# PRELIMINARY INFORMATION <br> ON THE <br> COMPENT OF REGISTERS OF THE <br> 1103A FLOATING POINT INSTRUCTIONS 

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The layout of the "Floating Point Content of Registers" is not the same as that of the fixed point instructions. There are more conditions affecting the final content of A. First, has the NEFF been set or cleared by instruction 05? Second, what is the relative size of (u) and (v)? For these reasons, only the Pack and Unpack commands are in the usual format. Since the arithmetic Floating Point comands do not change (u) and (v), $(u)_{f}$ and $(v)_{f}$ are not ineluded in the Contents of Registers of operations $64,65,66,67,01$, and 02.

The binary point of floating point numbers is usually between the twentyeighth and the twenty-seventh place. After the arithmetic pseudo-normalizing process, the mantissa is in $A_{L}$, and the binary point is between $A_{63}$ and $A_{62}$. It may or may not be normalized. The position of the most significant bit (MSB) indicates what has occurred. If normalized, the MSB will be in A $_{62}{ }^{\circ}$

The value of the significant bits depends upon whether rounding has occurred. Rounding in effect adds an extra bit to the value of (a) at $\mathbb{A}_{35}$ (unless the addition of the rounding bit carries into $A_{62}$, in which case the final left shift is omitted and the rounding bit remains added to the value of $A_{34}{ }^{\circ}$

The value of $(Q)_{f}$ will be either (1) the normalized rounded, and packed result (MRP), or (2) the pseudo-normalized result (PN).

HOTE: If $A$ or $Q$ is the $v$-address of any floating point command other than the pack or unpack command (A) or (Q) will be destroyed by the Unpack $(u)$ sequence before the unpack ( $v$ ) sequence is reached.

| Instruction: Floating Add (FAuv) |  |  |  | Operation: 64 |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Function: Form in $Q$ the normalized rounded and packed floating point sum of ( $u$ ) and ( $v$ ). |  |  |  |  |  |  |  |
| HES | Arithmetic Conditions |  | $(\mathrm{A})_{\underline{P}}$ |  |  | $(\mathrm{Q})_{1}$ |  |
| FF |  |  | MSB | Value of significant bits | Round | Norm | Value |
| 0 | $\begin{aligned} & (u) \geq(v) \\ & (u)<(v) \end{aligned}$ |  | $\begin{aligned} & A_{62} \\ & A_{62} \end{aligned}$ | $\begin{aligned} & \left(u_{m}\right) \cdot 2^{\left(v_{c}\right)-\left(v_{c}\right)_{+}\left(v_{m}\right)} \\ & \left(v_{m}\right) \cdot 2^{\left(v_{c}\right)-\left(u_{c}\right)}+\left(v_{m}\right) \end{aligned}$ | yes <br> yes | RRP <br> MRP | $\begin{aligned} & (u)+(v) \\ & (u)+(v) \end{aligned}$ |
| 1 | $(u) \geq(v)$ $(u)<(v)$ | $\left\|\begin{array}{l} \left(u_{c}\right)-\left(v_{c}\right) \geq 2 \\ \left(u_{c}\right)-\left(v_{c}\right)<2 \\ \left(v_{c}\right)-\left(u_{c}\right) \geq 2 \\ \left(v_{c}\right)-\left(u_{c}\right)<2 \end{array}\right\|$ | $\left\{\begin{array}{l} A_{61} \\ A_{61}-A_{33} \\ A_{61} \\ A_{61}-A_{33} \end{array}\right.$ | $\begin{aligned} & \left(v_{m}\right) .2^{\left(u_{c}\right)-\left(v_{c}\right)}+\left(v_{m}\right) \\ & \left(u_{m}\right) .2^{\left(u_{c}\right)-\left(v_{c}\right)_{+}\left(v_{m}\right)} \\ & \left(v_{m}\right) \cdot 2^{\left(v_{c}\right)-\left(u_{c}\right)_{+}}\left(u_{m}\right) \\ & \left(v_{m}\right) \cdot 2^{\left(v_{c}\right)-\left(v_{c}\right)}+\left(v_{m}\right) \end{aligned}$ | no <br> no <br> no <br> no | PII <br> PN <br> PN <br> PN | $\begin{aligned} & (u)+(v) \\ & (u)+(v) \\ & (u)+(v) \\ & (u)+(v) \end{aligned}$ |





| Instruction: Floating Point Polynomial Multiply (PPuv) Operation: 01 |  |  |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Function: Form in $Q$ the sum of $(v)$ and the product of $(Q)_{i}$. $(u)$ <br> (IRE FF should be cleared for the execution of this instruction. If it is not the product mantissa will be rounded not with one, but with ( $A_{L}$ ).) |  |  |  |  |  |  |  |  |  |
| NE | Arithmetic Condition |  | $(\mathrm{A})_{\mathrm{f}}$ |  |  |  |  | $(Q)_{f}$ |  |
| FF |  |  | MSB | Value of | ntissa |  | Hound | Norm | Value |
| 0 | (Q) $(u) \geq(\nabla)$ <br> (Q) $(u)<(v)$ |  | $A_{62}$ | $\begin{aligned} & (Q u)_{m} \cdot 2^{(Q u)_{c}-\left(v_{c}\right)_{+}+\left(v_{m}\right)} \\ & \left(v_{m}\right) \cdot 2^{\left(v_{c}\right)-(Q u)_{c}}+(Q u)_{m} \end{aligned}$ |  |  | yes <br> yes | NRP | $(Q)_{i}(u)+(v)$ |


| Instruction: Floating Point Inner Product (PIuv) |  |  |  | Operation: |  |  | 02 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Function: Form in $Q$ the normalized, rounded and packed sum of $\left(Q_{1}\right)$ and the product of ( $u$ ) and ( $v$ ). <br> (NE FF should be cleared for the execution of this instruction; if it is not, the product one, but with $\left(A_{L}\right)$ ) |  |  |  |  |  |  |  |
| HE | Arithmetic <br> Condition | $(\mathrm{A}){ }_{f}$ |  |  |  |  | $(\mathrm{Q})_{\mathrm{f}}$ |
| FF |  | MSB | Value of significant Round |  |  | Horm | Value |
| 0 | $\begin{aligned} & (Q)_{i} \geq(u)(v) \\ & (Q)<(u)(v) \end{aligned}$ | $\begin{aligned} & A_{62} \\ & A_{62} \end{aligned}$ | $\begin{aligned} & \left(Q m_{m}\right)_{i} 2^{(Q c)_{i}-(u v)_{c_{+}(u v)_{m}}} \\ & \left(u v_{m}\right) \cdot 2^{\left.(u v)_{c}-(Q c)_{i+(\text { (Qm }}\right)_{i}} \end{aligned}$ |  | $\begin{array}{r} \text { yes } \\ \text { yes } \end{array}$ |  | $\begin{aligned} & (Q)_{i}+(v \\ & (Q)_{i}+(v \end{aligned}$ |



Function: Replace (u) with the normalized rounded packed floating point number obtained from the possibiy unnormalized mantissa in (u) ${ }_{i}$ and the biased characteristic in $(v)_{c}$.

It is assumed that $(u)_{i}$ has the binary point between $u_{27}$ and $u_{26}$ $\left((u)_{i}\right.$ is scaled $\left.2^{-27}\right)$.

| Storage <br> Class |  | Contents of Registers \& Storage Position After Operation |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | $(\mathrm{MC})_{c}$ or | (MD) $^{1}$ |  | $(\mathrm{A}){ }_{\mathrm{f}}$ |  | $(Q)$ |
| $\mathfrak{u}$ | $\checkmark$ | u | $\nabla$ | MSB | Value of bits | Round |  |
| $\begin{aligned} & \text { MC } \\ & \text { or } \\ & \text { MD } \end{aligned}$ | MD or MC <br> A <br> $Q$ | $\begin{aligned} & \operatorname{RRP}(u)+\left(v_{c}\right) \\ & \operatorname{RRP}(u)+\left(v_{c}\right) \\ & \operatorname{RRP}(u)+\left(v_{c}\right) \end{aligned}$ | No change | ${ }^{A_{62}}$ <br> ${ }^{A_{62}}$ <br> ${ }^{A_{62}}$ | $\begin{aligned} & \left(u_{m}\right)_{f} \\ & \left(u_{m}\right)_{f} \\ & \left(n_{m}\right)_{f} \end{aligned}$ | yes <br> yes <br> yes | No change <br> Ho change <br> No change |
| A | MC <br> MC <br> A <br> Q |  | No change | $\begin{aligned} & A_{34} \\ & A_{34} \\ & A_{34} \end{aligned}$ | $\begin{aligned} & \operatorname{RRP}\left(A_{R}\right)_{i}+\left(\nabla_{c}\right) \\ & \operatorname{RRP}\left(A_{R}\right)_{i}+\left(A_{R c}\right)_{i} \\ & \operatorname{RRP}\left(A_{R}\right)_{i}+\left(Q_{C}\right) \end{aligned}$ | $\begin{aligned} & \text { yes } \\ & \text { yes } \\ & \text { yes } \end{aligned}$ | Ho change <br> Ho change <br> No change |
| Q | MD or MC <br> A <br> $Q$ |  | No change | $\begin{aligned} & A_{62} \\ & A_{62} \\ & A_{62} \end{aligned}$ | $\begin{aligned} & \left(u_{m}\right)_{f} \\ & \left(n_{m}\right)_{f} \\ & \left(u_{m}\right)_{f} \end{aligned}$ | $\begin{aligned} & \text { yes } \\ & \text { yes } \\ & \text { yes } \end{aligned}$ | $\begin{aligned} & \operatorname{RRP}(Q)_{i}+\left(\nabla_{c}\right) \\ & \operatorname{HRP}(Q)_{i}+\left(A_{R c}\right)_{i} \\ & \operatorname{HRP}(Q)_{i}+\left(Q_{c}\right)_{i} \end{aligned}$ |


| Instruction: | Floating Point Normalize Exit (HEJ) | Operation: | 05 |
| :---: | :---: | :---: | :---: |
| Function: | If $J=0$ clear the normalize exit flip-ilop (designated HFF); if $j-1$ set LIFF to 1 <br> (a) The results of setting NFF te 1 is set forth in the Contents of Registers ${ }^{n}$ <br> (b) Then EFF in set to $\pi$, it will remain set until cleared by another REJ - intaruction <br> (c) Hif must be cleared for FP, FI, and instructions |  |  |

