



Systems Reference Library

IBM 1130 Assembler Language

This publication contains the information necessary to write programs in the IBM 1130 Assembler language. Included are rules for statement writing, mnemonic codes and descriptions of operands, and descriptions of the instructions used to control the Assembler program.











PRE FACE

This manual describes the IBM 1130 Assembler language and defines the programming rules. It is intended as reference material for the writing of an assembler source program and the accomplishment of the steps required to produce the resulting object program. For those without programming experience or a knowledge of the principles involved, the IBM publication, <u>Introduction to IBM Data</u> <u>Processing Systems</u> (Form F22-5517), is suggested as preliminary reading.

For those without experience involving different number systems, i.e., binary and hexadecimal, the publication <u>IBM Student Text: Number</u> <u>Systems</u> (Form C20-1618) is recommended.

The reader should also be familiar with the following: <u>IBM 1130 Functional Characteristics</u> (Form A26-5918) and <u>IBM 1130 Computing System</u>, Input/Output Systems (Form A26-5890).

This publication is oriented towards the 1130 Card/Paper Tape Programming System which includes the basic assembler programs. The assembler language is also valid for the 1130 Monitor Programming system and a section on special Monitor statements is included. The Monitor System is described in the publication, <u>IBM 1130</u> Monitor System Specifications (Form C26-5940).

The operating procedures for the 1130 Card/ Paper Tape Assembler are described in the publication, <u>IBM 1130 Card/Paper Tape Programming</u> System Operator's Guide (Form C26-3629).

MACHINE REQUIREMENTS

The minimum machine configuration for assembling programs with the IBM 1130 Card/Paper Tape Assembler is as follows:

IBM 1131 Central Processing Unit, Model 1, with 4096 words of core storage

IBM 1442 Card Read Punch, or IBM 1134 Paper Tape Reader and IBM 1055 Paper Tape Punch.

Second Edition

This edition is a reprint of C26-5927-1 and incorporates changes released in Technical Newsletter N26-0554, dated March 10, 1966.

Significant changes or additions to the specifications contained in this publication will be reported in subsequent revisions or Technical Newsletters.

Requests for copies of IBM publications should be made to your IBM representative or to the IBM branch office serving your locality.

A form is provided at the back of this publication for reader's comments. If the form has been removed, comments may be addressed to IBM Corporation, Programming Publications, Department 452, San Jose, California 95114.

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INTRODUCTION

The IBM 1130 Assembler language, while similar in structure to machine language, replaces binary instruction codes with mnemonic symbols and uses labels for other fields of an instruction. Other features, such as pseudo-operations, expand the programming facilities of machine language. Thus, the programmer has available, through an assembler language, all of the flexibility and versatility of machine language, plus facilities that greatly reduce machine language programming effort.

COMPONENTS OF THE IBM 1130 ASSEMBLER

The IBM 1130 Assembler has two parts: the symbolic language used in writing a program and the assembler program that converts the symbolic language into machine language. An additional component is a library of relocatable input/output, arithmetic, and functional subroutines.

Symbolic Language

Symbolic language is the notation used by the programmer to write (code) the program. A program written in symbolic language is called a <u>source</u> program. It consists of systematically arranged mnemonic operation codes, special characters, address, and data, which symbolically describe the problem to be solved by the computer.

The use of symbolic language:

- 1. Makes a program independent of actual machine locations, thus allowing programs and routines to be relocated and combined as desired.
- 2. Allows routines within a program that can be written independently and that cause no loss of efficiency in the final program.
- 3. Permits instructions to be added to or deleted from a source program without the user having to reassign storage addresses.

Assembler Program

The assembler program (processor), supplied to the user by IBM in either paper tape or punched card form, converts (assembles) a symbolic-language source program into a machine-language program. The conversion is one for one — that is, the assembler produces one machine-language instruction for each symbolic-language instruction.

The IBM 1130 Assembler is a two-pass program. The processor is loaded into the computer and is followed by the first pass of the source program. During the first pass, the source statements are read and a symbol table is generated. During the second pass, the source program is read again and the object program and/or error indications are punched into the first 20 columns of each source card. If paper tape is used, the second pass results in the punching of a new tape that contains both source statements and corresponding object information. Both card and tape object programs must be compressed (via a Compressor Program supplied with the assembler) into a relocatable binary deck (or tape) before they can be loaded into core storage for execution. The output from the second pass is called the list deck (or tape) and can be used to obtain a program listing of source statements and corresponding object statements.

Subroutines

A library of input/output, arithmetic, and functional subroutines is available for use with the IBM 1130 Assembler.

The user can incorporate any subroutine into his program by simply writing a call statement referring to the subroutine name. The assembler generates the linkage necessary to provide a path to the subroutine and a return path to the user's program. The ability to use subroutines simplifies programming and reduces the time required to write a program.

A description of available subroutines is contained in the publication <u>IBM 1130 Subroutine Library</u> (Form C26-5929).

FEATURES OF THE ASSEMBLER

The significant features of the IBM 1130 Assembler are summarized below. More detailed explanations are given later in this bulletin.

<u>Mnemonic Operation Codes</u>: Mnemonic operation codes are used for all machine instructions instead of the more cumbersome internal binary operation codes of the machine. For example, the Subtract instruction can be represented by the mnemonic, S, instead of the machine operation code, 10010.

Symbolic References to Storage Addresses: Instructions, data areas, and other program elements can be referred to by symbolic names or actual machine addresses and designations.

Automatic Storage Assignment: The assembler assigns consecutive addresses to program elements as it encounters them. After processing each element, the assembler increments a counter by the number of words assigned to that element. This counter indicates the storage location available to the next element.

Convenient Data Representation: Constants can be specified as decimal digits, alphabetic characters, hexadecimal digits, and storage addresses. Conversion of the data into the appropriate machine format of the 1130 System is performed by the assembler. Data can be in a form suitable for use in fixed-point or floating-point arithmetic operations.

<u>Renaming Symbols</u>: A symbolic name can be equated to another symbol, so that both refer to the same

storage location. This makes it possible for the same program item to be referred to by different names in different parts of the program.

<u>Relocatable Programs</u>: The assembler can produce object programs in a relocatable format; that is, a format that enables programs to be loaded and executed at storage locations different from those assigned when the programs were assembled.

<u>Assembler Instructions</u>: A set of special instructions to the assembler is included in the language. Some of the features described in this section are implemented by these instructions.

<u>Program Listings</u>: For every assembly, the user can obtain a program listing. This listing can be produced either off-line on an IBM 407 or similar device, or on-line by the use of an IBM-supplied utility routine.

Error Checking: Source programs are examined by the assembler for errors arising from incorrect use of the language. Where an error is detected, a coded warning message appears in the program listing.

MNEMONIC CONCEPT

Symbolic programming may be defined as a method whereby names and symbols are used to write a program. The symbolic language includes a standard set of mnemonic operation codes. Mnemonic operation codes are easier to remember than machine language codes because they are usually abbreviations for actual instruction descriptions. For example:

Description	Mnemonic				
Add	А				
Execute I/O	OIX				

Each IBM 1130 machine instruction has a corresponding mnemonic operation code. In addition, there are some mnemonic codes that assign storage and others that allow the user to exercise control over the assembly process.

FORMAT OF STATEMENTS

A source program consists of a sequence of statements punched into cards or paper tape. These statements can be written on a standard coding form (X26-5994), provided by IBM. The information on each line of the form (Figure 1) is punched into one card or paper tape record. The position numbers shown on the form correspond to the card columns. If paper tape is used, the first character of each tape record corresponds to column 21. Space is provided at the top of the coding form to identify the program; however, none of this information is punched into the statement cards.

Statement Fields

An assembler statement is composed of one to seven fields: label field, operation field, format field, tag field, operand field, comments field, and identification-sequence field.

Label Field (Columns 21-25)

The label field represents the machine location of either data or instructions. The field may be left blank, may contain an asterisk in column 21, or may be filled with a symbolic address, left-justified in the field. Only data or instructions that are referred to elsewhere in the program need a label. A label that is not further referred to is not an error.

A label can consist of up to five alphameric characters, beginning at the leftmost position of the label field. A label is always a symbol and must therefore conform to the rules for symbols (see <u>Symbols</u>). The example below shows the symbol <u>ALPHA</u> used as a label.

	Label	Opera	tion	F	T			
21	25	27	30	32	33	35	40	. 45
A	L, P, H, A	ST	0,			$A_1 N_1$	EXPRE	SS10N

If the label field is left blank, it is ignored by the assembler and has no effect on the assembled program. If column 21 contains an asterisk (*), the entire statement is treated as comments and appears only in the listing. If the field contains a symbolic name (label), and the statement represents a standard machine language operation (Add, Store, etc.), the value assigned to the label is the address of the assembled instruction, which is equal to the value of the Location Assignment Counter (see Location <u>Assignment Counter</u>) at the time the statement is encountered by the assembler. Values assigned to labels of the various assembler instructions are specified in the section entitled <u>Assembler</u> Instructions.

The best labels to select are those that are mnemonically descriptive of the area or instruction to which they are assigned. Labels that have an obvious meaning provide easily remembered references for the original programmer and assist others who may assume responsibility for the program.

Operation Field (Columns 27-30)

Each machine instruction and assembler instruction has a unique mnemonic operation code associated with it. When a particular operation is to be represented, its mnemonic code must be punched, leftjustified, in columns 27–30 of the source statement card.

rogram					 	Date
rogrammed by					 	Page No of
Label		Operation		FT	Operands & Remarks	Identification
21 2	5	27 3		32 33	<u>35 40 45 50 55 60 65</u>	70 75
		L		+-	╶┫╴┚╶╽┈╿╶╽╼┨╶┧╼╢╴╽╶╢╼┨┈╽╼┨╸┨╼╢╴┸╼╂╴╢╼┨╴╢╾┨╸╢╸┥╸╢╺┪╸╢╺┪╸╢	
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- I - I - I - I						
					<u> </u>	

Figure 1. Coding Form

Format Field (Column 32)

The format field specifies the type of machine instruction being represented and, in the use of short (one-word) instructions, how the displacement field is to be handled. Any one of four entries is permitted: two for short instructions, one for a standard long (two-word) instruction, and one for an indirectly-addressed long instruction. For convenience, these formats are referred to by the character used to specify them, namely blank format, X format, L format, and I format.

Blank Format. A blank in the format field (column 32) signifies a one-word instruction. Bit 5 of the assembled instruction is set to zero. A blank

also indicates that any expression in the operand field be interpreted as the desired effective address for the statement.

During execution of certain short instructions, the effective address is the sum of the displacement and the contents of the Instruction Address Register (IAR). A blank format for such instructions causes the assembler to subtract the current value of the Location Assignment counter from the expression in the operand field. Thus, when this result is added to the IAR during execution of the instruction, the correct effective address is obtained.

The effective address of one-word Store Index (STX) instructions is <u>always</u> obtained by adding the displacement to the IAR.

The displacement of Load Index (LDX), Load Status (LDS), WAIT, all shift instructions, and all condition testing instructions is <u>never</u> added to the IAR.

The effective address of all other one-word instructions is obtained by adding the displacement to the IAR, if the instructions are not indexed; that is, if column 33 is blank or zero.

A short instruction assembled with a blank format field must not be relocated at execution time by an amount differing from the relocation increment of the data referred to by that instruction.

If a short instruction of the kind relatively addressed to the IAR is to be moved, under program control, to a location different from the location at which the instruction was assembled, the automatic subtraction of the address counter from the displacement operand value should be suppressed by use of the X format.

This requirement is not in conflict with the loadtime relocation process because the process shifts the whole program, including instructions and reference data, to a memory area different from that for which it was assembled. The relative distances between instructions and data remain the same, and the displacements remain correct.

In a relocatable assembly, the expression specifying an operand relatively addressed to the IAR must be relocatable so that the actual displacement is an absolute quantity (see <u>Expressions</u>). If this rule is not followed, a relocation error will be indicated. Also, since displacements must lie in the range -128_{10} to $+127_{10}$, the value of the displacementspecifying expression must not be more than 127_{10} greater, nor more than 128_{10} less than the address of the next location after the instruction in which it appears; otherwise, an addressing error will be indicated. An example illustrating the blank format is shown below:

Assume $A = \text{location } 1000_{10}$ B = location 1050_{10}

The value of the machine instruction counter will be 1001_{10} when instruction A is executed. Therefore, the value computed by the assembler for the displacement will be 49_{10} .

Label			Operation	F	T		•												
21	25		27 30	32	33	35				40					45				
A	,]	2	L.D			B,		1	.1		,		1			1	1		
							1		1	1	ı.	,	1_	1			1	1	
		ě.,	• , , ,				,	1		1	,		1	1	ł				1-
			• , , ,				,	,	,		,	,	,	,	1		1	1	,
1 1 1			•				,	,		1			,		1			,	
	.		•			.		,		,				,					,
		2	•				1		,		1						1	,	
B		1	D.C.			0,0	1,1	1,5	7	-,				1	1			,	
								_	1	1							1	1	

In the case of an instruction which is <u>not</u> relatively addressed to the IAR, the assembler interprets the expression in the operand field as the desired contents of the displacement field, without modification. In this case, the operand specifying the displacement must be absolute and must be in the range -128_{10} to $+127_{10}$, or relocation and addressing errors result.

<u>X Format.</u> An X in the format field indicates to the assembler that the related statement is to be assembled as a short instruction. It further indicates that any expression in the operand field is to be interpreted as the desired displacement value.

Consider the example illustrated in Figure 2; the purpose of this instruction sequence is to change the flow of a program by inserting a branch instruction in a location that previously contained an effective "no operation." If the branch instruction at BRCON were specified as MDX GO (i.e., blank format), the assembler would compute the displacement on the basis of the IAR value of 1101. (The IAR would have a value of 1101 if the BRCON instruction were executed where it was assembled.) However, the programmer, knowing the instruction will be executed at location SWTCH, computes the displacement himself and specifies the X format.

<u>L Format.</u> If column 32 contains the character L, it signifies a long (two-word) instruction. Bit 5 (F) of the assembled instruction is set to 1. The operand field expression, which may be relocatable or absolute, is used to fill the second word (bits 16-31) of the assembled instruction. A second operand may be present, separated from the first operand by a comma (,). This operand may be used in one of two ways:

- 1. To specify symbolic condition codes for use with BSC, BSI and BOSC instructions.
- 2. To specify an expression that has a value in the range of -128 to +127 that is not relocatable.

This second operand yields bits to fill bit positions 8-15 of the assembled instruction.

<u>I-Format.</u> If column 32 contains the character I, it signifies an indirectly addressed two-word instruction. Bit 5 and bit 8 are set to 1. In all other respects an indirect instruction is treated exactly as a long direct instruction. If a displacement operand is specified, its high-order bit (bit 8) will always be a one because this bit is also the indirect flag bit.

Label		Operation		F	T							Operan	ds & Remarks					
21 25		27 30	,	32	33		35	40		45		50	55	6	<u> </u>	65		70
		• · · · ·							J. L. L						.		<u> </u>	
	ана С	•••••						- 1									L	
		•							1 1_1		1				<u> </u>			I
		<u>•</u> • • • •					┟╌┛╌╌└╌┘	1 _1	1 1						1		<u> </u>	
SWITCH		S.L.T.	1	_			0, , E	FIFIE	<u>, C'1</u>	$I_1V_1E_1$	N_O_	OP_	ERATI	0, N, 1				
	81	•	2					الم ال	1						1.1			
	÷	••••	8					_1_1_		<u> </u>					1			
		•							1		1.1.1				11			
		LD					BIRIC	0, N,	CH	ANG	P.,	ROG	RAH F	LOW	AT	, <u>,</u> S,W,	TICI	1
		S.T.O.					S,W,T	C _ H _	<u> </u>						1	.		
		•				j.			<u> </u>		1		<u> </u>		1	بيب		l.
	Ċ.					્ર			1									
		·				201 1-3				<u> </u>				<u> </u>	<u> </u>	لسبب		
BIRICION		M.D.X.		X			6.0	6 W T	С.Н.	-1.1	1.1.1		<u> </u>		1.1			
Luu							L		<u>. </u>									1.1.1

Assume SWTCH = location 1000_{10}

 $GO = location 1050_{10}$ BRCON = location 1100_{10}

Figure 2. Use of X Format

Tag Field (Column 33)

Column 33 is used to specify an index register if one is required. The code in column 33 is the index register number; i.e., 1 = Index Register 1, 2 = IndexRegister 2, and 3 = Index Register 3. A zero or a blank indicates that no index register is to be used.

If no tag is specified in an LDX, MDX, or STX instruction, the IAR is used. The example below shows an add instruction in its long form addressing the core location whose address is zero plus the contents of Index Register 2.

Label		Operation	F	т		
21	25	27 30	32	33	35	40
S,U,M,		A	٤	2		
	J					

Operands and Remarks Field (Columns 35-71)

The operand field is used to specify subfields in instructions and constants. The content of the operand field for the various instruction formats are described under <u>Format Field</u>. Blanks must not appear within the operand(s) except as character values or in the EBC statements.

Some examples of one- and two-operand statements are shown in Figure 3.

Remarks Field

Remarks are for the convenience of the programmer. They permit lines or paragraphs of descriptive information about the program to be inserted in the program listing. Remarks appear only in the program listing; they have no effect on the assembled object program. Any valid characters (including blanks) can be used as remarks.

The Remarks field must appear to the right of the operand field and must be separated from it by at least one blank.

Comments Field

The entire statement field can be used for comments by placing an asterisk in column 21; the entire

SHORT	5,7,0,		A.C.C.U.1.
	-		
L.O.N.G.	M.D.X.	L	$A_1C_1C_1U_1I_1U_1O_1O_1 T_1W_1O_1 = 0 P_1E_1R_1A_1N_1D_1 S_1T_1A_1T_1E_1M_1E_1N_1T_1 + 1$
<u></u>			
LONG	5.7.0	14	$A_1C_1C_1U_1I_1 = A_1A_1A_1E_1A_1A_1A_1A_1A_1A_1A_1A_1A_1A_1A_1A_1A_$
ليبينا			

Figure 3. One- and Two-Operand Statements

statement is then treated as comments. The identification-sequence field (columns 73-80) should not be used for comments.

If it is necessary to continue comments on additional lines, each line must have an asterisk in column 21, as illustrated in Figure 4.

Identification-Sequence Field (Columns 73-80)

The identification-sequence field may be used for program identification and statement-sequence numbers. It is limited to columns 73-80. The information in this field normally is punched in every statement card. The assembler, however, does not check this field.

STATEMENT WRITING

Symbolic language statements are accepted by the assembler only if they conform to the rules of syntax presented in this section. Subsequent sections of this publication deal with the format and content of the specific types of assembler statements (machine instructions and assembler instructions). Instructions of both types are formed by using the basic elements described here. Many of the points introduced in this section are covered more extensively in subsequent sections.

Character Set

The following characters can be used in statements:

Monocase Alphabetics	A through Z, $, #, @,$
Numerics	0 through 9
Special Characters	/ * + - = & ¬ <>
	' . , :;()% - ?
	(blank)

The codes that the assembler accepts for these characters are listed in Appendix A which also contains additional codes which may be used in comments statements, as character values, and as alphameric constants. The + and & special characters may be used interchangeably as operators.

Symbols

Storage areas, instructions, and other elements may be given symbolic names for the purpose of referring to them in the program. The symbolic name is called a symbol. It can contain up to five characters. While the first character of a symbol must be alphabetic, the remainder may be alphabetic, numeric, or any combination of the two. No blanks or special characters may be used. Any violation of these rules is detected by the assembler and indicated as an error in the program listing.

The following	are valid symb	ols:
PUNCH	START	Ν
A2345	LOOP2	$\mathbf{BC}\#@$
The following	symbols are in	valid, for the
reasons noted:		

256B	First character is not
	alphabetic
RECORDAREA2	More than 5 characters
END 1	Contains a blank

If a symbol is to be used as an operand, it must be defined in the program by using it as the label of a statement. Two types of label assignments are allowed. In machine-instruction statements and certain assembler statements, the label is assigned an address equal to the current value of the location assignment counter. In the Equate Symbol statement (see Symbol Definition Statement), the label is assigned the value specified in the operand of the statement.

Symbol Table: For every program assembled, a table of the symbols in that program is created. This is the symbol table; each entry in the table records the value and relocation property of a symbol.

All symbols defined in the program are entered in the symbol table. Symbols that appear in the label field of assembler instructions which do not use labels (for example, ABS, END, ENT) are <u>not</u> placed in the symbol table.

General Restrictions on the Use of Symbols: The following restrictions are imposed on the use of symbols:

- 1. A symbol may appear only once in a program as the label of a statement. If a symbol is used as a label more than once, only the first usage is recognized. Each subsequent usage of the symbol as a label is ignored and is noted as an error in the program listing. In addition, any reference to such a symbol is noted as an error.
- 2. The number of symbols that can be defined in a program is restricted by the amount of core storage available to the assembler (see <u>1130</u> <u>Card/Paper Tape Programming System Opera-tor's Guide (C26-3629).</u>

LOCATION ASSIGNMENT COUNTER

The assembler maintains a counter to assign sequential storage addresses to program statements. This counter is called the Location Assignment Counter. It always indicates the next available address. As each machine instruction is processed, the counter

	Lobel				Оре	ation		F	T															(Oper	and	fs &	Ren	nark	:5							_					
21		2	۶ (2	<u>, </u>	3	•	3	3		35				40)				45					50				55				60					65				70
*	T_1H_1	E	6	18	<u>,</u> 7	E,1	2/1	1	: K	·	1	N.	(2,0	2,2	· 1	. 2	1		M	A	ĸ	Ē	s,	i	T,I	4,1	1,5		A		C, C	, м	,м	,E	, N,	Τ.	s,		L.,	٨. /	Ι.Ε.
*	A. N.		ls	7	' <u>, E</u>	R. 1	1	<u> </u>	1	1	s	Ĺ	R,l	5,6	2.4	41	R	E	, D	2	F	0	R	_	ξ"	A, (c ,/	Ч,	۷,	1	N	Ξ,	,0	F	,	, C ,	0,	M	4	E,1	1,7	S,
L			ľ	1		<u> </u>												ı	1	1		1			1	1	1			1			-									

Figure 4. Example of Comments Statement

is incremented by the number of words assigned to that instruction. Certain assembler instructions also cause the Location Assignment Counter to be set or incremented, whereas others do not affect it (see <u>Assembler Instructions</u>).

Location Assignment Counter Overflow: The maximum value of the Location Assignment Counter is 65535, a 16-bit value. If a program being assembled causes the counter to be incremented beyond 65535, the assembler retains only the rightmost 16 bits in the counter and continues the assembly, checking for any other source program errors. No usable object program is produced. The user can, however, still obtain a listing of the entire source program.

RELATIVE ADDRESSING

Once an instruction has been named by a symbol in the label field, it is possible for other instructions to refer to that instruction by using the same symbol. Moreover, it is possible to refer to instructions preceding or following the instruction named by indicating their positions relative to that instruction. This procedure is referred to as relative addressing. A relative address is, effectively, a type of expression (see Expressions).

For example, in the sequence

Label	Operation	Γ	F	T		
21 25	27 30		32	33	35 40	45
START	A				BIETA	
	S				STORE	<u></u>
Luci	S.T.O.		L		$A_1D_1D_1R_1$	
A, L, I, S,T	A		L		$L_1 S_T$	
	D_{1}				L0,C,2	

control can be transferred to the second instruction by either of the following instructions:

	Lobei		Operation		F	T			
21	25	12	27 30	18 C.	32	33	199	40	<u></u>
			B,S,C,		L			START+1	
			B,S,C		L			A, L, I, S, T, -13, 1	
			1.1.1.						

By using relative addressing, it is also possible to refer to a particular word within a block of reserved storage. For example, the instruction

Label 21 25	Operation 27 30	F T 32 33	35	40 45	
B.E.T.A.	BISIS,		50.		1_1_
					<u></u>

reserves a block of 50 words, in which BETA is assigned to the first word in the block. The address BETA +1 then refers to the second word, BETA +2 to the third word, and BETA +n to the (n+1)th word.

Relative addressing can also be effected by using the current value of the Location Assignment Counter in an operand. In symbolic language this value is denoted by an asterisk (*). (See <u>The Asterisk Used</u> <u>as an Element.</u>)

SELF-DEFINING VALUES

A self-defining value is an actual machine value or a bit configuration.

Self-defining values can be used to specify such program elements as data, masks, addresses, and address increments. The type of representation selected (decimal, hexadecimal, or character) depends on what is being specified.

Decimal Values

A machine decimal value is an absolute number from 0 to 65535. It is assembled as its binary equivalent. Some examples of decimal self-defining values are

500	003
17	52324
7230	1

If a number larger than 65535 is specified in address arithmetic, the value is truncated modulo 65536; that is, only the low order 16 bits of the binary value are retained.

Hexadecimal Values

A hexadecimal value is an unsigned hexadecimal number written as a sequence of digits. The digits must be preceded by a slash (/). The hexadecimal digits represent the 16 possible combinations of four bits.

Each hexadecimal digit is assembled as its four bit value. The hexadecimal digits and their bit patterns are as follows:

0 - 0000	4 - 0100	8 - 1000	C - 1100
1 - 0001	5 - 0101	9 - 1001	D - 1101
2 - 0010	6 - 0110	A - 1010	E - 1110
3 - 0011	7 - 0111	B - 1011	F - 1111

The following are examples of hexadecimal, self-defining values:

```
/FFFF (highest value)
/AB12
/F2
/379B
```

If more than four hexadecimal digits are specified in one sequence, only the four low-order digits are retained by the assembler. If less than four hexadecimal digits are specified, they are entered, right-justified.

A table for converting decimal values to hexadecimal values is provided in Appendix B.

Character Values

A character value is a single character, preceded by a period. A character value may be a blank, any combination of punches in a single card column, or a paper tape character that translates into the eightbit IBM Extended BCD Interchange Code. Appendix A is a table of these combinations, their interchange codes and, where applicable, their printer graphics. A period used as a character value is represented as two periods in sequence, (i. e., ..). Examples of character values are:

- . A . 1 . 2 . D
- (blank)

The same value can frequently be represented by any one of the three types of self-defining values. For example, the decimal value 196 can be expressed in hexadecimal as /C4 and as a character, .D. The selection of a particular type of value is left to the programmer. Decimal values can be used for actual addresses and input/output unit numbers, hexadecimal values for masks, and character values for data.

EXPRESSIONS

The term "expression" refers to symbols or selfdefining values used as operands, either singly or in arithmetic combinations. Expressions are used to specify the various fields of machine instructions. They are also used as the operands of assemblerinstruction statements.

An expression has three components: elements, terms, and operators.

Elements

The smallest component of an expression is an element. An element is either a single symbol or a single self-defining value. The following are valid elements:

TMP
/1A6
•В
A
*
4

The Asterisk Used As an Element

When used as an element the asterisk is relocatable and stands for the current value of the Location Assignment Counter for the instruction in which it appears (i.e., the rightmost word of the current instruction + 1). Thus, the asterisk as an element can have different values for different instructions.

Label		Operation		F	т	•	
21	25	27 30	28	32	33		35 40 45
		LD					A,B,C,
S,U,M,		A				1	$D_1 E_1 F_1$
		S					DATA
		B,S,C		L			S,U,M,,+,,+,,,,,,,,,,,,,,,,,,,,,,,,,,,,,

The last instruction is a conditional transfer to location SUM and can be written

21	Label 25	Operation	F 32	T 33	35 40	45
		B,S,C,	٢		*-4.+.	
Ŀ		1-1-1				

Be sure the asterisk refers to the proper word when it is used with a long instruction or in an area where long instructions are present. In the previous example, the BSC instruction will become two machine language words after assembly. Therefore, during assembly of the BSC instruction, the Location Assignment Counter contains a value one greater than if the BSC were a short instruction.

Terms

A term can consist of a single element, two elements separated by an asterisk (which denotes multiplication), or three elements separated by two asterisks, etc. A term must begin with an element and end with an element. It is not permissible to write two elements in succession. The following are valid terms:

TMP * FUNC * TAXY A * 4 X * Y * 5 6 * 4096 3

Operators

An operator is a character that denotes an arithmetic function. The recognized operators are + or & (plus or ampersand), - (minus), and * (asterisk), denoting addition, subtraction, and multiplication, respectively. An operator must be used between two terms. Two operators may not be used in succession.

There is no ambiguity between the use of the asterisk as an element and the use of the asterisk as an operator to denote multiplication because the position of the asterisk always makes clear what is meant. Thus, **10 means "the value of the Location Assignment Counter multiplied by 10."

Evaluation of Expressions

From a symbolically written operand, the evaluation procedure derives an integer value which can be used as (1) a displacement value in a short instruction, (2) an address in a long instruction, or (3) an absolute numeric quantity.

An expression is evaluated as follows:

- 1. Each element is replaced by its numeric value.
- 2. Each term is evaluated by performing the indicated multiplications from left to right, in the order in which they occur. In multiplication, the low-order 16 bits are retained.
- 3. The terms are combined from left to right, in the order in which they occur. If the result is negative, it is replaced by its 2's complement.

Grouping of terms, by parentheses or otherwise, is not permitted; however, this restriction can often be circumvented. For example, the product of 25 times the quantity B-C can be expressed as

25 * B - 25 * C

Types of Expressions

In addition to evaluating expressions, the assembler program must decide whether the expression is <u>absolute or relocatable</u>. Without this information the assembler would be unable to assign the proper relocation indicator bits for use by the loading routine.

Rules for Determining the Type of Expression

The rules by which the expression type is determined are:

- A symbol that is defined by means of the Location Assignment Counter is a relocatable element.
- Decimal and hexadecimal integers and character values are absolute elements.
- A relocatable element alone is a relocatable expression.
- A relocatable element, plus or minus an absolute element, is a relocatable expression.

- The difference of two relocatable elements is an absolute expression.
- A symbol that has been equated to an expression (by means of the EQU assembler instruction) assumes the same relocation property as that expression.

These rules are clarified by the following example:

Assume that a programmer wishes to incorporate a table into a relocatable program, and he knows that he may later wish to add or delete items without changing program references to the table. The first step is to assign symbols to the first (lowest addressed) word in the table and to the location immediately after the last (highest addressed) word of the table. These symbols could be BGTBL and ENTBL, respectively. Regardless of the number of items in the table or of the number of later additions or deletions, the number of words in the table is always equivalent to the value of the expression ENTBL-BGTBL. This illustrates the rule that the difference of two relocatable elements is an absolute expression.

Expanding this example, assume the programmer wishes to use a second table the same length as the first. The first (lowest addressed) word of the second table can be indicated by the symbol STBL. Then, the location following the last (highest addressed) word of the second table can be indicated by the expression

STBL + ENTBL - BGTBL

This address is subject to relocation; hence, the expression is relocatable, following the rule that a relocatable element plus or minus an absolute element is a relocatable expression.

Procedure for Determining the Type of Expression

The following paragraphs describe the procedure for determining expression type (absolute or relocatable):

- 1. Discard any term that contains only absolute elements.
- 2. Examine each term of the expression. If any term contains more than one relocatable element, the expression will yield a relocation error.

- 3. Replace each relocatable element by the symbol r, and replace each absolute element by its value. This yields a new expression which in-volves only numbers and the symbol r.
- 4. Rewrite the expression in simplest form by evaluating it according to the address arithmetic rules given above in the section, <u>Evaluation of</u> Expressions.

If the result is an integer, the operand is absolute. If the result is r, the expression is relocatable. If the result contains r to any power other than one or contains r with a coefficient other than one, the operand does not have a well-defined relocation property and will yield a relocation error. The following examples illustrate this procedure.

NOTE: When the terms absolute symbol and relocatable symbol are used in text, they mean symbols that refer to addresses.

Example 1: Consider the expression

4+3*TRANS-2*FUNC+COUNT

where TRANS and FUNC are relocatable symbols, and COUNT is an absolute symbol. Discarding the terms involving only absolute elements leaves

3*TRANS-2*FUNC

This does not contain any illegal terms. Replacing each symbol by the letter r results in

3*r-2*r

Evaluating this produces r; therefore, the expression is relocatable.

Example 2: Consider the expression,

2*3*TRANS-FUNC

This reduces to

2*3*r-r

 \mathbf{or}

5r

This is neither r nor a number; therefore, the expression will cause a relocation error.

Example 3: Consider the expression,

A*2*R-A*A*R+5

where A is an absolute symbol, and R is a relocatable symbol. The expression is absolute if the value of A is zero or two and relocatable if the value of A is 1. If the value of A is anything else, a relocation error will result.

In the following examples, A, B, C, and D are relocatable symbols, and J, K, L, M, and N are absolute symbols.

Relocatable expressions:

А	1*A
A+J	250*A-249*B
A+B+C - D-*	100*A+50*B-75*C-74*D

Absolute expressions:

12345	0*A
A-B+C-D+5	500*A-400*B-100*C

Relocation Errors

If a source program contains an expression having in it one or more of the following, that expression is flagged as a relocation error.

- The negative (complement) of a relocatable element
- An absolute element minus a relocatable element
- The sum of two relocatable elements

In the following examples, A, B, C, and D are relocatable symbols, and J, K, L, M, and N are absolute symbols.

A+B	(+2r)	A*B	(r ²)							
-A	(-1r)	2*A	(2r)							
15-*	(-1r)	5*A-6*A	(-1r)							
A+J+M+N+B-C+D+L(+2r)										

NOTE: In an absolute assembly headed by an ABS statement (described later), all symbols and asterisk values are defined as being absolute; therefore, no relocation errors are possible.

MACHINE-INSTRUCTION STATEMENTS

All machine instructions can be represented symbolically as assembler language statements. There are two basic formats: short and long. However, within each basic format, further variations are possible.

The symbolic format of a machine instruction parallels, but does not duplicate, its actual format. A mnemonic operation code is written in the operation field, and one or more operands are written in the operand field. Comments can be appended to a machine-instruction statement as previously explained.

Any machine-instruction statement can be named by a symbol, which other assembler statements can use as an operand. The value of the symbol is the address of the leftmost word assigned to the assembled instruction.

MNEMONICS

A list of all IBM 1130 machine language instructions and their associated mnemonics is given in Table 1.

Condition-Testing Instructions (BSC, BOSC, BSI)

The machine instructions Branch or Skip on Condition (BSC), Branch Out or Skip on Condition (BOSC), and the long form of Branch and Store Instruction counter (BSI) use bits 10–15 of the displacement to test any combination of six conditions associated with the accumulator. When coding these instructions, the user does not use an expression to specify the displacement field, but, instead, writes a series of unique characters, each of which represents one bit of the condition-testing mask. These character symbols may be written in any combination; the bits they represent are combined by the assembler in a logical OR fashion. The symbols and their representations are:

Unique			Bit Position
Character	Condition	Description	Set to 1
O (Alpha)	Overflow	Skip or do not branch if Overflow indicator <u>off</u>	15
с	Carry	Skip or do not branch if Carry indicator <u>off</u>	14
E	Even	Skip or do not branch if bit 15 of acc.=0	13
+ or &	Plus	Skip or do not branch if bit 0 of the acc. =0, but not all bits of acc. =0	12
-	Minus	Skip or do not branch if bit 0 of acc. =1	11
Z	Z ero	Skip or do not branch if all bits of acc. =0	10
Examples			
DSC			
BSC	+ +-	Skip on plus condition	
BSC	+ -	Skip on non-zero (plus or n	ninus)
DSC PSC	2-	Skip on non-plus (zero or n	ninusj
DSC I		Skip in Carry Indicator <u>on</u>	
DSC L	EA11, +	(zero or minus)	
BSC L	EXIT, + -	Branch to EXIT if zero (not minus)	plus or
BSC L	EXIT	Unconditional Branch to EX	IT
BSC L1	0, Z+	Branch to the contents of X minus (not zero or plus)	R1 on
BSI L	SUBR, O	Branch and Store instruction to SUBR if Overflow is <u>on</u>	n counter

Table 1. Machine Instruction Mnemonics

Mnemonic	OP Code (Hexadecimal Representation) ¹	Instruction
1 · · · · · · · · · · · · · · · · · · ·		
Load and Store		
LD LDD LDX LDX* STO STD STD STX STS	C00 C80 -600 200 D00 D80 680 280	Load Accumulator Load Double Load Index Load Status Store Accumulator Store Double Store Index Store Status
Arithmetic	-	
A AD S SD M D AND OR EOR	800 880 900 980 A00 A80 E00 E80 F00	Add Add Double Subtract Subtract Double Multiply Divide And Or Exclusive Or
Branch	· · · · ·	
BSC BOSC ² MDX	400 480 484 700	Branch and Store Instruction Counter Branch or Skip Conditionally Branch Out or Skip Conditionally Modify Index and Skip
Shift		•
SLA* SLT* SLC* SLCA* SRA* SRT* RTE*	100 108 10C 104 180 188 188	Shift Left Accumulator Shift Left Accumulator and Q Reg. Shift Left and Count Accumulator and Q Reg. Shift Left and Count Accumulator Shift Right Accumulator Shift Right Accumulator and Q Reg. Rotate Right
Input/Output		
XIO	080	Execute I/O
Miscellaneous ³	100	
WAIT*	300	No Operation Wait

*Valid in short format only.

1. The hexadecimal representation of the machine operation code is derived from the instruction format in the manner shown below. Bits 5, 6, 7, 10, and 11 are assumed to be zeros because they do not enter into the makeup of any operation codes.

2. Same as BSC with Bit 9 set to one.

3. An operand should not be specified.

Hexadecimal Characters

· .	
	1st 2nd 3rd
	0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15
OP Code	
Format (F)	
Index Tag bits (T)	
IA bit, part of displacement, or extension of OP code	
Part of displacement, or extension of OP code	
Condition indicators, or part of displacement	

Just as machine instructions are requests to the computer to perform a sequence of operations during program execution time, assembler instructions are requests to the assembler to perform certain operations during the assembly. In contrast to machineinstruction statements, assembler-instruction statements do not always cause machine instructions to be included in the assembled program. Some, such as BSS and BES, generate no instructions but do cause storage areas to be set aside for constants and other data. Others (e.g., EQU) are effective only at assembly time; they may or may not generate something in the assembled program. If nothing is generated, the Location Assignment Counter is not affected.

The following is a list of all assembler instructions permitted by the IBM 1130 Assembler.

Program Control

ABS	- Absolute Assembly
LIBR	- Transfer Vector Subroutine
SPR	- Standard Precision
EPR	- Extended Precision
ORG	– Define Origin
END	- End of Source Program
Data Definit	tion
DC	- Define Constant
DEC	- Decimal Data
XFLC	- Extended Floating Constant
EBC	- Extended Binary Coded Information
Storage Allo	ocation
BSS	- Block Started by Symbol
BES	- Block Ended by Symbol
Symbol Defi	inition
EQU	– Equate Symbol
Program Lin	king
ENT	- Define Subroutine Entry Point
ISS	- Define Interrupt Service Entry Point
ILS	- Define Interrupt Level Subroutine
CALL	- Call Subroutine (2-word call)
LIBF	- Call Subroutine (1-word call)

PROGRAM CONTROL STATEMENTS

Program control statements are used to set the Location Assignment Counter to a specific value, to define the end of a source program, or to specify whether a particular program is to be assembled as absolute or relocatable. None of these assembler statements generate machine-language instructions or constants in the object program.

ABS – Assemble Absolute

An ABS statement is used to specify that a main program is to be assembled as an absolute program. An absolute program is one in which the core locations used at execute time are the same as those specified by the programmer in the source program. The ABS statement is punched as shown below and is then used as the first statement of a source program.

Label			Operation		F	T													
21	25	Č.	27 3	0	32	33	1000	35			40)			. 4	5			
	, []		A,B,S,	100					1							1	1	1	
	. [1	1		1	1	ł	
	2									÷									

If the first (non-comment) statement of a source program is not an ABS statement, the program will be assembled as relocatable. In an absolute assembly headed by an ABS statement, all symbols and asterisk values are defined as absolute quantities; therefore, no relocation errors are possible. The significance of relocatable and absolute assemblies is explained in the following paragraphs.

Relocatable Assembly

Some programs assembled by the IBM 1130 Assembler are absolute; that is, the locations of assembled instructions are known at assembly time and the location on the listing is the actual location where a particular word is loaded. However, subroutines used by an absolute program must be in such a form that they may be loaded at various locations; otherwise, it would be necessary for the user to reassemble the subroutines each time he assembled a main program that required them. Therefore, all subroutines and certain main programs should be assembled relocatable.

Every relocatable program or subroutine produced by the IBM 1130 Assembler is assembled as though it begins at location zero. Since a job to be executed may contain several subroutines, it is obvious that they cannot all be loaded into locations starting with location zero. In fact, no relocatable program is ever loaded at location zero; instead, each program is relocated. The relocatable main program is loaded into the first available location. Subroutines are then loaded into successively higher locations of core storage, each beginning with the next even location after the last core storage location used by the preceding subroutine, as indicated on the end-of-program card. When a particular program has been loaded, the address of the first word is called the load address for that program.

Thus, the address in core storage actually occupied by an instruction of the program is the address assigned to that instruction during assembly, plus the load address of that program. To keep the program self-consistent, the load address must be added to the address of many (but not all) 2-word instructions, and those constants whose values are relocatable.

This process of conditionally adding the load address is performed by the loading program before execution and is called relocation. In relocating instructions, the loading program is guided by relocation indicator bits that are inserted into the binary object program cards when the program is compressed following assembly.

Absolute Assembly

The programmer uses the ORG assembler statement in his source program to specify the locations into which the binary object program resulting from an absolute assembly is loaded. The loader makes no adjustment to any field during loading. Subroutines are loaded into successively higher even-core locations following the end of the main program.

Remember that only main programs may be assembled absolute; subroutines must be assembled relocatable.

LIBR - Transfer Vector Subroutine

An LIBR statement is used as the first statement of a subroutine to specify that the subroutine is to be called by LIBF statements only (see <u>Program-</u> <u>Linking Statements</u>). The absence of an LIBR statement specifies that the subroutine is to be called by CALL statements only. LIBR statements are for subroutines only, as ABS statements are for main programs only. An LIBR statement needs no operands.

<u>SPR – Standard Precision and EPR – Extended</u> <u>Precision</u>

The SPR or EPR statement specifies that the program (main or subroutine) in which it appears uses standard precision or extended precision, respectively, for arithmetic operations. If these statements are included in the user's programs, the loader ensures that main programs and subroutines always match with regard to precision. Their use is optional, however. If used, the SPR or EPR statement must follow the ABS or LIBR statement. If no ABS or LIBR statement is used, the SPR or EPR statement is the first statement in the program.

ORG - Define Origin

This assembler instruction is used to set the Location Assignment Counter (i.e., the next location to be assigned) to any desired value. In this way the programmer is able to control the assignment of storage to instructions, constants, and data. If a Define Origin statement is not the first entry in an absolute source program, the processor begins the assignment of storage at a location compatible with the size of the applicable loader. A typical Define Origin statement is shown below.

21	Label 2	Operation 27 30	F 32	T 33	35	40	45
	1.1.1	$O_{R}G_{I}$			3,0,0,0	4	
	- 1 - 1 - 1				Lui		

The label, if used, is assigned a value equal to the value of the Location Assignment Counter at the time the statement is encountered in the source program. (This assignment is made <u>before</u> the counter is modified.) If any symbols are used in the expression, they must have been previously defined. In a relocatable assembly, an absolute expression in the oper-and field is considered a relocation error and the statement is ignored.

Some examples of Define Origin statements are given below:

Label		Operation	F	т				Ope
21 25		27 30	32	33	35	40	45	50
	Å.	O.R.G.			X.4.Z			
START		O,R,G,			XYZ	+ 50		
SITIART		O,R,G			¥1+.5	0, 140,C	$C_1T_1R_1+$	5,0, 1

If the label XYZ has been previously defined as 1000_{10} the first entry directs the assembler to begin the assignment of succeeding entries at location 1000. The second entry directs the assembler to begin the assignment of succeeding entries 50 core locations beyond the location that has been assigned to the symbol XYZ. The third entry directs the assembler to begin the assignment of succeeding entries at the

address specified by the current address of the Location Assignment Counter plus 50.

END - End of Source Program

An END statement is the last statement of a source program; it indicates to the assembler that all statements of the source program have been processed. An END statement is also used to cause execution of a main program to begin immediately after loading. To do this, the END statement requires an operand that represents the starting address of the program. At the completion of loading, execution begins at the address specified by the operand. For subroutines, all entry points are specified by ENT statements (described later); therefore, the operand of the END statement for a subroutine is blank.

The following statements illustrate both types of END statements.

	L	abe	el i		6	ре	rat	ion	F	T														Ope
21				25	27			30	32	33	35					40				45				50
					E	٨,	', l	2				,E	٦,	I,D), [–]	,0	<i>..</i> ,	J	2, R	2,0	G,	R,	A,/	И, ,
			,	,				,				,												
			1	1	ε	٨,	1,1	>,			G	0		,¢	R	A	N	CIA	1	17	0,		6,0	2, 1
						,	1.	-1					1	1				1	1				1	

DATA DEFINITION STATEMENTS

Data Definition statements are used to enter data constants into storage. The statements can be named by symbols so that other program statements can refer to the fields generated. Any type of data definition statement can be used in standard or extended precision program.

<u>DC</u> – Define Constant

The Define Constant statement is for generating constant data in main storage. Data can be specified as characters, hexadecimal numbers, decimal numbers, storage addresses or any valid expression. One 16bit word is generated for each DC statement. The format of this statement is shown below:

Label 21 25	Operation 27 30	F 32	T 33	35	40	45
L,A,B,E,L	DIC			A.N.	EXPRE	5,5,1,0,N,

If a label is used, the address assigned to it is the location of the generated data word and is equal to the current value of the Location Assignment Counter. Some examples of DC statements follow:

Labe I		Operation		F	T	Operande
21 25		27 30	, 🖉	32	33	35 40 45 50
$H_{1}E_{1}X_{1}$		D.C				V,F,F,F,F,F,H,E,X,CO,N,S,T
			1			
D,E,C,	ે	D.C.	2			-13,8,5, DEC 1,NTGER
		1-1-1				
A,L,P,H,A		D,C, ,				. B. C.H.A.R. C.O.N.S.T.
			ि			
A, D, D, R,S		D,C, ,				A.L.P.H.A.+ 5. A.D.D.R. C.O.N.

DEC - Decimal Data

The Decimal Data statement is used to enter binary data, expressed in decimal form, into a program. One DEC statement generates two 16-bit words of binary information. The format of the DEC statement is as follows:

Label 21 25	Operation 27 30	F 32	T 33	35	40	45
LA,BEL	D, E, C,			DEC	DATA	ITEM

If a label is used, its value is equal to the current value of the Location Assignment Counter if the current value is even; if the current value is odd, the label will be equal to the current value plus one. The label is assigned to the leftmost word of the generated constant. The types of data permitted in the operand field are described in the paragraphs entitled <u>Decimal</u> <u>Data Items</u>. An example of a DEC statement follows:

Label 21 25	Operation 27 30	F 32	T 33	35		40		45		
DATA.	$D_{1}E_{1}C_{1}$			+,	1.9	1	 1 1	1	1 1	
					L.11		 			1_1_

If the value of the Location Assignment Counter is 1000 when the DEC statement is encountered, the two words in storage look like this:

Location	Contents in Hexadecimal Form
01000	0000
01001	0013

Decimal Data Items

A decimal data item is used to specify, in decimal form, two or three words of data to be converted into binary form. Decimal data items are used in the

Ý

operand field of DEC assembler statements. Three types of decimal-data items are permitted: decimal integers, floating-point numbers, and fixed-point numbers. A floating-point decimal-data item can also be used as the operand of an XFLC statement that generates a 3-word constant.

<u>Decimal Integers</u>. A decimal integer is composed of a series of numeric digits with or without a preceding plus or minus sign. The allowable range of decimal integers is $-(2^{31}-1)$ to $2^{31}-1$.

Examples

Decimal Integer	Stored As	
50	000003216	
1535	000005FF16	
-3729	FFFF16F ₁₆	
	(2's complement)	

Floating-Point Numbers. A floating-point number has two components: a mantissa and an exponent.

- <u>Mantissa</u> The mantissa is a signed or unsigned decimal number, which can be written with or without a decimal point. The decimal point can appear at the beginning, at the end, or within the decimal number. If the exponent (see below) is present, the decimal point can be omitted, in which case it is assumed to be located at the right-hand end of the decimal number.
- Exponent The exponent consists of the letter E followed by a signed or unsigned decimal integer. The exponent part can be omitted if the mantissa contains a decimal point. If used, it must follow the mantissa.

A floating-point number is converted to a normalized, floating-point, binary number. The exponent part, if present, specifies a power of ten by which the mantissa is multiplied during conversion. For example, all of the following floating-point numbers are equivalent and will be converted to the same floating-point binary number.

4.500 45.00E-1 4500E-3 .4500E1

In standard precision, the above floating-point numbers are converted and stored in two consecutive storage locations as follows: $\frac{\text{Word 1}}{4800}$

Word 2 0083

The DEC assembler instruction stores floatingpoint numbers in the standard precision floatingpoint format described in the manual, <u>IBM 1130</u> Subroutine Library (Form C26-5929).

Fixed Point Numbers. A fixed-point number can have up to three components: a mantissa, an exponent, and a binary-point identifier.

- <u>Mantissa</u> The mantissa is the same as described for floating-point numbers.
- Exponent The exponent is the same as described for floating-point numbers.
- <u>Binary-Point Identifier</u> This identifier consists of the letter B, followed by a signed or unsigned decimal integer. The binary-point identifier must be present in a fixed-point number and must come after the mantissa. If the number has an exponent, the binary point identifier may precede or follow the exponent.

A fixed-point number is converted to a fixedpoint binary number that contains an understood binary point. The purpose of the binary-point identifier of the number is to specify the location of this understood binary point within the word. The number that follows the letter B specifies the number of binary places in the word to the left of the binary point (that is, the number of integral places in the word). The sign bit is not counted. Thus, a binary-point identifier of zero specifies a 31-bit binary fraction. B2 specifies two integral places and 29 fractional places. B31 specifies a binary integer. B-2 specifies a binary point located two places to the left of the leftmost bit of the word; that is, the word would contain the loworder 31 bits of binary fraction. As with floatingpoint numbers, the exponent, if present, specifies a power of ten by which the mantissa is multiplied during conversion.

A fixed-point number preceded by a minus sign is stored in 2's complement form.

The following fixed-point numbers all specify the same configuration of bits, but not all of them specify the same location for the understood binary point:

22.5B5 11.25B4 1125B4E-2

1125E-2B4 9B7E1

All of the above fixed-point numbers are converted to the same binary configuration, whose hexadecimal representation is:

Word 1	Word 2
5A00	0000

XFLC - Extended Floating Constant

The XFLC assembler instruction is used to introduce into a program an extended precision floating-point constant, expressed in three consecutive data words. When assembled, this instruction produces a format identical to the extended range floating-point format described in the manual, <u>IBM 1130 Subroutine</u> Library (Form C26-5929).

The format of the XFLC instruction is shown below:

	Label			Ope	ati	on		F	Ţ						Oper
21	2:	1	2	7		30	Ø.,	32	33	35			40	45	50
L	ABEL)	(_. F	L	,C				F	1,0	A,T	P.0		N,U,M
	1 1 1 1					1.		Γ			1 1.		1 1		1 1 1 1 1

The label is optional; if it is used, it is assigned to the location of the leftmost word generated.

Some examples of the XFLC instruction are shown below:

	Labe	I		Oper	ration		F	T				-			
23		25		27	30		32	33	35		40			45 -	
				X,F	, L ,C				0	5 . 3	5,1,2,	5		1 1 1	
						ціў.:			r				1.1	1 1 1	
F.	L.0	PT	्र	XF	LC	100			-10	5, 5	5,3,1	2,5		1.1.1	
					11					1.1				1 1 1	-
	1			X _. F	, L ,C				 51.	1,2	2, E, 2			1 1 1	

The data (in hexadecimal form) generated by each of these examples is

1.	Word 1	Word 2	Word 3
•	0080	4400	0000
2.	Word 1	Word 2	Word 3
	0080	BC00	0000
3.	Word 1	Word 2	Word 3
	008A	4000	0000

EBC – Extended Binary Coded Information

The EBC statement is used to generate data words, each consisting of two 8-bit characters in the Extended BCD Interchange Code (see <u>Appendix A</u>). Up to 18 sixteen-bit words can be generated with one EBC statement. The format of the statement is shown below:

Label 21 25	Operation 27 30	F 32	T 33 35	40	45
L.A.B.1.	E, B, C,			A,L,P,H,A	D,A,T,A
					1 1 1 1 1 1 1 1

If a label is present, it is assigned to the location of the leftmost word generated. The operand field contains the alphameric data to be represented in storage. This data must begin and end with a period. The data can be any valid character in the Extended BCD Interchange Code, including the period.

Examples

	Label		Operatio	'n		F	т	Γ										_
21		25	27	30		32	33	35		14		40			45			_
С	,0,N,S	٦ <u>,</u>	E,B,C						E_{1}	R ₁ R	°.0	RI.					<u>ц</u>	_
Γ													1.1	1				
A	LPH	I,A	E,B,C		•	1			C .	2,N	15	TIA	1, <i>N</i> ,	Τι.	Ļ		1.4	
	<u>і́ і і і</u>	1							_			I	<u> </u>			1		

The first example generates three words of data, with the location of the label CONST assigned to the leftmost location of the first word generated.

	Word 1	Word 2	Word 3
	C5D9	D9D6	D940
CONST	· ·		

Note that if the constant has an odd number of characters, as in the above example, the last word of data ends with the 8-bit equivalent of blank.

The second example generates four words of data:

Word 1	Word 2	Word 3	Word <u>4</u>
C3D6	D5E2	E3C1	D5E3

NOTE: A period may not appear in the comments field of an EBC instruction.

STORAGE ALLOCATION STATEMENTS

Storage allocation statements are used to reserve blocks of storage for data or work areas. Two such statements are available with the IBM 1130 Assembler: Block Started by Symbol and Block Ended by Symbol.

BSS - Block Started by Symbol

The BSS assembler instruction is used to reserve an area of core storage, within a program, for data storage or for working space. The format of the BSS instruction follows:

Label 21 25	Operation 2730	F 32	T 33	35 40 45
LA, B, E, L	B,S,S,			ABSOLUTE EXPR

The expression specifies the number of words to be reserved; the label, if specified, refers to the leftmost word reserved. The location of the block of storage within the object program is determined by the location of the BSS statement within the source program.

If the character E is punched in column 32, the assembler assigns the leftmost word of the reserved location to the next available <u>even</u> location. If a blank or any character other than E appears in column 32, the assembler assigns the leftmost word of the reserved area to the next available location regardless of whether that location is even or odd. This feature is useful when defining areas for use with double precision instructions.

A BSS statement with an E format and an operand value of zero causes the Location Assignment Counter to be made even (if necessary) before the next instruction is assembled.

A BSS instruction causes an area to be skipped, not cleared; therefore, it should not be assumed that an area reserved by a BSS instruction contains zeros.

Any symbols in the operand field of a BSS assembler instruction must have been previously defined. The expression in the operand field must be an absolute expression.

In the following example, the symbol AREA is equivalent to 3000; the next location assigned is 3028.

Label	Operation	F	T				
21 25	27 30	32	33		35	40	45
	O, R, G,				3,0,0,0		
A, R, E, A,	BSS				2,8, ,		
				ે			

BES - Block Ended by Symbol

The BES instruction is identical to the BSS instruction except that the address assigned to the label is the rightmost word in the area plus 1, i.e., the next location available for assignment.

In the previous example, the symbol AREA is equivalent to 3028.

SYMBOL DEFINITION STATEMENT

One symbol definition statement (EQU) is available in the IBM 1130 Assembler language.

EQU – Equate Symbol

The EQU statement is used to assign to a symbol a value other than the value of the Location Assignment Counter at the time the symbol is encountered. The format of the EQU statement is

Label		Oper	ation	F	т			
21	25	27	30	32	33	35	40	45
S, Y, M, L	9,2	EQ	U,			A.N.	EXPRE	S.S.1.0.N.
	, [

The symbol in the label field is made equivalent to the value of the expression. The expression may be absolute or relocatable. All symbols appearing in this expression must have appeared as a label in a previous statement. If an asterisk (*) is used as the expression, the value assigned to it is the next location to be assigned by the assembler.

Examples

Label	Operat	ion	F	T												
21 25	27	30	32	33	35				40			. 4	15			
N,A,M,E,	EQU	/.			2,0	6.	,			 ,	,	,			,	-
		,						,	,				,	1	,	
LOOP	E,Q,L	1			NJ	A ,/	M, C	5,7	1		-					-
									,							

In the first example, the symbol NAME is assigned a value of 26. In the second example, the symbol LOOP is assigned a value of 27.

PROGRAM-LINKING STATEMENTS

Program-linking statements are used to establish communication between a main program and its subroutines. Five statements are provided: three for entry point definition and two for subroutine calls.

ENT - Define Subroutine Entry Point

The ENT statement should be used to define the entry point(s) in all subroutines except ISS and ILS. Up to ten entry points may be defined for each subroutine (this would require ten ENT statements). The format of the ENT statement is shown below.

Label	Operation	FT	
21 25	27 30	32 33	35 40
	E,N,T		N,A,ME

NAME is a symbol which identifies an entry point for the associated subroutine. This symbol must be relocatable. All ENT statements for a given subroutine must be together and must precede all statements except LIBR, SPR, EPR, and comments statements. ENT, ISS, or ILS statements (see below) may not be used in the same subroutine.

ISS - Define Interrupt Service Entry Point

IBM provides interrupt service subroutines (ISS) for all devices; however, the user is given the option of replacing these subroutines with his own or adding a subroutine for interrupt run mode. The ISS statement is used to define an entry point in an interrupt service subroutine and to establish interrupt linkages to the subroutine at load time. Only one entry point may be defined for each subroutine. The format of the ISS statement is shown below.

Label	Operation		F	т			<u> </u>
21 25	27 30	Ì	32	33	35	40	45
	I,S,S,		N	N	N,A,M,E		<u> </u>
L							

Word 30 of the header record can be set for identification purposes as shown below. Word 30 is not used by any of the 1130 programs.

Label	ISS Header Word 30
blank	blank
1130	1
1800	2

NAME is as described for the ENT statement and NN (the ISS number) is a number from 01 to 20 used by the Relocating or core image loader to establish the linkage from the appropriate point in the corresponding ILS. The numbers and associated devices used in the routines provided by IBM are listed below:

<u>Number</u> *	Device or Function
01	1442 Card Read Punch
02	Input Keyboard; Console Printer
03	1134 Paper Tape Reader;
	1055 Paper Tape Punch
04	2310 Disk
06	1132 Printer
07	1627 Plotter

* Numbers 08 through 20 are assignable by the user.

L is a one-digit number used to indicate the interrupt level(s) associated with the subroutine. The level numbers (0-5) can be listed in any order in columns 45, 50, 55, 60, 65, and 70 and they must be left-justified in these columns (first in 45, second in 50 etc.).

An ISS statement must precede all statements except LIBR, SPR, EPR and comments statements.

Procedures for writing interrupt-service subroutines are provided in the publication <u>IBM 1130</u> Subroutine Library (C26-5929).

ILS - Define Interrupt Level Subroutine

IBM provides interrupt level subroutines for the various I/O devices and their associated interrupt levels; however, the user may replace these subroutines with his own. The ILS statement is used to define an interrupt level subroutine and to associate

the subroutine with a specific interrupt level. The format of the ILS statement is shown below:

Label	Operation	FT
21 25	27 30	32 33 35
	ILS.	NN

NN is the interrupt level number (00-05) associated with the interrupt level subroutine and is used by the relocating loader or core image converter. The devices associated with each interrupt level are shown below:

Device(s)
1442 Card Read Punch
1132 Printer
Disk Storage
1627 Plotter
Console Keyboard, Console
Printer, 1442 Card Read
Punch, 1134 Paper Tape
Reader, 1055 Paper Tape
Punch
PROGRAM STOP Key or
Interrupt Run Mode.

An ILS statement must precede all statements except LIBR, SPR, EPR, and comments statements.

Procedures for writing interrupt level subroutines are provided in the publication, <u>IBM 1130</u> <u>Subroutine Library</u> (Form C26-5929).

CALL - Call Direct Reference Subroutine

A CALL statement is used to call some of the subroutines in the IBM Subroutine Library or any userwritten subroutine written for the CALL statement. At execution time, this type of call takes the form of a long (two-word) BSI, to the entry point named in the CALL and the corresponding ENT or ISS statement.

When BSI is executed, the location of the first word following it is placed in the entry point location, and control is transferred to the first word following the entry point. The format of the CALL statement is:

	Label		Оре	ration	F	T		
21		25	27	30	32	33	35	40
L,	A,B,E	L	C,A	LL			N,A,M	Ε, , ,
1	. 1 . 1							
		1						

If used, the label is assigned to the current value of the Location Assignment Counter, which is the same as the leftmost word of the generated BSI instruction. The name of the called subroutine is assembled into the object deck, together with a unique code identifying the CALL. This code is recognized by the loader, which generates the BSI to the called subroutine.

<u>LIBF – Call TV (Transfer Vector) Reference</u> <u>Subroutine</u>

An LIBF statement is used to call any of the subroutines in the Subroutine Library (or any userwritten subroutine) written to utilize the transfer vector (see the following section, Subroutine Transfer Vector). The format of the LIBF statement is:

Labe	1	Oper	ration	F	T		
21	25	27	30	32	33	35	40
LAB	EL	LI	BF	Τ		N,A,M	ΙΕ, , ,

If used, the label is assigned to the current value of the Location Assignment Counter when the LIBF statement is encountered. The name of the called subroutine is assembled into the object deck, together with a unique code identifying the call as an LIBF call. This code is recognized by the loader, which generates a TV linkage to the called subroutine. During execution, the TV subroutine calls Index Register 3. Therefore, if Index Register 3 is used by any other instruction in the user's program, it must be saved and restored before it is needed by any TV subroutine calls. An LIBF Call is sometimes referred to as a One Word Call.

LIBF Subroutine Transfer Vector

To understand fully the use of the LIBF statement, the user should be familiar with the makeup of the subroutine transfer vector. A transfer vector allows main programs to communicate with relocatable subroutines (and relocatable subroutines to communicate with each other) without knowing where in core storage the subroutines are loaded. The transfer vector consists of three 16-bit words for each subroutine entry point referred to by an LIBF statement. The contents of the three words vary as the subroutine goes through the three phases of being called, loaded, and executed. The following paragraphs describe these three phases, and illustrate the contents of the transfer vector for each phase. Recognizing the Subroutine Call. All subroutines that utilize the transfer vector are called via LIBF statements. These statements take the following general form:

\mathbf{LIBF}	NAME
DC	Parameter
DC	Parameter
etc.	

When an LIBF call is recognized during the loading of an object program, the loader begins to build the transfer vector by setting aside a three-word record and placing in it the name of the called subroutine.

Name of Subroutine	Zeros

Subsequent LIBF statements produce additional records in the transfer vector, each containing a unique subroutine name. Calls to a subroutine previously listed in the transfer vector do not produce a new record. Each time an LIBF call is recognized, the loader replaces the associated LIBF NAME statement in the user's program with a one-word, indexed BSI instruction pointing to the first word of the associated transfer vector record. This instruction, generated by the loader, uses Index Register 3 and a computed displacement to refer to the proper transfer vector record.

Original Statement

21 25	Operation 27 30	F 32 3	T 13 35	40	45
	L,I,B,F		N _I A,M _I E		

Modified Statement

	Label			Operation		F	T			. <i>"</i>
21		Z5	673	27 30	878. I	32	33	382	35	40 45
				B,S,I			3		D,	1, S, P, L, A, C, E, M, E, N, T,

When this BSI instruction is encountered during execution of the main program, it causes a branch to the associated transfer vector record and from there to the entry address of the subroutine (see the following section, <u>Loading the Subroutine</u>). A BSI statement is generated for each LIBF statement encountered.

NOTE: Index Register 3 is reserved for LIBF subroutine calls. Therefore, if any instructions are to be tagged for register 3, register 3 should be restored prior to any LIBF subroutine call. Loading the Subroutine. After the main program has been read and all the called subroutine names are in the transfer vector, the loader loads the subroutines. As the header card of each subroutine in the library deck is read, the subroutine name is compared with the list of names previously placed in the transfer vector. When a match is found, the subroutine is loaded in an available core location, determined by the loader, and the subroutine entry address is placed in the third word of the associated transfer vector record. In addition, the address of the transfer vector record is inserted in the subroutine at entry point +2.



After the loader has read all of the subroutines in the subroutine deck and loaded those that were called, it modifies the transfer vector record of each subroutine that was loaded. This modification consists of deleting the subroutine names and placing, in the second word of each record, the first half of a long BSC instruction.

Undefined	BSC L	Entry address of Subroutine
-----------	-------	--------------------------------

Note that word 3 of the record is not affected and thus becomes the effective address of the long BSC instruction. This instruction is used to branch to a particular subroutine when that subroutine is called during execution of the main program.

NOTE: The preceding description of the subroutine loading process applies only to the card and paper tape systems and not to the Monitor system.

Executing the Subroutine. As the main program is executed, the subroutine calls are encountered one by one. Remember that each call is now in the form of a BSI instruction. The BSI instruction loads the contents of the IAR into the first word of the transfer vector record associated with the subroutine being called and then branches to the second word of the transfer vector record. As a result, the first word of the transfer vector record contains the address LIBF +1.

LIBF + 1 address	BSC L	Entry address of Subroutine
---------------------	-------	--------------------------------

Following execution of the BSI instruction, the BSC instruction in word 2 of the transfer vector record is executed, transferring control to the subroutine. In order to collect parameters and to return from a subroutine to the calling program, the programmer must know the address of the transfer vector record associated with the subroutine. This address was placed into the subroutine location, entry point +2, by the loader. Having this information, the programmer can find the address, LIBF +1. If parameters follow the statement, the address LIBF +1 (stored in the first word of the transfer vector record) is the address of the first parameter, otherwise it is the return address. A suggested method of computing the return address for a subroutine with one parameter is illustrated in Figure 5.

Size and Location of the Transfer Vector

The transfer vector consists of a maximum of 256 core locations and is located in the uppermost part of core storage (highest-numbered locations). If a

COMMON area is used for FORTRAN programs, the transfer vector is located below the COMMON area.

The location of the transfer vector in the Monitor system is different. The size of the transfer vector (256 locations) is sufficient to contain 85 three-word records. All but the first two of these records are available for use as subroutine communication areas. These two reserved records are used by certain arithmetic subroutines as a floatingpoint pseudo-accumulator, and as program error indicators.

Restrictions on the Use of the TV Reference Subroutines

- 1. No more than 83 unique entry point names referred to by LIBF statements can be loaded at one time.
- 2. In the card and paper tape systems, a subroutine called by an LIBF call must not contain any CALL (direct reference call) statements.

Label	Operation	F	т		11 181	- 1000 - 1 - 10 - 10 - 10 - 10 - 10 - 1	Operands	& Remarks			
21 25	27 30	3:	2 33	35	40	45	50	55	60	65	70
	L,1,8,F			S.U.B.	R	MALLN	P.R.O.G	RIAMI C	ALL		
	D,C,			AIRG	AD	DIRIEISIS	OF A	RIGIUME			
	•								I. J		
	•	•					<u></u>		1		
	•								J		
	•										
	END			STA	RT EN	$D_1 \circ_1 F_1$	MAJ/N	PROGK			
	.L.I.B.R									III I	
	$E_1N_1T_1$			S.U.B	$R_1 S_1 U_1 B_1$	$R_1O_1U_1T_1I_1$	NELHE	ADER	STAT	E, M, E, N, T	
S ₁ U ₁ B ₁ R ₁	5 ,T,X,		2	SAV	$E_1 + I_1 S_1$	$A_1V_1E_1 + I_1$	NDEX	RIEIG			
	LDX,	I	2	0	H_{1}	1,5, 1,0	C A T 1 0	N, WIII		BILANK	
					W,H	E_1N_1 S_1U	BRIIS	WRIT	Γ, <i>Τ</i> , <i>Ε</i> , <i>Ν</i> , , , Ε	$B_1 U_1 T_1 W_1 I_1$	LLI
					B,E,	$F_1/_1L_1L_1$	ED WI	THITH	I ENTI	e, 4, , A, D, D,	RESIS
					B.4.	$T_1H_1E_1$	LOADE	RULL			
	L.D.	1	2	0, 1	0, A, D, A,	RIGUME	NITLIN	T,0, A,	REG H	FORINS	ε
	• • • •			ר ב						-1	
	•• • •					المراجع المراجع					
	•			<u> </u>	AIIN B	0.0.4. 0	FUSUB	ROUTI	NE		
	•			سلسل		_1_!_I					
	•							<u></u>			
	MOX		2		A.D.D. 1	TO I	X 2 TO	OBT A	IIN RE	$\mathcal{F}_{1}\mathcal{T}_{1}\mathcal{U}_{1}\mathcal{R}_{1}\mathcal{N}_{1}$	ADDR
	SITIX		2	* +3	S.T.O.R.	$E_1 R_1 E_1 T_1$	URNIA	DDRES	$s_1 s_1 / N_1$	SAVET	3
$S_1A_1V_1E_1$	LDX	4	2	0	$R_1E_1S_1T_1O_1$	$\mathcal{R}_{1}\mathcal{E}_{1}$ $\mathcal{I}_{1}\mathcal{N}$	DEX	E1G1 121			
	B,S,C,	I		0	RETUR	N_1 , T_10_1	MAIN	PROGR			
	E _N D			بينا							

Figure 5. Example of a User-Written TV Reference Subroutine

The following assembler statements apply only to the IBM 1130 Monitor System. These statements are not valid for use with the card or paper tape assembler.

LINK - Load Link Program

In the assembler language, the LINK statement is used to load and execute another complete program. The program loaded and executed must be specified by name. The format of the LINK statement is:

- 1. A symbol or blanks in the label field
- 2. The mnemonic, LINK, in columns 27-30
- 3. A valid program name in the operand field

The label of the LINK pseudo-operation is defined as the current value of the Location Assignment Counter when the LINK statement is encountered; this value is the address of the first word generated by the LINK statement.

The operand field contains a valid program name (one to five alphameric characters), left-justified in the field. The name must be present in the Location Equivalence Table at execution time. The LINK statement causes four words to be generated in the object program. The first two words contain a two-word BSI instruction, which branches to a specified location within the Skeleton Supervisor program. The next two words contain the program name, left-justified, with blanks inserted in unused rightmost positions. The Supervisor Program uses the program name and begins the process required to load the new program.

EXIT - Return to Supervisor

In the assembler language, the EXIT statement is used to return control to the Supervisor. The format of the EXIT statement is:

- 1. A symbol or blanks in the label field
- 2. The mnemonic, EXIT, in columns 27-30

The label of the EXIT statement is defined as the current value of the Location Assignment Counter when the EXIT statement is encountered; this value is the address of the first word generated by an EXIT statement.

The operand field is ignored and can, therefore, be used for comments. The EXIT statement causes a one-word branch instruction to be generated in the object program. The instruction branches to a fixed location in the Skeleton Supervisor program. At object time the branch is executed and control is returned to the Supervisor. The EXIT statement should be the last logical statement in a program.

DSA – Define Sector Address

The DSA statement allows the programmer to refer symbolically to a disk-stored data file or core-image program without knowing the specific disk location of the data or program. The disk location of data files and programs can vary on disk because of deletions, but the DSA statement allows easy reference through the use of the symbolic name of the data file or program.

The format of the DSA statement is:

Label 21 25	Operation 27 30	F 32	T 33	35	40 45
L,A,B,E,L	D,S,A,			S	Y,M,B,O,L,I,C, N,A,M,E,

The label is defined as the current value of the Location Assignment Counter when the DSA statement is encountered. The symbol in the operand field must be the name of a data file or core-image program that is on disk at execution time.

The following statements illustrate the use of the DSA statement to read one sector of data. For a description of the disk calling sequences, see the publication <u>IBM 1130 Subroutine Library</u>, Form C26-5929.

Label	Operation	F	т	
1 25	27 30	32	33	35 40 45
	•			
	•			
· · · · ·	L, I, B, F			$D_I I_S K_I I_I$
	D.C.			/1,0,0,0,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,
	D.C.	•		$I \not 0 A R$
	D.C.			E.R.R.O.R.
T.Ø.A.R.	D.S.A.			$D_{A}T_{A}$
	B.S.S.			3,1,9, , , , , , , , , , , , , , , ,
	•		ŀ	
	•			

The Assembler reserves three words in the object program for each DSA statement. These words are filled in by the loader. For a data file they will contain:

Word 1 — Length (in words) Word 2 — Sector Address Word 3 — Not used

For a program they will contain:

Word 1 — Length (in words) Word 2 — Sector Address Word 3 — Execution Address of the Program

If the area corresponding to the DSA statement is used as the I/O area for a disk read operation, the execution address of the program must be saved prior to the disk call to bring in the program. (The contents of the third word are destroyed by the operation of the Disk I/O routine.)

The following statements illustrate the use of the DSA statement to supply the disk address of a one-sector program.

Labei		Operation	F	т			
21 25		27 30	32	33	35	40	45
		•, , ,	_		1. I		
		•					
		L _D ,			$I_{0}\phi_{A_{1}}R_{1}+$	121	
		S,T,O			BRNCH	1+1	
		L, I, B, F			$D_I I_S K_I$		
		DICI			1,1,0,0,0		
		D_1C_1			I.Ø.A.R.		
		DC			ERROK		
C, A, L, L, L	•	L, I, B, F			$D_I I_S K_1$		
		D,C, ,			10,0,0,0	2	
		D_1C_1			I, O, A, R		
		M,D,X,	L		CALL		
BRNCH		B,S,C,	2		0		1 1 1 1 1 1 1 1 1
		•	L				
		•				1_1_1	
		•				1.1.1	
$I_{0,A,R_{1}}$		D,S,A,			PRGRA	1	
		B,S,S,			3,1,9,		
		•			1 1 1 1	1 1 1	3 1 1 1 1 1 1
		•					

The following statements can be added to the previously shown program call to call a second pro-

gram and have it loaded to the same area as the first.

Label		Operation	F	т				
21 25		27 30	32	33		35	40	45
		L_D_{\perp}				A.D.R.	2	
		5,T,O,				110,A,	R	
		L.D.	L			A,D,R,	2,+,1	
		5,T,O,				I,0,A,	R+1	<u> </u>
	L	L.D.				A,D,R	2,+,2, ,	
		STO,		L		B,R,N,	$C_1H_1+_1I_1$	
		M,D,X			L	BRN	C,H,-,8,	
A, D, R, 2		$D_1S_1A_1$	L			PGR	M2 1	
				1				

The execution address of the second program can be different from the first, but the programs must be executable from the same locations. This requires a certain amount of planning before assembling the "overlay" programs.

Programming Considerations

The following considerations must be observed by the user who wishes to use the DSA statement to supply the disk address for programs.

- 1. The called programs must be in core-image format.
- 2. If the calling program is converted to core image format, the data for the DSA statement is filled in during the core image conversion and will be fixed for all subsequent executions. Thus, if the referenced program or data files are subsequently moved, incorrect results will occur.

HDNG - Heading

The HDNG statement is used to specify a one line page heading for a printed listing. The heading line consists of the data in the Operand-Remarks field (card columns 35-71, or the corresponding frames of an HDNG statement in paper tape format). Multiple HDNG statements may be used thus allowing different sections of a listing to have different page headings.

When the 1132 is the principal printer, the HDNG statement causes the listing to be ejected to a new page and the heading is printed. The same heading

is repeated at the top of each succeeding page until a new HDNG statement is encountered.

When the typewriter is the principal printer, the heading line is preceded and followed by a single line feed, and otherwise functions as a comments statement.

Hexadecimal Notation

In hexadecimal notation, each digit represents a four-bit binary value. This means that a 16-bit word in the Processor-Controller can be expressed as four hexadecimal digits. The binary — hexa-decimal — decimal correspondence is defined as follows:

Binary	Hexadecimal	Decimal
0000	0	0
0001	ů 1	1
0010	2	2
0011	3	3
0100	4	4
0101	5	5
0110	6	6
0111	7	7
1000	8	8
1001	9	9
1010	А	10
1011	В	11
1100	С	12
1101	D	13
1110	E	14
1111	\mathbf{F}	15

Extended Binary Coded Decimal Interchange Code (EBCDIC)

In the EBCDIC code, each character is represented by a unique configuration of eight binary bits. In the table that follows, each EBCDIC character is expressed as two hexadecimal digits.

IBM Card Code

In the IBM card code, each character represents a 12-bit card-column image. In the table that follows, each card code character is expressed as three hexa-decimal digits and as the card-column image.

Paper Tape Transmission Code, 8 Channel (PTTC/8)

In the PTTC/8 code, each character is represented by a unique configuration of a case shift, plus an eight-bit code. The case shift can be common to more than one character and need be inserted only when a case shift change is necessary. In the table that follows, each character is expressed as two hexadecimal digits, followed by the case shift in parentheses.

1132 Printer EBCDIC Subset Hex Code

In the 1132 Printer EBCDIC subset hex code, each character is represented by a unique configuration of eight bits. In the table that follows, each 1132 Printer character is expressed as two hexadecimal digits.

Console Printer Hex Code

In the Console Printer hexadecimal code each character is represented as two hexadecimal digits. * Recognized by all Conversion subroutines

NOTE: Codes that are not asterisked are recognized only by the SPEED subroutine.

	EBC	CDIC			IB	мc	ard (Code				1132	PTTC/8	Console
Ref No.	Binary 0123 45	567 H	lex	12	R 11	lows 09	8	7-1	Hex	Graph	ics and Control Names	Printer EBCDIC Subset Hex	Hex U-Upper Case L-Lower Case	Printer Hex _{Notes}
0 1 2 3 4 5* 6* 7* 8 9 10 11 12 13 14 15		000 0 001 0 010 0 011 0 100 0 101 0 101 0 0001 0 0001 0 0001 0 0001 0 0011 0 0100 0 1001 0 1001 0 1100 0 1110 0	00 01 02 03 04 05 06 07 08 07 08 09 0A 09 0A 0B 0C 0D 0E 0F	12 12 12 12 12 12 12 12 12 12 12 12 12 1		0 9 9 9 9 9 9 9 9 9 9 9 9 9 9 9 9 9 9 9	8 8 8 8 8 8 8 8 8 8 8 8	1 2 3 4 5 6 7 1 2 3 4 5 6 7	B030 9010 8810 8410 8210 8050 8050 8050 8030 9030 8830 8430 8230 8130 8080 8070	NUL PF HT LC DEL	Punch Off Horiz.Tab Lower Case Delete		6 D (S) 6 E (S) 7 F (S)	41 ①
16 17 18 19 20* 21* 22* 23 24 25 26 27 28 29 30 31	0001 00 000 00 00 00 00 00 00 00 00 00 00 00	000 001. 010 100 101 110 101 111 000 001 011 100 101 110 111	10 11 12 13 14 15 16 17 18 19 1A 1B 1C 1D 1E 1F	12	11 11 11 11 11 11 11 11 11 11 11	99999999999999999999	8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8	1 2 3 4 5 6 7 1 2 3 4 5 6 7	D030 5010 4810 4410 4210 4050 4050 4050 4030 5030 4830 4430 4230 4130 4080 4070	RES NL BS IDL	Restore New Line Backspace Idle		4C ⑤ DD ⑤ 5E ⑤	05 ② 81 ③ 11
32 33 34 35 36 37* 38* 39 40 41 42 43 44 45 46 47		000 2001 001 2010 010 2011 100 2011 101 2011 101 2000 001 2001 001 2010 010 2011 100 2010 101 2011 100 2010 111 2000	20 21 22 23 24 25 26 27 28 27 28 29 2A 28 29 2A 2B 2C 2D 2E 2F		11	0 9 9 0 9 9 0 9 9 0 9 9 0 9 9 0 9 9 0 9 9 0 9 9 0 9 9 0 9 9 0 9 9 9 0 9 9 9 0 0 9 9 0 0 9 9 0 0 9 9 0 0 0 9 9 0 0 0 9 9 0	8 8 8 8 8 8 8 8 8 8 8 8 8	1 2 3 4 5 6 7 1 2 3 4 5 6 7	7030 3010 2810 2410 2210 2050 2050 2050 2030 3030 2830 2430 2230 2230 22130 2080 2070	BYP LF EOB PRE	Bypass Line Feed End of Block Prefix		3D (S) 3E (S)	03
48 49 50 51 52 53* 54* 55 56 57 58 59 60 61 62 63		000 001 0000 0000 0000 0000 0000 0000	30 31 32 33 34 35 36 37 38 37 38 37 38 39 3A 3B 3C 3D 3E 3F	12	11	099999999999999999999999999999999999999	8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8	1 2 3 4 5 6 7 1 2 3 4 5 6 7	F030 1010 0810 0410 0210 0050 0050 0030 1030 0830 0430 0430 0430 0130 0080 0070	PN RS UC EOT	Punch On Reader Stop Upper Case End of Trans.		0D (S) 0E (S)	09 ④

NOTES: Typewriter Output ① Tabulate ② Shift to black

(5) The Same in Either Case

Carrier Return
 Shift to red

	EBCDIC		IBM Card Code			1132	PTTC/8	Console
Ref No.	Binary	Hex	Rows	Hex	Graphics and Control Names	Printer FBCDIC	Hex II-IInner Case	Printer
	0123 4567		12 11 0 9 8 7-1			Subset Hex	L-Lower Case	Hex
64* 65 66 67 68 69 70 71 72 73 74* 75* 76* 75* 78* 79*	0100 0000 0001 0010 0011 0100 0101 0111 1000 1001 1010 1011 1100 1101 1110	40 41 42 43 44 45 46 47 48 47 48 40 4A 4B 4C 4D 4E 4F	no punches 12 0 9 1 12 0 9 2 12 0 9 3 12 0 9 3 12 0 9 4 12 0 9 5 12 0 9 6 12 0 9 8 12 0 9 8 12 0 9 8 12 8 1 12 8 3 12 8 3 12 8 4 12 8 5 12 8 6 12 8 7	0000 B010 A810 A410 A110 A090 A050 A030 8820 8820 8820 8820 8120 80A0 8060	(space) ¢ . (period) < († I (logical OR)	* 4B 4D 4E	10 (5) 20 (U) 6B (L) 02 (U) 19 (U) 70 (U) 3B (U)	21 02 00 DE FE DA C6
80* 81 82 83 84 85 86 87 88 87 88 90* 91* 92* 93* 94* 95*	0101 0000 0001 0010 0010 0101 0101 0110 0111 1000 1001 1011 1100 1011 1110 1111	50 51 52 53 54 55 56 57 58 57 58 57 58 57 58 57 58 57 58 57 58 57 58	$\begin{array}{cccccccccccccccccccccccccccccccccccc$	8000 D010 C810 C410 C110 C090 C050 C030 5020 4820 4420 4420 4120 40A0 4060	& ! \$ * ; (logical NOT)	50 5B 5C 5D	5B (U) 5B (U) 5B (L) 08 (U) 1A (U) 13 (U) 6B (U)	44 42 40 D6 F6 D2 F2
96* 97* 98 99 100 101 102 103 104 105 106 107* 108* 109* 110* 111*	0110 0000 0001 0010 0011 0100 0101 0110 0111 1000 1001 1011 1100 1101 1110	60 61 62 63 64 65 66 67 68 67 68 60 66 66 66 66	$\begin{array}{cccccccccccccccccccccccccccccccccccc$	4000 3000 6810 6210 6110 6090 6050 6030 3020 2420 2220 2120 20A0 2060	- (dash) / (comma) % _ (underscore) ?	60 61 6B	40 (L) 31 (L) 38 (L) 15 (U) 40 (U) 07 (U) 31 (U)	84 BC 80 06 BE 46 86
112 113 114 115 116 117 118 119 120 121 122* 123* 124* 125* 126* 127*	0111 0000 0001 0010 0011 0100 0101 0110 0111 1000 1001 1010 1011 1100 1101 1100	70 71 72 73 74 75 76 77 78 79 7A 78 79 7A 7B 7C 7D 7E 7F	$\begin{array}{cccccccccccccccccccccccccccccccccccc$	E000 F010 E810 E410 E110 E050 E050 E030 E030 1020 0820 0420 0420 0120 00A0 0060	: # @ ' (apostrophe) = "	7D 7E	04 (U) 0B (L) 20 (L) 16 (U) 01 (U) 0B (U)	82 C0 04 E6 C2 E2

*Any code other than those defined will be interpreted by PRNT1 as a space.

			1	вм с	Card	Code			1132	PTTC/8	Console		
Ket No.	Bind	ary	Hex			Rows			Hex	Graphics and Control Names	EBCDIC	Hex U-Upper Case	Printer
	0123	4567		12	11	0 9	8	7-1			Subset Hex	L-Lower Case	пех
128 129 130 131 132 133 134 135 136 137 138 139 140 141 142 143	1000	0000 001 0010 0011 0100 0101 0110 0111 1000 1001 1010 1011 1100 1101 1110	80 81 82 83 84 85 86 87 88 87 88 88 88 88 88 80 88 88 88 88 88 88 88	12 12 12 12 12 12 12 12 12 12 12 12 12 1			8 8 8 8 8 8 8 8 8 8 8 8 8 8	1 2 3 4 5 6 7 2 3 4 5 6 7	B020 B000 A800 A400 A200 A100 A080 A020 A010 A820 A220 A120 A120 A120 A040 A040 A020 A020 A020 A000 A0200 A000 A0200 A000 A	a b c d e f ar i			
144 145 146 147 148 149 150 151 152 153 154 155 156 157 158 159	1001	0000 0001 0010 0011 0100 0101 0110 0111 1000 1001 1010 1001 1110 1110 1111	90 91 92 93 94 95 96 97 98 97 98 99 98 99 90 92 95 97	12 12 12 12 12 12 12 12 12 12 12 12 12 1	11 11 11 11 11 11 11 11 11 11 11 11	5	8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8	1 1 2 3 4 5 6 7 2 3 4 5 6 7	D020 D000 C800 C100 C080 C040 C020 C010 C820 C420 C420 C220 C120 C0A0 C060	j k I m n o P q r			
160 161 162 163 164 165 166 167 168 169 170 171 172 173 174 175	1010	0000 0001 0010 0011 0100 0101 0110 1001 1010 1011 1100 1101 1110 1111	A0 A1 A2 A3 A4 A5 A6 A7 A8 A7 A8 A7 A8 A0 AD AE AF))))))))))))))))))))		8 8 8 8 8 8 8 8 8 8 8 8 8 8	1 2 3 4 5 6 7 2 3 4 5 6 7	7020 7000 6800 6400 6100 6080 6040 6010 6820 6420 6420 6420 6120 60A0 6060	s t u v w x y z			
176 177 178 179 180 181 182 183 184 185 186 187 188 189 190 191	1011	0000 0001 0010 0011 0100 0101 0110 0111 1000 1001 1010 1011 1100 1101 1110	B0 B1 B2 B3 B4 B5 B6 B7 B8 B7 B8 B7 B8 B7 B8 B7 B8 B7 B8 B7 B8 B7 B8 B7 B8 B7 B8 B7 B8 B8 B7 B7 B8 B7 B7 B7 B7 B7 B7 B7 B7 B7 B7 B7 B7 B7	12 12 12 12 12 12 12 12 12 12 12 12 12 1	11 11 11 11 11 11 11 11 11 11 11 11 11	0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	8 8 8 8 8 8 8 8 8 8 8	1 1 2 3 4 5 6 7 2 3 4 5 6 7	F020 F000 E800 E400 E000 E000 E040 E020 E010 E820 E420 E420 E420 E120 E0A0 E060				

<u> </u>	EBCDIC					BM	Card	Code			1132	PTTC/8	Cancala
Ref	Binary	y	Hex			Rows			Hex	Graphics and Control	Printer		Printer
110.	0123	4567		12	11	09	8	7-1		i quilles	Subset Hex	L-Lower Case	Hex
192 193* 194* 195* 196* 197* 198* 199* 200* 201* 202 203 204 205 206 207		0000 0001 0010 0011 0100 0101 0110 0111 1000 1001 1010 1011 1100 1101 1110 1111	C0 C1 C2 C3 C4 C5 C7 C8 C7 C8 C7 C8 C7 CB CC CD CE CF	12 12 12 12 12 12 12 12 12 12 12 12 12 1		0 0 0 5 0 5 0 5 0 5 0 5 0 5 0 5 0 5 5 0 5 5 0 5 5 0 5	8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8 8	1 2 3 4 5 6 7 2 3 4 5 6 7	A000 9000 8800 8400 8200 8000 8010 A800 A8010 A830 A430 A430 A430 A430 A430 A080 A070	(+ zero) A B C D E F G H I	C1 C2 C3 C4 C5 C6 C7 C8 C9	61 (U) 62 (U) 73 (U) 64 (U) 75 (U) 76 (U) 67 (U) 68 (U) 79 (U)	3C or 3E 18 or 1A 1C or 1E 30 or 32 34 or 36 10 or 12 14 or 16 24 or 26 20 or 22
208 209* 210* 212* 213* 214* 215* 216* 217* 216* 217* 219 220 221 222 223		0000 0001 0010 0011 0100 0101 0110 0111 1000 1001 1011 1100 1101 1110 1111	D0 D1 D2 D3 D5 D6 D7 D8 D7 D7 D8 D7 D7 D7 D7 D7 D7 D7 D7 D7 D7 D7 D7 D7	12 12 12 12 12 12	11 11 11 11 11 11 11 11 11 11 11 11 11	0 9 9 9 9 9 9 9 9 9 9 9	8 8 8 8 8 8 8 8 8 8 8 8 8 8	1 2 3 4 5 6 7 2 3 4 5 6 7	6000 5000 4800 4200 4100 4080 4040 4020 4010 C830 C4300 C230 C130 C080 C070	(- zero) J K L M N O P Q R	D1 D2 D3 D4 D5 D6 D7 D8 D9	51 (U) 52 (U) 43 (U) 54 (U) 45 (U) 46 (U) 57 (U) 58 (U) 49 (U)	7C or 7 E 58 or 5A 5C or 5E 70 or 72 74 or 76 50 or 52 54 or 56 64 or 66 60 or 62
224 225 226* 227* 228* 230* 231* 232* 233* 233 235 236 237 238 239		0000 0001 0010 0011 0100 0111 0100 0111 1000 1011 1100 1101 1110 1111	E0 E1 E2 E3 E4 E5 E67 E89 EA ECD EE EF		11 11 11 11 11 11	0 9 0 9 0 0 0 0 0 9 0 9 0 9 0 9 0 9 0 9	8 8 8 8 8 8 8 8 8 8 8	2 1 2 3 4 5 6 7 2 3 4 5 6 7	2820 7010 2800 2400 2100 2080 2040 2040 2010 6830 6430 6430 6430 6430 6080 6070	S T U V W X Y Z	E2 E3 E4 E5 E6 E7 E8 E9	32 (U) 23 (U) 34 (U) 25 (U) 26 (U) 37 (U) 38 (U) 29 (U)	98 or 9A 9C or 9E B0 or B2 B4 or B6 90 or 92 94 or 96 A4 or A6 A0 or A2
240* 241* 242* 243* 244* 245* 246* 247* 249* 250 251 252 253 254 255		0000 0001 0010 0011 0100 0111 0100 0111 1000 1001 1010 1001 1100 1101 1110	F0 F1 F2 F3 F5 F6 F7 F8 F7 F8 F0 F0 F0 F0 F0 F0 F0 F0 F0 F1 F7 F7 F7 F7 F7 F7 F7 F7 F7 F7 F7 F7 F7	12 12 12 12 12 12	11 11 11 11 11 11	0 9 0 9 0 9 0 9 0 9 0 9 0 9 0 9 0 9	8 8 8 8 8 8 8 8 8 8	1 2 3 4 5 6 7 2 3 4 5 6 7	2000 1000 0800 0200 0100 0040 0040 0040 0040 0830 E830 E430 E430 E130 E080 E070	0 1 2 3 4 5 6 7 8 9	F0 F1 F2 F3 F4 F5 F6 F7 F8 F9	1A (L) 01 (L) 02 (L) 13 (L) 04 (L) 15 (L) 16 (L) 07 (L) 08 (L) 19 (L)	C4 FC D8 DC F0 F4 D0 D4 E4 E0

The tables printed below are used to convert decimal numbers to hexadecimal and hexadecimal numbers to decimal. In the descriptions that follow, the explanation of each step is followed by an example in parentheses. Decimal to Hexadecimal Conversion. Locate the decimal number (0489) in the body of the table. The two high-order digits (1E) of the hexadecimal number are in the left column on the same line, and the low-order digit (9) is at the top of the column. Thus, the hexadecimal number 1E9 is equal to the decimal number 0489.

<u>Hexadecimal to Decimal Conversion</u>. Locate the first two digits (1E) of the hexadecimal number (1E9) in the left column. Follow the line of figures across the page to the column headed by the low-order digit (9). The decimal number (0489) located at the junction of the horizontal line and the vertical column is the equivalent of the hexadecimal number.

Ç.	0	1	2	3	4	5	6	7	8	9	٨	В	с	D	E	F] [c.	- 0	1	2	3	4	5	6	7	8	9	٨	в	с	D	E	F
00 <u>-</u>	0000	0001	0002	0003	0004	0005	0006	0007	0008	0009	0010	0011	0012	0013	0014	0015		40 -	1024	1025	1026	1027	1028	1029	1030	1031	1032	1033	1034	1035	1036	1037	1038	1039
01 <u>-</u>	0016	0017	0018	0019	0020	0021	0022	0023	0024	0025	0026	0027	0028	0029	0030	0031		41 -	1040	1041	1042	1043	1044	1045	1046	1047	1048	1049	1050	1051	1052	1053	1054	1055
02 <u>-</u>	0032	0033	0034	0035	0036	0037	0038	0039	0040	0041	0042	0043	0044	0045	0046	0047		42 -	1056	1057	1058	1059	1060	1061	1062	1063	1064	1065	1066	1067	1068	1069	1070	1071
03 <u>-</u>	0048	0049	0050	0051	0052	0053	0054	0055	0056	0057	0058	0059	0060	0061	0062	0063		43 -	1072	1073	1074	1075	1076	1077	1078	1079	1080	1081	1082	1083	1084	1085	1086	1087
04 -	0064	0065	0066	0067	0068	0069	0070	0071	0072	0073	0074	0075	0076	0077	0078	0079		44 -	1088	1089	1090	1091	1092	1093	1094	1095	1096	1097	1098	1099	1100	1101	1102	1103
05 -	0080	0081	0082	0083	0084	0085	0086	0087	0088	0089	0090	0091	0092	0093	0094	0095		45 -	1104	1105	1106	1107	1108	1109	1110	1111	1112	1113	1114	1115	1116	1117	1118	1119
06 -	0096	0097	0098	0099	0100	0101	0102	0103	0104	0105	0106	0107	0108	0109	0110	0111		46 -	1120	1121	1122	1123	1124	1125	1126	1127	1128	1129	1130	1131	1132	1133	1134	1135
07 -	0112	0113	0114	0115	0116	0117	0118	0119	0120	0121	0122	0123	0124	0125	0126	0127		47 -	1136	1137	1138	1139	1140	1141	1142	1143	1144	1145	1146	1147	1148	1149	1150	1151
08 _	0128	0129	0130	0131	0132	0133	0134	0135	0136	0137	0138	0139	0140	0141	0142	0143		48 -	1152	1153	1154	1155	1156	1157	1158	1159	1160	1161	1162	1163	1164	1165	1166	1167
09 _	0144	0145	0146	0147	0148	0149	0150	0151	0152	0153	0154	0155	0156	0157	0158	0159		49 -	1168	1169	1170	1171	1172	1173	1174	1175	1176	1177	1178	1179	1180	1181	1182	1183
0A _	0160	0161	0162	0163	0164	0165	0166	0167	0168	0169	0170	0171	0172	0173	0174	0175		4A -	1184	1185	1186	1187	1188	1189	1190	1191	1192	1193	1194	1195	1196	1197	1198	1199
0B _	0176	0177	0178	0179	0180	0181	0182	0183	0184	0185	0186	0187	0188	0189	0190	0191		4B -	1200	1201	1202	1203	1204	1205	1206	1207	1208	1209	1210	1211	1212	1213	1214	1215
0C -	0192	0193	0194	0195	0196	0197	0198	0199	0200	0201	0202	0203	0204	0205	0206	0207		4C -	1216	1217	1218	1219	1220	1221	1222	1223	1224	1225	1226	1227	1228	1229	1230	1231
0D-	0208	0209	0210	0211	0212	0213	0214	0215	0216	0217	0218	0219	0220	0221	0222	0223		4D -	1232	1233	1234	1235	1236	1237	1238	1239	1240	1241	1242	1243	1244	1245	1246	1247
0E -	0224	0225	0226	0227	0228	0229	0230	0231	0232	0233	0234	0235	0236	0237	0238	0239		4E -	1248	1249	1250	1251	1252	1253	1254	1255	1256	1257	1258	1259	1260	1261	1262	1263
0F -	0240	0241	0242	0243	0244	0245	0246	0247	0248	0249	0250	0251	0252	0253	0254	0255		4F -	1264	1265	1266	1267	1268	1269	1270	1271	1272	1273	1274	1275	1276	1277	1278	1279
10 _	0256	0257	0258	0259	0260	0261	0262	0263	0264	0265	0266	0267	0268	0269	0270	0271		50 -	1280	1281	1282	1283	1284	1285	1286	1287	1288	1289	1290	1291	1292	1293	1294	1295
11 _	0272	0273	0274	0275	0276	0277	0278	0279	0280	0281	0282	0283	0284	0285	0286	0287		51 -	1296	1297	1298	1299	1300	1301	1302	1303	1304	1305	1306	1307	1308	1309	1310	1311
12 _	0288	0289	0290	0291	0292	0293	0294	0295	0296	0297	0298	0299	0300	0301	0302	0303		52 -	1312	1313	1314	1315	1316	1317	1318	1319	1320	1321	1322	1323	1324	1325	1326	1327
13 _	0304	0305	0306	0307	0308	0309	0310	0311	0312	0313	0314	0315	0316	0317	0318	0319		53 -	1328	1329	1330	1331	1332	1333	1334	1335	1336	1337	1338	1339	1340	1341	1342	1343
14 -	0320	0321	0322	0323	0324	0325	0326	0327	0328	0329	0330	0331	0332	0333	0334	0335		54 -	1344	1345	1346	1347	1348	1349	1350	1351	1352	1353	1354	1355	1356	1357	1358	1359
15 -	0336	0337	0338	0339	0340	0341	0342	0343	0344	0345	0346	0347	0348	0349	0350	0351		55 -	1360	1361	1362	1363	1364	1365	1366	1367	1368	1369	1370	1371	1372	1373	1374	1375
16 -	0352	0353	0354	0355	0356	0357	0358	0359	0360	0361	0362	0363	0364	0365	0366	0367		56 -	1376	1377	1378	1379	1380	1381	1382	1383	1384	1385	1386	1387	1388	1389	1390	1391
17 -	0368	0369	0370	0371	0372	0373	0374	0375	0376	0377	0378	0379	0380	0381	0382	0383		57 -	1392	1393	1394	1395	1396	1397	1398	1399	1400	1401	1402	1403	1404	1405	1406	1407
18 _	0384	0385	0386	0387	0388	0389	0390	0391	0392	0393	0394	0395	0396	0397	0398	0399		58 -	1408	1409	1410	1411	1412	1413	1414	1415	1416	1417	1418	1419	1420	1421	1422	1423
19 _	0400	0401	0402	0403	0404	0405	0406	0407	0408	0409	0410	0411	0412	0413	0414	0415		59 -	1424	1425	1426	1427	1428	1429	1430	1431	1432	1433	1434	1435	1436	1437	1438	1439
1A _	0416	0417	0418	0419	0420	0421	0422	0423	0424	0425	0426	0427	0428	0429	0430	0431		5A -	1440	1441	1442	1443	1444	1445	1446	1447	1448	1449	1450	1451	1452	1453	1454	1455
1B _	0432	0433	0434	0435	0436	0437	0438	0439	0440	0441	0442	0443	0444	0445	0446	0447		5B -	1456	1457	1458	1459	1460	1461	1462	1463	1454	1465	1466	1467	1468	1469	1470	1471
1C_	0448	0449	0450	0451	0452	0453	0454	0455	0456	0457	0458	0459	0460	0461	0462	0463		5C -	1472	1473	1474	1475	1476	1477	1478	1479	1480	1481	1482	1483	1484	1485	1486	1487
1D_	0464	0465	0466	0467	0468	0469	0470	0471	0472	0473	0474	0475	0476	0477	0478	0479		5D -	1488	1489	1490	1491	1492	1493	1494	1495	1496	1497	1498	1499	1500	1501	1502	1503
1E_	0480	0481	0482	0483	0484	0485	0486	0487	0488	0489	0490	0491	0492	0493	0494	0495		5E -	1504	1505	1506	1507	1508	1509	1510	1511	1512	1513	1514	1515	1516	1517	1518	1519
1F_	0496	0497	0498	0499	0500	0501	0502	0503	0504	0505	0506	0507	0508	0509	0510	0511		5F -	1520	1521	1522	1523	1524	1525	1526	1527	1528	1529	1530	1531	1532	1533	1534	1535
20 -	0512	0513	0514	0515	0516	0517	0518	0519	0520	0521	0522	0523	0524	0525	0526	0527		60 -	1536	1537	1538	1539	1540	1541	1542	1543	1544	1545	1546	1547	1548	1549	1550	1551
21 -	0528	0529	0530	0531	0532	0533	0534	0535	0536	0537	0538	0539	0540	0541	0542	0543		61 -	1552	1553	1554	1555	1556	1557	1558	1559	1560	1561	1562	1563	1564	1565	1566	1567
22 -	0544	0545	0546	0547	0548	0549	0550	0551	0552	0553	0554	0555	0556	0557	0558	0559		62 -	1568	1569	1570	1571	1572	1573	1574	1575	1576	1577	1578	1579	1580	1581	1582	1583
23 -	0560	0561	0562	0563	0564	0565	0566	0567	0568	0569	0570	0571	0572	0573	0574	0575		63 -	1584	1585	1586	1587	1588	1589	1590	1591	1592	1593	1594	1595	1596	1597	1598	1599
24 -	0576	0577	0578	0579	0580	0581	0582	0583	0584	0585	0586	0587	0588	0589	0590	0591		64 -	1600	1601	1602	1603	1604	1605	1606	1607	1608	1609	1610	1611	1612	1613	1614	1615
25 -	0592	0593	0594	0595	0596	0597	0598	0599	0600	0601	0602	0603	0604	0605	0606	0607		65 -	1616	1617	1618	1619	1620	1621	1622	1623	1624	1625	1626	1627	1628	1629	1630	1631
26 -	0608	0609	0610	0611	0612	0613	0614	0615	0616	0617	0618	0619	0620	0621	0622	0623		66 -	1632	1633	1634	1635	1636	1637	1638	1639	1640	1641	1642	1643	1644	1645	1646	1647
27 -	0624	0625	0626	0627	0628	0629	0630	0631	0632	0633	0634	0635	0636	0637	0638	0639		67 -	1648	1649	1650	1651	1652	1653	1654	1655	1656	1657	1658	1659	1660	1661	1662	1663
28 -	0640	0641	0642	0643	0644	0645	0646	0647	0648	0649	0650	0651	0652	0653	0654	0655		68 -	1664	1665	1666	1667	1668	1669	1670	1671	1672	1673	1674	1675	1676	1677	1678	1679
29 -	0656	0657	0658	0659	0660	0661	0662	0663	0664	0665	0666	0667	0668	0669	0670	0671		69 -	1680	1681	1682	1683	1684	1685	1686	1687	1688	1689	1690	1691	1692	1693	1694	1695
2A -	0672	0673	0674	0675	0676	0677	0678	0679	0680	0681	0682	0683	0684	0685	0686	0687		6A -	1696	1697	1698	1699	1700	1701	1702	1703	1704	1705	1706	1707	1708	1709	1710	1711
2B -	0688	0689	0690	0691	0692	0693	0694	0695	0696	0697	0698	0699	0700	0701	0702	0703		6B -	1712	1713	1714	1715	1716	1717	1718	1719	1720	1721	1722	1723	1724	1725	1726	1727
2C -	0704	0705	0706	0707	0708	0709	0710	0711	0712	0713	0714	0715	0716	0717	0718	0719		6C	1728	1729	1730	1731	1732	1733	1734	1735	1736	1737	1738	1739	1740	1741	1742	1743
2D -	0720	0721	0722	0723	0724	0725	0726	0727	0728	0729	0730	0731	0732	0733	0734	0735		6D	1744	1745	1746	1747	1748	1749	1750	1751	1752	1753	1754	1755	1756	1757	1758	1759
2E -	0736	0737	0738	0739	0740	0741	0742	0743	0744	0745	0746	0747	0748	0749	0750	0751		6E	1760	1761	1762	1763	1764	1765	1766	1767	1768	1769	1770	1771	1772	1773	1774	1775
2F -	0752	0753	0754	0755	0756	0757	0758	0759	0760	0761	0762	0763	0764	0765	0766	0767		6F	1776	1777	1778	1779	1780	1781	1782	1783	1784	1785	1786	1787	1788	1789	1790	1791
30 -	0768	0769	0770	0771	0772	0773	0774	0775	0776	0777	0778	0779	0780	0781	0782	0783		70 -	1792	1793	1794	1795	1796	1797	1798	1799	1800	1801	1802	1803	1804	1805	1806	1807
31 -	0784	0785	0786	0787	0788	0789	0790	0791	0792	0793	0794	0795	0796	0797	0798	0799		71 -	1808	1809	1810	1811	1812	1813	1814	1815	1816	-1817	1818	1819	1820	1821	1822	1823
32 -	0800	0801	0802	0803	0804	0805	0806	0807	0808	0809	0810	0811	0812	0813	0814	0815		72 -	1824	1825	1826	1827	1828	1829	1830	1831	1832	1833	1834	1835	1836	1837	1838	1839
33 -	0816	0817	0818	0819	0820	0821	0822	0823	0824	0825	0826	0827	0828	0829	0830	0831		73 -	1840	1841	1842	1843	1844	1845	1846	1847	1848	1849	1850	1851	1852	1853	1854	1855
34 -	0832	0833	0834	0835	0836	0837	0838	0839	0840	0841	0842	0843	0844	0845	0846	0847		74 -	1856	1857	1858	1859	1860	1861	1862	1863	1864	1865	1866	1867	1868	1869	1870	1871
35 -	0848	0849	0850	0851	0852	0853	0854	0855	0856	0857	0858	0859	0860	0861	0862	0863		75 -	1872	1873	1874	1875	1876	1877	1878	1879	1880	1881	1882	1883	1884	1885	1886	1887
36 -	0864	0865	0866	0867	0868	0869	0870	0871	0872	0873	0874	0875	0876	0877	0878	0879		76 -	1888	1889	1890	1891	1892	1893	1894	1895	1896	1897	1898	1899	1900	1901	1902	1903
37 -	0880	0881	0882	0883	0884	0885	0886	0887	0888	0889	0890	0891	0892	0893	0894	0895		77 -	1904	1905	1906	1907	1908	1909	1910	1911	1912	1913	1914	1915	1916	1917	1918	1919
38 -	0896	0897	0898	0899	0900	0901	0902	0903	0904	0905	0906	0907	0908	0909	0910	0911		78 -	1920	1921	1922	1923	1924	1925	1926	1927	1928	1929	1930	1931	1932	1933	1934	1935
39 -	0912	0913	0914	0915	0916	0917	0918	0919	0920	0921	0922	0923	0924	0925	0926	0927		79 -	1936	1937	1938	1939	1940	1941	1942	1943	1944	1945	1946	1947	1948	1949	1950	1951
3A -	0928	0929	0930	0931	0932	0933	0934	0935	0936	0937	0938	0939	0940	0941	0942	0943		7A -	1952	1953	1954	1955	1956	1957	1958	1959	1960	1961	1962	1963	1964	1965	1966	1967
3B -	0944	0945	0946	0947	0948	0949	0950	0951	0952	0953	0954	0955	0956	0957	0958	0959		7B -	1968	1969	1970	1971	1972	1973	1974	1975	1976	1977	1978	1979	1980	1981	1982	1983
3C -	0960	0961	0962	0963	0964	0965	0966	0967	0968	0969	0970	0971	0972	0973	0974	0975		7C -	1984	1985	1986	1987	1988	1989	1990	1991	1992	1993	1994	1995	1996	1997	1998	1999
3D -	0976	0977	0978	0979	0980	0981	0982	0983	0984	0985	0986	0987	0988	0989	0990	0991		7D -	2000	2001	2002	2003	2004	2005	2006	2007	2008	2009	2010	2011	2012	2013	2014	2015
3E -	0992	0993	0994	0995	0996	0997	0998	0999	1000	1001	1002	1003	1004	1005	1006	1007		7E -	2016	2017	2018	2019	2020	2021	2022	2023	2024	2025	2026	2027	2028	2029	2030	2031
3F -	1008	1009	1010	1011	1012	1013	1014	1015	1016	1017	1018	1019	1020	1021	1022	1023		7F -	2032	2033	2034	2035	2036	2037	2038	2039	2040	2041	2042	2043	2044	2045	2046	2047

34 1130 Assembler Language

ſ	0	1	2	3	4	5	6	7	8	9	A	В	С	D	Е	F	. Г	ſ	- 0	1	2	3	4	5	6	7	8	9	A	В	С	D	E	
80 - 81 - 82 - 83 -	2048 2064 2080 2096	2049 2065 2081 2097	2050 2066 2082 2098	2051 2067 2083 2099	2052 2068 2084 2100	2053 2069 2085 2101	2054 2070 2086 2102	2055 2071 2087 2103	2056 2072 2088 2104	2057 2073 2089 2105	2058 2074 2090 2106	2059 2075 2091 2107	2060 2076 2092 2108	2061 2077 2093 2109	2062 2078 2094 2110	2063 2079 2095 2111		C0- C1- C2- C3-	3072 3088 3104 3120	3073 3089 3105 3121	3074 3090 3106 3122	3075 3091. 3107 3123	3076 3092 3108	3077 3093 3109 3125	3078 3094 3110	3079 3095 3111	3080 3096 3112	3081 3097 3113	3082 3098 3114	3083 3099 3115	3084 3100 3116	3085 3101 3117	3086 3102 3118	
84 - 85 - 86 - 87	2112 2128 2144 2160	2113 2129 2145 2161	2114 2130 2146 2162	2115 2131 2147 2163	2116 2132 2148 2164	2117 2133 2149	2118 2134 2150	2119 2135 2151 2167	2120 2136 2152	2121 2137 2153 2169	2122 2138 2154	2123 2139 2155	2124 2140 2156	2125 2141 2157	2126 2142 2158	2127 2143 2159		C4 - C5 - C6 -	3136 3152 3168	3137 3153 3169	3138 3154 3170	3139 3155 3171	3140 3156 3172	3141 3157 3173	3142 3158 3174	3143 3159 3175	3144 3160 3176	3145 3161 3177	3146 3162 3178	3147 3163 3179	3148 3164 3180	3149 3165 3181	3150 3166 3182	3 3 3 3 3
88 _ 89 _ 8A _	2100 2176 2192 2208	2177 2193 2209	2178 2194 2210	2103 2179 2195 2211	2180 2196 2212	2103 2181 2197 2213	2182 2198 2214	2183 2199 2215	2103 2184 2200 2216	2103 2185 2201 2217	2186 2202 2218	2171 2187 2203 2219	2172 2188 2204 2220	2173 2189 2205 2221	2190 2206 2222	2173 2191 2207 2223		C8 - C9 - CA -	3200 3216 3232	3201 3217 3233	3202 3218 3234	3203 3219 3235	3188 3204 3220 3236	3189 3205 3221 3237	3190 3206 3222 3238	3191 3207 3223 3239	3192 3208 3224 3240	3193 3209 3225 3241	3194 3210 3226 3242	3195 3211 3227 3243	3196 3212 3228 3244	3197 3213 3229 3245	3198 3214 3230 3246	
8B 8C 8D 8E 8F	2224 2240 2256 2272 2288	2225 2241 2257 2273 2289	2226 2242 2258 2274 2290	2227 2243 2259 2275 2291	2228 2244 2260 2276 2292	2229 2245 2261 2277 2293	2230 2246 2262 2278 2294	2231 2247 2263 2279 2295	2232 2248 2264 2280	2233 2249 2265 2281 2297	2234 2250 2266 2282 2208	2235 2251 2267 2283	2236 2252 2268 2284 2300	2237 2253 2269 2285	2238 2254 2270 2286 2302	2239 2255 2271 2287		CB- CC- CD- CE-	3248 3264 3280 3296	3249 3265 3281 3297	3250 3266 3282 3298	3251 3267 3283 3299	3252 3268 3284 3300	3253 3269 3285 3301	3254 3270 3286 3302	3255 3271 3287 3303	3256 3272 3288 3304	3257 3273 3289 3305	3258 3274 3290 3306	3259 3275 3291 3307	3260 3276 3292 3308	3261 3277 3293 3309	3262 3278 3294 3310	
90 _ 91 _ 92 _	2304 2320 2336	2305 2321 2337	2306 2322 2338	2307 2323 2339	2308 2324 2340	2309 2325 2341	2310 2326 2342	2311 2327 2343	2312 2328 2344	2313 2329 2345	2314 2330 2346	2315 2331 2347	2316 2332 2348	2317 2333 2349	2318 2334 2350	2303 2319 2335 2351		D0 - D1 - D2 -	3328 3344 3360	3329 3345 3361	3330 3346 3362	3331 3347 3363	3832 3348 3364	3333 3349 3365	3334 3350 3366	3335 3351 3367	3336 3352 3368	3321 3337 3353 3369	3322 3338 3354 3370	3323 3339 3355 3371	3324 3340 3356 3372	3325 3341 3357 3373	3326 3342 3358 3374	
93 _ 94 _ 95 _ 96 _ 97 _	2352 2368 2384 2400 2416	2353 2369 2385 2401 2417	2354 2370 2386 2402 2418	2355 2371 2387 2403 2419	2356 2372 2388 2404 2420	2357 2373 2389 2405 2421	2358 2374 2390 2406 2422	2359 2375 2391 2407 2423	2360 2376 2392 2408 2424	2361 2377 2393 2409 2425	2362 2378 2394 2410 2428	2363 2379 2395 2411 2427	2364 2380 2396 2412 2428	2365 2381 2397 2413 2429	2366 2382 2398 2414 2430	2367 2383 2399 2415 2431		D3 - D4 - D5 - D6 - D7 -	3376 3392 3408 3424 3440	3377 3393 3409 3425 3441	3378 3394 3410 3426 3442	3379 3395 3411 3427 3443	3380 3396 3412 3428 3444	3381 3397 3413 3429 3445	3382 3398 3414 3430 3446	3383 3399 3415 3431 3447	3384 3400 3416 3432 3148	3385 3401 3417 3433 2449	3386 3402 3418 3434 2450	3387 3403 3419 3435 2451	3388 3404 3420 3436	3389 3405 3421 3437 2452	3390 3406 3422 3438	
98 - 99 - 9A - 9B -	2432 2448 2464 2480	2433 2449 2465 2481	2434 2450 2466 2482	2435 2451 2467 2483	2436 2452 2468 2484	2437 2453 2469 2485	2438 2454 2470 2486	2439 2455 2471 2487	2440 2456 2472 2488	2441 2457 2473 2489	2442 2458 2474 2490	2443 2459 2475 2491	2444 2460 2476 2492	2445 2461 2477 2493	2446 2462 2478 2494	2447 2463 2479 2495		D8 - D9 - DA - DB -	3456 3472 3488 3504	3457 3473 3489 3505	3458 3474 3490 3506	3459 3475 3491 3507	3460 3476 3492 3508	3461 3477 3493 3509	3462 3478 3494 3510	3463 3479 3495 3511	3464 3480 3496 3512	3465 3481 3497 3513	3466 3482 3498 3514	3467 3483 3499 3515	3468 3484 3500 3516	3469 3485 3501 3517	3470 3486 3502 3518	
9C _ 9D _ 9E _ 9F _	2496 2512 2528 2544	2497 2513 2529 2545	2498 2514 2530 2546	2499 2515 2531 2547	2500 2516 2532 2548	2501 2517 2533 2549	2502 2518 2534 2550	2503 2519 2535 2551	2504 2520 2536 2552	2505 2521 2537 2553	2506 2522 2538 2554	2507 2523 2539 2555	2508 2524 2540 2556	2509 2525 2541 2557	2510 2526 2542 2558	2511 2527 2543 2559		DC. DD. DE. DF.	3520 3536 3552 3568	3521 3537 3553 3569	3522 3538 3554 3570	3523 3539 3555 3571	3524 3540 3556 3572	3525 3541 3557 3573	3526 3542 3558 3574	3527 3543 3559 3575	3528 3544 3560 3576	3529 3545 3561 3577	3530 3546 3562 3578	3531 3547 3563 3579	3532 3548 3564 3580	3533 3549 3565 3581	3534 3550 3566 3582	
A0 _ A1 _ A2 _ A3 _	2560 2576 2592 2608	2561 2577 2593 2609	2562 2578 2594 2610	2563 2579 2595 2611	2564 2580 2596 2612	2565 2581 2597 2613	2566 2582 2598 2614	2567 2583 2599 2615	2568 2584 2600 2616	2569 2585 2601 2617	2570 2586 2602 2618	2571 2587 2603 2619	2572 2588 2604 2620	2573 2589 2605 2621	2574 2590 2606 2622	2575 2591 2607 2623		E0 _ E1 _ E2 _ E3 _	3584 3600 3616 3632	3585 3601 3617 3633	3586 3602 3618 3634	3587. 3603 3619 3635	3583 3604 3620 3636	3589 3605 3621 3637	3590 3606 3622 3638	3591 3607 3623 3639	3592 3608 3624 3640	3593 3609 3625 3641	3594 3610 3626 3642	3595 3611 3627 3643	3596 3612 3628 3644	3597 3613 3629 3645	3598 3614 3630 3646	
A4 A5 A6 A7	2624 2640 2656 2672	2625 2641 2657 2673	2626 2642 2658 2674	2627 2643 2659 2675	2628 2644 2660 2676	2629 2645 2661 2677	2630 2646 2662 2678	2631 2647 2663 2679	2632 2648 2664 2680	2633 2649 2665 2681	2634 2650 2666 2682	2635 2651 2667 2683	2636 2652 2668 2684	2637 2653 2669 2685	2638 2654 2670 2686	2639 2655 2671 2687		E4 - E5 - E6 - E7 -	3648 3664 3680 3696	3649 3665 3681 3697	3650 3666 3682 3698	3651 3667 3683 3699	3652 3668 3684 3700	3653 3669 3685 3701	3654 3670 3686 3702	3655 3671 3687 3703	3656 3672 3688 3704	3657 3673 3689 3705	3658 3674 3690 3706	3659 3675 3691 3707	3660 3676 3692 3708	3661 3677 3693 3709	3662 3678 3694 3710	
A8 _ A9 _ AA _ AB _	2688 2704 2720 2736	2689 2705 2721 2737	2690 2706 2722 2738	2691 2707 2723 2739	2692 2708 2724 2740	2693 2709 2725 2741	2694 2710 2726 2742	2695 2711 2727 2743	2696 2712 2728 2744	2697 2713 2729 2745	2698 2714 2730 2746	2699 2715 2731 2747	2700 2716 2732 2748	2701 2717 2733 2749	2702 2718 2734 2750	2703 2719 2735 2751		E8 - E9 - EA - EB -	3712 3728 3744 3760	3713 3729 3745 3761	3714 3730 3746 3762	3715 3731 3747 3763	3716 3732 3748 3764	3717 3733 3749 3765	3718 3734 3750 3766	3719 3735 3751 3767	3720 3736 3752 3768	3721 3737 3753 3769	3722 3738 3754 3770	3723 3739 3755 3771	3724 3740 3756 3772	3725 3741 3757 3773	3726 3742 3758 3774	
AC_ AD_ AE_ AF_	2752 2768 2784 2800	2753 2769 2785 2801	2754 2770 2786 2802	2755 2771 2787 2803	2756 2772 2788 2804	2757 2773 2789 2805	2758 2774 2790 2806	2759 2775 2791 2807	2760 2776 2792 2808	2761 2777 2793 2809	2762 2778 2794 2810	2763 2779 2795 2811	2764 2780 2796 2812	2765 2781 2797 2813	2766 2782 2798 2814	2767 2783 2799 2815		EC - ED - EE - EF -	3776 3792 3808 3824	3777 3793 3809 3825	3778 3794 3810 3826	3779 3795 3811 3827	3780 3796 3812 3828	3781 3797 3813 3829	3782 3798 3814 3830	3783 3799 3815 3831	3784 3800 3816 3832	3785 3801 3817 3833	3786 3802 3818 3834	3787 3803 3819 3835	3788 3804 3820 3836	3789 3805 3821 3837	3790 3806 3822 3838	00000
B0 _ B1 _ B2 _ B3 _	2816 2832 2848 2864	2817 2833 2849 2865	2818 2834 2850 2866	2819 2835 2851 2867	2820 2836 2852 2863	2821 2837 2853 2869	2822 2838 2854 2870	2823 2839 2855 2871	2824 2840 2856 2872	2825 2841 2857 2873	2826 2842 2858 2874	2827 2843 2859 2875	2828 2844 2860 2876	2829 2845 2861 2877	2830 2846 2862 2878	2831 2847 2863 2879		F0 - F1 - F2 - F3 -	3840 3856 3872 3888	3841 3857 3873 3889	3842 3858 3874 3890	3843 3859 3875 3891	3844 3860 3876 3892	3845 3861 3877 3893	3846 3862 3878 3894	3847 3863 3879 3895	3848 3864 3880 3896	3849 3865 3881 3897	3850 3866 3882 3898	3851 3867 3883 3899	3852 3868 3884 3900	3853 3869 3885 3901	3854 3870 3886 3902	3333
B4 - B5 - B6 - B7 -	2880 2896 2912 2928	2881 2897 2913 2929	2882 2898 2914 2930	2883 2899 2915 2931	2884 2900 2916 2932	2885 2901 2917 2933	2886 2902 2918 2934	2887 2903 2919 2935	2888 2904 2920 2936	2889 2905 2921 2937	2890 2906 2922 2938	2891 2907 2923 2939	2892 2908 2924 2940	2893 2909 2925 2941	2894 2910 2926 2942	2895 2911 2927 2943		F4 - F5 - F6 - F7 -	3904 3920 3936 3952	3905 3921 3937 3953	3906 3922 3938 3954	3907 3923 3939 3955	3908 3924 3940 3956	3909 3925 3941 3957	3910 3926 3942 3958	3911 3927 3943 3959	3912 3928 3944 3960	3913 3929 3945 3961	3914 3930 3946 3962	3915 3931 3947 3963	3916 3932 3948 3964	3917 3933 3949 3965	3918 3934 3950 3966	0000
B8 _ B9 _ BA _ BB _	2944 2960 2976 2992	2945 2961 2977 2993	2946 2962 2978 2994	2947 2963 2979 2995	2948 2964 2980 2996	2949 2965 2981 2997	2950 2966 2982 2998	2951 2967 2983 2999	2952 2968 2984 3000	2953 2969 2985 3001	2954 2970 2986 3002	2955 2971 2987 3003	2956 2972 2988 3004	2957 2973 2989 3005	2958 2974 2990 3006	2959 2975 2991 3007		F8 _ F9 _ FA _ FB _	3968 3984 4000 4016	3969 3985 4001 4017	3970 3986 4002 4018	3971 3987 4003 4019	3972 3988 4004 4020	3973 3989 4005 4021	3974 3990 4008 4022	3975 3991 4007 4023	3976 3992 4008 4024	3977 3993 4009 4025	3978 3994 4010 4026	3979 3995 4011 4027	3980 3996 4012 4028	3981 3997 4013 4029	3982 3998 4014 4030	3 3 4 4
BC _ BD _ BE _ BF _	3008 3024 3040 3056	3009 3025 3041 3057	3010 3026 3042 3058	3011 3027 3043 3059	3012 3028 3044 3060	3013 3029 3045 3061	3014 3030 3046 3062	3015 3031 3047 3063	3016 3032 3048 3064	3017 3033 3049 3065	3018 3034 3050 3066	3019 3035 3051 3067	3020 3036 3052 3068	3021 3037 3053 3069	3022 3038 3054 3070	3023 3039 3055 3071		FC - FD - FE - FF -	4032 4048 4064 4080	4033 4049 4065 4081	4034 4050 4066 4082	4035 4051 4067 4083	4038 4052 4068 4084	4037 4053 4069 4085	4038 4054 4070 4088	4039 4055 4071 4087	4040 4056 4072 4088	4041 4057 4073 4089	4042 4058 4074 4090	4043 4059 4075 4091	4044 4060 4076 4092	4045 4061 4077 4093	4046 4062 4078 4094	4444
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Dec	Bin	Hex	Dec	Bin	Hex
0	0000	0	8	1000	8
1	0001	1	9	1001	9
2	0010	2	10	1010	A
3	0011	3	11	1011	В
4	0100	4	12	1100	C
5	0101	5	13	1101	D
6	0110	6	14	1110	Е
7	0111	7	15	1111	F

The table to the left gives the decimal, binary, and hexadecimal coding for the full range of fourbinary bits, from zero through F_{16} and 15_{10} .

To convert a four-digit hexadecimal number to decimal, determine the decimal value of the three low-order hexadecimal digits in the main table, and add the value for the high-order digit, as shown in the

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extended chart to the right.

For conversion of decimal values beyond the main table, deduct the largest number in the table at the right that will yield a positive result. The related digit is the highorder hexadecimal digit. Determine the three remaining hexadecimal digits by converting the product of the above subtraction in the main table.

Hex	Dec	Hex	Dec
1000	4096	9000	36864
2000	8192	A000	40960
3000	12288	B000	45056
4000	16384	C000	49152
5000	20480	D000	53248
6000	24576	E000	57344
7000	28672	F000	61440
8000	32768		

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