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RRRRRRRR      MM      MM      333333      GGGGGGGG      EEEEEEEEEEE      TTTTTTTTTTT
RRRRRRRR      MM      MM      333333      GGGGGGGG      EEEEEEEEEEE      TTTTTTTTTTT
RR      RR      MMMM      MMMM      33      33      GG      GG      EE      TT
RR      RR      MMMM      MMMM      33      33      GG      GG      EE      TT
RR      RR      MM      MM      33      33      GG      GG      EE      TT
RR      RR      MM      MM      33      33      GG      GG      EE      TT
RRRRRRRR      MM      MM      33      33      GG      GG      EEEEEEEEE      TT
RRRRRRRR      MM      MM      33      33      GG      GG      EEEEEEEEE      TT
RR      RR      MM      MM      33      33      GG      GGGGGG      EE      TT
RR      RR      MM      MM      33      33      GG      GGGGGG      EE      TT
RR      RR      MM      MM      33      33      GG      GG      EE      TT
RR      RR      MM      MM      33      33      GG      GG      EE      TT
RR      RR      MM      MM      33      33      GG      GG      EE      TT
RR      RR      MM      MM      333333      GGGGGG      EEEEEEEEEEE      TT
RR      RR      MM      MM      333333      GGGGGG      EEEEEEEEEEE      TT

```

```

LL      111111      SSSSSSSS
LL      111111      SSSSSSSS
LL      11      SS
LL      11      SS
LL      11      SS
LL      11      SS
LL      11      SSSSSS
LL      11      SSSSSS
LL      11      SS
LL      11      SS
LL      11      SS
LL      11      SS
LLLLLLLLLLLL 111111  SSSSSSSS
LLLLLLLLLLLL 111111  SSSSSSSS

```

1  
2  
3  
4  
5  
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8  
9  
10  
11  
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57

```

0001 0 MODULE RM3GET (LANGUAGE (BLISS32) ,
0002 0             IDENT = 'V04-000'
0003 0             ) =
0004 1 BEGIN
0005 1
0006 1 *****
0007 1 *
0008 1 *   COPYRIGHT (c) 1978, 1980, 1982, 1984 BY
0009 1 *   DIGITAL EQUIPMENT CORPORATION, MAYNARD, MASSACHUSETTS.
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0018 1 *
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0020 1 *   AND SHOULD NOT BE CONSTRUED AS A COMMITMENT BY DIGITAL EQUIPMENT
0021 1 *   CORPORATION.
0022 1 *
0023 1 *   DIGITAL ASSUMES NO RESPONSIBILITY FOR THE USE OR RELIABILITY OF ITS
0024 1 *   SOFTWARE ON EQUIPMENT WHICH IS NOT SUPPLIED BY DIGITAL.
0025 1 *
0026 1 *
0027 1 *****
0028 1
0029 1 ++
0030 1
0031 1 FACILITY:      RMS32 INDEX SEQUENTIAL FILE ORGANIZATION
0032 1
0033 1 Abstract:
0034 1           This module implements the get and find record operations
0035 1           for the indexed file organization.
0036 1
0037 1
0038 1 ENVIRONMENT:
0039 1
0040 1           VAX/VMS OPERATING SYSTEM
0041 1
0042 1 --
0043 1
0044 1
0045 1 AUTHOR:      E. H. MARISON      CREATION DATE:      18-APR-78  13:11
0046 1
0047 1
0048 1 MODIFIED BY:
0049 1
0050 1           V03-025  JWT0193      Jim Teague      13-Aug-1984
0051 1           Fix bug in re-accessing records after they have been
0052 1           found to be locked.  If a process had done a $GET on
0053 1           a record, then a $RELEASE, and then had to wait to
0054 1           $GET the record a second time, too much context was
0055 1           still around from the first $GET.  This caused problems
0056 1           when the sought-after record had been deleted.  RMS
0057 1           treated the $GET + $RELEASE + $GET case just like

```

58 0058 1  
59 0059 1  
60 0060 1  
61 0061 1  
62 0062 1  
63 0063 1  
64 0064 1  
65 0065 1  
66 0066 1  
67 0067 1  
68 0068 1  
69 0069 1  
70 0070 1  
71 0071 1  
72 0072 1  
73 0073 1  
74 0074 1  
75 0075 1  
76 0076 1  
77 0077 1  
78 0078 1  
79 0079 1  
80 0080 1  
81 0081 1  
82 0082 1  
83 0083 1  
84 0084 1  
85 0085 1  
86 0086 1  
87 0087 1  
88 0088 1  
89 0089 1  
90 0090 1  
91 0091 1  
92 0092 1  
93 0093 1  
94 0094 1  
95 0095 1  
96 0096 1  
97 0097 1  
98 0098 1  
99 0099 1  
100 0100 1  
101 0101 1  
102 0102 1  
103 0103 1  
104 0104 1  
105 0105 1  
106 0106 1  
107 0107 1  
108 0108 1  
109 0109 1  
110 0110 1  
111 0111 1  
112 0112 1  
113 0113 1  
114 0114 1

a \$FIND + \$GET case, and would end up with the wrong record.

Also, improve \$GET/\$FIND performance. Leave the infinite GET\_RECORD loop immediately if GET\_RECORD returns an unqualified success status. Formerly, RMS was forced to grind through an unbelievably perverted IF test EVERY TIME it returned from GET\_RECORD.

V03-024 TSK0001 Tamar Krichevsky 15-Jun-1983  
Change addressing mode to long relative for RMSRU\_RECLAIM.

V03-023 MCN0015 Maria del C. Nasr 24-Mar-1983  
More linkages reorganization.

V03-022 TMK0015 Todd M. Katz 11-Mar-1983  
If RMS had to wait for a record lock, and it must re-position to the primary data record by calling RMS\$FIND\_BY\_RRV, then make sure the primary data bucket containing the record is locked exclusively if the possibility exists that some reclamation maybe done (the file is write accessed and RU Journallable).

V03-021 MCN0014 Maria del C. Nasr 24-Feb-1983  
Reorganize linkages

V03-020 TMK0014 Todd M. Katz 14-Jan-1983  
Add support for Recovery Unit Journalling and RU ROLLBACK Recovery of ISAM files. Support involves modifications to RMS\$GET3B and RMS\$GET\_RECORD.

The purpose of the routines within this module is to retrieve a non-deleted primary data record by the user specified access mode. If during its search for such a record RMS in its low-level routines encounters records that are marked RU\_DELETE, RMS will try and delete them for good at this time provided it has write access to the file and the Recovery Unit in which they were deleted has completed successfully.

If RMS is able to delete a primary data record marked RU\_DELETE in these low-level routines, then RMS proceeds to continue looking for a non-deleted primary data record just as if it had encountered a deleted record in the first place. Likewise, if RMS is unable to delete a record that is marked RU\_DELETE because it does not have write access to the file, it merely continues its search. However, if RMS is unable to delete the record for good in these low-level routines because the Recovery Unit in which it was marked RU\_DELETE has not successfully terminated, then RMS returns this record as if it was the non-deleted primary data record to be returned, and lets a higher-level routine decide whether or not to wait for the Recovery Unit in which the record was deleted to complete, or to return an error to the user.

The routines within this module are the high-level routines which decide what to do with RU\_DELETED records that are returned from the low-level positioning routines.

```

: 115 0115 1
: 116 0116 1
: 117 0117 1
: 118 0118 1
: 119 0119 1
: 120 0120 1
: 121 0121 1
: 122 0122 1
: 123 0123 1
: 124 0124 1
: 125 0125 1
: 126 0126 1
: 127 0127 1
: 128 0128 1
: 129 0129 1
: 130 0130 1
: 131 0131 1
: 132 0132 1
: 133 0133 1
: 134 0134 1
: 135 0135 1
: 136 0136 1
: 137 0137 1
: 138 0138 1
: 139 0139 1
: 140 0140 1
: 141 0141 1
: 142 0142 1
: 143 0143 1
: 144 0144 1
: 145 0145 1
: 146 0146 1
: 147 0147 1
: 148 0148 1
: 149 0149 1
: 150 0150 1
: 151 0151 1
: 152 0152 1
: 153 0153 1
: 154 0154 1
: 155 0155 1
: 156 0156 1
: 157 0157 1
: 158 0158 1
: 159 0159 1
: 160 0160 1
: 161 0161 1
: 162 0162 1
: 163 0163 1
: 164 0164 1
: 165 0165 1
: 166 0166 1
: 167 0167 1
: 168 0168 1
: 169 0169 1
: 170 0170 1
: 171 0171 1

```

1. If RMS is unable to lock such a record because another process currently has it locked, then an RLK error is returned.
2. If RMS is able to lock such a record, regardless of whether it had to wait for it or not, then if it finds that the record is not marked RU DELETE it will return it provided all other normal conditions have been met.
3. If on the other hand, RMS finds that the record is still marked RU DELETE after it has locked it, then it will delete the record for good at this time (if the stream has write access to the file), and continue the search for a non-deleted primary data record provided the access mode is not by RFA.

I have also made two other changes in support of RU Journalling and Recovery. First, the ROP bit RABSV\_NLK is totally ignored whenever a stream is currently within a Recovery Unit. Finally, it is also possible that a RU\_UPDATE marked record might be re-formatted before releasing the bucket in which it is found provided the stream has write access to the file. The record being re-formatted in this case can only be the record that is to be returned as the non-deleted primary data record.

I have made an additional change to RMSGET RECORD. If RMS is currently randomly positioning by key to what it thinks is the current record, then it query locks the current record to make sure that this record is in fact locked to avoid a window in which the record is deleted between the time the record lock is released, and the bucket in which the record is found is accessed. If the user has specified record waiting it is disabled for this query lock. Currently it is disabled by clearing the RABSV\_WAIT bit if it is set, and then re-establishing its state after the query lock. The state bit IRBSV\_NO\_Q\_WAIT maybe set to accomplish this same thing and it avoids modifying the user's control block.

I have created a routine RMSPOS\_RFA whose functionality parallels that of RMSPOS\_SEQ and RMSPOS\_KEY. That is, the routine RMSGET\_RECORD will call this routine whenever it is to position to the next primary data record by RFA instead of performing the positioning itself.

V03-019 TMK0013            Todd M. Katz            09-Nov-1982  
 Fix a bug in record unlocking. Whenever RMS must wait for a record lock (the RABSV\_WAIT ROP bit is set), and upon being granted the lock finds that the record it has been waited on has been deleted, RMS must perform a re-positioning. (There is one exception to this rule. If RMS was accessing the record by its RFA then the record deleted error is returned.) RMS must also perform a re-positioning whenever it is positioning by means of an alternate key and has had to wait for a record lock. As part of this re-positioning, RMS must release the lock it obtained during the prior positioning attempt. The problem is that RMS was using the wrong RFA when it went to release the record

```

: 172 0172 1
: 173 0173 1
: 174 0174 1
: 175 0175 1
: 176 0176 1
: 177 0177 1
: 178 0178 1
: 179 0179 1
: 180 0180 1
: 181 0181 1
: 182 0182 1
: 183 0183 1
: 184 0184 1
: 185 0185 1
: 186 0186 1
: 187 0187 1
: 188 0188 1
: 189 0189 1
: 190 0190 1
: 191 0191 1
: 192 0192 1
: 193 0193 1
: 194 0194 1
: 195 0195 1
: 196 0196 1
: 197 0197 1
: 198 0198 1
: 199 0199 1
: 200 0200 1
: 201 0201 1
: 202 0202 1
: 203 0203 1
: 204 0204 1
: 205 0205 1
: 206 0206 1
: 207 0207 1
: 208 0208 1
: 209 0209 1
: 210 0210 1
: 211 0211 1
: 212 0212 1
: 213 0213 1
: 214 0214 1
: 215 0215 1
: 216 0216 1
: 217 0217 1
: 218 0218 1
: 219 0219 1
: 220 0220 1
: 221 0221 1
: 222 0222 1
: 223 0223 1
: 224 0224 1
: 225 0225 1
: 226 0226 1
: 227 0227 1
: 228 0228 1

```

whenever it was re-positioning because the record it had to wait for had been deleted while it was waiting for it. The RFA it was using was the RFA of the current record. This record had been locked during the previous positioning operation, and had been unlocked during the first positioning attempt of the current operation. The lock RMS wants to release is for the record locked during the previous positioning attempt.

The fix for this problem is relatively straightforward. RMS never has to re-position unless it has had to wait for a record lock. Therefore, what I did was set the state bit OK\_WAT\_STATUS whenever a re-positioning has to be done instead of just setting it whenever the re-positioning is being done because RMS had to wait for a record lock while positioning along an alternate index. The setting of this state bit forces RMS to unlock the correct record during the re-positioning attempt.

V03-018 TMK0012 Todd M. Katz 29-Oct-1982  
Make sure that RMS has the index descriptor for the primary key before the size of the primary data record to be returned is determined, the record unpacked (if the file is a prologue 3 file), and the record moved into the user's buffer.

V03-017 TMK0011 Todd M. Katz 11-Oct-1982  
Fix a record locking bug. Whenever the ROP bit RAB\$V\_WAT is set the possibility exists that RMS might have to wait for a record lock. If RMS is positioning by means of an alternate index, and has to wait for such a record lock, then it had to give up the SIDR bucket while it was waiting. Because it gave up the SIDR bucket, the information which it has in order to update the NRP list can no longer be considered valid. Since there is no way for RMS to easily re-access the SIDR bucket, RMS must re-position to it by re-calling GET\_RECORD. Part of this re-positioning includes unlocking the very same primary data record which it had to wait for a record lock on. Unfortunately, GET\_RECORD uses the NRP information to unlock primary data records, and RMS of course, didn't get to the point where it updated the NRP! Therefore, RMS is either not unlocking any record, or it is unlocking the wrong record. Both cases represent errors.

To fix this what I have done is added an input parameter to GET\_RECORD. If it is set, RMS is re-positioning because of the above mentioned problem, and uses the RFA internally saved from the prior positioning attempt to unlock the record; otherwise, the RFA from the current record saved as part of the NRP context is used. Furthermore, whenever RMS does such a re-positioning, it now notes that it had to do so because of an OK\_WAT success status positioning on an alternate key of reference. If it is successful at re-positioning, it sets the status to OK\_WAT which represents the status that it would have returned if the re-positioning had not been necessary.

I have made two additional changes concerning when re-positioning is required. First, if RMS is positioning by key value, and after waiting for a record lock finds that the record it has been waited for has been deleted, then RMS will perform a

```

: 229 0229 1
: 230 0230 1
: 231 0231 1
: 232 0232 1
: 233 0233 1
: 234 0234 1
: 235 0235 1
: 236 0236 1
: 237 0237 1
: 238 0238 1
: 239 0239 1
: 240 0240 1
: 241 0241 1
: 242 0242 1
: 243 0243 1
: 244 0244 1
: 245 0245 1
: 246 0246 1
: 247 0247 1
: 248 0248 1
: 249 0249 1
: 250 0250 1
: 251 0251 1
: 252 0252 1
: 253 0253 1
: 254 0254 1
: 255 0255 1
: 256 0256 1
: 257 0257 1
: 258 0258 1
: 259 0259 1
: 260 0260 1
: 261 0261 1
: 262 0262 1
: 263 0263 1
: 264 0264 1
: 265 0265 1
: 266 0266 1
: 267 0267 1
: 268 0268 1
: 269 0269 1
: 270 0270 1
: 271 0271 1
: 272 0272 1
: 273 0273 1
: 274 0274 1
: 275 0275 1
: 276 0276 1
: 277 0277 1
: 278 0278 1
: 279 0279 1
: 280 0280 1
: 281 0281 1
: 282 0282 1
: 283 0283 1
: 284 0284 1
: 285 0285 1

```

re-positioning to the next record which matches the search key in keybuffer 2 according to the characteristics of the search. Formerly, RMS was just returning a record deleted error, but I believe the other approach has more merit. Second, whenever RMS is positioning by an alternate key of reference (sequentially or randomly by key value), and must wait for a record lock, then RMS must re-position to re-establish the NRP information for the SIDR. Formerly, this re-positioning was not done if RMS was performing a random \$FIND. However, since the stream which has the record locked can delete the SIDR array positioned to by the waiting stream without deleting the actual primary data record (by means of an \$UPDATE), then as the record eventually returned would not have the "correct" alternate key if re-positioning were not done, I believe that this requires this re-positioning to take place, even though the NRP is not going to be updated by this particular operation.

Finally, the last thing I did was make some changes on how the record unlocking is done when buffer errors are encountered during a \$GET/\$FIND. At this point the record has already been locked, and must be unlocked before control returns to the user. The routine GET\_RECORD returns information in AP to RMSGET3B as to whether any special action is required to unlock this record on buffer errors. Unfortunately, AP is used throughout the remainder of RMSGET3B as input to record unpacking and key extraction; thus, its contents should a buffer error be detected and the record need to be unlocked, are unreliable. To fix this problem, I now set a flag bit on return from GET\_RECORD if in fact special action will be required to unlock the record on buffer errors, and reference this bit in that circumstance rather than the AP.

V03-016 TMK0010 Todd M. Katz 29-Sep-1982  
If a file is a prologue 3 file with alternate keys, and RMS is positioning by means of an alternate key of reference, then RMS was not unpacking the record before returning it to the user because it assumed that the record had been unpacked during the positioning and there was no need to unpack it a second time. However, while this is true, RMS at this point does not know the unpacked record's size. Thus, for the time being RMS must always unpack the record before moving it into the user's record buffer if the file is a prologue 3 file.

V03-015 TMK0009 Todd M. Katz 09-Sep-1982  
The field IRB\$B\_SRCHFLAGS is now a word in size. Change all references to it.  
  
Whenever RMS is positioning by means of an alternate key of reference (IRB\$B\_RP\_KREF > 0), then there is never a need in the local routine GET\_RECORD to extract the alternate key of the record positioned to into keybuffer 2. This is because as part of positioning to the primary data record from the SIDR in the first place, the SIDR key has already been extracted into keybuffer 2.

Eliminate all references to the routine RMSKEY\_TYPE\_CONV, since this routine doesn't do anything anyway.

The only time it is necessary to check for a valid packed decimal key is when the key type is packed decimal. It is never necessary to check for a valid packed decimal type when there is more than one segment and the file is a prologue 3 file. The packed decimal verification routine no longer requires parameters.

V03-014 KBT0294 Keith B. Thompson 23-Aug-1982  
Reorganize psects

V03-013 TMK0008 Todd M. Katz 10-Aug-1982  
At the present time, when the accessing of a record by RFA fails, the error returned by RMS\$FIND\_BY\_RRV is the error that gets reported to the user. Change this so that if this routine returns an error of RMS\$\_EOF (because the RFA VBN is greater than the VBN of any primary data bucket), this error gets mapped into an error of RMS\$\_PNF.

V03-012 MCN0013 Maria del C. Nasr 10-Aug-1982  
Check for less than 0 on call to RMS\$COMPARE\_KEY so that LIM check is done correctly. This is to fix bug introduced by MCN0012.

V03-011 TMK0007 Todd M. Katz 19-Jun-1982  
Implement the RMS cluster NRP solution. Basically this involves removal of the NRP cells from system space, and the maintenance of the next record positioning context locally within the IFAB. Changes required to the routines in this module are as follows:

1. The routine SETUP\_NRP\_DATA now sets up the current record context in the process local IRAB instead of in the system-wide NRP cell.
2. The IRAB variables IRBSL\_NEXT\_VBN and IRBSW\_NEXT\_ID are used to temporarily hold the RFA address of the "next" primary data record until the updating of the local NRP context takes place. This is because nothing in the local NRP context maybe modified, until everything is modified!
3. The local routines must also be modified both to make use of the next record positioning context now saved within the IRAB instead of within a systemwide NRP cell.
4. If RMS encounters the end-of-file set the IRBSV\_EOF bit. This bit is also cleared after successfully positioning randomly by key value. The former function of this bit has now been taken over by the new bit IRBSV\_CON\_EOF.
5. Special processing is required for \$GETs following random \$FINDs. A random \$FIND does not change the notion of what the next record is although it does change the notion of what the current record is! Example with the record sequence 0 A B - sequential \$GET to A, random \$FIND to 0, \$DELETE 0, followed by a sequential \$GET returns B, the next record. The random \$FIND changed the current record to 0, but did not change the next record to 0! The RMS cluster solution

286 0286 1  
287 0287 1  
288 0288 1  
289 0289 1  
290 0290 1  
291 0291 1  
292 0292 1  
293 0293 1  
294 0294 1  
295 0295 1  
296 0296 1  
297 0297 1  
298 0298 1  
299 0299 1  
300 0300 1  
301 0301 1  
302 0302 1  
303 0303 1  
304 0304 1  
305 0305 1  
306 0306 1  
307 0307 1  
308 0308 1  
309 0309 1  
310 0310 1  
311 0311 1  
312 0312 1  
313 0313 1  
314 0314 1  
315 0315 1  
316 0316 1  
317 0317 1  
318 0318 1  
319 0319 1  
320 0320 1  
321 0321 1  
322 0322 1  
323 0323 1  
324 0324 1  
325 0325 1  
326 0326 1  
327 0327 1  
328 0328 1  
329 0329 1  
330 0330 1  
331 0331 1  
332 0332 1  
333 0333 1  
334 0334 1  
335 0335 1  
336 0336 1  
337 0337 1  
338 0338 1  
339 0339 1  
340 0340 1  
341 0341 1  
342 0342 1



```

: 343 0343 1
: 344 0344 1
: 345 0345 1
: 346 0346 1
: 347 0347 1
: 348 0348 1
: 349 0349 1
: 350 0350 1
: 351 0351 1
: 352 0352 1
: 353 0353 1
: 354 0354 1
: 355 0355 1
: 356 0356 1
: 357 0357 1
: 358 0358 1
: 359 0359 1
: 360 0360 1
: 361 0361 1
: 362 0362 1
: 363 0363 1
: 364 0364 1
: 365 0365 1
: 366 0366 1
: 367 0367 1
: 368 0368 1
: 369 0369 1
: 370 0370 1
: 371 0371 1
: 372 0372 1
: 373 0373 1
: 374 0374 1
: 375 0375 1
: 376 0376 1
: 377 0377 1
: 378 0378 1
: 379 0379 1
: 380 0380 1
: 381 0381 1
: 382 0382 1
: 383 0383 1
: 384 0384 1
: 385 0385 1
: 386 0386 1
: 387 0387 1
: 388 0388 1
: 389 0389 1
: 390 0390 1
: 391 0391 1
: 392 0392 1
: 393 0393 1
: 394 0394 1
: 395 0395 1
: 396 0396 1
: 397 0397 1
: 398 0398 1
: 399 0399 1

```

for NRP positioning handles this by keeping the current primary data record's RFA and the RFA of the primary data record for NRP positioning in separate fields. Most operations set all NRP fields and as a result the RFA address of the current primary data record and the RFA address of the primary data record used for NRP positioning are the same. However, a random \$FIND will set only the current primary data record's RFA field. If the random \$FIND is immediately followed by a sequential \$GET, then it is only at that moment that the local NRP context is setup to return the randomly found record as the next record.

Also, it is no longer necessary within GET\_RECORD to loop on calls to RMSPOS\_SEQ or RMSPOS\_KEY when these routines return RLK errors. An RLK error could occur only when positioning on an alternate index and signalled that re-positioning should be forced. This re-positioning is now handled at a much lower level, and there is no longer any need to force it.

During the performance optimization of TMK0005 one incorrect assumption was made: that no deleted records were encountered between the last record retrieved, whose key is in keybuffer 1, and the new record that has just been retrieved. If this is true, the optimization holds, but if it is not, we can not use the key of the last retrieved record to uncompress the key of the new record, because the compression of the key of the new record is based upon the intervening deleted records, and not the key of the last record. In such a situation, the key of the new record must be extracted, and re-expanded in the old way performing a bucket scan if necessary.

Finally, it will no longer be necessary to unpack the primary data record when the file is a prologue 3 file, and RMS is currently positioning by an alternate key since the record will have been already unpacked and is within the internal record buffer.

V03-010 MCN0012 Maria del C. Nasr 29-Jun-1982  
Allow keys of different data types other than string.  
Change all CH\$COMPARE calls to RMSCOMPARE\_KEY to compare keys taking into consideration the different data types.

V03-009 TMK0006 Todd M. Katz 26-May-1982  
I have changed how the ROP=LIM key comparison is performed. Formerly, the routine RMSCOMPARE\_REC was being called. It was being called because the (incorrect) assumption existed that the key of the next record might have to be extracted and re-expanded, if key compression was enabled, in order to make the comparison. As it turns out, at this point in the operation, RMS has already extracted (and re-expanded if necessary) the key of the next record into keybuffer 2. Thus, in order to make the comparison, only a CH\$COMPARE between the key in the user's key buffer and the key in keybuffer 2 need be made. Thus, this comparison has now been made prologue independent, it is a performance optimization for all prologue versions, and the performance realized for prologue 3 files is considerable because it eliminates the need for one more









```
628 0628 1 : V01-005 C. D. Saether 21-Sep-1978 16:44  
629 0629 1 : Clear SRCFLAGS always.  
630 0630 1 :  
631 0631 1 : V01-004 W. Koenig 21-Sep-1978 15:50  
632 0632 1 : Return the data to the user on any seccess, not just "suc".  
633 0633 1 :  
634 0634 1 : V01-003 C. D. Saether 20-Sep-1978 16:25  
635 0635 1 : Clear NRP update flags when storing NRP.  
636 0636 1 :  
637 0637 1 : V01-002 C. D. Saether 12-Sep-1978 15:21  
638 0638 1 : Remove NXTBDB setup on RFA access.  
639 0639 1 :  
640 0640 1 : *****  
641 0641 1 :  
642 0642 1 LIBRARY 'RMSLIB:RMS';  
643 0643 1 :  
644 0644 1 REQUIRE 'RMSSRC:RMSIDXDEF';  
645 0709 1 :  
646 0710 1 : Define default psects for code  
647 0711 1 :  
648 0712 1 PSECT  
649 0713 1 CODE = RMSRMS3(PSECT_ATTR),  
650 0714 1 PLIT = RMSRMS3(PSECT_ATTR);  
651 0715 1 :  
652 0716 1 : Linkages.  
653 0717 1 :  
654 0718 1 LINKAGE  
655 0719 1 L_COMPARE_KEY,  
656 0720 1 L_JSB,  
657 0721 1 L_JSB01,  
658 0722 1 L_PRESERVE1,  
659 0723 1 L_QUERY_AND_LOCK,  
660 0724 1 L_RABREG,  
661 0725 1 L_RABREG_67,  
662 0726 1 L_RABREG_7,  
663 0727 1 L_REC_OVRD,  
664 0728 1 :  
665 0729 1 : Local Linkages.  
666 0730 1 :  
667 0731 1 L_GET_RECORD = JSB () :  
668 0732 1 GLOBAL (COMMON_RABREG, R_REC_ADDR, R_IDX_DFN)  
669 0733 1 NOPRESERVE (2, 3, 4, 5),  
670 0734 1 L_SETUP_NRP = JSB () :  
671 0735 1 GLOBAL (COMMON_RABREG, R_IDX_DFN)  
672 0736 1 NOPRESERVE (2, 3, 4, 5);  
673 0737 1 :  
674 0738 1 : Forward Routines.  
675 0739 1 :  
676 0740 1 FORWARD ROUTINE  
677 0741 1 GET_RECORD : L_GET_RECORD;  
678 0742 1 :  
679 0743 1 : External Routines.  
680 0744 1 :  
681 0745 1 EXTERNAL ROUTINE  
682 0746 1 RMSCOMPARE_KEY : RL$COMPARE_KEY,  
683 0747 1 RMSFIND_BY_RRV : RL$RABREG_67,  
684 0748 1 RMSKEY_DEST : RL$RABREG_7,
```

:	685	0749	1	RMSLOCK	:	RL\$QUERY_AND_LOCK,
:	686	0750	1	RMSNOREAD_LONG	:	RL\$JSB,
:	687	0751	1	RMSNOWRT_LONG	:	RL\$JSB,
:	688	0752	1	RMS\$PCKDEC_CHECK	:	RL\$RABREG_7,
:	689	0753	1	RMS\$POS_KEY	:	RL\$RABREG-67,
:	690	0754	1	RMS\$POS_RFA	:	RL\$RABREG-67,
:	691	0755	1	RMS\$POS_SEQ	:	RL\$RABREG-67,
:	692	0756	1	RMS\$QUERY_LCK	:	RL\$QUERY_AND_LOCK,
:	693	0757	1	RMS\$RECORD_IC	:	RL\$RABREG_67,
:	694	0758	1	RMS\$RECORD_KEY	:	RL\$PRESERVE1,
:	695	0759	1	RMS\$RECORD_VBN	:	RL\$PRESERVE1,
:	696	0760	1	RMS\$REC_OVHD	:	RL\$REC_OVHD,
:	697	0761	1	RMS\$RLSBKT	:	RL\$PRESERVE1,
:	698	0762	1	RMS\$RU_RECLAIM	:	RL\$RABREG_67 ADDRESSING_MODE( LONG_RELATIVE ),
:	699	0763	1	RMS\$UNLOCK	:	RL\$QUERY_AND_LOCK,
:	700	0764	1	RMS\$UNPACK_REC	:	RL\$JSB01;
:	701	0765	1			

SETUP\_NRP\_DATA

```

0766 1 %SBTTL 'SETUP_NRP_DATA'
0767 1 ROUTINE SETUP_NRP_DATA : L_SETUP_NRP NOVALUE =
0768 1
0769 1 |**
0770 1
0771 1 FUNCTIONAL DESCRIPTION:
0772 1
0773 1     This routine saves the next record positioning data
0774 1     in the IRAB from the temporary IRAB locations filled
0775 1     in during the positioning to the primary data record.
0776 1
0777 1 CALLING SEQUENCE:
0778 1
0779 1     SETUP_NRP_DATA()
0780 1
0781 1 INPUT PARAMETERS:
0782 1     NONE
0783 1
0784 1 IMPLICIT INPUTS:
0785 1
0786 1     IRAB
0787 1     IRBSL_FIRST_ID - Current SIDR's first SIDR array element ID
0788 1     IRBSL_FIRST_VBN - Current SIDR's first SIDR array element VBN
0789 1     IRBSL_KEYBUF - Pointer to keybuffers (to access keybuffer 2)
0790 1     IRBSW_NEXT_ID - ID of current primary data record
0791 1     IRBSL_NEXT_VBN - VBN of current primary data record
0792 1     IRBSW_RFA_ID - ID of current record (SIDR/primary)
0793 1     IRBSL_RFA_VBN - VBN of current record (SIDR/primary)
0794 1     IRBSB_RP_RREF - Key of reference used to retrieve user data record
0795 1     IRBSW_SAVE_POS - Number of elements before current SIDR element
0796 1
0797 1     IFAB
0798 1     IFBSW_KBUFSZ - Size of keybuffer (to access keybuffer 2)
0799 1
0800 1 OUTPUT PARAMETERS:
0801 1     NONE
0802 1
0803 1 IMPLICIT OUTPUTS:
0804 1     IRAB
0805 1     IRBSW_CUR_COUNT - Number of elements before current SIDR element
0806 1     IRBSW_CUR_ID - ID of current record (SIDR/primary)
0807 1     IRBSB_CUR_KREF - Key of reference of current record (SIDR/primary)
0808 1     IRBSL_CUR_VBN - VBN of current record (SIDR/primary)
0809 1     IRBSV_EOF - clear indicating stream is not at end-of-file
0810 1     IRBSL_KEYBUF - Pointer to keybuffers (to access keybuffer 1)
0811 1     IRBSW_POS_ID - ID of primary data record for NRP positioning
0812 1     IRBSL_POS_VBN - VBN of primary data record for NRP positioning
0813 1     IRBSL_SIDR_VBN - Current SIDR's first SIDR array element VBN
0814 1     IRBSW_SIDR_ID - Current SIDR's first SIDR array element ID
0815 1     IRBSW_UDR_ID - ID of current primary data record
0816 1     IRBSL_UDR_VBN - VBN of current primary data record
0817 1
0818 1 ROUTINE VALUE:
0819 1     NONE
0820 1
0821 1 --
0822 2 BEGIN

```



```

: 760 0823 2
: 761 0824 2
: 762 0825 2
: 763 0826 2
: 764 0827 2
: 765 0828 2
: 766 0829 2
: 767 0830 2
: 768 0831 2
: 769 0832 2
: 770 0833 2
: 771 0834 2
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: 777 0840 2
: 778 0841 2
: 779 0842 2
: 780 0843 2
: 781 0844 2
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: 783 0846 2
: 784 0847 2
: 785 0848 2
: 786 0849 2
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: 796 0859 2
: 797 0860 2
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: 800 0863 2
: 801 0864 2
: 802 0865 2
: 803 0866 2
: 804 0867 2
: 805 0868 2
: 806 0869 2
: 807 0870 2
: 808 0871 2
: 809 0872 2
: 810 0873 1

EXTERNAL REGISTER
R_IDX DFN STR,
COMMON_RAB_STR;

! Indicate that this stream is no longer at the file's end of file.
IRAB[IRBSV_EOF] = 0;

! Move the VBN of the current record into the appropriate IRAB location
IRAB[IRBSL_CUR_VBN] = .IRAB[IRBSL_RFA_VBN];

! If the current record happens to also be the primary data record, then
! move its ID into the appropriate IRAB location.
IF .IRAB[IRBSB_RP_KREF] EQLU 0
THEN
    IRAB[IRBSW_CUR_ID] = .IRAB[IRBSW_RFA_ID]

! If the current record happens to be a SIDR, then it has no ID to save,
! and instead save the SIDR first array element's VBN and ID (this uniquely
! identifies the SIDR), and the number of array elements preceding the
! current element (which points to the primary data record that is being
! retrieved).
ELSE
    BEGIN
        IRAB[IRBSW_CUR_COUNT] = .IRAB[IRBSW_SAVE_POS];
        IRAB[IRBSL_SIDR_VBN] = .IRAB[IRBSL_FIRST_VBN];
        IRAB[IRBSW_SIDR_ID] = .IRAB[IRBSW_FIRST_ID];
    END;

! Move the RFA of the current primary data record from its temporary
! location into the local NRP context and make it both the current primary
! data record and the primary data record for NRP positioning.
IRAB[IRBSL_UDR_VBN] = .IRAB[IRBSL_NEXT_VBN];
IRAB[IRBSW_UDR_ID] = .IRAB[IRBSW_NEXT_ID];

IRAB[IRBSL_POS_VBN] = .IRAB[IRBSL_NEXT_VBN];
IRAB[IRBSW_POS_ID] = .IRAB[IRBSW_NEXT_ID];

! Setup up the key of reference of the current record, and move the key of
! the current record into keybuffer 1.
IRAB[IRBSB_CUR_KREF] = .IRAB[IRBSB_RP_KREF];
CH$MOVE(.IFAB[IFBSW_KBUFSZ], KEYBUF_ADDR(2), KEYBUF_ADDR(1));
RETURN;

END;
```

```

.TITLE RM3GET
.FINNT \V04-000\
.EXTRN RMSCOMPARE_KEY, RMSFIND_BY_RRV
```

```

.EXTRN RMSKEY_DESC, RMSLOCK
.EXTRN RMSNOREAD_LONG, RMSNOWRT_LONG
.EXTRN RMSPOCKDEC-CHECK
.EXTRN RMSPOS_KEY, RMSPOS_RFA
.EXTRN RMSPOS_SEQ, RMSQUERY_LCK
.EXTRN RMSRECORD_ID, RMSRECORD_KEY
.EXTRN RMSRECORD_VBN, RMSREC_OVHD
.EXTRN RMSRLSBKT, RMSRU_RECLAIM
.EXTRN RMSUNLOCK, RMSUNPACK_REC

.PSECT RMSRMS3, NOWRT, GBL, PIC, 2

```

```

04 A9 02 8A 00000 SETUP_NRP_DATA:
00A8 C9 70 A9 D0 00004 BITCB2 #2, 4(IRAB) : 0830
00C2 C9 95 0000A MOVL 112(IRAB), 168(IRAB) : 0834
08 12 0000E TSTB 194(IRAB) : 0839
00B8 C9 74 A9 B0 00010 BNEQ 1$ : 0841
13 11 00016 BRB 2$ : 0841
00C0 C9 76 A9 B0 00018 1$: MOVW 118(IRAB), 192(IRAB) : 0851
00B4 C9 7C A9 D0 0001E MOVL 124(IRAB), 180(IRAB) : 0852
00BE C9 0082 C9 B0 00024 MOVW 130(IRAB), 190(IRAB) : 0853
00B0 C9 78 A9 D0 0002B 2$: MOVL 120(IRAB), 176(IRAB) : 0860
00BC C9 0080 C9 B0 00031 MOVW 128(IRAB), 188(IRAB) : 0861
00AC C9 78 A9 D0 00038 MOVL 120(IRAB), 172(IRAB) : 0863
00BA C9 0080 C9 B0 0003E MOVW 128(IRAB), 186(IRAB) : 0864
00C3 C9 00C2 C9 90 00045 MOVB 194(IRAB), 195(IRAB) : 0869
50 00B4 CA 3C 0004C MOVZWL 180(IFAB), R0 : 0870
50 60 A9 C0 00051 ADDL2 96(IRAB), R0 : 0870
60 B9 60 00B4 CA 28 00055 MOVCL3 180(IFAB), (R0), @96(IRAB) : 0873
05 0005C RSB : 0873

```

; Routine Size: 93 bytes, Routine Base: RMSRMS3 + 0000

; 811 0874 1

```

813 0875 1 %SBTTL 'RM$GET3B'
814 0876 1 GLOBAL ROUTINE RM$GET3B : RL$RABREG =
815 0877 1
816 0878 1 :++
817 0879 1
818 0880 1 :FUNCTIONAL DESCRIPTION:
819 0881 1
820 0882 1 :       This routine implements the get/find operation for the
821 0883 1 :       indexed file organization.
822 0884 1
823 0885 1 :CALLING SEQUENCE:
824 0886 1
825 0887 1 :       RM$GET3()
826 0888 1
827 0889 1 :INPUT PARAMETERS:
828 0890 1
829 0891 1 :       R11      Impure area pointer
830 0892 1 :       R10      IFAB -- Pointer to IFAB
831 0893 1 :       R9       IRAB -- Pointer to IRAB
832 0894 1 :       R8       RAB -- pointer to users RAB
833 0895 1 :       ROP field options (NLK,ULK,RLK,LOC,NXR)
834 0896 1 :       RAC field = (SEQ, or KEY, or RFA)
835 0897 1 :       RFA field if RAC = RFA
836 0898 1 :       KBF,KSZ,KRF if RAC = KEY and KBF,KSZ if RAC = SEQ and LIM set
837 0899 1 :       UBF,USZ -- if a GET
838 0900 1
839 0901 1 :IMPLICIT INPUTS:
840 0902 1
841 0903 1 :       IRAB fields:
842 0904 1
843 0905 1 :       IRBSV_UNLOCK_RP - current record should be unlocked before
844 0906 1 :       accessing new record.
845 0907 1 :       IRBSV_FIND_LAST - last operation was a FIND.
846 0908 1 :       IRBSV_SKIP_NEXT - last operation was sequential, the record
847 0909 1 :       described by the nrp info is to be skipped
848 0910 1 :       and the record beyond it becomes the new
849 0911 1 :       record.
850 0912 1 :       IRBSL_KEYBUF    (key buffer or 2, maybe 3)
851 0913 1
852 0914 1 :       IFAB fields:
853 0915 1
854 0916 1 :       IFBSB_PLG_VER
855 0917 1 :       IFBSV_RU_RLK    - if set, perform pseudo record locking
856 0918 1 :       IFBSV_RUP      - if set, Recovery Unit is in progress
857 0919 1 :       IFBSW_KBUFSZ   - size of each keybuffer
858 0920 1 :       IFBSV_NORECLCK - record locking not required, i.e., not
859 0921 1 :       sharing the file and single stream only.
860 0922 1 :       IFBSV_WRTACC   - if accessed for other than read only.
861 0923 1
862 0924 1 :OUTPUT PARAMETERS:
863 0925 1
864 0926 1 :       RAB fields:
865 0927 1
866 0928 1 :       RFA of record found
867 0929 1 :       STV if io errors
868 0930 1 :       RBF,RSZ -- if a GET
869 0931 1

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```
IMPLICIT OUTPUTS:
  IRAB fields:
    IRBSV_UNLOCK RP
    IRBSV_FIND_LAST
    IRBSV_SKIP_NEXT
    IRBSB_RP KREF
    IRBSB_CUR KREF
    IRBSL_KEYBUF (key buffer 1 or 2, maybe 3)
    IRBSL_RBF User buffer address and size
    IRBSW_RSZ

ROUTINE VALUE:
  Internal RMS status code

SIDE EFFECTS:
  Retrieved record maybe locked, and next record context is modified.
--
BEGIN
BUILTIN
  AP,
  TESTBITSC;
EXTERNAL REGISTER
  COMMON_RAB_STR;
GLOBAL REGISTER
  R_REC_ADDR_STR,
  R_IDX_DFN_STR;
LOCAL
  FLAGS : BLOCK[1],
  STATUS;
MACRO
  AP_STATUS = 0,0,1,0 %;
  OK_WAT_STATUS = 0,1,1,0 %;
  RU_DEL_STATUS = 0,2,1,0 %;

: Continue to attempt to get the next record under the following
: circumstances:
: 1. The status returned from GET_RECORD indicates the next record has
: been deleted and RMS's access mode is sequential or random by key
: value. Any key of reference. This can only happen if RMS has had
: to wait for a record lock to be granted (status returned will be
: RMSS_DEL), or RMS has positioned to and managed to lock a primary
: data record that is marked RU_DELETE (status returned will be 0).
: 2. The status returned from GET_RECORD is an alternate success status
```

```

927 0989 2
928 0990 2
929 0991 2
930 0992 2
931 0993 2
932 0994 2
933 0995 2
934 0996 2
935 0997 2
936 0998 2
937 0999 2
938 1000 2
939 1001 2
940 1002 2
941 1003 2
942 1004 2
943 1005 2
944 1006 2
945 1007 2
946 1008 2
947 1009 2
948 1010 2
949 1011 2
950 1012 2
951 1013 2
952 1014 2
953 1015 2
954 1016 2
955 1017 2
956 1018 2
957 1019 2
958 1020 2
959 1021 2
960 1022 6
961 1023 5
962 1024 5
963 1025 4
964 1026 4
965 1027 3
966 1028 4
967 1029 4
968 1030 4
969 1031 3
970 1032 3
971 1033 3
972 1034 3
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```

: (OK_WAT), and the key of reference is NOT the primary key.
:
: 3. An IDX_DFN has been allocated for the key descriptor, indicating
: there are no problems with the key of reference.
:
: Force the key descriptor to be initially zero.
:
: IDX_DFN = 0;
: FLAGS = 0;
:
: WHILE 1
: DO
: BEGIN
:
: STATUS = GET_RECORD (.FLAGS[OK_WAT_STATUS]
: OR
: TESTBITSC(FLAGS[RU_DEL_STATUS]));
:
: : If an unqualified success, avoid the contorted mass of
: : logic below and exit immediately...
:
: IF .STATUS<0,16> EQLU RMSSUC()
: THEN
: EXITLOOP;
:
: : Check the key descriptor after return from GET_RECORD.
: : If still zero, something wrong with the key of
: : reference, so exit loop.
:
: IF .IDX_DFN EQLU 0
: THEN
: EXITLOOP;
:
: IF NOT ((.STATUS<0,16> EQLU RMSERR(DEL)
: OR
: .STATUS EQLU 0)
: AND
: .RAB[RAB$B_RAC] NEQU RAB$C_RFA)
: THEN
: IF NOT (.IDX_DFN[IDX$B_KEYREF] NEQU 0
: AND
: .STATUS<0,16> EQLU RMSSUC(OK_WAT))
: THEN
: EXITLOOP;
:
: : Let us back off from the radical position above concerning the
: : ambitious attempts to continue to get a record. We should
: : never go back for a record if all of the following are
: : true:
: : - status from GET_RECORD is RMS$DEL
: : - primary key of reference
: : - random access (keyed OR RFA access)
: : - no dups on primary key
: : - exact key match
:
: : Under these circumstances, it is at least useless to
: : to back after a record, and sometimes downright WRONG!

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!
IF (.STATUS<0,16> EQLU RMSERR(DEL)
    AND
    .IDX_DFN[IDX$B_KEYREF] EQLU 0
    AND
    .RAB[RAB$B_RAC] NEQU RAB$C_SEQ
    AND
    NOT .IDX_DFN[IDX$V_DUPKEYS]
    AND
    NOT (.RAB[RAB$V_KGE] OR .RAB[RAB$V_KGT]))
THEN
    BEGIN
    STATUS = RMSERR(RNF);
    EXITLOOP;
    END;

! If RMS has to perform a re-positioning, then either it had to have
! waited for a record lock, or it positioned to a RU_DELETE marked
! primary data record. Therefore, set the appropriate state bit either
! of which will cause the correct lock to be released during
! re-positioning, and so that the proper status will be returned if
! RMS is able to position to a record.
IF .STATUS NEQU 0
THEN
    FLAGS[OK_WAT_STATUS] = 1
ELSE
    FLAGS[RU_DEL_STATUS] = 1;

CHSMOVE (.IDX_DFN[IDX$B_KEYSZ], KEYBUF_ADDR(1), KEYBUF_ADDR(2));
END;

! NOTE: AP is 0 if no special action is needed to unlock the RP on errors
! produced due to user buffer/size errors. Otherwise it is 1. Save this
! status.
FLAGS[AP_STATUS] = .AP<0,1>;

! If RMS was successful at obtaining the next record, but at some earlier
! time was forced to do a re-positioning because it had to wait for a record
! lock, then change the status to an OK_WAT success. This can only happen
! when RMS is positioning by means of an alternate index.
IF .FLAGS[OK_WAT_STATUS]
    AND
    .STATUS<0,16> EQLU RMSSUC()
THEN
    STATUS = RMSSUC(OK_WAT);

IRAB[IRB$V_FIND_LAST] = 0;

! Obtain user buffer address and size for later probe.
IRAB [IRB$L_RBF] = .RAB [RAB$L_RBF];
IRAB [IRB$W_RSZ] = .RAB [RAB$W_RSZ];

```

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```

: If the user has set the RAB$V_LIM bit in the ROP field on a sequential
: $GET/$FIND RMS reports whether the specified key exceeds the key of the
: record found.
IF .STATUS
  AND
  (.RAB[RAB$B_RAC] EQL RAB$C_SEQ)
  AND
  .RAB[RAB$V_LIM]
THEN
  BEGIN
    LOCAL
      KBF_ADDR      : LONG,
      KEYSIZE;

    KEYSIZE = .RAB[RAB$B_KSZ];

    IF .KEYSIZE EQL 0
    THEN
      IF .IDX_DFN[IDX$B_DATATYPE] EQL IDX$C_STRING
      OR .IDX_DFN[IDX$B_SEGMENTS] GTR 1
      THEN
        STATUS = RMSERR(KSZ)
      ELSE
        KEYSIZE = .IDX_DFN[IDX$B_KEYSZ];

    BEGIN

    MAP
      KEYSIZE      : BYTE;

    IF .KEYSIZE GTRU .IDX_DFN[IDX$B_KEYSZ]
    THEN
      STATUS = RMSERR(KSZ)
    END;

    KBF_ADDR = .RAB [RAB$L_KBF];
    IFNORD(KEYSIZE, .KBF_ADDR, [RAB[IRB$B_MODE],
      STATUS = RMSERR(RBF));

    IF .STATUS
    THEN
      BEGIN
        AP = 3;          ! Contiguous key compare

        ! The key of the current record has been previously saved in
        ! keybuffer 2, and may now be used to determine whether the user
        ! specified key limit has been exceeded.

        IF RMSCOMPARE_KEY ( KEYBUF_ADDR(2), .KBF_ADDR, .KEYSIZE ) LSS 0
        THEN
          STATUS = RMSSUC(OK_LIM);
        END;
      END;

    END;

```

```

: 1098 1160 2
: 1099 1161 2
: 1100 1162 2
: 1101 1163 2
: 1102 1164 2
: 1103 1165 2
: 1104 1166 2
: 1105 1167 2
: 1106 1168 2
: 1107 1169 2
: 1108 1170 2
: 1109 1171 3
: 1110 1172 4
: 1111 1173 4
: 1112 1174 4
: 1113 1175 4
: 1114 1176 5
: 1115 1177 4
: 1116 1178 4
: 1117 1179 4
: 1118 1180 4
: 1119 1181 4
: 1120 1182 4
: 1121 1183 4
: 1122 1184 5
: 1123 1185 5
: 1124 1186 5
: 1125 1187 4
: 1126 1188 4
: 1127 1189 4
: 1128 1190 3
: 1129 1191 3
: 1130 1192 3
: 1131 1193 3
: 1132 1194 4
: 1133 1195 4
: 1134 1196 4
: 1135 1197 4
: 1136 1198 4
: 1137 1199 4
: 1138 1200 4
: 1139 1201 4
: 1140 1202 4
: 1141 1203 5
: 1142 1204 5
: 1143 1205 5
: 1144 1206 5
: 1145 1207 5
: 1146 1208 5
: 1147 1209 5
: 1148 1210 5
: 1149 1211 5
: 1150 1212 5
: 1151 1213 5
: 1152 1214 5
: 1153 1215 5
: 1154 1216 5

```

```

IF .STATUS
THEN
BEGIN
  IF .IRAB[IRBSV_FIND]
  THEN
    ! This is a find operation don't
    ! return record etc.
    BEGIN
      IRAB[IRBSV_FIND_LAST] = 1;
      ! Set up the next record context for non-random $FIND operations.
      IF (.RAB[RAB$B_RAC] EQL RAB$C_SEQ)
      THEN
        SETUP_NRP_DATA()
        ! If this is a random $FIND operation then save the RFA of the
        ! found primary data record as the current primary data record.
      ELSE
        BEGIN
          IRAB[IRB$L_UDR_VBN] = .IRAB[IRB$L_NEXT_VBN];
          IRAB[IRB$W_UDR_ID] = .IRAB[IRB$W_NEXT_ID];
        END;
      END
    ELSE
      END
  ELSE
    ! Return the user the data on the record.
    BEGIN
      LOCAL
        RSZ      : WORD;
      RMSKEY_DESC(0);
      ! Add record overhead, and calculate record's size.
      BEGIN
        LOCAL
          REC_SIZE,
          RECRD_OVHD;
        RECORD_OVHD = RMSREC_OVHD(0; REC_SIZE);
        ! If this primary data record had been updated within a Recovery
        ! Unit then retrieve its true size from the last two bytes of the
        ! reserved space.
      IF .REC_ADDR[IRCSV_RU_UPDATE]
      THEN

```

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: 1100 1162 2
: 1101 1163 2
: 1102 1164 2
: 1103 1165 2
: 1104 1166 2
: 1105 1167 2
: 1106 1168 2
: 1107 1169 2
: 1108 1170 2
: 1109 1171 3
: 1110 1172 4
: 1111 1173 4
: 1112 1174 4
: 1113 1175 4
: 1114 1176 5
: 1115 1177 4
: 1116 1178 4
: 1117 1179 4
: 1118 1180 4
: 1119 1181 4
: 1120 1182 4
: 1121 1183 4
: 1122 1184 5
: 1123 1185 5
: 1124 1186 5
: 1125 1187 4
: 1126 1188 4
: 1127 1189 4
: 1128 1190 3
: 1129 1191 3
: 1130 1192 3
: 1131 1193 3
: 1132 1194 4
: 1133 1195 4
: 1134 1196 4
: 1135 1197 4
: 1136 1198 4
: 1137 1199 4
: 1138 1200 4
: 1139 1201 4
: 1140 1202 4
: 1141 1203 5
: 1142 1204 5
: 1143 1205 5
: 1144 1206 5
: 1145 1207 5
: 1146 1208 5
: 1147 1209 5
: 1148 1210 5
: 1149 1211 5
: 1150 1212 5
: 1151 1213 5
: 1152 1214 5
: 1153 1215 5
: 1154 1216 5

```









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```

: these same fields would contain the RFA of the current record from
when it had been originally locked.
IF TESTBITSC(IRAB[IRBSV_UNLOCK_RP])
THEN
  RMSUNLOCK (.IRAB[IRBSL_NEXT_VBN], .IRAB[IRBSW_NEXT_ID]);
: If end-of-file has been reached, set the corresponding IRAB bit.
IF .STATUS EQL RMSERR(EOF)
THEN
  IRAB[IRBSV_EOF] = 1;
: There is no longer a current primary data record.
IRAB[IRBSL_UDR_VBN] = 0;
IRAB[IRBSW_UDR_ID] = 0;
END;
RAB[RAB$R_RFA0] = .IRAB[IRBSL_UDR_VBN];
RAB[RAB$R_RFA4] = .IRAB[IRBSW_UDR_ID];
IRAB[IRBSV_FIND] = 0;
BEGIN
  GLOBAL REGISTER
  R_BDB;
  IF (BDB = .IRAB[IRBSL_CURBDB]) NEQ 0
  THEN
    RMSRLSBKT(0);
  IRAB[IRBSL_CURBDB] = 0;
END;
RETURN .STATUS
END;

```

```

00FC 8F BB 00000 RM$GET3B::
      SE           10 C2 00004  PUSH R      #M<R2,R3,R4,R5,R6,R7>          : 0876
      04           57 D4 00007  SUBL2     #16, SP                               : 0996
      AE D4 00009  CLRL      IDX,DFN                             : 0997
      50 D4 0000C 1$:    CLRL      R0                               : 1005
      02 04 AE     02 E5 0000E  BBCC     #2, FLAGS, 2$
      50 D6 00013  INCL      R0
      51 EF 00015  2$:    EXTZV     #1, #1, FLAGS, R1
      0000V 51 C9 0001B  BISL3     R1, R0, -(SP)
      04 C0 00022  BSBW     GET_RECORD                               : 1004
      5E           04 C0 00022  ADDL2     #4, SP
      6E           50 D0 00025  MOVL      R0, STATUS
      01           6E B1 00028  CMPW     STATUS, #1                               : 1010
      6C 13 0002B  BEQL     10$

```



		05	1B	000E9		BLEQU	15\$					
	6E	85A4	8F	3C	000EB	MOVZWL	#34212, STATUS				1138	
	53	30	AB	DO	000F0	15\$:	MOV	48(RAB), KBF_ADDR			1141	
	50	0A	A9	OC	000F4	PROBER	10(IRAB), KEYSIZE, (KBF_ADDR)				1143	
			05	12	000F9	BNEQ	16\$					
	6E	858C	8F	3C	000FB	MOVZWL	#34188, STATUS					
	18		6E	E9	00100	16\$:	BLBC	STATUS, 17\$			1145	
	5C		03	DO	00103	MOVL	#3, AP				1148	
	51	00B4	CA	3C	00106	MOVZWL	180(IFAB), R1				1154	
	51	60	A9	CO	0010B	ADDL2	96(IRAB), R1					
			0000G	30	0010F	BSBW	RMSCOMPARE_KEY					
			50	D5	00112	TSTL	R0					
			05	18	00114	BGEQ	17\$					
	6E	8051	8F	3C	00116	MOVZWL	#32849, STATUS				1156	
	03		6E	E8	0011B	17\$:	BLBS	STATUS, 18\$			1161	
			010A	31	0011E	BRW	38\$					
1D	05	A9	01	E1	00121	18\$:	BBC	#1, 5(IRAB), 21\$			1165	
	04	A9	20	88	00126	BISB2	#32, 4(IRAB)				1172	
		1E	AB	95	0012A	TSTB	30(RAB)				1176	
			05	12	0012D	BNEQ	19\$					
			FE71	30	0012F	BSBW	SETUP_NRP_DATA				1178	
			0D	11	00132	BRB	20\$					
	00B0	C9	78	A9	DO	00134	19\$:	MOVL	120(IRAB), 176(IRAB)		1185	
	00BC	C9	00B0	C9	B0	0013A	MOVW	128(IRAB), 188(IRAB)			1186	
				72	11	00141	20\$:	BRB	29\$		1165	
				7E	D4	00143	21\$:	CLRL	-(SP)		1199	
			0000G	30	00145	BSBW	RMSKEY_DESC					
		5E		04	C0	00148	ADDL2	#4, SP				
				51	D4	0014B	CLRL	R1			1209	
			0000G	30	0014D	BSBW	RMSREC_OVHD					
0E		66		06	E1	00150	BBC	#6, (REC_ADDR), 22\$			1215	
52		56		50	C1	00154	ADDL3	RECORD_OVHD, REC_ADDR, R2			1217	
			FE	A142	9F	00158	PUSHAB	-2(REC_SIZE)[R2]				
	08	AE		9E	B0	0015C	MOVW	@(SP)+, RSZ				
				04	11	00160	BRB	23\$				
	08	AE		51	B0	00162	22\$:	MOVW	REC_SIZE, RSZ		1221	
		56		50	C0	00166	23\$:	ADDL2	RECORD_OVHD, REC_ADDR		1223	
		03	00B7	CA	91	00169	CMPB	183(IFAB), #3			1229	
				1C	12	0016E	BNEQ	26\$				
			00C2	C9	95	00170	TSTB	194(IRAB)			1240	
				05	12	00174	BNEQ	24\$				
		5C		02	DO	00176	MOVL	#2, AP			1242	
				02	11	00179	BRB	25\$				
				5C	D4	0017B	24\$:	CLRL	AP		1244	
	51	08	AE	3C	0017D	25\$:	MOVZWL	RSZ, R1			1246	
	50	68	A9	DO	00181	MOVL	104(IRAB), R0					
			0000G	30	00185	BSBW	RMSUNPACK_REC					
	08	AE		50	B0	00188	MOVW	R0, RSZ				
		27	06	AB	E9	0018C	26\$:	BLBC	6(RAB), 30\$		1249	
22	22	AA		03	E0	00190	BBS	#3, 34(IFAB), 30\$				
		50	20	A9	DO	00195	MOVL	32(IRAB), R0			1251	
19	0A	A0		04	E0	00199	BBS	#4, 10(R0), 30\$				
		03	00B7	CA	91	0019E	CMPB	183(IFAB), #3			1254	
				06	1E	001A3	BGEQU	27\$				
	28	A8		56	DO	001A5	MOVL	REC_ADDR, 40(RAB)			1256	
				05	11	001A9	BRB	28\$				
	28	A8	68	A9	DO	001AB	27\$:	MOVL	104(IRAB), 40(RAB)		1258	

22	A8	08	AE	B0	001B0	28\$:	MOVW	RSZ, 34(RAB)	1259
			71	11	001B5	29\$:	BRB	37\$	1249
	50	20	A8	B0	001B7	30\$:	MOVW	32(RAB), USZ	1267
			0F	12	001BB		BNEQ	32\$	1269
	04	04	AE	E9	001BD		BLBC	FLAGS, 31\$	1273
05	A9		20	88	001C1		BISB2	#32, 5(IRAB)	1275
	6E	86F4	8F	3C	001C5	31\$:	MOVZWL	#34548, STATUS	1277
			5C	11	001CA		BRB	37\$	1269
0C	AE	24	A8	D0	001CC	32\$:	MOVL	36(RAB), UBF_ADDR	1284
	50	08	AE	B1	001D1		CMPW	RSZ, USZ	1286
			0E	1B	001D5		BLEQU	33\$	
0C	A8	08	AE	3C	001D7		MOVZWL	RSZ, 12(RAB)	1289
08	AE		50	B0	001DC		MOVW	USZ, RSZ	1290
	6E	81A8	8F	3C	001E0		MOVZWL	#33192, STATUS	1291
	7E	0A	A9	9A	001E5	33\$:	MOVZBL	10(IRAB), -(SP)	1294
		10	AE	DD	001E9		PUSHL	UBF_ADDR	
	7E	10	AE	3C	001EC		MOVZWL	RSZ, -(SP)	
			0000G	30	001F0		BSBW	RMSNOWRT_LONG	
	5E		0C	C0	001F3		ADDL2	#12, SP	
	0F		50	E9	001F6		BLBC	R0, 35\$	
	04	04	AE	E9	001F9		BLBC	FLAGS, 34\$	1298
05	A9		20	88	001FD		BISB2	#32, 5(IRAB)	1300
	6E	86EC	8F	3C	00201	34\$:	MOVZWL	#34540, STATUS	1302
			20	11	00206		BRB	37\$	1294
22	A8	08	AE	B0	00208	35\$:	MOVW	RSZ, 34(RAB)	1306
28	A8	0C	AE	D0	0020D		MOVL	UBF_ADDR, 40(RAB)	1307
	03	00B7	CA	91	00212		CMPB	183(IFAB), #3	1308
			08	1E	00217		BGEQU	36\$	
OC	BE	66	AE	28	00219		MOV3	RSZ, (REC_ADDR), @UBF_ADDR	1310
			07	11	0021F		BRB	37\$	
OC	BE	68	AE	28	00221	36\$:	MOV3	RSZ, @104(IRAB), @UBF_ADDR	1312
		07	6E	E8	00228	37\$:	BLBS	STATUS, 39\$	1324
	81A8	8F	6E	B1	0022B	38\$:	CMPW	STATUS, #33192	1326
			4C	12	00230		BNEQ	43\$	
	6D	05	01	E0	00232	39\$:	BBS	#1, 5(IRAB), 46\$	1330
			FD69	30	00237		BSBW	SETUP_NRP_DATA	1336
	66	1C	A7	E8	0023A		BLBS	28(IDX_DF), 46\$	1343
	52		56	D0	0023E		MOVL	REC_ADDR, TMP_REC_ADDR	1350
	56	28	A8	D0	00241		MOVL	40(RAB), REC_ADDR	1351
	7E	0A	A9	9A	00245		MOVZBL	10(IRAB), -(SP)	1354
			56	DD	00249		PUSHL	REC_ADDR	1353
	7E	22	A8	3C	0024B		MOVZWL	34(RAB), -(SP)	
			0000G	30	0024F		BSBW	RMSNOREAD_LONG	
	5E		0C	C0	00252		ADDL2	#12, SP	
	05		50	E9	00255		BLBC	R0, 40\$	
	6E	8654	8F	3C	00258		MOVZWL	#34388, STATUS	1356
	5C		03	D0	0025D	40\$:	MOVL	#3, AP	1357
	07		6E	E8	00260		BLBS	STATUS, 41\$	1358
	81A8	8F	6E	B1	00263		CMPW	STATUS, #33192	1360
			0F	12	00268		BNEQ	42\$	
	50	00B4	CA	3C	0026A	41\$:	MOVZWL	180(IFAB), R0	1367
		60	B940	3F	0026F		PUSHAW	@96(IRAB)[R0]	
			0000G	30	00273		BSBW	RMSRECORD_KEY	
	5E		04	C0	00276		ADDL2	#4, SP	
	56		52	D0	00279	42\$:	MOVL	TMP_REC_ADDR, REC_ADDR	1369
			26	11	0027C		BRB	46\$	1324
OC	04	A9	0D	E5	0027E	43\$:	BBCC	#13, 4(IRAB), 44\$	1391





```

: 1365      1426 1 %SBTTL 'GET_RECORD'
: 1366      1427 1 ROUTINE GET_RECORD (REPOS_STATUS) : L_GET_RECORD =
: 1367      1428 1
: 1368      1429 1 :++
: 1369      1430 1
: 1370      1431 1 FUNCTIONAL DESCRIPTION:
: 1371      1432 1
: 1372      1433 1     This routine implements the actual retrieval of the
: 1373      1434 1     data record for internal RMS usage. The use request
: 1374      1435 1     is checked for valid input parameters and all internal state
: 1375      1436 1     information is setup to retrieve the record.
: 1376      1437 1     Then current record is unlocked if required and the
: 1377      1438 1     requested record retrieved and locked if required. All NRP
: 1378      1439 1     update data is saved in the IRAB but is not placed in the
: 1379      1440 1     NRP fields of the IRAB.
: 1380      1441 1
: 1381      1442 1 CALLING SEQUENCE:
: 1382      1443 1
: 1383      1444 1     GET_RECORD()
: 1384      1445 1
: 1385      1446 1 INPUT PARAMETERS:
: 1386      1447 1
: 1387      1448 1     REPOS_STATUS    - if 1, then RMS is performing a re-positioning.
: 1388      1449 1
: 1389      1450 1 IMPLICIT INPUTS:
: 1390      1451 1
: 1391      1452 1     Same as for RMSGET3 or RMSFIND3
: 1392      1453 1
: 1393      1454 1 OUTPUT PARAMETERS:
: 1394      1455 1
: 1395      1456 1     IRAB context setup for retrieved record:
: 1396      1457 1
: 1397      1458 1         CURBDB,RFA_VBN,RFA_ID,SAVE_POS,FIRST_VBN
: 1398      1459 1         FIRST_ID,REC_ADDR,NEXT_VBN,NEXT_ID
: 1399      1460 1
: 1400      1461 1     RAB The DCT field is cleared in all cases.
: 1401      1462 1
: 1402      1463 1     If the value of the routine is a success status then
: 1403      1464 1     the AP = 0 if no special action is needed to unlock the RP
: 1404      1465 1     and is 1 if special action is needed, on errors detected after
: 1405      1466 1     this routine.
: 1406      1467 1
: 1407      1468 1 IMPLICIT OUTPUTS:
: 1408      1469 1
: 1409      1470 1     IRBSV_UNLOCK_RP
: 1410      1471 1     IRBSV_FIND_LAST = 0
: 1411      1472 1     IRBSB_RP_KREF
: 1412      1473 1
: 1413      1474 1 ROUTINE VALUE:
: 1414      1475 1
: 1415      1476 1     Internal RMS status code
: 1416      1477 1
: 1417      1478 1 SIDE EFFECTS:
: 1418      1479 1
: 1419      1480 1     Retrieved record maybe locked.
: 1420      1481 1     Old current record may have been unlocked.
: 1421      1482 1     The data bucket for the retrieved record is accessed.

```

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```

NLK is ignored if the process is in a Recovery Unit.
If a primary data record is found to have been deleted within a Recovery
Unit it might be deleted for good.
If a primary data record is found to have been updated within a Recovery
Unit it might be re-formatted.

--

BEGIN

BUILTIN
  TESTBITCC,
  TESTBITSC,
  AP;

EXTERNAL REGISTER
  R_REC_ADDR_STR,
  R_IDX_DFN_STR,
  COMMON_RAB_STR;

LABEL
  UNLOCK,
  KEY;

IRAB[IRB$B_CACHEFLGS] = 0;
IRAB[IRB$W_SRCHFLAGS] = 0;
IRAB[IRB$V_DUP] = 0;

! Based on the record access mode (RAC) of this operation (GET/FIND)
! set up the IRAB RP fields, key buffer 2 etc to retrieve the record.

RAB[RAB$L_DCT] = 0;

! Get record block 1 --- set up the IRAB search context data to get the
! record the user is requesting, and unlock the current record if this is
! required.

UNLOCK :
BEGIN

CASE RAB[RAB$B_RAC] FROM RAB$C_SEQ TO RAB$C_RFA OF
SET
  [RAB$C_SEQ] :
    +
    Sequential Access:

    Setup to retrieve the record associated with the NRP if this
    is a GET and the last operation was not a FIND or if this is
    a FIND.

    If last operation was a FIND and this operation is a GET then
    retrieve record which is described by the NRP if that FIND was
    sequential.

```

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: 1479 1540 3
: 1480 1541 3
: 1481 1542 3
: 1482 1543 3
: 1483 1544 3
: 1484 1545 3
: 1485 1546 3
: 1486 1547 4
: 1487 1548 4
: 1488 1549 4
: 1489 1550 4
: 1490 1551 4
: 1491 1552 4
: 1492 1553 5
: 1493 1554 5
: 1494 1555 5
: 1495 1556 4
: 1496 1557 5
: 1497 1558 5
: 1498 1559 5
: 1499 1560 5
: 1500 1561 5
: 1501 1562 5
: 1502 1563 5
: 1503 1564 5
: 1504 1565 5
: 1505 1566 5
: 1506 1567 5
: 1507 1568 5
: 1508 1569 5
: 1509 1570 5
: 1510 1571 5
: 1511 1572 6
: 1512 1573 5
: 1513 1574 5
: 1514 1575 5
: 1515 1576 5
: 1516 1577 5
: 1517 1578 5
: 1518 1579 5
: 1519 1580 5
: 1520 1581 5
: 1521 1582 5
: 1522 1583 6
: 1523 1584 6
: 1524 1585 6
: 1525 1586 5
: 1526 1587 5
: 1527 1588 5
: 1528 1589 5
: 1529 1590 5
: 1530 1591 5
: 1531 1592 5
: 1532 1593 5
: 1533 1594 5
: 1534 1595 5
: 1535 1596 5

```

```

: If the FIND was random then change the NRP data for the record
: which was found and retrieve it.

: Note that a sequential FIND following a random FIND returns to
: the sequential next record (NRP). That is to say that the
: random operation will not change the NRP VBN and ID fields.

BEGIN
IF (.IRAB[IRBSV_CON_EOF ])
THEN
RETURN RMSERR(EOF);

IF (.IRAB[IRBSV_FIND_LAST]
AND
NOT (.IRAB[IRBSV_FIND]))
THEN
BEGIN
: NOTE: keybuffer 2 contains key value, RP_KREF has key of
: reference, RP_VBN and RP_ID contains record's RFA/RRV, and
: SAVE_DUP contains the duplicate position count.
IF TESTBITSC(IRAB[IRBSV_SKIP_NEXT])
THEN
: Last find was sequential so retrieve the record described
: by the NRP unless RMS is already positioned at the end
: of the file.
IF .IRAB[IRBSV_EOF]
THEN
RETURN .RMSERR (EOF)
ELSE
RMSKEY_DESC(.IRAB[IRBSB_CUR_KREF])
ELSE
: Last operation was a find random and this operation is a
: get sequential. Setup the local NRP context so that this
: same record will be retrieved. Since there has been no
: intervening operation, calling SETUP_NRP_DATA will
: accomplish this.
BEGIN
SETUP_NRP_DATA();
RMSKEY_DESC(.IRAB[IRBSB_RP_KREF]);
END;

IRAB[IRBSB_KEYSZ] = .IDX_DFN[IDX$B_KEYSZ];

: Unless no lock is desired on this record (and the process
: is not within a Recovery Unit), leave this block to avoid
: unlocking the current record. This avoids a potential
: window where the record is unlocked as it is reaccessed on
: this get operation.
IF NOT .RAB[RAB$V_NLK]

```

```

: 1536      1597      5
: 1537      1598      5
: 1538      1599      5
: 1539      1600      5
: 1540      1601      5
: 1541      1602      5
: 1542      1603      4
: 1543      1604      5
: 1544      1605      5
: 1545      1606      5
: 1546      1607      5
: 1547      1608      5
: 1548      1609      5
: 1549      1610      5
: 1550      1611      5
: 1551      1612      5
: 1552      1613      5
: 1553      1614      5
: 1554      1615      5
: 1555      1616      5
: 1556      1617      5
: 1557      1618      5
: 1558      1619      5
: 1559      1620      5
: 1560      1621      5
: 1561      1622      5
: 1562      1623      5
: 1563      1624      5
: 1564      1625      5
: 1565      1626      4
: 1566      1627      4
: 1567      1628      3
: 1568      1629      3
: 1569      1630      3
: 1570      1631      3
: 1571      1632      4
: 1572      1633      4
: 1573      1634      4
: 1574      1635      4
: 1575      1636      4
: 1576      1637      4
: 1577      1638      4
: 1578      1639      4
: 1579      1640      4
: 1580      1641      4
: 1581      1642      4
: 1582      1643      4
: 1583      1644      4
: 1584      1645      4
: 1585      1646      4
: 1586      1647      4
: 1587      1648      4
: 1588      1649      4
: 1589      1650      5
: 1590      1651      5
: 1591      1652      7
: 1592      1653      6

```

```

      OR
      .IFAB[IFBSV_RUP]
    THEN
      LEAVE UNLOCK;
    END
  ELSE
    BEGIN
      ! Return immediately if already at end-of-file.
      IF .IRAB[IRBSV_EOF]
    THEN
      RETURN RMSERR(EOF);

      IRAB[IRBSV_SKIP_NEXT] = 1;

      ! First time call after $CONNECT or $REWIND. Then what we
      ! want to retrieve is the very first record, so don't skip
      ! next record.
      IF .IRAB[IRBSL_CUR_VBN] EQL 0
    THEN
      IRAB[IRBSV_SKIP_NEXT] = 0;

      RETURN_ON_ERROR (RMSKEY_DESC(.IRAB[IRBSB_CUR_KREF]));

      IRAB[IRBSB_KEYSZ] = .IDX_DFN[IDXSB_KEYSZ];
      CH$MOVE(.IDX_DFN[IDXSB_KEYSZ], KEYBUF_ADDR(1), KEYBUF_ADDR(2));
      END;

    END;

  [RAB$C_KEY] :
  KEY :
    BEGIN
      LOCAL
        KEYSIZE : BYTE,
        KBF_ADDR : LONG;

      IRAB[IRBSV_CON_EOF] = 0;
      IRAB[IRBSV_SKIP_NEXT] = 0;

      RETURN_ON_ERROR (RMSKEY_DESC(.RAB[RAB$B_KRF]));
      KEYSIZE = .RAB[RAB$B_KSZ];

      ! Check and setup for user key value.
      IF .IDX_DFN[IDXSB_DATATYPE] EQL IDX$C_STRING
    OR .IDX_DFN[IDXSB_SEGMENTS] GTR 1
    THEN
      BEGIN
        IF ((.KEYSIZE EQL 0)
          OR

```



```

: 1650 1711 4
: 1651 1712 4
: 1652 1713 4
: 1653 1714 4
: 1654 1715 4
: 1655 1716 4
: 1656 1717 4
: 1657 1718 4
: 1658 1719 4
: 1659 1720 4
: 1660 1721 4
: 1661 1722 4
: 1662 1723 4
: 1663 1724 4
: 1664 1725 4
: 1665 1726 4
: 1666 1727 4
: 1667 1728 4
: 1668 1729 4
: 1669 1730 4
: 1670 1731 4
: 1671 1732 4
: 1672 1733 4
: 1673 1734 4
: 1674 1735 4
: 1675 1736 4
: 1676 1737 4
: 1677 1738 4
: 1678 1739 4
: 1679 1740 4
: 1680 1741 5
: 1681 1742 5
: 1682 1743 5
: 1683 1744 5
: 1684 1745 5
: 1685 1746 5
: 1686 1747 5
: 1687 1748 5
: 1688 1749 5
: 1689 1750 5
: 1690 1751 5
: 1691 1752 4
: 1692 1753 5
: 1693 1754 5
: 1694 1755 5
: 1695 1756 5
: 1696 1757 5
: 1697 1758 5
: 1698 1759 5
: 1699 1760 5
: 1700 1761 5
: 1701 1762 5
: 1702 1763 5
: 1703 1764 5
: 1704 1765 5
: 1705 1766 5
: 1706 1767 5

```

```

*
At this point we have determined that this a random access for an
exact match by key. Now try to find out if this is for the
current record, i.e., is it the same one we just got. This will
be checked only for primary key. The following conditions must be
met to take this optimization:

Previous operation was a GET.
This operation is for primary key.
Duplicates aren't allowed on primary key.
The full key size is being used.
The key value matches the saved primary key value of the
current record (in keybuffer 3).
There is a current record (rp_vbn neq 0).
The current record is already locked, if locking required.
The new record is to be locked, if locking required.

IF .IRAB[IRBSV_FIND_LAST]
OR
.IDX_DFN[IDX$B_KEYREF] NEQ 0
OR
.IDX_DFN[IDX$V_DUPKEYS]
THEN
LEAVE KEY;

IF .IRAB[IRBSB_KEYSZ] NEQ .IDX_DFN[IDX$B_KEYSZ]
THEN
LEAVE KEY;

BEGIN
LOCAL
SIZE;

SIZE = .IRAB[IRBSB_KEYSZ];

IF NOT CH$EQL(.SIZE, KEYBUF_ADDR(2), .SIZE, KEYBUF_ADDR(3))
THEN
LEAVE KEY;

END;
BEGIN
LOCAL
VBN;

IF (VBN = .IRAB[IRBSL_UDR_VBN]) EQL 0
THEN
LEAVE KEY;

! If record locking is required, make sure this record is already
! locked, otherwise it may be deleted or locked by another
! accessor. Also that the new record is to be locked also,
! otherwise there is potentially an obscure window where it could
! be deleted while reaccessing the bucket after the current lock is
! released.

```

1707	1768	5
1708	1769	5
1709	1770	5
1710	1771	6
1711	1772	6
1712	1773	6
1713	1774	5
1714	1775	6
1715	1776	5
1716	1777	6
1717	1778	6
1718	1779	6
1719	1780	6
1720	1781	6
1721	1782	6
1722	1783	6
1723	1784	6
1724	1785	6
1725	1786	6
1726	1787	6
1727	1788	6
1728	1789	6
1729	1790	6
1730	1791	6
1731	1792	6
1732	1793	6
1733	1794	6
1734	1795	6
1735	1796	6
1736	1797	6
1737	1798	6
1738	1799	7
1739	1800	6
1740	1801	6
1741	1802	6
1742	1803	5
1743	1804	5
1744	1805	4
1745	1806	4
1746	1807	4
1747	1808	4
1748	1809	4
1749	1810	4
1750	1811	4
1751	1812	4
1752	1813	4
1753	1814	4
1754	1815	4
1755	1816	4
1756	1817	4
1757	1818	4
1758	1819	4
1759	1820	4
1760	1821	4
1761	1822	4
1762	1823	4
1763	1824	4

```

:
IF NOT .IFAB[IFBSV_NORECLK]
  OR
  (.IFAB[IFBSV_RU_RLK]
   AND
   .IFAB[IFBSV_RUP])
THEN
  BEGIN
    LOCAL
      ST;

    IF .RAB[RABSV_NLK]
      AND
      NOT .IFAB[IFBSV_RUP]
    THEN
      LEAVE KEY;

    ! If the user has requested waiting for record locking,
    ! disable such waiting at this time. RMS does not want to
    ! wait because it is just interested in whether this stream
    ! has locked the record or there is a window. RMS does not
    ! want to wait for it here if it has to. The lock logic later
    ! on will do the waiting if it is necessary.
    IF .RAB[RABSV_WAT]
    THEN
      IRAB[IRBSV_NO_Q_WAIT] = 1;

    ST = RMSQUERY_LCK(.VBN, .IRAB[IRBSV_UDR_ID]);

    IF .ST<0, 16> NEQU RMSSUC(OK_ALK)
    THEN
      LEAVE KEY

  END;
END;

! If we are here we have determined that this is the same record
! that we already have locked as the current record. Flag that
! state by setting DUP. By leaving the unlock block, the current
! record is not unlocked.

      >>>====> NOTE <====<<<

Setting DUP will cause us to fall into an optimization below,
which will change our access to sequential in order to quickly
reaccess the record that we have been waiting on (and should
now have locked). There are some problems with this way of
thinking, especially when the following are true:

- we just had this record, and released it (which means
  that our record context is STILL INTACT)
- the record we were waiting on was deleted after we
  released it
- keyed access

```

1764 1825 4  
1765 1826 4  
1766 1827 4  
1767 1828 4  
1768 1829 4  
1769 1830 4  
1770 1831 4  
1771 1832 4  
1772 1833 4  
1773 1834 4  
1774 1835 4  
1775 1836 4  
1776 1837 4  
1777 1838 3  
1778 1839 3  
1779 1840 3  
1780 1841 3  
1781 1842 3  
1782 1843 3  
1783 1844 3  
1784 1845 4  
1785 1846 4  
1786 1847 4  
1787 1848 4  
1788 1849 4  
1789 1850 4  
1790 1851 4  
1791 1852 4  
1792 1853 4  
1793 1854 4  
1794 1855 4  
1795 1856 4  
1796 1857 4  
1797 1858 4  
1798 1859 3  
1799 1860 3  
1800 1861 3  
1801 1862 3  
1802 1863 3  
1803 1864 3  
1804 1865 3  
1805 1866 3  
1806 1867 3  
1807 1868 3  
1808 1869 3  
1809 1870 3  
1810 1871 3  
1811 1872 3  
1812 1873 3  
1813 1874 3  
1814 1875 3  
1815 1876 3  
1816 1877 3  
1817 1878 3  
1818 1879 3  
1819 1880 3  
1820 1881 3

```

- exact match by primary key
- no dups

The above conditions will give us the first non-deleted record,
regardless of key (since we temporarily used sequential access).
And since KEYBUFFER 3 still has the key from when we last had
this record (before we released it), it will match our search
key. I have attempted in RMSGET3B to prevent coming back here a
second time if these conditions are true.

IRAB[IRBSV_DUP] = 1;
LEAVE UNLOCK;

END;                                     ! of block KEY

[RAB$C_RFA] :
! RFA access -- check RFA for legality and setup for primary key
! access for next record pointer (NRP) data.
BEGIN
IRAB[IRBSV_CON_EOF] = 0;
IRAB[IRBSV_SKIP_NEXT] = 0;             ! flag random access
IF .RAB[RAB$R_RFA0] EQL 0
  OR
  .RAB[RAB$R_RFA4] EQL 0
THEN
  RETURN RMSERR(RFA);
RETURN_ON_ERROR (RMS$KEY_DESC(0));
END;

[OUTRANGE] :
RETURN RMSERR(RAC);

TES;
! The current record is now unlocked before accessing the new record,
! unless it has already been determined that the new record is the same as
! the old current record, in which case this block was left and this code
! is skipped.
IF TESTBITSC(IRAB[IRBSV_UNLOCK_RP])
THEN
! If RMS is performing a re-positioning then unlock the record
! positioned to during the previous positioning attempt; otherwise,
! unlock the current record (locked during the previous positioning
! operation) if there is one.
IF .REPOS_STATUS
THEN
  RMSUNLOCK (.IRAB[IRBSL_NEXT_VBN], .IRAB[IRBSW_NEXT_ID])

```



```

: 1821 1882 3
: 1822 1883 3
: 1823 1884 3
: 1824 1885 3
: 1825 1886 3
: 1826 1887 2
: 1827 1888 2
: 1828 1889 2
: 1829 1890 2
: 1830 1891 2
: 1831 1892 2
: 1832 1893 3
: 1833 1894 3
: 1834 1895 3
: 1835 1896 3
: 1836 1897 3
: 1837 1898 4
: 1838 1899 5
: 1839 1900 5
: 1840 1901 5
: 1841 1902 5
: 1842 1903 5
: 1843 1904 5
: 1844 1905 5
: 1845 1906 5
: 1846 1907 5
: 1847 1908 6
: 1848 1909 6
: 1849 1910 6
: 1850 1911 6
: 1851 1912 6
: 1852 1913 6
: 1853 1914 6
: 1854 1915 6
: 1855 1916 6
: 1856 1917 6
: 1857 1918 6
: 1858 1919 6
: 1859 1920 6
: 1860 1921 6
: 1861 1922 6
: 1862 1923 6
: 1863 1924 5
: 1864 1925 5
: 1865 1926 5
: 1866 1927 5
: 1867 1928 5
: 1868 1929 5
: 1869 1930 5
: 1870 1931 5
: 1871 1932 5
: 1872 1933 5
: 1873 1934 5
: 1874 1935 5
: 1875 1936 5
: 1876 1937 5
: 1877 1938 5

```

```

ELSE
  IF .IRAB[IRB$$_UDR_VBN] NEQU 0
  THEN
    RMSUNLOCK (.IRAB[IRB$$_UDR_VBN], .IRAB[IRB$$_UDR_ID]);
END;
IRAB[IRB$$_RP_KREF] = .IDX_DFN[IDX$$_KEYREF];
! of block UNLOCK
! Get record block 2 -- position and perform lock logic for record which
! the IRAB search context data describes.
BEGIN
LOCAL
  STATUS;
IF NOT (STATUS =
  BEGIN
  LOCAL
    RAC : BYTE;
  RAC = .RAB[RAB$$_RAC];
  IF .IRAB[IRB$$_DUP]
  THEN
    ! re-accessing current record
    BEGIN
      ! If next record info also for primary key, then use sequential
      ! positioning code and nrp info - it's faster. irb$$_skip_next
      ! will be clear in this case so that record itself is retrieved.
      ASSUME_C(RAB$$_SEQ, 0);
      ! sneaky way to set rac = rab$$_seq
      RAC = .IRAB[IRB$$_CUR_KREF];
      IF .RAC NEQ 0
      THEN
        ! when cur_kref = 0.
        RAC = RAB$$_RFA;
      END;
    CASE .RAC FROM RAB$$_SEQ TO RAB$$_RFA OF
    SET
      ! Sequential access.
      [RAB$$_SEQ] : STATUS = RMSPOS_SEQ();
      ! Random access by key.
      [RAB$$_KEY] : STATUS = RMSPOS_KEY();
      ! Random access by RFA.

```

```

1878
1879
1880
1881
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1885
1886
1887
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1889
1890
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1894
1895
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1900
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1934

```

```

[ RAB$C_RFA ] : STATUS = RMS$POS_RFA();
TES
END)
THEN
RETURN .STATUS;

! Setup record pointer (RP) to the RFA/RRV of retrieved record.
!
IRAB[IRB$W_NEXT_ID] = RMS$RECORD_ID();
AP = 3;
BEGIN
GLOBAL REGISTER
R_BDB;

IRAB[IRB$L_NEXT_VBN] = RMS$RECORD_VBN();
END;

! Move the key of the internally current record into keybuffer 2 in
! preparation for updating the local NRP context. It will not be necessary
! to extract the key, if positioning was done by means of an alternate key
! index, because as part of that positioning, the key would have been moved
! into keybuffer 2.
IF .IRAB[IRB$B_RP_KREF] EQLU 0
THEN
! If any of the following conditions holds, the key of the internally
! current record (which must be a primary data record) must be
! extracted from the record itself.
!
! 1. If the file is a prologue 1 or 2 file.
! 2. If the record was retrieved randomly.
! 3. If key compression is not enabled in this key of reference.
! 4. If deleted records were encountered during the positioning.
IF ((.IFAB[IFB$B_PLG_VER] LSSU PLG$C_VER_3)
OR
(.RAB[RAB$B_RAC] NEQU RAB$C_SEQ)
OR
NOT .IDX_DFN[IDX$V_KEY_COMPR]
OR
.IRAB[IRB$V_DEL_SEEN])
THEN
BEGIN
GLOBAL REGISTER
R_BDB;

AP = 0;
RMS$RECORD_KEY (KEYBUF_ADDR(2));
END

```

```

: 1935
: 1936
: 1937
: 1938
: 1939
: 1940
: 1941
: 1942
: 1943
: 1944
: 1945
: 1946
: 1947
: 1948
: 1949
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: 1951
: 1952
: 1953
: 1954
: 1955
: 1956
: 1957
: 1958
: 1959
: 1960
: 1961
: 1962
: 1963
: 1964
: 1965
: 1966
: 1967
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: 1970
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: 1978
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: 1984
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: 1988
: 1989
: 1990
: 1991

```

```

1996 4
1997 4
1998 4
1999 4
2000 4
2001 4
2002 3
2003 4
2004 4
2005 4
2006 4
2007 4
2008 4
2009 4
2010 4
2011 4
2012 4
2013 4
2014 4
2015 4
2016 4
2017 4
2018 3
2019 3
2020 3
2021 3
2022 3
2023 3
2024 3
2025 4
2026 4
2027 4
2028 3
2029 3
2030 3
2031 2
2032 2
2033 3
2034 3
2035 3
2036 3
2037 3
2038 4
2039 4
2040 4
2041 4
2042 4
2043 4
2044 4
2045 4
2046 4
2047 4
2048 4
2049 4
2050 4
2051 4
2052 4

```

```

: If this is a prologue 3 file, and RMS is positioning sequentially,
: then RMS may use the key of the last retrieved record, stored in
: keybuffer 1 to supply any characters front compressed off the current
: key provided no intervening records were encountered between the
: last retrieved, and the new record.
ELSE
  BEGIN
    LOCAL
      KEY      : REF BBLOCK;

    MACRO
      KEY_LEN = 0,0,8,0 %;
      CMP_CNT = 1,0,8,0 %;

    KEY = .REC_ADDR + RMSREC_OVHD(0);

    CH$COPY (.KEY[CMP_CNT], KEYBUF_ADDR(1),
             .KEY[KEY_LEN], .KEY+2,
             (.KEY + .KEY[KEY_LEN] + 1),
             .IDX_DFNC[IDX$B_KEYSZ], KEYBUF_ADDR(2));

    END;

: Don't do any record-locking if there aren't any writers of the file, and
: if pseudo record locking is not to be done.
IF .IFAB[IFB$V_NORECLK]
  AND
  NOT (.IFAB[IFB$V_RU_RLK]
       AND
       .IFAB[IFB$V_RUP])
THEN
  RETURN .STATUS;

END;                                ! of block defining local STATUS

BEGIN
  LOCAL
    STATUS;

  BEGIN
    LABEL
      OK_WAT;

: Flag no special action needed for unlocking the RP
AP = 0;

: only query_lock the record if:
: 1. The user has specified no locking ( NLK ) and is not in a Recovery
:    Unit.

```

```

1992 2053 4      ! 2. The file is opened for read ( not WRTACC )
1993 2054 4      !       and the user has not specified read-only locking ( not REA )
1994 2055 4      !
1995 2056 5      IF (.RAB[RAB$V_NLK]
1996 2057 5          AND
1997 2058 5          NOT .IFAB[IFB$V_RUP])
1998 2059 5
1999 2060 4      OR
2000 2061 6      ( (NOT .IFAB[IFB$V_WRTACC])
2001 2062 5          AND
2002 2063 5          (NOT .RAB[RAB$V_REA]) )
2003 2064 4      THEN
2004 2065 4          STATUS = RMSQUERY_LCK (.IRAB[IRB$S_NEXT_VBN], .IRAB[IRB$W_NEXT_ID])
2005 2066 4      ELSE
2006 2067 4          STATUS = RMSLOCK (.IRAB[IRB$S_NEXT_VBN], .IRAB[IRB$W_NEXT_ID]);
2007 2068 4
2008 2069 4
2009 2070 4      ! OK_WAT success status means we had to wait for someone else to unlock the
2010 2071 4      ! record. To wait, we deaccessed the bucket. Therefore, we must reaccess
2011 2072 4      ! it, and we can use the record pointer information for this. Deaccessing
2012 2073 4      ! the bucket also means that our NRP context updating information in the
2013 2074 4      ! IRAB cannot longer be considered to be valid.
2014 2075 4
2015 2076 5      IF .STATUS EQL RMSSUC(OK_WAT)
2016 2077 4      THEN
2017 2078 4      OK_WAT:
2018 2079 5          BEGIN
2019 2080 5
2020 2081 5          LOCAL
2021 2082 5          TEMP_STATUS;
2022 2083 5
2023 2084 5      ! Reposition to the record using record pointer contents. If it is
2024 2085 5      ! possible that some reclamation maybe done make sure the primary data
2025 2086 5      ! bucket is exclusively accessed.
2026 2087 5
2027 2088 5      IF .IFAB[IFB$V_WRTACC]
2028 2089 5          AND
2029 2090 5          .IFAB[IFB$V_RU]
2030 2091 5      THEN
2031 2092 5          IRAB[IRB$B_CACHEFLGS] = CSH$M_LOCK;
2032 2093 5
2033 2094 6      IF NOT (TEMP_STATUS = RMSFIND_BY_RRV (.IRAB[IRB$S_NEXT_VBN],
2034 2095 6          .IRAB[IRB$W_NEXT_ID],
2035 2096 6          0))
2036 2097 5      THEN
2037 2098 5          STATUS = .TEMP_STATUS;
2038 2099 5
2039 2100 5      ! If RMS after re-positioning to the record finds that it had been
2040 2101 5      ! deleted within a Recovery Unit, then RMS will have to re-position
2041 2102 5      ! after deleting the record for good if it has write access to the
2042 2103 5      ! file.
2043 2104 5
2044 2105 5      IF .STATUS
2045 2106 5          AND
2046 2107 5          .REC_ADDR[IRC$V_RU_DELETE]
2047 2108 5      THEN
2048 2109 5          LEAVE OK_WAT;

```

```

2049 2110 5
2050 2111 5 IF .STATUS
2051 2112 5 THEN
2052 2113 5
2053 2114 5 : If our key of reference is the primary key, then we can reclaim
2054 2115 5 : our NRP updating information from the primary data bucket's
2055 2116 5 : VBN and the record ID.
2056 2117 5
2057 2118 6 IF (.IDX_DFN[IDX$B_KEYREF] EQL 0)
2058 2119 5 THEN
2059 2120 6 BEGIN
2060 2121 6 IRAB[IRB$RFA_VBN] = .BBLOCK[IRAB[IRB$L_CURBDB], BDB$L_VBN];
2061 2122 6 IRAB[IRB$RFA_ID] = IRC$_ID(REC_ADDR);
2062 2123 6 END
2063 2124 6
2064 2125 6 : If the key of reference is not the primary key, then RMS has no
2065 2126 6 : easy way to reclaim the NRP list updating information which is
2066 2127 6 : for the SIDR bucket (long since released) and not the primary
2067 2128 6 : data bucket. Since the stream which has the record locked might
2068 2129 6 : delete the SIDR array positioned to but not the primary data
2069 2130 6 : record itself (by means of an $UPDATE), RMS must re-position
2070 2131 6 : in order to guarantee that the key of the SIDR array it
2071 2132 6 : positions to is actually represented in the primary data record
2072 2133 6 : to be returned. The alternate success status, and the fact that
2073 2134 6 : the key of reference is other than the primary will force another
2074 2135 6 : attempt to access the primary data bucket after accessing the
2075 2136 6 : necessary SIDR, and to lock the next record.
2076 2137 6
2077 2138 5 ELSE
2078 2139 6 BEGIN
2079 2140 6
2080 2141 6 GLOBAL REGISTER
2081 2142 6 R_BDB_STR;
2082 2143 6
2083 2144 6 RELEASE (IRAB[IRB$L_CURBDB]);
2084 2145 6 END
2085 2146 5 ELSE
2086 2147 5 IRAB[IRB$L_CURBDB] = 0;
2087 2148 5
2088 2149 4 END;
2089 2150 4
2090 2151 4 : If RMS finds that the current record has been modified within a Recovery
2091 2152 4 : Unit, then subject it to further processing before deciding whether to
2092 2153 4 : return it as the non-deleted primary data record to be returned to the
2093 2154 4 : user, or whether to return a status to force RMS to re-position.
2094 2155 4
2095 2156 4 IF .STATUS
2096 2157 4 AND
2097 2158 4 .IRAB[IRB$L_CURBDB] NEQU 0
2098 2159 4 AND
2099 2160 5 (.REC_ADDR[IRCSV_RU_DELETE]
2100 2161 5 OR
2101 2162 5 .REC_ADDR[IRCSV_RU_UPDATE])
2102 2163 4 THEN
2103 2164 5 BEGIN
2104 2165 5
2105 2166 5 LOCAL

```

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: 2106
: 2107
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2167 5
2168 5
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2170 5
2171 5
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2173 5
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2197 6
2198 6
2199 6
2200 7
2201 6
2202 7
2203 6
2204 6
2205 5
2206 4
2207 4
2208 4
2209 4
2210 5
2211 5
2212 5
2213 4
2214 5
2215 5
2216 5
2217 4
2218 4
2219 4
2220 3
2221 3
2222 3
2223 3

```

```

RECORD_ID : WORD;
RECORD_ID = .REC_ADDR[IRCSW_ID];
! If the file has been open for write access, then attempt to delete
! the record if it was deleted within a Recovery Unit, or re-format
! the record if it was updated within a Recovery Unit.
IF .IFAB[IFBSV_WRTACC]
THEN
RMSRU_RECLAIM();
! If the record had been deleted within a Recovery Unit, then RMS will
! not return this record to the user as a non-deleted primary data
! record. Therefore, release the primary data bucket, and change the
! return status to 0 if RMS did not have to wait for the record lock, or
! change the return status to RMS$DEL if RMS had to wait for the
! record lock. Returning a status of RMS$DEL in the latter case will
! allow the information that RMS had to wait for a record lock to
! eventually be returned to the user along with the non-deleted primary
! data record when such a record is eventually found.
IF .RECORD_ID NEQU .REC_ADDR[IRCSW_ID]
OR
.REC_ADDR[IRCSV_RU_DELETE]
THEN
BEGIN
GLOBAL REGISTER
R_BDB_STR;
RELEASE (IRAB[IRBSL_CURBDB]);
IF .STATUS<0,16> EQLU RMSSUC(OK_WAT)
THEN
STATUS = RMSERR(DEL)
ELSE
STATUS = 0;
END;
END;
! Return here if QUERY_LCK.
IF (.RAB[RABSV_NLK]
AND
NOT .IFAB[IFBSV_RUP])
OR
(NOT .IFAB[IFBSV_WRTACC]
AND
NOT .RAB[RABSV_REA])
THEN
RETURN .STATUS;
END;
! If UNLOCK_RP was set coming here, it can only mean that this was a
! reaccessing of a previously automatically locked record that was not

```

```

: 2163      2224      3      ! unlocked at the beginning of this operation to avoid a locking window.
: 2164      2225      3      ! It will get an ok_alk status (not suc) from rm$lock. It wants to release
: 2165      2226      3      ! the current record on potential buffer errors. The case where we don't
: 2166      2227      3      ! want to release the now current record lock is if the status from rm$lock
: 2167      2228      3      ! was rms$ok_alk (i.e., not suc) which meant that it had been previously
: 2168      2229      3      ! manually locked, and should remain that way even if this operation fails.
: 2169      2230      3
: 2170      2231      3      IF TESTBITC(IRAB[IRBSV_UNLOCK_RP])
: 2171      2232      3      THEN
: 2172      2233      4          IF .STATUS<0,16> EQLU RMSSUC()
: 2173      2234      5              OR
: 2174      2235      6              .STATUS<0,16> EQLU RMSSUC(OK_WAT)
: 2175      2236      7              OR
: 2176      2237      8              .STATUS<0,16> EQLU RMSSUC(OK_RULK)
: 2177      2238      9          THEN
: 2178      2239      9              AP = 1
: 2179      2240      9          ELSE
: 2180      2241      9
: 2181      2242      9              ! If it is necessary for us to release the record lock set
: 2182      2243      9              ! IRBSV_UNLOCK_RP. It will only be necessary to release the record
: 2183      2244      9              ! lock in some circumstances when we have had to stall waiting for
: 2184      2245      9              ! it, and whenever RMS has managed to position to a record that was
: 2185      2246      9              ! deleted within a Recovery Unit. In the former case if after
: 2186      2247      9              ! waiting for the lock, we have some problem reaccessing the bucket,
: 2187      2248      9              ! or we find that the record is deleted while we were waiting, then
: 2188      2249      9              ! we must release the record lock. If we are positioning by means
: 2189      2250      9              ! of an alternate key and we have had to stall waiting for the
: 2190      2251      9              ! record lock, and this is an operation where the NRP list must be
: 2191      2252      9              ! updated (any operation but a nonrandom $FIND), then the record
: 2192      2253      9              ! lock must also be released.
: 2193      2254      9
: 2194      2255      4          BEGIN
: 2195      2256      4
: 2196      2257      5              IF (.IRAB[IRBSL_CURBDB] EQL 0)
: 2197      2258      6                  AND
: 2198      2259      7                  (.STATUS<0,16> NEQ RMSERR(RLK))
: 2199      2260      8              THEN
: 2200      2261      8                  IRAB[IRBSV_UNLOCK_RP] = 1;
: 2201      2262      8
: 2202      2263      8              RETURN .STATUS;
: 2203      2264      8              END
: 2204      2265      3          ELSE
: 2205      2266      3              AP = 1;
: 2206      2267      3
: 2207      2268      3          IF NOT .RAB[RABSV_ULK]
: 2208      2269      3          THEN
: 2209      2270      3              IRAB[IRBSV_UNLOCK_RP] = 1;
: 2210      2271      3
: 2211      2272      3          RETURN .STATUS;
: 2212      2273      3
: 2213      2274      3          END
: 2214      2275      1          END;

```

```

! of local block defining STATUS
! of routine

```

			SE	18	C2	00000	GET_RECORD:						
				40	A9	94	00003	SUBL2	#24, SP	1427			
				42	A9	B4	00006	CLRFB	64(IRAB)	1507			
				04	A9	9E	00009	CLRWB	66(IRAB)	1508			
	14	AE		1000	8F	AA	0000E	MOVAB	4(IRAB), 20(SP)	1509			
	14	BE		38	A8	D4	00014	BICW2	#4096, @20(SP)				
02		00		1E	A8	8F	00017	CLRL	56(RAB)	1515			
01A3		0097		000D			0001C	CASEB	30(RAB), #0, #2	1524			
								.WORD	2\$-1\$ -				
									12\$-1\$ -				
									29\$-1\$ -				
			50	8644	8F	3C	00022	MOVZWL	#34372, R0	1862			
					4C	11	00027	BRB	8\$				
42	14	BE			17	E0	00029	BBS	#23, @20(SP), 7\$	1549			
38	14	BE			05	E1	0002E	BBC	#5, @20(SP), 6\$	1553			
33	14	BE			09	E0	00033	BBS	#9, @20(SP), 6\$	1555			
0C	14	BE			08	E5	00038	BBCC	#11, @20(SP), 3\$	1563			
2E	14	BE			01	E0	0003D	BBS	#1, @20(SP), 7\$	1570			
					00C3	C9	9A	00042	MOVZBL	195(IRAB), -(SP)	1574		
						08	11	00047	BRB	4\$			
					FC87	30	00049	RSBW	SETUP NRP_DATA	1584			
			7E	00C2	C9	9A	0004C	MOVZBL	194(IRAB), -(SP)	1585			
					0000G	30	00051	BSBW	RMSKEY_DE\$C				
			SE		04	C0	00054	ADDL2	#4, SP				
	00A6	C9		20	A7	90	00057	MOVB	32(IDX_DFN), 166(IRAB)	1588			
06	06	A8			04	E1	0005D	BBC	#4, 6(RAB), 5\$	1596			
48	00A2	CA			02	E1	00062	BBC	#2, 162(IFAB), 11\$	1598			
					01A0	31	00068	BRW	36\$	1600			
08	14	BE			01	E1	0006B	BBC	#1, @20(SP), 9\$	1608			
					50	827A	8F	3C	00070	MOVZWL	#33402, R0	1610	
						03B6	31	00075	BRW	71\$			
					14	BE	0800	8F	A8	00078	9\$: BISW2	#2048, @20(SP)	1612
						00A8	9	D5	0007E	TSTL	168(IRAB)	1618	
							06	12	00082	BNEQ	10\$		
					14	BE	0800	8F	AA	00084	BICW2	#2048, @20(SP)	1620
						7E	00C3	C9	9A	0008A	10\$: MOVZBL	195(IRAB), -(SP)	1622
							0000G	30	0008F	BSBW	RMSKEY_DE\$C		
			SE		04	C0	00092	ADDL2	#4, SP				
			DD		50	E9	00095	BLBC	STATUS, 8\$				
	00A6	C9		20	A7	90	00098	MOVB	32(IDX_DFN), 166(IRAB)	1624			
		51		20	A7	9A	0009E	MOVZBL	32(IDX_DFN), R1	1625			
		50		00B4	CA	3C	000A2	MOVZWL	180(IFAB), R0				
		50		60	A9	C0	000A7	ADDL2	96(IRAB), R0				
	60	60	B9		51	28	000AB	MOVCS	R1, @96(IRAB), (R0)				
					0135	31	000B0	BRW	33\$	1524			
14	BE	01			00	F0	000B3	INSV	#0, #23, #1, @20(SP)	1638			
					14	BE	0800	8F	AA	000B9	12\$: BICW2	#2048, @20(SP)	1640
						7E	35	A8	9A	000BF	MOVZBL	53(RAB), -(SP)	1642
							00U0G	30	000C3	BSBW	RMSKEY_DE\$C		
			SE		04	C0	000C6	ADDL2	#4, SP				
			A9		50	E9	000C9	BLBC	STATUS, 8\$				
			50		34	A8	90	000CC	MOVB	52(RAB), KEYSIZE	1643		
					1D	A7	95	000D0	TSTB	29(IDX_DFN)	1647		
						06	13	000D3	BEQL	13\$			
			01		1E	A7	91	000D5	CMPB	30(IDX_DFN), #1	1648		
						0C	1B	000D9	BLEQU	14\$			
						50	95	000DB	13\$: TSTB	KEYSIZE	1652		



			16	13	000DD	BEQL	16\$			
	20	A7	50	91	000DF	CMPB	KEYSIZE, 32(IDX_DFN)		1654	
			17	1B	000E3	BLEQU	17\$			
			0E	11	000E5	BRB	16\$		1656	
			50	95	000E7	14\$: TSTB	KEYSIZE		1662	
			04	12	000E9	BNEQ	15\$			
	20	50	20	A7	90	000EB	MOVB	32(IDX_DFN), KEYSIZE	1664	
		A7	50	91	000EF	15\$: CMPB	KEYSIZE, 32(IDX_DFN)		1666	
			07	13	000F3	BEQL	17\$			
		50	85A4	8F	3C	000F5	16\$: MOVZWL	#34212, R0	1668	
			44	11	000FA	BRB	20\$			
	00A6	C9	50	90	000FC	17\$: MOVB	KEYSIZE, 166(IRAB)		1672	
		52	30	A8	00	00101	MOVL	48(RAB), KBF_ADDR	1673	
62		50	0A	A9	0C	00105	PROBER	10(IRAB), KEYSIZE, (KBF_ADDR)	1675	
			07	12	0010A	BNEQ	18\$			
		50	858C	8F	3C	0010C	MOVZWL	#34188, R0		
			2D	11	00111	BRB	20\$			
		51	50	9A	00113	18\$: MOVZBL	KEYSIZE, R1		1679	
		50	00B4	CA	3C	00116	MOVZWL	180(IFAB), R0		
		50	60	A9	70	0011B	ADDL2	96(IRAB), R0		
60		62	51	28	0011F	MOV(C3	R1, (KBF_ADDR), (R0)			
		05	1D	A7	91	00123	CMPB	29(IDX_DFN), #5	1683	
			06	12	00127	BNEQ	19\$			
			0000G	30	00129	BSBW	RMS\$PCKDEC CHECK		1685	
		11	50	E9	0012C	BLBC	STATUS, 20\$			
12		54	04	A8	9E	0012F	19\$: MOVAB	4(RAB), R4	1689	
		64	15	E1	00133	BBC	#21, (R4), 22\$			
08		64	16	E1	00137	BBC	#22, (R4), 21\$		1692	
		50	867C	8F	3C	00138	MOVZWL	#34428, R0	1694	
			02EB	31	00140	20\$: BRW	71\$			
	42	A9	10	88	00143	21\$: BISB2	#16, 66(IRAB)		1697	
			08	11	00147	BRB	23\$		1698	
07		64	16	E1	00149	22\$: BBC	#22, (R4), 24\$		1703	
	42	A9	02	88	0014D	BISB2	#2, 66(IRAB)		1706	
			0094	31	00151	23\$: BRW	33\$		1707	
F8		04	A9	05	E0	00154	24\$: BBS	#5, 4(IRAB), 23\$	1729	
			21	A7	95	00159	TSTB	33(IDX_DFN)	1731	
			F3	12	0015C	BNEQ	23\$			
		EF	1C	A7	E8	0015E	BLBS	28(IDX_DFN), 23\$	1733	
	20	A7	00A6	C9	91	00162	CMPB	166(IRAB), 32(IDX_DFN)	1737	
			7E	12	00168	BNEQ	33\$			
		51	00A6	C9	9A	0016A	MOVZBL	166(IRAB), SIZE	1746	
		50	00B4	CA	3C	0016F	MOVZWL	180(IFAB), R0	1748	
			60	B940	3F	00174	PUSHAW	@96(IRAB)[R0]		
9E		60	B940	51	29	00178	CMPC3	SIZE, @96(IRAB)[R0], @ (SP)+		
			68	12	0017E	BNFQ	33\$			
		51	00B0	C9	00	00180	MOVL	176(IRAB), VBN	1758	
			61	13	00185	BEQL	33\$			
0C		06	AA	03	E1	00187	BRC	#3, 6(IFAB), 25\$	1769	
27		00A2	CA	03	E1	0018C	BBC	#3, 162(IFAB), 28\$	1771	
21		00A2	CA	02	E1	00192	BBC	#2, 162(IFAB), 28\$	1773	
06		64	14	E1	00198	25\$: BBC	#20, (R4), 26\$		1780	
46		00A2	CA	02	E1	0019C	BBC	#2, 162(IFAB), 33\$	1782	
04		64	11	E1	001A2	26\$: BBC	#17, (R4), 27\$		1793	
		07	A9	01	88	001A6	BISB2	#1, 7(IRAB)	1795	
			52	00BC	C9	3C	001AA	27\$: MOVZWL	188(IRAB), R2	1797
			0000C	30	001AF	BSBW	RMS\$QUERY_LCK			

	8039	8F		50	B1	001B2		CMPW	ST, #32825		1799	
				2F	12	001B7		BNEQ	33\$			
	05	A9		10	88	001B9	28\$:	BISB2	#16, 5(IRAB)		1835	
				4C	11	001BD		BRB	36\$		1836	
14	BE		01	00	F0	001BF	29\$:	INSV	#0, #23, #1, @20(SP)		1847	
	14	BE		0800	8F	AA	001C5	BICW2	#2048, @20(SP)		1849	
				10	A8	D5	001CB	TSTL	16(RAB)		1851	
					05	13	001CE	BEQL	30\$			
				14	A8	B5	001D0	TSTW	20(RAB)		1853	
					08	12	001D3	BNEQ	32\$			
				50	865C	8F	3C	001D5	30\$:	MOVZWL	#34396, R0	1855
					0251	31	001DA	31\$:	BRW	71\$		
					7E	D4	001DD	32\$:	CLRL	-(SP)		1857
					0000G	30	001DF		BSBW	RMSKEY_DESC		
		5E			04	C0	001E2	ADDL2	#4, SP			
		F2			50	E9	001E5	BLBC	STATUS, 31\$			
	1E	04			0D	E5	001E8	33\$:	BBCC	#13, 4(IRAB), 36\$		1871
		A9			0D	E9	001ED	BLBC	REPOS STATUS, 34\$			1879
		0B		1C	AE	E9	001F1	MOVZWL	128(IRAB), R2			1881
		52		0080	C9	3C	001F1	MOVL	120(IRAB), R1			
		51			78	A9	D0	001F6	BRB	35\$		
					0C	11	001FA	MOVL	176(IRAB), R1			1883
		51		00B0	C9	D0	001FC	34\$:	BEQL	36\$		
					08	13	00201	MOVZWL	188(IRAB), R2			1885
		52		00BC	C9	3C	00203	BSBW	RMSUNLOCK			
					0000G	30	00208	35\$:	MOVB	33(IDX DFN), 194(IRAB)		1888
	00C2	C9			21	A7	90	0020B	36\$:	MOVB	30(RAB), RAC	1904
		50			1E	A8	90	00211	BBC	#4, 5(IRAB), 37\$		1906
	0A	05			04	E1	00215	MOVB	195(IRAB), RAC			1918
		50		00C3	C9	90	0021A	BEQL	37\$			1920
					03	13	0021F	MOVB	#2, RAC			1922
		50			02	90	00221	CASEB	RAC, #0, #2			1926
	02	00			50	8F	00224	37\$:	.WORD	39\$-38\$,-		
	0010	000B			0006		00228	38\$:		40\$-38\$,-		
										41\$-38\$		
					0000G	30	0022E	39\$:	BSBW	RMSPOS_SEQ		1931
					08	11	00231	BRB	42\$			
					0000G	30	00233	40\$:	BSBW	RMSPOS_KEY		1935
					03	11	00236	BRB	42\$			
					0000G	30	00238	41\$:	BSBW	RMSPOS RFA		1939
		14	AE		50	D0	0023B	42\$:	MOVL	R0, STATUS		
		03			14	AE	E8	0023F	BLBS	STATUS, 43\$		1926
					00AE	31	00243	BRW	48\$			
					0000G	30	00246	43\$:	BSBW	RMSRECORD ID		1949
	0080	C9			50	B0	00249	MOVW	R0, 128(IRAB)			
		5C			03	D0	0024E	MOVL	#3, AP			1951
					0000G	30	00251	BSBW	RMSRECORD VBN			1958
		78	A9		50	D0	00254	MOVL	R0, 120(IRAB)			
				00C2	C9	95	00258	TSTB	194(IRAB)			1967
					27	12	0025C	BNEQ	45\$			
		03		00B7	CA	91	0025E	CMPB	183(IFAB), #3			1979
					0F	1F	00263	PLSU	44\$			
					1E	A8	95	00265	TSTB	30(RAB)		1981
					0A	12	00268	BNEQ	44\$			
	05	1C	A7		06	E1	0026A	BBC	#6, 28(IDX DFN), 44\$			1983
	13	43	A9		01	E1	0026F	BBC	#1, 67(IRAB), 46\$			1985
					5C	D4	00274	44\$:	CLRL	AP		1992

			50	00B4	CA	3C	00276		MOVZWL	180(IFAB), R0		1993
				60	B940	9F	0027B		PUSHAB	@96(IRAB)[R0]		
			5E		0000G	30	0027F		BSBW	RMSRECORD_KEY		
						04	00282		ADDL2	#4, SP		
						5C	00285	45\$:	BRB	47\$		1979
						51	00287	46\$:	CLRL	R1		2012
					0000G	30	00289		BSBW	RMSREC_OVHD		
6E			50			56	0028C		ADDL3	REC_ADDR, R0, KEY		
50			6E			01	00290		ADDL3	#1, KEY, R0		2014
	10		AE			60	00294		MOVZBL	(R0), 16(SP)		
	04		AE			00	00298		MOVZBL	@KEY, 4(SP)		2015
	0C		AE			20	0029D		MOVZBL	32(IDX DFN), 12(SP)		2017
			50	00B4	CA	3C	002A2		MOVZWL	180(IFAB), R0		
			08	AE		60	002A7		MOVAB	@96(IRAB)[R0], 8(SP)		
7E			04	AE		01	002AD		ADDL3	#1, 4(SP), -(SP)		
			6E	AE		04	002B2		ADDL2	KEY, (SP)		
OC	AE		9E	60	B9	14	002B6		MOVC5	20(SP), @96(IRAB), @ (SP)+, 12(SP), @8(SP)		
						08	002BE					
						21	002C0		BGEQ	47\$		
			08	AE		10	002C2		ADDL2	16(SP), 8(SP)		
			0C	AE		10	002C7		SUBL2	16(SP), 12(SP)		
7E			04	AE		01	002CC		ADDL3	#1, 4(SP), -(SP)		
			6E	AE		04	002D1		ADDL2	KEY, (SP)		
7E			04	AE		02	002D5		ADDL3	#2, KEY, -(SP)		
OC	AE		9E	9E		0C	002DA		MOVC5	12(SP), @ (SP)+, @ (SP)+, 12(SP), @8(SP)		
						08	002E1					
13	06	AA				05	002E3	47\$:	BBC	#3, 6(IFAB), 49\$		2023
06	00A2	CA				03	002E8		BBC	#3, 162(IFAB), 48\$		2025
07	00A2	CA				02	002EE		BBS	#2, 162(IFAB), 49\$		2027
		50				14	002F4	48\$:	MOVL	STATUS, R0		2029
						0133	002F8		BRW	71\$		
						5C	002FB	49\$:	CLRL	AP		2045
06	06	A8				04	002FD		BBC	#4, 6(RAB), 50\$		2056
09	00A2	CA				02	00302		BBC	#2, 162(IFAB), 51\$		2058
		13				06	00308	50\$:	BLBS	6(IFAB), 52\$		2061
0E		A8				02	0030C		BBS	#2, 4(RAB), 52\$		2063
		52				0080	00311	51\$:	MOVZWL	128(IRAB), R2		2065
		51				78	00316		MOVL	120(IRAB), R1		
						0000G	0031A		BSBW	RMSQUERY_LCK		
						0C	0031D		BRB	53\$		
			52			0080	0031F	52\$:	MOVZWL	128(IRAB), R2		2067
			51			78	00324		MOVL	120(IRAB), R1		
						0000G	00328		BSBW	RMSLOCK		
			52			50	0032B	53\$:	MOVL	R0, STATUS		
	00008061	8F				52	0032E		CMPL	STATUS, #32865		2076
						67	00335		BNEQ	60\$		
			0A			06	00337		BLBC	6(IFAB), 54\$		2088
			CA			01	0033B		BBC	#1, 160(IFAB), 54\$		2090
04	00A0	A9				01	00341		MOVB	#1, 64(IRAB)		2092
	40					7E	00345	54\$:	CLRL	-(SP)		2094
			7E			0080	00347		MOVZWL	128(IRAB), -(SP)		2095
						78	0034C		PUSHL	120(IRAB)		2094
						0000G	0034F		BSBW	RMSFIND_BY_RRV		
			5E			0C	00352		ADDL2	#12, SP		
			03			50	00355		BLBS	TEMP_STATUS, 55\$		
			52			50	00358		MOVL	TEMP_STATUS, STATUS		2098
			3D			52	0035B	55\$:	BLBC	STATUS, 59\$		2105

3C	66		05	E0	0035E	BBS	#5, (REC_ADDR), 60\$	2107
	36		52	E9	00362	BLBC	STATUS, 59\$	2121
		21	A7	95	00365	TSTB	33(IDX_DFN)	2118
			20	12	00368	BNEQ	58\$	
	50	20	A9	D0	0036A	MOVL	32(IRAB), R0	2121
70	A9	1C	A0	D0	0036E	MOVL	28(R0), 112(IRAB)	
	03	00B7	CA	91	00373	CMPB	183(IFAB), #3	2122
			06	1E	00378	BGEQU	56\$	
	50	01	A6	9A	0037A	MOVZBL	1(REC_ADDR), R0	
			04	11	0037E	BRB	57\$	
	50	01	A6	3C	00380	MOVZWL	1(REC_ADDR), R0	
74	A9		50	B0	00384	MOVW	R0, 1T6(IRAB)	
			14	11	00388	BRB	60\$	2118
	54	20	A9	D0	0038A	MOVL	32(IRAB), BDB	2144
		20	A9	D4	0038E	CLRL	32(IRAB)	
			7E	D4	00391	CLRL	-(SP)	
			0000G	30	00393	BSBW	RMSRLSBKT	
	5E		04	C0	00396	ADDL2	#4, SP	
			03	11	00399	BRB	60\$	2118
		20	A9	D4	0039B	CLRL	32(IRAB)	2147
	44		52	E9	0039E	BLBC	STATUS, 65\$	2156
		20	A9	D5	003A1	TSTL	32(IRAB)	2158
			3F	13	003A4	BEQL	65\$	
04	66		05	E0	003A6	BBS	#5, (REC_ADDR), 61\$	2160
37	66		06	E1	003AA	BBC	#6, (REC_ADDR), 65\$	2162
	53	01	A6	B0	003AE	MOVW	1(REC_ADDR), RECORD_ID	2169
	06	06	AA	E9	003B2	BLBC	6(IFAB), 62\$	2175
		00000000G	EF	16	003B6	JSB	RMSRU RECLAIM	2177
	01	A6	53	B1	003BC	CMPW	RECORD_ID, 1(REC_ADDR)	2189
			04	12	003C0	BNEQ	63\$	
1F	66		05	E1	003C2	BBC	#5, (REC_ADDR), 65\$	2191
	54	20	A9	D0	003C6	MOVL	32(IRAB), BDB	2198
		20	A9	D4	003CA	CLRL	32(IRAB)	
			7E	D4	003CD	CLRL	-(SP)	
			0000G	30	003CF	BSBW	RMSRLSBKT	
	5E		04	C0	003D2	ADDL2	#4, SP	
8061	8F		52	B1	003D5	CMPW	STATUS, #32865	2200
			07	12	003DA	BNEQ	64\$	
	52	8262	8F	3C	003DC	MOVZWL	#33378, STATUS	2202
			02	11	003E1	BRB	65\$	
			52	D4	003E3	CLRL	STATUS	2204
06	06	A8	04	E1	003E5	BBC	#4, 6(RAB), 66\$	2210
3B	00A2	CA	02	E1	003EA	BBC	#2, 162(IFAB), 70\$	2212
		05	AA	E8	003F0	BLBS	6(IFAB), 67\$	2214
32	04	A8	02	E1	003F4	BBC	#2, 4(RAB), 70\$	2216
21	04	A9	0D	E4	003F9	BBSC	#13, 4(IRAB), 68\$	2231
		01	52	B1	003FE	CMPW	STATUS, #1	2233
			1C	13	00401	BEQL	68\$	
	8061	8F	52	B1	00403	CMPW	STATUS, #32865	2235
			15	13	00408	BFQL	68\$	
	8071	8F	52	B1	0040A	CMPW	STATUS, #32881	2237
			0E	13	0040F	BEQL	68\$	
		20	A9	D5	00411	TSTL	32(IRAB)	2257
			15	12	00414	BNEQ	70\$	
	82AA	8F	52	B1	00416	CMPW	STATUS, #33450	2259
			0A	12	0041B	BNEQ	69\$	
			0C	11	0041D	BPB	70\$	2263

; R0

04	06	5C	01	DO	0041F	68\$:	MOVL	#1, AP	:	2266
	05	A8	02	EO	00422		BBS	#2, 6(RAB), 70\$	:	2268
		A9	20	88	00427	69\$:	BISB2	#32, 5(IRAB)	:	2270
		50	52	DO	0042B	70\$:	MOVL	STATUS, R0	:	2272
		5E	18	CO	0042E	71\$:	ADDL2	#24, SP	:	2275
			05	00431			RSB		:	

; Routine Size: 1074 bytes, Routine Base: RM\$RMS3 + 032D

; 2215           2276 1 END                           ! of module  
; 2216           2277 1  
; 2217           2278 0 ELUDOM

PSECT SUMMARY

Name	Bytes	Attributes
RM\$RMS3	1887	NOVEC, NOWRT, RD, EXE, NOSHR, GBL, REL, CON, PIC, ALIGN(2)

Library Statistics

File	Total	Symbols Loaded	Percent	Pages Mapped	Processing Time
_\$255\$DUA28:[RMS.OBJ]RMS.L32;1	3109	140	4	154	00:00.4

COMMAND QUALIFIERS

BLISS/CHECK=(FIELD,INITIAL,OPTIMIZE)/LIS=LIS\$:RM3GET/OBJ=OBJ\$:RM3GET MSRC\$:RM3GET/UPDATE=(ENH\$:RM3GET)

; Size: 1887 code + 0 data bytes  
; Run Time: 00:48.3  
; Elapsed Time: 01:24.0  
; Lines/CPU Min: 2829  
; Lexemes/CPU-Min: 17134  
; Memory Used: 399 pages  
; Compilation Complete



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- RM3MKTDX LIS
- RM3FNDRU LIS
- RM3FNDRFA LIS
- RM3GET LIS
- RM3LUDR LIS
- RM3JOURN LIS
- RM3KEYDSC LIS
- RM3MISC LIS
- RM3MISPLT LIS

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