



\*\*FILE\*\*ID\*\*RM3MISPUT

K 16

RRRRRRRR	MM	MM	333333	MM	MM	IIIIII	SSSSSSSS	PPPPPPPP	UU	UU	TTTTTTTTTT
RRRRRRRR	MM	MM	333333	MM	MM	IIIIII	SSSSSSSS	PPPPPPPP	UU	UU	TTTTTTTTTT
RR RR	RR	MMMM	MMMM	33	33	IIIIII	SS	PP	PP	UU	TT
RR RR	RR	MMMM	MMMM	33	33	IIIIII	SS	PP	PP	UU	TT
RR RR	RR	MM MM	MM	33	33	IIIIII	SS	PP	PP	UU	TT
RR RR	RR	MM MM	MM	33	33	IIIIII	SS	PP	PP	UU	TT
RRRRRRRR	MM	MM	33	MM	MM	IIIIII	SSSSSS	PPPPPPPP	UU	UU	TT
RRRRRRRR	MM	MM	33	MM	MM	IIIIII	SSSSSS	PPPPPPPP	UU	UU	TT
RR RR	RR	MM	MM	33	MM	IIIIII	SS	PP	UU	UU	TT
RR RR	RR	MM	MM	33	MM	IIIIII	SS	PP	UU	UU	TT
RR RR	RR	MM	MM	33	33	IIIIII	SS	PP	UU	UU	TT
RR RR	RR	MM	MM	333333	MM	IIIIII	SSSSSSSS	PP	UUUUUUUUUU	UUUUUUUUUU	TT
RR RR	RR	MM	MM	333333	MM	IIIIII	SSSSSSSS	PP	UUUUUUUUUU	UUUUUUUUUU	TT

LL	IIIIII	SSSSSSSS
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LL	II	SS
LLLLLLLL	IIIIII	SSSSSSSS
LLLLLLLL	IIIIII	SSSSSSSS

```
1 0001 0
2 0002 0 MODULE RM3MISPUT (LANGUAGE (BLISS32) ,
3 0003 0 IDENT = 'V04-000'
4 0004 0 ) =
5 0005 1 BEGIN
6 0006 1
7 0007 1 *****
8 0008 1 *
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26 0026 1 *
27 0027 1 *
28 0028 1 *****
29 0029 1
30 0030 1 ++
31 0031 1
32 0032 1 FACILITY: RMS32 INDEX SEQUENTIAL FILE ORGANIZATION
33 0033 1
34 0034 1 ABSTRACT:
35 0035 1 Miscellaneous put routines
36 0036 1
37 0037 1
38 0038 1 ENVIRONMENT:
39 0039 1
40 0040 1 VAX/VMS OPERATING SYSTEM
41 0041 1
42 0042 1 --
43 0043 1
44 0044 1
45 0045 1 AUTHOR: Christian Saether CREATION DATE: 6-JUL-78 10:56
46 0046 1
47 0047 1
48 0048 1 MODIFIED BY:
49 0049 1
50 0050 1 V03-012 MCN0010 Maria del C. Nasr 04-Apr-1983
51 0051 1 Change linkage of RMSCOMPRESS_KEY.
52 0052 1
53 0053 1 V03-011 MCN0009 Maria del C. Nasr 15-Mar-1983
54 0054 1 More linkages reorganization
55 0055 1
56 0056 1 V03-010 MCN0008 Maria del C. Nasr 22-Feb-1983
57 0057 1 RMSBLD_IDX_REC and RMSMOVE_IN_VBN are not global routines
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58       0058 1 | anymore, make them local. Also, make local linkages more  
59       0059 1 | general, and reorganize linkages.  
60       0060 1 |  
61       0061 1 | V03-009 TMK0006 Todd M. Katz 10-Sep-1982  
62       0062 1 | Add support for Prologue 3 SIDs. The routines which have been  
63       0063 1 | modified or rewritten are: RMSRECORD\_SIZE, RMSADD\_TO\_ARRAY,  
64       0064 1 | RMSBLD\_NEW\_SIDR, and RMSINS\_REC.  
65       0065 1 |  
66       0066 1 | I also reformatted the whole module, putting the routines in  
67       0067 1 | alphabetical order, and redid some of the linkages.  
68       0068 1 |  
69       0069 1 | V03-008 KBT0222 Keith B. Thompson 23-Aug-1982  
70       0070 1 | Reorganize psects  
71       0071 1 |  
72       0072 1 | V03-007 TMK0005 Todd M. Katz 02-Jun-1982  
73       0073 1 | Implement the RMS cluster solution for next record processing.  
74       0074 1 | There is no longer any need to reference the NRP cell to obtain  
75       0075 1 | the RFA address of the current primary data record because that  
76       0076 1 | information is now kept locally in the IRAB as part of the next  
77       0077 1 | record context instead of in the NRP cell which has been  
78       0078 1 | deleted. Also refer to the fields IRBS\_PUTUP\_VBN and  
79       0079 1 | IRBSW\_PUTUP\_ID when referring to the new/changed primary data  
80       0080 1 | which requires the SIDR entry.  
81       0081 1 |  
82       0082 1 | V03-006 KBT0066 Keith B. Thompson 17-Jun-1982  
83       0083 1 | Remove rm\$move\_sig\_cnt routine  
84       0084 1 |  
85       0085 1 | V03-005 TMK0004 Todd M. Katz 11-May-1982  
86       0086 1 | Added code for the support of two index bucket split cases  
87       0087 1 | which are currently not supported. At the present time when  
88       0088 1 | one of these split cases occur, the routines in this module are  
89       0089 1 | never called, but instead, a higher routine returns an  
90       0090 1 | alternate success status of index not updated. These two split  
91       0091 1 | cases are: two-pass two-bucket with empty bucket split case  
92       0092 1 | and two-pass multi-bucket with empty bucket split case. Support  
93       0093 1 | has been added for all prologue versions. The routines which  
94       0094 1 | have been modified are RMSRECORD\_SIZE, RMSSHFT VBNS,  
95       0095 1 | RMSBLD\_IDX\_REC, and RMSV3\_IDX\_REC. These routines have been  
96       0096 1 | modified so that they can perform their individual functions  
97       0097 1 | during each of the two individual passes required to correctly  
98       0098 1 | update the index during these two-pass empty bucket index  
99       0099 1 | bucket split cases. During the first pass, these routines  
100      0100 1 | function on the old (left) index bucket, and during the second  
101      0101 1 | pass these routines function on the new (right) index bucket.  
102      0102 1 |  
103      0103 1 | V03-004 TMK0003 Todd M. Katz 26-Apr-1982  
104      0104 1 | Redid the way all empty bucket splits are done for prologue  
105      0105 1 | 3 files. There were many errors in the way they were being done  
106      0106 1 | especially in those cases when a decision was made not to swing  
107      0107 1 | the level one index pointer from the old (left) bucket to a  
108      0108 1 | new (middle) bucket because the downpointer was not currently  
109      0109 1 | pointing to the left bucket. Changes were made to the routine  
110      0110 1 | RMSRECORD\_SIZE so that the size of the new index record(s)  
111      0111 1 | could be correctly computed for all empty bucket involving  
112      0112 1 | index updates. Changes were made to the routine RMSSHFT VBNS  
113      0113 1 | so that the spreading apart of VBNs during index updating was  
114      0114 1 | performed correctly for all index updates involving empty data

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126 0126 1 |  
127 0127 1 |  
128 0128 1 |  
129 0129 1 |  
130 0130 1 |  
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163 0163 1 |  
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167 0167 1 |  
168 0168 1 |  
169 0169 1 |  
170 0170 1 |  
171 0171 1 |

buckets. Finally, the routine RMSV3\_IDX\_REC was modified to correctly perform all index updates involving empty buckets, and most especially those updates were no pointer is to be swung because the level one index down pointer does not point to the old (leftmost) data bucket. These changes will be included as a patch on the V3.1 update floppy.

V03-003 TMK0002 Todd M. Katz 13-Apr-1982  
Fixed a bug in RMSRECORD SIZE. This bug only showed up when a multi-bucket data level split of a prolog 3 file with index compression required updating of the level 1 index with two index records. The size of these two index records was being incorrectly computed. I was referring to keybuffer 5 (containing the compressed key of the first record) in one instance when I should have been referring to keybuffer 4 (containing the compressed key of the second record), and thus the combined size of the two index records was being incorrectly computed.

V03-002 LJA0007 Laurie Anderson 25-Mar-1982  
Change KBUFSZ to reference a macro when computing buffer size and make IFBS\$B\_KBUFSZ a word, now: IFBS\$W\_KBUFSZ.

V03-001 TMK0001 Todd M. Katz 03-Mar-1982  
Fixed a bug in the determination of the record size of a two-bucket with empty bucket bucket split case involving prolog 3 index buckets with compressed keys. In such a case there is no key to be inserted, but due to insufficient checks the routine RMSV3KEY\_SZ was being called as if there was a key. This routine thinks there is a key in keybuffer 2 and attempts to move it into keybuffer 5 so it can compress it. Since there isn't a key, it ends up moving garbage into the keybuffer and into the RLB which follows it guaranteeing an access violation the next time RM\_UNLOCK is called. The fix is to make sure that this routine will never be called for this particular prolog 3 bucket split case.

V02-022 KPL0001 Peter Lieberwirth 2-Mar-1982  
Fix Linkage from v02-021

V02-021 TMK0005 Todd M. Katz 02-Mar-1982  
Made several changes to the routines in this module:

1. RMSV3KEY\_SZ is no longer a global routine.
2. Added a symbol to RMSFOOLED\_YUH so that reprobining of the RAB would not be neccessairy.
3. Added support for rear-end truncation of keys in prolog 3 buckets containing compressed keys. This included changes in the way the size of two compressed keys (from a multibucket split) is computed in RMSRECORD SIZE, changes in the way the size of a single key is determined in RMSV3KEY\_SZ, and changes in the way in which keys are inserted into prolog 3 buckets (with compressed keys) in RMSADD\_V3KEY. All these changes reflect the fact that compressed keys are now rear-end truncated as well as front-end compressed.

172 0172 1 |  
173 0173 1 |  
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177 0177 1 |  
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180 0180 1 |  
181 0181 1 |  
182 0182 1 |  
183 0183 1 |  
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188 0188 1 |  
189 0189 1 |  
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191 0191 1 |  
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208 0208 1 |  
209 0209 1 |  
210 0210 1 |  
211 0211 1 |  
212 0212 1 |  
213 0213 1 |  
214 0214 1 |  
215 0215 1 |  
216 0216 1 |  
217 0217 1 |  
218 0218 1 |  
219 0219 1 |  
220 0220 1 |  
221 0221 1 |  
222 0222 1 |  
223 0223 1 |  
224 0224 1 |  
225 0225 1 |  
226 0226 1 |  
227 0227 1 |  
228 0228 1 |

V02-020 TMK0004 Todd M. Katz 12-Feb-1982  
Add checks for a multibucket split in RMS\$RECORD SIZE and  
RMS\$BLD\_IDX REC in all instances when it isn't absolutely  
clear whether RMS is currently processing a two-pass two-bucket  
split case or a two-pass multibucket split case (where both  
records are to go in the same bucket).

V02-019 TMK0003 Todd M. Katz 08-Feb-1982  
Document what is going on in RMS\$RECORD SIZE so that someone  
besides me might possibly understand what is going on from  
the comments. The current comments are both confusing and  
wrong.

V02-018 TMK0002 Todd M. Katz 30-Jan-1982  
Made many changes to several of these miscellaneous routines  
inorder to fix a variety of bugs, and to add functionality  
that was lacking. I will list the reasons for the changes  
made, and how the routines were changed.

1. Added REC\_SZ as an arguement to both RMS\$BLD\_IDX\_REC  
and RMS\$V3\_IDX REC. It was require in the former to pass  
it to the latter, and in the latter to fix the prolog 3  
multibucket split trailing key recompression bug described  
below.
2. Fixed a prolog 3 multibucket split key recompression bug.  
When a multibucket split occurs, and the index must be  
updated with two keys, the key at the insertion point  
must have its front compression redetermined. However,  
the key recompression routine, RMS\$RECOMPR\_KEY, is unable  
to recompute the front compression of a key if its  
current front compression is not equal to the front  
compression of the key preceeding it. If it has greater  
compression the index is out-of-order, and if it has less,  
it assumes there is no need for recompression. Well, if  
the two keys BB AND BC were to be inserted between AA and  
BD, the trailing key BD was not being recompressed because  
its current fron compression (0) was less than the  
compression of the record preceeding it (1) even though  
it should be recompressed to have a new front compression  
value of 1. The fix was that on a multibucket split, the  
first key is inserted and used to recompress what will  
become the trailing key, and then the second key is inserted  
and used to recompress the trailing key. This required  
changes to RMS\$ADD\_V3KEY, so that recompression is NOT  
performed if the key being inserted is the first of two  
keys, and changes to RMS\$V3\_IDX REC, to recompress what  
will become the trailing key after the first of two keys  
is inserted, based on the first key that is inserted. In  
addition, I also added the change such that if the first  
of two compressed keys to be inserted is zero front  
compressed, then the pointer to the last noncompressed key  
is incremented to this key to facilitate the determination  
of the front compression of the next key when it is  
inserted.

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230 0230 1 |  
231 0231 1 |  
232 0232 1 |  
233 0233 1 |  
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251 0251 1 |  
252 0252 1 |  
253 0253 1 |  
254 0254 1 |  
255 0255 1 |  
256 0256 1 |  
257 0257 1 |  
258 0258 1 |  
259 0259 1 |  
260 0260 1 |  
261 0261 1 |  
262 0262 1 |  
263 0263 1 |  
264 0264 1 |  
265 0265 1 |  
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271 0271 1 |  
272 0272 1 |  
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277 0277 1 |  
278 0278 1 |  
279 0279 1 |  
280 0280 1 |  
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282 0282 1 |  
283 0283 1 |  
284 0284 1 |  
285 0285 1 |
3. Added code to restrict when trailing key recompression is required. If one or two keys are inserted into a prolog 3 index bucket containing compressed keys at the beginning of the bucket's freespace, there is no need to redetermine the front compression of the trailing key because there is no trailing key! Formerly, the front compression of the trailing key was blindly determined as part of the process of inserting any compressed key regardless of whether there was any key that trailed the key inserted or not.
  4. Changed the decision by which the keys and VBNs in prolog 3 index buckets are shifted in order to make room for the insertion of a new key-VBN pair. At the present time, if the key/VBN is to be added as the high order key in the bucket, neither the keys nor the VBNs are shifted since both "grow" inward towards the center. An additional check is added to also not shift the keys or VBNs, if the number of bytes by which they are to be shifted "inwards" is 0. This will occur when the split point of an index bucket matches the insertion point, and just the VBN of the new lower order key of the new (right) bucket is to be updated to its new value.
  5. Added code necessary for support of a Prolog 1 & 2 split case. When the insertion point of a multibucket split is found to be at the split point, but there is sufficient room in an index bucket for only two index records, and the insertion point happens to come at the bucket's beginning, both of the old index records must be moved out of the old bucket to make room for the two new index records which are both put in the old bucket. This index bucket split case required code to be added to RM\$RECORD\_SIZE and RM\$BLD\_IDX\_REC. As this split case is handled by making two passes, one for insertion of the two new index records in the old index bucket, and the second to update a VBN pointer in the low order key of the new bucket, both of these routines had to be modified, to take the proper action for each of the passes of this case. RM\$BLD\_IDX\_REC oversees addition of the both keys and a VBN during the first pass, and updates a VBN during the second, while RM\$RECORD\_SIZE makes sure to return the proper size of the index record to be added during each pass.
  6. There are three bucket split cases for when the insertion point and the bucket split point are the same and no empty buckets are involved. For all three cases, two passes must be made, and for all three cases no code existed for Prolog 3 buckets. Changes were made to RM\$RECORD\_SIZE, RM\$SHFT\_VBN and RM\$V3\_IDX\_REC to support these three bucket split cases. The changes made to RM\$V3\_IDX\_REC allow the routine to recognize these three two-pass bucket split cases, which of the two passes is in progress, and to take the appropriate action. As the VBNs in prolog 3 buckets are ordered from right to left and "grow inwards", if a index record with a key value less than the current high key value of the bucket is inserted, the VBNs must be shifted to make room for the

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298 0298 1 |  
299 0299 1 |  
300 0300 1 |  
301 0301 1 |  
302 0302 1 |  
303 0303 1 |  
304 0304 1 |  
305 0305 1 |  
306 0306 1 |  
307 0307 1 |  
308 0308 1 |  
309 0309 1 |  
310 0310 1 |  
311 0311 1 |  
312 0312 1 |  
313 0313 1 |  
314 0314 1 |  
315 0315 1 |  
316 0316 1 |  
317 0317 1 |  
318 0318 1 |  
319 0319 1 |  
320 0320 1 |  
321 0321 1 |  
322 0322 1 |  
323 0323 1 |  
324 0324 1 |  
325 0325 1 |  
326 0326 1 |  
327 0327 1 |  
328 0328 1 |  
329 0329 1 |  
330 0330 1 |  
331 0331 1 |  
332 0332 1 |  
333 0333 1 |  
334 0334 1 |  
335 0335 1 |  
336 0336 1 |  
337 0337 1 |  
338 0338 1 |  
339 0339 1 |  
340 0340 1 |  
341 0341 1 |  
342 0342 1 |

VBN associated with the just inserted key. The routine which performs this shifting (RMSSHFT VBNS) had to be modified to recognize these three bucket split cases, which of the two passes was current, and perform any required VBN shifting. Finally, RMSRECORD SIZE had to be similarly modified so that it would return the size of the index record to be inserted into the "current" bucket. The size, of course, as well as the current bucket depends upon the split case, and which of the two passes is currently under way. The three bucket split cases and their effect on these routines are as follows:

TWO-BUCKET SPLIT:  
During the first pass (when the index record contents of the old bucket are updated) just a key is added. The size of the record is the size of the key, and no VBNs need be shifted. During the second pass (when the index record contents of the new index bucket are appropriately modified), just a VBN is updated. The size of the record is the difference between the VBN size required to hold the old data level VBN and the size required to hold the new data level VBN. No VBN shifting is required unless the size is greater than 0.

THREE-BUCKET SPLIT - ONE KEY IN EACH INDEX BUCKET:  
During the first pass, just a key is added to the contents of the old index bucket so the record size is the size of the key and no VBNs need be shifted. During the second pass, a new key and VBN is added as the low order key in the new index bucket, and a VBN is also updated. In this case VBNs must be shifted to make room for the new VBN and any change in VBN size, and the record size includes the key added, the VBN added, and any difference in data level VBN size.

THREE-BUCKET SPLIT - BOTH KEYS IN OLD INDEX BUCKET:  
This split case only occurs when the insertion point and the split point are at the beginning of the old index bucket and both keys in the old bucket must be moved out to make room for the two new keys. During the first pass, both keys are added as is a VBN. The record size is the size of the two keys and the VBN added, and no VBN shifting need take place. During the second pass just a VBN is updated. The record size is computed to be any difference in sizes required to hold data level VBNs, and VBN shifting takes place only if there is a difference.

V02-017 TMK0001 Todd M. Katz 11-Jan-1982  
Make a change to RMSSHFT VBNS which affects three-way bucket splits where one of the resulting buckets is empty. The routine was adjusting the offset pointer to where the VBN shift should start. This is wrong, and in fact, caused index corruption so I removed it.

V02-016 PSK0007 Paulina S. Knibbe 14-Dec-1981  
Change record\_size to take BKT\_ADDR as an input

V02-015 PSK0006 Paulina S. Knibbe 26-Oct-1981

343 0343 1 Add support for compressed indexes to rm\$V3KEY\_SZ  
344 0344 1  
345 0345 1 V02-014 PSK0005 Paulina S. Knibbe 15-Aug-1981  
346 0346 1 Fix some problems w/large VBNs in index buckets.  
347 0347 1  
348 0348 1 V02-013 PSK0004 Paulina S. Knibbe 29-Jul-1981  
349 0349 1 Remove support for growing prologue three compressed  
350 0350 1 indexes. Add code to handle indexes with no chars  
351 0351 1 compressed off of the rear.  
352 0352 1  
353 0353 1 V02-012 PSK0003 Paulina S. Knibbe 24-Jul-1981  
354 0354 1 Fix off-by-one error in SHIFT\_VBNS; other problems  
355 0355 1 with growing VBNs  
356 0356 1  
357 0357 1 V02-011 PSK0002 Paulina S. Knibbe 13-Jul-1981  
358 0358 1 Add 'RMSSHFT\_VBNS' to handle spreading the VBN chain  
359 0359 1 in prologue three index buckets.  
360 0360 1  
361 0361 1 V02-010 MCN0007 Maria del C. Nasr 28-Jun-1981  
362 0362 1 Add RMSMOVE\_SIG\_CNT.  
363 0363 1  
364 0364 1 V02-009 PSK0001 Paulina S. Knibbe 08-Jun-1981  
365 0365 1 Add support for prologue three index buckets:  
366 0366 1 RMSRECORD\_SIZE, RMSNEW\_VBN\_BYTES, RMSV3KEY\_SZ  
367 0367 1 RMSV3\_IDX\_REC, RMSADD\_V3VBN, RMSADD\_V3KEY  
368 0368 1  
369 0369 1 V02-008 MCN0006 Maria del C. Nasr 16-Mar-1981  
370 0370 1 Increase size of record identifier to a word in the NRP.  
371 0371 1 Change linkage to RMSRECORD\_SIZE to include IFAB.  
372 0372 1  
373 0373 1 V02-007 REFORMAT C D Saether 01-Aug-1980 18:10  
374 0374 1  
375 0375 1 V0006 CDS0077 C D SAETHER 24-JAN-1980 13:50  
376 0376 1 Don't want to create dupes count field on continuation  
377 0377 1 SIDR arrays either.  
378 0378 1  
379 0379 1 V0005 CDS0071 C D SAETHER 15-JAN-1980 12:28  
380 0380 1 Fix bug creating dupes count field in SIDR records when  
381 0381 1 dupes aren't allowed (rm\$record\_size and rm\$bld\_new\_sidr).  
382 0382 1  
383 0383 1 Revision history:  
384 0384 1  
385 0385 1 Wendy Koenig, 24-OCT-78 14:02  
386 0386 1 X0002 - MAKE CHANGES CAUSED BY SHARING CONVENTIONS  
387 0387 1  
388 0388 1 Christian Saether, 26-JAN-79 9:01  
389 0389 1 X0003 - check down pointer before attempting to swing from empty bucket  
390 0390 1  
391 0391 1 Christian Saether, 1-july-79 11:00  
392 0392 1 X0004 - fix nxtrecid logic in RMSINS\_REC (caused re-use of ID's)  
393 0393 1 \*\*\*\*\*  
394 0394 1  
395 0395 1  
396 0396 1 LIBRARY 'RMSLIB:RMS';  
397 0397 1  
398 0398 1 REQUIRE 'RMSSRC:RMSIDXDEF';  
399 0463 1

```
400 0464 1 !  
401 0465 1 : define default psects for code  
402 0466 1 :  
403 0467 1 :  
404 0468 1 PSECT  
405 0469 1     CODE = RMSRMS3(PSECT_ATTR),  
406 0470 1     PLIT = RMSRMS3(PSECT_ATTR);  
407 0471 1 :  
408 0472 1 :  
409 0473 1 : Linkages  
410 0474 1 :  
411 0475 1 LINKAGE  
412 0476 1     L_COMPARE_KEY,  
413 0477 1     L_JSB,  
414 0478 1     L_JSB01,  
415 0479 1     L_PRESERVE1,  
416 0480 1     L_RABREG_567,  
417 0481 1     L_RABREG_67,  
418 0482 1     L_RABREG_7,  
419 0483 1     L_REC_OVHD,  
420 0484 1 :  
421 0485 1 : Local Linkages  
422 0486 1 :  
423 0487 1     RL$ADD_TO_ARRAY = JSB () :  
424 0488 1             GLOBAL (R_IDX_DFN, R_IRAB, R_IFAB, R_REC_ADDR),  
425 0489 1     RL$COMMON_LINK = JSB (STANDARD) :  
426 0490 1             GLOBAL (R_BKT_ADDR, R_IFAB, R_IRAB, R_IDX_DFN, R_REC_ADDR),  
427 0491 1     RL$MOVE_IN_VBN = JSB (REGISTER= 2) :  
428 0492 1             GLOBAL (R_REC_ADDR)  
429 0493 1             NOPRESERVE (2),  
430 0494 1     RL$NEW_VBN_BYTES = JSB () :  
431 0495 1             GLOBAL (R_BKT_ADDR, R_IDX_DFN);  
432 0496 1 :  
433 0497 1 : Forward Routines  
434 0498 1 :  
435 0499 1 FORWARD ROUTINE  
436 0500 1     RMSMOVE_IN_VBN      : RL$MOVE_IN_VBN NOVALUE,  
437 0501 1     RMSHFT_VBN      : RL$COMMON_LINK NOVALUE,  
438 0502 1     RMSV3_IDX_REC    : RL$COMMON_LINK NOVALUE,  
439 0503 1     RMSV3REY_5Z      : RL$COMMON_LINK,  
440 0504 1     RMSVBN_SIZE      : RL$PRESERVE1;  
441 0505 1 :  
442 0506 1 : External Routines  
443 0507 1 :  
444 0508 1 EXTERNAL ROUTINE  
445 0509 1     RMSCNTRL_ADDR    : RL$RABREG_567,  
446 0510 1     RMSCOMPARE_KEY   : RL$COMPARE_KEY,  
447 0511 1     RMSCOMPRESS_KEY  : RL$JSB01,  
448 0512 1     RMSGETNEXT_REC   : RL$RABREG_67,  
449 0513 1     RMSMOVE        : RL$PRESERVE1,  
450 0514 1     RMSNOREAD_LONG   : RL$JSB,  
451 0515 1     RMSRECOMPR_KEY  : RL$JSB01,  
452 0516 1     RMSRECORD_VBN    : RL$PRESERVE1,  
453 0517 1     RMSREC_OVHD     : RL$REC_OVHD,  
454 0518 1     RMSV3_VBN       : RL$RABREG_567;  
455 0519 1 :  
456 0520 1 MACRO
```

RM3MISPUT  
V04-000

H 1  
16-Sep-1984 01:51:28    VAX-11 Bliss-32 V4.0-742  
14-Sep-1984 13:01:29    [RMS.SRC]RM3MISPUT.B32;1

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(1)

: 457            0521 1    KEY\_LEN  
: 458            0522 1    CMPR\_CNT    = 0,0,8,0 %:  
: 459            0523 1

RM3  
V04

```
461 0524 1 %SBTTL 'RMSADD_TO_ARRAY'  
462 0525 1 ROUTINE RMSADD_TO_ARRAY : RL$ADD_TO_ARRAY NOVALUE =  
463 0526 1 !++  
464 0527 1 FUNCTIONAL DESCRIPTION:  
465 0528 1 This routine adds a new array to an existing SIDR array, and updates  
466 0529 1 the size field of the SIDR.  
467 0530 1 CALLING SEQUENCE:  
468 0531 1 RMSADD_TO_ARRAY()  
469 0532 1 INPUT PARAMETERS:  
470 0533 1 NONE  
471 0534 1 IMPLICIT INPUTS:  
472 0535 1 IFAB IFBSB_PLG_VER - address of IFAB  
473 0536 1 - prologue version of the file  
474 0537 1 IRAB IRBSL_LST_REC - address of IRAB  
475 0538 1 IRBSW_PUTUPD_ID - address of beginning of SIDR  
476 0539 1 IRBSL_PUTUPD_VBN - ID of UDR to put into record pointer  
477 0540 1 - VBN of UDR to put into record pointer  
478 0541 1 REC_ADDR - address of where new array element is to go  
479 0542 1 OUTPUT PARAMETERS:  
480 0543 1 NONE  
481 0544 1 IMPLICIT OUTPUTS:  
482 0545 1 REC_ADDR points - address of SIDR's size field  
483 0546 1 ROUTINE VALUE:  
484 0547 1 NONE  
485 0548 1 SIDE EFFECTS:  
486 0549 1 NONE  
487 0550 1 !--  
488 0551 1 BEGIN  
489 0552 1 EXTERNAL REGISTER  
490 0553 1 R_IDX_DFN,  
491 0554 1 R_IFAB_STR,  
492 0555 1 R_IRAB_STR,  
493 0556 1 R_REC_ADDR_STR;  
494 0557 1 GLOBAL REGISTER  
495 0558 1 R_RAB,  
496 0559 1 R_IMPURE,  
497 0560 1 R_BDB;
```

```

518      0581 2 LOCAL
519      0582 2     VBN_SIZE : BYTE,
520      0583 2     TOTAL_SIZE : BYTE;
521
522      0585 2     | Determine what will be the size of the record pointer's VBN.
523      0586 2
524      0587 2     VBN_SIZE = RM$VBN_SIZE (.IRAB[IRB$L_PUTUP_VBN]);
525
526      0589 2     | Set up the pointer size field in the new array element's control byte.
527      0590 2
528      0591 2     (.REC_ADDR)<0, 8> = .VBN_SIZE - 2;
529      0592 2     REC_ADDR = .REC_ADDR + 1;
530
531      0593 2     | Insert the record pointer's ID into the new SIDR array element, and
532      0594 2     compute the overall size of the new SIDR array element. Both of these
533      0595 2     steps are prologue dependent.
534      0596 2
535      0597 2     IF .IFAB[IFBSB_PLG_VER] LSSU PLG$C_VER_3
536      0599 2     THEN
537      0600 2       BEGIN
538      0601 2         (.REC_ADDR)<0, 8> = .IRAB[IRB$W_PUTUP_ID];
539      0602 2         REC_ADDR = .REC_ADDR + 1;
540      0603 2         TOTAL_SIZE = .VBN_SIZE + IRC$C_DATOVHDSZ;
541      0604 2       END
542      0605 2     ELSE
543      0606 2       BEGIN
544      0607 2         (.REC_ADDR)<0, 16> = .IRAB[IRB$W_PUTUP_ID];
545      0608 2         REC_ADDR = .REC_ADDR + 2;
546      0609 2         TOTAL_SIZE = .VBN_SIZE + IRC$C_DATOVHSZ;
547      0610 2       END:
548
549      0611 2     | Insert the record pointer's VBN into the new SIDR array element.
550      0612 2
551      0614 2     (.REC_ADDR)<0, .VBN_SIZE*8> = .IRAB[IRB$L_PUTUP_VBN];
552
553      0615 2     | Position to the size field of the SIDR, and update it to reflect
554      0616 2     the addition of the new array element.
555
556      0618 2
557      0619 2     REC_ADDR = .IRAB[IRB$L_LST_REC];
558      0620 2     REC_ADDR = .REC_ADDR + RMSREC_OVHD(-1) - IRC$C_DATSZFLD;
559      0621 2     (.REC_ADDR)<0, T6> = (.REC_ADDR)<0, 16> + .TOTAL_SIZE;
          0622 1     END:

```

```

.TITLE RM3MISPUT
.IDENT \V04-000\

.EXTRN RMSCNTRL_ADDR, RMSCOMPARE_KEY
.EXTRN RMSCOMPRESS_KEY
.EXTRN RMSGETNEXT_REC, RMSMOVE
.EXTRN RMSNOREAD_LONG, RMSRECOMPR_KEY
.EXTRN RMSRECORD_VBN, RMSREC_OVHD
.EXTRN RMSV3_VBN

.PSECT RMSRMS3,NOWRT, GBL, PIC,2

```

0910 8F BB 00000 RMSADD\_TO\_ARRAY:

RM3MISPUT  
V04-000

RMSADD\_TO\_ARRAY

K 1  
16-Sep-1984 01:51:28    VAX-11 Bliss-32 V4.0-742  
14-Sep-1984 13:01:29    [RMS.SRC]RM3MISPUT.B32;1

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				PUSHR #^M<R4,R8,R11>	: 0525
				PUSHL 120(IRAB)	: 0587
				BSBW RMSVBN_SIZE	:
				ADDL2 #4, SP	:
				MOVB R0, VBN_SIZE	:
				SUBB3 #2, VBN_SIZE, (REC_ADDR)+	0591
				CMPB 183(IFAB), #3	0598
				BGEQU 1\$	:
				MOVBL 128(IRAB), (REC_ADDR)+	0601
				ADDB3 #2, VBN_SIZE, TOTAL_SIZE	0603
				BRB 2\$	0598
				MOVW 128(IRAB), (REC_ADDR)+	0607
				ADDB3 #3, VBN_SIZE, TOTAL_SIZE	0609
				MOVZBL VBN_SIZE, R0	0614
				MULL2 #8, R0	:
				INSV 120(IRAB), #0, R0, (REC_ADDR)	0619
				MOVL 76(IRAB), REC_ADDR	0620
				MNEGL #1, R1	:
				BSBW RM\$REC_OVHD	0621
				MOVAB -2(R0)[REC_ADDR], REC_ADDR	0621
				MOVZBL TOTAL_SIZE, R0	0621
				ADDW2 R0, (REC_ADDR)	0622
				POPR #^M<R4,R8,R11>	0622
				RSB	:

; Routine Size: 85 bytes,    Routine Base: RMSRMS3 + 0000

561 0623 1 %SBTTL 'RMSADD\_V3KEY'  
562 0624 1 ROUTINE RMSADD\_V3KEY (KEY\_ADDR) : RL\$COMMON\_LINK NOVALUE =  
563 0625 1 ++  
564 0626 1  
565 0627 1 RM\$ADD\_V3KEY  
566 0628 1  
567 0629 1 This routine inserts the key at KEY\_ADDR into the bucket at REC\_ADDR  
568 0630 1 The contents of the bucket have already been spread apart. The  
569 0631 1 trailing record is recompressed if necessary.  
570 0632 1  
571 0633 1 CALLING SEQUENCE:  
572 0634 1  
573 0635 1 RM\$ADD\_V3KEY()  
574 0636 1  
575 0637 1 IMPLICIT INPUTS:  
576 0638 1 IFAB  
577 0639 1 REC\_ADDR  
578 0640 1 BKT\_ADDR  
579 0641 1 IDX\_DFN  
580 0642 1 IRAB[IRB\$V\_BIG\_SPLIT]  
581 0643 1 IRAB[IRB\$V\_EMPTY\_BKT]  
582 0644 1  
583 0645 1 OUTPUT PARAMETERS:  
584 0646 1  
585 0647 1 NONE  
586 0648 1  
587 0649 1 IMPLICIT OUTPUTS:  
588 0650 1  
589 0651 1 REC\_ADDR points past the inserted key  
590 0652 1 IRAB[IRB\$L\_LST\_NCMP] may possibly be updated  
591 0653 1  
592 0654 1 ROUTINE VALUE:  
593 0655 1  
594 0656 1 NONE  
595 0657 1  
596 0658 1 --  
597 0659 1  
598 0660 2 BEGIN  
599 0661 2  
600 0662 2 MAP  
601 0663 2 KEY\_ADDR : REF BBLOCK;  
602 0664 2  
603 0665 2 EXTERNAL REGISTER  
604 0666 2 R\_IDX\_DFN\_STR,  
605 0667 2 R\_IFAB\_STR,  
606 0668 2 R\_REC\_ADDR\_STR,  
607 0669 2 R\_BKT\_ADDR\_STR,  
608 0670 2 R\_IRAB\_STR;  
609 0671 2  
610 0672 2 GLOBAL REGISTER  
611 0673 2 R\_BDB,  
612 0674 2 R\_IMPURE,  
613 0675 2 R\_RAB;  
614 0676 2  
615 0677 2 IF NOT .IDX\_DFN [IDX\$V\_IDX\_COMPR]  
616 0678 2 THEN  
617 0679 2

```

618 0680 2 ! Fixed length index key.
619 0681 2
620 0682 2 BEGIN
621 0683 3 REC_ADDR = RMSMOVE (.IDX_DFN [IDX$B_KEYSZ], .KEY_ADDR, .REC_ADDR);
622 0684 3 RETURN
623 0685 2 END;
624 0686 2
625 0687 2 ! Compressed index key
626 0688 2 First fill in the length and compression count
627 0689 2
628 0690 2 RMSCOMPRESS_KEY (.KEY_ADDR);
629 0691 2
630 0692 2 ! Now move in the new key including the two bytes of compression overhead.
631 0693 2
632 0694 2 RMSMOVE (.KEY_ADDR[KEY_LEN] + 2,
633 0695 2 .KEY_ADDR,
634 0696 2 .REC_ADDR);
635 0697 2
636 0698 2 ! Fix up next record (front compression might change). This will be done
637 0699 2 only when a single a key is to be added, or when the second of two keys
638 0700 2 is being added. A key must also trail the key just inserted for this
639 0701 2 recompression to be performed.
640 0702 2
641 0703 3 IF (NOT .IRAB[IRBSV_BIG_SPLIT] OR .IRAB[IRBSV_EMPTY_BKT])
642 0704 2 AND
643 0705 3 (.REC_ADDR + .REC_ADDR[KEY_LEN] + 2
644 0706 3 LSSU .BKT_ADDR + .BKT_ADDR[BKT$W_FREESPACE])
645 0707 2 THEN
646 0708 2 RMSRECOMPR_KEY (.REC_ADDR, .REC_ADDR + .REC_ADDR [KEY_LEN] + 2);
647 0709 2
648 0710 2 ! If the key to be added is the first of two keys (from a multibucket split)
649 0711 2 and the front compression of the key just added is 0, increment the pointer
650 0712 2 to the last noncompressed key to this key to make the determination of the
651 0713 2 front compression of the second of the two keys easier.
652 0714 2
653 0715 3 IF .IRAB[IRBSV_BIG_SPLIT] AND (.REC_ADDR[CMPR_CNT] EQLU 0)
654 0716 2 THEN
655 0717 2 IRAB[IRBSL_LST_NCMP] = .REC_ADDR;
656 0718 2
657 0719 2 ! Reset REC_ADDR to point to the key which follows the key(s) just inserted
658 0720 2
659 0721 2 REC_ADDR = .REC_ADDR + .REC_ADDR [KEY_LEN] + 2;
660 0722 2
661 0723 2 RETURN
662 0724 2
663 0725 1 END;

```

14	1C	A7	0910	8F BB 00000 RMSADD_V3KEY:		
				PUSHR	#^M<R4,R8,R11>	: 0624
				BBS	#3, 28(IDX_DFN), 1\$	: 0677
				PUSHL	REC_ADDR	: 0683
	7E		14	AE DD 00004		
				56 DD 00009		
				PUSHL	KEY_ADDR	
				MOVZBL	32(IDX_DFN), -(SP)	
			20	A7 9A 0000E		

			0000G 30 00012	BSBW	RMSMOVE	
		5E	0C C0 00015	ADDL2	#12, SP	
		56	50 D0 00018	MOVL	R0, REC_ADDR	
			5C 11 0001B	BRB	5\$	0682
		50	10 AE D0 0001D	MOVL	KEY_ADDR, R0	0690
			0000G 30 00021	BSBW	RMSCOMPRESS_KEY	
			56 DD 00024	PUSHL	REC_ADDR	0696
		7E	14 AE DD 00026	PUSHL	KEY_ADDR	0695
		6E	18 BE 9A 00029	MOVZBL	@KEY_ADDR, -(SP)	0694
			02 C0 0002D	ADDL2	#2, TSP	
			0000G 30 00030	BSBW	RMSMOVE	
		05	5E 0C C0 00033	ADDL2	#12, SP	
	22	44	A9 02 E1 00036	BBC	#2, 68(IRAB), 2\$	0703
		A9 06 E1 0003B	BBC	#6, 68(IRAB), 3\$		
		50 66 9A 00040	MOVZBL	(REC_ADDR), R0	0705	
		51 02 A046 9E 00043	MOVAB	2(R0)[REC_ADDR], R1		
		50 04 A5 3C 00048	MOVZWL	4(BKT_ADDR), R0	0706	
		50 55 C0 0004C	ADDL2	BKT_ADDR, R0		
		50 51 D1 0004F	CMPL	R1, R0		
			0E 1E 00052	BGEQU	3\$	
		50 66 9A 00054	MOVZBL	(REC_ADDR), R0	0708	
		51 02 A046 9E 00057	MOVAB	2(R0)[REC_ADDR], R1		
		50 56 D0 0005C	MOVL	REC_ADDR, R0		
			0000G 30 0005F	BSBW	RMSRECOMP KEY	
	0A	44	A9 02 E1 00062	BBC	#2, 68(IRAB), 4\$	0715
		01 A6 95 00067	TSTB	1(REC_ADDR)		
			05 12 0006A	BNEQ	4\$	
	0098	C9 56 D0 0006C	MOVL	REC_ADDR, 152(IRAB)	0717	
		50 66 9A 00071	MOVZBL	(REC_ADDR), R0	0721	
		56 02 A046 9E 00074	MOVAB	2(R0)[REC_ADDR], REC_ADDR		
		0910 8F BA 00079	POPR	#^M<R4,R8,R11>	0725	
			05 0007D	RSB		

: Routine Size: 126 bytes, Routine Base: RMSRMS3 + 0055

; 665 0726 1 XSBTTL 'RM\$ADD\_V3VBN'  
666 0727 1 ROUTINE RMSADD\_V3VBN (VBN) : RL\$COMMON\_LINK NOVALUE =  
667 0728 1 ++  
668 0729 1  
669 0730 1 RMSADD\_V3VBN  
670 0731 1  
671 0732 1 This routine adds a VBN to the chain at the end of the bucket.  
672 0733 1 If the VBN is longer than the current VBN size all VBNs grow  
673 0734 1 by the appropriate number of bytes.  
674 0735 1  
675 0736 1 CALLING SEQUENCE:  
676 0737 1  
677 0738 1 RM\$ADD\_V3VBN(.VBN)  
678 0739 1  
679 0740 1 IMPLICIT INPUTS:  
680 0741 1 IRAB  
681 0742 1 BKT\_ADDR  
682 0743 1 IDX\_DFN  
683 0744 1  
684 0745 1 OUTPUT PARAMETERS:  
685 0746 1  
686 0747 1 NONE  
687 0748 1  
688 0749 1 IMPLICIT OUTPUTS:  
689 0750 1  
690 0751 1  
691 0752 1 ROUTINE VALUE:  
692 0753 1  
693 0754 1 NONE  
694 0755 1  
695 0756 1 --  
696 0757 1  
697 0758 2 BEGIN  
698 0759 2  
699 0760 2 MACRO  
700 0761 2 OFFSET = 0,0,16,0 %;  
701 0762 2  
702 0763 2 EXTERNAL REGISTER  
703 0764 2 R\_IDX DFN\_STR,  
704 0765 2 R\_IRAB STR,  
705 0766 2 R\_REC ADDR\_STR,  
706 0767 2 R\_IFAB STR,  
707 0768 2 P\_BKT\_ADDR\_STR;  
708 0769 2  
709 0770 2 GLOBAL REGISTER  
710 0771 2 R\_BDB,  
711 0772 2 R\_RAB,  
712 0773 2 R\_IMPURE;  
713 0774 2  
714 0775 2 LOCAL  
715 0776 2 NEW\_SIZE,  
716 0777 2 OLD\_SIZE;  
717 0778 2  
718 0779 2 NEW\_SIZE = RM\$VBN\_SIZE (.VBN);  
719 0780 2 OLD\_SIZE = .BKT\_ADDR [BKT\$V\_PTR\_SZ] + 2;  
720 0781 2  
721 0782 2 IF .NEW\_SIZE GTRU .OLD\_SIZE

```
722      0783 2 THEN
723      0784 2
724      0785 2      ! We need to shuffle all of the VBNs down
725      0786 2      | First figure out where to move them.
726      0787 2
727      0788 2      BEGIN
728      0789 2
729      0790 2      LOCAL
730      0791 3          OLD_END      : REF BBLOCK,
731      0792 3          NUM_RECS,
732      0793 3          NEW_END,
733      0794 3          SAVE;
734      0795 3
735      0796 3          OLD_END = .BKT_ADDR + (.IDX_DFN [IDX$B_IDXBKTSZ] * 512) - BKT$C_ENDOVHD;
736      0797 3          NUM_RECS = (.OLD_END - (.BKT_ADDR+.OLD_END [OFFSET])) / .OLD_SIZE;
737      0798 3          NEW_END = .OLD_END - (.NUM_RECS * .NEW_SIZE);
738      0799 3
739      0800 3          IF .BKT_ADDR [BKT$B_LEVEL] EQLU 0
740      0801 3          AND .IDX_DFN [IDX$V_DUPKEYS]
741      0802 3          THEN
742      0803 3              NEW_END = .NEW_END - 4;
743      0804 3
744      0805 3          OLD_END [OFFSET] = .NEW_END - .BKT_ADDR - 1;
745      0806 3
746      0807 3          ! Now actually move them
747      0808 3
748      0809 3          SAVE = .IRAB [IRB$L_REC_COUNT];
749      0810 3          IRAB [IRB$L_REC_COUNT] = .NUM_RECS - 1;
750      0811 3
751      0812 3          WHILE .IRAB [IRB$L_REC_COUNT] GEQ 0
752      0813 3          DO
753      0814 4              BEGIN
754      0815 4                  (.NEW_END) <0,8 * .NEW_SIZE> = .(RMS$CNTRL_ADDR())<0,8*.OLD_SIZE>;
755      0816 4                  IRAB [IRB$L_REC_COUNT] = .IRAB [IRB$L_REC_COUNT] - 1;
756      0817 4                  NEW_END = .NEW_END + .NEW_SIZE
757      0818 3                  END;
758      0819 3
759      0820 3          IRAB [IRB$L_REC_COUNT] = .SAVE;
760      0821 3
761      0822 3          ! Update the VBN size
762      0823 3
763      0824 3          BKT_ADDR [BKT$V_PTR_SZ] = .NEW_SIZE - 2
764      0825 3
765      0826 2          END;
766      0827 2
767      0828 2          ! Actually add the new VBN
768      0829 2
769      0830 3          BEGIN
770      0831 3
771      0832 3          LOCAL
772      0833 3              ADDR;
773      0834 3
774      0835 3              ADDR = RMS$CNTRL_ADDR();
775      0836 3              (.ADDR) <0,.NEW_SIZE * 8> = .VBN
776      0837 3
777      0838 2          END; ! Of local ADDR
778      0839 2
```

			091C	8F BB 00000 RMSADD_V3VBN:			
			5E	24 0C C2 00004	PUSHR	#^M<R2,R3,R4,R8,R11>	0727
				AE DD 00007	SUBL2	#12, SP	0779
				0000V 30 0000A	PUSHL	VBN	
			5E	04 C0 0000D	BSBW	RMSVBN_SIZE	
			6E	50 D0 00010	ADDL2	#4, SP	
			52	03 EF 00013	MOVL	R0, NEW_SIZE	
			52	02 C0 00019	EXTZV	#3, #2, 13(BKT_ADDR), OLD_SIZE	0780
			52	6E D1 0001C	ADDL2	#2, OLD_SIZE	
				7D 1B 0001F	CMPL	NEW_SIZE, OLD_SIZE	0782
					BLEQU	4S	
52	OD A5		50	16 A7 9A 00021	MOVZBL	22(IDX_DFN), R0	0796
			50	09 78 00025	ASHL	#9, R0, R0	
			50	FC A045 9E 00029	MOVAB	-4(R0)[BKT_ADDR], OLD_END	
			51	60 3C 0002E	MOVZWL	(OLD_END), R1	0797
			51	55 C0 00031	ADDL2	BKT_ADDR, R1	
			51	51 C3 00034	SUBL3	R1, OLD_END, R1	
			53	52 C6 00038	DIVL2	OLD_SIZE, NUM_RECS	
			53	51 C5 0003B	MULL3	NEW_SIZE, NUM_RECS, R3	0798
			54	53 C3 0003F	SUBL3	R3, OLD_END, NEW_END	
				OC A5 95 00043	TSTB	12(BKT_ADDR)	
				07 12 00046	BNEQ	1S	
				03 1C A7 E9 00048	BLBC	28(IDX_DFN), 1S	0801
			53	04 C2 0004C	SUBL2	#4, NEW_END	0803
			54	55 C3 0004F	1\$: SUBL3	BKT_ADDR, NEW_END, R3	0805
			60	01 A3 00053	SUBW3	#1-R3, (OLD_END)	
			53	0094 C9 D0 00057	MOVL	148(IRAB), SAVE	0809
			08	C9 FF A1 9E 0005D	MOVAB	-1(R1), 148(IRAB)	0810
			0094	53 C9 9E 00063	MOVAB	148(R9), R3	0812
			04 AE	6E 03 78 00068	ASHL	#3, NEW_SIZE, 4(SP)	0815
				52 08 C4 0006D	MULL2	#8, R2	
				63 D5 00070	2\$: TSTL	(R3)	0812
				1A 19 00072	BLSS	3S	
				0000G 30 00074	BSBW	RMSCTRL_ADDR	0815
51	04 60		52	00 EF 00077	EXTZV	#0, R2, TRO, R1	
64	04 AE		00	51 F0 0007C	INSV	R1, #0, 4(SP), (NEW_END)	
			53	0094 C9 9E 00082	MOVAB	148(R9), R3	0816
				63 D7 00087	DECL	(R3)	
			54	6E C0 00089	ADDL2	NEW_SIZE, NEW_END	0817
				E2 11 0008C	BRB	2S	
			0094	C9 08 AE D0 0008E	3\$: MOVL	SAVE, 148(IRAB)	0820
			50	6E 02 C3 00094	SUBL3	#2, NEW_SIZE, R0	0824
			02	03 50 F0 00098	INSV	R0, #3, #2, 13(BKT_ADDR)	
OD A5	02		03	51 50 D0 000A1	BSBW	RMSCTRL_ADDR	0835
				8E 03 78 000A4	MOVL	R0, ADDR	
			50	00 20 AE F0 000A8	ASHL	#3, NEW_SIZE, R0	0836
				5E 08 C0 000AE	INSV	VBN, #0, R0, (ADDR)	
			091C	8F BA 000B1	ADDL2	#8, SP	0840
				05 000B5	POPR	#^M<R2,R3,R4,R8,R11>	
					RSB		

RM3MISPUT  
V04-000

RMSADD\_V3VBN

: Routine Size: 182 bytes, Routine Base: RMSRMS3 + 00D3

E 2  
16-Sep-1984 01:51:28  
14-Sep-1984 13:01:29

VAX-11 Bliss-32 V4.0-742  
[RMS.SRC]RM3MISPUT.B32;1

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(4)

RM3  
V04

781 0841 1 %SBTTL 'RMSBKT\_SORT'  
782 0842 1 GLOBAL ROUTINE RMSBKT\_SORT (BKT\_ADDR) : RL\$RABREG\_7 =  
783 0843 1  
784 0844 1 !++  
785 0845 1  
786 0846 1 FUNCTIONAL DESCRIPTION:  
787 0847 1  
788 0848 1 this routine scans the bucket, makes a table of all used ID's,  
789 0849 1 and then updates the NXTRECID and LSTRECID fields to reflect the existence  
790 0850 1 of the lowest range of un-used ID's available.  
791 0851 1  
792 0852 1 CALLING SEQUENCE:  
793 0853 1 RMSBKT\_SORT (BKT\_ADDR)  
794 0854 1  
795 0855 1 INPUT PARAMETERS:  
796 0856 1 BKT\_ADDR - pointer to bucket to be searched for ID's  
797 0857 1  
798 0858 1 IMPLICIT INPUTS:  
799 0859 1 IDX\_DFN - pointer to index descriptor  
800 0860 1 IRAB - pointer to internal RAB  
801 0861 1 IFAB - IFAB  
802 0862 1  
803 0863 1 OUTPUT PARAMETERS:  
804 0864 1 NONE  
805 0865 1  
806 0866 1 IMPLICIT OUTPUTS:  
807 0867 1 NONE  
808 0868 1  
809 0869 1 ROUTINE VALUE:  
810 0870 1 low bit set if success (at least one ID was found)  
811 0871 1 low bit clear if all ID's in use  
812 0872 1  
813 0873 1 SIDE EFFECTS:  
814 0874 1 NXTRECID is set to lowest ID available  
815 0875 1 LSTRECID is set to highest contiguous ID available after NXTRECID  
816 0876 1  
817 0877 1 --  
818 0878 1  
819 0879 2 BEGIN  
820 0880 2  
821 0881 2 MAP  
822 0882 2 BKT\_ADDR : REF BBLOCK;  
823 0883 2  
824 0884 2 EXTERNAL REGISTER  
825 0885 2 R\_IDX\_DFN,  
826 0886 2 R\_IFAB,  
827 0887 2 R\_IRAB;  
828 0888 2  
829 0889 2 GLOBAL REGISTER  
830 0890 2 R\_REC\_ADDR\_STR;  
831 0891 2  
832 0892 2 LOCAL  
833 0893 2 BIT\_NUM,  
834 0894 2 BIT\_LEN;  
835 0895 2  
836 0896 2 ! create 8 longword array for table of used ID's  
837 0897 2

```
838 0898 2
839 0899 2
840 0900 2
841 0901 2
842 0902 2
843 0903 2
844 0904 2
845 0905 2
846 0906 2
847 0907 2
848 0908 3
849 0909 3
850 0910 3
851 0911 3
852 0912 3
853 0913 3
854 0914 3
855 0915 3
856 0916 3
857 0917 3
858 0918 3
859 0919 3
860 0920 4
861 0921 4
862 0922 4
863 0923 4
864 0924 4
865 0925 4
866 0926 4
867 0927 4
868 0928 3
869 0929 3
870 0930 3
871 0931 3
872 0932 2
873 0933 2
874 0934 2
875 0935 2
876 0936 2
877 0937 2
878 0938 2
879 0939 2
880 0940 2
881 0941 2
882 0942 2
883 0943 2
884 0944 2
885 0945 2
886 0946 3
887 0947 3
888 0948 3
889 0949 3
890 0950 4
891 0951 4
892 0952 4
893 0953 4
894 0954 4

     STACKLOCAL
     USED_IDS      : VECTOR [8];
     LOCAL
     EOB;
     DECR I FROM 7 TO 0 DO
       USED_IDS[I] = 0;
     BEGIN
     MAP
       USED_IDS      : BITVECTOR [256];
       REC_ADDR = .BKT_ADDR + BKT$C_OVERHDSZ;
       EOB = .BKT_ADDR[BKT$W_FREESPACE] + .BKT_ADDR;
       ! scan records in bucket.  set bit in bitvector for those id values found.
       WHILE .REC_ADDR LSSA .EOB
       DO
         BEGIN
         GLOBAL REGISTER
           R_RAB,
           R_IMPURE;
         USED_IDS[REC_ADDR[IRC$B_ID]] = 1;
         RMSGETNEXT_REC();
         END;
         USED_IDS[0] = 0;                                ! flag set when a free id found
         USED_IDS[255] = 1;                              ! make sure last one always set
         END;                                            ! of USED_IDS as bitvector
         BIT_LEN = 31;
         BIT_NUM = 1;
         ! now search the bit array for the first clear bit (that corresponds to
         ! an unused id value) one longword at a time.  when found, check the
         ! remaining bits in that longword looking for a used id.  if none found
         ! there, continue searching the remaining longwords in the array to
         ! determine the extent of the unused id values.
         INCR I FROM 0 TO 7 DO
         IF NOT .USED_IDS
         THEN
           BEGIN
           IF NOT FFC(BIT_NUM, BIT_LEN, USED_IDS[I], BIT_NUM)
           THEN
             BEGIN
             BKT_ADDR[BKT$B_NXTRECID] = .I*32 + .BIT_NUM;
             USED_IDS<0,1> = 1;
             BIT_NUM = 32 - .BIT_NUM;
           
```

```

895      0955  4      IF NOT FFS(BIT_NUM, BIT_LEN, USED_IDS[.I], BIT_NUM)
896      0956  4      THEN
897      0957  5      BEGIN
898      0958  5      BKT_ADDR[BKT$B_LSTRECID] = .I*32 + .BIT_NUM - 1;
899      0959  5      EXITLOOP;
900      0960  5
901      0961  5      END
902      0962  4      ELSE
903      0963  4      BIT_LEN = 32;
904      0964  4
905      0965  3      END;
906      0966  3
907      0967  3      BIT_NUM = 0;
908      0968  3      END
909      0969  2      ELSE
910      0970  2
911      0971  2      IF NOT FFS(BIT_NUM, BIT_LEN, USED_IDS[.I], BIT_NUM)
912      0972  2      THEN
913      0973  3      BEGIN
914      0974  3      BKT_ADDR[BKT$B_LSTRECID] = .I*32 + .BIT_NUM - 1;
915      0975  3      EXITLOOP;
916      0976  3
917      0977  3      END
918      0978  2      ELSE
919      0979  3      BEGIN
920      0980  3      BIT_LEN = 32;
921      0981  3      BIT_NUM = 0;
922      0982  2      END;
923      0983  2
924      0984  2      ! if this is clear, no free id's were found. if set, at least one
925      0985  2      is available.
926      0986  2
927      0987  2      RETURN .USED_IDS[0];
928      0988  2
929      0989  1      END;

```

	094C	8F BB 00000 RMSBKT_SORT::		
	5E	20 C2 00004	PUSHR	#^M<R2,R3,R6,R8,R11>
	50	07 D0 00007	SUBL2	#32, SP
	6E40	D4 0000A 1\$:	MOVL	#7, I
	FA	50 F4 0000D	CLRL	USED_IDS[I]
	52	38 AE D0 00010	SOBGEQ	I, 1\$
	56	0E A2 9E 00014	MOVL	BKT_ADDR, R2
	53	04 A2 3C 00018	MOVAB	14(R2), REC_ADDR
	53	52 C0 0001C	MOVZWL	4(R2), EOB
	53	56 D1 0001F 2\$:	ADDL2	R2, EOB
		0D 1E 00022	CMPL	REC_ADDR, EOB
	00	50 A6 9A 00024	BGEQU	4\$
	6E	01 6E 00028	MOVZBL	1(REC_ADDR), R0
		0000G 30 0002C 3\$:	BBSS	R0, USED_IDS, 3\$
		EE 11 0002F	BSBW	RM\$GETNEXT_REC
	6E	01 8A 00031 4\$:	BRB	2\$
			BICB2	#1, USED_IDS

0842  
0905  
0906  
0913  
0914  
0918  
0926  
0927  
0918  
0930

			1F	AE	80	8F	88 00034	BISB2	#128, USED_IDS+31	0931
			58		1F	D0	00039	MOVL	#31, BIT_LEN	0933
			56		01	D0	0003C	MOVL	#1, BIT_NUM	0934
					50	D4	0003F	CLRL	I	0942
			58		6E40	DE	00041 5\$:	MOVAL	USED_IDS[I], R11	0948
			20		6E	E8	00045	BLBS	USED_IDS, 6\$	0944
	56	68	58		56	EB	00048	FFC	BIT_NUM, BIT_LEN, (R11), BIT_NUM	0948
					32	13	0004D	BEQL	9\$	
			51		05	78	0004F	ASHL	#5 I, R1	0951
	06	A2	51		56	81	00053	ADDB3	BIT_NUM, R1, 6(R2)	0952
			6E		01	88	00058	BISB2	#1 USED_IDS	0953
	56	68	58		56	C3	0005B	SUBL3	BIT_NUM, #32, BIT_LEN	0955
					56	EA	0005F	FFS	BIT_NUM, BIT_LEN, (R11), BIT_NUM	
					18	13	00064	BEQL	8\$	
	56	68	58		0B	11	00066	BRB	7\$	0958
					56	EA	00068 6\$:	FFS	BIT_NUM, BIT_LEN, (R11), BIT_NUM	0971
			51		0F	13	0006D	BEQL	8\$	
					05	78	0006F	ASHL	#5 I, R1	0974
			51		53	A641	9E 00073 7\$:	MOVAB	-1(BIT_NUM)[R1], R3	0973
			07	A2	53	90	00078	MOVB	R3, 7(R2)	
					09	11	0007C	BRB	10\$	
					58	20	D0 0007E 8\$:	MOVL	#32, BIT_LEN	0980
			BA		56	D4	00081 9\$:	CLRL	BIT_NUM	0981
					07	F3	00083	AOBLEQ	#7, I, 5\$	0944
					50	8E	D0 00087 10\$:	MOVL	USED_IDS, R0	0987
					5E	10	C0 0008A	ADDL2	#28, SP	0989
					094C	81	BA 0008D	POPR	#^M<R2,R3,R6,R8,R11>	
							05 00091	RSB		

; Routine Size: 146 bytes, Routine Base: RMSRMS3 + 0189

971  
972 0990 1 %SBTTL 'RMSBLD\_IDX\_REC'  
973 0991 1 ROUTINE RMSBLD\_IDX\_REC (REC\_SZ) : RL\$COMMON\_LINK NOVALUE =  
974 0992 1 !++  
975 0993 1  
976 0994 1  
977 0995 1 FUNCTIONAL DESCRIPTION:  
978 0996 1  
979 0997 1 This routine builds an index record in the buffer. It will do one of the  
980 0998 1 following:  
981 1000 1 0) Call RMSV3\_IDX\_REC to do everything for prologue three index  
982 1001 1 buckets.  
983 1002 1  
984 1003 1 1) replace key value of existing down pointer with the new high key  
985 1004 1 value (in keybuffer 2) of the original bucket (described as vbn\_left)  
986 1005 1 and use that key value along with the value of vbn\_right to create  
987 1006 1 a new index record pointing to the right hand bucket of the split.  
988 1007 1  
989 1008 1 2) same as the above except create another record between the two  
990 1009 1 described using the contents of keybuffer 3 and vbn\_right. this  
991 1010 1 record points to the middle record of a 3 or 4 bucket split.  
992 1011 1  
993 1012 1 3) as an option on the first situation (1), change the down  
994 1013 1 pointer to be the value of vbn\_mid if the current down pointer matches  
995 1014 1 the value in vbn\_left.  
996 1015 1  
997 1016 1 4) build only the left hand part of the possible pieces described above  
998 1017 1 and make the break just before the potential new middle bucket pointer.  
999 1018 1  
1000 1019 1 5) build the part remaining after only doing 4).  
1001 1020 1  
1002 1021 1 Any modifications to this code should be done only after thoroughly  
1003 1022 1 understanding the interactions between this routine, RM\$RECORD\_SIZE,  
1004 1023 1 and RM\$SPLIT\_EM.  
1005 1024 1  
1006 1025 1 CALLING SEQUENCE:  
1007 1026 1 RMSBLD\_IDX\_REC (REC\_SZ)  
1008 1027 1  
1009 1028 1 INPUT PARAMETERS:  
1010 1029 1 REC\_SZ - total size of record to be inserted  
1011 1030 1  
1012 1031 1 IMPLICIT INPUTS:  
1013 1032 1 REC\_ADDR - points to location in buffer where record is to be built  
1014 1033 1  
1015 1034 1 IRAB [vbn\_left] - left hand (original) bucket in split. zero if  
1016 1035 1 only right hand part of record to be built.  
1017 1036 1 [vbn\_mid] - possible middle bucket in 3 or 4 bucket split  
1018 1037 1 [vbn\_right] - right hand bucket of split  
1019 1038 1 [empty\_bkt] - left hand bucket is dead, i.e., no data records  
1020 1039 1 [big\_split] - more than 1 new bucket in split  
1021 1040 1 [spl\_idx] - only build left hand part of index record  
1022 1041 1 [pos\_ins] - position of insert of the new record  
1023 1042 1 [keybuf] - address of keybuffers  
1024 1043 1  
1025 1044 1 IFAB [kbufsz] - size of keybuffers  
1026 1045 1  
1027 1046 1 IDX\_DFN [keysz] - size of this key

```
988 1047 1
989 1048 1 | OUTPUT PARAMETERS:
990 1049 1 |   NONE
991 1050 1
992 1051 1 | IMPLICIT OUTPUTS:
993 1052 1 |   REC_ADDR - points to first byte beyond record.
994 1053 1
995 1054 1 | ROUTINE VALUE:
996 1055 1 |   NONE
997 1056 1
998 1057 1 | SIDE EFFECTS:
999 1058 1 |   NONE
1000 1059 1
1001 1060 1 | --
1002 1061 1
1003 1062 2 | BEGIN
1004 1063 2
1005 1064 2 | EXTERNAL REGISTER
1006 1065 2 |   R_BKT_ADDR_STR,
1007 1066 2 |   R_REC_ADDR_STR,
1008 1067 2 |   R_IRAB_STR,
1009 1068 2 |   R_IDX_DFN_STR,
1010 1069 2 |   R_IFAB_STR;
1011 1070 2
1012 1071 2 | GLOBAL REGISTER
1013 1072 2 |   R_BDB,
1014 1073 2 |   R_IMPURE,
1015 1074 2 |   R_RAB;
1016 1075 2
1017 1076 2 | IF .IFAB [IFB$B_PLG_VER] EQLU PLG$C_VER_3
1018 1077 2 | THEN
1019 1078 2
1020 1079 2 |   Prologue three index records. They basically
1021 1080 2 |   have nothing in common with pre-prologue three
1022 1081 2 |   index records. Just call the routine that
1023 1082 2 |   handles them and get out.
1024 1083 2
1025 1084 3 | BEGIN
1026 1085 3
1027 1086 3 | RMSV3_IDX_REC (.REC_SZ);
1028 1087 3 | RETURN
1029 1088 2 | END;
1030 1089 2
1031 1090 2 | first key we have to build if present is the new key for the left
1032 1091 2 | hand bucket this will be zero only when splitting up index record when
1033 1092 2 | two new key values are required - currently 'empty' buckets (i.e.,
1034 1093 2 | buckets which end up without any records in them after a split) are not
1035 1094 2 | removed from the index
1036 1095 2
1037 1096 2
1038 1097 2 | IF .IRAB[IRBSL_VBN_LEFT] NEQ 0
1039 1098 2 | THEN
1040 1099 2
1041 1100 2 |   the bucket has already been spread apart but the index record we
1042 1101 2 |   positioned to still exists where we found it, so using the existing
1043 1102 2 |   pointer, change the key value to the high key of VBN_LEFT
1044 1103 2
```

```
1045 1104 3 BEGIN
1046 1105 3
1047 1106 3 ! check for empty bucket situation, compare down pointers before
1048 1107 3 deciding whether to modify down pointer or not
1049 1108 3
1050 1109 3
1051 1110 3 IF .IRAB[IRB$V_EMPTY_BKT]
1052 1111 3 THEN BEGIN
1053 1112 4
1054 1113 4
1055 1114 4 BUILTIN
1056 1115 4 AP;
1057 1116 4
1058 1117 4 AP = 1; ! set for index level
1059 1118 4
1060 1119 4 IF .IRAB[IRB$L_VBN_LEFT] EQLA RMSRECORD_VBN()
1061 1120 4
1062 1121 4 ! current down pointer matches so modify down pointer
1063 1122 4
1064 1123 4
1065 1124 4 THEN RMSMOVE_IN_VBN(.IRAB[IRB$L_VBN_MID])
1066 1125 4 ELSE REC_ADDR = .REC_ADDR + .REC_ADDR[IRC$V_PTRSZ] + IRC$C_IDXPTRBAS
1067 1126 4 + IRC$C_IDXOVHDSZ; ! advance past pointer
1068 1127 4
1069 1128 4
1070 1129 4 IF NOT .IRAB[IRB$V_BIG_SPLIT]
1071 1130 4 THEN RETURN; ! no more to do, quit
1072 1131 4
1073 1132 4
1074 1133 4 END
1075 1134 3 ELSE REC_ADDR = .REC_ADDR + .REC_ADDR[IRC$V_PTRSZ] + IRC$C_IDXPTRBAS +
1076 1135 3 + IRC$C_IDXOVHDSZ;
1077 1136 3
1078 1137 3
1079 1138 3 ! move in key from keybuffer 2
1080 1139 3
1081 1140 3 REC_ADDR = RMSMOVE(.IDX_DFN[IDX$B_KEYSZ], KEYBUF_ADDR(2), .REC_ADDR);
1082 1141 3
1083 1142 3 ! We move in the VBN from IRB$L_VBN_RIGHT if and only if this is a
1084 1143 3 one-pass multi-bucket empty bucket split case.
1085 1144 3
1086 1145 4 IF (.IRAB[IRB$V_EMPTY_BKT]
1087 1146 4 AND
1088 1147 4 .IRAB[IRB$V_BIG_SPLIT]
1089 1148 4 AND
1090 1149 4 NOT .IRAB[IRB$V_SPL_IDX])
1091 1150 3 THEN RMSMOVE_IN_VBN(.IRAB[IRB$L_VBN_RIGHT]);
1092 1151 3
1093 1152 3 ! We are done if this was a empty bucket split case
1094 1153 3
1095 1154 3
1096 1155 3 IF .IRAB[IRB$V_EMPTY_BKT]
1097 1156 3 THEN RETURN;
1098 1157 3
1099 1158 3
1100 1159 2
1101 1160 2 END;
```

```
1102      1161 2 | At this point all bucket splits involving empty buckets have completed
1103      1162 2 | except for the second pass of the two-pass multi-bucket with empty
1104      1163 2 | bucket split case.
1105      1164 2 |
1106      1165 2 | We continue for all other cases with one exception. When the insertion
1107      1166 2 | point is found to be at the split point, two passes must be made - one
1108      1167 2 | to update the old bucket, and the second to update the new bucket. If
1109      1168 2 | this is the first pass to update the old bucket components we do not
1110      1169 2 | continue unless a multibucket split had occurred, two and only two
1111      1170 2 | index records could fit in an index bucket, and the split point requires
1112      1171 2 | both of them to be moved to the new bucket to make room for the two new
1113      1172 2 | index records. In all other "first pass" situations we do not continue.
1114      1173 2 |
1115      1174 2 | IF NOT .IRAB[IRB$V_SPL_IDX]
1116          OR
1117          (.IRAB[IRB$V SPL IDX]
1118              AND .IRAB[IRBSV_BIG_SPLIT]
1119              AND .IRAB[IRBSW_POS_INS] EQLU BKT$C_OVERHDSZ)
1120      THEN
1121          BEGIN
1122          |
1123          | If a multibucket split occurred, the middle VBN and its corresponding
1124          | key are inserted into the current index bucket with two exceptions.
1125          | If this is the second pass of a two-pass multibucket split case where
1126          | both records had to be put in the old index bucket, or of a two-pass
1127          | multi-bucket empty bucket split case, then this bucket updating does
1128          | not take place.
1129          1188 3 |
1130          1189 3 |
1131          1190 3 |
1132          1191 4 | IF .IRAB[IRB$V_BIG_SPLIT]
1133          1192 3 |     AND
1134          1193 4 |     (.IRAB[IRBSL_VBN_MID] NEQU 0)
1135          1194 4 |     THEN
1136          1195 4 |         BEGIN
1137          1196 4 |             RM$MOVE_IN_VBN(.IRAB[IRBSL_VBN_MID]);
1138          1197 4 |             |
1139          1198 4 |             | move in key from keybuffer 3
1140          1199 4 |             REC_ADDR = RM$MOVE(.IDX_DFN[IDX$B_KEYSZ],
1141          1200 3 |                         KEYBUF_ADDR(3), .REC_ADDR);
1142          1201 3 |             END;
1143          1202 3 |
1144          1203 3 |
1145          1204 3 |
1146          1205 3 |
1147          1206 3 |
1148          1207 3 |
1149          1208 3 |
1150          1209 3 |
1151          1210 3 |
1152          1211 2 |     IF NOT .IRAB[IRB$V_SPL_IDX]
1153          1212 2 |         THEN
1154          1213 2 |             RM$MOVE_IN_VBN(.IRAB[IRBSL_VBN_RIGHT]);
1155          1214 1 |             END;
```

0914 BF BB 00000 RMSBLD_IDX_REC:									
						PUSHR	#^M<R2,R4,R8,R11>		0991
		03	00B7	CA	91 00004	CMPB	183(IFAB), #3		1076
			14	OB	12 00009	BNEQ	1\$		
				AE	DD 0000B	PUSHL	REC SZ		1086
				0000V	30 0000E	BSBW	RMSV3_IDX_REC		
		5E		04	C0 00011	ADDL2	#4, SP		
				30	11 00014	BRB	4\$		1084
		51	0088	C9	D0 00016	1\$: MOVL	136(IRAB), R1		1097
				69	13 0001B	BEQL	8\$		
27	44	A9		06	E1 0001D	BBC	#6, 68(IRAB), 5\$		1110
		5C		01	D0 00022	MOVL	#1, AP		1117
				0000G	30 00025	BSBW	RMSRECORD_VBN		1119
		50		51	D1 00028	CMPL	R1, R0		
				0A	12 0002B	BNEQ	2\$		
		52	0090	C9	D0 0002D	MOVL	144(IRAB), R2		1124
				0000V	30 00032	BSBW	RMSMOVE_IN_VBN		
				0A	11 00035	BRB	3\$		
50	66	02		00	EF 00037	2\$: EXTZV	#0, #2, (REC_ADDR), R0		1126
		56	03	A046	9E 0003C	MOVAB	3(R0)[REC_ADDR], REC_ADDR		1127
	0D	44	A9		02 E0 00041	3\$: BBS	#2, 68(IRAB), 6\$		1129
50	66	02		0083	31 00046	4\$: BRW	11\$		1131
		56	03	A046	9E 0004E	EXTZV	#0, #2, (REC_ADDR), R0		1135
				56	DD 00053	MOVAB	3(R0)[REC_ADDR], REC_ADDR		
		50	00B4	CA	3C 00055	PUSHL	REC ADDR		1140
				60	B940 9F 0005A	MOVZWL	180(IFAB), R0		
		7E	20	A7	9A 0005E	PUSHAB	@96(IRAB)[R0]		
				0000G	30 00062	MOVZBL	32(IDX DFN), -(SP)		
		5E		0C	C0 00065	BSBW	RMSMOVE		
		56		50	D0 00068	ADDL2	#12, SP		
16	44	A9		06	E1 0006B	MOVL	R0, REC ADDR		1145
OC	44	A9		02	E1 00070	BBC	#6, 68(IRAB), 8\$		1147
	08		44	A9	E8 00075	BBC	#2, 68(IRAB), 7\$		1149
	52	008C		C9	D0 00079	BLBS	68(IRAB), 7\$		1151
				0000V	30 0007E	MOVL	140(IRAB), R2		
						BSBW	RMSMOVE IN VBN		
46	44	A9		06	E0 00081	7\$: BBS	#6, 68(IRAB), 11\$		1155
	0B		44	A9	E9 00086	8\$: BLBC	68(IRAB), 9\$		1174
3D	44	A9		02	E1 0008A	BBC	#2, 68(IRAB), 11\$		1177
	0E		48	A9	B1 0008F	CMPW	72(IRAB), #14		1178
				37	12 00093	BNEQ	11\$		
26	44	A9		02	E1 00095	9\$: BBC	#2, 68(IRAB), 10\$		1189
				C9	D5 0009A	TSTL	144(IRAB)		1191
				20	13 0009E	BEQL	10\$		
		52	0090	C9	D0 000A0	MOVL	144(IRAB), R2		1194
				0000V	30 000A5	BSBW	RMSMOVE IN_VBN		
				56	DD 000A8	PUSHL	REC ADDR		1199
		50	00B4	CA	3C 000AA	MOVZWL	180(IFAB), R0		
				60	B940 3F 000AF	PUSHAB	@96(IRAB)[R0]		
		7E	20	A7	9A 000B3	MOVZBL	32(IDX DFN), -(SP)		1198
				0000G	30 000B7	BSBW	RMSMOVE		
		5E		0C	C0 000BA	ADDL2	#12, SP		
		56		50	D0 000BD	MOVL	R0, REC ADDR		
	08	44	A9	E8 000C0	10\$: BLBS		68(IRAB), 11\$		1209
	52	008C	C9	D0 000C4	MOVL		140(IRAB), R2		1211

: R



1157 1215 1 %SBTTL 'RMSBLD\_NEW\_SIDR'  
1158 1216 1 ROUTINE RMSBLD\_NEW\_SIDR (SIDR\_SIZE) : RL\$COMMON\_LINK NOVALUE =  
1159 1217 1  
1160 1218 1 !++  
1161 1219 1  
1162 1220 1 FUNCTIONAL DESCRIPTION:  
1163 1221 1  
1164 1222 1 This routine builds a new SIDR using the key found in keybuffer 2.  
1165 1223 1  
1166 1224 1 CALLING SEQUENCE:  
1167 1225 1  
1168 1226 1 RMSBLD\_NEW\_SIDR()  
1169 1227 1  
1170 1228 1 INPUT PARAMETERS:  
1171 1229 1  
1172 1230 1 SIDL\_SIZE - full size of the new SIDL  
1173 1231 1  
1174 1232 1 IMPLICIT INPUTS:  
1175 1233 1  
1176 1234 1 BKT\_ADDR - address of SIDL bucket  
1177 1235 1 BKT\$W\_FREESPACE - offset to first free byte in bucket  
1178 1236 1 BKT\$B\_NXTRECID - next available ID in bucket  
1179 1237 1  
1180 1238 1 IDX\_DFN - address of index descriptor  
1181 1239 1 IDX\$V\_DUPKEYS - if set, duplicate alternate keys are allowed  
1182 1240 1 IDX\$V\_KEY\_COMPR - if set, SIDL key compression is enabled  
1183 1241 1 IDX\$B\_KEYSZ - size of alternate key  
1184 1242 1  
1185 1243 1 IFAB - address of IFAB  
1186 1244 1 IFB\$W\_KBUFSZ - size of each of the contiguous keybuffers  
1187 1245 1 IFB\$B\_PLG\_VER - prologue version of file  
1188 1246 1  
1189 1247 1 IRAB - address of IRAB  
1190 1248 1 IRB\$V\_CONT\_BKT - if set, bucket is a continuation bucket  
1191 1249 1 IRB\$V\_DUP\_KEY - if set, not first SIDL with this key value  
1192 1250 1 IRB\$L\_KEYBUF - address of contiguous keybuffers  
1193 1251 1  
1194 1252 1 REC\_ADDR - address of point of insertion of new SIDL  
1195 1253 1  
1196 1254 1 OUTPUT PARAMETERS:  
1197 1255 1 NONE  
1198 1256 1  
1199 1257 1 IMPLICIT OUTPUTS:  
1200 1258 1  
1201 1259 1 IRAB[IRB\$L\_LST\_REC] - address of new SIDL  
1202 1260 1 REC\_ADDR - address of first byte following new SIDL  
1203 1261 1  
1204 1262 1 ROUTINE VALUE:  
1205 1263 1 NONE  
1206 1264 1  
1207 1265 1 SIDE EFFECTS:  
1208 1266 1  
1209 1267 1 If the file is a prologue 1 or 2 file, BKT\_ADDR[BKT\$B\_NXTRECID] will  
1210 1268 1 have been incremented by 1.  
1211 1269 1 If key compression is enabled, and a SIDL follows the new SIDL being  
1212 1270 1 added, then its key will be re-compressed, and the SIDL bucket's  
1213 1271 1 freespace offset pointer updated appropriately.

```
1214 1272 1 !--  
1215 1273 1 !--  
1216 1274 1  
1217 1275 2 BEGIN  
1218 1276 2 EXTERNAL REGISTER  
1219 1277 2 R_BKT_ADDR_STR,  
1220 1278 2 R_IFAB_STR,  
1221 1279 2 R_IRAB_STR,  
1222 1280 2 R_IDX_DFN_STR,  
1223 1281 2 R_REC_ADDR_STR;  
1224 1282 2  
1225 1283 2  
1226 1284 2 LOCAL  
1227 1285 2 FIRST_ARRAY_ADDR : REF BBLOCK;  
1228 1286 2  
1229 1287 2 ! Save the address of the point of insertion of the new SIDR.  
1230 1288 2  
1231 1289 2 IRAB[IRB$L_LST_REC] = .REC_ADDR;  
1232 1290 2  
1233 1291 2 ! Create the record control byte, DUP_CNT (if appropriate), and record ID  
1234 1292 2 fields for the prologue 1 or 2 SIDR. It is not necessary to do this  
1235 1293 2 for prologue 3 SIDRs because they don't have any of these fields. After  
1236 1294 2 this step, REC_ADDR will be pointing to the size field of the new SIDR,  
1237 1295 2 regardless of the file's prologue version.  
1238 1296 2  
1239 1297 2 IF .IFAB[IFB$B_PLG_VER] LSSU PLG$C_VER_3  
1240 1298 2 THEN  
1241 1299 3 BEGIN  
1242 1300 3  
1243 1301 3 ! If duplicate keys are allowed on this index, and this bucket is not a  
1244 1302 3 continuation bucket, then a DUP_CNT field will be created as part  
1245 1303 3 of the overhead of the new SIDR even though this field will not be  
1246 1304 3 maintained. The control byte, ID, and DUP_CNT fields are initialized.  
1247 1305 3  
1248 1306 3 IF .IDX_DFN[IDX$V_DUPKEYS]  
1249 1307 3 AND  
1250 1308 3 NOT .IRAB[IRB$V_CONT_BKT]  
1251 1309 3 THEN  
1252 1310 4 BEGIN  
1253 1311 4 (.REC_ADDR)<0, 8> = 1;  
1254 1312 4 REC_ADDR = .REC_ADDR + 1;  
1255 1313 4  
1256 1314 4 (.REC_ADDR)<0, 8> = .BKT_ADDR[BKT$B_NXTRECID];  
1257 1315 4 REC_ADDR = .REC_ADDR + 1;  
1258 1316 4  
1259 1317 4 (.REC_ADDR)<0, 32> = 0;  
1260 1318 4 REC_ADDR = .REC_ADDR + IRC$C_DCNTSZFLD;  
1261 1319 4 END  
1262 1320 4  
1263 1321 4 ! Either duplicate alternate keys are not allowed, or if they are, this  
1264 1322 4 bucket is a continuation bucket. In either case, it is not necessary  
1265 1323 4 to allocate and initialize a DUP_CNT field. Just the control byte  
1266 1324 4 and the ID field of the new record are initialized.  
1267 1325 4  
1268 1326 3 ELSE  
1269 1327 4 BEGIN  
1270 1328 4 (.REC_ADDR)<0, 8> = IRC$M_NODUPCNT;
```

```
: 1271      1329 4      REC_ADDR = .REC_ADDR + 1;
1272      1330 4
1273      1331 4      (.REC_ADDR)<0, 8> = .BKT_ADDR[BKT$B_NXTRECID];
1274      1332 4      REC_ADDR = .REC_ADDR + 1;
1275      1333 3      END;
1276      1334 3
1277      1335 3      ! Increment the next record ID field in the bucket.
1278      1336 3
1279      1337 3      BKT_ADDR[BKT$B_NXTRECID] = .BKT_ADDR[BKT$B_NXTRECID] + 1;
1280      1338 2      END;
1281      1339 2
1282      1340 2      If key compression is not enabled, copy the new SIDR's key from keybuffer
1283      1341 2      2 into its place in the SIDR bucket, and initialize the new SIDR's record
1284      1342 2      size to be the size of the key. This size field will be updated to its
1285      1343 2      correct value (key size + SIDR array size) by the routine which will
1286      1344 2      create the first element in the SIDR array.
1287      1345 2
1288      1346 3      BEGIN
1289      1347 3
1290      1348 3      GLOBAL REGISTER
1291      1349 3          R_BDB,
1292      1350 3          R_IMPURE,
1293      1351 3          R_RAB;
1294      1352 3
1295      1353 3      IF NOT .IDX_DFN[IDX$V_KEY_COMPR]
1296      1354 3      THEN
1297      1355 4      BEGIN
1298      1356 4          (.REC_ADDR)<0, 16> = .IDX_DFN[IDX$B_KEYSZ];
1299      1357 4          REC_ADDR = .REC_ADDR + IRCS_C_DATSZF[D];
1300      1358 4
1301      1359 4          REC_ADDR = RMSMOVE(.IDX_DFN[IDX$B_KEYSZ], KEYBUF_ADDR(2), .REC_ADDR);
1302      1360 4      END
1303      1361 4
1304      1362 4      If key compression is enabled, the key in keybuffer 2 must first be
1305      1363 4      compressed before it is moved into the SIDR bucket. The key of the
1306      1364 4      following SIDR (if there is one) must also be re-compressed using the
1307      1365 4      new SIDR's key. Finally, the size field must be initialized to the size of
1308      1366 4      the new SIDR's compressed key plus the two bytes of key compression
1309      1367 4      overhead. This size field will be updated to its correct value (key size +
1310      1368 4      SIDR array size) by the routine which will create the first element in
1311      1369 4      the SIDR array.
1312      1370 4
1313      1371 3      ELSE
1314      1372 4      BEGIN
1315      1373 4
1316      1374 4      LOCAL
1317      1375 4          KEYBUF : REF BBLOCK;
1318      1376 4
1319      1377 4      ! Compress the new SIDR's key within its keybuffer - keybuffer 2.
1320      1378 4
1321      1379 4          KEYBUF = KEYBUF_ADDR(2);
1322      1380 4          RMSCOMPRESS_KEY~(.KEYBUF);
1323      1381 4
1324      1382 4      ! Fix up the front compression of the following record because it might
1325      1383 4      change due to the insertion of this new SIDR. Of course, this step
1326      1384 4      is unnecessary, if there is no following record.
1327      1385 4
```

```

1328    1386 5      IF (.REC_ADDR + .SIDR_SIZE LSSU .BKT_ADDR + .BKT_ADDR[BKT$W_FREESPACE])
1329    1387 4      THEN
1330    1388 4          RMSRECOMPRESS_KEY (.KEYBUF, .REC_ADDR + IRC$C_DATSZFLD + .SIDR_SIZE);
1331    1389 4
1332    1390 4      ! Initialize the size field of the new SIDR to the size of the
1333    1391 4      compressed key plus the two bytes of key compression overhead, and
1334    1392 4      move the new SIDR key into the bucket including the two bytes of
1335    1393 4      compression overhead.
1336    1394 4
1337    1395 4      (.REC_ADDR)<0,16> = .KEYBUF[KEY_LEN] + IRC$C_KEYCMPOVH;
1338    1396 4      REC_ADDR = .REC_ADDR + IRC$C_DATSZFLD;
1339    1397 4
1340    1398 4      REC_ADDR = RMSMOVE (.KEYBUF[KEY_LEN] + 2, .KEYBUF, .REC_ADDR);
1341    1399 3      END;
1342    1400 2      END;
1343    1401 2
1344    1402 2      ! Save the address of what will be the first array element in this SIDR
1345    1403 2      before creating the array.
1346    1404 2
1347    1405 2      FIRST_ARRAY_ADDR = .REC_ADDR;
1348    1406 2      RM$ADD_TO_ARRAY();
1349    1407 2
1350    1408 2      ! If this is a prologue 3 file, and this is the first such SIDR with this
1351    1409 2      key value in the file, then set a bit indicating this in the control byte
1352    1410 2      of the SIDR array's first element.
1353    1411 2
1354    1412 2      IF .IFAB[IFB$B_PLG_VER] GEQU PLG$C_VER_3
1355    1413 2          AND
1356    1414 2          NOT .IRAB[IRB$V_DUP_KEY]
1357    1415 2      THEN FIRST_ARRAY_ADDR[IRC$V_FIRST_KEY] = 1;
1358    1416 2
1359    1417 1      END;

```

0914 8F BB 00000 RMSBLD_NEW_SIDR:							
	4C	A9	03	00B7	56 D0 00004	PUSHR #^M<R2,R4,R8,R11>	1216
					CA 91 00008	MOVL REC ADDR, 76(IRAB)	1289
					1E 1E 0000D	CMPB 183(IFAB), #3	1297
OB	44	10		1C	A7 E9 0000F	BGEQU 3\$	1306
		A9			04 E0 00013	BLBC 28(IDX_DFN), 1\$	1308
		86			01 90 00018	BBS #4, 68(IRAB), 1\$	1311
		86	06		A5 90 0001B	MOV B #1, (REC_ADDR)+	1314
					86 D4 0001F	MOV B 6(BKT_ADDR), (REC_ADDR)+	1317
					07 11 00021	CLRL (REC_ADDR)+	1306
						BRB 2\$	1328
		86			10 90 00023 1\$:	MOV B #16, (REC_ADDR)+	1331
		86	06		A5 90 00026	MOV B 6(BKT_ADDR), (REC_ADDR)+	1337
					06 A5 96 0002A 2\$:	INC B 6(BKT_ADDR)	1353
15	1C	A7			06 E0 0002D 3\$:	BBS #6, 28(IDX_DFN), 4\$	1356
		86	20		A7 9B 00032	MOVZBW 32(IDX_DFN), (REC_ADDR)+	1359
					56 DD 00036	PUSHL REC ADDR	
		50	00B4		CA 3C 00038	MOVZWL 180(IFAB), R0	
					60 B940 9F 0003D	PUSHAB @96(IRAB)[R0]	
		7E			20 A7 9A 00041	MOVZBL 32(IDX_DFN), -(SP)	

RM3MISPUT  
V04-000

RMSBLD\_NEWSIDR

G 3  
16-Sep-1984 01:51:28  
14-Sep-1984 13:01:29VAX-11 Bliss-32 V4.0-742  
[RMS.SRC]RM3MISPUT.B32;1Page 34  
(7)RM3  
V04

52	00B4	37	11	00045		BRB	6\$		1379	
52	60	CA	3C	00047	4\$:	MOVZWL	180(IFAB), KEYBUF		1	
50		A9	C0	0004C		ADDL2	96(IRAB), KEYBUF		1	
		52	D0	00050		MOVL	KEYBUF, R0		1	
		0000G	30	00053		BSBW	RMS\$COMPRESS KEY		1	
50	56	14	AE	C1	00056	ADDL3	SIDR SIZE, REC_ADDR, R0		1	
	51	04	A5	3C	0005B	MOVZWL	4(BKT_ADDR), RT		1	
	51		55	C0	0005F	ADDL2	BKT_ADDR, R1		1	
	51		50	D1	00062	CMPL	R0, R1		1	
			0A	1E	00065	BGEQU	5\$		1	
	51	02	A0	9E	00067	MOVAB	2(R0), R1		1	
	50		52	D0	0006B	MOVL	KEYBUF, R0		1	
			0000G	30	0006E	BSBW	RMS\$REC\$COMPRESS KEY		1	
	50		62	9A	00071	5\$::	MOVZBL	(KEYBUF), R0		1
	50		02	C0	00074	ADDL2	#2, R0		1	
	86		50	B0	00077	MOVW	R0, (REC_ADDR)+		1	
		0045	8F	BB	0007A	PUSHR	#^M<R0,R2,R6>		1	
			0000G	30	0007E	6\$::	BSBW	RMS\$MOVE	1	
	5E		0C	C0	00081	ADDL2	#12, SP		1	
	56		50	D0	00084	MOVL	R0, REC_ADDR		1	
	52		56	D0	00087	MOVL	REC_ADDR, FIRST_ARRAY_ADDR		1	
	03	00B7	FC87	30	0008A	BSBW	RMS\$ADD_TO_ARRAY		1	
			CA	91	0008D	CMPB	183(IFAB), #3		1	
			08	1F	00092	BLSSU	7\$		1	
	04	43	A9	E8	00094	BLBS	67(IRAB), 7\$		1	
	62	80	8F	88	00098	BISB2	#128, (FIRST_ARRAY_ADDR)		1	
		0914	8F	BA	0009C	POPR	#^M<R2,R4,R8,R11>		1	
				05	000A0	RSB			1	

: Routine Size: 161 bytes, Routine Base: RMS\$RMS3 + 02EC

```
1361 1418 1 %SBTTL 'RMSFOOLED_YUH'  
1362 1419 1 GLOBAL ROUTINE RMSFOOLED_YUH : RL$RABREG_7 =  
1363 1420 1  
1364 1421 1 !++  
1365 1422 1  
1366 1423 1 FUNCTIONAL DESCRIPTION:  
1367 1424 1  
1368 1425 1 routine to do all the probing, etc. to make sure the user hasn't  
1369 1426 1 been fooling around w/ rsz, rbf or the buffer describing his record  
1370 1427 1 between the time i started to position for insert and i go to insert  
1371 1428 1  
1372 1429 1  
1373 1430 1 CALLING SEQUENCE:  
1374 1431 1 bsbw rm$fooled_yuh  
1375 1432 1  
1376 1433 1 INPUT PARAMETERS:  
1377 1434 1 NONE  
1378 1435 1  
1379 1436 1 IMPLICIT INPUTS:  
1380 1437 1 irab -- mode, address of internal key buffer to check against  
1381 1438 1 IFAB -- kbufsz  
1382 1439 1 rab -- rsz, rbf  
1383 1440 1 idx_dfn -- minrecsz, keysz  
1384 1441 1  
1385 1442 1 OUTPUT PARAMETERS:  
1386 1443 1 NONE  
1387 1444 1  
1388 1445 1 IMPLICIT OUTPUTS:  
1389 1446 1 NONE  
1390 1447 1  
1391 1448 1 ROUTINE VALUE:  
1392 1449 1 rmssuc if the user has been considerate  
1393 1450 1 rsz or rbf if not  
1394 1451 1  
1395 1452 1 SIDE EFFECTS:  
1396 1453 1 ap is clobbered  
1397 1454 1  
1398 1455 1 --  
1399 1456 1  
1400 1457 2 BEGIN  
1401 1458 2  
1402 1459 2 EXTERNAL REGISTER  
1403 1460 2 R_IFAB_STR,  
1404 1461 2 R_IRAB_STR,  
1405 1462 2 R_RAB_STR,  
1406 1463 2 R_IDX_DFN_STR;  
1407 1464 2  
1408 1465 2 LOCAL  
1409 1466 2 RBF_ADDR,  
1410 1467 2 STATUS;  
1411 1468 2  
1412 1469 2 BUILTIN  
1413 1470 2 AP;  
1414 1471 2  
1415 1472 2 ! do all the probing and comparing to make sure the user hasn't been  
1416 1473 2 fooling around w/ the record buffer  
1417 1474 2
```

```

1418    1475  2
1419    1476  2 IF .RAB[RAB$W_RSZ] LSSU .IDX_DFN[IDX$W_MINRECSZ]
1420    1477  2 THEN
1421    1478  2     RETURN RMSERR(RSZ);
1422    1479  2
1423    1480  2 RBF_ADDR = .RAB[RAB$L_RBF];
1424    1481  2 IF RMSNOREAD_LONG(.RAB[RAB$W_RSZ], .RBF_ADDR, .IRAB[IRB$B_MODE])
1425    1482  2 THEN
1426    1483  2     RETURN RMSERR(RBF);
1427    1484  2
1428    1485  2 AP = 2;
1429    1486  2
1430    1487  2 ! If keys do not match, return error
1431    1488  2 !
1432    1489  2
1433    1490  2 IF RMSCOMPARE_KEY(.RBF_ADDR,
1434          KEYBUF_ADDR(3),
1435          .IDX_DFN[IDX$B_KEYSZ])
1436    1493  2 THEN
1437    1494  2     RETURN RMSERR(RBF);
1438    1495  2
1439    1496  3 RETURN RMSSUC()
1440    1497  3
1441    1498  1 END;

```

OC BB 00000 RMS\$FOOLED_YUH::					
22	A7	22	A8	B1 00002	PUSHR #^M<R2,R3>
		07	07	1E 00007	CMPW 34(RAB), 34(IDX_DFN)
50	86A4	8F	3C 00009	BGEQU 1\$	
		3B	11 0000E	MOVZWL #34468, R0	
52	28	A8	D0 00010	BRB 4\$	
7E	0A	A9	9A 00014	MOVL 40(RAB), RBF_ADDR	
		52	DD 00018	MOVZBL 10(IRAB), -(SP)	
7E	22	A8	3C 0001A	PUSHL RBF_ADDR	
		0000G	30 0001E	MOVZWL 34(RAB), -(SP)	
5E	0C	CO 00021	BSBW RMSNOREAD_LONG		
1A	50	E8 00024	ADDL2 #12, SP		
5C	02	D0 00027	BLBS R0, 2\$		
50	00B4	CA 3C 0002A	MOVL #2, AP		
53	60	B940 3E 0002F	MOVZWL 180(IFAB), R0		
50	20	A7 9A 00034	MOVAW @96(IRAB)[R0], R3		
51		52 D0 00038	MOVZBL 32(IDX_DFN), R0		
		0000G 30 0003B	MOVL RBF_ADDR, R1		
07		50 E9 0003E	BSBW RMSCOMPARE_KEY		
50	8654	8F 3C 00041	BLBC R0, 3\$		
		03 11 00046	MOVZWL #34388, R0		
50		01 D0 00048	BRB 4\$		
		0C BA 0004B	MOVL #1, R0		
		05 0004D	POPR #^M<R2,R3>		
			RSB		

; Routine Size: 78 bytes, Routine Base: RMS\$RMS3 + 038D

RM3MISPUT  
V04-000

RMSFOOLED\_YUH

J 3  
16-Sep-1984 01:51:28  
14-Sep-1984 13:01:29      VAX-11 Bliss-32 v4.0-742  
[RMS.SRC]RM3MISPUT.B32;1

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(8)

RM3  
V04

1443 1499 1 %SBTTL 'RMSINS REC'  
1444 1500 1 GLOBAL ROUTINE RMSINS\_REC (BUCKET, REC\_SZ) : RL\$RABREG\_67 =  
1445 1501 1  
1446 1502 1 !++  
1447 1503 1  
1448 1504 1 FUNCTIONAL DESCRIPTION:  
1449 1505 1  
1450 1506 1 Insert new record into bucket. Check for ID availability if required,  
1451 1507 1 make room by spreading existing records apart, and call appropriate  
1452 1508 1 routine to build the record.  
1453 1509 1  
1454 1510 1 CALLING SEQUENCE:  
1455 1511 1  
1456 1512 1 RMSINS\_REC()  
1457 1513 1  
1458 1514 1 INPUT PARAMETERS:  
1459 1515 1  
1460 1516 1 BKT\_ADDR - address of bucket in which to insert record  
1461 1517 1  
1462 1518 1 REC\_SZ - total size of record to be inserted  
1463 1519 1  
1464 1520 1 IMPLICIT INPUTS:  
1465 1521 1  
1466 1522 1 REC\_ADDR - is where to build the record  
1467 1523 1  
1468 1524 1 IRAB - internal stream context  
1469 1525 1 [dups\_seen] - this is a duplicate record  
1470 1526 1 [spl\_idx] - left hand part only of index record  
1471 1527 1 [pos\_ins] - offset to position of insert  
1472 1528 1  
1473 1529 1 IFAB - internal file context, used by called routines  
1474 1530 1  
1475 1531 1 IDX\_DFN - index descriptor, used by called routines  
1476 1532 1  
1477 1533 1 OUTPUT PARAMETERS:  
1478 1534 1 NONE  
1479 1535 1  
1480 1536 1 IMPLICIT OUTPUTS:  
1481 1537 1 NONE  
1482 1538 1  
1483 1539 1 ROUTINE VALUE:  
1484 1540 1 SUCCESS  
1485 1541 1 low bit clear if no id's available  
1486 1542 1  
1487 1543 1  
1488 1544 1 --  
1489 1545 1  
1490 1546 2 BEGIN  
1491 1547 2  
1492 1548 2 EXTERNAL REGISTER  
1493 1549 2 COMMON RAB\_STR,  
1494 1550 2 R\_REC\_ADDR\_STR,  
1495 1551 2 R\_IDX\_DFN\_STR;  
1496 1552 2  
1497 1553 2 GLOBAL REGISTER  
1498 1554 2 R\_BKT\_ADDR\_STR;  
1499 1555 2

```
1500      1556 2 LOCAL
1501      1557 2     TEMP   : BBLOCK [6];
1502      1558 2
1503      1559 2 MACRO
1504      1560 2     OLD_OVHD    = TEMP[0,0,16,0] %;
1505      1561 2     LENGTH      = TEMP[2,0,32,0] %;
1506      1562 2     KEY_SZ      = REC_SZ<0,16>%;
1507      1563 2     VBN_SZ      = REC_SZ<16,16>%;
1508      1564 2
1509      1565 2     BKT_ADDR = .BUCKET;
1510      1566 2
1511      1567 2     OLD_OVHD = 0;
1512      1568 2
1513      1569 2 IF .IFAB [IFB$B_PLG_VER] LSSU PLG$C_VER_3
1514      1570 2 THEN
1515      1571 2
1516      1572 2     BEGIN
1517      1573 2
1518      1574 2     IF .BKT_ADDR [BKT$B_LEVEL] EQL 0
1519      1575 2 THEN
1520      1576 4     BEGIN
1521      1577 4
1522      1578 4     IF NOT .IRAB[IRB$V_DUPS_SEEN]
1523      1579 4     AND
1524      1580 5         (.BKT_ADDR [BKT$B_NXTRECID] EQL 0
1525      1581 5         OR
1526      1582 5         (.BKT_ADDR [BKT$B_NXTRECID] GTRU .BKT_ADDR [BKT$B_LSTRECID]))
1527      1583 4 THEN
1528      1584 4     RETURN_ON_ERROR (RM$BKT_SORT(.BKT_ADDR));
1529      1585 4
1530      1586 4 END
1531      1587 3 ELSE
1532      1588 3     IF NOT .IRAB[IRB$V_SPL_IDX]
1533      1589 3     THEN
1534      1590 3     OLD_OVHD = .REC_ADDR[IRC$V_PTRSZ] + IRC$C_IDXPTRBAS +
1535      1591 3           IRC$C_IDXOVHDSZ;
1536      1592 2 END;
1537      1593 2
1538      1594 2 ! Figure out length of data to move and do it if non-zero - it will
1539      1595 2 be zero when SDR record is being inserted at the end of the bucket
1540      1596 2
1541      1597 2
1542      1598 2 LENGTH = .BKT_ADDR[BKT$W_FREESPACE] - .IRAB[IRB$W_POS_INS] - .OLD_OVHD;
1543      1599 2
1544      1600 3 IF (.LENGTH GTR 0)
1545      1601 3     AND
1546      1602 3     (.KEY_SZ GTR 0)
1547      1603 2 THEN
1548      1604 3     BEGIN
1549      1605 3
1550      1606 3     GLOBAL REGISTER
1551      1607 3         R_BDB;
1552      1608 3
1553      1609 3     RMSMOVE(.LENGTH, .REC_ADDR + .OLD_OVHD,
1554      1610 3           .REC_ADDR + .KEY_SZ + .OLD_OVHD);
1555      1611 2 END;
1556      1612 2
```

```

1557 1613 2 | Update the freespace pointer to reflect new size of data area
1558 1614 2
1559 1615 2 BKT_ADDR [BKT$W_FREESPACE] = .BKT_ADDR [BKT$W_FREESPACE] + .KEY_SZ;
1560 1616 2
1561 1617 2 IF (.IFAB [IFBSB_PLG_VER] GEQU PLG$C_VER_3)
1562 1618 2 AND
1563 1619 2 (.BKT_ADDR[BKT$B_LEVEL] GTRU 0)
1564 1620 2 THEN
1565 1621 2
1566 1622 2 | We must spread apart the VBN list for prologue 3 index records, also.
1567 1623 2
1568 1624 2 RMSSHFT_VBNS();
1569 1625 2
1570 1626 2 | Insert new index record and return success.
1571 1627 2
1572 1628 2 IF .BKT_ADDR [BKT$B_LEVEL] NEQU 0
1573 1629 2 THEN
1574 1630 2 BEGIN
1575 1631 2 RMSBLD_IDX_REC (.REC_SZ);
1576 1632 3 RETURN 1
1577 1633 2
1578 1634 2
1579 1635 2 | If duplicates have been seen, then RMS merely has to add a new SIDR
1580 1636 2 array element to an existing array, otherwise, it must create a new
1581 1637 2 SIDR.
1582 1638 2
1583 1639 2 IF .IRAB[IRBSV_DUPS_SEEN]
1584 1640 2 THEN
1585 1641 2 RMSADD_TO_ARRAY()
1586 1642 2 ELSE
1587 1643 2 RMSBLD_NEW_SIDR (.REC_SZ);
1588 1644 2
1589 1645 2 RETURN 1
1590 1646 1 END;

```

		30 BB 00000 RMSINS_REC::		
5E	08	C2 00002	PUSHR	#^M<R4,R5>
55	14	AE D0 00005	SUBL2	#8, SP
03	00B7	6E B4 00009	MOVL	BUCKET, BKT_ADDR
		CA 91 0000B	CLRW	TEMP
		30 1E 00010	CMPB	183(IFAB), #3
	0C	A5 95 00012	BGEQU	3\$
		1E 12 00015	TSTB	12(BKT_ADDR)
	44	A9 95 00017	BNEQ	2\$
		26 19 0001A	TSTB	68(IRAB)
	06	A5 95 0001C	BLSS	3\$
		07 13 0001F	TSTB	6(BKT_ADDR)
07	A5	06 A5 91 00021	BEQL	1\$
		1A 1B 00026	CMPB	6(BKT_ADDR), 7(BKT_ADDR)
		55 DD 00028 1\$: FD81	BLEQU	3\$
	5E	30 0002A	PUSHL	BKT_ADDR
		04 C0 0002D	BSBW	RMSBKT_SORT
			ADDL2	#4, SP

: 1500  
: 1565  
: 1567  
: 1569  
: 1574  
: 1578  
: 1580  
: 1582  
: 1584

RM3I  
V04-

		0F	50	E8 00030	BLBS	STATUS, 3\$	
		09	77	11 00033	BRB	10\$	
50	66	02	A9	E8 00035	BLBS	68(IRAB), 3\$	1588
		50	00	EF 00039	EXTZV	#0, #2, (REC_ADDR), R0	1590
		50	03	A1 0003E	ADDW3	#3, R0, TEMP	
		50	04	A5 3C 00042	MOVZWL	4(BKT_ADDR), R0	1598
		51	48	A9 3C 00046	MOVZWL	72(IRAB), R1	
		50	51	C2 0004A	SUBL2	R1, R0	
02	AE	54	6E	3C 0004D	MOVZWL	TEMP, R4	
		50	54	C3 00050	SUBL3	R4, R0, TEMP+2	1600
			1E	15 00055	BLEQ	4\$	1602
			18	AE B5 00057	TSTW	REC_SZ	
		50	18	AE 3C 0005C	BEQL	4\$	1610
		50	56	C0 00060	MOVZWL	REC_SZ, R0	
		51	6E	3C 00063	ADDL2	REC_ADDR, R0	
			6140	9F 00066	MOVZWL	TEMP, R1	
			6146	9F 00069	PUSHAB	(R1)[R0]	
		0A	AE	DD 0006C	PUSHAB	(R1)[REC_ADDR]	1609
			0000G	30 0006F	PUSHL	TEMP+2	
			OC	C0 00072	BSBW	RMSMOVE	
04	5E	04	A5	A0 00075	ADDL2	#12, SP	
		03	00B7	CA 91 0007A	ADDW2	REC_SZ, 4(BKT_ADDR)	1615
			08	1F 0007F	CMPB	183TIFAB), #3	1617
			0C	A5 95 00081	BLSSU	5\$	
			03	13 00084	TSTB	12(BKT_ADDR)	1619
			0000V	30 00086	BEQL	5\$	
			0C	A5 95 00089	BSBW	RMSHFT_VBNS	1624
			08	13 0008C	TSTB	12(BKT_ADDR)	1628
		18	AE	DD 0008E	BEQL	6\$	
			FDAC	30 00091	PUSHL	REC_SZ	1631
			10	11 00094	BSBW	RMSBLD_IDX_REC	
		44	A9	95 00096	BRB	8\$	
			05	18 00099	TSTB	68(IRAB)	1639
			FB87	30 0009B	BGEQ	7\$	
			09	11 0009E	BSBW	RMSADD_TO_ARRAY	1641
		18	AE	DD 000A0	BRB	9\$	
		5E	FE6B	30 000A3	PUSHL	REC_SZ	1643
			04	C0 000A6	BSBW	RMSBLD_NEW_SIDR	
		50	01	D0 000A9	ADDL2	#4, SP	1645
		5E	08	C0 000AC	9\$:	#1, R0	
			30	BA 000AF	MOVL	#8, SP	1646
			05	000B1	ADDL2	#^M<R4,R5>	
					POPR		
					RSB		

; Routine Size: 178 bytes, Routine Base: RMSRMS3 + 03DB

```

1592    1647 1 %SBTTL 'RMSMOVE_IN_VBN'
1593    1648 1 ROUTINE RMSMOVE_IN_VBN (VBN) : RL$MOVE_IN_VBN NOVALUE =
1594    1649 1
1595    1650 1 ++
1596    1651 1
1597    1652 1 FUNCTIONAL DESCRIPTION:
1598    1653 1
1599    1654 1 set pointer size field and move in VBN, updating REC_ADDR
1600    1655 1
1601    1656 1 CALLING SEQUENCE:
1602    1657 1     RMSMOVE_IN_VBN(VBN)
1603    1658 1
1604    1659 1 INPUT PARAMETERS:
1605    1660 1     VBN - VBN to move in
1606    1661 1
1607    1662 1 IMPLICIT INPUTS:
1608    1663 1     REC_ADDR - address to move value to
1609    1664 1
1610    1665 1 OUTPUT PARAMETERS:
1611    1666 1     NONE
1612    1667 1
1613    1668 1 IMPLICIT OUTPUTS:
1614    1669 1     REC_ADDR - points one byte beyond where value placed
1615    1670 1
1616    1671 1 ROUTINE VALUE:
1617    1672 1     NONE
1618    1673 1
1619    1674 1 SIDE EFFECTS:
1620    1675 1     NONE
1621    1676 1
1622    1677 1 --
1623    1678 1
1624    1679 2 BEGIN
1625    1680 2
1626    1681 2 EXTERNAL REGISTER
1627    1682 2     R_REC_ADDR;
1628    1683 2
1629    1684 2 GLOBAL REGISTER
1630    1685 2     COMMON_RABREG,
1631    1686 2     R_BDB,
1632    1687 2     R_IDX_DFN;
1633    1688 2
1634    1689 3 BEGIN
1635    1690 3
1636    1691 3 REGISTER
1637    1692 3     SIZE = 1      : BYTE;
1638    1693 3
1639    1694 3     ! get vbn size and set pointer field
1640    1695 3
1641    1696 3     SIZE = RMSVBN_SIZE(.VBN);
1642    1697 3     (.REC_ADDR)<0, 8> = .SIZE - 2;
1643    1698 3     REC_ADDR = .REC_ADDR + 1;
1644    1699 3
1645    1700 3     ! stick in bucket pointer
1646    1701 3
1647    1702 3     (.REC_ADDR)<0, .SIZE*8> = .VBN;
1648    1703 3     REC_ADDR = .REC_ADDR + .SIZE;

```

RM3MISPUT  
V04-000

RM\$MOVE\_IN\_VBN

C 4  
16-Sep-1984 01:51:28  
14-Sep-1984 13:01:29 VAX-11 Bliss-32 V4.0-742  
[RMS.SRC]RM3MISPUT.B32;1

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(10)

: 1649 1704 2 END:  
: 1650 1705 1 END:

		OF94	8F BB 00000 RM\$MOVE_IN_VBN:			
			0000V 30 00004	PUSHR	#^M<R2,R4,R7,R8,R9,R10,R11>	1696
		5E	04 C0 00007	BSBW	RMSVBN_SIZE	
		51	50 90 0000A	ADDL2	#4, SP-	
	86	51	02 83 0000D	MOVB	R0, SIZE	
		50	51 9A 00011	SUBB3	#2, SIZE, (REC_ADDR)+	1697
		50	08 C4 00014	MOVZBL	SIZE, R0	1702
66	50	00	52 F0 00017	MULL2	#8, R0	
		50	51 9A 0001C	INSV	VBN, #0, R0, (REC_ADDR)	
		56	50 C0 0001F	MOVZBL	SIZE, R0	1703
			OF90 8F BA 00022	ADDL2	R0, REC ADDR	
			05 00026	POPR	#^M<R4,R7,R8,R9,R10,R11>	1705
				RSB		

; Routine Size: 39 bytes, Routine Base: RM\$RMS3 + 048D

```
: 1652      1706 1 %SBTTL 'RMSNEW_VBN_BYTES'  
1653      1707 1 ROUTINE RMSNEW_VBN_BYTES (VBN_SIZE) : RL$NEW_VBN_BYTES =  
1654      1708 1 ++  
1655      1709 1  
1656      1710 1 RMSNEW_VBN_BYTES  
1657      1711 1  
1658      1712 1 This routine returns the number of additional bytes which will  
1659      1713 1 be needed if all the VBNs grow to VBN_SIZE  
1660      1714 1  
1661      1715 1 CALLING SEQUENCE:  
1662      1716 1  
1663      1717 1 RMSNEW_VBN_BYTES (VBN_SIZE)  
1664      1718 1  
1665      1719 1 INPUT PARAMETERS:  
1666      1720 1  
1667      1721 1 VBN_SIZE - potential number of bytes per VBN  
1668      1722 1  
1669      1723 1 IMPLICIT INPUTS:  
1670      1724 1  
1671      1725 1 BKT_ADDR - address of current bucket  
1672      1726 1 IDX_DFN - index descriptor structure  
1673      1727 1  
1674      1728 1 OUTPUT PARAMETERS:  
1675      1729 1  
1676      1730 1 NONE  
1677      1731 1  
1678      1732 1 IMPLICIT OUTPUTS:  
1679      1733 1 NONE  
1680      1734 1  
1681      1735 1 ROUTINE VALUE:  
1682      1736 1  
1683      1737 1 R0 - number of new bytes which will be needed if all the  
1684      1738 1 VBNs in the bucket grow to be VBN_SIZE long  
1685      1739 1  
1686      1740 1 --  
1687      1741 1  
1688      1742 2 BEGIN  
1689      1743 2  
1690      1744 2 MACRO  
1691      1745 2     FREESPC = 0,0,16,0 %;  
1692      1746 2  
1693      1747 2 GLOBAL REGISTER  
1694      1748 2     R_BKT_ADDR_STR,  
1695      1749 2     R_IDX_DFN_STR;  
1696      1750 2  
1697      1751 2 LOCAL  
1698      1752 2     CUR_VBN_SZ,  
1699      1753 2     FIRST_VBN      : REF BBLOCK,  
1700      1754 2     NUM_VBNS;  
1701      1755 2  
1702      1756 2     CUR_VBN_SZ = .BKT_ADDR [BKT$V_PTR_SZ] + 2;  
1703      1757 2  
1704      1758 2     ! If VBNs didn't grow return  
1705      1759 2  
1706      1760 2     IF .VBN_SIZE LEQU .CUR_VBN_SZ  
1707      1761 2     THEN  
1708      1762 2     RETURN 0;
```

```

: 1709    1763  2
: 1710    1764  2 ! Number of new bytes = (new_size - old_size) * num_vbns
: 1711    1765  2
: 1712    1766  2 FIRST_VBN = .BKT_ADDR + (.IDX_DFN [IDX$B_IDXBKTSZ] * 512) - BKT$C_ENDOVHD;
: 1713    1767  2 NUM_VBNS = (.FIRST_VBN - (.BKT_ADDR + .FIRST_VBN [FREESPC]))/ (.BKT_ADDR [BKT$V_PTR_SZ] + 2);
: 1714    1768  2 RETURN (.VBN_SIZE = .CUR_VBN_SZ) * .NUM_VBNS
: 1715    1769  2
: 1716    1770  1 END:
INFO#250      L1:1756
Referenced REGISTER symbol BKT_ADDR is probably not initialized
INFO#250      L1:1765
Referenced REGISTER symbol IDX_DFN is probably not initialized

```

			00A4 8F BB 00000 RMSNEW_VBN_BYTES:		
51	OD A5		PUSHR #^M<R2,R5,R7>		1707
		02	EXTZV #3, #2, 13(BKT_ADDR), R1		1756
		51	ADDL2 #2, R1		
		52	MOVL R1, CUR_VBN_SZ		
		52	CMPL VBN_SIZE, CUR_VBN_SZ		1760
		10	BLEQU 1\$		
57		57	MOVZBL 22(IDX_DFN), R7		1766
		57	ASHL #9, R7, R7		
		50	MOVAB -4(R7)[BKT_ADDR], FIRST_VBN		
		57	MOVZWL (FIRST_VBN), R7		1767
		60	ADDL2 R7, BKT_ADDR		
		55	SUBL2 R5, R0		
55	50	50	DIVL3 R1, R0, NUM_VBNS		
	10	AE	SUBL3 CUR_VBN_SZ, VBN_SIZE, R0		1768
		50	MULL2 NUM_VBNS, R0		
		51	BRB 2\$		
		52	CLRL R0		1770
		53	POPR #^M<R2,R5,R7>		
		02	RSB		
		50			
		54			
		00A4 8F BA 0003C 1\$: 05 00040 2\$:			

: Routine Size: 65 bytes, Routine Base: RM\$RMS3 + 04B4

1718 1771 1 %SBTTL 'RMSRECORD\_SIZE'  
 1719 1772 1 GLOBAL ROUTINE RMSRECORD\_SIZE : RL\$RABREG\_567 =  
 1720 1773 1  
 1721 1774 1 ++  
 1722 1775 1  
 1723 1776 1 FUNCTIONAL DESCRIPTION:  
 1724 1777 1 Calculate record size of SIDR or index records  
 1725 1778 1  
 1726 1779 1  
 1727 1780 1 CALLING SEQUENCE:  
 1728 1781 1  
 1729 1782 1 RMSRECORD\_SIZE()  
 1730 1783 1  
 1731 1784 1 INPUT PARAMETERS:  
 1732 1785 1 NONE  
 1733 1786 1  
 1734 1787 1 IMPLICIT INPUTS:  
 1735 1788 1 IFAB - for RMSRECORD VBN  
 1736 1789 1 IRAB - pointer to internal RAB  
 1737 1790 1 [ STOPLEVEL ] - current level in tree  
 1738 1791 1 [ DUPS\_SEEN ] - flag set if duplicates seen  
 1739 1792 1 [ CONT\_BKT ] - this is continuation bucket  
 1740 1793 1 [ BIG\_SPLIT ] - flag set if more than two bucket split  
 1741 1794 1 [ VBN\_LEFT ] - VBN of leftmost bucket in split  
 1742 1795 1 [ VBN\_RIGHT ] - VBN of rightmost bucket in split  
 1743 1796 1 [ VBN\_MID ] - VBN of middle bucket in 3-4 bucket split  
 1744 1797 1 [ REC\_COUNT ] - which record this is in the bucket  
 1745 1798 1 [ POS\_INS ] - position of insert of this record  
 1746 1799 1 [ KEYBUF ] - pointer to the 5 contiguous keybuffers  
 1747 1800 1 [ PUTUP\_VBN ] - VBN of new primary data record (for SIDR ptr)  
 1748 1801 1 [ PUTUP\_ID ] - ID of new primary data record (for SIDR ptr)  
 1749 1802 1 IDX\_DFN - pointer to index descriptor structure  
 1750 1803 1 [ DUPKEYS ] - duplicate keys are allowed  
 1751 1804 1 [ KEYSZ ] - size of key  
 1752 1805 1 REC\_ADDR - position of insert (used on index levels only)  
 1753 1806 1 BKT\_ADDR - pointer to current bucket  
 1754 1807 1  
 1755 1808 1 OUTPUT PARAMETERS:  
 1756 1809 1 NONE  
 1757 1810 1  
 1758 1811 1 IMPLICIT OUTPUTS:  
 1759 1812 1 NONE  
 1760 1813 1  
 1761 1814 1 ROUTINE VALUE:  
 1762 1815 1 size of record required  
 1763 1816 1  
 1764 1817 1 PROLOG 1 & 2: Size is returned as a single quantity  
 1765 1818 1  
 1766 1819 1 PROLOG 3 Index: Size is returned as two contiguous words.  
 1767 1820 1 High order word contains number of VBN bytes  
 1768 1821 1 Low order word contains number of key bytes  
 1769 1822 1  
 1770 1823 1 PROLOG 3 SIDR: Size is returned as a single quantity  
 1771 1824 1  
 1772 1825 1 SIDE EFFECTS:  
 1773 1826 1  
 1774 1827 1 sets AP = 1 when checking down pointer on empty bucket cases

```
1775 1828 1 | (index level).
1776 1829 1 | Keybuffer 5 will be used if index/key compression is enabled.
1777 1830 1 | Keybuffer 4 will be used if index compression is enabled and a big
1778 1831 1 | split occurred.
1779 1832 1 |
1780 1833 1 | --
1781 1834 1 | BEGIN
1782 1835 2 | BUILTIN
1783 1836 2 | AP:
1784 1837 2 | EXTERNAL REGISTER
1785 1838 2 | COMMON RAB_STR,
1786 1839 2 | R_BKT_ADDR_STR,
1787 1840 2 | R_REC_ADDR_STR,
1788 1841 2 | R_IDX_DFN_STR;
1789 1842 2 |
1790 1843 2 |
1791 1844 2 |
1792 1845 2 |
1793 1846 2 | GLOBAL REGISTER
1794 1847 2 | R_BDB;
1795 1848 2 |
1796 1849 2 | LOCAL
1797 1850 2 | SIZE;
1798 1851 2 |
1799 1852 2 | ! RMS is to determine the size of a new SIDR or SIDR array element.
1800 1853 2 |
1801 1854 2 | IF .IRAB[IRB$B_STOPLEVEL] EQ 0
1802 1855 2 | THEN
1803 1856 3 | BEGIN
1804 1857 3 |
1805 1858 3 | ! Calculate the size of the record pointer part of record (which will
1806 1859 3 | always be present).
1807 1860 3 |
1808 1861 3 | SIZE = RM$VBN_SIZE(.IRAB[IRB$L_PUTUP_VBN]);
1809 1862 3 |
1810 1863 3 | IF .IFAB[IFB$B_PLG_VER] LSSU PLG$C_VER_3
1811 1864 3 | THEN
1812 1865 3 | SIZE = .SIZE + IRC$C_DATOVHDSZ
1813 1866 3 | ELSE
1814 1867 3 | SIZE = .SIZE + IRC$C_DATOVHSZ;
1815 1868 3 |
1816 1869 3 |
1817 1870 3 | ! If a new SIDR is required, then add in the size of the key, and the
1818 1871 3 | overhead associated with a SIDR.
1819 1872 3 |
1820 1873 3 | IF NOT .IRAB[IRB$V_DUPS_SEEN]
1821 1874 3 | THEN
1822 1875 3 |
1823 1876 3 | ! Prologue 1 and 2 SIDR.
1824 1877 3 |
1825 1878 3 | IF .IFAB[IFB$B_PLG_VER] LSSU PLG$C_VER_3
1826 1879 3 | THEN
1827 1880 4 | BEGIN
1828 1881 4 |
1829 1882 4 | ! If duplicates are allowed, then the overhead for the SIDR
1830 1883 4 | includes a four byte duplicate count field which will never
1831 1884 4 | be used but is there for compatibility. This field is
```

1832 1885 4 ! unnecessary, if the new SIDR is to go into a continuation  
1833 1886 4 bucket.  
1834 1887 4  
1835 1888 4 IF .IDX\_DFN[IDX\$V\_DUPKEYS]  
1836 1889 4 AND  
1837 1890 4 NOT .IRAB[IRBSV\_CONT\_BKT]  
1838 1891 4 THEN  
1839 1892 4 SIZE = .SIZE + IRCSC\_DATOVHDSZ + IRCSC\_DATSZFLD  
1840 1893 4 + IRCSC\_DCNTSZFLD  
1841 1894 4 ELSE  
1842 1895 4 SIZE = .SIZE + IRCSC\_DATOVHDSZ + IRCSC\_DATSZFLD;  
1843 1896 4  
1844 1897 4 SIZE = .SIZE + .IDX\_DFN[IDX\$B\_KEYSZ];  
1845 1898 4 END  
1846 1899 4  
1847 1900 4 ! Prologue 3 SIDR.  
1848 1901 4  
1849 1902 3 ELSE  
1850 1903 4 BEGIN  
1851 1904 4 SIZE = .SIZE + IRCSC\_SDROVHSZ3 + RMSV3KEY\_SZ (KEYBUF\_ADDR(2));  
1852 1905 4  
1853 1906 4 IF .IDX\_DFN[IDX\$V\_KEY\_COMPR]  
1854 1907 4 THEN  
1855 1908 4 SIZE = .SIZE + IRCSC\_KEYCMPOVH;  
1856 1909 3 END;  
1857 1910 3 END ! of SIDR data level

1859 1911 3  
1860 1912 2 ELSE  
1861 1913 2  
1862 1914 2 IF .IFAB [IFB\$B\_PLG\_VER] LSSU PLG\$C\_VER\_3  
1863 1915 2 THEN  
1864 1916 2  
1865 1917 2 This is a Prolog 1 or 2 index bucket. The size to be returned is  
1866 1918 2 given below for each of the possible split cases.  
1867 1919 2  
1868 1920 2 1) two-bucket split no empty buckets:  
1869 1921 2 size of new key + size of VBN\_RIGHT  
1870 1922 2  
1871 1923 2  
1872 1924 2  
1873 1925 2  
1874 1926 2  
1875 1927 2 size of VBN\_MID - size of old VBN in record  
1876 1928 2 (only update VBN pointer)  
1877 1929 2  
1878 1930 2  
1879 1931 2  
1880 1932 2 size of two new keys + size of first new VBN (VBN\_MID)  
1881 1933 2 + size of second new VBN (VBN\_RIGHT)  
1882 1934 2  
1883 1935 2  
1884 1936 2  
1885 1937 2  
1886 1938 2  
1887 1939 2  
1888 1940 2  
1889 1941 2  
1890 1942 2  
1891 1943 2  
1892 1944 2  
1893 1945 2  
1894 1946 2  
1895 1947 2  
1896 1948 2  
1897 1949 2  
1898 1950 2  
1899 1951 2  
1900 1952 2  
1901 1953 2  
1902 1954 2  
1903 1955 2  
1904 1956 2  
1905 1957 2  
1906 1958 2  
1907 1959 2  
1908 1960 2  
1909 1961 2  
1910 1962 2  
1911 1963 2  
1912 1964 2  
1913 1965 2  
1914 1966 2  
1915 1967 2  
Two-pass Split Cases Without Empty Buckets.  
1943 2 5) first pass of two-pass two-bucket split case:  
1944 2 size of new key + size of old VBN in record  
1945 2  
1946 2  
1947 2  
1948 2  
1949 2  
1950 2  
1951 2  
1952 2  
1953 2  
1954 2  
1955 2  
1956 2  
1957 2  
1958 2  
1959 2  
1960 2  
1961 2  
1962 2  
1963 2  
1964 2  
1965 2  
1966 2  
1967 2  
size of first new key + size of old VBN in record  
8) second pass of two-pass multibucket split case -  
one of the new keys in each bucket:  
size of second new key + size of new VBN (VBN\_MID)  
+ size of VBN\_RIGHT  
- size of old VBN in record  
(update one VBN pointer)  
9) first pass of two-pass multibucket split case -  
both of the new keys in the same bucket:

```

1916    1968 2      size of the two new keys + size of new VBN (VBN_MID)
1917    1969 2      + size of old VBN in record
1918    1970 2
1919    1971 2      10) second pass of two-pass multibucket split case -
1920    1972 2      both of the new keys in the same bucket:
1921    1973 2      size of VBN_RIGHT - size of old VBN in record
1922    1974 2      (only update VBN pointer)
1923    1975 2
1924    1976 2
1925    1977 2      Two-pass Bucket Split Cases with Empty Buckets
1926    1978 2
1927    1979 2      11) second pass of two-pass two-bucket split case:
1928    1980 2      size of VBN_MID - size of old VBN in record
1929    1981 2      (only update VBN pointer)
1930    1982 2
1931    1983 2
1932    1984 2      11) first pass of two-pass multi-bucket split case:
1933    1985 2      size of new key + size of VBN_MID
1934    1986 2      - size of old VBN in record
1935    1987 2      (only update VBN in pointer)
1936    1988 2
1937    1989 2
1938    1990 2      13) second pass of two-pass multi-bucket split case:
1939    1991 2      size of VBN_RIGHT - size of old VBN in record
1940    1992 2      (update VBN pointer)
1941    1993 2
1942    1994 2
1943    1995 2      sizes are always computed so as to include any control bytes
1944    1996 2      required
1945    1997 2
1946    1998 3      BEGIN
1947    1999 3      SIZE = 0;
1948    2000 3
1949    2001 3      ! Handle all empty bucket split cases
1950    2002 3
1951    2003 3      IF .IRAB[IRBV_EMPTY_BKT]
1952    2004 3      THEN
1953    2005 4      BEGIN
1954    2006 4      AP = 1;                      ! set for index level
1955    2007 4
1956    2008 4      ! In all empty bucket split cases the existing down pointer
1957    2009 4      must match VBN_LEFT before we modify it. If it doesn't
1958    2010 4      match we can't change it or we may be causing crossed
1959    2011 4      pointers down to the level below.
1960    2012 4
1961    2013 4      IF .IRAB[IRBSL_VBN_LEFT] EQLA RMSRECORD_VBN()
1962    2014 4      THEN
1963    2015 4
1964    2016 4      ! If this is the first pass of the two-pass multi-bucket
1965    2017 4      with empty bucket split case then if we are going to
1966    2018 4      swing the pointer the size is IRBSL_VBN_MID.
1967    2019 4
1968    2020 4      IF .IRAB[IRBV_SPL_IDX]
1969    2021 4      THEN
1970    2022 4      SIZE = RMSVBN_SIZE (.IRAB[IRBSL_VBN_MID])
1971    2023 4
1972    2024 4      ! Otherwise we are just going to change the VBN pointer in

```

RMSRECORD\_SIZE

1973 2025 4  
1974 2026 4  
1975 2027 4  
1976 2028 4  
1977 2029 4  
1978 2030 4  
1979 2031 4  
1980 2032 4  
1981 2033 4  
1982 2034 4  
1983 2035 4  
1984 2036 4  
1985 2037 4  
1986 2038 4  
1987 2039 4  
1988 2040 4  
1989 2041 4  
1990 2042 4  
1991 2043 5  
1992 2044 4  
1993 2045 4  
1994 2046 4  
1995 2047 4  
1996 2048 4  
1997 2049 4  
1998 2050 4  
1999 2051 4  
2000 2052 4  
2001 2053 4  
2002 2054 4  
2003 2055 5  
2004 2056 5  
2005 2057 5  
2006 2058 5  
2007 2059 5  
2008 2060 5  
2009 2061 4  
2010 2062 4  
2011 2063 4  
2012 2064 3  
2013 2065 3  
2014 2066 3  
2015 2067 3  
2016 2068 3  
2017 2069 3  
2018 2070 3  
2019 2071 3  
2020 2072 3  
2021 2073 3  
2022 2074 3  
2023 2075 3  
2024 2076 3  
2025 2077 3  
2026 2078 3  
2027 2079 3  
2028 2080 3  
2029 2081 3

| place and the size is the difference between the current  
| size, and the size of the VBN that will replace it  
| (found in IRB\$L\_VBN\_MID).  
|  
ELSE  
SIZE = RMSVBN\_SIZE(.IRAB[IRB\$L\_VBN\_MID]) -  
.REC\_ADDR[IRC\$V\_PTRSZ] = 2;  
|  
IF NOT .IRAB[IRBSV\_BIG\_SPLIT]  
THEN  
RETURN .SIZE;  
|  
If this is a sizing for the second pass of a two-pass  
multi-bucket empty bucket split case, then the size of  
the index record to be inserted is only the difference  
between the size of the current record's VBN and the size  
of the VBN that is to replace it.  
|  
IF (.IRAB[IRB\$L\_VBN\_LEFT] EQLA 0)  
THEN  
SIZE = RMSVBN\_SIZE(.IRAB[IRB\$L\_VBN\_RIGHT])  
.REC\_ADDR[IRC\$V\_PTRSZ] - 2  
|  
For all other multi-bucket empty bucket split cases, the  
size of the index record(s) to be inserted includes the  
key size, and in the case of the one-pass multi-bucket  
empty bucket split case includes the size of the VBN in  
IRB\$L\_VBN\_RIGHT.  
|  
ELSE  
BEGIN  
SIZE = .SIZE + .IDX\_DFN[IDX\$B\_KEYSZ] + 1;  
|  
IF NOT .IRAB[IRBSV\_SPL\_IDX]  
THEN  
SIZE = .SIZE + RMSVBN\_SIZE(.IRAB[IRB\$L\_VBN\_RIGHT]);  
END;  
|  
RETURN .SIZE;  
END;  
|  
Determine size of index records to be added for all bucket split  
cases not involving empty buckets.  
|  
If this is not the first pass of a two-pass bucket split case,  
compute the difference between the current record VBN pointer  
size and the size of the VBN to swing (VBN\_RIGHT). This will  
facilitate computation of the record size during the second  
pass of all two-pass bucket split cases when a VBN pointer must  
be updated (in fact becomes the size for two of the two-pass  
bucket split cases), and it becomes part of the computation  
of the VBN pointer size required for VBN\_RIGHT for all one-pass  
bucket split cases.  
|  
IF NOT .IRAB[IRBSV\_SPL\_IDX]  
THEN

2030 2082 3  
2031 2083 3  
2032 2084 3  
2033 2085 3  
2034 2086 3  
2035 2087 3  
2036 2088 3  
2037 2089 3  
2038 2090 3  
2039 2091 3  
2040 2092 3  
2041 2093 3  
2042 2094 4  
2043 2095 4  
2044 2096 3  
2045 2097 4  
2046 2098 4  
2047 2099 4  
2048 2100 3  
2049 2101 3  
2050 2102 3  
2051 2103 3  
2052 2104 3  
2053 2105 3  
2054 2106 3  
2055 2107 3  
2056 2108 3  
2057 2109 3  
2058 2110 3  
2059 2111 3  
2060 2112 3  
2061 2113 3  
2062 2114 3  
2063 2115 3  
2064 2116 3  
2065 2117 3  
2066 2118 3  
2067 2119 3  
2068 2120 3  
2069 2121 3  
2070 2122 3  
2071 2123 3

SIZE = RMSVBN\_SIZE(.IRAB[IRBSL\_VBN\_RIGHT])  
- .REC\_ADDR[IRC\$V\_PTRSZ] = 2;

The size of the middle key key and its corresponding VBN must be included in the computation of record size if this is a one-pass multibucket case, the first pass of the two-pass multibucket split case where both keys go in the same (old/left) index bucket, or the second pass of the two-pass multibucket split case where one of the new keys goes in each of the old and new index buckets and we are currently updating the new index bucket with the second of the two keys and its corresponding VBN.

IF (.IRAB[IRBSV\_BIG\_SPLIT] AND NOT .IRAB[IRBSV\_SPL\_IDX] AND  
.IRAB[IRBSL\_VBN\_MID] NEQU 0)  
OR  
(.IRAB[IRBSV\_SPL\_IDX]  
AND .IRAB[IRBSV\_BIG\_SPLIT]  
AND (.IRAB[IRBSQ\_POS\_INS] EQLU BKT\$C\_OVERHDSZ))  
THEN  
SIZE = .SIZE + RMSVBN\_SIZE(.IRAB[IRBSL\_VBN\_MID]) +  
.IDX\_DFN[IDX\$B\_KEYSZ] + 1;

Finally, the size of the record must include the size of the first key on all one-pass bucket split cases not involving empty buckets, and on all two-pass bucket split cases during the first pass when the contents of the old index buckets are being updated.

Note that the current record size is also added in. If this is a one-pass bucket split cases this quantity plus the difference between the VBN pointer size in the old record and the VBN pointer size required for VBN RIGHT (already included within SIZE) yield the size of VBN RIGHT. If this is one of the two-pass bucket split cases, it is at this point that we include in the size determination the size of the old record's VBN pointer.

IF .IRAB[IRBSL\_VBN\_LEFT] NEQ 0  
THEN  
SIZE = .SIZE + .REC\_ADDR[IRC\$V\_PTRSZ] + 2 +  
.IDX\_DFN[IDX\$B\_KEYSZ] + 1;

END

2073 2124 3  
2074 2125 2 ELSE  
2075 2126 2  
2076 2127 2  
2077 2128 2  
2078 2129 2  
2079 2130 2  
2080 2131 2  
2081 2132 2  
2082 2133 2  
2083 2134 2  
2084 2135 2  
2085 2136 2  
2086 2137 2  
2087 2138 2  
2088 2139 2  
2089 2140 2  
2090 2141 2  
2091 2142 2  
2092 2143 2  
2093 2144 2  
2094 2145 2  
2095 2146 2  
2096 2147 2  
2097 2148 2  
2098 2149 2  
2099 2150 2  
2100 2151 2  
2101 2152 2  
2102 2153 2  
2103 2154 2  
2104 2155 2  
2105 2156 2  
2106 2157 2  
2107 2158 2  
2108 2159 2  
2109 2160 2  
2110 2161 2  
2111 2162 2  
2112 2163 2  
2113 2164 2  
2114 2165 2  
2115 2166 2  
2116 2167 2  
2117 2168 2  
2118 2169 2  
2119 2170 2  
2120 2171 2  
2121 2172 2  
2122 2173 2  
2123 2174 2  
2124 2175 2  
2125 2176 2  
2126 2177 2  
2127 2178 2  
2128 2179 2  
2129 2180 2

This is a Prolog 3 index bucket. The size to be returned is given below for each of the possible split cases. Note that it will always be possible for the VBN pointer size to change in a bucket. This requires that the additional bytes needed to increase the VBN pointer size of every currently existing index record in the current bucket to the new value be included in the size determination.

1) two-bucket split no empty buckets:

size of new key + VBN pointer size of bucket  
+ number of new VBN bytes

2) two-bucket split empty bucket:

number of new VBN bytes  
(only update VBN pointer)

3) multibucket split no empty buckets:

size of two new keys + 2 \* VBN pointer size of bucket  
+ number of new VBN bytes

4) multibucket split with empty buckets:

size of two new keys + VBN pointer size of bucket  
+ number of new VBN bytes  
(update one VBN pointer)

Two-pass Bucket Split Cases Without Empty Buckets

5) first pass of two-pass two-bucket split case:

size of new key + VBN pointer size of bucket  
number of new VBN bytes

6) second pass of two-pass two-bucket split case:

number of new VBN bytes  
(only update VBN pointer)

7) first pass of two-pass multibucket split case -  
one of the new keys in each bucket:

size of first new key + VBN pointer size of bucket  
+ number of new VBN bytes

8) second pass of two-pass multibucket split case -  
one of the new keys in each bucket:

size of second new key + VBN pointer size of bucket  
+ number of new VBN bytes  
(update one VBN pointer)

```

2130      2181 2           9) first pass of two-pass multibucket split case -
2131      2182 2           both of the new keys in the same bucket:
2132      2183 2           size of the two new keys + 2 * VBN pointer size of bucket
2133      2184 2           + number of new VBN bytes
2134      2185 2
2135      2186 2
2136      2187 2           10) second pass of two-pass multibucket split case -
2137      2188 2           both of the new keys in the same bucket:
2138      2189 2           number of new VBN bytes
2139      2190 2           (only update VBN pointer)
2140      2191 2
2141      2192 2
2142      2193 2           Two-pass Bucket Split Cases With Empty Buckets
2143      2194 2
2144      2195 2           11) second pass of two-pass two-bucket split case:
2145      2196 2           number of new VBN bytes
2146      2197 2           (only update VBN pointer)
2147      2198 2
2148      2199 2
2149      2200 2           12) first pass of two-pass multi-bucket split case:
2150      2201 2           size of new key + number of new VBN bytes
2151      2202 2           (update VBN pointer)
2152      2203 2
2153      2204 2
2154      2205 2           13) second pass of two-pass multi-bucket split case:
2155      2206 2           number of new VBN bytes
2156      2207 2           (only update VBN pointer)
2157      2208 2
2158      2209 2
2159      2210 2           Sizes are always computed so as to include any control bytes
2160      2211 2           required. For cases 3, 4, and 9, key size include compression
2161      2212 2           of the second record if compression is enabled.
2162      2213 2
2163      2214 2           BEGIN
2164      2215 2
2165      2216 2           ! First set up some useful constants
2166      2217 2
2167      2218 2
2168      2219 2           LOCAL
2169      2220 2           CONTEXT : REF BBLICK,   Address of block containing context
2170      2221 2           for potential update of an existing
2171      2222 2           index rec
2172      2223 2           KEY_SZ,          Number of bytes needed to store the
2173      2224 2           key of the index record(s)
2174      2225 2           (compression index is added in and
2175      2226 2           front compression of the second
2176      2227 2           record is subtracted)
2177      2228 2           VBN_SZ,          Size of all the VBNs after the
2178      2229 2           insertion (they could all grow by
2179      2230 2           one or two bytes)
2180      2231 2           NEW_VBN_BYTES; Number of bytes which will be added
2181      2232 2           to the existing VBNs because of
2182      2233 2           changing to a larger VBN size
2183      2234 2           (may be 0)
2184      2235 2
2185      2236 2           ! Determine the size (VBN_SZ) of the VBNs after insertion of the
2186      2237 2           index records to be inserted, and the number of VBN bytes
2187      2238 2           (NEW_VBN_BYTES) which will have to be added in order adjust the

```

```

2188      2239 2
2189      2240 2
2190      2241 2
2191      2242 2
2192      2243 2
2193      2244 2
2194      2245 2
2195      2246 2
2196      2247 2
2197      2248 2
2198      2249 2
2199      2250 2
2200      2251 2
2201      2252 2
2202      2253 2
2203      2254 2
2204      2255 2
2205      2256 2
2206      2257 2
2207      2258 2
2208      2259 2
2209      2260 2
2210      2261 2
2211      2262 2
2212      2263 2
2213      2264 2
2214      2265 2
2215      2266 2
2216      2267 2
2217      2268 2
2218      2269 2
2219      2270 2
2220      2271 2
2221      2272 2
2222      2273 2
2223      2274 2
2224      2275 2
2225      2276 2
2226      2277 2
2227      2278 2
2228      2279 2
2229      2280 2
2230      2281 2
2231      2282 2
2232      2283 2
2233      2284 2
2234      2285 2
2235      2286 2
2236      2287 2
2237      2288 2

```

```
: 2187      2238 3           | VBN pointer size of the pre-existing index records to this value.  
: 2188      2239 3           | Determination of VBN_SZ will directly depend upon the bucket  
: 2189      2240 3           | split case of the current invocation of this routine.  
: 2190      2241 3           NEW_VBN_BYTES = 0;  
: 2191      2242 3           | For all one-pass bucket split cases and two-pass two-bucket  
: 2192      2243 3           empty bucket split cases the VBN pointer size is the maximum of  
: 2193      2244 3           the number of bytes need to store the current bucket pointer  
: 2194      2245 3           size, VBN_RIGHT, and VBN_MID.  
: 2195      2246 3           IF (NOT .IRAB[IRBSV_SPL_IDX])  
: 2196      2247 3           AND  
: 2197      2248 3           (.IRAB[IRBSL_VBN_LEFT] NEQU 0)  
: 2198      2249 4           THEN  
: 2199      2250 3           VBN_SZ = MAXU (RM$VBN_SIZE(.IRAB[IRBSL_VBN_MID]),  
: 2200      2251 4           RM$VBN_SIZE(.IRAB[IRBSL_VBN_RIGHT]),  
: 2201      2252 3           .BKT_ADDR[BKT$V_PTR_SZ]+2)  
: 2202      2253 3           ELSE  
: 2203      2254 3           | Determine VBN_SZ for the first pass of all two-pass  
: 2204      2255 3           bucket split cases.  
: 2205      2256 3           IF .IRAB[IRBSV_SPL_IDX]  
: 2206      2257 3           THEN  
: 2207      2258 3           | For case 9 (two-pass multi-bucket split with both keys  
: 2208      2259 3           in the same bucket) and case 12 (multi-pass multi-bucket  
: 2209      2260 3           empty bucket split) VBN_SZ is the maximum number of  
: 2210      2261 3           bytes needed to store VBN_MID and the current bucket  
: 2211      2262 3           VBN pointer size.  
: 2212      2263 3           IF .IRAB[IRBSV_BIG_SPLIT]  
: 2213      2264 3           AND  
: 2214      2265 3           ((.IRAB[IRBSW_POS_INS] EQLU BKT$C_OVERHDSZ)  
: 2215      2266 3           OR  
: 2216      2267 3           .IRAB[IRBSV_EMPTY_BKT])  
: 2217      2268 3           THEN  
: 2218      2269 3           VBN_SZ = MAXU (RM$VBN_SIZE(.IRAB[IRBSL_VBN_MID]),  
: 2219      2270 3           .BKT_ADDR[BKT$V_PTR_SZ]+2)  
: 2220      2271 3           | For all other two-pass bucket split cases, the VBN_SZ  
: 2221      2272 5           is the number of bytes need to store the current VBN  
: 2222      2273 4           pointer bucket size.  
: 2223      2274 4           ELSE  
: 2224      2275 3           VBN_SZ = .BKT_ADDR[BKT$V_PTR_SZ]+2  
: 2225      2276 3           | Determine VBN_SZ for the second pass of all two-pass  
: 2226      2277 3           bucket split cases.  
: 2227      2278 3           ELSE  
: 2228      2279 3           | For case 8 (two-pass multibucket split with a key  
: 2229      2280 3           inserted in each of the buckets) VBN_SZ is the  
: 2230      2281 3           number of bytes needed to store the maximum of the  
: 2231      2282 3           current VBN bucket pointer size and VBN_MID.  
: 2232      2283 3           ; Ro  
: 2233      2284 3           ; 26
```

```
: 2244 RMSRECORD_SIZE
: 2245 2295 3
: 2246 2296 3
: 2247 2297 3
: 2248 2298 3
: 2249 2299 3
: 2250 2300 3
: 2251 2301 3
: 2252 2302 3
: 2253 2303 3
: 2254 2304 3
: 2255 2305 3
: 2256 2306 3
: 2257 2307 3
: 2258 2308 3
: 2259 2309 3
: 2260 2310 3
: 2261 2311 3
: 2262 2312 3
: 2263 2313 3
: 2264 2314 3
: 2265 2315 3
: 2266 2316 3
: 2267 2317 3
: 2268 2318 3
: 2269 2319 3
: 2270 2320 3
: 2271 2321 3
: 2272 2322 3
: 2273 2323 3
: 2274 2324 4
: 2275 2325 3
: 2276 2326 4
: 2277 2327 3
: 2278 2328 3
: 2279 2329 3
: 2280 2330 3
: 2281 2331 3
: 2282 2332 3
: 2283 2333 3
: 2284 2334 3
: 2285 2335 3
: 2286 2336 3
: 2287 2337 3
: 2288 2338 3
: 2289 2339 4
: 2290 2340 3
: 2291 2341 3
: 2292 2342 3
: 2293 2343 3
: 2294 2344 3
: 2295 2345 3
: 2296 2346 3
: 2297 2347 3
: 2298 2348 3
: 2299 2349 3
: 2300 2350 3
: 2301 2351 3

        ! IF .IRAB[IRBSL_VBN_MID] NEQU 0
        THEN
          VBN_SZ = MAXU (RMSVBN_SIZE(.IRAB[IRBSL_VBN_MID]),
                         RMSVBN_SIZE(.IRAB[IRBSL_VBN_RIGHT]))
        ! For all other two-pass bucket split cases, VBN_SZ becomes
        ! the number of bytes needed to store VBN_RIGHT.
        ELSE
          VBN_SZ = RMSVBN_SIZE(.IRAB[IRBSL_VBN_RIGHT]);
        NEW_VBN_BYTES = RMSNEW_VBN_BYTES(.VBN_SZ);

        ! Determine the key size of the first key to be inserted.
        ! The two bytes required as overhead for key compression are
        ! included in the key size if index compression is enabled.

        IF .IDX_DFN [IDX$V_IDX_COMPR]
        THEN
          KEY_SZ = 2
        ELSE
          KEY_SZ = 0;

        ! For all one-pass bucket split cases where there are keys to be
        ! inserted, and for the first pass of all two-pass bucket split
        ! cases, the first key to be inserted is the key found in key
        ! buffer 2.

        IF (.IRAB[IRBSL_VBN_LEFT] NEQU 0)
          AND
          (NOT (.IRAB[IRBSV_EMPTY_BKT] AND NOT .IRAB[IRBSV_BIG_SPLIT]))
        THEN
          KEY_SZ = .KEY_SZ + RMSV3KEY_SZ (KEYBUF_ADDR(2))
        ELSE
          ! For the second pass of the two-pass multibucket split case
          ! where a key goes in each of the two index buckets (case 8)
          ! the first and only key to be inserted is actually the second
          ! of the two keys to be inserted, and it is found in key
          ! buffer 3.

          IF NOT .IRAB[IRBSV_EMPTY_BKT]
            AND
            (.IRAB[IRBSL_VBN_MID] NEQU 0)
          THEN
            KEY_SZ = .KEY_SZ + RMSV3KEY_SZ (KEYBUF_ADDR(3))

          ! For the second pass of all other two-pass split cases,
          ! the key size is zero.

          ELSE
            KEY_SZ = 0;

        IF .IRAB [IRBSV_EMPTY_BKT]
        THEN
```

: 264  
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: 300

```
: 2301      2352 3           ! We have an empty bucket.  
: 2302      2353 3           !  
: 2303      2354 4           BEGIN  
: 2304      2355 4           SIZE = 0;  
: 2305      2356 4           !  
: 2306      2357 4           If there was a multiple bucket split, add new key and VBN  
: 2307      2358 4           unless this is the second pass of a two-pass multi-bucket  
: 2308      2359 4           empty bucket split case.  
: 2309      2360 4           !  
: 2310      2361 5           IF (.IRAB[IRB$V_BIG_SPLIT]  
: 2311          AND  
: 2312          (.IRAB[IRB$L_VBN_LEFT] NEQA 0))  
: 2313      THEN  
: 2314          BEGIN  
: 2315          SIZE = .KEY_SZ;  
: 2316          !  
: 2317      2368 5           If this is not a two-pass split case, the number of  
: 2318      2369 5           new VBN bytes is VBN size plus the number of bytes  
: 2319      2370 5           which will have to be added because of a possible  
: 2320      2371 5           VBN size change within the index bucket (will be  
: 2321      2372 5           zero if no VBN size change occurred).  
: 2322      2373 5           !  
: 2323      2374 5           IF NOT .IRAB[IRB$V_SPL_IDX]  
: 2324      THEN  
: 2325          NEW_VBN_BYTES = .NEW_VBN_BYTES + .VBN_SZ  
: 2326          !  
: 2327      2378 5           If this is the first pass of a two-pass multi-bucket  
: 2328      2379 5           empty bucket split case then the number of new VBN bytes  
: 2329      2380 5           will be the number of bytes which will have to be added  
: 2330      2381 5           because of a possible VBN size change within the index  
: 2331      2382 5           bucket provided we are allowed to update the VBN of the  
: 2332      2383 5           current record because it is the same as the VBN of the  
: 2333      2384 5           leftmost primary data bucket. If this is not the case,  
: 2334      2385 5           then because we can't update the VBN (this could cause  
: 2335      2386 5           crossed pointers down to the data level), the number  
: 2336      2387 5           of new VBN bytes is set to zero.  
: 2337      2388 5           !  
: 2338      2389 5           ELSE  
: 2339      2390 6           IF (.IRAB[IRB$L_VBN_LEFT] NEQA RM$V3_VBN())  
: 2340      THEN  
: 2341          NEW_VBN_BYTES = 0;  
: 2342      END  
: 2343      2394 5           !  
: 2344      2395 5           If there wasn't a multibucket split, and we are not going  
: 2345      2396 5           to swing a pointer because this could be causing crossed  
: 2346      2397 5           pointers down to the level below, then there is need to add  
: 2347      2398 5           new VBN bytes.  
: 2348      2399 5           !  
: 2349      2400 4           ELSE  
: 2350      2401 5           IF (NOT .IRAB[IRB$V_BIG_SPLIT]  
: 2351          AND  
: 2352          (.IRAB [IRB$L_VBN_LEFT] NEQU RM$V3_VBN()))  
: 2353      THEN  
: 2354          NEW_VBN_BYTES = 0;  
: 2355          !  
: 2356      2406 4           END  
: 2357      2407 4           !  
: 2358      2408 4           END
```

2358 2409 3 ELSE  
2359 2410 3  
2360 2411 3 No empty buckets. For now, assume that entire record  
2361 2412 3 will fit into one bucket and compute its size.  
2362 2413 3 Always add VBN and key unless this is the second pass of one  
2363 2414 3 of the two-pass bucket split cases where all we are going  
2364 2415 3 to do is update a VBN pointer. In such cases, the size  
2365 2416 3 returned does not include the size of a VBN and a key.  
2366 2417 3  
2367 2418 4 BEGIN  
2368 2419 4 SIZE = .KEY\_SZ;  
2369 2420 4  
2370 2421 5 IF (.IRAB[IRBSL\_VBN\_LEFT] NEQU 0)  
2371 2422 4 OR  
2372 2423 5 (.IRAB[IRBSL\_VBN\_MID] NEQU 0)  
2373 2424 4 THEN  
2374 2425 4 NEW\_VBN\_BYTES = .NEW\_VBN\_BYTES + .VBN\_SZ;  
2375 2426 4  
2376 2427 4 If this is a one-pass multibucket split or the first pass  
2377 2428 4 of a two-pass multibucket split where both index records  
2378 2429 4 are to go in the same (old) index bucket, the size of the  
2379 2430 4 index record must include (in addition to the key size of the  
2380 2431 4 first index record, the new VBN size of the bucket, and the  
2381 2432 4 number of new VBN bytes which must be added to all  
2382 2433 4 pre-existing VBN pointers to make them all the same size) the  
2383 2434 4 key size of the second index record (less its front  
2384 2435 4 compression if index compression is enabled) and the size of  
2385 2436 4 another new VBN. If index compression is enabled, then the  
2386 2437 4 key size will not include any rear-end truncated characters.  
2387 2438 4  
2388 2439 4 IF .IRAB[IRBSV\_BIG\_SPLIT]  
2389 2440 4 AND  
2390 2441 6 ((NOT .IRAB[IRBSV\_SPL\_IDX]  
2391 2442 6 AND .IRAB[IRBSL\_VBN\_LEFT] NEQU 0)  
2392 2443 5 OR  
2393 2444 6 (.IRAB[IRBSV\_SPL\_IDX]  
2394 2445 5 AND .IRAB[IRBSW\_POS\_INS] EQLU BKT\$C\_OVERHDSZ))  
2395 2446 4 THEN  
2396 2447 5 BEGIN  
2397 2448 5 NEW\_VBN\_BYTES = .NEW\_VBN\_BYTES + .VBN\_SZ;  
2398 2449 5  
2399 2450 5 IF NOT .IDX\_DFN [IDX\$V\_IDX\_COMPR]  
2400 2451 5 THEN  
2401 2452 5 SIZE = .SIZE + .IDX\_DFN [IDX\$B\_KEYSZ]  
2402 2453 5 ELSE  
2403 2454 6 BEGIN  
2404 2455 6  
2405 2456 6 LOCAL  
2406 2457 6 KEYBUF3 : REF BBLOCK,  
2407 2458 6 KEYBUF4 : REF BBLOCK,  
2408 2459 6 KEYBUF5 : REF BBLOCK,  
2409 2460 6 SAVE\_BKT\_ADDR;  
2410 2461 6  
2411 2462 6 ! To determine the size of both keys, we will  
2412 2463 6 ! need key buffers 3, 4, and 5.  
2413 2464 6  
2414 2465 6 KEYBUF3 = KEYBUF\_ADDR(3);

```

2415    2466 6      KEYBUF4 = KEYBUF_ADDR(4);
2416    2467 6      KEYBUF5 = KEYBUF_ADDR(5);
2417    2468 6
2418    2469 6
2419    2470 6      | Move the second key into key buffer 4 and compress
2420    2471 6      it as if it was to go at the point of insertion
2421    2472 6      instead of following the first key (which is to
2422    2473 6      be inserted there). Then recompress the second key
2423    2474 6      based upon the first key (which is still in key
2424    2475 6      buffer 5).
2425    2476 6      RMSMOVE (.KEYBUF3[KEY_LEN] + 2,
2426    2477 6          .KEYBUF3,
2427    2478 6          .KEYBUF4);
2428    2479 6      RMSCOMPRESS_KEY (.KEYBUF4);
2429    2480 6
2430    2481 6      SAVE_BKT_ADDR = .BKT_ADDR;
2431    2482 6      BKT_ADDR = 0;
2432    2483 6      RMSRECOMPR_KEY (.KEYBUF5, .KEYBUF4);
2433    2484 6      BKT_ADDR = .SAVE_BKT_ADDR;
2434    2485 6
2435    2486 6      | Add the size of the second key to the size of the
2436    2487 6      first key to obtain the key contribution to the
2437    2488 6      size of the record.
2438    2489 6
2439    2490 6      SIZE = .SIZE + .KEYBUF4[KEY_LEN] + 2;
2440    2491 5      END;
2441    2492 4      END;
2442    2493 3      END;
2443    2494
2444    2495 3      | Return the size of the record for prolog 3 index buckets as two
2445    2496 3      contiguous words. The high order word contains the number of VBN
2446    2497 3      bytes required while the low order word contains the number of
2447    2498 3      bytes required to hold the key(s).
2448    2499
2449    2500 3      RETURN (.NEW_VBN_BYTES)^16 + .SIZE;
2450    2501 2      END;
2451    2502 2
2452    2503 2      | Return the size of the prologue 1 & 2 record or Pprologue 3 SIDR as
2453    2504 2      a single unified quantity.
2454    2505 2
2455    2506 2      RETURN .SIZE;
2456    2507 1      END;

```

	1C BB 00000 RMSRECORD_SIZE::			
5E	0C	C2 00002	PUSHR	#^M<R2,R3,R4>
	41	A9 95 00005	SUBL2	#12 SP
		5F 12 00008	TSTB	65(IRAB)
	78	A9 DD 0000A	BNEQ	8\$
		0000V 30 0000D	PUSHL	120(IRAB)
5E	04	C0 00010	BSBW	RMSVBN_SIZE
52	50	D0 00013	ADDL2	#4, SP
	50	D4 00016	MOVL	R0, SIZE
			CLRL	R0

: 1772  
: 1854  
: 1861  
: 1863

		03	0087	CA	91 00018	CMPB	183(IFAB), #3	
				07	1E 0001D	BGEQU	1\$	
		52		50	D6 0001F	INCL	R0	
				02	C0 00021	ADDL2	#2, SIZE	1865
		52		03	11 00024	BRB	2\$	
			44	03	C0 00026	ADDL2	#3, SIZE	1867
				A9	95 00029	TSTB	68(IRAB)	1873
				39	19 0002C	BLSS	7\$	
		1A		50	E9 0002E	BLBC	R0, 6\$	1878
		0A	1C	A7	E9 00031	BLBC	28(IDX_DFN), 3\$	1888
		A9		04	E0 00035	BBS	#4, 68(IRAB), 3\$	1890
		52		08	C0 0003A	ADDL2	#8, SIZE	1893
				03	11 0003D	BRB	4\$	1892
		52		04	C0 0003F	ADDL2	#4, SIZE	1895
			20	A7	9A 00042	MOVZBL	32(IDX_DFN), R0	1897
		50		50	C0 00046	ADDL2	R0, SIZE	
				1C	11 00049	BRB	7\$	
		50	00B4	CA	3C 0004B	MOVZWL	180(IFAB), R0	1878
				60	B940 9F 00050	PUSHAB	096(IRAB)[R0]	1904
				0000V	30 00054	BSBW	RMS\$V3KEY_SZ	
		5E		04	C0 00057	ADDL2	#4, SP	
		52	02	A042	9E 0005A	MOVAB	2(R0)[SIZE], SIZE	
		A7		06	E1 0005F	BBC	#6, 28(IDX_DFN), 12\$	1906
		52		02	C0 00064	ADDL2	#2, SIZE	1908
				68	11 00067	BRB	12\$	1878
		03	0087	CA	91 00069	CMPB	183(IFAB), #3	1914
				03	1F 0006E	BLSSU	9\$	
				00E2	31 00070	BRW	21\$	
				52	D4 00073	CLRL	SIZE	
		73	44	A9	06 E1 00075	BBC	#6, 68(IRAB), 14\$	1999
				5C	01 D0 0007A	MOVL	#1, AP	2003
				0000G	30 0007D	BSBW	RMS\$RECORD_VBN	2006
				50	0088 C9 D1 00080	CMPL	136(IRAB), R0	2013
				29	12 00085	BNEQ	11\$	
			OF	44	A9 E9 00087	BLBC	68(IRAB), 10\$	2020
				0090	C9 DD 0008B	PUSHL	144(IRAB)	2022
				0000V	30 0008F	BSBW	RMS\$VBN_SIZE	
		5E		04	C0 00092	ADDL2	#4, SP	
		52		50	DD 00095	MOVL	R0, SIZE	
				16	11 00098	BRB	11\$	
				0090	C9 DD 0009A	PUSHL	144(IRAB)	2030
				0000V	30 0009E	BSBW	RMS\$VBN_SIZE	
		51	66	5E	04 C0 000A1	ADDL2	#4, SP	
				02	EF 000A4	EXTZV	#0, #2, (REC_ADDR), R1	2031
				50	51 C2 000A9	SUBL2	R1, R0	
			1C	52 FE A0 9E 000AC	MOVAB	-2(R0), SIZE		
				02	E1 000B0	TSTL	#2, 68(IRAB), 12\$	2033
				0088	C9 D5 000B5	BNEQ	13\$	2043
				18	12 000B9	PUSHL	140(IRAB)	2045
				008C	C9 DD 000BB	BSBW	RMS\$VBN_SIZE	
		51	66	5E	04 C0 000C2	ADDL2	#4, SP	
				02	EF 000C5	EXTZV	#0, #2, (REC_ADDR), R1	2046
				50	51 C2 000CA	SUBL2	R1, R0	
				52 FE A0 9E 000CD	MOVAB	-2(R0), SIZE		
				7F	11 000D1	BRB	20\$	2045
			50	20 A7 9A 000D3	MOVZBL	32(IDX_DFN), R0	2056	

			52	01 A042 9E 000D7	MOVAB	1(R0)[SIZE]	SIZE		
			72	44 A9 E8 000DC	BLBS	68(IRAB)	20\$	2058	
				008C C9 DD 000E0	PUSHL	140(IRAB)		2060	
				0000V 30 000E4	BSBW	RMSVBN_SIZE			
			5E	04 C0 000E7	ADDL2	#4, SP			
				FF59 31 000EA	BRW	5\$			
			16	44 A9 E8 000ED 14\$:	BLBS	68(IRAB), 15\$		2080	
				008C C9 DD 000F1	PUSHL	140(IRAB)		2082	
				0000V 30 000F5	BSBW	RMSVBN_SIZE			
			5E	04 C0 000F8	ADDL2	#4, SP			
			02	00 EF 000FB	EXTZV	#0, #2, (REC_ADDR), R1		2083	
			50	51 C2 00100	SUBL2	R1, R0			
			52	FE A0 9E 00103	MOVAB	-2(R0), SIZE			
51	66		0A	02 E1 00107 15\$:	BBC	#2, 68(IRAB), 16\$		2094	
			44	44 A9 E8 0010C	BLBS	68(IRAB), 17\$			
			0A	0090 C9 D5 00110	TSTL	144(IRAB)		2095	
				0F 12 00114	BNEQ	18\$			
			21	44 A9 E9 00116 16\$:	BLBC	68(IRAB), 19\$		2097	
			44	02 E1 0011A 17\$:	BBC	#2, 68(IRAB), 19\$		2098	
1C	44		OE	48 A9 B1 0011F	CMPW	72(IRAB), #14		2099	
				16 12 00123	BNEQ	19\$			
			0090	C9 DD 00125 18\$:	PUSHL	144(IRAB)		2101	
				0000V 30 00129	BSBW	RMSVBN_SIZE			
			5E	04 C0 0012C	ADDL2	#4, SP			
			50	52 C0 0012F	ADDL2	SIZE, R0			
			51	20 A7 9A 00132	MOVZBL	32(IDX_DFN), R1		2102	
			52	01 A140 9E 00136	MOVAB	1(R1)[R0], SIZE			
			0088	C9 D5 0013B 19\$:	TSTL	136(IRAB)		2118	
				11 13 0013F	BEQL	20\$			
50	66		02	00 EF 00141	EXTZV	#0, #2, (REC_ADDR), R0		2120	
			50	52 C0 00146	ADDL2	SIZE, R0			
			51	20 A7 9A 00149	MOVZBL	32(IDX_DFN), R1		2121	
			52	03 A140 9E 0014D	MOVAB	3(R1)[R0], SIZE			
				01C7 31 00152 20\$:	BRW	51\$		1914	
			37	08 AE D4 00155 21\$:	CLRL	NEW_VBN_BYTES		2242	
				44 A9 E8 00158	BLBS	68(IRAB), 24\$		2249	
			0088	C9 D5 0015C	TSTL	136(IRAB)		2251	
				2D 13 00160	BEQL	23\$			
			0090	C9 DD 00162	PUSHL	144(IRAB)		2253	
				0000V 30 00166	BSBW	RMSVBN_SIZE			
			51	50 D0 00169	MOVL	R0, R1			
			6E	008C C9 D0 0016C	MOVL	140(IRAB), (SP)		2254	
				0000V 30 00171	BSBW	RMSVBN_SIZE			
			5E	04 C0 00174	ADDL2	#4, SP			
			50	51 D1 00177	CMPL	R1, R0			
				03 1E 0017A	BGEQU	22\$			
			51	50 D0 0017C 22\$:	MOVL	R0, R1			
50	OD A5		02	03 EF 0017F	EXTZV	#3, #2, 13(BKT_ADDR), R0		2255	
			50	02 C0 00185	ADDL2	#2, R0			
			50	51 D1 00188	CMPL	R1, R0			
				5F 1F 0018B	BLSSU	28\$			
				60 11 0018D	BRB	29\$		2253	
			3A	44 A9 E9 0018F 23\$:	BLBC	68(IRAB), 27\$		2261	
28	44	A9	0E	02 E1 00193 24\$:	BBC	#2, 68(IRAB), 26\$		2270	
			48	A9 B1 00198	CMPW	72(IRAB), #14		2272	
				05 13 0019C	BEQL	25\$			
			1D	44 A9 06 E1 0019E	BBC	#6, 68(IRAB), 26\$		2274	

			0090	C9	DD 001A3 25\$:	PUSHL	144(IRAB)	2276
			0000V	30 001A7		BSBW	RMS\$VBN_SIZE	
51	OD A5		02	04	CO 001AA	ADDL2	#4, SP	
			5E	03	EF 001AD	EXTZV	#3, #2, 13(BKT_ADDR), R1	
			02	02	CO 001B3	ADDL2	#2, R1	
			51	50	D1 001B6	CMPL	R0, R1	
				44	1E 001B9	BGEQU	31\$	
				50	D0 001BB	MOVL	R1, R0	
				51	11 001BE	BRB	31\$	
04	AE	OD A5	04	02	EF 001C0 26\$:	EXTZV	#3, #2, 13(BKT_ADDR), VBN_SZ	2276
			AE	02	CO 001C7	ADDL2	#2, VBN_SZ	2284
				36	11 001CB	BRB	32\$	
				50	D0 001CD 27\$:	MOVL	144(IRAB), R0	2270
				21	13 001D2	BEQL	30\$	2296
				50	DD 001D4	PUSHL	R0	2298
				0000V	30 001D6	BSBW	RMS\$VBN_SIZE	
			51	50	DO 001D9	MOVL	R0, R1	
			6E	008C	C9 DO 001DC	MOVL	140(IRAB), (SP)	2299
				0000V	30 001E1	BSBW	RMS\$VBN_SIZE	
				5E	04 CO 001E4	ADDL2	#4, SP	
				50	51 D1 001E7	CMPL	R1, R0	
				03	1E 001EA	BGEQU	29\$	
			04	51	D0 001EC 28\$:	MOVL	R0, R1	
			AE	51	D0 001EF 29\$:	MOVL	R1, VBN_SZ	2298
				0E	11 001F3	BRB	32\$	
				008C	C9 DD 001F5 30\$:	PUSHL	140(IRAB)	2305
				0000V	30 001F9	BSBW	RMS\$VBN_SIZE	
			04	5E	04 CO 001FC	ADDL2	#4, SP	
			AE	50	D0 001FF 31\$:	MOVL	R0, VBN_SZ	
				04	AE DD 00203 32\$:	PUSHL	VBN SZ	2307
				FDB6	30 00206	BSBW	RMS\$NEW_VBN_BYTES	
			05	08	5E 04 CO 00209	ADDL2	#4, SP	
			1C	A7	50 D0 0020C	MOVL	R0, NEW VBN BYTES	2313
				6E	03 E1 00210	BBC	#3, 28(IDX_DFN), 33\$	2315
					02 D0 00215	MOVL	#2, KEY_SZ	
					02 11 00218	BRB	34\$	
					6E D4 0021A 33\$:	CLRL	KEY SZ	2317
					C9 D5 0021C 34\$:	TSTL	1367(IRAB)	2324
			05	44	A9 06 E1 00222	BBC	#6, 68(IRAB), 35\$	2326
			0B	44	A9 02 E1 00227	BBC	#2, 68(IRAB), 36\$	
				50	00B4 CA 3C 0022C 35\$:	MOVZWL	180(IFAB), R0	2328
					60 B940 9F 00231	PUSHAB	@96(IRAB)[R0]	
			1A	44	A9 14 11 00235	BRB	37\$	
				0090	06 E0 00237 36\$:	BBS	#6, 68(IRAB), 38\$	2337
				C9 14 13 00240	TSTL	144(IRAB)	2339	
				50 00B4 CA 3C 00242	BEQL	38\$		
				60 B940 3F 00247	MOVZWL	180(IFAB), R0	2341	
				0000V 30 0024B 37\$:	PUSHAW	@96(IRAB)[R0]		
				5E 04 CO 0024E	BSBW	RMS\$V3KEY_SZ		
				6E 50 CO 00251	ADDL2	#4, SP		
				02 11 00254	ADDL2	R0, KEY_SZ		
			2C	53 44 A9 6E D4 00256 38\$:	BRB	39\$		
			63	06 E1 00258 39\$:	CLRL	KEY SZ	2347	
				52 06 E1 0025C	MOVAB	68(IRAB), R3	2349	
				52 D4 00260	BBC	#6, (R3), 42\$		
					CLRL	SIZE	2355	

17	63	0088	02	E1 00262	BBC	#2 (R3), 41\$	2361	
			C9	D5 00266	TSTL	136(IRAB\$)	2363	
			0D	13 0026A	BEQL	40\$	2362	
	52		6E	D0 0026C	MOVL	KEY_SZ, SIZE	2366	
	0B		63	E8 0026F	BLBS	(R3), 41\$	2374	
	AE	04	AE	C0 00272	ADDL2	VBN_SZ, NEW_VBN_BYTES	2376	
			51	11 00277	BRB	48\$	2377	
4D	63	0000G	02	E0 00279	40\$: BBS	#2 (R3), 48\$	2401	
	50	0088	30	0027D	41\$: BSBW	RM\$V3 VBN	2403	
			C9	D1 00280	CMPL	136(IRAB), R0	2402	
			43	13 00285	BEQL	48\$	2404	
		08	AE	D4 00287	CLRL	NEW_VBN_BYTES	2405	
	52		3E	11 0028A	BRB	48\$	2349	
			6E	D0 0028C	42\$: MOVL	KEY_SZ, SIZE	2419	
		0088	50	D4 0028F	CLRL	R0	2421	
			C9	D5 00291	TSTL	136(IRAB)	2422	
			04	13 00295	BEQL	43\$	2423	
			50	D6 00297	INCL	R0	2424	
			06	11 00299	BRB	44\$	2425	
		0090	C9	D5 0029B	43\$: TSTL	144(IRAB)	2426	
			05	13 0029F	BEQL	45\$	2427	
68	08	AE	04	AE	ADDL2	VBN_SZ, NEW_VBN_BYTES	2428	
	63		02	E1 002A1	44\$: BBC	#2 (R3), 50\$	2439	
	06		63	E8 002A6	BLBS	(R3), 46\$	2441	
	09		50	E8 002AA	BLBS	R0, 47\$	2442	
	5F		63	E9 002B0	BLBC	(R3), 50\$	2444	
	0E	48	A9	B1 002B3	46\$: CMPW	72(IRAB), #14	2445	
			59	12 002B7	BNEQ	50\$	2446	
09	08	AE	04	AE	ADDL2	VBN_SZ, NEW_VBN_BYTES	2448	
	1C	A7	04	AE	002B9	47\$: ADDL2	32(IDX_DFN), 49\$	2450
	50		03	E0 002BE	BBS	#3, -28(IDX_DFN), 49\$	2451	
	52		20	A7 9A 002C3	MOVZBL	32(IDX_DFN), R0	2452	
			50	C0 002C7	ADDL2	R0, SIZE	2453	
			46	11 002CA	BRB	50\$	2454	
	51	0084	CA	3C 002CC	48\$: MOVZWL	180(IFAB), R1	2465	
	50		60	B941	3E 002D1	MOVAW	a96(IRAB)[R1], KEYBUF3	2466
	51		03	C5 002D6	MULL3	#3, R1, R3	2467	
	53	60	A9	CO 002DA	ADDL2	96(IRAB), KEYBUF4	2468	
	04	AE	60	B941	DE 002DE	MOVAL	a96(IRAB)[R1], KEYBUF5	2469
			09	BB 002E4	PUSHR	#^M<R0, R3>	2470	
	7E		60	9A 002E6	MOVZBL	(KEYBUF3), -(SP)	2471	
	6E		02	CO 002E9	ADDL2	#2 (SP)	2472	
			0000G	30 002EC	BSBW	RM\$MOVE	2473	
	5E		0C	CO 002EF	ADDL2	#12, SP	2474	
	50		53	DO 002F2	MOVL	KEYBUF4, R0	2475	
			0000G	30 002F5	BSBW	RM\$COMPRESS_KEY	2476	
	54		55	DO 002F8	MOVL	BKT_ADDR, SAVE_BKT_ADDR	2477	
	51		55	D4 002FB	CLRL	BKT_ADDR	2478	
	50	04	AE	DO 00300	MOVL	KEYBUF4, R1	2479	
			0000G	30 00304	BSBW	KEYBUF5, R0	2480	
	55		54	DO 00307	MOVL	RMSRECCMPR_KEY	2481	
	50		63	9A 0030A	MOVZBL	SAVE_BKT_ADDR, BKT_ADDR	2482	
	52	02	A042	9E 0030D	MOVAB	(KEYBUF4), R0	2483	
			10	78 00312	50\$: ASHL	2(R0)[SIZE], SIZE	2484	
	50		52	C0 00317	ADDL2	#16, NEW_VBN_BYTES, R0	2485	
			03	11 0031A	BRB	SIZE, R0	2486	
	50		52	DO 0031C	51\$: MOVL	52\$	2487	
						SIZE, R0	2506	

RM3MISPUT  
V04-000

RMSRECORD\_SIZE

K 5  
16-Sep-1984 01:51:28 VAX-11 Bliss-32 V4.0-742  
14-Sep-1984 13:01:29 [RMS.SRC]RM3MISPUT.B32;1

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(14)

5E

0C C0 0031F 52\$: ADDL2 #12, SP  
1C BA 00322 POPR #^M<R2,R3,R4>  
05 00324 RSB

; 2507

; Routine Size: 805 bytes, Routine Base: RMS\$RMS3 + 04F5

RM3  
V04

; R

2458 1 %SBTTL 'RMSSHFT\_VBNS'  
2459 1 ROUTINE RMSSHFT\_VBNS : RLSCOMMON\_LINK NOVALUE =  
2460 1  
2461 1 !++  
2462 1  
2463 1 RMSSHFT\_VBNS  
2464 1  
2465 1 Spread the VBNs at the end of the bucket to make room for  
2466 1 the new VBN(s).  
2467 1  
2468 1 This routine handles the following conditions:  
2469 1  
2470 1 1) two bucket split, no empty buckets  
2471 1 2) two bucket split, empty bucket  
2472 1 3) multi-bucket split (three or four), no empty buckets  
2473 1 4) multi-bucket split, empty buckets  
2474 1  
2475 1 The two-pass non-empty bucket bucket split cases which occur when the  
2476 1 insertion point is found to be the same as the bucket split point are  
2477 1 handled as follows:  
2478 1  
2479 1 5) first pass two bucket split - treated as 1)  
2480 1 6) second pass two bucket split - treated as 2)  
2481 1 first pass multi-bucket split  
2482 1 7) one in each bucket - treated as 4)  
2483 1 8) both in old bucket - treated as 3)  
2484 1 second pass multi-bucket split  
2485 1 9) one in each bucket - treated as 4)  
2486 1 10) both in old bucket - treated as 2)  
2487 1  
2488 1 The two-pass empty bucket bucket split cases which occur when the  
2489 1 insertion point is found to be the same as the bucket split point are  
2490 1 handled as follows:  
2491 1  
2492 1 11) second pass two bucket split - treated as 2)  
2493 1 12) first pass multi-bucket split - treated as 4)  
2494 1 13) second pass multi-bucket split - treated as 2)  
2495 1  
2496 1 In all cases it assumes that the new VBNs will be the same size  
2497 1 as the current VBNs. If they actually are larger, the routine that  
2498 1 inserts them will have to grow each one individually.  
2499 1  
2500 1 INPUT  
2501 1  
2502 1 none  
2503 1  
2504 1 IMPLICIT INPUT  
2505 1  
2506 1 IRAB  
2507 1 [EMPTY BUCKET]  
2508 1 [BIG SPLIT]  
2509 1 [VBN-LEFT]  
2510 1 [VBN-MID]  
2511 1 [SPL-IDX]  
2512 1 [POS-INS]  
2513 1 BKT\_ADDR  
2514 1

2515 2565 1 | OUTPUT  
2516 2566 1 | none  
2517 2567 1 | IMPLICIT OUTPUT  
2518 2568 1 |  
2519 2569 1 |  
2520 2570 1 |  
2521 2571 1 | The VBN chain is spread apart and ready to insert new VBNs  
2522 2572 1 |  
2523 2573 1 | ---  
2524 2574 1 |  
2525 2575 2 | BEGIN  
2526 2576 2 |  
2527 2577 2 | EXTERNAL REGISTER  
2528 2578 2 | R\_REC\_ADDR\_STR,  
2529 2579 2 | R\_BKT\_ADDR\_STR,  
2530 2580 2 | R\_IRAB\_STR,  
2531 2581 2 | R\_IDX\_DFN\_STR,  
2532 2582 2 | R\_IFAB\_STR;  
2533 2583 2 |  
2534 2584 2 | GLOBAL REGISTER  
2535 2585 2 | R\_RAB,  
2536 2586 2 | R\_BDB,  
2537 2587 2 | R\_IMPURE;  
2538 2588 2 |  
2539 2589 2 | MACRO  
2540 2590 2 | OFFSET = 0,0,16,0 %;  
2541 2591 2 |  
2542 2592 2 | LOCAL  
2543 2593 2 | FREE\_SPACE : REF BBLOCK, | Address of offset to VBN freespace  
2544 2594 2 | LENGTH, | Number of bytes to move  
2545 2595 2 | SWING\_FLG, | Logical - so we don't have to check swing conditions  
2546 2596 2 | every time around  
2547 2597 2 | INS\_ADDR, | Address where the new VBN will be inserted  
2548 2598 2 | NUM\_BYTES; | Number of bytes which will be inserted  
2549 2599 2 |  
2550 2600 2 | The swing flag is set under the following bucket split cases:  
2551 2601 2 |  
2552 2602 2 | 1. All empty bucket split cases  
2553 2603 2 | 2. First pass multibucket split cases where one key goes  
2554 2604 2 | in each of the two index buckets  
2555 2605 2 | 3. All second pass bucket split cases  
2556 2606 2 |  
2557 2607 2 | SWING\_FLG = 0;  
2558 2608 2 |  
2559 2609 2 | IF .IRAB[IRBSV\_EMPTY\_BKT]  
2560 2610 2 | OR  
2561 2611 3 | (.IRAB[IRBSL\_VBN\_LEFT] EQLU 0)  
2562 2612 2 | OR  
2563 2613 3 | (.IRAB[IRBSV\_SPL\_IDX]  
2564 2614 3 | AND .IRAB[IRBSV\_BIG\_SPLIT]  
2565 2615 3 | AND .IRAB[IRBSW\_POS\_INS] NEQU BKT\$C\_OVERHDSZ)  
2566 2616 2 | THEN  
2567 2617 2 | SWING\_FLG = 1;  
2568 2618 2 |  
2569 2619 2 | Assume a two bucket split with no empties  
2570 2620 2 |  
2571 2621 2 | INS\_ADDR = RMSCNTRL\_ADDR();

```
2572    2622 2     NUM_BYTS = .BKT_ADDR [BKT$V_PTR_SZ] + 2;  
2573    2623 2  
2574    2624 2     | Consider all nonmulti-bucket cases, the second pass of the multibucket  
2575    2625 2     | case where both index records go in the same (old) index bucket, and the  
2576    2626 2     | second pass of the multi-bucket with empty bucket split case.  
2577    2627 2  
2578    2628 3     IF (NOT .IRAB [IRB$V_BIG_SPLIT])  
2579    2629 3     OR  
2580    2630 3     (.IRAB[IRBSL_VBN_MID] EQLA 0)  
2581    2631 2     THEN  
2582    2632 2     BEGIN  
2583    2633 2  
2584    2634 2     IF .SWING_FLG  
2585    2635 2     THEN  
2586    2636 2  
2587    2637 3     | For cases 2, 6, 10, 11, and 13 - just update the VBN in place  
2588    2638 3  
2589    2639 3     NUM_BYTS = 0;  
2590    2640 2  
2591    2641 2  
2592    2642 2  
2593    2643 3     | For cases 1, and 5 - shift the VBNs to make room for 1 new VBN  
2594    2644 3  
2595    2645 3     | Consider all multibucket cases except case 10 and 13 handled above.  
2596    2646 3  
2597    2647 2  
2598    2648 2  
2599    2649 2  
2600    2650 2  
2601    2651 2  
2602    2652 2  
2603    2653 2  
2604    2654 2     | For cases 3, and 8 - shift the VBNs to make room for 2 new VBNs  
2605    2655 2  
2606    2656 2  
2607    2657 2     | For cases 4, 7, 9, and 12 - shift the VBNs to make room for 1 new VBN  
2608    2658 2     Cases 4, and 9 will also update one VBN in place  
2609    2659 2  
2610    2660 2     Now actually shift the VBN chain provided at least one VBN is being added  
2611    2661 2     and it is not being added to the end of the chain (ie the index record(s)  
2612    2662 2     to be added do not have key values higher than all index records currently  
2613    2663 2     in the bucket)  
2614    2664 2     FREE_SPACE = .BKT_ADDR + (.IDX_DFN[IDX$B_IDXBKTSZ] * 512) - BKT$C_ENDOVHD;  
2615    2665 2     LENGTH = .INS_ADDR - (.BKT_ADDR + .FREE_SPACE [OFFSET]);  
2616    2666 2  
2617    2667 3     IF (.LENGTH GTR 0)  
2618    2668 2     AND  
2619    2669 3     (.NUM_BYTS GTR 0)  
2620    2670 2     THEN  
2621    2671 2     RMSMOVE (.LENGTH, .BKT_ADDR + .FREE_SPACE [OFFSET],  
2622    2672 2     .BKT_ADDR + .FREE_SPACE [OFFSET] - .NUM_BYTS);  
2623    2673 2  
2624    2674 2  
2625    2675 2     | Update the free space pointer.  
2626    2676 2     FREE_SPACE [OFFSET] = .FREE_SPACE [OFFSET] - .NUM_BYTS  
2627    2677 2  
2628    2678 1     END;
```

0914 8F BB 00000 RM\$SHFT_VBNS:							
15	44	A9	0088	54 D4 00004	PUSHR	#^M<R2,R4,R8,R11>	2509
				06 E0 00006	CLRL	SWING FLG	2607
				C9 D5 0000B	BBS	#6, 68(IRAB), 1\$	2609
				0F 13 0000F	TSTL	136(IRAB)	2611
09	44	0E	44	A9 E9 00011	BEQL	1\$	
	44	A9	48	02 E1 00015	BLBC	68(IRAB), 2\$	2613
		0E		A9 B1 0001A	BBC	#2, 68(IRAB), 2\$	2614
				03 13 0001E	CMPW	72(IRAB), #14	2615
				54 01 D0 00020	BEQL	2\$	
52	0D	A5	02	0000G 30 00023	MOVL	#1, SWING FLG	2617
			52	03 EF 00026	BSBW	RM\$CNTRL_ADDR	2621
	06	44	A9	02 C0 0002C	EXTZV	#3, #2, T3(BKT_ADDR), NUM_BYTES	2622
			02	E1 0002F	ADDL2	#2, NUM_BYTES	
			02	C9 D5 00034	BBC	#2, 68(IRAB), 3\$	2628
			07	12 00038	TSTL	144(IRAB)	2630
			0A	54 E9 0003A	BNEQ	4\$	
				52 D4 0003D	BLBC	SWING FLG, 5\$	2634
				06 11 0003F	CLRL	NUM_BYTES	2639
			03	54 E8 00041	BRB	5\$	2628
			52	02 C4 00044	BLBS	SWING FLG, 5\$	2649
51			51	16 A7 9A 00047	MULL2	#2, NUM_BYTES	2654
			51	09 78 0004B	MOVZBL	22(IDX_DFN), R1	2664
			51	FC A145 9E 0004F	ASHL	#9, R1-R1	
			54	61 3C 00054	MOVAB	-4(R1)[BKT_ADDR], FREE_SPACE	
			54	55 C0 00057	MOVZWL	(FREE SPACE), R4	2665
			50	54 C2 0005A	ADDL2	BKT_ADDR, R4	
				10 15 0005D	SUBL2	R4, LENGTH	
				52 D5 0005F	BLEQ	6\$	2667
				0C 15 00061	TSTL	NUM_BYTES	2669
7E			54	52 C3 00063	BLEQ	6\$	
				11 BB 00067	SUBL3	NUM_BYTES, R4, -(SP)	2672
			0000G	30 00069	PUSHR	#^M[R0,R4]	2671
			5E	OC C0 0006C	BSBW	RM\$MOVE	
			61	52 A2 0006F	ADDL2	#12, SP	
				0914 8F BA 00072	SUBL2	NUM_BYTES, (FREE SPACE)	2676
				05 00076	POPR	#^M[R2,R4,R8,R11]	2678
					RSB		

; Routine Size: 119 bytes, Routine Base: RMSRMS3 + 081A

: 2629 2679 1

```
: 2631      2680 1 %SBTTL 'RMSV3_IDX_REC'  
: 2632      2681 1 ROUTINE RMSV3_IDX_REC (REC_SZ) : RL$COMMON_LINK NOVALUE =  
: 2633      2682 1 !++  
: 2634      2683 1  
: 2635      2684 1 RMSV3_IDX_REC  
: 2636      2685 1  
: 2637      2686 1 This routine builds a new prologue three index record. It does  
: 2638      2687 1 one of the following  
: 2639      2688 1  
: 2640      2689 1 2) Change a VBN (swing pointer because of empty bucket)  
: 2641      2690 1 3) Add a key and VBN  
: 2642      2691 1 4) Add an additional key and VBN  
: 2643      2692 1  
: 2644      2693 1 CALLING SEQUENCE:  
: 2645      2694 1  
: 2646      2695 1 RMSV3_KEY_REC (REC_SZ)  
: 2647      2696 1  
: 2648      2697 1 INPUT PARAMETERS:  
: 2649      2698 1 REC_SZ - total size of record to be inserted  
: 2650      2699 1  
: 2651      2700 1 IMPLICIT INPUTS:  
: 2652      2701 1 IRAB  
: 2653      2702 1 REC_ADDR  
: 2654      2703 1 BKT_ADDR  
: 2655      2704 1 IDX_DFN  
: 2656      2705 1 IRAB[IRBV_SPL_IDX]  
: 2657      2706 1 IRAB[IRBV_EMPTY_BKT]  
: 2658      2707 1 IRAB[IRBSB_BIG_SPLIT]  
: 2659      2708 1 IRAB[IRBSW_POS_INS]  
: 2660      2709 1 IRAB[IRBSL_VBN_LEFT]  
: 2661      2710 1 IRAB[IRBSL_VBN_MID]  
: 2662      2711 1  
: 2663      2712 1 OUTPUT PARAMETERS:  
: 2664      2713 1  
: 2665      2714 1 NONE  
: 2666      2715 1  
: 2667      2716 1 IMPLICIT OUTPUTS:  
: 2668      2717 1  
: 2669      2718 1 REC_ADDR is updated to point past the record(s) inserted  
: 2670      2719 1 REC_COUNT is incremented by the number of VBN's inserted  
: 2671      2720 1  
: 2672      2721 1  
: 2673      2722 1 ROUTINE VALUE:  
: 2674      2723 1  
: 2675      2724 1 NONE  
: 2676      2725 1  
: 2677      2726 1 --  
: 2678      2727 1  
: 2679      2728 2 BEGIN  
: 2680      2729 2  
: 2681      2730 2 EXTERNAL REGISTER  
: 2682      2731 2 R_IRAB_STR,  
: 2683      2732 2 R_IFAB_STR,  
: 2684      2733 2 R_IDX_DFN_STR,  
: 2685      2734 2 R_REC_ADDR_STR,  
: 2686      2735 2 R_BKT_ADDR_STR;  
: 2687      2736 2
```

```

2688 2737 2 LOCAL
2689 2738 2 SAVE_REC_ADDR;
2690 2739 2
2691 2740 2 MACRO
2692 2741 2 KEY_SZ      = REC_SZ<0,16> %,
2693 2742 2 VBN_SZ       = REC_SZ<16,16>%;
2694 2743 2
2695 2744 2 | Handle all empty bucket cases
2696 2745 2
2697 2746 2 IF .IRAB [IRB$V_EMPTY_BKT]
2698 2747 2 THEN
2699 2748 2
2700 2749 3 BEGIN
2701 2750 3
2702 2751 3 GLOBAL REGISTER
2703 2752 3   R_RAB,
2704 2753 3   R_IMPURE;
2705 2754 3
2706 2755 3 | We will only swing the VBN, if the VBN of the current index record
2707 2756 3 is the same as the VBN of the leftmost primary data bucket. This may
2708 2757 3 not be the case because of a previous index update failure, system
2709 2758 3 crash, or during certain continuation bucket splits. In such cases
2710 2759 3 if we swing the VBN we could be causing crossed pointers to the
2711 2760 3 level below.
2712 2761 3
2713 2762 3 | The only time this VBN swing is not done (besides when the
2714 2763 3 downpointer of the current index record is not equal to the leftmost
2715 2764 3 data VBN) is during the second pass of the two-pass multi-bucket
2716 2765 3 empty bucket split case.
2717 2766 3
2718 2767 3 IF .IRAB [IRB$L_VBN_LEFT] EQA RMSV3_VBN()
2719 2768 3 THEN
2720 2769 3   RMSADD_V3VBN(.IRAB [IRB$L_VBN_MID]);
2721 2770 3
2722 2771 3 | If there was a multi-bucket split involving empty buckets, RMS
2723 2772 3 continues the processing here.
2724 2773 3
2725 2774 3 IF .IRAB[IRB$V_BIG_SPLIT]
2726 2775 3 THEN
2727 2776 4 BEGIN
2728 2777 4
2729 2778 4 | Unless this is the second pass of the two-pass multi-bucket split
2730 2779 4 case insert the key that is in keybuffer 2. If this is a one-pass
2731 2780 4 split, also increment IRB$L_REC_COUNT, the counter of the number
2732 2781 4 of VBNs succeeding the position where the VBN corresponding to
2733 2782 4 the just inserted key is to be inserted.
2734 2783 4
2735 2784 5 IF (.IRAB[IRB$L_VBN_LEFT] NEQA 0)
2736 2785 4 THEN
2737 2786 5   BEGIN
2738 2787 5     RMSADD_V3KEY (KEYBUF_ADDR(2));
2739 2788 5
2740 2789 5     IF NOT .IRAB[IRB$V_SPL_IDX]
2741 2790 5     THEN
2742 2791 5       IRAB[IRB$L_REC_COUNT] = .IRAB[IRB$L_REC_COUNT] + 1
2743 2792 5
2744 2793 4 END;

```

```
: 2745      2794 4
: 2746      2795 4 | Provided this is not the first pass of the two-pass multi-bucket
: 2747      2796 4 | with empty bucket split case, RMS now updates the VBN down pointer
: 2748      2797 4 | associated with the index record representing the high key of the
: 2749      2798 4 | old (left) data bucket before the data bucket split, and its high key
: 2750      2799 4 | became the high key of the rightmost data bucket.
: 2751      2800 4
: 2752      2801 4 | IF NOT .IRAB[IRBSV_SPL_IDX]
: 2753      2802 4 | THEN
: 2754      2803 4 |     RMSADD_V3VBN (.IRAB[IRBSL_VBN_RIGHT])
: 2755      2804 4
: 2756      2805 3 | END:
: 2757      2806 3
: 2758      2807 3 | RETURN
: 2759      2808 3
: 2760      2809 2 | END:
: 2761      2810 2
: 2762      2811 2 | Handle all split cases not involving empty buckets.
: 2763      2812 2
: 2764      2813 2 | We will always add at least one key provided this is not the second pass
: 2765      2814 2 | of one of the three two-pass bucket split cases. If this is one of those
: 2766      2815 2 | three cases, the key we always add was added during the first pass when the
: 2767      2816 2 | old index bucket contents were updated, and should not be again added during
: 2768      2817 2 | this update of the new index bucket contents.
: 2769      2818 2
: 2770      2819 2 | SAVE_REC_ADDR = .REC_ADDR;
: 2771      2820 2
: 2772      2821 2 | IF .IRAB[IRBSL_VBN_LEFT] NEQU 0
: 2773      2822 2 | THEN
: 2774      2823 2 |     RMSADD_V3KEY (KEYBUF_ADDR(2));
: 2775      2824 2
: 2776      2825 2 | On a multibucket split, the key which will trail the two keys added must
: 2777      2826 2 | be recompressed based on the first of the two keys added. This recompression
: 2778      2827 2 | is not done if the keys being added are being added to the end of the index
: 2779      2828 2 | bucket (no trailing key), if the multibucket split case is the two-pass case
: 2780      2829 2 | where a key is added to each of the two index buckets and this is either of
: 2781      2830 2 | the two passes, or if index compression is not enabled (of course).
: 2782      2831 2
: 2783      2832 3 | IF (.IRAB[IRBSV_BIG_SPLIT]
: 2784      2833 3 | AND NOT .IRAB[IRBSV_SPL_IDX]
: 2785      2834 3 | AND .IRAB[IRBSL_VBN_LEFT] NEQU 0)
: 2786      2835 2 | AND
: 2787      2836 3 |     (.IDX_DFN[IDX$V_IDX_COMPR]
: 2788      2837 4 |     AND (.SAVE_REC_ADDR + .KEY_SZ LSSA .BKT_ADDR +
: 2789      2838 3 |                           .BKT_ADDR[BKT$W_FREESPACE]))
: 2790      2839 2 | THEN
: 2791      2840 2 |     RMSRECOMPRESS_KEY (.SAVE_REC_ADDR , .SAVE_REC_ADDR + .KEY_SZ);
: 2792      2841 2
: 2793      2842 2 | A key and VBN must be added for all multibucket split cases. For the one-
: 2794      2843 2 | pass multibucket split case and the two-pass multibucket split case this
: 2795      2844 2 | key is the second of two keys to be added during this current pass and the
: 2796      2845 2 | first VBN. For the latter case, both keys and the VBN get added during the
: 2797      2846 2 | first pass and none get added during the second pass. For the two-pass
: 2798      2847 2 | multibucket split case where a key is inserted in each of the old and new
: 2799      2848 2 | index buckets, this key and its corresponding VBN are inserted as the low
: 2800      2849 2 | order key in the new index bucket during the second pass having inserted
: 2801      2850 2 | the first key in the old index bucket during the first pass.
```

```

2802      2      !
2803      2      IF .IRAB [IRBSV_BIG_SPLIT]
2804      2      AND
2805      4      ((NOT .IRAB[IRBSV_SPL_IDX]
2806      4      AND .IRAB[IRBSL_VBN_MID] NEQU 0)
2807      3      OR
2808      4      (.IRAB[IRBSV_SPL_IDX]
2809      3      AND .IRAB[IRBSW_POS_INS] EQLU BKTSC_OVERHDSZ))
2810      2      THEN
2811      2      BEGIN
2812      2
2813      3      IF .IRAB[IRBSL_VBN_LEFT] NEQU 0
2814      3      THEN
2815      3      IRAB[IRBSL_REC_COUNT] = .IRAB[IRBSL_REC_COUNT] + 1;
2816      3
2817      3      RMSADD_V3VBN (.IRAB[IRBSL_VBN_MID]);
2818      3
2819      3      IRAB[IRBSV_BIG_SPLIT] = 0;
2820      3      RMSADD_V3KEY (KEYBUF_ADDR(3));
2821      3      IRAB[IRBSV_BIG_SPLIT] = 1
2822      2
2823      2
2824      2
2825      2      ! A VBN update is always required for one-pass bucket split cases. This is
2826      2      because the VBN pointer of at least one of the index records on the current
2827      2      level has changed because of the split at the level below it. For all
2828      2      two-pass bucket split cases, this VBN updating must also take place, for the
2829      2      same reason. However, it always takes place during the second of the two
2830      2      passes when the contents of the new index bucket are being modified since it
2831      2      is the VBN pointer of one of those index records that has changed.
2832      2
2833      2      IF NOT .IRAB[IRBSV_SPL_IDX]
2834      2      THEN
2835      2      BEGIN
2836      2
2837      3      IF (.IRAB[IRBSL_VBN_LEFT] NEQU 0)
2838      3      OR
2839      4      (.IRAB[IRBSL_VBN_MID] NEQU 0)
2840      3      THEN
2841      3      IRAB [IRBSL_REC_COUNT] = .IRAB [IRBSL_REC_COUNT] + 1;
2842      3
2843      3      RMSADD_V3VBN (.IRAB [IRBSL_VBN_RIGHT])
2844      3
2845      2
2846      2
2847      1      END;

```

				0900 8F BB 00000 RMSV3_IDX REC:		
40	44	A9		06 E1 00004	PUSHR	#^M<R8,R11>
				0000G 30 00009	BBC	#6 68(IRAB), SS
	50	0088		C9 D1 0000C	BSBW	RM\$V3 VBN
				0A 12 00011	CMPL	136(IRAB), R0
					BNEQ	1\$

2681  
2746  
2767

			0090	C9	DD 00013	PUSHL	144(IRAB)	2769	
			F828	30	00017	BSBW	RMSADD_V3VBN		
			04	C0	0001A	ADDL2	#4, SP-		
			02	E0	0001D	1\$: ADDL2	#2, 68(IRAB), 3\$	2774	
			00D1	31	00022	2\$: BBS	15\$		
			0088	C9	D5 00025	3\$: BRW	136(IRAB)	2784	
			17	13	00029	TSTL	4\$		
			50	00B4	CA 3C 0002B	BEQL	180(IFAB), R0	2787	
			60	B940	9F 00030	MOVZWL	#96(IRAB)[R0]		
			F78D	30	00034	PUSHAB	RMSADD_V3KEY		
			5E	04	C0 00037	BSBW	#4, SP-		
			E4	44	A9 E8 0003A	ADDL2	68(IRAB), 2\$	2789	
			0094	C9	D6 0003E	BLBS	148(IRAB)	2791	
			DC	44	A9 E8 00042	INCL	68(IRAB), 2\$	2801	
			00A3	31	00046	BLBS	14\$	2803	
			58	56	D0 00049	BRW	REC_ADDR, SAVE_REC_ADDR	2819	
			0088	C9	D5 0004C	MOVL	136(IRAB)	2821	
			0F	13	00050	TSTL	6\$		
			50	00B4	CA 3C 00052	BEQL	180(IFAB), R0	2823	
			60	B940	9F 00057	MOVZWL	#96(IRAB)[R0]		
			F766	30	0005B	PUSHAB	RMSADD_V3KEY		
			04	C0 0005E	BSBW	ADDL2	#4, SP-		
			72	44	A9 2C	02 E1 00061	BBC	68(IRAB), 11\$	2832
			44	A9	E8 00066	BLBS	68(IRAB), 7\$	2833	
			0088	C9	D5 0006A	TSTL	136(IRAB)	2834	
			26	13	0006E	BEQL	7\$		
			21	1C	A7	03 E1 00070	BBC	#3, 28(IDX_DFN), 7\$	2836
			50	0C	AE 3C 00075	MOVZWL	REC_SZ, R0	2837	
			50	58	C0 00079	ADDL2	SAVE REC_ADDR, R0		
			51	04	A5 3C 0007C	MOVZWL	4(BKT ADDR), R1	2838	
			51	55	C0 00080	ADDL2	BKT_ADDR, R1		
			51	50	D1 00083	CMPL	RO, R1	2837	
			50	0E	1E 00086	BGEQU	7\$		
			51	50	AE 3C 00088	MOVZWL	REC_SZ, R0	2840	
			58	50	C1 0008C	ADDL3	RO, SAVE_REC_ADDR, R1		
			50	58	D0 00090	MOVL	SAVE REC_ADDR, R0		
			0000G	30	00093	BSBW	RMSRECOMPR KEY		
			3D	44	A9 0A	02 E1 00096	BBC	68(IRAB), 11\$	2852
			44	A9	E8 0009B	BLBS	68(IRAB), 8\$	2854	
			0090	C9	D5 0009F	TSTL	144(IRAB)	2855	
			33	0A	12 000A3	BNEQ	9\$		
			0E	44	A9 E9 000A5	BLBC	68(IRAB), 12\$	2857	
			48	A9	B1 000A9	CMPW	72(IRAB), #14	2858	
			0088	29	12 000AD	BNEQ	11\$		
			0094	C9	D5 000AF	TSTL	136(IRAB)	2862	
			0090	C9	D6 000B5	BEQL	10\$		
			F782	DD	000B9	INCL	148(IRAB)	2864	
			44	A9	8A 000C0	PUSHL	144(IRAB)	2866	
			50	00B4	CA 3C 000C4	BSBW	RMSADD_V3VBN		
			6E	60	B940 3E 000C9	BICB2	#4, 68(IRAB)	2868	
			F6F3	30	000CE	MOVZWL	180(IFAB), R0	2869	
			44	5E	04 C0 000D1	MOVAW	#96(IRAB)[R0], (SP)		
			1A	A9	88 000D4	BSBW	RMSADD_V3KEY		
			0088	04	E8 000D8	ADDL2	#4, SP-	2870	
			11\$:	C9	D5 000DC	BISB2	68(IRAB), 15\$	2882	
			12\$:			BLBS	136(IRAB)	2886	
						TSTL			

RM3MISPUT  
V04-000

RMSV3\_IDX\_REC

H 6  
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0090	06	12	000E0	BNEQ	13\$		
	C9	D5	000E2	TSTL	144(IRAB)	2888	
	04	13	000E6	BEQL	14\$		
0094	C9	D6	000E8	13\$:	INCL	148(IRAB)	
008C	C9	DD	000EC	14\$:	PUSHL	140(IRAB)	
	F74F	30	000F0	BSBW	RMSADD V5VBN	2890	
5E	04	C0	000F3	ADDL2	#4, SP	2892	
	0900	8F	BA	000F6	15\$:	POPR	#^M<R8,R11>
			05	000FA		RSB	2896

; Routine Size: 251 bytes,    Routine Base: RMSRMS3 + 0891

2849 2897 1 XSBTTL 'RMSV3KEY\_SZ'  
2850 2898 1 ROUTINE RMSV3KEY\_SZ (KEY\_ADDR) : RL\$COMMON\_LINK =  
2851 2899 1 !++  
2852 2900 1 FUNCTIONAL DESCRIPTION:  
2853 2901 1 This routine computes the number of chars needed to store the  
2854 2902 1 prologue three key in KEY\_ADDR at REC\_ADDR.  
2855 2903 1 CALLING SEQUENCE:  
2856 2904 1 RMSV3KEY\_SZ()  
2857 2905 1 INPUT PARAMETERS:  
2858 2906 1 KEY\_ADDR - Address of the key to be inserted  
2859 2907 1 IMPLICIT INPUTS:  
2860 2908 1 BKT\_ADDR - Pointer to the current bucket  
2861 2909 1 IDX\_DFN - Pointer to index descriptor structure  
2862 2910 1 IFAB - Pointer to internal FAB  
2863 2911 1 IRAB - Pointer to Internal RAB  
2864 2912 1 REC\_ADDR - Address to insert key  
2865 2913 1 OUTPUT PARAMETERS:  
2866 2914 1 NONE  
2867 2915 1 IMPLICIT OUTPUTS:  
2868 2916 1 NONE  
2869 2917 1 ROUTINE VALUE:  
2870 2918 1 R0 - number of bytes needed to store the key after front compression  
2871 2919 1 and rear-end truncation (doesn't include two bytes of overhead).  
2872 2920 1 --  
2873 2921 1 BEGIN  
2874 2922 1 MAP  
2875 2923 1 KEY\_ADDR : REF BBLOCK;  
2876 2924 1 EXTERNAL REGISTER  
2877 2925 1 R\_BKT\_ADDR\_STR,  
2878 2926 1 R\_IDX\_DFN\_STR,  
2879 2927 1 R\_REC\_ADDR\_STR,  
2880 2928 1 R\_IFAB\_STR,  
2881 2929 1 R\_IRAB\_STR;  
2882 2930 1 GLOBAL REGISTER  
2883 2931 1 R\_IMPURE,  
2884 2932 1 R\_RAB,  
2885 2933 1 R\_BDB;  
2886 2934 1 LOCAL  
2887 2935 1  
2888 2936 2  
2889 2937 2  
2890 2938 2  
2891 2939 2  
2892 2940 2  
2893 2941 2  
2894 2942 2  
2895 2943 2  
2896 2944 2  
2897 2945 2  
2898 2946 2  
2899 2947 2  
2900 2948 2  
2901 2949 2  
2902 2950 2  
2903 2951 2  
2904 2952 2  
2905 2953 2

```

2906    2954 2      KEYBUF : REF BBLOCK;
2907    2955 2
2908    2956 2      ! Data level.
2909    2957 2
2910    2958 2      IF .BKT_ADDR [BKT$B_LEVEL] EQLU 0
2911    2959 2      THEN
2912    2960 2          BEGIN
2913    2961 2
2914    2962 2          IF NOT .IDX_DFN [IDX$V_KEY_COMPR]
2915    2963 2          THEN
2916    2964 2              RETURN .IDX_DFN [IDX$B_KEYSZ];
2917    2965 2
2918    2966 2
2919    2967 3      ! Index level.
2920    2968 3
2921    2969 2
2922    2970 2      ELSE
2923    2971 2          BEGIN
2924    2972 3
2925    2973 3          IF NOT .IDX_DFN [IDX$V_IDX_COMPR]
2926    2974 3          THEN
2927    2975 3              RETURN .IDX_DFN [IDX$B_KEYSZ];
2928    2976 2
2929    2977 2      ! Must be a compressed key. Move key into keybuffer 5, compress it, and
2930    2978 2      return its size (not including the two byte of compression overhead).
2931    2979 2
2932    2980 2      KEYBUF = KEYBUF_ADDR(5);
2933    2981 2      RM$MOVE (.KEY_ADDR[KEY_LEN] + 2, .KEY_ADDR, .KEYBUF);
2934    2982 2
2935    2983 2      RM$COMPRESS_KEY (.KEYBUF);
2936    2984 2
2937    2985 2      RETURN .KEYBUF[KEY_LEN]
2938    2986 2
2939    2987 1      END;

```

			0914	8F BB 00000 RMSV3KEY_SZ:			
			0C	A5 95 00004	PUSHR	#^M<R2,R4,R8,R11>	2898
				07 12 00007	TSTB	12(BKT_ADDR)	2958
0D	1C	A7		06 E0 00009	BNEQ	1\$	
				05 11 0000E	BBS	#6, 28(IDX_DFN), 3\$	2962
06	1C	A7	50	03 E0 00010 1\$:	BRB	2\$	2964
			20	A7 9A 00015 2\$:	BBS	#3, 28(IDX_DFN), 3\$	2972
				25 11 00019	MOVZBL	32(IDX_DFNT, R0	2974
			50	CA 3C 0001B 3\$:	BRB	4\$	
			52	00B4 60 B940 DE 00020	MOVZWL	180(IFAB), R0	2980
				52 DD 00025	MOVAL	@96(IRAB)[R0], KEYBUF	
				18 AE DD 00027	PUSHL	KEYBUF	2981
			7E	1C BE 9A 0002A	PUSHL	KEY ADDR	
			6E	02 C0 0002E	MOVZBL	@KEY ADDR, -(SP)	
				0000G 30 00031	ADDL2	#2, TSP)	
			5E	OC C0 00034	BSBW	RM\$MOVE	
			50	52 D0 00037	ADDL2	#12, SP	
					MOVL	KEYBUF, R0	2983

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RMSV3KEY\_SZ

K 6  
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RM3  
V04

50      0000G 30 0003A      BSBW      RMS\$COMPRESS KEY  
      62 9A 0003D      MOVZBL      (KEYBUF), R0  
0914 8F BA 00040 4\$:      POPR      #^M<R2,R4,R8,R11>  
      05 00044      RSB

; 2985  
; 2987

; Routine Size: 69 bytes,    Routine Base: RMSRMS3 + 098C

```
2941 2988 1 %SBTTL 'RMSVBN_SIZE'
2942 2989 1 GLOBAL ROUTINE RMSVBN_SIZE (VBN) : RL$PRESERVE1 =
2943 2990 1 ++
2944 2991 1 FUNCTIONAL DESCRIPTION:
2945 2992 1 Calculate number of bytes required to describe input VBN
2946 2993 1 CALLING SEQUENCE:
2947 2994 1 RMSVBN_SIZE (VBN)
2948 2995 1 INPUT PARAMETERS:
2949 2996 1 VBN - VBN whose size is to be determined
2950 2997 1 IMPLICIT INPUTS:
2951 2998 1 NONE
2952 2999 1 OUTPUT PARAMETERS:
2953 3000 1 NONE
2954 3001 1 IMPLICIT OUTPUTS:
2955 3002 1 NONE
2956 3003 1 ROUTINE VALUE:
2957 3004 1 number of bytes needed to describe VBN
2958 3005 1 SIDE EFFECTS:
2959 3006 1 NONE
2960 3007 1 --
2961 3008 1 BEGIN
2962 3009 1 LOCAL
2963 3010 1 SIZE;
2964 3011 1 MAP
2965 3012 1 VBN : VECTOR [4, BYTE];
2966 3013 1 SIZE = 3;
2967 3014 1 DO
2968 3015 1 IF .VBN[.SIZE] NEQ 0
2969 3016 1 THEN
2970 3017 1 BEGIN
2971 3018 1 SIZE = .SIZE + 1;
2972 3019 1 RETURN .SIZE
2973 3020 1 END
2974 3021 1 ELSE
2975 3022 1 SIZE = .SIZE - 1
2976 3023 1 UNTIL .SIZE EQL 1;
2977 3024 1 SIZE = .SIZE + 1
```

RM3MISPUT  
V04-000

RMSVBN\_SIZE

: 2998

3045 1 END;

M 6  
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RM3  
V04

50	03	DO 00000 RMSVBN_SIZE::				
	04 AE40	95 00003 1\$:	MOVL	#3, SIZE		3028
	07	12 00007	TSTB	VBN[SIZE]		3032
	50	D7 00009	BNEQ	2\$		
01	50	D1 0000B	DECL	SIZE		3040
	F3	12 0000E	CMPL	SIZE, #1		3042
	50	D6 00010 2\$:	BNEQ	1\$		
	05	00012	INCL	SIZE		3044
			RSB			3045

: Routine Size: 19 bytes, Routine Base: RMSRMS3 + 09D1

: 2999 3046 1  
: 3000 3047 1 END  
: 3001 3048 1  
: 3002 3049 0 ELUDOM

#### PSECT SUMMARY

Name	Bytes	Attributes
RMSRMS3	2532	NOVEC,NOWRT, RD , EXE,NOSHR, GBL, REL, CON, PIC,ALIGN(2)

#### Library Statistics

File	-----	Symbols	-----	Pages	Processing
	Total	Loaded	Percent	Mapped	Time
\$_\$255\$DUA28:[RMS.OBJ]RMS.L32;1	3109	81	2	154	00:00.4

: Information: 2  
: Warnings: 0  
: Errors: 0

#### COMMAND QUALIFIERS

: BLISS/CHECK=(FIELD,INITIAL,OPTIMIZE)/LIS=LISS:RM3MISPUT/OBJ=OBJ\$:RM3MISPUT MSRC\$:RM3MISPUT/UPDATE=(ENH\$:RM3MISPUT)

RM3MISPUT  
V04-000

RMSVBN\_SIZE

N 6  
16-Sep-1984 01:51:28  
14-Sep-1984 13:01:29 VAX-11 Bliss-32 V4.0-742  
[RMS.SRC]RM3MISPUT.B32;1

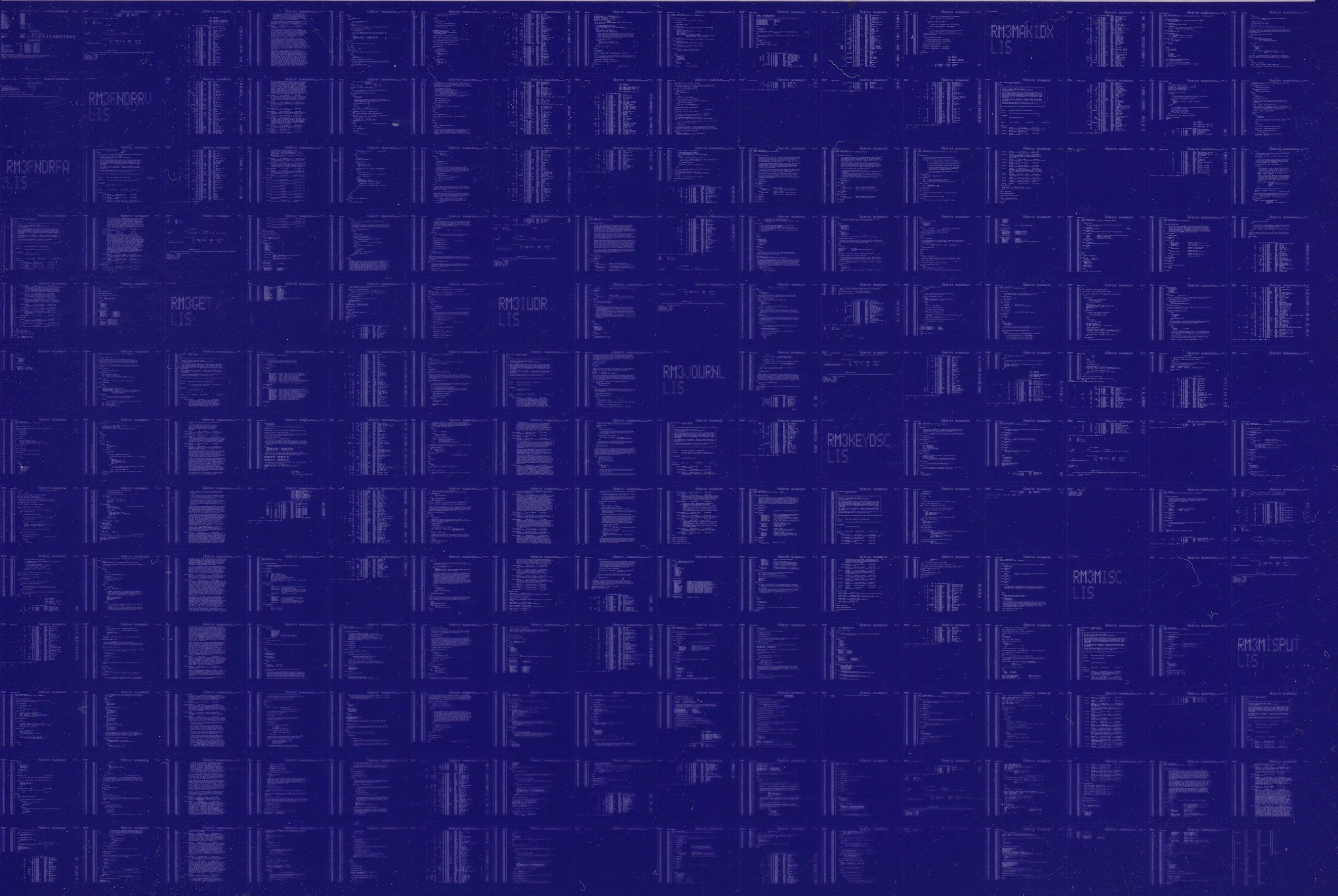
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; 3003 3050 0  
; Size: 2532 code + 0 data bytes  
; Run Time: 00:58.1  
; Elapsed Time: 02:02.7  
; Lines/CPU Min: 3148  
; Lexemes/CPU-Min: 15870  
; Memory Used: 287 pages  
; Compilation Complete

RM3  
V04

0325 AH-BT13A-SE  
VAX/VMS V4.0

DIGITAL EQUIPMENT CORPORATION  
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0326 AH-BT13A-SE  
VAX/VMS V4.0

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CONFIDENTIAL AND PROPRIETARY

RM3NEXTRE  
LIS

RM3OPEN  
LIS

RM3POSRFA  
LIS

RM3PCKUP  
LIS

RM3POSKEY  
LIS

RM3POSSEQ  
LIS