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IBM System/360 Operating System: Programmer's Guide to Debugging

OS Release 21.7





Seventh Edition (November, 1972)

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GN28-2520 (dated April 15, 1972)

This edition with Technical Newsletter GN28-2545 applies to release 21.7 and component release 360S-OS-586 of IBM System/360 Operating System, and to all subsequent releases until otherwise indicated in new editions or Technical Newsletters. Changes are continually made to the information herein; before using this publication in connection with the operation of IBM systems, consult the latest IBM System/360 and System/370 Bibliography, GA22-6822, for the editions that are applicable and current.

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This publication is intended to help you use the debugging facilities provided with the IBM System/360 Operating System. It describes, in assembler language terms, the major debugging facilities provided with the System/360 Operating System, and is directed towards the programmer who deals with system and application program problems.

The publication is divided into three principal parts: "Section 1: Operating System Concepts;" "Section 2: Interpreting Dumps;" and "Section 3: Tracing Aids," plus an Introduction and a set of Appendixes that provide specific debugging information.

The Introduction provides a brief survey of the material presented in the balance of the publication.

Section 1 deals with internal aspects of the operating system that are pertinent to debugging. A working knowledge of this information will provide you with the means of determining the status of the system at the time of failure, and the course of events which led up to that failure. The general precedure for debugging with an operating system dump (Appendix A) assumes knowledge of this control flow.

Section 2 includes instructions for invoking, reading, and interpreting storage dumps of systems with MFT or MVT control programs. The material is intended to aid you in interpreting dumps and isolating errors.

Section 3 deals with the save area chain, the Trace Option, and the Generalized Trace Facility. Output from the Generalized Trace Facility is discussed.

Before reading this publication, you should have a general knowledge of operating system features and concepts as presented in the prerequisite publications. Occasionally, the text refers you to other publications for detailed discussions beyond the scope of this book. For information on debugging facilities provided within higher languages, consult the programmers' guides associated with the respective languages. Other System/360 Operating System publications, such as <u>Messages and Codes</u>, describe additional debugging aids provided for the assembler language programmer.

Notice: Coding level information presented in this publication must not be used for coding purposes or exposure to changes in implementation may result. The information is presented for debugging purposes only.

PREREQUISITE PUBLICATIONS

IBM System/360: Principles of Operation, GA22-6821

IBM System/360 Operating System:

Supervisor Services and Macro Instructions, GC28-6646

Data Management Services, GC26-3746

REFERENCE PUBLICATIONS

IBM System/360 Operating System:

System Control Blocks, GC28-6628

Messages and Codes, GC28-6631

Data Management Macro Instructions, GC26-3794

Service Aids, GC28-6719

TCAM Programmer's Guide and Reference, GC30-2024.

TCAM Serviceability Aids, GY30-2027.

TCAM, GY30-2029.

TSO Control Program, GY27-7199.

4 Programmer's Guide to Debugging (Release 21)

Contents

SUMMARY OF AMENDMENTS FOR GC28-6670-5	
OS RELEASE 21	9
SUMMARY OF AMENDMENTS FOR GC28-6670-4	
AS UPDATED BY GN28-2457 AND GN28-2472	
	LO
SUMMARY OF AMENDMENTS FOR GC28-6670-3	
	10
INTRODUCTION	11
	13
	13
	13
	13
Active RB Queue	16
	L6
Job Pack Area Queue (MFT With	
Subtasking Only)	17
Effects of LINK, ATTACH, XCTL, and	
	18
	19
Systems With MFT (Without	
	19
,, _,, _	21
Main Scorage Supervision	6 J
Storage Control in Systems with MFT	14
	21
Storage Control in Systems with	~ ~
	22
Storage Control for a Region in	
	22
Storage Control for a Subpool in	
	24
Storage Control for a Load Module	
in Systems With MVT	
System Control Blocks and Tables	26
	26
Task Input/Output Table (TIOT)	26
	27
Event Control Block (ECB)	
Input/Output Block (IOB)	57
Data Control Block (DCB)	27
Data Extent Block (DEB)	
Summary of Control Block	21
Relationships	21
SECTION 2: INTERPRETING DUMPS	20
	29
	29
Invoking an ABEND/SNAP Dump (MFT) .	29
Contents of an ABEND/SNAP Dump	
(MFT) Guide to Using an ABEND/SNAP Dump	32
Guide to Using an ABEND/SNAP Dump	
	44
	46
	46
Contents of an ABEND/SNAP Dump	
	46
(MVT)	-0
(MUT)	
	ເລ
Indicative Dump	63 65
Indicative Dump	63 65 65

Guide to Using an Indicative Dump	. 67
Guide to Using an Indicative Dump Storage Dumps	- 68
Damage Assessment Routine (DAR)	. 68
Console Dump	. 68
Console Dump	. 68
System Failure	. 68
The SYS1.DUMP Data Set	. 68
Tape	. 68
Direct Access	. 69
IMDPRDMP Output Queue Control Block Traces Link Pack Area Maps Major System Control Block Formats .	. 70
Queue Control Block Traces	• 70
Link Pack Area Maps	. 71
Major System Control Block Formats .	. 75
MVT Control Block Formatting	• 75
MFT Control Block Formatting	• 92
TSO System Block Formatting Task Control Block Summaries	.108
Task Control Block Summarles	.124
The General Format	.128
Cuide to Charge Dump	-130
Guide to Storage Dumps	120
Tack Structure	120
Task Structure	130
MFT System (With Subtasking)	.139
MVT System	.139
MVT System	.141
Main Storage Contents	.142
Load List (MFT)	.142
Load List (MVT)	.142
Load List (MFT) Load List (MVT) Job Pack Area Queue (MFT With	
Subtasking, MVT)	.142
Subtasking, MVT)	.143
Free Areas in MFT Systems	.143
Gotten Subtask Areas (MFT)	.143
Region Structure in MVT System	.143
I/O Control Blocks	.143
I/O Control Blocks	.143
UCBs	.144
DCB and TIOT	.144
IOB	.144
ECB	.145
TSO Control Blocks	.145
ECB TSO Control Blocks	.145
RCB	.146
UMSM	.146
SWAP DCB	.146
	.146
	.146
	.146
	.146
TAXE	.146
CROMTON 2. MOLOTNO LING	147
	.147
	.147
	.148 .148
	.148
	.150
IO and PCI/IO Minimal Trace Record	- 151
SIO Minimal Trace Record	
	.153
EXT Minimal Trace Record	.154

PGM Minimal Trace Record	
SVC Minimal Trace Record	
SSM Minimal Trace Record	
GTF Comprehensive Trace Records158	A
IO and PCI/IO Comprehensive Trace	C
Record	
SIO Comprehensive Trace Record161	
DSP Comprehensive Trace Record162	
EXT Comprehensive Trace Record163	7
PGM Comprehensive Trace Records165	-
SSM Comprehensive Trace Record166	Į
Time and Lost Event Records	•
Hexadecimal Format Record	7
	Ĩ
GTF SVC Comprehensive Trace Records .169	1
SVC Comprehensive Trace Records	
Group 1 Basic Fields	Į
SVC Comprehensive Trace Records	
Group 1 Basic Fields	7
SVC Comprehensive Trace Records	
Group 2 - Basic Fields Plus DDNAME	1
Field	
SVC Comprehensive Trace Records;	1
Group 3 - Basic Fields Plus	
Parameter List Field	1
SVC Comprehensive Trace Records;	-
Group 4 - Basic Fields Plus	7
Variable Fields	•
	1
	•

IMDPRDMP Output Comments - GTF Processing
APPENDIX A: DEBUGGING WITH AN
OPERATING SYSTEM DUMP
Specialized Program Checks
Debugging Procedure Summary207
APPENDIX B: SVCs
APPENDIX C: COMPLETION CODES
APPENDIX D: SYSTEM MODULE NAME
PREFIXES
APPENDIX E: LIST OF ABBREVIATIONS221
APPENDIX F: ECB COMPLETION CODES223
APPENDIX G: UCB SENSE BYTES
APPENDIX H: SERVICE AIDS
APPENDIX J: TCAM DEBUGGING AIDS231
APPENDIX K: CONTROL BLOCK POINTERS233
APPENDIX L: OPEN/CLOSE/EOV DEBUGGING
AIDS
INDEX

Figure 1. Control Information	Figure 27. Sample of General Format
Available Through the TCB	Dump
Figure 2. RB Formats 15	Figure 28. Permanently Assigned
Figure 3. Active RB Queue 16	Hardware Control Words
Figure 4. Load List (MFT) 16	Figure 29. Finding the Partition TCBs
Figure 5. Job Pack Area Queue 18	in MFT
Figure 6. Main Storage Snapshot (MFT	Figure 30. Finding the TCB
Without Subtasking)	Figure 31. IMDPRDMP TCB Summary141
Figure 7. Partition (MFT Without	Figure 32. Determining Module From
Subtasking)	CDE in MVT
Figure 8. Main Storage Snapshot (MFT	Figure 33. Subpool Descriptions in
With Subtasking) 20	MVT - IMDPRDMP Storage Print 144
Figure 9. Main Storage Snapshot	Figure 34. I/O Control Blocks 145
(MVT)	Figure 35. Save Area Trace
Figure 10. Storage Control for a	Figure 36. Trace Table Entries (MFT) .148
Partition (MFT Without Subtasking) 22	Figure 37. Trace Table Entries (MVT) .148
Figure 11. Storage Control for Subtask	Figure 38. Trace Table Entries (MVT)
Storage (MFT With Subtasking) 22	with Model 65 multiprocessing)149
Figure 12. Storage Control for a	Figure 39. Sample Trace Table Entries
Region (MVT)	(MFT)
Figure 13. Storage Control for a	Figure 40. Sample Trace Table Entries
Subpool (MVT)	(MVT)
Figure 14. Storage Control for a Load	Figure 41. IO and PCI/IO Minimal
Module (MVT)	Trace Record
Figure 15. Control Block	Figure 42. SIO Minimal Trace Record .152
Relationships	
Figure 16. Sample of an ABEND Dump	
(MFT) (Part 1 of 2)	Figure 45. PGM Minimal Trace Record .155 Figure 46. SVC Minimal Trace Record .156
Figure 17. SYSABEND DD Statements 33	
Figure 18. Sample of Complete ABEND	Figure 47. SSM Minimal Trace Record .157
Dump (MVT) (Part 1 of 2) 47	Figure 48. IO and PCI/IO
Figure 19. Contents of an Indicative	Comprehensive Trace Record
Dump	Figure 49. SIO Comprehensive Trace
Figure 20. Queue Control Block Trace	Record
Sample	Figure 50. DSP Comprehensive Trace
Figure 21. Link Pack Area Map Sample . 73	Record
Figure 22. Sample of MVT Major	Figure 51. EXT Comprehensive Trace
Control Block Format	Record
Figure 23. Sample of MFT Control	Figure 52. PGM Comprehensive Trace
Block Format	Record
Figure 24. Sample of TSO Control	Figure 53. SSM Comprehensive Trace
Block Format (Part 1 of 3)	Record
Figure 25. TSB Summary Sample for	Figure 54. Hexadecimal Format Record .168
System That Operated Under MVT or MFT	Figure 55. Basic SVC Comprehensive
With Subtasking	Trace Record
Figure 26. TCB Summary Sample for	Figure 56. Control Block Flow 237
Systems that Operated Under MFT	Figure 57. MVT Storage Control Flow .239
Without Subtasking	

8 Programmer's Guide to Debugging (Release 21)

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Summary of Amendments

Summary of Amendments for GC28-6670-5,-6 as Updated by GN28-2545 OS Release 21.7

<u>Model 91 or 195 Imprecise Interruption</u> An explanation of the bits set by an imprecise interruption on the Model 91 or 195 has been added.

AMDPRDMP Format Changes Some additional fields have been added to the JOB information and the TCB.

AMDPRDMP Messages Several new messages can be put out by AMDPRDMP to aid in debugging.

GTF Records

The DSP, EXT, PRM, and SVC comprehensive trace records have had a parameter added to the MODN field.

Summary of Amendments for GC28-6670-5 OS Release 21

PCP REMOVAL

References to the PCP version of Operating System/360 have been deleted from the publication.

TESTRAN REMOVAL

References to the TESTRAN testing facility of Operating System/360 have been deleted from the publication.

IMDPRDMP SERVICE AID OUTPUT

Storage dumps as formatted and displayed by the IMDPRDMP service aid are now discussed in this publication. This material was formerly in the Service Aids publication, GC28-6719.

GENERALIZED TRACE FACILITY (GTF) OUTPUT

GTF trace records, as processed by the EDIT function of the IMDPRDMP service aid are illustrated and discussed in Section 3 of the publication.

DEVICE SUPPORT

The sense byte information given in Appendix G is updated to include information for the:

IBM 3420 Magnetic Tape Unit and 3803 Tape Control The explanation for SVC 79 (STATUS) has been changed.

<u>New Emulator</u> The 7074 emulator can be run on the Model 155.

<u>New I/O Devices</u> Sense byte information has been added for the 3213, 3215, 3272, 3277, and 3333 as well as additional support for the 3410/3411.

<u>Miscellaneous Corrections</u> Several additional changes have been made to correct errors.

IBM 2596 Card Read Punch IBM 3505 Card Reader IBM 3525 Card Punch IBM 3410 Magnetic Tape Unit IBM 3411 Magnetic Tape Unit and Control

PROBLEM DETERMINATION

Addition of an Appendix discussing problem determination aids for OPEN/CLOSE/EOV processing.

Updating of the completion codes and service aids Appendixes to reflect release 21 changes.

The Console Dump facility, used to obtain a storage dump for later processing by IMDPRDMP, is briefly described in the storage dump and IMDPRDMP formatting section of the publication.

MISCELLANEOUS

Editorial improvements and corrections to existing material have been made throughout the publication. Summary of Amendments for GC28-6670-4 as Updated by GN28-2457 and GN28-2472 OS Release 20.1

TCAM

3330, 2305, 2319

new devices.

Appendix F

Section 2: ABEND/SNAP Dump (PCP and MFT) ABEND/SNAP Dump (MVT) Appendix A Appendix H A brief description of TCAM debugging Aids and a new SVC.

<u>TSO</u>

MISCELLANEOUS

Section 2: TSO Control Blocks <u>Appendix A</u> The addition of new SVCs and a summary of the control blocks formatted by IMDPRDMP. <u>Appendix C</u> 1. Addition of module name prefixes for emulator programs.

Appendix G

2. New features of service aid program IMAPTFLE.

Additional of sense byte information fo

Summary of Amendments for GC28-6670-3 OS Release 20

IMDPR DMP

<u>TSO</u>

"Guide to Using a Storage Image Dump" IMDPRDMP is used instead of IEAPRINT to print MFT and MVT dumps. <u>Appendix A</u> New SVCs in Appendix A. This information is for planning purposes only. To debug efficiently, you should be familiar with the system control information reflected in dumps. This control information, in the form of control blocks and traces, tells you what has happened up to the point of error and where key information related to the program is located. To provide an insight into the IBM System/360 Operating System and its complex aspects of task management and storage supervision, Section 1 of this publication provides an orientation in the control functions of the operating system.

The IBM System/360 Operating System provides extensive debugging facilities to aid you in locating errors and determining the system state quickly. Some debugging aids, such as console messages, provide limited information that may not always help you identify the error. This manual discusses those debugging facilities that provide you with the most extensive information:

- a. Abnormal termination (ABEND) and snapshot (SNAP) dumps.
- b. Indicative dumps.
- c. Storage image dumps.
- d. Tracing facilities.

Dumps are discussed in Section 2 and tracing facilities in Section 3.

<u>ABEND and SNAP Dumps</u> are invoked by ABEND and SNAP macro instructions, respectively. They are grouped in a single category because they provide identical information. In addition to a hexadecimal dump of main storage, they can contain converiently edited control information and displays of the operating system nucleus and trace table.

<u>Indicative dumps</u> contain control information useful in isolating the instruction that caused an abnormal end of task situation. The information is similar to that given in an ABEND/SNAP dump, but does not include a dump of main storage.

<u>Storage dumps</u> are produced by either the system dump facility at the time of a system failure, or by a dump program created through use of the IMDSADMP service aid. IMDSADMP programs must be loaded into storage through use of the IPL facilities and are intended for use in situations in which the system is not operative, e.g., a disabled wait state or an unending system loop.

The system dump facility writes to the SYS1.DUMP data set. The IMDPRDMP service aid is used to format and print the SYS1.DUMP data set. IMDPRDMP output is described in this publication. The IMDSADMP programs write to tape (high-speed dump) or to tape or printer (low-speed dump). The output tape produced by the high-speed dump must be processed by the IMDPRDMP program; low-speed output to tape may be processed by IMDPRDMP, IEBPTPCH or the IEBGENER utility program.

Storage dumps taken by the system dump facility consist of control information followed by a display of printable storage from location 00 to the capacity of storage. Storage words are displayed in both hexadecimal and EBCDIC notation. Storage dumps taken by an IMDSADMP program consist of register contents followed by a display of storage from location 00 to the capacity of storage. Notation is in both hexadecimal and EBCDIC.

Tracing facilities consist of the save area chain trace, the Trace Option and the Generalized Trace Facility.

The save area chain enables tracing of the save areas for each level of load module in a task. The save area trace is displayed in ABEND/SNAP and storage dumps.

The Trace Option, if installed in the system, provides records of system interruptions (IO, SIO, etc.) that are displayed in ABEND/SNAP and storage dumps.

The Generalized Trace Facility (GTF) enables selective tracing of system and application program events and records the information internally, in a table which is displayed in printouts of ABEND dumps and storage dumps, or externally in a data set which is processed by the IMDPRDMP service aid to provide edited and formatted GTF trace records. (For complete information on GTF see the Service Aids publication.) The GTF output, as processed by IMDPRDMP, is discussed in Section 3 of this publication.

Introduction 11

General Notes:

- Displacements and addresses shown in the text and illustrations of this publication are given in decimal numbers, followed by the corresponding hexadecimal number in parentheses, e.g., TCB+14(E); location 28(1C); SVC 42(2A). All other numbers in the text are decimal, e.g., the <u>seventeenth</u> word of the TCB; a <u>4</u>-word control block; <u>15</u> job steps.
- Control block field names referred to are those used in the <u>IBM System/360</u> <u>Operating System: System Control</u> <u>Blocks</u> manual, GC28-6628.
- Wherever possible, diagrams, and reproductions of dumps have been included to aid you during the debugging process.

Section 1: Operating System Concepts

This section introduces you to the control information that you must know to interpret dumps. It is divided into three topics:

- task management
- main storage supervision
- system control blocks and tables

The first two topics deal with those aspects of task management and main storage management, respectively, that are represented in dumps. The third topic describes the remaining system control blocks and tables helpful in pinpointing errors.

Note: The descriptions of system control blocks and tables in this section emphasize function rather than byte-by-byte contents. Appendix K summarizes the contents of those control blocks most useful in debugging.

For a more detailed description of system control blocks and tables, refer to the <u>System Control Blocks</u> publication, GC28-6628.

Task Management

The task management control information most useful in debugging with a dump includes the task control block and its associated request blocks and elements. The functions, interactions, and relationships to other system features of

these items are discussed in this topic. A summary of how task supervision differs at each system level concludes the topic.

Task Control Block

The operating system keeps pointers to all information related to a task in a task control block (TCB). For the most part, the TCB contains pointers to other system control blocks. By using these pointers, you can learn such facts as what I/O devices were allocated to the task, which data sets were open, and which load modules were requested.

Figure 1 shows some of the control information that can be located by using the pointers in the TCB. Later, in the discussion of system control blocks and tables, Figure 1 is expanded to show the actual block names and pointer addresses.

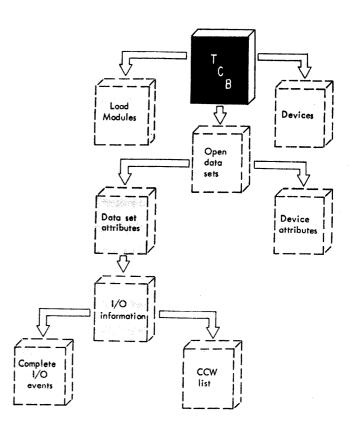


Figure 1. Control Information Available Through the TCB

Request Blocks

Frequently, the routines that comprise a task are not all brought into main storage with the first load module. Instead, they are requested by the task as it requires them. This dynamic loading capability necessitates another type of control block to describe each load module associated with a task -- a request block (RB). An RB is created by the control program when it receives a request from the system or from a problem program to fetch a load module for execution, and at other times, such as when a type II supervisor call (SVC) is issued. By looking at RBs, you can determine which load modules have been executed, why each lost control, and, in most cases, which one was the source of an error condition.

There are seven types of RBs created by the control program:

- Program request block (PRB)
- Supervisor request block (SVRB)
- Interrupt request block (IRB)

- Supervisor interrupt request block (SIRB)
- Loaded program request block (LPRB)
- Loaded request block (LRB)
- Finch request block (FRB)

Of these, you will most often encounter the PRB and SVRB in dumps. The type of RB created depends on the routine or load module with which it is associated.

<u>PRB (Systems with MFT)</u>: A PRB is created whenever an XCTL, LINK, or ATTACH macro instruction is issued. It is located ' immediately before the load module with which it is associated.

<u>PRB (Systems with MVT)</u>: A PRB is created whenever an XCTL or LINK macro instruction is issued. It is located in a fixed area of the operating system.

SVRB: An SVRB is created each time a type II, III, or IV supervisor call is issued. (Type I SVC routines are resident, but run disabled; they do not require a request block.) This block is used to store information if an interruption occurs during execution of these SVC routines. A list of SVCs, including their numbers and types, appears in Appendix A.

IRB: An IRB is created each time an asynchronous exit routine is executed. It is associated with an event that can occur at an unpredictable time during program execution, such as a timing routine initiated by an STIMER macro instruction. The IRB is filled at the time the event occurs, just before control is given to the exit routine.

<u>SIRB</u>: An SIRB is similar to an IRB, except that it is associated only with IBM-supplied input/output error routines. Its associated error routine is fetched from the SYS1.SVCLIB data set.

<u>LPRB (MFT only)</u>: An LPRB is created when a LOAD macro instruction is issued unless the LOAD macro instruction specifies:

- A routine that has already been loaded.
- A routine that is being loaded in response to a LOAD macro instruction previously issued by a task in the partition (MFT with subtasking).
- A routine that is "only loadable" (see LRB).

An LPRB is located immediately before the load module with which it is associated. Routines for which an LPRB is created can also be invoked by XCTL, LINK, and ATTACH macro instructions. LRB (MFT only): The LRB is a shortened form of an LPRB. Routines associated with LRBs can be invoked only by a LOAD macro instruction. This attribute is assigned to a routine through the OL (only loadable) subparameter in the PARM parameter of the EXEC statement that executes the linkage editor. The most common reason for assigning this attribute is that linkage conventions for XCTL, LINK, and ATTACH are not followed. This request block is located immediately before the load module with which it is associated.

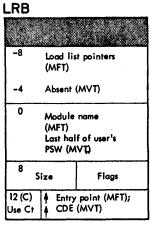
FRB (MFT with subtasking only): An FRB is created and attached to the job pack area queue, during LOAD macro instruction processing, if the requested module is not already in the job pack area. The FRB describes a module being loaded in response to a LOAD macro instruction. Any subsequent requests for the same module, received while it is still being loaded, are deferred by means of wait list elements (WLEs) queued to the FRB. When the module is fully loaded, an LRB or an LPRB is created, the FRB is removed from the job pack area queue, and any requests, represented by wait list elements, are reinitiated.

Figure 2 shows the relative size of the seven types of RBs and the significant fields in each.

In Figure 2, the "size" field tells the number of doublewords in both the RB and its associated load module. The PSW contained in the "resume PSW" field reflects the reason that the associated load module lost control. Other fields are discussed in succeeding topics.

This far, the characteristics of the TCB and its associated RBs have been discussed. With the possibility of many RBs subordinate to one task, it is necessary that queues of RBs be maintained. In systems with MFT without subtasking, two queues are maintained by the system -- the active RB queue and the load list. In MFT systems with subtasking, a job pack area queue, containing FRBs, and LRBs and LPRBs that represent reenterable modules is also maintained. MVT systems maintain an active RB queue and a contents directory. The contents directory is made up of three separate queues: the link pack area control queue (LPAQ); the job pack area control queue (JPAQ); and the <u>load list</u>.

LPRB				
-12		3 address th subtasking)		
-8	Load list pointers (MFT)			
-4	Absent (м∨т)		
0 Module name (MFT) Last half of user's PSW (MVT)				
8 Size Flags				
12(C) Use Ct		ry point (MFT); E (MVT)		
16 (10)				
Resume PSW				
28(IC) Wait Ct	14 N	Vext RB		



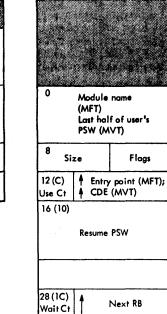
Program Extent List

hiearchy 1

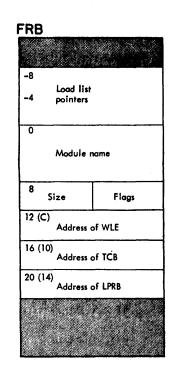
Length of extent in hiearchy 0

Length of extent in

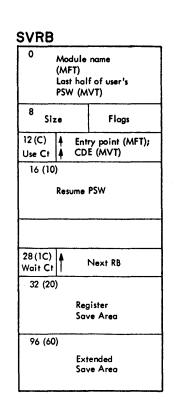
Address of extent in hiearchy 0 + 12(C) Address of extent in hiearchy 1



PRB



Note: Program extent list is added to LPRB, LRB, or PRB if the program described was hiearchy block loaded.

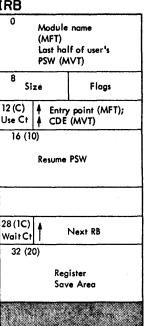




L±0.

+ 4

+ 8



SIRB 0 Module name (MFT) Last half of user's PSW (MVT) 8 Size Flags Entry point (MFT);
 CDE (MVT) 12 (C) Use Ct 16 (10) **Resume PSW** 28(1C) Next RB WaitCt 32 (20) Register Save Area

Figure 2. RB Formats

Active RB Queue

The active RB queue is a chain of request blocks associated with active load modules and SVC routines. This queue can contain PRBs, SVRBs, IRBs, SIRBs, and under certain circumstances, LPRBs. Figure 3 illustrates how the active RB queue links together the TCB and its associated RBs.

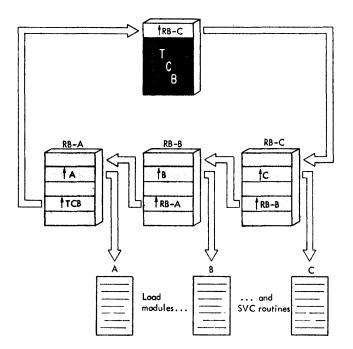


Figure 3. Active RB Queue

The request blocks in the active RB queue in Figure 3 represent three load modules. Load module A invokes load module B, and B, in turn, invokes C. When execution of A began, only one RB existed. When the first invoking request was encountered, a second RB was created, the TCB field that points to the most recent RB was changed, and A's status information was stored in RB-A. A similar set of actions occurred when the second invoking request was encountered. As each load module is executed and control is returned to the next higher level load module, its RB is removed from the chain and pointers are updated accordingly.

Load List

The load list is a chain of request blocks or elements associated with load modules invoked by a LOAD macro instruction. The load list differs from the active RB queue in that RBs and associated load modules are not deleted automatically. They remain intact until they are deleted with a DELETE macro instruction or job step termination occurs. By looking at the load list, you can determine which system and problem program routines were loaded before the dump was taken. The format of the load list differs with control program levels.

Systems with MFT (without subtasking): At this control program level, the load list associated with a TCB contains LRBs and LPRBs. RBs on the load list are linked together somewhat differently from those on the active RB queue because of the characteristics of the LOAD macro instruction. Because RBs may be deleted from a load list in a different order than they were created (depending on the order of DELETE macro instructions), they must have both forward and backward pointers. Figure 4 illustrates how a load list links together a TCB and three RBs.

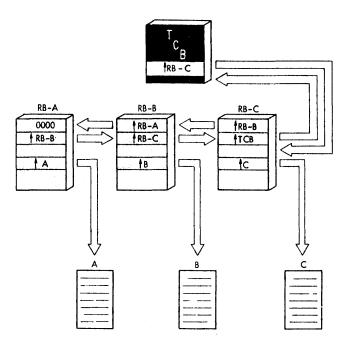


Figure 4. Load List (MFT)

Here, each RB contains a pointer both to the previous RB and the next most recent RB in the list. If there is no previous or more recent RB, these fields contain zeros and a pointer to the TCB, respectively.

Another field of a load list RB that merits consideration is the use count. Whenever a LOAD macro instruction is issued, the load list is searched to see if the routine is already loaded. If it is loaded, the system increments the use count by one and passes the entry point address to the requesting routine.

Each time a DELETE macro instruction is issued for the routine, the use count is decremented by one. When it reaches zero, the RB is removed from the load list and storage occupied by the associated routine is freed. Systems With MFT (With Subtasking): At this control program level, the load list is used as described for MFT without subtasking, with the following exceptions:

- The LRBs and LPRBs queued on the load list represent modules that are not reenterable. LRBs and LPRBs representing reenterable modules are queued on the job pack area queue.
- 2. When a LOAD macro instruction is issued, the system searches the job pack area queue before searching the load list.

Systems With MVT: Instead of LRBs and LPRBs created as a result of LOAD macro instructions, the load list maintained by a system with MVT contains elements representing load modules. Load list elements (LLEs) are associated with load modules through another control medium called the <u>contents directory</u>.

The contents directory is made up of three separate queues: the <u>link pack area</u> <u>control queue (LPAQ)</u>, the job pack area <u>control queue</u> (JPAQ), and the <u>load list</u>.

The LPAQ is a record of every program in the system link pack area. This area contains reenterable routines specified by the control program or by the user. The routines in the system link pack area can be used repeatedly to perform any task of any job step in the system. The entries in the LPAQ are contents directory entries (CDEs).

There is a JPAQ for each job step in the system that uses a program not in the link pack area. The JPAQ, like the LPAQ, is made up of CDEs. It describes routines in a job step region. The routines in the job pack area can be either reenterable or not reenterable. These routines however, cannot be used to perform a task that is not part of the job step.

The load list represents routines that are brought into a job pack area or found in the link pack area by the routines that perform the Load function. The entries in the load list are load list elements, not CDEs. Each load list element is associated with a CDE in the JPAQ or the LPAQ; the programs represented in the load list are thus also represented in one of the other contents directory queues.

Load list elements also contain a count field that corresponds to the use count in a LPRB or LRB. Each time a LOAD macro instruction is issued for a load module already represented on the load list, the count is incremented by one. As corresponding DELETE macro instructions are issued, the count is decremented until it reaches zero. An LLE has the following format:

Res	Next LLE	Count		CDE
0	1	4	5	

Byte 0: Reserved (RES).

Bytes 1-3:	Pointer to	the next more	recent
	LLE on the	load list.	

Byte 4: Count.

Bytes 5-7: Pointer to the corresponding CDE.

More will be said about CDEs in the next topic of Section 1, titled "Main Storage Supervision."

Job Pack Area Queue (MFT With Subtasking Only)

In an MFT system with subtasking, the job pack area queue is a chain of request blocks associated with load modules invoked by a LOAD macro instruction. The queue contains FRBs, and those LRBs and LPRBs that represent reenterable modules. FRBs are queued on the job pack area queue until the requested module is completely loaded. When the module is completely loaded into main storage, the FRB is removed from the job pack area queue and replaced with an LBR or an LPR queue on the job pack area queue if the loaded module is reenterable, and on the load list if it is not.

In the MFT with subtasking configuration, the load list represents non-reenterable modules, while the job pack area queue represents cnly reenterable modules within the partition. These RBs on the job pack area queue are not deleted automatically, but remain intact until they are deleted by a DELETE macro instruction, or until job step termination occurs. Reenterable load modules are therefore retained in the partition for use by the job step task or any subtasks which may be created.

Whenever a LOAD macro instruction is issued, the job pack area queue is searched. If the routine is already fully loaded and represented by an LRB or an LPRB on the JPAQ (the routine is reenterable), the system increments the use count by one and passes the module entry point address to the requesting routine. If an FRB for the requested module is found, a wait list element (WLE) representing the deferred request is queued to the FRB, and the request is placed in a wait. When the requested routine is fully loaded, the system releases the request from the wait condition, and the request is re-initiated. If no RB for the requested routine is found, an FRB is created and queued on the The system then searches the load JPAO. list of the requesting task for an RB for the requested routine. If an RB for that routine is found on the load list (the routine is not reenterable), the use count is incremented by one, the entry point address of the module is passed to the requesting routine, and the FRB is dequeued from the JPAQ. If no RB is found on the load list, the FRB remains on the JPAQ and the system begins loading the requested module.

Each time a DELETE macro instruction is issued for the routine, the use count is decremented by one (the DELETE routine ignores FRBs). When the use count reaches zero, the RB is removed from the queue.

Figure 5 illustrates how the job pack area queue is chained to a TCB.

In Figure 5, each RB contains a pointer to the previous RB and a pointer to the next RB on the queue. If there is no previous RB on the queue, that pointer will contain zero; if there is no next RB on the queue (this RB is the most recent on the JPAQ), the next RB pointer will point back to the job pack area queue pointer in the PIB.

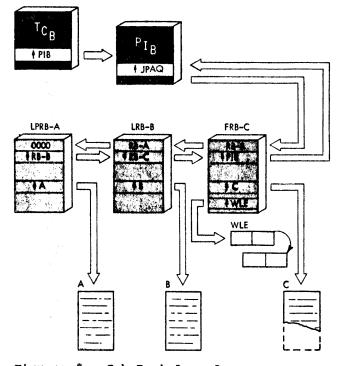


Figure 5. Job Pack Area Queue

Two wait list elements (WLEs) are queued to FRB-C representing deferred requests waiting until the initial loading of the module is completed. The last WLE contains zero in its forward pointer, indicating that it is the last element on the WLE queue.

Effects of LINK, ATTACH, XCTL, and LOAD

LINK, ATTACH, XCTL, and LOAD, though similar, have some distinguishing characteristics and system dependencies worth mentioning. By knowing what happens when these macro instructions are issued, you can make more effective use of the active RB queue and the load list.

LINK: A LINK results in the creation of a PRB chained to the active RB queue. Upon completion of the invoked routine, control is returned to the invoking routine. In systems with MFT, the RB is removed from the queue. The storage occupied by the invoked routine is freed unless the routine is also represented on the load list, or on the job pack area queue in MFT systems with subtasking. In systems with MVT, the use count in the CDE is decremented by one; if it is then zero, the RB and the storage occupied by the routine are marked for deletion. A LINK macro instruction generates an SVC 6.

ATTACH: An ATTACH is similar to the other three macro instructions in systems with MFT without subtasking. In systems with MFT with subtasking or MVT, ATTACH is the means for dynamically creating a separate but related task -- a subtask.

At the MFT without subtasking level, ATTACH effectively performs the same functions as LINK with two notable additions:

- You can request an exit routine to be given control upon normal completion of the attached routine.
- You can request the posting of an event control block upon the routine's completion.

Exit routines are represented by additional RBs on the active RB queue. The ATTACH macro instruction generates an SVC 42(2A).

XCTL: An XCTL also results in the creation of a PRB and immediate transfer of control to the invoked routine. However, XCTL differs from the other macro instructions in that, upon completion of the invoked routine, control is passed to a routine other than the invoking routine. In fact, an XCTL does not result in the creation of a lower level RB. Instead, the invoking routine and its associated RBs are deleted when the XCTL is issued. In effect, the RB for the invoked routine replaces the invoking routine's RB. The XCTL macro instruction generates an SVC 7.

LOAD: The LOAD macro instruction was treated previously in the discussion of the load list. To summarize: the system responds to a LOAD by fetching the routine into main storage and passing the entry point address to the requesting routine in register 0. Because the system does not have an indication of when the routine is no longer needed, a LOAD must be accompanied by a corresponding DELETE macro instruction. If not, the routine and its RB remain intact until the job step is terminated. The LOAD macro instruction generates an SVC 8.

System Task Control Differences

Thus far, this topic has dealt with the aspects of task supervision that are similar for MFT and MVT. There are, however, some major differences:

- 1. The number of tasks that can be known to the system concurrently.
- 2. The layout of main storage.
- 3. The additional main storage control information in systems with MVT.

The first two subjects are discussed here, by system. The third subject, because of its volume, is discussed in the next topic of Section 1.

Systems With MFT (Without Subtasking)

Figure 6 is a snapshot of main storage in a system with MFT without subtasking.

The <u>fixed area</u> contains the nucleus (including TCB queue, transient area loading task, communications task, and master scheduler task), and the system queue area. Optionally it may contain access methods and SVC routine which are normally nonresident, a list of absolute addresses for routines which reside on direct access devices, and a reenterable load module area.

One TCB exists for each task. All TCBs are linked by dispatching priority in a TCB queue, beginning with the three resident tasks.

The <u>dynamic area</u> is divided into a maximum of 52 partitions. Each partition contains one task. The dynamic area can contain as many as 3 reading tasks, 36 writing tasks, and 15 job step tasks, providing that the total number of tasks does not exceed 52. Partition sizes and attributes are defined during system generation. Figure 7 shows the contents of an MFT partition.

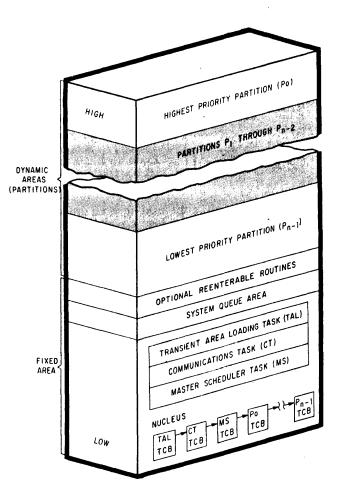


Figure 6. Main Storage Snapshot (MFT Without Subtasking)

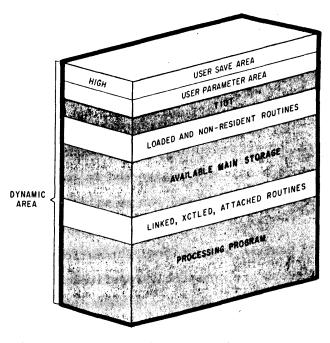


Figure 7. Partition (MFT Without Subtasking)

Jobs are processed sequentially in a partition, one job step at a time. An ATTACH macro instruction does not create a subtask.

Systems with MFT (With Subtasking): Operating Systems that provide multiprogramming with a fixed number of tasks with the subtasking option (MFT with subtasking) differ from MFT systems without subtasking in the following major areas:

- MFT with subtasking has an ATTACH 1. facility similar to the ATTACH facility in MVT. While the number of job step TCBs still may not exceed 15, the number of tasks in any partition, and therefore the total number of tasks in the system, is now variable. Job step task TCBs reside in the nucleus. They are queued, following the system task TCBs, in the same manner as in MFT without subtasking. When subtasks are created, however, the subtask TCBs are placed in the system queue area and queued to the job step TCBs according to dispatching priority (TCBTCB field), and according to subtask relationships (TCBNTC, TCBOTC, TCBLTC fields).
- 2. MFT with subtasking provides the ability to change the dispatching priority of any task within a partition through the use of the CHAP macro instruction.

Figure 8 is a snapshot of main storage in an MFT system with subtasking. Note here that the TCBs in the nucleus are all job step TCBs, while those residing in the sytem queue area are the subtask TCBs.

Systems with MVT: In Operating Systems that provide multiprogramming with a variable number of tasks (MVT), as many as 15 job steps can be executed concurrently. Each job step requests an area of main storage called a <u>region</u> and is executed as a job step task. In addition, <u>system tasks</u> request regions and can be executed concurrently with job step tasks.

Regions are assigned automatically from the dynamic area when tasks are initiated. Regions are constantly redefined according to the main storage requirements of each new task.

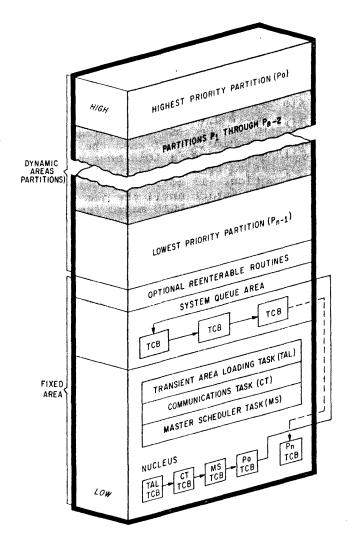


Figure 8. Main Storage Snapshot (MFT With Subtasking)

With the facility of attaching subtasks available to each task through the ATTACH macro instruction, the number of TCBs in the system is variable. Tasks gain control of the CPU by priority. To keep track of the priority and status of each task in the system, TCBs are linked together in a TCB queue.

Figure 9 is a snapshot of main storage in a system with MVT. The <u>fixed area</u> is occupied by the resident portion of the control program loaded at IPL. The <u>system</u> <u>queue space</u> is reserved for control blocks and tables built by the control program. The <u>dynamic area</u> is divided into variable-sized regions, each of which is allocated to a job step task or a system task. Finally, the <u>link pack area</u> contains selected reenterable routines, loaded at IPL. If an IBM 2361 Core Storage device and Main Storage Hierarchy Support are included in the system, a secondary link pack area may be created in hierarchy 1 to contain other reenterable routines.

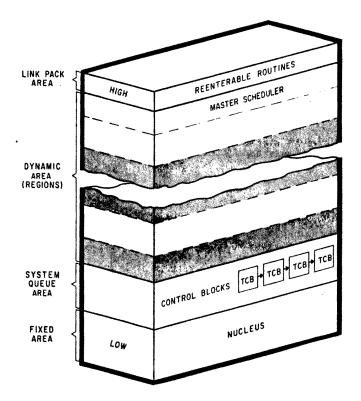


Figure 9. Main Storage Snapshot (MVT)

Main Storage Supervision

Storage control information is kept in a series of control blocks called <u>queue</u> <u>elements</u>. In systems with MFT without subtasking, queue elements reflect areas of main storage that are unassigned. In MFT systems with subtasking, a <u>gotten subtask</u> <u>area queue element</u> (GQE) is introduced to record storage obtained for a subtask by a supervisor issued GETMAIN macro instruction. In systems with MVT, more elaborate storage control is maintained; at any given time, queue elements reflect the distribution of main storage in regions, subpools, and load modules.

The dynamic area may be significantly expanded by including IBM 2361 Core Storage in the system. Main Storage Hierarchy Support for IBM 2361 Models 1 and 2 permits selective access to either processor storage (hierarchy 0) or 2361 Core Storage (hierarchy 1). If IBM 2361 Core Storage is not included, requests for storage from hierarchy 1 are obtained from hierarchy 0. If 2361 Core Storage is not present in an MVT system and a region is defined to exist in two hierarchies, a two-part region is established within processor storage. The two parts are not necessarily contiguous.

Storage Control in Systems with MFT (Without Subtasking)

The chain of storage control information in an MFT system without subtasking begins at a table called the main storage supervisor (MSS) boundary box, located in the system nucleus. There is one MSS boundary box for each partition. It is pointed to by the TCB (TCBMSS field) for the partition.

Each boundary box contains 3 words. The first word points to the Free Queue Element (FQE) associated with the highest free area in the partition. The second word points to the lowest limit of the partition. The third word contains the highest address in the partition plus 1.

If Main Storage Hierarchy Support is included, the first half of each expanded boundary box describes the processor storage (hierarchy 0) partition segment, and the second half describes the 2361 Core Storage (hierarchy 1) partition segment. Any partition segment not currently assigned storage in the system has the applicable boundary box pointers set to zero. If the partition is established entirely within hierarchy 0, or if 2361 Core Storage is not included in the system, the hierarchy 1 pointers in the second half of the expanded boundary box are set to zero. If a partition is established entirely within hierarchy 1, the hierarchy 0 pointers in the first half of the expanded boundary box are set to zero.

<u>FQE</u>: Each free area in a partition is described by an FQE. FQEs are chained beginning with the FQE associated with the free area having the highest address in the partition. If Main Storage Hierarchy Support is present, one FQE chain exists for each hierarchy specified. Each FQE occupies the first 8 bytes of the area it describes. It has the following format:

Next Fo	2E Number of bytes
Bytes 0-3:	Pointer to FQE associated with
	next lower free area or, if this is the last FQE, zeros.
Bytes 4-7:	Number of bytes in the free area.
	ummarizes storage control in h MFT without subtasking.

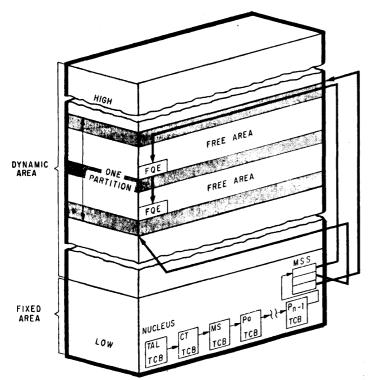


Figure 10. Storage Control for a Partition (MFT Without Subtasking)

<u>Storage Control in Systems with MFT (With</u> <u>Subtasking)</u>

Storage control information for the job step or partition TCB in MFT systems with subtasking is handled in the same way as in MFT systems without subtasking. However, when subtasks are created, the supervisor builds another control block, the <u>gotten</u> <u>subtask area queue element</u> (GQE). The GQEs associated with each subtask originate from a one word pointer addressed by the TCBMSS field of the subtask TCB.

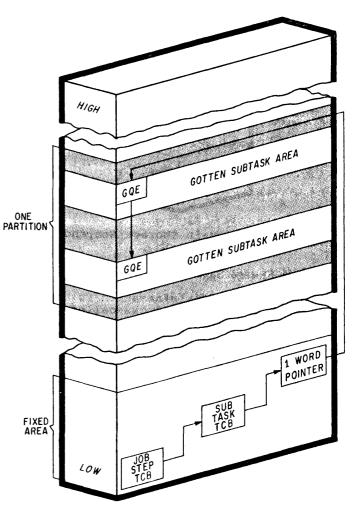
<u>GQE</u>: Each area in main storage belonging to a subtask, and obtained by a supervisor issued GETMAIN macro instruction, is described by a <u>gotten subtask area queue</u> <u>element</u> (GQE). GQEs are chained in the order they are created. The TCBMSS field of the subtask TCB contains the address of a word which points to the most recently created GQE.

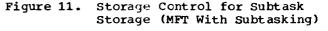
If Main Storage Hierarchy Support is present in the system, the GQE chain can span from hierarchy 0 to hierarchy 1 and back in any order. Each GQE occupies the first eight bytes of the area it describes, and has the following format:



- Bytes 0-3: Pointer to the Previous GQE or, if zero, this is the last GQE on the chain.
- Bytes 4-7: Number of bytes in the gotten subtask area.

Figure 11 summarizes the chaining of GQEs to a subtask TCB.





Storage Control for a Region in Systems with MVT

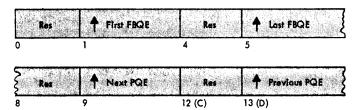
Unassigned areas of main storage within each region of a system with MVT are reflected in a queue of partition queue elements (PQEs) and a series of free block queue elements (FBQEs). <u>PQE</u>: The partition queue associated with a region resides in the system queue space. It is connected to the TCBs for all tasks in the job step through a dummy PQE located in the system queue space. A dummy PQE has the following format:

fint PQE	Lost PQE
٥	4

Bytes 0-3: Pointer to the first PQE in the partition queue.

Bytes 4-7: Pointer to the last PQE in the partition queue.

In systems that do not include the rollout/rollin feature or Main Storage Hierarchy Support for IBM 2361 Models 1 and 2, there is one PQE for each job step. If the rollout feature is used, additional PQEs are added each time a job step borrows storage space from existing steps or acquires unassigned free space to satisfy an unconditional GETMAIN request. These additional PQEs are removed from the queue as the rollin feature is used. If Main Storage Hierarchy Support is present, one PQE exists for each hierarchy used by the job step. A PQE has the following format:



{	1 Owning		Region Size 2	
16 (10)	17 (11)	20 (14)	21 (15)	

Res	Region	Flags	Rea
24 (18)	25 (19)	28 (1C)	29 (1D)

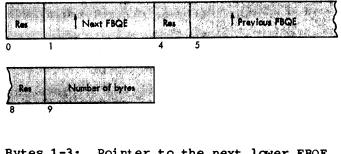
- Bytes 1-3: Pointer to the first FBQE or, if there are no FBQEs, a pointer to the PQE itself.
- Bytes 5-7: Pointer to the last FBQE or, if there are no FBQEs, a pointer to the PQE itself.
- Bytes 9-11(B): Pointer to the next PQE or, if this is the last PQE, zeros.

- Bytes 13-15(D-F): Pointer to the previous PQE or, if this is the first PQE, zeros.
- Bytes 17-19(11-13): Pointer to the TCB of the owning job step.
- Bytes 21-23(15-17): Size of the region, in 2K (2048) bytes.
- Bytes 25-27(19-1B): Pointer to the first byte of the region.

Byte 28(1C): Rollout flags.

<u>FBQE</u>: The FBQEs chained to a PQE reflect the total amount of free space in a region. Each FBQE is associated with one or more contiguous 2K blocks of free storage area. FBQEs reside in the lowest part of their associated area. As area distribution within the region changes, FBQEs are added to and deleted from the free block queue.

An FBQE has the following format:



- Bytes 1-3: Pointer to the next lower FBQE or, if this is the last FBQE, a pointer to the PQE.
- Bytes 5-7: Pointer to the preceding FBQE, or, if this is the first FBQE, a pointer to the PQE.
- Bytes 8-11(B): Number of bytes in the free block.

The remaining main storage in a region is used by problem programs and system programs. For convenience in referring to storage areas, the total amount of space assigned to a task represents one or more numbered <u>subpools</u>. (Subpools can also be shared by tasks.) Subpools are designated by a number assigned to the area through a GETMAIN macro instruction. Subpool numbers available for problem program use range from 0 through 127. Subpool numbers 128 through 255 are either unavailable or used by system programs.

Storage control elements and queues for a region are summarized in Figure 12.

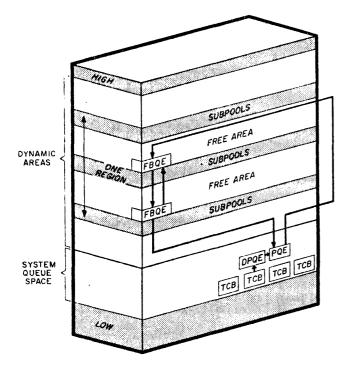


Figure 12. Storage Control for a Region (MVT)

Storage Control for a Subpool in Systems with MVT

Main storage distribution within each subpool is reflected in a subpool queue element (SPQE) and queues of descriptor queue elements (DQEs) and free queue elements (FQEs).

SPQE: SPQEs are associated with the subpools created for a task. SPQEs reside in the system queue space and are chained to the TCB(s) that use the subpool. They serve as a link between the TCB and the descriptor queue, and may be part of a subpool queue if the task uses more than one subpool. If a subpool is used by more than one task, only one SPQE is created. An SPQE has the following format:

Flags	Next SPQE SP / First DQE
0	1 4 5
Byte	0:
-	Bit 0 - Subpool is owned by this task if zero; shared, and owned by another task, if one.
	Bit 1 - This SPQE is the last on the queue, if one.
	Bit 2 - Subpool is shared and owned by this task, if one.
	Bits 3-7 - Reserved.

24 Programmer's Guide to Debugging (Release 21)

- Bytes 1-3: Pointer to next SPQE or, in last SPQE, zero.
- Byte 4: Subpool number.
- Bytes 5-7: Pointer to first DQE or, if the subpool is shared, a pointer to the "owning" SPQE.

DQE: DQEs associated with each SPQE reflect the total amount of space assigned to a subpool. Each DQE is associated with one or more 2K blocks of main storage set aside as a result of a GETMAIN macro instruction. Each DQE is also the starting point for the free queue. A DQE has the following format:

Res	1 First FQE	Res	Next DQE
0	1	4	3
Ret	First 2K block	Res	Length of area

Bytes 1-3: Pointer to the FQE associated with the first free area.

- Bytes 5-7: Pointer to the next DQE or, if this is the last DQE, zeros.
- Bytes 9-11(B): Pointer to first 2K block described by this DQE.
- Bytes 13-15(D-F): Length in bytes of area described by this DQE.

FQE: The FOE describes a free area within a set of 2K blocks described by a DQE. It occupies the first eight bytes of that free area. Since the FQE is within the subpool, it has the same protect key as the task active within that subpool. Extreme care should be exercised to see that FOEs are not destroyed by the problem program. If an FQE is destroyed, the free space that it describes is lost to the system and cannot be assigned through a GETMAIN. As area distribution within the set of blocks changes, FQEs are added to and deleted from the free queue. An FQE has the following format:

Re	8	Next FQE	Re	3	Number of bytes	
0	1		4	5		

- Bytes 1-3: Pointer to the next lower FQE or, if this is the last FQE, zeros.
- Bytes 5-7: Number of bytes in the free area.

Storage control for a subpool is summarized in Figure 13.

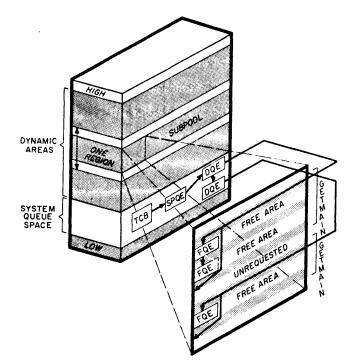


Figure 13. Storage Control for a Subpool (MVT)

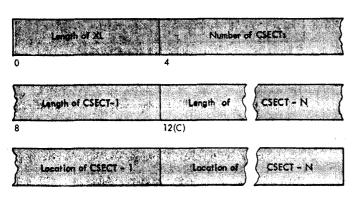
Storage Control for a Load Module in Systems With MVT

Each load module in main storage is described by a contents directory entry (CDE) and an extent list (XL) that tells how much space it occupies.

<u>CDE</u>: The contents directory is a group of queues, each of which is associated with an area of main storage. The CDEs in each queue represent the load modules residing in the associated area. There is a CDE queue for the link pack area and one for each region, or job pack area. The TCB for the job step task that requested the region Contents directory queues reside in the system queue space. A CDE has the following format:

Flige Next CDE Res 1 RB 0 1 4 5
ESCDIC program name
Count tentry point Flags tXL
16(10) 17(11) 20(14) 21(15)
Byte 0: Flag bits, when set to one, indicate: Bit 0 - Module was loaded by NIP. Bit 1 - Module is in process of being loaded. Bit 2 - Module is reenterable. Bit 3 - Module is serially reusable.
Bit 4 - Module may not be reused. Bit 5 - This CDE reflects an alias name (a minor CDE). Bit 6 - Module is in job pack area. Bit 7 - Module is not only-loadable.
Bytes 1-3: Pointer to next CDE.
Bytes 5-7: Pointer to the RB.
Bytes 8-15(F): EBCDIC name of load module.
Byte 16(10): Use count.
Bytes 17-19(11-13): Entry point address of load module.
Byte 20: Flag bits, when set to one, indicate: Bit 0 - Reserved. Bit 1 - Module is inactive. Bit 2 - An extent list has been built for the module. Pit 3 - This CDE contains a relocated alias entry point address. Bit 4 - The module is refreshable. Bits 5, 6, 7 - Reserved.
Bytes 21-23(15-17): Pointer to the XL for this module or, if this is a minor CDE, pointer to the major CDE.

<u>XL</u>: The total amount of main storage occupied by a load module is reflected in an extent list (XL). XLs are located in the system queue space. An XL has the following format:



Bytes 0-3: Length of XL in bytes.

- Bytes 4-7: Number of scattered control sections. If the control sections are block-loaded, 1.
- Remaining Length in bytes of each bytes: control section in the module (4 bytes for each control section) and starting location of each control section (4 bytes for each control section).

Storage control elements and queues for load modules are summarized in Figure 14.

System Control Blocks and Tables

In addition to the key task management control blocks (TCB and RB), several other control blocks containing essential debugging information are built and maintained by data management and job management routines. Although some of these blocks are not readily identifiable on a storage dump, they can be located by following chains of pointers that begin at the TCB.

The control blocks discussed here have the same basic functions at each control program level. The precise byte-by-byte contents of the blocks can be found in the publication <u>System Control Blocks</u>. Block contents useful in debugging are listed in Appendix K.

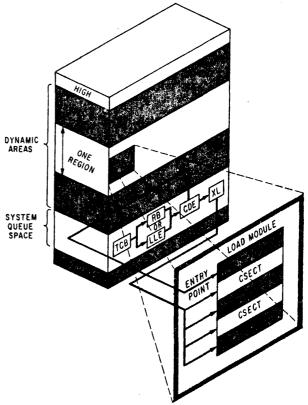


Figure 14. Storage Control for a Load Module (MVT)

Communications Vector Table (CVT)

The CVT provides a means of communication between nonresident routines and the control program nucleus. Its most important role in debugging is its pointer to two words of TCB addresses. These words enable you to locate the TCB of the active task, and from there to find other essential control information. Storage locations 16(10) and 76(4c) contain a pointer to the CVT.

Task Input/Output Table (TIOT)

A TIOT is constructed by job management for each task in the system. It contains primarily pointers to control blocks used by I/O support routines. It is usually located in the highest part of the main storage area occupied by the associated task (in systems with MVT, TIOTs are in the system queue space.) Through the TIOT, you can obtain addresses of unit control blocks allocated to the task, the job and step name, the ddnames associated with the step, and the status of each device and volume used by the data sets.

Unit Control Block (UCB)

The UCB describes the characteristics of an I/O device. One UCB is associated with each I/O device configured into a system. The UCB's most useful debugging aid is the sense information returned by the last sense command issued to the associated device.

Event Control Block (ECB)

The ECB is a 1-word control block created when a READ or WRITE macro instruction is issued, initiating an asynchronous I/O operation. At the completion of the I/O operation, the access method routine posts the ECB. By checking this ECB, the completion status of an I/O operation can be determined. In all access methods but QTAM, the ECB is the first word of a larger block, the data event control block.

Input/Output Block (IOB)

The IOB is the source of information required by the I/O supervisor. It is filled in with information taken from an I/O operation request. In debugging, it is useful as a source of pointers to the DCB associated with the I/O operation and the channel commands associated with a particular device.

Data Control Block (DCB)

The DCB is the place where the operating system and the problem program store all pertinent information about a data set. It may be completely filled by operands in the DCB macro instruction, or partially filled in and completed when the data set is opened, with subparameters in a DD statement and/or information from the data set label. The format of DCBs differs slightly for each of the various access methods and device types. The DCB's primary debugging aids are its pointers to the DEB and current IOB associated with its data set, and the offset value of the ddname in the TIOT.

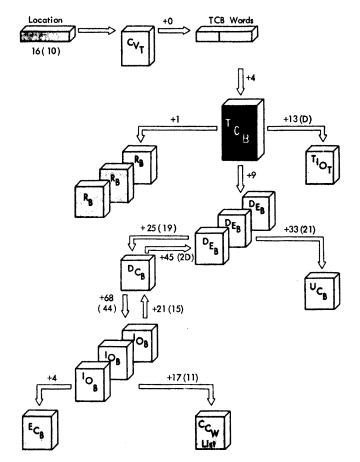
Data Extent Block (DEB)

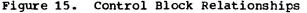
A DEB describes a data set's auxiliary storage assignments and contains pointers

to some other control blocks. The DEB is created and queued to the TCB at the time a data set is opened. Each TCB contains a pointer to the first DEB on its chain. Through this pointer you can find out which data sets are opened for the task at a given time, what extents are occupied by open data sets, and where the DCB and UCB are located.

Summary of Control Block Relationships

Figure 15, an expansion of Figure 1, shows the relationships among the principal control blocks and tables in the System/360 Operating System.





Topics composing Section 2 are:

- ABEND/SNAP dumps issued by systems with MFT.
- ABEND/SNAP dumps issued by systems with MVT.
- Indicative dumps.
- Storage dumps.

Each topic includes instructions for invoking the dump, a detailed description of the dump's contents, and a guide to using the dump.

ABEND/SNAP Dump (MFT)

ABEND/SNAP storage dumps are issued whenever the control program or problem program issues an ABEND or SNAP macro instruction, or the operator issues a CANCEL command requesting a dump, and proper dump data sets have been defined. However, in the event of a system failure, if a SYS1.DUMP data set has been defined and is available, a full storage dump will be provided, as explained in the section "Storage Dumps."

Since, in an MFT with subtasking system, subtasks may be created, you may receive one or more partial dumps in addition to the complete dump of the task that caused the abnormal termination. A complete dump includes a printout of all control information related to the terminating task, and the nucleus and all allocated storage within the partition in which the abending task resided. A partial dump of a task related to the terminating task includes only control information. The partial dump is identified by either ID=001 or ID=002 printed in the first line of the dump. Figure 16 is a copy of the first few pages of a complete ABEND dump of an MFT system with subtasking. It illustrates some of the key areas on an ABEND dump, as issued by systems with MFT. Those portions of the dump that would only appear on a dump of a subtasking system are noted in the later discussions as appearing only in a dump of an MFT with subtasking system.

For a discussion of a formatted ABEND dump using the telecommunications acces method (TCAM) in an MFT environment, see <u>IBM System/360 Operating System: TCAM</u> <u>Program Logic Manual</u>, GY30-2029. References to other TCAM debugging aids are found in Appendix J.

Invoking an ABEND/SNAP Dump (MFT)

ABEND dumps are produced as a result of an ABEND macro instruction, issued either by a processing program or an operating system routine. The macro instruction requires a DD statement in the input stream for each job step that is subject to abnormal termination. This DD statement must be identified by one of the special ddnames SYSABEND or SYSUDUMP. SYSABEND results in edited control information, the system nucleus, the trace table, and a dump of main storage; SYSUDUMP excludes the nucleus and the trace table. In the event of a system failure, the Damage Assessment routine (DAR) attempts to write a storage image dump to the SYS1.DUMP data set. A full explanation of storage dumps may be found in the section "Storage Dumps."

<u>SNAP Dumps</u> result from a problem program issuing a SNAP macro instruction. The contents of a SNAP dump vary according to the operands specified in the SNAP macro instruction. SNAP dumps also require a DD statement in the input stream. This DD statement has no special characteristics except that its ddname must not be SYSABEND or SYSUDUMP. The processing program must define a DCB for the snapshot data set. The DCB macro instruction must contain, in addition to the usual DCB requirements, the operands DSORG=PS, RECFM=VBA, MACRF=(W), BLKSIZE=882 or 1632, and LRECL=125. In addition, the DCB must be opened before the first SNAP macro instruction is issued.

Main Storage Considerations: Three BSAM modules (IGG019BA, IGG019BB, and the device-dependent EOB module) are required to process dumps. These modules should be made resident in the Resident Access Method (RAM) area by specifying RESIDNT=ACSMETH in the SUPRVSOR macro instruction during system generation. If these modules are not resident, as much as 1352 bytes of main storage within the partition are required to contain them.

In addition to the area required for the BSAM modules, 2784 bytes must be available in the partition. 1344 of these bytes are required for EOV processing should the initial space specification for a direct access device be exceeded by the dump requirements.

	4 ³¹	· ·							an a			in i			
	* ABDL	IMP REQU	ESTED	•											·
		EOT24		STEP S	TEP	TIM	E 000737	DATE 9	9366					PAG	E 0001
C	OMPLET		E	USER =	0123										
ł	INTERAL	UPT AT C	6EF 5A												
					0000 4006										
	TCB	01C820	MSS ESAC	0007FC5 0001CC5 1506E8F 0000000	8 PK/FL B TCB	0000000 G 1081040 0001D0A 10000000 0000000	8 FLG O TME	0007F78C 000001F8 0001C808	LLS P18	0007FDB0 00000000 F0012420 00000000 00000000	JLB	8000007 0007FF7 0000000 F805000 0000000	B JST O OTC O Resv	00000000 00005508 0001CDEC 000000000 000000000	jê Î
	ACTIVE	RBS													
1	PRB	06EE28	NM TA	THBLOG	SZ/STAB	00302000	USE/E	P 0106EE4		00150000				/LNK 000	
:	SVRB	07FD20	NM SV RG 0- 8-15-		00240	00120062 80000078 0007FF78	USE/E 00000 0007F	P 0000787 000 00 FB0 00	8 PSW 080000 07FFF8	FF040033 0007F6 4006E6	48 00	0 Q 90 0000098 006EE60	0390 WT 0000550 0000984		6EE28 7FC30 00000
	SVR B	07FC58	NM SV RG 0- 8-15-		7F7E8	000CD062 0007FD80 0006F296	USE/E 40007 0001C		8 PSW 0097F8 00225C	FF040008 0001C8 0001C8	320 00	C Q F8 007FD20 006F230	03F8 WT 0006F23 90007C8		7FD20 055D8 1E7C8
C. A. A	LPRB Q	CL ARBA DECAD OTERN D STOLD S STOLD S MODEL S S S S S S S S S S S S S S S S S S S	n Pati Pati Pati	ALIOS ALIOS CLOS CLOS CLOS	SZZSTAD SZZSTAD SZZSTAD	002F2010 00302000 00122090 00122090 00132090	that we get	0106EC60 D1D4EE44 0104F031 d104F030 0104F0300	PSV PSV PSV	No where	4006EF5 40006AE 4006F16	0 000 0 0 000 0 0 000 1 0 000	000		CDE0 C820 CC80 D0AQ1
	P/P ST	ORAGE BO	UNDAR	ES 0006	E800 TO 0	0080000									
	FREE A	REAS	SIZE												
	06EB 06EC		000000												
	06F5	68 (000FC	8											
	07F7 07F8 07F8	40 (0000001	28											
. 20	07FE		000000												
	OTTEN		-61. 26												
	06F2 0765 0765	29 90 0 F0 0													
	07750 0777 0777 07774 07774	00 60 . A													
												5 a. 19			

Figure 16. Sample of an ABEND Dump (MFT) (Part 1 of 2)

	SAVE	AREA TR	ACE												PAGE	0002	
	TATH	BLOG WAS	ENTE	RED													
	SA		WD1 R1	0606 EAC8 0001 CC 80	R2	00000100	L SA R3	0006EE60) R4	00009848 0007FE48	85	4006EE48	R6	000095CE 000055D8			
	SA	06EE60	WD1	0007FC30	HSA	0006ECE0		0007FF7) RET	0007FFB0	EPA	00007FFF8	RO				
				000000000000000000000000000000000000000		000000000000000000000000000000000000000	R 3 R 9	00000000		00000000		00000000		00000000			
	PROCI	EEDING B	ACK \	IA REG 1	5												
	SA	068660	R1	000000000000000000000000000000000000000	R2	0006EBF8 00000000 00000000		0000000) R4	00000000 00000000 00000000	R5	00000000	R6	00000000			
	TATH	B10G WAS	ENTE	RED													
	SA	06E8F8	R1	0606EAC8 0001CC80 0007FC30	R2	00000100 00000000 0006ECE0		0006EE60 00080000 0007FF78	R4	00009848 0007FE48 0007FE80	R S	4006EE48 00000098 0007FFF8	R6				
·	DATA	SETS															
	SNAP		UCE	192	002250	DE	B 07F	780	DCB 0	6EF84							
	DUMDO		UCE		002250		8 07F		008 0								
	JOBLI		UCE		002180												
	SYSPE		UCE	_	002250												
	SYSAE		UCE		002250												
	SNAP		UCB		002180												4
	REGS	AT ENTR	Y TO	ABEND													
		.REGS O	-6	00.00	0000 0	0000000	0	0.000000	000000	00 00	0.000	000000 000	00	00.000000	0000000		
	REGS REGS			000002A 0006EE6			00000 0007F		80000 7FFF8		0007F 4006E		00098 6EE60	000055D8 00009848	0007FC30 00000000		
	NUCLE	EUS															
	00002		0000	00000510	FOFOF	5C1 00000	000	0000976	8 0001	3440 01040 3000 FF040	0080 8	003ACD4			8		
	00004	40 100	7F5E8	50000000 0000033A	00001	480 00009	7F 8	6008500	0 0000	0000 00040 0000 00040 00040	0000 0	0000282	*5Y				
	00006	000 06	15380	000000000	00000	000 0000	000	0000000	0000 00		0 0000	0000000	*		• • • • • • • • • • • • • • • • • • •	*	
		INES I	00000	0~000140	SAN	E AS ABOV	£			A7A0 0000					• • • • • • • • • • • • • • • • • • •		
	00016	80 000	16820	00007E91	0006F	465 80007	D16	0000008	0006	F491 0000	0001 0	006F4A8	*	4			
		INE O	00100) SAME) 00006888	AS ABO	VE				AD42 9000					• • • • • • • • • • • • • • • • • • •		
	00020	000 000	08460	00008364	00006	780 00006	942	0000100	00000	0F28 0000 5840 3000	9730 0	0013350	*		• • • • • • • • • • • • • • • • • • •	••••	
	00024	401	00038	94FD4011 003847F0	90A13	030 58900	210	0589585	0 0210	5890 02140 4780 0448)7F9 9	OALOLEO	*		0 9	• • • • *	
	00028	0 044	09029	018091F0	02384	780 02909	041	01E0D20	7 0440	0018 47F0	0282 5	589006C4	*	0	K0	•••D*	
	00020	0 918	00018	47800208	58200	208 05224	7F0	026A000	1000 00	5388 0000	870A 0	A0390A9	*			*	
	00032	0000 000	01244	47C00332	90C2B	004 18185	880	0218928	1000	98F0 A000	8900 C	0001200	*****	P	D	*	
	00032		5 7 0F U	. JU2C+IEU		CAD OTAOS	2.00	0020181	0000	VE10 0178	,		+•••U		• • • • • • • • • • • • • • •	••••	

Figure 16. Sample of an ABEND Dump (MFT) (Part 2 of 2)

Device and Space Considerations: DD statements for ABEND/SNAP dumps, must contain parameters appropriate for a basic sequential (BSAM) data set. Data sets can be allocated to any device supported by the basic sequential access method. There are several ways to code these DD statements depending on what type of device you choose and when you want the dump printed.

If you wish to have the dump printed immediately, code a DD statement defining a printer data set.

//SYSABEND DD UNIT=1443,DCB=(...

A printer is associated with the SYSOUT class, you can also obtain immediate printing by routing the data set through the output stream.

//SNAPDUMP DD SYSOUT=A, DCB=(...

This type of request is the easiest, most economical way to provide for a dump. All other DD statements result in the tying up of an output unit or delayed printing of the dump.

If you wish to retain the dump, you can keep or catalog it on a direct access or tape unit. The last step in the pertinent job can serve several functions: to print out key data sets in steps that have been abnormally terminated, to print an ABEND or SNAP dump stored in an earlier step, or to release a tape volume or direct access space acquired for dump data sets. Conditional execution of the last step can be established through proper use of the COND parameter and its subparameters, EVEN and ONLY, on the EXEC statement. Direct access space should be requested in units of average block size rather than in cylinders (CYL) or tracks (TRK). If abnormal termination occurs and the data set is retained, the tape volume or direct access space should be released (DELETE in the DISP parameter) at the time the data set is printed.

<u>Sample DD Statements</u>: Figure 17 shows a set of job steps that include DD statements for ABEND dump data sets.

The SYSABEND DD statement in STEP2 takes advantage of the direct access space acquired in STEP1 by indicating MOD in the DISP parameter. Note that the space request in STEP1 is large so that the dumping operation is not inhibited due to insufficient space. The final SYSABEND DD statement in the job should indicate a disposition of DELETE to free the space acquired for dumping.

Contents of an ABEND/SNAP Dump (MFT)

This explanation of the contents of ABEND/SNAP dumps for systems with MFT is interspersed with sample sections taken from an ABEND dump. Capital letters represent the headings found in all dumps, and lowercase letters, information that varies with each dump. The lowercase letter used indicates the mode of the information, and the number of letters indicates its length:

- h represents 1/2 byte of hexadecimal information
- d represents 1 byte of decimal information
- c represents a 1-byte character

You may prefer to follow the explanation on your own ABEND or SNAP dump.

									1	13			
						-				ż.			
												 	•
-												 	
-			+									 	
175	TEPI	EXEC	PGN=PRO	GRAMI									
// \$	YSABEN	D DD	DSNAME =	DUMP, UN1 SER=1234	T=231	1, DIS	P=(,K	EEP,K	EEP),			X	
11			VOLUME=	SER=1234	, SPAC	E= (TR	K,(11	0,10))			 	
	othe	rvu	statemen									 	
115	TEP2	EXEC	PGM= PRO	GRAM2								 	
//s	YSABEN	D DD	DSN AME =	*.STEPI.	SYSAB	END,D	ISP=(MOD, D	ELETE	, KEEP),	 X	
11			VOL=REF	**.STEPI.	. SYSA	BEND						 	
												 _	
												 	*

Figure 17. SYSABEND DD Statements

* * A B D U M P R E Q U E S T E D * * * identifies the dump as an ABEND or SNAP dump.

*ccccccc....
is omitted or is one or more of the
following:

*CORE NOT AVAILABLE, LOC. hhhhhhhhhh TAKEN... indicates that the ABDUMP routine confiscated storage locations hhhhh through hhhhhh because not enough storage was available. This area is printed under P/P STORAGE, but can be ignored because the problem program originally in it was overlaid during the dumping process.

- *MODIFIED, /SIRB/DEB/LLS/ARB/MSS...
 indicates that the one or more
 queues listed were destroyed or
 their elements dequeued during
 abnormal termination:
 - SIRB -- system interruption request block queue. One or more SIRB elements were found in the active RB queue: these elements are always dequeued during dumping.
 - DEB -- DEB queue. If the first message also appeared, either a DEB or an associated DCB was overlaid.
 - LLS -- load list. If the first message also appeared, one or more loaded RBs were overlaid.
 - ARB -- active RB queue. If the first message also appeared, one or more RBs were overlaid.
 - MSS -- boundary box queue. One or more MSS elements were dequeued, but an otherwise valid control block was found

in the free area specified by an MSS element.

- *FOUND ERROR IN /DEB/LLS/ARB/MSS... indicates that one or more of the following contained an error:
 - DEB: data extent block
 - LLS: load list
 - ARB: active RB
 - MSS: boundary box

This message appears with either the first or second message above. The error could be: improper boundary alignment, control block not within storage assigned to the program being dumped, or an infinite loop (300 times is the maximum for this test). For an MSS block, 4 other errors could also be found: incorrect descending sequence (omitting loop count), overlapping free areas, free area not entirely within the storage assigned to the program being dumped, or count in count field not a multiple of 8.

- JOB cccccccc is the job name specified in the JOB statement.
- STEP cccccccc is the step name specified in the EXEC statement for the problem program being dumped.
- TIME dddddd is the hour (first 2 digits), minute (second 2 digits), and second (last 2 digits) when the ABDUMP routine began processing.
- DATE ddddd is the year (first 2 digits) and day of the year (last 3 digits). For example, 67352 would be December 18, 1967.

PAGE dddd is the page number. Appears at the top of each page.

COMPLETION CODE SYSTEM=hhh or COMPLETION CODE USER=dddd

is the completion code supplied by the control program (SYSTEM=hhh) or the problem program (USER=dddd). Either SYSTEM=hhh or USER=dddd is printed, but not both. Common completion codes are explained in Appendix B.

ccccc...

explains the completion code or, if a program interruption occurred: PROGRAM INTERRUPTION ccccc... AT LOCATION hhhhhh,

where ccccc is the program interruption cause -- OPERATION, PRIVILEGED OPERATION, EXECUTE, PROTECTION, ADDRESSING, SPECIFICATION, DATA,FIXED-POINT OVERFLOW, FIXED-POINT DIVIDE, DECIMAL OVERFLOW, DECIMAL DIVIDE, EXPONENT OVERFLOW, EXPONENT UNDERFLOW, SIGNIFICANCE, or FLOATING-POINT DIVIDE; and hhhhhh is the starting address of the instruction being executed when the interruption occurred.

INTERRUPT AT hhhhhh

is the address of next instruction to be executed in the problem program. It is obtained from the resume PSW of the PRB or LPRB in the active RB queue at the time abnormal termination was requested.

PSW AT ENTRY TO ABEND hhhhhhhh hhhhhhhh or PSW AT ENTRY TO SNAP hhhhhhhh hhhhhhh is the PSW for the problem or control program that had control when abnormal termination was requested or when the SNAP macro instruction was executed.

TCB	hhhhhh	RB MSS	hhhhhhhh hhhhhhhh	PIE PK/F	hhhhhhhh LG hhhhhhhh	DEB FLG	hhhhhhhh hhhhhhhh	TIOT LLS	հհհհհհհ հհհհհհհ	CMP JLB	հհհհհհհհ հհհհհհհհ	TRN JST	հհհհհհհ հհհհհհհ
		RG 0-			hhhhhhhh hhhhhhhh	hhhhhl hhhhhl		nhhhh nhhhh	hhhhhhh		hhhhhh	hhhhhh	
		FSA	15 hhhhhh hhhhhhhh	TCB	hhhhhhh	TME	nnn nnn hhhhhhhh	PIB	hhhhhhhh hhhhhhhh	NTC	hhhhhh hhhhhhhh	hhhhhhi OTC	hh hhhhhhhh hhhhhhhh
		LTC	hhhhhhh	IQE	hhhhhhh	ECB	hhhhhhh				hhhhhhh		hhhhhhh
		STAE	hhhhhhh	TCT	hhhhhhhh	USER	hhhhhhh	DAR	hhhhhhh	RESV	հհհհհհհ	JSCB	hhhhhhh

TCB hhhhhh

is the starting address of the TCB.

ŧ

RB hhhhhhhh

is the TCBRBP field (bytes 0 through 3): starting address of the active RB queue and, consequently, the most recent RB on the queue (usually ABEND's RB).

PIE hhhhhhhh is the TCBPIE field (bytes 4 through 7): starting address of the program interruption element (PIE) for the task.

- DEB hhhhhhhh is the TCBDEB field (bytes 8 through 11): starting address of the DEB queue.
- TIOT hhhhhhhh is the TCBTIO field (bytes 12 through 15): starting address of the TIOT.

CMP hhhhhhhh is the TCBCMP field (bytes 16 through 19): task completion code in hexadecimal. System codes are shown in the third through fifth digits and user codes in the sixth through eighth.

TRN hhhhhhhh

is the TCBTRN field (bytes 20 through 23): starting address of control core (table) for controlling testing of the task by TESTRAN.

MSS hhhhhhhh

is the TCBMSS field (bytes 24 through 27): starting address of the main storage supervisor's boundary box.

PK/FLG hhhhhhhh

contains, in the first 2 digits, the TCBPKF field (byte 28): protection key.

FLG hhhhhhhh
 contains, in the first 4 digits, the
 last 2 bytes of the TCBFLGS field
 (bytes 32 and 33): last 2 flag bytes.

contains, in the next 2 digits, the TCBLMP field (byte 34): number of resources on which the task is queued. contains, in the last 2 digits, the TCBDSP field (byte 35):

- Reserved in MFT without subtasking; both digits are zero.
- In MFT with subtasking, this field contains the dispatching priority of the TCB.
- LLS hhhhhhhh
 - is the TCBLLS field (bytes 36 through 39): starting address of the RB most recently added to the load list.
- JLB hhhhhhhh
 - is the TCBJLB field (bytes 40 through
 43): starting address of the DCB
 for the JOBLIB data set.
- JST hhhhhhh
 - is the TCBJST field (bytes 44 through 47). Not currently used in MFT without subtasking. In MFT with subtasking - the starting address of the TCB for the job step task.
- RG 0-7 and RG 8-15 is the TCBGRS field (bytes 48 through 111): contents of general registers 0 through 7 and 8 through 15, as stored in the save area of the TCB when a task switch occurred. These 2 lines appear only in TCBs of tasks other than the task in control when the dump was requested.
- FSA hhhhhhhh contains. in the

contains, in the first 2 digits, the TCBIDF field (byte 112): TCB identifier field.

contains, in the last 6 digits, the TCBFSA field (bytes 113 through 115): starting address of the first problem program save area. This save area was set up by the control program when the job step was initiated.

TCB hhhhhhhh

is the TCBTCB field (bytes 116 through 119): starting address of the next TCB of lower priority or, if this is the last TCB, zeros. TME hhhhhhhh

is the TCBTME field (bytes 120 through 123): starting address of the timer element created when an STIMER macro instruction is issued by the task. This field is not printed if the computer does not contain the timer option.

- PIB hhhhhhhh is the TCBPIB field (bytes 124 through 127): starting address of the program information block.
- NTC hhhhhhhh

is the TCBNTC field (bytes 128 through 131):

MFT without subtasking: zeros.

MFT with subtasking: the starting address of the TCB for the previous subtask on this subtask TCB queue. This field is zero both in the job step task, and in the TCB for the first subtask created by a parent task.

OTC hhhhhhhh

is the TCBOTC field (bytes 132 through 135): starting address of the TCB for the parent task. Both in the TCB for the job step task, and in MFT systems without subtasking this field is zero.

LTC hhhhhhhh

is the TCBLTC field (bytes 136 through 139): starting address of the TCB for the most recent subtask created by this task. This field is zero in the TCB for the last subtask of a job step, or in the TCB for a task that does not create subtasks. <u>This field</u> is always zero in an MFT system without subtasking.

hhhhhhhh is the TCBIQE field (bytes 140 through 143).

MFT without subtasking: zero.

MFT with subtasking: starting address of the interruption queue element (IQE) for the ETXR exit routine. This routine is specified by the ETXR operand of the ATTACH macro instruction that created the TCB being dumped. The routine is to be entered when the task terminates.

ECB hhhhhhhh is the TCBECB field (bytes 144 through 147).

MFT without subtasking: zero. TCT hhhhhhhh is the TCBTCT field (bytes 164 through MFT with subtasking: starting address 167): of the ECB field to be posted by the control program at task termination. This field is zero if the task was Address of the Timing Control Table attached without an ECB operand. (TCT): zeros if the System Management Facilities option is not present in XTCB hhhhhhhh the system. reserved for future use. LP/FL hhhhhhhh **USER** hhhhhhhh MFT without subtasking: reserved. is the TCBUSER field (bytes 168 through 171): to be used as the user MFT with subtasking: contains in the chooses. first byte, the limit priority of the task (byte 152). contains, in the last three bytes the field TCBFTFLG DAR hhhhhhhh contains, in the first 2 digits, (bytes 153 through 155) - flag bytes. Damage Assessment Routine (DAR) flags (byte 172); **RESV** hhhhhhhh reserved for future use. contains, in the last 6 digits, the secondary non-dispatchability bits STAE hhhhhhhh (bytes 173 through 175). contains, in the first 2 digits, STAE RESV hhhhhhhh flags (byte 160). reserved for future use. contains, in the last 6 digits, the TCBNSTAE field (bytes 161 through JSCB hhhhhhhh 163): starting address of the current is the TCBJSCB field (bytes 180 STAE control block for the task. This through 183): the last three bytes

CTIVE RBS								
ccc hhhhhh	NM CCCCCCC SZ/S	TAB hhhhhhhh	USE/EP hh	hhhhhh PSW	hhhhhhhh	hhhhhhh Q	hhhhhh WT/I	NK hhhhhhh
	RG 0-7 hhhhhhh	hhhhhhh	հհհհհհհ	hhhhhhhh	հհհհհհհ	հհհհհհհ	hhhhhhh	hhhhhhh
	RG 8-15 hhhhhhhh	hbhhhhh	hhhhhhh	hhhhhhh	hhhhhhh	hhhhhhh	hhhhhhh	hhhhhhhh

field is zero if STAE has not been

issued.

ACTIVE RBS Note: Three SVRBs for ABEND identifies the next lines as the processing exist in the nucleus. They are used when there is insufficient contents of the active RBs queued to the TCB. space in the partition to create an SVRB. cccc hhhhhh indicates the RB type and its starting address. NM XXXXXXX is the XRBNM field (bytes 0 through 7): in PRB, LRB, and LPRB, the The RB types are: program name; in IRB, the first byte contains flags for the timer or, if PRB Program request block the timer is not being used, contains SIRB Supervisor interrupt request no meaningful information; in SVRB for block a type 2 SVC routine, the first 4 bytes contain the TTR of the load LPRB Loaded program request block module in the SVC library, and the last 4 bytes contain the SVC number in IRB Interruption request block signed, unpacked decimal. SVRB Supervisor request block

contain the address of the Job Step

Control Block.

SZ/STAB hhhhhhhh (bytes 25 through 27): in PRB and LPRB, starting address of an LPRB for contains in the first 4 digits, the XRBSZ field (bytes 8 and 9): number an entry identified by an IDENTIFY of contiguous doublewords in the RB, macro instruction; in IRB, starting the program (if applicable), and address of a request element; in SVRB associated supervisor work areas. for a type 3 or 4 SVC, size of the program in bytes. contains in the last 4 digits, the XSTAB field (bytes 10 and 11): flag WT/LNK hhhhhhh contains, in the first 2 digits, the bytes. XRBWT field (byte 28): wait count. USE/EP hhhhhhhh contains, in the first 2 digits, the contains, in the last 6 digits, the XRBLNK field (bytes 29 through 31): XRBUSE field (byte 12): use count. primary queuing field. It is the contains, in the last 6 digits, the starting address of the previous RB XRBEP field (bytes 13 through 15): for the task or, in the first RB to be address of entry point in the placed on the queue, the starting address of the TCB. associated program. PSW hhhhhhhh hhhhhhh RG 0-7 and RG 8-15 is the XRBPSW field (bytes 16 through is the XRBREB field (bytes 32 through 23): resume PSW. 95 in IRBs and SVRBs): contents of general registers 0 through 15 stored Q hhhhhh in the RB. These 2 lines do not appear for PRBs, LPRBs, and LRBs.

cccc hhhhhh NM ccccccc SZ/STAB hhhhhhhh USE/EP hhhhhhhh PSW hhhhhhhh hhhhhhhh Q hhhhhh

is the last 3 bytes of the XRBQ field

contains, in the last 4 digits, the

LOAD LIST identifies the next lines as the contents of the load list queued to the TCB.

cccc hhhhhh indicates the RB type and its starting address.

The RB types are:

LOAD LIST

- Loaded request block LRB LPRB Loaded program request block D-LPRB Dummy loaded program request block. (Present if the resident reenterable load module option was selected).
- NM ccccccc is the XRBNM field (bytes 0 through 7): program name.
- SZ/STAB hhhhhhhh contains, in the first 4 digits, the XRBSZ field (bytes 8 and 9): number of contiguous doublewords for the RB, the program (if applicable), and associated supervisor work areas.

XSTAB field (bytes 10 and 11): flag bytes.

WT/LNK hhhhhhhh

USE/EP hhhhhhhh contains, in the first 2 digits, the XRBUSE field (byte 12): use count.

> contains, in the last 6 digits, the XRBEP field (bytes 12 through 15): address of entry point in the program.

- PSW hhhhhhhh hhhhhhh is the XRBPSW field (bytes 16 through 23): resume PSW.
- 0 hhhhhh is the last 3 bytes of the XRBQ field (bytes 25 through 27): in LPRB, starting address of an LPRB for an entry identified by an IDENTIFY macro instruction; in LRB, unused.
- WT/LNK hhhhhhhh contains, in the first 2 digits, the XRBWT field (byte 28): wait count.

contains, in the last 6 digits, the XRBLNK field (bytes 29 through 31):

primary queuing field for LRBs and LPRBs also on the active RB queue. It points to the previous RB for the task or, in the oldest RB in the queue, back to the TCB.

JOB PACK AREA QUEUE

cccc hhhhhh NM cccccccc SZ/STAB hhhhhhhh USE/EP hhhhhhhh PSW hhhhhhhh hhhhhhhh Q hhhhhh WT/LNK hhhhhhhh cccc hhhhhh NM cccccccc SZ/STAB hhhhhhhh WTL hhhhhhhh REQ hhhhhhhh TLPRB hhhhhhhh cccc hhhhhh NM cccccccc SZ/STAB hhhhhhhh USE/EP hhhhhhhh PSW hhhhhhhh hhhhhhhh Q hhhhhh WT/LNK hhhhhhhh

JOB PACK AREA QUEUE (MFT with subtasking only) identifies the next lines as the contents of the job pack area queue originating in the partition information block (PIB). cccc hhhhhh indicates the RB type and its starting address. The RB types are: FRB Finch request block LRB Loaded request block LPRB Loaded program request block NM ccccccc is the XRBNM field (bytes 0 through 7): Program name. SZ/STAB hhhhhhhh contains, in the first 4 digits, the XRBSZ field (bytes 8 and 9): number of contiguous doublewords for the RB, the program (if applicable), and associated supervisor work areas. contains, in the last 4 digits, the XSTAB field (bytes 10 and 11): flag bytes. USE/EP hhhhhhhh (LPRB, LRB Only) contains, in the first 2 digits, the XRBUSE field (byte 12): use count.

contains, in the last 6 digits, the XRBEP field (bytes 13 through 15): address of entry point in the program.

WTL hhhhhhhh (FRB Only) is the XRWTL field of the FRB (bytes 12 through 15): address of the most recent wait list element (WLE) on the WLE queue.

PSW hhhhhhhh hhhhhhhh (LPRB, LRB Only)
is the XRBPSW field (bytes 16 through
23): resume PSW.

REQ hhhhhhhh (FRB Only)
 is the XRREQ field of the FRB (bytes
 16 through 19): address of the TCB of
 the requesting task.

TLPRB hhhhhhhh (FRB Only) is the XRTLPRB field of the FRB (bytes 20 through 23): address of the LPRB built by the Finch routine for the requested program.

Q hhhhhh (LRB, LPRB Only) is the last 3 bytes of the XRBQ field (bytes 25 through 27):

- in an LPRB, the starting address of an LPRB for an entry identified by an IDENTIFY macro instruction.
- in an LRB, unused.

WT/LNK hhhhhhhh (LRB, LPRB Only) contains, in the first 2 digits, the XRBWT field (byte 28): wait count.

> contains, in the last 6 digits (bytes 29 through 31): primary queuing field for RBs. These RBs may be queued either on the job pack area queue or on the active RB queue. It points to the previous RB for the task or, in the oldest RB on the queue, back to the TCB.

P/P STORAGE BOUNDARIES hhhhhhhh TO hhhhhhhh FREE AREAS SIZE hhhhhh hhhhhhh GOTTEN CORE SIZE hhhhhhh hhhhh SAVE AREA TRACE CCCCCCCC WAS ENTERED VIA LINK (CALL) ddddd AT EP CCCCC... RO hhhhhhhh R6 hhhhhhhh HSA hhhhhhhh LSA hhhhhhhh RET hhhhhhhh EPA hhhhhhhh 5 2 hhhhhh WD1 hhhhhhhh Rl hhhhhhhh R2 hhhhhhhh R8 hhhhhhhh R4 hhhhhhhh R5 hhhhhhhh R3 hhhhhhhh R7 hhhhhhhh R9 hhhhhhhh R10 hhhhhhhh Rll hhhhhhhh R12 hhhhhhhh INCORRECT BACK CHAIN PROCEEDING BACK VIA REG 13

- P/P STORAGE BOUNDARIES hhhhhhhh TO hhhhhhhh gives the addresses of the lower and upper boundaries of a main storage area assigned to the task. This heading is repeated for every noncontiguous block of storage owned by the task.
- FREE AREAS SIZE

.

- hhhhhh hhhhh
- hhhhhh hhhhhh are the starting addresses of free areas and the size, in bytes, of each area contained within the P/P STORAGE BOUNDARIES field listed above.
- GOTTEN CORE SIZE
- hhhhhh hhhhhhh
- •
- -
- hhhhhh hhhhhhh (Printed only in a dump of a system with the MFT with subtasking option). These figures represent the starting addresses of the gotten areas (those areas obtained for a subtask through a supervisor issued GETMAIN macro instruction), and the size, in bytes, of each area contained within the P/P STORAGE BOUNDARIES field listed above. If main storage hierarchy support is included in the system, the values in this field can address storage in either hierarchy 0 or hierarchy 1, or both.
- SAVE AREA TRACE identifies the next lines as a trace of the save areas for the program.

cccccccc WAS ENTERED is the name of the program that stored register contents in the save area. This name is obtained from the RB.

- VIA LINK (CALL) ddddd indicates the macro instruction (LINK or CALL) used to give control to the next lower level module, and is the ID operand, if it was specified, of the LINK or CALL macro instruction.
- AT EP ccccc... is the entry point identified, which appears only if it was specified in the SAVE macro instruction that filled the save area.
- SA hhhhhh is the starting address of the save area.
- WD1 hhhhhhhh is the first word of the save area: use of this word is optional.
- HSA hhhhhhhh is the second word of the save area: starting address of the save area in the next higher level module. In the first save area in a job step, this word contains zeros. In all other save areas, this word must be filled.
- ISA hhhhhhhh is the third word of the save area (register 13): starting address of the save area in the next lower level module.
- RET hhhhhhhh is the fourth word of the save area (register 14): return address. Optional.

EPA hhhhhhhh is the fifth word of the save area (register 15): entry point to the invoked module. Optional.

- R0 hhhhhhhh R1 hhhhhhhh ... R12 hhhhhhhh are words 6 through 18 of the save area (registers 0 through 12): contents of registers 0 through 12 immediately after the linkage for the module containing the save area.
- INCORRECT BACK CHAIN indicates that the following lines may not be a save area because the second

word in this area does not point back to the previous save area in the chain.

PROCEEDING BACK VIA REG 13

indicates that the next 2 save areas are (1) the save area in the lowest level module, followed by (2) the save area in the next higher level module. The lowest save area is assumed to be the save area pointed to by register 13. These 2 save areas appear only if register 13 points to a full word boundary and does not contain zeros.

DATA SETS ***** N O T F O R M A T T E D ***** cccccccc UCB ddd hhhhhh DEB hhhhhh DCB hhhhhh **D/S FORMATTING TERMINATED**

DATA SETS

indicates that the next lines present information about the data sets for the task. For unopened data sets, only the ddname and UCB information are printed.

NOT FORMATTED

indicates that the abnormal termination dump routine confiscated storage (indicated by *CORE NOT AVAILABLE, LOC. hhhhhh-hhhhhh TAKEN); because DCBs may have been overlaid, or that the dump is for an OLTEP task. Data set information is not presented.

cccccccc is the name field (ddname) of the DD statement.

UCB ddd hhhhhh is the unit to which the data set was assigned, and the starting address of the UCB for that unit. If the data set was assigned to several units, the additional units are identified on following lines.

DEB hhhhhh

is the starting address of the DEB for the data set. Appears only for open data sets.

DCB hhhhhh

is the starting address of the DCB for the data set. Appears only for open data sets.

D/S FORMATTING TERMINATED

indicates that no more data set information is presented because a DCB is incorrect, possibly because a program incorrectly modified it.

TRACE T	TRACE TABLE - STARTING WITH OLDEST ENTRY											
dddd dddd dddd	I/O ddd SIO ddd SVC ddd	PSW հհհհհհհ հհհհհհհ CC = d PSW հհհհհհհ հհհհհհհ	CAW hhhhhhhh RG 0 hhhhhhhh	CSW OLD CSW RG 1	հհհհհհհհ հեհհհհ հեհհհհհհ հհհհհհհհ		ATUS hhhh)					

TRACE TABLE -- STARTING WITH OLDEST ENTRY identifies the next lines as the contents of the trace table. Each entry is presented on one line. The types of entries are:

I/O Input/output interruption entry

SIO Start input/output (SIO) entry

SVC Supervisor call (SVC) interruption entry

dddd

is the number assigned to each entry. The oldest entry receives the number 0001.

I/O ddd is the channel a

is the channel and unit that caused the input/output interruption.

- PSW hhhhhhh hhhhhhh is the program status word that was stored when the input/output interruption occurred.
- CSW hhhhhhh hhhhhhh is the channel status word that was stored when the input/output interruption occurred.
- SIO ddd is the device specified in the SIO instruction.

CC=d is the condition code resulting from execution of the SIO instruction. Zero indicates a successful start.

- CAW hhhhhhhh is the channel address word used by the SIO instruction.
- OLD CSW hhhhhhhh hhhhhhhh is the channel status word stored during execution of an SIO operation. It appears when CC is not equal to 1.
- CSW STATUS hhhh is the status portion of the channel status word stored during execution of an SIO instruction. Appears when CC is equal to 1.

SVC ddd is the SVC instruction's operand.

PSW hhhhhhh hhhhhhh is the PSW stored during the SVC interruption. An F in the fifth digit of the first word identifies the entry as representing a task switch.

RG 0 hhhhhhhh is the contents of register 0 as passed to the SVC routine.

RG 1 hhhhhhhh is the contents of register 1 as passed to the SVC routine.

REGS AT ENTRY TO ABEND (SNAP)											
FLTR 0-6	հհհհհհհհհհհհհ	հիհիհիհիհիհի	hhhhhhhhhhhhhhh	հեհեհեհեհեհեհ							
REGS 0-7 REGS 8-15	հհհհհհհ հհհհհհհ հհհհհհհ հհհհհհհ	հեններին հեններին հենների հենների	իհիհիհի հիհիհիհ հիհիհիհ հիհիհիհ	հհհհհհհհ հհհ հ հհհհ հհհհհհհ հհհհհ հ							

REGS AT ENTRY TO ABEND OF REGS AT ENTRY TO SNAP

identifies the next 3 lines as the contents of the floating point and general registers when the abnormal termination routine received control in response to an ABEND macro instruction or when the SNAP routine received control in response to a SNAP macro instruction. FLTR 0-6 is the contents of floating point registers 0, 2, 4, and 6.

REGS 0-7 is the contents of general registers 0 through 7.

REGS 8-15 is the contents of general registers 8 through 15.

NUCLEUS							
hhhhh	հհհհհհհհ հհհհհհհ	hhhhhhh hhhhhhhh	հհհհհհհ	hhhhhhh	hhhhhhh	hhhhhhh	***************************************
hhhhhh LINE	hhhhhhhh hhhhhhhh hhhhhh SAME A	hhhhhhhh hhh <mark>hhhhh</mark> As above	hhhhhhhh	hhhhhhh	hhhhhhh	hhhhhhh	***************************************
hhhhhh LINES	hhhhhhhh hhhhhhhh hhhhhh-hhhhhh	hhhhhhhh hhhhhhhh SAME AS ABOVE	հհհհհհհ	hhhhhhh	hhhhhhh	hhhhhhh	*****************
hhhhhh P/P STORA	hhhhhhhh hhhhhhhh NGE	hhhhhhhh hhhhhhhh	hhhhhhh	հհհհհհհ	հհհհհհհ	hhhhhhhh	***************************************
hhhhh	հհհհհհհհ հհհհհհհ	հհհհհհհ հհհհհհհ	hhhhhhh	hhhhhhh	հհհհհհհ	hhhhhhh	***************************************
hhhhh	hhhhhhhh hhhhhhhh	hhhhhhhh hhhhhhhh	hhhhhhh	hhhhhhh	hhhhhhh	hhhhhhh	*ccccccccccccccccccccccccccccccccc
hhhhhh LINES	hhhhhhhh hhhhhhh hhhhhh-hhhhhh	hhhhhhhh hhhhhhhh SAME AS ABOVE	հհհհհհհ	hhhhhhh	hhhhhhh	հհհհհհհ	***************************************
hhhhh	hhhhhhhh hhhhhhh END OF DUMP	հհհհհհհհ հհհհհհհ	hhhhhhhh	hhhhhhhh	հհհհհհհ	hhhhhhh	*======================================

The content of main storage is given under 2 headings: NUCLEUS and P/P STORAGE. Under these headings, the lines have the following format:

- First entry: the address of the initial byte of main storage contents presented on the line.
- Next 8 entries: 8 full words (32 bytes) of main storage in hexadecimal.
- Last entry (surrounded by asterisks): the same 8 full words of main storage in EBCDIC. Only A through Z, 0 through 9, and blanks are printed; a period is printed for anything else. An exception occurs in the printed lines representing the ABDUMP work area. The contents of the ABDUMP work area during the printing of EBCDIC characters

differs from the contents during printing of the hexadecimal characters because a portion of the work area is used to write lines to the printer. This exception should not create any problems since the contents of the ABDUMP work area is of little use in debugging.

The following lines may also appear:

- LINES hhhhhhhh-hhhhhhh SAME AS ABOVE are the starting addresses of the first and last line of a group of lines that are identical to the line immediately preceding.
- LINE hhhhhh SAME AS ABOVE is the starting address of a line that is identical to the line immediately preceding.

NUCLEUS identifies the next lines as the contents of the control program nucleus.

P/P STORAGE

identifies the next lines as the contents of the main storage area assigned to the task (problem program).

END OF DUMP indicates that the dump or snapshot is completed.

Guide to Using an ABEND/SNAP Dump (MFT)

<u>Cause of Abnormal Termination</u>: Evaluate the user (USER Decimal code) or system (SYSTEM=hex code) completion code using Appendix C or the publication <u>Messages and</u> <u>Codes</u>.

Active RB Queue: The first RB shown on the dump represents the oldest RB on the queue. The RB representing the load module that had control when the dump was taken is third from the bottom. The last RB represents the ABDUMP routine, and the second from last, the ABEND routine. The names of load modules represented in the active RB queue are given in the RB field labeled NM in the dump. Names of load modules in SVC routines are presented in the format:

r	
NM	SVC-mnnn
L	Jan 1997 - J

where m is the load module number (minus 1) in the routine and nnn is the signed decimal SVC number. The last two RBs on an ABEND/SNAP dump will always be SVRBs with edited names <u>SVC-105A</u> (ABDUMP--SVC 51) and <u>SVC-401C</u> (ABEND--SVC 13).

Resume PSW: The resume PSW field is the fourth entry in the first line of each RB printout. It is identified by the subheading PSW. For debugging purposes, the resume PSW of the third RB from the bottom, on the dump, is most useful. The last three characters of the first word give the SVC number or the I/O device address, depending on which type of interruption caused the associated routine to lose control. It also provides the CPU state at the time of the interruption (bit 15), the length of the last instruction executed in the program (bits 32,33), and the address of the next instruction to be executed (bytes 5-8). Load List and Job Pack Area Queue: The load module that had control at the time of abnormal termination may not contain the instruction address pointed to by the resume PSW. In that case, look at the RBs on the load list and on the job pack area queue (MFT with subtasking). Compare the instruction address with the entry points of each load module (shown in the last 3 bytes of the field labeled USE/EP). The module which contains the instruction pointed to by the resume PSW is the one in which abnormal termination occurred. The name of the load module is indicated in the field labeled NM.

Trace Table: Entries in the trace table reflect SIO, I/O, and SVC interruptions and task switching. <u>SIO</u> entries can be used to locate the CCW (through the CAW), which reflects the operation initiated by an SIO instruction. If the SIO operation was not successful, the CSW STATUS portion of the entry will show you why it failed.

<u>I/O</u> entries reflect the I/O old PSW and the CSW that was stored when the interruption occurred. From the PSW, you can learn the address of the device on which the interruption occurred (bytes 2 and 3), the CPU state at the time of interruption (bit 15), and the instruction address where the interruption occurred (bytes 5-8). The CSW provides you with the unit status (byte 4), the channel status (byte 5), and the address of the previous CCW plus 8 (bytes 0-3).

<u>SVC</u> entries provide the SVC old PSW and the contents of registers 0 and 1. The PSW offers you the hexadecimal SVC number (bits 20-31), the CPU mode (bit 15), and the address of the SVC instruction (bytes 5-8). The contents of registers 0 and 1 are useful in that many system macro instructions use these registers for parameter information. Contents of registers 0 and 1 for each SVC interruption are given in Appendix A.

A task switch entry is similar to an SVC entry, except that words 3 and 4 of the entry contain the address of the TCBs for the "new" and "old" tasks being performed, respectively. The trace table entries for one particular task are contained between sets of two task switch entries. Word 3 of the beginning task switch entry and word 4 of the ending task switch entry point to the TCB for that task. Task switch entries are identified by a fifth digit of 'F'.

Notes: If an ABEND macro instruction is issued by the system when a program check interruption causes abnormal termination, an SVC entry does not appear in the trace table, but is reflected in the PSW at entry to ABEND. Dumps issued by systems with MFT contain only the last four characters of the module name in the RB APSW field. You cannot distinguish between IFG0xxxx and IGG0xxxx. After an SVC 19 has been issued, the OPEN where-to-go table should be checked for the module name. Free Areas: ABEND/SNAP dumps do not print out areas of main storage that are available for allocation. Since the ABEND routine uses some available main storage, the only way you can determine the amount of free storage available when abnormal termination occurred is to re-create the situation and take a stand-alone dump.

ABEND/SNAP Dump (MVT)

MVT dumps differ from PCP and MFT dumps in the addition of detailed main storage control information, the omission of a complete main storage dump, and the omission of a trace table in ABEND dumps. MVT dumps occur immediately after an abnormal termination, provided an ABEND or SNAP macro instruction was issued and proper dump data sets were defined. However, if a system failure has occurred and a SYS1.DUMP data set has been defined and is available, a full storage image dump is provided, as explained in the section headed "Storage Image Dump."

With MVT's subtask creating capability, you may receive one or more partial dumps in addition to a complete dump of the task that caused abnormal termination. A complete dump includes all control information associated with the terminating task and a printout of the load modules and subpools used by the task. A partial dump of a task related to the terminating task includes only control information. A partial dump is identified by either ID=001 or ID=002 printed in the first line of the dump. Figure 18 shows the key areas of a complete dump.

In systems with MVT, you can effect termination of a job step task upon abnormal termination of a lower level task. To do this, you must either terminate each task upon finding an abnormal termination completion code issued by its subtask or pass the completion code on to the next higher level task.

For a discussion of a formatted ABEND dump using the telecommunications access method (TCAM) in an MVT environment, see <u>IBM System/360 Operating System: TCAM</u> <u>Program Logic Manual</u>, GY30-2029. References to other TCAM debugging aids are found in Appendix J.

Invoking an ABEND/SNAP Dump (MVT)

ABEND/SNAP dumps issued by systems with MVT are invoked in the same manner as those under systems with PCP and MFT. They result from an ABEND or SNAP macro instruction in a system or user program, accompanied by a properly defined data set. In the case of a system failure, the damage assessment routine (DAR) attempts to write a storage image dump to the SYS1.DUMP data set. A full explanation of storage image dumps may be found in the section headed "Storage Image Dump." The instructions that invoke an ABEND/SNAP dump in MVT environment are the same as those given in the preceding topic for systems with MFT. However, some additional considerations must be made in requesting main storage and direct access space.

MVT Considerations: In specifying a region size for a job step subject to abnormal termination, you must consider the space requirements for opening a SYSABEND or SYSUDUMP data set (if there is one), and loading the ABDUMP routine and required data management routines. This space requirement can run as high as 6000 bytes.

Direct access devices are used frequently for intermediate storage of dump data sets in systems with MVT. To use direct access space efficiently, the space for the dump data set should be varied, depending on whether or not abnormal termination is likely. A small quantity should be requested if normal termination is expected. To prevent termination of the dump due to a lack of direct access space, always specify an incremental (secondary) quantity when coding a SPACE parameter for a dump data set. You can obtain a reasonable estimate of the direct access space required for an ABEND/SNAP dump by adding, (1) the number of bytes in the nucleus, (2) the part of the system queue space required by the task (9150 bytes is a sufficient estimate), and (3) the amount of region space occupied by the task. Multiply the sum by 4, and request this amount of space in 1024-byte blocks.

This formula gives the space requirements for one task. Request additional space if partial dumps of subtasks and invoking tasks will be included.

Contents of an ABEND/SNAP Dump (MVT)

This explanation of the contents of ABEND/SNAP dumps issued by systems with MVT is interspersed with sample sections from an ABEND dump. Capital letters represent the headings found in all dumps, and lowercase letters, information that varies with each dump. The lowercase letter used indicates the mode of the information and the number of letters indicates its length:

- h represents 1/2 byte of hexadecimal information
- d represents 1 byte of decimal information
- c represents a 1-byte character

You may prefer to follow the explanation on your own ABEND or SNAP dump.

PSW 41 F41EV IN LARMAN FF040007 S0000-08 TG D7778 WF 00076767 AFF 0000000 PF 0000000 PF 0000000 PF 0000000 PF 0000000 PF 0000000 PF 00000000		JOB 1	PCT41		STEP F	XSTEP		TIME	007404	> n4	TF 993	66					PAGE	0001
TCB 077/72 VFF. 077/72 VFF. 077/72 VFF. 07000000 JFF. 00000000 JFF. JFF. 00000000 JFF. 00000000 JFF. JFF. <t< td=""><td></td><td>COMPL</td><td>ETTON C</td><td>JUE</td><td>SYSTEM</td><td>= 837</td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td></t<>		COMPL	ETTON C	JUE	SYSTEM	= 837												
WS 0 101176 PF-FL C FORTAGE FL C DONDORON LIS DOSTRON JLA DOSTRON TO DOSTRON DIA DOSTRON TO DOSTR		PSW A	T ENTRY	TO ABE	ND FF04	0007 5	0000408											
P#B 030789 WTV 00000000 WTV 00000000 WTVV 00000000 WTVV 00000000 WTVVV 00000000 WTVVVVVVVVVVVVVVVVVVVVVVVVVVVVVVVVVVVV		TCB	025028	455 FSA LTC	01031738 01060768 00000000	РК- ТСВ ТОЕ	FLG F085 0000 0000	0409 0000 0000	FLG TME FCR	00000 00000 00030	000 000 484	LLS JST STA	00030980 0002F028 00000000	JLB NTC D-PQ	0000000 0000000 F 0003266	0 JPQ 0 OTC 8 SQS	000301E8 00030508 0002EAA0	
Office Optimization Optimization Optimization PMR 030904 Style Optimization Optimization Optimization Optimization Style Optimization Optimization Optimization Optimization Style Optimization Optimization Optimization Optimization Style Optimization Optimization Optimization Optimization Optimization Optimization Optimization Optimization Opti		ACTIV	FRAS															
O/TIS OPTOBOO WT-LKE ONINGORF SV#R 02F0F0 TATLE MOSTRAGE WT-LKE ONINGOFF WT-LKE ONINGOF		PRB	030NF8						WC-9	5Z-STA	8 0004	0082	FL-CDE	00031	290 PSW	FFF50006	7003553E	
O/TTF 00030FF VT-LNE 00030FF 000000FF 000000FF 00030FF4 00030FF6		PPB	030988						WC-1	57-STA	8 0004	0002	FL-CDE	00030	ERO PSW	FFF50037	5207EC4A	
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02ED20 00000050 00000000 <			032380 032280 032250 032380 0321F0 032380	SZ 00 SZ 00 S7 00 S7 00 SZ 00 SZ 00	000010 000010 000010 000010 000010 000010	NO 00 NO 00 NO 00 NO 00 NO 00 NO 00	000001 000001 000001 000001 000001 000001	A 8 8 8 8 8	000001 000021 000018 0000021 0000021 0000021 0000021	10 10 10 10	0006C9 0007FA 0007F4 0007F8 0007F8 0007E3 0007E3	80 00 A0 80 00 A0 10		- 0		0140030	-	
02ED20 00000050 0000000 00002020A 000028E0 0F000000 0002F028 0402EED4 98000000 +		DFB																
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				× .														
												۰, ^۲						

Figure 18. Sample of Complete ABEND Dump (MVT) (Part 1 of 2)

DER PAGE 0002 00003050 07000750 0000050 0000070 24003000 03027028 0400000 8800000 04027ER0 18002648 0000037 0009038 00080004 18002648 0000037 00090040 00083004 18002648 0000043 00090044 00033004 18002648 0000045 00090046 00083004 18002648 0000045 00090046 00083004 18002648 0000047 00090046 00083004 18002648 0000047 00090046 00083004 18002648 0000047 00090046 00083004 18002648 0000047 00090046 00083004 18002648 0000048 00090046 00083004 18002648 0000048 00090046 02FEA0 02FEC0 02FFE0 02FF00 00000050 0000000 0000020F 000114F0 0F00000 1000000 0000000 FF0394F4 00080032 18002448 0000003E 0000003F 00080004 18002648 00000040 00090041 *..... 02EE20 028F40 028F40 028F60 028F80 0008000A 18002648 00000040 00090041 0004000A 18002648 00000042 00090043 0008000A 18002648 00000046 00090045 0008000A 18002648 00000046 00090047 0008000A 18002648 00000046 00090048 0008000A 18002648 00000046 00090048 *.......... *..... 02FFA0 02FFC0 02FFE0 *.....ARAJCD6.* Sec. 1 TIOT STEP EXSTEP JOB IPCT41 PGH=+.00 00 14040101 00230E00 80002648 SYSABEND F TO6F001 F TNL IN 00240900 00240000 00250100 00 14040100 R0002648 80002648 00 14040100 80003984 00 00 14000000 SY SPUNCH 00250800 00000000 14040101 SYSIN 00250A00 80002648 MSS ******* FOF ******* ********** ***************** DQF ************** FLGS NSPOE FOE BLK LN NDQE NFQE LN 031738 60 031740 251 031250 00035000 00035000 00000800 00031050 00000000 00000508 00035800 00035800 00017000 00000000 0000000 00000108 031740 00 031488 252 031400 00060800 00060800 00000800 00030878 00000000 00000588 0006C000 0006C000 00000800 00030308 0006C800 0006C800 00000800 0002F388 00000000 00000480 00068800 00068800 00000800 0000000 00000000 000001A0 031488 031400 031400 031488 C 0 60 000000 000 00060000 00000020 000000 000 00050000 00060748 00000800 00000000 FIRST 00031460 LAST 00031460 0004C800 LFB 0004C800 NPQ 0 00030508 PSI 00039000 RAD 0 D-POE 00032668 031460 FFB 0004C800 TCB 00030508 NPQ 00000000 RAD 00035000 PPQ 00000000 FLG 0000 POF **FROF 04C800** NFR 00031460 PFR 00031460 SZ 0001F000 OCR TRACE MA'J 0311C8 NMAJ 00030100 PMAJ 0001C6A0 FMIN 00031088 NM SYSDSN MIN 031088 FOEL 00031698 PMIN 000311C8 NHIN 00000000 NM FF SYS1.MACLIB NOEL 0000000 POEL 80031088 TCR 00030508 SVR8 00030100 MAJ 030100 NMAJ 0000000 PMAJ 000311C8 FMIN 00030140 SYSIEA01 MIN 030140 FOFI 00030190 PHIN 00030100 NMIN DODDOOD NH FO IFA NOFL 00000000 POEL 00030140 TC8 0002F028 SVRB 0002EBE8 SAVE AREA TRACE 5A 060768 W01 00000000 HSA 00000000 LSA 00000000 RET 00000000 EPA 00000000 RO 0000000 00000000 R2 R8 00000000 R3 R9 00000000 R4 00000000 R10 00000000 R5 00000000 R11 00000000 R6 0000000 R12 0000000 00000000 00000000 INTERRUPT AT OTEC44 PROCEEDING BACK VIA REG 13 SA 039500 W01 957095FF HSA 70004780 LSA 95789180 REJ 80064710 EPA 958C1811 RO 5203936E R1 9207E3A0 R7 00060688 82 00060570 00060780 R3 R9 000396F4 00000FD9 0003966 00060570 R6 7F0605CC R12 0007EC10 RA R10 0007EC10 R11 5207F434 R0 FF000000 R6 00000000 R12 FF000000 S۵ 004780 WD1 47900000 HSA EE000000 LSA 00000000 RET 00000000 EPA 47400000 00000000 47C00000 82 98 00000000 FF000000 83 89 47800000 R4 FF000000 R10 00000000 R5 00000000 R11 47000000 NUCLEUS 000000 0000000 0000000 0000000 0000000 00000868 0000000 FF040080 80038724 0000FF00 00000000 FF060336 8000000 0836E88C 0001389C 00040000 0000F678 *..... 000020 FF050001 4007EC3C FFF50001 02036CF2 0000A7C8 0000000 000725A0 00000B68 . - 1 2

Figure 18. Sample of Complete ABEND Dump (MVT) (Part 2 of 2)

JOB CCCCCCC STEP CCCCCCC TIME dddddd DATE COMPLETION CODE SYSTEM = hhh (or USER = dddd) PSW AT ENTRY TO ABEND (SNAP) hhhhhhhh hhhhhhhh	ddddd ID - ddd PAGE dddd
JOB ccccccc is the job name specified in the JOB statement.	task being abnormally terminated, up to and including the job step task.
STEP cccccccc is the step name specified in the EXEC statement for the problem program associated with the task being dumped.	PAGE dddd is the page number. Appears at the top of each page. Page numbers begin at 0001 for each task or subtask dumped.
TIME dddddd is the hour (first 2 digits), minute (next 2 digits), and second (last 2 digits) when the abnormal termination dump routine began processing.	COMPLETION CODE SYSTEM=hhh or COMPLETION CODE USER=dddd is the completion code supplied by the control program (SYSTEM=hhh) or the
DATE ddddd is the year (first 2 digits) and day of the year (last 3 digits). For example, 67352 would be December 18, 1967.	problem program (USER=dddd). For a system completion code of 0C0, denoting an imprecise interruption on a Model 91 or Model 195, bits 20 through 29 of the CMP field of the TCD are significant for a Model 91, while all the bits (20 through 31) are used
 ID=ddd is an identification of the dump. For dumps requested by an ABEND macro instruction, this identification is: Absent if the dump is of the task being abnormally terminated. 	by a Model 195. For further discussion on interpreting the imprecise interruption configuration refer to the publication <u>IBM</u> <u>System/360 Operating</u> <u>System: Supervisor Services</u> , Order No. GC28-6646.
 001 if the dump is of a subtask of the task being abnormally terminated. (Note that, when a task is abnormally terminated, its subtasks are also abnormally 	PSW AT ENTRY TO ABEND hhhhhhhh hhhhhhh or PSW AT ENTRY TO SNAP hhhhhhhh hhhhhhh is the PSW for the problem program or control program routine that had control when abnormal termination was

• 002 if the dump is of a task that directly or indirectly created the

terminated.)

ABEND/SNAP Dump (MVT) 49

requested, or when the SNAP macro instruction was executed. It is not necessarily the PSW at the time the error condition occurred.

TCB	hhhhhh	RBP	hhhhhhh	PIE	hhhhhhh	DEB	hhhhhhh	TIO	հհհհհհհ	CMP	hhhhhhh	TRN	hhhhhhh	ı
		MSS	հհհհհհհ	PK-FL	🙃 hhhhhhhh	FLG	hhhhhhh	LLS	հհհհհհհ	JLB	hhhhhhh	JPQ	hhhhhhh	า
		RG 0-	7 hhhhł	hhh	hhhhhhhh	hhhhi	hhhh hl	hhhhhh	hhhhhi	nhh	hhhhhhh	hhhh	hhhh	hhhhhhh
		RG 8-	15 hhhhł	hhh	հհհհհհհ	hhhh	hhhh h	hhhhhh	hhhhh!	nhh	hhhhhhh	hhhh	հհհհ	hhhhhhh
		FSA	հհհհհհհ	TCB	hhhhhhh	TME	hhhhhhh	JST	hhhhhhh	NTC	հհհհհհհ	OTC	hhhhhhh	1
		LTC	hhhhhhh	IQE	hhhhhhh	ECB	հհհհհհհ	STA	hhhhhhh	D-PQE	hhhhhhh	sqs	hhhhhhh	ı
		NSTAE	hhhhhhh	TCT	hhhhhhh	USER	hhhhhhh	DAR	հհհհհհհ	RESV	հհհհհհհ	JSCB	hhhhhhh	`

TCB hhhhhh

is the starting address of the TCB.

RBP hhhhhhhh

is the TCBRBP field (bytes 0 through 3): starting address of the active RB queue and, consequently, the most recent RB on the queue.

PIE hhhhhhhh

is the TCBPIE field (bytes 4 through 7): starting address of the program interruption element (PIE) for the task; however, in an abnormal termination dump for the task causing the abnormal termination, zeros. The field is zeroed by the ABEND routine to prevent interruptions during dumping.

DEB hhhhhhhh

is the TCBDEB field (bytes 8 through 11): starting address of the DEB queue. Under the heading DEB in the dump, the prefix section for the first DEB in the queue is presented in the first 8-digit entry on the first line. The 6-digit entry at the left of each line under DEB is the address of the second column on the line, whether or not the column is filled. The contents of the TCBDEB field may differ in the main storage printout from what appears in the TCBDEB field of the formatted section. This occurs when the number of extents specified in the DEB for the dump data set is not sufficient to complete ABDUMP processing. When the dump of main storage is given, the END OF VOLUME routine may have built another DEB having additional extents for the dump data set and dequeued the original DEB. Therefore, the TCBDEB field in the main storage printout may contain the address of the new DEB built by END OF VOLUME.

TIO hhhhhhhh

is the TCBTIO field (bytes 12 through 15): starting address of the TIOT.

CMP hhhhhhhh

is the TCBCMP field (bytes 16 through 19): task completion code or contents of register 1 when the dump was requested. System codes are given in the third through fifth digits and user codes in the sixth through eight digits.

For a system completion code of 0C0, denoting an imprecise interruption on a Model 91 or Model 195, digits six through eight contain the imprecise interruption configuration. These bit settings represent bits 16 to 27 of the program old PSW (location X'28') at the time of the program exception. Therefore, in the TCB CMP field, bits 20 to 29 are significant for a Model 91 while a Model 195 makes use of all the bits (20 to 31). For further discussion on interpreting the imprecise interruption configuration refer to the publication IBM System/360 Operating System: Supervisor Services, Order No. GC28-6646.

TRN hhhhhhhh

is the TCBTRN field (bytes 20 through 23): starting address of the control core (table) for controlling testing of the task by TESTRAN.

MSS hhhhhhhh is the TCBMSS field (bytes 24 through 27): starting address of SPQE most recently added to the SPQE queue.

contains, in the last 6 digits, the first 3 bytes of the TCBFLGS field (bytes 29 through 31): first 3 flag bytes.

FLG hhhhhhhh
 contains, in the first 4 digits, the
 last 2 bytes of the TCBFLGS (bytes 32
 and 33): last 2 flag bytes.

contains, in the next 2 digits, the

TCBLMP field (byte 34): limit priority (converted to an internal priority, 0 to 255).

contains, in the last 2 digits, the TCBDSP field (byte 35): dispatching priority (converted to an internal priority, 0 to 255).

LLS hhhhhhhh

is the TCBLLS field (bytes 36 through 39): starting address of the load list element most recently added to the load list.

JLB hhhhhhhh

is the TCBJLB field (bytes 40 through 43): starting address of the DCB for the JOBLIB data set.

JPQ hhhhhhhh

is the TCBJPQ field (bytes 41 through
47): when translated into binary
bits:

- Bit 0 is the purge flag.
- Bits 1 through 7 are reserved for future use and are zeros.
- Bits 8 through 31 are the starting address of the queue of CDEs for the job pack area control queue, which is for programs acquired by the job step.

The TCBJPQ field is used only in the first TCB in the job step; it is zeros for all other TCBs.

RG 0-7 and RG 8-15

is the TCBGRS field (bytes 48 through 111): contents of general registers 0 through 7 and 8 through 15, as stored in the save area of the TCB when a task switch occurred. These 2 lines appear only in dumps of tasks other than the task in control when the dump was requested.

FSA hhhhhhhh

contains, in the first 2 digits, the TCBQEL field (byte 112): count of enqueue elements.

contains, in the last 6 digits, the TCBFSA field (bytes 113 through 115): starting address of the first problem program save area. This save area was set up by the control program when the job step was initiated.

TCB hhhhhhhh

is the TCBTCB field (bytes 116 through 119): starting address of the next lower priority TCB on the TCB queue or, if this is the lowest priority TCB, zeros. Page of GC28-6670-5,6, Revised April 16, 1973, By TNL: GN28-2545

TME hhhhhhh is the TCBTME field (bytes 120 through 123): starting address of the timer element created when an STIMER macro instruction is issued by the task.

JST hhhhhhhh

is the TCBJSTCB field (bytes 124 through 127): starting address of the TCB for the job step task. For tasks with a protection key of zero, this field contains the starting address of the TCB.

NTC hhhhhhhh

is the TCBNTC field (bytes 128 through 131): the starting address of the TCB for the previous subtask on this subtask queue. This field is zero in the job step task, and in the TCB for the first subtask created by a parent task.

OTC hhhhhhhh

is the TCBOTC field (bytes 132 through 135): starting address of TCB for the parent task. In the TCB for the job step task, this field contains the address of the initiator.

LTC hhhhhhhh

is the TCBLTC field (bytes 136 through 139): starting address of the TCB for the most recent subtask created by this task. This field is zero in the TCB for the last subtask of a job step, or in a TCB for a task that does not create subtasks.

IQE hhhhhhhh

is the TCBIQE field (bytes 140 through 143): starting address of the interruption queue element (IQE) for the ETXR exit routine. This routine is specified by the ETXR operand of the ATTACH macro instruction that created the TCB being dumped. The routine is to be entered when the task terminates.

ECB hhhhhhhh

is the TCBECB field (bytes 144 through 147): starting address of the ECB to be posted by the control program at task termination. This field is zero if the task was attached without an ECB operand.

TSF hhhhhhhh

is the TCBSFLG field (bytes 94 through 97). The first two digits, byte 94, contain internal TSO flags, the next two digits, byte 95, contain the stop count, and the last two bytes, 96 and 97, contain the limit and dispatching priority of the TSO task. D-PQE hhhhhhhh

is the TCBPQE field (bytes 152 through 155): starting address minus 8 bytes of the dummy PQE. This field is passed by the ATTACH macro instruction to each TCB in a job step.

SQS hhhhhhhh is the TCBAQE field (bytes 156 through 159): starting address of the allocation queue element (AQE).

NSTAE hhhhhhhh contains, in the first 2 digits, STAE flags (byte 160).

contains, in the last 6 digits, the TCBNSTAE field (bytes 161 through 163): starting address of the current STAE control block for the task. This field is zero if STAE has not been issued.

- TCT hhhhhhhh is the TCBTCT field (bytes 164 through 167): address of the Timing Control Table (TCT).
- USER hhhhhhhh is the TCBUSER field (bytes 168 through 171): to be used as the user chooses.
- DAR hhhhhhh contains, in the first two digits, Damage Assessment Routine (DAR) flags (byte 172).
- RESV hhhhhhhh reserved for future use.

JSCB hhhhhhh is the TCBJSCB field (bytes 180 through 183): the last three bytes contain the address of the Job Step Control Block.

ACTIVE RBS

			APSW hhhhhhh WT-LNK hhhhhhh		FAB hhhhhhhh	ccccc hh	hhhhhh PSW	հհհհհհհ	հհհհհհ
R	G 0-7 G 8-15 XTSA	hhhhhhh hhhhhhhh hhhhhhhh hhhhhhh	հհիհհհհ հհհհհհհ հհհհհհհ հհհհհհհ	hhhhhhh hhhhhhh hhhhhhh hhhhhhh	հհհհհհհ հհհհհհհ հհհհհհհ հհհհհհհ	հհհհհեհ հհհհհհհ հհհհհհհ	հևհհհհհհ դհհհհհհհ հծհհհհհհ	հհհհհհհ հհհհհհհ հհհհհհհ	հհհհհհհհ հհհհհհհ հհհհհհհ

ACTIVE RBS

identifies the next lines as the contents of the active RBs queued to the TCB, beginning with the oldest RB first.

cccc hhhhhh indicates the RB type (cccc) and starting address (hhhhhh).

The RB types are:

PRB program request block IRB interruption request block SVRB supervisor request block

cccccc hhhhhhhh

indicates the RB's function (cccccc)
and bytes 0 through 3 of the RB
(hbhhhhhh):

- RESV hhhhhhh indicates PRB or SVRB for resident routines. Bytes 0 through 3 are reserved for later use and contain zeros.
- TAB-LN hhhhhhh indicates SVRB for transient routines. The first 4 digits contain the RBTABNO field (bytes 0 and 1): displacement from the beginning of the transient area control table (TACT) to the entry for the module represented by the RB. The last 4 digits contain the RBRTLNTH field (bytes 2 and 3): length of the SVC routine.
- FL-PSA hhhhhhh indicates IRB. The first 2 digits contain the RBTMFLD field (byte 0): indicators for the timer routines. This byte contains zeros when the IRB does not represent a timer routine. The last 6 digits contain the RBPSAV field (bytes 1 through 3): starting address of the problem program register save area (PSA).

APSW hhhhhhhh

is the RBABOPSW field (bytes 4 through 7):

- In PRB, right half of the problem program's PSW when the interruption occurred.
- In IRB or SVRB for type II SVC routines, right half of routine's PSW during execution of ABEND or ABTERM, or zeros.
- In SVRB for type III or IV SVC routines, right half of routine's PSW during execution of ABEND or ABTERM, or the last four characters of the name of the requested routine. (The last two characters give the SVC number.)
- WC-SZ-STAB hhhhhhhh contains, in the first 2 digits, the RBWCSA field (byte 8): wait count in effect at time of abnormal termination of the program.

contains, in the second 2 digits, the RBSIZE field (byte 9): size of the RB in doublewords.

contains, in the last 4 digits, the RBSTAB field (bytes 10 and 11): status and attribute bits.

cccccc hhhhhhhh

indicates the RB's function (cccccc) and bytes 12 through 15 of the RB (hhhhhhhh):

• FL-CDE hhhhhhhh indicates SVRB for resident routines, or PRB. The first 2 digits contain the RBCDFLGS field (byte 12): control flags. The last 6 digits contain the RBCDE field (bytes 13 through 15): starting address of the CDE for the module associated with this RB.

- EPA hhhhhhh is the RBEP field of an IRB (bytes 12 through 15): entry-point address of asynchronously executed routine.
- TQN hhhhhhh indicates SVRB for transient routines. Is the RBSVTQN field (bytes 12 through 15): address of the next RB in the transient control queue.
- PSW hhhhhhh hhhhhhh
 is the RBOPSW field (bytes 16 through
 23): resume PSW.
- Q/TTR hhhhhhh
 - In PRBs and SVRBs for resident routines, contains zeros in the first 2 digits. The last 6 digits contain the RBPGMQ field (bytes 25 through 27): queue field for serially reusable programs (also called the secondary queue).
 - In IRBs, contains the RBUSE field in the first 2 digits (byte 24): count of requests for the same exit (ETXR). The RBIQE field in last 6 digits (bytes 25 through 27): starting address of the queue of interruption queue elements (IQE), or zeros in the first 4 digits and the RBIQE field in the last 4 digits (bytes 26 and 27): starting address of the request queue elements.

- In SVRBs for transient routines the first 2 digits contain the RBTAWCSA field (byte 24): number of requests (used if transient routine is overlaid) and the last 6 digits, the RBSVTTR field (bytes 25 through 27): relative track address for the SVC routine.
- WT-LNK hhhhhhhh contains, in the first 2 digits, the RBWCF field (byte 28): wait count.

contains, in the last 6 digits, the RBLINK field (bytes 29 through 31): starting address of the previous RB on the active RB queue (primary queuing field) or, if this is the first or only RB, the starting address of the TCB.

RG 0-7 and RG 8-15 is the RBGRSAVE field (bytes 32 through 95): in SVRBs and IRBs, contents of registers 0 through 15.

EXTSA

- In IRBs, contains the RBNEXAV field in the first 8 digits (bytes 96 through 99): address of next available interruption queue element (IQE), and in the remaining digits, the interruption queue element work space (up to 1948 bytes).
- In SVRBs, contains the RBEXSAVE field (bytes 96 through 143): extended save area for SVC routine.

LOAD LIST					
NE hhhhhhh	RSP-CDE hhhhhhhh	NE hhhhhhhh	RSP-CDE hhhhhhhh	NE hhhhhhhh	RSP-CDE hhhhhhhh

LOAD LIST

identifies the next lines as the contents of the load list elements (LLEs) queued to the TCB by its TCBLLS field. The contents of 3 load list elements are presented per line until all elements in the queue are shown.

NE hhhhhhhh contains, in the first 2 digits, LLE byte 0: zeros.

> contains, in the last 6 digits, LLE bytes 1 through 3: starting address of the next element in the load list.

RSP-CDE hhhhhhhh

contains, in the first 2 digits, LLE byte 4: the count of the number of requests made by LOAD macro instructions for the indicated load module. This count is decremented by DELETE macro instructions.

contains, in the last 6 digits, LLE bytes 5 through 7: starting address of the CDE for the load module.

CDE									
hhhhhhh	ATR1 hh	NCDE hhhhhh	ROC-RB hhhhhhhh	NM ccccccc	USE hh	EPA hhhhhh	ATR2 hh	XL/MJ hhhhhh	

CDE

identifies the next lines as the contents directory addressed by an LLE or RB. One entry is presented per line.

hhhhhhh

is the starting address of the entry given on the line.

ATR1 hh

is the attribute flags.

NCDE hhhhhh is the starting address of the next entry in the contents directory.

ROC-RB hhhhhhhh

contains, in the first 2 digits, zeros.

contains, in the last 6 digits, the starting address of the RB for the load module represented by this entry. NM cccccccc is the name of the entry point to the load module represented by this entry.

USE hh

is the count of the uses (through ATTACH, LINK, and XCTL macro instructions) of the load module, and of the number of LOAD macro instructions executed for the module.

EPA hhhhhh is the entry point address associated with the name in the NM field.

ATR2 hh is the attribute flags.

XL/MJ hhhhhh

is the starting address of the extent list (XL) for a major CDE, or the starting address of the major CDE for a minor CDE. (Minor CDEs are for aliases.)

XL				LN	ADR	LN	ADR	LN	ADR	
	hhhhhh	SZ hhhhhhhh	NO hhhhhhhh	hhhhhhh	hhhhhhh	hhhhhhh	hhhhhhh			

XL

indicates the next lines are entries in the extent list, which is queued to the major contents directory entry. Each extent list entry is given in one or more lines. Only the first line for an entry contains the left 3 columns; additional lines for an entry contain information only in the right 6 columns.

hhhhh

is the starting address of the entry.

SZ hhhhhhhh is the total length, in bytes, of the entry. NO hhhhhhhh is the number of scattered control sections in the load module described by this entry. If this number is 1, the load module was loaded as one block.

LN hhhhhhh gives the length, in bytes, of the control sections in the load module described by this entry. Bit 0 is set to 1 in the last, or only, LN field to signal the end of the list of lengths.

ADR hhhhhhhh gives the starting addresses of the control sections. Each ADR field is paired with the LN field to its left.

TOI	JOB DD	ccccccc STEP hhhhhhhh	ccccccc cccccccc	PROC ccccc hhhhhhh					
				D DOC					
	hhhhhh	hhhhhhh	hhhhhhhh						
	hhhhhh	hhhhhhh	hhhhhhhh	hhhhhhh	հհհհհհհ	հհհհհհհ	հհհհհհհ	հհհհհհհ	hhhhhhhh
	hhhhh	hhhhhhh	hhhhhhh	hhhhhhh	hhhhhhh	hhhhhhh	hhhhhhh	hhhhhhh	hhhhhhhh
	hhhhh			hhhhhhh	hhhhhhh	hhhhhhh	hhhhhhh	hhhhhhh	hhhhhhhh
EB									

DEB

identifies the next lines as the contents of the DEBs and their prefix sections. The first 6 digits in each line give the address of the DEB contents shown on the line, beginning with the second column. The first six digits of the first line contains the prefix section for the first DEB on the queue.

Note: DEBs are not formatted if the dump is for an OLTEP task. If a dump of the DEB chain is desired, use a SYSABEND DD card so that the nucleus will be dumped.

TIOT

identifies the next lines as the contents of the TIOT.

- JOB cccccccc is the name of the job whose task is being dumped.
- STEP cccccccc is the name of the step whose task is being dumped.
- PROC cccccccc is the name for the job step that called the cataloged procedure. This field appears if the job step whose task is being dumped was part of a cataloged procedure.

DD

identifies the line as the contents of the DD entry in the TIOT.

155		****	****** SP	QE ***	********	•			DQE ******	*******	****** FQE	*******
		FLGS	NSPQE	SPID	DQE		BLK	FQE	LN	NDQE	NFOE	LN
1	hhhhh	hh	hhhhh	aaa	hhhhhh		hhhhh	hhhhh	հհհհհ	hhhhh	հհհհհհհ	hhhhhhh
D-PQE	hhhhh	FIRST	hhhhhhh	LAST	hhhhhhhh							
PQE	hhhhh	FFB	hhhhhhhh	LFB	հհհհհհհ	NPO	hhhhhh	h PPO	հհհհհհհ			
		TCB	hhhhhhhh	RSI	hhhhhhhh	RAD	hhhhhh	h FLC	հհհհհհհ			
FBQE	hhhhh	NFB	հհհհհհհ	PFB	hhhhhhh	SZ	հհհհհհ	h				
•		•		•								
•		•		•		•						
•		•		•		•						
PQE	hhhhhh	FFB	hhhhhhhh	LFB	hhhhhhh	NPO	hhhhhhh	h PPO	hhhhhhh			
		TCB	hhhhhhh	RSI	hhhhhhhh	RAD	hhhhhh	h FLG	hhhhhhh			
FBOE	hhhhhh	NFB	hhhhhhh	PFB	hhhhhhh	SZ	hhhhhh	h				

MSS

identifies the next lines as the contents of the main storage supervisor queue. This queue includes subpool queue elements (SPQE), descriptor queue elements (DQE), and free queue elements (FQE). hhhhhh

is the starting address of the first element shown on the line.

SPQE

identifies the 4 columns beneath it as the contents of SPQEs.

FLGS hh is the SPQE flag byte. NSPQE hhhhhh is the starting address of the next SPQE in the queue. SPID ddd is the subpool number. DQE hhhhhh for a subpool owned by the task being dumped: the starting address of the first DQE for the subpool. for a subpool that is shared: the starting address of the SPQE for the task that owns the subpool. DOE identifies the 4 columns beneath it as the contents of DQEs. BLK hhhhhh is the starting address of the allocated 2K block of main storage or set of 2K blocks. FQE hhhhhh is the starting address of the first FQE within the allocated blocks. LN hhhhhh is the length, in bytes, of the allocated blocks. NDQE hhhhhh is the starting address of the next DQE. FOE identifies the 2 columns beneath it as the contents of FQEs. NFQE hhhhhhhh is the starting address of the next FOE. LN hhhhhhhh indicates the number of bytes in the free area. D-PQE hhhhhh is the TCBPQE field (bytes 152 through 155): starting address minus 8 bytes of the dummy PQE shown on the line. FIRST hhhhhhhh is the starting address of the first PQE. LAST hhhhhhhh is the starting address of the last PQE. PQE hhhhhh is the starting address of the PQE shown on the line.

FFB hhhhhhhh is bytes 0 through 3 of the PQE: starting address of the first FBQE. If no FBQEs exist, this field is the starting address of this PQE LFB hhhhhhhh is bytes 4 through 7 of the PQE: starting address of the last FBQE. Tf no FBQEs exist, this field is the starting address of this PQE. NPQ hhhhhhhh is bytes 8 through 11 of the element: starting address of the next PQE or, if this is the last PQE, zeros. PPQ hhhhhhhh is bytes 12 through 15 of the element: starting address of the preceding PQE or, if this is the first PQE, zeros. TCB hhhhhhhh is bytes 16 through 19 of the element: starting address of the TCB for the job step to which the space belongs or, if the space was obtained from unassigned free space, zeros. **RSI** hhhhhhhh is bytes 20 through 23 of the element: size of the region described by this PQE (a multiple of 2048). RAD hhhhhhhh is bytes 24 through 27 of the element: starting address of the region described by this PQE. FLG hhhhhhhh is byte 28 of the element: bit 0 when 0, indicates space described by this PQE is owned; when 1, indicates space is borrowed. bit 1 when 1, indicates region has been rolled out (meaningful only when bit 0 is 0). bit 2 when 1, indicates region has been borrowed. bit 3-7, reserved for future use. Note: POE information is contained in two lines on the dump. When the rollout/rollin feature or Main Storage Hierarchy Support is included in the system, PQE information (with associated FBQEs) appears once in the dump for each region segment of the job step. (Each PQE on the partition queue defines a region segment. A job step's region contains more than one segment only

when the step has rolled out another step

is present.)

or steps, or Main Storage Hierarchy Support

- FBQE hhhhhh is the starting address of the FBQE shown on the line.
- NFB hhhhhhhh is bytes 0 through 3 of the element: starting address of the next FBOE. In the highest or only FBQE, this field contains the address of the PQE.
- PFB hhhhhhhh is bytes 4 through 7 of the element: starting address of the previous FBQE. In the lowest or only FBQE, the field contains the address of the PQE.

SZ hhhhhhhh is bytes 8 through 11 of the element: size, in bytes, of the free area.

QCB	TRACE				
MAJ	hhhhh	NMAJ hhhhhhhh	PMAJ hhhhhhhh	FMIN hhhhhhhh	NM ccccccc
MIN	hhhhh	FQEL hhhhhhhh	PMIN hhhhhhhh	NMIN hhhhhhhh	NM xx xxxxxxx
		NQEL hhhhhhhh	PQEL hhhhhhhh	TCB hhhhhhhh	SVRB hhhhhhh

QCB TRACE

- identifies the next lines as a trace of the queue control blocks (QCB) associated with the job step. Lines beginning with MAJ show major QCBs, lines beginning with MIN show minor QCBs, and lines beginning with NQEL show queue elements (QEL).
- MAJ hhhhhh is the starting address of the major QCB whose contents are given on the line.
- NMAJ hhhhhhhh is the starting address of the next major QCB for the job step.
- PMAJ hhhhhhhh is the starting address of the previous major QCB for the job step.
- FMIN hhhhhhh is the starting address of the first minor QCB associated with the major QCB given on the line.
- NM CCCCCCC is the name of the serially reusable resource represented by the major QCB.
- MIN hhhhhh is the starting address of the minor QCB whose contents are given on the line.
- FQEL hhhhhhhh is the starting address of the first queue element (QEL), which represents a request to gain access to a serially reusable resource or set of resources.

PMIN hhhhhhhh is the starting address of the

- previous minor QCB.
- NMIN hhhhhhh is the starting address of the next minor QCB.
- NM XX XXXXXXX

indicates, in the first 2 digits, the scope of the name or address of the minor QCB being dumped. If the scope is hexadecimal FF, the name is known to the entire operating system. If the scope is hexadecimal 00 or 10 through FO, the name is known only to the job step; in this case, the scope is the protection key of the TCB enqueuing the minor QCB.

Also contains, in the last 8 digits, the name or the starting address of the minor QCB.

NOEL hhhhhhhh indicates, by hexadecimal 10 in the first 2 digits, that the queue element on the line represents a request for step-must-complete; by 00, ordinary request; and by 20, a set-must-complete request.

> Also contains, in the last 6 digits, the starting address of the next queue element in the queue, or for the last queue element in the queue, zeros.

POEL hhhhhhhh

indicates, by hexadecimal 80 in the first 2 digits, that the queue element represents a shared request or, by hexadecimal 00, that the element represents an exclusive request. (If

the shared DASD option was selected, hexadecimal 40 in the first 2 digits indicates an exclusive RESERVE request and 00 indicates a shared RESERVE request.)

TCB hhhhhhhh

is the starting address of the TCB under which the ENQ macro instruction was issued. SVRB hhhhhhhh

is the starting address, of the SVRB under which the routine for the ENQ macro instruction is executed, or, after the requesting task receives control of the resource, the UCB address of a device being reserved through a RESERVE macro instruction (the latter value occurs only when the shared DASD option was selected).

cccc	CCCC WAS	ENT	ERED VIA LI	INK (C	(ALL) dadad	AT	EP CCCCC	•				
52	հհհհհ	WD1 R1 R7	հհհհհհհ հհհհհհհ հհհհհհհ	R2	hhhhhhhh hhhhhhhh hhhhhhhh		հհհհհհհ հհհհհհհ հհհհհհհ	R4	հհհհհհհհ հհհհհհհ հհհհհհհ	R5	hhhhhhhh hhhhhhhh hhhhhhh	 hhhhhhhi hhhhhhhi hhhhhhhi
INCO	RRECT BA	CK C	HAIN									
_	RRUPT AT											

SAVE AREA TRACE

identifies the next lines as a trace of the save areas for the program. Each save area is presented in 3 or 4 lines. The first line gives information about the linkage that last used the save area. This line will not appear when the RB for the linkage cannot be found. The second line gives the contents of words 0 through 5 of the save area. The third and fourth lines give the contents of words 6 through 18 of the save area; these words are the contents of registers 0 through 12. Save areas are presented in the following order:

- 1. The save area pointed to in the TCBFSA field of the TCB. This save area is the first one for the problem program; it was set up by the control program when the job step was initiated.
- 2. If the third word of the first save area was filled by the problem program, then the second save area shown is that of the next lower level module of the task. However, if the third word of the first area points to a location whose second word does not point back to the first area, the message INCORRECT BACK CHAIN appears, followed by possible contents of the second save area.

3. The third, fourth, etc. save areas are then shown, provided the third word in each higher save area was filled and the second word of each lower save area points back to the next higher save area. This process is continued until the end of the chain is reached (the third word in a save area contains zeros) or INCORRECT BACK CHAIN appears.

Following the forward trace, the message INTERRUPT AT hhhhhh appears, followed by the message PROCEEDING BACK VIA REG 13. Then, the save area in the lowest level module is presented, followed by the save area in the next higher level. The lowest save area is assumed to be the 76 bytes beginning with the byte addressed by register 13. These two save areas appear only if register 13 points to a full word boundary and does not contain zeros.

- cccccccc WAS ENTERED is the name of the module that stored register contents in the save area. This name is obtained from the RB.
- VIA LINK ddddd or VIA CALL ddddd indicates the macro instruction (LINK or CALL) used to give control to the next lower level module, and is the ID

operand, if it was specified, of the LINK or CALL macro instruction.

AT EP CCCCC...

is the entry point identifier, which appears only if it was specified in the SAVE macro instruction that filled the save area.

- SA hhhhhh is the starting address of the save area.
- WD1 hhhhhhhh is the first word of the save area (optional).
- HSA hhhhhhhh

is the second word of the save area: starting address of the save area in the next higher level module. In the first save area in a job step, this word contains zeros. In all other save areas, this word must be filled.

LSA hhhhhhhh

is the third word of the save area (register 13): starting address of the save area in the next lower level (called) module. If the module containing this save area did not fill the word, it contains zeros.

RET hhhhhhhh

is the fourth word of the save area (register 14): return address (optional); if the called module did not fill the word, it contains zeros.

EPA hhhhhhhh

is the fifth word of the save area

(register 15): entry point to the called module. Use of this word is optional; if the called module did not fill the word, it contains zeros.

- RO hhhhhhh R1 hhhhhhh ... R12 hhhhhhh are words 6 through 18 of the save area (registers 0 through 12): contents of registers 0 through 12 for the module containing the save area immediately after the linkage. Use of these words is optional; if the called module did not fill these words, they contain zeros.
- INCORRECT BACK CHAIN indicates that the following lines may not be a save area because the second word in this area does not point back to the previous save area in the trace.
- INTERRUPT AT hhhhhh is the address of the next instruction to be executed in the problem program. It is obtained from the resume PSW word of the last PRB or LPRB in the active RB queue.

PROCEEDING BACK VIA REG 13

indicates that the next 2 save areas are (1) the save area in the lowest level module, followed by (2) the save area in the next higher level module. The lowest save area is the save area pointed to by register 13. These 2 save areas appear only if register 13 points to a fullword boundary and does not contain zero.

CPUX PSA				
hhhhhh hhhhhhh hhhhhh hhhhhhh	hhhhhhh hhhhhhhh hhhhhhh hhhhhhh hhhhhhh	հհհհհհհհ հհհհհհհհ հհհհհ հհհհհհհհ հհհհհհհ		**************************************
NUCLEUS				
	հհհհհհհհ հհհհհհհհ հհհհհհհ	hhhhhhh hhhhhhhh hhhhh		*22222222222222222222222222222222222222
hhhhhh hhhhhhh	հհհհհհհհ հհհհհհհհ հհհհհհհ	հհհհհհհհ հհհհհհհհհհհհ	hh hhhhhhhh *ccccc	ccccccccccccccccccccccccccc
NUCLEUS CONT.				
	հհհհհհհհ հհհհհհհ հհհհհհհ հհհհհհհհ հհհհհհ	hhhhhhh hhhhhhhh hhhhh hhhhhhh hhhhhhh hhhhh		*22222222222222222222222222222222222 cccccc
REGS AT ENTRY TO A	ABEND (SNAP)			
FLTR 0-6	հհհհհհհհհհհհհհհ հհհհհ	հիհիհիհի լ	ոհհհհհհհհհհհհ	հհհհհհհհհհհհհ
REGS 0-7 REGS 8-15	հհհհհհհհ հհհհհհհհ հհհհհ հհհհհհհ հհհհհհհ հհհհհ		ոհհհհհհհ հհհհհհհ հհհհհհհ հհհհհհհ	հհհհհհհհ հհհհհհհհ հհհհհհհ հհհհհհհ
LOAD MODULE cccccc	ccc			
hhhhhh hhhhhhh	hhhhhhh hhhhhhhh hhhhhhh hhhhhhh hhhhhhh	հհհհհհհհ հհհհհհհհ հհհհհ հհհհհհհ հհհհհհհհ		cccccccccccccccccccccccccce*
hhhhhh hhhhhhh	hhhhhhh hhhhhhhh hhhhhhhh hhhhhhhh hhhhhh	hhhhhhhh hhhhhhhh hhhhh hhhhhhhh hhhhhhh		ccccccccccccccccccccccccccc*
LINE hhhhhh	SAME AS ABOVE			
CSECT dd OF ccccco	200			
	հհհհհհհհ հհհհհհհհ հհհհհհհ հհհհհհհհ հհհհհհ	հհհհհհհհ հհհհհհհհ հհհհհ հհհհհհհ հհհհհհհ հհհհհ		**************************************

The contents of main storage are given under 6 headings: CPUx PSA, NUCLEUS, NUCLEUS CONT., LOAD MODULE cccccccc, CSECT dd OF cccccccc, and in the trace table, SP ddd BLK hh. Under these headings, the lines have the following format:

- First entry: the address of the initial bytes of the main storage presented on the line.
- Next 8 entries: 8 full words (32 bytes) of main storage in hexadecimal.
- Last entry (surrounded by asterisks): the same 8 full words of main storage in EBCDIC. Only A through Z, 0 through 9, and blanks are printed; a period is printed for anything else.

The following lines may also appear:

- LINES hhhhhh-hhhhhh SAME AS ABOVE are the starting addresses of the first and last lines for a group of lines that are identical to the line immediately preceding.
- LINE hhhhhh SAME AS ABOVE is the starting address of a line that is identical to the line immediately preceding.

CPUx PSA (Model 65 Multiprocessing dumps only)

identifies the next lines as the contents of the prefixed storage area (PSA) -- 0 through 4095 (FFF). If the system is operating in partitioned mode (1 CPU), x is the CPU identification. If the system is operating in a 2 CPU multisystem mode, both PSAs are printed, the first under the heading CPUA PSA and the second under CPUB PSA.

NUCLEUS

identifies the next lines as the contents of the nucleus of the control program.

NUCLEUS CONT.

identifies the next lines as the contents of the part of the nucleus that lies above the trace table.

REGS AT ENTRY TO ABEND OF REGS AT ENTRY TO SNAP

identifies the next 3 lines as the contents of the floating point and general registers when the abnormal termination routine received control in response to an ABEND macro instruction or when the SNAP routine received control in response to a SNAP macro instruction. These are not the registers for the problem program when the error occurred.

- FLTR 0-6 indicates the contents of floating point registers 0, 2, 4, and 6.
- REGS 0-7 indicates the contents of general registers 0 through 7.
- REGS 8-15 indicates the contents of general registers 8 through 15.
- LOAD MODULE cccccccc identifies the next lines as the contents of the main storage area occupied by the load module cccccccc addressed by an LLE or RB. All the modules for the job step are dumped under this type of heading. Partial dumps do not contain this information.
- CSECT hhhh OF ccccccc identifies the next lines as the contents of the main storage area occupied by the control section (CSECT) indicated by hhhh. This control section belongs to the scatter-loaded load module cccccccc.

TRACI	E TABLE													
DSP	NEW PSW	hhhhhhh	hhhhhhh	R15/R0	hhhhhhh	hhhhhhh	Rl	hhhhhhh	SW	hhhhhhh	тсв	հհհհհհհ	TME	hhhhhhh
1/0	OLD PSW	hhhhhhh	hhhhhhhh	R15/R0	hhhhhhh	hhhhhhh	R1	հհհհհհհ	PES	հհհհհհհ	TCB	hhhhhhh	TME	hhhhhhh
SIO	CC/DEV/CAW	hhhhhhh	hhhhhhh	CSW	hhhhhhh	hhhhhhh	RES	hhhhhhh	RES	hhhhhhh	TCB	hhhhhhh	TME	hhhhhhh
svc	OLD PSW	hhhhhhh	hhhhhhh	R15/R0	hhhhhhhh	hhhhhhh	R1	hhhhhhh	RES	hhhhhhh	TCB	hhhhhhh	TME	hhhhhhh
PGM	OLD PSW	hhhhhhh	hhhhhhh	R15/R0	hhhhhhh	hhhhhhh	R 1	hhhhhhh	PES	hhhhhhh	TCB	hhhhhhh	TME	hhhhhhh
EXT	OLD PSW	hhhhhhh	hhhhhhh	R15/R0	hhhhhhh	hhhhhhh	R1	hhhhhhh	RES	հհհհհհհ	TCB	hhhhhhhh	TME	hhhhhhh

TRACE TABLE (SNAP dumps only)
 identifies the next lines as the
 contents of the trace table. Each
 trace table entry is presented on one
 line; the name at the beginning of
 each line identifies the type of entry
 on the line:

• DSP Dispatcher entry

- I/O Input/output interruption entry
- SIO Start input-output (SIO) entry
- SVC Supervisor call (SVC) interruption entry
- PGM Program interruption entry
- EXT External interruption entry
- OLD PSW hhhhhhhh hhhhhhhh is the PSW stored when the interruption represented by the entry occurred.
- NEW PSW hhhhhhhh hhhhhhhh is the new PSW stored in the entry.
- CC/DEV/CAW hhhhhhhh hhhhhhhh contains, in the first 2 digits: completion code.

contains, in the next 6 digits: device type. contains, in the last 8 digits: address of the channel address word (CAW) stored in the entry.

R15/RO hhhhhhh hhhhhhh contains, in the first 8 digits: contents of register 15 stored in the entry.

> contains, in the last 8 digits: contents of register 0 stored in the entry.

- CSW hhhhhhh hhhhhhhh is the channel status word (CSW) stored in the entry.
- R1 hhhhhhhh is the contents of register 1 stored in the entry.
- RES hhhhhhhh is reserved for future use; all digits are zeros.
- SW hhhhhhhh is reserved for future use; all digits are zeros.
- TCB hhhhhhhh is the starting address of the TCB associated with the entry.
- TME hhhhhhhh is a representation of the timer element associated with the entry.

TRT														
X DSP	NEW PSW	hhhhhhh	hhhhhhh	R15/R0	hhhhhhh	hhhhhhh	R1	hhhhhhh	NUA	hhhhhhh	NUB	hhhhhhh		hhhhhh
X I/O	OLD PSW	hhhhhhh	hhhhhhh	CSW	hhhhhhh	hhhhhhh	R1	hhhhhhh	OLA	hhhhhhh		hhhhhhhh		hhhhhh
X S10	CC/DEV/CAW	hhhhhhh	hhhhhhh	CSW	hhhhhhh	hhhhhhh	TCB	hhhhhhh	OLA	hhhhhhh	OLB	hhhhhhhh		hhhhhh
X SVC	OLD PSW	hhhhhhh	hhhhhhh	R15/R0	hhhhhhh	hhhhhhh	R1	hhhhhhh	OLA	hhhhhhh	OLB	hhhhhhh	TME	hhhhhh
X PGM	OLD PSW	hhhhhhh	hhhhhhh	R15/R0	hhhhhhh	hhhhhhh	R1	hhhhhhh		hhhhhhhh		hhhhhhh		hhhhh
X EXT	OLD PSW	hhhhhhhh	hhhhhhh	R15/R0	hhhhhhh	հհհհհհհ	R1	hhhhhhh	MSK	hhhhhhhh	TOE	hhhhhhhh		hhhhh
X SSM	OLD PSW	hhhhhhh	hhhhhhh	R15/R0	hhhhhhh	hhhhhhh	R1	hhhhhhh	AFF	yyhhhhhh	OLB	hhhhhhhh	TME	hhhhhh

TRT (MVT with Model 65 multiprocessing dumps only)

- identifies the next lines as the contents of the trace table. Each trace table entry is presented on one line; the letter and name at the beginning of each line identify the CPU and the type of entry, respectively:
- DSP Dispatcher entry.
- I/O Input/output interruption entry.
- SIO Start input/output entry.
- SVC Supervisor call interruption entry.
- PGM Program interruption entry.
- EXT External interruption entry.
- SSM Set system mask entry.
- OLD PSW hhhhhhhh hhhhhhh is the PSW stored when the interruption represented by the entry occurred.
- NEW PSW hhhhhhhh hhhhhhhh is the new PSW stored in the entry.
- CC/DEV/CAW hhhhhhhh hhhhhhh contains, in the first 2 digits: completion code; in the next 6 digits: device type; in the last 8 digits: address of the channel address word stored in the entry.
- R15/R0 hhhhhhh hhhhhhh contains, in the first 8 digits: contents of register 15; in the last 8 digits: contents of register 0, both as stored in the entry.

- CSW hhhhhhhh hhhhhhhh is the channel.status word stored in the entry.
- R1 hhhhhhh is the contents of register 1 as stored in the entry.
- TCB hhhhhhhh is the starting address of the TCB associated with the entry.
- NUA hhhhhhhh is the starting address of the new TCB for CPU A, as stored in the entry.
- OLA hhhhhhhh is the starting address of the old TCB for CPU A, as stored in the entry.
- MSK hhhhhhhh is the STMASK of the other CPU as stored in the entry.
- NUB hhhhhhhh is the starting address of the new TCB for CPU B, as stored in the entry.
- OLB hhhhhhhh is the starting address of the old TCB for CPU B, as stored in the entry.
- TQE hhhhhhhh is the first word of the timer queue element stored in the entry, provided a timer interrupt occurred.
- TME hhhhhhhh is a representation of the timer element associated with the entry.
- AFF yyhhhhh contains, in the first 2 digits: the ID of the locking CPU at the time of the interrupt; in the last 6 digits: starting address of the old TCB for CPU A, as stored in the entry.

SP ddd

 END OF DUMP

SP ddd

identifies the next lines as the contents of a block of main storage obtained through a GETMAIN macro instruction, and indicates the subpool number (ddd). The part of subpool 252 that is the supervisor work area is presented first, followed by the entire contents of any problem program subpools (0 through 127) in existence during the dumping.

END OF DUMP

indicates that the dump or snapshot is completed. If this line does not appear, the ABDUMP routine was abnormally terminated before the dump was completed, possibly because enough space was not allocated for the dump data set.

Guide to Using an ABEND/SNAP Dump (MVT)

<u>Cause of Abnormal Termination</u>: Evaluate the user (USER=decimal code) or system (SYSTEM=hex code) completion code using Appendix B or the publication <u>Messages and</u> <u>Codes</u>.

<u>Dumped Task</u>: Check the ID field for an indication of which task is being dumped in relation to the task that was abnormally terminated:

- 001 indicates a partial dump of a subtask
- 002 indicates a partial dump of the invoking task

If the ID field is absent, the dump contains a full dump of the task that was abnormally terminated.

Active RB Queue: The first RB shown on the dump represents the oldest RB on the queue. The RB representing the load module that had control when the dump was taken is third from the bottom. The last RB represents the ABDUMP routine and the second from last, the ABEND routine. The load module name and entry point (for a PRB) are given in a contents directory entry, the address of which is shown in the last 3 bytes of the FL/CDE field. <u>Program Check PSW</u>: The program check old PSW is the fifth entry in the first line of each RB printout. It is identified by the subheading APSW. For debugging purposes, the APSW of the third RB from the bottom of the dump is most useful. It provides the length of the last instruction executed in the program (bits 32,33), and the address of the next instruction to be executed (bytes 5-8).

Load List: Does the resume PSW indicate an instruction address outside the limits of the load module that had control at the time of abnormal termination? If so, look at the LLEs on the load list. Each LLE contains the CDE address in the dump field labeled RSP-CDE.

<u>CDEs</u>: The entries in the contents directory for the region are listed under the dump heading CDE. The printouts for each CDE include the load module and its entry point. If you have a complete dump, each load module represented in a CDE is printed in its entirety following the NUCLEUS section of the dump.

Trace Table (SNAP dumps only): Entries on an MVT SNAP dump, if valid, represent occurrences of SIO, external, SVC, program, I/O, and dispatcher interruptions. <u>SIO</u> entries can be used to locate the CCW (through the CAW), which reflects the operation initiated by an SIO instruction. If the SIO operation was not successful, the CSW STATUS portion of the entry will show you why it failed. <u>EXT</u> and <u>PGM</u> entries are useful for locating the instruction where the interruption occurred (bytes 5-8 of the PSW).

SVC trace table entries provide the SVC old PSW and the contents of registers 0, 1, and 15. The PSW offers you the hexadecimal SVC number (bits 20-31), the CPU mode (bit 15), and the address of the SVC instruction (bytes 5-8). The contents of registers 0 and 1 are especially useful in that many system macro instructions pass key information in these registers. (See Appendix A.)

<u>I/O entries</u> reflect the I/O old PSW and the CSW that was stored when the interruption occurred. From the PSW, you can learn the

address of the device that caused the interruption (bytes 2 and 3), the CPU state at the time of interruption (bit 15), and the instruction address where the interruption occurred (bytes 5-8). The CSW provides you with the unit status (byte 4), the channel status (byte 5), and the address of the previous CCW plus 8 (bytes 0-3).

You can use the <u>DSP</u> entry to delimit the entries in the trace table. To find all entries for the terminated task, scan word 7 of each trace table entry for the TCB address in a DSP entry. The lines between this and the next DSP entry represent interruptions that occurred in the task.

<u>Region Contents</u>: Free areas for the region occupied by the dumped task are identified under headings PQE and FBQE. The field labeled SZ gives the number of bytes in the free area represented by the FBQE.

<u>Subpool Contents</u>: Free and requested areas of the subpools used by the dumped task are described under the dump heading MSS. Subpool numbers are given under the SPID column in the list of SPQEs. If a GETMAIN macro instruction was issued without a subpool specification, space is assigned from subpool 0. Thus, two SPQEs may exist for subpool 0. The sizes of the requested areas and free areas are given under the LN column in the lists of DQEs and FQEs, respectively.

Load Module Contents: The contents of each load module used by the job step are given under the heading XL. Each entry includes the sizes (LN) and starting addresses (ADR) of the control sections in the load module.

Indicative Dump

An indicative dump is issued when a task is abnormally terminated by an ABEND macro instruction, and a dump is requested, but a dump data set is not available, due either to omission or incorrect specification of a SYSABEND or SYSUDUMP DD statement. An indicative dump is issued automatically on the system output (SYSOUT) device.

Systems with MVT do not issue indicative dumps.

Contents of an Indicative Dump

This explanation of indicative dumps utilizes capital letters for the headings found in all dumps, and lowercase letters for information that varies with each dump. The lowercase letter used indicates the mode of the information, and the number of letters indicates its length:

- h represents 1/2 byte of hexadecimal information
- d represents 1 byte of decimal information
- c represents a 1-byte character

Figure 19 shows the contents of an indicative dump. You may prefer to follow the explanation on your own indicative dump.

CONTROL BYTE=hh

describes the contents of the indicative dump.

First digit:

<u>Bit</u>	<u>Setting</u>	Meaning
0	0	Instruction image not
		present
	1	Instruction image present
1	0	Floating-point registers
		not present
	1	Floating-point registers present
2	0	One general register set
		present
	1	Two general register sets
		present
3	0	All active RBs present
•	1	All active RBs not present
	-	
Last	t digit:	

st algit:

1

2

Digit in			
Hexadecimal	Meaning		
0	All loaded	RBs	present

- 8 All loaded RBs not present
- TCB FLAGS=hh is the first byte of TCBFLGS field (byte 29 in the TCB for the program being dumped): task end flag byte:
 - Bit Setting Meaning
 - Abnormal termination in 1 0 process
 - 1 Normal termination in process
 - Abnormal termination was 1 initiated by the resident ABTERM routine

19							
×.*							
14							
	CONTROL BYTE=hh	TCB FLAGS=hh	NO. ACTI	VE RB=dd	NC. LOAD F	B=dd	
	COMPLETION CODE	- CVCTEMabbb	11SFR#dddd				
		- SISIER-MAN	00211-0000				
	cccccc						
	REGISTER SET 1						
	hbhhhhhh hhhhhh	hh hbhhhhhh	bbbbbbbbb	hhhhhhhh	hhhhhhhh	hhhhhhhh	hhhhhhhh
			hhhhhhh	hhhhhhhh	hhhhhhh	hhhhhhh	hhhhhhhh
	hhhhhhhh hhhhhh	hh hhhhhhhh	mmmm	[]]]]]]]]]]]]]]]]]]]]]]]]]]]]]]]]]]]]]]			
	REGISTER SET 2						
	hhhhhhhh hhhhhh	իի իիիիիիի	nhhhhhhh	hhhhhhhh	hhhhhhhh	hhhhhhh	hhhhhhhh
			hhhhhhhh	hhhhhhh	hhbbhbh	hbhbhbhb	hhhhhhh
	hhhhhhhh hhhhhh						
	INSTRUCTION IMAG	iE=hhhhhhhhhhhh	hhhhhhhhhhh	ממח			
1.	հերհերհերհերհերհերհերհերհերհերհերհերհերհ			hhhhhhhh	hhhhhhhhh	hhhhhhhh	hhhhhhhh
				POINT=hhhl	հհհ		
• •	PROGRAM ID=ccccc					•hhhhhh	
	RESUME PSW SM=h	h K=h AMWP=				•000000	
	PROGRAM ID=ccccc	CCC RB TYPE=	hh ENTRY	POINT=hhhl	hhh		
	TROOMIN ID COOCC						
5.							
· · ·							

Figure 19. Contents of an Indicative Dump

- 3 1 ABTERM routine entered because of program interruption
- 4 1 Reserved for future use
- 5 1 Data set closing initiated by the ABTERM routine
- 6 1 The ABTERM routine overlaid some or all of the problem program
- 7 1 The system prohibited queuing of asynchronous exit routines for this task
- NO. ACTIVE RB=dd is the number of active RBs presented in the dump.
- NO. LOAD RB=dd is the number of RBs in the load list presented in the dump.
- COMPLETION CODE SYSTEM=hhh USER=dddd is the completion code supplied by the control program (SYSTEM=hhh) or the problem program (USER=dddd). Both SYSTEM=hhh and USER=dddd are printed; however, one of them is always zero.
- cccccc...

1

- explains the completion code or, if a program interruption occurred:
- PROGRAM INTERRUPTION ccccc.. AT LOCATION hhhhh where ccccc is the program interruption cause: OPERATION, PRIVILEGED OPERATION, EXECUTE, PROTECTION, ADDRESSING, SPECIFICATION, DATA, FIXED-POINT OVERFLOW, FIXED-POINT DIVIDE, DECIMAL OVERFLOW, DECIMAL DIVIDE, EXPONENT OVERFLOW, DECIMAL DIVIDE, EXPONENT OVERFLOW, EXPONENT UNDERFLOW, SIGNIFICANCE, or FLOATING-POINT DIVIDE; and hhhhhh is the address of the instruction being executed when the interruption occurred.
- REGISTER SET 1
 - indicates that the next 2 lines give the contents of general registers 0 through 7 and 8 through 15 for a program being executed under control of an RB when it:
 - Passed control to a type I SVC routine through an SVC instruction and the routine terminated abnormally.

- Lost control to the input/output interruption handler, which subsequently terminated abnormally.
- Was abnormally terminated by the control program because of a program interruption.
- Issued an ABEND macro instruction to request an abnormal termination.

If REGISTER SET 2 also appears in the dump, the lines under REGISTER SET 1 give the general register contents for a type II, III, or IV SVC routine operating under an SVRB.

REGISTER SET 2

indicates that the next 2 lines give the contents of general registers 0 through 7 and 8 through 15 for a program being executed under control of an RB other than an SVRB when the program last passed control to a type II, III, or IV SVC routine.

PROGRAM ID=ccccccc

is the XRBNM field (bytes 0 through 7): in PRB, LRBs, and LPRBs, the program name; in IRBs, the first character contains flags for the timer or, if the timer is not being used, contains no meaningful information; in SVRBs for a type II SVC routine, contains no meaningful information; in SVRBs for a type III or IV SVC routine, the first 4 bytes contain the relative track address (TTR) of the load module in the SVC library and the last 4 bytes contain the SVC number in signed, unpacked decimal; in SIRBs, the name of the error routine currently occupying the 400-byte input/output supervisor transient area.

RB TYPE=hh

- indicates the type of active RB
- hh Type of RB 00 PRB that does not contain entry points identified by IDENTIFY macro instructions
- 10 PRB that contains one or more entry points identified by IDENTIFY macro instructions
- 20 LPRB that does not contain entry points identified by IDENTIFY macro instructions
- 30 LPRB that contains one or more entry points identified by IDENTIFY macro instructions
- 40 IRB
- 80 SIRB
- CO SVRB for a type II SVC routine
- D0 SVRB for a type III or IV SVC routine
- E0 LPRB for an entry point identified by an IDENTIFY macro instruction
- FO LRB
- ENTRY POINT=hhhhhh is the XRBEP field (bytes 13 through 15): address of entry point in the program.

RESUME PSW

XRBPSW field (bytes 16 through 23): is the contents of the resume PSW.

SM=hh

is bits 0 through 7 of PSW: system mask.

K=h is bits 8 through 11 of PSW: protection key.

AMWP=h

is bits 12 through 15 of PSW: indicators.

IC=hhhh

is bits 16 through 31 of PSW: interruption code.

IL.CC=h

is bits 32 through 35 of PSW: instruction length code (bits 32 and 33) and condition code (bits 34 and 35). PM=h
 is bits 36 through 39 of PSW: program
 mask.
IA=hhhhhh
 is bits 40 through 63 of PSW:
 instruction address.
PROGRAM ID=cccccccc

- is the XRBNM field (bytes 0 through 7): program name.
- RB TYPE=hh indicates the type of RB:
 - hh Type of RB
 - 20 LPRB that does not contain entry points identified by IDENTIFY macro instructions.
 - 30 LPRB that contains one or more entry points identified by IDENTIFY macro instructions.
 - E0 LPRB for an entry point identified by an IDENTIFY macro instruction.
 F0 LRB.

ENTRY POINT=hhhhhh

is the XRBEP field (bytes 13 through 15): address of entry point in the program.

Guide to Using an Indicative Dump

<u>Completion Code</u>: Evaluate the user (USER=decimal code) or system (SYSTEM=hex code) completion code using either Appendix C of this publication or the publication <u>Messages and Codes</u>. The line under the completion code gives a capsule explanation of the code or the type of program interruption that occurred.

<u>Instruction Address</u>: If a program interruption occurred, get the address of the erroneous instruction in the last 3 bytes of the field labeled INSTRUCTION IMAGE.

Active RB Queue: RBs are shown in the first group of two-line printouts labeled PROGRAM ID and RESUME PSW, with the most recent RB shown first. There are two lines for as many RBs indicated by NO. ACTIVE RB=dd.

Register Contents: General register contents at the time a program last had control are given under the heading REGISTER SET 2 or, if this heading is not present, under REGISTER SET 1. Register contents, particularly those of register 14, may aid you in locating the last instruction executed in your program.

Storage Dumps

Storage dumps record the contents of main storage from location 00 to the end of printable storage.

Storage dumps are produced by the damage assessment routine (DAR) or other system recovery routines, the Console Dump facility, or the stand-alone service aid program IMDSADMP.

DAMAGE ASSESSMENT ROUTINE (DAR)

The damage assessement routine produces a storage dump when a system task fails and is designed to provide increased system availability in the event of system failure. The storage dump is written to the SYS1.DUMP data set.

If a system routine fails, DAR attempts to reinitialize the failing task, thereby permitting the system to continue operation without interruption. DAR permits the system to continue processing in a degraded condition if it encounters a system failure that does not permit total reinstatement of the affected task or region. The operator will be informed, via a WTO, that the system is in an unpredictable state; he then must decide whether or not already-scheduled jobs should be allowed to attempt completion.

Note: If TSO is installed in the system and a failure occurs in the TSO subsystem or in the operating system the TSO SWAP data set must be recorded for use in diagnosis if needed. The system recovery routines do not do this. The IMDPRDMP service aid can be used as a high-performance dumping program for this purpose by directing its output to tape. Refer to the Service Aids publication for details of this usage of the IMDPRDMP program.

CONSOLE DUMP

The Console Dump function is designed to meet the requirements for a dynamic main storage dumping tool in the operating system. The operator initiates the Console Dump from the primary console via a DUMP command. Execution of the function allows a dump to be taken to the SYS1.DUMP data set of all or selective portions of main storage. The dump operation is performed during system operation and requires no The storage dump may then be IPL. formatted and printed by the IMDPRDMP Service Aid program. Refer to the Operator's Guide publication for details of the DUMP command.

IMDSADMP SERVICE AID

In situations where the system is not operative, an IMDSADMP program is loaded into storage through use of the IPL facilities. The storage dump taken may be written in a high-speed version to tape or disk, and in a low speed version to tape or printer. The high-speed IMDSADMP dump must be processed by the IMDPRDMP program. The low-speed tape output may be processed by a program such as the IEBGENER utility program. The format of the low-speed IMDSADMP output is similar to the general format listing produced by the IMDPRDMP program and therefore is not illustrated in this publication. A sample IMDSADMP listing and a discussion of the program are contained in the Service Aids publication.

SYSTEM FAILURE

If a system failure occurs, the damage assessment routine immediately attempts to write a storage dump to the SYS1.DUMP data set. A system failure may be caused by a failure in any of the following system tasks:

MFT:

Communications Task Master Scheduler Task Log Task

MVT:

System Error Task Rollout/Rollin Task Communications Task Master Scheduler Task Transient Area Fetch Task

A system failure is also caused by an ABEND recursion in other than OPEN, CLOSE, ABDUMP, or STAE; by a failure of a task in 'must complete' status; or, in MFT only, by a failure in the scheduler if no SYSABEND or SYSUDUMP DD card is provided.

THE SYS1.DUMP DATA SET

The SYS1.DUMP data set may reside on tape or on a direct access device.

<u>Tape</u>

If you wish to have the SYS1.DUMP data set reside on tape, you may specify the tape drive during IPL. If the drive has not been made ready prior to IPL, a MOUNT message is issued to the console, specifying the selected device. The device should be mounted with an unlabeled tape. After writing a storage image dump, the damage assessment routine writes a tape mark and will position the tape to the next file. The tape drive will remain in a ready state to receive another storage image dump.

Direct Access

If you wish to have the SYS1.DUMP data set placed on a direct access device, you may preallocate the data set at system generation or prior to any IPL of the system. The following restrictions apply:

- The data set name must be SYS1.DUMP.
- The data set must be cataloged on the IPL volume.

- The data set may be preallocated on any volume that will be online during system operation.
- The data set must be sequential.
- Sufficient space must be allocated to receive a storage image dump for all of main storage.

When a direct access device is used for the SYS1.DUMP data set, the data set can hold only one storage dump. If additional failures occur, and if the SYS1.DUMP data set is occupied, DAR does not attempt to write another storage image dump.

Use the IMDPRDMP service aid to format and list the SYS1.DUMP data set.

IMDPRDMP Output

Main storage information processed by the IMDPRDMP program is presented in six different output formats. The output format used is determined by the function of the particular area of the dumped system's main storage that is being printed. Two of these formats, the queue control block trace and the link pack area map, are invoked by specific format statements. A third format is used to print the major system control blocks. TWO formats are used for TSO; one for system control blocks and the other for user control blocks. Any areas of the dumped system's main storage that do not fall into any of the aforementioned functional categories are processed in the general format.

<u>Dump List Headings</u>: Each page of output listing contains a heading. This heading has the optional user specified title, the name of the module that invoked the dump, the date and the time the dump was taken except when processing Generalized Trace Facility output when the heading will be "EXTERNAL TRACE - DD ddname." <u>Note</u>: If the dump was produced by IMDSADMP on a system with the time-of-day (TOD) clock, IMDPRDMP can not determine the time at which the dump was taken; the time is replaced by "TOD CLK."

<u>Dump Header</u>: If the dump was produced by SVC DUMP, IMDPRDMP will print the title taken from the dump header record. A maximum of 100 characters are printed on the second line of the first page of the output listing.

<u>Output Comments</u>: While formatting the dump, the IMDPRDMP program occasionally is unable to locate, format and print a control block. On those occasions IMDPRDMP prints a comment explaining why the control block could not be formatted and printed. These comments are printed within the body of the formatted dump and are part of the IMDPRDMP output. A complete list of these output comments along with further explanations is contained at the end of this chapter.

<u>Summary Information</u>: In addition to formats, the following summary information is printed at the end of each execution of IMDPRDMP:

- The number of entries to the read routine;
- The number of times that the required address was not found in a buffer;
- The number of blocks read from the dump data set;

- The number of permanent I/O errors encountered during the execution;
- The average number of buffers used for each operation performed during this execution;
- The number of blocks read from the TSO swap data sets;
- The ratio of the number of times the read routine was called to the number of times the requested address was not in a buffer.
- When processing Generalized Trace Facility output, the number of trace records processed.

QUEUE CONTROL BLOCK TRACES

In a multiprogramming environment, requests for system resources are enqueued. This process is accomplished through the use of queue control blocks (QCBs).

Certain system failures, such as task contention deadlocks, become evident to the user upon examination of a queue control block trace. When requested through the use of the QCBTRACE statement, the QCB trace appears on a separate page of the IMDPRDMP program dump listing. The trace, a sample of which appears in Figure 20, contains a listing of all queue control blocks that were present in the dumped system, and is available to users who are processing main storage information gathered from an MVT or MFT system.

(For more information on system resource queuing, see <u>IBM System/360 Operating</u> <u>System: MVT Supervisor</u>, GY28-6659.)

The page of the IMDPRDMP listing containing the Queue Control Block trace is identified by two heading lines. The first line contains an optional title, the name of the module that invoked the dump, and the date and time that the information was gathered from the dumped system. The second line of the heading identifies the page as containing a Queue Control Block trace. The individual QCBs are then listed for each Task Control Block. Each Queue Control Block is formatted as follows:

MAJOR hhhhhh

The starting address of a major queue control block, the contents of which are given, indented, on the line or lines below.

NAME cccccccc The name of the system resource represented by the major QCB.

SAMPLE QCB TRACE MODULE IMDSADMP DATE 7/04/70 TIME 0.10 PAGE 2 QUEUE CONTROL BLOCK TRACE * * * * MAJOR 024100 NAME SYSDN MINOR 0239AO NAME FF SYS1.LINKLIB **OEL 024068** TCB 023488 SHARED NAME FF SYS1.MACLIB MINOR 023838 TCB 023448 QEL 023ED8 SHARED MAJOR 0235E8 NAME SYSIEFSD MINOR 0235C8 NAME FF Q5 QEL 023208 TCB 023480 EXCLUSIVE QEL 023C10 TCB 0238E0 EXCLUSIVE

Figure 20. Queue Control Block Trace Sample

MINOR hhhhhh The starting address of the minor queue control block. Contents are given on this line or the lines below.

NAME hh ccccccc The first two characters appearing after the NAME field identifier indicate the scope of the minor QCB being dumped. If the scope is given as hexadecimal FF, the name of the QCB is known to the entire operating system. If the scope indicator is hexadecimal 00 or 10 through F0, the name of the QCB is known only to the job step. The scope indicator shows the storage protection key of the TCB that enqueued this minor QCB. The NAME field also contains the name of the specific system resource represented by the minor QCB.

QEL hhhhhh The address of a queue element (QEL) associated with the minor QCB described on the line above. A QEL line appears for each resource requested by the task associated with the minor QCB.

TCB hhhhhh The starting address of the task control block of the requesting task. This task requests a specific system resource through the use of the QEL indicated on this line.

SHARED or EXCLUSIVE This indicator tells whether the system resource is available to one task (EXCLUSIVE) or several tasks (SHARED).

LINK PACK AREA MAPS

Information on routines residing in either the MVT link pack area or the MFT resident reenterable load module area of the dump system is available to the user through use of the LPAMAP (link pack area map) format statement.

For users who are processing an MVT dump, the IMDPRDMP program produces a listing of all routines loaded into the link pack area by the nucleus initialization program (NIP). For MFT dumps, this list contains information pertaining to all resident reentrant routines loaded into the reenterable load module area by NIP.

The IMDPRDMP user will find the link pack area map, for MVT, or the reenterable load module area map, for MFT, to be a useful tool in isolating system failures that occurred in program modules that reside outside the user's partition or region. If requested, the applicable map appears on a separate page of the IMDPRDMP program dump listing. A sample Link Pack Area map is shown in Figure 21.

The dump listing page containing the link pack area map is identified by two heading lines. The first line contains the optional title supplied by the user, the name of the module that invoked the dump, and the date and time that the information was gathered from the dumped system. The second line of the heading identifies the page as containing a link pack area map. Information on each module contained in the link pack area or reenterable load module area, is given in the following format:

- NAME cccccccc The name of the load module represented by this entry.
- EPA hhhhhh

The entry point address of the module identified on the corresponding line in the NAME column.

STA hhhhhh

The starting address of the named module's control section.

LNGH hhhhhh The length, in bytes, of the control sections in the load module described on this line.

TYPE CCCCC

The attributes of the control block associated with the module being described on this line. Under MVT, the type of the contents directory entry (CDE) associated with the module is given. The type may be either MAJOR or MINOR. Under MFT, the type is shown as either a loaded request block (LRB) or a loaded program request block (LPRB).

IMDPRDMP Output Formatting: Link Pack Area Maps

· 73

					MODULE	INDSADMP	DATE 11/12/70	TIME 00.
	* * *	* *	LINK	РАСК	AREA	МАР	* * * *	
NAME	EPA	STA	LNGH	TYPE				
IEELWAIT	072418 (072418	0003E8	MAJOR				
I GGO 209 Z	C748C0 (C74800	000400	MAJOR				
16G0201Z		C74C00	000400	MAJOR				
IGG0201¥		C75000	000400	MAJOR				
1 G G C 200 Z		075400	000400	MAJOR				
IGG0200Y		C75800	000400	MAJOR				
IGG0200H		C75CCO	000400	MAJOR				
I GGO 200G		C760C0	000400	MAJOR				
1 GG O 2 O O F		076400	000400	MAJOR				
IGGO200A		076800	000400	MAJOR				
IGG0199M		076000	000400	MAJDR				
I GGO 196B		077000	000400	MAJOR				
IGG0196A		C77400	C00400	MAJOR				
16601917		C778CO	000400	MAJOR				
IGG01911		077C00	C00400	MAJOR				
IGG01910		078000	C00400	MAJOR				
16601910		C78400	000400	MAJOR				
IGG0191G		C78800	000400	MAJOR				
IGG0191D		C78CCO	000400	MAJOR				
IGGC191B		000079	000400	MAJOR				
IGGC191A		C79400	000400	MAJOR				
IGG0190S		079800	000400	MAJOR				
IGG0190N		002673	000400	MAJOR				
IGG0190M		C7A000	000400	MAJOR				
IGG0190L		C7A400	000400	MAJOR				
IGC0005E		C74800	000400	MAJOR				
1900002		C7AC00	000400	MAJOR Major				
IGCOCOLI		C7B360 C7CA00	C00400 000060	MAJOR				
IGG019CK		C7CA60	0000E8	MAJOR				
IGG019BC		C7C848	000128	MAJOR				
IGG019BD		C7CC70	000000	MAJOR				
IGGO19AD IGGO19AL		070030	000158	MAJOR				
IGCO19AC		070848	0000E8	MAJOR				
IGG019CA		C7D930	000088	MAJOR				
IGG019CB		C7D988	000098	MAJOR				
IGGO19AG		C7DA50	000090	MAJOR				
IGG019BE		C7DAE0	000188	MAJOR				
IGG019AM		070068	000078	MAJOR				
IGGOL9AN		C7DCE0	000008	MAJOR				
IGG019AV		C7DDB8	000058	MAJOR				
1GG019M0		C7DE10	0000F0	MAJOR				
IGGC19MB		078760	0010A0	MAJOR				
IGGC19MA		07CE88	000978	MAJOR				
IGG019CL		07E820	000040	MAJOR				
IGG019CF		C7DFC0	000100	MAJOR				
IGG019CE		07E038	000088	MAJOR				
IGGC19AJ		076000	000120	MAJOR				
IGG019AI		C7E1E0	000080	MAJOR				
IGG019BB		C7E860	000058	MAJOP				
IGGO19BA		C7E260	000180	MAJOP				

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PACE 0001

Figure 21. Link Pack Area Map Sample

PAGE 0020 MODULE INDSADMP DATE 11/12/70 TIME 00.15 PROCSTEP STEP1 **JOB JOB4** STEP GO CURRENT TASK **** ***** TCB 02D400 RBP 0002E410 PIE 0000000 DEB 0002DABC TI0 0002E1F0 CMP 00000000 TRN 00000000 JLB 00000000 JPQ 0002E3E8 MSS 0002E770 PK-FLG FC000000 FLG 00001B1B LLS OCO2E3E0 RG 0-7 0C0C00C0 00000066 0002DFBC 00000000 00020660 0002D1E8 0002E234 00020BA8 RG 8-15 0C02CFA0 00000000 0002DFC8 0005DF08 400 5DE 56 0005DF08 6007F060 60008342 FSA 00068F68 TME 00000000 JST 0002D400 NTC 00000000 OTC 0002D1E8 TCB 00000000 SQS 0002DA90 LTC 0000000 IQE 00000000 ECB 0002DFC4 TSPR 0C000000 D-POE 0002E770 USR 0000000 DAR 00000000 RES 0000000 JSCB 0002E33C 0002CF28 STA CCCCCOOO TCT ACTIVE RBS PSW FFF50009 AC05DEF9 PRB 02E410 RESV 00000000 APSW 00000000 WC-SZ-STAB 00040082 FL-CDE 0002E5E8 EPA 05DE50 STA 05DE50 LN 0001B0 WT-LNK 00020400 ATR1 OB Q/TTR 000C0000 NM GO MAIN STORAGE D-POE 0002E770 FIRST 0002E688 LAST 0002E688 NPQ 00000000 PPQ 0000000 PQE 02E688 FFB 0005ECC0 LFB CCC5E000 TCB 0002D1E8 RSI COOOFOOO RAD 0005D800 FLG 0000 LOAD LIST EPA 050DC8 STA 050DC8 LN 000088 CDE 02E3E8 NM RETURNS USE OL RESP 01 ATR1 08 STA 07E928 LN CODOD8 ATR1 BO EPA 07E928 028850 NM IGGC19CC **USE 03** RESP 01 CDF ATR1 BO EPA 07E8B8 STA 07E888 LN 000070 NM IGG019CH RESP 01 CDE 028820 LSE 03 NM IGG019AC RESP 01 ATR1 BO EPA 070848 STA 070848 LN 0000E8 CDE 028730 LSE 02 RESP 01 ATR1 BO EPA 07F020 STA 07F020 LN 000078 CDE 0288F0 NM IGG019AQ **USE 03** JOB PACK QUEUE NM RETURNS ATR1 OB EPA 05DDC8 STA 0500C8 LN 000088 CDE 02E3E8 USE 01 RESP NA LN 0001BC CDE 02E5E8 NM GO USE 01 RESP NA ATR1 OB EPA 05DE50 STA 05DE50 PCI 000D72 CH END 000D72 JEB 02DABC APPENDAGES END OF EXT 07E8B8 SID 000D72 AB END 000072 C5000006 000108E0 11000000 PFX 0000000 ASYN F8000000 SPRG 00000000 UPRG 01068E18 PLST 18000000 DCB FF05DFA0 TCB 0402D400 NDEB 1000000 AVT 0402CA98 FM-UCB TRKS START END 580026AC 00020003 COC20003 0001 TIOT 02E1F0 JOB JOB4 STEP GO PROC STEP1 OFFSET CONAME TTR-STC STB-UCB LN-STA PGM=*.DD 800026AC 0018 14040101 00271500 002C 14040101 DUMMY 00271900 800026AC

Figure 22. Sample of MVT Major Control Block Format

74 Programmer's Guide to Debugging (Release 21)

MAJOR SYSTEM CONTROL BLOCK FORMATS

Formatting of the major system control blocks associated with a task is a function of either a FORMAT control statement, or one of the several noted parameters associated with the PRINT control statement. The control blocks of several tasks may be printed during one execution of IMDPRDMP. When more than one task is printed, the associated task control blocks (TCBs) are grouped into a TCB summary, listed following the printing of all requested tasks. This summary provides an index to the formatted TCBs by jobname. see the discussion "Task Control Block Summaries."

For ease of identifying various dump printouts, specific headings are printed on each dump; such as FORMAT, DAR AND F03 TASKS, PRINT CURRENT, and PRINT JOBNAME.

Each task being printed begins on a new page, identified by two heading lines. The first heading line contains the optional title supplied by the user, the name of the module that invoked the dump, and the date and time that the information was gathered from the dumped system, and a page number. The second line of the heading identifies the particular task being printed. This task information is broken down into the following named fields:

UNKNOWN JOB cccccccc The JOB field displays the eight-character name that was specified in the label field of the JOB statement.

UNKNOWN indicates that the job name is unavailable.

STEP cccccccc The STEP field shows the eightcharacter step name of the problem program associated with the task being dumped. This name was supplied in the label field of the EXEC statement.

PROCSTEP ccccccc If the job step being displayed was invoked from a cataloged procedure, the step name of the cataloged procedure, as contained in the cataloged procedure's EXEC statement, is displayed in this field.

If the task being printed was in control at the time the dump was taken, a third heading line follows the two previously described. The line "**** CURRENT TASK ****" identifies the TCB associated with the task in control when the dump was taken.

While formatting the dumped control blocks, IMDPRDMP may issue various output comments to assist the person who analyzes the printout. The output comments are discussed following the control block discussion.

Specific formatting of the major system control blocks is dependent upon the operating system option under which the dumped system was operating. To allow the reader to concentrate on the particular operating system with which he is concerned, the discussion of control block formatting is divided into three parts: MVT, MFT, and the TSO option of MVT.

MVT Control Block Formatting

The formats described below are repeated for each requested task that is printed. A sample of the major system control blocks, as formatted from an MVT dump, is shown in Figure 22.

<u>MVT TASK CONTROL BLOCK (TCB) FORMATTING</u>: The task control block (TCB) contains information that pertains to the specific task named in the heading lines that appear at the top of the page. Each TCB is formatted as follows:

TCB hhhhhh

The address of the task control block being displayed is given in this first field.

RBP hhhhhhhh

The address of the request block (RB) that was currently associated with the task represented by this TCB.

PIE hhhhhhhh

The address of the first program interrupt element (PIE) enqueued by this TCB.

DEB hhhhhhhh

The address of the beginning of the data extent block (DEB) queue that was associated with this task. Information on the contents of each DEB in the queue is given in a separate portion of this MVT dump listing.

TIO hhhhhhhh

The address of the task input output table (TIOT) that was constructed during device allocation for the task represented by this TCB. The contents of this table are displayed in a later portion of this task's display.

CMP hhhhhhh

This word contains ABEND indicators and user and system completion codes. The usage of this field is as follows:

byte 0

1	Bit 0 indicates that a
	dump had been requested.
.1	Bit 1 set indicates that
	a step ABEND had been
	requested.
xx xxxx	Bits 2 through 7 are
	reserved for future use.

bytes 1-3

The first 12 bits contain a system completion code. These codes and their meanings are explained in the publication <u>IBM System/360 Operating</u> <u>System: Messages and Codes</u>, GC28-6631 under the heading "System Completion Codes." A user completion code is contained in the last 12 bits.

For a system completion code of 0C0, denoting an imprecise interruption on a Model 91 or Model 195, digits six through eight contain the imprecise interruption configuration. These bit settings represent bits 16 through 27 of the program old PSW (location X'28') at the time of the program exception. Therefore, in the TCB CMP field, bits 20 through 29 are significant for a Model 91, while a Model 195 makes use of all the bits (20 through 31). For further discussion on interpreting the imprecise interruption configuration refer to the publication IBM System/360 Operating System: Supervisor Services, Order No. GC28-6646.

TRN hhhhhhhh

Contains flags and TESTRAN indicators as follows:

byte 0

1	Bit 0 set indicates that both TESTRAN and decimal
	simulator programs were
	being used on a
	System/360 Model 91
	machine.
.1	Bit 1 set indicates that
	checkpoints were not
	taken for this step.
	Bit 2 set indicates that
	the TCB being displayed

1	belonged to either a graphics foreground or the graphic job processor. Bit 3 set indicates that the TCB being displayed was associated with a 7094 emulator task that was being run on a System/360 Model 85 machine. Bit 4 is reserved for
	future use.
1	Bit 5 set indicates that this is a time shared task under control of the TEST command
1.	processor. Bit 6 set indicates that the OLTEP functions require cleanup before abnormal termination can
···· ··· x	be invoked. Bit 7 is reserved for future use.

bytes 1-3

The address of the control core table that was used by TESTRAN.

- MSS hhhhhhhh Main storage supervision information as follows:
 - byte 0 This byte determined roll-out eligibility for the job step associated with this TCB.
 - 00 in this byte indicated that the job step may be rolled out.
 - nz (nonzero) in this byte indicated that the job step may not be rolled out.

bytes 1-3 These bytes contain the starting address of the last subpool queue element (SPQE).

PK-FLG hhhhhhhh The storage protection key of the task and a series of flags. This word is divided into several subfields. These are:

byte 0

XXXX	The storage protection key of the task
0000	represented by this TCB. Always contain zeros.
byte 1	

1... Bit 0 set indicates thatan abnormal termination was in progress at the time the

.1	dump was taken. Bit 1 set indicates that a normal termination was	1	Bit 3 set causes the Purge routine in ABEND to enter when ABEND is
1	in progress at the time the dump was taken. Bit 2 set causes the Erase routine in ABEND to enter when ABEND is in control again.	1	in control again. Bit 4 set indicates that the Graphics Abnormal Termination routine was in control of the task associated with this TCB

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76.2 Programmer's Guide to Debugging (Release 21.7)

	at the time the dump was	huto 3	
	at the time the dump was taken. Bit 7 in byte 3 of this word must also	byte 3 1	Bit 0 set indicates that a PSW associated with
	be on.		the task represented by
1	Bit 5 set indicates that the top task in the TCB		this TCB was in the supervisor state.
	chain (usually the job	.1	Bit 1 set is applicable
	step TCB) was in the process of being		to job step TCBs. Setting of this bit
	terminated when the dump		indicates that the job
1.	was taken. Bit 6 set indicates that		step had invoked rollouts that were still
	an abnormal dump has		in effect at the time
1	been completed. Bit 7 indicates that		the dump was taken. Bit 2 set indicates that
	asynchronous exits could		ABEND was processing in
	not be scheduled.		such a manner as to prevent multiple ABENDS
byte 2	Bit 0 set indicates that		from occurring in the dumped system.
1	the SYSABEND (or	1x	Bit 3 set indicates that
	SYSUDUMP) data set for		the SYSABEND (or SYSUDUMP) data set is
	the job step is being opened. Operands of		being opened by this
	ABEND macro instruction have been saved in	···· 1··×	task. (See also bit 7.) Bit 4 set indicates that
	TCBCMP field.		an ABDUMP was in process
.1	Bit 1 set indicates that if this is an initiator		for the task associated with this TCB at the
	TCB, the second job step		time the dump was taken.
1	interval has expired. Bit 2 set indicates that		(See bit 7 of this byte.)
	for a job step TCB, the	1	Bit 5 set is applicable
	job step can cause rollout.		only for job step TCBs. With this bit set, no
1	Bit 3 set indicates that		abnormal termination
	the current task had a forced completion		dumps could have been provided within the job
	imposed upon it. Other tasks in the system		step represented by this TCB.
	could not have been	1x	Bit 6 set indicates that
	performed until the current task had been		a CLOSE had been issued during ABEND processing.
	completed.		(See bit 7 of this
1	Bit 4 set indicates that the job step had a	x x.x1	byte.) Bit 7 set, in
	forced completion		conjunction with bits 3,
	imposed upon it. Other tasks in the job step		4, or 6 of this byte or bit 4 in byte 1 of this
	could not have been		word indicates that, had
	performed until the present job step had		the dumped system been allowed to continue
1	been completed. Bit 5 set indicates that		processing without interruption by the
	the SYSABEND (or		IMDSADMP dump program, a
	SYSUDUMP) data set has been opened for the job		valid reentry to ABEND would have been
	step.		effected.
1.	Bit 6 set indicates that an EXTR exit was	FIG hhhhhhhh	
	requested by an	This field di	splays a further series
1	attaching task. Bit 7 set indicates that		certain priority This word is divided into
	the task associated with this TCB was a member of	subfields as	follows:
	a time-sliced group.		

byte 0 If any one of the flags comprising this byte were set at the time the dump was taken, the task represented by this TCB was considered to be non-dispatchable.

1	Bit 0 was set by ABDUMP					
•X•• ••••	Bit 1 is reserved for					
	future use.					
	Bit 2 set indicates that					
	the supply of I/O					
	request queue elements					
	(RQEs) had been					
	exhausted.					
•••X XX••	Bits 3 through 5 are					
	reserved for future use.					
1.	Bit 6 is applicable only					
	to M65 multiprocessing					
	situations. The setting					
	of this bit indicates					
	that the task					
	represented by this TCB					
	had been flagged					
	non-dispatchable by one					
	CPU to prevent any CPU					
	from working on it.					
1	Bit 7 set indicates that					
	the task associated with					
	this TCB entered the					
	ABEND routine while the					
	data control block					

representing the SYSABEND data set was being opened for another task.

byte 1

If any one of the flags comprising this byte were set at the time the dump was taken, the task represented by this TCB was considered to be non-dispatchable.

1	Bit 0 set indicates that the task represented by this TCB was terminated prior to the time the
.1	dump was taken. Bit 1 set indicates that the task represented by this TCB was a candidate for termination by
	ABEND.
1	Bit 2 set indicates that a routine of the task represented by this TCB issued an unconditional GETMAIN that could only have been satisfied by the rolling out of another job step.
1	Bit 3 indicates that the job step associated with this TCB was rolled out.
1	Bit 4 set indicates that another task was in system-must-complete status.

1	Bit 5 set indicates that another task in this job step was in step-must-complete status at the time the dump was taken.
••••	Bit 6 is applicable only for an initiator task.
	Setting of this bit
	indicates that a request
	for a region could not
	be satisfied.
1	Bit 7 is the primary
	non-dispatchability
	indicator. Setting of
	this bit indicates that
	one or more of the
	secondary
	non-dispatchability bits
	(bytes 1-3 of the DAR
	field) was set at the
	time the dump was taken.

byte 2

The dispatching priority limit for the task represented by this TCB.

- byte 3 The dispatching priority of the task represented by this TCB.
- LLS hhhhhhhh The load list element (LLE) for the program that was loaded by means of the LOAD macro instruction.
- JLB hhhhhhhh The address of the data control block associated with the JOBLIB associated with the task.
- JPQ hhhhhhhh Contains information pertaining to a job step TCB as follows:

byte 0

byce u	
1	Bit 0 set indicates that if the associated job
	5
	step had been allowed to
	continue processing
	without being
	interrupted by the dump
	program, the job step
	would have been purged.
•XXX XXXX	Bits 1 through 7 are
	reserved for future use.

bytes 1-3 The address of the last contents directory entry for a job pack area (JPA) control queue.

RG 0-7 and RG 8-15 The register save area of the TCB being displayed. The general registers were stored in this area upon entry to the first routine invoked in the task. On entry to any task, register 13 points to this TCB's register save area. This pointer is useful in locating the entry points of first routines and in tracing the save area chains.

- FSA hhhhhhhh The address of the first problem program save area.
- TCB hhhhhhhh The address of the TCB that had the next lowest priority on the ready queue at the time the dump was taken.
- TME hhhhhhhh The address of the timer element.
- JST hhhhhhh The address of the first TCB for a job step. For tasks with a storage protection key of zero (as shown in the first byte of the PK-FLG field), this word contains the address of this TCB.
- NTC hhhhhhh The address of the previous TCB that existed on the originating task's queue of subtask TCBs (sister). If this TCB was the first on the queue, this field contains zeros.
- OTC hhhhhhhh The address of the TCB representing the originating task (mother).

LTC hhhhhhh The address of the last TCB that existed on the originating task's queue of subtask TCBs at the time the dump was taken (daughter). If this TCB was the last on the queue, this field contains zeros.

IQE hhhhhhhh The address of the interruption queue element (IQE) that was used in scheduling the EXTR routine on the originating task.

ECB hhhhhhh The address of the event control block (ECB) that would have been posted by the supervisor's task termination routines had either normal or abnormal task termination been allowed to occur.

TSPR hhhhhhhh

byte 0 This field contains flags that indicate the status of the time sharing (TSO). Without TSO or when TSO has not been started, this field contains zeros.

- 1... Bit 0 set indicates that this task is a time sharing task.
- .1.....Bit 1 set indicates that the time sharing task should be set non-dispatchable. This bit was set by the TCBSTP routine while the routine was not executing as a privileged program.
- ..1. Bit 2 set indicates that the system is executing and requires that the time sharing task must not be interrupted by the attention exit or by the STATUS SVC.
- ...1 Bit 3 set indicates that a terminal I/O purge is required.
- xxxx Bits 4 through 7 are reserved for future use.

byte 1

This field contains the number of SET STATUS starts required to make this time sharing task dispatchable.

byte 2 This field contains the limit priority of the time sharing task.

byte 3 This field contains the dispatching priority of the time sharing task.

D-PQE hhhhhhhh The address of the region dummy partition queue element minus 8 (DPQE-8).

SQS hhhhhhhh

The address of an allocated queue element (AQE) which contains the amount of available bytes assigned to this task in the system queue area (SQA), and a pointer to the next AQE for this task.

STA hhhhhhhh

Internal STAE routine flags and the address of the STAE control block that was in effect at the time the dump was taken.

TCT hhhhhhhh This word contains information pertaining to the dumped system's timing control table (TCT). The TCT field is divided into the following two subfields:

byte 0 Reserved for future use.

option was pr	management facilities esent in the dumped	1	Bit 0 set indicates that the task represented by this TCB is temporarily
	bytes contain the e dumped system's timing •	.1	non-dispatchable. Bit 1 set indicates that the task represented by this TCB is permanently
USR hhhhhhhh This word is available to the user of the dumped system. It contains any information placed in it by the user.		••*** ••• <i>•</i>	non-dispatchable. Bits 2 and 3 are recovery management support and system error recovery flags. Their
DAR hhhhhhhh The contents	of this field were used	1	meanings are: Bit 2 set indicates that
(DAR). Certa	assessment routines in subfields displayed in e also used to control		the task represented by this TCB is temporarily non-dispatchable.
the dispatcha	bility of the dumped	1	Bit 3 set indicates that
task. The DA the following	R field is divided into subfields:		the task represented by the TCB is permanently non-dispatchable.
byte 0		···· X···	Bit 4 is reserved for
	e of the DAR field	1	future use.
as follows:	flags. These flags are	•••• • 1 ••	Bit 5 set indicates that this task is temporarily non-dispatchable. Timer
1	Bit 0 set indicates that		services have been
	primary DAR recursion		requested and the
	occurred in the dumped system. The damage		time-of-day clock is still inoperative.
	assessment routine	•••• ••XX	
	failed while writing a		reserved for future use.
	main storage image dump.		
.1	Bit 1 set indicates that	byte 2	
	secondary DAR recursion occurred in the dumped system. The damage	follows:	ings for byte 2 are as
	assessment routine	x	Bit 0 is reserved for
	failed while attempting to reinstate a failing	.1	future use. Bit 1 set indicates that
••XX ••••	region or partition. Bits 2 and 3 are		this task has been stopped by a STATUS
	reserved for future use.		stop.
1	Bit 4 set indicates that the system error task is	1	Bit 2 set indicates that task is
	failing. The DAR dump		non-dispatchable. An
	should not request any		SVC dump is executing
	error recovery procedure (ERP) processing.	1	for another task. Bit 3 set indicates that
···· ·XX.	Bits 5 and 6 are	1	this task is being
	reserved for future use.		swapped out by the time
1	Bit 7 set indicates that		sharing (TSO).
	an SVC dump is executing	1	Bit 4 set indicates that
byte 1	for this task.		this task is in an input wait state.
-	gh 3 of the DAR field are	1	Bit 5 set indicates that
used to store			this task is in an
	bility flags. If any of		output wait state.
	in this subfield were	•••• ••XX	Bits 6 and 7 are
	in this subfield were ary non-dispatchability	•••• ••XX	Bits 6 and 7 are reserved for future use.
	in this subfield were ary non-dispatchability bit in the FIG field)	xx byte 3	
will also hav The bit setti	in this subfield were ary non-dispatchability bit in the FLG field) e been non-dispatchable. ngs that may appear in		reserved for future use.
will also hav	in this subfield were ary non-dispatchability bit in the FLG field) e been non-dispatchable. ngs that may appear in	byte 3	reserved for future use.

routines. Their meanings are:

JSCB hhhhhhhh The address of the job step control block.

RES hhhhhhhh Reserved for future use.

IOBRC hhhhhhhh Contains the address of the IOB restore chain for I/O quiesced by end-of-task.

<u>MVT ACTIVE REQUEST BLOCK (RB) FORMATTING</u>: Request blocks (RBs) were used by the lines at the top of the dump page and in the preceding TCB display, are listed in the portion of the dump listing labeled "ACTIVE RBS." Information on each RB associated with the task is formatted as shown below:

PRB

IRB hhhhhh SVRB

STRB

Each RB display is preceded by a field that indicates the type and address of the RB being displayed. The four types of RBs that may be displayed under an MVT task are:

PRB

program request block

IRB

interruption request block SVRB

supervisor request block (SVRBs may be divided into two categories; type 2 for resident routines and type 3 or 4 for transient routines)

SIRB

system interruption request block.

The type acronym for each RB is displayed in the first portion of the field. The starting address of the indicated request block appears in the last portion of the field. The contents of certain fields in the body of the formatted display are dependent upon the type of RB being displayed. Variations in display field usage are noted in the descriptions of the fields in which they occur.

RESV

TAB-LN hhhhhhhh

FL-PSA

This field shows both the function and the first word of the request block being displayed. The meanings of the function indicators and the values that follow them are:

RESV indicates that the request block is either a PRB or an SVRB for resident routines. The first word of these particular RBs is reserved for future use and contains zeros.

TAB-LN

indicates that the request block being displayed is used as an SVRB for transient routines. The value field is divided into two subfields of two bytes each. The first two bytes show the displacement of the entry point of the module represented by this SVRB from the beginning of the transient area control table (TACT). The second subfield shows the length, in bytes, of the SVC routine.

FL-PSA

indicates that the RB being displayed is an IRB. The value portion of this field is divided into two subfields. The first subfield has a length of one byte and contains indicators for the timer routines. When there were no timer routines, this field contains zeros. The timer routine indicators set at the time the dump was taken are shown as:

1	••••	indicates that the timer element was not
.1	••••	on queue. indicates that the local time-of-day
00	••••	option was used. indicates that the time interval was
01		requested in timer units. indicates that the time interval was
		requested in binary units.
11	••••	indicates that the time interval was
		requested in decimal form.
• • • •	1	indicates that the time interval had
••••	.000	expired. indicates a task
••••	.100	request indicates a task
		request with an exit specified.
• • • •	.001	indicates a wait
	.011	request. indicates a real
		request.
••••	.111	indicates a real request with an exit specified.

The second subfield is three bytes long and contains the starting address of the problem program register save area (PSA).

- APSW hhhhhhhh
 - The APSW field displays information pertaining to the program status word that was active at the time the dump was taken. The functional variations associated with the usage of this field are:
 - PRBs being formatted contain the right half (bytes 4 through 7) of the problem program's PSW when an ABTERM interruption occurred.
 - IRBs, SIRBs, and SVRBs for resident routines use this field to display the right half (bytes 4 through 7) of the PSW that was active, in the dumped system, during the execution of an ABEND or ABTERM routine. If no ABEND or ABTERM routine was envoked in the dumped system, this field contains zeros.
 - SVRBs for transient routines use this field in much the same way as SVRBs for resident routines. If an ABEND or ABTERM routine was invoked in the dumped system, bytes 4 through 7 of the associated PSW are displayed in this field. If an ABEND or ABTERM routine was not invoked, this field contains the last four characters of the name of the requested routine. (The last two characters of the name represent the SVC number.)
- WC-SZ-STAB hhhhhhhh This field contains information pertaining to wait conditions, request block sizes, and RB status and attribute characteristics. This field is divided into three subfields, as follows:
 - byte 0
 The wait count that was in effect at
 the time of the dump.
 byte 1
 The size of this request block. This
 RB size is expressed as the number of
 doublewords comprising the block.
 byte 2
 The last two bytes of the WC-SZ-STAB
 field contain bit settings that
 reflect the status and attributes of
 the request block. The settings that

may appear in byte 2 are:

xx..... Bits 0 and 1 indicate the type of RB being displayed. The possible

	settings for these two bits and their meanings
	are:
00	This is a program
	request block (PRB).
01	This is an interrupt
	request block (IRB).
10	This is a system
	interrupt request block
	(SIRB).
11	This is a supervisor
	request block (SVRB).
••x• x•xx	Bits 2, 4, 6 and 7 are
	reserved for future use.
1	Bit 3 set indicates that
	this request block is an
	SVRB for a transient
	routine.
1	Bit 5 is applicable only
	if the request block
	being displayed is an
	SVRB. If this bit is
	set, a checkpoint could
	have been taken in a
	user exit from the SVC
	routine associated with
	this RB.
e 3	CHIS RD.
	of the WC-SZ-STAB field
ontains more	status and attribute

The last byte of the WC-SZ-STAB field contains more status and attribute flags. The possible settings for this subfield and their meanings are:

byte

1	Bit 0 set indicates that
	the WT-LNK field in this
	RB display, contains in
	its last three bytes,
	the address of the TCB
	to which this request
	block is linked.
.1	Bit 1 applies only to
	IRBs and SIRBs. If this
	bit is set, the
	indication is that at
	the time the dump was
	taken, the program
	associated with this RB
	was active.
••X• ••••	Bit 2 is reserved for
÷	future use.
1	Bit 3 is applicable only
	to IRBs. The setting of
	this bit is an
	indication that the IRB
	was associated with an
	ETXR exit routine.
XX	Bits 4 and 5 concern
	interruption queue
	elements (IQEs) and
	request queue elements
	(RQEs). This flag is
	used as follows:
00	This setting indicates
•••• VV••	
	that the request queue
	element was not to be
	returned to the free
	list when the exit was
	taken.

- 01.. This setting indicates that the IRB had queue elements for asynchronously executed routines that were RQEs. This setting is applicable only if the RB being displayed is an IRB.
- 10.. This setting indicates
 that the IQE was not to
 have been returned at
 EXIT.
 11.. This setting is
- 11.. This setting is applicable only to IRBs. If this setting appears, the indication is that the IRB had queue elements for asynchronously executed routines that were IQEs. Bit 6 set indicates that request block storage could be freed at the time of exit.
 - wait conditions. The meanings of the two possible settings for this bit are:
- Bit 7 not set indicates that the request had to wait for a single event or all of a number of events.
- Bit 7 set indicates that the request had to wait for a number of events. This number of events was less than the total number of events that were waiting.

FL-CDE

EPA hhhhhhhh

TQN

This field shows both the function and the fourth word of the request block being displayed. The meaning of the function indicator and the value following it is given below:

FL-CDE

the request block being displayed is either a PRB or an SVRB for a resident routine. The value field is divided into two elements. The first subfield has a length of one byte and contains control flag settings.

These control flags are as follows:

xxxx x... Bits 0 through 4 are reserved for future use.

-1.. indicates that a SYNC macro instruction was requested.1. indicates that an XCTL macro instruction was
- requested. indicates that a LOAD macro instruction was requested.

The second subfield is three bytes long and contains the address of the contents directory entry (CDE) representing the module that this request block was associated with.

EPA

The request block being displayed is an IRB. The value field contains the entry point address of a routine that was asynchronously executed.

TQN The request block being displayed represents a transient routine SVRB. The value field contains the address of the next request block that was on the queue of transient routines.

PSW hhhhhhhh hhhhhhh

The resume program status word. This PSW represents the status of the program represented by the RB being displayed when a <u>new</u> RB was created. Had the dumped system been allowed to continue processing without being interrupted by the dump program, operation would have resumed on this PSW.

Q/TTR hhhhhhhh

This word is used to display various data, depending upon the type of request block being displayed. Usage of the Q/TTR value field is used by each type of request block as follows:

- PRBs and SVRBs that represented resident routines do not use this field; the first byte always contains zeros. Bytes 1 through 3 of the field show the address of a request block that requested the use of the same serially reusable program.
- IRBs utilize this field in one of two ways, to show either the three-byte link-field segment or the two-byte link-field segment, depending upon the IRB usage. The three-byte link-field segment appears in the Q/TTR value field as follows:

byte 0 Contains a count of the number of requests for the same exit (ETXR). This use count is utilized by the ATTACH macro instruction.

byte 1-3 Contains the starting address of the queue of interruption queue elements (IQEs).

Alternately, the Q/TTR value field may be formatted to show the two-byte link-field segment. In this instance, the field is used thusly:

byte 0-1 Reserved for future use.

- bytes 2-3 The starting address of the queue of request queue elements (RQEs).
- SVRBs that represented transient routines display two data elements in this field. The first subfield has a length of one byte and shows the number of requests if the transient routine was overlaid. The last three bytes of the Q/TTR field contain the relative direct access device address for the associated supervisor routine in the form TTR.

WT-LNK hhhhhhhh

This field displays information pertaining to wait counts and request block linkages. In the case of a transient SVC, if this field contains x'FF', either the routine represented by the SVRB is currently being brought into the transient area, or this routine has been displaced in the transient area by a routine requested by a higher priority task. To tell what has happened, compare the APSW and NM field contents as described under NM below. This field is divided into two subfields, one with a length of one byte and the other with a length of three bytes. These subfields show the following:

byte O

The number of requests that were pending at the time the dump was taken (wait count).

byte 1-3

The address of the next request block on the RB queue. In the last RB on the queue, this field contains the address of the task control block (TCB). NM ccccccc

The eight character name of the load module represented by the request block being displayed with a possible exception for transient SVRBs.

If byte 0 of the WT-LNK field contains x'FF', it is possible that the module represented by this SVRB has been overlaid in the transient area by a module requested by a higher priority task. Compare the APSW field, (providing it contains the four low-order bytes of a module name) with the last four characters (the hexadecimal should be translated to EBCDIC) of the module name in the NM field. No match indicates the user of the transient area has been pre-empted by a higher priority task. NM therefore represents the module currently in the transient area, not the module represented by this SVRB.

If a match results, NM correctly identifies the module name requested by this SVRB.

EPA hhhhhh

The address of the entry point of the module named in the NM field of this RB display.

STA hhhhhh The starting address of the module identified in the NM field of this

RB's display.

The length, in bytes, of the load module that is represented by this request block.

ATR1 hh

This one byte field displays the attributes of the described module. These attributes are taken from the contents directory entry associated with the module. The meanings of the attribute flag settings are given below:

1	Bit 0 set indicates that the module was resident
.1	in the link pack area. Bit 1 set indicates that
	at the time the dump was
	taken, the module
	represented by this
	request block was in the
	process of being
	fetched.
1	Bit 2 set indicates that
	the module was
	reenterable.
1	Bit 3 set indicates that

...1 Bit 3 set indicates that the module was serially reusable.

Bit 4 set indicates that 1... the module could not have been reused. This flag setting is not applicable if either bit 2 or 3 is set.1.. Bit 5 set indicates that the contents directory entry associated with this module reflects the use of an alias name. This information applies only to minor CDEs. Bit 6 set indicates that1. the module was in the job pack area. Bit 7 set indicates that1 the module was considered not only-loadable.

<u>MVT MAIN STORAGE INFORMATION</u>: Each task operating under the MVT option of the operating system was dynamically assigned a region of main storage that consisted of one or more 2K-byte subpool areas. To keep track of main storage allocations, the MVT supervisor maintained a partition queue associated with each region. Composed of partition queue elements (PQEs), and residing in the system queue area, this partition queue was connected to the TCBs for each task in a job step through a dummy partition queue element (DPQE).

Information on the areas of main storage allocated to each task, is presented to the user in a separate portion of each task's dump listing headed "MAIN STORAGE." This main storage information is formatted as shown below:

- D-PQE hhhhhhhh The address minus eight bytes of the dummy partition queue element (DPQE-8) connecting the partition queue to this task's TCB.
- FIRST hhhhhhhh The starting address of the first partition queue element (PQE) on this region's partition queue.
- LAST hhhhhhhh The starting address of the last PQE on the partition queue.

PQE hhhhhh The starting address of one of the partition queue elements on the partition queue bounded by the addresses given on the line above.

FFB hhhhhhhh The starting address of the first free block queue element (FBQE) on the free block queue associated with this PQE. If no FBQEs exist, this field contains the address of the PQE being displayed

LFB hhhhhhhh

The starting address of the last free block queue element (FBQE) on the free block queue associated with this PQE. If no FBQEs exist, this field shows the starting address of this PQE.

NPQ hhhhhhh The starting address of the next partition queue element on the partition queue. If the PQE being displayed was the last PQE on the queue, this field contains zeros.

PPQ hhhhhhhh

The starting address of the partition queue element on the partition queue that preceded this PQE. If this PQE was the first on the queue, this field contains zeros.

- TCB hhhhhhh The starting address of the TCB of the job step to which the described region is assigned. If this field contains zeros, the indication is that the area of main storage was obtained from unassigned free space.
- RSI hhhhhhhh The size of the region being described. This number is a multiple of 2K (2048).
- RAD hhhhhhhh The starting address of the region being described by this PQE.
- FLG hhh The FLG field shows the settings of several PQE flags whose meanings are given below:

x	Bit 0 indicates region
	ownership. The meanings
	of the settings are:
0	indicates that the space
	described by this PQE
	was owned by the
	associated task.
1	indicates that the space
	described by this PQE
	was borrowed.
.1	The setting of bit 1 is
	meaningful only if bit 0
	was not set. If this
	bit is set and bit 0 is
	not set, the indication
	is that the region had
	been rolled out.
· .	
1	Bit 2 set indicates that
	the region described by
	this PQE was borrowed by
	another task.
X XXXX	Bits 3 through 7 are
	reserved for future use.

MVT LOAD LIST FORMATTING: A load list was maintained by the dumped system's supervisor in order to keep track of the load modules that were in main storage and the area of main storage each occupied. The load list maintained by a system operating under the MVT option of the operating system contained a series of load list elements (LLEs), each of which was associated with a particular load module through the use of a control block called a contents directory entry (CDE). A formatted listing of the dumped system's MVT load list appears as follows:

CDE hhhhhh

The starting address of the contents directory entry associated with this load list item.

NM ccccccc

The eight-character name of the entry point to the load module represented by this entry.

USE hh

The count of the number of uses (through the ATTACH, LINK and XCTL macro instructions) of the load module, and the number of times a LOAD macro instruction was issued for the module.

RESP hh

The responsibility count contained in the load list entry associated with the load module. This count indicates the number of requests made by the LOAD macro instruction for the indicated load module. This count was decremented by one for each occurrence of the DELETE macro instruction.

ATR1 hh

The attributes of the load module described in this load list entry. These attributes are taken from the contents directory entry associated with the module. The meanings of the attribute flag settings are given below:

1... Bit 0 set indicates that the module was resident in the link pack area. Bit 1 set indicates that .1.. at the time the dump was taken, the load module represented by this load list element was in the process of being loaded. Bit 2 set indicates that the load module was reenterable. ...1 Bit 3 set indicates that the load module was serially reusable.

1	Bit 4 set indicates that the load module could
	not have been reused.
	This flag setting is not
	applicable if either bit
	2 or 3 is set.
1	Bit 5 set indicates that
	the contents directory
	entry associated with
	this load module
	reflects the use of an
	alias name. If this bit
	is set, this line of the
	load list display
	reflects information
	taken from a minor CDE.
	Bit 6 set indicates that
	the load module was in
	the job pack area.
1	Bit 7 set indicates that
	the load module was
	considered not
	only-loadable.

EPA hhhhhh

The address of the entry point of the load module named in the NM field of this load list display line.

STA hhhhhh

This field contains the starting address of the load module identified in the NM field of this load list display line.

LN hhhhhh The LN field supplies the user with the length, in bytes, of the load module represented by this load list entry (LLE).

<u>MVT JOB PACK QUEUE FORMAT</u>: A job pack area control queue (JPACQ) exists for each job step in the dumped system that used a program not in the link pack area. The job pack queue, like the link pack area, is made up of contents directory entries (CDEs). This area describes routines in a job step region that were brought into main storage by contents supervision routines to perform a task in the job step. The IMDPRDMP program displays the contents of the dumped MVT system's job pack queue as follows:

CDE hhhhhh

The starting address of the contents directory entry associated with this job pack queue element.

NM ccccccc

The eight-character name of the entry point to the load module represented by this entry. USE hh The count of the number of uses (through the ATTACH, LINK and XCTL macro instructions) of the load module, and the number of times a LOAD macro instruction was issued for the module.

RESP NA

This responsibility count field is flagged 'NA' to indicate that the information is not applicable to modules displayed in the job pack queue.

ATR1 hh

The attributes of the load module described in this job pack queue entry. These attributes are taken from the contents directory entry associated with the module. The meanings of the attribute flag settings are:

1... Bit 0 set indicates that the module was resident in the link pack area. .1.. Bit 1 set indicates that at the time the dump was taken, the load module represented by this job pack queue entry was in the process of being loaded.

..1. Bit 2 set indicates that the load module was reenterable.

...1 Bit 3 set indicates that the load module was serially reusable.

2 or 3 is set.

Bit 4 set indicates that

applicable if either bit

Bit 5 set indicates that

the load module could

not have been reused. This flag setting is not

.... 1...

.... .1..

the contents directory entry associated with this load module reflects the use of an alias name. If this bit is set, this line of the job pack queue display reflects information taken from a minor CDE. Bit 6 set indicates that the load module was in the job pack queue area. Bit 7 set indicates that1 the load module was considered not only-loadable.

EPA hhhhhh

The address of the entry point of the load module named in the NM field of this job pack queue entry display line. STA hhhhhh

This field contains the starting address of the load module identified in the NM field of this job pack queue entry display line.

LN hhhhhh

The LN field supplies the user with the length, in bytes, of the load module represented by this job pack queue entry.

MVT DATA EXTENT BLOCK (DEB) FORMATTING: Data extent blocks (DEBs), describing a data set's external storage requirements, were queued to those task control blocks (TCBs) that represented tasks requiring auxiliary storage input/output processing. External storage information, taken from each DEB, is formatted as shown below:

DEB hhhhhh The starting address of the basic section of the DEB being displayed.

APPENDAGES

The word "appendages" informs the user that the five named fields on this line contain information taken from the appendage vector table preceding the DEB being displayed. The named fields appearing on the rest of this line are:

END OF EXT hhhhhh The entry point of the end-of-extent appendage routine.

SIO hhhhhh The entry point of the start I/O appendage routine.

PCI hhhhhh The entry point of the program-controlled-interruption appendage routine.

CH END hhhhhh The entry point of the channel-end appendage routine.

AB END hhhhhh The entry point of the abnormal-end appendage routine.

PFX hhhhhhh hhhhhhh hhhhhhh The second line of a DEB display contains information taken from the prefix section of the DEB being displayed. The area is subdivided as follows:

byte 0 The first byte of the prefix area contain the contents of the I/O support work area. This area is used only by DEBs dealing with direct access storage devices.

bytes 1-7 The next seven bytes of the DEB prefix section are used by DEBs associated with direct access storage device functions. This subfield displays the data set control block's (DSCB) address used by I/O support. The address is expressed in the following format: bytes 1 and 2 the bin (cell) number. bytes 3 and 4 the cylinder address. bytes 5 and 6 the track address. the record number. byte 7 bytes 8-11 The third word of the PFX field contains the data control block (DCB) modification mask that was used by I/O support. byte 12 The length of the DEB in doublewords . bytes 13-15 The remainder of the DEB prefix section is reserved for future use. **TCB** hhhhhh This field marks the beginning of the basic section of the data extent block. The TCB field is divided into two subfields as follows: byte 0 The number of subroutines for which a LOAD macro instruction was issued during the execution of the OPEN executor routines. bytes 1-3 The starting address of the task control block to which this DEB was enqueued. NDEB hhhhhh The NDEB field is also used to display two data elements. It is subfielded as follows: bvte 0 The overall length of a data extent block includes the length of a variable length access method dependent section. The first byte of the NDEB field, expresses the length of the access method dependent section in bytes. If the access method was BDAM, this indicator is expressed as a number of fullwords. bytes 1-3 The last portion of the NDEB field displays the starting address of the basic section of the next DEB on the task's queue. If this DEB was the last on the queue, the contents of this field are the starting address of the TCB that enqueued this DEB.

ASYN hhhhhhhh This field contains fata set status flags and the address of the associated IRB. This field is used as follows: byte 0 The first byte of the ASYN field contains data set status flags. These flags have the following meanings: xx..... Bits 0 and 1 indicate the data set's disposition. The possible settings are: 01.. This setting indicates that the disposition was OLD. 10.. This setting indicates that the disposition of the data set was MOD (modify). 11.. This setting indicates that the disposition was NEW. ..1. Bit 2 set indicates that an end-of-volume (EOV) or end-of-file (EOF) condition had been encountered.1 The setting of bit 3 has one of two meanings depending upon the external storage medium. For disk this indicator reflects a release of unused external storage. For tape, the meaning of this indicator is that an emulator tape with second generation format was being used. Bit 4 set is a data 1... control block (DCB) modification indicator.1.. Bit 5 set has two meanings, depending upon the auxiliary storage recording medium. For disk, the setting of bit 5 indicates that a split cylinder was encountered. For tape, this flag indicates that an emulator tape with possible mixed parity records was used. Bit 6 set indicates the1. use of nonstandard labels.1 Bit 7 set indicates that reduced error recovery procedures were used on magnetic tapes containing the data set represented by this DEB.

bytes 1-3 The last portion of the ASYN field shows the starting address of the IRB that was associated with asynchronous appendage exit scheduling.		specified in this DEB.	f extents that were the DSCBs associated with
	ontains information on I/O		of the first input/output In the user's purge chain.
	thods and the system 2. The usage of this collows:		y and supervisor purge are contained in this
	e of this field contain dicate the method of	field. This follows:	field is formatted as
disposition of to have been	processing and the of the data set that was performed when an end-of-	byte 0 The priority this DEB was	of the task under which enqueued.
settings are:	ion occurred. These flag	bytes 1-3	
		-	address of a parameter
1	Bit 0 was set by ABEND.		s used to locate the purge
	The setting of this bit		l block (ECB) for a
	indicates that the data set associated with this	supervisor pu	irge request.
	DEB was a SYSABEND or	DCB hhhhhhhh	
	SYSUDUMP data set.		d contains three data
.0	Bit 1 is always zero.		nese are displayed in the
•••×× •••••	Bits 2 and 3 show the end-of-volume disposition procedure	format given byte 0	pelow:
	disposition procedure. The values for this flag		The storage protection
	are:		key that was associated
01	REREAD		with the task under
	LEAVE		which this DEB was
XXXX	The last half of this byte contains flags that	1111	enqueued. A hexadecimal 'F' in
	indicate the type of		bits 4 through 7 of this
	input/output processing		field identify this
	that was performed on		control block as a data
	the data set represented		extent block (DEB).
	by this DEB. The values for this flag are:	bytes 1-3	
0000	INPUT	-	address of the data
1111	OUTPUT		k (DCB) that was
0011	INOUT		ith this DEB.
0111	OUTIN		
0001	RDBACK	AVT hhhhhhhh	
0100	UPDAT		d displays two DEB data is subfielded as follows:
byte 1		Cremenco and	- Superseded as LOREOWS.
The quiesce of	count. The byte is	byte O	
	ith the system PURGE		nt scale that is used to
	C 16) and indicates the		e size of the device
	kiliary storage devices ecuting the user's channel		ction of this DEB. For s devices, a 4 is
programs.	ecucing the user's chamler		this subfield. For a
£		nondirect acc	cess device or a
bytes 2-3			n device, a 2 is
Reserved for	future use.	displayed.	
UPRG hhhhhhhh		bytes 1-3	
	ld contains extent		s the last portion of the
information a	and data used by the	AVT field she	ows the starting address
user's purge	routines. This field is	of the append	dage vector table
	the following two		is DEB. This table of
subfields:			utine addresses appears on ne of this DFB's display.

the first line of this DEB's display.

OP-UCB hhhhhhh byte 0 The contents of this field have The device modifier showing the file meaning only when the DEB being mask. displayed describes a data set that bytes 1-3 was assigned to a unit record or magnetic tape device. This The starting address of the unit control block (UCB) that was information is formatted from the device dependent section of the DEB. associated with the data extent. The OP-UCB field is subfielded as follows: START hhhhhhhh The address of the beginning of the direct access device extent. The byte 0 first four characters represent the This first subfield is applicable only to data sets assigned to magnetic tape cylinder address and the last four devices and shows the SET MODE characters represent the track operation code. For a data set that address. was assigned to a unit record device, this subfield is reserved. END hhhhhhhh The address of the end of the data bytes 1-3 extent. Cylinder and track references The starting address of the unit are formatted as in the extent control block (UCB) associated with beginning address, described above. the data set described by the DEB being displayed. TRKS hhhh The number of direct access tracks The following four fields are present only bounded by the starting and ending for data sets assigned to the IBM 3525 Card addresses shown in the previous two Punch for multi-function. The information columns. is formatted as shown below: MVT TASK INPUT/OUTPUT TABLE (TIOT) UCB hhhhhhhh FORMATTING: A task input output table (TIOT) was constructed for each task in the byte 0 The device modifier field (not used dumped system by MVT jcb management for the 3525). routines. Residing in the system queue area, this table contained primary pointers bytes 1-3 to control blocks used by I/O support The starting address of the unit control block (UCB) associated with routines. As the functions of several TIOT the data set described by the DEB fields were dependent upon the state of being displayed. associated external storage devices, multiple definitions may apply. The TIOT that was constructed in an MVT system is RDRDCB hhhhhhhh The starting address of the data formatted as shown. control block (DCB) for the read associated data set. TIOT hhhhhh The starting address of the task PCHDCB hhhhhhhh input/output table being displayed. The starting address of the data control block (DCB) for the punch JOB ccccccc The eight-character name of the job associated data set. for which this TIOT was constructed. WTRDCB hhhhhhhh The starting address of the data STEP ccccccc control block (DCB) for the print The eight-character name specified in associated data set. the label field of the EXEC JCL statement associated with this job The final portion of a DEB display shows step. information pertaining to a data set that was assigned to a direct access device. PROC ccccccc This information, taken from the DEB's If the job step for which this TIOT device dependent section, is arranged in was constructed was invoked from a columnar format with a line for each cataloged procedure, the procedure extent. The information is formatted as name, as contained in the EXEC JCL shown below: statement, is displayed in this field. Each data set associated with the indicated FM-UCB hhhhhhhh The first column displays two data task is represented by a separate DD entry elements and is formatted as follows: that is included in the TIOT. Each TIOT

entry is displayed on a separate line in

colum	anar fo	rmat	. The	use	and	meaning	of
each	column	is	given l	below	/:		

- OFFSET hhhh The offset of this DD entry from the beginning of the TIOT in hexadecimal.
- LN-STA hhhhhhhh Four bytes of length and status information, described below:

byte 0

The total length (including all device entries) in bytes of the DD entry being displayed on this line.

byte 1

Status byte A, one of three status bytes in a TIOT entry. The meanings of the status byte settings are:

xx	Bits 0 and 5 indicate
	the tape label
	processing that was to
	have been performed.
	The meanings of the
	settings are:
0	Nonlabeled tape or an
	indication to bypass
	label processing.
01	Standard labels and
	standard user labels.
1	Nonstandard labels.
.1	The setting of status
	bit 1 has two meanings,
	depending upon the
	processing phase that
	had been reached at the
	time the system was

dumped. During

allocation processing, the setting of this bit indicates that this entry represents a split cylinder primary space allocation DD. If the dump was taken during

..1.

step termination processing, the setting of this bit indicated that no unallocation of space was necessary. The setting of status bit 2 works under the same philosophy as status bit 1. During allocation processing, the setting of this bit indicates that this entry represents a split cylinder secondary space allocation DD. If the dump was taken during step termination processing, the indication was one of rewinding with no unload.

Bit 3 set indicates that ...1 this DD entry represents a JOBLIB. 1... Bit 4 set indicates that direct access device space management was deemed necessary. The setting of bit 61. specifies that the tape volume was to have been rewound and unloaded. The setting of bit 7 •••• specifies that the tape volume was to have been rewound.

byte 2

The third byte of this column has meaning only during the allocation phase. This displays the number of devices that were requested by the data set represented by the TIOT entry displayed on this line.

byte 3

The last byte of the LN-STA field displays a TIOT field that had meaning at two points during the processing of this task. During the allocation process, this field contained a link to the appropriate prime split, unit affinity, volume affinity or suballocate TIOT entry. After CLOSE processing, this byte was used thusly:

1	The setting of bit 0
	indicates that the data
	set represented by this
	DD entry was a SYSOUT
	data set that contained
	data.
•XXX XXXX	Bits 1 through 7 are
	reserved for future use.

DDNAME cccccccc The eight-character DD name associated with the TIOT entry being displayed.

TTR-STC hhhhhhh The first three bytes of this column display the relative track address (TTR) of the job file control block (JFCB) associated with this entry.

STB-UCB hhhhhhhh The last column in a TIOT display contains information taken from the one-word device entries that are appended to each TIOT entry. One TIOT device entry exists for each allocated device. This display field shows this information in the following format:

byte O	
Status byte B.	. The status bits have
the following	meanings:
1	Bit 0 set indicates that
	the data set associated
	with this line of the
	TIOT display was present
	on the device
	represented by this TIOT
	device entry.
.1	Bit 1 set indicates that
	the data set associated
	with this line of the
	TIOT display would have
	used the device
	represented by this TIOT
	device entry.
1	Bit 2 set indicates that
****	the device represented
	by this device entry
1	violated separation. Bit 3 set indicates that
	a volume serial number
	was present.
···· 1···	Bit 4 set indicates that
	a setup message was
	required.
•••• •X••	Bit 5 indicates the
	device disposition that
	would have taken place
	had the dumped system
	been allowed to continue
	processing this task.
	The settings for this
	bit are:
0	Indicates t it if the
	volume was required to
	be unloaded, the volume
	was to have been
	deleted.
1	Indicates that if the
	volume was requires to
	be unloaded, the
	unloaded volume was to
	have been retained.
1.	Bit 6 indicates that an
	unload requirement had
	been made.
1	Bit 7 set indicates that
	a load or label
	verification requirement
	had been made.

bytes 1-3

The address of the UCB that was used in all cases except when the device was a 2321 data cell drive. For a 2321, this address is that of the description in the UCB of the cell in the bin.

MFT Control Block Formatting

The formats described below are repeated for each requested task that is printed. A sample of the major system control blocks, as formatted from an MFT dump, is shown in Figure 23.

92 Programmer's Guide to Debugging (Release 21)

MFT TASK CONTROL BLOCK (TCB) FORMATTING: The task control block (TCB) contains information pertaining to the specific task identified in the heading lines at the top of the dump listing page. It is formatted as follows:

TCB hhhhhh The address of the task control block being displayed is given in this first display field.

- RBP hhhhhhhh The starting address of the request block (RB) that was currently associated with the task represented by this TCB.
- PIE hhhhhhhh The address of the first program interrupt element (PIE) enqueued by this TCB.
- DEB hhhhhhh The address of the beginning of the data extent block (DEB) queue that was associated with this task. Information on the contents of each DEB in the queue is given in a separate portion of this MFT task's dump listing.
- TIO hhhhhhh The starting address of the task input/output table (TIOT) that was constructed during device allocation for the task represented by this TCB. The contents of this table are displayed in a later portion of this task's display.
- CMP hhhhhhhh

This word contains ABEND indicators and user and system completion codes as follows:

byte 0

1	Bit 0 set indicates that a dump had been
	requested.
.1	Bit 1 is reserved for
	future use but is set
	for MVT compatibility.
1	Bit 2 set indicates that
	a portion of the problem
	program's main storage
	area was overlaid by a
	second load of ABEND.
	A first load overlay is
	indicated by the setting
	of bit 14 of the PK-FLG
	field.
X	Bit 3 is reserved for
	future use.
1	Bit 4 set indicates that
	a double ABEND occurred

in the dumped task.

MUDULE INDSADMP DATE 11/12/70 TIME 00.50 PAGE 0011 MFT DUMP LISTING PROCSTEP STEP1 JUB J085 STEP GO ***** CURRENT TASK ***** CMP 0000000 TRN COCCOCO TC3 009148 R8P 00009228 PIE 000000000 DEB 00071634 TIO JJ071728 JLB CCOOCCOC JST 00009148 PK-FLG 1000008 FLG 000001E3 LLS 000712F8 MSS 00009210 56007FD2 CCC0C0CC 000C011A RG 10-1 00071780 0C02A910 5002A826 9BC712B0 4002A896 0002C304 0007176C 000004C 00009148 000717F8 0CC71778 00000000 RG 2-9 30000000 THE 00009228 PIS E0019A88 NTC 00000CCC GTC 00000000 00009348 FSA 08071730 TC B 00000000 ECB OCCODOOC XTCB 0000000 LP/FL E300000C RES 0000000 LTC 00000000 IQE USR 00000000 CAR 0000000 RES C0000C00 JSCB 00021284 8A6C2000 STA COUCUUDO TC T ALTIVE RBS USE/EP C002A820 PS# FF150C80 9002AE7A 0 0000000 WT-LNK 00009148 PRA DZA 800 NM GO SZ/STAB 0C2C00C0 WT-LNK CCC24800 USE/EP 0002A87E PS# FF150193 8002A8AA Q CCCC5288 IR3 009228 NA SGKJ AKY SZ/STAB OCDE404C 00000040 0CC05148 C00717F8 RG 10-1 FA000048 00009228 00020304 00071760 00000000 JC02A510 13C0G0CC 4C0122EA 00071778 00000000 000717B0 0602A910 5002A826 KG 2-9 00071280 00009228 00009148 EXTSA 0000000 P/P BOUNDRIES HIER 0 00024800 TO 00071800 HIER 1 00000000 TO 00000000 LOAD LIST USE/EP 01071310 LRB 071300 NM DUM4YUL SZ CCOC88 SZ C000A8 USE/EP 01071380 LPRB 071390 NM RETURNS JUB PACK QUEUE NUTHING IN JOB PACK PCI DU 3FF4 CH END CC3FF4 AB END 0C3FF4 END CF EXT 0229C0 SID 003FF4 DEB 07163+ APPENDAGES C5C00005 00010BE0 11000000 PFX DCJCCOOO CCB 1F02A8BC UPRG C107144C PLST E3CCOCOO NDEB 1007150C ASYN F8CC0000 SPRG 00000000 TLB 34009148 AVT 04071610 FM-UCS START END TRKS 53001560 00020003 00020003 0001 ENC OF EXT 0138F0 SIG 013922 PCI 0136F8 CH END G13864 AB END C13922 DEB 07150C APPENDAGES PFX 00000000 05C0C007 000007E0 0F000000 NDEB 0000000 ASYN A8000000 SPRG 0000000 UPRG 01000000 PLST E3CCCC00 CCB 0F071778 TCB CCU09148 AVT 040136E4 START END TRKS FM--- UCa 0 CC 50 00 9 0011 5 8001520 00040003 PRCC STEP1 TIGT 071728 JOB JOB 5 STEP GO TTR-STC ST6-UCB **JFFSET** LN-STA DDNAME 00700000 800015EC PGM=+.DD 0018 14040100 DUMMY 007F0300 80001500 14640100 002C

Figure 23. Sample of MFT Control Block Format

IMDPRDMP Output Formatting: MFT -- TCB 93

CMP hhhhhhhh -- byte 0 -- (continued)1.. Bit 5 set indicates that a dump message (WTO) was to have been issued.1. Bit 6 set indicates that the dumped system's scheduler was to have printed an indicative dump. Bit 7 set indicates that1 an ABEND message, to be printed by the ABDUMP routine, was provided. bytes 1-3 The first 12 bits contain a system completion code. These codes and their meanings are explained in the publication IBM System/360 Operating System: Messages and Codes, GC28-6631 under the heading "System Completion Messages." A user completion code is contained in the last 12 bits. TRN hhhhhhh Contains flags as follows: byte 0 1.... Bit 0 set indicates that decimal simulator programs were being used on a System/360 model 91 machine. Bit 1 set indicates that .1.. checkpoints were not taken for this step. Bit 2 set indicates that the TCB being displayed was associated with either a graphics foreground job or the graphic job processor. ...1 Bit 3 set indicates that the TCB being displayed was associated with a 7094 emulator task that was being run on a System/360 model 85 machine. xxxx Bits 4 through 7 are reserved for future use. bytes 1-3 Reserved. MSS hhhhhhhh Main storage supervision as follows: byte 0 This byte is reserved for future use. bytes 1-3 This subfield displays one of two addresses. If the TCB being displayed represents a job step, this subfield contains the address of the boundary box. If this TCB represents a

subtask, this field displays the address of the gotten queue element (GQE). GQEs are preset only if the dumped system issued a GETMAIN macro instruction for the space.

PK-FLG hhhhhhhh The storage protection key and a series of flags associated with the task being displayed. This field is divided into several subfields. Thes e are: byte 0 xxxx The storage protection key associated with the task represented by this TCB. 0000 Always contain zeros. byte 1 Bit 0 set indicates that 1... an abnormal termination was in progress at the time the dump was taken. Bit 1 set indicates that .1.. a normal termination was in progress at the time the dump was taken. Bit 2 set indicates that ..1. ABEND was initiated by the resident abnormal termination routine. ...1 Bit 3 set indicates that recursion through ABEND was permitted. 1... Bit $\overline{4}$ set indicates that the graphics abnormal termination routine had been entered for the task represented by the TCB being displayed. Bit 5 set indicates that1.. the CLOSE routine was initiated by ABEND. Bit 6 set indicates that1. a portion of the problem program's main storage area was overlaid in order to process ABEND routines. (See also bit 2 of the CMP display field.) Bit 7 set indicates that1 the queueing of asynchronous exits for the task represented by the TCB being displayed, was prohibited. byte 2 Bit 0 set indicates that 1... ABEND was prohibited for this task. The setting of this bit has meaning only if the TCB being displayed represents a system task.

•XX• ••X•	Bits 1, 2 and 6 are reserved for future use.
1	Bit 3 set indicates that the task represented by
	the TCB being displayed
	had a forced completion imposed upon it. Other
	tasks in the dumped
	system could not have
	been performed until this task had been
	completed.
1	Bit 4 set indicates that the job step had a
	forced completion
	imposed upon it. Other tasks in the dumped
	system could not have
	been performed until
	this job step had been completed.
1	Bit 5 indicates that
	dump processing had been initiated in ABEND.
1	Bit 7 set indicates that
	the task represented by
	the TCB being displayed was a member of a time
	sliced group.
byte 3 xx.xx	Bits 0, 1, 3 and 7 are
	reserved for future use.
1	Bit 2 is an exit
	effector indicator. The setting of this bit
	indicates that at the
	time the dump was taken, system error routines
	were operating on this
1	task. Bit 4 set indicates that
**** 7***	floating point registers
	existed in the dumped
1	system. Bit 5 set indicates that
	at the time the dump was
	taken, job scheduler routines were
	processing.
1.	Bit 6 set indicates that
	at the time the dump was taken, an XCTL routine
	was changing the storage
	protection key in the PSW from zero to the one
	used by the problem
	program.
FLG hhhhhhhh	
This field dis	splays a further series
	certain priority This word is formatted as
follows:	

byte 0

Reserved for future use.

byte 1

XXXX XXX.

reserved for future use. Bit 7 is the primary1 non-dispatchability indicator. Setting of this bit indicates that one or more of the secondary non-dispatchability bits (bytes 1-3 of the DAR field) was set at the time the dump was taken. If this bit is set, the task represented by this TCB was considered to be non-dispatchable.

Bits 0 through 6 are

byte 2

This byte contains the number of resources for which the task represented by this TCB was enqueued.

byte 3 This byte displays the dispatching priority of the task represented by this TCB.

LLS hhhhhhhh

The address of the last request block (RB) that was created by the loading of a module that used the LOAD macro instruction.

JLB hhhhhhhh The address of the data control block (DCB) representing the JOBLIB associated with this task.

JST hhhhhhhh

Job step information. The contents of this field have meaning only when the dumped MFT system was operating with the subtasking option. If this was the case, this field shows the address of the first TCB for a job step.

- RG 0-7 and RG 8-15 The register save area of the TCB being displayed. This pointer is useful in locating the entry points of first routines and in tracing the save area chains.
- FSA hhhhhhhh This field displays two data elements and is formatted as follows:
 - byte 0 The TCB identification code.
 - byte 1-3 The address of the first problem program save area.

IMDPRDMP Output Formatting: MFT -- TCB 95

TCB hhhhhhhh		
The address of next lowest p	the TCB that had the ciority on the ready time the dump was taken.	
_	-	ord
TME hhhhhhhh The address of	the timer element.	
PIB hhhhhhh		
The PIB field	displays two items of	
information in	the following format:	
byte 0		LTO
This byte cont	ains flags that identify	
	attributes. These flags	
are: xx	Bits 0 and 1 indicate	
	the function of the	
	partition. The possible	
	functions are given	
00	below: System task partition.	
01	Reader partition.	
10	Writer partition.	IQI
11	Processing program	
••X• ••••	partition. Bit 2 gives the	
•••••	partition size. The	
	meanings of the possible	
<u>^</u>	settings are:	
0 1	Small partition. Large partition.	
	Bit 3 set indicates that	EC
	CPU timing was stopped	
	by FINCH until a	
	transient routine was loaded.	
xx	Bits 4 and 5 are	
	reserved for future use.	
1.	Bit 6 set indicates that	
	the partition associated with this task was a	
	writer partition. This	XT
	bit is used by ABEND,	
	transient writers and	
1.	resident writers. Bit 7 set indicates that	LP.
**** ****	at the time the system	
	was dumped, the	
	scheduler was in	
	control. Had this task's TIOT been written	
	to SYS1. SYSJOBQE, this	
	bit would not be set.	
huton 1.2		
bytes 1-3 The last port	ion of the PIB field	
shows the add:	ress of the partition	
	lock (PIB) that was	
associated with	th this task's partition.	
NTC hhhhhhhh		
The address of	f the previous TCB that	
existed on the	e originating task's	
queue of subta the TCR was +1	ask TCBs (sister). If he first on the queue,	
this field co	ntains zeros. The	
contents of t	he NTC field have meaning	

only if the dumped system was operating with the MFT subtasking option.

OTC hhhhhhhh

The OTC field is applicable only when the dumped system was operating under MFT subtasking option. If this was the case, this field displays the address of the TCB representing the originating task (mother).

LTC hhhhhhhh

The address of the last TCB that existed on the originating task's queue of subtask TCBs (daughter) at the time the dump was taken. If this TCB was the last on the queue, this field contains zeros. This field is applicable only if the dumped system was operating under the MFT subtasking option.

IQE hhhhhhhh

The address of the interruption queue element (IQE) that was used in scheduling the ETXR routine on the originating task. The contents of this field have no meaning unless the dumped system was operating under the MFT subtasking option.

ECB hhhhhhhh

If the dumped system was operating under the MFT subtasking option, this field displays the address of the event control block (ECB) that would have been posted by the supervisor's task termination routines had either normal or abnormal task termination been allowed to occur.

- TCB hhhhhhhh The XTCB field in this TCB display is reserved for future use.
- P/FL hh hhhhhh Priority and dump information on tasks that were operating under the subtasking option of MFT. The LP/FL field displays its data as follows:
 - byte 0 The limit priority of the task represented by the TCB being displayed.

byte 1 Dump information flags.

> xxxx x... Bits 0 through 4 are reserved for future use. 1... Bit 5 set indicates that the task represented by the TCB being displayed was the top task in the tree of abnormally terminating tasks.

	Bit 6 set indicates that an abnormal termination dump had been completed. Bit 7 set indicates that		e bytes contain the ne dumped system's timing e (TCT).
	the task represented by this TCB was enqueued on a dump data set. tains more dump	the dumped sy	available to the user of ystem. It contains any placed in it by the user.
information f of these bits	lag bits. The meanings are:		$\mathbf{of}^{'}$ this field were used
1	Bit 0 set indicates that at the time the system was dumped, an OPEN was in process for the dump data set.	(DAR). Certa this word wer the dispatcha task. The DA	e assessment routine in subfields displayed in te also used to control ability of the dumped AR field is divided into
•xxx x••x	Bits 1 through 4 and bit 7 are reserved for future use.	the following byte 0	
1	Bit 5 set indicates that the dump data set was open for the job step.		e of the DAR field flags. The flags are as
···· ··X.	Bit 6 indicates the type of dump data set. The possible setting are:	1	Bit 0 set indicates that primary DAR recursion
	SYSUDUMP data set. SYSABEND data set.		occurred in the dumped system. The damage assessment routine
shows abnorma	e of the LP/FL field l termination flags as	.1	failed while writing a main storage image dump. Bit 1 set indicates that
follows: xxx. x.xx	Bits 0, 1, 2, 4, 6 and 7		secondary DAR recursion occurred in the dumped system. The damage
1	are reserved for future use. Bit 3 set indicates that		assessment routine failed while attempting to reinstate a failing
	a valid message recursion occurred in ABEND.	1	partition. Bit 2 set indicates that only the dump capability
•••• •1••	Bit 5 set indicates that no abnormal termination dumps could be provided	····x ····	of the damage assessment routine was requested. Bit 3 is reserved for
	within the job step associated with the TCB being displayed.	1	future use. Bit 4 set indicates that the system error task is failing. The DAR dump
RES hhhhhhhh This field is	reserved for future use.		should not request any error recovery procedure (ERP) processing.
address of th	routine flags and the STAE control block that at the time the dump was	···· · ××. ···· · · · 1	Bits 5 and 6 are reserved for future use. Bit 7 set indicates that an SVC dump is executing for this task.
system's timi The TCT field following two	ertaining to the dumped ng control table (TCT). is divided into the subfields:	field are use secondary nor bits. If any subfield were non-dispatcha	agh 3 of the DAR display ed to show the settings of a-dispatchability flags y of the flags in this e set, the primary ability flag (the last bit
byte 1-3 If the system	reserved for future use. management facilities esented in the dumped	set and the t TCB will have The bit setti	ield) will also have been task represented by this the been non-dispatchable. Ings that may appear in their meanings are:

XX	Bits 0 and 1 were set by	byte 3	
	the damage assessment	1	The setting of this
	routines. Their		first bit has meaning
	meanings are:		only if the dumped
1	Bit 0 set indicates that		system was operating
	the task represented by		with the MFT subtasking
	the TCB being displayed		option. If this bit is
	was flagged temporarily		set, the indication is
	non-dispatchable. Bit 1 set indicates that		that the task represented by the TCB
.1	the task represented by		being displayed was
	this TCB was deemed		terminated.
	permanently	.1	Bit 1 set indicates that
	non-dispatchable.		had the dumped MFT
• • XX • • • •	Bits 2 and 3 are		system, operating with
	recovery management		the subtasking option,
	support and system error		been allowed to continue
	recovery flags. Their		processing without
	meanings are:		intervention by the dump
1	Bit 2 set indicates that		program, the task
	the task represented by this TCB was flagged		represented by this TCB would have been
	temporarily		terminated by ABEND.
	non-dispatchable.	••xx xxxx	Bits 2 through 7 are
1	Bit 3 set indicates that		reserved for future use.
	the task represented by	RES hhhhhhhh	
	the TCB being displayed	Reserved for	future use.
	was deemed permanently		
	non-dispatchable.	JSCB hhhhhhhh	
•••• X•••	Bit 4 is reserved for		address of the job step
	future use.	control block	•
1	Bit 5 set indicates that	RES hhhhhhhh	
	this task is temporarily non-dispatchable. Time	Reserved for	future use
	services have been	Reserved for	incure ase.
	requested and the	IOBRC hhhhhhhh	
	time-of-day clock is	Contains the	address of the IOB
	still inoperative.	restore chain	for I/O quiesced by
•••• ••XX	Bits 6 and 7 are	end-of-task.	
	reserved for future use.		
h			BLOCK (RB) FORMATTING:
byte 2	Dit 0 indiant on that at		s) were used by the
1	Bit 0 indicates that at the time the dumped		pervisor to maintain
	system was active,	information concer	ning a task. RBs he task identified in the
	ABDUMP was processing.		he top of the dump page
	The setting of this flag		ng TCB display, are
	bit has meaning only if		ion of the dump listing
	the dumped system was	labeled "ACTIVE RB	
	operating with the		the task is formatted as
	subtasking option of	shown below:	
	MFT. Bit 1 is reserved for		
•X•• ••••	future use.	PRB LPRB	
	Bit 2 set indicates that	SVRB hhhhhh	
	this task is	SIRB	
	non-dispatchable. An	IRB	
	SVC dump is executing		ay is preceded by a field
	for another task.		s the type and starting
x xxx.	Bits 3 through 6 are		e RB being displayed.
-	reserved for future use.		s of RBs that may be
•••• •••1	Bit 7 set indicates that	displayed und	er an MFT task are:
	at the time the system	PRB	
	was dumped, the dump data set was in the		request block
	process of being opened.	Program	TOTACOL NICOV
		LPRB	
		loaded p	rogram request block

SVRB

supervisor request block (SVRBs may be divided into two categories; type 2 for resident routines and type 3 or 4 for transient routines).

SIRB

system interrupt request block

IRB

interruption request block

The type acronym for each RB is displayed in the first portion of the field. The starting address of the indicated request block appears in the last portion of the field. The contents of certain fields in the body of the formatted RB display are dependent upon the type of RB being displayed. Variations in display field usage are noted in the descriptions of the fields in which they occur.

NM ccccccc

The variations associated with the usage of this field are:

- PRBs and LPRBs use this field to display the name of the program they represented.
- SVRBs display the SVRB type in this field.
- SIRBS use this field to present the eight-character name of the error routine that was occupying the supervisor transient area at the time the dump was taken.
- IRBs display meaningful information in this field only if the timer was being used. If this was the case, the first character in this field represents the setting of the timer flags. The remainder of the NM field is meaningless.

SZ/STAB hhhhhhh

This field displays two data elements; RB size information and STAB flag bit settings. This field is subfielded as follows:

bytes 0-1 The number of contiguous doublewords that were occupied by the request block, the associated program (if applicable), and associated supervisor work areas. If a program extent list was present, the program size is not included in this figure. byte 2 STAB flag bit settings. The meaning of these flags are depends upon the type of request block being displayed. These flags are presented, by RB type, below:

- PRB The following bit settings are applicable to program request block displays:
 - indicates that the 0000 program represented by this PRB was not loaded by a LOAD macro instruction; nor did it have minor entries identified by an IDENTIFY macro instruction. 0001 indicate that the program represented by this PRB was not loaded by a LOAD macro instruction but did have minor entries identified by an IDENTIFY macro instruction. xx.. Bit 4 and 5 have no meaning in PRB displays. ..1. indicates that the program represented by this PRB was hierarchy block loaded and that a program extent list existed. ...1 indicates that the program module represented by this PRB was refreshable.
- LPRB

Loaded program request blocks being displayed may have the following bit settings in this byte:

0010 indicates that the program represented by this LPRB was not loaded by a LOAD macro instruction; nor did it have minor entries identified by an IDENTIFY macro instruction. 0011 indicates that the program represented by this LPRB was not loaded by a LOAD macro instruction but did have minor entries identified by an IDENTIFY macro instruction.

- 1110 indicates that this LPRB describes a minor entry identified by an IDENTIFY macro instruction.
- Bits 4 and 5 have no XX meaning in LPRB displays. indicates that the ..1. program represented by this LPRB was hierarchy block loaded and that a program extent list existed. ...1 indicates that the
 - program module represented by this LPRB was refreshable.

SVRB

- Supervisor request blocks display the following bit settings in this subfield:
- 1100 indicates that the program represented by this SVRB is a type 2 SVC routine that had not been loaded at the time the dump was taken. 1101 indicates that the program represented by this SVRB is a type 3 or SVC routine that had been loaded. indicates that the 1... type 3 or 4 SVC routine was resident. indicates that while .1.. the dumped system was active, a checkpoint could have been taken in a user exit from the SVC routine represented by this SVRB. ..xx bits 6 and 7 have no meaning in SVRB
- SIRB The flag bit setting applicable to supervisor interrupt request block displays is as follows:

displays.

1000 indicates that the RB being displayed is a supervisor interrupt request block (SIRB).
.... xxxx bits 4 through 7 have no meaning in SIRB displays.

IRB

- Interrupt request block displays use these flag bits in the following manner.
 - 0100 indicates that the RB
 being displayed is an
 interrupt request
 block (IRB).
 xxxx bits 4 through 7 have
 no meaning in IRB
 displays.
- byte 3 The last byte of the SZ/STAB field displays more status and attribute flags. The possible settings for this subfield and their meanings are:
 - Bit 0 set indicates that 1... the WT-LNK field in this RB display contains, in its last three bytes, the address of the TCB to which this request block is linked. Bit 1 set indicates that .1.. at the time the dumped system was active, the program associated with the RB being displayed was active. ..1. Bit 2 set indicates that had the dumped system been allowed to continue processing without intervention by the dump program, general registers 2 through 14 would have been restored from this RB's general register save area, displayed on the following two lines. The setting of this bit is valid only for IRB, SIRB and SVRB displays. Bit 3 set indicates that1 the program module represented by this request block was reenterable or reusable. xx.. Bits 4 and 5 are used only in IRB or LPRB displays. The settings of these bits and their meanings are: 00... This setting indicates that the IRB being displayed had no interrupt queue elements (IQEs) associated with

it.

.... 01.. This setting indicates that the IRB being displayed had associated with it interrupt queue elements that were

.... 10...

.... ..1.

request queue elements (RQEs). This setting indicates that the request block being displayed is a dummy LPRB, in a partition that represents a program in the reenterable load module area. The LPRB for the program is in the reenterable load module area. 11.. This setting indicates that the IRB being

displayed had interrupt queue elements associated with it that were not request queue elements (RQEs).

- Bit 6 set indicates that when the dumped system was active, request block storage was to have been freed when the program returned.
- Bit 3 indicates waitX request conditions. The meanings of the two possible settings for this bit are:
- Bit 7 not set indicates0 that the request had to wait for a single event or for all of a number of events.
- Bit 7 set indicates that1 the request had to wait for a number of events. This number of events was less than the total number of events that were waiting.

USE/EP hhhhhhhh The USE/EP field, as indicated by the field identifier, displays two data elements. These are shown in the following format:

byte 0

The first byte of this field contains the use count that was applied to the program module represented by the request block being displayed. This use count was calculated by subtracting the number of invocations of the DELETE macro instruction from the number of times the LOAD macro instruction was used.

byte 1-3

The second portion of the USE/EP field displays the address of the entry point of the module represented by this request block.

PSW hhhhhhhh hhhhhhh

The two words of the PSW field display to the user the dumped system's old program status word. If the dumped system had been allowed to continue processing without interruption by the dump program, operation would have resumed on this PSW.

Q hhhhhhhh

The information displayed in this field depends upon type of RB being displayed. The contents of this display field are described below, by **RB** type:

- PRBs and LPRBs use this field to display the address of an LPRB describing an entry that was identified via the IDENTIFY macro instruction.
- SVRBs representing type 3 or 4 SVCs use this field to indicate the size of the program they represent in bytes.
- SIRBs and IRBs display in this field the address of a 12- or 16-byte request element.

WT-LNK hhhhhhhh

This field displays information pertaining to wait counts and request block linkages. The field is divided into the following two subfields:

byte 0 The number of requests that were pending at the time the dump was taken (wait count).

byte 1-3 The address of the next request block on the RB queue. If the RB being displayed was the last request block on the queue, this field shows the address of the task control block (TCB) that enqueued this RB.

RG 0-7 and RG 8-15 The sixteen-word register save area appears only after IRB, SIRB or SVRB displays. These two lines display the contents of general registers 0 through 15 as they were stored in the request block.

MFT PROBLEM PROGRAM BOUNDARIES INFORMATION: Each task operating under the MFT option of the operating system was assigned a main storage partition in which to operate. If

the system configuration included 2361 Large Core Storage, partitions may have included area from both hierarchy 0 (main storage) and hierarchy 1 (low speed main storage). If 2361 Large Core Storage was not available or was not used, hierarchy 1 pointers were set to zero. Each MFT task displays in its dump listing the limits of the partition in which it operated. This display is presented under the heading "P/P BOUNDARIES" (problem program boundaries) in the following format:

- HIER 0 hhhhhhhh
 - The starting address of the problem program's hierarchy 0 partition.
- TO hhhhhhhh

The ending address of the problem program's hierarchy 0 main storage partition.

- HIER 1 hhhhhhhh
 - The starting address of the problem program's hierarchy 1 partition. If this field contains zeros, the indication is that 2361 Large Core Storage was either not available or not utilized by this task.
- TO hhhhhhhh

This last field indicates the high limit of the problem program's hierarchy 1 partition if one was used. If this field contains zeros, either 2361 Large Core Storage was not available or it was not used by this task.

MFT LOAD LIST FORMATTING: A load list was maintained by the dumped system's supervisor in order to keep track of the load modules that were in main storage and the area of main storage each occupied. A load list created by an MFT supervisor is composed of loaded request blocks (LRBs) and loaded program request blocks (LPRBs). A formatted listing of the dumped MFT system's load list appears as follows:

LRB

LPRB hhhhhhh The type of request block being displayed and its starting address.

NM ccccccc

The eight-character name of the program module represented by the request block being displayed.

SZ hhhhhh

The number of contiguous double words that were occupied by the request block, the associated program (if applicable) and associated supervisor work areas. If a program extent list was present, the program size is not included. USE/EP hhhhhhhh

Use count and entry point address as follows:

byte 0

The use count that was applied to the program module represented by the request block being displayed. This use count was calculated by subtracting the number of times the DELETE macro instruction was issued from the number of times the LOAD macro instruction was used.

byte 1-3

The address of the entry point of the program module named in the NM field of this RB display line.

MFT JOB PACK QUEUE FORMATTING: A job pack area queue was maintained by the dumped system's supervisor for each job step that used a program not in the resident reenterable load module area. A job pack queue created by an MFT supervisor consists of loaded request blocks (LRBs), loaded program request blocks (LPRBs) and FINCH request blocks (FRBs). A formatted job pack area queue display appears as follows:

LRB

LPRB hhhhhh

FRB

The type of request block being displayed and its starting address.

NM cccccccc The eight-character name of the module represented by the request block being displayed.

SZ hhhhhh

The number of contiguous doublewords that were occupied by the request block, the associated program (if applicable) and associated supervisor work areas. If a program extent list was present, the program size is not included.

USE/EP hhhhhhhh

XRWTL

The usage of this display field is dependent upon the type of request block being displayed:

USE/EP is used for LRBs and LPRBs and displays the use count and entry point address as follows:

byte 0 The use count that was applied to the program module represented by the request block being displayed. This use count was calculated by subtracting the number of times the DELETE macro instruction was issued from the number of times the LOAD macro instruction was used.

bytes 1-3

The address of the entry point of the program module named in the NM field of this display line.

XRWTL

is used for FRBs and shows the starting address of the wait list element.

XRREQ hhhhhhhh

This field appears only in FRB displays, and shows the address of the TCB representing the task on whose behalf this FRB was constructed.

XRTLPRB hhhhhhhh

This field appears only in FRB displays and shows the starting address of the area of main storage that was acquired by the FETCH routine for the module identified by the NM field of this line.

MFT DATA EXTENT BLOCK (DEB) FORMATTING: Data extent blocks (DEBs), describing a data set's external storage requirements, were queued to those task control blocks (TCBs) that represented tasks requiring auxiliary storage input/output processing. External storage information, taken from each DEB, is formatted as shown below:

DEB hhhhhh

The starting address of the basic section of the DEB being displayed.

APPENDAGES

The word "appendages" informs the user that the five named fields on this line contain information taken from the appendage vector table preceding the DEB being displayed. The named fields appearing on the rest of this line are:

END OF EXT hhhhhh The entry point of the end-of-extent appendage routine.

SIO hhhhhh The entry point of the start I/O appendage routine.

PCI hhhhhh The entry point of the program-controlled-interruption appendage routime.

CH END hhhhhh The entry point of the channel-end appendage routine. AB END hhhhhh The entry point of the abnormal-end appendage routine.

PFX hhhhhhh hhhhhhh hhhhhhh The second line of a DEB display contains information taken from the prefix section of the DEB being displayed. The area is subdivided as follows:

byte 0 The first byte of the prefix area contains the contents of the I/O support work area. This area is used only by DEBs dealing with direct access storage devices.

bytes 1-7

The next seven bytes of the DEB prefix section are used by DEBs associated with direct access storage device functions. This subfield displays the data set control block's (DSCB) address used by I/O support. The address is expressed in the following format:

bytes 1 and 2 the bin (cell) number. bytes 3 and 4 the cylinder address. bytes 5 and 6 the track address. byte 7 the record number.

bytes 8-11 The third word of the PFX field contains the data control block (DCB) modification mask that was used by I/O support.

byte 12 The length of the DEB in double words.

bytes 13-15 The remainder of the DEB prefix section is reserved for future use.

TCB hhhhhhh This field marks the beginning of the basic section of the data extent block. The TCB field is divided into two subfields as follows:

- byte 0 The number of subroutines for which a LOAD macro instruction was issued during the execution of the OPEN executor routines.
- bytes 1-3 The starting address of the task control block to which this DEB was enqueued.

NDEB hhhhhhhh

byte 0 The overall length of a data extent block includes the length of a variable length access method dependent section. The first byte of the NDEB field expresses the length of the access method dependent section in bytes. If the access method was BDAM, this indicator is expressed as a number of full words.

bytes 1-3

The last portion of the NDEB field displays the starting address of the basic section of the next DEB on the task's queue. If this DEB was the last on the queue, the content of this field is the starting address of the TCB that enqueued this DEB.

ASYN hhhhhhhh

This field contains data set status flags and the address of the associated IRB:

byte 0

The first byte of the ASYN field contains data set status flags. These flags have the following meanings:

xx	Bits 0 and 1 indicate the data set's disposition. The
	possible settings are:
01	This setting indicates
	that the disposition was
	OLD.
10	This setting indicates
	that the disposition of
	the data set was MOD
11	(modify). This setting indicates
17	that the disposition was
	NEW.
	Bit 2 set indicates that
	an end-of-volume (EOV)
	or end-of-file (EOF)
	condition had been
	encountered.
1	The setting of bit 3 has
	one of two meanings
	depending upon the
	external storage medium. For disk, this indicator
	reflects a release of
	unused external storage.
	For tape, this indicator
	means that an emulator
	tape with second
	generation format was
	being used.
1	Bit 4 set is a data
	control block (DCB)
1	modification indicator. Bit 5 set has two
	meanings, depending upon
	the auxiliary storage
	recording medium. For
	disk, the setting of bit
	5 indicates that a split
	cylinder was

encountered. For tape, this flag indicates that an emulator tape with possible mixed parity records was used. Bit 6 set indicates the1. use of nonstandard labels. Bit 7 set indicates that1 reduced error recovery procedures were used on magnetic tapes containing the data set represented by this DEB.

bytes 1-3

The last portion of the ASYN field shows the starting address of the IRB that was associated with asynchronous appendage exit scheduling.

SPRG hhhhhhhh

This field contains information on I/O processing methods and the system PURGE routine.

byte 0

The first byte of this field contains flags that indicate the method of input/output processing and the disposition of the data set that was to have been performed when an end-ofvolume condition occurred. These flag settings are:

1 .0 xx	Bit 0 was set by ABEND. The setting of this bit indicates that the data set associated with this DEB was a SYSABEND or SYSUDUMP data set. Bit 1 is always zero. Bit 2 and 3 show the end-of-volume disposition procedure. The values for this flag
01 11 xxxx	are: REREAD LEAVE The last half of this byte contains flags that indicate the type of input/output processing that was performed on the data set represented by this DEB. The values
0001	for this flag are: INPUT OUTPUT INOUT OUTIN RDBACK UPDAT

byte 1

The quiesce count. The byte is associated with the system PURGE routines (SVC 16), and indicates the

number of auxiliary storage devices that were executing the user's channel programs. bytes 2-3 Reserved for future use. UPRG hhhhhhhh The UPRG field contains extent information and data used by the user's purge routines. This field is divided into the following two subfields: byte 0 The number of extents that were specified in the DSCBs associated with this DEB. bytes 1-3 The address of the first input/output block (IOB) in the user's purge chain. PLST hhhhhhhh Task priority and supervisor purge information are contained in this field. This field is formatted as follows: byte 0 The priority of the task under which this DEB was enqueued. bytes 1-3 The starting address of a parameter list that was used to locate the purge event control block (ECB) for a supervisor purge request. DCB hhhhhhh The DCB field contains three data elements. These are displayed in the format given below: byte 0 XXXX The storage protection key that was associated with the task under which this DEB was enqueued. A hexadecimal "F" in 1111 bits 4 through 7 of this field identify this control block as a data extent block (DEB). bytes 1-3 The starting address of the data control block (DCB) that was associated with this DEB. AVT hhhhhhhh The AVT field displays two DEB data elements and is subfielded as follows: byte 0 The DEB extent scale that is used to determine the size of the device

dependent section of this DEB. For direct access devices, a 4 is displayed in this subfield. For a nondirect access device or a communication device, a 2 is displayed.

bytes 1-3 In most cases, the last portion of the AVT field shows the starting address of the appendage vector table preceding this DEB. This table of appendage routine addresses appears on the first line of this DEB's display.

OP-UCB hhhhhhhh

The contents of this field have meaning only when the DEB being displayed describes a data set that was assigned to a unit record or magnetic tape device. This information is formatted from the device dependent section of the DEB. The OP-UCB field is subfielded as follows:

byte 0 This first subfield is applicable only to data sets assigned to magnetic tape devices, and shows the SET MODE operation code. For a data set that was assigned to a unit record device, this subfield is reserved.

bytes 1-3 The starting address of the unit control block (UCB) associated with the data set described by the DEB being displayed.

The following four fields are present only for data sets assigned to the IBM 3525 Card Punch for multi-function. The information is formatted as shown below:

UCB hhhhhhh byte 0 The device modifier field (not used for the 3525). bytes 1-3 The starting address of the unit control block (UCB) associated with the data set described by the DEB being displayed. RDRDCB hhhhhhh The starting address of the data

The starting address of the data control block (DCB) for the read associated data set.

PCHDCB hhhhhhhh The starting address of the data control block (DCB) for the punch associated data set.

WTRDCB hhhhhhhh The starting address of the data control block (DCB) for the print associated data set.	STEP c T t s s
The final portion of a DEB display shows information pertaining to a data set that was assigned to a direct access device. This information, taken from the DEB's device dependent section, is arranged in columnar format with a line for each extent. The information is formatted as shown below:	PROC C I W C n S
FM-UCB hhhhhhhh The first column displays two data elements and is formatted as follows: byte 0	Each da task i that i entry column each c
The device modifier showing the file mask.	OFFSET T
bytes 1-3 The starting address of the unit control block (UCB) that was associated with the data extent.	LN-STA
START hhhhhhhh The address of the beginning of the direct access device extent. The first four characters represent the cylinder address and the last four characters represent the track address.	byt T e b byt S b
END hhhhhhhh The address of the end of the data extent. Cylinder and track references are formatted as in the extent beginning address, described above.	o
TRKS hhhh The number of direct access tracks bounded by the starting and ending addresses shown in the previous two columns.	
MFT TASK INPUT/OUTPUT TABLE (TIOT) FORMATTING: A task input/output table (TIOT) was constructed for each task in the dumped system by MFT job management routines. This table contained primary pointers to control blocks used by I/O support routines. As the functions of several TIOT fields were dependent upon the state of associated external storage devices, multiple definitions may apply. The TIOT that was constructed in the dumped MFT system is formatted as shown.	
TIOT hhhhhh The starting address of the task input/output table being displayed.	
JOB cccccccc The eight-character name of the job for which this TIOT was constructed.	

---ccccccc The eight-character name specified in the label field of the EXEC JCL statement associated with this job step. ccccccc If the job step for which this TIOT vas constructed was invoked from a cataloged procedure, the procedure name, as contained in the EXEC JCL statement, is displayed in this field. lata set associated with the indicated is represented by a separate DD entry is included in the TIOT. Each TIOT is displayed on a separate line in nar format. The use and meaning of column is given below: r hhhh The offset of this DD entry from the beginning of the TIOT in hexadecimal. A hhhhhhhh te O The total length (including all device entries) in bytes of the DD entry being displayed on this line. te 1 Status byte A, one of three status oytes in a TIOT entry. The meanings of the status byte settings are: Bits 0 and 5 indicate x... .x.. the tape label processing that was to have been performed. The meanings of the settings are: 0.... Nonlabeled tape or an indication to bypass label processing. Standard labels or 0....1... standard user labels. 1...... Nonstandard labels. The setting of status .1.. bit 1 has two meanings, depending upon the processing phase that had been reached at the time the system was dumped. During allocation processing, the setting of this bit indicates that this entry represents a split

> cylinder primary space allocation DD. If the dump was taken during step termination

processing, the setting of this bit indicates that no unallocation of space was necessary.

1	The setting of status bit 2 works under the same philosophy as status bit 1. During allocation processing, the setting of this bit indicates that this entry represents a split cylinder secondary space allocation DD. If the dump was taken during step termination processing, the indication was one of rewinding with no	STB-UCB hhl The la contai one-wa device device inform byte 0 Status the fo 1
1	unload. Bit 3 set indicates that this DD entry represents a JOBLIB. Bit 4 set indicates that direct access device	.1.
1.	space management was deemed necessary. The setting of bit 6 specifies that the tape volume was to have been rewound and unloaded. The setting of bit 7 specifies that the tape	••1
byte 2	volume was to have been rewound.	•••
The third byt meaning only phase. This devices that data set rep	e of this column has during the allocation displays the number of were requested by the cesented by the TIOT entry	•••
displays a T at two points this task. I process, this to the approp affinity, vo suballocate T	e of the LN-STA field OT field that had meaning s during the processing of During the allocation s field contained a link priate prime split, unit lume affinity or TIOT entry. After CLOSE this byte was used as	
1	The setting of bit 0 indicates that the data set represented by this DD entry was a SYSOUT data set that contained data. Bits 1 through 7 are	
	reserved for future use.	• • •
with the TIO TTR-STC hhhhhhhh The first the display the r (TTR) of the	aracter DD name associated I entry being displayed. The entry being displayed. The entry being displayed. The entry being displayed. The entry being displayed by the entry by the entry by th	bytes 1 The au in al was a 2321, descr the b

B hhhhhhhh	
	mn in a TIOT display
ontains info	rmation taken from the
	ce entries that are
ppended to ea	ach TIOT entry. One TIOT
levice entry e	exists for each allocated display field shows this
	n the following format:
	in one routowing rounder
e 0	
	. The status bits have
he following	
1	Bit 0 set indicates that the data set associated
	with this line of the
	TIOT display was present
	on the device
	represented by this TIOT
4	device entry.
.1	Bit 1 set indicates that the data set associated
	with this line of the
	TIOT display would have
	used the device
	represented by this TIOT
1	device entry.
1	Bit 2 set indicates that the device represented
	by this device entry
	violated separation.
1	Bit 3 set indicates that
	a volume serial number
	was present.
1	Bit 4 set indicates that
	a setup message was required.
x	Bit 5 indicates the
	device disposition that
	would have taken place
	had the dumped system
	been allowed to continue processing this task.
	The settings for this
	bit are:
0	Indicates that if the
	volume was required to
	be unloaded, the volume
	was to have been deleted.
1	Indicates that if the
	volume was required to
	be unloaded, the
	unloaded volume was to have been retained.
	Bit 6 indicates that an
	unload requirement had
	been made.
••••	Bit 7 set indicates that
	a load or label
	verification requirement
tes 1-3	had been made.
	f the UCB that was used
	except when the device

The address of the UCB that was used in all cases except when the device was a 2321 data cell drive. For a 2321, this address is that of the description in the UCB of the cell in the bin.

TSO System Block Formatting

The TSO control blocks are divided into two groups: system and user. The control blocks are discussed in the order in which they appear when both groups are requested. Some control blocks are formatted and printed when either group is requested.

An example of a TSO system and user dump listing is shown in Figure 24.

TIME SHARING COMMUNICATIONS VECTOR TABLE (TSCVT) FORMATTING: The time sharing communications vector table is a secondary CVT to meet the time sharing requirements. The time sharing CVT resides in the time sharing region; therefore, it exists only while the time sharing region is active. When time sharing does not exist in the system, the MVT CVT pointer to the TSCVT is zero.

TSCVT hhhhhh The address of this time sharing communications vector table.

TJB hhhhhhhh

The address of the time-sharing job block (TJB) table. This table contains all of the TJBs allowed TSO users. The first TJB is for the terminal job identification (TJID) equal to zero.

RCB hhhhhhhh

The address of the region control block (RCB) table. It is an indexed table containing one RCB for each possible time sharing region; therefore, the table contains the maximum number of RCBs that may be used by time sharing. The first RCB is for region one.

RPT hhhhhhhh

The address of the reference point table (RPT). It is used by the terminal input output coordinator (TIOC).

FLG hhhh

These flags indicate functions requested from the time sharing control task (TSC).

byte 0

1... TSCSWPND: Bit 0 set indicates that a swap has ended.

- .1.. ... TSCSWPBG: Bit 1 set indicates that a swap should be started.
- ..1. TSCLOGON: Bit 2 set indicates that a logon is required.
- ...1 TSCDISC: Bit 3 set indicates that a disconnect is required.
- xxxx Bits 4 through 7 are reserved for future use.
- byte 1

Reserved for future use.

FL1 hhhh

These flags indicate atypical functions required by the time sharing control task (TSC).

- byte 0
 - 1... TSCSSTOP: Bit 0 indicates that a system stop has been requested and the time sharing system is in the process of stopping.
 - .1.. ... TSCRSTOP: Bit 1 indicates that a region stop has been requested.
 - ..1. ... TSCASTOP: Bit 2 is the ABEND-STOP flag. When set, it indicates to the time sharing control task (TSC) that time sharing should be stopped. This flag is set by (1) the TSO/RMS interface return when a machine check occurs in TCAM or (2) the TCAM STAE exit when TCAM abnormally terminates.
 - ...x xxxx Bits 3 through 7 are reserved for future use.
- byte 1 Reserved for future use.
- SDC hhhhhhhh The address of the first data control block (DCB) for swap data sets.
- CUS hhhh A count of the current TSO users logged onto the system. For additional users to be logged onto the system, this number must be less than the value in LUS.

NMBR 01 PKEY E0 UMSMN 04 FLG 40 FLG2 20 FBQE 0 UTTMQ 0000 CUSE 0004 EXTNT 000A7F63 UMSM 000DDFA8 SDCB 000DE120 PQE 0001AC20					MODULE INDSADMP 1	DATE 11/12/70 TIME 00	12 PAGE 0006
NMBR 01 PKEY EO UMSMN 04 FLG 40 FLG2 20 FBQE 00 UTTMQ 0000 CUSE 0004 EXTNT 000A7F69 UMSM 000DFA8 SDCB 000DE120 PQE 0001AC20 CONID 00 RESV 000000 PRG2 000A7F1C QPL 000A7F10 STECB 00000000 RCDVR 0B00FF0 UMSM 0DDFA8 ADDR-LN C458020 ADDR-LN 0000000 ADDR-LN 00000000 ADDR-LN 0000000 ADDR-LN 00000000 ADDR-LN 00000000 ADDR-	TSCVT ODDA90	CUS 0C04 SAV 0C0DDB20 I02 000D3B50 SLF 0C0CDF10 SVQ CCCC0000	LUS 000A ECB 000DDB14 I03 000D3E46 TSC 0001ACD0 ABN 000D1C20	NTJ 000A SIA 000DDCDC D02 000D28C8 SPL 0001B4E8	SZU 0030 C' ICB 000DDC34 I LCQ 00000000 TI RSZ 0028 R	TR 0001 MUS 01 000D38C4 TQE 000 RB 00000000 LPA 000 SV 0000 SVT 000	000A 14674 000000 000000
SWAP DCB DDC0D0 CA5800 STDRAGE KEY 0 0A5800 0.0000000 D00A58C8 0000000 D00A58C2 000A5820 *	RCB ODDFB8	NMBR 01 UTTMQ 0000 PRG 0E000000	PKEY E0 CUSE 0004 PRG1 000A79D0	UMSMN 04 EXTNT 000A7F6	FLG 40 UMSM 000DDFA8	FLG2 20 F SDCB 000DE120 F	BQE 01
CA580J STURAGE KEY 0 0A5800 00000000 000A58C8 00000000 000A58C8 0000000 000A58C3 000C0C0 ************************************	UMSM ODDFA8	ADDR-LN CÅ580060	ADDR-LN OCE	880020 ADDR-1	DCA 0000000 ADD	R-LN 000000C0	
0A5800 0 0000000 000A58C8 00CA7260 00000000 00000000 ************************************	SWAP DCB 000	0000					
	0A5800 0 0000 0A5820 0 0000 0A584C 0 0000 0A5880 0 0000 0A5880 0 TO N 0A6620 0 0012 0A6640 0 0000 0A6680 0 000A	00000 000458C8 000A7 00000 000A7000 00000 00000 00001468 00000 IEXT LINE ADDRESS SA 00000 CCCC0000 FFC40 0000 CCC0000 FFC40 0000 0000028C 000A6 17788 00010100 00009	000 0000000 000 0000000 000 0000000 ME AS ABOVE 000 0000CAF8 006 000A7700 DF4 0000000 000 0000000	0001A7B9 000280 00000000 000000 00000000 000000 000000	000 00045800 000000 000 0000000 000000 000 0000000 000000 700 0000000 000000 001 0000000 000000 000 0000000 000000 001 0000000 000000 000 0000000 0000000 000 0000000 0000000	************************************	***************************************
	046000 0 0004	0000 0000000 0000	000 0000000	00000000000000	000000 00000000 000	00	••••••••••

MODULE IMDSADNP DATE 11/12/70 TIME 00.12 PAGE 0007 TSO USER CONTROL BLOCKS ***** **** USER KGNO1 TJID=0001 ****** **** TJB 0DDD18 TSB 000090F4 ATTN STAX STAT2 00 EXTNT 00047F68 00 01 STAT 00 RSTOR RCB 000DCFB8 UMSM 000DDF08 SDCB 000DE120 UTTMQ 0002 48 UMSMN 04 USER KGN01 IPPB 00000000 NEWID 00 FLUSL 00 TJID 0001 MONI 00 RSV 000000 UMSM ODDF08 ADDR-LN 0A580C38 ADDR-LN 0A980058 ADDR-LN OCBOOO28 ADDR-LN 00000000 TSB 0D9DF4 STAT 81 TJB ODCD18 FLG1 00 WT SB 000000 LNSZ 78 OTBFP 000000 CBFP 000000 IBFP NOBF BPKFL 00 **ITBFP** 000000 NITR 01 ODAOFO 00 QCB OE1CCO CLEAR 00 ECB 00000000 TJID 0001 STCC 0000 ATNLC 0016 ATNTC 0000 AUTOS 0000000 LNNO 00 BLNK 00 ASRCE 0000 ATNCC 0003 AUTOI 0000000 ERSDS 0000000 **** THE FOLLOWING TJBX, JAXE, PSCB, TCB'S AND STORAGE ARE FROM THE SWAPPED DATA SET ***** TJBX 0A7F68 XEST COOA7DAC XSVRB 00047700 XLAST 000A6D68 XDSE 000A7320 XRQE 0000000 XIQE 0000000 TAXE 000A6CB0 XLECB 0000000 00000000 XAIQE 0000000 XPSWD XQPL 000A7F10 RSV XNOPE 000A XNTC B 0002 XLOPL 0054 HBFL 0000 XACT 00000000 XAECB 00014534 XKEYA COCA7FBO JOB KGN01 STEP KGN01 PROCSTEP STARTING TCB OA7DAO RBP 000A7D18 FIE 0000000 DEB 0000000 TIO 000A7864 CMP 00000000 TRN 00000000 MSS 030A79A0 PK-FLG E000000 FLG 00018888 LLS OCOA7EAO JLB 00000000 JPQ 000A7EB0 RG 0-7 00000001 FFF58C74 0001A534 00014500 000A 7510 COOA7DAO 00000000 00000001 RG 8-15 000A7370 FFFFFF9 COOA7564 000A6D68 600FEAB2 000A7534 400FE930 600062FA TME 00000000 FSA 03000000 TCB 000A6D68 JST 000A7DA0 NTC 00000000 OTC 0001A788 ECB 000DDFBC TSPR 8000B828 D-PQE 000A5810 LTC 000A6D68 1 QE 0000000 SOS 000A6040 STA 200CC498 TCT 00CA73D8 USR 00000000 DAR 00001000 RES 00000000 JSCB 000A7E00 ACTIVE RBS PRB 0A7D18 RESV 00000000 APSW 00000000 WC-SZ-STAB 00040083 FL-CDE 00010580 PSW FF050001 500FEC8A Q/TTR 0C0C0000 WT-LNK 010A7DA0 NM IEFSD263 EPA OFEABO STA OFEABO LN 000550 ATR1 B9 MAIN STORAGE D-PQE 000A5810 FIRST 000A5820 LAST 00045820 PQE 0A5820 PPQ 0000000 FFB OCCCCOOO LFB 000AF000 NPQ 00000000 TCB COC1A788 RSI 00028000 RAD 00045800 FLG 0000

Figure 24. Sample of TSO Control Block Format (Part 2 of 3)

110 Programmer's Guide to Debugging (Release)

MODULE INDSADMP DATE 11/12/70 TIME 00.12 PAGE 0010 DEB 0A74A4 APPENDAGES END OF EXT 01516E SIO 01516C PCI 0151DC CH END 0151A0 AB END 01516C PFX 0000000 C2C00C0B 00003FE2 11000000 TCB 050A6D68 NDEB 01000000 ASYN 6900000 SPRG 0000000 UPRG 0200000 PLST 88000000 DCB EFOCCE64 AVT 04015158 FM-UCB START END TRKS 50002AF0 00610000 00920013 03E8 50002AB0 009F0000 00C60013 0320 TIOT 0A6E28 JOB KGNO1 STEP TMP PROC KGNPO1 OFFSET LN-STA DDNAME TTR-STC ST8-UC8 0018 14040100 S'SPRINT 00491600 80002570 002C 14040140 SYSCEMD 00480A00 80002AF0 80002AB0 0040 14040100 C0481000 0054 14040100 SYSUCUMP 00491800 80002530 14040100 0068 SYSUT1 00481200 80002530 0070 14040100 SYSUT2 004B0600 80002570 0090 14040100 00491A00 BSLOUT 800025F0 0044 14040100 SNAPTAPE 004C1100 80002530 COB8 14000010 DD1 C04B0800 00000000 0000 14000010 DD2 004B0C00 00000000 00E0 14000010 DD3 00480E00 00000000 00F4 14000010 0C4E0100 00000000 DD4 0108 14000010 DD5 004E0300 00000000 011C 14000010 004E0500 00000000 006 0130 14000010 DD7 004E0900 00000000 0144 14000010 004E0B00 DD8 0000000 PSCB 0A7B88 USER KGN01 USRL 05 GPNM SYSDA ATR1 E000 ATR2 0000 CPU 00018800 SWP CO4C33FD LTIM 008A0560 TCPU 00000000 TCON 00000000 TSWP 0000000 TC01 00000000 RLGB OCCA8700 RSV2 0000000 UPT 000A86F0 UPTL USE1 00000000 0010 PSV1 0000 USE2 CCCCODCO TAXE 046CB0 TMFLD 00 PPSAV 0CD710 ABOPSV 0000000 WCSA 00 SIZE 12 STAB 4034 00000302 EΡ 00008834 LCPSW 00040000 ROPSW USE 00 TOE 000000 WCF 00 C000000C GR1 0000000 GR2 00000000 G93 LINK 000000 GRO 00000000 GR4 00000000 00000000 GR7 00000000 00000000 GR 9 GR10 00000000 GR5 00000000 GR6 GR8 00000000 00000000 **GR13** 0000000 SR14 00000000 00000000 GR12 GR15 00000000 NIGE 0000000C GR11 IRB 000A6CB0 PRM1 00000000 000A6CB0 LNK 000A6014 TCB TLNK 000A6D68 XPSW 0000000 PARM 000ABBF8 EXIT 00000000 STAT 00000000 TAIE COOCCF7C IBUF 00000000 USER 000CCDB4

Figure 24. Sample of TSO Control Block Format (Part 3 of 3)

IMDPRDMP Output Formatting: TSO -- TSCVT 111

LUS hhhh

The maximum number of TSO users that may be logged onto the system. For additional users to be logged onto the system, the value of LUS must be greater than the value in CUS. LUS cannot exceed the value in NTJ. LUS is set by the time sharing control task (TSC). This field is initially set to the same value as MUS; however, if TSO encounters I/O errors while swapping users in and out, the time sharing control task reduces this value to limit the number of TSO users.

NTJ hhhh

The number of time-sharing job blocks (TJBs) and terminal status blocks (TSBs) allocated when TSO was started. The dummy TJB for the terminal job identification (TJID) equal zero is not included. The value of LUS cannot exceed this number.

SZU hhhh

The number of bytes in the time sharing job block (TJB).

CTR hhhh

Contains the number of region control blocks (RCBs) allocated when TSO was started. This number cannot be increased after the TSO system is started.

MUS hhhh

The maximum number of users that may be logged onto a TSO system. This field is set by the START and MODIFY commands issued by the operator.

- SAV hhhhhhh The beginning address of three 18-word save areas used by the time sharing control task (TSC), the time sharing interface program (TSIP), and the time sharing dispatcher.
- ECB hhhhhhh The address of the table control block (TSECBTAB) which contains the event control blocks (ECBs) used to post the time sharing control task (TSC), the region control tasks (RCTs), and the terminal input output coordinator (TIOC).
- SIA hhhhhhhh The address of the time sharing interface area (TSIA).
- ICB hhhhhhhh The address of the time sharing interface control block (TSICB).

101 hhhhhhh The address of the branch entry point IKJEAI01 in the time sharing interface program (TSIP).

- TQE hhhhhhhh The address of the timer queue element (TQE) used by TSO for time slicing.
- 102 hhhhhhhh The address of the entry point IKJEAI02 in the time sharing dispatcher.
- 103 hhhhhhhh The address of the entry point IKJEAI03 in the time sharing dispatcher.
- D02 hhhhhhh The address of the entry point to the TSO driver routine (IKJEAD02), or the equivalent entry in a user written routine.
- LCQ hhhhhhhh The address of the first element in the logon communications queue.
- TRB hhhhhhh The first address in the trace control block chain. This address is established and used by the statistics collection routine. It is set to zero by the time sharing control task (TSC).
- LPA hhhhhhh The address of the first contents directory entry (CDE) in the time sharing link pack area.
- SLF hhhhhhhh The address of the system-initiated logoff routine.
- TSC hhhhhhhh The address of the task control block (TCB) for the time sharing control task (TSC).
- SPL hhhhhhhh The address of the start parameter list.
- RSZ hhhh The minimum number of 2K blocks for a region during logon.
- RSV hhhh Reserved for future use.
- SVT hhhhhhhh The contents of the SVC table entry used by the time sharing interface program (TSIP).

- SVQ hhhhhhh The contents of the SVC table entry used by the TCAM/TIOC interface program.
- ABN hhhhhhh The address of the out-of-main storage abnormal termination routine (IKJEAT07). The routine is resident in main storage.
- D03 hhhhhhh The entry point address to the TSO driver MODIFY routine (IKJEAD03), or the equivalent entry point address in a user written routine.
- FLM hhhhhhhh The entry point address IKJEFLM for the system initiated logoff routine.
- QTP hhhhhhhh The entry point address IKJGGQT1 for the branch entry to the TCAM interface program (QTIP).
- T08 hhhhhhh The entry point address to the TSO command routine (IJEAT08) for TSO dumps taken by the time sharing control task (TSC) TSO dumps.
- DMP hhhhhhh The address of the TSO dump control block.
- T06 hhhhhhh The TCB address of the TSO dump routine (IKJEAT06) for the time sharing control task (TSC) modify routine.

TIME SHARING REGION CONTROL BLOCK (RCB) FORMATTING: A region control block (RCB) contains information that is unique to a time sharing region. There is one RCB for each time sharing region. The RCBs reside in the time sharing control tasks region, they are contiguous, and they are created during initialization of the time sharing controller.

RCB hhhhhh The address of the RCB.

RCT hhhhhhh The address of the task control block (TCB) for this region control task (RCT). The TCB contains the address of the partition queue element (PQE) that defines the region.

ECB hhhhhhh The event control block (ECB) on which this region control block (RCB) waits. This ECB must be posted before this region control task (RCT) can perform one of its functions. DIECB hhhhhhh The event control block (ECB) that is posted upon completion of this region control task (RCT). The time sharing control task (TSC) waits for this ECB to be posted. TJID hhhh The terminal job identification (TJID)

The terminal job identification (TJID) for the time sharing job currently executing in this region.

RSIZE hhhh

The number of 2K blocks in this region. It is set by the time sharing control task (TSC) when the time sharing system is started.

LSQSZ hhhh

The number of 2K blocks in the local system queue space (LSQS) for this region. It is set by the time sharing control task (TSC) when the time sharing system is started.

NMBR hh

The identification number assigned to this region.

PKEY hh

The protect key (PKEY) for the time sharing job currently executing in this region.

UMSMN hh

The number of entries in the main storage map which describes the main storage image that was initialized during logon.

FLG hh

This field contains the first byte of the region control block (RCB) flags. The flags indicate various functions to be performed by the region control task (RCT) and time sharing control task (TSC). These flags are set by the time sharing interface program (TSIP), the time sharing control task (TSC), and the terminal input/output coordinator (TIOC). These flags are tested and reset by the region control task (RCT) and the time sharing control task (TSC).

- 1... RCBFQO: Bit 0 is the quiesce flag. When set, this flag indicates that the current user of this region should be quiesced.
- .1... RCBFSO: Bit 1 is the swap out flag. When set, this flag indicates that the current user should be swapped out.

- ..1. RCBFSI: Bit 2 is the swap in flag. When set, this flag indicates that the current user of this region should be swapped in. The user's terminal job identification (TJID) is in the region control block (RCB).
- ...1 RCBFRS: Bit 3 is the restore flag. When set, this flag indicates that the user, whose terminal job identification (TJID) is in the region control block (RCB), should be restored by the region control task (RCT).
- 1... RCBOCAB: Bit 4 set indicates that the out-of-main storage abnormal termination routine was invoked.
-x.. Bit 5 is reserved for future use.
-1. RCBFAT: Bit 6 is the attention exit flag. When set, this flag indicates that an attention exit has been requested for one or more users.
- RCBFND: Bit 7 is the END region control task (RCT) flag. When set, this flag indicates that the region control task (RCT) should terminate normally and return control to the time sharing control task (TSC).

FLG2 hh

- This field contains the second byte of the region control block (RCB) flags. See FLG.
 - 1... RCBFSE: Bit 0 is the swap end flag. When set, this flag indicates that the swap-in operation for the current user of this region is complete.
 - .1.. RCBSTOP: Bit 1 is the region stop flag. When set, this flag indicates that a request has been made to stop the region. Every user of this region will be logged off.
 - ..1. RCBACTV: Bit 2 indicates the active status of the region control task (RCT).

The flag is set to one when the region control task is initialized; it is set to zero when the region control task is terminated.

- ...1 RCBSTR1: Bit 3 indicates that a region start has been requested, and the region control task should be attached.
- 1... RCBSTR2: Bit 4 indicates that a region start has been requested, and a swap logon image should be created.
-xxx Bits 5 through 7 are reserved for future use.

FBQE hh

The number of free block queue elements (FBQEs) for this region.

UTTMQ hhhh

The relative track address (TT) of the map queue pointer. The map queue pointer describes the location of the region's initialized logon image on the swap data set.

CUSE hhhh

The number of users logged on to use this region. The time sharing control task (TSC) increments the count before disconnect (DISC) and decrements the count during logon.

EXTNT hhhhhhhh

The address of the initialized time sharing job block extension (TJBX). The TJBX is created during the logon initialization for this region.

UMSM hhhhhhhh

The address of the user main storage map. This map describes the initialized logon main storage image for this region.

SDCB hhhhhhhh

The address of the swap data set control block (SDCB). This block points to the location of the initialized logon image on the swap data set for this region.

PQE hhhhhhhh

The address of the partition queue element (PQE) pointer in the system queue space (SQS). The PQE describes the main storage space assigned to this region. The PQE pointer is used to manipulate main storage when (1) this region control task's (RCT's) region is obtained during start time sharing initialization and (2) this region control task's (RCT's) region

is freed during region control task termination. PRG hhhhhhhh PRG1 hhhhhhhh PRG2 hhhhhhhh These three words constitute the SVC I/O purge parameter list. For further information, see the "Purge Macro Instruction" in the publication <u>IBM</u> <u>System/360:</u> System Programmer's Guide, GC28-6550. QPL hhhhhhhh The address of the quiesce I/O parameter list. STECB hhhhhhhh An event control block (ECB). During a subsystem recovery, the time sharing control task (TSC) waits for this ECB to be posted by the region control task (RCT). The posting is done during end processing. RCOVR hhhhhhhh These bits indicate the current recovery status of the region control task (RCT) in the event of a subsystem failure. byte 0 1.... RCBRCOVR: Bit 0 set indicates that the status bits in the following 3 bytes are valid. .xxx xx.. Reserved for future use. processing complete. **X RCBTACMP:** Transient area restore processing complete. byte 1 -- RCBRSFLG 1.... RCBRSTRT: Bit 0 set indicates a restore. .1....RCBTCBDN: Bit 1 set indicates that the task control blocks (TCBs) have been requeued. ..1. RCBQELCM: Bit 2 set indicates that the queue element (QEL) restore processing is complete. ...1 RCBTQECM: Bit 3 set indicates that the timer queue element (TQE) restore processing is complete. 1... RCBRQIQC: Bit 4 set indicates that both the request queue element (RQE) and the interrupt queue element (IQE) restore

processing is complete.

.... .1.. RCBIORSC: Bit 5 set indicates that the I/O restore processing is complete.

.... ..xx Bits 6 and 7 are reserved for future use.

- byte 2 -- RCBQUFLG 1... RCBQUSTR: Bit 0 set indicates that quiesce has started.
 - .1.. ... RCBIOSTR: Bit 1 set indicates that the first entry into the I/O purge routine is complete.
 - ..1. RCBTADON: Bit 2 set indicates that the transient area quiesce is complete.
 - ...1 RCBWTORD: Bit 3 set indicates that the write to operator with reply (WTOR) quiesce is complete.
 - 1... RCBQELDN: Bit 4 set indicates that the queue element (QEL) quiesce is complete.
 -1.. RCBIODON: Bit 5 set indicates that the second entry I/O purge is complete.
 -1. RCBTQEDN: Bit 6 set indicates that the timer queue element (TQE) quiesce is complete.
 - RCBRQIQD: Bit 7 set indicates that both the request queue element (RQE) and the interrupt queue element (IQE) are complete.
- byte 3 1... RCBSWTCH: Bit 0 indicates the method of search used by various subroutines in IKJEAT07. When equal to zero, all system users are purged according to the terminal job identification (TJID). When equal to one, all users in this region are purged as indicated by the region control block addresses.
 - .1.. RCBSWTON: When bit 1 is set along with bit 0 being set, all system users are purged. A search is made according to the terminal job identification (TJID) and the request control block (RCB).

..xx xxxx Bits 2 through 7 are reserved for future use.

CONID hh

The routing code of the console that issued the last START, MODIFY, or STOP command.

RESV hhhhhh

Reserved for future use.

<u>RAIN STORAGE MAP (UMSM) FORMATTING:</u>

UMSM is used in the swap operation. One user main storage map exists for each possible time sharing user. The UMSM contains a series of consecutive one-word extent fields (ADDR-LN). Each one-word extent contains a halfword address field (ADDR) and a halfword length field (LN) that describe the main storage space allocated to the time sharing user. The number of UMSM extents has established defaults that can be modified by the operator when he starts the time sharing system. The number of extent entries is stored in the time sharing job block (TJB) at TJBUMSMN. Unused extent fields contain zeros.

- UMSM hhhhhh The address of the user main storage map.
- ADDR-LN hhhhhhhh
 - bytes 0 and 1 Begin Address: This field contains the two high order bytes of the beginning address of the main storage segment allocated to the time sharing user. Since main storage is allocated in 2K blocks, the low order byte is always zero and, therefore, need not be kept in a control block.
 - bytes 2 and 3 This field contains the two high-order bytes designating the length of the main storage space allocated to the time sharing user. Since main storage is allocated in 2K blocks, the low-order byte is always zero and, therefore, need not be kept in a control block.

SWAP DATA CONTROL BLOCK (SWAP DCB) FORMATTING: The swap data control block (SWAP DCB) is used whenever a time sharing user's region is swapped into or out of main storage. Each region control task (RCT) has one swap data control block. Following the address of the swap data control block is the contents of the main storage data that was written on the swap data set.

SWAP DCB hhhhhhhh The address of the swap data control block. TIME SHARING JOB BLOCK (TJB) FORMATTING: The time sharing job block (TJB) contains status information about the time sharing user. The TJB is retained in main storage while the user is swapped out. One time sharing job block exists for each possible simultaneous time sharing user. The space for the TJB is obtained from the time sharing control task (TSC) region during time sharing initialization. Status information about terminals related to this TJB is contained in the terminal status block (TSB). The address of the terminal status block is the first word of the TJB.

```
TJB hhhhhh
The address of this TJB.
```

TSB hhhhhhhh The address of the terminal status block (TSB) that owns this terminal job. If zero, this job was started by an operator command.

ATTN hh

A count of the unprocessed attention interrupts for this job.

STAX hh

The number of scheduled specify terminal attention exits (STAXs).

- STAT hh This field contains flags that indicate the status of the time sharing job.
 - 1.... TJBNJB: Bit 0 set indicates that this TJB is currently unused.
 - .1.. TJBINCOR: Bit 1 set indicates that this user is currently in main storage.
 - ..1. TJBLOGON: Bit 2 set indicates that the logon start has been set by the terminal input output coordinator (TIOC) during a dialup to request a logon. This bit is reset by the time sharing control task (TSC).
 - ...1 TJBIWAIT: Bit 3 set indicates that the terminal job is in an input wait state.
 - 1... TJBOWAIT: Bit 4 set indicates that the terminal job is in an output wait state.

-1.. TJBSILF: Bit 5 set indicates that the user is to be logged off the system. This bit is set by the IKJSILF subroutine and tested by the region control task (RCT) restore routine that posts the logon ECB. This bit is tested and reset by the logon/logoff routine.
-1. TJBDISC: Bit 6 set indicates that a request has been made to the terminal input output coordinator (TIOC) to disconnect the line.
- Bit 7 is reserved for future use.
- STAT2 hh These flags indicate the status of the time sharing job.
 - 1.... TJBHUNG: Bit 0 set indicates that the user's communication line disconnected.
 - .1....TJBHOLD: Bit 1 set indicates that an output wait (OWAIT) exists because of a hold option.
 - ..1. TJBOCAB: Bit 2 set indicates an out-of-main storage abnormal termination has occurred for this user.
 - ...1 TJBRNAV: Bit 3 set indicates that the user cannot be logged onto the time sharing system because (1) a machine check occurred in the user's region or (2) the region is too small for the user.
 - 1... TJBSURSV: Bit 4 set indicates that on the next swap in the swap unit is not marked as available for the user.
 -xxx Bits 5 through 7 are reserved for future use.
- EXTNT hhhhhhhh The address of the terminal job block extension (TJBX) when it is in main storage.
- RCB hhhhhhhh The address of the region control block (RCB) for this job.

UMSM hhhhhhhh The address of the user main storage map (UMSM) for this job.

- SDCB hhhhhhhh The address of the swap data control block (DCB) for this job.
- UTTMO hhhh
 - 1... TJBUTTMP: Bit 0 of byte 0 set indicates a parallel swap.
 - .111 1111 Bits 1 through 7 of byte 0 along with byte 1 contain the offset into the map queue. The map queue contains a chain of allocation units for this user on the swap data set. The address of the queue is in the UTTMQ field of the TSO region control block (RCB).
- RSTOR hh This field contains the status flags used by the region control task (RCT) restore operation.
 - 1... TJBOWP: Bit 0 set indicates to the terminal input output coordinator (TIOC) to end the output wait (OWAIT) condition.

 - ..x. Bit 2 is reserved for future use.
 - ...1 TJBLOGP: Bit 3 set indicates that the event control block (ECB) waited for by the logon image should be posted. This flag is set by the time sharing control task (TSC) logon routine and by the IKJSILF subroutine.
 - 1... TJBLWAIT: Bit 4 set indicates that if the user is not made ready by restore processing, he should be swapped out again.
 -x.. Bit 5 is reserved for future use.
 -1. TJBFAT: Bit 6 set indicates that an attention exit is requested for this user's job.

.... Bit 7 is reserved for future use.

- UMSMN hh The number of entries in the user main storage map (UMSM).
- USER CCCCCCC The userid of the user who owns this job. This field may have trailing blanks when the user identification contains less than eight characters.
- IPPB hhhhhhh An address pointer to the beginning of a chain of inter-partition post blocks that indicate the event control blocks (ECBs) to be posted by the restore operation.

NEWID hh

Identifies the region where the user should be logged on. When this field is zero, the TSO driver should select the region. When this field is set by the end-of-routine for logon/logoff, it identifies the new region to which the user will be shifted.

FLUSL hh

Reserved for future use.

TJID hhhh

This field contains the terminal job identification (TJID) for this time sharing job.

MONI hh

These flags indicate various processing functions that cause operator messages to be sent to this terminal. The flags are set and reset when the terminal user issues the MONITOR subcommand of the OPERATOR command.

- 1... TJBMDSN: Bit 0 set indicates that the first non-temporary data set allocated to a new volume should be displayed as part of the mount and keep messages.
- .1.. ... TJBMJBN: Bit 1 set indicates that the name of each job is to be displayed on the console when each job is initiated and terminated, and that the unit record allocations are to be displayed when a job step is initiated.

- ..1. TJBMSES: Bit 2 set indicates that, when a terminal session is initiated or terminated a message is displayed on the operator console.
- ...1 TJBMSPA: Bit 3 set indicates that the available space on a direct access device is to be displayed on the operator console as part of the demount message.
- 1... TJBMSTA: Bit 4 set indicates that, at the end of a job or job step, certain data set disposition information should be printed with the demount messages. These dispositions are: KEEP, CATLG, or UNCATLG.
-xxx Bits 5 through 7 are reserved for future use.
- RSV hhhhhh Reserved for future use.

TERMINAL STATUS BLOCK (TSB) FORMATTING:

Each terminal status block (TSB) contains status information about one terminal user. The terminal input output coordinator (TIOC) uses this information. During system initialization, one TSB is created for each possible user. The main storage space is obtained in one contiguous block for all of the TSBs in the region of the time sharing control task (TSC); this contiguous string of TSBs is called the TSB table. The origin pointer to the TSB table is the TIOCTSB field in the TIOCRPT.

- TSB hhhhhh The address of this terminal status block (TSB).
- STAT hh This field contains the terminal status indicator flags.
 - 1... TSBINUSE: Bit 0 set indicates that this TSB is being used.
 - .1.. TSBLWAIT: Bit 1 set indicates that the terminal keyboard is locked due to a lack of input buffer space.
 - ..1. TSBDSPLY: Bit 2 set indicates that this TSB represents a terminal which is a graphic device.

- ...1 TSBNOBUF: Bit 3 set indicates that TPUT found no time sharing buffers.
- 1... TSBITOFF: Bit 4 set indicates that this user wishes to prevent inter-terminal communications.
-1.. TSBDISC: Bit 5 set indicates that this TSB has been processed by logoff.
-x. Bit 6 is reserved for future use.
- I TSBATNLD: Bit 7 set indicates an attention for an input line deletion.
- TJB hhhhhh The address of the time sharing job block (TJB) currently used by this terminal. This field contains zeros when this terminal is not associated with a time sharing job block.
- FLG1 hh This field contains terminal status flags.
 - 1... TSBANSR: Bit 0 set indicates that an attention simulation is requested.
 - .1.. ... TSBOFLSH: Bit 1 set indicates that the output trailer queue is to be flushed. This bit is set by TCLEARQ.
 - ..1. TSBOWIP: Bit 2 set indicates that a TPUT operation is in progress.
 - ...1 TSBWOWIP: Bit 3 set indicates that a task is waiting for another task to complete a TPUT operation.
 - 1... TSBIFLSH: Bit 4 set indicates that an input queue flush is in progress.
 - 1.. TSBTJOW: Bit 5 set indicates that this user is already using the maximum number of output buffers that can be allocated. This TSB waits on event control block (ECB) for this TCB. This bit is set by a TPUT macro instruction with a terminal job identification (TJID).

- Bit 6 is reserved for future use.
- I TSBTJBF: Bit 7 set indicates that no time sharing buffers were available when the SVC for TPUT with the terminal job identification (TJID) was issued. The system waits for the TJB event control block (ECB) to be posted.

WTSB hhhhhh Reserved for future use.

LNSZ hh The number of characters that can be printed on one line for this terminal. This field is set by either logon or STSIZE.

- OTBFP hhhhhh The address of the trailer buffer if the heading buffer for a message has been removed from the message queue. This field is reset to zeros when the message has been completely moved to the TCAM buffers.
- NOBF hh The number of buffers on the output queue.
- OBFP hhhhhh The address of the first buffer on the output buffer queue.
- BRKFL hh These flags indicate the status of the communication line.
 - 1... TSBBIPI: Bit 0 set indicates to the TSINPUT that a partial line exists for prompting. Set by TSOUTPUT.

.1.. ... TSBAUTON: Bit 1 set indicates that automatic input line numbering is requested.

- ..1. TSBBRKIN: Bit 2 set indicates that TPUT is using the breakin option and a partial line was assigned to this function. This bit is set by TSINPUT. TSINPUT is a TCAM subtask.
- ...1 TSBAULST: Bit 3 set indicates that automatic line numbering has started.
- 1... TSBAUTOC: Bit 4 set indicates that automatic character prompting is used.

IMDPRDMP Output Formatting: TSO -- TSB 119

-1.. TBSTAUT: Bit 5 set indicates that the user is being prompted with the next line number.
-11 TSBSATN1: Bits 6 and 7 contain a count of the number of characters used to simulate attention.
- ITBFP hhhhhh The address of the first buffer in the trailer input buffer chain.
- NIBF The number of buffers on the input queue.
- IBFP hhhhhh The address of the first buffer in the input buffer queue chain.
- CLEAR hh
 - This field contains terminal status flags.
 - 1.... TSBATTN: Bit 0 set indicates that an attention from this terminal has been ignored.
 - .1.. TSBTJMSG: Bit 1 set indicates that TSOUTPUT is processing a terminal job identification (TJID) message.
 - TSBSPIT: Bit 2 set indicates that breakin prompt and automatic prompt are suppressed.
 - ...1 TSBNBKSP: Bit 3 set indicates that the next character in the user's buffer is a backspace character.
 - xxxx Bits 4 through 7 are reserved for future use.
- QCB hhhhhh The address of the queue control block (QCB) that contains the destination for the message being sent.
- ECB hhhhhhhh The event control block (ECB) at which the inter-terminal communication (TPUT with TJID) waits (1) when there are no time sharing buffers, (2) when the TSBOWIP bit is set, or (3) when the TSBOQHLD bit is set.
- TJID hhhh The terminal job identification (TJID) of the task waiting on this TCB's event control block (ECB).

These two bytes define special purpose characters that may be redefined by the terminal user. byte 0 TSBLNDCC: This byte contains the line delete character. byte 1 TSBKSPCC: This byte contains the character delete character. ATNLC hhhh The number of successive lines of printed output between attention simulation reads. ATNTC hhhh The number of seconds between attention simulation reads. LNNO hh When a graphic terminal device is used, this is the number of line that can be displayed. BLNK hh Reserved for future use. ASRCE hhhh This field contains the same information as the PRFSRCE field in the TCAM buffer prefix. ATNCC CCCC This field contains from one to four characters that are used to simulate attention. Some of the character positions may contain blanks. AUTOS hhhhhhhh This field initially contains the starting line number for the first line number. AUTOI hhhhhhhh This field contains the value that is

STCC hhhh

input line. While the line of input information is being received from the terminal user, this field is updated to contain the value of the current

used to automatically increment the value of the input line numbers. This field can be modified by the terminal user.

ERSDS cccc

When a graphic terminal device is used, this word contains the characters used to erase the display screen.

TIME SHARING JOB BLOCK EXTENSION (TJBX) FORMATTING: The time sharing job block extension (TJBX) contains user job information that can be rolled out to the swap data set with the user's job. The

TJBX resides in the local system queue space (ISQS) for the region. The TJBX location is pointed to by the third word of the time sharing job block (TJB). The space for the TJBX is obtained by the region control task (RCT) during initialization.

TJBX hhhhhh The address of the TJBX.

- XFST hhhhhhhh The address of the logon TCB. The logon TCB is the first TCB on the user's ready queue.
- XLAST hhhhhhhh The address of the last TCB on the user's ready queue.
- XDSE hhhhhhhh The address of the data set extension (DSE) used by TSO dynamic allocation.
- XSVRB hhhhhhhh The address of the first supervisor request block (SVRB) purged from the transient area queue.
- XRQE hhhhhhhh The address of the first request queue element (RQE) purged from the asynchronous exit queue.
- XIQE hhhhhhhh The address of the first interrupt queue element (IQE) purged from the asynchronous exit queue.
- TAXE hhhhhhhh The address of the queue of terminal attention exit elements (TAXEs) used to schedule the attention exits.

XLECB hhhhhhhh The logon event control block (ECB) that was posted by the region control task (RCT) to activate logon/logoff.

XPSWD cccccccc The password entered by the terminal user during logon. If the password contains less than eight characters, the field is padded to the right with blanks. The entire field contains blanks when the user is not required to enter a password.

XATTR hhhhhhhh The address of the chain of attribute control blocks (ATRCBs).

XAIQE hhhhhhhh The address of the attention interrupt queue element (IQE) currently being processed by the attention proloque. XQPL hhhhhhhh The address of the quiesce parameter list (QPL).

XNQPE hhhh The number of entries in the quiesce parameter list (QPL).

- XNTCB hhhh The number of task control blocks (TCBs) active in the user's job step. When the value in XNTCB exceeds XNQPE, the quiesce parameter list is enlarged.
- XLQPL hhhh The number of bytes in the quiesce parameter list.
- RSV hhhh Reserved for future use.
 - XACT hhhhhhhh The relative track and record address (TTR) for the account control table (ACT) on SYSJOBQE.
 - XAECB hhhhhhhh This field contains either: (1) The address of the logon/logoff event control block (ECB) when logon processing begins. (2) The address of the command scheduling block (CSCB's) cancel event control block (ECB) after the CSCB is created.

XKEYA hhhhhhhh The address of the storage key save area.

<u>PROTECTED STEP CONTROL BLOCK (PSCB)</u>: The protected step control block (PSCB) contains accounting information related to a single user. All timing information is in software timer units. A software timer unit is equal to 26.04166 microseconds.

PSCB hhhhhh The address of this PSCB.

- USER ccccccc These seven bytes contain the userid entered by the terminal user during logon. If necessary, it is padded to the right with blanks. This field uniquely identifies each terminal user in the time sharing system.
- USRL hh The number of nonblank characters in the userid.

GPNM cccccccc An eight-byte group name initialized by logon from the user attribute data set (UADS). When a name is not available from UADS, the unit name used by the dynamic allocation

IMDPRDMP Output Formatting: TSO -- TJBX 121

interface routine (DAIR) is used, if a name is required.

ATR1 hhhh Sixteen bits used to define terminal user attributes.

byte 0

- 1... PSCBCTRL: Bit 0 set indicates that the user may use the OPERATOR command.
- .1.. ... PSCBACCT: Bit 1 set indicates that the user may use the ACCOUNT command.
- ..1. PSCBJCL: Bit 2 set indicates that the user may use the SUBMIT, STATUS, CANCEL, and OUTPUT commands.
- ...x xxxx Bits 3 through 7 are reserved for future use.

byte 1 Reserved for future use.

ATR2 hhhh

bytes 0 and 1 Reserved for use by IBM customers.

- CPU hhhhhhh The cumulative CPU time used by this terminal user during this session. The CPU field is set to zero during logon.
- SWP hhhhhhh The cumulative time that this terminal user has been resident in the region. The SWP field is set to zero during logon.
- LTIM hhhhhhh The actual time of day that this user logged on to the time sharing system for this session.
- TCPU hhhhhhh The total CPU time used by this terminal user, excluding the current session.
- TSWP hhhhhhh The total time that the terminal user has been resident in the region during this accounting period, excluding the current session.

TCON hhhhhhhh

TCO1 hhhhhhh TCON and TCO1 are a single eight byte field. This field contains the total connect time for this terminal user during this accounting period, excluding the current session. RLGB hhhhhhhh The address of the re-logon buffer block used by logon as a pointer to the re-logon command buffer.

- UPT hhhhhhhh The address of the user profile table (UPT).
- UPTL hhhh The number of bytes in the user profile table.
- RSV1 hhhh

RSV2 hhhhhhhh RSV1 and RSV2 are a single six byte field that is reserved for future use.

- USE1 hhhhhhhh
- USE2 hhhhhhhh USE1 and USE2 are a single eight byte field reserved for use by IBM customers.

TERMINAL ATTENTION EXIT ELEMENT (TAXE) FORMATTING: The TSO terminal attention exit element (TAXE) consists of a regular 24 word interrupt request block (IRB) plus a TSO addendum. It is used to schedule an attention exit resulting from a terminal attention interruption. It is created, queued, and dequeued by the specify terminal attention exit (STAX) macro instruction. The main storage space for the TAXE is obtained in the local system queue space (LSQS) of the terminal user's region.

TAXE hhhhhh The address of this TAXE when it is in main storage.

TMFLD hh

This field contains indicators for the time routines.

- 1... Bit 0 set indicates that the timer element was not queued.
- .1.. Bit 1 set indicates that the local time-of-day option is used.
- ..00 Bits 2 and 3 set to zero-zero indicate that the time interval was requested in timer units (26.04166 microseconds).
- ..01 Bits 2 and 3 set to zero-one indicate that the time interval was requested in binary units.

..10 Reserved for future use.

- ..11 Bit 2 and 3 set to one-one indicate that the time interval was requested in decimal digits.
- 1... Bit 4 set indicates that the time interval has expired.
-000 Bits 5 through 7 set to zero-zero-zero indicate an STIMER task time request.
-001 Bits 5 through 7 set to zero-zero-one indicate an STIMER wait request.
-011 Bits 5 through 7 set to zero-one-one indicate an STIMER REAL time request.
-100 Bits 5 through 7 set to one-zero-zero indicate an STIMER task time request with a specified exit.
-111 Bits 5 through 7 set to one-one-one indicate an STIMER REAL time request with a specified exit.
- Other combinations of bits 5 through 7 are reserved for future use.
- PPSAV hhhhhh The starting address of the register save area for the problem program.
- ABOPSW hhhhhhhh

This field displays the right half (bytes 4 through 7) of the program status word (PSW) that was active in the dump system during the execution of an ABEND or ABTERM routine. If these routines have not been invoked, then this field contains zeros.

- WCSA hh The number of requests waiting when termination occurred.
- SIZE hh The number of doublewords in this request block.
- STAB hhhh This field contains two bytes of status and attribute information.

byte 0 The TAXE is a type of interrupt request block (IRB). Byte zero identifies the type of request block; however, for the TAXE, only the IRB identification is used. 01....Bits 0 and 1 set to zero-one indicate that this is an interrupt request block (IRB).

- byte 1 This byte contains various request block indicators.
 - 1... Bit 0 set indicates that the RBLINK field points to the task control block (TCB).
 - .1... Bit 1 set indicates that the program related to the interrupt request block (IRB) is active.
 - ..1. Bit 2 set indicates that this interrupt request block (IRB) is for an exit routine (ETXR).
 - ...x Bit 3 is reserved for future use.
 - 00.. Bits 4 and 5 set to zero-zero indicate that the request queue element (IQE) is not to be returned.
 - 01.. Bits 4 and 5 set to zero-one indicate that the interrupt request block (IRB) has queue elements for asynchronously executed routines that are request queue elements (RQEs).
 - 10.. Bits 4 and 5 set to one-zero indicate that an interrupt queue element (IQE) is not to be returned at EXIT.
 - 11.. Bits 4 and 5 set to one-one indicate that the interrupt request block (IRB) has queue elements for asynchronously executed routines that are interrupt queue elements (IQEs).
 -1. Bit 6 set indicates that the request block storage can be freed at exit.
 - Bit 7 set to zero indicates a wait for a single event or all of a number of events.
 - Bit 7 set to one indicates a wait for a number of events that is less than the total number of events that are waiting.
- EP hhhhhhh The address of the routine that was asynchronously executed.

IMDPRDMP Output Formatting: TSO -- TAXE 123

ROPSW hhhhhhhh (Right half of PSW) This program status word (PSW) contains the status of the program represented by the request block being displayed when a new request block was created. Had the dumped system been allowed to continue processing normally, the operation would have been resumed with this PSW. USE hh This field contains the use count as used by ATTACH. IQE hhhhhh The address of the list origin for the interrupt queue element (IQE). WCF hh The number of requests that were pending when this dump was taken. LINK hhhhhh The address of the next request block (RB) on this RB queue. If this is the last request block on the queue, then this field contains the address of the task control block (TCB). GR0 hhhhhhhh GR15 hhhhhhhh The general register save area used by the supervisor. NIQE hhhhhhhh The address of the next available interrupt queue element (IQE). LNK hhhhhhhh The address of the next interrupt queue element (IQE). PRM1 hhhhhhhh The address of the parameter list for the asynchronous exit routine. IRB hhhhhhhh The address of the interrupt request block (IRB) to be scheduled next. TCB hhhhhhhh The address of the task control block (TCB) for this TAXE. TLNK hhhhhhhh The address of the next TAXE on this queue. XPSW hhhhhhhh The left half (bytes 0 through 3) of the program status word (PSW) for the user attention exit routine.

LOPSW hhhhhhhh (Left half of PSW)

EXIT hhhhhhhh The address of the user attention exit routine.

- STAT hhhhhhhh This field contains status flags for this TAXE.
 - byte 0 1... TAXEFKEY: Bit 0 set indicates that the task issuing the specify terminal attention exit (STAX) macro instruction is a problem program.
 - .1.. TAXEMOD: Bit 1 set indicates that the task issuing the specify terminal attention exit (STAX) macro instruction is in problem program mode.
 - ..1. TAXEFFREQ: Bit 2 set indicates that the requested TAXE is not available for scheduling.
 - ...x xxxx Bits 3 through 7 are reserved for future use.
 - bytes 1-3 Reserved for future use.
- PARM hhhhhhh The address of the parameter list for the specify terminal attention exit (STAX) macro instruction.
- TAIE hhhhhhhh The address of the terminal attention interrupt element.
- IBUF hhhhhhhh The address of the user input buffer.
- USER hhhhhhhh The address of the user parameter list from the specify terminal attention exit (STAX) macro instruction.

Task Control Block Summaries

If, during the course of program execution, the IMDPRDMP program formatted the major system control blocks of more than one MVT or MFT task, a summary of each displayed task's TCB is presented at the end of the control block portion of the dump listing. Depending upon the operating system option under which the dumped task was operating, either the MVT/MFT-with-subtasking TCB summary format (Figure 25), or the abridged MFT-without-subtasking TCB summary format (Figure 26) is presented.

Both summary formats are identified by two lines of heading information. The

first heading line displays the optional dump listing title, the name of the module that invoked the dump, and the date and time that the information was captured from the dumped system. The second line of heading displays the identifying phrase "**** TCB SUMMARY ****."

The individual TCB summaries contain the following information:

<u>MVT or MFT with Subtasking TCBs</u>: Are summarized in the two-line array illustrated in Figure PROUT-9 and described below:

JOB ccccccc The JOB field in the first line of each task control block array displays to the user the eight-character name of the job associated with the TCB.

STEP cccccccc The STEP field shows the eightcharacter name of the job step as it appeared on the label field of the EXEC JCL statement associated with the step.

TCB hhhhhh The starting address of the task control block.

CMP hhhhhhhh

This field shows the ABEND indicators and user and system completion codes associated with this TCB. (See the relevant TCB discussion for the contents of this field.)

NTC hhhhhhhh This word contains the address of the TCB that occurred previous to this one

on the originating task's subtask queue. If the TCB being summarized was the first on the queue, this field displays zeros.

OTC hhhhhhhh

The OTC field displays the address of the TCB representing the originating task.

LTC hhhhhhhh This field c

This field contains the address of the TCB that occurred last on the originating task's subtask queue at the time the dump was taken. If the TCB being summarized was the last on the subtask queue, this field contains zeros.

PAGE ddd The page of the dump listing on which the formatted control blocks associated with this TCB, may be found.

MFT Without Subtasking TCBs: Are summarized in the two line abridged array illustrated in Figure PROUT-10 and described below:

JOB cccccccc The JOB field in the first line of each task control block array, displays to the user the eight-character name of the job associated with the TCB being summarized.

STEP cccccccc The STEP field shows the eight-character name of the job step as it appeared in the label field of the EXEC JCL statement associated with the step.

TCB hhhhhh The starting address of the task control block being summarized is given in the first field of this second line.

CMP hhhhhhh This field shows the ABEND indicators and user and system completion codes associated with the TCB. (See the MFT TCB discussion for a description of the contents of this field.)

PAGE ddd The page of the dump listing on which the formatted control blocks associated with this TCB, are found.

			MODULE INDSADMP DATE 11/12/70 TIME 00.15 PAGE 0032
		* * * *	TCBSUMMARY * * * *
JOB		NTC C0000000	OTC 00009CA0 LTC 00000000 PAGE 0004
108		NTC 0000000	0TC 00009CA0 LTC 00000000 PAGE 0005
. JJB	STEP TCB 008868 CMP 0C000000	NTC 0000000	OTC 00009CA0 LTC 00000000 PAGE 0006
	STEP TCB 0089A8 CMP C00C0000	NTC 0000000	OTC 00009CA0 LTC 0000000 PAGE 0007
JOB		NTC 00000000	OTC 00009CA0 LTC 00000000 PAGE 0008
JOB	STEP TCB 008C28 CMP 00000C00	NTC 0000000	DTC 00009CA0 LTC 0000000 PAGE 0009
JOB		NTC 0000000	OTC 00009CA0 LTC 00000000 PAGE 0010
JJB		NTC 0000000	OTC 00009CA0 LTC 00000000 PAGE 0011
J05	STEP TCB 008FE8 CMP 00000000	NTC C0000000	DTC 00009CA0 LTC 00000000 PAGE 0012

		MODULE IMDSADMP DATE 11/12/70 TIME 00.15 PAGE 0033
	* * * *	TCBSUMMARY * * * *
JOB		OTC 00000000 LTC 0002E268 PAGE 0022
JGB	MASTER STEP SCHEDULR	
JOB	TCB 0288C8 CMP 00000000 NTC 00009BA8	OTC 00009CA0 LTC 00000000 PAGE 0025
	TCB 02E0F8 CMP 0000000 NTC 000288C8	DTC 00009CA0 LTC 0002D1E8 PAGE 0027
	TCB 0201E8 CMP C0000000 NTC 0000000	DTC 0002E0F8 LTC 00020400 PAGE 0028
	TCB 32D400 CMF 0000000 NTC 0000000	OTC 0002D1E8 LTC 00000000 PAGE 0025
J08	WTR STEP CCE	
	TCB 02E268 CMP 0C000000 NTC C002E0F8	DTC C0009CA0 LTC 00C2D108 PAGE 0030
	TCB 020108 CMP 00000000 NTC 00000000	OTC C002E268 LTC 0000000 PAGE 0031
1		

Figure 25. TSB Summary Sample for System That Operated Under MVT or MFT With Subtasking

126 Programmer's Guide to Debugging (Release 21)

AFT	DUMP LISTING				MJOULÉ IM	JSADMP	DATE 11/12/7	TIME CO.5C	PAGE 0C22
		* * * *	ТСВ	SU AIM	ARY	* * * *	r.		
90 1	STEP TCB 008778 CMP 0000CCCC	PAGE 0001							
9 OR	MASTER STEP SCHEDULR TCB JO3353 CMP JCCCODOO	PAGE 0002							
109	STEP TCB 008938 CMP 0000CC0C	PAGE JC04							
109	STEP TCB J08418 CMP OCOCOCOC	PAGE 0005							
109	MASTER STEP SCHEDULR TCB 008368 CMP 00COCCOC	PAGE OCO6							
1 08	MASTER STEP SCHEDULR TCB 008340 CMP JG00CCOC	PAGE 0007							
109	WTR STEP PU TCB 038148 CMP OCDOCCOC	PAGE CCO9							·
103	STEP TCB 038F40 CMP 30000000	PAGE 0010							
109	JOB5 STEP GJ TCB JJ9148 CMP J0000000	PAGE COLL							
109	STÉP TCB 009348 - ÚMP úcccococ	PAGE 0013							
119	STEP TCB J09348 CMP JOCODCCC	PAGE 0014							
108	STEP TCB 009748 CMP 000CC00C	PAGE 0015							
109	STEP TCB 009548 CMP 0000COC	PAGE 0016							
103	STEP - TCB 009348 CMP 0000000	PAGE 0017							
103	STÉP TEB 009348 CMP 00000000	PAGE CO18							
1 08	STEP TCB 009F48 CMP D0000000	PAGE JO19							
109	STÉP TCB JOA148 CMP JC00000C	PAGE 0020							
1 08	STEP TCB JJA348 CMP JOCCCJOC	PAGE 0021							

Figure 26. TCB Summary Sample for Systems that Operated Under MFT Without Subtasking

IMDPRDMP Output Formatting: General Format 127

THE GENERAL FORMAT

The IMDPRDMP program uses a general format to display the hexadecimal contents of main storage. The particular areas of main storage displayed are determined by the parameters entered after the PRINT user control verb.

To identify various dump printouts, IMDPRDMP prints specific headings on each dump, such as ALLOCATED STORAGE, PRINT STORAGE, and NUCLEUS and SQA PRINT. A sample of a general format dump is shown in Figure 27.

The IMDPRDMP program also reverts to the general format if it is unable to format control block information because it encountered either a control block error or one of several user control statement format errors.

Each page of an IMDPRDMP program dump listing containing information displayed in the general format is identified by a heading line. This heading line shows the optional title supplied by the user followed by the date and time that the information was taken from the dumped system. A sequential page number also appears in each heading line.

Listings being produced under control of the PRINT ALL, PRINT CURRENT, or PRINT STORAGE (no operands) format control statement display the contents of the sixteen general purpose registers. If the dump was obtained from a multiprocessing system and both sets of registers were obtained, then the contents of both sets of registers are displayed. Where applicable, the beginning of each main storage region is noted by a line that gives the job, step and procedure step name of the owning task, followed by the status of the region (BORROWED, ROLL-OUT, OWNED).

Then, starting at an address requested by the user, as specified in a PRINT user control statement, (or location zero if no address was specified) the contents of main storage are displayed. Each line of the general format displays eight words of main storage. Preceding each line of information is the address of the first byte displayed followed by a one-character storage protection key indicator representing the key associated with the area of main storage being displayed on this line. Following each line of information, a 32-character translation field is printed. This field gives the EBCDIC translation of the translatable characters in the eight hexadecimal words. Untranslatable bytes are represented by positional periods.

Printing of any line that duplicates the contents of the line printed previously, is suppressed. Duplicate lines are indicated by the phrase "TO NEXT LINE ADDRESS SAME AS ABOVE" following the line duplicated.

1							MO	DULE IMDSA	DATE	11/12/70 TIME 00.15. PAGE 0001
	R 0-7	00000000	00002208	00000000	80002144	00002280	A000000A	00000000	00000000	*••••••
1	R 8-15		00000000			00000000	00000000	00000000	40002 JB4	*
1	000000	00000191	00001000	40002084	600002B	08000080	40000001	FFE 50000	900432B6	*••••V••••*
	000020	FFC40C01	5000BBB2	FFF50004	4006E7C2	0000FF00	00000000	FF060009	80 00 00 00	*
	000040	000022E8	00000000	00002280	00005508	5A64336D	48100002	41200000	50200048	*•••Y•••••
	000060	98240008	90001000	00020000	0000003	90001000	47700070	91030044	4750007C	*•••H•••••••**
	000030	310000A6	4000005	0800080	40000001	05001C00	40000500	06001000	000004B0	*
1	000040	00000000	00000000	00000450	00D20650	44500088	47F0006C	02002000	00084040	*Q *
1	000000	02000008	2000048	C2C5D5C4	40404040			40404040		*HEND *
1	03 00 00	40404040	40404040	40404040	404 040C 6	FOF8C1D7	D9F7F040	FOF04BF1	F140F1F4	★ F08APR70 00.11 14*
1	000100		F9F94040					00000000		*•01•99 ••••••••*
	000120		00000000					00000000		*•••••
1	000160		00000000					00000000		*
	000180		8000000					00000001		*
1	0001A0		00009AF4					00009AD0		*4
1	000100		00009BB4					00000030		*•••••••
1	C 201E0		00072900					A006E740		*XXX
1	000200		00067594					00000000		*•••••
	000220		00000000					50881804		*•••••
	000240		D5022015					0E681B99		*••••N••••••**************************
ł	000260		02724873					002089A0		** *
1	000280		700247E0					00004100		*M*
	C002A0		43A7C00A 2004501B					20082008 02544700		**
1	0002C0 2002E0		58E00FC8					05901211		*HKO
1	003220		47100352					D300100C		*K.
1	000320		40105088					47708008		*
	000340		91107006					47800308		*
1	000360		50A0C024					5090A000		*
	000380		471003E6					48400044		*WKU*
1	000340		58AC7030					478003C0		**
1	000300		478003DA					701850A0		*
1	C003E0		06249104			-		05929101		*0
	000400		20019110			70064710	04D241A0	703140A0	50 3AD 203	*X
	000420	50007031	91012000	47100432	D2077030	20201BAA	43A70030	89A00004	41AA3020	*
	000440	91082001	47100490	05037033	A0064740	053ED503	7033A00A	4720053E	91027013	*
	000460	471004D2	D5017031	A0044770	053E9104	30084780	04900501	70354008	4740048A	*KN
1	000480	D5017035	A00C47C0	04900201	7035A008	41A05038	41800578	45C 00 5E 8	47700688	*N K
1	000440	90006000	47BC04A0	48A00044	54405058			D2062009		* K*
1	000400		8008945F					89B00002		*••••••••••••
	0004E0	4BC 05 096	41AC7031	40A0C002	43B07030	89800004	43983020	4290C00D	D202C 011	*K*
	000500		20004780					9618C00D		*KK*
1	000520		47800530					700647F0		*0
	00054 0		000005EF					47F00564		*0
1	000560		94FE7C06					D5037031		*•••••••*•••••••••
1	000580		70064060					CDD69148		*
1	000540		70064780					C40E9110		*
1	000500		00189101					50082018		*•D••••••*
1	000550		41B00624					47800624		`••••••••••••••••••••••••••••••••
1	000600		58F3001C					096447F0		,.3*
1	000620		92000048					03000048		······*
1	000640		88A00018					40107014		· · · · · · · · · · · · · · · · · · ·
	000660		4770068E					700C509A		*
	000680		075207F8 471006E4					A00850A0 00444780		*••••••8K••••••••••••••••• K•* *••••••U••••
	0006A0 0006C0		00444710					47808008		*••••••••••••*************************
1	000000	10007120	00444710	5004 9008	10003140	00440115	1040044		41000004	***************************************

Figure 27. Sample of General Format Dump

129

OUTPUT COMMENTS

The following output comments are printed within the body of a formatted dump whenever IMDPRDMP is unable to locate, format and print a control block. These comments explain why the referenced control block is not printed within the dump listing, these output comments are separated from the main storage information by a blank line both before and after each output comment. Note: Output Comments produced when IMDPRDMP is processing GTF output are shown in Section 3 of this publication under the heading 'IMDPRDMP Output Comments - GTF Processing'.

DUPLICATE PREFIX FOLLOWS - ID 'A'

Explanation: While processing a dump from a Model 65 multiprocessing system, IMDPRDMP has determined that the CPU prefixes (CPUIDs) are the same. If the task that performs the dump is initiated on one CPU, interrupted, dispatched to the second CPU, and completed the rest of the processing on the second CPU, then the prefix shown in this output comment is that of the first CPU to which the task was assigned. Processing continues.

END OF FILE ON DUMP TAPE

Explanation: While trying to locate a block of main storage on the dump tape, IMDPRDMP reached the end of the tape. This message is printed only if IMDPRDMP is either trying to extract the CVT pointer or trying to extract an area of storage for printing.

Processing terminates. If IMDPRDMP did no formatting and the tape does not contain a low-speed dump produced by IMDSADMP, the job may be rerun using the CVT control statement to direct IMDPRDMP to the CVT in this dump. Low-speed dumps produced by IMDSADMP can not be formatted by IMDPRDMP.

ERROR FINDING PARTITION BOUNDARIES FOR TCB aaaaaa

Explanation: IMDPRDMP found (1) a TCB family chain pointer, (2) a partition boundary box pointer, or (3) a pointer within the partition boundary box that was one of the following:

 Addressed an area that was not on a full-word boundary.
 Addressed an area that was higher than the highest address in the dump.

3. Could not be extracted from the dump because either an I/O error was encountered while attempting to read the block containing the pointer or the block containing the pointer was missing from the dump. A possible cause for a missing block is that the routine that produced the dump encountered an I/O error while attempting to write the block.

Processing continues.

ERROR FINDING REGION BOUNDARIES FOR TCB aaaaaa.

Explanation: While attempting to determine the region boundaries for the family of TCBs attached to the job step TCB at address aaaaaa, one of the following conditions occurred:

- IMDPRDMP found a chain with more than fifty partition queue elements (PQEs); or,
- IMDPRDMP found (1), a TCB family chain pointer, (2) a partition queue element (PQE) pointer (TCB + X'98'), or (3) a pointer within a PQE that:
 - Addressed an area that was not on a word boundary; or,
 - 2. Addressed an area that was higher than the highest address in the dump; or,
- 3. Could not be extracted from the dump decause either an I/\tilde{O} error was encountered while attempting to read the block containing the pointer or the block containing the pointer was missing from the dump; a possible cause for a missing block is that the routine that produced the dump encountered an I/O error while attempting to write the block. Processing continues.

ERROR FORMATTING TCB

Explanation: One of the fields in the TCB required for formatting could not be extracted from the dump because:

• IMDPRDMP encountered an I/O error while attempting to read the block that contains the required data; or, • The block containing the required data was missing from the dump; a possible cause is that the routine that produced the dump encountered an I/O error while attemtping to write the block. Processing continues.

ERROR IN DEB CHAIN

Explanation: The routine that formats the data extent block (DEB) found one of the following errors:

- A DEB chain pointer:
 - 1. Was not on a word boundary; or,
 - 2. Addressed an area of main storage higher than the highest address in the dump; or,
- The address of the DEB was invalid causing the address of the DEB prefix (DEB - 16) to be zero or negative; or,
- A DEB chain pointer or one of the fields necessary to format the DEB could not be extracted from the dump because:
 - IMDPRDMP encountered an I/O error attempting to read the block that contained the data; or,
 - 2. The block containing the data was missing from the dump; a possible cause is that the routine that produced the dump encountered an I/O error while attempting to write the block.

Processing continues.

ERROR IN EXTENT LIST

Explanation: While formatting the load list or job pack area of an MVT dump, IMDPRDMP encountered a contents directory entry (CDE) that had a block extent list with a relocation factor (extent list + 4) of zero or greater than twenty-five. A relocation factor of zero is an error; however, a value greater than twenty-five can be valid. The value of twenty-five was established by IMDPRDMP as a reasonable limit; it is improbable that a normal task would have a program that has more than twenty-five CSECTs causing it to get an extent list with a relocation factor greater than twenty-five. Processing continues with the next CDE. ERROR IN JOB PACK QUEUE

Explanation: The routine that formats the job pack area encountered one of the following errors:

- A job pack queue chain pointer addressed an area that:
 - 1. Was not on a word boundary, or,
 - Was greater than the highest address in the dump.
- A job pack queue chain pointer or one of the fields in a job pack area control block could not be extracted from the dump because:
 - IMDPRDMP encountered an I/O error attempting to read the block containing the needed data, or,
 - 2. The block containing the needed data was missing from the dump; a possible cause is that the routine that produced the dump encountered an I/O error while attempting to write the block.

Processing continues.

ERROR IN LOAD LIST

Explanation: The load list print routine encountered one of the following errors:

- A pointer in the load list control block chain referenced an area of main storage that:
 - 1. Was not on a word boundary, or
 - 2. Was greater than the highest address in the dump.
- A field in a load list queue control block could not be extracted from the dump because:
 - IMDPRDMP encountered an I/O error attempting to read the block that contained the data needed to format the load list; or,
 - 2. The block containing the data was missing from the dump; a possible cause is that the routine that produced the dump encountered an I/O error while attempting to write the block.

Processing continues.

ERROR IN TCB CHAIN DURING PRINT ALL FUNCTION

Explanation: The PRINT ALL routine encountered a TCB chain pointer that was not on a full-word boundary or addressed an area that could not be extracted from the dump decause:

- The pointer addressed an area higher than the highest address in th dump.
- 2. IMDPRDMP encountered an I/O error while trying to read the record that contained the area addressed by the pointer.
- 3. The record containing the addressed area was missing from the dump, possibly because the routine that produced the dump encountered an I/O error while attempting to write the record.
 No more TCBs are processed, but all tasks encountered up to the point are printed. The last TCB printed may contain the error described in 1 above or point to the unavailable area described in 2 and 3 above.

ERROR IN TCB CHAIN TCB aaaaaa

Explanation: The routine that formats the TCBs encountered one of the errors given below; the address of the TCB associated with the error replaces the aaaaaa field of the output comment.

- A TCB pointer for one of the TCBs in the TCB family chain addressed an area not on a word boundary; or,
- A TCB pointer or the TIOT pointer in the TCB at location aaaaaa points to an area that could not be extracted from the dump because:
 - IMDPRDMP encountered an I/O error while attempting to read the block that contains the pointer; or,
 - 2. The routine that produced the dump encountered an I/O error while writing the block that contains the pointer; therefore, the block is missing from

the dump. ERROR IN TIOT

Explanation: The format routine found one of the following errors:

- The task input output table (TIOT) pointer (TCB + X'C') was not on a word boundary; or,
- One of the fields required to format the TIOT could not be extracted from the dump because:
 - IMDPRDMP encountered an I/O error while attempting to read the block that contains the required data, or,
 - 2. The block containing the required data was missing from the dump; a possible cause is that the routine that produced the dump encountered an I/O error while attempting to write the block.

ERROR WHILE FORMATTING CONTROL BLOCKS CONTINUING

> Explanation: While building a list of job step TCB's for all partition regions in the dump data set, PRDMP encountered one of the following conditions:

- One of the TCB chain pointers was greater than the highest address in the dump.
- 2. One of the TCB chain pointers addressed an area that was missing from the dump data set.

PRDMP will attempt to use the partial list and continue with its formatting.

ERROR WHILE FORMATTING PSCB

Explanation: One of the following errors occurred while IMDPRDMP was formatting the protected step control block (PSCB):

- The address of the PSCB in the time sharing job block extension (TJEX) was greater than the highest main storage address in dump; or,
- An I/O error occurred while reading the block of dump information that contained the needed data; or,
- A block of dump information containing part of the PSCB was not found on either the dump or swap data sets.

Processing continues. IMDPRDMP attempts to format the control blocks for the next TSO user.

ERROR WHILE FORMATTING RCB

Explanation: One of the following errors occurred while IMDPRDMP was formatting the time sharing region control blocks (RCBs):

• The address of the RCB table in the time sharing

communication vector table (TSCVT) was greater than the highest main storage address in the dump; or,

- An I/O error occurred while reading the block of dump information that contained the needed data; or,
- A block of dump information containing part of an RCB was not found on the dump data set. This happens when an I/O error occurred while the dump routine was writing the data onto the dump data set.
 Processing continues. IMDPRDMP attempts to format the next entry in the RCB table.

ERROR WHILE FORMATTING SWAP CONTROL BLOCK

Explanation: One of the following errors occurred while IMDPRDMP was formatting the swap control block (SWAP DCB):

• The address of the SWAP DCB in the time sharing communication vector table (TSCVT) was greater than the highest main storage address in the dump; or,

132.2 Programmer's Guide to Debugging (Release 21.7)

- An I/O error occurred while reading the block of dump information that contained the data; or,
- A block of dump information containing part of the SWAP DCB was not found on the dump data set. This happens when an I/O error occurred while the dump routine was writing the data onto the dump data set.

Processing continues. IMDPRDMP attempts to continue formatting the time sharing job block (TJB).

ERROR WHILE FORMATTING TAXE

Explanation: One of the following errors occurred while IMDPRDMP was formatting the terminal attention exit element (TAXE):

- The address of the TAXE in the time sharing job block extension (TJBX) was not aligned in a fullword boundary; or,
- The address of the TAXE in the TJBX was greater than the highest main storage address in the dump; or,
- An I/O error occurred while reading the block of dump information that contained the needed data; or,
- A block of dump information containing part of the TJBX was not found on the dump or swap data sets.

Processing continues. IMDPRDMP attempts to format the control blocks for the next TSO user.

ERROR WHILE FORMATTING TJB

Explanation: One of the following errors occurred while IMDPRDMP was formatting the time sharing job block (TJB):

- The address of the TJB table in the time sharing communication vector table (TSCVT) was greater than the highest main storage address in the dump; or,
- An I/O error occurred while reading the block of dump information that contained the needed data; or,

• A block of dump information containing part of the TJB was not found on the dump data set. This happens when an I/O error occurred while the dump routine was writing the data onto the dump data set.

Processing continues. IMDPRDMP attempts to format the next active TJB.

ERROR WHILE FORMATTING TJBX

Explanation: One of the following errors occurred while formatting the time sharing job block extension (TJEX):

- The terminal job block (TJB), that contained the address of the TJB was not aligned on a fullword boundary; or,
- The address of the TJBX in the TJB was greater than the highest address in the system dump; or,
- An I/O error occurred while reading the block of dump information that contained the needed data; or,
- A block of dump information containing part of the TJBX was not found on either the dump or swap data sets.

Processing continues. IMDPRDMP attempts to format the control blocks associated with the next TSO user.

ERROR WHILE FORMATTING TSB

Explanation: One of the following errors occurred while IMDPRDMP was formatting the terminal status block (TSB):

- The address of the TSB table in the time sharing communication vector table (TSCVT) was greater than the highest main storage address in the dump; or,
- An I/O error occurred while reading the block of dump information that contained the needed data; or,
- A block of dump information containing part of the TSB was not found on the dump data set. This happens when a I/O error occurred while the dump routine was writing the data onto the dump data set.

Processing continues. IMDPRDMP attempts to format the associated time sharing job block extension (TJBX).

ERROR WHILE FORMATTING TSCVT

Explanation: One of the following errors occurred while IMDPRDMP was formatting the time sharing communication vector table (TSCVT):

- The address of the TSCVT in the communication vector table (CVT) was greater than the highest main storage address in the dump; or,
- An I/O error occurred while reading the block of dump information that contained the needed data; or,
- A block of the dump information containing part of the TSCVT was not found on the dump data set. This happens when a I/O error occurred while the dump routine was writing the data onto the dump data set.

Processing continues. IMDPRDMP attempts to format the time sharing region control blocks (RCBs).

ERROR WHILE FORMATTING USER MAIN STORAGE MAP

Explanation: One of the following errors occurred while IMDPRDMP was formatting the user main storage map (UMSM):

- The address of the UMSM associated with the region control task (RCT) or time sharing job block (TJB) was greater than the highest main storage address in the dump; or,
- An L/O error occurred while reading the block of dump information that contained needed data; or,
- A block of dump information containing part of the UMSM was not found in the dump data set. This happens when an I/O error occurred while the dump routine was writing the data onto the dump data set.

Processing continues. IMDPRDMP attempts to continue formatting with the terminal status block (TSB).

FORMAT ERROR DURING TCB SUMMARY

Explanation: The routine that prints the TCB summary must extract a TCB completion code (TCB + X'16') or a TCB family chain pointer from the dump. In this case, IMDPRDMP was unable to do so because:

- IMDPRDMP encountered an I/O error while attempting to read the block containing the completion code or pointer; or,
- The block containing the completion code or pointer was missing from the dump; a possible cause is that the routine that produced the dump encountered an I/O error while attempting to write the block.

Processing for the current control statement is terminated.

FORMAT ERROR IN MAIN STORAGE BLOCKS

Explanation: While formatting main storage control blocks, IMDPRDMP encountered one of the following errors:

- A pointer in a main storage control block addressed an area that:
 - 1. Was not on a word boundary; or,
 - Was higher than the maximum address in the dump; or,
- One of the fields in a main storage control block could not be extracted from the dump because:
 - IMDPRDMP encountered an I/O error while attempting to read the block that contains the required field; or,
 - 2. The block containing the required field is missing from the dump; a possible cause is that the routine that produced the dump encountered an I/O error while attempting to write the block.

Processing continues.

INFINITE LOOP IN DEB CHAIN

Explanation: While formatting the data extent blocks (DEBs), IMDPRDMP found more than 200 DEBs chained to the TCB. The limit of 200 DEBs prevents IMDPRDMP from looping. When the limit is exceeded, a loop is assumed which causes this comment to be printed. Processing continues after the first 200 DEBs are printed.

INFINITE LOOP IN JOB PACK QUEUE

Explanation: In MVT, IMDPRDMP found more than 255 CDEs on the job pack queue associated with the TCB. In MFT, IMDPRDMP found more than 255 RBs on the job pack queue associated with the TCB. A limit of 255 job pack queue control blocks has been established by IMDPRDMP to prevent a possible looping condition. When the limit is exceeded, a loop is assumed and this comment is issued. The first 255 job pack queue control blocks are printed and then processing continues.

INFINITE LOOP IN LOAD LIST

Explanation: In MVT, IMDPRDMP found more than 255 load list elements (LLEs) on the load list chained to the TCB. In MFT, IMDPRDMP found more than 255 RBs on the load list. A limit of 255 elements on the load list has been established by IMDPRDMP to prevent a possible looping condition. When the limit is exceeded, a loop is assumed and this comment is issued. The first 255 load list elements are printed and then processing continues.

INFINITE LOOP IN PQES

Explanation: The main storage print routine found more than 50 partition queue elements (PQEs) chained to the TCB. A limit of 50 PQEs has been established by IMDPRDMP to prevent a possible looping condition. When the limit is exceeded, a loop is assumed and this comment is issued. The first 50 PQEs are printed and then processing continues.

INFINITE LOOP IN QCB CHAIN

Explanation: The QCB print routine found more than 100 queue control blocks (QCBs) on the major QCB chain or more than 100 QCBs on a minor QCB chain. A limit of 100 QCBs has been established within IMDPRDMP to prevent a possible looping condition. When the limit is exceeded, a loop is assumed and this comment is issued. The first 100 QCBs are printed and then processing continues.

INFINIT LOOP IN QEL CHAIN

Explanation: The QCB print routine found more than 100 queue elements (QELs) on a QEL chain. A limit of 100 QELs has been established within IMDPRDMP to prevent a possible looping condition. When the limit is exceeded, a loop is assumed and this comment is issued. The first 100 QELs are printed and then processing continues.

INFINITE LOOP IN RB CHAIN

Explanation: The RB print routine found more than 50 request blocks (RBs) on the RB chain. A limit of 50 RBs has been established within IMDPRDMP to prevent a possible looping condition. When the limit is exceeded, a loop is assumed and this comment is issued. The first 50 RBs are printed and then processing continues.

INFINITE LOOP IN TIOT

Explanation: While formatting the TIOT, IMDPRDMP found more than 255 DD entries in the TIOT. Since only 255 DD statements are allowed per job step, a valid TIOT cannot have more than 255 DD entries. A loop is assumed because the limit was exceeded and the TIOT end indicator could not be found. Processing continues.

INVALID TIOT

Explanation: While formatting the task input output table (TIOT), the FORMAT routine found an invalid job name in the TIOT. To be valid, the first character of the job name must be A through Z, or \$, #, @ or a blank (X'40'). Processing continues.

NO ELEMENTS ON LOAD LIST

Explanation: The load list pointer in the TCB (displacement X'24') is zero. The zero pointer indicates that (1) no programs were loaded by the LOAD macro instruction or (2) the load list pointer was overlaid with zero. Processing continues.

NO EXTENT LIST

Explanation: While formatting the load list and job pack queue for an MVT dump, IMDPRDMP encountered zeros in the extent list pointer (CDE + '20') in a major contents directory entry (CDE). This zero pointer usually indicates an error condition in which the extent list pointer was overlaid with zeros. Processing continues with the next CDE.

NO LINK PACK AREA QUEUE

Explanation: In MFT, an LPAMAP was requested but the link pack area queue pointer (CVT + X'BC') was zero. Processing continues.

NO MAJOR QCBS

Explanation: The QCB TRACE routine found zeros as the pointer to the first major queue control block (QCB). This indicates that no resources have been enqueued at the time of the dump or that the pointer to the QCB queue has been overlaid with zeros. Processing continues.

NO RSVC MODULES IN LPA RAM

> Explanation: In MVT, either the CDE or LLE chain, pointing to the modules in the link pack area, was empty. RSVC modules are found by following the LLE pointer in the master scheduler TCE (TCBLLEP). Rpocessing continues.

RAM modules are found by following the CDE chain from the CVT (CVTQLPAQ). Processing for this IMDPRDMP verb ends.

NOTHING IN JOB PACK

Explanation: In MVT, the job pack queue field of the TCB (TCB + X'2C') is zero. In MFT, the partition information block (PIB) field (TCB + X'7C') or the job pack queue pointer (PIB + X'24') is zero. PCP does not have a job pack pointer; therefore, this comment does not appear in a PCP dump. A zero job pack queue pointer is usually a normal condition, especially for a system task. Processing continues.

RB FORMAT ERROR

Explanation: While formatting a request block (RB), the RB print routine found that the request block (RB) chain pointer addressed an area of main storage that:

- Was not on a word boundary; or,
- Was higher than the highest address in main storage; or,
- Could not be extracted from the dump because:
 - IMDPRDMP encountered a I/O error while attempting to read the block that contained the pointer; or,
 - 2. The block that contained the pointer was missing from the dump. One possible cause for this is that the program that produced the dump may have encountered an I/O error while writing the block.
- A field in the RB, or a contents directory entry (CDE) associated with the RB, necessary to formatting the RB could not be extracted from the dump. Either IMDPRDMP encountered an I/O error while trying to read the block, or the block that contained the pointer is missing from the dump.

REGISTERS FROM OTHER CPU ARE INVALID-NOT FORMATTED

Explanation: Multiprocessing systems only. Only the registers for the CPU in which the dump program was executed will be displayed on the dump listing. This can occur when the dump is taken on a multiprocessing system either when the NOMP option of IMDSADMP is used or when the direct control feature is not operational.

TASK HAS NO OPEN DATA SETS

Explanation: IMDPRDMP found the data extent block (DEB) pointer in the TCB (TCB + X'8') to be zero. This situation indicates that there were actually no open data sets or the DEB pointer in the TCB was overlaid with zeros. Processing continues.

TASK HAS NO TIOT

Explanation: While attempting to format the task input output table (TIOT), IMDPRDMP found that the TIOT pointer (TCB + T'C') was either zero or larger than the highest address in the dump. The zero TIOT pointer could be a normal condition for a system communication task, but for a problem program task this is an error condition. Processing continutes.

Explanation: After formatting a TCB, this comment is printed below the TCB if the first bit (the terminated bit) of the flag byte at X'21' of the TCB is set. Processing continues with the next TCB.

TCB CHAIN ERROR IN F03 PRINT ROUTINE TCB aaaaaa...CONTINUING WITH NEXT TCB

> Explanation: The Print F03 routine encountered a TCB chain pointer that:

- Was not on a word boundary; or,
- Addressed an area that could not be extracted from the dump because:
 - The pointer addressed an area higher than the maximum address in the dump; or,

- IMDPRDMP encountered an I/O error trying to read the record containing the area addressed by the pointer; or,
- 3. The block containing the addressed area was missing from the dump, probably because the routine that produced the dump encountered an I/O error while attempting to write the block.

The address of the TCB associated with the error replaces the aaaaaa field in the message. Processing continues with the next TCB.

UNABLE TO OBTAIN JOB STEP TCB FOR CURRENT TASK

Explanation: In an MFT system with ATTACH, the job step TCB is used to get the partition boundaries. If the address of this TCB is zero in the current TCB, then the above comment is issued. Processing continues.

XXXXXX THROUGH XXXXXX COULD NOT BE PRINTED

Explanation: The block(s) of storage between the two addresses could not be printed because IMDPRDMP encountered an I/O error while attempting to read the block(s) or the block(s) that were missing from the dump. Processing continues.

TASK HAS TERMINATED

Page of GC28-6670-5,6, Revised April 16, 1973, By TNL: GN28-2545

136.2 Programmer's Guide to Debugging (Release 21.7)

Guide to Storage Dumps

The purpose of this section is to suggest debugging procedures that you may use with a storage dump. This discussion applies to the output of the following programs:

- IMDSADMP- The low speed version that creates an unformatted dump of main storage.
- IMDPRDMP- Reads, formats, and prints storage dumps from MFT or MVT systems and the high speed version of IMDSADMP.

These programs produce hexadecimal dumps of the contents of main storage from location zero to the highest machine address.

The IMDPRDMP program provides formatting capabilities which can be used to display the important system control blocks for easy examination. The IMDPRDMP program does most of the procedures described in this section automatically. The cases in which the IMDPRDMP program does not provide formatting are identified. A complete description of the services provided by the IMDPRDMP program is found in the publication, <u>IBM System/360 Operating</u> System: Service Aids, GC28-6719.

Since the formatting for the IMDPRDMP program depends on the contents of the dump, it is not always possible to provide complete formatting. For example, if the CVT of the system to be dumped has been overlaid, the IMDPRDMP program can provide only a hexadecimal dump of main storage.

DETERMINING THE CAUSE OF THE DUMP

Main storage dumps are invoke by system routines and these routines can be identified by module names appearing in the most recent request block (RB) for the failing task. The main storage dump is invoked by SVC 51. This SVC PSW appears as the resume PSW in the second most recent RB of some task in the system. The module name in the current RB for that task must be 201C.

Main storage locations from zero to 128 (hexadecimal 80) are permanently assigned and contain hardware control words. Figure 28 shows these fields, their location, their length, and their purpose.

[
		Length In Bytes	Purpose	
0	0	8	IPL PSW	
8	8	8	IPL CCW1	
16	10	8	IPL CCW2	
24	18	8	External old PSW	
32	20	8	Supervisor call old PSW	
40	28	8	Program old PSW	
48	30	8	Machine check old PSW	
56	38	8	I/O old PSW	
64	40	8	Channel Status Word	
72	48	4	Channel Address Word	
76	4C	4	Unused	
80	50	4	Timer	
84	54	4	Unused	
88	58	8	External new PSW	
96	60	8	Supervisor call new PSW	
104	68	8	Program new PSW	
112	70	8	Machine check new PSW	
120	78	8	I/O new PSW	
Figure 28. Permanently Assigned Hardware				

Figure 28. Permanently Assigned Hardware Control Words

<u>Cause of the Dump</u>: Evaluate the PSWs that appear in the formatted section of the dump (the first four lines) to find the cause of

the dump. The PSW has the following format:

Program Status Word

	System	n Mask	Key	AMWP	Interruption Code	
0		7	8 11	12 15	16	31
ILC	сс	Program Mask		İnstr	ruction Address	
32 33	34 35	36 39	40			63

- Does the instruction address field of the old machine check PSW show either the value E2 or E02? If so, a hardware error has occurred.
- Does the instruction address field of the old program check PSW have a value other than zero? If so, a program check at the instruction preceding that address caused the interruption.

TASK STRUCTURE

MFT_System (Without Subtasking)

There is a TCB associated with each partition of main storage there are also TCBs for critical system tasks such as the master scheduler task and the transient area loading task. Figure 28 shows location 76 (4C) unused for hardware control words. The control program uses this word to contain a pointer to the CVT. Use this CVT pointer to locate the first byte of the CVT, then the CVTIXAVL field (offset 124) in the CVT. The address contained at CVTIXAVL is a pointer to the IOS freelist. At offset 4 in the IOS freelist is a pointer to the first address in a list of TCB addresses. You can look through this list of TCB addresses, and, keeping your system options in mind, find the TCBs for each partition. The TCB addresses are listed in the following order:

- Transient area loading task.
- System error task (MFT with subtasking).
- Multiple console support write-to-log task (optional).
- I/O recovery management support task (optional).
- Communications task.
- Master scheduler task.
- System management facilities task (optional).
- Partition 0 task.

- Partition 1 task.
- •
- Partition n task.

Figure 29 shows how to locate the partition TCBs in sample output from the IMDPRDMP program.

MFT System (With Subtasking)

For MFT with subtasking (and for MVT), a task may create a subtask. The partition TCBs for MFT with subtasking are referred to as job step TCBs. The task structure for a job step may be reconstructed in a main storage dump by using the information in Figure 30.

For MFT with subtasking, the job step TCB may be found using the method described for MFT without subtasking or by a more direct method. CVT offset 245 (F5) contains a pointer to the partition 0 job step TCB address in this address table.

To recreate the task structure within any partition, simply locate the job step TCB, and follow the TCB pointers - as explained in the previous section.

MVT System

To find the current TCB, look at location 76 (4C) for a pointer to the CVT. The first word of the CVT contains a pointer to a doubleword of TCB addresses, which contains pointers to the next TCB to be dispatched (first word) and the current TCB (second word). Beginning with the current TCB, you can recreate the task structure for the job step using the methods in Figure 30.

If the first word of the current TCB points to itself, there are no ready tasks to be dispatched, and the system has been placed in an enabled wait state. This TCB, now in control, is called the <u>system wait</u> <u>TCB</u>.

All TCBs in the system are maintained in a queue called the CVT ready queue. These TCBs are queued according to their dispatching priority. The CVTHEAD field, offset +160 (A0) in the CVT, contains the address of the highest priority TCB in the system. Offset +116 (74) in the TCB points to the TCB with the next lowest priority. Figure 30 shows how to locate all of the TCBs in the system.

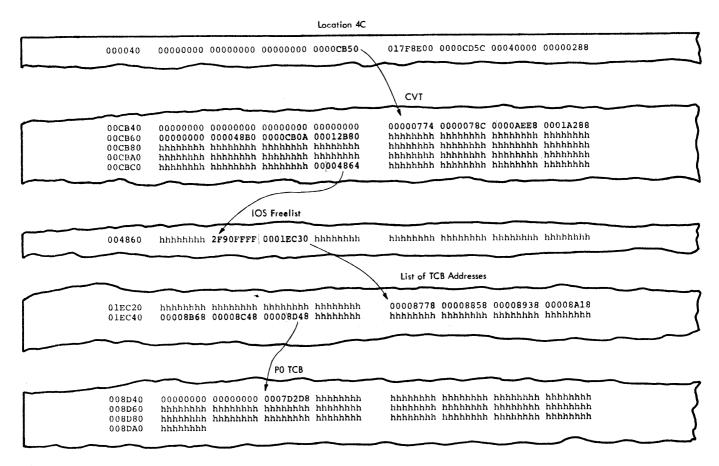
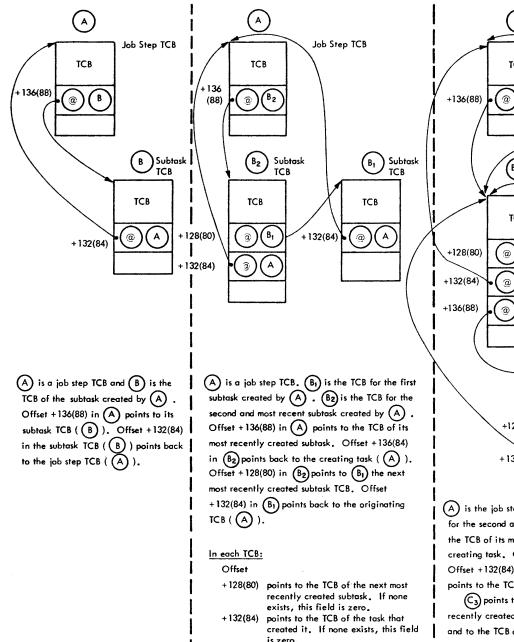
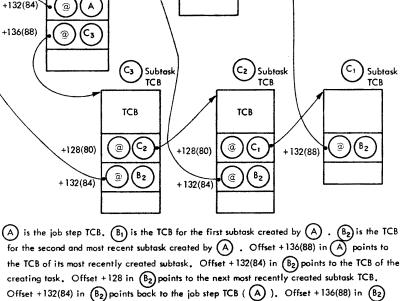


Figure 29. Finding the Partition TCBs in MFT





(B,

TCB

@ A

Subtask TCB

Α

TCB

B2

TCB

B,

Subtask

TCB

Job Step TCB

points to the TCB of its most recently created subtask ((C_3)). (C_3) points to the TCB of its creating task (B_2) and to the TCB of the subtask most

recently created by (B_2) . (C_2) contains pointers to the TCB of the originating task ((B_2)) and to the TCB of the task most recently created by (B_2) . (C_2) contains only a pointer to the TCB of the invoking task $((B_2))$.

- is zero. +136(88) points to the TCB of the most recent subtask created by this task. If none exists, this field is zero.
- Figure 30. Finding the TCB

Keep in mind that all TCBs in the system appear on this queue. Therefore, not only does a particular job step TCB appear on the ready queue, but all of its subtask TCBs also appear.

You can find the job step TCB associated with any TCB by using the TCBJSTCB field of the TCB, offset +124 (7C). This field contains the address of the job step TCB for the TCB you are examining.

In response to the FORMAT control statement, the IMDPRDMP program will do most of this work for you. It will recreate the task structure, format all TCBs in the system, and provide a TCB summary. The TCB summary shows the task structure. Figure 31 shows a portion of the TCB summary information from an MVT system. TCBs associated with a particular job are grouped together under the job name | and step name. The TCB summary contains the TCB address, the completion code, and, when applicable, the address of the originating TCB and the addresses of created TCBs.

TASK STATUS - ACTIVE RB QUEUE

The first word of the TCB contains a one-word pointer to the first word of the most recent RB added to the queue. In its eighth word, RB+28(1C), each RB contains a pointer to the next most recent RB. The last RB points back to the TCB.

You can determine the idenity of the load module by looking either in the first and/or second words of the RB for its EBCDIC name or in the last 3 digits of the resume PSW in the previous RB for its SVC number. The entry point to the module is in the last 3 bytes of the fourth word in the RB, RB-13(D).

In an MVT system, the name and entry point of the associated load module are not always contained in the RB associated with the module. Instead, they are found in a contents directory entry (CDE).

The address of the contents directory entry for a particular load module is given in the fourth word of the RB, RB+12(C). The CDE gives the address of the next entry in the directory (bytes 1-3), the name of the load module, bytes 8-15(F); the entry points of the module, bytes 17-19(11-13).

Figure 31 shows the formatting that the IMDPRDMP program does for a task in an MVT system. Notice the connection between the RB and the CDE. The IMDPRDMP program extracts the CDE information and displays this information with the RB.

The wait-count field of the RB is particularly important when locating the TCB by using the CVT ready queue (CVTHEAD). The high-order byte of the RB link field, RB-28(1C), of the most recent RB for a TCB contains a count of the number of events for which the task is waiting. Tasks that have a zero wait count are ready to be dispatched (unless marked non-dispatchable). Such a task will be dispatched or become the current task when all TCBs of higher priority are waiting for

		* * * * T	CB SUMM2	A R Y * * * *		
JOB	MASTER TCBhhhhhh	STEP SCHEDULER CMPhhhhhhh	NTChhhhhhh	OTChhhhhhhh	LTChhhhhhhh	PAGE hhhh
JOB	MASTER TCBhhhhhh	STEP SCHEDULER CMPhhhhhhh	NTChhhhhhhh	OTChhhhhhhh	LTChhhhhhh	PAGE hhhh
JOB		STEP 00E CMPhhhhhhh CMPhhhh h hhh	NTChhhhhhhh NTChhhhhhhh	OTChhhhhhh h OTChhhhhhhh	LTChhhhhhhh LTChhhhhhhh	PAGE hhhh PAGE hhhh
JOB	TCBhhhhhh	STEP GO CMPhhhhhhh CMPhhhhhhhh CMPhhhhhhhh	NTChhhhhhh NTChhhhhhhh NTChhhhhhhh	OTChhhhhhhh OTChhhhhhhh OTChhhhhhhh	LTChhhhhhhh LTChhhhhhhh LTChhhhhhhh	PAGE hhhh PAGE hhhh PAGE hhhh
JOB	JOB12 TCBhhhhhh	STEP GO CMPhhhh h hh	NTChhhhhhhh	OTChhhhhhhh	LTChhhhhhhh	PAGE hhhh

Figure 31. IMDPRDMP TCB Summary

the completion of an event. To determine the events for which a task is waiting, use the instruction address field in the resume PSW to locate the WAIT macro instruction in the source program. This will point you to the operation being executed at the time of the dump.

MAIN STORAGE CONTENTS

Load List (MFT)

The load list is a chain of request blocks associated with load modules invoked by a LOAD macro instruction. By looking at the load list, and at the job pack area queue described below, you can determine which system and problem program routines were loaded before the dump was taken. construct the load list associated with the task in control, look at the tenth word in the TCB, TCB+36(24), for a pointer to the most recent RB entry on the load list, minus 8 bytes (RB-8). This word, in turn, points to the next most recent entry (minus 8), and so on. If this is the last RB, RB-8 will contain zeroes. The word preceding the most recent RB on the list (RB-4) points back to the TCB's load list pointer.

Load List (MVT)

To construct the load list associated with the task in control, look at the tenth word in the TCB, TCB+36(24), for a pointer to the most recent load list entry (LLE). Each LLE contains the address of the next most recent entry (bytes 0-3), the count (byte 4), and the address of the CDE for the associated load module (bytes 5-7). If this is the last LLE in the list, TCB+36(24) will contain zeroes.

Job Pack Area Queue (MFT With Subtasking, MVT)

In systems with MFT with subtasking or MVT control programs, the job pack area queue is used to maintain reenterable modules within a partition or region. The complete description of this queue is found under the topic "Task Status-Active RB Queue".

MFT System: To reconstruct the job pack area queue in an MFT system with subtasking, look at TCB+125(7D) for a three byte pointer to the partition information block (PIB). The twelfth word of the PIB, PIB+44(2C), points to the most recent RB on the job pack area queue minus 8 bytes (RB-8). This word in turn points to the next most recent RB minus 8, and so on. The last RB will have zero in this field. The word preceding the most recent RB on the queue (RB-4) points back to the job pack area queue pointer in the PIB. You can determine the identity of the load module by looking either in the first and/or second word of the RB for its EBCDIC name, or in the last three digits of the resume PSW in the previous RB for the SVC number. The entry point of the module is given in the last three bytes of the fourth word in the RB, RB+29(1D), unless it is an FRB.

The first five words of an FRB (beginning at offset minus 8) are identical in content to those of other RBs. The XRWTL field, offset 12(C), contains the address of a wait list element. The first word of the WLE points to the next WLE, or

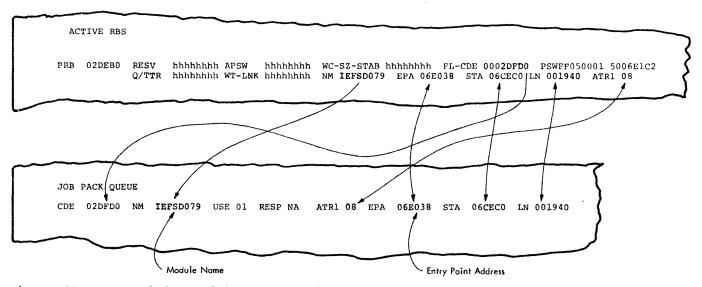


Figure 32. Determining Module From CDE in MVT

142 Programmer's Guide to Debugging (Release 21.7)

contains zeros if the WLE is the last one. The second word points to the waiting SVRB. You can determine the number of deferred requests for the module by tracing the chain of WLES.

The XRREQ field of an FRB, offset 16(10), contains a pointer to the TCB of the requesting task. The next word, CRTLPRB, offset 20(14), points to an LPRB built by the Finch routine for the requested program. The FRB for the requested program is removed from the job pack area queue by the Finch routine when the program is fully loaded.

<u>MVT System</u>: In MVT, the job pack area queue is maintained in the same manner as the load list. The distinction between the two queues is that the job pack area queue contains reenterable programs. There are no FRBs in MVT.

MAIN STORAGE SUPERVISION

Free Areas in MFT Systems

Areas of main storage that are available for allocation at the time the dump was taken are described by the MSS boundary box and a series of free queue elements (FQEs). The seventh word of the TCB for the task, TCB+24(18), points to a six-word MSS boundary box. The first word of the MSS boundary box points to the FQE with the highest processor storage address in the partition (hierarchy 0), and the fourth word, to the highest 2361 Core Storage address in the partition (hierarchy 1). The first word of each FQE points to the next lower FQE; the second word of the FQE gives the length of the area it describes. FQEs occupy the first 8 bytes of the area they describe.

Gotten Subtask Areas (MFT)

In MFT with subtasking, areas of a partition allocated by the system to a subtask within the partition are described by gotten subtask area queue elements (GQEs). The seventh word of the subtask TCB, TCB+24(18), points to a one word pointer to the most recently created GQE on the GQE queue. Bytes 0 through 3 of the GQE contain a pointer to the previous GQE or, if zero, indicate that the GQE is the last one on the queue. Bytes 4 through 7 of the GQE contain the length of the gotten subtask area. Each GQE occupies the first eight bytes of the gotten subtask area it describes.

Region Structure in MVT System

The region associated with a particular task in an MVT system is described by

partition queue elements (PQEs). The thirty-ninth word of the TCB, offset +152 (98) contains a pointer to the dummy PQE (D-PQE) for the region. The first word of the dummy PQE points to the first PQE and the second word, to the last PQE. The first and second words of each PQE point to the first and last free block queue elements (FBQEs), respectively, associated with the PQE. Separate PQEs are used to describe parts of a region in different storage hierarchies or part of a region that was obtained by another task which has been rolled out.

FBQEs describe free areas in the region that have a a length which is a multiple of 2048 bytes. These free areas are available for allocation to a specific subpool.

Subpool Descriptions (SPQEs) (MVT): The seventh word of the TCB, TCB+24(18), points to the SPQE representing the first subpool used by the task. Each SPQE contains the address of the next SPQE (bytes 1-3), the subpool number (byte 4), and the address of the first descriptor queue element (DQE) for the subpool (bytes 5-7) or, if the subpool is owned by another task (bit 0 is 1), the address of the SPQE that describes it (bytes 5-7).

Storage within a subpool is described by a descriptor queue element. Each DQE contains the number of bytes of main storage in the subpool. This count is always a multiple of 2048 bytes. If a request for space from a subpool cannot be satisfied with the space described by an existing DQE the GETMAIN routine builds another DQE and links the new DQE to the chain of existing DQE's. Each DQE contains a pointer to the FQE that represents the free area with the highest main storage address in the subpool (bytes 1-3), a pointer to the next DQE (bytes 5-7), and the length of the area described by the DQE, bytes 13-15(D-F).

Figure 33 shows the control blocks used to describe the subpools for a task in an MVT system.

I/O CONTROL BLOCKS

Queue of DEBs

To find the queue of DEBs for the task, look at the third word in the TCB (TCB+8). The address given here points to the first word of the most recent entry on the DEB queue. There is a DEB on this queue for each data set opened to the task at the time of the dump. DEBs are enqueued in the same order as the data sets are opened. The last three bytes of the second word in each DEB (DEB+5) points to the next most recent DEB on the queue. The queue contains one DEB for each open data set.

UCBs

You can find unit information for each device in your system in the unit control block (UCB) for that device. The address of the UCB is contained in the last 3 bytes of the ninth word of the DEB, DEB+33(21). If the DEB queue is empty, scan the dump around location 4096(1000) for words whose fifth and sixth digits are FF. These are the first words of the UCBs for the system; UCBs are arranged in numerical order by device address. (You may find it easier to locate UCBs by looking for the device address in the EBCDIC printout to the right of each page.) The first two bytes of the second word of each UCB give the device address. The device type and class are given in the third and fourth bytes of the fifth word, UCB+18(12), respectively. The sense bytes, with the exception of those for devices with extended sense, begin in the last two bytes of the sixth UCB word, UCB+22(16), and continue from 1 to 6 bytes depending on the device type. For the extended sense devices, UCB+22 and UCB+23 are ignored. UCB+24(18) in this case contains the number of bytes of sense

information to be found starting at the address specified in UCB+25(19). Sense bytes are given in Appendix G of this publication.

DCB and TIOT

The address of the DCB, a control block that describes the attributes of an open data set, is located in the last 3 bytes of the seventh DEB word, DEB+25(19). The first two bytes of the ninth word of the DCB, offset 40(28), contains the offset in the task input/output table (TIOT) of the DD name entered for the data set. Therefore, the address of the DD name for a particular data set may be found by adding the TIOT offset in the DCB to the TIOT address in the TCB (TCB+12), plus 24(16) bytes for the TIOT header.

IOB

If a data set is being accessed by a sequential access method with normal scheduling, the address of the input/output block (IOB) prefix (IOB-8) is located in the seventeenth word of the DCB, DCB-68(44). The first word of the IOB prefix points to the next IOB (if more than one IOB exits for the data set). Each IOB

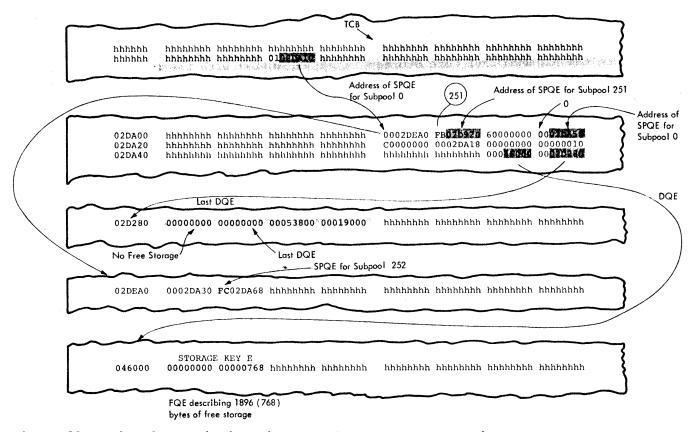


Figure 33. Subpool Descriptions in MVT - IMDPRDMP Storage Print

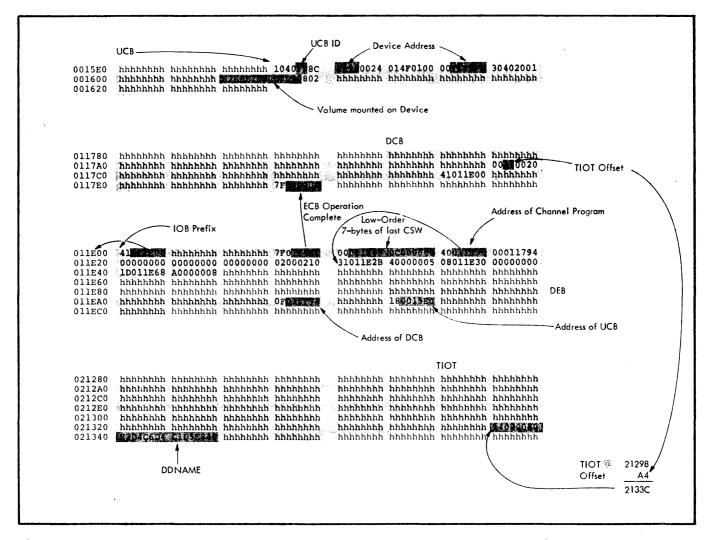


Figure 34. I/O Control Blocks

for an open data set contains a pointer to the CCW list in the last three bytes of the fifth word, IOB+17(11).

ECB

The Completion code for an I/O operation is posted in the first byte of the event control block (ECB). ECB completion codes are explained in Appendix F. If the I/O event is not complete and an SVC I (WAIT) has been issued, the high-order bit of the ECB is on, and bytes one through three contain the address of the associated RB. For the sequential and basic partition access methods the second word of an IOB points to its associated ECB.

Figure 34 shows the DEB, UCB, DCB, and IOB for a BSAM data set.

TSO CONTROL BLOCKS

The time sharing (TSO) control blocks are obtained from the IMDPRDMP service aid program by specifying the TSO control statement in the input stream. The first part of the TSO dump is the same as the normal MVT dump. The control blocks that IMDPRDMP formats are divided into two group: system and user.

TSCVT

The time sharing communications vector table (TSCVT) is a secondary CVT for the MVT CVT. The time sharing CVT resides in the time sharing region; therefore, it exists only while the time sharing region is active. When time sharing does not exist in the system, the MVT CVT pointer to the TSCVT (CVT+229) is zero.

A region control block (RCB) contains information that is unique to a time sharing region. There is one RCB for each time sharing region. The RCBs reside in the time sharing controller's region, they are contiguous, and they are created during initialization of the time sharing controller.

The TSCVT points to a region control block table. The RCB table is an indexed table containing one RCB address for each possible time sharing region, therefore, the table contains the maximum number of RCBs that may be used by time sharing. The first RCB is for region one, the second for region two, etc. The time sharing job block (TJB) of a job points to the RCB associated with that job.

UMSM

One user main storage map (UMSM) exists for each possible time sharing user. The UMSM contains a series of consecutive one-word extent fields (ADDR-LN). Each one-word extent contains a halfword address field (ADDR) and a halfword length field (LN) that describes the main storage allocated to the time sharing user. The UMSM contains the address and length of a storage block (a multiple of 2K bytes) that has been allocated to the user; only this allocated storage will be swapped out for the user. The time sharing job block (TJB) points to the UMSM.

SWAP DCB

The swap data control block (SWAP DCB) is used whenever a time sharing user's region is swapped into or out of main storage. It describes a swap data set that contains an IOB, area for channel programs, and the track map queue. The TJB points to the swap DCB.

<u>TJB</u>

The time sharing job block (TJB) contains status information about a time sharing user. The TJB is retained in main storage while the user is swapped out. One time sharing job block exists for each possible simultaneous time sharing user. The space for the TJB is obtained from the time sharing control task (TSC) region during time sharing initialization. Status information about the terminal related to this TJB is contained in the terminal status block (TSB). The address of the terminal status block is the first word of the TJB. The first word of the TSCVT points to the TJB.

<u>TSB</u>

Each terminal status block (TSB) contains status information about one terminal. The terminal input/output coordinator (TIOC) uses this information. During system initialization, one TSB is created for each possible user. The main storage space is obtained in one contiguous block for all of the TSBs in the region of the time sharing control task (TSC); this contiguous string of TSBs is called the TSB table. The origin pointer to the TSB table is the TIOCTSB field of the TIOCRPT.

TJBX

The time sharing job block extension (TJBX) contains user job information that can be rolled out to the swap data set with the user's job. The TJBX resides in the local system queue space (LSQS) for the region. The TJBX location is pointed to by the third word of the time sharing job block (TJB). The space for the TJBX is obtained by the region control task (RCT) during initialization.

PSCB

The protected step control block (PSCB) contains accounting information related to a single user. All timing information is in software timer units. A software timer unit is equal to 26.04166 micro seconds. The job step control block (JSCB), offset 268, points to the PSCB.

TAXE

The TSO terminal attention exit element (TAXE) is a physical addendum to a regular 24 word interrupt request block (IRB). It is used to schedule an attention exit resulting from a terminal attention interruption. It is created, queued, and dequeued by the specify terminal attention exit (STAX) macro instruction. The main storage space for the TAXE is obtained in the local system queue space (LSQS) of the terminal user's region.

For a more detailed description of the TSO control blocks formatted by the IMDPRDMP program, see the Control Block and/or TSO Control Program PLM publications.

<u>RCB</u>

Section 3: Tracing Aids

Tracing aids available are the save area chain, trace option, and Generalized Trace Facility (GTF). This section provides a description of each tracing aid, and, for GTF, describes its output after processing by the IMDPRDMP service aid.

Save Area Chain

The save area chain is edited and clearly identified in ABEND/SNAP dumps, and can be located easily in storage dumps produced by system dump facilities or the IMDSADMP service aid.

A save area is a block of 72 bytes containing chain pointers and register contents. It has the following format:

0	1 Next higher sove grea
) † N• 8	xf lower save area Register 14
} .	Register 15 Contents of Fregisters 0-12
16(10)	20(14)
Bytes	4-7: Pointer to the next higher level save area or, if this is the highest level save area, zeros.
Bytes	8-11(B): Pointer to the next lower level save area or, if this is the lowest level save area, unused.
Bytes	12-15(C-F): Contents of register 14 (optional)
Bytes	16-19(10-13): Contents of register 15 (optional)
Bytes	20-71(14-3F): Contents of registers 0 to 12

The save area for the first or highest level load module in a task (save area 1) is provided by the control program. The address of this area is contained in register 13 when the load module is first entered. It is the responsibility of the highest level module to:

- Save registers 0-12 in bytes 20-71(14-3F) of save area 1 when it is entered.
- Establish a new save area (save area 2).
- 3. Place the contents of register 13 into bytes 4-7 of save area 2.
- 4. Place the address of save area 2 into register 13.
- 5. Place the address of save area 2 into bytes 8-11(B) of save area 1.

At this point, the save areas appear as shown in Figure 35.

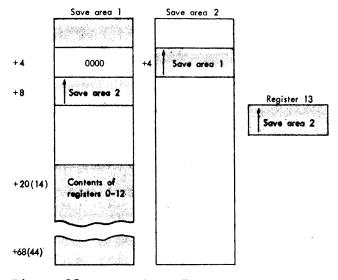


Figure 35. Save Area Trace

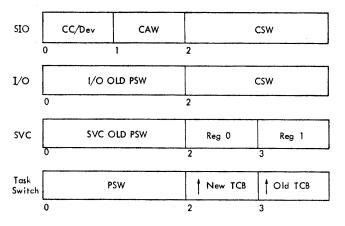
If a module requests a lower level module, it must perform actions 1 through 4 to ensure proper restoration of registers when it regains control. (Action 5 is not required, but must be performed if the dump printout of the field is desired.) A module that does not request a lower level module need only perform the first action.

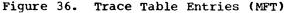
ABEND and SNAP dumps include edited information from all save areas associated with the dumped task under the heading "SAVE AREA TRACE". In a stand-alone dump, the highest level save area can be located through a field of the TCB. Subsequent save areas can be located through the save area chain.

TRACE OPTION

The tracing routine is an optional feature specified during system generation. This routine places entries, each of which is associated with a certain type of event, into a trace table. When the table is filled, the routine overlays old entries with new entries, beginning at the top of the table (the entry having the lowest storage address). The contents and size of a trace table are highly system-dependent.

Systems With MFT: Trace table entries for systems with MFT are 4 words long and represent SIO, I/O, SVC and dispatcher task-switching interruptions. Figure 36 shows the word contents of each type of entry.





Systems with MVT: The trace table in a system with MVT is expanded to include more entries and more information in each entry. Trace table printouts occur only on SNAP dumps and stand-alone dumps. Entries are eight words long and represent occurences of SIO, external, SVC, program, and I/O interruptions, and dispatcher loaded PSWs.

Figure 37 shows the word contents of trace table entries for SNAP dumps and stand-alone dumps. Figure 38 shows the contents of trace table entries as filled by MVT with Model 65 multiprocessing. (SSM -- set system mask -- entries are optional.)

INTERPRETING TRACE TABLE ENTRIES

Location 84(54) in main storage contains the address of the first word of the three word trace table control block. The trace table control block immediately preceeds the table. The trace table control block describes the bounds of the table and the most recent entry at the time of the dump.

Current	Entry	First	Entry	Last	Entry
L	•	L		L	

0 4 8 You can locate the trace table by scanning the contents of main storage between locations 16,384(4000) and 32,768(8000) for trace table entries. Entries are four words long and begin at addresses ending with zero. To find the table boundaries and current entry, scan the table in reverse until you reach the trace table control block.

Trace Table Entries in MFT: Trace table entries for systems with MFT are 4 words long and represent occurrences of SIO, I/O, SVC, and task-switching interruptions. Figure 39 gives some sample entries and their contents.

SIO entries can be used to locate the CCW (through the CAW), which reflects the operation initiated by an SIO instruction. If the SIO operation was not successful, the CSW STATUS portion of the entry will show you why it failed.

<u>I/O</u> entries reflect the I/O old PSW and the CSW that was stored when the interruption occurred. From the PSW, you can learn the address of the device on which the interruption occurred (bytes 2 and 3), the CPU state at the time of interruption (bit 15), and the instruction address where the interruption occurred (bytes 5-8). The CSW provides you with the unit status (byte 4), the channel status (byte 5), and the address of the previous CCW plus 8 (bytes 0-3).

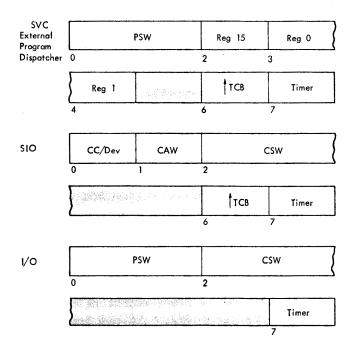


Figure 37. Trace Table Entries (MVT)

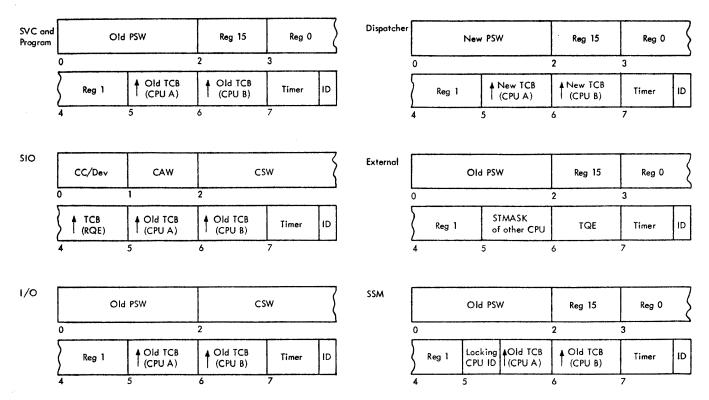


Figure 38. Trace Table Entries (MVT with Model 65 multiprocessing)

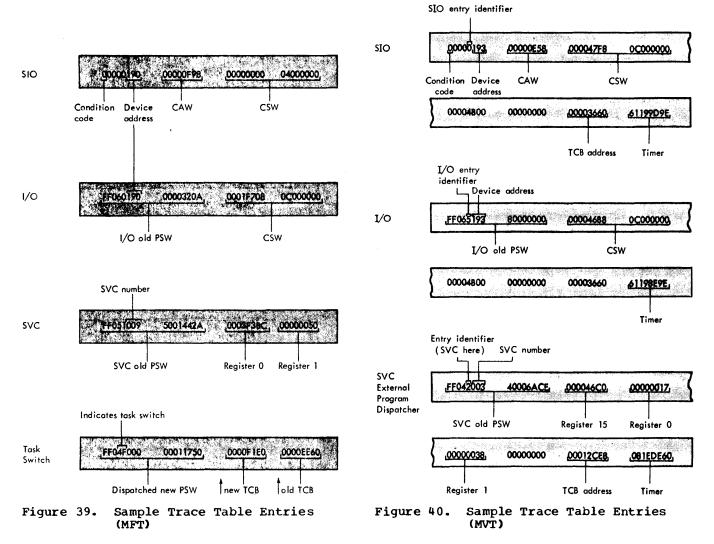
SVC entries provide the SVC old PSW and the SIO 5th word address of TCB. contents of registers 0 and 1. The PSW 6th word: address of old TCB offers you the hexadecimal SVC number (bits 20-31), the CPU mode (bit 15), and the address of the SVC instruction (bytes 5-8). The contents of registers 0 and 1 are useful in that many system macro instructions use these registers for I. parameter information. Contents of registers 0 and 1 for each SVC interruption are given in Appendix B. Trace Table Entries in MVT and M65MP: S Entries in an MVT trace table are 8 words P long and represent occurrences of SIO, external, SVC, program, I/O, and dispatcher interruptions. You can identify what type of interruption caused an entry by looking at the fifth digit: D

0	=	SIO
1 :	=	External
2	-	SVC
3	z	Program
5	=	1/0
D	=	Dispatcher

Figure 40 gives some sample entries and their contents.

In dumps of Model 65 Multiprocessing system, trace table entries differ as follows:

	ocn	woru.	for CPU A.
	7th	word:	address of old TCB
			for CPU B.
	8th	word	CPU identification
			(last byte).
I/0	3rd	word:	contents of register
			15.
	4th	word	contents of register
			0.
	8th	word	CPU identification
			(last byte).
SVC and	6th	word:	address of old TCB
Program			for APU A.
	7th	word:	address of old TCB
		_	for CPU B.
	8th	word	CPU identification
Dimminal	C + b		(last byte).
Dispatcher	6 t n	wora:	address of new TCB
	7+6	word:	for CPU A.
	/tn	word:	address of new TCB for CPU B.
	0+h	word:	CPU identification
	ocn	woru:	(last byte).
External	6+h	word:	STMASK of other CPU.
DACCTIMI		word:	TQE if timer inter-
	/ СП	woru.	rupt occurred.
	8+h	word:	CPU identification
	0.011	word.	(last byte). If so, a
			program check at the
			instruction preceding
			that address caused
			the interruption.
			-



Generalized Trace Facility

The Generalized Trace Facility (GTF) traces system and application program events and records information about these events. Trace records can be stored internally -in a table similiar to the trace table of the Trace Option -- or they can be recorded externally in a data set that becomes input to the IMDPRDMP service aid program. (When stored internally the trace table is formatted in ABEND/SNAP dumps.) The IMDPRDMP service aid edits and formats the GTF external trace records as specified in an EDIT control statement.

This section describes the output of GTF; it does not tell how to use GTF. For a description of the functions performed by GTF and IMDPRDMP refer to the Service Aids publication. System events traced by GTF in MFT, MVT, and MVT-M65MP systems are:

IO interrupts SIO operations SVC interrupts Program interrupts External interrupts Task Switches by the system dispatcher SSM interrupts in multi-processing systems

GTF MINIMAL TRACE RECORDS

The following material describes the records produced under the minimal trace option (SYSM) of GTF. The formats described appear in both ABEND/SNAP dumps (under the heading GTF TRACE TABLE) and in IMDPRDMP output. Minimal trace records are produced for IO and PCI/IO, SIO, SVC, PGM, EXT, DSP, and SSM events. ${A \\ B} {IO \\ PCI}$ OLD PSW hhhhhhhh hhhhhhhh CSW hhhhhhhh hhhhhhhh RQE TCB ${******** \\ hhhhhhhh \\ N/A}$ OLD TCB hhhhhhhh

Figure 41. IO and PCI/IO Minimal Trace Record

- {A }
 {B }
 appears in MVT-M65MP system records;
 identifies the CPU associated with the
 event.
- IO PCI identifies the type of trace record.
- OLD PSW hhhhhhhh hhhhhhh the program status word that was current at the time the IO or PCI/IO interrupt occurred.
- CSW hhhhhhh hhhhhhh the channel status word associated with the IO or PCI/IO interrupt being traced.

RQE TCB

indicates that an error occurred while gathering the information. hhhhhhh is the address of the TCB of the task for which this I/O operation is being performed.

- N/A indicates the interrupt was unsolicited: either the I/O supervisor did not issue an SIO instruction to the device; or there is no valid UCB for the device.
- OLD TCB hhhhhhhh in MFT and MVT system trace records, the address of the TCB for the task that was in control when the interrupt occurred.

in MVT-M65MP systems the OLA and OLB fields replace the OLD TCB field and contain the address of the TCB for the task in control of CPU A and CPU B respectively, at the time the interrupt occurred.

SIO Minimal Trace Record

${A \\ B}$ SIO CC/DEV/CAW hhhhhhhh hhhhhhh	n CSW hhhhhhhhhhhhhhhhhh	$ \begin{array}{c} & \left\{ \begin{array}{c} ******** \\ hhhhhhhh \\ hhhhhhhh \\ N/A \end{array} \right\} \left\{ \begin{array}{c} OLA \\ OLD \\ TCB \\ hhhhhhhh \\ OLD \\ TCB \\ hhhhhhhh \\ O \\ O \end{array} \right\} $	hhhhhhh
--	--------------------------	--	---------

Figure 42. SIO Minimal Trace Record

 $\left\{ \begin{array}{c} \mathbf{A} \\ \mathbf{B} \end{array} \right\}$

appears in MVT-M65MP system records; identifies the CPU associated with the event.

SIO

identifies the type of trace record.

CC/DEV/CAW hhhhhhhh hhhhhhh displays the SIO condition code, the device address, and the CAW (channel address word) for the I/O operation just initiated.

> The first four digits represent the condition code returned from the SIO operation; the next four digits represent the device address; and the last eight digits represent the CAW.

CSW hhhhhhhh hhhhhhhh the channel status word associated with this event.

RQE TCB

indicates that an error occurred while gathering the information.

hhhhhhh

is the address of the TCB of the task for which this I/O operation is being performed.

N/A

indicates the interrupt was unsolicited, i.e., the I/O supervisor did not issue an SIO instruction to the device; or, there is no valid UCB for the device.

OLD TCB hhhhhhhh

in MFT and MVT system trace records, the address of the TCB for the task that was in control when the interrupt occurred.

In MVT-M65MP systems the OLA and OLB fields replace the OLD TCB field and contain the address of the TCB for the task in control of CPU A and CPU B respectively, at the time the interrupt occurred.

DSP Minimal Trace Record

 ${A \\ B}$ DSP ${RES PSW \\ NEW PSW}$ hhhhhhhh hhhhhhh R15/R0 hhhhhhhh hhhhhhhh R1 hhhhhhhh ${NUB hhhhhhhh}$ ${NUB hhhhhhhh}$ ${NUB hhhhhhhh}$

Figure 43. DSP Minimal Trace Record

${A \\ B}$

appears in MVT-M65MP records; identifies the CPU associated with the event.

DSP

identifies the type of record.

NEW PSW hhhhhhhh hhhhhhhh the PSW for the task about to be dispatched.

In a record obtained from a MVT-M65MP system this field will be labeled RES PSW.

R15/R0 hhhhhhhh hhhhhhh the contents of general purpose registers 15 and 0 as they will be when the task being dispatched is given control.

- R1 hhhhhhhh the contents of general purpose register 1 as it will be when the task being dispatched is given control.
- NEW TCB hhhhhhhh the address of the TCB for the task about to be dispatched.

In a record obtained from a MVT-M65MP system this field is replaced by the NUA and NUB fields containing the addresses of the tasks to be dispatched on CPU A and CPU B when processing resumes.

EXT Minimal Trace Record

A EXT OLD PSW hhhhhhhh hhhhhhh R15/R0 hhhhhhhh hhhhhhhh R1 hhhhhhhh STMSK hhhhhhhh TQE TCB hhhhhhhhh (N/A

Figure 44. EXT Minimal Trace Record

 $\left\{ \begin{array}{c} \mathbf{A} \\ \mathbf{B} \end{array} \right\}$

appears in MVT-M65MP records; identifies the CPU associated with the event.

EXT

identifies the type of trace record.

- OLD PSW hhhhhhhh hhhhhhh the program status word that was current at the time the external interrupt occurred.
- R15/R0 hhhhhhh hhhhhhh the contents of general purpose registers 15 and 0 at the time the interrupt occurred.
- R1 hhhhhhhh the contents of general purpose register 1 at the time the interrupt

occurred.

STMSK hhhhhhhh appears in MVT-M65MP records only; displays the SHOULDER TAP MASK at the time the interrupt occurred.

.

- ******* indicates that an error occurred while gathering the information.
- hhhhhhh is the address of the TCB of the task that requested this timer interrupt.

N/A

indicates the interrupt was other than a timer interrupt.

PGM Minimal Trace Record

 A
 PGM
 OLD
 PSW
 hhhhhhhh
 Nhhhhhhh
 R15/R0
 hhhhhhhh
 hhhhhhhh
 OLD
 TCB
 hhhhhhhhh
 A
 B
 B
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Figure 45. PGM Minimal Trace Record

$\left\{ \begin{array}{c} \mathbf{A} \\ \mathbf{B} \end{array} \right\}$

appears in MVT-M65MP system records; identifies the CPU associated with the event.

- PGM identifies the type of trace record.
- OLD PSW hhhhhhhh hhhhhhh the program status word that was current at the time the program interrupt occurred.
- R15/R0 hhhhhhh hhhhhhh the contents of general purpose registers 15 and 0 at the time the interrupt occurred.

R1 hhhhhhhh the contents of general purpose register 1 at the time the interrupt occurred.

OLD TCB hhhhhhhh the address of the TCB for the task that was in control when the interrupt occurred.

> In MVT-M65MP trace records this field is replaced by the OLA and OLB fields that contain, respectively, the address of the TCB for the tasks in control of CPU A and CPU B at the time the interrupt occurred.

 ${\Lambda \\ B}$ SVC OLD PSW hhhhhhh hhhhhhh R15/R0 hhhhhhh hhhhhhhh R1 hhhhhhhh ${
m OLD}$ TCB hhhhhhhh ${
m OLB}$ hhhhhhhh

Figure 46. SVC Minimal Trace Record

${A \\ B}$

appears in MVT-M65MP system records; identifies the CPU associated with the event.

SVC

identifies the type of trace record.

- OLD PSW hhhhhhhh hhhhhhh the program status word that was current at the time the interrupt occurred. The SVC number, e.g., SVC 51, is represented by the last two hexadecimal digits in the first word.
- R15/R0 hhhhhhh hhhhhhh the contents of general purpose

registers 15 and 0 at the time the interrupt occurred.

R1 hhhhhhh the contents of general purpose register 1 at the time the interrupt occurred.

OLD TCB hhhhhhhh the address of the TCB for the task that issued the SVC.

> In MVT-M65MP systems the OLA and OLB fields replace the OLD TCB field and contain the address of the TCB for the task in control of CPU A and CPU B respectively, at the time the interrupt occurred.

SSM Minimal Trace Record

 ${A \ B}$ SSM LK C OPSW hhbhhhhh hhhhhhhh R15/R0 hhhhhhhh hhhhhhhh R1 hhhhhhhh OLA hhhhhhhh OLB hhhhhhhh

Figure 47. SSM Minimal Trace Record

{ ^A _B }	indicates the CPU associated with the event.	OPSW hhhhhhhh hhhhhhhh the program status word that was current at the time the interrupt occurred. Obtained from the CPU on which the interrupt occurred.
SSM IK c	identifies the type of trace record. CPU affinity byte:	R15/R0 hhhhhhhh hhhhhhhh R1 hhhhhhhh The contents of general purpose registers 15, 0, and 1 from the CPU on which the interrupt occurred, at the time the interrupt occurred.
	A indicates CPU A executing disabled. B indicates CPU B executing disabled. 0 Neither CPU executing disabled.	OLA hhhhhhhh OLB hhhhhhhh the addresses of the TCBs of the tasks in control in CPU A and CPU B respectively at the time the interrupt occurred.

GTF COMPREHENSIVE TRACE RECORDS

.

The following material describes the records produced when comprehensive tracing is specified at the invoking of GTF (MODE=EXT). The formats described appear in the output from IMDPRDMP service aid processing of the data recorded by GTF.

Comprehensive trace records are produced for IO, PCI/IO, SIO, DSP, EXT, PGM, SSM, and SVC events.

Figure 48. IO and PCI/IO Comprehensive Trace Record

{A B This field appears only in MVT-M65MP system I/O or PCI trace records and identifies the computer associated with the event.

${IO \\ PCI}$

This field identifies the type of trace record -- input/output (IO) or program controlled interrupt (PCI).

cuu

This field displays the device address for the device associated with the interrupt in channel/unit form.

OLD PSW hhhhhhh hhhhhhh This field displays the program status word that was current at the time the IO or PCI interrupt being traced, occurred.

This field has three possible entries, as follows:

ccccccc

is the one to eight character name of the job associated with the interrupt being traced.

asterisks indicate that a bad control block chain prevented the jobname from being obtained.

N/A

in PCI trace records N/A indicates that the interrupt was issued by the system and there is no associated jobname; in IO interrupt trace records N/A indicates either a system issued interrupt as for PCI or an interrupt issued without a valid UCB for the device issuing the interrupt.

DDNM {ccccccc ********* N/A

This field has three possible entries, as follows:

- cccccccc is the name of the DD statement associated with the interrupt being traced.
- *******

asterisk indicate that a bad control block chain prevented the data definition name from being obtained.

N/A

N/A appears in the DDNM field for one of the following reasons:

- An interrupt was issued without a valid UCB for the device issuing the interrupt.
- The post bit in the UCB is 'off.'
- The data event block (DEB) pointer to the TCB is set to 0.
- The DCB is not opened.
- The DCB TIOT offset is outside the valid range.
- The TCB TIOT pointer is set to 0.
- The DDNAME in the TIOT is not recorded in EBCDIC characters.

OLTCB hhhhhhhh

In MFT and MVT system trace records this field displays the address of the TCB that was current at the time the IO or PCI interrupt being traced, occurred.

In MVT-M65MP system IO and PCI trace records the following fields replace the OLTCB field:

OLA hhhhhhh This field displays the address of the A computer TCB that was current when the IO or PCI interrupt occurred.

OLB hhhhhhh This field displays the address of the B computer TCB that was current when the IO or PCI interrupt occurred.

CSW hhhhhhhh hhhhhhhh This field displays the channel status word from permanent storage location 64.

```
RQE (hhhhhhhh hhhhhhhh hhhhhhhh
N/A
```

This field has three possible entries as follows:

hhhhhhh hhhhhhh hhhhhhh is the content of the first three words of the Request Queue Element associated with the IO or PCI interrupt.

******* ******* ******* asterisks indicate that a bad control block chain prevented the RQE information from being obtained.

N/A

indicates that the interrupt was issued without a valid UCB for the device issuing the interrupt. RQE TCB

This field has three possible entries as follows:

hhhhhhhh

is the address of the TCB associated with the Request Queue Element

asterisks indicate that a bad control block chain prevented the TCB address from being obtained.

N/A

indicates that the interrupt was issued without a valid UCB for the device issuing the interrupt.

SENS (hhhhhhhh)

l n/a

This field has two possible entries as follows:

hhhhhhh

is the content of the four sense bytes in the UCB beginning at UCB + 22 which describe the IO or PCI interrupt being traced. For more information about the sense bytes see Appendix G.

N/A

indicates that the interrupt was issued without a valid UCB for the device issuing the interrupt.

Figure 49. SIO Comprehensive Trace Record

 $\left\{ \begin{array}{c} A \\ B \end{array} \right\}$ OLTCB hhhhhhhh in MFT/MVT systems the address of the appears in MVT-M65MP system trace TCB that was current when the SIO was records; identifies the computer issued. associated with the event. SIO in MVT-M65MP systems the OLA and OLB the type of trace record. fields replace the OLTCB field. cuu the device address in channel/unit OLA hhhhhhhh form for the device associated with is the A computer address of the the record. TCB that was current when the SIO was issued. CC hh hh - is the condition code set by the OLB hhhhhhhh SIO event. is the B computer address of the CAW hhhhhhhh TCB that was current when the SIO the channel address word associated was issued. with this event -- taken from permanent storage location 72. CSW hhhhhhhh hhhhhhh JOBN (CCCCCCC) the channel status word associated) N/A with this event -- taken from permanent storage location 64. ccccccc is the one to eight character RQE hhhhhhhh hhhhhhhh jobname of the job associated the first three words of the Request Queue Element associated with the SIO with this event. operation. N/A indicates the SIO was issued by RQE TCB hhhhhhhh the system and there is no the address of the TCB associated with associated jobname. the request queue element.

DSP Comprehensive Trace Record

{ A B	DSP RES PSW hhhhhhhh hhhhhhhh JOBN (ccccccc) MOE	WAITTCB SVC-cccc SVC-RES **IRB*** cccccccc Iccccccc 999999999
Figu	re 50. DSP Comprehensive Trace Record	
$ \left\{ \begin{matrix} \mathbf{A} \\ \mathbf{B} \end{matrix} \right\} $	MVT-M65MP systems only. Identifies the computer associated with the event.	<pre>**IRB*** an asynchronous routine is about to be dispatched and the module name is not available.</pre>
DSP	the type of trace record. PSW hhhbhbhb hhbbbbhb	cccccccc in MVT systems the eight character module name from the CDE associated with the task to be dispatched; or, the name of an error exit routine from the SIRB
	the PSW for the task about to be dispatched. If this task was interrupted at some previous point in time, then this was the current PSW at the interrupt.	associated with the task. in MFT systems the eight character name from the LRB, LPRB, PRB or FRB associated with
JOBN	{ccccccc} N/A	the task being dispatched; or an error exit routine name from the SIRB associated with the task.
	ccccccc is the eight character name of the job associated with the task being dispatched.	Iccccccc indicates that error fetch is in the process of loading an error recovery module. The last seven characters of the module name are shown.
	indicates the task switch is for a system task; no jobname is available.	999999999 In MFT (with subtasking) an ATTACH was issued, but the module requested did not receive
MO DN	SVC-cccc SVC-RES ++IRB+++ cccccccc Iccccccc	<pre>control. NUTCB hhhhhhh the address of the new TCB the TCB of the next-to-be-dispatched task.</pre>
1	(99999999) WAITTCB the WAIT task is about to be given control.	in MVT-M65MP systems the following fields replace the NUTCB field: NUA hhhhhhh the address of the TCB of the next-to-be-dispatched task in the
	SVC-cccc indicates a type 3 or 4 SVC is about to get control; cccc is the last four characters in the module name.	A computer. NUB hhhhhhhh the address of the TCB of the next-to-be-dispatched task in the B computer.
	SVC-RES indicates a resident type SVC routine is about to be given control.	PRTY hh hh the dispatching priority of the next-to-be-dispatched task.

EXT Comprehensive Trace Record

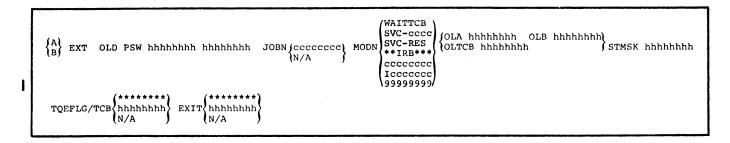


Figure 51. EXT Comprehensive Trace Record

 ${A \\ B}$

This field appears only in MVT-M65MP system EXT trace records and identifies the computer associated with the event.

EXT

This field identifies the trace record as an EXT trace record.

OLD PSW hhhhhhh hhhhhhh This field displays the program status word that was current at the time the external interrupt occurred.

JOBN (CCCCCCCC)

This field has two possible entries as follows:

ccccccc

is the one to eight character name of the job associated with the event.

N/A indicates that the interrupt was issued by the system and there is not associated job name.

WAITTCB SVC-cccc MODN SVC-RES **IRB*** cccccccc Iccccccc 99999999

> WAITTCB The WAIT task was interrupted.

> SVC-cccc A type 3 or 4 SVC routine was interrupted; cccc is the last four characters of the routine name.

SVC-RES a resident SVC routine was interrupted.

IRB*
 the EXT interrupt occurred during
 execution of an asynchronous
 routine with an associated IRB.

ccccccc

in MVT systems the eight character name of the module that was interrupted - taken from the CDE associated with the task; or the name of an error routine taken from the SIRB associated with the task.

in MFT systems the eight character name of the module that was interrupted - taken from either the LRB, LPRB, PRB, or FRB; or the name of an error routine - taken from the SIRB associated with the task.

Iccccccc indicates that error fetch was in the process of loading an error recovery routine when the interrupt occurred. The last seven characters of the module name are shown.

999999999

In MFT (with subtasking) an ATTACH was issued, but the module requested did not receive control.

OLTCB hhhhhhhh In MFT/MVT systems the address of the TCB that was current when the interrupt occurred.

In MVT-M65MP systems the OLA and OLB fields replace the OLTCB field.

OLA	hhhhhhh
	is the address of the TCB in the
	A computer that was current when
	the interrupt occurred.

OLB hhhhhhhh is the address of the TCB in the B computer that was current when the interrupt occurred.

STMSK hhhhhhhh

In MVT-M65MP systems only - the 'shoulder tap' mask from location X'2BC' in the other computers prefix.

TQEFLG/TCB

hhhhhhh

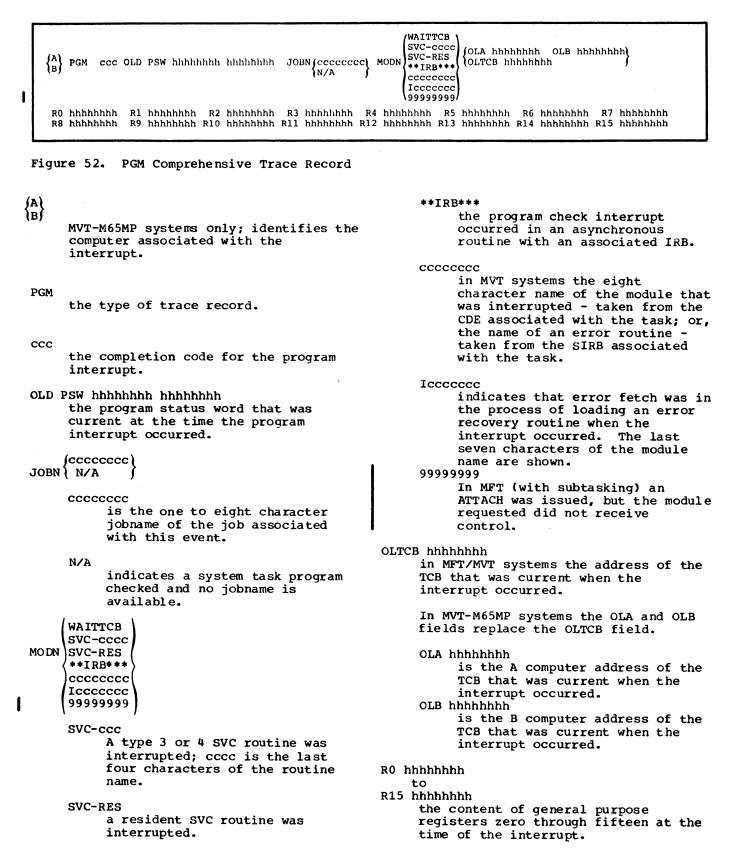
is the first word of the timer queue element (TQE). The first byte of the word is the TQEFLGS and the remaining three bytes the TQETCB, which is the address of the TCB for the task in which this timer element is being used. *******
asterisks indicate that a bad
control block chain prevented the
information from being obtained.
N/A
indicates that this EXT interrupt
was not caused by the timer.
EXIT
{
hhhhhhhh
h
is the address of the exit
routine - taken from the eighth
word of the TQE.

N/A

indicates that this EXT interrupt was not caused by the timer.

asterisks indicate that a bad control block chain prevented the information from being obtained.

PGM Comprehensive Trace Records



SSM Comprehensive Trace Record

	${A \\ B}$ SSM OLD PSW hhhhhhhh	JOBN {cccccccc} {N/A }	WAITTCB SVC-cccc MODN SVC-RES **IRB*** cccccccc Iccccccc	OLA hhhhhhhh	OLB hhhhhhhh	TKID C	
--	---------------------------------	---------------------------	--	--------------	--------------	--------	--

Figure 53. SSM Comprehensive Trace Record

۶	A	l
l	в	ſ

identifies the computer associated with the SSM interrupt.

SSM

- identifies this trace record as an SSM trace record.
- OLD PSW hhhhhhh hhhhhhh
 - the program status word that was current at the time the set system mask instruction was issued.
- JOBN (CCCCCCCC)

A\N *f*

- cccccccc is the one to eight character name of the job associated with SSM interrupt.
- N/A
 - indicates that the system originated the interrupt and there is no associated jobname.
- MODN WAITTCB SVC-cccc SVC-RES ** IRB*** ICCCCCCC
 - WAITTCB

the WAIT task was interrupted.

SVC-cccc a type 3 or 4 SVC routine was interrupted; cccc is the last four characters of the routine name.

SVC-RES a resident SVC routine was interrupted.

the SSM interrupt occurred during execution of an asynchronous routine with an associated IRB.
ccccccc
the eight character name of the module that was interrupted - taken from the content directory element (CDE) for the task; or the name of an error routine - taken from the SIRB associated with the task.
Icccccc
indicates that error fetch was in the process of loading an error recovery routine when the interrupt occurred. The last seven characters of the module name are shown.
OLA hbhhhbh
is the A computer address of the TCB
that was current when the interrupt

OLB hhhhhhhh is the B computer address of the TCB that was current when the interrupt occurred.

LKID C

occurred.

IRB*

- CPU affinity byte:
 - A indicates CPU A executing disabled.
 - B indicates CPU B executing disabled.
 - 0 Neither CPU executing disabled.

TIME AND LOST EVENT RECORDS

GTF produces two types of time records and a lost event record as follows:

TIME ddddddddd

appears on the last line of every event record if TIME=YES was specified in the GTF start command, and designates in decimal the number of seconds and microseconds since the last midnight.

***DATE: DAY ddd YEAR dddd TIME dd.dd.dd

This timestamp record appears at the beginning of the printout of each buffer filled by GTF and represents the time the first record was placed in the buffer.

DAY ddd is the Julian date.

YEAR dddd is the year. TIME dd.dd.dd is the time since midnight in a twenty-four hour format (hours.minutes.seconds).

*** LOST EVENTS: NUM ddddddddd TIME dd.dd.dd (GTF DISABLED)

> The lost event record appears whenever GTF loses records, whether it is because the GTF buffers overflowed or because GTF was temporarily disabled by ABEND. The record is not produced if GTF terminates when the buffers are full.

NUM ddddddddd is the number of records that were lost; one to ten decimal digits.

TIME dd.dd.dd is the time GTF resumed recording; 24-hour format starting at midnight.

GTF DISABLED

appears only if the events were lost because GTF was temporarily disabled, e.g., ABEND temporarily disables GTF in order to format GTF output for an ABEND dump. HEXFORMAT USER SYSTEM SUBSYS

Figure 54. Hexadecimal Format Record

Under some circumstances IMDPRDMP formats and prints GTF records in hexadecimal notation. The conditions under which GTF records are formatted and printed in hexadecimal format by IMDPRDMP are presented in the discussion of the hexformat record that follows:

HEXFORMAT

This label identifies a record dumped in hex format at the request of the user on a GTRACE macro. This request was made by not specifying a format appendage, that is FID=00.

USR

This label identifies this record as dumped in hexformat because the user requested a format appendage on the GTRACE macro that could not be found. This format appendage was identified by FID=hh, and therefore its name is IMDUSR hh.

SYSTEM

This label identifies a record that was dumped in hex format because either it is a GTF error record or the format appendage for it has been scratched by the user. If relative bytes 0, 1 or 8, 9 contain X'EEEE', then this is an error record produced by GTF. This error record was produced as a result of an unrecoverable error in a GTF data gathering routine. When the error was encountered message IHL118I was written on the master system console indentifying the error and the action taken. This message is not issued if the error occurred while building a comprehensive SVC trace record.

Except for comprehensive SVC records, this was the last record of its type produced during the run of GTF that produced it. If the X'EEEE' were not in the record, then it was dumped in hexformat because the IMDPRDMP format appendage that formats this type of record was not found by IMDPRDMP.

SUBSYS

This label identifies this record as dumped in hexformat because the subsystem format appendage requested by the subsystem on a GTRACE was not found by IMDPRDMP. The request was made via FID=hh, and therefore, it's name is IMDUSRhh.

```
AID hh
```

This field contains the AID of this record, and should always be X'FF'. The AID is the application identifier, and GTF's is always X'FF'.

FID hh

This field contains the FID, or format identifier. It is appended to 'IMDSYS' or 'IMDUSR' to obtain the name of the format appendage that was to have formatted this record.

EID hhhh This field contains the EID, or event identifier, for this record. The EID uniquely identifies the event that produced this record.

hhhhhhhh hhhhhhh hhhhhhh up to 64 words (256 bytes) of record in the GTF internal format. The internal format of GTF records is available in the Service Aids PLM.

GTF SVC COMPREHENSIVE TRACE RECORDS			<u>svc #</u> 45	<u>Group</u> 3	<u>Page</u> 185
There are four groups of GTF SVC		45	1	172	
Comprehensive Trace records.		47	3	185	
compron			48	3	185
Grou	р 1 Т	hose with Basic Fields	49	ĭ	172
Group 2 Those with Basic Fields plus			50 null		
	NAME Fie		51	4	197
Grou	р 3 Т	hose with Basic Fields plus	52	1	172
		List Field	53	2	181
Grou	ир 4 T	hose with Basic Fields plus	54	4	198
Vari	able Fie	ld (s)	55	2	181
			56	3	186
The following sub-index lists the SVCs in			57	2	181
sequence, identifies the group to which			58	2	181
they belong, and gives the page where			59	1	172
register contents and other variable fields are noted.			60	3	186
are not	eu.		61 62	1 4	172 198
SVC #	Group	Page	63	3	198
0	4	190	64	3	187
1	3	182	65	4	198
2	1	170	66	4	199
3	1	171	67	4	199
4	3	182	68	1	172
5	3	183	69	2	181
6	4	190	70	3	187
7	4	190	71	4	199
8	4	190	72	1	173
9	4	191	73	3	187
10	3	171	74	3	187
11	1	171	75	4	199
12	1	171	76	1	173
13	4	191	77	3	187
14	4	191	78	4	200
15	4	191	79	1	173
16	4	191	80	3	188
17 18	4	192	81 82	4	200
18	3 3	183 183	82 83	4 1	200 173
20	3	183	84	1	173
21	4	192	85	1	173
22	3	183	86	4	201
23	3	183	87	3	188
24	2	181	88	4	202
25	4	192	89	4	202
26	4	193	90	3	188
27	4	194	91	1	173
28	4	195	92	1	173
29	4	195	93	1	174
30	4	195	94	1	175
31	2	181	95	1	177
32	4	195	96	3	188
33 34	4	195	97 97	1	178
34	1 3	171	98 99	4 3	202 189
36	1	183 171	99 100	3 1	178
37	3	184	100	1	178
38	1	171	101	3	189
39	3	184	102	1	179
40	3	184	104	1	179
41	4	196	105	1	180
42	4	196	109	1	180
43	1	171	116	ī	180
44	4	197	117	1	180

SVC Comprehensive Trace Records Group 1 -- Basic Fields

OLA hhhhhhhh OLB hhhhhhhh A SVC ddd OLD PSW hhhhhhhh hhhhhhhh JOBN ccccccc MODN ccccccc OLTCB hhhhhhhh R15/R0 hhhhhhhh hhhhhhhh R1 hhhhhhhh Figure 55. Basic SVC Comprehensive Trace Record SVC-RES {A B indicates the SVC was issued by a this field appears only in MVT-M65MP resident SVC with an associated SVRB. records and identifies the CPU associated with the event. SVC-nnn indicates the SVC was issued by a SVC ddd transient SVC module with an the decimal number of the SVC associated SVRB. nnnn denotes the last four characters of the module name. OLD PSW hhhhhhh hhhhhhh *cccccc the program status word that was current at the time the SVC interrupt indicates that error fetch is in occurred. When SVC processing is the process of loading an error completed, operation is resumed under recovery module. ccccccc is the control of this PSW. last seven characters of the module name. ******* ccccccc JOBN ccccccc is, in MVT systems, the eight N/A character name of the module issuing the SVC -- taken from the ******* CDE associated with the task; or the name of an error routine -indicate an error occurred while attempting to retrieve the taken from the SIRB associated jobname, e.g., an incorrect TIOT with the task. address in the TCB could result in asterisks being placed in this In MFT systems the module name is taken from the LRB, LPRB, PRB, or field. FRB and the error routine name is ccccccc taken from the SIRB associated with the task. is the eight character jobname of the job issuing the SVC. N/A N/A indicates the RB CDE pointer was indicates that the SVC was issued zero. by the system and there is no ******* associated jobname. indicates that an error occurred **IRB*** while attempting to retrieve the SVC-RES module name. MODN\SVC-nnnn *cccccc 99999999 ccccccc In MFT (with subtasking) an N/A ATTACH was issued, but the module ******* requested did not receive I 999999999 control. **IRB*** indicates the SVC was issued by an asychronously executed routine with an associated IRB.

OLTCB hhhhhhhh the address of the TCB that was current when the SVC was issued.

> In MVT-M65MP systems the OLA and OLB fields replace the OLTCE field and indicate the addresses of the TCBs that were current in CPU A and CPU B when the SVC was issued.

R15/R0 hhbhbhhh hhbhbhhh R1 hhbhbhhh the contents of registers 15, 0, and 1 when the SVC was issued. SVC Comprehensive Trace Records Group 1 --Basic Fields

SVC 2 (POST)

R15 contains no applicable information. R0 contains the completion code to be placed in the ECB. R1 contains the address of the ECB to be posted.

170.2 Programmer's Guide to Debugging (Release 21.7)

<u>SVC 3 (EXIT)</u> registers contain no applicable information.

SVC 10 (REGMAIN)

R15 contains no applicable information. R0 contains the number of the subpool requested in the high order byte, and the number of bytes requested in the low order three bytes. R1 contains any negative value if the request is for a GETMAIN; contains the address of the storage to be freed if the request is for a FREEMAIN; contains zero value if the request is for a FREEMAIN for an entire subpool.

SVC 11 (TIME)

R15 contains no applicable information. R0 contains no applicable information. R1 contains flag bits in the low order byte that designate how the time is to be returned in Register 0.

If the low order byte is:

x'00'

register 0 is to contain a 32 bit unsigned binary number representing the number of timer units that have elapsed. (A timer unit is 26.04 micro-seconds).

x'01'

register 0 is to contain elapsed time in hundredths of a second.

x*02*

register 0 is to contain packed decimal digits representing elapsed time in hours, minutes, seconds, tenths of a second, and hundredths of a second (HHMMSSth).

SVC 12 (SYNCH)

R15 contains the address of the entry point for the processing program that is to be given control. R0 contains no applicable information. R1 contains no applicable information.

SVC 34 (MGCR)

R15 contains no applicable information.

R0 and R1 contents are as follows: R1, if positive, contains a pointer to the command buffer of the command to be processed. R0 is not used in this case. If R1 is negative and R0 is zero, then R1 contains a pointer to the CSCB that is to be either added to the chain or deleted from the chain.

If R1 is negative and R0 is positive, then R1 contains a pointer to the CIB that is to be added to or deleted from the chain. R0 contains a pointer to the beginning of the chain.

If R1 is negative and R0 is negative, then R0 contains a pointer to the CIB in which the CIB count is to be set and R1 contains the value to which the CIB count is to be set.

SVC 36 (WTL)

R15 contains no applicable information. R0 contains no applicable information. R1 if positive, contains a pointer to the user record that is to be written to the system log dataset.

> If negative, contains a pointer to the LCA indicating either initialization, (both data sets have to be opened), or data set switching is required.

SVC 38 (TTROUTER)

Registers 15, 0, and 1 do not contain any applicable information.

SVC 43 (CIRB)

R15 contains no applicable information. R0 contains the entry point address of the user's asynchronous exit routine. R1 contains option bit flags in the high order halfword and the size of the work area requested (in double words) in the low order halfword.

Flag settings are:

....1..

<u>flaq byte 1</u> 1	DIRB
0	CIRB
.1000	bits 1-4 always set as shown
1	problem program key
0	supervisor key
1.	problem program state
	supervisor state
1	save area for
0	registers requested no save area requested
flag byte 2	
XXXXXX	reserved
1	do not return IQEs at exit

Section	3:	Tracing	Aids	171

return IQEs at exit

SVC 46 (TTIMER)
R15 contains no applicable
information.
R0 contains no applicable information.
R1 the low order three bytes carry
code determining how TTIMER should
work, as follows:

- x'00'
 the time remaining in the
 current tasks time interval
 is to be returned in
 register 0; the interval
 timer is not to be canceled.
- x'01' the current task's time interval is to be canceled.
- x'02' the time interval of a related task is to be canceled.
- <u>SVC 49 (TTOPEN)</u> Registers 15, 0, and 1 do not contain any applicable information.
- <u>SVC</u> 52 (Restart/SMB Reader) Registers 15 and 0 have no applicable information.
 - R1 contents are as follows: If SVC 52 is issued by the Initiator for the purpose of reading SMBs (containing JCL) for an automatic step or checkpoint restart, register 1 points to a job queue DCB, SMB buffer, and general work space.

If SVC 52 is issued from module IEFRSTRT to initiate a check point restart, register 1 contains a pointer to a parameter list.

SVC 59 (OLTEP)

R15 contains no applicable information. R0 contains a pointer to a three word parameter list, which, in turn contains pointers as follows:

Word 1 -- pointer to UCB Word 2 -- pointer to DEB Word 3 -- pointer to IECIOLTS (I/O interrupt handler). R1 contains a call code used to locate the particular OLTEP function requested. The value will be greater than x*00' and equal or less than x'94'. <u>SVC 61 (TSAV)</u> Registers 15 and 0 have no applicable information.

> R1 contains zeroes if the routine is being entered from the Overlay Supervisor.

R1 contains the address of the DCB used to fetch the module (set to a negative value) if the routine is being entered from the Contents Supervisor.

SVC 68 (SYNADAF/SYNADRLS)

Entry from SYNADAF:

R15 contains a flag byte in the high-order position and three bytes of user data or an address of an entry point to the SYNAD routine.

The flag byte contains codes as follows:

00 EXCP request BPAM request 01 02 BSAM request 03 QSAM request 04 **BDAM request** 05 **BISAM request** 06 QISAM request 07 BTAM request 80 QTAM request 09 GAM request

R0 contains, in the three low order bytes, the address of the DECB (BSAM, BPAM, BDAM) or the address of the IOB (BISAM, QISAM, QSAM).

Additionally, when a QSAM request is made, the high-order byte contains the offset of the first CCW in the IOB.

R1 contains a flag byte and the address of the DCB in the high-order byte and the three low-order bytes respectively.

The flag byte bit settings are:

0000000	BISAM and QISAM
1	error caused by input
.1	error caused by
	output
1	error caused by BSP,
	CNTRL, Or POINT
1	record had been
,	successfully read
1	INVALID request
1	PT conversion -
	invalid character
1.	BDAM only - hardware
	error
1	BDAM only - no space
	for record

Entry from SYNADRLS: Registers 0 and 1 have no applicable information.

R15 contains x'FF' in the high-order byte, indicating the SVC routine is being entered from the SYNADRLS macro instruction and three bytes of user data.

SVC 72 (CHATR) Registers 15 and 0 have no applicable information.

> R1 contains the address of a parameter list with the following structure:

<u>Offset</u>

- 0 address of parameter list+8
- 4 address of DCB
- 8 module name for XCTL 16 code for OPEN/CLOSE (1 byte); address of UCM entry (3
- 20 address of UCM 24 address of return

SVC 76 (IFBSTAT)

R15 contains no applicable information.

The content and applicability of Registers 0 and 1 vary with the presence or absence of RDE (Reliability Data Extractor) routines in the control program.

If RDE is present: R0 contains a positive 0 or 8. R1 has no applicable information.

> A positive 0 in R0 indicates that EOD recording is requested; a positive 8 indicates that IPL recording is requested.

If RDE is not present: R0 contains a negative number representing the length in bytes of a record to be placed in the SYS1.LOGREC data set. R1 contains the address of the record to be written.

SVC 79 (STATUS)

R15 has no applicable information.

R0 If ND was specified, the two high-order bytes contain the bits indicating dispatchability. The two low-order bytes contain the function code: 0 cannot be rolled out 1 must complete - step 2 must complete - system 3 non-dispatchable - step 4 non-dispatchable - system 5 non-dispatchable - TCB address 6 stop non-dispatchable 7 start dispatchable R1 the high-order bit indicates SET or RESET: 0 SET(SS)

1 RESET(RS)

The three low-order bytes contain the TCB address if it was specified, or, if not, zero.

SVC 83 (SMFWTM)

Registers 15 and 0 contain no applicable information.

R1 contains a pointer as follows: If positive a pointer to the record that is to be written to the SMF data set.

> if negative a pointer to the SMCA indicating either initialization or processing for a SWITCH command to switch SMF data sets.

SVC 84 (Restart Address Routine)

SVC 84 is issued by the GPS Graphic I/O Control Routine to have the buffer restart address stored in the UCB associated with the display unit for which the routine builds a channel program.

R15 contains no applicable information.

R0 contains the buffer restart address to be stored in the UCB in the high order two bytes. The low order two bytes point to the UCB. R1 contains a zero

SVC 85 (SWAP)

Registers 15, 0, and 1 do not contain any applicable information.

SVC 91 (VOLSTAT)

R15 contains no applicable information. R0 when negative, contains the address of the UCB. Note: If device type is disk go to SVC 91 load 2. R0 when positive, contains the address of the DCB. R1 contents are as follows: if zero, the SVC was issued by CLOSE if X'32', the SVC was issued by DDR if X'33' the SVC was issued by EOD if X'63', the SVC was issued by EOD if any other than the above, the SVC was issued by UNALLOCATION

SVC 92 (TCBEXCP)

R15 contains no applicable information R0 contains the address of the TCB for the <u>issuers</u> task. R1 contains the address of the IOB.

SVC 93 (TGET/TPUT)

Entry from TGET

R15 contains no applicable information

R0 the two high-order bytes are reserved. The two low-order bytes contain the buffer size in bytes.

- R1 contains a flag byte and an address as follows:
 - the high order byte is a flag byte with these bit settings.

1	Denotes "TGET"
	specified
0	Denotes "TPUT"
	specified
.1	Reserved.
	Reserved for TPUT
1	Denotes "NOWAIT"
	specified means that
	control should be
	returned to the
	program that issued
	the TGET whether or
	not an input line is available from the
	available from the
	terminal if no input
1	line is obtained, a
	return code of 4 will
	be found in register
•	15.
0	Denotes "WAIT"
	specified means that
	control will not be
	returned to the
	program that issued
	the TGET until an
	input line has been
	put into the
	program's buffer if
	an input line is not
	available from the
	terminal, the issuing
	program is put into a
	wait state until a
	line does become
	available and is
	placed in the
	program's buffer
1	Reserved for TPUT
1	Reserved for TPUT
10	Reserved for TPUT
01	Denotes "ASIS"
	specified means that
	normal or minimal
	editing will be
	performed.
00	Denotes "EDIT"
	specified means that
	in addition to the
	normal ("ASIS")
	editing, further
	editing will be
	performed.
	-

the low-order three bytes contain the address of the buffer that is to receive the input line. Entry from TPUT R15 contains no applicable information. R0 the two high-order bytes contain the Terminal Job Identifier number; the two low-order bytes contain the size of the input buffer in bytes. R1 contains a flag byte and an address as follows: the high-order byte is a flag byte with these bit settings: 1..... Denotes "TGET" specified 0..... Denotes "TPUT" specified .1..... Reserved ..1.... Denotes "LOWP" specified means that the terminal will not receive any inter-terminal messages if TSBITOFF is on even if a key-zero task is sending the messages may only be specified on a TPUT with TJID. Denotes "HIGHP" specified means that the terminal will receive inter-terminal messages even if TSBITOFF is on if a key-zero task is sending the messages may only be specified on a TPUT with TJID. Denotes "NOWAIT" ...1.... specified means that control should be returned to the program that issued the TPUT whether or not system output buffers are available for the output line if no buffers are available, a return code of 4 will be found in register 15. Denotes "WAIT"

...0....

specified means that control will not be returned to the program that issued the TPUT until the output line has been placed in a system cutput buffer if no

buffers are

....1...

....0...

....1..

....0...

.....10

.....01

available, the issuing program will be put into a wait state until buffers do become available and the output line is placed in them. Denotes "HOLD" specified means that the program that issued the TPUT cannot continue its processing until this output line has been either written to the terminal or deleted. Denotes "NOHOLD" specified means that control should be returned to the program that issued the TPUT as soon as the output line has been placed on the output queue. Denotes "BREAKIN" specified means that output has precedence over input; that is, if the user at the terminal is transmitting, he is interrupted, and this output line is sent any data that was received before the interruption is kept and displayed at the terminal following this output line. Denotes "NOBREAK" specified means that input has precedence over output; that is, the output message will be placed on the output queue to be printed at some future time when the terminal user is not entering a line. Denotes "CONTROL" specified means that this line is composed of terminal control characters and will not print or move the carriage on the terminal. Denotes "ASIS" specified; means that normal or minimal editing will be performed. Denotes "EDIT" specified; means that in addition to the normal ("ASIS")

editing, further editing will be performed.

the low-order three bytes contain the address of the buffer that is to hold the line of output.

SVC 94 (TERMCTL)

Entry from TCLEARQ:

R15 contains no applicable information.

R0 Contents:

Bytes 0 01 -- Entry code 1-3 0 -- Reserved

R1 Contents:

Bytes		
0	80	"INPUT" specified
	00	"OUTPUT" specified
1-3	0	Reserved

Entry from STBREAK:

R15 contains no applicable information.

R0 Contents:

<u>Bytes</u>			
0	04	 Entry	code
1-3	0	 Reserv	red

R1 Contents:

<u>Bytes</u>		
0	80	"YES" specified
	00	"NO" specified
1-3	0	Reserved

Entry from STCOM:

R15 contains no applicable information.

R0 Contents:

<u>Bytes</u>			
0	05	 Entry	code
1-3	0	 Reserv	zed

R1 Contents:

Bytes	
0	80 YES specified
	00 NO specified
1-3	0 Reserved

Entry from STTIMEOU:

R15 contains no applicable information.

<u>Bytes</u>			
0	06	 Entry	code
1-3	0	 Reserv	ved

R1 Contents:

<u>Bytes</u>	
0	80 "YES" specified
	00 "NO" specified
1-3	0 Reserved

Entry from STCC:

R15 contains no applicable information.

R0 Contents:

<u>Bytes</u>		
0	07	 Entry code
1-3	0	 Reserved

R1 Contents:

<u>Bytes</u> 0	Flag byte as follows:
	1 first operand specified
	.1 ATTN specified
	LD specified
	1 CD specified
	00000000 no operands specified,
	retain
	previously-used
	characters.
1	0 Reserved
2	hh line delete control
	character. The
	hexadecimal
	representation of any
	EBCDIC character on the
	terminal keyboard except the new line (NL) and
	carriage return (CR)
	control characters.
	c the character
	representation of any
	EBCDIC character on the
	terminal keyboard.
3	hh character delete
	control character. The
	hexadecimal
	representation of any
	EBCDIC character on the
	terminal keyboard except
	the new line (NL) and
	carriage return (CR)
	characters. c the character
	c the character representation of any
	EBCDIC character on the
	terminal keyboard.
	constant volvourus

Entry from STATTN:

R15 contains no applicable information.

R0 Contents:

Bytes	
0	08 Entry code
1	00 Reserved
2	hh Lines byte. The number of consecutive
	lines of output that can
	be directed to the
	terminal before the
	keyboard will unlock.
	00 Output line
	counting is not used.
3	hh Tens byte. The
	tens of seconds that can
	elapse before the
	keyboard will unlock.
	00 Locked keyboard
	timing is not used.
	criting is not used.

R1 Contents:

00000000 no o spec resu	ES specified 5 specified at address cified pperands cified, alts in a
NOP 1-3 hhhhhh Cha string addres 000000 no string was sp	ss. character

Entry from STAUTOLN:

R15 contains no applicable informtion.

R0 Contents

<u>Bytes</u> 0 1-3	09 Entry code hhhhhh the address	
-		the

R1 Contents:

Bytes	
0	00 Reserved
1-3	hhhhhh the address of
	a fullword containing the
	increment value used in
	assigning line numbers.

Entry from STSIZE:

R15 contains no applicable information.

R0 Contents:

<u>Bytes</u>	0A Entry code
0	0000 Reserved.
1,2	hh lines byte. The
3	number of lines (depth)
	number of lines (depth) that can appear on the screen.

R1 Contents:

<u>Bytes</u>	
0-2	000000 Reserved
3	hh size byte. The logical line size (width) in characters of the terminal.

Entry from GTSIZE, STAUTOCP, SPAUTOPT, RTAUTOPT

R15 contains no applicable information.

R0 Contents:

<u>Bytes</u> 0	Entry codes as follows:		
1-3	0B GTSIZE 0C STAUTOCP 0D SPAUTOPT 0E RTAUTOPT 000000 Reserved		
R1 Contents: No applicable information, will be zeroed.			
Entry from STCLEAR:			
R15 contains information.	no applicable		
R0 Contents:			
<u>Bytes</u> 0 1-3	10 Entry code 000000 Reserved		
R1 Contents:			
Bytes 0	00 Reserved.		

1=3 hhhhhh -- erasure character string address. Entry from TCABEND

R15 contains no applicable information.

R0 Contents:

Bytes			
0	00	 Entry	code
1-3	0	 Reserv	ved

R1 Contents: No applicable information will be zeroed.

Entry from TSABEND

R15 contains no applicable information.

R0 Contents:

Bytes 0 OF -- Entry code 0 -- Reserved 1-3

R1 Contents: No applicable information will be zeroed.

SVC 95 (TSIP)

02

03

04

05

06

07

08

0A

0B

0C

R15 contains no applicable information.

R0 Contents:

<u>Bytes</u>	
0,1	zero or Terminal Job
-	Identifier (TJID) or not
	applicable.
2	00 Reserved
3	Entry code as follows:

Calling Routine Problem Program (TMP) Entry Code 00 01 Timer Second - Level

Interruption Handler

TGET/TPUT

Region Control Task

Dequeue, TIOC (Attention, TSINPUT, TSOUTPUT), Timer SLIH, WTOR

- Region Control Task
 - Enqueue
 - Dequeue
- TSO Dispatcher
- 09 TSO Dispatcher
 - TSO Dispatcher
 - TSO Dispatcher
 - Region Control Task (Quiesce)

0 D	Region Control Task (Quiesce)
0 E	Time Sharing Control Task (Swap)
0 F	Time Sharing Control Task (Swap)
10	Time Sharing Control Task (Swap)
11	Time Sharing Control Task (Swap)
12	Region Control Task (Restore)
13	Region Control Task (Restore)
14-18	Reserved

R1 Contents:

Bytes 0,1,2,3 variable as follows:

Entry Code	Content
00	Address of 8-character
	command name sign-bit:
	0-ended
	1-beginning

01 not applicable

02 Sign-bit: 0-Input 1-Output

> Bytes 3&4: Number of free buffers

- 03-05 not applicable
- 06 Estimated must complete time
- 07-0C not applicable
 - 0D Number of FBQEs
 - 0E Byte 0: Swap Units Byte 1: Swap device code (0,4,8,c)

Bytes 263: Swap size in 2K blocks

0F-13 not applicable

SVC 97 (TEST(TSO))

Entered from: Any module of the tested program, when used as a breakpoint handler. If used as a breakpoint handler the TCBTCP bit is '1' in the current TCB and registers 15, 0, and 1 contain no applicable information.

Any module of the TSO Test Command Processor when used as a subroutine of TSO TEST. In this case the current TCBTCP bit is '0' and registers are as follows: R15 contains no applicable information. R0 Contents: Bytes õ Entry code as follows: 40 -- Set TCBTCP bit to '1' 20 -- Set TCBTCP bit to **'0'** 10 -- Alter TCBTRN field 08 -- Alter second word of RBOPSW field 04 -- Alter specific register in SVC 97's SVRB register save area 04 -- Alter all registers in SVC 97's SVRB register save area 02 -- Alter floating-point register in TCB save area 01 -- Set RB wait count to 0 (zero). 1-3 Address of target TCB, PRB, or IRB R1 Contents: Register 1 contents are variable as follows: **Bvtes** Entry code 0123 not applicable **4**0 entry code 0123 not applicable 20 entry code 0 not applicable 1,2,3 **TCBTRN** value entry code 0 instruction length, 08 completion code program mask 80 address of value for 1,2,3 second word of RBOPSW field. entry code 0 07 1 register number 1,2,3 address of new value x'FF' entry code 0 04 1,2,3 address of 64-byte value entry code 0 floating-point 02 1,2,3 register number address of new value for register entry code 0,1, not applicable 01 2,3

SVC 100

SVC 100 is used by the SUBMIT, OUTPUT, OPERATOR, and CANCEL/STATUS processors.

R1 Content: Contact your FE programming representative for information concerning the content of General Purpose Registers 15, 0, and 1 upon entry to SVC 100. <u>Bytes</u> σ 0 -- zeroed. 1-3 hhhhhh -- variable by SVC 101 (QTIP) SVC 101 is used only by the TSO entry code in R0 as sub-system and the MCP and provides an follows: interface between them for inter-region communication and data 00 -- address of movement. savearea within AVT **R15** Contents: 03 -- not applicable 04-0D -- address of savearea within AVT Bytes OE -- (without entry 0 0 -- zeroed. by entry code in R0 address in R15; 1-3 hhhhhh -- variable by address of savearea in entry code in R0 as AVT) (with entry follows: address in R15; not applicable) **OF-11 --** address of 00 -- not applicable 03 -- entry address of savearea within AVT 12-16 -- not applicable 17 -- zeroed; OTIP0030 within IEDAYAA 04-0D -- not applicable indicates no savearea OE -- (with savearea is being passed 18 -- not applicable 19-1A -- address of address in R1) not applicable. (Without savearea within AVT savearea address in R1) entry address of 1C -- not applicable QTIP0140 within 1D -- address of IEDAYOO savearea within OF-11 -- not applicable TIOCRPT 12-16 -- entry address of IKJGGQT1, branch entry to OTIP SVC 17 -- address of TSB SVC 103 (XLATE) being logged off 18 -- (same as 12-16) R15 contains no applicable 19-1A -- not applicable information. 1C -- entry address of QTIP0280 within R0 contains the length of the field to IEDAYII be translated. 1D -- not applicable. **R1 Contents: R0 Content: Bytes** hh Bytes action byte as 0 -- zeroed. follows: 0 hh -- entry codes as 1-3 follows 80-translate from EBCIDIC to ASCII 00 -- invokes IEDAYAA 00-translate from 03 -- invokes IEDAYAA ASCII to EBCDIC 04 -- invokes IEDAYHH 05-09 -- invokes IEDAYII; 1-3 hhhhhh address of field to 0A -- invokes IEDAYLL; be translated **OB-11 -- invokes IEDAYOO** 12-14 -- invokes IEDAYGP 15-16 -- invokes IEDAYAA; SVC 104 (TCAM) 17 -- invokes IKJGG088 18 -- invokes IEDAYOO; R15 contains no applicable information 19-1A -- IEDAYZZ invoked 1C -- invokes IEDAYII

1D -- IEDAYGP invoked;

R0 indicates the subroutine to be executed as follows:

Bytes 0-3

00000001 IGC0010D entry point routine 00000002 GTFIELDA decode routine 00000003 STTNME operator command addressing routine 00000004 IEDQCA02 scan routine

R1 contains the address of the operator control work area

SVC 105 (IMGLIB)

R15 contains no applicable information

R0 contains no applicable information

R1 indicates actions to be taken as follows:

<u>Bytes</u>

0-3

0000000 construct a DCB and DEB for SYS1.IMGLIB hhhhhhh delete DCB at this address and also the DEB pointed to by the DCB.

SVC 109

Type 3 and type 4 SVC routing routine.

R15 contains an index value, converted to 3 digit EBCDIC number and appended to name IGC00. This routine is then called.

R0/R1 contain no applicable information for SVC 109, contents are to be used by called routine IGX00.

<u>SVC 116</u>

Type 1 SVC routing routine.

R15 contains an index value, used in binary form to index into a table to call other SVC routines.

R0/R1 contain no applicable information for SVC 116, contents are to be used by called routines.

<u>SVC 117</u>

Type 2 SVC routing routine.

R15 contains an index value, used in binary form to index into a table to call other SVC routines.

R0/R1 contain no applicable information for SVC 117, contents are to be used by called routines. SVC Comprehensive Trace Records Group 2 -Basic Fields Plus DDNAME Field

Group 2 SVC comprehensive trace records add a DDNAME field to the fields composing the basic record. The format is:

- *******
 asterisks indicate an error
 occurred while gathering the
 information.
- cccccccc the name of the associated DD statement.

N/A

indicates that the DD name could not be obtained for the following reasons:

The DCB was not opened The DCB TIOT offset was outside the valid range The DEB TCB pointer was set to 0 The TCB TIOT pointer was set to 0 The DD name in the TIOT was not in EBCDIC notation

Following are descriptions of register 15, 0, and 1 content for the Group 2 SVCs.

SVC 24 (DEVTYPE)

R15 contains no applicable information.

R0 contains the address of the output area or the two's compliment of the output area address.

R1 contains the address of the DD name, or the two's compliment of the DD name address.

When control returns from the DEVTYPE SVC routine, the output area will contain 8, 20, or 24 bytes of device data, depending on the value (+ or -) of R0 and R1, and the device type associated with the DDNAME as follows.

	Output	Area	<u>Size</u> (Bytes)
	RPS-DA	DA	Non-DA
R0 and R1 positive R0 negative	20	20	8
and R1 positive R0 and R1	20	20	8
negative	24	20	8

SVC 31 (FEOV)

R15 and R0 contain no applicable information

R1 contains the address of the DCB

SVC 53 (RELEX)

R15 contains no applicable information

R0 contains the address of a parameter list which contains either:

hhhhhhh relative block or TTR MBBCCHHR actual address

R1 contains the address of the DCB

SVC 55 (EOV)

R15 contains no applicable information

R0 contains the IOB address if the following are true:

DCBOFLAGS = ...1.... DCBMACRF = 0..... and R0 is not equal to x'00001000'

R1 contains the DCB address

SVC 57 (FREEDBUF)

R15 contains no applicable information

R0 contains the address of the DECB

R1 contains the address of the DCB

SVC 58 (REQBUF/RELBUF)

R15 contains no applicable information

R0 contains the request count or release address

R1 contains the DCB address

SVC 69 (BSP)

R15 and R0 contain no applicable information

R1 contains the address of the DCB

<u>SVC Comprehensive Trace Records; Group 3 -</u> Basic Fields Plus Parameter List Field

Group 3 SVC comprehensive trace records add a parameter list field to the fields composing the basic record. The parameter list field displays all or a portion of the parameter list being passed to the SVC routine by the caller. The format is:

PLIST { hhhhhhh hhhhhhh hhhhhhh ...}

- N/A indicates that there is no applicable information
- hhhhhhh hhhhhhh ... parameter list display. Content and amount varies with the SVC being traced.
- ******* ********
 indicates that an error occurred
 while gathering the information.

Following are descriptions of register 15, 0, and 1 content, and PLIST content for the Group 3 SVCs.

SVC 1 (WAIT)

R15 contains no applicable information

R0 contains the count of the events being waited on. If zero the wait is treated as a NOP.

R1 if positive, contains the address of the ECB being waited on. If negative, contains the address of a list of ECBs, in two's complement form.

PLIST may contain up to 40 bytes of information. It consists of a list of ECB addresses up to a maximum of 10.

SVC 4 (GETMAIN)

R15 and R0 contain no applicable information.

R1 contains the address of the parameter list passed when the SVC was called. (If R1 is zero there is no parameter list and the PLIST field will not be present.)

料本時半 in ten bytes in length and breaks down as follows:

<u>Bytes</u>

0-3 hhhhhhhh

- a. For a single area request - the length requested.
- b. For a variable request - the address of a doubleword containing the minimum and maximum length requested as shown below:

<u>Bytes</u>

- 0 zero 1,2,3 minimum length 4 zero 5,6,7 maximum length
- c. For a list request the address of a list of GETMAIN length requests (1 word per request) the last word containing x'80' in byte 0.

4

5-7

Hierarchy identifier (optional)

hhhhhhh

hh

- a. For a single area request - the address of a word GETMAIN will initialize as the beginning allocated core area.
- b. For a variable area request - the address of a doubleword which GETMAIN will initialize with the address of the GETMAINed area and the actual length allocated.
- c. For a list area request - the address of a list of words which GETMAIN will initialize with the address of allocated areas.

8 hh Flag byte as follows:

- 00 unconditional single area request
- 20 conditional single area request
- C0 unconditional variable request
- E0 conditional variable request

80 unconditional list request A0 conditional list request

Subpool identification 9 hh

SVC 5 (FREEMAIN)

R15 and R0 contains no applicable information.

R1 contains the address of the parameter list passed when the SVC was called. (If R1 is zero, no list passed, and PLIST will not appear.)

PLIST is 10 bytes in length and breaks down as follows:

Bytes

4-7

9 hh

- 0 3a. For a single area request the length to be freed.
 - b. For a list area request -the address of a list of FREEMAIN length requests (1 word per request), the list word containing x'80' in byte 0.
 - a. For a single area request -- the address of an area to be freed.
 - b. For a list area request -the address of a list of addresses of the areas to be freed.
- 8 hh Flag byte as follows:
 - 00 unconditional single area request
 - 20 conditional single area request
 - 80 unconditional list area request
 - A0 conditional list area request
 - Subpool identification.

SVC 18 (BLDL/FIND - Type D)

R15 contains no applicable information.

R0 contains the address of the parameter list.

R1 contains the address of the DCB and indicates the macro instruction that issued the SVC call; if R1 is positive -- BLDL; if R1 is negative -- FIND.

PLIST The BLDL parameter list is 12 bytes in length:

Bytes

- the numbering entries 0,1 2,3 entry length

- 4-11 the hexadecimal representation of the member name for which the BLDL was issued.
 - The FIND parameter list is 8 bytes in length:

Bytes

0-7 the hexadecimal representation of the member name for which the FIND was issued.

SVC 19,20,22,23 (OPEN, CLOSE, OPENJ, TCLOSE)

R15 and R0 contain no applicable data.

R1 contains the address of the parameter list.

PLIST is up to 40 bytes in length and consists of a series of 4-byte entries (up to 10). Each entry breaks down as follows:

below:

Option byte as shown

Bytes ō hh

Bits

1		Last Entry	indicator
.011		LEAVE	
.001		REREAD	
.100		REWIND	
.010		IDLE	
.000		DISP	
	0000	INPUT	
	1111	OUTPUT	
	0011	INOUT	
	0111	OUTIN	
	0100	UPDAT	
	0001	RDBACK	

1-3 hhhhhh DCB address

SVC 35 (WTO/WTOR)

R15 contains no applicable information.

R0 contains console source ID.

R1 contains the address of the parameter list being passed to the svc.

PLIST is 12 bytes in length for WTO and 20 bytes in length for WTOR.

The PLIST field for WTO breaks down as follows:

Bytes

00-- indicates WTO 0 parameter list.

1	hh	message	length	plus
	four.		-	_

- 2.3 hhhh-- MCS flag bytes; bit settings as follows:
 - Byte 2

1	Invalid entry Message is to be queued to the console whose source ID is passed in Register 0.
	the WTO is an immediate command response.
1	the WTO macro instruction is a reply to a WTOR macro
1.	instruction. Message should be broadcast to all active consoles.
1	Message queued for hard copy only.
1	Message queued unconditionally to the console whose source ID is passed in register 0.

Byte 3

1	time is not appended
	to the message.
.1111	Invalid entry
••••1••	message is not queued
	for hard copy
11	invalid entry

4-11 First eight bytes of message

The PLIST field for WTOR breaks down as follows:

<u>Bytes</u>	
0	hhlength of reply
1-3	hhhhhhaddress of
	reply buffer
4-7	hhhhhhhaddress of
	reply ECB
8	00zeroed
9	hhmessage length plus four
10,11	hhhhMCS flag bytes, see WTO PLIST
12-19	first eight bytes of message.

SVC 37 (SEGLD/SEGWT)

R15 contains no applicable information.

R0 if zero, entry was from SEGLD; non-zero indicates entry from SEGWT. R1 contains the address of the parameter list.

PLIST is 12 bytes in length and breaks down as follows:

<u>Bytes</u>

0-3	hhhhhhh branch instruction
	(to SVC 45)
4-7	hhhhhhh address of
	Referred-to Symbol
8	hh "To" segment number
9-11	hhhhhh Previous caller or 0

SVC 39 (LABEL)

R15 and R0 contain no applicable information.

R1 contains the address of the parameter list.

PLIST is 20 bytes in length and breaks down as follows:

Bytes

0-2	c00004 REWIND option c00006 UNLOAD option
3	hh relative UCB in TIOT
	to use for mounting purposes.
4-7	hhhhhhh address of 8 byte DDNAME for DD card
	that allocates devices
	for mounting tapes.
8-11	hhhhhhhaddress of
	volume label set.
12,13	hhhh length of one
	volume label.
14	hh number of labels
	in volume label set
15	hh command byte of
	control CCW
16-19	hhhhhhh address of
	the first 10 bytes of
	volume header label.
	volume neader label.

SVC 40 (Extract)

R15 and R0 contain no applicable information.

R1 contains the address of the parameter list.

PLIST is 12 bytes in length and breaks down as follows:

<u>Bytes</u>

0		Reserved
1-3	hhhhhh	address of list area in which the extracted information will be stored.
4	00	Reserved

5-7

8

hhhhhh

flags byte; indicates hh the fields to be extracted as follows: Bits 1..... address of the general register save area .1..... address of floating point register save area ..1.... reserved1.... address of end-oftask exit routine1... limit priority & dispatching priority1.. task completion code1. address of TIOT1 address of the command scheduler communication list in the CSCB

0^0000 EXTRACT will obtain

information from the current TCB and/or its related control blocks.

address of TCB from

which EXTRACT is to get requested information.

9

TSO only flags byte; indicates the TSO

fields to be extracted

hh

Bits 1..... address of timesharing flags in TCB .1..... address of protected storage control block ..1.... terminal job identifier for task ...xxxxx reserved

as follows:

10,11 0000 reserved

SVC 45 (OVLYBRCH)

R15 contains the address of the Entry Table entry which caused the SVC to be issued.

R0 and R1 contain no applicable information.

PLIST is 12 bytes in length and breaks down as follows:

0-3	hhhhhhh	Branch (inst. to
4-7	hhhhhhh	SVC 45) address of
8	hh	Referred-to-Symbol "To" segment number
9-11	hhhhh	Previous caller or 0

SVC 47 (STIMER)

R0 contents:

Bytes		
0	hh	STIMER option byte
		as follows:
		x'40' TOD option
		x'30' DINTVL option
		x'10' BINTVL option
		x'00' TUINTVL option
1-3	hhhhhh	exit address

R15 contains the address of the time value

PLIST is four or eight bytes in length depending on the option in force:

- a. For the DINTVL and TOD options PLIST is eight bytes in length and represents the time value.
- b. For the BINTVL and TUNINTVL options PLIST is 4 bytes in length and represents the time value.

SVC 48 (DEQ)

R15 and R0 contain no applicable information.

R1 contains the address of the parameter list.

PLIST is 16 bytes in length and breaks down as follows:

<u>Bytes</u>		
0	hh	if set to x'FF'
		indicates the last
		element in the
		parameter list.
		Otherwise no
		meaning.
1	hh	the length of the
		minor name whose
		address is in bytes
		8, 9, 10 and 11 of
		this element.
	00	the length of the
		minor name is in the
		first byte of the
		minor name field
		whose address is in
		bytes 8, 9, 10, and
		11 of this element
		(does not include
		length byte itself).
2	hh	DEQ parameters byte
		as follows:
	Dit Octi	

<u>Bit Settings</u> 0..... Exclusive request 1..... Shared request .0..... MINOR name is known only to job step

Section 3: Tracing Aids 185

		.1	the scope of minor name is SYSTEM		.0	MINOR name is known only to job step
			Set must complete		.1	the scope of minor
			equal to SYSTEM			name is SYSTEM
		••• . • • • •	Set must complete equal to STEP		••1••••	Set must complete equal to SYSTEM
			RET=NONE		1	Set must complete
			RET=HAVE			equal to STEP
		010	RET=CHNGE			RET=NONE RET=HAVE
		-	RET=TEST			RET=CHNGE
		1	RELEASE		011	RET=USE
	3	hh	return code field for codes returned			RET=TEST
			to the issuer by DEQ	3	1 hh	return code field
	4-7	hhhhhhh	address of major	-		for codes returned
			resource name	. 7	LLLLLLL	to the issuer by ENQ
	8-11	hhhhhhh	(QNAME) address of minor	4-7	hhhhhhh	address of major resource name
	V 11		resource name			(QNAME)
			(RNAME)	8-11	hhhhhhh	address of minor
	12-15	hhhhhhh	if the DEQ			resource name
			parameters byte bit 4 (RELEASE) is set	12-15	hhhhhhh	(RNAME) if the ENO
			on this word			parameters byte bit
			contains the UCB			4 (RESERVE) is set
			address; otherwise the content of this			on, this word contains the UCB
			word is			address; otherwise
			unpredictable.			the content of this
SVC	56 (ENQ)	•				word is unpredictable.
510			n no applicable			unprediceubie.
	informa			() (00)		
	TULOTING	ition	SVC	60 (STA	E/STAI)	
					_	colicable information
	R1 cont		ddress of the	R15 co R0 con	ntains no a tents:	pplicable information
	R1 cont paramet	tains the ad ter list	ddress of the	R15 co R0 con 0	ntains no aj tents: 0 Create	pplicable information
	R1 cont paramet PLIST :	tains the ad ter list		R15 co R0 con 0 0	ntains no a tents:	
	R1 cont paramet PLIST : down as	tains the adter list	ddress of the	R15 co R0 con 0 0 0	ntains no aj tents: 0 Create 4 Cancel 8 Overlaj	Ŷ
	R1 cont paramet PLIST :	tains the ad ter list is 16 bytes s follows:	ddress of the	R15 co R0 con 0 0 R1 con	ntains no aj tents: 0 Create 4 Cancel 8 Overlag tains the ag	Y ddress of the
	R1 cont paramet PLIST : down as <u>Bytes</u>	tains the adter list	ddress of the in length and breaks	R15 co R0 con 0 0 R1 con parame	ntains no aj tents: 0 Create 4 Cancel 8 Overlaj tains the ac ter list.	Ŷ
	R1 cont paramet PLIST : down as <u>Bytes</u>	tains the ad ter list is 16 bytes s follows:	ddress of the in length and breaks if set to x'FF' indicate the last element in the	R15 co R0 con 0 0 R1 con parame	ntains no a tents: 0 Create 4 Cancel 8 Overlay tains the ac ter list.	Y ddress of the The high-order <u>bit</u> is
	R1 cont paramet PLIST : down as <u>Bytes</u>	tains the ad ter list is 16 bytes s follows:	ddress of the in length and breaks if set to x'FF' indicate the last element in the parameter list.	R15 co R0 con 0 0 R1 con parame set to was co	ntains no aj tents: 0 Create 4 Cancel 8 Overlag tains the ac ter list. 5 one if the ded.	Y ddress of the The high-order <u>bit</u> is XCTL=YES parameter
	R1 cont paramet PLIST : down as <u>Bytes</u>	tains the ad ter list is 16 bytes s follows:	ddress of the in length and breaks if set to x'FF' indicate the last element in the	R15 co R0 con 0 0 R1 con parame set to was co PLIST	ntains no aj tents: 0 Create 4 Cancel 8 Overlag tains the ac ter list. 5 one if the ded.	Y ddress of the The high-order <u>bit</u> is XCTL=YES parameter tes in length and
	R1 cont paramet PLIST : down as <u>Bytes</u>	tains the ad ter list is 16 bytes s follows:	ddress of the in length and breaks if set to x'FF' indicate the last element in the parameter list. Otherwise no meaning. the length of the	R15 co R0 con 0 0 R1 con parame set to was co PLIST breaks	ntains no a tents: 0 Create 4 Cancel 8 Overlay tains the ac ter list. one if the ded. is eight by	Y ddress of the The high-order <u>bit</u> is XCTL=YES parameter tes in length and
	R1 cont paramet PLIST : down as <u>Bytes</u> 0	tains the ad ter list is 16 bytes 5 follows: hh	ddress of the in length and breaks if set to x'FF' indicate the last element in the parameter list. Otherwise no meaning. the length of the minor name whose	R15 co R0 con 0 0 R1 con parame set to was co PLIST breaks	ntains no a tents: 0 Create 4 Cancel 8 Overlay tains the ac ter list. one if the ded. is eight by	Y ddress of the The high-order <u>bit</u> is XCTL=YES parameter tes in length and llows:
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	R1 cont paramet PLIST : down as <u>Bytes</u> 0	tains the ad ter list is 16 bytes 5 follows: hh	in length and breaks if set to x'FF' indicate the last element in the parameter list. Otherwise no meaning. the length of the minor name whose address is in bytes 8, 9, 10, and 11 of this element.	R15 co R0 con 0 0 R1 con parame set to was co PLIST breaks	ntains no a tents: 0 Create 4 Cancel 8 Overlay tains the ac ter list. one if the ded. is eight by	Y ddress of the The high-order <u>bit</u> is XCTL=YES parameter tes in length and llows: flag byte as follows: x '80' for STAI
	R1 cont paramet PLIST : down as <u>Bytes</u> 0	tains the ad ter list is 16 bytes 5 follows: hh	ddress of the in length and breaks if set to x'FF' indicate the last element in the parameter list. Otherwise no meaning. the length of the minor name whose address is in bytes 8, 9, 10, and 11 of this element. the length of the	R15 co R0 con 0 0 R1 con parame set to was co PLIST breaks	ntains no a tents: 0 Create 4 Cancel 8 Overlay tains the ac ter list. one if the ded. is eight by	Y ddress of the The high-order <u>bit</u> is XCTL=YES parameter tes in length and llows: flag byte as follows: x '80' for STAI processing
	R1 cont paramet PLIST : down as <u>Bytes</u> 0	tains the ad ter list is 16 bytes 5 follows: hh	in length and breaks if set to x'FF' indicate the last element in the parameter list. Otherwise no meaning. the length of the minor name whose address is in bytes 8, 9, 10, and 11 of this element. the length of the minor name is in the	R15 co R0 con 0 0 R1 con parame set to was co PLIST breaks	ntains no a tents: 0 Create 4 Cancel 8 Overlay tains the ac ter list. one if the ded. is eight by	Y ddress of the The high-order <u>bit</u> is XCTL=YES parameter tes in length and llows: flag byte as follows: x '80' for STAI processing x '20' for STAE
	R1 cont paramet PLIST : down as <u>Bytes</u> 0	tains the ad ter list is 16 bytes 5 follows: hh	ddress of the in length and breaks if set to x'FF' indicate the last element in the parameter list. Otherwise no meaning. the length of the minor name whose address is in bytes 8, 9, 10, and 11 of this element. the length of the	R15 co R0 con 0 0 R1 con parame set to was co PLIST breaks	ntains no a tents: 0 Create 4 Cancel 8 Overlay tains the ac ter list. one if the ded. is eight by	y ddress of the The high-order <u>bit</u> is XCTL=YES parameter tes in length and llows: flag byte as follows: x '80' for STAI processing x '20' for STAE processing If zero, the
	R1 cont paramet PLIST : down as <u>Bytes</u> 0	tains the ad ter list is 16 bytes 5 follows: hh	in length and breaks if set to x'FF' indicate the last element in the parameter list. Otherwise no meaning. the length of the minor name whose address is in bytes 8, 9, 10, and 11 of this element. the length of the minor name is in the first byte of the minor name field whose address is in	R15 co R0 con 0 0 R1 con parame set to was co PLIST breaks <u>Bytes</u> 0	ntains no aj tents: 0 Create 4 Cancel 8 Overlay tains the ac ter list. 7 one if the ded. is eight by down as fo	y ddress of the The high-order <u>bit</u> is XCTL=YES parameter tes in length and llows: flag byte as follows: x '80' for STAI processing x '20' for STAE processing If zero, the 'CAMCE:' operand is
	R1 cont paramet PLIST : down as <u>Bytes</u> 0	tains the ad ter list is 16 bytes 5 follows: hh	in length and breaks if set to x'FF' indicate the last element in the parameter list. Otherwise no meaning. the length of the minor name whose address is in bytes 8, 9, 10, and 11 of this element. the length of the minor name is in the first byte of the minor name field whose address is in bytes 5, 9, 10, and	R15 co R0 con 0 0 R1 con parame set to was co PLIST breaks <u>Bytes</u> 0	ntains no aj tents: 0 Create 4 Cancel 8 Overlay tains the ac ter list. 7 one if the ded. is eight by down as fo	y ddress of the The high-order <u>bit</u> is XCTL=YES parameter tes in length and llows: flag byte as follows: x '80' for STAI processing x '20' for STAE processing If zero, the 'CAMCE:' operand is in effect; otherwise
	R1 cont paramet PLIST : down as <u>Bytes</u> 0	tains the ad ter list is 16 bytes 5 follows: hh	in length and breaks if set to x'FF' indicate the last element in the parameter list. Otherwise no meaning. the length of the minor name whose address is in bytes 8, 9, 10, and 11 of this element. the length of the minor name is in the first byte of the minor name field whose address is in bytes 1, 9, 10, and 11 of this element (does not include	R15 co R0 con 0 0 R1 con parame set to was co PLIST breaks <u>Bytes</u> 0	ntains no aj tents: 0 Create 4 Cancel 8 Overlay tains the ac ter list. 7 one if the ded. is eight by down as fo	y ddress of the The high-order <u>bit</u> is XCTL=YES parameter tes in length and llows: flag byte as follows: x '80' for STAI processing x '20' for STAE processing If zero, the 'CAMCE:' operand is
	R1 cont paramet PLIST : down as <u>Bytes</u> 0	tains the ad ter list is 16 bytes s follows: hh hh	in length and breaks if set to x'FF' indicate the last element in the parameter list. Otherwise no meaning. the length of the minor name whose address is in bytes 8, 9, 10, and 11 of this element. the length of the minor name is in the first byte of the minor name field whose address is in bytes 1, 9, 10, and 11 of this element (does not include length byte itself).	R15 co R0 con 0 0 R1 con parame set to was co PLIST breaks <u>Bytes</u> 0	ntains no a tents: 0 Create 4 Cancel 8 Overlay tains the ac ter list. 7 one if the ded. is eight by down as for	y ddress of the The high-order <u>bit</u> is XCTL=YES parameter tes in length and llows: flag byte as follows: x '80' for STAI processing x '20' for STAE processing If zero, the 'CAMCE:' operand is in effect; otherwise this is the address of the STAE/STAI exit routine.
	R1 cont paramet PLIST : down as <u>Bytes</u> 0	tains the ad ter list is 16 bytes 5 follows: hh	in length and breaks if set to x'FF' indicate the last element in the parameter list. Otherwise no meaning. the length of the minor name whose address is in bytes 8, 9, 10, and 11 of this element. the length of the minor name is in the first byte of the minor name field whose address is in bytes 1, 9, 10, and 11 of this element (does not include length byte itself). ENQ parameters byte	R15 co R0 con 0 0 R1 con parame set to was co PLIST breaks <u>Bytes</u> 0	ntains no aj tents: 0 Create 4 Cancel 8 Overlay tains the ac ter list. 7 one if the ded. is eight by down as fo	y ddress of the The high-order <u>bit</u> is XCTL=YES parameter tes in length and llows: flag byte as follows: x '80' for STAI processing x '20' for STAE processing If zero, the 'CAMCE:' operand is in effect; otherwise this is the address of the STAE/STAI exit routine. address of the exit
	R1 cont paramet PLIST : down as <u>Bytes</u> 0	hh	in length and breaks if set to x'FF' indicate the last element in the parameter list. Otherwise no meaning. the length of the minor name whose address is in bytes 8, 9, 10, and 11 of this element. the length of the minor name is in the first byte of the minor name field whose address is in bytes 5, 9, 10, and 11 of this element (does not include length byte itself). ENQ parameters byte as follows:	R15 co R0 con 0 0 R1 con parame set to was co PLIST breaks <u>Bytes</u> 0	ntains no a tents: 0 Create 4 Cancel 8 Overlay tains the ac ter list. 7 one if the ded. is eight by down as for	y ddress of the The high-order <u>bit</u> is XCTL=YES parameter tes in length and llows: flag byte as follows: x '80' for STAI processing x '20' for STAE processing If zero, the 'CAMCE:' operand is in effect; otherwise this is the address of the STAE/STAI exit routine. address of the exit routine parameter list; if zero no
	R1 cont paramet PLIST : down as <u>Bytes</u> 0	hh <u>Bit Sett</u>	in length and breaks if set to x'FF' indicate the last element in the parameter list. Otherwise no meaning. the length of the minor name whose address is in bytes 8, 9, 10, and 11 of this element. the length of the minor name is in the first byte of the minor name field whose address is in bytes 1, 9, 10, and 11 of this element (does not include length byte itself). ENQ parameters byte as follows: ings	R15 co R0 con 0 0 R1 con parame set to was co PLIST breaks <u>Bytes</u> 0	ntains no a tents: 0 Create 4 Cancel 8 Overlay tains the ac ter list. 7 one if the ded. is eight by down as for	y ddress of the The high-order <u>bit</u> is XCTL=YES parameter tes in length and llows: flag byte as follows: x '80' for STAI processing x '20' for STAE processing If zero, the 'CAMCE:' operand is in effect; otherwise this is the address of the STAE/STAI exit routine. address of the exit routine parameter list; if zero no exit routine
	R1 cont paramet PLIST : down as <u>Bytes</u> 0	tains the address the second ter list is 16 bytes follows: hh hh 00 $\frac{\text{Bit Sett:}}{0}$	in length and breaks if set to x'FF' indicate the last element in the parameter list. Otherwise no meaning. the length of the minor name whose address is in bytes 8, 9, 10, and 11 of this element. the length of the minor name is in the first byte of the minor name field whose address is in bytes 5, 9, 10, and 11 of this element (does not include length byte itself). ENQ parameters byte as follows:	R15 co R0 con 0 0 R1 con parame set to was co PLIST breaks <u>Bytes</u> 0	ntains no a tents: 0 Create 4 Cancel 8 Overlay tains the ac ter list. 7 one if the ded. is eight by down as for	y ddress of the The high-order <u>bit</u> is XCTL=YES parameter tes in length and llows: flag byte as follows: x '80' for STAI processing x '20' for STAE processing If zero, the 'CAMCE:' operand is in effect; otherwise this is the address of the STAE/STAI exit routine. address of the exit routine parameter list; if zero no

SVC 63 (CHKPT)

R15 and R0 contain no applicable info.

R1 contents: a. the address of the parameter list b. Zero if a CANCEL request

PLIST is eight bytes in length and breaks down as follows:

Bytes

0	00	check ID address
		provided via the second
		parameter of CHKPT
		macro instruction
	80	No check ID address
		provided
1-3	hhhhhh	address of checkpoint
		DCB
4	00	check ID address not
		provided
	01 to 1	Ocheck ID length
		provided via third
		parameter of the CHKPT
		macro instruction
	FF	"S" specified as third
		manage at an af our DO

- parameter of CHKPT macro instruction; the system generated check ID is to be placed at the address specified in bytes 5-7
- 5-7 hhhhhh address for storing system generated check ID or address of user provided check ID

SVC 64 (RDJFCB)

R15 and R0 contain no applicable information

R1 contains the address of the parameter list

PLIST is up to forty bytes in length and consists of a series of 4-byte entries containing the DCB address. The high-order byte has bit 0 set to one to indicate the last entry.

SVC 70 (GSERV)

R15 and R0 contain no applicable information.

R1 contents:

Bytes

0	hh	is a mask indicating which bits in the
1-3	hhhhhh	Graphic Control Byle (GCB) should be reset. the address of a
1 5		fullword field that identifies the DCB related to the GCB in
		which bits are to be reset.

PLIST is four bytes in length and displays the fullword pointed to by R1. Byte 0 is a unit index factor used to locate the UCB address in the DEB associated with the DCB. (The GCB to be reset is in the UCB).

SVC 73 (SPAR)

R15 and R0 contain no applicable information

R1 contains the address of the parameter list

PLIST is up to 40 bytes in length and consists of a series of 4-byte entries. The first entry breaks down as follows:

Bytes

0	hh	the priority specified
		for the attention
		routine by the SPAR
		macro instruction.
1	hh	Reserved
2,3	hhhh	the number of words in
		the parameter list.
Each a	dditiona	1 entry contains a GACB
addres	s as spe	cified by the SPAR

SVC 74 (DAR)

macro.

R15 and R0 contain no applicable information.

R1 contains the address of the parameter list.

PLIST is up to forty bytes in length, consisting of 4-byte entries. The first entry breaks down as follows:

Bytes			
0,1	hh	Reserved	
2,3	hh	the number of words	in
		the parameter list.	

Each additional entry contains a GACB address specified by the DAR macro.

SVC 77 (ONLT)

R15 contains the address of the UCB of the line for the terminal being tested.

R0 contains the address of the first of five '9's in the test request buffer for ONLT (five '9's' indicate a request for an online test).

R1 contains the address of the parameter list.

PLIST is 14 bytes in length and breaks down as follows:

Bytes		
0-3	hhhhhhh	address of the ECB
		and the prefix of the
		request buffer.
4-7	hhhhhhh	address of the
		GETMAIN parameters
		and terminal test
		pattern table.
8-11	hhhhhhh	address of special
		line control
		characters
12	hh	00 means test is
		valid
		01 means test is
		invalid and not set
		up
13	thh	00 means no answer on
		dial line
		01 means answer on
		dial line

SVC 80 (GJP/GFX)

(The SVC 80 Processing Routine serves as a communication link between GJP routines and the GFX Task, and between the GFX task and ABEND Hook routine.)

R15 contains no applicable information.

R0 contains the address of the parameter list.

R1 contains the address of the console control table.

PLIST is eight bytes in length and breaks down as follows:

Bytes 0 - 3

cccc	indicates which		
	routine passed to SVC		
	80 as follows:		

PLOG -- Log Off PBEG -- Begin Job Processor ABDH -- Abend Hook Routine IERR -- Internal Error Routine NPRO -- Initial Processor

4-7 hhhhhhh the 2250 unit address that indicates which graphic job processor is using the SVC 80 routine.

SVC 87 (DOM) 17-19 hhhhhh R15 contains no applicable information.

RO the value (positive or negative) of R0 determines the content of R1.

R1 If R0 is not negative, R1 contains a message ID word (which is also displayed in the PLIST field).

If R0 is negative, R1 contains the address of a list of message ID words.

PLIST is up to 40 bytes in length, consisting of 4-byte entries. Each entry is a message ID word. The last entry is identified by the 0 bit in the high-order byte being set to 1.

SVC 90 (XQMNGR)

R15 and R0 contain no applicable information.

R1 contains the address of the QMPA.

PLIST is 36 bytes in length and contains the OMPA fields. The OMPA and its associated control blocks are described in the MVT Job Management PLM, Order No. GY28-6660.

SVC 96 (STAX)

R15 and R0 contain no applicable information.

R1 contains the address of the parameter list.

PLIST is 20 bytes in length and breaks down as follows:

Bytes		
0-3	hhhhhhh	address of user
		program to get
		control at attention
		interrupt.
4,5	hhhh	size of input buffer
-		(max 4095)
6.7	hhhh	size of output
•		buffer (max 4095)
8-11	hhhhhhh	address of output
•		buffer
12-15	hhhhhhh	address of input
		buffer
16	hh	STAX option flag
10	1111	~ 7
		byte as follows:
		Bits

1.....Reserved .0....replace=YES .1....replace=NO ..1....defer=YES ...l....defer=NO1111Reserved

address of user

program.

parameters for user

		1 3		mand an her	antin sea
SVC 99 (TSO Dynamic		1-3	nnnnn	as follow	action code
information.	tain no applicable			Action Co	
	e address of the			20,02,10	ECB address
parameter list	•			08,04	Data Address
PLIST is up to Consult your F	40 bytes in length.	4	hh	varies by as follow	action code
representative	e for information				
concerning the field.	e data displayed in this			20	x'80', last four bytes
				80,40,01,	-
SVC 102 (TCAM)				08,04,02, 10	x'00'
R15 and R0 con information.	ntain no applicable				reserved
		5-7	hhhhhh		action code
R1 contains th parameter list	he address of the			as follow	'S :
-				Action Co	de
depending on t) 12 bytes in length the function and breaks			20,02,10	
down as follow	/S :			80.40.01.	address TCB address
Bytes 0 hh	Pation and but for			08,04	Target
0 hh	Action code byte for SVC 102 as follows:				address (for enqueuing an
	1 Flag issuing				element the target
	task not				address is
	eligible for rollout				the address of the
	.1 Post rollout/				disabled
	rollin ECB				ready queue
	complete				in the TCAM AVT).
	or TSO ECB	8	hh	wari og vi	th action code
	complete	0	1111	as follow	
	task not eligible for			Action Co	de
	swap				
	1 Move data across			80,40,20,	10,08, x'80', last
	partition			04,02,01	four bytes
	boundary	9-11	hhhhh	varies wi	th action code
	element on	/ ==		as follow	
	disabled ready queue			Action Co	ode
	and post MCP ECB complete			80,40,01	DEB address
	1. Flag issuing			08,04	Length
	task eligible for swap			10,20,02	address TCB address
	1 Flag issuing				
	task eligible for rollout				

SVC Comprehensive Trace Records; Group 4 -Basic Fields Plus Variable Fields

GTF Group 4 SVC comprehensive trace records have a variety of fields -- differing from SVC to SVC -- added to the fields composing the basic SVC record (Group 1). Format and content of the additional fields for each SVC are discussed in the following material.

SVC 0 (EXCP)

Additional fields -- DDNAME, DCB, DEB.

Register 15, 0, and 1 content, and DDNAME DCB, and DEB format and content follow:

R15 and R0 contain no applicable information.

R1 contains the address of the IOB associated with this request.

- DDNAME ccccccc N/A
 - See explanation of DDNAME field under Group 2.
- DCB hhhhhhhh

address of the DCB associated with this I/O request.

DEB hhhhhhhh

address of the DEB associated with this I/O request.

SVC 6 (LINK)

Additional fields -- PLIST, NAME

Register 15, 0, and 1 content, and PLIST and NAME format and content follow:

R15 contains the address of the parameter list.

R0 and R1 contain no applicable information.

PLIST hhhhhhhh hhhhhhhh is eight bytes in length and breaks down as follows:

> Bytes 0 hh flag byte as follows:

> > 80 DE form of macro instruction 00 EP and EPLOC form of macro instruction

- 1-3 hhhhhh If byte 0 is 80; the address of the directory entry list.
 - If byte 0 is 00; the address of the entry point name.
- 4 hh hierarchy ID as follows:
 - 00 -- no hierarchy 01 -- hierarchy 0 02 -- hierarchy 1
- 5 hhhhhh address of DCB or zero.

NAME cccccccc is the entry point/directory entry (EP/DE) name of the module to be linked to or control transferred to.

SVC 7 (XCTL)

SVC 8 (LOAD)

Additional field -- NAME

R15 contains no applicable information.

R0 Content: If byte 0 contains x'00', bytes 1, 2, and 3 contain the address of the entry point name.

> If byte 0 contains x'80', bytes 1, 2, and 3 contain the address of the directory entry list.

R1 Content: In LCS systems, byte 0 contains the hierarchy ID as follows:

> 00 -- no hierarchy 01 -- hierarchy 0 02 -- hierarchy 1

In systems without LCS byte 0 contains no significant information.

Bytes 1, 2, and 3 contain the DCB address or zero if the default for DCB was specified.

NAME cccccccc is the entry point/directory entry name of the module to be loaded.

SVC 9 (Delete)

Additional field -- NAME

R15 and R1 contain no applicable information.

R0 contains the address of the entry point name.

NAME cccccccc is the entry point name of the module to be deleted.

SVC 13 (ABEND)

Additional field -- CMP CODE

R15 and R0 contain no applicable information.

R1 contains significant information only if SVC 13 was <u>not</u> called by the ABTERM routines. In this case R1 contains the following:

Bytes
0hhFlag byte as follows:Bits1.......1.....STEP option..xxxxxxxreserved

1-3 hhhhhh ABEND completion code

CMP CODE hhhhhhhh

is the ABEND completion code if SVC 13 was called by the ABTERM routines. It is the content of the TCBCMP field of the current TCB at the time the SVC interrupt occurred. If ABEND recursion has occurred this field will contain the recursive completion code.

SVC 14 (SPIE)

Additional field -- PICA

R15 and R0 contain no applicable information.

R1 contains the address of the program interrupt control area (PICA).

PICA hhhhhhh hhhh displays the program interrupt control area from the associated SPIE macro instruction.

SVC 15 (ERREXCP)

Additional fields -- DDNAME, RQE, RQE TCB, CUU hhhh

R15 and R0 contain no applicable information.

R1 contains the address of the Request Queue Element (RQE) which was assigned to this I/O request by IOS.

DDNAME cccccccc is the name of the DD statement associated with this I/O request.

RQE hhhhhhhh hhhhhhh hhhhhhh is the first 12 bytes of the RQE assigned to this request by IOS. The breakdown is:

Bytes		
0,1	hhhh	not applicable
2,3	hhhh	address of the UCB
4	hh	TCB ID for MFT
5,6,7	hhhhhh	address of IOB
8	hh	priority byte
9	hhhhh	address of DEB

RQE TCB hhhhhhhh is the address of the TCB associated with the I/O request.

CUU hhhh device address in channel-unit form of the device associated with this I/O request.

SVC 16 (PURGE)

Additional fields -- DDNAME, DCB, PLIST

R15 and R0 contain no applicable information.

R1 address of the purge parameter list.

 $DDNAME \left\{ \begin{array}{l} N/A \\ cccccccc \\ ******** \end{array} \right\}$

ccccccc

is the name of the DD statement associated with the requests being purged.

- DCB hhhhhhhh is the address of the DCB associated with the purge request.
- PLIST hhhhhhhh hhhhhhh hhhhhhh displays the PURGE parameter list which breaks down as follows:

Bytes			
0	hh	option byt	e as
		follows:	

0... Purge request elements in complete DEB chain starting with DEB specified in address field.

	1	Purge the requests associated with the DEB specified in address field.
	.1	
		Allow the active
	1	request to quiesce. Halt the I/O
	0	operations.
		Purge all requests.
	1	
	0	requests. Purge AEQ, RB and IOS
		logical channel
		queue.
	1	Purge AEQ and IOS
		logical channel
		queue.
		Purge by DEB
	1.	Purge by TCB
1 2		
1-3	hhhhhh	address of DEB.
4	hh	completion code
5-7	hhhhh	address of TCB
8	hh	quiesce indicator:
		01 if one or more
		requests are
		quiescing.

SVC 17 (RESTORE)

9-13 hhhhhh

Additional fields -- DDNAME, DCB, DEB

address of IOB.

R15 and R0 contain no applicable information.

R1 contains the address of a pointer to the chain of IOBs to be restarted.

N/A DDNAME ccccccc *******

ccccccc is the name of the DD statement associated with this IOB.

DCB hhhhhhhh is the address of the DCB associated with the IOB.

DEB hhhhhhhh is the address of the DEB associated with the IOB.

SVC 21 (STOW)

Additional fields -- DDNAME, PLIST

R15 contains no applicable information.

R0 contains the address of the parameter list.

R1 contains the address of the associated DCB.

The values, positive or negative, of R0 and R1, indicate the directory action STOW is to take as follows:

- R0 <u>R1</u> Action +
- ADD + ŧ REPLACE
 - ÷ DELETE
 - CHANGE

N/A DDNAME ccccccc *******

> ccccccc is the name of the associated DD statement.

PLIST

-

hhhhhhh ... (2 or 4 words)

is eight or 16 bytes in length, depending on the directory action being performed:

> For ADD, REPLACE, or DELETE actions the PLIST field is eight bytes long and contains, the member name or alias of the PDS directory entry being acted upon.

For CHANGE the PLIST field is 16 bytes long, the first eight bytes containing the old member name or alias, and the second eight bytes contain the new member name or alias.

SVC 25 (TRKBAI)

Additional fields -- DDNAME, DCBFDAD, DCBTRBAL

R15 and R0 contain no applicable information.

R1 contains the address of the associated DCB. Note: If R1 is negative, the address is in complement form and the DCBFDAD and DCBTRBAL fields are meaningless.

N/A DDNAME { ccccccc

> is the name of the associated DD statement.

DCBFDAD hhhhhhhh hhhhhhhh is the full direct access address (MBBCCHHR) from the DCB pointed to by R1.

	BAL hhhh			2	00	
1	is the track balance the number of bytes remaining on the current track after a write. The field is negative if no bytes			3	00	
t				4	00	Reserved
	remain. FALOG/INI	DEX/LOCATE)		5 -7	hhhhhh	address of the area that contains the data set name
Addit	ional fie	elds PLIS	T, DSN	8	00	Reserved
	nd R0 con mation.	ntain no app	licable	9-11	hhhhhh	the address of the CVOL ID, or zeroed.
param			of the OG or INDEX	12 13 -1 5	00 hbbbbb	address of the volume
			of CAMLST as	15 15		list
genera instru	ated by t	the CAMLST m			from INI	DEX :
call.				Bytes 0	hh	option byte as follows:
:		ata set name				Bits 1 Search is to
	is the pa the SVC i macro ins varies, d	routine by t	st passed to the calling Its content the macro			start on specified CVOL 0 Search is to start on SYS.RES
	from CA	_		1	hh	option byte as follows:
<u>Bytes</u> 0	hh	option byte <u>Bits</u> 1 0 1. 1 1 1	e as follows: Search is to start on specified CVOL Search is to start on SYS.RES Catalog a data set Recatalog a data set Uncatalog a data set Es as follows: Build all missing index	2 3	hh	Bits .1. Build an index 1. Build a generation index 1. Build an alias 1 Delete an index 1 Delete an alias option byte as follows: Bits 1 Disconnect CVOLs Disconnect DELETE option Indicate DELETE option Indicate EMPTY option size of generation data group Bits
		1	levels Delete all unneeded index levels except the high level Indicate presence of DSCB TIR	4 5-7	00 hhhhhh	 a. address of the index name. b. address of an eight byte area that contains a high-level index name.

c.	address of an area
	that contains an
	alias to be deleted.

- 8 00
- 9-11 hhhhhh the address of the area that contains the CVOL ID, or zeroed.
- 12 00
- 13-15 hhhhh a. address of an eight-byte area that contains an alias for a high-level index.
 b. address of a ten-byte area that contains the 4-byte device code of the CVOL to be connected followed by its 6-byte volume serial number.

Entry from LOCATE:

<u>Bytes</u> 0	hh	option byte as follows:
		Bits 1 Search is to start on specified
		CVOL 0 Search is to start on SYS.RES
		1. Read a block by TTR.
		00 0.0. LOCATE a name
1	hh	option byte as follows:
		<u>Bits</u> .000 0000 LOCATE a name
2	hh	option byte as follows:
		<u>Bits</u> 0 LOCATE a name
3	00	
4	00	
5-7	hhhhhh	address of the data set name or the relative track address (TTR) of the desired block in the catalog.
8	00	
9-11	hhhhhh	address of the CVOL ID or zeroes.
12	00	

13-15, hhhhhh address of a 265 byte workarea which must be on a doubleword boundary. If the issuer of LOCATE has a non-zero protect key, then the workarea must have a matching storage protect key.

SVC 27 (OBTAIN)

Additional fields -- PLIST, VOLSER, DSN/CCHHR

R15 and R0 contain no applicable information.

R1 contains the address of the parameter list.

- PLIST hhhhhhhh ... (4 words) displays the OBTAIN parameter list which breaks down as follows:
 - Bytes 0-3 hhhhhhhh operation code as follows:
 - C1000000 SEARCH for DSNAME C1800000 SEEK for track address
 - 4-7 hhhhhhh address of data set name or address of track address of DSCB, CCHHR depending on operation code.
 - 8-11 hhhhhhhh address of the volume serial number
 - 12-15 hhhhhhhh address of 14-byte workarea.

VOLSER CCCCCC

cccccc is the volume serial number of the associated volume.

N/A indicates that the volser pointer in the parameter list was zero.

(nnnnn DSN/CCHHR (ccccccccc ...)

> nnnnn is the track address in EBCDIC notation and is displayed when the operation code in Word 1 of the parameter list indicates SEEK.

cccccc ... is the data set name and is displayed when the operation code in word 1 of the parameter list indicates SEARCH. N/A if the name is unavailable.

SVC 28 (OPENEXT)

Additional fields -- content of R13

R15 contains no applicable information.

RO contains zeroes, or the DCB address of the SYSCTLG to be processed.

R1 contains the UCB address of the volume whose SYSTCLG is to be opened, <u>if R0 contains zeroes</u>.

SVC 29 (SCRATCH)

Additional fields -- PLIST, DSN

R15 contains no applicable information.

R0 contains zeroes; or, the address of a UCB or a SUBUCB (for a 2321 device) for the device upon which unmounted volumes may be mounted.

PLIST hhhhhhhh ... (4 words) displays the SCRATCH parameter list which breaks-down as follows:

> Bytes 0-3 hhhhhhh operation code as follows:

> > 41004000 -- check purge date 41005000 -override purge date

4-7 hhhhhhhh address of data set name

8-11 not used

12-15 address of the volume list

DSN ccccccccc ... is the data set name. N/A if the name is unavailable.

SVC 30 (RENAME)

Additional fields -- PLIST, OLD DSN, NEW DSN

R15 contains no applicable information.

R0 contains the address of the UCB for the device on which unmounted volumes should be mounted, or zero.

R1 contains the address of the parameter list.

PLIST hhhhhhhh (4 words)

displays the RENAME parameter list which breaks-down as follows:

Bytes 0-3 x*41002000* 5-7 hhhhhhhh address of old data set name 8-11 hhhhhhhh 12-15 hhhhhhhh volume list

- OLD DSN ccccc ... is the fully qualified name of the data set to be renamed. N/A if the name is unavailable.
- NEW DSN ccccc ... is the new name for the data set being renamed. N/A if the name is unavailable.

SVC 32 (ALLOCATE)

Additional fields -- CUU, DSN

R15 contains no applicable information

R0 when positive, contains the address of the associated job file control block; when negative (not complemented--high-order bit is set on), contains the address of the associated model DSCB.

R1 contains the address of the UCB list.

CUU ccc is the unit address from the UCB pointed to by R1.

DSN cccccccc ...
is the data set name from the DSN
field of the JFCB or DSCB pointed
to by R0. N/A if the DSN field
was blank.

SVC 33 (IOHALT)

Additional fields -- CUU

R15 and R0 contain no applicable information

R1 contains the address of the UCB associated with the request to be halted.

cu	UU hhhh						01 bytes 13-15
is the device address associated with the device being halted.							contain the address of a listing of subpool
<u>SVC 41</u>	SVC 41 (IDENTIFY)				3-15	hhhhh	numbers. a subpool number
Ad	dditional	fields -	- EPNAME				or address of subpool list
	15 contair nformatior		blicable	1	.6	hh	(determined by byte 12) SHSP flag byte:
	0 contains ddress	s the ent	ry point name				00 bytes 17-19 contain a subpool number 01 bytes 17-19
			n storage address name being added.				contain the address of a list of subpool
EI	PNAME cccc is the added.	e entry p	ooint name being	1	7-19	hhhhh	numbers. à subpool number or address of a subpool list
SVC 42	(ATTACH)						(determined by byte 16)
Ad	ditional	fields -	- SUPRVLIST, PPLIST	2	20	hh	Roll-In/Roll-Out flag:
pa			dress of the ag passed to the SVC				00 new task may not be rolled-out and cannot invoke roll-out.
	0 contains nformation		nificant				01 new task may not be rolled-out but can invoke
pa Ca	R1 contains the address of the parameter list being passed to the called program, or zero (no parameter list being passed).						roll-out 02 new task may be rolled-out but cannot invoke
SUPRVLIST hhhhhhhh (36 bytes) is the parameter list being passed to the SVC routine and breaksdown as follows:						roll-out 03 new task may be rolled-out and can invoke roll-out	
	<u>Bytes</u> 0	hh	EP/DE flag byte:	2	1-23	hhhhhh	address of the end-of-task exit
			00 EP or EPLOC specified 80 DE specified	2	4,25	hhhh	routine dispatching priority number
	1-3	hhhhhh	address of the EP name or directory entry (determined by byte 0).	2	6	hh	limit priority number
	4	hh	hierarchy flag (used if option chosen):	2	7	hh	Key Flags byte as follows:
			00 no hierarchy specified 01 hierarchy 0 02 hierarchy 1				 Reserved Propagate the JSCB field from
	5-7	hhhhhh	address of the DCB; or zero.				the originating task
	8 9-11 12	hh hhhhhh hh	Reserved. address of the ECB GSP flag byte: 00 bytes 13-15 contain subpool number			.1	 If the origina- ting task has a protect key of 0, move the specified JSCB address into the

	0.	••••	attached TCB; otherwise, propagate the originating task's TCBJSCB field Subpools 251 and 252 and the job pack queue pointer of the originating task
	1.		are not given to the attached task. Subpools 251 and 252 of the job pack queue pointer are given to the
	1	••••	attached task. the attached task is to have a protect key
	••••	0	to be shared with other
	••••	1	tasks. Subpool zero is not to be shared
	• • • •	•0••	A save area of 72 bytes is to be obtained for
	••••	.1	the task. No save area is to be obtained.
	••••	• • 0.	TCBJSTCB field from the
	••••	1.	originating task. The TCBJSTCB of the new task is to point to the
	••••	0	new task. The new task is to operate in problem program
	••••	1	mode. The new task is to operate in supervisor mode.
28 - 35	hhhh	na bi	supervisor mode. ne entry point ame for EP; or lank for EPLOC or E specification.

PPLIST hhhhhhhh hhhhhhh hhhhhhh ... (up to 40 bytes) is the parameter list being passed to the called program and consists of a series of four-byte entries, each entry having it's high-order byte reserved, and an address in the low-order three bytes. SVC 44 (CHAP)

Additional fields -- CHAP TCB

R15 contains no applicable information.

R0 contains a signed value to be added to the dispatching priority of the specified task. A negative value will be in two's-complement form.

R1 contains the address of an area containing the address of the TCB whose priority is to be changed; or zero. If zero, it indicates that the active task's priority is to be changed.

CHAP TCB hhhhhhhh is the address of the TCB of the active task at the time the SVC interrupt occurred.

SVC 51 (SNAP)

Additional fields -- PLIST, MODN

R15 and R0 contain no applicable information.

R1 contains the address of the parameter list.

PLIST:

- The PLIST field when SVC 51 is called by the SNAP macro instruction is 12 bytes in length and breaksdown as follows:
- PLIST hhhhhhhh hhhhhhhh hhhhhhh displays three words of the parameter list passed to SVC 51 by SNAP.

Bytes

0	hh	ID number to be printed in the identification
1	00	heading of the dump.

2 hh option flag bytes as follows:

Bits 0... ... ABEND request 1... ... SNAP request .1. ... TCB address given .1. ... Display all supervisor data ...1 ... Display trace table ... 1... Display nucleus 1... Snapshot list is given

.... ..1. ID given 1 Display QCBs 3 hh option flag byte as follows: <u>Bits</u> 1... Save area (see next flags) .0.. Display entire save area .1.. Display heading only ..1. Display registers on entry to ABEND or SNAP ...1 Display link pack area 1... Display job pack area1.. Display PSW on entry to ABEND or SNAP1. Display all subpools less than subpool 128 X Reserved 4 00 5-7 hhhhhh address of DCB 8 00 9-11 hhhhhh address of the TCB specified in the SNAP macro instruction; or zero. If zero, the dump is for the current task. Certain calls for SVC 51 may result in a 16 byte PLIST field being recorded. If there is a problem in this area please contact your FE programming representative for programming support. N/A MODN CCCCCCCC ccccccc is the name of the module calling SVC 51. N/A appears if no module name is available. SVC 54 (DISABLE) Additional fields -- DDNAME, DCB, DEB R15 and R0 contain no applicable information. R1 contains the address of the associated DCB

(N/A DDNAME { ccccccc ******* is the name of the DD statement associated with this request. DCB hhhhhhhh is the address of the associated DCB-DEB hhhhhhhh is the address of the associated DEB. SVC 62 (DETACH) Additional fields -- DETACH TCB R15 and R0 contain no applicable information. R1 contains the address of an area containing the address of the TCB to be detached. Note: If R1 contains zero the DETACH TCB field is meaningless. DETACH TCB hhhhhhhh is the address of the TCB to be detached. SVC 65 (QWAIT) Additional fields -- R2, QCB R15, R0 and R1 contain no applicable information. R2 contains the address of the QCB for the element being waited on. QCB hhhhhhhh hhhhhhh hhhhhhh is the queue control block pointed to by R2, and breakdown as follows: **Bytes** 0 hh queue status: $\overline{01}$ -- not on ready queue 02 -- not waiting 03 -- waiting hhhhhh address of first 1-3 element on the queue. Ш hh priority of the queue when linked onto the ready queue. 5-7 hhhhhh address of the next item on the ready queue. hh reserved. 8 9-11 hhhhhh address of the STCB for the subtask to

be activated.

SVC 66 (BTAM TEST)

Additional fields -- IOBERINF

R15 and R0 contain no applicable information.

R1 contains the address of the IOB pointed to when the SVC was issued.

IOBERINF hhhhhhhh ... (4 words) is the error information field used by BTAM error recovery routines.

SVC 67 (OPOST)

Additional fields -- R2, QCB

R15 and R0 contain no applicable information.

R1 contains the address of the element being posted.

R2 contains the address of the QCB to which the element is being posted.

QCB hhhhhhhh hhhhhhh hhhhhhh is the queue control block pointed to by R2 and breaksdown as follows:

<u>Bytes</u> 0	hh	queue status: 01 not on Ready queue 02 not waiting
		03 waiting
1-3	hhhhhh	address of first

element on the queue. 4 hh priority of the queue when linked onto the ready queue. 5-7 hhhhhh address of the next 7 item on the ready queue. 8 hh reserved. 9-11 hhhhhh address of the STCB for the subtask to be activated.

SVC 71 (ASGNBFR/RLSEBFR/BUFINQ) Additional fields -- DDNAME, PLIST

R15 and R0 contain no applicable information

R1 contains the address of the parameter list.

DDNAME cccccccc is the name of the DD statement associated with the DCB specified by the macro instruction. PLIST hhhhhhhh hhhh ... (up to 12 bytes) displays the parameter list pointed to by R1. The content varies according to the macro instruction calling the SVC. Entry from ASGNBFR: <u>Bytes</u> 04 request byte, 04 0 indicates ASGNBFR the DCB address 1-3 hhhhhh 4-7 hhhhhhh the address of a half-word field containing the number of bytes of buffer to be assigned. Entry from RLSEBFR: Bytes hh request byte: 08 indicates RLSEBFR 0C indicates RLSEBFR ALL 1-3 hhhhhh the DCB address 4-7 hhhhhhh the address of a half-word field containing the number of bytes of buffer to be released. Entry from BUFINO: Bytes 10 request byte, 10 n indicates BUFINQ 1-3 hhhhhh the DCB address 4-7 hhhhhh address of the table of buffer

addresses (must be on a fullword boundary) 8-11 hhhhhhh the number of bytes specified to be available for the table of buffer addresses

SVC 75 (Dequeue Routine)

Additional fields -- IQE

R15 contains no applicable information

R0 contains the address of the next IQE on the IRB active list for the attention routine when ATTNINQ has specified clear mode; otherwise, contains zero.

R1 Content:	<u>Bytes</u> 0-3 hhhhhhhh address of the DCB 4-7 hhhhhhhh EBCDIC character
Bytes 0 hh is a unit index to identify a particular 2260 display station; or	set image ID 8 hh LOAD MODE indicator:
00 for a 2250 station. 1-3 hhhhhh the GACB address	.0 no fold .1 fold x.xx xxxx reserved
N/A IQE hhhhhhhh hhhhhhhh hhhhhhhh	9 hh verification indicator:
when ATTNINQ specifies clear mode this field displays the first 3 words of the IQE pointed to by R0:	1 verify 0 don't verify xxx. xxxx reserved
$\frac{Bytes}{0-3}$ hhhhhhhh the address of the	10 hh data check indicator:
 4-7 hhhhhhhh not meaningful 8-11 hhhhhhh IRB associated with the IQE. 	<pre>1 block .1 unblock 00 as DCB specifies 1 unfold UCS 3211 fold UCS 3211xxxx reserved</pre>
N/A	11-14 hhhhhhh EBCDIC FCB
will appear in this field whenever the ATTNINQ macro instruction did not specify clear mode.	image ID 15 hh FCB parameter options:
	1 verify FCB align .xxx xxx. reserved
SVC_78 (LSPACE)	
Additional fields CUU R15 and R1 contain no applicable information	<u>SVC</u> 82 (DISKANAL) Entered from modules: IEHDANAL, IEHOGETA, IEHDCELL, IEHDLABL, IEHDREST, IEHDDUMP
R0 contains the address of the associated UCB	Additional fields VOLSER, DA-ADDR, PLIST
CUU hhhh is the unit address	R15 and R0 contain no applicable information
SVC 81 (SETPRT)	R1 contains the address of the parameter list.
Additional fields DDNAME, PLIST R15 and R0 contain no applicable information	VOLSER cccccc is the volume serial number
R1 contains the address of the parameter list.	DA-ADDR N/A hhhhhhh hhhhhhhh
DDNAME cccccccc is the name of the DD statement associated with the data set being printed.	displays a six or eight byte track address or N/A, dependent on the options in effect for the SVC routine. The breakdown is:
PLIST hhhhhhhh (four words) is four words of the parameter list being passed to SVC 81 and breaks down as follows:	Option analyze or formatDA-ADDR Content six-byte track addresspost UCBeight-byte trac address

address of N/A alternate track ССНН eight-byte unlabeled volume track address new volume N/A PLIST hhhhhhhh ... (16 bytes maximum) is either 8, 12, or 16 bytes of the parameter list pointed to by The first four bytes always R1. consist of a flag byte, defining the function to be performed, and a 3-byte UCB address. The fifth, ninth, and thirteenth bytes, when present, will contain a flag indicating the last element (4-bytes) in the list. The breakdown is as follows: Bytes 0 hh function byte as follows: 8F -- new volume 1F -- address of alternate track ССНН 00 -- ANALYZE or FORMAT 08 -- POST UCB 88 -- unlabeled volume 1-3 hhhhhh address of UCB (function 8F) 80 flag byte -- last 4 element 5-7 hhhhhh address of DCB (function 1F) 80 flag byte -- last Ш. element 5-7 hhhhhh address of alternate track CCHH (function 00) 4-70 hhhhhhh address of alternate track ССНН 8 80 flag byte -- last element 9-11 hhhhhh address of alternate track information (function 08) 4-7 hhhhhhh address of serial number 8 80 flag byte -- last element 9-11 hhhhhh address of VTOC address of VTOC (function 88) 4-7 hhhhhhh address of serial number

12 flag byte -- last 80 element address of DEB 13-15 hhhhhh SVC 86 (ATLAS) Additional fields -- PLIST, CCHHR R15 and R0 contain no applicable information R1 contains the address of the parameter list PLIST hhhhhhhh hhhhhhh is the parameter list passed to SVC 86 and breaks down as follows: Bytes 0 hh flag byte as follows: 1.... User's channel program can not be re-executed. .xxx xxxx reserved 1 - 3hhhhhh address of IOB 4 hh flag byte as follows: 1.... IEHATLAS is the calling program .1.. a partial count (CCHH only) has been passed by the calling program ..1. a write special CCW is required for a track overflow record ...1 a write special CCW is not required xxxx reserved 5-7 hhhhhh address of count (CCHHR) or partial count (CCHH) field CCHHR hhhhhhhhhh is the five-byte track address of the complete (CCHHR) or partial count (CCHH) field passed by the

hhhhhhh address of VTOC

8-11

calling program. Note: If entry to SVC 86 is from the IEHATLAS program (byte 4, bit 0 in parameter list) this address points to the CCHH part of the count field.

SVC 88 (MOD 88)

Additional fields -- DEB, DSSTAT FLGS, DEVMOD

R15 and R0 contain no applicable information.

R1 contains the address of the DCB associated with the current task at the time the SVC was issued.

- DEB hhhhhhh is the address of the data extent block (taken from DCB pointed to by R1)
- DSSTAT hh the data set status flags field (taken from the DEB)
- DEVMOD hh the device modifier field (taken from the DEB)

SVC 89 (EMSERV)

Additional fields -- PLIST, RESMCW

R15 and R0 contain no applicable information

R1 contains the address of the parameter list

PLIST hhhhhhhh displays four bytes from the parameter list being passed to the SVC routine. The breakdown is:

Bytes

hh flag byte:

CO -- enter emulator mode AO -- leave emulator mode 1-3 hhhhhh address of control storage lead name

RESMCW hhhhhhhh hhhhhhh displays dight bytes of the RESMCW field from the RMS common area.

SVC 98 (TSO PROTECT)

Additional fields -- PLIST, DSN

R15 and R0 contain no applicable information

R1 contains the address of the parameter list

PLIST hhhhhhhh displays the first four bytes of the parameter list as follows:

01	entry code for the
	add function
02	entry code for the
	replace function
03	entry code for the
	delete function
04	entry code for the
	list function
	02 03

1-3 hhhhhh varies by function as follows:

000000 -- add function 000000 -- replace function 000000 -- delete function hhhhhh -- 80 byte buffer address

DSN ccccccc ... is the data set name IMDPRDMP OUTPUT COMMENTS - GTF PROCESSING

The following comments may appear in the listing of GTF trace records.

I/O ERROR ON ddname - CONTINUE

Explanation: The EDIT function of IMDPRDMP is being used to process a GTF external trace data set. An I/O error was encountered while attempting to read the trace data set identified by ddname. Fewer than three consecutive I/O errors have occurred for this data set, so EDIT continues processing, ignoring the current block that caused the I/O error.

I/O ERROR ON ddname - EDIT PROCESSING TERMINATED

Explanation: The EDIT function of IMDPRDMP is being used to process a GTF external trace data set. Three consecutive I/O errors have been encountered while attempting to read the trace data set identified by ddname. EDIT processing terminates.

ERROR IN GTF BUFFER CHAIN

Explanation: The EDIT function of IMDPRDMP is being used to process an internal (dump) trace data set. While attempting to locate the GTF trace buffers, IMDPRDMP encountered one of the following errors:

- A buffer pointer was not on a word boundary.
- A buffer pointer addressed an area of main storage that could not be extracted from the dump for one of the following reasons:
 - The pointer addressed an area higher than the highest address in the dump.
 - IMDPRDMP encountered an I/O error while attempting to read the record containing the area addressed by the pointer.
 - 3. The block containing the addressed area was missing from the dump, perhaps because the program that produced the dump encountered an I/O error while attempting to write the block. EDIT processing is terminated.

ERROR IN GTF BUFFER - CONTINUING WITH NEXT BUFFER

Explanation: The EDIT function of IMDPRDMP is being used to process an internal (dump) trace data set. EDIT has encountered a GTF trace record with a length that does not lie within the acceptable range of 4 to 272 bytes. EDIT continues processing with the next GTF buffer.

GTF NOT ACTIVE AT TIME OF DUMP

Explanation: The Edit function of IMDPRDMP is being used to process an internal (dump) trace data set. EDIT has determined that GTF was not active at the time that the dump was taken. EDIT processing is terminated.

TRACE RECORD LL INVALID, DD ddname BLOCK NO xxxyyy - EDIT PROCESSING TERMINATED

> Explanation: The EDIT function of IDMPRDMP is being used to process a GTF external trace data set. EDIT has encountered a GTF trace record with a length that does not lie within the acceptable range of 4 to 272 bytes. Ddname identifies the GTF external data set being processed; xxxyyy identifies the number of the block containing the faulty record. EDIT processing is terminated.

EDIT TERMINATED UPON USER'S REQUEST

Explanation: A user exit has requested EDIT termination by returning to EDIT with a return code of 24.

EXIT DELETED UPON USER'S REQUEST

Explanation: A user exit has requested that it no longer be invoked during the current EDIT execution. This is the result of a user exit routine return code of 16 or 20.

GTF OPTIONS IN EFFECT - option

Explanation: The input trace data set was created by GTF with trace options in effect as indicated by 'option'. The Service Aids publication describes the options available.

204 Programmer's Guide to Debugging (Release 21)

Appendix A: Debugging With an Operating System Dump

The first facts you must determine in debugging with an operating system dump are the cause of the abnormal termination and whether it occurred in a system routine or a problem program. To aid you in making these determinations, ABEND, SNAP, and indicative dumps provide two vital pieces of information -- the completion code and the active RB queue. Similar information can be obtained from a storage image dump or a stand-alone dump by analyzing PSWs and re-creating an active RB queue.

A completion code is printed at the top of ABEND, SNAP, and indicative dumps. It consists of a system code and a user code. The system code is supplied by the control program and is printed as a 3-digit hexadecimal number. The user code is the code you supplied when you issued your own ABEND macro instruction; it is printed as a 4-digit decimal number. If the dump shows a user code, the error is in your program, and the completion code should lead you directly to the source of error. Normally, however, a system code will be listed; this indicates that the operating system issued the ABEND. Often the system completion code gives enough information for you to determine the cause of the error. The explanations of system completion codes, along with a short explanation of the action to be taken by the programmer to correct the error, are contained in the publication IBM System/360 Operating System: Messages and Codes, GC28-6631.

To locate the load module that had control at the time the dump was issued, find the RB associated with the module. If the dump resulted from an ABEND or SNAP macro instruction, the third most recent RB on the queue represents the load module that had control. The most recent and second most recent RBs represent the ABDUMP and ABEND routines, respectively. Storage image dumps and stand-alone dumps contain PSW information that can be used to identify the load module in control.

Once you have located the RB or load module, look at its name. If it does not have a name, it is probably an SVRB for an SVC routine, such as one resulting from a LINK, ATTACH, XCTL or LOAD macro instruction. To find the SVC number, look at the last three digits of the resume PSW in the previous RB on the queue. If a previous RB does not exist, the RB in question is an SVRB for a routine invoked by an XCTL macro instruction. Register 15 in the extended save area of the RB gives a pointer to a parameter list containing the name of the routine that issued the XCTL.

If the RB does not bear the name of one of your load modules, either an RB was overlaid or termination occurred during execution of a system routine. The first three characters of the name identify the system components.

If the RB bears the name of one of your load modules, you can be reasonably certain that the source of the abnormal termination lies in your object code. However, an access method routine may be at fault. This possibility arises because your program branches to access method routines instead of invoking them through a supervisor-assisted linkage. Thus, an access method routine is not represented on the active RB queue. To ascertain whether an access method routine was the source of the abnormal termination, you must examine the resume PSW field in the RB. If the last 3 bytes in this field point to a main storage address outside your program, check the load list to see if an access method routine is loaded at that address. If it is, you can assume that it, and not your program, was the source of abnormal termination.

Abnormal Termination in System Routines: By analyzing the RB's name field or the SVC number in the previous RB, you can determine which system load module requested the termination. If the RB has a system module name, the first three characters tell you the name of the system component. The remaining characters in the name identify the load module in error.

Remember, although a system routine had control when the dump was taken, a problem program error may indirectly have been at fault. Such a situation might result from an incorrectly specified macro instruction, an FQE modified inadvertently, a request for too much storage space, a branch to an invalid storage address, etc. To determine the function of the load module that had control, consult Appendix C. With its function in mind, the completion code together with an examination of the trace table may help you to uncover which instruction in the problem program incorrectly requested a system function.

Program Check Interruptions in Problem Programs: If you have determined from the completion code or PSWs and evaluation of the RB queue that the dump resulted from a program check in your problem program, examine the status of your program in main storage. (If you have received only an indicative dump, you must obtain either an ABEND/SNAP dump or a stand-alone dump at this point.) Locate your program using pointers in the RB. If its entry point does not coincide with the lower boundary of the program, you can find the lower boundary by adding 32(20) to the address of the RB (systems with MFT). The RB's size field gives the number of doublewords occupied by the RB, the program, and associated supervisor work areas. ABEND/SNAP dumps with MFT have the storage boundaries of the problem program calculated and printed.

Next, locate the area within your program that was executed immediately prior to the dump. To do this, you must examine the program check old PSW. Pertinent information in this PSW includes:

- Bits 12-15: AMWP bits
- Bits 32,33: Instruction length in halfwords.
- Bits 40-63: Instruction address

A useful item of information in the PSW is the P bit of the AMWP bits (bits 12-15). If the P bit is on, the PSW was stored while the CPU was operating in the problem program state. If it is off, the CPU was operating in the supervisor state.

Find the last instruction executed before the dump was taken by subtracting the instruction length from the instruction address. This gives you the address of the instruction that caused the termination. If the source program was written in a higher level language, you must evaluate the instructions that precede and follow the instruction at fault to determine their function. You can then relate the function to a statement in the source program.

Other Interruptions in Problem Programs: If the completion code or PSWs and the active RB queue indicate a machine check interruption, a hardware error has occurred. Call your IBM Field Engineering representative and show him the dump.

If an external interruption is indicated, with no other type of interruption, the dump probably was taken by the operator. Check with him to find out why the dump was taken at this point. The most likely reasons are an unexpected wait or a program loop. If a trace table exists, examine it for the events preceding the trouble or, if the trace table was made ineffectual by a program loop, resubmit the job and take a dump at an earlier point in the program.

The remaining causes of a dump are an error during either execution of an SVC or an I/O interruption. In either case, examine the trace table. Entries in the table tell you what events occurred leading up to termination. From the sequence of events, you should be able to determine what caused a dump to be taken. From here, you can turn to system control blocks and save areas to get specific information. For example, you can find the sense information issued as a result of a unit check in the UCB, a list of the open data sets from the DEB chain, the CCW list from the IOB, the reason for an I/O interrupt in the status portion of the CSW, etc.

Specialized Program Checks

In addition to the error program checks (1-15), other system events cause program checks which are normally transparent to the user. They could, however, if seen in a dump (except ABEND dumps where they do not appear, result in some confusion. One such event is the <u>monitor call interrupt</u>. On 360 CPU's, the monitor call appears as a 01 (operation) interrupt code in the program old PSW. To verify that a simulated monitor call occurred, check the address in the program old PSW. A monitor call occurred if:

- The address (-4) points to an execution instruction ('44');
- The execute is operating on an x'AF00' in low core;
- 3. A NOP (x'470') follows the execute.

370 CPU's support the real monitor call interrupt. The code in the program old PSW is a x'40', and the PSW address (-4) points directly at an x'AF' instruction.

On 360 CPU's, the x'AF' opcode is simulated as follows:

- The first time an x'AF' instruction is encountered, an execute instruction is substituted for the x'AF'.
- 2. The execute is of an instruction in a low-core table (Class Mask Table).
- If the monitor call should occur, the instruction in the Class Mask Table is an x'AF00'; if it should not occur, the instruction is a x'0700' (NOP).

 Required class and ID information for the monitor call are contained in the x'470' NOP following the execute.

On 370 CPU's, the monitor call occurs under control of a mask in Control Register 8.

The Generalized Trace Facility (GTF) is a user of the monitor call interrupt. For more detailed information, refer to the <u>Service Aids Logic PLM</u>, GY28-6721.

Debugging Procedure Summary

- Look at the completion code or PSW printouts to find out what type of error occurred. Common completion codes and causes are explained in Appendix C.
- 2. Check the name of the load module that had control at the time the dump was taken by looking at the active RB's.
- 3. If the name identifies a system routine, proceed to step 4. If the name identifies a problem program and the completion code or PSW indicates a program check, proceed to step 6. If the name identifies a problem program, and the completion code or PSW indicates other than a program check, proceed to step 10.
- 4. Find the function of the system routine using Appendix D.
- 5. If the dump contains a trace table, begin at the most recent entry and

proceed backward to locate the most recent SVC entry indicating the problem state. From this entry, proceed forward in the table, examining each entry for an error that could have caused the system routine to be terminated.

- 6. If the name identifies one of your load modules, check the instruction address and the load list to see if an access method routine last had control. If so, return to step 4.
- 7. Locate your program in the dump.
- 8. Locate the last instruction executed before the dump.
- 9. Examine the instruction and, if the program was written in a high-level language, the instructions around it for a possible error in object code.
- If a machine check interruption is indicated, call your IBM Field Engineering representative.
- 11. If only an external interruption is indicated, ask the operator why he took the dump. Resubmit the job and take a dump at the point where trouble first occurred.
- 12. Examine the trace table, if one is present, for events leading up to the termination. Use trace table entries and/or information in system control blocks and save areas to isolate the cause of the error.

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Register contents at entry to an SVC routine are often helpful in finding pointers and control information. The table below lists SVC numbers in decimal and hexadecimal, and gives the type, associated macro instruction, and significant contents of registers 0 and 1 at entry to each SVC routine.

Decimal No.	Hex. No.	Туре	Macro	Register 0	Register 1
0	0	I	EXCP		IOB address
0	0	I	XDAP		
1	1	I	WAIT	Event count	ECB address
1	1	I	WAITR	Event count	2's complement of ECB address
1	1	I	PRTOV		
2	2	I	POST	Completion code	ECB address
3	3	I	EXIT		
4	4	I	GETMAIN		Parameter list address
5	5	I	FREEMAIN		Parameter list address
6	6	II	LINK		Parameter list address
7	7	II	XCTL		Parameter list address
8	8	II	LOAD	Address of entry point address	DCB address
9	9	I, II	DELETE	Address of program name	
10	A	I	FREEMAIN	Subpool number (byte 0) Length (bytes 1-3)	Address of area to be freed
10	A	I	(R Operand) FREEPOOL		
11	В	I, III	TIME		Time units code
12	с	II	SYNCH		
13	D	IV	ABEND		Completion code
14	Е	II, III	SPIE		PICA address
15	F	I	ERREXCP		Address of request queue

(Part 1 of 5)

Decimal No.	Hex. No.	Туре	Macro	Register 0	Register 1
16	10	III	PURGE	a an an an die der der der der der der der der der de	
17	11	III	RESTORE		IOB chain address
18	12	II	BLDL	Address of build list	DCB address
18	12	11	FIND		
19	13	VI	OPEN		Address of parameter list of DCB addresses
20	14.	VI	CLOSE		Address of parameter list of DCB addresses
21	15	111	STOW	Parameter list address	DCB address
22	16	IV	OPEN TYPE=J		Address of parameter list of DCB addresses
23	17	IV	CLOSE TYPE=T		Address of parameter list of DCB addresses
24	18	III	DEVTYPE		ddname address
25	19	III	TRKBAL		DCB address
26	1A	IV	CATALOG		Parameter list address
26	1A	IV	INDEX		Parameter list address
26	1A	111	LOCATE		Parameter list address
27	1B	III	OBTAIN		Parameter list address
28	1C	IV	CVOL		
29	1D	IV	SCRATCH	UCB address	Parameter list address
30	1E	IV	RENAME	UCB address	Parameter list address
31	1F	IV	FEOV		DCB address
32	20	IV	ALLOC		Address of UCB list
33	21	111	IOHALT		UCB address
34	22	IV	MGCR (MAST CMD EXCP)		
35	23	IV	WTO		Message address
35	23	IV	WIOR		Message address
36	24	IV	WTL		Address of message
37	25	II	SEGLD		Segment name address
37	25	II	SEGWT		Segment name address
38	26	II	TTROUTER		
39	27	111, IV	LABEL		Parameter list address

(Part 2 of 5)

Decimal No.	Hex. No.	Туре	Macro	Register 0	Register 1
40	28	I, II, III	EXTRACT		Parameter list address
41	29	11, 111	IDENTIFY	 Entry point name address	Entry point address
42	2A	II, III	АТТАСН		
43	2B	11, 111	CIRB	Entry point address	Size of work area in doublewords
44	2C	I	СНАР	+ Increase priority - Decrease priority	TCB address
45	2D	11	OVLYBRCH		
46	2E	I	TTIMER		1: Cancel
47	2F	11	STIMER	Exit address	Timer interval address
48	30	I, II	DEQ		QCB address
49	31	111	TEST		
50	32				
51	33	IV	SNAP		Parameter list address
52	34	IV	RESTART		DCB address
53	35	III	RELEX	Key address	DCB address
54	36	II	DISABLE		
55	37	IV	EOV	EOB address	DCB address
56	38	I, II	ENQ	QEL address	QCB address
56	38	I, II	RESERVE		
57	39	III	FREEDBUF	DECB address	DCB address
58	3A	I	RELBUF		DCB address
58	3A	I	REQBUF		DCB address
59	3B	III	OLTEP		
60	3C	III	STAE	 0 Create SCB 4 Cancel SCB 8 0	Parameter list address
61	3D	III	TTSAV		Parameter list address
62	3E	11	DETACH		TCB address
63	3F	IV	СНКРТ		DCB address
64	40	III	RDJFCB		Address of parameter lis of DCB addresses
65	41	II	QWAIT		Parameter list address
66	42	IV	BTAMTEST		

(Part 3 of 5)

Decimal No.	Hex. No.	Type	Macro	Register 0	Register 1
67	43	II	ENDREADY		QPOST
68	44	IV	SYNADAF	Same as register 0 on entry to SYNAD	Same as register 1 on entry to SYNAD
68	44	IV	SYNADRLS		
69	45	III	BSP		DCB address
70	46	II	GSERV		Parameter list address
71	47	III	RLSEBFR		Parameter list address
71	47	III	ASGNBFR		Parameter list address
71	47	III	BUFINQ		Parameter list address
72	48	IV	CHATR		Parameter list address
73	49	III	SPAR		Parameter list address
74	4A	III	DAR		Parameter list address
75	4B	III	DQUEUE		Parameter list address
76	4C	IV	IFBSTAT		
77	4D	IV	QTAMTEST		
78	4E	III	WSCAN		
79	4F	I	STATUS		
80	50	III	IKASVC		
81	51	IV	SETPRT		
. 82	52	IV	DASDR		
83	53	III	SMFWTM		Message address
84	54	I	GRAPHICS	UCB address and buffer restart address	
85	55	IV	DDRSWAP		
86	56	IV	ATLAS		Parameter list address
87	57	III	DOM	If zero If negative 	A DOM message I.D. A pointer to a list of DOM message I.Ds
88	58	III	MOD88	Routine code	DCB address
89	59	III	EMSRV		Parameter list address
90	5A	IV	XQMNGR	Address of list of ECB/IOB pointers (optional)	QMPA address
91	5B	111	VOLSTAT	DCB address	zero: issued by CLOSE Non-zero: issued by EOV

(Part 4 of 5)

Decimal Hex. No. No. 92 5C 93 5D 94 5E 95 5F	Type I IV IV	Macro TCBEXCP TGET/TPUT STERMINAL STATUS	Register 0 TJID & buffer Size	Register 1 Address of User's Buffer
93 5D 94 5E	IV IV	TGET/TPUT STERMINAL	TJID & buffer Size	Address of User's Buffer
94 5E	IV	STERMINAL	TJID & buffer Size	Address of User's Buffer
95 5F	-	OTUTO	Entry code	
1 1 1	I	TSEVENT	TJID/Entry Code or 0	Not Always Applicable
96 60	111	STAX		Parameter List Address
97 61	III	TEST-TSO		
98 62	IV	PROTECT		Parameter List Address
99 63	IV	none		
100 64	III	FIB		
101 65	I	QTIP	Entry code	Parameter List Address
102 66	I	AQCTL		Parameter List Address
103 67		XLATE	Field length	Action byte and field address
104 68	IV	TOPCTL	Subroutine indicator	Address of operator control word area
105 69	111	IMAGLIB		Action indication
109 6D	IV		contents used by calle	ed routine
116 74	I	АТ	contents used by calle	ed routines
117 75	11		contents used by calle	ed routines

(Part 5 of 5)

Appendix C: Completion Codes

Completion codes issued by operating system routines are often caused by problem program errors. This appendix includes the most common system completion codes, their probable causes, and how to correct the error or locate related information using a dump. For a more comprehensive coverage of completion codes, see the publication <u>Messages and Codes</u>.

- 0Cx A program check occurred without a recovery routine. If bit 15 of the old program PSW (PSW at entry to ABEND) is on, the problem program had control when the interruption occurred; "x" reflects the type of error that causes the interruption:
 - <u>x Cause</u> 1 Operation 2 Privileged operation 3 Execute 4 Protection 5 Addressing 6 Specification 7 Data 8 Fixed-point overflow Fixed-point divide 9 Α Decimal overflow Decimal divide B C Exponent overflow D Exponent underflow E Significance
 - F Floating-point

The correct register contents are reflected under the heading "REGS AT ENTRY TO ABEND" in an ABEND/SNAP dump. In a stand-alone dump, register contents can be found in the register save area for ABEND'S SVRB.

0F1 A program check occurred in the interruption handling part of the input/output supervisor. The applicable program check PSW can be found at location 40(28). (In systems with MFT, this PSW is valid only if the first four digits are 0004).

The problem program can be responsible for this code if:

- An access method routine in the problem program storage area has been overlaid.
- An IOB, DCB, or DEB has been modified after an EXCP has been issued, but prior to the completion of an event.

If a trace table exists (trace option was specified at system generation), the instruction address in the new program check PSW, location 104(68), contains the address of a field of register contents. This field includes registers 10 through 9 on an ABEND/SNAP dump, or 10 through 1 on a stand-alone dump.

If no trace table exists, the above field contains registers 10 through 1 on both ABEND/SNAP (MFT only) and stand-alone dumps.

- 0F2 Most frequently caused by incorrect parameters passed to a type I SVC routine.
- 100 A device has been taken off-line without informing the system, or a device is not operational.

If a trace table exists, the most current entry is an SIO entry beginning with 30. The last 3 digits of the first word give the device address.

If a trace table does not exist, register 1 (in the SVRB for the ABEND routine) contains a pointer to the IOB associated with the device.

- 101 The wait count, contained in register 0 when a WAIT macro instruction was issued, is greater than the number of ECBs being waited upon.
- 102 An invalid ECB address has been given in a POST macro instruction.

If a POST macro instruction has been issued by the problem program, the ECB address is given in register 1 of either the trace table entry or the SVRB for the ABEND routine.

If the POST was issued by an I/O interruption handler, the ECB address can be found in the IOB associated with the event.

106 During a transient area load or a dynamic load resulting from a LINK, LOAD, XCTL, or ATTACH macro instruction, the fetch routine found an error. A description of the error is contained in register 15 of ABEND's SVRB register save area:

Appendix C: Completion Codes 215

- 0D The control program found an invalid record type.
- OE The control program found an invalid address. The problem program may contain a relocatable expression that specifies a location outside the partition boundaries.
- 0F A permanent I/O error has occurred. This error can probably be found in the trace table prior to the ABEND entry.

Register 6 of ABEND'S SVRB register save area points to the work area used by the fetch routine. This area contains the IOB, channel program, RLD buffer, and the BLDL directory entry associated with the program being loaded.

- 122 The operator cancelled the job and requested a dump.
- 155 An unauthorized user (a user other than dynamic device reconfiguration) has issued SVC 85. The user's task has been abnormally terminated by dynamic device recognition.
- 200 The error was detected when an I/O operation was requested and the storage protection keys of the IOB, ECB, and DCB were not the same as the key in the DEB. (checked for MVT only)
- 201 This completion code is identical to 102, but applies to the WAIT macro instruction instead of POST.
- 202 An invalid RB address was found in an ECB. The RB address is placed in the ECB when a WAIT macro instruction is issued.
- 213 The error occurred during execution of an OPEN macro instruction for a data set on a direct-access device. Either:
 - The data set control block (DSCB) could not be found on the direct access device.
 - An uncorrectable input/output error occurred in reading or writing the data set control block.

Register 4 contains the address of a combined work and control block area. This address plus x'64' is the address of the data set name in the JFCBDSNM field of the job file control block (JFCB).

- 222 The operator cancelled the job without requesting a dump. The cancellation was probably the result of a wait state or loop.
- 301 A WAIT macro instruction was issued, specifying an ECB which has not been posted complete from a previous event. Either:
 - The ECB has been reinitialized by the problem program prior to a second WAIT on the same ECB, or
 - 2. The high order bit of the ECB has been inadvertently turned on.
- 308 The problem program requested the loading of a module using an entry point given to the control program by an IDENTIFY macro instruction.

Register 0 of LOAD'S SVRB register save area contains the address (or its complement) of the name of the module being loaded.

- 400 The control program found an invalid IOB, DCB, or DEB. Check the following blocks for the indicated information:
 - IOB a valid DCB address.
 - DCB a valid DEB address.
 - DEB ID of OF and a valid UCB address.
 - UCB a valid identification of FF.

Note: In systems with MVT, this code may appear instead of a 200 code, for the reasons given under 200.

406 A program has the "only loadable" attribute or has an entry point given to the control program by an IDENTIFY macro instruction. In either case, the program was invoked by a LINK, XCTL, or ATTACH macro instruction.

> Register 15 of the LINK, XCTL, or ATTACH SVRB register save area contains the address of the name of the program being loaded.

506 The error occurred during execution of a LINK, XCTL, ATTACH, or LOAD macro instruction in an overlay program or in a program that was being tested using the TESTRAN interpreter.

The program name can be found as follows:

 If a LOAD macro instruction was issued, register 0 in the trace table SVC entry or in the SVRB register save area contains the address (or its complement) of the program name.

2. If a LINK, XCTL, or ATTACH was issued, register 15 of the associated SVRB register save area contains the address of a pointer to the program name.

Note: Programs written in an overlay structure or using TESTRAN should not reside in the SVC library.

- 604 During execution of a GETMAIN macro instruction, the control program found one of the following:
 - A free area exceeds the boundaries of the main storage assigned to the task. This can result from a modified FQE.
 - 2. The A-operand of the macro instruction specified an address outside the main storage boundaries assigned to the task.
- 605 During execution of a FREEMAIN macro instruction, the control program found that part of the area to be freed is outside the main storage boundaries assigned to the task, possibly resulting from a modified FQE.

Item 1 under the 604 completion code is also applicable to 605.

606 During execution of a LINK, XCTL, ATTACH, or LOAD macro instruction, a conditional GETMAIN request was not satisfied because of a lack of available main storage for a fetch routine work area. Consequently, the request was not satisfied.

> The name of the load module can be found as described under completion code 506.

- 60A Results from the same situations described under 604 and 605 for R-form GETMAIN and FREEMAIN macro instructions.
- 613 The error occurred during execution of an OPEN macro instruction for a data set on magnetic tape. An uncorrectable input/output error occurred in tape positioning or in label processing.
- 700 A unit check resulted from an SIO issued to initiate a sense command.

The defective device can be determined from the SIO trace table entry that

reflects a unit check in the CSW status.

704 A GETMAIN macro instruction requested a list of areas to be allocated. This type of request is valid only for systems with MVT.

The applicable SVC can be found in a trace table entry or in the PSW at entry to ABEND.

- 705 Results from the same situations described under 704 for FREEMAIN macro instructions.
- 706 During execution of a LINK, LOAD, XCTL, or ATTACH macro instruction, the requested load module was found to be not executable.

The name of the module can be found as described under the completion code 506.

- 804 The error occurred during execution of a GETMAIN macro instruction with a mode operand of EU or VU. More main storage was requested than was available.
- 806 The error occurred during execution of a LINK, XCTL, ATTACH, or LOAD macro instruction.

An error was detected by the control program routing for the BLDL macro instruction. This routine is executed as a result of these macro instructions if the problem program names the requested program in an EP or EPLOC operand. The contents of register 15 indicate the nature of the error:

- X'04' The requested program was not found in the indicated source.
- X'08' An uncorrectable input/output error occurred when the BLDL control program routine attempted to search the directory of the library indicated as containing the requested program.

Register 12 contains the address of the BLDL list used by the routine. In systems with MFT this address plus 4 is the location of the 8-byte name of the requested program that could not be loaded. In systems with MVT, registers 2 and 3 contain the name of the requested module.

- 80A The error occurred during execution of an R-form GETMAIN macro instruction. More main storage was requested than was available.
- 905 The address of the area to be freed (given in a FREEMAIN macro instruction) is not a multiple of eight. The contents of register one in either the trace table entry or ABEND'S SVRB register save area reflect the invalid address.
- 90A Results from the same situations described under 905 for R-form FREEMAIN macro instructions.
- A05 The error occurred during execution of a FREEMAIN macro instruction. The area to be freed overlaps an already existing free area. This error can occur if the address or the size of the area to be freed were incorrect or modified.

The contents of registers 0 and 1 in either the SVC trace table entry or ABEND's SVRB register save area reflect the size and address.

- AOA Results from the same situations described under AO5 for R-form of GETMAIN and FREEMAIN macro instructions.
- B04 This error occurred during execution of a GETMAIN macro instruction. A subpool number greater than 127 was specified. The problem program is restricted to using subpools 0-127. This error can occur if the subpool number was either incorrectly specified or modified.

A displacement of nine bytes from the list address passed to GETMAIN in register 1 contains the subpool number. Register 1 can be found in either the SVC trace table entry or ABEND's SVRB register save area.

- B05 Results from the same situation described under B04 for FREEMAIN macro instructions.
- BOA Results from the same situations described under BO4 and BO5 for R-form of GETMAIN and FREEMAIN macro instructions.

The subpool number can be found in the high order bytes of register 0 in either the SVC trace table entry or ABEND's SVRB register save area.

B37 The error occurred at an end of volume. The control program found that all space on the currently mounted volumes was allocated, that more space was required, and that no volume was available for demounting.

> Either allocate more devices or change the program so that a device will be free when a volume must be mounted.

Fnn An SVC instruction contained an invalid operand; nn is the hexadecimal value of the SVC.

> This error can occur if either an invalid instruction was issued by the problem program or an operand referring to an optional function was not included during system generation.

Appendix D: System Module Name Prefixes

All load modules associated with a specific operating system component have a common prefix on their module names. This appendix lists the module name prefixes and the associated system component(s).

<u>Prefix</u>	Component	<u>Prefix</u>	Component
IBC	Independent utility programs	IFF	Graphic programming support
IEA	Supervisor, I/O supervisor, and NIP	IFG	Close, open, and related routines
IEB	Data set utility programs	IGC	Transient SVC routines
IEC	Input/output supervisor	IGE	I/O error routines
IEE	Master scheduler	IGF	Machine check handler program
IEF	Job scheduler	IHA	System control blocks
IEG	TESTRAN	IHB	Assembler during expansion of supervisor and data management
IEH	System utility programs		macro instructions
IEI	Assembler program during system generation	IHC	FORTRAN library subroutines
IEJ	FORTRAN IV E compiler	IHD	COBOL library subroutines
IEK	FORTRAN IV H compiler	IHE	PL/I library subroutines
IEM	PL/I F compiler	IHF	PL/I library subroutines
IEP	COBOL E compiler	IHG	Update analysis program
IEQ	COBOL F compiler	IHI	Object program originally coded in ALGOL language
IER	Sort/Merge program	IHJ	Checkpoint/restart
IES	Report program generator	IHL	Generalized Trace Facility
IET	Assembler E	IHK	Remote job entry
IEU	Assembler F	IIN	7094 emulator program for the Model 85
IEW	Linkage editor/overlay supervisor/program fetch	110	7074 emulator program on the Models 155 and 165
IEX	ALGOL compiler	IIP	
IEY	FORTRAN IV G compiler	11r	7080 emulator program on the Mod el 165
IEZ	System Interfaces	IIQ	1401/1440/1460 emulator program on Models 135, 145, and 155
IFB	Environment recording routines	IIR	1440/7010 emulator program on
IFC	Environment recording and print routines		Models 145 and 155
IFD	Online test executive program	IIT	709/7090/7094/7094 II emulator program on the Model 165

IIU	7074 emulator program on the Model 155	IKF	USAS COBOL compiler
IKA	Graphic Job Processor	IKJ	Time Sharing
IKD	Satellite graphic job processor messages	ILB	USAS COBOL subroutines

Appendix E: List of Abbreviations

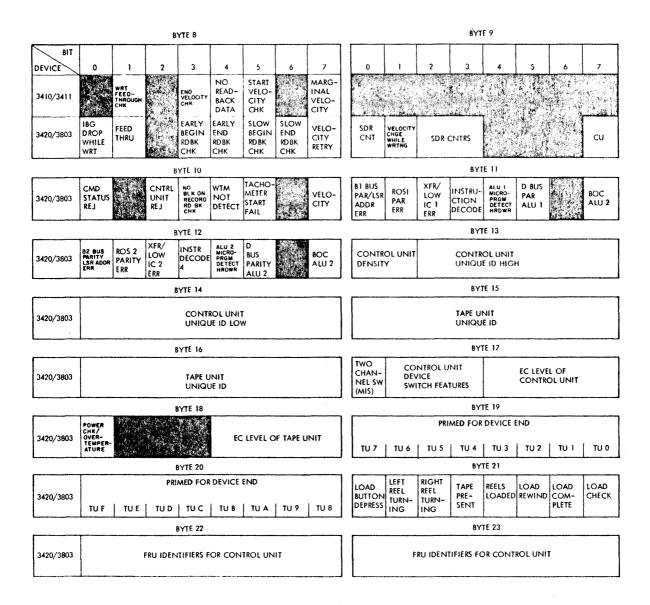
ABEND	abnormal end-of-task	MFT	multiprogramming with a fixed number of tasks
APR	alternate path retry		
CCW	channel command word	MVT	multiprogramming with a variable number of tasks
CDE	contents directory entry	NIP	nucleus initialization program
CPU	central processing unit	PIB	partition information block
CSW	channel status word		
CVT	communications vector table	PQE	partition queue element
DAR	damage assessment routine	PRB	program request block
DCB	data control block	PSA	prefixed storage area
DD	data definition	PSW	program status word
DDR	dynamic device reconfiguration	QCB	queue control block
DEB	data extent block	QEL	queue element
		RB	request block
DPQE	dummy partition queue element	SCB	STAE control block
DQE	descriptor queue element	SIO	start input/output
ECB	event control block	SIRB	supervisor interrupt request block
FBQE	free block queue element	SPOE	subpool queue element
FQE	free queue element	SVC	supervisor call
FRB	finch request block	SVRB	supervisor request block
GQE	gotten subtask area queue element		
IOB	input/output block	SYSOUT	system output
IPL	initial program loading	TCB	task control block
IRB	interrupt request block	TIOT	task input/output table
LLE	load list element	UCB	unit control block
LPRB	loaded program request block	WLE	wait list element
_		XCTL	transfer control
LRB	loaded request block	XL	extent list

Hexadecimal Code	Meaning
7F000000	Channel program has terminated without error. (CSW contents can be useful.)
41000000	Channel program has terminated with permanent error. (CSW contents can be useful.)
42000000	 Channel program has terminated because a direct access extent address has been violated. (CSW contents do not apply.)
44000000	Channel program has been intercepted because of permanent error associated with device end of previous request. You may reissue the intercepted request. (CSW contents do not apply.)
48000000	Request element for channel program has been made available after it has been purged. (CSW contents do not apply.)
4F000000	Error recovery routines have been entered because of direct access error but are unable to read home address of record 0. (CSW contents do not apply.)

Appendix G: UCB Sense Bytes

			B	YTE O								BYT	re 1			
BIT	0	1	2	3	4	5,	6	7	0	1	2	3	4	5	6	7
2400	CMD REJ	INT REQ	BUS OUT	EQ CHK	DATA CHK	OVER- RUN	WORD CNT ZERO	DATA CNVTT CHK	NOISE	00-NON-X 01-NOT RI 10-RDY & 11-RDY &	ADY NO RWD	7 TRK	AT LOAD POINT	WRT STATUS	FILE PROT- ECT	NOT CAP- ABLE
2311, 2841	CMD RE J	INT REQ	BUS OUT	EQ CHK	DATA CHK	OVER- RUN	TRK CCND CHK	SEEK CHK	DATA CHK FLD	TRK OVER- RUN	END OF CYL	IN- VALID SEQ	NO REC FOUND	FILE PROT	MISSING ADR MRKR	OVER FLOW
2301, 2302 2303, 2314 2319, 2820	CMD REJ	INT REQ	BUS OUT	EQ CHK	ДАТА СНК	OVER- RUN		INVAL ADDR	DATA CHK IN COUNT	TRK OVER- RUN	END OF CYL	INVAL SEQ	NO REC FOUND	FILE PROT	SERVICE OVER- RUN	OVER FLOW
2250	CMD REJ	SHOULD NOT OCCUR	BUS OUT	SHOULD NOT OCCUR	DATA CHK	SHOULD NOT OCCUR	BUFFER RUN- NING	SHOULD NOT OCCUR	LIGHT PEN DETECT	END ORDER SEQ	CHAR MODE	j ¢	annet det			
2280	CMD REJ	INT REQ	BUS OUT	EQ CHK	DATA СНК		SHOULD NOT OCCUR		READ COUNT CHK	FILM	RECRDR FORCED GAP	NOT	SHOULD NOT OCCUR		2840 INPUT CHK	GRAPH
2282	CMD RE J	INT REQ	BUS OUT	EQ CHK	DATA CHK		SHOULD NOT OCCUR	ILLGL SEGN	READ COUNT CHK	FILM LOW	RECRDR FORCED GAP	FILM	SHOULD NOT OCCUR		2840	GRAPH
1052, 2150	CMD REJ	INT REQ	BUS OUT	EQ CHK		<u>loccor</u> II	<u> </u>	All and a second second				5. S. F . N 1997 - S.			28,18 2	
1285	CMD RE J	INT REQ	BUS OUT	EQ CHK	DATA CHK	OVER- RUN	NON	KYBD CORR								n Shi A H H Miret Miret
1287	CMD REJ	INT REQ	BUS OUT	EQ CHK	DATA CHK	OVER- RUN	NON RCVY	KYBD CORR	TAPE MODE	LATE STKR SELECT	NO DOC FOUND	SHOULD NOT OCCUR	INVAL OP	NOT	SHOULD NOT OCCUR	NOT
1288	CMD REJ	INT REQ	BUS OUT	EQ CHK	DATA CHK	OVER- RUN	NON RCVY	SHOULD NOT OCCUR	SHOULD NOT OCCUR	END OF PAGE	NO DOC	SHOULD NOT OCCUR	INVAL OP	SHOULD NOT	SHOULD NOT OCCUR	
2495	CMD REJ	INT REQ	BUS OUT	EQ CHK	DATA CHK	SHOULD NOT OCCUR	POSN CHK	SHOULD NOT OCCUR	e in jegejej		1				.	
2540, 2021	CMD REJ	INT REQ	BUS OUT	EQ CHK	DАТА СНК		UN- USUAL CMD				· · · · · ·	a an Ar				
3505, 3525	CMD REJ	INT REQ	BUS OUT	EQ CHK	DATA CHK	NOT USED	ABNOR- MAL FORMAT RESET	PERMAN- ENT ERROR (BYPASS KEY	PERM- ANENT ERROR	AUTO- MATIC RETRY	MOTION MALFUNC- TION	RETRY AFTER INT REQ COMPLETE		iner Le		
3211	CMD RE J	INT REQ	BUS OUT	EQ CHK	DATA CHK	BUFFER PARITY CHK	LOAD CHK	сн 9	CMD RETRY	PRINT CHK	PRINT QUAL- ITY	LINE POS	FORMS CHK	CMD SUP	MECHAN- ICAL MOTION	
1403, 1443	CMD REJ	INT REQ	BUS OUT	EQ CHK	TYPE BAR			сн 9	, geologie , geologie				•	~ 1	•	
1442,2596 2501,2520	CMD RE J	INT REQ	BUS OUT	EQ CHK	DATA CHK	OVER- RUN	en la composita de la composita La composita de la composita de	9.999 1997	t në s të	el Ale	1945					e yi - ".
2671, 2822	CMD REJ	INT REQ	BUS OUT	EQ CHK	DATA CHK						ngg Su	5100 - 14 - 14				ing ang Sh
2260	CMD REJ	INT REQ	BUS OUT	EQ CHK	SHOULD NOT OCCUR	SHOULD NOT OCCUR	SHOULD NOT OCCUR	SHOULD NOT OCCUR				n an	ta tula a tu		ند * ۱۹۹۹ - ۱۹۹۹ - ۱۹۹۹ -	
2701, 2702	CMD REJ	INT REQ	BUS OUT	EQ CHK	DATA CHK	OVER- RUN	LOST DATA	TIME		titelet Setter						
1419/1275 PCU	CMD RE J	INT REQ	BUS OUT	NOT USED	DATA CHK	OVER- RUN	AUTO SELECT	NOT USED	NOT USED		DOC UNDER READ HEAD	AMT FIELD VALID	PROCESS CNTRL FIELD VALID	ACCT # FIELD VALID	TRANSIT FIELD VALID	SERIAL FIELD VALID
1419/1275 SCU	CMD REJ	INT REQ	BUS OUT CHK	NOT USED	NOT USED	LATE STKR SELECT	AUTO SELECT	OP ATT	- 1. J. (18) 20	1. 1923 (1937)		Allor:				1111212
	CMD RE J	INT REQ	BUS OUT	EQ CHK	DATA CHK	OVER- RUN		12.160	PERM ERR	INVLD TRK FORMAT	END OF CYL	STATE VAR PRES	NO REC FOUND	FILE	WRITE	OPER- ATION
3330, 3333		+		EQ	DATA	OVER-	WORD	DATA CNVTT	NOISE	TU STATUS	TU STATUS	7 TRK	AT LOAD	WRT	FILE PROT-	NOT
	CMD RE J	INT REQ	BUS	СНК	СНК	RUN	ZERO	CHK !		IA	18	110	POINT	1314103	ECT	ABLE
3333					CHK DATA CHK	OVER- RUN	ZERO	СНК	PERM ERR	A INVLD TRK FORMAT	B END OF CYL	TU	POINT NO REC FOUND	FILE	ECT	ABLE OPER- ATION

2211, 2211, 2201, 2002 SERAL- LINE TAG CHK ALU UNSE CHK UNSE CHK VRC CHK VRC CHK <t< th=""><th></th><th></th><th></th><th></th><th></th><th></th><th></th><th></th><th></th><th></th><th></th><th></th><th></th><th></th><th></th><th></th><th></th><th></th></t<>																		
DEVICE 0 1 2 3 4 5 6 7 2400 BITS 0-7 INDICATE A TRACK IS IN ERROR MOLENTER OF MOLENTER SEE V CRC SEE V	_				B	YTE 2	.	-	.			-		BYT	E 3	.	<u></u>	
2200 BITS 0-7 INDICATE A TRACK IS IN ERROR 6.6 7 INDICATE MULTI-ERROR V/V URC SKEW C.C. SKEW C.C. C.C. KVD C.C.	ſ		0	1	2	3	4	5	6	7	0	1	2	3	4	5	6	7
2211, 2211, 2301, 200 UN- CHK SFEAL- LINE ALU UNSEL CHK ALU UNSEL CHK EAU EAU <td>1</td> <td></td> <td></td> <td>L</td> <td>L</td> <td>.</td> <td>1</td> <td>1</td> <td>6 & 7 IN NO ERR</td> <td>OROR</td> <td></td> <td>LRCR</td> <td>SKEW</td> <td>CRC</td> <td>REQ</td> <td></td> <td>BKWD STATUS</td> <td>COM- PARE</td>	1			L	L	.	1	1	6 & 7 IN NO ERR	OROR		LRCR	SKEW	CRC	REQ		BKWD STATUS	COM- PARE
2010, 2021 2010, 2024 2010, 2024 2010, 2024 UN-F FAIL SKEW CHK CTR FAIL COMP CHK CTR CHK COMP CHK Comp CHK <td>ŀ</td> <td></td> <td></td> <td></td> <td>IZER</td> <td>LINE</td> <td></td> <td></td> <td></td> <td></td> <td>READY</td> <td></td> <td></td> <td></td> <td></td> <td>OF</td> <td></td> <td>SEEK INCMPL</td>	ŀ				IZER	LINE					READY					OF		SEEK INCMPL
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2230 BIT 15 BIT 14 BIT 12 BIT 12 BIT 11 BIT 10 BIT 9 2280 BIT 15 BIT 14 BIT 12 BIT 12 BIT 11 BIT 10 BIT 9 2280 BIT 15 BIT 14 BIT 13 BIT 12 BIT 11 BIT 10 BIT 9 2280 BIT 15 BIT 14 BIT 13 BIT 12 BIT 11 BIT 10 BIT 9 2282 BIT 15 BIT 14 BIT 13 BIT 12 BIT 11 BIT 10 BIT 9 2282 BIT 15 BIT 44 BIT 3 BIT 20 BIT 3 BIT 2 BIT 4 BIT 3 BIT 2 BIT 3 BIT 4 BIT 3 BIT 2 BIT 4 BIT 3 BIT 3 BIT 4 BIT 3 BIT 3 BIT 4 BIT 3 BIT 4 BIT 3 BIT 4	F				1	BUFFER A	DDRESS I	REGISTER				L	l	BUFFER A	DDRESS R	EGISTER		
2280 BIT 15 BIT 14 BIT 13 BIT 12 BIT 11 BIT 10 BIT 9 BIT 7 BIT 6 BIT 7 BIT 6 BIT 7		2250		BIT 15	BIT 14	BIT 13	BIT 12	BIT 11	BIT 10	BIT 9	BIT 8	BIT 7	BIT 6	BIT 5	BIT 4	BIT 3	BIT 2	BIT 1
2282 BIT 15 BIT 14 BIT 13 BIT 12 BIT 11 BIT 10 BIT 9 BIT 7 BIT 7 BIT 7 BIT 3		2280		BIT 15	BIT 14	1	1	1	BIT 10	BIT 9	BIT 8	BIT 7	1	I	1	1	BIT 2	BIT 1
BIT 13 BIT 14 BIT 13 BIT 12 BIT 12<	ſ	2282				BUFFER A	DDRESS	REGISTER					1	BUFFER A	DDRESS R	EGISTER		
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I 3333 330, 330, 2305, 3325 ARECT- LOG NECT- FULL DATA PRESENT PATA PRESENT 2305 BJF COR- RECT- FULL COR- RECT- ARLE PRESENT RESTART COMMAND 3505, 3525 USED FOR DIAGNOSTIC PURPOSES ONLY USED FOR DIAGNOSTIC PURPOSES ONLY USED FOR DIAGNOSTIC PURPOSES ONLY 3410/3411 TRACK IN ERROR VRC MTE/ LACR SKEW ENV CRC PPT ENV BKWD 3420/3803 TRACK IN ERROR TRACK IN ERROR R/W MTE/ VRC SKEW ENV CRC ENV ENV PPT ENV BKWD 2400 ECHO ERR TRACK IN ERROR SEQ UNIT SEQ LOCK SEQ COMMAND IN PROGRESS WHEN OVERFLOW INCOMPLETE ERR ENV ERR ERR ENV ERR ERR ERR <td>ł</td> <td>3211</td> <td>10</td> <td>SEQ CHK</td> <td>STOP</td> <td>FAILED</td> <td>TO</td> <td></td> <td>MO-</td> <td>OVER-</td> <td></td> <td></td> <td></td> <td>PROT</td> <td>MER</td> <td></td> <td>SYNC</td> <td>SYNC CHK</td>	ł	3211	10	SEQ CHK	STOP	FAILED	TO		MO-	OVER-				PROT	MER		SYNC	SYNC CHK
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USED FOR DIAGNOSTIC PURPOSES ONLY UNIT TRACK IN ERROR READ <td></td> <td>2305</td> <td>LOG</td> <td>RECT-</td> <td></td> <td>a serena Ng sa</td> <td>243 (m 243 (m)</td> <td>化滑机</td> <td>\$</td> <td>n seren en e</td> <td></td> <td>tas sur tue tourise</td> <td>arts and all rates in</td> <td></td> <td></td> <td></td> <td></td> <td></td>		2305	LOG	RECT-		a serena Ng sa	243 (m 243 (m)	化滑机	\$	n seren en e		tas sur tue tourise	arts and all rates in					
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2400 ECHO ERR TAPE UNIT CLOCK ERR CLOCK ERR CNTR ERR IND C IND B IND A IND A IND C IND B IND C IND B IND C IND	-		·····			······								BYT	E 5			
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3333 PHYSICAL DRIVE IDENTIFICATION CYLINDER ADDRESS 2305 CYLINDER ADDRESS 3410/3411 TAPE POS CHK TAPE UNIT CHK DIAG TRK UNIT CHK TAPE TRK DIAG TRK TAPE UNIT CHK PE SUB- SUB- SYSTEM PE TAPE ID CHK PARITY CHK TACHO- FALSE END MARK BURST PARE CHK 3420/3803 AU HARD- WARF ERROR REJ TAPE UNIT TAPE TOR PRGM TAPE CHK UNIT CHK SPARE CHK NEW SUB- SYSTEM NEW SUB- SYSTEM NEW TAPE CHK NEW BURST PARE CHK PE BURST PARE CHK FALSE END MARK BURST PARE CHK FALSE END MARK BURST PARE CHK ENCOM- PE SYSTEM 3420/3803 TAPE WR PE VRC TAPE TOR ERROR LWR UNIT ERROR TAPE CHK NEW UNIT CHK NEW SUB- SYSTEM NEW SUB- SYSTEM START READ READ READ READ READ RECORD PARIAL POST RECORD PARIAL POST PARE OR TM BYTE 6 BYTE 7 FORMAT OF REMAINING SENSE BYTES (8-23)		2303,2314	IND	IND	IND	IND	IND	IND	IND	IND							· .	
3410/3411 TAPE UNIT CHK TAPE UNIT REJ TAPE UNIT CHK TAPE TRK DIAG TRF TAPE UNIT CHK NEW SUB- SUB- SUB- SUB- SUB- SUB- SUB- SUB-	1				PHYSICA	L DRIVE II	DENTIFIC	ATION				da et de en en dijiliji fan de	CY	LINDER	ADD RESS			
3410/3411 UNIT CHK UNIT REJ EOT TRK UNIT CHK SPARE NEW NEW TAPE ID COM- SUB- SYSTEM ID COM- SYSTEM ID COM- SYSTEM ID COM- SYSTEM ID COM- SYSTEM ID		2305			· · · .	Que e	- Arange	i n ken	₋				CY	LINDER	ADDRESS	• • • • • • • • • • • • • • • • • • • •		
3420/3803 ALU MARD- WARE EAROR REJ TAPE UNIT TAPE TGR WRITE PGB PGB ERROR TAPE UNIT NEW UNIT NEW SUB- SYSTEM NEW SUB- SYSTEM WAT TAPE CHK PE TAPE ID SYSTEM START READ CHK PARTIAL READ BURST PARTIAL POST- CHK NEW SUB- SYSTEM SYSTEM NEW SUB- SYSTEM NEW SUB- S		3410/3411	UNIT	UNIT	EOT	n an		TRK	UNIT	SPARE	SUB-	SUB-	TAPE MARK	ID	COM-	METER	END	RPQ
BYTE 6 BYTE 7 3330,* RE- 3333, VERSE HIGH HIGH HEAD ADDR FORMAT OF REMAINING ENCODED ERROF SENSE BYTES (8-23) MESSAGE		3420/3803	ALU HARD- WARE	TAPE	INDI-	TGR	DETECT		UNIT	22 - 192 22 - 4 - 193 24 - 4 - 193	NEW SUB-	NEW SUB-	WRT TAPE MARK	PE ID	START READ	PARTIAL	EXCESSIVE POST- AMPLE	RPQ
I 3333 VERSE HIGH HIGH HEAD ADDR SENSE BYTES (8-23) MESSAGE	-					1	.		4	ليستحضينا	Distant	10101011	1.0			1	1	1
	1						HE	AD ADD	2					ING				
2305* CURRENT HEAD ADDR ENCODED ERROR MESSAGE	ſ	2305*			•	CURREN	T HEAD A	DDR					EN		ERROR ME	ESSAGE		
7 SHORT DUAL ALTER- 3410/3411 TRK GAP DEN- NATE TU MODE SITY TU MODEL		3410/3411	TRK	GAP	DEN-	ALTER- NATE DENSITY		TU MO	ODEL		FAILURE	TOM LEFT	TOM RIGHT	RESET	ERASE			
		3420/3803	7 TRK	WRT	DUAL DEN-	DEN-	TAP	E UNIT M	ODEL DE	FINED	LAMP	TAPE	TAPE		DATA		BEARING PRES-	LOAD FAILURE



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Appendix H: Service Aids

In addition to the debugging facilities discussed in this manual, IBM provides the following service aid programs to aid you in debugging. A complete description of each of these service aids and instructions for their use are found in the publication <u>IBM System/360 Operating System</u> <u>Service Aids</u>, GC28-6719.

Functional Description

IMDSADMP A stand-alone program, assembled with user-selected options, that dumps the contents of main storage onto a tape or a printer. The program has two versions:

Program Name

- A high speed version that dumps the contents of main . storage to a tape.
- A low speed version that formats and dumps the contents of main storage either to a tape or directly to a printer.
- IMDPRDMP A problem program that allows the user to format and print IMSADMP output data sets, the SYS1.DUMP data set, the TSO dump data set and its associated swap data sets, and Generalized Trace Facility output data sets. IMDPRDMP can also be used to transfer a system dump from a SYS1.DUMP data set on a direct access device to another data set for later formatting and printing.
- IMCJQDMP A stand-alone program that reads, formats, and prints either the entire operating system data set SYS1.SYSJOBQE, or selects and prints information related to a specific job in that data set. Because it operates independently of the operating system, IMCJQDMP can print the contents of the job queue as it appeared at the time of abnormal termination.
- IMCOSJQD A problem program that reads, formats, and prints the contents of the system job queue data set (SYS1.SYSJOBQE). Either the entire job queue or information related to a specific job may be printed.

Because the program can be run under OS, it is not necessary to re-IPL the operating system as with IMCJQDMP.

- IMBLIST A problem program that produces formatted listings of object modules, load modules, module cross references, CSECT identification records (IDRs), and PTFs.
- IMBMDMAP A problem program that produces a map of the system nucleus, any load module, the resident reenterable load module area of an MFT system, or the link pack area of an MVT system. The listing produced by this program shows the locations of CSECTS, external references, and entry points within a load module.
- IMASPZAP A problem program that can inspect and modify either data records or load modules located on a direct access storage device.

- IMAPTFLE A problem program that generates job control language (JCL) statements necessary to add a PTF to the Operating System in a later step, or applies PTFs to the Operating System by dynamically invoking the linkage editor.
- IFCDIP00 A problem program that initializes the SYS1.LOGREC data set.
- IFCEREP0 A problem program that edits, writes, and accumulates environment records on the SYS1.LOGREC data set.

In addition to the debugging facilities described in this publication, the telecommunications access method provides the following aids to debugging:

- I/O error recording procedures.
- I/O interrupt trace table (line trace).
- A dispatcher subtask trace table (STCB trace).
- Sequential listings of buffers and message queue data sets.

Optional formatted listings of the line and STCB traces are available with TCAM. These debugging aids are described in the publications <u>IBM</u> <u>System/360 Operating System: TCAM Programmer's Guide and Reference</u> <u>Manual, GC30-2024, and IBM System/360 Operating System: TCAM</u> <u>Serviceability Aids Program Logic Manual, GY30-2027. A discussion of</u> the TCAM formatted ABEND dump is given in the publication <u>IBM System/360</u> <u>Operating System: TCAM Program Logic Manual</u>, GY30-2029.

Appendix K: Control Block Pointers

This appendix summarizes the contents of the control blocks that are most useful in debugging. Control blocks are presented in alphabetical order, with displacements in decimal, followed by the hexadecimal counterpart in parentheses. Figure 56 illustrates control block relationships in the System/360 Operating System. Figure 57 shows relationships between storage control elements in a system with MVT.

CVT - Con	munications Vector Table	<u> RB - Requ</u>	est Block (MVT)
+0	Address of TCB control words	+4	Last half of user's PSW
+53(35)	Address of entry point of ABTERM	+13(D)	CDE address
+193(C1)		+16(10)	
	only with Model 65	+29(1D)	Address of previous RB
	Multiprocessing systems and TSO)		
		MT (M) - Mo	ok Innut (Out out Mable
		$\frac{1101 - 1a}{+0}$	sk Input/Output Table Job name
DCP - Dat	a Control Block	+8	
$\frac{DCB}{+40(28)}$	<u>a Control Block</u> ddname (before open); offset to	+24(18)	Step name DD entries begin (one variable-
140(20)	ddname in TIOT (after open)	•24(10)	length entry for each DD
+45(2D)	DEB address		statement)
+69(45)	IOB address	+0	Length of DD entry
	20D address	+4	ddname
		+16(10)	Device entries begin (one 4-byte
			entry for each device)
		+20(14)	Next device entry (if there is
DEB - Dat	a Extent Block		one)
+1	TCB address		
+5	Address of next DEB		
+25(19)		-	(Next DD entry begins at 24(18)
+33(21)			plus length of first DD entry)
+38(26)			1
+42(2A)	Address of end of extent		
		,	
			k Control Block (PCP and MFT)
		+1	Address of most recent RB
		+1 +9	Address of most recent RB Address of most recent DEB
	ent Control Block	+1 +9 +13(D)	Address of most recent RB Address of most recent DEB TIOT address
<u>ECB - Eve</u> +1	ent Control Block RB address or completion code	+1 +9 +13(D) +16(10)	Address of most recent RB Address of most recent DEB TIOT address Completion code
		+1 +9 +13(D) +16(10) +25(19)	Address of most recent RB Address of most recent DEB TIOT address Completion code MSS boundary box address
		+1 +9 +13(D) +16(10)	Address of most recent RB Address of most recent DEB TIOT address Completion code MSS boundary box address Address of most recent RB on load
+1	RB address or completion code	+1 +9 +13(D) +16(10) +25(19) +37(25)	Address of most recent RB Address of most recent DEB TIOT address Completion code MSS boundary box address Address of most recent RB on load list
+1 10B - Ing	RB address or completion code	+1 +9 +13(D) +16(10) +25(19) +37(25) +113(71)	Address of most recent RB Address of most recent DEB TIOT address Completion code MSS boundary box address Address of most recent RB on load list Address of first save area
+1	RB address or completion code out/Output Block Address of next IOB (BSAM, QSAM,	+1 +9 +13(D) +16(10) +25(19) +37(25) +113(71) +161(A1)	Address of most recent RB Address of most recent DEB TIOT address Completion code MSS boundary box address Address of most recent RB on load list Address of first save area Address of STAE control block
+1 <u>IOB - Inp</u> -7	RB address or completion code <u>out/Output Block</u> Address of next IOB (BSAM, QSAM, and BPAM)	+1 +9 +13(D) +16(10) +25(19) +37(25) +113(71)	Address of most recent RB Address of most recent DEB TIOT address Completion code MSS boundary box address Address of most recent RB on load list Address of first save area Address of STAE control block Address of the job step control
+1 10B - Ing	RB address or completion code out/Output Block Address of next IOB (BSAM, QSAM,	+1 +9 +13(D) +16(10) +25(19) +37(25) +113(71) +161(A1)	Address of most recent RB Address of most recent DEB TIOT address Completion code MSS boundary box address Address of most recent RB on load list Address of first save area Address of STAE control block
+1 <u>IOB - Inp</u> -7 +2	RB address or completion code <u>out/Output Block</u> Address of next IOB (BSAM, QSAM, and BPAM) Sense bytes ECB address	+1 +9 +13(D) +16(10) +25(19) +37(25) +113(71) +161(A1)	Address of most recent RB Address of most recent DEB TIOT address Completion code MSS boundary box address Address of most recent RB on load list Address of first save area Address of STAE control block Address of the job step control
+1 <u>IOB - Inp</u> -7 +2 +5 +9	RB address or completion code <u>out/Output Block</u> Address of next IOB (BSAM, QSAM, and BPAM) Sense bytes ECB address CSW	+1 +9 +13(D) +16(10) +25(19) +37(25) +113(71) +161(A1)	Address of most recent RB Address of most recent DEB TIOT address Completion code MSS boundary box address Address of most recent RB on load list Address of first save area Address of STAE control block Address of the job step control
+1 <u>IOB - Ing</u> -7 +2 +5	RB address or completion code <u>out/Output Block</u> Address of next IOB (BSAM, QSAM, and BPAM) Sense bytes ECB address CSW CCW list address	+1 +9 +13(D) +16(10) +25(19) +37(25) +113(71) +161(A1) +181(B5)	Address of most recent RB Address of most recent DEB TIOT address Completion code MSS boundary box address Address of most recent RB on load list Address of first save area Address of STAE control block Address of the job step control block
+1 <u>IOB - Inp</u> -7 +2 +5 +9 +17(11)	RB address or completion code <u>out/Output Block</u> Address of next IOB (BSAM, QSAM, and BPAM) Sense bytes ECB address CSW CCW list address	+1 +9 +13(D) +16(10) +25(19) +37(25) +113(71) +161(A1) +181(B5)	Address of most recent RB Address of most recent DEB TIOT address Completion code MSS boundary box address Address of most recent RB on load list Address of first save area Address of STAE control block Address of the job step control block
+1 <u>IOB - Inp</u> -7 +2 +5 +9 +17(11)	RB address or completion code <u>out/Output Block</u> Address of next IOB (BSAM, QSAM, and BPAM) Sense bytes ECB address CSW CCW list address	+1 +9 +13(D) +16(10) +25(19) +37(25) +113(71) +161(A1) +181(B5)	Address of most recent RB Address of most recent DEB TIOT address Completion code MSS boundary box address Address of most recent RB on load list Address of first save area Address of STAE control block Address of the job step control block
+1 <u>IOB - Inp</u> -7 +2 +5 +9 +17(11)	RB address or completion code <u>out/Output Block</u> Address of next IOB (BSAM, QSAM, and BPAM) Sense bytes ECB address CSW CCW list address	+1 +9 +13(D) +16(10) +25(19) +37(25) +113(71) +161(A1) +181(B5) <u>TCB - Tas</u> (MFT) wit	Address of most recent RB Address of most recent DEB TIOT address Completion code MSS boundary box address Address of most recent RB on load list Address of first save area Address of STAE control block Address of the job step control block <u>Address of TCB for job step task</u> Address of TCB for next subtask
+1 <u>IOB - Inr</u> -7 +2 +5 +9 +17(11) +21(15)	RB address or completion code <u>out/Output Block</u> Address of next IOB (BSAM, QSAM, and BPAM) Sense bytes ECB address CSW CCW list address	+1 +9 +13(D) +16(10) +25(19) +37(25) +113(71) +161(A1) +181(B5) <u>TCB - Tas</u> (MFT) wit +45(2D)	Address of most recent RB Address of most recent DEB TIOT address Completion code MSS boundary box address Address of most recent RB on load list Address of first save area Address of STAE control block Address of the job step control block <u>Address of TCB for job step task</u> Address of TCB for next subtask attached by same parent task
+1 <u>IOB - Inr</u> -7 +2 +5 +9 +17(11) +21(15)	RB address or completion code <u>out/Output Block</u> Address of next IOB (BSAM, QSAM, and BPAM) Sense bytes ECB address CSW CCW list address DCB address	+1 +9 +13(D) +16(10) +25(19) +37(25) +113(71) +161(A1) +181(B5) <u>TCB - Tas</u> (MFT) wit +45(2D)	Address of most recent RB Address of most recent DEB TIOT address Completion code MSS boundary box address Address of most recent RB on load list Address of first save area Address of STAE control block Address of the job step control block <u>Address of TCB for job step task</u> Address of TCB for next subtask attached by same parent task
+1 <u>IOB - Inp</u> -7 +2 +5 +9 +17(11) +21(15) <u>RB - Requ</u>	RB address or completion code <u>out/Output Block</u> Address of next IOB (BSAM, QSAM, and BPAM) Sense bytes ECB address CSW CCW list address DCB address DCB address	+1 +9 +13(D) +16(10) +25(19) +37(25) +113(71) +161(A1) +161(A1) +181(B5) <u>TCB - Tas</u> (MFT) wit +45(2D) +129(81)	Address of most recent RB Address of most recent DEB TIOT address Completion code MSS boundary box address Address of most recent RB on load list Address of first save area Address of STAE control block Address of the job step control block <u>Address of TCB for job step task</u> Address of TCB for next subtask attached by same parent task Address of TCB for parent task
+1 <u>IOB - Inr</u> -7 +2 +5 +9 +17(11) +21(15) <u>RB - Requ</u> -8 -4	RB address or completion code <u>out/Output Block</u> Address of next IOB (BSAM, QSAM, and BPAM) Sense bytes ECB address CSW CCW list address DCB address <u>DCB address</u> <u>DCB address</u>	+1 +9 +13(D) +16(10) +25(19) +37(25) +113(71) +161(A1) +161(A1) +181(B5) <u>TCB - Tas</u> (MFT) wit +45(2D) +129(81) +133(85) +137(89)	Address of most recent RB Address of most recent DEB TIOT address Completion code MSS boundary box address Address of most recent RB on load list Address of first save area Address of STAE control block Address of the job step control block <u>Address of TCB for job step task</u> Address of TCB for next subtask attached by same parent task Address of TCB for most recent subtask
+1 <u>IOB - Inr</u> -7 +2 +5 +9 +17(11) +21(15) <u>RB - Requ</u> -8 -4 +0	RB address or completion code <u>out/Output Block</u> Address of next IOB (BSAM, QSAM, and BPAM) Sense bytes ECB address CSW CCW list address DCB address DCB address <u>ddress</u> <u>ddress</u> of previous RB on load list Address of next RB on load list Module name	+1 +9 +13(D) +16(10) +25(19) +37(25) +113(71) +161(A1) +161(A1) +181(B5) <u>TCB - Tas</u> (MFT) wit +45(2D) +129(81) +133(85)	Address of most recent RB Address of most recent DEB TIOT address Completion code MSS boundary box address Address of most recent RB on load list Address of first save area Address of STAE control block Address of the job step control block <u>Address of the job step control</u> block <u>Address of TCB for job step task</u> Address of TCB for next subtask attached by same parent task Address of TCB for parent task Address of TCB for most recent subtask Address of ECB to be posted at
+1 <u>IOB - Inr</u> -7 +2 +5 +9 +17(11) +21(15) <u>RB - Requ</u> -8 -4	RB address or completion code <u>out/Output Block</u> Address of next IOB (BSAM, QSAM, and BPAM) Sense bytes ECB address CSW CCW list address DCB address <u>DCB address</u> <u>DCB address</u>	+1 +9 +13(D) +16(10) +25(19) +37(25) +113(71) +161(A1) +161(A1) +181(B5) <u>TCB - Tas</u> (MFT) wit +45(2D) +129(81) +133(85) +137(89)	Address of most recent RB Address of most recent DEB TIOT address Completion code MSS boundary box address Address of most recent RB on load list Address of first save area Address of STAE control block Address of the job step control block <u>Address of the job step control</u> block <u>Address of TCB for job step task</u> Address of TCB for next subtask attached by same parent task Address of TCB for most recent subtask Address of ECB to be posted at task completion

+16(10) Resume PSW

+29(1D) Address of previous RB

+181(B5) Address of the job step control block

TCB - Task Control Block (MVT)						
+1	Address of most recent RB					
+9	Address of most recent DEB					
+13(D)	TIOT address					
+16(10)						
+25(19)	Address of most recent SPQE					
+33(21)	Bit 7 Non-dispatchability bit					
+37(25)	Address of most recent LLE					
+113(71)	Address of first save area					
+125(7D)	Address of TCB for job step task					
+129(81)	Address of TCB for next subtask					
	attached by same parent task					
+133(85)	Address of TCB for parent task					
+137(89)	Address of TCB for most recent					
	subtask					
+145(91)	Address of ECB to be posted at					
	task completion					
+153(99)	Address of dummy PQE minus 8					
	bytes					
+161(A1)	Address of STAE control block					

+181(B5) Address of the job step control block

UCB - Unit Control Block						
-4	CPU ID (used only with Model 65					
	Multiprocessing systems)					
+2	FF (UCB identification)					
+4	Device address					
+13(D)	Unit name					
+18(12)	Device class					
+19(13)	Device type					
+22(16)	Sense bytes (except devices with					
	extended sense)					
+24(18)	Number of sense bytes (devices					
	with extended sense)					
+25(19)	Address of sense bytes (devices					
	with extended sense)					
+40(28)	Number of outstanding RESERVE					

+40(28) Number of outstanding RESERVE requests (shared DASD only)

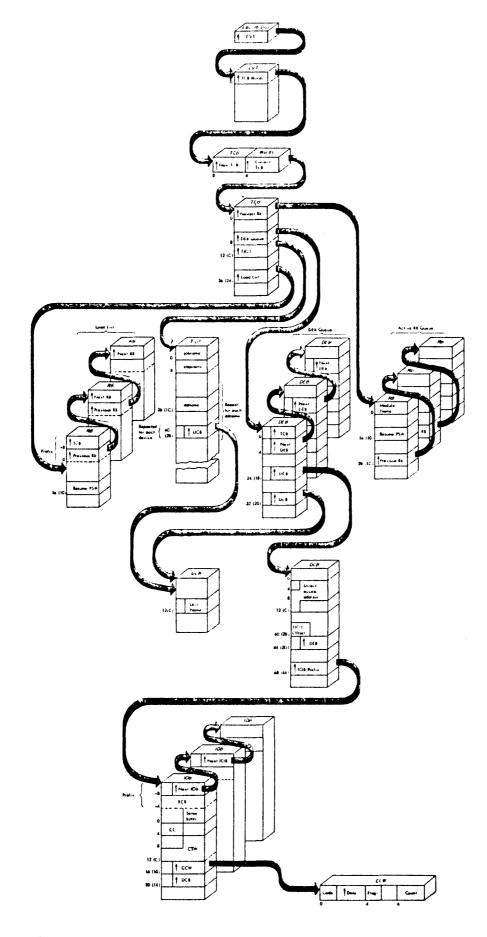


Figure 56. Control Block Flow

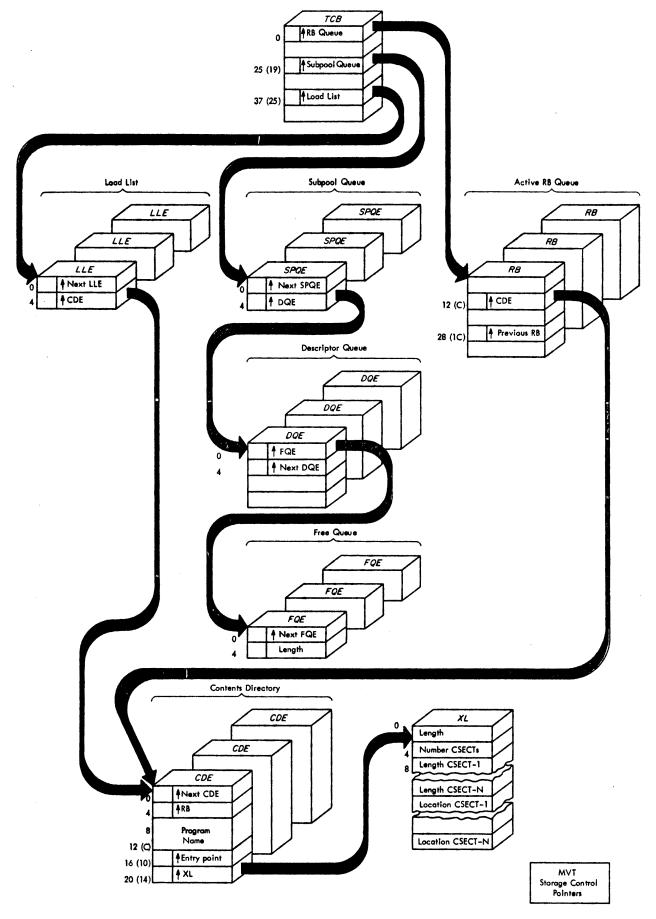


Figure 57. MVT Storage Control Flow

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Appendix L: OPEN/CLOSE/EOV Debugging Aids

There are two types of traces that may be performed during OPEN/CLOSE/EOV processing, provided that GTF is active.

- ABEND trace A trace performed before an OPEN/CLOSE/EOV problem determination module calls an ABEND routine.
- Optional work area trace A trace performed when an OPEN/CLOSE/EOV module has finished execution. This trace is made only if DCB=DIAGNS=TRACE is specified in the DD statement of the data set for which the trace is desired.

Further information on requesting these traces is contained in <u>IBM System/360</u> <u>Operating System: Data Management</u> <u>Services</u>, GC26-3746.

The format of both types of OPEN/CLOSE/EOV trace output is as follows:

r			ده دوره همه ميده هاله خلي معرد تادي هيرة كران		
USRFF	FFF	CCC	control	block	fields
L					

USRFF

is the name (excluding the IMD prefix) of the IMDPRDMP appendage which formats the control block and work area information collected by OPEN/CLOSE/EOV and included in the GTF output data set. FF is the format ID for OPEN/CLOSE/EOV. FFF

is the event ID which defines the event which caused the trace entry. Everything traced by OPEN/CLOSE/EOV has an event ID of FF.

ccc

is the control block that was traced to provide the problem program with OPEN/CLOSE/EOV data for debugging purposes.

When the OPEN/CLOSE/EOV ABEND trace occurs, only those control blocks meaningful to an ABEND condition will be traced. The selection of these control blocks is described in <u>IBM</u> <u>System/360 Input/Output Support</u> (OPEN/CLOSE/EOV) PLM, GY28-6609.

If the optional work area trace has been requested, the OPEN/CLOSE/EOV work area and the user's DCB will be traced after the execution of each OPEN/CLOSE/EOV module.

control block fields

are the contents of fields in control block ccc. For descriptions of the fields shown, refer to <u>IBM System/360</u> <u>Operating System: System Control</u> <u>Blocks, GC28-6628 or IBM System/360</u> <u>Operating System: Input/Output</u> <u>Support (OPEN/CLOSE/EOV) PLM,</u> GY28-6609.

Index

Indexes to systems reference library manuals are consolidated in the publication IBM System/360 Operating System: Systems Reference Library Master Index, GC28-6644. For additional information about any subject listed below, refer to other publications listed for the same subject in the Master Index. When more than one page reference is given, the major reference is first. ABEND dumps contents of (MVT) 46-64 contents of (MFT) 29-45 guide to using (MVT) 63-64 guide to using (MFT) 44-45 how to invoke (MVT) 46 how to invoke (MFT) 29 introduction to 11 samples of (MVT) 47-48 samples of (MFT) 30-31 Abnormal termination, cause of in an ABEND/SNAP dump (MVT) 63 in an ABEND/SNAP dump (MFT) 44 Abnormal termination dumps (see ABEND dumps) Active RB queue description of 16 instructions for using 205 in a storage dump 141 in an ABEND/SNAP dump (MVT) 52,63 in an ABEND/SNAP dump (MFT) 37,44 in an indicative dump 67 AMWP bits in an indicative dump 67 meaning of 206 APSW field, in an ABEND/SNAP dump (MVT) 52,63 ATTACH macro instruction, effects of 18 Attaching subtasks 20 Boundary problem program 40,206 Catalog dump 32 CDE as used with the load list (1.7)format of 25 in an ABEND/SNAP dump in a storage dump 141 Communications vector table (see CVT) Complete dump (MVT) description of 46 sample of 47,48 Completion codes description of common 215 explanation of 205 in an ABEND/SNAP dump (MVT) 49 in an ABEND/SNAP dump (MFT) 35 in an indicative dump 67

Console dump facility 68 COND parameter, to regulate job step execution 32 Contents directory description of 17,25 entries (see CDE) Control blocks descriptions of 26,27 pointers in 233 relationships between 27 use in debugging 205 Control block displays (IMDPRDMP output) 74-122 MFT DEB format 103 MFT job pack queue format 102 MFT load list format 102 MFT problem program boundaries 101-102 MFT RB format 98 MFT TCB format 92-98 MFT TIOT format 107 MVT DEB format 87-90 MVT job pack queue format 86-87 MVT load list format 86 MVT main storage information 85 MVT RB format 81-85 MVT TCB format 75-80 MVT TIOT format 90-92 TSO PSCB format 121 TSO RCB format 113 TSO SWAP DCB format 116 TSO TAXE format 123 TSO TJB format 116 TSO TJBX format 120-121 TSO TSB format 118 TSO TSCVT format 108 TSO UMSM format 116 Control information 13 Control program nucleus ABEND/SNAP (MVT) 60 ABEND/SNAP (MFT) 43-44 CVT description of 26 in a storage image dump 138,139 pointers in 233 ŧ Data control block (see DCB) Data event control block 27 Data extent block (see DEB) Damage assessment routine (DAR) 68 DCB description of 27 in a storage dump 144 pointers in 233 DD statements required with ABEND/SNAP dumps 29,32 sample of SYSABEND 32 DEB description of 27 in a storage dump 143 in an ABEND/SNAP dump (MVT) 55 in an ABEND/SNAP dump (MFT) 41 L pointers in 233

DEB queue in a storage dump 143 50 in an ABEND/SNAP dump (MVT) in an ANDED/SNAP dump (MFT) 35 Debugging procedure description of 205 summary 207 DELETE macro instruction 16 Dequeued elements 34 Descriptor queue element (see DQE) Destroyed queues 32 Device considerations, for ABEND/SNAP dumps 29,32 Dispatcher trace table entry (MVT) format of 148,149 in a SNAP dump 61,64 in a storage image dump 146 Dispatching priority 19,20 DOE format of 24 in a storage dump 143 in an ABEND/SNAP dump 56,64 Dump (see individual type of dump, e.g., ABEND, indicative) Dump data set MVT 46 MFT 29 Dump list heading (IMDPRDMP output) 70 Dynamic area in systems with MVT 20 in systems with MFT 19

ECB completion codes, list of 223 description of 27 in a storage dump 145 pointers in 233 posting of, using ATTACH 18 Event control block (see ECB) Extent list (see XL) External interruption 206 External trace table entry format of 148 in a SNAP dump 61,63,64 in a storage dump 148,149

FBQE format of 22-23 in a storage dump 146 in an ABEND/SNAP dump 57,64 FINCH request block 14 Finding the partition TCB 141 FRB 14 Fixed area in systems with MFT 19 in systems with MVT 20 FOE format of (MFT) 21 format of (MVT) 24 Free areas in an ABEND/SNAP dump (MFT) 44 Free block queue element (see FBQE) Free queue element (see FQE)

General debugging procedure description of 205-207 summary 207 General format (IMDPRDMP output) 128 Generalized trace faility (GTF) 150-203 comprehensive trace records 158-166,169-202 DSP 162 EXT 163 IO 159 PCI/IO 159 PGM 165 SIO 161 SSM 166 SVC 170 description of 15 hexadecimal format record 168 lost event record 167 minimal trace records 150-157 DSP 153 EXT 154 IO 151 PCI/IO 151 PGM 155 SIO 152 SSM 157 SVC 156 SVC comprehensive trace records 169-202 sub-index 169 time record 167 GETMAIN macro instruction 21-22 Gotten subtask area queue element 21-22 GQE 21-22 GTF (see Generalized Trace Facility) GTF trace table in ABEND/SNAP dumps 150 in IMDPRDMP output 150 Guide to using storage image dump 137 Hardware error 205,206 Hierarchy, main storage 20-22 IFCDIP00 230 IFCEREP0 230 IMAPTFLE 230 IMA SP ZAP 229 IMBLIST 229 IMBMDMAP 229 IMCJQDMP 229 IMCO SJQD 229 229 IMDPRDMP IMDSADMP 229 Indicative dumps contents of 65-67 description of 65 guide to using 67 introduction 11 Input/output block (see IOB) Interrupt request block 14 Interruptions 205,206 IOB description of 27 in a storage dump 144-145

pointers in 233

1

I/O trace table entry format of 148 in a storage dump (MFT) 148-149 in a storage dump (MVT) 148 in a SNAP dump (MVT) 61,63-64 in an ABEND/SNAP dump (MFT) 42,44 IRB 14 Job pack area 14,17 Job pack area queue 17 Job step 19 Job step task (MVT) 20,46 JPAQ 17 LINK macro instruction, effects of 18 Link pack area (MVT) 20-21 Link pack area maps (IMDPRDMP output) 71-73 LLE count field 17 description of 17 in an ABEND/SNAP dump (MVT) 50 Load list description of 16 instruction for using 205,207 in a storage dump 142 in an ABEND/SNAP dump (MVT) 53,63 in an ABEND/SNAP dump (MFT) 38,44 in an indicative dump 66 in systems with MVT 17 in systems with MFT 16-17 Load list element (see LLE) LOAD macro instruction, effects of 18-19 Load module, storage control for in an ABEND/SNAP dump (MVT) 53-54,64 in systems with MVT 25-26 Loaded program request block 14 LPRB 14 LRB 14 Main storage hierarchy support inclusion of 20-22 effects on MSS boundary box 21 effects on partition queue 21 Main storage layout in systems with MFT with subtasking 19-20 in systems with MFT without subtasking 19 in the systems with MVT 20 Main storage supervisor's boundary box (see MSS) Machine check interruption 205-207 MFT, systems with considerations in using an ABEND/SNAP dump of 44-45 contents of an ABEND/SNAP dump of 34-45 guide to using a storage dump of 137 how to invoke an ABEND/SNAP dump of 29 main storage layout in 20 storage control in 21-23 task control characteristics of 19-20 trace table entries in 148

Model 65 Multiprocessing system trace table formats 149 prefixed storage area, as shown in an ABEND/SNAP dump (MVT) 60 trace table entries in a SNAP dump 62 Module name prefixes, list of 219 description of (MFT) 21-22 in an ABEND/SNAP dump (MVT) 55-56 starting address (MFT) 35 Multiprogramming with a fixed number of tasks (see MFT, systems with) Multiprogramming with a variable number of tasks (see MVT, system with) MVT, systems with complete ABEND/SNAP dump of 47-48 contents of an ABEND/SNAP dump 46-64 guide to using a storage image dump of 137-146 guide to using an ABEND/SNAP dump of 63-64 how to invoke an ABEND/SNAP dump of 46 load list 16 main storage layout in 20 storage control in 22-26 task control characteristics in 20 trace table entries in 148 Nucleus contents of 19-20 in an ABEND/SNAP dump (MVT) 60 in an ABEND/SNAP dump (MFT) 44 Only loadable (OL) 14 OPEN/CLOSE/EOV debugging 241 Output comments (IMDPRDMP output) 130-136 from GTF processing 203 Overlaid problem program 34 Partition (MFT) 19-20 Partition queue element (see PQE) Partition TCBs 138-141 PIE 35,49 Pointers, control block 231 PQE format of 23 in a storage dump 143 in an ABEND/SNAP dump 56-64 PRB 14 Priority 19-20 Problem program, how to locate in a dump 205-207 Problem program storage boundaries, in an ABEND/SNAP dump (MFT) 40 Program check interruption 205,206 Program check old PSW in an ABEND/SNAP dump (MVT) 52,63 information in 205,206 Program check trace table entry format of 148 in a SNAP dump 61-62 in a storage dump 145-146 Program interruption element (see PIE) Program request block 14 Protection key 35 PSCB 146

PSW at entry to ABEND in an ABEND/SNAP dump (MVT) 49 in an ABEND/SNAP dump nMFT) 35 PSW, program check old (see program check old PSW) PSW, resume (see resume PSW) OCB 57 Queue control block trace (IMDPRDMP output) 70-71 Queue elements (MVT) 21,22-26 Queues destroyed 34 RB as affected by LINK, ATTACH, XCTL and LOAD 18 formats of 13-16 in an ABEND/SNAP dump (MVT) 52-54 in an ABEND/SNAP dump (MFT) 37,38,44 in an indicative dump 66 most recent 35,49 name field, in a dump 205 purpose of 13 pointers in 233 pointers to, in a storage dump 141-142 queue (see active RB queue) sizes of 15 types of 13-14 when created 13 which ones appear in a dump 205 RCB 145 Re-creating the task structure MFT with subtasking 139 MVT 139 Reenterable load module area (MFT) 19 Region (MVT) contents of, in an ABEND/SNAP dump 64 description of 20 storage control for 22-24 Register contents in a save area 147 in an ABEND/SNAP dump (MVT) 60 - 61in an ABEND/SNAP dump (MFT) 43 in an indicative dump 66 Request block (see RB) Resume PSW in an ABEND/SNAP dump (MVT) 53,62 in an ABEND/SNAP dump (MFT) 38,44 in an indicative dump 67 Retain dump 32 Rollout/rollin effects on partition queue 22-23 Save areas format of 147 in an ABEND/SNAP dump (MVT) 58 in an ABEND/SNAP dump (MFT) 40 Save area chain 147 Sense bytes, UCB table of 225 Sequential partitioned system (see MFT, systems with) Service aids 229 Set system mask trace table entry

I

SIO trace table entry format of (MFT) 148 format of (MVT) 148 in a SNAP dump (MVT) 62-63 in an ABEND-SNAP dump (MFT) 42,44-45 SIRB 14 SNAP dumps contents of (MVT) 46-63 contents of (MFT) 32-44 guide to using (MVT) 63-64 guide to using (MFT) 44-45 how to invoke (MVT) 46 how to invoke 5MFT) 29 introduction to 11 Snapshot dumps (see SNAP dumps) Space considerations, for ABEND/SNAP dumps 32 SPQE format of 24 in a storage dump 143 in an ABEND/SNAP dump 55,64 SQS (see system queue space) SSM (see set system mask trace table entry) Storage control in systems with MFT with subtasking 22 in systems with MFT without subtasking 21 in systems with MVT 25,26 Storage dumps discussion of 68 guide to using 137 introduction to 11 Subpool definition of 24 in a storage dump 143 in an ABEND/SNAP dump 55,64 queue elements (see SPQE) Subtask, as created by ATTACH 18 Supervisor calls, list of 209 Supervisor interrupt request block 14 Supervisor request block 14 SVC trace table entries format of (MFT) 148 format of (MVT) 148 in a SNAP dump (MVT) 61-62 in an ABEND/SNAP dump (MFT) 42-44 SVCs, list of 209 SVRB 14 SYSABEND DD statement description of 32 samples of 33 SYSOUT, as a dump data set 32 System control blocks (see control blocks) System differences in task control 18-20 System failure 68 System queue space (MVT) 20 System tasks 18-20 System wait TCB 75 SYS1.DUMP data set 68,29 SYSUDUMP DD statement 29 Task completion code (see completion codes) Task control block (see TCB) Task control differences, by system 18-20 Task dispatching priority 20

Task input/output table (see TIOT)

format of

149

in an ABEND/SNAP dump 62

Task management 13-21 Task supervision 13-21 Task structure, recreating the, using a storage dump (MVT) 139 Task switch trace table entry (MFT) format of 148 in an ABEND/SNAP dump 44 Task switching (MFT) 20 **TAXE 146** TCAM Debugging Aids 231 TCB description of 13 in an ABEND/SNAP dump (MVT) 49-51 in an ABEND/SNAP dump (MFT) 35-37 information available through 13 locating, in a storage dump 138-141 pointers in 233 pointers to, in a storage dump (MFT) 138-139 queue (MFT) 19,20 queue (MVT) 20 relationships 19-20 summary display (IMDPRDMP output) 124-127 MFT without subtask TCBs 125 MFT or MVT with subtask TCBs 125 TCBLCT 51,233 TCBNTC 51, 233 TCBOTC 51,233 TCBTCB 50,233 Telecommunications Access Method (see TCAM) Termination, abnormal (see abnormal termination) TIOT description 26 in ABEND dump (MVT) 55 in IMDPRDMP output (MVT) 90 pointers in 233

TJB 146 **TJBX 146** Traces 147-150 Trace table control block 148 delimiting entries, in an ABEND/SNAP dump (MFT) 44 description of 148 format of entries (MFT) 148 format of entries (MVT) 149 format of entries (Mod 65 multiprocessing systems) 149 in a SNAP dump (MVT) 61-62 in a storage dump 148-149 in an ABEND/SNAP dump (MFT) 42 samples of entries 148-149 usefulness in debugging 206 TSB 146 TSCVT 145 TSO Control Blocks 145-146 UCB description of 27 in a storage dump 144 in an ABEND/SNAP dump (MFT) 41 pointers in 234 UMSM 146 Unit control block (see UCB) Use count 17

Wait list 17 Wait list element 17,18 WLE 17

XCTL macro instruction, effects of 18 XL description of 25 in a ABEND/SNAP dumps 54,64

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International Business Machines Corporation Data Processing Division 1153 Westbibister Avenue, White Plains, New York 10604 (U.S.A. only)

IBM World Trade Corporation 821 United Nations Place, New York, New York \$1317 (International)