TIITE
TYPE
NUMBER OF WORDS
TEMPORARY STORAGE
DURATION
ACCURACY
ENITY

PRESET PARAMETERS

DESCRIPTION

Chi-Squared
Closed with two program parameters
23
0,1 , and 2
( $2.4 \mathrm{~m}+3$ ) milliseconds
$\pm 2^{-40}$
When this routine is located at $q$, entry is made by

| p | -mF |
| :--- | :--- |
|  | 50 pF |
| $\mathrm{p}+1$ | 26 qF |
| $\mathrm{p}+2$ | 00 nF |

At the end of this routine control is transferred to the left hand order at location $p+2$.
When this routine is read in the following preset parameters must have been stored in locations 3, 4 and 5, respectively.

S3 $00 \mathrm{~F} \quad \mathrm{~s}$ is a scaling factor and is normally 00 sF chosen to be $1,10,100$, or 1000 depending on whether none, one, two, or three decimal places of accuracy are desired after the decimal point.

S4 $00 \mathrm{~F} \quad$ a is the location at which the quantities
$00 \mathrm{aF} \quad p_{i}, i=0,1, \ldots, m-1$, are stored prior to entering this routine.
S5 $00 \mathrm{~F} \quad \mathrm{~b}$ is the location at which the quantities
$00 \mathrm{bF} \quad \mathrm{f}_{\mathrm{i}}$, $i=0,1, \ldots, m-1$, are stored prior to entering this routine.
In its most frequent application chi-squared is given by the formula

$$
X^{2}=\sum_{i=0}^{m-1} \frac{\left(E_{i}-\theta_{i}\right)^{2}}{E_{i}}
$$

where each $E_{i}$ is the expected number of members in the $i^{\text {th }}$ of $m$ classes for a given sample size and the $\theta_{i}$ are the observed values. If we let

$$
\left.p_{i}=\frac{E_{i}}{n} ; \quad f_{i}=\frac{\theta_{i}}{n} \quad 2 a, b\right)
$$

and multiply both sides of equation l) by a number $s$, the resulting equation is

$$
s \chi^{2}=\operatorname{sn} \sum_{i=0}^{m-1} \frac{\left(p_{i}-f_{i}\right)^{2}}{p_{i}}
$$

This last equation corresponds to the quantity computed by this routine. $s \chi^{2}$ is computed as an integer and is placed in the A register at the end of the routine. The quantities $p_{i}$ and $f_{i}$ are fractions and must be in the ranges

$$
\begin{aligned}
& 0<p_{i}<1 \\
& 0 \leq f_{i}<I
\end{aligned}
$$

A value of 1 for one of the $f_{i}$ may be represented in the machine as -1. Each value $p_{i}$ is stored at location $a+i$ before this routine is entered, the first address a being given by preset parameter S4. Similarly the $f_{i}$ are stored at $b+i, b$ being specified by preset parameter $S 5$. The number of values $m$ of the $p_{i}$ or $f_{i}$ is specified by a program parameter (See Entry). The number $s$ is a positive integer specified by preset parameter S3 and serves as a scaling factor for $\chi^{2}$. Normally $s$ will be chosen to be 1, 10, 100
or 1000 depending upon the number of decimal places of accuracy required in the value of $\chi^{2}$. For example if $\mathcal{Z}^{2}$ is $2.531 \ldots$ and $s$ is 100 , the number in the $A$ register at the end of this routine will be $253 \times 2^{-39}$.*

The program parameter $n$ is also a positive integer and is chosen so as to put the $p_{i}$ and $f_{i}$ in the required range given by $4 \mathrm{a}, \mathrm{b}$ ) as determined by equations $2 \mathrm{a}, \mathrm{b}$ ). A logical choice for $n$ is the sample size. The $p_{i}$ are then the predicted probabilities and the $f_{i}$ are the corresponding observed values. The requirements $4 a, b$ ) will then be satisfied automatically. If the values $E_{i}$ and $\theta_{i}$ are known directly it may be more convenient to chose $n$ to be the smallest power of 10 which is greater than or equal to the sample size. In some applications a power of 2 may be more convenient. In any case the values $p_{i}$ and $f_{i}$ are determined by equations $\left.2 a, b\right)$, and n must be specified upon entering the routine (See Entry).

If it is much more convenient to produce the values $p_{i}$ and $f_{i}$ one at a time, it is possible to enter this routine $m$ times with the new values of $p_{i}$ and $f_{i}$ in locations $a$ and $b$, respectively, each time. If this procedure is used the program parameter m must always be 1 , and the necessary summation must be carried on outside this routine. Other things being equal, this method will be slower and less accurate.

In addition to $4 \mathrm{a}, \mathrm{b}$ ) there are certain other moderate requirements on the quantities involved. The first of these is that the following inequality must hold:

$$
\frac{\chi^{2}}{n}=\sum_{i=0}^{m-1} \frac{\left(p_{i}-f_{i}\right)^{2}}{p_{i}}<256=2^{8}
$$

* If Library Routine P 1 is used the results may be printed in such a way that the position of the decimal point is specified.

Since each term in this summation is non-negative, each term must also be less than 256. If the latter does not hold a division hang-up will occur. If the summation is too large the answer will be in error by a negative integral multiple of 5l2s. Any danger that the requirement 5) might be violated can usually be avoided by using a larger value of n and making the appropriate changes in the $p_{i}$ and $f_{i}$ as determined by 2a, b). Secondly, in order to obtain the stated accuracy of $\pm 2^{-40}$ the product mns should be small compared to $2^{30} \approx 10^{9}$. Although this generally means that values of $\chi^{2}$ accurate to a large number of decimal places are obtainable, printing out results to greater accuracy than actually required or justified by the data should be assiduously avoided.

|  |
| :---: |
| Rt: $7 / 22 / 59$ <br> RETYPED <br> DATE 6/2/55: $5 / 31 / 56$ |
| PROGRAMMED BY C.Farrington APPROVED BY |


| LOCATION | ORDER |  | NOTES PAGE 1 |
| :---: | :---: | :---: | :---: |
| 0 | 41 F |  |  |
| 1 | K5 F |  |  |
|  | 425 L |  | Plant address of $n$ |
|  | 46 F |  | Store m |
| 2 | L5 19L |  | Set addresses |
|  | 408 L |  |  |
| 3 | $4611 L$ |  |  |
|  | L4 F |  | Set test constant |
| 4 | 46.20 L |  |  |
|  | F5 5L |  | Plant link |
| 5 | 42 18L |  |  |
|  | 50 F |  | Extract and store $n$ |
| 6 | 001043 F |  |  |
|  | 01.20 F |  |  |
| 7 | 40 F |  |  |
|  | 41 F |  | Clear $\Sigma$ |
| 8 | L5 S4 | from 15 | $p_{i}-f_{i}$ |
|  | L0. 55 |  |  |
| 9 | 402 F |  | $\left(p_{i}-f_{i}\right)^{2} \times 2-8$ |
|  | 502 F |  | $\frac{1}{p_{i}} \times 2$ |
| 10 | 752 F |  |  |
|  | 1085 |  |  |
| 11 | 66.54 |  |  |
|  | S5 F |  | Summation |
| 12 | L4 1 F |  |  |
|  | $401 F$ |  |  |
| 13 | I.5 8L |  | Step addresses |
|  | $\text { F4 } 211$ |  |  |
| 14 | 408 L |  |  |
|  | 46112 |  |  |
| 15 | L0 20L |  | Test for end |
|  | 368 L |  |  |
| 16 | 50 F |  | ns $\times 2^{\text {m }} 39$ in $Q^{4}$ |
|  | 7522.4 |  |  |



