

RRRRRRRRRRRR		MMM		MMM	SSSSSSSSSSSS	
RRRRRRRRRRRR		MMM		MMM	SSSSSSSSSSSS	
RRRRRRRRRRRR		MMM		MMM	SSSSSSSSSSSS	
RRR	RRR	MMMMMM	MMMMMM	SSS		
RRR	RRR	MMMMMM	MMMMMM	SSS		
RRR	RRR	MMMMMM	MMMMMM	SSS		
RRR	RRR	MMM	MMM	MMM	SSS	
RRR	RRR	MMM	MMM	MMM	SSS	
RRR	RRR	MMM	MMM	MMM	SSS	
RRRRRRRRRRRR		MMM		MMM	SSSSSSSSSS	
RRRRRRRRRRRR		MMM		MMM	SSSSSSSSSS	
RRRRRRRRRRRR		MMM		MMM	SSSSSSSSSS	
RRR	RRR	MMM		MMM		SSS
RRR	RRR	MMM		MMM		SSS
RRR	RRR	MMM		MMM		SSS
RRR	RRR	MMM		MMM		SSS
RRR	RRR	MMM		MMM		SSS
RRR	RRR	MMM		MMM		SSS
RRR	RRR	MMM		MMM		SSS
RRR	RRR	MMM		MMM	SSSSSSSSSSSS	
RRR	RRR	MMM		MMM	SSSSSSSSSSSS	
RRR	RRR	MMM		MMM	SSSSSSSSSSSS	

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[illegible]


```
1 0001 0 MODULE RM3SPLUDR (LANGUAGE (BLISS32) ,
2 0002 0 IDENT = 'V04-000'
3 0003 0 ) =
4 0004 1 BEGIN
5 0005 1
6 0006 1 *****
7 0007 1 *
8 0008 1 * COPYRIGHT (c) 1978, 1980, 1982, 1984 BY
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25 0025 1 *
26 0026 1 *
27 0027 1 *****
28 0028 1
29 0029 1 ++
30 0030 1
31 0031 1 FACILITY: RMS32 INDEX SEQUENTIAL FILE ORGANIZATION
32 0032 1
33 0033 1 ABSTRACT:
34 0034 1 split user data record buckets
35 0035 1
36 0036 1
37 0037 1 ENVIRONMENT:
38 0038 1
39 0039 1 VAX/VMS OPERATING SYSTEM
40 0040 1
41 0041 1 --
42 0042 1
43 0043 1
44 0044 1 AUTHOR: Wendy Koenig CREATION DATE: 5-JUL-78 14:46
45 0045 1
46 0046 1
47 0047 1 MODIFIED BY:
48 0048 1
49 0049 1 V03-013 JWT0157 Jim Teague 23-Feb-1984
50 0050 1 When RMS attempted to calculate whether a series of
51 0051 1 duplicate records (including the new record) would
52 0052 1 fit within a single bucket, it neglected to account
53 0053 1 for the fact that the first record in the chain will
54 0054 1 undergo full expansion when it ends up as the first
55 0055 1 record in the new bucket. If it is currently partially
56 0056 1 compressed based on the previous key, then that could
57 0057 1 (and sometimes DID) cause bucket overflow when the
```


58 0058 1
59 0059 1
60 0060 1
61 0061 1
62 0062 1
63 0063 1
64 0064 1
65 0065 1
66 0066 1
67 0067 1
68 0068 1
69 0069 1
70 0070 1
71 0071 1
72 0072 1
73 0073 1
74 0074 1
75 0075 1
76 0076 1
77 0077 1
78 0078 1
79 0079 1
80 0080 1
81 0081 1
82 0082 1
83 0083 1
84 0084 1
85 0085 1
86 0086 1
87 0087 1
88 0088 1
89 0089 1
90 0090 1
91 0091 1
92 0092 1
93 0093 1
94 0094 1
95 0095 1
96 0096 1
97 0097 1
98 0098 1
99 0099 1
100 0100 1
101 0101 1
102 0102 1
103 0103 1
104 0104 1
105 0105 1
106 0106 1
107 0107 1
108 0108 1
109 0109 1
110 0110 1
111 0111 1
112 0112 1
113 0113 1
114 0114 1

duplicate chain is moved merrily into the new bucket.
Keep track of the compression count for the first
record in the dup chain, and add it to the total
size of the chain before comparing to bucket size.

V03-012 JWT0142 Jim Teague 16-Dec-1983
Correct incorrect bucket VBN comparison.

V03-011 MCN0008 Maria del C. Nasr 22-Mar-1983
More changes in the linkages

V03-010 MCN0007 Maria del C. Nasr 28-Feb-1983
Reorganize linkages

V03-009 TMK0004 Todd M. Katz 10-Nov-1982
At the present time, under certain circumstances, the number
of RRVs which will be required to be created when a simple
two-bucket split is done is being incorrectly calculated. This
will happen only during SUPDATES when the record being updated
is to go into the old (left) bucket and prior to the split is
in its original bucket. Even then it does not happen under all
possible circumstances, but only when duplicate records are
involved. It is possible that the number of RRVs calculated to
be required will be several less than the actual number which
will be needed. Under certain circumstances, the number of RRVs
needed may actually be calculated as a negative number - an
impossibility. Much depends upon the bucket composition. While
this does not influence the actual creation of RRVs, what it
does affect is where the bucket split point is calculated to be
since RRVs to be created do take up space in the old (left)
bucket. In fact, this problem came to my attention because of
the occurrence of a bucket split which resulted in the right
bucket, the new bucket, being empty, and the old (left) bucket
containing all the records even though there was no room for
them (or the bucket split would not have been required in the
first place). This split was caused by the number of RRVs
required being calculated as -1 instead of 0 such that, instead
of having the RRV spacial requirements added to the left bucket
size requirements, they were subtracted.

To fix this problem I have adjusted how the number of needed
RRVs are to be calculated. To start, the number of needed RRVs
is calculated to be the number of records (including the record
being updated which is not currently in the bucket) whose
original bucket is the bucket splitting. Then, as the split
point of the bucket is adjusted from left to right, this number
is decremented as records (which are in their original bucket)
are designated to stay in the left or old bucket. This is where
my change comes in. Previously, that the updated record was to
stay in the old bucket was determined at several different
points, and each time the count of the number of needed RRVs
was decremented, as long as the other conditions were met.
Unfortunately, this allowed for this determination to take place
more than once, and for the RRV count to be decremented multiple
times for the same record. My fix prevents this from occurring.
While it is still determined in several places that the updated
record is to go in the old bucket, I have made sure that those

115 0115 1
116 0116 1
117 0117 1
118 0118 1
119 0119 1
120 0120 1
121 0121 1
122 0122 1
123 0123 1
124 0124 1
125 0125 1
126 0126 1
127 0127 1
128 0128 1
129 0129 1
130 0130 1
131 0131 1
132 0132 1
133 0133 1
134 0134 1
135 0135 1
136 0136 1
137 0137 1
138 0138 1
139 0139 1
140 0140 1
141 0141 1
142 0142 1
143 0143 1
144 0144 1
145 0145 1
146 0146 1
147 0147 1
148 0148 1
149 0149 1
150 0150 1
151 0151 1
152 0152 1
153 0153 1
154 0154 1
155 0155 1
156 0156 1
157 0157 1
158 0158 1
159 0159 1
160 0160 1
161 0161 1
162 0162 1
163 0163 1
164 0164 1
165 0165 1
166 0166 1
167 0167 1
168 0168 1
169 0169 1
170 0170 1
171 0171 1

places are orthogonal to one another, so that the RRV count is not decremented more than once for the same record, the record whose update is causing the split.

V03-008 KBT0234 Keith B. Thompson 23-Aug-1982
Reorganize psects

V03-007 TMK0003 Todd M. Katz 02-Jul-1982
Implement RMS cluster solution for next record positioning. The next record positioning context is now kept in the IRAB, where it may be retrieved from, instead of in the NRP list which has been eliminated. When referring to the RFA address of the new/changed primary data record use the subfields IRBSL_PUTUP_VBN and IRBSW_PUTUPD_ID.

V03-006 MCN0006 Maria del C. Nasr 29-Jun-1982
Allow keys of different data types other than string in prologue 3 files.
Change all CH\$COMPARE calls to RM\$COMPARE_KEY to compare keys taking into consideration the different data types.

V03-005 MCN0005 Maria del C. Nasr 11-Jun-1982
Eliminate overhead at end of data bucket that was to be used for duplicate continuation bucket processing.

V03-004 TMK0002 Todd M. Katz 31-May-1982
Performance enhancements. I have made four changes to the routine RM\$SPLIT_UDR_3 which should cut down on the length of the bucket scans required at various times to re-expand keys. The enhancements involve setting the IRAB field IRBSL_LST_NCMP to the current record if key compression is enabled and the key of the current record is zero front compressed during various bucket scans required in the determination of the split point(s). The first of these scans is made to position to and extract the key of the last record in the bucket. The second and third scans are when the split code has decided that the best split point is one record previous to the current position, and must scan the bucket to obtain the record previous to that position so its key can be extracted and used during the index update. The fourth scan is the right-to-left record-by-record scan made to decide whether a two-bucket split is possible, and if so, where is the best place to split. In all four cases, I have added code to set the last noncompressed record pointer before continuing the scan with the next record, if the current record was zero front compressed.

V03-003 TMK0001 Todd M. Katz 10-May-1982
The algorithm for determining the split point of a prologue three data bucket with compressed keys first determines whether a two-bucket split can be done by scanning the old bucket from left-to-right record-by-record determining whether the lefthand sides and righthand sides of each possible split point will fit into a bucket. This size determination must take into account the position of insertion of the new (or updated) record, and the size determination of the righthand side must take into account the number of characters currently front

172 0172 1
173 0173 1
174 0174 1
175 0175 1
176 0176 1
177 0177 1
178 0178 1
179 0179 1
180 0180 1
181 0181 1
182 0182 1
183 0183 1
184 0184 1
185 0185 1
186 0186 1
187 0187 1
188 0188 1
189 0189 1
190 0190 1
191 0191 1
192 0192 1
193 0193 1
194 0194 1
195 0195 1
196 0196 1
197 0197 1
198 0198 1
199 0199 1
200 0200 1
201 0201 1
202 0202 1
203 0203 1
204 0204 1
205 0205 1
206 0206 1
207 0207 1
208 0208 1
209 0209 1
210 0210 1
211 0211 1
212 0212 1
213 0213 1
214 0214 1
215 0215 1
216 0216 1
217 0217 1
218 0218 1
219 0219 1
220 0220 1
221 0221 1
222 0222 1
223 0223 1
224 0224 1
225 0225 1
226 0226 1
227 0227 1
228 0228 1

compressed of what will become its low-order (and thus non-compressed key) record. What was missing, and what this change rectifies, is that what may become the low-order record of the righthand bucket is in fact the new (updated) record whose insertion is forcing this split to take place. In this case, the number of front compressed characters to be added to the righthand side total must come from the compressed key in keybuffer 5, if this is an \$UPDATE, or from the compressed key in the record buffer whose address is stored in IRB\$L_RECBUF, if this is a \$PUT. This change will be included as a patch on the V3.1 update floppy.

V03-002 MCN0004 Maria del C. Nasr 31-Mar-1982
Do not count records that will not need rrv's when moved out of the bucket. Their id's cannot be recycled in plg 3 files.

V03-001 MCN0003 Maria del C. Nasr 25-Mar-1982
Use macro to calculate keybuffer address.

V02-016 DJD0001 Darrell Duffy 1-March-1982
Fix references to RBF for better probing

V02-015 MCN0002 Maria del C. Nasr 09-Jul-1981
Fix a problem with update of the first record in a duplicate chain, in both old code, and new code. Also fix problem in new code with non-compressed keys.

V02-014 MCN0001 Maria del C. Nasr 02-Jun-1981
Add the routine to split prologue 3 data buckets.

V02-013 REFORMAT Ron Schaefer 23-Jul-1980 14:10
Reformat the source

V02-012 CDS0000 Christian Saether, 01-Jan-1980 15:00
FIX PROBLEM WHEN SPLITTING BECAUSE OUT OF ID'S.

REVISION HISTORY:

Wendy Koenig, 18-SEP-78 16:53
X0002 - FIX BUG IN BACKING UP PAST NEW RECORD

Wendy Koenig, 19-SEP-78 10:52
X0003 - DO SPLIT AT POINT OF INSERT IF ASCENDING ORDER DETECTED

Wendy Koenig, 12-OCT-78 13:21
X0004 - CHANGES FOR UPDATE

Wendy Koenig, 18-OCT-78 14:03
X0005 - IF WE PASS BY POS_INSERT WHILE SKIPPING OVER DUPS, NOTE IT

Wendy Koenig, 18-OCT-78 14:37
X0006 - FIX SOME PROBLEMS W/ 4-BKT SPLIT (\$UPDATE ONLY)

Wendy Koenig, 24-OCT-78 14:03
X0007 - MAKE CHANGES CAUSED BY SHARING CONVENTIONS

Wendy Koenig, 7-NOV-78 8:58


```

229 0229 1 | X0008 - FIX EMPTY_BKT BUG, NOT BEING SET WHEN SHOULD BE
230 0230 1 |
231 0231 1 | Wendy Koenig, 22-JAN-79 17:03
232 0232 1 | X0009 - IF LOA TRIES TO FORCE US TO SPLIT ALL DUPS, SPLIT AT POS_INS
233 0233 1 |
234 0234 1 | Wendy Koenig, 24-JAN-79 9:51
235 0235 1 | X0010 - CONDITION HOLDS EVEN IF LOA NOT SET
236 0236 1 |
237 0237 1 | Wendy Koenig, 29-JAN-79 15:58
238 0238 1 | X0011 - FIX PROBLEM W/ DUPLICATE ENTRIES IN INDEX
239 0239 1 |
240 0240 1 | *****
241 0241 1 |
242 0242 1 | LIBRARY 'RMSLIB:RMS';
243 0243 1 |
244 0244 1 | REQUIRE 'RMSSRC:RMSIDXDEF';
245 0309 1 |
246 0310 1 | ! define default psects for code
247 0311 1 |
248 0312 1 | PSECT
249 0313 1 |     CODE = RMSRMS3(PSECT_ATTR);
250 0314 1 |     PLIT = RMSRMS3(PSECT_ATTR);
251 0315 1 |
252 0316 1 | ! Linkages
253 0317 1 |
254 0318 1 | LINKAGE
255 0319 1 |     L_COMPARE_KEY,
256 0320 1 |     L_PRESERVE1,
257 0321 1 |     L_RABREG_4567,
258 0322 1 |     L_RABREG_67,
259 0323 1 |     L_REC_OVRD,
260 0324 1 |
261 0325 1 | ! Local linkages
262 0326 1 |
263 0327 1 |     RL$BUILD_KEY          = JSB () :
264 0328 1 |                             GLOBAL (R_IDX_DFN) PRESERVE(1,2,3,4,5),
265 0329 1 |     RL$MOVE_KEY          = JSB (REGISTER = 0, REGISTER = 6) :
266 0330 1 |                             GLOBAL (R_RAB, R_IRAB, R_IFAB, R_IDX_DFN, R_BKT_ADDR);
267 0331 1 |
268 0332 1 | ! Forward Routine
269 0333 1 |
270 0334 1 |
271 0335 1 | FORWARD ROUTINE
272 0336 1 |     RMSBUILD_KEY          : RL$BUILD_KEY NOVALUE,
273 0337 1 |     RMSMOVE_KEY           : RL$MOVE_KEY NOVALUE;
274 0338 1 |
275 0339 1 | ! External Routines
276 0340 1 |
277 0341 1 |
278 0342 1 | EXTERNAL ROUTINE
279 0343 1 |     RMSMOVE               : RL$PRESERVE1,
280 0344 1 |     RMSRECORD_VBN         : RL$PRESERVE1,
281 0345 1 |     RMSRECORD_KEY         : RL$PRESERVE1,
282 0346 1 |     RMSREC_OVRD           : RL$REC_OVRD,
283 0347 1 |     RMSVBN_SIZE           : RL$PRESERVE1,
284 0348 1 |     RMSCOMPARE_KEY        : RL$COMPARE_KEY,
285 0349 1 |     RMSCOMPARE_REC        : RL$RABREG_67,
```

RM3SPLUDR
V04-000

M 15
16-Sep-1984 02:03:28
14-Sep-1984 13:01:40

VAX-11 Bliss-32 V4.0-742
[RMS.SRC]RM3SPLUDR.B32;1

Page 6
(1)

: 286 0350 1 RM\$GETNEXT_REC : RL\$RABREG_67;


```
288 0351 1 ++
289 0352 1
290 0353 1
291 0354 1
292 0355 1
293 0356 1
294 0357 1
295 0358 1
296 0359 1
297 0360 1
298 0361 1
299 0362 1
300 0363 1
301 0364 1
302 0365 1
303 0366 1
304 0367 1
305 0368 1
306 0369 1
307 0370 1
308 0371 1
309 0372 1
310 0373 1
311 0374 1
312 0375 1
313 0376 1
314 0377 1
315 0378 1
316 0379 1
317 0380 1
318 0381 1
319 0382 1
320 0383 1
321 0384 1
322 0385 1
323 0386 1
324 0387 1
325 0388 1
326 0389 1
327 0390 1
328 0391 1
329 0392 1
330 0393 1
331 0394 1
332 0395 1
333 0396 1
334 0397 1
335 0398 1
336 0399 1
337 0400 1
338 0401 1
339 0402 1
340 0403 1
341 0404 1
342 0405 1
343 0406 1
344 0407 1
```

ALGORITHM FOR A TWO-BUCKET 50/50 SPLIT

GIVEN: that the record will not fit in the bucket.
i.e., we must split the bucket in some form.

INPUTS: the bucket, the record size and the position
to insert the record in the bucket

GOALS: to make the split as efficient as possible:

- 1) to create the fewest number of new buckets possible
- 2) to use the space in the available buckets efficiently --
i.e., the bucket with the most available space should contain
the most data after the split.

ALGORITHM IN A NUTSHELL:

- 1) A two-bucket split will occur IF AND ONLY IF there is a point in
the bucket at which all records to the left of the point and
necessary rrv's fit in a single bucket and all records to the right
of the point fit in a single bucket. This point must be on a
record boundary and must not be in the middle of a chain of
duplicates.
- 2) Given that such a point exists, the most optimal point for a
2-bucket split is the point at which the actual data records
are divided evenly between the available space in the original
bucket and the available space in the new (previously empty) bucket.

In theory, therefore, the idea is to find a point in the bucket such that
the point is on a boundary between duplicate records and that

- 1) records in the left hand side / space in the left hand bucket
=
- 2) records in the right hand side / space in the right hand bucket.

In practice, the idea is to minimize the absolute difference between
ratio 1) and ratio 2). Just to make it clearer, "records in the left
hand side" means the total size of the data records left of this point
(not including rrv's of any kind) and "space in the left hand bucket"
means the bucketsize of the data bucket minus the total size of existing
rrv's and the total size of rrv's which would have to be generated.

IMPLEMENTATION:

This algorithm needs two scans of the bucket. The first scan is very
quick and determines the total size of the existing rrv's. It also
counts the number of rrv's that would have to be generated in a worst
case situation (i.e., all records would be moved out). Thus, as the
second scan proceeds, all information needed to calculate the above
ratios EXACTLY is available.

In order for there to be a 2-bucket split, there must be a point
in the bucket such that the right hand side fits in a single bucket.
Scanning from the left (beginning) of the bucket, we can find the
first point at which the right hand side will fit. Since as we
continue scanning to the right we are decreasing the right hand side,
the righthand side will continue to fit as we scan rightward.

If at this point, the left hand side will not fit, we can not possibly

: 345 0408 1 :
: 346 0409 1 :
: 347 0410 1 :
: 348 0411 1 :
: 349 0412 1 :
: 350 0413 1 :
: 351 0414 1 :
: 352 0415 1 :
: 353 0416 1 :
: 354 0417 1 :
: 355 0418 1 :
: 356 0419 1 :
: 357 0420 1 :
: 358 0421 1 :
: 359 0422 1 :
: 360 0423 1 :
: 361 0424 1 :
: 362 0425 1 :
: 363 0426 1 :
: 364 0427 1 :
: 365 0428 1 :
: 366 0429 1 :
: 367 0430 1 :
: 368 0431 1 :
: 369 0432 1 :
: 370 0433 1 :
: 371 0434 1 :
: 372 0435 1 :
: 373 0436 1 :
: 374 0437 1 :
: 375 0438 1 :
: 376 0439 1 :
: 377 0440 1 :
: 378 0441 1 :
: 379 0442 1 :
: 380 0443 1 :
: 381 0444 1 :
: 382 0445 1 :
: 383 0446 1 :
: 384 0447 1 :
: 385 0448 1 :
: 386 0449 1 :
: 387 0450 1 :
: 388 0451 1 :
: 389 0452 1 :
: 390 0453 1 :
: 391 0454 1 :
: 392 0455 1 :
: 393 0456 1 :
: 394 0457 1 :
: 395 0458 1 :
: 396 0459 1 :
: 397 0460 1 :
: 398 0461 1 :
: 399 0462 1 :
: 400 0463 1 :
: 401 0464 1 :

have a 2-bucket split, since continuing our scan would only make the left hand side larger (or it may stay the same size). Once we have found a point at which we can do a 2-bucket split we can always return to it, if in our search for a more optimal split point we leave the range in which the left hand side will fit. This can occur if the records in the bucket are of minimal size, that is to say that the records are the same size as rrv's and therefore no additional space for data is gained by scanning to the right.

At this point (the first point at which the right hand side will fit), ratio 1 is less than ratio 2. As we proceed to the right, ratio 1 will increase and ratio 2 will decrease. This is due to the fact that the size of the right hand side (the numerator of ratio 2) decreases as we move rightward and the available space in the right bucket is a constant (the denominator of ratio 2). In ratio 1, both the numerator and denominator are increasing, but the numerator is increasing at a faster rate. As soon as we reach a point where ratio 1 is greater than or equal to ratio 2, we can stop the scan. Now we have a choice of split points available. We can use this point or the one immediately before it (if such a point exists). The decision is made by minimizing the absolute difference between the ratios and we have an optimal split point.

Things become complicated by the presence of duplicate records. When duplicate records occupy more than one bucket, the subsequent buckets are termed continuation buckets. In prologue version 1 and 2 files, there is a pointer from the index to the first bucket only, and the continuation buckets are found only from the horizontal links in the buckets. At one point, it was thought that disaster would ensue if the continuation buckets ever had a record with a key value other than that of the duplicates. Normally, this will not happen because the key value of the index pointer to the first bucket will be the same as that of the duplicate records in the chain and a record with a higher key value will follow the next index pointer down when positioning for insert. This will place it in the next bucket beyond the chain of continuation buckets. However, a bucket in which the record with the highest value has been deleted that subsequently receives a series of duplicates creating a continuation chain will generate a situation where a record with a key value between that of the duplicate chain and the original high key value of the bucket will be inserted at the end of the duplicate chain. A far more common situation is created by RMS-11 (at least thru v1.5) when loading a file in ascending primary key sequence will pack the buckets 100% (or the load factor) full, including records of non-dupe key values at the end of continuation buckets.

At any rate, the fact that the situation exists notwithstanding, much of the code that follows is there to keep duplicates together when splitting, and to put only records with duplicate key values in continuation buckets. It appears to be a good thing to do from an overall space efficiency standpoint over a period of time, but the code could probably be considerably simplified if it wasn't necessary. With all that in mind, the split situation with all possible record 'partitions' within the bucket prior to splitting is as follows:

402 0465 1
403 0466 1
404 0467 1
405 0468 1
406 0469 1
407 0470 1
408 0471 1
409 0472 1
410 0473 1
411 0474 1
412 0475 1
413 0476 1
414 0477 1
415 0478 1
416 0479 1
417 0480 1
418 0481 1
419 0482 1
420 0483 1
421 0484 1
422 0485 1
423 0486 1
424 0487 1
425 0488 1
426 0489 1
427 0490 1
428 0491 1
429 0492 1
430 0493 1
431 0494 1
432 0495 1
433 0496 1
434 0497 1
435 0498 1
436 0499 1
437 0500 1
438 0501 1
439 0502 1
440 0503 1
441 0504 1
442 0505 1
443 0506 1
444 0507 1
445 0508 1
446 0509 1
447 0510 1
448 0511 1
449 0512 1
450 0513 1
451 0514 1
452 0515 1
453 0516 1
454 0517 1
455 0518 1
456 0519 1
457 0520 1
458 0521 1

! low set ! low dupes !! high dupes ! high set !

^
point of insert (new record)

From the point of view of the split code, an update operation in which the record is growing and causes a split is identical (almost) to a new record being inserted. The original record is removed from the bucket after determining that the updated record will cause a split and the updated record is more or less treated as a new record. One of the most important differences is that in an update situation, the 'new' record gets the id of the old record, rather than a new id. Another is that because duplicate records are always inserted at the end of a chain of duplicates, some split cases can only occur on an update operation.

In fact, the situation postulated above can happen only in an update situation, and may cause 3 new buckets to be generated on the split operation. This will occur when the updated record is in the middle of a group of duplicate records and grows to the extent that no other records will fit in the bucket with it anymore. Using 1 byte key values to make this easier to visualize, the bucket above prior to the update may look like this (the artificial partitioning of the bucket corresponds to the breakdown above):

! A B C ! D D D ! D ! D D ! E F G !

^\
this record gets updated

The record being updated changes size and grows such that it needs an entire bucket for itself. To keep all the duplicates together, the situation after the split looks like this:

----- ----- ----- -----
! A B C D D D ! -> ! D ! -> ! D D ! -> ! E F G !
----- ----- ----- -----

this is the these two are
original bucket continuation buckets

The original bucket probably had an index pointer with the value 'G' pointing to it (or some previous bucket if there was a previous index update failure). After the split, the key value for that pointer will be updated to have the key value 'D', and the key value that used to point to it (probably 'G'), will now point to the right hand bucket (with 'E', 'F', and 'G' in it). The continuation buckets never have an index pointer to them.

All other split situations are a variation of this one, with one or more of the 'partitions' not present, dependent on the key value and position of insert within the bucket of the record being inserted or updated. For example, if there are no duplicates, there are no 'low dupes' or 'high dupes'. Or if the position of insert is at the end of the bucket, there is no 'high set'.

Now that I've started on it, may as well try to document some other

interesting split situations. Note that a '2 bucket split' means that there are 2 buckets after the split, i.e., 1 new bucket is added. The situation described above is a 4 bucket split.

The most interesting split from an index updating point of view is the 3 bucket split where a record is being inserted in the middle of the bucket and doesn't fit in a bucket with either the low set or the high set. Again with 1 byte key values to illustrate:

G (this is supposed to represent an index
pointer to this bucket with key value 'g')

! A B C !! E F G !

/\

new record with key value 'D' inserted, but is so large
that it has to have bucket of its own.

After split (with new index pointers):

 C D G
 | | |
 v v v

! A B C ! -> ! D ! -> ! E F G !

The new pointer 'C' is the bucket pointer from the original index record 'G' with the new key value 'C'. The 'D' pointer is an entirely new record (i.e., key value 'D' and bucket pointer). The pointer 'G' is the key value from the original record 'G' with a new bucket pointer. The bucket pointer for the 'D' bucket comes from `irb$1_vbn_mid` and the bucket pointer for the 'G' bucket comes from `irb$1_vbn_right`. Remember that all of this stuff works correctly if the index update failed and we got to the bucket that's splitting by following the horizontal bucket links at the data level. For example, consider the following case where prior index corruption exists:

G (index update failed when right hand bucket split off
during a previous insert operation)

 |
 v

! A B C ! -> ! D !! F G !

/\

new record 'E' will be inserted here and cause split

After split:

 E G
 | |
 v v

$$\begin{array}{c} \text{V} \\ \hline ! \text{ A B C } ! \end{array} \rightarrow \begin{array}{c} \text{---} \\ \hline ! \text{ D E } ! \end{array} \rightarrow \begin{array}{c} \text{V} \\ \hline ! \text{ F G } ! \end{array}$$

After split:

Note that the index pointer 'I' was not moved to point to the new bucket. If it had been, the bucket containing 'D E F' would have been 'lost' by random access from the index. This condition is detected by setting `irb$l_vbn_left` to the `vbn` of the `rrv` only bucket. During the index update procedure, the pointer will be moved to point to the new bucket only if the existing down pointer points to the bucket that was split, i.e., `irb$l_vbn_left` (this


```

: 573      0636 1  is normally the case as index corruption is not normal). Note that
: 574      0637 1  an empty left hand bucket may also be present in a 3 bucket split
: 575      0638 1  situation.
: 576      0639 1
: 577      0640 1  Following is a list of the specific split cases handled in the
: 578      0641 1  code. They are basically variations of the above cases.
: 579      0642 1
: 580      0643 1  these are all the cases of 3 and 4 bucket splits that i can think of
: 581      0644 1  any or all of these cases can have the empty left-hand bucket
: 582      0645 1  -- this would occur if the first split point is at the beginning
: 583      0646 1  -- of the bucket and all data records got moved out
: 584      0647 1
: 585      0648 1  low dups exist -- no high dups
: 586      0649 1  low dups fit w/ rec
: 587      0650 1  3 bkt split low, low dups w/ rec, hi set -- rec goes w/ lo
: 588      0651 1  ( SPLIT TYPE 1 )
: 589      0652 1
: 590      0653 1  low dups don't fit w/ rec
: 591      0654 1  3 bkt split w/ rec in its own continuation bucket
: 592      0655 1  ( SPLIT TYPE 2, W/ DUPS SEEN )
: 593      0656 1
: 594      0657 1  hi dups exist -- no low dups
: 595      0658 1  hi dups fit w/ rec
: 596      0659 1  3 bkt split low, hi dups w/ rec, hi set
: 597      0660 1  ( SPLIT TYPE 1 )
: 598      0661 1
: 599      0662 1  hi dups don't fit w/ rec
: 600      0663 1  if no more hi, 3 bkt split low, rec, hi = hi dups is a cont. bkt
: 601      0664 1  ( SPLIT TYPE 2 )
: 602      0665 1
: 603      0666 1  if there is more hi, 4 bkt split low, rec, hi dups, hi
: 604      0667 1  ( SPLIT TYPE 2B )
: 605      0668 1
: 606      0669 1  no dups at all
: 607      0670 1  record goes in its own bucket, 3 bkt split
: 608      0671 1  (SPLIT TYPE 3 )
: 609      0672 1
: 610      0673 1  low dups and hi dups
: 611      0674 1  all dups fit together
: 612      0675 1  3 bkt split w/ dups in middle bkt
: 613      0676 1  ( SPLIT TYPE 1 )
: 614      0677 1
: 615      0678 1  no dups fit w/ record
: 616      0679 1  if no more hi, 3 bkt split low, rec = cont. bkt, hi = hi dups = cont. bkt
: 617      0680 1
: 618      0681 1  ( SPLIT TYPE 2B )
: 619      0682 1
: 620      0683 1  if there is more hi, 4 bkt split low, rec = cont. bkt, hi dups, hi
: 621      0684 1  ( SPLIT TYPE 2B )
: 622      0685 1
: 623      0686 1
: 624      0687 1  hi dups fit w/ record
: 625      0688 1  if no more hi, 2 bkt split low, rec w/ hi dups = cont. bkt
: 626      0689 1  ( this is a 2 bkt split case that the previous alg. wouldn't handle)
: 627      0690 1  ( SPLIT TYPE 4B )
: 628      0691 1
: 629      0692 1  if there is more hi, 3 bkt split low, rec w/ hi dups = cont. bkt, hi
```


:	630	0693	1	:	(SPLIT TYPE 4)
:	631	0694	1	:	
:	632	0695	1	:	low dups fit w/ record
:	633	0696	1	:	if lo and hi, 4 bkt split low, low dups w/ rec, hi dups, hi
:	634	0697	1	:	(SPLIT TYPE 5)
:	635	0698	1	:	
:	636	0699	1	:	if no lo and no hi, 3 bkt split (rrv's), low dups w/ rec, hi dups = cont.
:	637	0700	1	:	
:	638	0701	1	:	bkt
:	639	0702	1	:	(SPLIT TYPE 5, w/ empty original bkt and no high)
:	640	0703	1	:	
:	641	0704	1	:	if lo but no hi, 3 bkt split lo, low dups w/ rec, hi dups = cont. bkt
:	642	0705	1	:	(SPLIT TYPE 5, w/ no high)
:	643	0706	1	:	
:	644	0707	1	:	if hi but no low, 4 bkt split (rrv's), low dups w/ rec, hi dups, hi
:	645	0708	1	:	(SPLIT TYPE 5, w/ empty bkt)
:	646	0709	1	:	--

RMSMOVE_KEY

```

: 648 0710 1 %SBTTL 'RMSMOVE_KEY'
: 649 0711 1 ROUTINE RMSMOVE_KEY (ADDRESS, CUR_REC_ADDR) : RL$MOVE_KEY NOVALUE =
: 650 0712 1
: 651 0713 1 ++
: 652 0714 1
: 653 0715 1 FUNCTIONAL DESCRIPTION:
: 654 0716 1
: 655 0717 1 Routine to move the key from wherever it is desired into
: 656 0718 1 key buffer 2.
: 657 0719 1
: 658 0720 1 CALLING SEQUENCE:
: 659 0721 1 bsbw rm$move_key (address,cur_rec_addr)
: 660 0722 1
: 661 0723 1 INPUT PARAMETERS:
: 662 0724 1 address from which to get the key from
: 663 0725 1 the current value of rec_addr
: 664 0726 1
: 665 0727 1 IMPLICIT INPUTS:
: 666 0728 1 BKT_ADDR,
: 667 0729 1 RAB -- user's buffer address
: 668 0730 1 IRAB -- pos_ins, rec_w_lo, keybuf
: 669 0731 1 IFAB -- kbufsz, prologue version
: 670 0732 1 IDX_DFN -- for call to record_key, and compression flags
: 671 0733 1
: 672 0734 1 OUTPUT PARAMETERS:
: 673 0735 1 none
: 674 0736 1
: 675 0737 1 IMPLICIT OUTPUTS:
: 676 0738 1 key is moved into key buffer 2
: 677 0739 1
: 678 0740 1 ROUTINE VALUE:
: 679 0741 1 none
: 680 0742 1
: 681 0743 1 SIDE EFFECTS:
: 682 0744 1 key is moved into key buffer 2
: 683 0745 1 AP is clobbered
: 684 0746 1
: 685 0747 1 --
: 686 0748 1
: 687 0749 2 BEGIN
: 688 0750 2
: 689 0751 2 BUILTIN
: 690 0752 2 AP;
: 691 0753 2
: 692 0754 2 GLOBAL REGISTER
: 693 0755 2 R_BDB,
: 694 0756 2 R_IMPURE,
: 695 0757 2 R_REC_ADDR_STR;
: 696 0758 2
: 697 0759 2 EXTERNAL REGISTER
: 698 0760 2 R_IFAB_STR,
: 699 0761 2 R_RAB,
: 700 0762 2 R_IRAB_STR,
: 701 0763 2 R_IDX_DFN_STR,
: 702 0764 2 R_BKT_ADDR_STR;
: 703 0765 2
: 704 0766 2 IF .CUR_REC_ADDR - .BKT_ADDR EQLU .IRAB[IRBSW_POS_INS]
```



```

: 705      0767 2      AND
: 706      0768 2      .IRAB[IRB$V_REC_W_LO]
: 707      0769 2      . THEN
: 708      0770 2      BEGIN
: 709      0771 2      AP = 3; ! no overhead, not compressed
: 710      0772 2      REC_ADDR = .IRAB[IRB$L_RBF];
: 711      0773 2      END
: 712      0774 2      ELSE
: 713      0775 2      BEGIN
: 714      0776 2      AP = 0;
: 715      0777 2      REC_ADDR = .ADDRESS;
: 716      0778 2
: 717      0779 2      ! In prologue 3 version files, if the key is compressed, it must be
: 718      0780 2      ! rebuilt. Make sure that the last non-compressed pointer, is before
: 719      0781 2      ! the record we are looking at.
: 720      0782 2
: 721      0783 2
: 722      0784 2      IF .IDX_DFN[IDX$V_KEY_COMP]
: 723      0785 2      THEN
: 724      0786 4      BEGIN
: 725      0787 4
: 726      0788 4      IF .IRAB[IRB$L_LST_NCMP] GTRU .ADDRESS
: 727      0789 4      THEN
: 728      0790 5      BEGIN
: 729      0791 5
: 730      0792 5      IF .(.ADDRESS + RM$REC_OVHD() + 1)<0,8> EQLU 0
: 731      0793 5      THEN
: 732      0794 5      IRAB[IRB$L_LST_NCMP] = .ADDRESS
: 733      0795 5      ELSE
: 734      0796 5      IRAB[IRB$L_LST_NCMP] = .BKT_ADDR + BKT$C_OVERHDSZ;
: 735      0797 4      END;
: 736      0798 3      END;
: 737      0799 2      END;
: 738      0800 2
: 739      0801 2      ! We are storing in key buffer 2 the possible key to be inserted at the
: 740      0802 2      ! index level.
: 741      0803 2
: 742      0804 2
: 743      0805 2      RM$RECORD_KEY ( KEYBUF_ADDR(2) );
: 744      0806 2
: 745      0807 2      RETURN;
: 746      0808 2
: 747      0809 1      END;
```

```

.TITLE RM3SPLUDR
.IDENT \V04-000\

.EXTRN RM$MOVE, RM$RECORD_VBN
.EXTRN RM$RECORD_KEY, RM$REC_OVHD
.EXTRN RM$VBN_SIZE, RM$COMPARE_KEY
.EXTRN RM$COMPARE_REC, RM$GETNEXT_REC

.PSECT RM$RMS3,NOWRT, GBL, PIC,2
```

```
0850 8F BB 00000 RM$MOVE_KEY:
      PUSH  #^M<R4,R6,R11>
```

: 0711

RM3SPLUDR
V04-000

RMSMOVE_KEY

J 16
16-Sep-1984 02:03:28
14-Sep-1984 13:01:40

VAX-11 Bliss-32 V4.0-742
[RMS.SRC]RM3SPLUDR.B32;1

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56	48	A9	5B	50	D0	00004	MOVL	R0, R11	:		
			56	55	C2	00007	SUBL2	BKf_ADDR, R6	:	0766	
			10	00	ED	0000A	CMPZV	#0, #16, 72(IRAB), R6	:		
				0E	12	00010	BNEQ	1\$:		
		09	44	A9	03	E1	00012	BBC	#3, 68(IRAB), 1\$:	0768
			5C	03	D0	00017	MOVL	#3, AP	:	0771	
			56	58	A9	D0	0001A	MOVL	88(IRAB), REC_ADDR	:	0772
				26	11	0001E	BRB	3\$:	0766	
				5C	D4	00020	CLRL	AP	:	0776	
			56	5B	D0	00022	MOVL	ADDRESS, REC_ADDR	:	0777	
		1C	1C	A7	06	E1	00025	BBC	#6, 28(IDX_DFN), 3\$:	0784
			54	0098	C9	9E	0002A	MOVAB	152(IRAB), R4	:	0788
			5B		64	D1	0002F	CPL	(R4), ADDRESS	:	
					12	1B	00032	BLEQU	3\$:	
					0000G	30	00034	BSBW	RMSREC_OVHD	:	0792
				01	A04B	95	00037	TSTB	1(R0)[ADDRESS]	:	
					05	12	0003B	BNEQ	2\$:	
			64		5B	D0	0003D	MOVL	ADDRESS, (R4)	:	0794
					04	11	00040	BRB	3\$:	
			64	0E	A5	9E	00042	MOVAB	14(R5), (R4)	:	0796
			50	00B4	CA	3C	00046	MOVZWL	180(IFAB), R0	:	0805
				60	B940	9F	0004B	PUSHAB	@96(IRAB)[R0]	:	
					0000G	30	0004F	BSBW	RMSRECORD_KEY	:	
			5E		04	C0	00052	ADDL2	#4, SP	:	
				0850	8F	BA	00055	POPR	#^M<R4,R6,R11>	:	0809
					05	00059	RSB		:		

; Routine Size: 90 bytes, Routine Base: RMSRMS3 + 0000

; 748 0810 1

RMSBUILD_KEY

```
750 0811 1 %SBTTL 'RMSBUILD_KEY'
751 0812 1 ROUTINE RMSBUILD_KEY (ADDRESS, KEYBUF) : RL$BUILD_KEY NOVALUE =
752 0813 1
753 0814 1 !++
754 0815 1
755 0816 1 FUNCTIONAL DESCRIPTION:
756 0817 1
757 0818 1 This routine builds a compressed key from the record into the
758 0819 1 given buffer, knowing that the front characters are valid from
759 0820 1 the previous expansion.
760 0821 1
761 0822 1 CALLING SEQUENCE:
762 0823 1 bsbw rm$build_key (address,keybuf)
763 0824 1
764 0825 1 INPUT PARAMETERS:
765 0826 1 - address in bucket which points to key compression overhead
766 0827 1 - key output buffer
767 0828 1
768 0829 1 IMPLICIT INPUTS:
769 0830 1 IDX_DFN - index definition for key size
770 0831 1
771 0832 1 OUTPUT PARAMETERS:
772 0833 1 none
773 0834 1
774 0835 1 IMPLICIT OUTPUTS:
775 0836 1 key is moved into appropriate key buffer
776 0837 1
777 0838 1 ROUTINE VALUE:
778 0839 1 none
779 0840 1
780 0841 1 SIDE EFFECTS:
781 0842 1 key is moved into appropriate key buffer
782 0843 1
783 0844 1 --
784 0845 1
785 0846 2 BEGIN
786 0847 2
787 0848 2 EXTERNAL REGISTER
788 0849 2 R_IDX_DFN_STR;
789 0850 2
790 0851 2 LOCAL
791 0852 2 TRUN_CHAR,
792 0853 2 LENGTH;
793 0854 2
794 0855 2 BIND
795 0856 2 REC_KEY = ADDRESS : REF BBLOCK;
796 0857 2
797 0858 2 MACRO
798 0859 2 KEY_LEN = 0,0,8,0 %;
799 0860 2 CMP_CNT = 1,0,8,0 %;
800 0861 2
801 0862 2 KEYBUF = .KEYBUF + .REC_KEY[ CMP_CNT ]; ! skip characters already moved
802 0863 2 TRUN_CHAR = .REC_KEY + .REC_KEY[ KEY_LEN ] + 1;
803 0864 2 LENGTH = .IDX_DFN[ IDX$B_KEYSZ ] - .REC_KEY[ CMP_CNT ];
804 0865 2 CH$COPY ( .REC_KEY[ KEY_LEN ], .REC_KEY + 2,
805 0866 2 ..TRUN_CHAR,
806 0867 2 .LENGTH, .KEYBUF );
```


RM3SPLUDR
V04-000

RMSBUILD_KEY

L 16
16-Sep-1984 02:03:28
14-Sep-1984 13:01:40

VAX-11 Bliss-32 V4.0-742
[RMS.SRC]RM3SPLUDR.B32;1

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: 807
: 808
: 809
: 810
0868 2
0869 2 RETURN;
0870 2
0871 1 END;

				3E	BB	00000	RMSBUILD_KEY:			
							PUSHR	#^M<R1,R2,R3,R4,R5>		: 0812
							MOVL	REC KEY, R1		: 0862
							MOVZBL	1(RT), R0		:
							ADDL2	R0, KEYBUF		:
							MOVZBL	(R1), R0		: 0863
							MOVAB	1(R1)[R0], TRUN_CHAR		:
							MOVZBL	32(IDX_DFN), LENGTH		: 0864
							MOVZBL	1(R1), R0		:
							SUBL2	R0, LENGTH		:
							MOVZBL	(R1), R0		: 0865
							MOVCS	R0, 2(R1), (TRUN_CHAR), LENGTH, @KEYBUF		: 0867
							POPR	#^M<R1,R2,R3,R4,R5>		: 0871
							RSB			:

51 18 AE D0 00002
50 01 A1 9A 00006
1C AE 50 C0 0000A
50 61 9A 0000E
53 01 A140 9E 00011
52 20 A7 9A 00016
50 01 A1 9A 0001A
52 50 C2 0001E
50 61 9A 00021
52 50 2C 00024
50 A1 1C BE 0002A
3E BA 0002C
05 0002E

; Routine Size: 47 bytes, Routine Base: RMSRMS3 + 005A

: 811 0872 1


```
RM$SPLIT_UDR
: 813 0873 1 %SBTTL 'RM$SPLIT_UDR'
: 814 0874 1 GLOBAL ROUTINE RM$SPLIT_UDR : RL$RABREG_4567 NOVALUE =
: 815 0875 1
: 816 0876 1 ++
: 817 0877 1
: 818 0878 1 FUNCTIONAL DESCRIPTION:
: 819 0879 1
: 820 0880 1
: 821 0881 1 CALLING SEQUENCE:
: 822 0882 1     BSBW RM$SPLIT_UDR()
: 823 0883 1
: 824 0884 1 INPUT PARAMETERS:
: 825 0885 1     none
: 826 0886 1
: 827 0887 1 IMPLICIT INPUTS:
: 828 0888 1     BDB pointer, BUFFER pointer, REC_ADDR = point of insert, IDX_DFN
: 829 0889 1     in IRAB -- curbdb, associated w/ bdb and bkt_addr
: 830 0890 1     pos_ins corresponding to rec_addr
: 831 0891 1     in RAB -- rsz of record
: 832 0892 1     in IFAB -- rfm
: 833 0893 1     BKT$B_NXTRECID = 0 in original bucket signals that this is
: 834 0894 1     a split due to a lack of id's in the bucket
: 835 0895 1
: 836 0896 1 OUTPUT PARAMETERS:
: 837 0897 1     none
: 838 0898 1
: 839 0899 1 IMPLICIT OUTPUTS:
: 840 0900 1     in IRAB --
: 841 0901 1     if 2 bkt split --
: 842 0902 1         IRB$W_SPLIT, offset to split point
: 843 0903 1         IRB$V_REC_W_LO -- set if split point is pos_insert and
: 844 0904 1         record goes w/ lo set
: 845 0905 1         new high key for original bucket in keybuffer 2
: 846 0906 1         number of new buckets = 1
: 847 0907 1         if original bucket was all rrv's, set IRB$V_EMPTY_BKT flag
: 848 0908 1         if new bucket is a continuation bkt., set IRB$V_CONT_BKT flag
: 849 0909 1     if 3 bkt split --
: 850 0910 1         same as above w/ these changes:
: 851 0911 1         IRB$W_SPLIT_1, offset to second split point
: 852 0912 1         number of new buckets = 2
: 853 0913 1         if right bucket is a continuation bkt, set IRB$V_CONT_R flag
: 854 0914 1     if 4 bkt split --
: 855 0915 1         same as above w/ these changes:
: 856 0916 1         IRB$W_SPLIT_2, offset to third split point
: 857 0917 1         number of new buckets = 3
: 858 0918 1
: 859 0919 1 ROUTINE VALUE:
: 860 0920 1     rmssuc
: 861 0921 1
: 862 0922 1 SIDE EFFECTS:
: 863 0923 1     AP is clobbered
: 864 0924 1
: 865 0925 1 --
: 866 0926 1
: 867 0927 2 BEGIN
: 868 0928 2
: 869 0929 2 EXTERNAL REGISTER
```



```

: 870      0930      2      COMMON_RAB_STR,
: 871      0931      2      R_REC_ADDR_STR,
: 872      0932      2      R_IDX_DFN_STR,
: 873      0933      2      COMMON_IO_STR;
: 874      0934      2
: 875      0935      2      LOCAL
: 876      0936      2      SAVE_REC_W_LO,
: 877      0937      2      NUM_RRVs,
: 878      0938      2      POS_INSERT,
: 879      0939      2      EOB,
: 880      0940      2      RRV,
: 881      0941      2      RHS,
: 882      0942      2      LHS,
: 883      0943      2      LAST      : REF BBLOCK,
: 884      0944      2      LAST_DIFF,
: 885      0945      2      BKT_SIZE,
: 886      0946      2      REC_SIZE,
: 887      0947      2      DIFFERENCE;
: 888      0948      2
: 889      0949      2      MACRO
: 890      0950      2      NEED_RRV = NUM_RRVs<0,16> %,
: 891      0951      2      NOT_NEED_RRV = -NUM_RRVs<16,16> %;
: 892      0952      2
: 893      0953      2      LABEL
: 894      0954      2      DO_IT,
: 895      0955      2      HALF,
: 896      0956      2      NEXT;
: 897      0957      2
: 898      0958      2      DO_IT :
: 899      0959      2
: 900      0960      2      BEGIN
: 901      0961      2
: 902      0962      2      ! define a block so that we can have some common checks before returning
: 903      0963      2      ! successfully
: 904      0964      2      !
: 905      0965      2      HALF :
: 906      0966      2
: 907      0967      2      BEGIN
: 908      0968      2
: 909      0969      2      !+
: 910      0970      2      ! define a block so that we can simulate a go-to (naughty, naughty)
: 911      0971      2      ! if we have decided that we are positioning at the end of the bucket
: 912      0972      2      ! & we're in somewhat of an ascending order, where the last record
: 913      0973      2      ! inserted is a duplicate of the new record, skip over the 50-50 code
: 914      0974      2      ! and go to the code to take duplicates into account
: 915      0975      2      !
: 916      0976      2      ! scan 1 -- calculate
: 917      0977      2      ! size of existing rrv's and total number of rrv's needed to move the whole
: 918      0978      2      ! bucket out ( worst case) as a side effect, adjust eob ptr to pt to the
: 919      0979      2      ! rrv's instead of freespace assume not empty bucket until showed otherwise
: 920      0980      2      !-
: 921      0981      2
: 922      0982      2      IRAB[IRBSV_EMPTY_BKT] = 0;
: 923      0983      2
: 924      0984      2      ! new rec is tried 1st w/ hi set, then w/ lo set
: 925      0985      2      !
: 926      0986      2      IRAB[IRBSV_REC_W_LO] = 0;
```



```

: 927      0987 4      IRAB[IRB$V_NEW_BKTS] = 1;      ! assume 2-bkt split until showed otherwise
: 928      0988 4      NUM_RRVs = 0;                ! this zeroes NEED_RRV and NOT_NEED_RRV
: 929      0989 4      POS_INSERT = .REC_ADDR;
: 930      0990 4      REC_ADDR = .BKT_ADDR + BKT$C_OVERHDSZ;
: 931      0991 4      EOB = .BKT_ADDR + .BKT_ADDR[BKT$W_FREESPACE];
: 932      0992 4      LAST = 0;
: 933      0993 4
: 934      0994 4      DO
: 935      0995 5          BEGIN
: 936      0996 5
: 937      0997 5          BUILTIN
: 938      0998 5              AP;
: 939      0999 5
: 940      1000 5          IF .REC_ADDR[IRC$V_RRV]
: 941      1001 5              THEN
: 942      1002 5              EXITLOOP;
: 943      1003 5
: 944      1004 5          AP = 3;
: 945      1005 5
: 946      1006 5          IF .BDB[BDB$L_VBN] EQLU RM$RECORD_VBN()
: 947      1007 5              THEN
: 948      1008 5              NEED_RRV = .NEED_RRV + 1
: 949      1009 5
: 950      1010 5              ! the records not requiring rrv's are counted also because in the
: 951      1011 5              ! case where we're splitting due to lack of id's, the lhs side will
: 952      1012 5              ! fit with the new record if any of the record being moved to the
: 953      1013 5              ! new bucket doesn't require an rrv. this will be checked when we
: 954      1014 5              ! check to see if the lhs will fit after the first point that the
: 955      1015 5              ! rhs fits.
: 956      1016 5
: 957      1017 5          ELSE
: 958      1018 5              NOT_NEED_RRV = .NOT_NEED_RRV + 1;
: 959      1019 5
: 960      1020 5          LAST = .REC_ADDR;
: 961      1021 5          RM$GETNEXT_REC()
: 962      1022 5          END
: 963      1023 4      UNTIL .REC_ADDR GEQU .EOB;
: 964      1024 4
: 965      1025 4      ! set split_2 and split_1 to be eob, so if there's less than 3 new buckets
: 966      1026 4      ! bkt_spl can use the value w/o having to recalculate it also set up the
: 967      1027 4      ! bucket size and the record size
: 968      1028 4
: 969      1029 4      IRAB[IRB$W_SPLIT_1] = IRAB[IRB$W_SPLIT_2] = .REC_ADDR - .BKT_ADDR;
: 970      1030 4      BKTSIZE = .IDX_DFNC[IDX$B_DTBKTSZ]*512 - BKT$C_OVERHDSZ - 1;
: 971      1031 4
: 972      1032 4      REC_SIZE = .RAB[RAB$W_RSZ] + IRC$C_FIXOVHDSZ;
: 973      1033 4
: 974      1034 4      IF .IFAB[IFB$B_RFMORG] NEQ FAB$C_FIX
: 975      1035 4          THEN
: 976      1036 4          REC_SIZE = .REC_SIZE + 2;
: 977      1037 4
: 978      1038 4      ! if this is an update, may have to count in an rrv for the existing record
: 979      1039 4
: 980      1040 4      IF .IRAB[IRB$V_UPDATE]
: 981      1041 4          THEN
: 982      1042 5          BEGIN
: 983      1043 5
```



```
: 984      1044  5      IF .BDB[BDB$L_VBN] EQLU .IRAB[IRB$L_PUTUP_VBN]
: 985      1045  5      THEN
: 986      1046  5          NEED_RRV = .NEED_RRV + 1
: 987      1047  5
: 988      1048  4      END;
: 989      1049  4
: 990      1050  4      RRV = .EOB - .REC_ADDR;          ! size of existing rrv's
: 991      1051  4      EOB = .REC_ADDR;              ! adjust eob
: 992      1052  4
: 993      1053  4      ! special case it, if the bucket was all rrv's
: 994      1054  4      !
: 995      1055  4
: 996      1056  4      IF .REC_ADDR EQLU .BKT_ADDR + BKT$C_OVERHDSZ
: 997      1057  4      THEN
: 998      1058  5          BEGIN
: 999      1059  5              ! bkt is all rrv's yet the record wouldn't fit so we need to
1000      1060  5              ! allocate another bkt ( 2 bkt split) yet special case it so as not
1001      1061  5              ! to make another idx entry only to update the existing one by
1002      1062  5              ! setting empty bucket flag
1003      1063  5              !
1004      1064  5              IRAB[IRB$W_SPLIT] = .REC_ADDR - .BKT_ADDR;
1005      1065  5              LEAVE DO_IT
1006      1066  5
1007      1067  5          END;
1008      1068  4              ! { of special case an all-rrv bucket }
1009      1069  4
1010      1070  4      +
1011      1071  4      ! special case -- if we can detect a possible ascending order to these
1012      1072  4      ! records it probably will be better to do a straight point of insert split
1013      1073  4      ! this would put the new record in a bucket all by itself.
1014      1074  4      ! do this kind of split if and only if all the following conditions are met:
1015      1075  4      ! 1) the record is being inserted at the end of bucket
1016      1076  4      ! 2) the last record physically in the bkt is the last record to have
1017      1077  4      !     been inserted
1018      1078  4      ! 3) the last record and the new record do not have duplicate key values
1019      1079  4
1020      1080  4      ! note that if they are duplicates, we can still make an optimization by
1021      1081  4      ! skipping the 50-50 split code
1022      1082  4
1023      1083  4      ! note that last cannot be zero, since if it were we
1024      1084  4      ! would have an all rrv bkt
1025      1085  4      !
1026      1086  4
1027      1087  4      IF .POS_INSERT EQLU .REC_ADDR
1028      1088  4      AND
1029      1089  5          (((.LAST[IRC$B_ID] + 1) AND %X'FF') EQLU .BKT_ADDR[BKT$B_NXTRECID])
1030      1090  4      THEN
1031      1091  5          BEGIN
1032      1092  5              REC_ADDR = .LAST;
1033      1093  5
1034      1094  5              IF RM$COMPARE_REC(KEYBUF_ADDR(3), .IDX_DFN[IDX$B_KEYSZ], 0)
1035      1095  5              THEN
1036      1096  6                  BEGIN
1037      1097  6
1038      1098  6                  ! since we have detected a possible ascending order in the input
1039      1099  6                  ! let's try to optimize a little and split at the point of insert
1040      1100  6                  ! send the record by itself into the new bucket have to set up the
```



```
1041      1101  6      ! key value and the split point and that's it
1042      1102  6      !
1043      1103  6      RM$MOVE KEY(.REC_ADDR, .REC_ADDR);
1044      1104  6      IRAB[IRB$W_SPLIT] = .IRAB[IRB$W_POS_INS];
1045      1105  6      LEAVE DO_IT;
1046      1106  6
1047      1107  6      END
1048      1108  5      ELSE
1049      1109  5          LEAVE HALF;
1050      1110  5
1051      1111  5      ! { end of trying to special case insertion of records in ascending
1052      1112  5      ! order }
1053      1113  5      !
1054      1114  4      END;
1055      1115  4
1056      1116  4      REC_ADDR = .BKT_ADDR + BKT$C_OVERHDSZ;
1057      1117  4      LAST_DIFF = 'X'7FFFFFFF;
1058      1118  4      LAST = 0;
1059      1119  4      SAVE_REC_W_LO = 0;
1060      1120  4
1061      1121  4      ! start from the beginning of the bucket and scan rightward. first find the
1062      1122  4      ! 1st place the rhs will fit in 1 bkt then, as long as the lhs will fit in
1063      1123  4      ! a bkt, try to find an optimal point if there is no point where the rhs
1064      1124  4      ! and lhs will both fit. we can't do a 2-bkt split and this case will fall
1065      1125  4      ! out
1066      1126  4      !
1067      1127  4
1068      1128  4      WHILE 1
1069      1129  4      DO
1070      1130  5          BEGIN
1071      1131  5          RHS = .EOB - .REC_ADDR;
1072      1132  5
1073      1133  5          IF .REC_ADDR LEQU .POS_INSERT
1074      1134  5              AND
1075      1135  5              NOT .IRAB[IRB$V_REC_W_LO]
1076      1136  5          THEN
1077      1137  5              RHS = .RHS + .REC_SIZE;
1078      1138  5
1079      1139  5          ! the right hand side fits if there is enough room and there are id's
1080      1140  5          ! available. id's are always available in the new bucket in the update
1081      1141  5          ! situation, or if we're leaving at least 1 record behind in the old
1082      1142  5          ! bucket. note that nxtrecid is always zeroed if this is a split due to
1083      1143  5          ! lack of id's.
1084      1144  5          !
1085      1145  5          IF .RHS LSSU .BKTSIZE
1086      1146  5              AND
1087      1147  5              (.BKT_ADDR[BKT$B_NXTRECID] NEQ 0
1088      1148  6              OR
1089      1149  6              .IRAB[IRB$V_UPDATE]
1090      1150  6              OR
1091      1151  6              .REC_ADDR NEQA (.BKT_ADDR + BKT$C_OVERHDSZ)
1092      1152  7              OR
1093      1153  6              .IRAB[IRB$V_REC_W_LO])
1094      1154  6          THEN
1095      1155  5              BEGIN
1096      1156  6              LHS = .REC_ADDR - (.BKT_ADDR + BKT$C_OVERHDSZ);
1097      1157  6
```



```

: 1098      1158 6
: 1099      1159 6
: 1100      1160 6
: 1101      1161 6
: 1102      1162 6
: 1103      1163 6
: 1104      1164 6
: 1105      1165 6
: 1106      1166 6
: 1107      1167 6
: 1108      1168 6
: 1109      1169 6
: 1110      1170 6
: 1111      1171 6
: 1112      1172 6
: 1113      1173 6
: 1114      1174 6
: 1115      1175 6
: 1116      1176 6
: 1117      1177 6
: 1118      1178 6
: 1119      1179 6
: 1120      1180 7
: 1121      1181 7
: 1122      1182 7
: 1123      1183 7
: 1124      1184 7
: 1125      1185 7
: 1126      1186 7
: 1127      1187 6
: 1128      1188 7
: 1129      1189 7
: 1130      1190 7
: 1131      1191 7
: 1132      1192 7
: 1133      1193 7
: 1134      1194 7
: 1135      1195 7
: 1136      1196 7
: 1137      1197 7
: 1138      1198 7
: 1139      1199 7
: 1140      1200 7
: 1141      1201 7
: 1142      1202 7
: 1143      1203 7
: 1144      1204 8
: 1145      1205 8
: 1146      1206 8
: 1147      1207 8
: 1148      1208 8
: 1149      1209 8
: 1150      1210 8
: 1151      1211 8
: 1152      1212 8
: 1153      1213 8
: 1154      1214 8

IF .REC_ADDR GEQU .POS_INSERT
AND
.IRAB[IRB$V_REC_W_LO]
THEN
LHS = .LHS + .REC_SIZE;

! will lhs fit ? lhs doesn't fit if there is no space in the
! bucket, or if there won't be any id's available in the bucket.
! if not & if there is no previous point at which it fit, goto 3-bkt
! split code if there is a previous place where we could have had a
! 2-bkt split, use it

IF .LHS + .RRV + (7*.NEED_RRV) GTRU .BKTSIZE
! id's will be available in the original bucket if we aren't
! out of id's to begin with, this is an update, any record
! being moved out doesn't need an rrv, or the new record is
! going in the new bucket
OR
(.BKT_ADDR[BKT$B_NXTRECID] EQL 0
AND
NOT .IRAB[IRB$V_UPDATED]
AND
.NOT_NEED_RRV EQL 0
AND
.IRAB[IRB$V_REC_W_LO])
THEN
BEGIN
IF .LAST EQL 0
THEN
EXITLOOP;

REC_ADDR = .LAST;

IF NOT .SAVE_REC_W_LO
THEN
IRAB[IRB$V_REC_W_LO] = 0;

! 2 bkt split is possible rec_addr points to the most
! optimal place since we had to back up, reset last to point
! to the record immediately before the split point
BEGIN
LOCAL
TMP;

TMP = .REC_ADDR;
REC_ADDR = .BKT_ADDR + BKT$C_OVERHDSZ;
LAST = .REC_ADDR;

WHILE .REC_ADDR NEQU .TMP
DO
```



```
: 1155      1215  9      BEGIN
: 1156      1216  9      LAST = .REC_ADDR;
: 1157      1217  9      RMSGETNEXT_REC();
: 1158      1218  8      END;
: 1159      1219  8
: 1160      1220  7      END;
: 1161      1221  7      RMSMOVE KEY(.LAST, .REC_ADDR);
: 1162      1222  7      IRAB[IRB$W_SPLIT] = .REC_ADDR - .BKT_ADDR;
: 1163      1223  7
: 1164      1224  7      ! treat another exception case of the new record going off into
: 1165      1225  7      ! a cont. bkt all by itself
: 1166      1226  7
: 1167      1227  7
: 1168      1228  7      IF .IRAB[IRB$W_SPLIT] EQLU .IRAB[IRB$W_POS_INS]
: 1169      1229  7      THEN
: 1170      1230  7
: 1171      1231  7          IF .IRAB[IRB$W_SPLIT] EQLU .IRAB[IRB$W_SPLIT_1]
: 1172      1232  7          THEN
: 1173      1233  7
: 1174      1234  7              IF NOT .IRAB[IRB$V_REC_W_LO]
: 1175      1235  7              THEN
: 1176      1236  8                  BEGIN
: 1177      1237  8
: 1178      1238  8                      BUILTIN
: 1179      1239  8                      AP;
: 1180      1240  8
: 1181      1241  8                      AP = 3;
: 1182      1242  8
: 1183      1243  8                      IF NOT RMSCOMPARE KEY(KEYBUF_ADDR(2),
: 1184      1244  8                      KEYBUF_ADDR(3),
: 1185      1245  8                      .IDX_DFNC[IDX$B_KEYSZ])
: 1186      1246  8                      THEN
: 1187      1247  8                          IRAB[IRB$V_CONT_BKT] = 1;
: 1188      1248  8
: 1189      1249  7                      END;
: 1190      1250  7
: 1191      1251  7      LEAVE DO_IT
: 1192      1252  7
: 1193      1253  6      END;
: 1194      1254  6      ! { end of lhs doesn't fit anymore }
: 1195      1255  6      ! lhs fits also, calculate the magic ratio
: 1196      1256  6
: 1197      1257  8      DIFFERENCE = (.LHS*.BKTSIZE) - (.RHS*(.BKTSIZE - (7*.NEED_RRV) -
: 1198      1258  6      .RRV));
: 1199      1259  6
: 1200      1260  6      IF .DIFFERENCE GEQ 0
: 1201      1261  6      THEN
: 1202      1262  7          BEGIN
: 1203      1263  7
: 1204      1264  7              ! found the 1st point at which the magic ratio is positive
: 1205      1265  7              ! was the last point more optimal, if so use it
: 1206      1266  7
: 1207      1267  7
: 1208      1268  7          IF ABS(.DIFFERENCE) GTRU ABS(.LAST_DIFF)
: 1209      1269  7          THEN
: 1210      1270  8              BEGIN
: 1211      1271  8
```



```

: 1212      1272  8      IF .REC_ADDR EQLU .LAST
: 1213      1273  8      THEN
: 1214      1274  8          IRAB[IRB$V_REC_W_LO] = 0
: 1215      1275  8      ELSE
: 1216      1276  9          (REC_ADDR = .LAST;
: 1217      1277  9              IF .REC_ADDR LSSU .POS_INSERT
: 1218      1278  9              THEN
: 1219      1279  9                  IRAB[IRB$V_REC_W_LO] = 0);
: 1220      1280  8
: 1221      1281  8      LAST = 0;
: 1222      1282  8      END;
: 1223      1283  7
: 1224      1284  7      ! 2-bkt split is possible rec_addr points to the most
: 1225      1285  7      ! optimal place
: 1226      1286  7
: 1227      1287  7
: 1228      1288  7
: 1229      1289  7      IF .LAST EQL 0
: 1230      1290  7      THEN      ! just backed up rec_addr, need to recalc last
: 1231      1291  8          BEGIN
: 1232      1292  8              LOCAL
: 1233      1293  8                  TMP;
: 1234      1294  8
: 1235      1295  8              TMP = .REC_ADDR;
: 1236      1296  8              REC_ADDR = .BKT_ADDR + BKT$C_OVERHDSZ;
: 1237      1297  8              LAST = .REC_ADDR;
: 1238      1298  8
: 1239      1299  8              WHILE .REC_ADDR NEQU .TMP
: 1240      1300  8                  DO
: 1241      1301  8                      BEGIN
: 1242      1302  9                          LAST = .REC_ADDR;
: 1243      1303  9                          RM$GETNEXT_REC();
: 1244      1304  9                          END;
: 1245      1305  8
: 1246      1306  8              END;
: 1247      1307  7
: 1248      1308  7      RM$MOVE KEY(.LAST, .REC_ADDR);
: 1249      1309  7      IRAB[IRB$W_SPLIT] = .REC_ADDR - .BKT_ADDR;
: 1250      1310  7
: 1251      1311  7      ! treat another exception case of the new record going off into
: 1252      1312  7      ! a cont. bkt all by itself
: 1253      1313  7
: 1254      1314  7
: 1255      1315  7
: 1256      1316  7      IF .IRAB[IRB$W_SPLIT] EQLU .IRAB[IRB$W_POS_INS]
: 1257      1317  7      THEN
: 1258      1318  7
: 1259      1319  7          IF .IRAB[IRB$W_SPLIT] EQLU .IRAB[IRB$W_SPLIT_1]
: 1260      1320  7          THEN
: 1261      1321  7              IF NOT .IRAB[IRB$V_REC_W_LO]
: 1262      1322  7              THEN
: 1263      1323  7                  BEGIN
: 1264      1324  8                      BUILTIN
: 1265      1325  8                          AP;
: 1266      1326  8
: 1267      1327  8
: 1268      1328  8
```



```
: 1269      1329  8      AP = 3;
: 1270      1330  8
: 1271      1331  8      IF NOT RM$COMPARE_KEY(KEYBUF_ADDR(2),
: 1272      1332  8              KEYBUF_ADDR(3),
: 1273      1333  8              .IDX_DFN[IDX$B_KEY$Z])
: 1274      1334  8      THEN
: 1275      1335  8          IRAB[IRB$V_CONT_BKT] = 1;
: 1276      1336  8
: 1277      1337  7      END;
: 1278      1338  7
: 1279      1339  7          LEAVE DO_IT
: 1280      1340  7
: 1281      1341  6          END;
: 1282      1342  6
: 1283      1343  6      ! the magic ratio isn't positive yet, so save all the context and
: 1284      1344  6      ! move on to the next record
: 1285      1345  6
: 1286      1346  6      LAST_DIFF = .DIFFERENCE;
: 1287      1347  6      LAST = .REC_ADDR;
: 1288      1348  6
: 1289      1349  6      IF .IRAB[IRB$V_REC_W_LO]
: 1290      1350  6      THEN
: 1291      1351  6          SAVE_REC_W_LO = 1;
: 1292      1352  6
: 1293      1353  5      END;
: 1294      1354  5          ! { end of rhs fits, is this a good point? }
: 1295      1355  5      ! go on to the next record
: 1296      1356  5
: 1297      1357  5      NEXT :
: 1298      1358  6      BEGIN
: 1299      1359  6
: 1300      1360  6      IF .REC_ADDR EQLU .POS_INSERT
: 1301      1361  6      AND
: 1302      1362  6      NOT .IRAB[IRB$V_REC_W_LO]
: 1303      1363  6      THEN
: 1304      1364  7          BEGIN
: 1305      1365  7
: 1306      1366  7      ! if this is an update and we pass the record, check to see if it
: 1307      1367  7      ! needed an rrv
: 1308      1368  7
: 1309      1369  7      IF .IRAB[IRB$V_UPDATE]
: 1310      1370  7      THEN
: 1311      1371  8          BEGIN
: 1312      1372  8
: 1313      1373  8          IF .BDB[BDB$L_VBN] EQLU .IRAB[IRB$L_PUTUP_VBN]
: 1314      1374  8          THEN
: 1315      1375  8              NEED_RRV = .NEED_RRV - 1;
: 1316      1376  8
: 1317      1377  7          END;
: 1318      1378  7
: 1319      1379  7      IRAB[IRB$V_REC_W_LO] = 1;
: 1320      1380  7      RM$MOVE_KEY(.REC_ADDR, .REC_ADDR);
: 1321      1381  7
: 1322      1382  7      IF .REC_ADDR EQLU .EOB
: 1323      1383  7      THEN
: 1324      1384  7          LEAVE NEXT
: 1325      1385  7
```



```
: 1326      1386 7      ELSE IF RM$COMPARE_REC(KEYBUF_ADDR(2), .IDX_DFN[IDX$B_KEYSZ], 0)
: 1327      1387 7      THEN
: 1328      1388 7          LEAVE NEXT;
: 1329      1389 7
: 1330      1390 6      END;          ! { end of at position for insert for the 1st time }
: 1331      1391 6
: 1332      1392 6      ! fool move key a little by always clearing rec_w_lo to always get the
: 1333      1393 6      ! key associated w/ the record at pos_ins. (I think it is the key of
: 1334      1394 6      ! the record we are pointing to, not the one at pos_ins...)
: 1335      1395 6
: 1336      1396 7      BEGIN
: 1337      1397 7
: 1338      1398 7      LOCAL
: 1339      1399 7          TMP : BYTE;
: 1340      1400 7
: 1341      1401 7      TMP = .IRAB[IRB$B_SPL_BITS];
: 1342      1402 7      IRAB[IRB$V_REC_W[0]] = 0;
: 1343      1403 7      RM$MOVE KEY(.REC_ADDR, .REC_ADDR);
: 1344      1404 7      IRAB[IRB$B_SPL_BITS] = .TMP
: 1345      1405 6      END;
: 1346      1406 6
: 1347      1407 6      DO
: 1348      1408 7          BEGIN
: 1349      1409 7
: 1350      1410 7          BUILTIN
: 1351      1411 7              AP;
: 1352      1412 7
: 1353      1413 7          IF .REC_ADDR EQLU .EOB
: 1354      1414 7          THEN
: 1355      1415 7              EXITLOOP;
: 1356      1416 7
: 1357      1417 7          AP = 3;
: 1358      1418 7
: 1359      1419 7          IF .BDB[BDB$L_VBN] EQLU RM$RECORD_VBN()
: 1360      1420 7          THEN
: 1361      1421 7              NEED_RRV = .NEED_RRV - 1
: 1362      1422 7          ELSE
: 1363      1423 7              NOT_NEED_RRV = .NOT_NEED_RRV - 1;
: 1364      1424 7
: 1365      1425 7          RM$GETNEXT_REC();
: 1366      1426 7
: 1367      1427 7          IF .REC_ADDR EQLU .EOB
: 1368      1428 7          THEN
: 1369      1429 7              EXITLOOP;
: 1370      1430 7
: 1371      1431 7          END
: 1372      1432 7
: 1373      1433 7      ! compare_rec returns 0 if a match
: 1374      1434 7      !
: 1375      1435 6      UNTIL RM$COMPARE_REC(KEYBUF_ADDR(2), .IDX_DFN[IDX$B_KEYSZ], 0);
: 1376      1436 6
: 1377      1437 6      ! if the key compares brought us up to the pos of insert, see if the
: 1378      1438 6      ! key of the new record matches. if it does, have to include it w/ the
: 1379      1439 6      ! lhs
: 1380      1440 6      !
: 1381      1441 6
: 1382      1442 6      IF .REC_ADDR EQLU .POS_INSERT
```



```
: 1383      1443  6      THEN
: 1384      1444  7      BEGIN
: 1385      1445  7      BUILTIN
: 1386      1446  7      AP;
: 1387      1447  7
: 1388      1448  7
: 1389      1449  7      AP = 3;
: 1390      1450  7
: 1391      1451  7      IF NOT RM$COMPARE_KEY(KEYBUF_ADDR(2),
: 1392      1452  7      KEYBUF_ADDR(3),
: 1393      1453  7      .IDX_DFN[IDX$B_KEYSZ])
: 1394      1454  7      THEN
: 1395      1455  8      BEGIN
: 1396      1456  8      IRAB[IRB$V_REC_W_LO] = 1;
: 1397      1457  8
: 1398      1458  8      IF .IRAB[IRB$V_UPDATE]
: 1399      1459  8      AND
: 1400      1460  8      .BDB[BDB$L_VBN] EQLU .IRAB[IRB$L_PUTUP_VBN]
: 1401      1461  8      THEN
: 1402      1462  8      NEED_RRV = .NEED_RRV - 1;
: 1403      1463  7      END;
: 1404      1464  7
: 1405      1465  6      END;
: 1406      1466  6
: 1407      1467  6      IF .REC_ADDR GTRU .POS_INSERT
: 1408      1468  6      THEN
: 1409      1469  7      BEGIN
: 1410      1470  7      IRAB[IRB$V_REC_W_LO] = 1;
: 1411      1471  7
: 1412      1472  7      IF .IRAB[IRB$V_UPDATE]
: 1413      1473  7      AND
: 1414      1474  7      .BDB[BDB$L_VBN] EQLU .IRAB[IRB$L_PUTUP_VBN]
: 1415      1475  7      THEN
: 1416      1476  7      NEED_RRV = .NEED_RRV - 1;
: 1417      1477  6      END;
: 1418      1478  6
: 1419      1479  5      END;
: 1420      1480  4      END;
: 1421      1481  4      ! { end of scanning to find optimal split point }
: 1422      1482  3      END;
: 1423      1483  3      ! { end of half }
: 1424      1484  3      ! define a new block here so local storage can be redefined
: 1425      1485  3      !
: 1426      1486  4      BEGIN
: 1427      1487  4
: 1428      1488  4      MACRO
: 1429      1489  4      BEG_CHAIN = LHS %,
: 1430      1490  4      END_CHAIN = RHS %,
: 1431      1491  4      NUM_DUPS = NUM_RRVs %,
: 1432      1492  4      DUPS = RRV %;
: 1433      1493  4
: 1434      1494  4      BUILTIN
: 1435      1495  4      AP;
: 1436      1496  4
: 1437      1497  4      ! must be a 3 or 4 bucket split or we detected ascending order and the new
: 1438      1498  4      ! record was a dupe. we'll optimize here to the extent of trying to keep a
: 1439      1499  4      ! dup chain around the new record together and in the middle bucket
```



```
: 1440      1500  4      ! note that in all the cases that follow the new record is going into the
: 1441      1501  4      ! middle bucket. therefore, the "lhs" will always fit, since it can only
: 1442      1502  4      ! get smaller ( or stay the same size, in the degenerate case). also note
: 1443      1503  4      ! that in any of these case, the left hand bucket may be empty of data
: 1444      1504  4      ! records (have only rrv's in it) if the first split point is at the
: 1445      1505  4      ! beginning and all data records get moved out
: 1446      1506  4
: 1447      1507  4      IRAB[IRBSV_NEW_BKTS] = 2;      ! assume 3-bkt split until shown otherwise
: 1448      1508  4      IRAB[IRBSV_REC_W_LO] = 0;
: 1449      1509  4
: 1450      1510  4      ! initialize key buffer 2 with the contents of key buffer 3 (the value
: 1451      1511  4      ! of the primary key of the record being inserted). This is necessary
: 1452      1512  4      ! when the new record is at the beginning of the bucket and is going into
: 1453      1513  4      ! a bucket all by itself and there were already 255 records in their
: 1454      1514  4      ! original bucket and they all need rrv's therefore they all move into the
: 1455      1515  4      ! next bucket. At any rate, that seems to be the only case where key buffer
: 1456      1516  4      ! 2 is not correct coming into here and will be set correctly before
: 1457      1517  4      ! leaving.
: 1458      1518  4
: 1459      1519  4      RM$MOVE(.IDX_DFN[IDX$B_KEYSZ], KEYBUF_ADDR(3), KEYBUF_ADDR(2));
: 1460      1520  4
: 1461      1521  4      ! find beginning and end of this possible dups chain equal to the key value
: 1462      1522  4      ! of the record being inserted.
: 1463      1523  4
: 1464      1524  4      REC_ADDR = .BKT_ADDR + BKT$C_OVERHDSZ;
: 1465      1525  5      BEGIN
: 1466      1526  5
: 1467      1527  5      LOCAL
: 1468      1528  5          STATUS;
: 1469      1529  5
: 1470      1530  5      WHILE STATUS = RM$COMPARE_REC(KEYBUF_ADDR(3), .IRAB[IRBSB_KEYSZ], 0)
: 1471      1531  5      DO
: 1472      1532  6          BEGIN
: 1473      1533  6
: 1474      1534  6          IF .REC_ADDR LSSU .POS_INSERT
: 1475      1535  6          THEN
: 1476      1536  7              BEGIN
: 1477      1537  7                  AP = 0;
: 1478      1538  7                  RM$RECORD_KEY(KEYBUF_ADDR(2));
: 1479      1539  6                  END;
: 1480      1540  6
: 1481      1541  6          IF .REC_ADDR EQLU .EOB
: 1482      1542  6          OR
: 1483      1543  6          .STATUS LSS 0
: 1484      1544  6          THEN
: 1485      1545  7              BEGIN
: 1486      1546  7
: 1487      1547  7                  ! !!!! SPLIT TYPE 3 !!!! no duplicates found for simplicity, do a
: 1488      1548  7                  ! 3-bkt split at the point of insert w/ new record in its own
: 1489      1549  7                  ! bucket
: 1490      1550  7
: 1491      1551  7                  IRAB[IRBSW_SPLIT] = IRAB[IRBSW_SPLIT_1] = .IRAB[IRBSW_POS_INS];
: 1492      1552  7                  LEAVE DO_IT
: 1493      1553  7
: 1494      1554  7                  ! { end of didn't find a duplicate, put record in its own bucket }
: 1495      1555  7
: 1496      1556  6          END;
```



```
: 1497      1557  6
: 1498      1558  6      RM$GETNEXT_REC();
: 1499      1559  5      END;                                ! { end of while no duplicate has been found }
: 1500      1560  5
: 1501      1561  4      END;                                ! { end of block defining status for while loop }
: 1502      1562  4
: 1503      1563  4      ! found the beginning of the dups chain, now find the end
: 1504      1564  4      !
: 1505      1565  4      NUM_DUPS = 0;
: 1506      1566  4      BEG_CHAIN = .REC_ADDR;
: 1507      1567  4
: 1508      1568  4      DO
: 1509      1569  5          BEGIN
: 1510      1570  5              NUM_DUPS = .NUM_DUPS + 1;
: 1511      1571  5              RM$GETNEXT_REC();
: 1512      1572  5
: 1513      1573  5              IF .REC_ADDR EQLU .EOB
: 1514      1574  5              THEN
: 1515      1575  5                  EXITLOOP;
: 1516      1576  5
: 1517      1577  5              END
: 1518      1578  4      UNTIL RM$COMPARE_REC(KEYBUF_ADDR(3), .IRAB[IRB$B_KEYSZ], 0);
: 1519      1579  4                  ! compare_rec returns 0 if keys match
: 1520      1580  4
: 1521      1581  4      END_CHAIN = .REC_ADDR;
: 1522      1582  4
: 1523      1583  4      ! found the beginning and the end of the chain calculate its size if we got
: 1524      1584  4      ! here via an update, we never called rm$srch_by_key to set dups_seen
: 1525      1585  4      ! for us. so let us do that now if necessary
: 1526      1586  4      !
: 1527      1587  4
: 1528      1588  4      IF .POS_INSERT GTRU .BEG_CHAIN
: 1529      1589  4      THEN
: 1530      1590  4          IRAB[IRB$V_DUPS_SEEN] = 1;
: 1531      1591  4
: 1532      1592  4      DUPS = .END_CHAIN - .BEG_CHAIN;
: 1533      1593  4      DUPS = .DUPS + .REC_SIZE;
: 1534      1594  4
: 1535      1595  4      IF .DUPS LSSU .BKTSIZE
: 1536      1596  4          !
: 1537      1597  4          ! if there are 255 dups on a put, there won't be enough id's in the
: 1538      1598  4          ! new bucket even if there is enough space for them.
: 1539      1599  4          !
: 1540      1600  4          AND
: 1541      1601  5              (.IRAB[IRB$V_UPDATE]
: 1542      1602  5              OR
: 1543      1603  5              .NUM_DUPS<0, 8> LEQU 254)
: 1544      1604  4      THEN
: 1545      1605  5          BEGIN
: 1546      1606  5              !+
: 1547      1607  5              !!!!! SPLIT TYPE 1 !!!!!
: 1548      1608  5              ! duplicates found and fortunately, they all fit
: 1549      1609  5              ! in one bucket so 3-bkt split w/ all of the dups in the middle bucket
: 1550      1610  5              ! because of the optimization used for dups being inserted "in order"
: 1551      1611  5              ! this can still be a 2-bkt split if the new record is being inserted
: 1552      1612  5              ! at the end of the bucket
: 1553      1613  5
```



```
: 1554      1614  5      | 22-jan-79 if loa forced us to think that a
: 1555      1615  5      | bkt w/ all dups had to be split ( only on put) be smart and just put
: 1556      1616  5      | new record by itself a better solution would be not to split at all,
: 1557      1617  5      | but at this date it's rather inconceivable
: 1558      1618  5      | 23-jan-79 it's not only loa
: 1559      1619  5      | that can fool us, the bkt might have had a lot of rrv's
: 1560      1620  5      |
: 1561      1621  5      |
: 1562      1622  5      | IRAB[IRB$W_SPLIT] = .BEG_CHAIN - .BKT_ADDR;
: 1563      1623  5      | IRAB[IRB$W_SPLIT_1] = .END_CHAIN - .BKT_ADDR;
: 1564      1624  5      |
: 1565      1625  5      | IF .END_CHAIN EQLU .EOB
: 1566      1626  5      | THEN
: 1567      1627  6      |     BEGIN
: 1568      1628  6      |         IRAB[IRB$V_NEW_BKTS] = 1;
: 1569      1629  6      |
: 1570      1630  7      |         IF .BEG_CHAIN EQLU (.BKT_ADDR + BKT$C_OVERHDSZ)
: 1571      1631  6      |         THEN
: 1572      1632  7      |             BEGIN
: 1573      1633  7      |                 IRAB[IRB$W_SPLIT_1] = .IRAB[IRB$W_SPLIT_2];
: 1574      1634  7      |                 IRAB[IRB$W_SPLIT] = .IRAB[IRB$W_POS_INS];
: 1575      1635  7      |                 IRAB[IRB$V_CONT_BKT] = 1;
: 1576      1636  7      |             END
: 1577      1637  7      |
: 1578      1638  6      |         END
: 1579      1639  5      |     ELSE
: 1580      1640  6      |         BEGIN
: 1581      1641  6      |
: 1582      1642  6      |             IF .IRAB[IRB$W_SPLIT] EQLU BKT$C_OVERHDSZ<0, 16>
: 1583      1643  6      |             THEN
: 1584      1644  6      |                 IRAB[IRB$V_EMPTY_BKT] = 1;
: 1585      1645  6      |
: 1586      1646  6      |             ! Only force the record into the low bucket if it is not the
: 1587      1647  6      |             ! first one in the duplicate chain.
: 1588      1648  6      |             !
: 1589      1649  6      |
: 1590      1650  6      |             IF .END_CHAIN GEQU .POS_INSERT
: 1591      1651  6      |             AND .IRAB[IRB$W_SPLIT] NEQU .IRAB[IRB$W_POS_INS]
: 1592      1652  6      |             THEN
: 1593      1653  6      |                 IRAB[IRB$V_REC_W_LO] = 1;
: 1594      1654  5      |             END;
: 1595      1655  5      |
: 1596      1656  5      |         LEAVE DO_IT
: 1597      1657  5      |
: 1598      1658  4      |     END;      ! { end of duplicates found and they fit in one bucket }
: 1599      1659  4      |
: 1600      1660  4      | ! if we had 255 dups above we dropped thru to here and this next test
: 1601      1661  4      | ! will fail because it can only happen on an update so the all dups case
: 1602      1662  4      | ! will fall thru to split type 2, which will put the new record by itself.
: 1603      1663  4      | ! consider oddball update case in which there are dups before and after
: 1604      1664  4      | ! position of insert. ( note that if this case doesn't apply, the duplicates
: 1605      1665  4      | ! were only before or after -- and didn't fit w/ record -- so new record
: 1606      1666  4      | ! will end up by itself. for code flow purposes, leave that till later).
: 1607      1667  4      |
: 1608      1668  4      |
: 1609      1669  4      | IF .IRAB[IRB$V_DUPS_SEEN]
: 1610      1670  4      |     AND
```



```
1611      1671  4      .END_CHAIN GTRU .POS_INSERT
1612      1672  4      THEN
1613      1673  5      BEGIN
1614      1674  5
1615      1675  5      IF .DUPS - (.POS_INSERT - .BEG_CHAIN) LSSU .BKTSIZE
1616      1676  5      THEN
1617      1677  5          ! if high dups will fit w/ record, put them in a bucket together
1618      1678  5          !
1619      1679  5          BEGIN
1620      1680  6
1621      1681  6          +
1622      1682  6          !!!!! SPLIT TYPE 4 !!!!!
1623      1683  6          ! 3 bkt split where middle bkt is a continuation bkt containing
1624      1684  6          ! new record and dups following it
1625      1685  6          !
1626      1686  6          !!!!! AND SPLIT TYPE 4B !!!!! however, if the hi set consists
1627      1687  6          ! solely of duplicates, we can still have a 2-bkt split case that
1628      1688  6          ! would not have been picked up by the previous algorithm ( since
1629      1689  6          ! it won't divide dups).
1630      1690  6          !
1631      1691  6          -
1632      1692  6
1633      1693  6          IRAB[IRB$V_CONT_BKT] = 1;
1634      1694  6          IRAB[IRB$W_SPLIT] = .IRAB[IRB$W_POS_INS];
1635      1695  6
1636      1696  6          IF .END_CHAIN EQLU .EOB
1637      1697  6          THEN
1638      1698  6              IRAB[IRB$V_NEW_BKTS] = 1
1639      1699  6          ELSE
1640      1700  6              IRAB[IRB$W_SPLIT_1] = .END_CHAIN - .BKT_ADDR;
1641      1701  6
1642      1702  6          REC_ADDR = .BEG_CHAIN;
1643      1703  6          AP = 0;
1644      1704  6          RMS$RECORD_KEY(KEYBUF_ADDR(2));
1645      1705  6          LEAVE DO_IT
1646      1706  6
1647      1707  5          END;
1648      1708  5
1649      1709  5          ! try to fit new record w/ before-dups in middle bucket
1650      1710  5          !
1651      1711  5
1652      1712  5      IF .DUPS - (.END_CHAIN - .POS_INSERT) LSSU .BKTSIZE
1653      1713  5      THEN
1654      1714  6          BEGIN
1655      1715  6
1656      1716  6          +
1657      1717  6          !!!!! SPLIT TYPE 5 !!!!!
1658      1718  6          ! 3 or 4 bkt split ( depending on status of
1659      1719  6          ! high set) where left-middle bkt is new record w/ before-dups
1660      1720  6          ! and right-middle bkt, if it is needed, is a continuation bkt
1661      1721  6          ! w/ the after-dups. it is needed iff the dups aren't the whole hi
1662      1722  6          ! set it still is a continuation bkt.
1663      1723  6          !
1664      1724  6          !***** NOTE FROM NOV-7-78
1665      1725  6          ! This case doesn't take into account the fact that the
1666      1726  6          ! whole bucket may be dups. In the case of all dups, we could
1667      1727  6          ! end up generating an empty bucket when we don't have to (if
```



```
: 1668      1728 6      ! no RRV's) or a relatively useless bucket (some RRV's). In any
: 1669      1729 6      ! event we could end up generating an extra bucket when we
: 1670      1730 6      ! don't have to
: 1671      1731 6      !-
: 1672      1732 6
: 1673      1733 6      IRAB[IRB$W_SPLIT] = .BEG_CHAIN - .BKT_ADDR;
: 1674      1734 6      IRAB[IRB$W_SPLIT_1] = .IRAB[IRB$W_POS_INS];
: 1675      1735 6
: 1676      1736 6      IF .IRAB[IRB$W_SPLIT] EQLU BKT$C_OVERHDSZ<0, 16>
: 1677      1737 6      THEN
: 1678      1738 6          IRAB[IRB$V_EMPTY_BKT] = 1;
: 1679      1739 6
: 1680      1740 6      IRAB[IRB$V_REC_W_LO] = 1;
: 1681      1741 6
: 1682      1742 6      IF .END_CHAIN LSSU .EOB
: 1683      1743 6      THEN
: 1684      1744 7          BEGIN
: 1685      1745 7              IRAB[IRB$V_NEW_BKTS] = 3;
: 1686      1746 7              IRAB[IRB$W_SPLIT_2] = .END_CHAIN - .BKT_ADDR;
: 1687      1747 7          END
: 1688      1748 6      ELSE
: 1689      1749 6          IRAB[IRB$V_CONT_R] = 1;
: 1690      1750 6
: 1691      1751 6      LEAVE DO_IT
: 1692      1752 6
: 1693      1753 5      END;
: 1694      1754 5
: 1695      1755 5      ! { end of oddball update case w/ dups on both sides of new record }
: 1696      1756 5
: 1697      1757 4      END;
: 1698      1758 4
: 1699      1759 4      !+
: 1700      1760 4      !!!!! SPLIT TYPE 2 !!!!!
: 1701      1761 4      ! the new record must go all by itself therefore,
: 1702      1762 4      ! this is a 3-bkt split if there are no after-dups or no hi set and a 4-bkt
: 1703      1763 4      ! split if both of those exist even more exceptional, this can still be a
: 1704      1764 4      ! 2-bkt split if there is no hi set at all ---- i.e., eob = end of the dups
: 1705      1765 4      ! chain
: 1706      1766 4      !-
: 1707      1767 4
: 1708      1768 4      IRAB[IRB$W_SPLIT] = IRAB[IRB$W_SPLIT_1] = .IRAB[IRB$W_POS_INS];
: 1709      1769 4
: 1710      1770 4      IF .IRAB[IRB$V_DUPS_SEEN]
: 1711      1771 4      THEN
: 1712      1772 5          BEGIN
: 1713      1773 5              IRAB[IRB$V_CONT_BKT] = 1;
: 1714      1774 5              REC_ADDR = .BEG_CHAIN;
: 1715      1775 5              AP = 0;
: 1716      1776 5              RMS$RECORD_KEY(KEYBUF_ADDR(2));
: 1717      1777 4          END;
: 1718      1778 4
: 1719      1779 4      IF .POS_INSERT EQLU .EOB
: 1720      1780 4      THEN
: 1721      1781 4          IRAB[IRB$V_NEW_BKTS] = 1
: 1722      1782 4      ELSE
: 1723      1783 4
: 1724      1784 4          IF .POS_INSERT LSSU .END_CHAIN
```



```
: 1725      1785  4      THEN
: 1726      1786  5      BEGIN
: 1727      1787  5
: 1728      1788  5      IF .END_CHAIN LSSU .EOB
: 1729      1789  5      THEN
: 1730      1790  5          IRAB[IRB$V_NEW_BKTS] = 3
: 1731      1791  5      ELSE
: 1732      1792  5          IRAB[IRB$V_CONT_R] = 1;
: 1733      1793  5
: 1734      1794  5      IRAB[IRB$W_SPLIT_2] = .END_CHAIN - .BKT_ADDR;
: 1735      1795  4      END;
: 1736      1796  4
: 1737      1797  3      END;          ! { end of block defining local symbols }
: 1738      1798  3
: 1739      1799  2      END;          ! { end of do_it }
: 1740      1800  2
: 1741      1801  2      ! if the first split point is at the beginning of the data, this means that
: 1742      1802  2      ! all data records will be moved out and only rrv's will be left in the
: 1743      1803  2      ! original bucket ..... therefore, we can mark this bucket as empty
: 1744      1804  2      !
: 1745      1805  2
: 1746      1806  2      IF .IRAB[IRB$W_SPLIT] EQLU BKT$C_OVERHDSZ<0, 16>
: 1747      1807  2      AND
: 1748      1808  2      NOT .IRAB[IRB$V_REC_W_LO]
: 1749      1809  2      THEN
: 1750      1810  2          IRAB[IRB$V_EMPTY_BKT] = 1;
: 1751      1811  2
: 1752      1812  2      RETURN;
: 1753      1813  2
: 1754      1814  1      END;          ! { end of routine }
```

				OC	BB	00000	RMSSPLIT_UDR::			
							PUSHR	#^M<R2,R3>	: 0874	
							SUBL2	#32, SP		
							BICB2	#72, 68(IRAB)	: 0986	
44	A9	02	44	5E	20	C2	00002	INSV	#1, #1, #2, 68(IRAB)	: 0987
				A9	8F	8A	00005	CLRL	NUM_RRVs	: 0988
				01	01	F0	0000A	PUSHL	REC_ADDR	: 0989
					7E	D4	00010	MOVAB	14(R5), R2	: 0990
					56	DD	00012	MOVL	R2, REC_ADDR	
				52	0E	A5	9E	MOVZWL	4(BKT_ADDR), R0	: 0991
				56	52	D0	00018	PUSHAB	(R0)[BKT_ADDR]	
				50	04	A5	3C	CLRL	LAST	: 0992
					6045	9F	0001F	BBS	#3, (REC_ADDR), 4\$: 1000
					7E	D4	00022	MOVL	#3, AP	: 1004
		20		66	03	E0	00024	BSBW	RM\$RECORD_VBN	: 1006
				5C	03	D0	00028	CMPL	28(BDB), R0	
					0000G	30	0002B	BNEQ	2\$	
				50	1C	A4	D1	INCW	NUM_RRVs	: 1008
					05	12	00032	BRB	3\$	
					0C	AE	B6	INCW	NUM_RRVs+2	: 1018
					03	11	00037	MOVL	REC_ADDR, LAST	: 1020
				6E	0E	AE	B6	BSBW	RM\$GETNEXT_REC	: 1021
					56	D0	0003C			
					0000G	30	0003F			

		04	AE		56	D1	00042		CMPL	REC_ADDR, EOB	1023
					DC	1F	00046		BLSSU	1\$	
50		56			55	C3	00048	4\$:	SUBL3	BKT_ADDR, REC_ADDR, R0	1029
	4E	A9			50	B0	0004C		MOVW	R0, 78(IRAB)	
	4C	A9			50	B0	00050		MOVW	R0, 76(IRAB)	
51		51		17	A7	9A	00054		MOVZBL	23(IDX_DFN), R1	1030
		51			09	78	00058		ASHL	#9, R1, R1	
	1C	AE		F1	A1	9E	0005C		MOVAB	-15(R1), BKTSIZE	
	2C	AE		22	A8	3C	00061		MOVZWL	34(RAB), REC_SIZE	1032
	2C	AE			07	C0	00066		ADDL2	#7, REC_SIZE	
		01		50	AA	91	0006A		CMPB	80(IFAB), #1	1034
					04	13	0006E		BEQL	5\$	
	2C	AE			02	C0	00070		ADDL2	#2, REC_SIZE	1036
0A		A9			03	E1	00074	5\$:	BBC	#3, 6(IRAB), 6\$	1040
	78	A9		1C	A4	D1	00079		CMPL	28(BDB), 120(IRAB)	1044
					03	12	0007E		BNEQ	6\$	
				0C	AE	B6	00080		INCW	NUM_RRVS	1046
20	AE	04	AE		56	C3	00083	6\$:	SUBL3	REC_ADDR, EOB, RRV	1050
		04	AE		56	D0	00089		MOVL	REC_ADDR, EOB	1051
		52			56	D1	0008D		CMPL	REC_ADDR, R2	1056
					03	12	00090		BNEQ	7\$	
					02FD	31	00092		BRW	47\$	
		56		08	AE	D1	00095	7\$:	CMPL	POS_INSERT, REC_ADDR	1087
					3B	12	00099		BNEQ	9\$	
51		6E			01	C1	0009B		ADDL3	#1, LAST, R1	1089
		50			61	9A	0009F		MOVZBL	(R1), R0	
					50	D6	000A2		INCL	R0	
	06	A5			50	91	000A4		CMPB	R0, 6(BKT_ADDR)	
					2C	12	000A8		BNEQ	9\$	
		56			6E	D0	000AA		MOVL	LAST, REC_ADDR	1092
					7E	D4	000AD		CLRL	-(SP)	1094
		7E		20	A7	9A	000AF		MOVZBL	32(IDX_DFN), -(SP)	
		50		00B4	CA	3C	000B3		MOVZWL	180(IFAB), R0	
				60	B940	3F	000B8		PUSHAW	296(IRAB)[R0]	
					0000G	30	000BC		BSBW	RM\$COMPARE_REC	
		5E			0C	C0	000BF		ADDL2	#12, SP	
		03			50	E8	000C2		BLBS	R0, 8\$	
					025B	31	000C5		BRW	43\$	
		50			56	D0	000C8	8\$:	MOVL	REC_ADDR, R0	1103
					FEA9	30	000CB		BSBW	RM\$MOVE_KEY	
	4A	A9		48	A9	B0	000CE		MOVW	72(IRAB), 74(IRAB)	1104
					044B	31	000D3		BRW	72\$	1105
		56		0E	A5	9E	000D6	9\$:	MOVAB	14(R5), REC_ADDR	1116
	24	AE	7FFFFFFF		8F	D0	000DA		MOVL	#2147483647, LAST_DIFF	1117
					6E	D4	000E2		CLRL	LAST	1118
				28	AE	D4	000E4		CLRL	SAVE_REC_W_LO	1119
10	AE	04	AE		56	C3	000E7	10\$:	SUBL3	REC_ADDR, EOB, RHS	1131
		08	AE		56	D1	000ED		CMPL	REC_ADDR, POS_INSERT	1133
					0A	1A	000F1		BGTRU	11\$	
05		44	A9		03	E0	000F3		BBS	#3, 68(IRAB), 11\$	1135
		10	AE	2C	AE	C0	000F8		ADDL2	REC_SIZE, RHS	1137
		1C	AE	10	AE	D1	000FD	11\$:	CMPL	RHS, BKTSIZE	1146
					03	1F	00102		BLSSU	13\$	
					0139	31	00104	12\$:	BRW	32\$	
				06	A5	95	00107	13\$:	TSTB	6(BKT_ADDR)	1148
					13	12	0010A		BNEQ	14\$	
0E		06	A9		03	E0	0010C		BBS	#3, 6(IRAB), 14\$	1150

		50	0E	A5	9E	00111	MOVAB	14(R5), R0	1152
		50		56	D1	00115	CMPL	REC_ADDR, R0	
				05	12	00118	BNEQ	14\$	
	E5	44	A9	03	E1	0011A	BBC	#3, 68(IRAB), 12\$	1154
	50		56	55	C3	0011F	SUBL3	BKT_ADDR, REC_ADDR, R0	1157
			52	A0	9E	00123	MOVAB	-14(R0), LHS	
		08	AE	56	D1	00127	CMPL	REC_ADDR, POS_INSERT	1159
				09	1F	0012B	BLSSU	15\$	
	04	44	A9	03	E1	0012D	BBC	#3, 68(IRAB), 15\$	1161
			52	AE	C0	00132	ADDL2	REC_SIZE, LHS	1163
	50		52	20	AE	C1	00136	RRV, LHS, R0	1172
			51	0C	AE	3C	0013B	MOVZWL	NUM_RRV, R1
14	AE		51	07	C5	0013F	MULL3	#7, R1, 20(SP)	
			50	14	AE	C0	00144	ADDL2	20(SP), R0
		1C	AE	50	D1	00148	CMPL	R0, BKT_SIZE	
				14	1A	0014C	BGTRU	16\$	
			06	A5	95	0014E	TSTB	6(BKT_ADDR)	1180
				38	12	00151	BNEQ	20\$	
	33	06	A9	03	E0	00153	BBS	#3, 6(IRAB), 20\$	1182
				0E	AE	B5	00158	TSTW	NUM_RRV+2
				2E	12	0015B	BNEQ	20\$	1184
	29	44	A9	03	E1	0015D	BBC	#3, 68(IRAB), 20\$	1186
				6E	D5	00162	TSTL	LAST	1190
				03	12	00164	BNEQ	17\$	
				01BA	31	00166	BRW	43\$	
		56		6E	D0	00169	MOVL	LAST, REC_ADDR	1194
		04	28	AE	E8	0016C	BLBS	SAVE_REC_W_LO, 18\$	1196
		44	A9	08	8A	00170	BICB2	#8, 68(IRAB)	1198
				56	D0	00174	MOVL	REC_ADDR, TMP	1209
				0E	A5	9E	00177	MOVAB	14(R5), REC_ADDR
				56	D0	0017B	MOVL	REC_ADDR, LAST	1210
				56	D1	0017E	CMPL	REC_ADDR, TMP	1211
				53	6C	13	00181	BEQL	28\$
				6E	56	D0	00183	MOVL	REC_ADDR, LAST
					0000G	30	00186	BSBW	RM\$GETNEXT_REC
					F3	11	00189	BRB	19\$
	50		1C	AE	C5	0018B	MULL3	BKT_SIZE, LHS, R0	1213
	51		14	AE	C3	00190	SUBL3	20(SP), BKT_SIZE, R1	1257
	51	1C	AE	51	C3	00196	SUBL3	R1, RRV, R1	1258
		20	AE	51	C4	0019B	MULL2	RHS, R1	1257
18	AE		50	51	C1	0019F	ADDL3	R1, R0, DIFFERENCE	
				03	18	001A4	BGEQ	21\$	1260
				0086	31	001A6	BRW	31\$	
		51	18	AE	D0	001A9	MOVL	DIFFERENCE, R1	1268
				03	18	001AD	BGEQ	22\$	
		51		51	CE	001AF	MNEGL	R1, R1	
		50	24	AE	D0	001B2	MOVL	LAST_DIFF, R0	
				03	18	001B6	BGEQ	23\$	
		50		50	CE	001B8	MNEGL	R0, R0	
		50		51	D1	001BB	CMPL	R1, R0	
				14	1B	001BE	BLEQU	26\$	
		6E		56	D1	001C0	CMPL	REC_ADDR, LAST	1272
				09	13	001C3	BEQL	24\$	
		56		6E	D0	001C5	MOVL	LAST, REC_ADDR	1276
		08	AE	56	D1	001C8	CMPL	REC_ADDR, POS_INSERT	1278
				04	1E	001CC	BGEQU	25\$	
		44	A9	08	8A	001CE	BICB2	#8, 68(IRAB)	1280

			6E	D4	001D2	25\$:	CLRL	LAST	1282
			6E	D5	001D4	26\$:	TSTL	LAST	1289
		53	17	12	001D6		BNEQ	28\$	
		56	56	D0	001D8		MOVL	REC_ADDR, TMP	1296
		56	0E	A5	9E 001DB		MOVAB	14(R5), REC_ADDR	1297
		6E	56	D0	001DF		MOVL	REC_ADDR, LAST	1298
		53	56	D1	001E2	27\$:	CMPL	REC_ADDR, TMP	1300
			08	13	001E5		BEQL	28\$	
		6E	56	D0	001E7		MOVL	REC_ADDR, LAST	1303
			0000G	30	001EA		BSBW	RMSGETNEXT_REC	1304
			F3	11	001ED		BRB	27\$	1300
		50	6E	D0	001EF	28\$:	MOVL	LAST, R0	1309
			FD82	30	001F2		BSBW	RMSMOVE_KEY	
4A	A9	56	55	A3	001F5		SUBW3	BKT_ADDR, REC_ADDR, 74(IRAB)	1310
		48	A9	4A	A9 B1 001FA		CMPW	74(IRAB), 72(IRAB)	1316
			0C	12	001FF		BNEQ	29\$	
		4C	A9	4A	A9 B1 00201		CMPW	74(IRAB), 76(IRAB)	1319
			05	12	00206		BNEQ	29\$	
	03	44	A9	03	E1 00208		BBC	#3, 68(IRAB), 30\$	1322
			0311	31	0020D	29\$:	BRW	72\$	
		5C	03	D0	00210	30\$:	MOVL	#3, AP	1329
		50	00B4	CA	3C 00213		MOVZWL	180(IFAB), R0	1332
		53	60	B940	3E 00218		MOVAV	@96(IRAB)[R0], R3	
	51	50	60	A9	C1 0021D		ADDL3	96(IRAB), R0, R1	1331
		50	20	A7	9A 00222		MOVZBL	32(IDX_DFN), R0	
			0000G	30	00226		BSBW	RMSCOMPARE_KEY	
		E1	50	E8	00229		BLBS	R0, 29\$	
			01EF	31	0022C		BRW	55\$	1335
		24	AE	18	AE D0 0022F	31\$:	MOVL	DIFFERENCE, LAST_DIFF	1346
		6E	56	D0	00234		MOVL	REC_ADDR, LAST	1347
	04	44	A9	03	E1 00237		BBC	#3, 68(IRAB), 32\$	1349
		28	AE	01	D0 0023C		MOVL	#1, SAVE_REC_W_LO	1351
		08	AE	56	D1 00240	32\$:	CMPL	REC_ADDR, POS_INSERT	1360
			3F	12	00244		BNEQ	36\$	
	3A	44	A9	03	E0 00246		BBS	#3, 68(IRAB), 36\$	1362
	0A	06	A9	03	E1 0024B		BBC	#3, 6(IRAB), 33\$	1369
		78	A9	1C	A4 D1 00250		CMPL	28(BDB), 120(IRAB)	1373
			03	12	00255		BNEQ	33\$	
			0C	AE	B7 00257		DECW	NUM_RRVS	1375
		44	A9	08	88 0025A	33\$:	BISB2	#8, 68(IRAB)	1379
		50	56	D0	0025E		MOVL	REC_ADDR, R0	1380
			FD13	30	00261		BSBW	RMSMOVE_KEY	
		04	AE	56	D1 00264		CMPL	REC_ADDR, EOB	1382
			03	12	00268		BNEQ	35\$	
			FE7A	31	0026A	34\$:	BRW	10\$	
			7E	D4	0026D	35\$:	CLRL	-(SP)	1386
		7E	20	A7	9A 0026F		MOVZBL	32(IDX_DFN), -(SP)	
		50	00B4	CA	3C 00273		MOVZWL	180(IFAB), R0	
			60	B940	9F 00278		PUSHAB	@96(IRAB)[R0]	
			0000G	30	0027C		BSBW	RMSCOMPARE_REC	
		5E	0C	C0	0027F		ADDL2	#12, SP	
		E5	50	E8	00282		BLBS	R0, 34\$	
		53	44	A9	90 00285	36\$:	MOVB	68(IRAB), TMP	1401
		44	A9	08	8A 00289		BICB2	#8, 68(IRAB)	1402
		50	56	D0	0028D		MOVL	REC_ADDR, R0	1403
			FCE4	30	00290		BSBW	RMSMOVE_KEY	
		44	A9	53	90 00293		MOVB	TMP, 68(IRAB)	1404

04	AE	56	D1	00297	37\$:	CMPL	REC_ADDR, EOB	1413	
		35	13	00298		BEQL	40\$		
	5C	03	D0	0029D		MOVL	#3, AP	1417	
		0000G	30	002A0		BSBW	RM\$RECORD_VBN	1419	
	50	1C	A4	D1	002A3	CMPL	28(BDB), R0		
		05	12	002A7		BNEQ	38\$		
		0C	AE	B7	002A9	DECW	NUM_RRVS	1421	
		03	11	002AC		BRB	39\$		
		0E	AE	B7	002AE	38\$:	DECW	NUM_RRVS+2	1423
		0000G	30	002B1	39\$:	BSBW	RM\$GETNEXT_REC	1425	
04	AE	56	D1	002B4		CMPL	REC_ADDR, EOB	1427	
		18	13	002B8		BEQL	40\$		
		7E	D4	002BA		CLRL	-(SP)	1435	
	7E	20	A7	9A	002BC	MOVZBL	32(IDX_DFN), -(SP)		
	50	00B4	CA	3C	002C0	MOVZWL	180(IFAB), R0		
		60	B940	9F	002C5	PUSHAB	@96(IRAB)[R0]		
		0000G	30	002C9		BSBW	RM\$COMPARE_REC		
	5E	0C	C0	002CC		ADDL2	#12, SP		
	C5	50	E9	002CF		BLBC	R0, 37\$		
08	AE	56	D1	002D2	40\$:	CMPL	REC_ADDR, POS_INSERT	1442	
		2F	12	002D6		BNEQ	41\$		
	5C	03	D0	002D8		MOVL	#3, AP	1449	
	50	00B4	CA	3C	002DB	MOVZWL	180(IFAB), R0	1452	
	53	60	B940	3E	002E0	MOVAW	@96(IRAB)[R0], R3		
51	50	60	A9	C1	002E5	ADDL3	96(IRAB), R0, R1	1451	
	50	20	A7	9A	002EA	MOVZBL	32(IDX_DFN), R0		
		0000G	30	002EE		BSBW	RM\$COMPARE_KEY		
	13	50	E8	002F1		BLBS	R0, 41\$		
	44	08	88	002F4		BISB2	#8, 68(IRAB)	1456	
0A	06	03	E1	002F8		BBC	#3, 6(IRAB), 41\$	1458	
	78	1C	A4	D1	002FD	CMPL	28(BDB), 120(IRAB)	1460	
		03	12	00302		BNEQ	41\$		
		0C	AE	B7	00304	DECW	NUM_RRVS	1462	
	08	56	D1	00307	41\$:	CMPL	REC_ADDR, POS_INSERT	1467	
		13	1B	0030B		BLEQU	42\$		
	44	08	88	0030D		BISB2	#8, 68(IRAB)	1470	
0A	06	03	E1	00311		BBC	#3, 6(IRAB), 42\$	1472	
	78	1C	A4	D1	00316	CMPL	28(BDB), 120(IRAB)	1474	
		03	12	0031B		BNEQ	42\$		
		0C	AE	B7	0031D	DECW	NUM_RRVS	1476	
		FDC4	31	00320	42\$:	BRW	10\$	1128	
		02	F0	00323	43\$:	INSV	#2, #1, #2, 68(IRAB)	1507	
	44	08	8A	00329		BICB2	#8, 68(IRAB)	1508	
	50	00B4	CA	3C	0032D	MOVZWL	180(IFAB), R0	1519	
		60	B940	9F	00332	PUSHAB	@96(IRAB)[R0]		
		60	B940	3F	00336	PUSHAW	@96(IRAB)[R0]		
	7E	20	A7	9A	0033A	MOVZBL	32(IDX_DFN), -(SP)		
		0000G	30	0033E		BSBW	RM\$MOVE		
	5E	0C	C0	00341		ADDL2	#12, SP		
28	AE	0E	A5	9E	00344	MOVAB	14(R5), 40(SP)	1524	
	56	28	AE	D0	00349	MOVL	40(SP), REC_ADDR		
		7E	D4	0034D	44\$:	CLRL	-(SP)	1530	
	7E	00A6	C9	9A	0034F	MOVZBL	166(IRAB), -(SP)		
	50	00B4	CA	3C	00354	MOVZWL	180(IFAB), R0		
		60	B940	3F	00359	PUSHAW	@96(IRAB)[R0]		
		0000G	30	0035D		BSBW	RM\$COMPARE_REC		
	5E	0C	C0	00360		ADDL2	#12, SP		

			53		50	D0	00363	MOVL	R0, STATUS		
			35		53	E9	00366	BLBC	STATUS, 50\$		
		08	AE		56	D1	00369	CMPL	REC_ADDR, POS_INSERT	1534	
					11	1E	0036D	BGEQU	45\$		
					5C	D4	0036F	CLRL	AP	1537	
		50		00B4	CA	3C	00371	MOVZWL	180(IFAB), R0	1538	
				60	B940	9F	00376	PUSHAB	@96(IRAB)[R0]		
					0000G	30	0037A	BSBW	RMSRECORD_KEY		
					04	C0	0037D	ADDL2	#4, SP		
		04	5E		56	D1	00380	CMPL	REC_ADDR, EOB	1541	
			AE		04	13	00384	BEQL	46\$		
					53	D5	00386	TSTL	STATUS	1543	
					0F	18	00388	BGEQ	49\$		
		50		48	A9	3C	0038A	MOVZWL	72(IRAB), R0	1551	
		4C	A9		50	B0	0038E	MOVW	R0, 76(IRAB)		
		4A	A9		50	B0	00392	MOVW	R0, 74(IRAB)		
					0188	31	00396	BRW	72\$	1552	
					0000G	30	00399	BSBW	RMSGETNEXT_REC	1558	
					AF	11	0039C	BRB	44\$	1530	
				0C	AE	D4	0039E	CLRL	NUM_RRVS	1565	
		52			56	D0	003A1	MOVL	REC_ADDR, LHS	1566	
				0C	AE	D6	003A4	INCL	NUM_RRVS	1570	
					0000G	30	003A7	BSBW	RMSGETNEXT_REC	1571	
		04	AE		56	D1	003AA	CMPL	REC_ADDR, EOB	1573	
					19	13	003AE	BEQL	52\$		
					7E	D4	003B0	CLRL	-(SP)	1578	
		7E		00A6	C9	9A	003B2	MOVZBL	166(IRAB), -(SP)		
		50		00B4	CA	3C	003B7	MOVZWL	180(IFAB), R0		
				60	B940	3F	003BC	PUSHAW	@96(IRAB)[R0]		
					0000G	30	003C0	BSBW	RMSCOMPARE_REC		
					0C	C0	003C3	ADDL2	#12, SP		
			5E		50	E9	003C6	BLBC	R0, 51\$		
		10	DB		56	D0	003C9	MOVL	REC_ADDR, RHS	1581	
			AE		08	AE	003CD	CMPL	POS_INSERT, LHS	1588	
			52		05	1B	003D1	BLEQU	53\$		
		44	A9	80	8F	88	003D3	BISB2	#128, 68(IRAB)	1590	
	20	AE			52	C3	003D8	SUBL3	LHS, RHS, RRV	1592	
		10	AE	2C	AE	C0	003DE	ADDL2	REC_SIZE, RRV	1593	
		20	AE	20	AE	D1	003E3	CMPL	RRV, BKTSIZE	1595	
		1C	AE		59	1E	003E8	BGEQU	58\$		
					03	E0	003EA	BBS	#3, 6(IRAB), 54\$	1601	
		07			0C	AE	003EF	CMPB	NUM_RRVS, #254	1603	
		06	A9		4D	1A	003F4	BGTRU	58\$		
		FE	8F		55	A3	003F6	SUBW3	BKT_ADDR, LHS, 74(IRAB)	1622	
	4A	A9			55	A3	003FB	SUBW3	BKT_ADDR, RHS, 76(IRAB)	1623	
	4C	A9		10	AE	D1	00401	CMPL	RHS, EOB	1625	
				04	AE	1C	00406	BNEQ	56\$		
44	A9				01	F0	00408	INSV	#1, #1, #2, 68(IRAB)	1628	
		02			52	D1	0040E	CMPL	LHS, 40(SP)	1630	
				28	AE	82	00412	BNEQ	48\$		
		4C	A9	4E	A9	B0	00414	MOVW	78(IRAB), 76(IRAB)	1633	
		4A	A9	48	A9	B0	00419	MOVW	72(IRAB), 74(IRAB)	1634	
		44	A9		10	88	0041E	BISB2	#16, 68(IRAB)	1635	
					6F	11	00422	BRB	62\$	1630	
			0E	4A	A9	B1	00424	CMPW	74(IRAB), #14	1642	
					05	12	00428	BNEQ	57\$		
				44	A9	8F	0042A	BISB2	#64, 68(IRAB)	1644	

		08	AE	10	AE	D1	0042F	57\$:	CMPL	RHS, POS_INSERT	1650
					5D	1F	00434		BLSSU	62\$	1651
		48	A9	4A	A9	B1	00436		CMPW	74(IRAB), 72(IRAB)	1653
					56	13	0043B		BEQL	62\$	1656
		44	A9		08	88	0043D		BISB2	#8, 68(IRAB)	1669
					50	11	00441		BRB	62\$	1671
				44	A9	95	00443	58\$:	TSTB	68(IRAB)	1675
					03	19	00446		BLSS	59\$	1693
					0080	31	00448		BRW	65\$	1694
		08	AE	10	AE	D1	0044B	59\$:	CMPL	RHS, POS_INSERT	1696
					79	1B	00450		BLEQU	65\$	1698
	50		52	08	AE	C3	00452		SUBL3	POS_INSERT, LHS, R0	1700
			50	20	AE	C0	00457		ADDL2	RRV, R0	1702
		1C	AE		50	D1	0045B		CMPL	R0, BKTSIZE	1703
					34	1E	0045F		BGEQU	63\$	1704
		44	A9		10	88	00461		BISB2	#16, 68(IRAB)	1705
		4A	A9	48	A9	B0	00465		MOVW	72(IRAB), 74(IRAB)	1712
		04	AE	10	AE	D1	0046A		CMPL	RHS, EOB	1733
					08	12	0046F		BNEQ	60\$	1734
44	A9		02		01	F0	00471		INSV	#1, #1, #2, 68(IRAB)	1736
					06	11	00477		BRB	61\$	1738
	4C	A9		10	AE	A3	00479	60\$:	SUBW3	BKT_ADDR, RHS, 76(IRAB)	1740
					56	D0	0047F	61\$:	MOVL	LHS, REC_ADDR	1742
					5C	D4	00482		CLRL	AP	1749
				50	00B4	CA	3C	00484	MOVZWL	180(IFAB), R0	1751
					60	B940	9F	00489	PUSHAB	296(IRAB)[R0]	1751
						0000G	30	0048D	BSBW	RM\$RECORD_KEY	1768
						04	C0	00490	ADDL2	#4, SP	1770
					6C	11	00493	62\$:	BRB	67\$	1773
	50		08	AE	10	AE	C3	00495	SUBL3	RHS, POS_INSERT, R0	1774
					20	AE	C0	0049B	ADDL2	RRV, R0	1775
		1C	AE		50	D1	0049F		CMPL	R0, BKTSIZE	1776
					26	1E	004A3		BGEQU	65\$	1777
					55	A3	004A5		SUBW3	BKT_ADDR, LHS, 74(IRAB)	1779
	4A	A9		4C	A9	B0	004AA		MOVW	72(IRAB), 76(IRAB)	1781
					0E	A9	B1	004AF	CMPW	74(IRAB), #14	1783
					05	12	004B3		BNEQ	64\$	1788
		44	A9	40	8F	88	004B5		BISB2	#64, 68(IRAB)	1790
		44	A9		08	88	004BA	64\$:	BISB2	#8, 68(IRAB)	1792
		04	AE	10	AE	D1	004BE		CMPL	RHS, EOB	1794
					4C	1F	004C3		BLSSU	69\$	1796
		44	A9		20	88	004C5		BISB2	#32, 68(IRAB)	1799
					56	11	004C9		BRB	72\$	1801
				50	A9	3C	004CB	65\$:	MOVZWL	72(IRAB), R0	1803
		4C	A9		50	B0	004CF		MOVW	R0, 76(IRAB)	1805
		4A	A9		50	B0	004D3		MOVW	R0, 74(IRAB)	1807
					44	A9	95	004D7	TSTB	68(IRAB)	1809
					18	18	004DA		BGEQ	66\$	1811
		44	A9		10	88	004DC		BISB2	#16, 68(IRAB)	1813
					52	D0	004E0		MOVL	LHS, REC_ADDR	1815
					5C	D4	004E3		CLRL	AP	1817
				50	00B4	CA	3C	004E5	MOVZWL	180(IFAB), R0	1819
					60	B940	9F	004EA	PUSHAB	296(IRAB)[R0]	1821
						0000G	30	004EE	BSBW	RM\$RECORD_KEY	1823
						04	C0	004F1	ADDL2	#4, SP	1825
		04	AE	08	AE	D1	004F4	66\$:	CMPL	POS_INSERT, EOB	1827
					08	12	004F9		BNEQ	68\$	1829

RM3SPLUDR
V04-000

RMSSPLIT_UDR

K 2
16-Sep-1984 02:03:28
14-Sep-1984 13:01:40

VAX-11 Bliss-32 V4.0-742
[RMS.SRC]RM3SPLUDR.B32;1

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44	A9	02	01	01	F0	004FB	INSV	#1, #1, #2, 68(IRAB)	: 1781
				1E	11	00501	BRB	72\$: 1784
		10	AE	08	AE	D1 00503	CMPL	POS_INSERT, RHS	: 1788
				17	1E	00508	BGEQU	72\$: 1790
		04	AE	10	AE	D1 0050A	CMPL	RHS, EOB	: 1792
				06	1E	0050F	BGEQU	70\$: 1794
		44	A9		06	88 00511	BISB2	#6, 68(IRAB)	: 1806
				04	11	00515	BRB	71\$: 1808
		44	A9		20	88 00517	BISB2	#32, 68(IRAB)	: 1810
4E	A9	10	AE		55	A3 0051B	SUBW3	BKT_ADDR, RHS, 78(IRAB)	: 1814
			OE	4A	A9	B1 00521	CMPL	74(IRAB), #14	:
					0A	12 00525	BNEQ	73\$:
	05	44	A9		03	E0 00527	BBS	#3, 68(IRAB), 73\$:
		44	A9	40	8F	88 0052C	BISB2	#64, 68(IRAB)	:
			5E		30	C0 00531	ADDL2	#48, SP	:
					0C	BA 00534	POPR	#^M<R2,R3>	:
					05	00536	RSB		:

; Routine Size: 1335 bytes, Routine Base: RM\$RMS3 + 0089

; 1755 1815 1


```
: 1757 1816 1 %SBTTL 'RMSSPLIT UDR 3'
: 1758 1817 1 GLOBAL ROUTINE RMSSPLIT_UDR_3(RECSZ) : RL$RABREG_4567 NOVALUE =
: 1759 1818 1
: 1760 1819 1 ++
: 1761 1820 1
: 1762 1821 1 FUNCTIONAL DESCRIPTION:
: 1763 1822 1 This routine calculates bucket splits for prologue 3 version files.
: 1764 1823 1
: 1765 1824 1 CALLING SEQUENCE:
: 1766 1825 1 BSBW RMSSPLIT_UDR_3(RECSZ)
: 1767 1826 1
: 1768 1827 1 INPUT PARAMETERS:
: 1769 1828 1 RECSZ - packed record size including overhead
: 1770 1829 1
: 1771 1830 1 IMPLICIT INPUTS:
: 1772 1831 1 BDB pointer
: 1773 1832 1 BUFFER pointer
: 1774 1833 1 REC_ADDR -- point of insert
: 1775 1834 1 RAB -- to be passed to RMSMOVE_KEY
: 1776 1835 1 IDX_DFN
: 1777 1836 1 in IRAB -- CURBDB, associated with bdb and bkt_addr
: 1778 1837 1 POS_INS corresponding to REC_ADDR
: 1779 1838 1 key buffer address
: 1780 1839 1 in IFAB -- key buffer size
: 1781 1840 1 BKT$B_NXTRECID = 0 in original bucket signals that this is
: 1782 1841 1 a split due to a lack of id's in the bucket
: 1783 1842 1
: 1784 1843 1 OUTPUT PARAMETERS:
: 1785 1844 1 none
: 1786 1845 1
: 1787 1846 1 IMPLICIT OUTPUTS:
: 1788 1847 1 in IRAB --
: 1789 1848 1 if 2 bkt split --
: 1790 1849 1 IRB$W_SPLIT, offset to split point
: 1791 1850 1 IRB$V_REC_W_LO -- set if split point is pos_insert and
: 1792 1851 1 record goes with lo set
: 1793 1852 1 key buffer 2 - new high key for original bucket, i.e. key to be
: 1794 1853 1 inserted at the index level
: 1795 1854 1 key buffer 4 - old high key
: 1796 1855 1 number of new buckets = 1
: 1797 1856 1 if original bucket was all rrv's, set IRB$V_EMPTY_BKT flag
: 1798 1857 1 if new bucket is a continuation bkt., set IRB$V_CONT_BKT flag
: 1799 1858 1 if 3 bkt split --
: 1800 1859 1 same as above with these changes:
: 1801 1860 1 key buffer 3 - implicitly it contains second key to be inserted
: 1802 1861 1 at the index level
: 1803 1862 1 IRB$W_SPLIT_1, offset to second split point
: 1804 1863 1 number of new buckets = 2
: 1805 1864 1 if right bucket is a continuation bkt, set IRB$V_CONT_R flag
: 1806 1865 1 if 4 bkt split --
: 1807 1866 1 same as above with these changes:
: 1808 1867 1 IRB$W_SPLIT_2, offset to third split point
: 1809 1868 1 number of new buckets = 3
: 1810 1869 1
: 1811 1870 1 ROUTINE VALUE:
: 1812 1871 1 rmssuc
: 1813 1872 1
```



```
: 1814      1873  1  | SIDE EFFECTS:
: 1815      1874  1  |   AP is clobbered
: 1816      1875  1  |
: 1817      1876  1  |  --
: 1818      1877  1  |
: 1819      1878  2  | BEGIN
: 1820      1879  2  |
: 1821      1880  2  | EXTERNAL REGISTER
: 1822      1881  2  |   COMMON_RAB_STR,
: 1823      1882  2  |   R_REC_ADDR_STR,
: 1824      1883  2  |   R_IDX_DFN_STR,
: 1825      1884  2  |   COMMON_IO_STR;
: 1826      1885  2  |
: 1827      1886  2  | LOCAL
: 1828      1887  2  |   SAVE_REC_W_LO,
: 1829      1888  2  |   NEED_RRV,
: 1830      1889  2  |   POS_INSERT,
: 1831      1890  2  |   EOB,
: 1832      1891  2  |   RRV,
: 1833      1892  2  |   RMS,
: 1834      1893  2  |   LHS,
: 1835      1894  2  |   LAST : REF BBLOCK,
: 1836      1895  2  |   LAST_DIFF,
: 1837      1896  2  |   BKT_SIZE,
: 1838      1897  2  |   DIFFERENCE;
: 1839      1898  2  |
: 1840      1899  2  | LITERAL
: 1841      1900  2  |   RRV_SIZE = 9;
: 1842      1901  2  |
: 1843      1902  2  | LABEL
: 1844      1903  2  |   DO_IT,
: 1845      1904  2  |   HALF,
: 1846      1905  2  |   NEXT;
: 1847      1906  2  |
: 1848      1907  2  | DO_IT :
: 1849      1908  2  |
: 1850      1909  3  | BEGIN
: 1851      1910  3  |
: 1852      1911  3  |   ! define a block so that we can have some common checks before returning
: 1853      1912  3  |   ! successfully
: 1854      1913  3  |   !
: 1855      1914  3  | HALF :
: 1856      1915  3  |
: 1857      1916  4  | BEGIN
: 1858      1917  4  |
: 1859      1918  4  |   !+
: 1860      1919  4  |   ! Define a block so that we can simulate a go-to (naughty, naughty),
: 1861      1920  4  |   ! if we have decided that we are positioning at the end of the bucket
: 1862      1921  4  |   ! & we're in somewhat of an ascending order, where the last record
: 1863      1922  4  |   ! inserted is a duplicate of the new record, skip over the 50-50 code
: 1864      1923  4  |   ! and go to the code to take duplicates into account.
: 1865      1924  4  |
: 1866      1925  4  |   ! scan 1 -- Calculate size of existing rrv's and total number of rrv's
: 1867      1926  4  |   ! needed to move the whole bucket out ( worst case). As a side effect,
: 1868      1927  4  |   ! adjust eob pointer to point to the rrv's instead of freespace. Assume
: 1869      1928  4  |   ! not empty bucket until showed otherwise.
: 1870      1929  4  |   !-
```



```
: 1871      1930  4      IRAB[IRB$V_EMPTY_BKT] = 0;
: 1872      1931  4
: 1873      1932  4
: 1874      1933  4      ! new rec is tried 1st with hi set, then with lo set
: 1875      1934  4
: 1876      1935  4      IRAB[IRB$V_REC_W_LO] = 0;
: 1877      1936  4      IRAB[IRB$V_NEW_BRTS] = 1;      ! assume 2-bkt split until showed otherwise
: 1878      1937  4      NEED_RRV = 0;
: 1879      1938  4      POS_INSERT = .REC_ADDR;
: 1880      1939  4      REC_ADDR = .BKT_ADDR + BKT$C_OVERHDSZ;
: 1881      1940  4      EOB = .BKT_ADDR + .BKT_ADDR[BKT$W_FREESPACE];
: 1882      1941  4      LAST = 0;
: 1883      1942  4
: 1884      1943  4      DO
: 1885      1944  5          BEGIN
: 1886      1945  5
: 1887      1946  5          BUILTIN
: 1888      1947  5              AP;
: 1889      1948  5
: 1890      1949  5          IF .REC_ADDR[IRC$V_RRV]
: 1891      1950  5          THEN
: 1892      1951  5              EXITLOOP;
: 1893      1952  5
: 1894      1953  5          AP = 3;
: 1895      1954  5
: 1896      1955  5          IF .BDB[BDB$L_VBN] EQLU RM$RECORD_VBN()
: 1897      1956  5          THEN
: 1898      1957  5              NEED_RRV = .NEED_RRV + 1;
: 1899      1958  5
: 1900      1959  5          LAST = .REC_ADDR;
: 1901      1960  5
: 1902      1961  5          ! If the front compression of the current record is zero, save its
: 1903      1962  5          ! address as the last noncompressed key. This may prevent a bucket
: 1904      1963  5          ! scan when it comes time to extract and re-expand the key of the
: 1905      1964  5          ! last record in the bucket.
: 1906      1965  5
: 1907      1966  5          IF .IDX_DFN[IDX$V_KEY_COMPR]
: 1908      1967  5          THEN
: 1909      1968  6              BEGIN
: 1910      1969  6
: 1911      1970  6                  IF .(.REC_ADDR + RM$REC_OVHD() + 1) < 0.8 > EQLU 0
: 1912      1971  6                  THEN
: 1913      1972  6                      IRAB[IRB$L_LST_NCMP] = .REC_ADDR;
: 1914      1973  5                  END;
: 1915      1974  5
: 1916      1975  5          RM$GETNEXT_REC()
: 1917      1976  5          END
: 1918      1977  4      UNTIL .REC_ADDR GEQU .EOB;                      ! end of first scan
: 1919      1978  4
: 1920      1979  4      ! Now that we have the address of the last record in the bucket, store
: 1921      1980  4      ! the key of that record in key buffer 4, to be used by index updating.
: 1922      1981  4
: 1923      1982  4      IF .LAST NEQU 0
: 1924      1983  4      THEN
: 1925      1984  5          BEGIN
: 1926      1985  5
: 1927      1986  5          LOCAL
```



```
: 1928      1987  5      TMP_ADDR;
: 1929      1988  5
: 1930      1989  5      BUILTIN
: 1931      1990  5      AP;
: 1932      1991  5
: 1933      1992  5      TMP_ADDR = .REC_ADDR;
: 1934      1993  5      REC_ADDR = .LAST;
: 1935      1994  5      AP = 0;          ! overhead and compressed form
: 1936      1995  5      RM$RECORD_KEY(KEYBUF_ADDR(4));
: 1937      1996  5      REC_ADDR = .TMP_ADDR;
: 1938      1997  4      END;
: 1939      1998  4
: 1940      1999  4      ! Set SPLIT_2 and SPLIT_1 to be EOB, so if there are less than 3 new
: 1941      2000  4      ! buckets BKT_SPL can use the value without having to recalculate it.
: 1942      2001  4
: 1943      2002  4      IRAB[IRB$W_SPLIT_1] = IRAB[IRB$W_SPLIT_2] = .REC_ADDR - .BKT_ADDR;
: 1944      2003  4
: 1945      2004  4      ! Set up the bucket size
: 1946      2005  4
: 1947      2006  4      BKTSIZE = .IDX_DFN[IDX$B_DATBKTSZ]*512 - BKT$C_OVERHDSZ - BKT$C_DATBKTOVH;
: 1948      2007  4
: 1949      2008  4      ! If this is an update, may have to count in an rrv for the existing record
: 1950      2009  4      !
: 1951      2010  4      IF .IRAB[IRB$V_UPDATE]
: 1952      2011  4      THEN
: 1953      2012  4      BEGIN
: 1954      2013  5
: 1955      2014  5
: 1956      2015  5      IF .BDB[BDB$L_VBN] EQLU .IRAB[IRB$L_PUTUP_VBN]
: 1957      2016  5      THEN
: 1958      2017  5      NEED_RRV = .NEED_RRV + 1;
: 1959      2018  4      END;
: 1960      2019  4
: 1961      2020  4      RRV = .EOB - .REC_ADDR;          ! size of existing rrv's
: 1962      2021  4      EOB = .REC_ADDR;              ! adjust eob
: 1963      2022  4
: 1964      2023  4      ! special case it, if the bucket was all rrv's
: 1965      2024  4      !
: 1966      2025  4
: 1967      2026  4      IF .REC_ADDR EQLU .BKT_ADDR + BKT$C_OVERHDSZ
: 1968      2027  4      THEN
: 1969      2028  5      BEGIN
: 1970      2029  5
: 1971      2030  5      ! Bkt is all rrv's yet the record wouldn't fit so we need to
: 1972      2031  5      ! allocate another bkt ( 2 bkt split). Yet special case it so as not
: 1973      2032  5      ! to make another idx entry, only to update the existing one by
: 1974      2033  5      ! setting empty bucket flag.
: 1975      2034  5
: 1976      2035  5      IRAB[IRB$W_SPLIT] = .REC_ADDR - .BKT_ADDR;
: 1977      2036  5      LEAVE DO_IT
: 1978      2037  5
: 1979      2038  4      END;          ! end { of special case an all-rrv bucket }
: 1980      2039  4
: 1981      2040  4      !+
: 1982      2041  4      ! * BLOCK 1 *
: 1983      2042  4      ! Special Case -- If we can detect a possible ascending order to these
: 1984      2043  4      ! records it probably will be better to do a straight point of insert split
```



```
: 1985      2044  4  ! which would put the new record in a bucket all by itself.
: 1986      2045  4  ! Do this kind of split if and only if all the following conditions are met:
: 1987      2046  4  !   1) the record is being inserted at the end of bucket
: 1988      2047  4  !   2) the last record physically in the bkt is the last record to have
: 1989      2048  4  !       been inserted
: 1990      2049  4  !   3) the last record and the new record do not have duplicate key values
: 1991      2050  4  !
: 1992      2051  4  ! Note that if they are duplicates, we can still make an optimization by
: 1993      2052  4  ! skipping the 50-50 split code.
: 1994      2053  4  !
: 1995      2054  4  ! Note that LAST cannot be zero, since if it were we would have an all
: 1996      2055  4  ! rrv bkt.
: 1997      2056  4  !
: 1998      2057  4  !
: 1999      2058  4  IF .POS_INSERT EQLU .REC_ADDR
: 2000      2059  4  AND
: 2001      2060  5  (((.LAST[IRC$W_ID] + 1) AND %X'FFFF') EQLU .BKT_ADDR[BKT$W_NXTRECID])
: 2002      2061  4  THEN
: 2003      2062  5  BEGIN
: 2004      2063  5  REC_ADDR = .LAST;
: 2005      2064  5  !
: 2006      2065  5  ! Check for duplicates:
: 2007      2066  5  ! If the key is compressed, and the new key has a length of zero, then
: 2008      2067  5  ! we know it is a duplicate of the previous one.
: 2009      2068  5  ! If the key is not compressed, then compare the new key (key buffer 3)
: 2010      2069  5  ! with the previous key.
: 2011      2070  5  !
: 2012      2071  5  !
: 2013      2072  5  !
: 2014      2073  5  IF .IDX_DFN[IDX$V_KEY_COMPR]
: 2015      2074  5  THEN
: 2016      2075  6  BEGIN
: 2017      2076  6  !
: 2018      2077  6  ! IF (.IRAB[IRB$L_RECBUF]) < 0,8> NEQU 0
: 2019      2078  6  ! THEN
: 2020      2079  7  ! BEGIN
: 2021      2080  7  !
: 2022      2081  7  ! Since we have detected a possible ascending order in the
: 2023      2082  7  ! input, let's try to optimize a little and split at the point
: 2024      2083  7  ! of insert. Send the record by itself into the new bucket
: 2025      2084  7  ! and store the new high key of the old bucket in keybuf2,
: 2026      2085  7  ! the high key of the new bucket in keybuf4, and split point.
: 2027      2086  7  !
: 2028      2087  7  !
: 2029      2088  7  ! RMSMOVE(.IDX_DFN[IDX$B_KEYSZ],
: 2030      2089  7  ! KEYBUF_ADDR(4),
: 2031      2090  7  ! KEYBUF_ADDR(2));
: 2032      2091  7  ! RMSMOVE(.IDX_DFN[IDX$B_KEYSZ],
: 2033      2092  7  ! KEYBUF_ADDR(3),
: 2034      2093  7  ! KEYBUF_ADDR(4));
: 2035      2094  7  ! IRAB[IRB$W_SPLIT] = .IRAB[IRB$W_POS_INS];
: 2036      2095  7  ! LEAVE DO_IT;
: 2037      2096  6  ! END
: 2038      2097  6  ! ELSE
: 2039      2098  6  ! LEAVE HALF
: 2040      2099  5  ! END
: 2041      2100  6  ! ELSE
: 2041      2100  6  ! BEGIN
```



```
2042 2101 6
2043 2102 6
2044 2103 6
2045 2104 6
2046 2105 6
2047 2106 6
2048 2107 6
2049 2108 6
2050 2109 6
2051 2110 6
2052 2111 6
2053 2112 6
2054 2113 6
2055 2114 6
2056 2115 7
2057 2116 7
2058 2117 7
2059 2118 7
2060 2119 7
2061 2120 7
2062 2121 7
2063 2122 7
2064 2123 7
2065 2124 7
2066 2125 6
2067 2126 6
2068 2127 6
2069 2128 6
2070 2129 6
2071 2130 6
2072 2131 4
2073 2132 4
2074 2133 4
2075 2134 4
2076 2135 4
2077 2136 4
2078 2137 4
2079 2138 4
2080 2139 4
2081 2140 4
2082 2141 4
2083 2142 4
2084 2143 4
2085 2144 4
2086 2145 4
2087 2146 4
2088 2147 4
2089 2148 4
2090 2149 5
2091 2150 5
2092 2151 5
2093 2152 5
2094 2153 5
2095 2154 5
2096 2155 5
2097 2156 5
2098 2157 5

LOCAL
  REC_OVHD;

BUILTIN
  AP;

REC_OVHD = RM$REC_OVHD(0);
AP = 3; ! Contiguous compare of keys

IF RM$COMPARE_KEY ( .REC_ADDR + .REC_OVHD,
                    KEYBUF_ADDR(3),
                    .IDX_DFN[IDX$B_KEYSZ] )
THEN
  BEGIN
    RM$MOVE(.IDX_DFN[IDX$B_KEYSZ],
            KEYBUF_ADDR(4),
            KEYBUF_ADDR(2));
    RM$MOVE(.IDX_DFN[IDX$B_KEYSZ],
            KEYBUF_ADDR(3),
            KEYBUF_ADDR(4));
    IRAB[IRB$W_SPLIT] = .IRAB[IRB$W_POS_INS];
    LEAVE DO_IT;
  END
ELSE
  LEAVE HALF
END

! * end of BLOCK 1 *
END;

REC_ADDR = .BKT_ADDR + BKT$C_OVERHDSZ;
IRAB[IRB$L_LST_NCMP] = .REC_ADDR;
LAST_DIFF = %X7FFFFFFF;
LAST = 0;
SAVE_REC_W_LO = 0;

! * BLOCK 2 *
! Start from the beginning of the bucket and scan rightward. First find the
! 1st place the rhs will fit in 1 bkt then, as long as the lhs will fit in
! a bkt, try to find an optimal point. If there is no point where the rhs
! and lhs will both fit, we can't do a 2-bkt split and this case will fall
! out.

WHILE 1
DO
  BEGIN
    RHS = .EOB - .REC_ADDR;

    IF .REC_ADDR LEQU .POS_INSERT
    AND
    NOT .IRAB[IRB$V_REC_W_LO]
    THEN
      RHS = .RHS + .RECSZ;
```



```
: 2099      2158 5
: 2100      2159 5
: 2101      2160 5
: 2102      2161 5
: 2103      2162 5
: 2104      2163 5
: 2105      2164 5
: 2106      2165 5
: 2107      2166 5
: 2108      2167 5
: 2109      2168 5
: 2110      2169 5
: 2111      2170 5
: 2112      2171 5
: 2113      2172 5
: 2114      2173 6
: 2115      2174 5
: 2116      2175 5
: 2117      2176 5
: 2118      2177 5
: 2119      2178 5
: 2120      2179 5
: 2121      2180 5
: 2122      2181 5
: 2123      2182 5
: 2124      2183 5
: 2125      2184 5
: 2126      2185 5
: 2127      2186 5
: 2128      2187 5
: 2129      2188 5
: 2130      2189 5
: 2131      2190 5
: 2132      2191 5
: 2133      2192 5
: 2134      2193 5
: 2135      2194 5
: 2136      2195 5
: 2137      2196 5
: 2138      2197 5
: 2139      2198 5
: 2140      2199 5
: 2141      2200 5
: 2142      2201 5
: 2143      2202 5
: 2144      2203 5
: 2145      2204 5
: 2146      2205 6
: 2147      2206 6
: 2148      2207 6
: 2149      2208 6
: 2150      2209 7
: 2151      2210 6
: 2152      2211 6
: 2153      2212 5
: 2154      2213 6
: 2155      2214 6

! If the primary key is compressed, then the righthand side total must
! include the count of characters currently front compressed off the
! key of the record which will be first in the right bucket.
IF .IDX_DFN[IDX$V_KEY_COMP]
THEN
! If the point of insertion of the new (updated) record is the same
! as that of the current split point, and the new (updated) record
! is to go in the new (right) bucket, the the number of front
! compressed characters to be added to the righthand total comes
! from the currently compressed key of the new (updated) record.
! This key will be found in keybuffer 5, if the current operation
! is an $UPDATE, or in a record buffer, if the current operation is
! a $PUT.
IF (.REC_ADDR EQA .POS_INSERT)
AND
NOT .IRAB[IRB$V_REC_W_LO]
THEN
IF .IRAB[IRB$V_UPDATE]
THEN
RHS = .RHS + .(KEYBUF_ADDR(5) + 1)<0,8>
ELSE
RHS = .RHS + .(.IRAB[IRB$L_RECBUF] + 1)<0,8>

! If the current split point is not at the point of insertion of
! the new (updated) record, or if it is but the new (updated)
! record is to go in the old (left) bucket, then the first record
! in the new (right) bucket will be the current record, and the
! number of characters currently front compressed off its key is
! added to the righthand side total.
ELSE
IF .REC_ADDR LSSA .EOB
THEN
RHS = .RHS + .(.REC_ADDR + RM$REC_OVHD(0) + 1)<0,8>;

* BLOCK 3 *
! The right hand side fits if there is enough room and there are id's
! available. Id's are always available in the new bucket in the update
! situation, or if we're leaving at least 1 record behind in the old
! bucket. note that nxtrecid is always zeroed if this is a split due to
! lack of id's.

IF .RHS LSSU .BKTSIZE
AND
(.BKT_ADDR[BKT$W_NXTRECID] NEQ 0
OR
.IRAB[IRB$V_UPDATE]
OR
.REC_ADDR NEQA (.BKT_ADDR + BKT$C_OVERHDSZ)
OR
.IRAB[IRB$V_REC_W_LO])
THEN
BEGIN
LHS = .REC_ADDR - (.BKT_ADDR + BKT$C_OVERHDSZ);
```



```
2156 2215 6
2157 2216 6
2158 2217 6
2159 2218 6
2160 2219 6
2161 2220 6
2162 2221 6
2163 2222 6
2164 2223 6
2165 2224 6
2166 2225 6
2167 2226 6
2168 2227 6
2169 2228 6
2170 2229 6
2171 2230 6
2172 2231 6
2173 2232 6
2174 2233 6
2175 2234 6
2176 2235 6
2177 2236 6
2178 2237 7
2179 2238 7
2180 2239 7
2181 2240 7
2182 2241 7
2183 2242 6
2184 2243 7
2185 2244 7
2186 2245 7
2187 2246 7
2188 2247 7
2189 2248 7
2190 2249 7
2191 2250 7
2192 2251 7
2193 2252 7
2194 2253 7
2195 2254 7
2196 2255 7
2197 2256 7
2198 2257 7
2199 2258 7
2200 2259 8
2201 2260 8
2202 2261 8
2203 2262 8
2204 2263 8
2205 2264 8
2206 2265 8
2207 2266 8
2208 2267 8
2209 2268 8
2210 2269 8
2211 2270 9
2212 2271 9

IF .REC_ADDR GEQU .POS_INSERT
AND
.IRAB[IRBSV_REC_W_LO]
THEN
LHS = .LHS + .RECSZ;

! * BLOCK 4 *
! will lhs fit ? lhs doesn't fit if there is no space in the
! bucket, or if there won't be any id's available in the bucket,
! if not & if there is no previous point at which it fit, goto 3-bkt
! split code if there is a previous place where we could have had a
! 2-bkt split, use it

IF .LHS + .RRV + (RRV_SIZE * .NEED_RRV) GTRU .BKTSIZE
! Id's will be available in the original bucket if we aren't
! out of id's to begin with, if this is an update,
! or if the new record is going in the new bucket
OR
(.BKT_ADDR[BKT$W_NXTRECID] EQL 0
AND
NOT .IRAB[IRBSV_UPDATE]
AND
.IRAB[IRBSV_REC_W_LO])
THEN
BEGIN
IF .LAST EQL 0
THEN
EXITLOOP;

REC_ADDR = .LAST;

IF NOT .SAVE_REC_W_LO
THEN
IRAB[IRBSV_REC_W_LO] = 0;

! 2 bkt split is possible rec_addr points to the most
! optimal place since we had to back up, reset last to point
! to the record immediately before the split point
BEGIN
LOCAL
TMP;

TMP = .REC_ADDR;
REC_ADDR = .BKT_ADDR + BKT$C_OVERHDSZ;
LAST = .REC_ADDR;

WHILE .REC_ADDR NEQU .TMP
DO
BEGIN
LAST = .REC_ADDR;
```



```

: 2213      2272  9
: 2214      2273  9
: 2215      2274  9
: 2216      2275  9
: 2217      2276  9
: 2218      2277  9
: 2219      2278  9
: 2220      2279  9
: 2221      2280  9
: 2222      2281 10
: 2223      2282 10
: 2224      2283 10
: 2225      2284 10
: 2226      2285 10
: 2227      2286  9
: 2228      2287  9
: 2229      2288  9
: 2230      2289  8
: 2231      2290  8
: 2232      2291  7
: 2233      2292  7
: 2234      2293  7
: 2235      2294  7
: 2236      2295  7
: 2237      2296  7
: 2238      2297  7
: 2239      2298  7
: 2240      2299  7
: 2241      2300  7
: 2242      2301  7
: 2243      2302  7
: 2244      2303  7
: 2245      2304  7
: 2246      2305  7
: 2247      2306  7
: 2248      2307  8
: 2249      2308  8
: 2250      2309  8
: 2251      2310  8
: 2252      2311  8
: 2253      2312  8
: 2254      2313  8
: 2255      2314  8
: 2256      2315  8
: 2257      2316  8
: 2258      2317  8
: 2259      2318  8
: 2260      2319  8
: 2261      2320  8
: 2262      2321  8
: 2263      2322  8
: 2264      2323  8
: 2265      2324  7
: 2266      2325  7
: 2267      2326  7
: 2268      2327  7
: 2269      2328  6

      ! If the front compression of the current record is zero,
      ! save its address as the last noncompressed key. This may
      ! prevent a bucket scan when it comes time to extract and
      ! re-expand the key of the last record in the bucket
      ! immediately before the split point.
      IF .IDX_DFN[IDX$V_KEY_COMPR]
      THEN
        BEGIN
          IF .(.REC_ADDR + RM$REC_OVHD() + 1) < 0.8> EQLU 0
          THEN
            IRAB[IRB$L_LST_NCMP] = .REC_ADDR;
          END;
        RM$GETNEXT_REC();
      END;
END;
RM$MOVE KEY(.LAST, .REC_ADDR);
IRAB[IRB$W_SPLIT] = .REC_ADDR - .BKT_ADDR;
! treat another exception case of the new record going off into
! a cont. bkt all by itself
IF .IRAB[IRB$W_SPLIT] EQLU .IRAB[IRB$W_POS_INS]
THEN
  IF .IRAB[IRB$W_SPLIT] EQLU .IRAB[IRB$W_SPLIT_1]
  THEN
    IF NOT .IRAB[IRB$V_REC_W_LO]
    THEN
      BEGIN
        BUILTIN
        AP;
        AP = 3;
        ! If the new last key in the bucket equals the key
        ! to be inserted in the new bucket, then we have a
        ! continuation bucket.
        IF NOT RM$COMPARE_KEY ( KEYBUF_ADDR(2),
                                KEYBUF_ADDR(3),
                                .IDX_DFN[IDX$B_KEYSZ] )
        THEN
          IRAB[IRB$V_CONT_BKT] = 1;
        END;
      LEAVE DO_IT
    END;
  ! end of * BLOCK 4 * (LHS does not fit)
```



```
: 2270      2329  6
: 2271      2330  6      ! lhs fits also, calculate the magic ratio
: 2272      2331  6
: 2273      2332  6      DIFFERENCE = (.LHS * .BKTSIZE) -
: 2274      2333  6      (.RHS * (.BKTSIZE - (RRV_SIZE * .NEED_RRV) - .RRV));
: 2275      2334  6
: 2276      2335  6      ! * BLOCK 5 *
: 2277      2336  6
: 2278      2337  6
: 2279      2338  6      IF .DIFFERENCE GEQ 0
: 2280      2339  6      THEN
: 2281      2340  7          BEGIN
: 2282      2341  7
: 2283      2342  7          ! found the 1st point at which the magic ratio is positive
: 2284      2343  7          ! was the last point more optimal, if so use it
: 2285      2344  7
: 2286      2345  7
: 2287      2346  7      IF ABS(.DIFFERENCE) GTRU ABS(.LAST_DIFF)
: 2288      2347  7      THEN
: 2289      2348  8          BEGIN
: 2290      2349  8
: 2291      2350  8          IF .REC_ADDR EQLU .LAST
: 2292      2351  8          THEN
: 2293      2352  8              IRAB[IRB$V_REC_W_LO] = 0
: 2294      2353  8          ELSE
: 2295      2354  9              BEGIN
: 2296      2355  9                  REC_ADDR = .LAST;
: 2297      2356  9
: 2298      2357  9                  IF .REC_ADDR LSSU .POS_INSERT
: 2299      2358  9                  THEN
: 2300      2359  9                      IRAB[IRB$V_REC_W_LO] = 0;
: 2301      2360  8                  END;
: 2302      2361  8
: 2303      2362  8          LAST = 0;
: 2304      2363  7          END;
: 2305      2364  7
: 2306      2365  7      ! 2-bkt split is possible rec_addr points to the most
: 2307      2366  7      ! optimal place
: 2308      2367  7
: 2309      2368  7
: 2310      2369  7      IF .LAST EQL 0
: 2311      2370  7      THEN          ! just backed up rec_addr, need to recalc last
: 2312      2371  8          BEGIN
: 2313      2372  8
: 2314      2373  8          LOCAL
: 2315      2374  8              TMP;
: 2316      2375  8
: 2317      2376  8          TMP = .REC_ADDR;
: 2318      2377  8          REC_ADDR = .BKT_ADDR + BKT$C_OVERHDSZ;
: 2319      2378  8          LAST = .REC_ADDR;
: 2320      2379  8
: 2321      2380  8          WHILE .REC_ADDR NEQU .TMP
: 2322      2381  8          DO
: 2323      2382  9              BEGIN
: 2324      2383  9                  LAST = .REC_ADDR;
: 2325      2384  9
: 2326      2385  9                  ! If the front compression of the current record is
```



```
: 2327      2386  9      : zero, save its address as the last noncompressed key.
: 2328      2387  9      : This may prevent a bucket scan when it comes time to
: 2329      2388  9      : extract and re-expand the key of the last record in
: 2330      2389  9      : the bucket immediately before the split point.
: 2331      2390  9      :
: 2332      2391  9      : IF .IDX_DFN[IDX$V_KEY_COMP]
: 2333      2392  9      : THEN
: 2334      2393 10      : BEGIN
: 2335      2394 10      :
: 2336      2395 10      : IF .(.REC_ADDR + RM$REC_OVHD() + 1) < 0,8> EQLU 0
: 2337      2396 10      : THEN
: 2338      2397 10      :     IRAB[IRB$L_LST_NCMP] = .REC_ADDR;
: 2339      2398  9      : END;
: 2340      2399  9      :
: 2341      2400  9      : RM$GETNEXT_REC();
: 2342      2401  8      : END;
: 2343      2402  8      :
: 2344      2403  7      : END;
: 2345      2404  7      :
: 2346      2405  7      : RM$MOVE KEY(.LAST, .REC_ADDR);
: 2347      2406  7      : IRAB[IRB$W_SPLIT] = .REC_ADDR - .BKT_ADDR;
: 2348      2407  7      :
: 2349      2408  7      : ! treat another exception case of the new record going off into
: 2350      2409  7      : ! a cont. bkt all by itself
: 2351      2410  7      :
: 2352      2411  7      :
: 2353      2412  7      : IF .IRAB[IRB$W_SPLIT] EQLU .IRAB[IRB$W_POS_INS]
: 2354      2413  7      : THEN
: 2355      2414  7      :
: 2356      2415  7      : IF .IRAB[IRB$W_SPLIT] EQLU .IRAB[IRB$W_SPLIT_1]
: 2357      2416  7      : THEN
: 2358      2417  7      :
: 2359      2418  7      : IF NOT .IRAB[IRB$V_REC_W_LO]
: 2360      2419  7      : THEN
: 2361      2420  8      : BEGIN
: 2362      2421  8      :
: 2363      2422  8      : BUILTIN
: 2364      2423  8      : AP;
: 2365      2424  8      :
: 2366      2425  8      : AP = 3;
: 2367      2426  8      :
: 2368      2427  8      : IF NOT RM$COMPARE_KEY ( KEYBUF_ADDR(2),
: 2369      2428  8      :                               KEYBUF_ADDR(3),
: 2370      2429  8      :                               .IDX_DFN[IDX$B_KEYSZ] )
: 2371      2430  8      : THEN
: 2372      2431  8      :     IRAB[IRB$V_CONT_BKT] = 1;
: 2373      2432  8      :
: 2374      2433  7      : END;
: 2375      2434  7      :
: 2376      2435  7      : LEAVE DO_IT
: 2377      2436  7      :
: 2378      2437  6      : END;          ! end of * BLOCK 5 *
: 2379      2438  6      :
: 2380      2439  6      : ! the magic ratio isn't positive yet, so save all the context and
: 2381      2440  6      : ! move on to the next record
: 2382      2441  6      :
: 2383      2442  6      : LAST_DIFF = .DIFFERENCE;
```



```
: 2384      2443  6      LAST = .REC_ADDR;
: 2385      2444  6
: 2386      2445  6      IF .IRAB[IRB$V_REC_W_LO]
: 2387      2446  6      THEN
: 2388      2447  6          SAVE_REC_W_LO = 1;
: 2389      2448  6
: 2390      2449  5      END;                      ! end of * BLOCK 3 *
: 2391      2450  5
: 2392      2451  5      ! Go get the next record, but special case when we are at the position
: 2393      2452  5      ! of insert.
: 2394      2453  5
: 2395      2454  5      NEXT :
: 2396      2455  6          BEGIN
: 2397      2456  6
: 2398      2457  6          IF .REC_ADDR EQLU .POS_INSERT
: 2399      2458  6              AND
: 2400      2459  6              NOT .IRAB[IRB$V_REC_W_LO]
: 2401      2460  6          THEN
: 2402      2461  7              BEGIN
: 2403      2462  7
: 2404      2463  7              ! If this is an update, check to see if it needed an rrv, since
: 2405      2464  7              ! the record will go in the left bucket.
: 2406      2465  7
: 2407      2466  7              IF .IRAB[IRB$V_UPDATE]
: 2408      2467  7              THEN
: 2409      2468  8                  BEGIN
: 2410      2469  8
: 2411      2470  8                  IF .BDB[BDB$L_VBN] EQLU .IRAB[IRB$L_PUTUP_VBN]
: 2412      2471  8                  THEN
: 2413      2472  8                      NEED_RRV = .NEED_RRV - 1;
: 2414      2473  8
: 2415      2474  7                  END;
: 2416      2475  7
: 2417      2476  7              ! Force record to low bucket, and put in key buffer 2 the key
: 2418      2477  7              ! of the record we are inserting (currently in keybuffer 3).
: 2419      2478  7
: 2420      2479  7              IRAB[IRB$V_REC_W_LO] = 1;
: 2421      2480  7              RM$MOVE(.IDX_DFN[IDX$B_KEYSZ], KEYBUF_ADDR(3), KEYBUF_ADDR(2));
: 2422      2481  7
: 2423      2482  7              ! If we are inserting at the end of the bucket, or if the record
: 2424      2483  7              ! at position of insert has a different key from that to be inserted,
: 2425      2484  7              ! leave NEXT so that no other record goes to the left bucket (so far).
: 2426      2485  7              ! If the key is a duplicate, then keep them together in the left
: 2427      2486  7              ! bucket.
: 2428      2487  7
: 2429      2488  7              IF .REC_ADDR EQLU .EOB
: 2430      2489  7              THEN
: 2431      2490  7                  LEAVE NEXT
: 2432      2491  7              ELSE
: 2433      2492  8                  BEGIN
: 2434      2493  8
: 2435      2494  8                  BUILTIN
: 2436      2495  8                      AP;
: 2437      2496  8
: 2438      2497  8                  LOCAL
: 2439      2498  8                      CURR_KEY,
: 2440      2499  8                      REC_OVHD;
```



```
: 2441      2500      8
: 2442      2501      8      REC_OVHD = RM$REC_OVHD(0);
: 2443      2502      8
: 2444      2503      8      ! When the key is compressed, we must build it first in key
: 2445      2504      8      ! buffer 5, and then compare. This build is easy because we
: 2446      2505      8      ! can take the front chars from the key to be inserted.
: 2447      2506      8
: 2448      2507      8
: 2449      2508      8      IF .IDX_DFN[IDX$V_KEY_COMPR]
: 2450      2509      8      THEN
: 2451      2510      9          BEGIN
: 2452      2511      9              CURR_KEY = KEYBUF_ADDR(5);
: 2453      2512      9              RM$MOVE ( (.REC_ADDR + .REC_OVHD + 1)<0,8>,
: 2454      2513      9                  KEYBUF_ADDR(2),
: 2455      2514      9                  .CURR_KEY );
: 2456      2515      9              RM$BUILD_KEY ( .REC_ADDR + .REC_OVHD, .CURR_KEY );
: 2457      2516      9              END
: 2458      2517      8      ELSE
: 2459      2518      8          CURR_KEY = .REC_ADDR + .REC_OVHD;
: 2460      2519      8          AP = 3;      ! Contiguous compare of keys
: 2461      2520      8
: 2462      2521      8          IF RM$COMPARE_KEY ( .CURR_KEY,
: 2463      2522      8              KEYBUF_ADDR(2),
: 2464      2523      8              .IDX_DFN[IDX$B_KEYSZ] )
: 2465      2524      8          THEN
: 2466      2525      8              LEAVE NEXT;
: 2467      2526      7          END;
: 2468      2527      7
: 2469      2528      6          END;      ! end of { at position for insert for the 1st time }
: 2470      2529      6
: 2471      2530      6      ! Now RMS will scan the bucket starting from the current record
: 2472      2531      6      ! position and keeping duplicates together, since RMS does not want to
: 2473      2532      6      ! split the bucket in the middle of a duplicate chain. Before scanning
: 2474      2533      6      ! RMS obtains the size of the current record, saves its address in
: 2475      2534      6      ! IRB$LST_NCMP, if its key is zero front compressed, and saves the
: 2476      2535      6      ! key of the current record in keybuffer 2
: 2477      2536      6
: 2478      2537      7      BEGIN
: 2479      2538      7
: 2480      2539      7      LOCAL
: 2481      2540      7          REC_OVHD,
: 2482      2541      7          S_REC_SIZE,
: 2483      2542      7          NOT_DUP;
: 2484      2543      7
: 2485      2544      7      NOT_DUP = 0;      ! assume duplicates
: 2486      2545      7
: 2487      2546      7      ! Determine the size of the current record.
: 2488      2547      7
: 2489      2548      7      REC_OVHD = RM$REC_OVHD(0; S_REC_SIZE);
: 2490      2549      7
: 2491      2550      7      ! Save the address of the current record if its key is zero front
: 2492      2551      7      ! compressed.
: 2493      2552      7
: 2494      2553      7      IF .IDX_DFN[IDX$V_KEY_COMPR]
: 2495      2554      7          AND
: 2496      2555      7          (.REC_ADDR + .REC_OVHD + 1)<0,8> EQLU 0
: 2497      2556      7      THEN
```



```
: 2498      2557 7      IRAB[IRB$L_LST_NCMP] = .REC_ADDR;
: 2499      2558 7
: 2500      2559 7      ! Move the key of the current record into keybuffer 2. Fool RM$MOVE_KEY
: 2501      2560 7      ! a little by always clearing REC_W_LO so that we get in key buffer-2
: 2502      2561 7      ! the key associated with the record we are pointing to.
: 2503      2562 7
: 2504      2563 8      BEGIN
: 2505      2564 8
: 2506      2565 8      LOCAL
: 2507      2566 8      TMP : BYTE;
: 2508      2567 8
: 2509      2568 8      TMP = .IRAB[IRB$B SPL BITS];
: 2510      2569 8      IRAB[IRB$V REC_W_CO] = 0;
: 2511      2570 8      RM$MOVE_KEY(.REC_ADDR, .REC_ADDR);
: 2512      2571 8      IRAB[IRB$B SPL BITS] = .TMP
: 2513      2572 7      END;
: 2514      2573 7
: 2515      2574 7      ! Position to the next record which does not contain a key duplicate to
: 2516      2575 7      ! that of the current record (whose key has been saved in keybuffer 2).
: 2517      2576 7
: 2518      2577 7      DO
: 2519      2578 8      BEGIN
: 2520      2579 8
: 2521      2580 8      BUILTIN
: 2522      2581 8      AP;
: 2523      2582 8
: 2524      2583 8      IF .REC_ADDR EQLU .EOB
: 2525      2584 8      THEN
: 2526      2585 8      EXITLOOP;
: 2527      2586 8
: 2528      2587 8      AP = 3;
: 2529      2588 8
: 2530      2589 8      IF .BDB[BDB$L_VBN] EQLU RM$RECORD_VBN()
: 2531      2590 8      THEN
: 2532      2591 8      NEED_RRV = .NEED_RRV - 1;
: 2533      2592 8
: 2534      2593 8      REC_ADDR = .REC_ADDR + .REC_OVHD + .S_REC_SIZE;      ! get next rec
: 2535      2594 8
: 2536      2595 8      IF .REC_ADDR EQLU .EOB
: 2537      2596 8      THEN
: 2538      2597 8      EXITLOOP;
: 2539      2598 8
: 2540      2599 8      REC_OVHD = RM$REC_OVHD(0; S_REC_SIZE);
: 2541      2600 8
: 2542      2601 8      IF .IDX_DFN[IDX$V_KEY_COMPR]
: 2543      2602 8      THEN
: 2544      2603 9      BEGIN
: 2545      2604 9
: 2546      2605 9      IF .(.REC_ADDR + .REC_OVHD)<0,8> NEQU 0
: 2547      2606 9      THEN
: 2548      2607 9      NOT_DUP = 1;
: 2549      2608 9      END
: 2550      2609 8      ELSE
: 2551      2610 9      BEGIN
: 2552      2611 9      AP = 3;      ! Contiguous compare of keys
: 2553      2612 9
: 2554      2613 9      IF RM$COMPARE_KEY ( .REC_ADDR + .REC_OVHD,
```



```
: 2555      2614  9      KEYBUF_ADDR(2),
: 2556      2615  9      .IDX_DFN[IDX$B_KEYSZ] )
: 2557      2616  9      THEN
: 2558      2617  9      NOT_DUP = 1;
: 2559      2618  8      END;
: 2560      2619  8
: 2561      2620  8      ! If RMS is currently positioned to the point of insertion of the
: 2562      2621  8      ! updated record, and if the key of the next record matches the
: 2563      2622  8      ! key of the previous record, then the updated record must go
: 2564      2623  8      ! into the old (left) bucket.
: 2565      2624  8
: 2566      2625  8      IF .REC_ADDR EQLU .POS_INSERT
: 2567      2626  8      AND
: 2568      2627  8      NOT .NOT_DUP
: 2569      2628  8      AND
: 2570      2629  8      .IRAB[IRB$V_UPDATE]
: 2571      2630  8      THEN
: 2572      2631  9      BEGIN
: 2573      2632  9
: 2574      2633  9      IRAB[IRB$V_REC_W_LO] = 1;
: 2575      2634  9
: 2576      2635  9      IF .BDB[BDB$L_VBN] EQLU .IRAB[IRB$L_PUTUP_VBN]
: 2577      2636  9      THEN
: 2578      2637  9      NEED_RRV = .NEED_RRV - 1;
: 2579      2638  8      END;
: 2580      2639  8
: 2581      2640  8      END
: 2582      2641  8
: 2583      2642  8      ! Loop until a non-duplicate record is found
: 2584      2643  8
: 2585      2644  7      UNTIL .NOT_DUP;
: 2586      2645  7
: 2587      2646  6      END;
: 2588      2647  6      ! end of block defining NOT_DUP
: 2589      2648  6
: 2590      2649  6      ! If the key compares brought us up to the pos of insert, see if the
: 2591      2650  6      ! key of the new record matches the key of the record before the
: 2592      2651  6      ! position of insert. If it does, have to include the new record with
: 2593      2652  6      ! the lhs.
: 2594      2653  6
: 2595      2654  6      IF .REC_ADDR EQLU .POS_INSERT
: 2596      2655  6      THEN
: 2597      2656  7      BEGIN
: 2598      2657  7
: 2599      2658  7      BUILTIN
: 2600      2659  7      AP;
: 2601      2660  7
: 2602      2661  7      AP = 3;
: 2603      2662  7
: 2604      2663  7      IF NOT RM$COMPARE_KEY( KEYBUF_ADDR(2),
: 2605      2664  7      KEYBUF_ADDR(3),
: 2606      2665  7      .IDX_DFN[IDX$B_KEYSZ] )
: 2607      2666  7      THEN
: 2608      2667  8      BEGIN
: 2609      2668  8      IRAB[IRB$V_REC_W_LO] = 1;
: 2610      2669  8
: 2611      2670  8      IF .IRAB[IRB$V_UPDATE]
```



```
: 2612      2671 8      AND
: 2613      2672 8      .BDB[BDB$L_VBN] EQLU .IRAB[IRB$L_PUTUP_VBN]
: 2614      2673 8      THEN
: 2615      2674 8      NEED_RRV = .NEED_RRV - 1;
: 2616      2675 7      END;
: 2617      2676 7
: 2618      2677 6      END;
: 2619      2678 6
: 2620      2679 5      END;          ! end of NEXT
: 2621      2680 4      END;          ! end of * BLOCK 2 *
: 2622      2681 4
: 2623      2682 3      END;          ! end of HALF
: 2624      2683 3
: 2625      2684 3      ! define a new block here so local storage can be redefined
: 2626      2685 3      !
: 2627      2686 4      BEGIN
: 2628      2687 4
: 2629      2688 4      MACRO
: 2630      2689 4      BEG_CHAIN = LHS %,
: 2631      2690 4      END_CHAIN = RHS %,
: 2632      2691 4      DUPS = RRV %;
: 2633      2692 4
: 2634      2693 4      LOCAL
: 2635      2694 4      FIRST_KEY_EXPANSION;
: 2636      2695 4
: 2637      2696 4      BUILTIN
: 2638      2697 4      AP;
: 2639      2698 4
: 2640      2699 4      ! If we end up with a duplicate chain here, we need to account for the
: 2641      2700 4      ! the fact that the first record which would end up in a new bucket
: 2642      2701 4      ! will have it's first key expanded fully. Initialize the expansion
: 2643      2702 4      ! amount to 0.
: 2644      2703 4
: 2645      2704 4      FIRST_KEY_EXPANSION = 0;
: 2646      2705 4
: 2647      2706 4      ! Must be a 3 or 4 bucket split or we detected ascending order and the new
: 2648      2707 4      ! record was a dupe. We'll optimize here to the extent of trying to keep a
: 2649      2708 4      ! dup chain around the new record together and in the middle bucket.
: 2650      2709 4      ! Note that in all the cases that follow the new record is going into the
: 2651      2710 4      ! middle bucket. Therefore, the "lhs" will always fit, since it can only
: 2652      2711 4      ! get smaller (or stay the same size, in the degenerate case). Also note
: 2653      2712 4      ! that in any of these cases, the left hand bucket may be empty of data
: 2654      2713 4      ! records (have only rrv's in it) if the first split point is at the
: 2655      2714 4      ! beginning and all data records get moved
: 2656      2715 4
: 2657      2716 4      IRAB[IRB$V_NEW_BKTS] = 2; ! assume 3-bkt split until shown otherwise
: 2658      2717 4      IRAB[IRB$V_REC_W_LO] = 0;
: 2659      2718 4
: 2660      2719 4      ! Initialize key buffer 2 with the contents of key buffer 3 (the value
: 2661      2720 4      ! of the primary key of the record being inserted). This is necessary
: 2662      2721 4      ! when the new record is at the beginning of the bucket and is going into
: 2663      2722 4      ! a bucket all by itself so that all the records in the bucket need rrv's
: 2664      2723 4      ! since they all move into the next bucket.
: 2665      2724 4      ! At any rate, that seems to be the only case where key buffer 2 is not
: 2666      2725 4      ! correct coming into here and will be set correctly before leaving.
: 2667      2726 4
: 2668      2727 4      RM$MOVE(.IDX_DFN[IDX$B_KEYSZ], KEYBUF_ADDR(3), KEYBUF_ADDR(2));
```



```
: 2669      2728  4
: 2670      2729  4      ! Find beginning and end of this possible dups chain equal to the key value
: 2671      2730  4      ! of the record being inserted.
: 2672      2731  4
: 2673      2732  4      REC_ADDR = .BKT_ADDR + BKT$C_OVERHDSZ;
: 2674      2733  5      BEGIN
: 2675      2734  5
: 2676      2735  5      LOCAL
: 2677      2736  5          STATUS,
: 2678      2737  5          REC_OVHD,
: 2679      2738  5          S_REC_SIZE,
: 2680      2739  5          CORR_KEY;
: 2681      2740  5
: 2682      2741  5      WHILE 1
: 2683      2742  5      DO
: 2684      2743  6          BEGIN
: 2685      2744  6
: 2686      2745  6          REC_OVHD = RM$REC_OVHD(0; S_REC_SIZE);
: 2687      2746  6
: 2688      2747  6          ! If the key is compressed, it must be rebuilt into keybuffer 5 first
: 2689      2748  6          !
: 2690      2749  6
: 2691      2750  6          IF .IDX_DFNC[IDX$V_KEY_COMPR]
: 2692      2751  6          THEN
: 2693      2752  7              BEGIN
: 2694      2753  7                  CURR_KEY = KEYBUF_ADDR(5);
: 2695      2754  7                  RM$BUILD_KEY ( .REC_ADDR + .REC_OVHD, .CURR_KEY );
: 2696      2755  7                  END
: 2697      2756  6          ELSE
: 2698      2757  6              ! Otherwise, we are already pointing to the beginning of the key
: 2699      2758  6              !
: 2700      2759  6              CURR_KEY = .REC_ADDR + .REC_OVHD;
: 2701      2760  6          AP = 3;      ! Contiguous compare of keys
: 2702      2761  6          STATUS = RM$COMPARE_KEY ( .CURR_KEY,
: 2703      2762  6              KEYBUF_ADDR(3),
: 2704      2763  6              .IDX_DFNC[IDX$B_KEYSZ] );
: 2705      2764  6
: 2706      2765  6          IF NOT .STATUS      ! If key matched, found beginning of chain
: 2707      2766  6          THEN
: 2708      2767  6              EXITLOOP;
: 2709      2768  6
: 2710      2769  6          IF .REC_ADDR LSSU .POS_INSERT
: 2711      2770  6          THEN
: 2712      2771  6              RM$MOVE( .IDX_DFNC[IDX$B_KEYSZ], .CURR_KEY, KEYBUF_ADDR(2) );
: 2713      2772  6
: 2714      2773  6          IF .REC_ADDR EQLU .EOB
: 2715      2774  6              OR
: 2716      2775  6              .STATUS LSS 0
: 2717      2776  6          THEN
: 2718      2777  7              BEGIN
: 2719      2778  7
: 2720      2779  7              ! !!!! SPLIT TYPE 3 !!!!
: 2721      2780  7              ! No duplicates found. For simplicity, do a 3-bkt split at the
: 2722      2781  7              ! point of insert with the new record in its own bucket.
: 2723      2782  7              !
: 2724      2783  7              IRAB[IRB$W_SPLIT] = IRAB[IRB$W_SPLIT_1] = .IRAB[IRB$W_POS_INS];
: 2725      2784  7              LEAVE DO_IT
```



```
: 2726      2785 7
: 2727      2786 7      ! { end of didn't find a duplicate, put record in its own bucket }
: 2728      2787 7
: 2729      2788 6      END;
: 2730      2789 6
: 2731      2790 6      REC_ADDR = .REC_ADDR + .REC_OVHD + .S_REC_SIZE;
: 2732      2791 5      END;      ! { end of while no duplicate has been found }
: 2733      2792 5
: 2734      2793 4      END;      ! { end of block defining status for while loop }
: 2735      2794 4
: 2736      2795 4      ! Found the beginning of the dups chain, now find the end.
: 2737      2796 4
: 2738      2797 4      BEG_CHAIN = .REC_ADDR;
: 2739      2798 4
: 2740      2799 5      BEGIN
: 2741      2800 5
: 2742      2801 5      LOCAL
: 2743      2802 5          NOT_DUP,
: 2744      2803 5          REC_OVHD,
: 2745      2804 5          S_REC_SIZE;
: 2746      2805 5
: 2747      2806 5      NOT_DUP = 0;      ! assume more duplicates
: 2748      2807 5      REC_OVHD = RM$REC_OVHD(0; S_REC_SIZE);
: 2749      2808 5
: 2750      2809 5      ! Ok, keep track of how much the first key would expand if placed
: 2751      2810 5      ! at the beginning of a new bucket.
: 2752      2811 5
: 2753      2812 5      IF .IDX_DFN[IDX$V_KEY_COMP]
: 2754      2813 5      THEN
: 2755      2814 5          FIRST_KEY_EXPANSION = .(.REC_ADDR + .REC_OVHD + 1)<0,8>;
: 2756      2815 5
: 2757      2816 5      DO
: 2758      2817 6          BEGIN
: 2759      2818 6
: 2760      2819 6          REC_ADDR = .REC_ADDR + .REC_OVHD + .S_REC_SIZE;
: 2761      2820 6          IF .REC_ADDR EQU .EOB
: 2762      2821 6          THEN
: 2763      2822 6              EXITLOOP;
: 2764      2823 6
: 2765      2824 6          REC_OVHD = RM$REC_OVHD(0; S_REC_SIZE);
: 2766      2825 6
: 2767      2826 6          IF .IDX_DFN[IDX$V_KEY_COMP]
: 2768      2827 6          THEN
: 2769      2828 7              BEGIN
: 2770      2829 7
: 2771      2830 7                  IF .(.REC_ADDR + .REC_OVHD)<0,8> NEQU 0
: 2772      2831 7                  THEN
: 2773      2832 7                      NOT_DUP = 1
: 2774      2833 7                  END
: 2775      2834 6          ELSE
: 2776      2835 7              BEGIN
: 2777      2836 7
: 2778      2837 7                  AP = 3;      ! Contiguous compare of keys
: 2779      2838 7
: 2780      2839 7                  IF RM$COMPARE_KEY ( .REC_ADDR + .REC_OVHD,
: 2781      2840 7                      KEYBUF_ADDR(3),
: 2782      2841 7                      .IDX_DFN[IDX$B_KEYSZ] )
```



```
: 2783      2842  7      THEN
: 2784      2843  7      NOT_DUP = 1
: 2785      2844  6      END;
: 2786      2845  6
: 2787      2846  6      END
: 2788      2847  6
: 2789      2848  5      UNTIL .NOT_DUP
: 2790      2849  4      END;                      ! end of found end of dups chain
: 2791      2850  4
: 2792      2851  4      END_CHAIN = .REC_ADDR;
: 2793      2852  4
: 2794      2853  4      ! found the beginning and the end of the chain. Calculate its size.
: 2795      2854  4      ! If we got here via an update, we never called RMSSRCH_BY_KEY to set
: 2796      2855  4      ! DUPS_SEEN for us, so let us do that now if necessary. Also be sure
: 2797      2856  4      ! to factor in the amount of key expansion that the first key would
: 2798      2857  4      ! undergo if placed first in a new bucket. If the keys aren't
: 2799      2858  4      ! compressed, don't sweat it -- FIRST_KEY_EXPANSION was initialized
: 2800      2859  4      ! to zero, and only changed if key compression is in effect.
: 2801      2860  4
: 2802      2861  4      IF .POS_INSERT GTRU .BEG_CHAIN
: 2803      2862  4      THEN
: 2804      2863  4          IRAB[IRB$V_DUPS_SEEN] = 1;
: 2805      2864  4
: 2806      2865  4      DUPS = .END_CHAIN - .BEG_CHAIN;
: 2807      2866  4      DUPS = .DUPS + .RECSZ + .FIRST_KEY_EXPANSION;
: 2808      2867  4
: 2809      2868  4      IF .DUPS LSSU .BKTSIZE
: 2810      2869  4      THEN
: 2811      2870  5          BEGIN
: 2812      2871  5              +
: 2813      2872  5              !!!!! SPLIT TYPE 1 !!!!!
: 2814      2873  5              Duplicates found and fortunately, they all fit in one bucket,
: 2815      2874  5              so do a 3-bkt split with all of the dups in the middle bucket.
: 2816      2875  5              Because of the optimization used for dups being inserted "in order"
: 2817      2876  5              this can still be a 2-bkt split if the new record is being inserted
: 2818      2877  5              at the end of the bucket .
: 2819      2878  5
: 2820      2879  5              22-jan-79 If LOA forced us to think that a bkt with all dups had to
: 2821      2880  5              be split ( only on put) be smart and just put new record by itself.
: 2822      2881  5              A better solution would be not to split at all, but at this date
: 2823      2882  5              it's rather inconceivable.
: 2824      2883  5
: 2825      2884  5              23-jan-79 It's not only LOA that can fool us, the bkt might have
: 2826      2885  5              had a lot of rrv's.
: 2827      2886  5              -
: 2828      2887  5
: 2829      2888  5
: 2830      2889  5          IRAB[IRB$W_SPLIT] = .BEG_CHAIN - .BKT_ADDR;
: 2831      2890  5          IRAB[IRB$W_SPLIT_1] = .END_CHAIN - .BKT_ADDR;
: 2832      2891  5
: 2833      2892  5          IF .END_CHAIN EQLU .EOB
: 2834      2893  5          THEN
: 2835      2894  6              BEGIN
: 2836      2895  6                  IRAB[IRB$V_NEW_BKTS] = 1;
: 2837      2896  6
: 2838      2897  7                  IF .BEG_CHAIN EQLU (.BKT_ADDR + BKT$C_OVERHDSZ)
: 2839      2898  6                  THEN
```



```
: 2840      2899 7      BEGIN
: 2841      2900 7      IRAB[IRBSW_SPLIT_1] = .IRAB[IRBSW_SPLIT_2];
: 2842      2901 7      IRAB[IRBSW_SPLIT] = .IRAB[IRBSW_POS_INS];
: 2843      2902 7      IRAB[IRBSV_CONT_BKT] = 1;
: 2844      2903 7      END
: 2845      2904 7
: 2846      2905 6      END
: 2847      2906 5      ELSE
: 2848      2907 6      BEGIN
: 2849      2908 6
: 2850      2909 6      IF .IRAB[IRBSW_SPLIT] EQLU BKTSC_OVERHDSZ<0, 16>
: 2851      2910 6      THEN
: 2852      2911 6          IRAB[IRBSV_EMPTY_BKT] = 1;
: 2853      2912 6
: 2854      2913 6      ! Only force record into the low bucket if it is not the first
: 2855      2914 6      ! one in a duplicate chain.
: 2856      2915 6
: 2857      2916 6
: 2858      2917 6      IF .END_CHAIN GEQU .POS_INSERT
: 2859      2918 6      AND .IRAB[IRBSW_SPLIT] NEQU .IRAB[IRBSW_POS_INS]
: 2860      2919 6      THEN
: 2861      2920 6          IRAB[IRBSV_REC_W_LO] = 1;
: 2862      2921 5      END;
: 2863      2922 5
: 2864      2923 5      LEAVE DO_IT
: 2865      2924 5
: 2866      2925 4      END;      ! { end of duplicates found and they fit in one bucket }
: 2867      2926 4
: 2868      2927 4      ! This next test can only happen on an update so the all dupes case
: 2869      2928 4      ! will fall thru to split type 2, which will put the new record by itself.
: 2870      2929 4      ! Consider oddball update case in which there are dups before and after
: 2871      2930 4      ! position of insert. ( note that if this case doesn't apply, the duplicates
: 2872      2931 4      ! were only before or after -- and didn't fit with record -- so new record
: 2873      2932 4      ! will end up by itself. For code flow purposes, leave that till later).
: 2874      2933 4
: 2875      2934 4
: 2876      2935 4      IF .IRAB[IRBSV_DUPS_SEEN]
: 2877      2936 4      AND
: 2878      2937 4      .END_CHAIN GTRU .POS_INSERT
: 2879      2938 4      THEN
: 2880      2939 5      BEGIN
: 2881      2940 5
: 2882      2941 5      IF .DUPS - (.POS_INSERT - .BEG_CHAIN) LSSU .BKTSIZE
: 2883      2942 5      THEN
: 2884      2943 5
: 2885      2944 5      ! if high dups will fit with record, put them in a bucket together
: 2886      2945 5
: 2887      2946 6      BEGIN
: 2888      2947 6
: 2889      2948 6      !+
: 2890      2949 6      !!!!! SPLIT TYPE 4 !!!!!
: 2891      2950 6      ! 3 bkt split where middle bkt is a continuation bkt containing
: 2892      2951 6      ! new record and dups following it
: 2893      2952 6
: 2894      2953 6      !!!!! AND SPLIT TYPE 4B !!!!! however, if the hi set consists
: 2895      2954 6      ! solely of duplicates, we can still have a 2-bkt split case that
: 2896      2955 6      ! would not have been picked up by the previous algorithm ( since
```



```
: 2897      2956 6      ! it won't divide dups).
: 2898      2957 6      !-
: 2899      2958 6
: 2900      2959 6      IRAB[IRB$V_CONT_BKT] = 1;
: 2901      2960 6      IRAB[IRB$W_SPLIT] = .IRAB[IRB$W_POS_INS];
: 2902      2961 6
: 2903      2962 6      IF .END_CHAIN EQLU .EOB
: 2904      2963 6      THEN
: 2905      2964 6          IRAB[IRB$V_NEW_BKTS] = 1
: 2906      2965 6      ELSE
: 2907      2966 6          IRAB[IRB$W_SPLIT_1] = .END_CHAIN - .BKT_ADDR;
: 2908      2967 6
: 2909      2968 6      REC_ADDR = .BEG_CHAIN;
: 2910      2969 6      RM$MOVE ( .IDX_DFN[IDX$B_KEYSZ], KEYBUF_ADDR(3), KEYBUF_ADDR(2) );
: 2911      2970 6      LEAVE DO_IT
: 2912      2971 6
: 2913      2972 5      END;
: 2914      2973 5
: 2915      2974 5      ! try to fit new record with before-dups in middle bucket
: 2916      2975 5      !
: 2917      2976 5
: 2918      2977 5      IF .DUPS - (.END_CHAIN - .POS_INSERT) LSSU .BKTSIZE
: 2919      2978 5      THEN
: 2920      2979 6          BEGIN
: 2921      2980 6
: 2922      2981 6          !+
: 2923      2982 6          !!!!! SPLIT TYPE 5 !!!!!
: 2924      2983 6          ! 3 or 4 bkt split ( depending on status of
: 2925      2984 6          ! high set) where left-middle bkt is new record with before-dups
: 2926      2985 6          ! and right-middle bkt, if it is needed, is a continuation bkt
: 2927      2986 6          ! with the after-dups. it is needed if the dups aren't the whole hi
: 2928      2987 6          ! set it still is a continuation bkt.
: 2929      2988 6
: 2930      2989 6          ***** NOTE FROM NOV-7-78
: 2931      2990 6          ! This case doesn't take into account the fact that the
: 2932      2991 6          ! whole bucket may be dups. In the case of all dups, we could
: 2933      2992 6          ! end up generating an empty bucket when we don't have to (if
: 2934      2993 6          ! no RRV's) or a relatively useless bucket (some RRV's). In any
: 2935      2994 6          ! event we could end up generating an extra bucket when we
: 2936      2995 6          ! don't have to
: 2937      2996 6          !-
: 2938      2997 6
: 2939      2998 6      IRAB[IRB$W_SPLIT] = .BEG_CHAIN - .BKT_ADDR;
: 2940      2999 6      IRAB[IRB$W_SPLIT_1] = .IRAB[IRB$W_POS_INS];
: 2941      3000 6
: 2942      3001 6      IF .IRAB[IRB$W_SPLIT] EQLU BKT$C_OVERHDSZ<0, 16>
: 2943      3002 6      THEN
: 2944      3003 6          IRAB[IRB$V_EMPTY_BKT] = 1;
: 2945      3004 6
: 2946      3005 6      IRAB[IRB$V_REC_W_LO] = 1;
: 2947      3006 6
: 2948      3007 6      IF .END_CHAIN LSSU .EOB
: 2949      3008 6      THEN
: 2950      3009 7          BEGIN
: 2951      3010 7              IRAB[IRB$V_NEW_BKTS] = 3;
: 2952      3011 7              IRAB[IRB$W_SPLIT_2] = .END_CHAIN - .BKT_ADDR;
: 2953      3012 7          END
```



```
2954 3013 6 ELSE
2955 3014 6 IRAB[IRB$V_CONT_R] = 1;
2956 3015 6
2957 3016 6 LEAVE DO_IT
2958 3017 6
2959 3018 6 END;
2960 3019 5
2961 3020 5 ! { end of oddball update case with dups on both sides of new record }
2962 3021 5 !
2963 3022 4 END;
2964 3023 4
2965 3024 4
2966 3025 4 !+
2967 3026 4 !!!!! SPLIT TYPE 2 !!!!!
2968 3027 4 the new record must go all by itself therefore,
2969 3028 4 this is a 3-bkt split if there are no after-dups or no hi set and a 4-bkt
2970 3029 4 split if both of those exist even more exceptional, this can still be a
2971 3030 4 2-bkt split if there is no hi set at all ---- i.e., eob = end of the dups
2972 3031 4 chain
2973 3032 4 -
2974 3033 4 IRAB[IRB$W_SPLIT] = IRAB[IRB$W_SPLIT_1] = .IRAB[IRB$W_POS_INS];
2975 3034 4
2976 3035 4 IF .IRAB[IRB$V_DUPS_SEEN]
2977 3036 4 THEN
2978 3037 5 BEGIN
2979 3038 5 IRAB[IRB$V_CONT_BKT] = 1;
2980 3039 5 REC_ADDR = .BEG_CHAIN;
2981 3040 5 RMSMOVE ( .IDX_DFN[IDX$B_KEYSZ], KEYBUF_ADDR(3), KEYBUF_ADDR(2) );
2982 3041 4 END;
2983 3042 4
2984 3043 4 IF .POS_INSERT EQLU .EOB
2985 3044 4 THEN
2986 3045 4 IRAB[IRB$V_NEW_BKTS] = 1
2987 3046 4 ELSE
2988 3047 4 IF .POS_INSERT LSSU .END_CHAIN
2989 3048 4 THEN
2990 3049 4 BEGIN
2991 3050 5 IF .END_CHAIN LSSU .EOB
2992 3051 5 THEN
2993 3052 5 IRAB[IRB$V_NEW_BKTS] = 3
2994 3053 5 ELSE
2995 3054 5 IRAB[IRB$V_CONT_R] = 1;
2996 3055 5
2997 3056 5 IRAB[IRB$W_SPLIT_2] = .END_CHAIN - .BKT_ADDR;
2998 3057 5 END;
2999 3058 5
3000 3059 4 END;
3001 3060 4
3002 3061 3 END;
3003 3062 3 ! { end of block defining local symbols }
3004 3063 2 END;
3005 3064 2 ! { end of do_it }
3006 3065 2
3007 3066 2 ! if the first split point is at the beginning of the data, this means that
3008 3067 2 all data records will be moved out and only rrv's will be left in the
3009 3068 2 original bucket ..... therefore, we can mark this bucket as empty
3010 3069 2
```



```

: 3011      3070      2      IF .IRAB[IRBSW_SPLIT] EQLU BKT$C_OVERHDSZ<0, 16>
: 3012      3071      2      AND
: 3013      3072      2      NOT .IRAB[IRBSV_REC_W_LO]
: 3014      3073      2      THEN
: 3015      3074      2      IRAB[IRBSV_EMPTY_BKT] = 1;
: 3016      3075      2
: 3017      3076      2      RETURN;
: 3018      3077      2
: 3019      3078      1      END;                                     ! { end of routine }

```

44	A9	02	44	5E A9 01	48	24 8F 01 7E 56	0C BB 00000	RM\$SPLIT	UDR 3::		
									PUSHR	#^M<R2,R3>	1817
									SUBL2	#36, SP	
									BICB2	#72, 68(IRAB)	1935
									INSV	#1, #1, #2, 68(IRAB)	1936
									CLRL	NEED_RRV	1937
									PUSHL	REC_ADDR	1938
									MOVAB	14(R5), 40(SP)	1939
									MOVL	40(SP), REC_ADDR	
									MOVZWL	4(BKT_ADDR), R0	1940
									PUSHAB	(R0)[BKT_ADDR]	
									CLRL	LAST	1941
									BBS	#3, (REC_ADDR), 4\$	1949
									MOVL	#3, AP	1953
									BSBW	RM\$RECORD_VBN	1955
									CMPL	28(BDB), R0	
									BNEQ	2\$	
									INCL	NEED_RRV	1957
									MOVL	REC_ADDR, LAST	1959
									BBC	#6, -28(IDX_DFN), 3\$	1966
									BSBW	RM\$REC_OVHD	1970
									TSTB	1(R0)[REC_ADDR]	
									BNEQ	3\$	
									MOVL	REC_ADDR, 152(IRAB)	1972
									BSBW	RM\$GETNEXT_REC	1975
									CMPL	REC_ADDR, EOB	1977
									BLSSU	1\$	
									TSTL	LAST	1982
									BEQL	5\$	
									MOVL	REC_ADDR, TMP_ADDR	1992
									MOVL	LAST, REC_ADDR	1993
									CLRL	AP	1994
									MOVZWL	180(IFAB), R0	1995
									MULL2	#3, R0	
									PUSHAB	@96(IRAB)[R0]	
									BSBW	RM\$RECORD_KEY	
									ADDL2	#4, SP	
									MOVL	TMP_ADDR, REC_ADDR	1996
									SUBL3	BKT_ADDR, REC_ADDR, R0	2002
									MOVW	R0, -78(IRAB)	
									MOVW	R0, 76(IRAB)	
									MOVZBL	23(IDX_DFN), R1	2006
									ASHL	#9, R1, R1	

	1C	AE	F0	A1	9E	0008D	MOVAB	-16(R1), BKTSIZE			
	06	A9		03	E1	00092	BBC	#3, 6(IRAB), 6\$	2011		
0A	78	A9	1C	A4	D1	00097	CMPL	28(BDB), 120(IRAB)	2015		
				03	12	0009C	BNEQ	6\$			
			0C	AE	D6	0009E	INCL	NEED RRV	2017		
18	AE	04	AE	56	C3	000A1	SUBL3	REC_ADDR, EOB, RRV	2020		
		04	AE	56	D0	000A7	MOVL	REC_ADDR, EOB	2021		
		30	AE	56	D1	000AB	CMPL	REC_ADDR, 48(SP)	2026		
				03	12	000AF	BNEQ	7\$			
			04	9E	31	000B1	BRW	68\$			
		56	08	AE	D1	000B4	CMPL	POS_INSERT, REC_ADDR	2058		
				0D	12	000B8	BNEQ	8\$			
51		6E		01	C1	000BA	ADDL3	#1, LAST, R1	2060		
		50		61	3C	000BE	MOVZWL	(R1), R0			
				50	D6	000C1	INCL	R0			
	06	A5		50	B1	000C3	CMPW	R0, 6(BKT_ADDR)			
				03	13	000C7	BEQL	9\$			
			00	80	31	000C9	BRW	14\$			
		56		6E	D0	000CC	MOVL	LAST, REC_ADDR	2064		
1F	1C	A7		06	E1	000CF	BBC	#6, 28(IDX_DFN), 12\$	2073		
			68	B9	95	000D4	TSTB	@104(IRAB)	2077		
				03	12	000D7	BNEQ	11\$			
			03	C3	31	000D9	BRW	62\$			
		50	00	B4	CA	3C	000DC	MOVZWL	180(IFAB), R0	2089	
			60	B9	40	9F	000E1	PUSHAB	@96(IRAB)[R0]		
		50	00	B4	CA	3C	000E5	MOVZWL	180(IFAB), R0	2088	
		50		03	C4	000EA	MULL2	#3, R0			
			60	B9	40	9F	000ED	PUSHAB	@96(IRAB)[R0]		
				2B	11	000F1	BRB	13\$	2087		
				51	D4	000F3	CLRL	R1	2108		
			00	00	G	30	000F5	BSBW	RM\$REC_OVHD		
		5C		03	D0	000F8	MOVL	#3, AP	2109		
		52	00	B4	CA	3C	000FB	MOVZWL	180(IFAB), R2	2112	
		53	60	B9	42	3E	00100	MOVAV	@96(IRAB)[R2], R3		
51		56		50	C1	00105	ADDL3	REC_OVHD, REC_ADDR, R1	2111		
		50	20	A7	9A	00109	MOVZBL	32(IDX_DFN), R0			
			00	00	G	30	0010D	BSBW	RM\$COMPARE_KEY		
		C6		50	E9	00110	BLBC	R0, 10\$			
			60	B9	42	9F	00113	PUSHAB	@96(IRAB)[R2]	2118	
		52		03	C4	00117	MULL2	#3, R2	2117		
			60	B9	42	9F	0011A	PUSHAB	@96(IRAB)[R2]		
		7E	20	A7	9A	0011E	MOVZBL	32(IDX_DFN), -(SP)	2116		
			00	00	G	30	00122	BSBW	RM\$MOVE		
		5E		08	C0	00125	ADDL2	#8, SP			
		51	00	B4	CA	3C	00128	MOVZWL	180(IFAB), R1	2121	
50		51		03	C5	0012D	MULL3	#3, R1, R0			
		6E	60	B9	40	9E	00131	MOVAB	@96(IRAB)[R0], (SP)		
			60	B9	41	3F	00136	PUSHAW	@96(IRAB)[R1]	2120	
		7E	20	A7	9A	0013A	MOVZBL	32(IDX_DFN), -(SP)	2119		
			00	00	G	30	0013E	BSBW	RM\$MOVE		
		5E		0C	C0	00141	ADDL2	#12, SP			
	4A	A9	48	A9	B0	00144	MOVW	72(IRAB), 74(IRAB)	2122		
			05	EF	31	00149	BRW	96\$	2123		
		56	30	AE	D0	0014C	MOVL	48(SP), REC_ADDR	2133		
0098		C9		56	D0	00150	MOVL	REC_ADDR, 152(IRAB)	2134		
28	AE	7FFFFFFF		8F	D0	00155	MOVL	#2147483647, LAST_DIFF	2135		
				6E	D4	0015D	CLRL	LAST	2136		

10	AE	04	AE	2C	AE	D4	0015F		CLRL	SAVE_REC_W_LO	2137
		08	AE		56	C3	00162	15\$:	SUBL3	REC_ADDR, EOB, RHS	2150
					56	D1	00168		CMPL	REC_ADDR, POS_INSERT	2152
					0A	1A	0016C		BGTRU	16\$	
	05	44	A9		03	E0	0016E		BBS	#3, 68(IRAB), 16\$	2154
		10	AE	40	AE	C0	00173		ADDL2	RECSZ, RHS	2156
	3E	1C	A7		06	E1	00178	16\$:	BBC	#6, 28(IDX DFN), 20\$	2162
		08	AE		56	D1	0017D		CMPL	REC_ADDR, POS_INSERT	2173
					24	12	00181		BNEQ	19\$	
	1F	44	A9		03	E0	00183		BBS	#3, 68(IRAB), 19\$	2175
	0C	06	A9		03	E1	00188		BBC	#3, 6(IRAB), 17\$	2177
		50		00B4	CA	3C	0018D		MOVZWL	180(IFAB), R0	2179
		50		60	B940	DE	00192		MOVAL	296(IRAB)[R0], R0	
					04	11	00197		BRB	18\$	
		50		68	A9	D0	00199	17\$:	MOVL	104(IRAB), R0	2181
		51		01	A0	9A	0019D	18\$:	MOVZBL	1(R0), R1	
		10	AE		51	C0	001A1		ADDL2	R1, RHS	
					14	11	001A5		BRB	20\$	2177
		04	AE		56	D1	001A7	19\$:	CMPL	REC_ADDR, EOB	2191
					0E	1E	001AB		BGEQU	20\$	
					51	D4	001AD		CLRL	R1	2193
				0000G	30	001AF			BSBW	RMSREC OVHD	
		53		01	A046	9A	001B2		MOVZBL	1(R0)[REC_ADDR], R3	
		10	AE		53	C0	001B7		ADDL2	R3, RHS	
		1C	AE	10	AE	D1	001BB	20\$:	CMPL	RHS, BKTSIZE	2203
					03	1F	001C0		BLSSU	22\$	
				0157	31	001C2		21\$:	BRW	44\$	
				06	A5	B5	001C5	22\$:	TSTW	6(BKT_ADDR)	2205
					13	12	001C8		BNEQ	23\$	
	0E	06	A9		03	E0	001CA		BBS	#3, 6(IRAB), 23\$	2207
		50		0E	A5	9E	001CF		MOVAB	14(R5), R0	2209
		50			56	D1	001D3		CMPL	REC_ADDR, R0	
					05	12	001D6		BNEQ	23\$	
	E5	44	A9		03	E1	001D8		BBC	#3, 68(IRAB), 21\$	2211
	50	56			55	C3	001DD	23\$:	SUBL3	BKT_ADDR, REC_ADDR, R0	2214
		52		F2	A0	9E	001E1		MOVAB	-14(R0), LHS	
		08	AE		56	D1	001E5		CMPL	REC_ADDR, POS_INSERT	2216
					09	1F	001E9		BLSSU	24\$	
	04	44	A9		03	E1	001EB		BBC	#3, 68(IRAB), 24\$	2218
		52		40	AE	C0	001F0		ADDL2	RECSZ, LHS	2220
	50	52		18	AE	C1	001F4	24\$:	ADDL3	RRV, LHS, R0	2230
30	AE	0C	AE		09	C5	001F9		MULL3	#9, NEED RRV, 48(SP)	
		50		30	AE	C0	001FF		ADDL2	48(SP), R0	
		1C	AE		50	D1	00203		CMPL	R0, BKTSIZE	
					0F	1A	00207		BGTRU	25\$	
				06	A5	B5	00209		TSTW	6(BKT_ADDR)	2237
					46	12	0020C		BNEQ	30\$	
	41	06	A9		03	E0	0020E		BBS	#3, 6(IRAB), 30\$	2239
	3C	44	A9		03	E1	00213		BBC	#3, 68(IRAB), 30\$	2241
					6E	D5	00218	25\$:	TSTL	LAST	2245
					03	12	0021A		BNEQ	26\$	
				0280	31	0021C			BRW	62\$	
		56			6E	D0	0021F	26\$:	MOVL	LAST, REC_ADDR	2249
		04		2C	AE	E8	00222		BLBS	SAVE_REC_W_LO, 27\$	2251
		A9			08	8A	00226		BICB2	#8, 68(IRAB)	2253
		53			56	D0	0022A	27\$:	MOVL	REC_ADDR, TMP	2264
		56		0E	A5	9E	0022D		MOVAB	14(R5), REC_ADDR	2265

		6E	56	D0	00231	MOVL	REC_ADDR, LAST	2266	
		53	56	D1	00234	CMPL	REC_ADDR, TMP	2268	
			75	13	00237	BEQL	38\$		
		6E	56	D0	00239	MOVL	REC_ADDR, LAST	2271	
	OE	A7	06	E1	0023C	BBC	#6, -28(IDX_DFN), 29\$	2279	
			0000G	30	00241	BSBW	RM\$REC_OVHD	2283	
			01	A046	95	00244	TSTB	1(R0)[REC_ADDR]	
			05	12	00248	BNEQ	29\$		
		0098	56	D0	0024A	MOVL	REC_ADDR, 152(IRAB)	2285	
			0000G	30	0024F	BSBW	RM\$GETNEXT_REC	2288	
			E0	11	00252	BRB	28\$	2268	
	51	52	1C	AE	C5	00254	MULL3	BKTSIZE, LHS, R1	2332
	50	AE	30	AE	C3	00259	SUBL3	48(SP), BKTSIZE, R0	2333
	50	AE		50	C3	0025F	SUBL3	R0, RRV, R0	
		50	10	AE	C4	00264	MULL2	RHS, R0	
24	AE	51	50	C1	00268	ADDL3	R0, R1, DIFFERENCE		
			03	18	0026D	BGEQ	31\$	2338	
			0099	31	0026F	BRW	43\$		
		51	24	AE	D0	00272	MOVL	DIFFERENCE, R1	2346
			03	18	00276	BGEQ	32\$		
		51		51	CE	00278	MNEGL	R1, R1	
		50	28	AE	D0	0027B	MOVL	LAST_DIFF, R0	
			03	18	0027F	BGEQ	33\$		
		50		50	CE	00281	MNEGL	R0, R0	
		50		51	D1	00284	CMPL	R1, R0	
				14	1B	00287	BLEQU	36\$	
		6E		56	D1	00289	CMPL	REC_ADDR, LAST	2350
				09	13	0028C	BEQL	34\$	
		56		6E	D0	0028E	MOVL	LAST, REC_ADDR	2355
	08	AE		56	D1	00291	CMPL	REC_ADDR, -POS_INSERT	2357
				04	1E	00295	BGEQU	35\$	
	44	A9		08	8A	00297	BICB2	#8, 68(IRAB)	2359
				6E	D4	0029B	CLRL	LAST	2362
				6E	D5	0029D	TSTL	LAST	2369
				2A	12	0029F	BNEQ	40\$	
		53		56	D0	002A1	MOVL	REC_ADDR, TMP	2376
		56	OE	A5	9E	002A4	MOVAB	14(R5), REC_ADDR	2377
		6E		56	D0	002A8	MOVL	REC_ADDR, LAST	2378
		53		56	D1	002AB	CMPL	REC_ADDR, TMP	2380
				1B	13	002AE	BEQL	40\$	
		6E		56	D0	002B0	MOVL	REC_ADDR, LAST	2383
	OE	A7		06	E1	002B3	BBC	#6, -28(IDX_DFN), 39\$	2391
				0000G	30	002B8	BSBW	RM\$REC_OVHD	2395
			01	A046	95	002BB	TSTB	1(R0)[REC_ADDR]	
				05	12	002BF	BNEQ	39\$	
		0098		56	D0	002C1	MOVL	REC_ADDR, 152(IRAB)	2397
				0000G	30	002C6	BSBW	RM\$GETNEXT_REC	2400
				E0	11	002C9	BRB	37\$	2380
		50		6E	D0	002CB	MOVL	LAST, R0	2405
				F76F	30	002CE	BSBW	RM\$MOVE_KEY	
4A	A9	56		55	A3	002D1	SUBW3	BKT_ADDR, REC_ADDR, 74(IRAB)	2406
		A9	4A	A9	B1	002D6	CMPL	74(IRAB), 72(IRAB)	2412
				0C	12	002DB	BNEQ	41\$	
		A9	4A	A9	B1	002DD	CMPL	74(IRAB), 76(IRAB)	2415
				05	12	002E2	BNEQ	41\$	
	03	A9		03	E1	002E4	BBC	#3, 68(IRAB), 42\$	2418
				044F	31	002E9	BRW	96\$	

51	5C	03	D0	002EC	42\$:	MOVL	#3, AP	2425	
	50	00B4	CA	3C	002EF	MOVZWL	180(IFAB), R0	2428	
	53	60	B940	3E	002F4	MOVAV	@96(IRAB)[R0], R3		
	50	60	A9	C1	002F9	ADDL3	96(IRAB), R0, R1	2427	
	50	20	A7	9A	002FE	MOVZBL	32(IDX_DFN), R0		
			0000G	30	00302	BSBW	RM\$COMPARE_KEY		
	E1	50	E8	00305		BLBS	R0, 41\$		
		0321	31	00308		BRW	78\$	2431	
04	28	AE	24	AE	D0	0030B	43\$:	2442	
	6E	56	D0	00310		MOVL	DIFFERENCE, LAST_DIFF	2443	
	44	A9	03	E1	00313	MOVL	REC_ADDR, LAST	2445	
	2C	AE	01	D0	00318	BBC	#3, 68(IRAB), 44\$	2447	
	08	AE	56	D1	0031C	44\$:	MOVL	#1, SAVE_REC_W_LO	2457
			03	13	00320	44\$:	CPL	REC_ADDR, POS_INSERT	
			008D	31	00322	45\$:	BEQL	46\$	
F8	44	A9	03	E0	00325	46\$:	BRW	52\$	2459
OA	06	A9	03	E1	0032A		BBS	#3, 68(IRAB), 45\$	2466
	78	A9	1C	A4	D1	0032F	BBC	#3, 6(IRAB), 47\$	2470
				03	12	00334	CPL	28(BDB), 120(IRAB)	
			0C	AE	D7	00336	BNEQ	47\$	
	44	A9	08	88	00339	47\$:	DECL	NEED_RRV	2472
	50		00B4	CA	3C	0033D	BISB2	#8, 68(IRAB)	2479
			60	B940	9F	00342	MOVZWL	180(IFAB), R0	2480
			60	B940	3F	00346	PUSHAB	@96(IRAB)[R0]	
	7E		20	A7	9A	0034A	PUSHAW	@96(IRAB)[R0]	
				0000G	30	0034E	MOVZBL	32(IDX_DFN), -(SP)	
			0C	C0	00351		BSBW	RM\$MOVE	
	04	AE	56	D1	00354		ADDL2	#12, SP	
			03	12	00358		CPL	REC_ADDR, EOB	2488
			FE05	31	0035A	48\$:	BNEQ	49\$	
			51	D4	0035D	49\$:	BRW	15\$	
			0000G	30	0035F		CLRL	R1	2501
53	30	AE	50	D0	00362		BSBW	RM\$REC_OVHD	
29	1C	A7	30	AE	C1	00366	MOVL	R0, REC_OVHD	
		50	06	E1	0036B		ADDL3	REC_OVHD, REC_ADDR, R3	2512
		51	00B4	CA	3C	00370	BBC	#6, 28(IDX_DFN), 50\$	2508
			60	B940	DE	00375	MOVZWL	180(IFAB), R0	2511
				51	DD	0037A	MOVAL	@96(IRAB)[R0], CURR_KEY	
			60	B940	9F	0037C	PUSHL	CURR_KEY	2514
			01	A3	9A	00380	PUSHAB	@96(IRAB)[R0]	2513
	7E			0000G	30	00384	MOVZBL	1(R3), -(SP)	2512
			08	C0	00387		BSBW	RM\$MOVE	
	5E		51	D0	0038A		ADDL2	#8, SP	
	6E		34	BE46	9F	0038D	MOVL	CURR_KEY, (SP)	2515
			F706	30	00391		PUSHAB	@REC_OVHD[REC_ADDR]	
				08	C0	00394	BSBW	RM\$BUILD_KEY	
				03	11	00397	ADDL2	#8, SP	
			53	D0	00399	50\$:	BRB	51\$	2508
	5C		03	D0	0039C	51\$:	MOVL	R3, CURR_KEY	2518
	53		00B4	CA	3C	0039F	MOVL	#3, AP	2519
	53		60	A9	C0	003A4	MOVZWL	180(IFAB), R3	2522
	50		20	A7	9A	003A8	ADDL2	96(IRAB), R3	
				0000G	30	003AC	MOVZBL	32(IDX_DFN), R0	2521
			50	E8	003AF		BSBW	RM\$COMPARE_KEY	
	A8		20	AE	D4	003B2	BLBS	R0, 48\$	
				51	D4	003B5	CLRL	NOT_DUP	2544
			0000G	30	003B7		CLRL	R1	2548
							BSBW	RM\$REC_OVHD	

	14	AE	50	D0	003BA	MOVL	R0, REC_OVHD	...	
	30	AE	51	D0	003BE	MOVL	R1, 48(SP)	...	
OF	1C	A7	06	E1	003C2	BBC	#6, 28(IDX_DFN), 53\$...	2553
50	14	AE	01	C1	003C7	ADDL3	#1, REC_OVHD, R0	...	2555
			6046	95	003CC	TSTB	(R0)[REC_ADDR]	...	
			05	12	003CF	BNEQ	53\$...	
	0098	C9	56	D0	003D1	MOVL	REC_ADDR, 152(IRAB)	...	2557
		53	44	A9	90 003D6	MOVB	68(IRAB), TMP	...	2568
	44	A9	08	8A	003DA	BICB2	#8, 68(IRAB)	...	2569
		50	56	D0	003DE	MOVL	REC_ADDR, R0	...	2570
			F65C	30	003E1	BSBW	RMSMOVE_KEY	...	
	44	A9	53	90	003E4	MOVB	TMP, 68(IRAB)	...	2571
	04	AE	56	D1	003E8	CMPL	REC_ADDR, EOB	...	2583
			79	13	003EC	BEQL	60\$...	
		5C	03	D0	003EE	MOVL	#3, AP	...	2587
			0000G	30	003F1	BSBW	RMSRECORD_VBN	...	2589
		50	1C	A4	D1 003F4	CMPL	28(BDB), R0	...	
			03	12	003F8	BNEQ	55\$...	
50			0C	AE	D7 003FA	DECL	NEED_RRV	...	2591
56		56	14	AE	C1 003FD	ADDL3	REC_OVHD, REC_ADDR, R0	...	2593
		50	30	AE	C1 00402	ADDL3	S_REC_SIZE, R0, REC_ADDR	...	
	04	AE	56	D1	00407	CMPL	REC_ADDR, EOB	...	2595
			5A	13	0040B	BEQL	60\$...	
			51	D4	0040D	CLRL	R1	...	2599
			0000G	30	0040F	BSBW	RMSREC_OVHD	...	
	14	AE	50	D0	00412	MOVL	R0, REC_OVHD	...	
	30	AE	51	D0	00416	MOVL	R1, 48(SP)	...	
08	1C	A7	06	E1	0041A	BBC	#6, 28(IDX_DFN), 56\$...	2601
			14	BE46	95 0041F	TSTB	@REC_OVHD[REC_ADDR]	...	2605
			21	13	00423	BEQL	58\$...	
			1B	11	00425	BRB	57\$...	2607
		5C	03	D0	00427	MOVL	#3, AP	...	2611
		53	00B4	CA	3C 0042A	MOVZWL	180(IFAB), R3	...	2614
		53	60	A9	C0 0042F	ADDL2	96(IRAB), R3	...	
51		56	14	AE	C1 00433	ADDL3	REC_OVHD, REC_ADDR, R1	...	2613
		50	20	A7	9A 00438	MOVZBL	32(IDX_DFN), R0	...	
			0000G	30	0043C	BSBW	RMSCOMPARE_KEY	...	
		04	50	E9	0043F	BLBC	R0, 58\$...	
	20	AE	01	D0	00442	MOVL	#1, NOT_DUP	...	2617
	08	AE	56	D1	00446	CMPL	REC_ADDR, POS_INSERT	...	2625
			17	12	0044A	BNEQ	59\$...	
		17	20	AE	E8 0044C	BLBS	NOT_DUP, 60\$...	2627
		A9	03	E1	00450	BBC	#3, 6(IRAB), 59\$...	2629
0E	06	A9	08	88	00455	BISB2	#8, 68(IRAB)	...	2633
	44	A9	1C	A4	D1 00459	CMPL	28(BDB), 120(IRAB)	...	2635
	78	A9	03	12	0045E	BNEQ	59\$...	
			0C	AE	D7 00460	DECL	NEED_RRV	...	2637
		81	20	AE	E9 00463	BLBC	NOT_DUP, 54\$...	2644
	08	AE	56	D1	00467	CMPL	REC_ADDR, POS_INSERT	...	2654
			2F	12	0046B	BNEQ	61\$...	
		5C	03	D0	0046D	MOVL	#3, AP	...	2661
		50	00B4	CA	3C 00470	MOVZWL	180(IFAB), R0	...	2664
		53	60	B940	3E 00475	MOVAV	@96(IRAB)[R0], R3	...	
51		50	60	A9	C1 0047A	ADDL3	96(IRAB), R0, R1	...	2663
		50	20	A7	9A 0047F	MOVZBL	32(IDX_DFN), R0	...	
			0000G	30	00483	BSBW	RMSCOMPARE_KEY	...	
		13	50	E8	00486	BLBS	R0, 61\$...	

44	A9	0A	44	A9	08	88	00489	BISB2	#8, 68(IRAB)	2668	
			06	A9	03	E1	0048D	BBC	#3, 6(IRAB), 61\$	2670	
			78	A9	A4	D1	00492	CMPL	28(BDB), 120(IRAB)	2672	
					03	12	00497	BNEQ	61\$		
					0C	AE	D7	00499	DECL	NEED_RRV	2674
					FCC3	31	0049C	BRW	15\$	2147	
					24	AE	D4	0049F	CLRL	FIRST_KEY_EXPANSION	2704
44	A9	02	44	01	02	F0	004A2	INSV	#2, #1, #2, 68(IRAB)	2716	
			50	A9	08	8A	004A8	BICB2	#8, 68(IRAB)	2717	
				50	CA	3C	004AC	MOVZWL	180(IFAB), R0	2727	
					60	B940	9F	004B1	PUSHAB	@96(IRAB)[R0]	
					60	B940	3F	004B5	PUSHAW	@96(IRAB)[R0]	
				7E	20	A7	9A	004B9	MOVZBL	32(IDX_DFN), -(SP)	
					0000G	30	004BD	BSBW	RM\$MOVE		
				5E	0C	C0	004C0	ADDL2	#12, SP		
			20	AE	0E	A5	9E	004C3	MOVAB	14(R5), 32(SP)	2732
			56		20	AE	D0	004C8	MOVL	32(SP), REC_ADDR	
					51	D4	004CC	CLRL	R1	2745	
					0000G	30	004CE	BSBW	RM\$REC_OVHD		
			28	AE	50	D0	004D1	MOVL	R0, REC_OVHD		
			30	AE	51	D0	004D5	MOVL	R1, 48(SP)		
				56	28	AE	C1	004D9	ADDL3	REC_OVHD, REC_ADDR, R1	2754
51			1C	A7	06	E1	004DE	BBC	#6, 28(IDX_DFN), 64\$	2750	
18			50		00B4	CA	3C	004E3	MOVZWL	180(IFAB), R0	2753
			14	AE	60	B940	DE	004E8	MOVAL	@96(IRAB)[R0], CURR_KEY	
					14	AE	DD	004EE	PUSHL	CURR_KEY	2754
					51	DD	004F1	PUSHL	R1		
					F5A4	30	004F3	BSBW	RM\$BUILD_KEY		
				5E	08	C0	004F6	ADDL2	#8, SP		
					04	11	004F9	BRB	65\$	2750	
			14	AE	51	D0	004FB	MOVL	R1, CURR_KEY	2759	
			5C		03	D0	004FF	MOVL	#3, AP	2760	
			0C	AE	00B4	CA	3C	00502	MOVZWL	180(IFAB), 12(SP)	2762
			50		0C	AE	D0	00508	MOVL	12(SP), R0	
			53		60	B940	3E	0050C	MOVAV	@96(IRAB)[R0], R3	
			50		20	A7	9A	00511	MOVZBL	32(IDX_DFN), R0	2761
			51		14	AE	D0	00515	MOVL	CURR_KEY, R1	
					0000G	30	00519	BSBW	RM\$COMPARE_KEY		
			2C	AE	50	D0	0051C	MOVL	R0, STATUS		
			42		2C	AE	E9	00520	BLBC	STATUS, 70\$	2765
			08	AE	56	D1	00524	CMPL	REC_ADDR, POS_INSERT	2769	
					15	1E	00528	BGEQU	66\$		
			50		0C	AE	D0	0052A	MOVL	12(SP), R0	2771
					60	B940	9F	0052E	PUSHAB	@96(IRAB)[R0]	
					18	AE	DD	00532	PUSHL	CURR_KEY	
				7E	20	A7	9A	00535	MOVZBL	32(IDX_DFN), -(SP)	
					0000G	30	00539	BSBW	RM\$MOVE		
				5E	0C	C0	0053C	ADDL2	#12, SP		
			04	AE	56	D1	0053F	CMPL	REC_ADDR, EOB	2773	
					05	13	00543	BEQL	67\$		
					2C	AE	D5	00545	TSTL	STATUS	2775
					0F	18	00548	BGEQ	69\$		
				50	48	A9	3C	0054A	MOVZWL	72(IRAB), R0	2783
			4C	A9	50	B0	0054E	MOVW	R0, 76(IRAB)		
			4A	A9	50	B0	00552	MOVW	R0, 74(IRAB)		
					01E2	31	00556	BRW	96\$	2784	
50			56		28	AE	C1	00559	ADDL3	REC_OVHD, REC_ADDR, R0	2790

56	50	30	AE	C1	0055E	ADDL3	S_REC_SIZE, R0, REC_ADDR	2741		
	52	FF66	31	00563	BRW	63\$		2797		
		56	D0	00566	70\$:	MOVL	REC_ADDR, LHS	2806		
		2C	AE	D4	00569	CLRL	NOT_DUP	2807		
		51	D4	0056C	CLRL	R1				
		0000G	30	0056E	BSBW	RMSREC_OVHD				
	14	AE	50	D0	00571	MOVL	R0, REC_OVHD			
	30	AE	51	D0	00575	MOVL	R1, 48(SP)			
0A	1C	A7	06	E1	00579	BBC	#6, 28(IDX_DFN), 71\$	2812		
50	14	AE	01	C1	0057E	ADDL3	#1, REC_OVHD, R0	2814		
	24	AE	6046	9A	00583	MOVZBL	(R0)[REC_ADDR], FIRST_KEY_EXPANSION			
	28	AE	14	BE46	9E	00588	71\$: MOVAB	@REC_OVHD[REC_ADDR], 40(SP)	2819	
56	28	AE	30	AE	C1	0058E	72\$: ADDL3	S_REC_SIZE, 40(SP), REC_ADDR	2820	
	04	AE	56	D1	00594	CMPL	REC_ADDR, EOB		2824	
			42	13	00598	BEQL	76\$			
			51	D4	0059A	CLRL	R1			
			0000G	30	0059C	BSBW	RMSREC_OVHD			
	14	AE	50	D0	0059F	MOVL	R0, REC_OVHD			
	30	AE	51	D0	005A3	MOVL	R1, 48(SP)			
	28	AE	14	BE46	9E	005A7	MOVAB	@REC_OVHD[REC_ADDR], 40(SP)	2830	
07	1C	A7	06	E1	005AD	BBC	#6, 28(IDX_DFN), 73\$	2826		
			28	BE	95	005B2	TSTB	@40(SP)	2830	
			21	13	005B5	BEQL	75\$			
			1B	11	005B7	BRB	74\$		2832	
			03	D0	005B9	73\$:	MOVL	#3, AP	2837	
	5C		00B4	CA	3C	005BC	MOVZWL	180(IFAB), R0	2840	
	50		60	B940	3E	005C1	MOVAB	@96(IRAB)[R0], R3		
	53		20	A7	9A	005C6	MOVZBL	32(IDX_DFN), R0	2839	
	50		28	AE	D0	005CA	MOVL	40(SP), R1		
	51		0000G	30	005CE	BSBW	RMSCOMPARE_KEY			
			50	E9	005D1	BLBC	R0, 75\$			
	2C	AE	01	D0	005D4	74\$:	MOVL	#1, NOT_DUP	2843	
	B2		2C	AE	E9	005D8	75\$:	BLBC	NOT_DUP, 72\$	2848
	10	AE	56	D0	005DC	76\$:	MOVL	REC_ADDR, RHS	2851	
	52		08	AE	D1	005E0	CMPL	POS_INSERT, LHS	2861	
			05	1B	005E4	BLEQU	77\$			
	44	A9	80	8F	88	005E6	BISB2	#128, 68(IRAB)	2863	
18	10	AE	52	C3	005EB	77\$:	SUBL3	LHS, RHS, RRV	2865	
50	18	AE	40	AE	C1	005F1	ADDL3	RECSZ, RRV, R0	2866	
	18	AE	24	BE40	9E	005F7	MOVAB	@FIRST_KEY_EXPANSION[R0], RRV		
	1C	AE	18	AE	D1	005FD	CMPL	RRV, BRTSIZE	2868	
			4D	1E	00602	BGEQU	82\$			
			55	A3	00604	SUBW3	BKT_ADDR, LHS, 74(IRAB)	2889		
4A	10	AE	55	A3	00609	SUBW3	BKT_ADDR, RHS, 76(IRAB)	2890		
4C	04	AE	10	AE	D1	0060F	CMPL	RHS, EOB	2892	
			1C	12	00614	BNEQ	79\$			
44	02		01	F0	00616	INSV	#1, #1, #2, 68(IRAB)	2895		
			20	AE	52	D1	0061C	CMPL	LHS, 32(SP)	2897
			2D	12	00620	BNEQ	81\$			
	4C	A9	4E	A9	B0	00622	MOVW	78(IRAB), 76(IRAB)	2900	
	4A	A9	48	A9	B0	00627	MOVW	72(IRAB), 74(IRAB)	2901	
	44	A9	10	88	0062C	78\$:	BISB2	#16, 68(IRAB)	2902	
			75	11	00630	BRB	86\$		2897	
		0E	4A	A9	B1	00632	79\$:	CMPW	74(IRAB), #14	2909
			05	12	00636	BNEQ	80\$			
	44	A9	40	8F	88	00638	BISB2	#64, 68(IRAB)	2911	
	08	AE	10	AE	D1	0063D	80\$:	CMPL	RHS, POS_INSERT	2917

				63	1F	00642	BLSSU	86\$		
	48	A9	4A	A9	B1	00644	CMPW	74(IRAB), 72(IRAB)	2918	
				5C	13	00649	BEQL	86\$		
	44	A9		08	88	0064B	BISB2	#8, 68(IRAB)	2920	
				56	11	0064F	BRB	86\$	2923	
			44	A9	95	00651	TSTB	68(IRAB)	2935	
				03	19	00654	BLSS	83\$		
				0086	31	00656	BRW	89\$		
	08	AE	10	AE	D1	00659	CMPL	RHS, POS_INSERT	2937	
				7F	1B	0065E	BLEQU	89\$		
50		52	08	AE	C3	00660	SUBL3	POS_INSERT, LHS, R0	2941	
		50	18	AE	C0	00665	ADDL2	RRV, R0		
	1C	AE		50	D1	00669	CMPL	R0, BKTSIZE		
				3A	1E	0066D	BGEQU	87\$		
	44	A9		10	88	0066F	BISB2	#16, 68(IRAB)	2959	
	4A	A9	48	A9	B0	00673	MOVW	72(IRAB), 74(IRAB)	2960	
	04	AE	10	AE	D1	00678	CMPL	RHS, EOB	2962	
				08	12	0067D	BNEQ	84\$		
44	A9	02	01	01	F0	0067F	INSV	#1, #1, #2, 68(IRAB)	2964	
				06	11	00685	BRB	85\$		
	4C	A9	10	AE	A3	00687	SUBW3	BKT_ADDR, RHS, 76(IRAB)	2966	
		56		52	D0	0068D	MOVL	LHS, REC_ADDR	2968	
		50	00B4	CA	3C	00690	MOVZWL	180(IFABT), R0	2969	
			60	B940	9F	00695	PUSHAB	@96(IRAB)[R0]		
			60	B940	3F	00699	PUSHAW	@96(IRAB)[R0]		
		7E	20	A7	9A	0069D	MOVZBL	32(IDX DFN), -(SP)		
				0000G	30	006A1	BSBW	RMSMOVE		
		5E		0C	C0	006A4	ADDL2	#12, SP		
				72	11	006A7	BRB	91\$	2970	
50	08	AE	10	AE	C3	006A9	SUBL3	RHS, POS_INSERT, R0	2977	
		50	18	AE	C0	006AF	ADDL2	RRV, R0		
	1C	AE		50	D1	006B3	CMPL	R0, BKTSIZE		
				26	1E	006B7	BGEQU	89\$		
4A	A9	52		55	A3	006B9	SUBW3	BKT_ADDR, LHS, 74(IRAB)	2998	
	4C	A9	48	A9	B0	006BE	MOVW	72(IRAB), 76(IRAB)	2999	
		0E	4A	A9	B1	006C3	CMPW	74(IRAB), #14	3001	
				05	12	006C7	BNEQ	88\$		
44	A9	40	8F	88	006C9		BISB2	#64, 68(IRAB)	3003	
44	A9		08	88	006CE	88\$:	BISB2	#8, 68(IRAB)	3005	
04	AE	10	AE	D1	006D2		CMPL	RHS, EOB	3007	
				52	1F	006D7	BLSSU	93\$		
44	A9		20	88	006D9		BISB2	#32, 68(IRAB)	3014	
				5C	11	006DD	BRB	96\$	3016	
		50	48	A9	3C	006DF	MOVZWL	72(IRAB), R0	3033	
4C	A9			50	B0	006E3	MOVW	R0, 76(IRAB)		
4A	A9			50	B0	006E7	MOVW	R0, 74(IRAB)		
			44	A9	95	006EB	TSTB	68(IRAB)	3035	
				1E	18	006EE	BGEQ	90\$		
44	A9			10	88	006F0	BISB2	#16, 68(IRAB)	3038	
	56			52	D0	006F4	MOVL	LHS, REC_ADDR	3039	
	50	00B4		CA	3C	006F7	MOVZWL	180(IFABT), R0	3040	
		60	B940	9F	006FC		PUSHAB	@96(IRAB)[R0]		
		60	B940	3F	00700		PUSHAW	@96(IRAB)[R0]		
		7E	20	A7	9A	00704	MOVZBL	32(IDX DFN), -(SP)		
				0000G	30	00708	BSBW	RMSMOVE		
		5E		0C	C0	0070B	ADDL2	#12, SP		
04	AE	08	AE	D1	0070E	90\$:	CMPL	POS_INSERT, EOB	3043	

RM3SPLUDR
V04-000

RM\$SPLIT_UDR_3

D 5
16-Sep-1984 02:03:28
14-Sep-1984 13:01:40

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[RMS.SRC]RM3SPLUDR.B32;1

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44	A9	02	01	08	12	00713	BNEQ	92\$:	3045
				01	F0	00715	INSV	#1, #2, 68(IRAB)	:	
				1E	11	0071B	BRB	96\$:	
		10	AE	08	AE	D1 0071D	CMPL	POS_INSERT, RHS	:	3048
				17	1E	00722	BGEQU	96\$:	
		04	AE	10	AE	D1 00724	CMPL	RHS, EOB	:	3052
				06	1E	00729	BGEQU	94\$:	
		44	A9		06	88 0072B	BISB2	#6, 68(IRAB)	:	3054
					04	11 0072F	BRB	95\$:	
		44	A9		20	88 00731	BISB2	#32, 68(IRAB)	:	3056
4E	A9	10	AE		55	A3 00735	SUBW3	BKT_ADDR, RHS, 78(IRAB)	:	3058
			0E	4A	A9	B1 0073B	CMPL	74(IRAB), #14	:	3070
					0A	12 0073F	BNEQ	97\$:	
	05	44	A9		03	E0 00741	BBS	#3, 68(IRAB), 97\$:	3072
		44	A9	40	8F	88 00746	BISB2	#64, 68(IRAB)	:	3074
			5E		34	C0 0074B	ADDL2	#52, SP	:	3078
					0C	BA 0074E	POPR	#^M<R2,R3>	:	
					05	00750	RSB		:	

; Routine Size: 1873 bytes, Routine Base: RM\$RMS3 + 05C0

: 3020	3079	1
: 3021	3080	1 END
: 3022	3081	1
: 3023	3082	0 ELUDOM

PSECT SUMMARY

Name	Bytes	Attributes
RM\$RMS3	3345	NOVEC,NOWRT, RD , EXE,NOSHR, GBL, REL, CON, PIC,ALIGN(2)

Library Statistics

File	----- Total	Symbols Loaded	----- Percent	Pages Mapped	Processing Time
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COMMAND QUALIFIERS

; BLISS/CHECK=(FIELD,INITIAL,OPTIMIZE)/LIS=LIS\$:RM3SPLUDR/OBJ=OBJ\$:RM3SPLUDR MSRC\$:RM3SPLUDR/UPDATE=(ENH\$:RM3SPLUDR)

RM3SPLUDR
V04-000

RMSSPLIT_UDR_3

E 5
16-Sep-1984 02:03:28
14-Sep-1984 13:01:40

VAX-11 Bliss-32 V4.0-742
[RMS.SRC]RM3SPLUDR.B32;1

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; Size: 3345 code + 0 data bytes
; Run Time: 01:25.9
; Elapsed Time: 02:43.7
; Lines/CPU Min: 2152
; Lexemes/CPU-Min: 13733
; Memory Used: 634 pages
; Compilation Complete

RM
V0

0327 AH-BT13A-SE
VAX/VMS V4.0

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