

Variable	Descriptive statistics	Statistical tests
Age	Mean = 34.5, SD = 10.2	ANOVA
Gender	Male = 65, Female = 35	Chi-square
Education	High school = 40, College = 40, Graduate = 20	ANOVA
Income	Low = 30, Middle = 40, High = 30	ANOVA
Marital status	Married = 50, Single = 30, Divorced = 20	Chi-square
Occupation	Professional = 30, Managerial = 30, Service = 40	ANOVA
Religious affiliation	Catholic = 40, Protestant = 30, Jewish = 20, Muslim = 10	Chi-square
Political affiliation	Democrat = 40, Republican = 30, Independent = 30	ANOVA
Health status	Good = 40, Fair = 30, Poor = 30	ANOVA
Life satisfaction	High = 30, Medium = 40, Low = 30	ANOVA
Stress level	Low = 30, Medium = 40, High = 30	ANOVA
Self-esteem	High = 30, Medium = 40, Low = 30	ANOVA
Resilience	High = 30, Medium = 40, Low = 30	ANOVA
Optimism	High = 30, Medium = 40, Low = 30	ANOVA
Gratitude	High = 30, Medium = 40, Low = 30	ANOVA
Forgiveness	High = 30, Medium = 40, Low = 30	ANOVA
Empathy	High = 30, Medium = 40, Low = 30	ANOVA
Compassion	High = 30, Medium = 40, Low = 30	ANOVA
Kindness	High = 30, Medium = 40, Low = 30	ANOVA
Generosity	High = 30, Medium = 40, Low = 30	ANOVA
Patience	High = 30, Medium = 40, Low = 30	ANOVA
Humility	High = 30, Medium = 40, Low = 30	ANOVA
Modesty	High = 30, Medium = 40, Low = 30	ANOVA
Shyness	High = 30, Medium = 40, Low = 30	ANOVA
Introversion	High = 30, Medium = 40, Low = 30	ANOVA
Extroversion	High = 30, Medium = 40, Low = 30	ANOVA
Social skills	High = 30, Medium = 40, Low = 30	ANOVA
Communication skills	High = 30, Medium = 40, Low = 30	ANOVA
Interpersonal skills	High = 30, Medium = 40, Low = 30	ANOVA
Teamwork skills	High = 30, Medium = 40, Low = 30	ANOVA
Leadership skills	High = 30, Medium = 40, Low = 30	ANOVA
Problem-solving skills	High = 30, Medium = 40, Low = 30	ANOVA
Decision-making skills	High = 30, Medium = 40, Low = 30	ANOVA
Conflict resolution skills	High = 30, Medium = 40, Low = 30	ANOVA
Emotional regulation skills	High = 30, Medium = 40, Low = 30	ANOVA
Stress management skills	High = 30, Medium = 40, Low = 30	ANOVA
Time management skills	High = 30, Medium = 40, Low = 30	ANOVA
Organization skills	High = 30, Medium = 40, Low = 30	ANOVA
Planning skills	High = 30, Medium = 40, Low = 30	ANOVA
Goal setting skills	High = 30, Medium = 40, Low = 30	ANOVA
Self-discipline skills	High = 30, Medium = 40, Low = 30	ANOVA
Perseverance skills	High = 30, Medium = 40, Low = 30	ANOVA
Resilience skills	High = 30, Medium = 40, Low = 30	ANOVA
Optimism skills	High = 30, Medium = 40, Low = 30	ANOVA
Gratitude skills	High = 30, Medium = 40, Low = 30	ANOVA
Forgiveness skills	High = 30, Medium = 40, Low = 30	ANOVA
Empathy skills	High = 30, Medium = 40, Low = 30	ANOVA
Compassion skills	High = 30, Medium = 40, Low = 30	ANOVA
Kindness skills	High = 30, Medium = 40, Low = 30	ANOVA
Generosity skills	High = 30, Medium = 40, Low = 30	ANOVA
Patience skills	High = 30, Medium = 40, Low = 30	ANOVA
Humility skills	High = 30, Medium = 40, Low = 30	ANOVA
Modesty skills	High = 30, Medium = 40, Low = 30	ANOVA
Shyness skills	High = 30, Medium = 40, Low = 30	ANOVA
Introversion skills	High = 30, Medium = 40, Low = 30	ANOVA
Extroversion skills	High = 30, Medium = 40, Low = 30	ANOVA
Social skills	High = 30, Medium = 40, Low = 30	ANOVA
Communication skills	High = 30, Medium = 40, Low = 30	ANOVA
Interpersonal skills	High = 30, Medium = 40, Low = 30	ANOVA
Teamwork skills	High = 30, Medium = 40, Low = 30	ANOVA
Leadership skills	High = 30, Medium = 40, Low = 30	ANOVA
Problem-solving skills	High = 30, Medium = 40, Low = 30	ANOVA
Decision-making skills	High = 30, Medium = 40, Low = 30	ANOVA
Conflict resolution skills	High = 30, Medium = 40, Low = 30	ANOVA
Emotional regulation skills	High = 30, Medium = 40, Low = 30	ANOVA
Stress management skills	High = 30, Medium = 40, Low = 30	ANOVA
Time management skills	High = 30, Medium = 40, Low = 30	ANOVA
Organization skills	High = 30, Medium = 40, Low = 30	ANOVA
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Shyness skills	High = 30, Medium = 40, Low = 30	ANOVA
Introversion skills	High = 30, Medium = 40, Low = 30	ANOVA
Extroversion skills	High = 30, Medium = 40, Low = 30	ANOVA
Social skills		

FILEID**DBGLIB

```
DDDDDDDD  BBBB BBBB  GGGGGGGG  LL      IIIIII  BBBB BBBB
DDDDDDDD  BBBB BBBB  GGGGGGGG  LL      IIIIII  BBBB BBBB
DD      DD  BB      BB  GG      LL      II      BB      BB
DD      DD  BB      BB  GG      LL      II      BB      BB
DD      DD  BB      BB  GG      LL      II      BB      BB
DD      DD  BB      BB  GG      LL      II      BB      BB
DC      DD  BBBB BBBB  GG      LL      II      BBBB BBBB
DD      DD  BBBB BBBB  GG      LL      II      BBBB BBBB
DD      DD  BB      BB  GG  GGGGGG  LL      II      BB      BB
DD      DD  BB      BB  GG  GGGGGG  LL      II      BB      BB
DD      DD  BB      BB  GG      GG  LL      II      BB      BB
DD      DD  BB      BB  GG      GG  LL      II      BB      BB
DDDDDDDD  BBBB BBBB  GGGGGG  LLLLLLLLLL  IIIIII  BBBB BBBB
DDDDDDDD  BBBB BBBB  GGGGGG  LLLLLLLLLL  IIIIII  BBBB BBBB
....
....
....
....
```

```
RRRRRRRR  EEEEEEEEE  QQQQQQ
RRRRRRRR  EEEEEEEEE  QQQQQQ
RR      RR  EE      QQ      QQ
RR      RR  EE      QQ      QQ
RR      RR  EE      QQ      QQ
RR      RR  EE      QQ      QQ
RRRRRRRR  EEEEEEEEE  QQ      QQ
RRRRRRRR  EEEEEEEEE  QQ      QQ
RR  RR  EE      QQ  QQ  QQ
RR  RR  EE      QQ  QQ  QQ
RR      RR  EE      QQ      QQ
RR      RR  EE      QQ      QQ
RR      RR  EEEEEEEEE  QQQQ  QQ
RR      RR  EEEEEEEEE  QQQQ  QQ
```

DBGLIB -- COMMON DEFINITION FILE FOR THE VAX DEBUGGER

Version: 'V04-000'

```
*****
*
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```

WRITTEN BY
Bert Beander June, 1980.MODIFIED BY
John Francis
Ping Sager
Rich Title
Vicki Holt
Sid MaxwellMODULE FUNCTION:
This REQUIRE file contains all literal, macro, and data structure definitions used by new Bliss modules in the Debugger. Old Bliss modules do not use this file for historical reasons.

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GENERAL DEBUGGER DEFINITIONS

The declarations in this section define symbols of general utility throughout the Debugger. This includes widely used literals and various utility macros.

DATA STRUCTURE DEFINITION AND ACCESS

Here we declare all the macros used to define and access BLISS data structures. These are the names L, W, B, W0, W1, and so forth. The actual definitions are in a separate REQUIRE file, but they are required for the LIBRARY compilation of the present REQUIRE file.

LIBRARY 'LIBS:STRUCDEF.L32';

BLISS-32 LANGUAGE EXTENSIONS

Here we declare various minor extensions to the Bliss-32 language.

LITERAL

```
TRUE      = 1,           ! Define TRUE and FALSE
FALSE     = 0;
```

MACRO

```
REPEAT      = WHILE 1 DO%, ! Infinite Loop
```

! ZEROCOR zeroes an area of memory. The two parameters are ADDRESS and COUNT, which are the address of the area to be zeroed and the number of contiguous longwords to be zeroed, respectively.

```
ZEROCOR (ADDRESS, COUNT) =
  CH$FILL (0, (COUNT) * 4, CH$PTR (ADDRESS))%,
```

```
ONEOF (X) [] =
  (MASK (%REMAINING)) ^ (X) LSS 0%,
```

```
OR_OP (X) [] =
  OR%,
```

```
MASK (X) [] =
  1 ^ (31-X) OR_OP (%REMAINING) MASK (%REMAINING)%,
```

CH\$SEQUENCE (N) =
VECTOR [CH\$ALLOCATION (N)]%,
CONSECUTIVE [X] = LITERAL X = %COUNT%;

ERROR REPORTING

The following macro is used for reporting internal DEBUG coding errors symbolically. The macro is used as follows:

```
$DBG_ERROR('MODNAME\ROUTNAME number');
```

This signals the internal DEBUG error message DBGS_INTERR which prints the text of the macro parameter. The convention is that this parameter should include the module name and the routine name of the signal location so that the point where the error was signalled can be easily determined by DEBUG developers. In addition, a number may be included in the signal text to make the text unique when there is more than one \$DBG_ERROR invocation in the same routine.

```
MACRO $DBG_ERROR (STRING) =  
  SIGNAL (DBGS_INTERR, 1, UPLIT BYTE (%ASCIC STRING)  
  %IF %LENGTH GTR 1  
  %THEN  
    %REMAINING)  
  %ELSE  
    )  
  %FI %;
```

DEBUG OUTPUT MACROS

The following MACROs are used in several places in the debugger to do formatted ASCII output (\$FAO).

MACRO

! \$FAO_STG_COUNT makes a counted byte string out of an ASCII string. This macro is useful to transform an FAO control string into the address of such a string, whose first byte contains the length of the string in bytes.

\$FAO_STG_COUNT (STRING) =
 DPLIT BYTE (%CHARCOUNT (STRING), %ASCII STRING)%,

! \$FAO_TT_OUT constructs a call to FAO with a control string, and some arguments to the control string. This formatted string is then output to the output device.

\$FAO_TT_OUT (CTL_STRING) =
 DBG\$FAO_OUT (\$FAO_STG_COUNT (CTL_STRING)
 %IF %LENGTH GTR 1
 %THEN
 , %REMAINING
 %FI
)%;


```

DBG$K_MIN_DESCR_TYPE = 120,  ! Used in FROM clause of CASE
DBG$K_LITERAL        = 120,  ! Literal value
DBG$K_PRIMARY_DESC   = 121,  ! Primary descriptor
DBG$K_VALUE_DESC     = 122,  ! Value descriptor
DBG$K_PERM_DESC      = 123,  ! Permanent Symbol descriptor
DBG$K_INSTRUCTION    = 124,  ! Instruction
DBG$K_NC_INSTRUCTION = 125,  ! Named constant, instruction type
DBG$K_NC_OTHER       = 126,  ! Named constant, not instruction
DBG$K_OTHER          = 127,  ! Language-specific data type
DBG$K_NOTYPE         = 128,  ! Virtual address without any type
DBG$K_EXTERNAL_DESC  = 129,  ! External (printable) format
DBG$K_VAX_DESC       = 130,  ! VAX standard descriptor
DBG$K_V_VALUE_DESC   = 131,  ! Volatile Value descriptor
DBG$K_AED            = 132,  ! This code appears in the DBG$B_DHDR_TYPE
                             ! field of the new primary descriptor
                             ! and also in the corresponding place
                             ! in the AED to tell that the descriptor
                             ! is an AED and not a new-style descriptor.

DBG$K_MAX_DESCR_TYPE = 132;  ! Used in TO clause of CASE.

```

```
! Define the Radix Literals.
```

```

LITERAL
DBG$K_DEFAULT = 1,  ! Default source language radix
DBG$K_BINARY  = 2,  ! Binary radix
DBG$K_OCTAL   = 8,  ! Octal radix
DBG$K_DECIMAL = 10, ! Decimal radix
DBG$K_HEX     = 16;  ! Hexadecimal radix

```

```
! Define Pseudo-Symbol codes.
```

```

LITERAL
DBG$K_CURRENT_LOC = 220,  ! Current location ""
DBG$K_PREDECESSOR = 221,  ! Previous location ""
DBG$K_SUCCESOR    = 222,  ! Next location <CR>
DBG$K_LAST_VALUE  = 223,  ! Last value ""
DBG$K_INDIRECTION = DBG$K_CURRENT_LOC;
                             ! An alternative meaning of ""
                             ! is indirection

```

```
! Token type literals. These are used by the Pathname Parser and scanners.
```

```

LITERAL
DBG$K_TOK_NULL = 0,  ! Null or EOL token
DBG$K_TOK_INVALID = 1,  ! Invalid token
DBG$K_TOK_LINE = 2,  ! '%LINE' token
DBG$K_TOK_LABEL = 3,  ! '%LABEL' token
DBG$K_TOK_BS = 4,  ! '\' token
DBG$K_TOK_ID = 5,  ! ID token
DBG$K_TOK_INT = 6,  ! integer token
DBG$K_TOK_DOT = 7,  ! '.' token
DBG$K_TOK_REG = 8,  ! '%REGISTER' token
DBG$K_TOK_QNAME = 9,  ! '%NAME' token
DBG$K_TOK_LOWEST = DBG$K_TOK_NULL,

```

DBG\$K_TOK_HIGHEST =DBG\$K_TOK_QNAME;

! Return state values for the Pathname Parser

LITERAL

DBG\$K_PN	= 0,	! Data reference or lexical reference
DBG\$K_REG	= 1,	! Register reference
DBG\$K_LINE	= 2,	! '%LINE' reference
DBG\$K_LABEL	= 3;	! '%LABEL' reference

! Command verb literals. These are used by the Command Line Interpreter.

LITERAL

DBG\$K_ALLOCATE_VERB	= 22,	! ALLOCATE
DBG\$K_AT_SIGN_VERB	= 1,	! @
DBG\$K_ATTACH_VERB	= 26,	! ATTACH
DBG\$K_CALL_VERB	= 2,	! CALL
DBG\$K_CANCEL_VERB	= 3,	! CANCEL
DBG\$K_DECLARE_VERB	= 20,	! DECLARE
DBG\$K_DEFINE_VERB	= 4,	! DEFINE
DBG\$K_DELETE_VERB	= 32,	! DELETE
DBG\$K_DEPOSIT_VERB	= 5,	! DEPOSIT
DBG\$K_DISPLAY_VERB	= 28,	! DISPLAY
DBG\$K_DUMP_VERB	= 27,	! DUMP (Only for developer use)
DBG\$K_EDIT_VERB	= 33,	! EDIT
DBG\$K_EVALUATE_VERB	= 6,	! EVALUATE
DBG\$K_EXAMINE_VERB	= 7,	! EXAMINE
DBG\$K_EXIT_VERB	= 8,	! EXIT
DBG\$K_EXITLOOP_VERB	= 19,	! EXITLOOP
DBG\$K_FOR_VERB	= 25,	! FOR
DBG\$K_GO_VERB	= 9,	! GO
DBG\$K_HELP_VERB	= 13,	! HELP
DBG\$K_IF_VERB	= 16,	! IF
DBG\$K_REPEAT_VERB	= 18,	! DO
DBG\$K_SAVE_VERB	= 31,	! SAVE
DBG\$K_SCROLL_VERB	= 29,	! SCROLL
DBG\$K_SEARCH_VERB	= 15,	! SEARCH
DBG\$K_SELECT_VERB	= 30,	! SELECT
DBG\$K_SET_VERB	= 10,	! SET
DBG\$K_SHOW_VERB	= 11,	! SHOW
DBG\$K_SPAWN_VERB	= 21,	! SPAWN
DBG\$K_STEP_VERB	= 12,	! STEP
DBG\$K_SYMBOLIZE_VERB	= 23,	! SYMBOLIZE
DBG\$K_TYPE_VERB	= 14,	! TYPE
DBG\$K_UNDEFINE_VERB	= 24,	! UNDEFINE
DBG\$K_WHILE_VERB	= 17,	! WHILE
DBG\$K_FIRST_VERB	= DBG\$K_AT_SIGN_VERB,	
DBG\$K_LAST_VERB	= DBG\$K_EDIT_VERB,	
EVENT\$K_SET_BREAK	= 1,	! Also SET_BREAK in DBGNSET
EVENT\$K_SET_BREAK_EXC	= 3,	! Also SET_EXCEPTION_BREAK in DBGNSET
EVENT\$K_SET_TRACE	= 14,	! Also SET_TRACE in DBGNSET
EVENT\$K_SET_WATCH	= 17,	! Also SET_WATCH in DBGNSET

```

EVENTSK_SET_STEP      = 11,   ! Also SET_STEP in DBGNSSET
EVENTSK_STEP          = 12,   ! Also DBGSK_STEP_VERB

EVENTSK_SHOW_BREAK    = 1,    ! Also SHOW_BREAK in DBGNSHOW
EVENTSK_SHOW_TRACE    = 11,   ! Also SHOW_TRACE in DBGNSHOW
EVENTSK_SHOW_WATCH    = 14,   ! Also SHOW_WATCH in DBGNSHOW

EVENTSK_CANCEL_BREAK  = 2,    ! Also CANCEL_BREAK in DBGNCANCL
EVENTSK_CANCEL_BREAK_EXC = 4,  ! Also CANCEL_EXCEPTION_BREAK in DBGNCANCL
EVENTSK_CANCEL_TRACE  = 9,    ! Also CANCEL_TRACE in DBGNCANCL
EVENTSK_CANCEL_WATCH  = 14;   ! Also CANCEL_WATCH in DBGNCANCL

```

! Value kind literals. These values are returned as 'value kinds' by routines
 DBG\$STA_SYMVALUE and DBG\$STA_VALSPEC.

```

LITERAL
DBG$K_VAL_NOVALUE      = 0,    ! Symbol is a type and has no value
DBG$K_VAL_LITERAL      = 1,    ! Value is a literal
DBG$K_VAL_ADDR         = 2,    ! Value is an address
DBG$K_VAL_DESCR        = 3,    ! Value is a descriptor pointer
DBG$K_VAL_UNALLOC      = 4;    ! Symbol was never allocated and
                                ! hence has no value

```

! Debug specific DTYPE codes which are used in DBGNEXMNE.
 ! These codes are TEMPORARY!

```

LITERAL
DBG$K_DTYPE_AD         = 56;   ! ASCII string

```

! Maximum number of digits in a packed decimal number.

```

LITERAL
DBG$K_LARGEST_PACKED   = 31;

```

! PL/I specific data type codes. These are referenced in DBGPERMOP and in
 DBGEVALOP in calls to the PL/I RTL during type conversions and calculations.

```

LITERAL
DBG$K_PLI_PIC          = 1,
DBG$K_PLI_FIX_BIN      = 2,
DBG$K_PLI_FLO_BIN      = 3,
DBG$K_PLI_FIX_DEC      = 4,
DBG$K_PLI_FLO_DEC      = 5,
DBG$K_PLI_CHAR         = 10,
DBG$K_PLI_CHAR_VAR     = 11,
DBG$K_PLI_UBIT         = 12,
DBG$K_PLI_ABIT         = 14;

```

! Debug Print Control Code.

```

LITERAL

```



```
DBG$K_PRTSET_LMARGIN = 1,  ! Set left margin for DBG$PRINT
DBG$K_PRTSET_RLMARGIN = 2,  ! Set relative left margin
DBG$K_PRTBRK_ON BLANKS = 3,  ! Set break on blanks flag
DBG$K_PRTSET_CONTINUE = 4,  ! Set indentation length for continuation
DBG$K_PRT_RESET       = 5;  ! Reset to default settings
```

!...-

! Input string descriptor definition. This definition is compatible with the
! use of the DSC\$B_DTYPE, DSC\$B_CLASS, DSC\$W_LENGTH, and DSC\$A_POINTER macros.

LITERAL

DBG\$K_STG_DESC_SIZE = 12; ! Length of descriptor in bytes

MACRO

DBG\$STG_DESC = BLOCK [DBG\$K_STG_DESC_SIZE, BYTE] %;

ADDRESS DESCRIPTORS

The Address Descriptor is used to define an address where a bit offset is needed in addition to a byte address.

```

3 3 2 2 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1 1
1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0

```

0	----- DBG\$L_ADDRESS_BYTE_ADDR -----
1	----- DBG\$L_ADDRESS_BIT_OFFSET -----

A pointer to an Address Descriptor is declared as follows:

```
ADDPTR: REF DBG$ADDRESS_DESC;
```

Define the fields in the Address Descriptor. Also define the declaration macro.

```

***
FIELD DBG$ADDRESS_DESC_FIELDS =
    SET
    DBG$L_ADDRESS_BYTE_ADDR = [ 0, L_ ], ! Byte address
    DBG$L_ADDRESS_BIT_OFFSET = [ 1, L_ ], ! Bit offset
    TES;

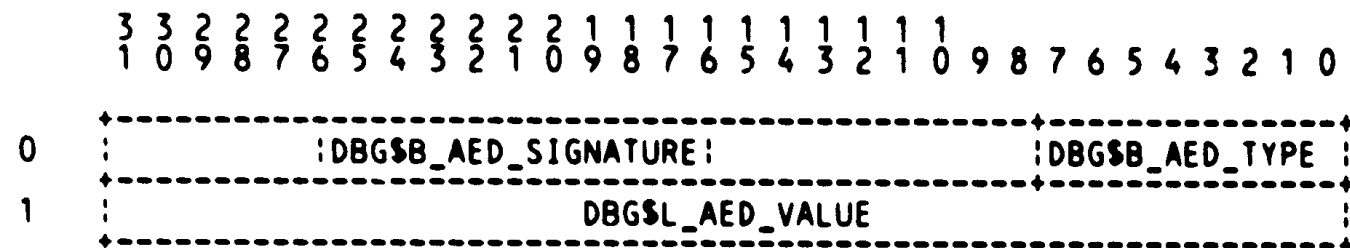
LITERAL
DBG$K_ADDRESS_DESC_SIZE = 2;           ! Size of descriptor in longwords

MACRO
DBG$ADDRESS_DESC = BLOCK(DBG$K_ADDRESS_DESC_SIZE)
    FIELD(DBG$ADDRESS_DESC_FIELDS) %;
***

```

ADDRESS EXPRESSION DESCRIPTORS

The Address Expression Descriptor is used to define an address expression. It has the following structure:



A pointer to an Address Expression Descriptor is declared as follows:

AEDPTR: REF DBGSAED

Define the fields in the Address Expression Descriptor. Also define the declaration macro.

```

FIELD DBG$AED_FIELDS =
    SET
    DBG$B_AED_TYPE = [ 0, B0_ ], ! The type of the value field
    DBG$B_AED_SIGNATURE = [ 0, B2_ ], ! Always contains a code DBG$K_AED
                                     ! saying this is an AED.
    DBG$L_AED_VALUE = [ 1, L_ ] ! L-value or pointer to descriptor
    TES;

```

```
LITERAL    DBGSK_AED_SIZE    = 2;                ! Size of descriptor in longwords
```

```
MACRO DBG$AED = BLOCK[DBG$K_AED_SIZE] FIELD(DBG$AED_FIELDS) %;
```

The following are the legal values for the DBGSB_AED_TYPE field.

DBG\$K_PRIMARY_DESC	:	The value field contains the address of a primary descriptor
DBG\$K_PERM_DESC	:	The value field contains the address of a permanent symbol descriptor
DBG\$K_INSTRUCTION	:	The value field contains a PC value
DBG\$K_NOTYPE	:	The value field contains an untyped L-value

C O M M A N D E X E C U T I O N T R E E S

This section contains the definitions for the nodes that are part of a command execution tree. A command execution tree is built when a debug command is parsed. The tree contains a verb node representing the command (e.g., EXAMINE, DEPOSIT, etc.) as the header node. Linked to the verb node are chains of adverb nodes and noun nodes. The noun node chain hangs off the OBJECT_PTR field and the adverb node chain hangs off the ADVERB_PTR field. The exact structure of the tree depends on the command.

VERB NODES

0	!SB_VERB_COMPOSITE!SB_VERB_LITERAL!
1	DBG\$L_VERB_ADVERB_PTR
2	DBG\$L_VERB_OBJECT_PTR

FIELD

DBG\$VERB_NODE_FIELDS =
SET

DBG\$B_VERB_LITERAL	= [0, 0, 8, 0],	! First level verb
DBG\$B_VERB_COMPOSITE	= [0, 8, 8, 0],	! Second or third level verb
DBG\$L_VERB_ADVERB_PTR	= [1, 0, 32, 0],	! Pointer to adverb node
DBG\$L_VERB_OBJECT_PTR	= [2, 0, 32, 0]	! Pointer to object (noun node)

TES;

!

LITERAL

DBG\$K_VERB_NODE_SIZE = 3; ! Length in longwords

MACRO

DBG\$VERB_NODE = BLOCK [DBG\$K_VERB_NODE_SIZE] FIELD (DBG\$VERB_NODE_FIELDS) %;

ADVERB NODES

0	!_ADVERB_LITERAL!
1	DBG\$L_ADVERB_VALUE
2	DBG\$L_ADVERB_LINK

FIELD

DBG\$ADVERB_NODE_FIELDS =
SET

DBG\$B_ADVERB_LITERAL	= [0, 0, 8, 0],	! Adverb id literal
DBG\$L_ADVERB_VALUE	= [1, 0, 32, 0],	! Value (integer)
DBG\$L_ADVERB_LINK	= [2, 0, 32, 0]	! Link to next adverb node

TES;

LITERAL DBG\$K_ADVERB_NODE_SIZE = 3; ! Length in longwords

MACRO DBG\$ADVERB_NODE = BLOCK [DBG\$K_ADVERB_NODE_SIZE] FIELD (DBG\$ADVERB_NODE_FIELDS) %;

NOUN NODES

0	DBG\$NOUN_VALUE
1	DBG\$ADJECTIVE_PTR
2	DBG\$NOUN_LINK
3	DBG\$NOUN_VALUE2

FIELD DBG\$NOUN_NODE_FIELDS =

SET
DBG\$NOUN_VALUE = [0, L_] ! Noun value
DBG\$ADJECTIVE_PTR = [1, L_] ! Pointer to adjective (unused)
DBG\$NOUN_LINK = [2, L_] ! Pointer to next noun
DBG\$NOUN_VALUE2 = [3, L_] ! Additional Noun value
DBG\$NOUN_VALUE3 = [4, L_] ! More Noun value (long nodes only)
DBG\$NOUN_VALUE4 = [5, L_] ! More Noun value (long nodes only)
TES;

LITERAL DBG\$K_NOUN_NODE_SIZE = 4; ! Length of normal Noun Node in longwords
DBG\$K_NOUN_NODE_SIZE_LONG = 6; ! Length of long Noun Node in longwords

MACRO DBG\$NOUN_NODE = BLOCK [DBG\$K_NOUN_NODE_SIZE] FIELD (DBG\$NOUN_NODE_FIELDS) %;

C O M M A N D I N P U T S T R E A M

The Command Input Stream (CIS) is a linked list that the debugger maintains to describe where it is currently taking input from.

The format of a Command Input Stream Descriptor is:

```

3 3 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1 1
1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0

```

0	Unused	SB_INPUT_TYPE	CISW_LENGTH
1	CISW_INPUT_PTR		
2	CISW_NEXT_LINK		
3	CISW_INIT_ADDR		
4	Unused	Flags	CISW_INIT_LENGTH
5	CISW_WHILE_CLAUSE		
6	CISL_FOR_UPPER_BOUND or CISL_REPEAT_COUNT		
7	CISW_FOR_LOOP_VAR		
8	CISL_FOR_LOOP_INCR		
9	CISL_SCREEN_OUTPUT		
10	CISL_SCREEN_SOURCE		
11	CISL_SCREEN_ERROR		
12	CISW_BUFLIST		
13	CISW_WHILE_LENGTH		

A pointer to a command input stream descriptor block is declared with:

```
CIS_PTR: REF CISLINK
```

Define the fields, the descriptor size, and the declaration macro.

```
FIELD CIS_FIELDS =
```

```

SET
CISW_INPUT_TYPE = [ 2, B_ ], ! Type of CIS link
CISW_LENGTH     = [ 0, W_ ], ! Length of input buffer
CISW_INPUT_PTR  = [ 4, L_ ], ! Pointer to input buffer

```

```

CIS$A_NEXT_LINK    = [ 8, L_ ]  : Ptr to next cis link
CIS$A_INIT_ADDR    = [ 12, L_ ] : Ptr to start of input buffer
CIS$W_INIT_LENGTH  = [ 16, W_ ] : Allocated buffer size
CIS$V_REM_FLAG     = [ 18, V_(0) ] : Is link flagged for removal?
CIS$V_WHILE_FLAG   = [ 18, V_(1) ] : TRUE when condition is true
CIS$V_SCREEN_NOGO  = [ 18, V_(2) ] : Saved NOGO flag to turn off STEPs
                                   and GOs in CIS_SCREEN entries
CIS$A_WHILE_CLAUSE = [ 20, L_ ]  : Pointer to WHILE command
CIS$L_REPEAT_COUNT = [ 24, L_ ]  : Count for REPEAT loops
CIS$L_FOR_UPPER_BOUND = [ 24, L_ ] : Upper bound for FOR loops
CIS$A_FOR_LOOP_VAR = [ 28, L_ ]  : Pointer to the FOR loop var name
CIS$L_FOR_LOOP_INCR = [ 32, L_ ] : Increment for FOR loops
CIS$L_SCREEN_OUTPUT = [ 36, L_ ] : Reset value for DBG$GL_SCREEN_OUTPUT
CIS$L_SCREEN_SOURCE = [ 40, L_ ] : Reset value for DBG$GL_SCREEN_SOURCE
CIS$L_SCREEN_ERROR  = [ 44, L_ ] : Reset value for DBG$GL_SCREEN_ERROR
CIS$A_BUFLIST      = [ 48, L_ ]  : List of buffers to free in CIS_REMOVE
CIS$W_WHILE_LENGTH = [ 52, W_ ]
TES;

```

```

LITERAL
CIS_ELEMENTS      = 54;          : CIS block size (bytes)

```

```

MACRO
CIS$LINK = BLOCK [CIS_ELEMENTS, BYTE] FIELD(CIS_FIELDS) %;

```

```

: The following are legal types for CIS links.

```

```

LITERAL
DBG$K_CIS_MINIMUM      = 0.      : Lowest type CIS
DBG$K_CIS_DBG$INPUT    = 0.      : Type DBG$INPUT
DBG$K_CIS_RAB          = 1.      : Type RAB
DBG$K_CIS_INPBUF       = 2.      : Type input buffer
DBG$K_CIS_ACBUF        = 3.      : Type action buffer
DBG$K_CIS_REPEAT       = 4.      : Type repeat buffer
DBG$K_CIS_WHILE        = 5.      : Type while buffer
DBG$K_CIS_IF           = 6.      : Type if buffer
DBG$K_CIS_FOR          = 7.      : Type for buffer
DBG$K_CIS_SCREEN       = 8.      : Type for Screen Display
DBG$K_CIS_MAXIMUM      = 8.      : Highest type CIS

```

```

: The definitions below should eventually be eliminated when all uses of
: them are converted to the DBG$K_ form.

```

```

LITERAL
CIS_DBG$INPUT    = 0.      : Type DBG$INPUT
CIS_RAB          = 1.      : Type RAB
CIS_INPBUF       = 2.      : Type input buffer
CIS_ACBUF        = 3.      : Type action buffer
CIS_REPEAT       = 4.      : Type repeat buffer
CIS_WHILE        = 5.      : Type while buffer
CIS_IF           = 6.      : Type if buffer
CIS_FOR          = 7.      : Type for buffer
CIS_SCREEN       = 8.      : Type for Screen Display

```


C O N T R O L F L A G S

Control flags are global DEBUG status flags used to control the overall flow of DEBUG execution. They represent the major state of DEBUG.

MACRO DBG\$CONTROL_FLAGS = BITVECTOR[]%;

L I T E R A L

DBG\$V_CONTROL_TDBG	= 0.	! Set if this is a testable DEBUG
DBG\$V_CONTROL_SDBG	= 1.	! Set if this is SUPERDEBUG
DBG\$V_CONTROL_KDBG	= 2.	! Reserved for future development
DBG\$V_CONTROL_URUN	= 3.	! Set if user program has been run to stop DEBUG allocating memory
DBG\$V_CONTROL_EXIT	= 4.	! Set if DEBUG is about to EXIT
DBG\$V_CONTROL_FAIL	= 5.	! Set by DEBUG internal errors
DBG\$V_CONTROL_DONE	= 6.	! Set if user execution complete
DBG\$V_CONTROL_ALLOCATE	= 7.	! Set if OK to allocate more memory (e.g., SET MODULE/ALLOCATE)
DBG\$V_CONTROL_USER	= 8.	! Set if user program is running
DBG\$V_CONTROL_STOP	= 9.	! Set by ^Y,DEBUG sequence to abort processing DEBUG command files
DBG\$V_CONTROL_TBIT	= 10.	! Used by DBGEVENT and DBGSTART to control AST/TBIT interactions.
DBG\$V_CONTROL_SCREEN	= 11.	! Set if screen displays must be updated because user program has run
DBG\$V_CONTROL_VERSION_4	= 12;	! Set if VMS 3B or 4.0 is running

! We also define a macro \$ABORT_ON_CONTROL_Y. The purpose of this macro is to allow premature termination of commands which can take a large amount of time to complete but to allow the individual command routines to make sure that they leave all DEBUGs internal data structures in a consistent state. Execution of this macro will have no effect if the Control-Y flag is clear, but will SIGNAL back to DEBUG command level if the flag is set

MACRO \$ABORT_ON_CONTROL_Y =

BEGIN

EXTERNAL DBG\$GV_CONTROL : BITVECTOR[];

IF .DBG\$GV_CONTROL[DBG\$V_CONTROL_STOP] THEN SIGNAL(DBG\$_ABORTED);

END%;


```
! LITERAL
DEFINE_LOWEST      = 1,      ! Lowest value of DEFINE codes
DEFINE_ADDRESS     = 1,
DEFINE_COMMAND     = 2,
DEFINE_PROCEDURE   = 3,
DEFINE_STRING      = 4,
DEFINE_VALUE       = 5,
DEFINE_PARAMETER   = 6,      ! This is used to temporarily store strings
                              ! representing actual parameters, until
                              ! they are bound to formals by the
                              ! DECLARE command.

DEFINE_KEY         = 7,
DEFINE_HIGHEST     = 7;      ! Highest value of DEFINE codes
```

```
! Macros for declaration of DEFINE entries and headers.
```

```
! MACRO
DEFINE$ENTRY =
  BLOCK [DBG$K_DEFINE_ENTRY_SIZE_B, BYTE] FIELD (DEFINE_ENTRY_FIELDS)%,
DEFINE$HEADER =
  BLOCK [DBG$K_DEFINE_HEADER_SIZE_B, BYTE] FIELD (DEFINE_HEADER_FIELDS)%;
```

E R R O R S T A T U S C O D E S

Debugger status codes are defined in the MDL file DBGMES.MDL--new ones should therefore be entered in that file. All that is found here are definitions relating to components of status codes.

Define the possible values of the status severity field.

LITERAL SYSSK_INFO = 3; ! Informational status severity


```

4 : [6] EVENT$V_OVRD_SOURCE
5 : [6] EVENT$V_STEP_SOURCE
6 : [6] EVENT$V_OVRD_NOSYSTEM
7 : [6] EVENT$V_STEP_NOSYSTEM
8 : [6] EVENT$V_OVRD_NOSILENT
9 : [6] EVENT$V_STEP_NOSILENT

```

```

A : [6] EVENT$V_STEP_BPT
B : [6] EVENT$V_STEP_TICK
C : [6] EVENT$V_STEP_LINE_OK

```

```

D : [6] EVENT$V_BREAK_ADDRESS

```

A pointer to an Event Table Entry is declared as follows:

```

EVENTPTR: REF EVENT$EVENT_DESCRIPTOR

```

Define the fields of the Event Table Entry. Also declare the maximum size of an Event Table Entry and define the declaration macro.

```

FIELD EVENT$EVENT_DESCRIPTOR_FIELDS =

```

```

SET

```

```

EVENT$A_EVENT_DESCRIPTOR = [0,A_], ! Event Descriptor address

```

```

EVENT$SL_CMD_FLINK = [0,L_], ! Command queue forward link
EVENT$SL_CMD_BLINK = [1,L_], ! Command queue backward link

```

```

EVENT$SL_EXC_FLINK = [2,L_], ! Exception queue forward link
EVENT$SL_EXC_BLINK = [3,L_], ! Exception queue backward link

```

```

EVENT$SL_EXCEPTION = [4,L_], ! Exception code we expect to see

```

```

EVENT$SB_CMD_TYPE = [5,B0_], ! Command type
EVENT$SB_CMD_KIND = [5,B1_], ! Command kind
EVENT$SB_SUB_KIND = [5,B2_], ! Command sub-kind

```

```

EVENT$SB_CMD_FLAGS = [5,B3_], ! Command flags

```

```

EVENT$V_SILENT = [5,V3_(0)], ! Flag set if /SILENT qualifier set
EVENT$V_THREADED = [5,V3_(1)], ! Flag set if this is threaded code
EVENT$V_ONCE_ONLY = [5,V3_(2)], ! Flag set for temporary event entry
EVENT$V_HAS_VAL_DSCR = [5,V3_(3)], ! Flag set if there is an Event

```

```

Value Descriptor
EVENT$V_ROUTINE = [5,V3_(4)], ! Flag set if ROUTINE EventPoint
EVENT$V_RET_AT_PC = [5,V3_(5)], ! Flag set if SET BREAK/RET (or TRACE)
address is at current PC

```

```

EVENT$V_DELETED = [5,V3_(7)], ! Flag set if this entry is deleted

```

```

EVENT$SL_STEP_FLAGS = [6,L_], ! 'Stepping' flags

```

```

EVENT$V_OVRD_LINE = [6,V_(0)], ! Flag set if override step/line
EVENT$V_STEP_LINE = [6,V_(1)], ! Flag set if step/line
EVENT$V_OVRD_OVER = [6,V_(2)], ! Flag set if override step/over
EVENT$V_STEP_OVER = [6,V_(3)], ! Flag set if step/over
EVENT$V_OVRD_SOURCE = [6,V_(4)], ! Flag set if override step/source

```

```

EVENTSV_STEP_SOURCE = [6,V_( 5)], ! Flag set if step/source
EVENTSV_OVRD_NOSYSTEM = [6,V_( 6)], ! Flag set if override step/nosystem
EVENTSV_STEP_NOSYSTEM = [6,V_( 7)], ! Flag set if step/nosystem
EVENTSV_OVRD_NOSILENT = [6,V_( 8)], ! Flag set if override step/nosilent
EVENTSV_STEP_NOSILENT = [6,V_( 9)], ! Flag set if step/nosilent

EVENTSV_STEP_BPT = [6,V_(10)], ! Flag set if step over BPT active
EVENTSV_STEP_TICK = [6,V_(11)], ! Flag set if new step
EVENTSV_STEP_NOLINE = [6,V_(12)], ! Flag set if line not found

EVENTSV_BREAK_ADDRESS = [6,V_(13)], ! Flag set if the prolog address is
                                     ! used
EVENTSB_SAVED_SUB_KIND = [6,B3_], ! Saved value of SUB_KIND

EVENTSL_AFTER_COUNT = [7,L_], ! After count
EVENTSL_STEP_COUNT = [7,L_], ! Step count
EVENTSL_SSV_COUNT = [7,L_], ! Address of a count which keeps
                             ! track of the number of SSV calls
                             ! originated from one SSV

EVENTSL_WHEN = [8,L_], ! WHEN expression pointer
EVENTSL_USERS_PC = [8,L_], ! User's PC for STEP/RETURN

EVENTSL_DO = [9,L_], ! DO command list pointer
EVENTSL_USERS_FP = [9,L_], ! User's FP for STEP/RETURN and
                             ! [BREAK|TRACE]/RETURN skips

EVENTSL_PRIMARY = [10,L_], ! Primary pointer
EVENTSL_OPCODE_LIST = [10,L_], ! Instruction opcode list

EVENTSL_ADDRESS = [11,L_], ! Byte address
EVENTSL_FP_RET_CNT = [11,L_], ! Address of a count which keeps
                              ! track of RET for SSV return levels

EVENTSL_VALDESC = [12,L_], ! Value descriptor pointer
EVENTSB_OPCODE = [12,B0_], ! Opcode byte
EVENTSL_THREAD = [12,L_], ! Thread
EVENTSA_VMSDESC = [13,A_], ! VMS Descriptor
EVENTSL_STEP_LO_PC = [13,L_], ! Step line lo PC
EVENTSL_STEP_HI_PC = [14,L_], ! Step line hi PC
EVENTSL_CALL_FRAME = [15,L_] ! Pointer to Saved value from FP
TES;

```

LITERAL

```
EVENTSK_EVENT_DESCRIPTOR_SIZE = 16; ! Descriptor size in longwords
```

MACRO

```

EVENTSEVENT_DESCRIPTOR =
  BLOCK [EVENTSK_EVENT_DESCRIPTOR_SIZE]
  FIELD (EVENTSEVENT_DESCRIPTOR_FIELDS) %;

```


! Macro to assign an increasing value to each member of a <name> list,
! starting with <constant>.

```
MACRO
  ENUMERATE (CONSTANT, NAME) [] =
    NAME = CONSTANT
    %IF %NULL (%REMAINING)
    %THEN %EXITMACRO
    %FI
    , ENUMERATE (CONSTANT + 1, %REMAINING) %;
```

! The following are the allowed values of the EVENT\$B_CMD_TYPE field.

```
LITERAL
  ENUMERATE(0,
    EVENT$K_TYPE_BREAK,      ! Starting with 0...
    EVENT$K_TYPE_TRACE,      ! Break
    EVENT$K_TYPE_WATCH,      ! Trace
    EVENT$K_TYPE_STEPS,      ! Watch
    EVENT$K_TYPE_SKIPS,      ! Steps
    EVENT$K_TYPE_IGNORE,     ! Skips
    EVENT$K_TYPE_SSINT,      ! Ignored activating an event
  );                          ! System service Intercept Event
```

! The following are the allowed values of the EVENT\$B_CMD_KIND field.

```
LITERAL
  ENUMERATE(0,
    EVENT$K_KIND_ACC,        ! Starting with 0...
    EVENT$K_KIND_INS,        ! Access (/READ, etc.)
    EVENT$K_KIND_EXC,        ! Instruction
    EVENT$K_KIND_SKP,        ! Exception
  );                          ! Skips
```

! The following are the allowed values of the EVENT\$B_SUB_KIND field.

```
LITERAL
  ENUMERATE(0,
    EVENT$K_ACC_READ,        ! Starting with 0...
    EVENT$K_ACC_WRIT,        ! Access Read
    EVENT$K_ACC_MDFY,        ! Access Write
    EVENT$K_ACC_EXEC,        ! Access Modify
    EVENT$K_ACC_RTRN,        ! Access Execute
    EVENT$K_INS_CALL,        ! Access Return
    EVENT$K_INS_BRAN,        ! Instruction Call
    EVENT$K_INS_LINE,        ! Instruction Branch
    EVENT$K_INS_EVR,        ! Instruction Line
    EVENT$K_INS_USER,        ! Instruction Every
    EVENT$K_SKP_READ,        ! Instruction User
    EVENT$K_SKP_WRIT,        ! Skipping Read
    EVENT$K_SKP_MDFY,        ! Skipping Write
  );                          ! Skipping Modify
```

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EVENTSK_SKP_EXEC,
EVENTSK_SKP_RTRN,
EVENTSK_EXC_EXC
);

! Skipping Execute
! Skipping Return

EVENT TABLE ENTRY FOR TREADED CODE EVENT POINTS

This form of the Event Table Entry is used for breakpoints and tracepoints in threaded code. Threaded code is only used by COBOL-74 and is thus obsolescent, but we still support it.

3 3 2 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1
1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0

0	EVENTSL_CMD_FLINK
1	EVENTSL_CMD_BLINK
2	EVENTSL_EXC_FLINK
3	EVENTSL_EXC_BLINK
4	EVENTSL_EXCEPTION
5	7: - 13:2:1:0: B_SUB_KIND : B_CMD_KIND : B_CMD_TYPE
6	B_SAVED_SUB_KIND: - unused - :C:B:A:9:8:7:6:5:4:3:2:1:0:
7	EVENTSL_AFTER_COUNT
8	EVENTSL_WHEN
9	EVENTSL_DO
10	EVENTSL_PRIMARY
11	EVENTSL_ADDRESS
12	EVENTSL_THREAD
13	EVENTSA_VMSDESC
14	
15	

EVENT TABLE ENTRY FOR /READ, /WRITE, /MODIFY AND WATCHPOINTS

This form of the Event Table Entry is used in all cases when data locations are watched for read accesses, write accesses, or modifying accesses. It is thus used both for SET BREAK/MODIFY, for example, and for the SET WATCH command.

3 3 2 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1 1 1
1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0

0	EVENT\$L_CMD_FLINK
1	EVENT\$L_CMD_BLINK
2	EVENT\$L_EXC_FLINK
3	EVENT\$L_EXC_BLINK
4	EVENT\$L_EXCEPTION
5	7! - 13!2!1!0! B_SUB_KIND B_CMD_KIND B_CMD_TYPE
6	B_SAVED_SUB_KIND! - unused - 1C!B!A!9!8!7!6!5!4!3!2!1!0!
7	EVENT\$L_AFTER_COUNT
8	EVENT\$L_WHEN
9	EVENT\$L_DO
10	EVENT\$L_PRIMARY
11	EVENT\$L_ADDRESS
12	EVENT\$L_VALDESC
13	EVENT\$A_VMSDESC
14	
15	

EVENT TABLE ENTRY FOR STEP

This form of Event Table Entry is used for the STEP command.

3 3 2 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1
1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0

0	EVENT\$L_CMD_FLINK
1	EVENT\$L_CMD_BLINK
2	EVENT\$L_EXC_FLINK
3	EVENT\$L_EXC_BLINK
4	EVENT\$L_EXCEPTION
5	7: - 13:2:1:0: B_SUB_KIND B_CMD_KIND B_CMD_TYPE
6	B_SAVED_SUB_KIND: - unused - 1C:1B:1A:9:8:7:6:5:4:3:2:1:0
7	EVENT\$L_STEP_COUNT
8	EVENT\$L_USERS_PC
9	EVENT\$L_USERS_FP
10	EVENT\$L_OPCODE_LIST
11	EVENT\$L_ADDRESS
12	Unused EVENT\$B_OPCODE
13	EVENT\$L_STEP_LO_PC
14	EVENT\$L_STEP_HI_PC
15	EVENT\$L_CALL_FRAME

EVENT STEPPING DESCRIPTOR

This structure defines the current step type(s).

```

      3 3 2 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1
      1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0
0  +-----+-----+-----+-----+-----+
   |XXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXX|^!%$#@|*|
   +-----+-----+-----+-----+-----+
1  |                                EVENT$SL_STEPPING_OP_LIST
   +-----+-----+-----+-----+-----+

```

```

* : [0] EVENT$B_STEPPING_KIND
@ : [0] EVENT$V_STEPPING_LINE
# : [0] EVENT$V_STEPPING_OVER
$ : [0] EVENT$V_STEPPING_SOURCE
% : [0] EVENT$V_STEPPING_NOSYSTEM
^ : [0] EVENT$V_STEPPING_NOSILENT

```

A pointer to an Step Descriptor Entry is declared as follows:

```
STEPPINGPTR: REF EVENT$STEPPING_DESCRIPTOR
```

Define the fields of the Step Descriptor Entry. Also declare the maximum size of an Step Descriptor Entry and define the declaration macro.

```
FIELD EVENT$STEPPING_DESC_FIELDS =
```

```

SET
EVENT$A_STEPPING_DESCRIPTOR = [0,A ], ! Event Descriptor address
EVENT$B_STEPPING_KIND = [0,B0 ], ! Step kind (as EVENT$B_SUB_KIND)
EVENT$V_STEPPING_LINE = [0,V (8)], ! Stepping /LINE
EVENT$V_STEPPING_OVER = [0,V (9)], ! Stepping /OVER
EVENT$V_STEPPING_SOURCE = [0,V (10)], ! Stepping /SOURCE
EVENT$V_STEPPING_NOSYSTEM = [0,V (11)], ! Stepping /NOSYSTEM
EVENT$V_STEPPING_NOSILENT = [0,V (12)], ! Stepping /NOSILENT
EVENT$SL_STEPPING_OP_LIST = [1,L_] ! /CALL, /BRANCH, etc.
TES;

```

```
LITERAL
```

```
EVENT$K_STEPPING_DESC_SIZE = 2; ! Descriptor size in longwords
```

```
MACRO
```

```

EVENT$STEPPING_DESCRIPTOR =
  BLOCK [EVENT$K_STEPPING_DESC_SIZE]
  FIELD (EVENT$STEPPING_DESC_FIELDS) %;

```


EVENT DO LIST DESCRIPTOR

An Event DO List Descriptor is created for each event-point which has an associated DO-action. This descriptor points to the text of the DO-action itself and is thus what allows the DO-action to be accessed when the event-point is taken. This is the format of the Event DO List Descriptor:

```

3 3 2 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1
1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0

```

0	EVENT\$DO_LIST_FLINK
1	EVENT\$DO_LIST_BLINK
2	EVENT\$DO_LIST_COUNT
3	EVENT\$DO_LIST_POINT

A pointer to an Event DO List Descriptor is declared as follows:

```
DOPTR: REF EVENT$DO_LIST_DESCRIPTOR
```

Define the fields of the Event DO List Descriptor. Also declare the descriptor size and define the declaration macro.

```

FIELD  EVENT$DO_LIST_DESCRIPTOR_FIELDS =
SET
EVENT$A_DO_LIST_DESCRIPTOR = [0,A_], ! DO List Descriptor address

EVENT$L_DO_LIST_FLINK      = [0,L_], ! DO queue forward
EVENT$L_DO_LIST_BLINK      = [1,L_], ! and backward links

EVENT$L_DO_LIST_COUNT      = [2,L_], ! DO reference count

EVENT$L_DO_LIST_POINT      = [3,L_] ! DO list pointer
TES;

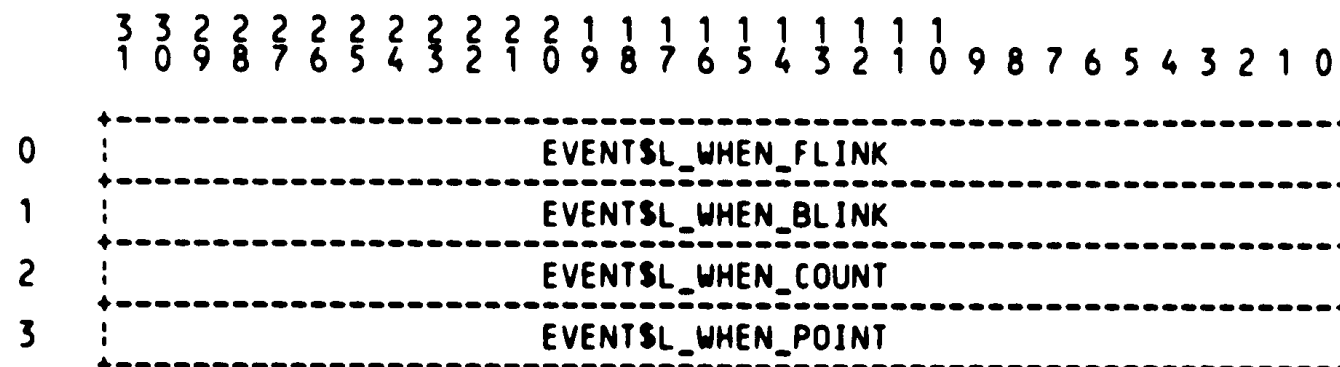
LITERAL
EVENT$K_DO_LIST_DESCRIPTOR_SIZE = 4; ! Descriptor size (longwords)

MACRO
EVENT$DO_LIST_DESCRIPTOR =
BLOCK [EVENT$K_DO_LIST_DESCRIPTOR_SIZE]
FIELD (EVENT$DO_LIST_DESCRIPTOR_FIELDS) %;

```

EVENT WHEN DESCRIPTOR

An Event WHEN Descriptor is created each time a SET BREAK, SET TRACE, or SET WATCH statement has a WHEN clause. The Event WHEN Descriptor points to the location of the boolean WHEN-expression and allows this expression to be evaluated each time the X-point is reached. This is the format of the Event WHEN Descriptor:



A pointer to an Event WHEN Descriptor is declared as follows:

```
WHENPTR: REF EVENT$WHEN_DESCRIPTOR
```

Define the fields of the Event WHEN Descriptor, its size, and its declaration macro.

```
FIELD EVENT$WHEN_DESCRIPTOR_FIELDS =
```

```
SET
EVENT$A_WHEN_DESCRIPTOR = [0,A_], ! WHEN Descriptor address
EVENT$SL_WHEN_FLINK      = [0,L_], ! WHEN queue forward link
EVENT$SL_WHEN_BLINK      = [1,L_], ! WHEN queue backward link
EVENT$SL_WHEN_COUNT      = [2,L_], ! WHEN reference count
EVENT$SL_WHEN_POINT      = [3,L_], ! WHEN list pointer
TES;
```

```
LITERAL
```

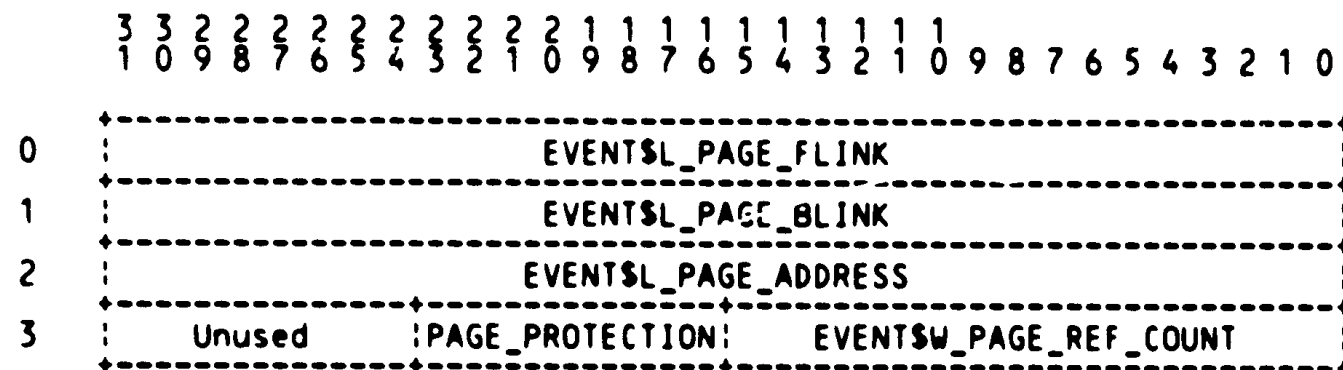
```
EVENT$K_WHEN_DESCRIPTOR_SIZE = 4; ! Descriptor size (longwords)
```

```
MACRO
```

```
EVENT$WHEN_DESCRIPTOR =
BLOCK [EVENT$K_WHEN_DESCRIPTOR_SIZE]
FIELD (EVENT$WHEN_DESCRIPTOR_FIELDS) %;
```


EVENT PAGE DESCRIPTOR

An Event Page Descriptor is created for each page which is write-protected to watch-point one or more variables on that page. Event Page Descriptors are linked together on a doubly linked list which is pointed to by variable XXXXXX. This is the format of an Event Page Descriptor:



A pointer to an Event Page Descriptor is declared as follows:

```
PAGEDESC: REF EVENT$PAGE_DESCRIPTOR
```

Define the fields of the Event Page Descriptor. Also define the descriptor size and the declaration macro.

```
FIELD EVENT$PAGE_DESCRIPTOR_FIELDS =
```

```
SET
EVENT$PAGE_DESCRIPTOR = [0,A_] : Page Descriptor address
EVENT$PAGE_FLINK      = [0,L_] : Page queue forward link
EVENT$PAGE_BLINK      = [1,L_] : Page queue backward link
EVENT$PAGE_ADDRESS     = [2,L_] : Page byte address
EVENT$PAGE_REF_COUNT   = [3,W0_] : Page reference count
EVENT$PAGE_PROTECTION  = [3,B2_] : Page old protection
TES;
```

```
LITERAL
```

```
EVENT$PAGE_DESCRIPTOR_SIZE = 4; ! Descriptor size (longwords)
```

```
MACRO
```

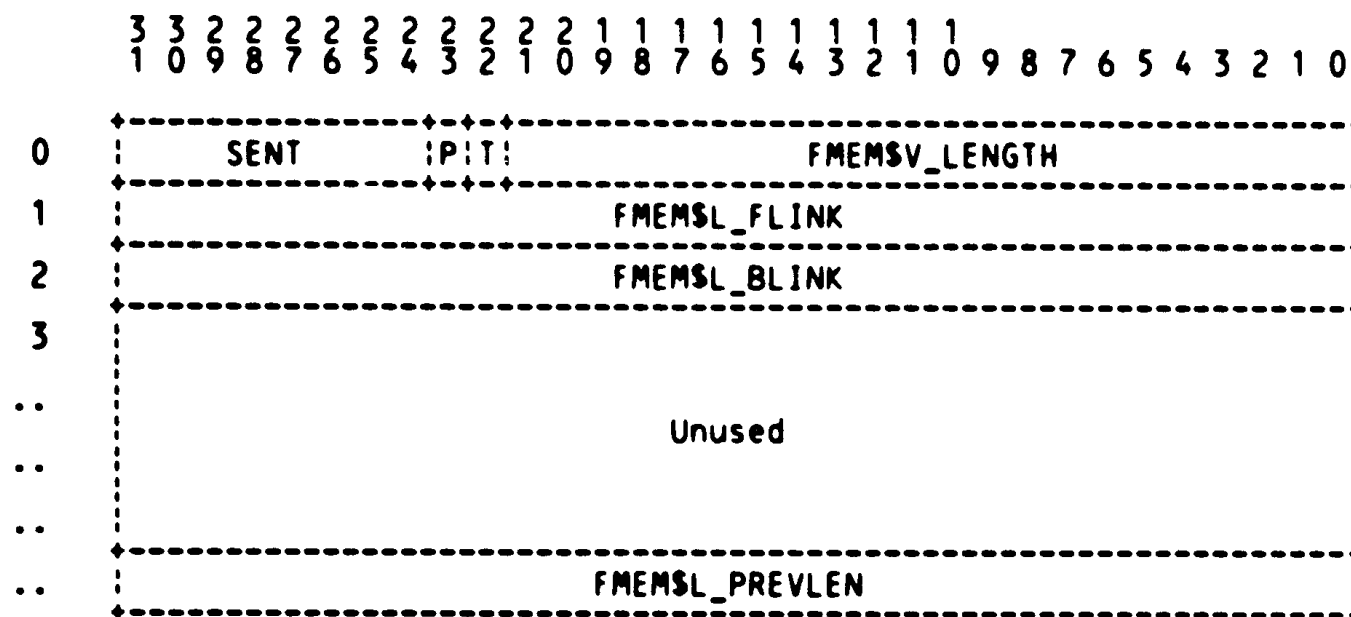
```
EVENT$PAGE_DESCRIPTOR =
BLOCK [EVENT$PAGE_DESCRIPTOR_SIZE]
FIELD (EVENT$PAGE_DESCRIPTOR_FIELDS) %;
```

FREE MEMORY MANAGEMENT

The Free Memory Manager manages the Debugger's free memory pool. Its routines are thus called throughout the Debugger to allocate and de-allocate memory blocks used for RST entries, Static Address Table entries, and numerous other kinds of descriptors and records. The Free Memory Manager has two kinds of blocks in its memory pool: free blocks and allocated blocks. The formats of both kinds of blocks and the overall structure of the memory pool is illustrated below.

FORMAT OF A FREE MEMORY BLOCK

The format of a free memory block is shown here:

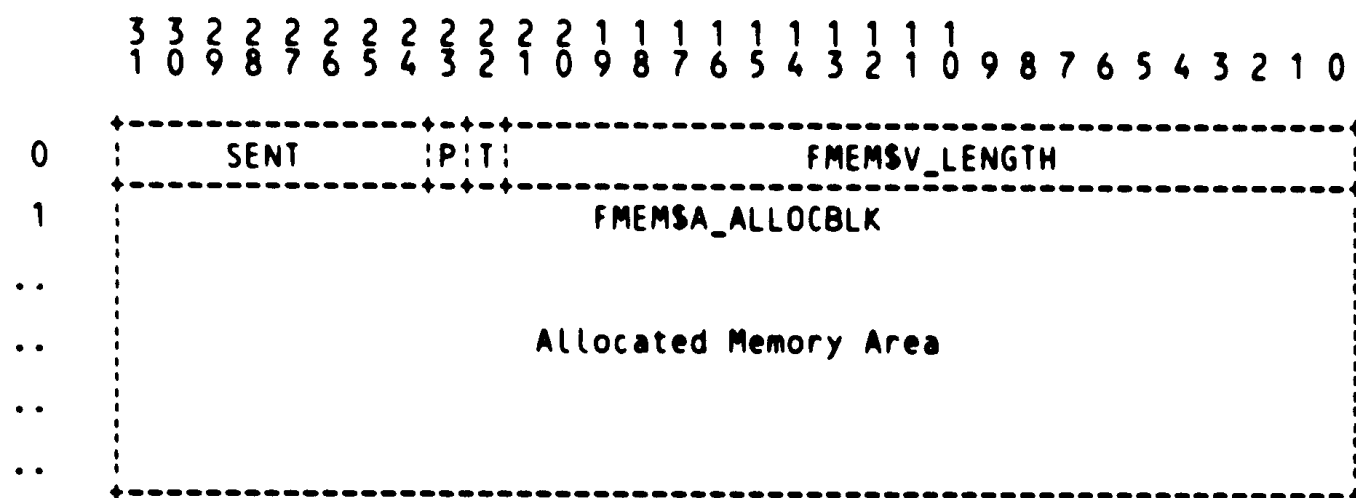


Here SENT represents FMESB_SENTINEL, P is FMESV_PREVALLOC, and T is FMESV_THISALLOC as defined on the next page.

A free block has its length (in longwords) stored both at the beginning of the block (FMESV_LENGTH) and at the end (FMESL_PREVLEN). The minimum length of a free block is four longwords. Free blocks are chained together on a doubly linked free-list. This list has a list head which is never allocated--the list is thus never empty. The address of the free-list list head is always stored in the OWN variable DBG\$FREE_LIST.

FORMAT OF AN ALLOCATED MEMORY BLOCK

The format of an allocated block is as follows:



Here SENT represents FMEMSB_SENTINEL, P is FMEMSV_PREVALLOC, and T is FMEMSV_THISALLOC as defined below. In an allocated block, the address returned to the requestor is represented by FMEMSA_ALLOCBLK.

A pointer to a memory block (free or allocated) is declared as follows:

BLKPTR: REF FMEMSBLOCK

This declaration is used in the Free Memory Manager only.

Define the fields in the two kinds of memory blocks.

FIELD FMEM\$FLD_DEF =

SET		
FMEMSV_LENGTH	= [0, V_(0,22)],	! The length of the block in longwords
FMEMSV_THISALLOC	= [0, V_(22)],	Set to TRUE if this block is allocated
FMEMSV_PREVALLOC	= [0, V_(23)],	Set to TRUE if the previous block is allocated (block at smaller addr)
FMEMSB_SENTINEL	= [0, B3_],	Sentinel value--always FMEMSK_SENTINEL
FMEMSL_FLINK	= [1, L_],	Forward link on free list
FMEMSL_BLINK	= [2, L_],	Backward link on free list
FMEMSL_PREVLEN	= [-1, L_],	Length of this free block relative to the start of the next block
FMEMSA_ALLOCBLK	= [1, A_],	Address returned to caller requesting a new memory block
FMEMSA_HEADER	= [-1, A_],	Address of block's header relative to the FMEMSA_ALLOCBLK location
TES;		

! Declare the sentinel value and the declaration macro.

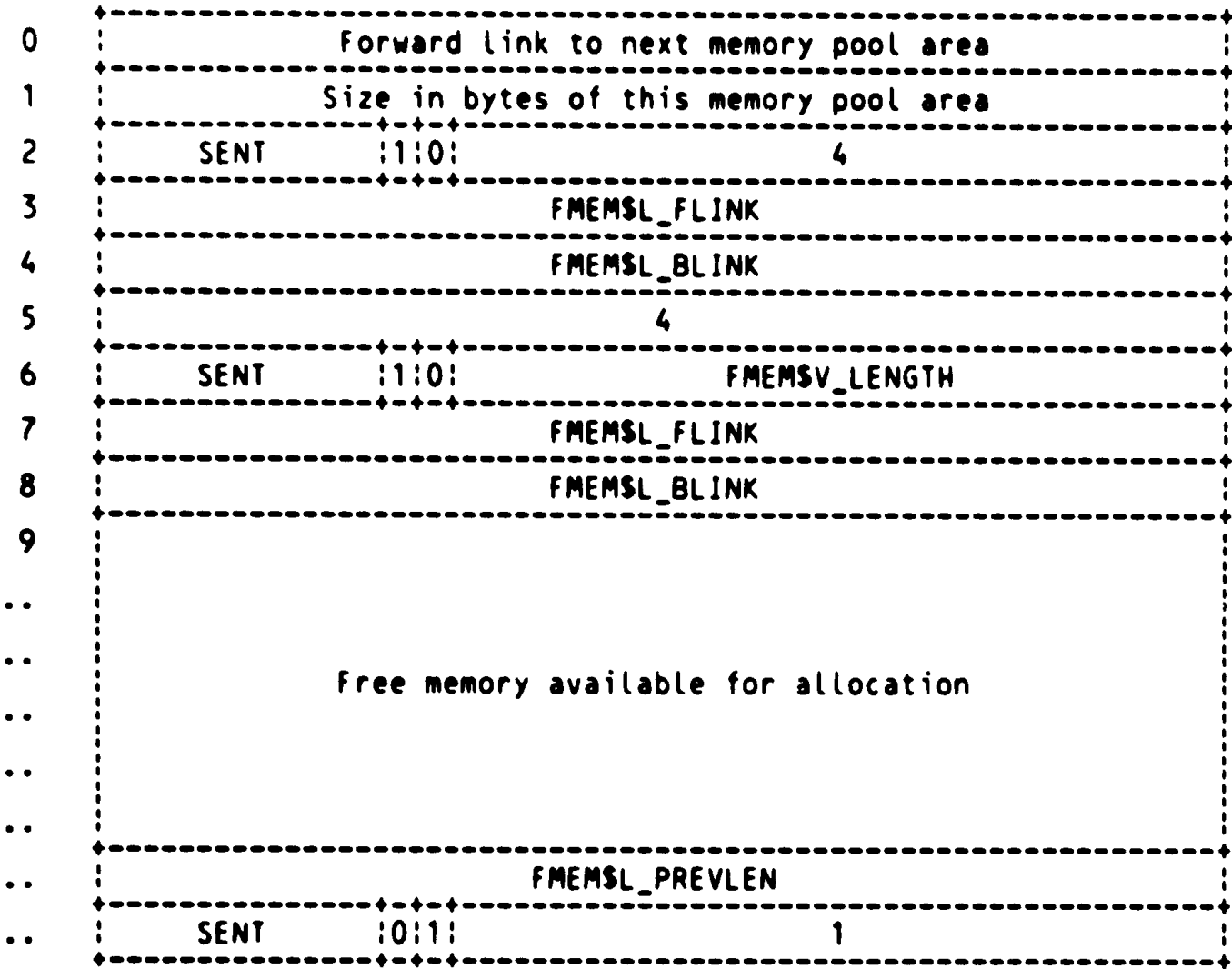
LITERAL FMEMSK_SENTINEL = %X'B2'; ! Magic sentinel value--used for error
! checking only

MACRO FMEMSBLOCK = BLOCK[,LONG] FIELD(FMEM\$FLD_DEF) %;

OVERALL STRUCTURE OF MEMORY POOL

The memory pool consists of one or more memory pool areas. The first such area is initialized to have the format shown here:

3 3 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1 1
1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0



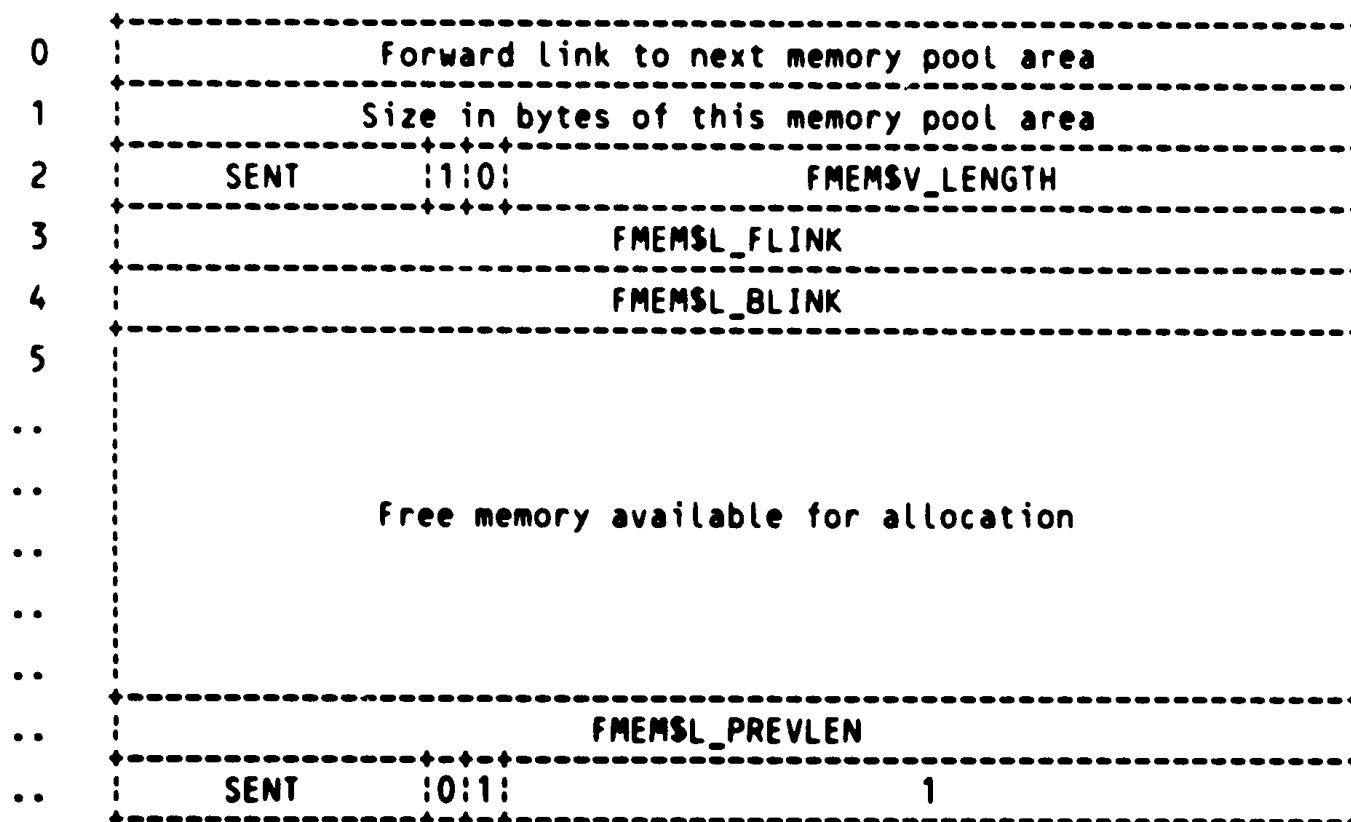
Here longword 0 contains a pointer to the next memory pool area (for this initial area the pointer is always zero) and longword 1 contains the size of this area in bytes. These two fields are present for internal error checking purposes only.

Longwords 2 - 5 contain the free-list list head. This is a free block which is always at the start of the doubly linked free list and which

is never coalesced with any other free block--the code thus never has to worry about a completely empty free list. The very last longword in the memory pool area is the header of an allocated block. It prevents any attempt to coalesce free blocks off the end of the area. The rest of the area (longword 6 to the next to last longword) is initially one big free block from which allocation can take place. Note how the P and T flags (FMEMSV_PREVALLOC and FMEMSV_THISALLOC) are set to prevent free block coalescing off either end of the memory pool area.

Additional memory pool areas may be acquired from LIB\$GET_VM during the Debugger's initialization phase. Each such area is initialized to have the same format as the first one except that the free-list list head is absent--only one is needed for the whole memory pool. The initial format of each additional memory pool area is illustrated here:

3 3 2 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1
1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0



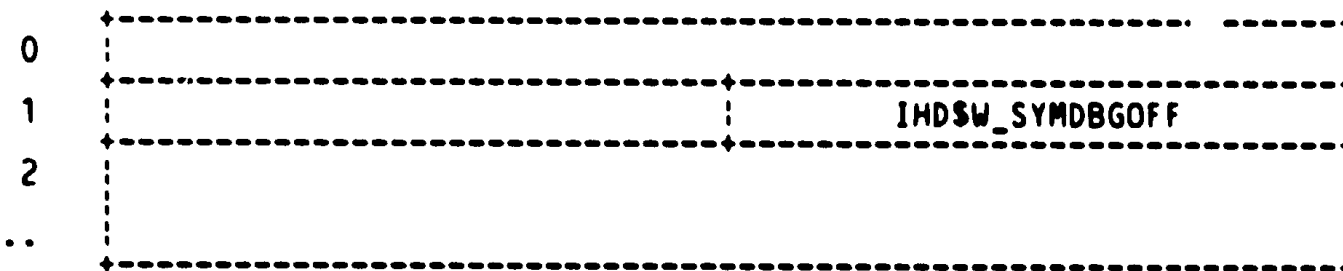
I M A G E F I L E H E A D E R D E F I N I T I O N S

The executable image file header block contains a pointer in a fixed location which points to a small block later in the header which gives the size and location of the Debug Symbol Table (DST) and the Global Symbol Table (GST). The first part of the executable image file header looks as follows:

```

3 3 2 2 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1 1
1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0

```



Here IHDSW_SYMDBGOFF contains the byte offset relative to the start of the header of an Image Header Symbol Table descriptor.

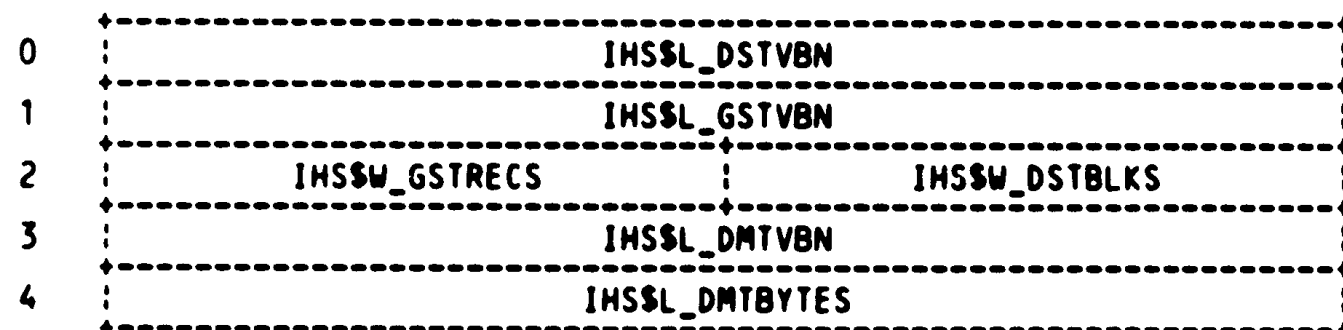
I M A G E H E A D E R S Y M B O L T A B L E D E S C R I P T O R

The Image Header Symbol Table Descriptor (IHS) pointed to by the IHDSW_SYMDBGOFF field in the header has the following format:

```

3 3 2 2 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1 1
1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0

```



Here IHSSW_DSTBLKS and IHSSL_DSTVBN give the size (in blocks) and loca-

tion (Virtual Block Number) of the Debug Symbol Table (DST) within the executable image file. The fields IHSSW GSTRECS and IHSSL GSTVBN give the size (in GST records) and start location (Virtual Block Number) of the Global Symbol Table (GST). Finally, the fields IHSSL DMTBYTES and IHSSL DMTVBN give the size (in bytes) and start location (Virtual Block Number) of the Debug Module Table (DMT). The DMT is described below. These field names are declared by macros in SYSSLIBRARY:LIB.L32.

The symbol IHD\$W SYMDBGOFF is defined in SYSSLIBRARY:LIB.L32 and is therefore not defined here. A pointer to the image file header is declared as follows:

IHDPTR: REF BLOCK[,BYTE]

A pointer to the image header symbol descriptor is declared like this:

IHSPTR: REF IH\$ENTRY

Define the IH\$ENTRY macro.

MACRO

IH\$ENTRY = BLOCK[IH\$K_LENGTH,BYTE] %; ! IH\$K_LENGTH is defined in
! SYSSLIBRARY:LIB.L32

IMAGE FILE DEBUG MODULE TABLE

The Image File Header in an executable image file points to the Image Header Symbol Table descriptor as described above. If bit 5 of field IHSSL_LNKFLAGS in the image header is set, this is a "new" image, i.e. one produced by the VMS V4.0 or later linker, and the IHSSL_DMTVBN and IHSSL_DMTBYTES fields exist in the Image Header Symbol Table descriptor. (If bit 5 is not set, this is an "old" image and those fields do not exist.) If non-zero, IHSSL_DMTVBN gives the Virtual Block Number in the image file of the Debug Module Table (the DMP). IHSSL_DMPBYTES then gives the size of the DMT in bytes. The DMT is only built if the user did a LINK/DEBUG; if he did not, IHSSL_DMTVBN and IHSSL_DMTBYTES are zero.

The Debug Module Table contains one entry per module in the Debug Symbol Table (the DST). This is the format of each such DMT entry:

3 3 2 2 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1 1
1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0

0	DST address of Module Begin DST Record
1	Size in bytes of module's DST
2	Number of PSECTs for this module
3	Start address of first PSECT in module
4	Length of first PSECT in module in bytes
..	
..	(Two longwords per PSECT)
..	
..	Start address of last PSECT in module
..	Length of last PSECT in module in bytes

Longword 0 gives the address relative to the start of the DST of the Module Begin DST Record for this module. Longword 1 gives the size of the DST in bytes for the same module. Longword 2 gives the number of PSECTs in the module (i.e., the number of statically allocated program sections), and this is followed by that number of two-longword pairs which give the start address and length (in bytes) of each such PSECT. Since the number of PSECTs cannot exceed 65K, the upper two bytes of longword 2 are available for future expansion.

The DMT is used during DEBUG initialization to initialize the RST and the Program Static Address Table. Using the DMT is much faster than the alternative procedure, namely reading through the entire DST to pick up the needed information. The information in the DMT entry is enough to build a Module RST Entry for each module in the DST and the PSECT information is used to build the Program SAT. The amount of RST symbol table space needed per module is not computable from the DMP information, but is estimated by multiplying the DST size of each module by an appropriate scale factor.

L E X I C A L S C A N N E R C H A R A C T E R T A B L E

The Lexical Scanner (routine DBGSLEXICAL_SCANNER) operates on a Character Table which gives lexical information about each ASCII character. This information consists of bits and small integer values which characterize the character from the Lexical Scanner's point of view--they indicate whether a given character can be part of an identifier, a numeric constant, an operator, or whatever. The Character Table is thus a vector of 256 elements, indexed by the ASCII character code, where each element contains the information defined here.

A Lexical Scanner Character Table is declared as follows:

CHARTBL: CHRTBL\$TABLE

Define the fields of the Character Table element. Also define the declaration macro.

FIELD CHRTBL\$FLD_DEF =

SET		
CHRTBL\$V_IDENT_START	= [0, V_(0)],	Identifier start character
CHRTBL\$V_IDENT_MIDDLE	= [0, V_(1)],	Identifier middle character
CHRTBL\$V_IDENT_END	= [0, V_(2)],	Identifier end character
CHRTBL\$V_NUMBER_START	= [0, V_(3)],	Number start character
CHRTBL\$V_NUMBER_CLASS	= [0, V_(4,4)],	Numeric constant character class
CHRTBL\$V_ALPHABETIC	= [0, V_(8)],	Alphabetic character (A-Z)
CHRTBL\$V_DIGIT	= [0, V_(9)],	Numeric digit (0-9)
CHRTBL\$V_STRING_QUOTE	= [0, V_(10)],	String quote character
CHRTBL\$V_OPCHAR	= [0, V_(11)],	Operator character
CHRTBL\$V_OPCHAR_PREFIX	= [0, V_(12)],	Operator prefix character
CHRTBL\$V_OPCHAR infix	= [0, V_(13)],	Operator infix character
CHRTBL\$V_OPCHAR_POSTFIX	= [0, V_(14)],	Operator postfix character
CHRTBL\$V_ADDRESS_OP	= [0, V_(15)],	Address Expression operator
CHRTBL\$V_SPECIAL_SYMBOL	= [0, V_(16)],	Special DEBUG symbol (. \ ^)
CHRTBL\$V_SPACE	= [0, V_(17)],	Space or Tab character
CHRTBL\$V_TERMINATOR	= [0, V_(18)],	May start Terminator Token
CHRTBL\$V_SPECIAL_CASE	= [0, V_(19)],	Special case in lex scanner
CHRTBL\$V_WHOLE_ENTRY	= [0, L_]	The whole entry for a char.

TES;

MACRO

CHRTBL\$TABLE = BLOCKVECTOR[256,1, LONG] FIELD(CHRTBL\$FLD_DEF) %;

Define mask values for all the bits in the Character Table element. These values are used to initialize Character Table elements.

LITERAL

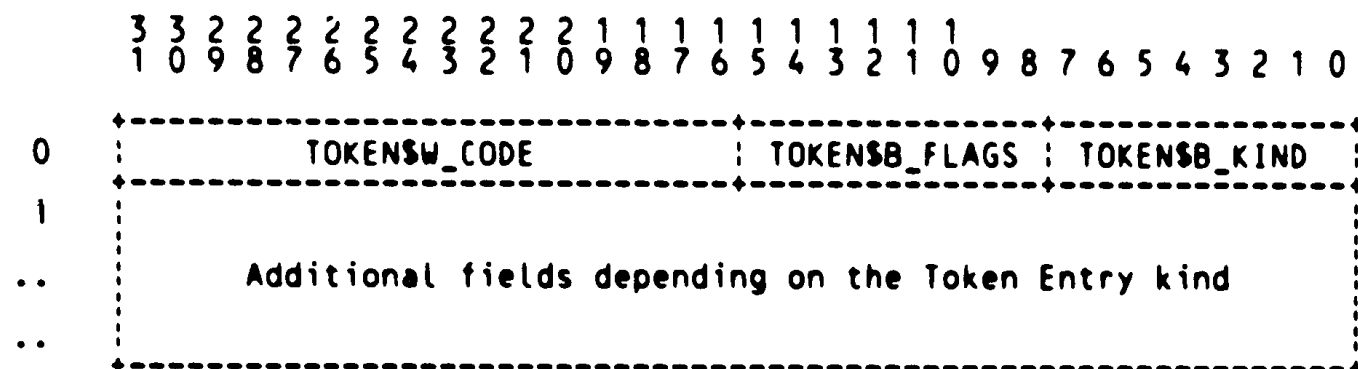
CHRTBL\$M_NOTHING	= 0,	No bits set
CHRTBL\$M_IDENT_START	= 1<0,	Identifier start character

CHRTBLSM_IDENT_MIDDLE	= 1^1,	Identifier middle character
CHRTBLSM_IDENT_END	= 1^2,	Identifier end character
CHRTBLSM_IDENT_ANYWHERE	= 7^0,	Anywhere in an identifier
CHRTBLSM_NUMBER_START	= 1^3,	Number start character
CHRTBLSM_ALPHABETIC	= 1^8,	Alphabetic character (A-Z)
CHRTBLSM_DIGIT	= 1^9,	Numeric digit (0-9)
CHRTBLSM_STRING_QUOTE	= 1^10,	String quote character
CHRTBLSM_OPCHAR	= 1^11,	Operator character
CHRTBLSM_OPCHAR_PREFIX	= 1^12,	Operator prefix character
CHRTBLSM_OPCHAR infix	= 1^13,	Operator infix character
CHRTBLSM_OPCHAR_POSTFIX	= 1^14,	Operator postfix character
CHRTBLSM_ADDRESS_OP	= 1^15,	Address Expression operator
CHRTBLSM_SPECIAL_SYMBOL	= 1^16,	Special DEBUG symbol (. \ ^)
CHRTBLSM_SPACE	= 1^17,	Space or tab character
CHRTBLSM_TERMINATOR	= 1^18,	May start terminator token
CHRTBLSM_SPECIAL_CASE	= 1^19,	Special case in lexical scanner

: Note that the possible values of the CHRTBLSV NUMBER CLASS field are defined
: under the description of the Number Scanner State Table. They are the same
: values as are allowed in the NUMST\$B_CHAR_CLASS field and their names all
: have the form NUMST\$K_CLASS_xxx.

LEXICAL TOKEN ENTRY

A Lexical Token Entry is returned by the Lexical Scanner every time it is called. Operator Token Entries are found in predefined tables specific to each language supported by DEBUG while Operand Token Entries are built as needed by the Lexical Scanner (DBG\$LEXICAL_SCANNER). The format of a Lexical Token Entry depends on the kind of Token Entry it is, but all Lexical Token Entries have this general format:



A pointer to a Lexical Token Entry is defined as follows:

TOKENPTR: REF TOKENSENTRY

Define the fields of the Lexical Token Entry.

FIELD TOKENSFLD_DEF =

SET			
TOKEN\$B_KIND	= [0, B0_],	:	Token kind
TOKEN\$B_FLAGS	= [0, B1_],	:	Token Entry flag bits
TOKEN\$V_PRIMARY	= [0, V1_(0)],	:	Token is Primary Symbol operator (used in Operator entries only)
TOKEN\$V_LEXICAL	= [0, V1_(1)],	:	Token is a lexical operator (used in Operator entries only)
TOKEN\$V_SGNEXT	= [0, V1_(2)],	:	Sign extension S in X<P,S,E> operation (used in Operator Lexical Token for the bit-select operation)
TOKEN\$V_BALANCED_PARENS	= [0, V1_(0)],	:	Token is Terminator Token only if parentheses are balanced (used in Terminator entries only)
TOKEN\$V_MUST_BE_SINGLE	= [0, V1_(1)],	:	Token is Terminator Token only if this character is not followed by it- self (: is term but not ::) (used in Terminator entries only)
TOKEN\$V_ARGUMENT_LIST	= [0, V1_(3)],	:	Argument list follows this Ada tick operator.
TOKEN\$W_CODE	= [0, W1_],	:	Operand or operator code for token
TOKEN\$B_BIF	= [1, B0_],	:	Built-in Function Code

```

TOKEN$B_LENGTH = [ 2, B0 ], ! Length of operand name in bytes
TOKEN$A_NAME    = [ 2, A1 ], ! Operand name in ASCII

TOKEN$B_R_PREC  = [ 1, B0 ], ! Operator's Right Precedence
TOKEN$B_L_PREC  = [ 1, B1 ], ! Operator's Left Precedence
TOKEN$W_SUBCODE = [ 1, W1 ], ! Sub-code field
TOKEN$W_BIT_OFFSET = [ 2, W0 ], ! Offset P in X<P,S,E> operation
TOKEN$W_BIT_LENGTH = [ 2, W1 ], ! Length S in X<P,S,E> operation
TOKEN$B_OPLEN    = [ 3, B0 ], ! Length of operator name in bytes
TOKEN$A_OPNAME   = [ 3, A1 ], ! Operator name in ASCII
TES;

```

LITERAL

```

TOKEN$K_ENTSIZE = 2, ! Size of fixed part of an Operand
                  ! Lexical Token Entry
TOKEN$K_ENTSIZE_OPERATOR = 3; ! Size of fixed part of an Operator
                  ! Lexical Token Entry

```

MACRO

```

TOKEN$ENTRY = BLOCK[,LONG] FIELD(TOKEN$FLD_DEF) %; ! Declaration macro

```

```

! Define the valid values of the TOKEN$B_KIND field. This field indicates
! what kind of token this is.

```

LITERAL

```

TOKEN$K_OPERAND      = 1, ! Operand
TOKEN$K_PREFIX_OP    = 2, ! Prefix Operator
TOKEN$K infix_OP     = 3, ! Infix Operator
TOKEN$K_POSTFIX_OP   = 4; ! Postfix Operator

```

```

! Define mask values for the bits in the TOKEN$B_FLAGS field.

```

LITERAL

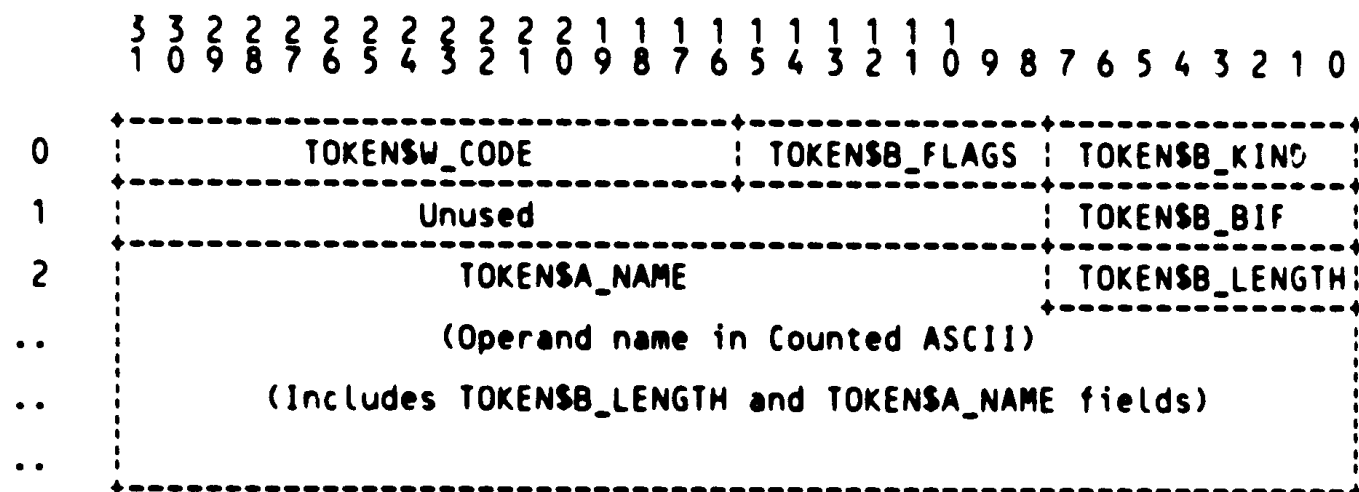
```

TOKEN$M_PRIMARY      = 1^0, ! Token is Primary Symbol operator
TOKEN$M_LEXICAL      = 1^1, ! Token is lexical operator
TOKEN$M_BALANCED_PARENS = 1^0, ! Terminator requires balanced parens
TOKEN$M_MUST_BE_SINGLE = 1^1; ! Terminator character must be single

```


OPERAND LEXICAL TOKEN ENTRY

Operands are identifiers, numeric constants, string or character constants, or special constructs such as %LINE nnn which generate identifiers. The Lexical Token Entry for an operand has the following format:



The TOKENSB_KIND field of an Operand Lexical Token Entry always has the value TOKENSK_OPERAND.

Define the valid values of the TOKENSW_CODE field for operands. This field says what kind of operand this Token Entry represents.

LITERAL

TOKENSK_MIN_OPERAND	= 1.	---	Minimum value for TOKENSW_CODE
TOKENSK_IDENTIFIER	= 1.		Identifier
TOKENSK_STRING	= 2.		Character String Constant
TOKENSK_CHARACTER	= 3.		Single Character Constant
TOKENSK_INTEGER	= 4.		Integer Constant
TOKENSK_HEX_INTEGER	= 5.		Hexadecimal Integer Constant
TOKENSK_FLOATING	= 6.		Floating Point Constant
TOKENSK_EXP_E_FLOAT	= 7.		Floating Constant with E Exponent
TOKENSK_EXP_D_FLOAT	= 8.		Floating Constant with D Exponent
TOKENSK_EXP_Q_FLOAT	= 9.		Floating Constant with Q Exponent
TOKENSK_BIN_INTEGER	= 10.		Binary Constant
TOKENSK_OCT_INTEGER	= 11.		Octal Constant
TOKENSK_COMPLEMENT	= 12.		1's Complement
TOKENSK_BIT_STRING	= 13.		Bit String Constant (as in PL/I)
TOKENSK_PACK_DECIMAL	= 14.		Packed Decimal Constant
TOKENSK_EXP_G_FLOAT	= 15.		Floating Constant with G Exponent
TOKENSK_BUILTIN_FUNCTION	= 16.		Built-in Function
TOKENSK_MAX_OPERAND	= 16.	---	Maximum value for TOKENSW_CODE

OPERATOR LEXICAL TOKEN ENTRY

The Lexical Token Entry for an operator contains additional information not needed for operands, namely the left and right precedences of the operator. This is the format of an Operator Lexical Token Entry:

3 3 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1 1
1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0

0	TOKEN\$W_CODE	TOKEN\$B_FLAGS	TOKEN\$B_KIND
1	TOKEN\$W_SUBCODE	TOKEN\$B_L_PREC	TOKEN\$B_R_PREC
2	TOKEN\$W_BIT_LENGTH	TOKEN\$W_BIT_OFFSET	
3	TOKEN\$A_OPNAME	TOKEN\$B_OPLEN	
..	(Operator name string in Counted ASCII)		
..	(Includes TOKEN\$B_OPLEN and TOKEN\$A_OPNAME fields)		
..			

The TOKEN\$B_KIND field of an Operator Lexical Token Entry always has one of the three values TOKEN\$K_PREFIX_OP, TOKEN\$K infix_OP, or TOKEN\$K postfix_OP.

Define some size literals used for allocating memory.

```
LITERAL
  TOKEN$K_FIXED_SIZE_BYTE = 13,
  TOKEN$K_FIXED_SIZE_LONG = 4;
```

Fixed size (not including the name) of an operator token. Actual size in bytes is this value plus the OPLEN field.

Fixed size (not including the name) of an operator token. Actual size in longwords is this value plus OPLEN / 4.

Define the valid values of the TOKEN\$W_CODE field for operators. This field indicates which operation is to be performed. Not all languages allow all operators, but all operators allowed in at least one language are listed in this table. Note that the same operator may be represented differently in different languages--thus "=" and ".EQ." are both considered to be the same

! comparison operator.

LITERAL

TOKENSK_MIN_PRIMARY_OP	= 1,	--- Minimum value of TOKENSW_CODE
TOKENSK_PRIMARY_TERM	= 1,	--- for Primary Symbol operators
TOKENSK_GLOBAL_SLASH	= 2,	Terminator operator for primary
TOKENSK_BACKSLASH	= 3,	Prefix backslash for GST scope (\X)
		Infix backslash for pathname qualifi-
		cation (EXAM M\R\X)
TOKENSK_SUBSCRIPT	= 4,	Subscript open parenthesis operator
TOKENSK_DOT	= 5,	Infix data qualification operator
		as in A.B(2).C in PL/I
TOKENSK_COBOL_OF	= 6,	Infix data qualification operator
		as in C OF B OF A(2) in COBOL
TOKENSK_PASCAL_DEREF	= 7,	Postfix dereference operator such as
		^ in PASCAL (as in P^X)
TOKENSK_PLI_DEREF	= 8,	Infix dereference operator such as
		-> in PL/I (as in P->X)
TOKENSK_INVOCNUM	= 9,	Invocation Number operator
TOKENSK_BIF_OP	= 10,	Built-in Function (parameter list)
TOKENSK_ADA_TICK	= 11,	The ADA tick operator for symbol
		attribute qualification.
TOKENSK_MAX_PRIMARY_OP	= 11;	--- Maximum value for TOKENSW_CODE
		--- for Primary Symbol operators

LITERAL

TOKENSK_MIN_OPERATOR	= 1,	--- Minimum value for TOKENSW_CODE
		--- for language expression operators
TOKENSK_INITIATOR	= 1,	Expression initiator pseudo-operator
TOKENSK_TERMINATOR	= 2,	Expression terminator pseudo-operator
TOKENSK_INDIRECT	= 3,	Indirection operator (prefix . or @)
TOKENSK_UNARY_PLUS	= 4,	Unary plus operator (+A)
TOKENSK_UNARY_MINUS	= 5,	Unary minus operator (-A)
TOKENSK_ADD	= 6,	Addition operator (A+B)
TOKENSK_SUBTRACT	= 7,	Subtraction operator (A-B)
TOKENSK_MULTIPLY	= 8,	Multiplication operator (A*B)
TOKENSK_DIVIDE	= 9,	Division operator (A/B)
TOKENSK_POWER_OF	= 10,	Exponentiation operator (A**B)
TOKENSK_OPENPAREN	= 11,	Open expression parenthesis "("
TOKENSK_CLOSEPAREN	= 12,	Close expression parenthesis ")"
TOKENSK_EQUAL	= 13,	Equal comparison operator
TOKENSK_NOT_EQUAL	= 14,	Not equal comparison operator
TOKENSK_GTR_THAN	= 15,	Greater than comparison operator
TOKENSK_GTR_THAN_U	= 16,	Unsigned Greater than comparison operator
TOKENSK_GTR_EQUAL	= 17,	Greater than or equal comparison oper.
TOKENSK_GTR_EQUAL_U	= 18,	Unsigned Greater than or equal to comparison operator
TOKENSK_LSS_THAN	= 19,	Less than comparison operator
TOKENSK_LSS_THAN_U	= 20,	Unsigned Less than comparison operator
TOKENSK_LSS_EQUAL	= 21,	Less than or equal comparison oper.
TOKENSK_LSS_EQUAL_U	= 22,	Unsigned Less than or equal comparison operator
TOKENSK_NOT	= 23,	Boolean NOT operator
TOKENSK_AND	= 24,	Boolean AND operator
TOKENSK_OR	= 25,	Boolean OR operator
TOKENSK_XOR	= 26,	Boolean XOR (exclusive OR) operator
TOKENSK_EQV	= 27,	Boolean EQV (equivalence) operator

TOKENSK_SHORT_AND	= 28,	Short-circuited boolean AND operator
TOKENSK_SHORT_OR	= 29,	Short-circuited boolean OR operator
TOKENSK_BIT_NOT	= 30,	Bitwise NOT operator
TOKENSK_BIT_AND	= 31,	Bitwise AND operator
TOKENSK_BIT_OR	= 32,	Bitwise OR operator
TOKENSK_BIT_XOR	= 33,	Bitwise XOR operator
TOKENSK_BIT_EQV	= 34,	Bitwise EQV operator
TOKENSK_CONCATENATE	= 35,	String concatenation
TOKENSK_MODULUS	= 36,	Modulus operator
TOKENSK_REMAINDER	= 37,	Remainder operator
TOKENSK_LEFT_SHIFT	= 38,	Left shift operator
TOKENSK_RIGHT_SHIFT	= 39,	Right shift operator
TOKENSK_ADDRESS_OF	= 40,	Address-of operator (as in C)
TOKENSK_SIZEOF	= 41,	Size-of operator (as in C)
TOKENSK_SET_MEMBER	= 42,	Set membership operator
TOKENSK_ABSOLUTE	= 43,	Absolute value operator
TOKENSK_INFIX_NOT	= 44,	Lexical operator (as in COBOL's NOT =)
TOKENSK_SUCCESOR	= 45,	Successor operator (BIF)
TOKENSK_PREDECESSOR	= 46,	Predecessor operator (BIF)
: Unused	= 47,	Unused
TOKENSK_INT_DIVIDE	= 48,	Lexical operator for integer divide (in PASCAL, A DIV B must be distinguished from A / B)
TOKENSK_BITSELECT	= 49,	Bit selection in BLISS and address expressions (<pos,size,ext>)
TOKENSK_DEPOSIT	= 50,	= for Deposit Command (acting as lang. dependent assignment)
TOKENSK_CONVERT	= 51,	Used as to indicate subscript conversion operator
TOKENSK_RADIX_DEC	= 52,	Set current radix to decimal
TOKENSK_RADIX_HEX	= 53,	Set current radix to hexadecimal
TOKENSK_RADIX_OCT	= 54,	Set current radix to octal
TOKENSK_RADIX_BIN	= 55,	Set current radix to binary
TOKENSK_IDENTITY	= 56,	Identity operator
TOKENSK_OPENSET	= 57,	Set constant open [
TOKENSK_BIT_IMP	= 58,	Bitwise IMP (implication) operator
TOKENSK_PRE_INCR	= 59,	Pre-increment (++X)
TOKENSK_POST_INCR	= 60,	Post-increment (X++)
TOKENSK_PRE_DECR	= 61,	Pre-decrement (--X)
TOKENSK_POST_DECR	= 62,	Post-decrement (X--)
TOKENSK_PREFIX_GTR	= 63,	Prefix > as in COBOL, ie. NOT >
TOKENSK_PREFIX_LSS	= 64,	Prefix < as in COBOL, ie. NOT <
TOKENSK_PREFIX_EQL	= 65,	Prefix = as in COBOL, ie. NOT =
TOKENSK_NEGCONST	= 66,	Negative constant
TOKENSK_POSCONST	= 67,	Positive constant
TOKENSK_MAX_OPERATOR	= 67:	--- Maximum value of TOKENSW_CODE --- for language expression operators

Define the possible values for the TOKENSK_SUBCODE field.
 This field is used to store additional information about the operator.
 At present the one use of this field is for the ADA attribute operators.
 For these, the TOKENSW_CODE field contains "TOKENSK_ADA_TICK" and the
 subcode field identifies which tick operator is present ("FIRST",
 "LAST", etc.
 The reason for doing it this way is because there may be a large number
 of tick operators, all of which are treated the same for purposes of

! operator precedence parsing. Tables thus are smaller if there is only
! one code for "tick" and a subcode for each tick operator.
!

LITERAL

TOKENSK_TICK_MIN = 1,

TOKENSK_TICK_CONSTRAINED= 1,

TOKENSK_TICK_FIRST = 2,

TOKENSK_TICK_LAST = 3,

TOKENSK_TICK_LENGTH = 4,

TOKENSK_TICK_POS = 5,

TOKENSK_TICK_PRED = 6,

TOKENSK_TICK_SIZE = 7,

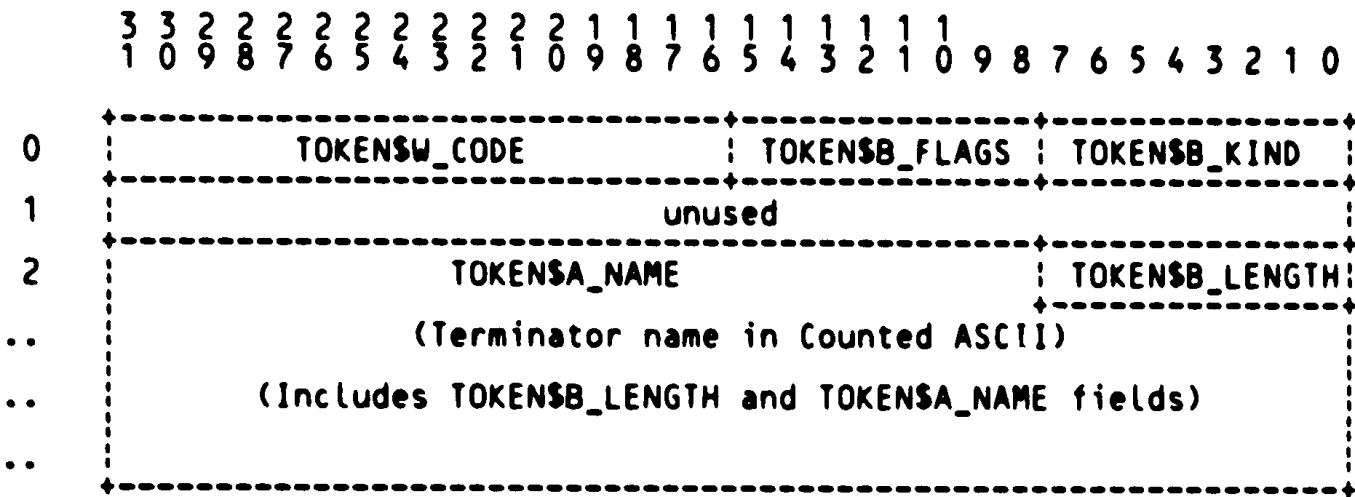
TOKENSK_TICK_SUCC = 8,

TOKENSK_TICK_VAL = 9,

TOKENSK_TICK_MAX = 9;

TERMINATOR LEXICAL TOKEN ENTRY

Terminators are lexical tokens which mark the end of an expression. For example, in subscript expressions, "." and ")" mark the ends of the expressions as in "ARR(K+2, 1*3-4)". Similarly, in the command SET BREAK X+4 DO(...), the keyword "DO" marks the end of the expression as long as it does not appear within parentheses. Terminator Lexical Token Entries define tokens which act as expression terminators in specified contexts. The Terminator Lexical Token Entry has the following format:



The TOKEN\$B_KIND field of a Terminator Lexical Token Entry always has the value zero (0).

The TOKEN\$B_FLAGS field can contain the flag TOKEN\$V_BALANCED_PARENS which, if set, indicates that this token acts as a terminator only if parentheses are balanced when it is encountered. For example, ")" terminates a subscript expression only the subscript expression's parentheses are balanced so far. Similarly, DO terminates the address expression in SET BREAK expr DO ... only if it is not enclosed in parentheses; otherwise there would be no way to include the identifier "DO" (as allowed in many languages) in the address expression.

Define the valid values of the TOKEN\$W_CODE field for terminators. This field says what kind of terminator this Token Entry represents.

LITERAL

TOKEN\$K_MIN_TERMINATOR	=	0,	---	Minimum value for TOKEN\$W_CODE
TOKEN\$K_TERM_NONE	=	0,		No terminator (just end of line)
TOKEN\$K_TERM_COMMA	=	1,		Comma ","
TOKEN\$K_TERM_CLOSE	=	2,		Close paren or bracket (")" or "]"
TOKEN\$K_TERM_COLON	=	3,		Colon ":"

TOKENSK_TERM_EQUAL	= 4.	Equal sign "="
TOKENSK_TERM_DO	= 5.	DO keyword "DO" (as in SET BREAK)
TOKENSK_TERM_THEN	= 6.	THEN keyword "THEN" (as in IF - THEN)
TOKENSK_TERM_WHEN	= 7.	WHEN keyword "WHEN" (as in SET BREAK)
TOKENSK_TERM_COMCOL	= 8.	Allow both comma "," and colon ":"
TOKENSK_TERM_CMWHDO	= 9.	Allow comma, WHEN, and DO (SET BREAK)
TOKENSK_TERM_GTRTHAN	= 10.	Greater-than sign ">" (as in <p,s,e>)
TOKENSK_TERM_OPEN	= 11.	Open paren "("
TOKENSK_TERM_COMPAREN	= 12.	"(" or ")" in CALL
TOKENSK_TERM_TO	= 13.	"to" in FOR i=1 TO n
TOKENSK_TERM_BY	= 14.	"BY"
TOKENSK_TERM_DOT	= 15.	Dot "." (as in Set constant, ie. 1..4)
TOKENSK_MAX_TERMINATOR	= 15.	--- Maximum value for TOKENSW_CODE

L I N K E D L I S T N O D E

Linked lists using this definition are generated by language specific routines to indicate:

1. The pages on which a symbol's R-value is contained.
(DBG\$NXXX_GET_PAGES)
2. Primary Descriptors which represent a range specification.
(DBG\$NXXX_RANGE_VAL)

```

3 3 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1
1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0

```

```

0  +-----+
   |                DBG$L_LINK_NODE_LINK                |
1  +-----+
   |                DBG$L_LINK_NODE_VALUE                |
   +-----+

```

A pointer to a Link Node is declared as follows:

```
LINKNODE: REF DBG$LINK_NODE
```

```
!***
```

Define the Link Node fields and declaration macro.

```
FIELD DBG$LINK_NODE_FIELDS =
```

```

SET
DBG$L_LINK_NODE_LINK    = [ 0, L_ ],    ! Pointer to next node
DBG$L_LINK_NODE_VALUE   = [ 1, L_ ],    ! Node value field
TES;

```

```

LITERAL
DBG$K_LINK_NODE_SIZE = 2;                ! Size of link node in
                                           ! longwords

```

```

MACRO
DBG$LINK_NODE = BLOCK [DBG$K_LINK_NODE_SIZE]
FIELD (DBG$LINK_NODE_FIELDS) %;

```

```
!***-
```


MOVED DST ENTRY DEFINITIONS

To support DST continuation records, every continued DST record is merged with its continuation records into one record and is moved to DEBUG permanent memory. A list of moved DST records is kept to allow access to these moved records. Each entry on that list has the following structure:

```

3 3 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1 1
1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0

```

0	-----
	MDESL_FLINK
1	-----
	MDESL_ORIGINAL_DSTPTR
2	-----
	MDESL_REAL_DSTPTR
3	-----
	MDESL_REAL_LENGTH
4	-----
	MDESL_NEXT_DSTPTR

As the RST for a module is built, continuation records are merged and the new 'moved' DST pointer is placed in the RST entry. Thus, moved DSTs are transparent for those routines that access the RST.

Those routines that obtain DST pointers from the DST itself are:

- o Routines that deal with SEPTYP records
- o Routines that access indirect type specs
- o Routines that access novel length type specs
- o Routines that access the next DST record

These routines must access those pointers through routines that search the structure to make sure that the DST hasn't been moved. The routines currently provided are:

- o DBG\$RST_DST_PTR - Return the 'real' DST pointer
- o DBG\$RST_NEXT_DST - Return the 'next' DST pointer

The structure is a singly linked list with its head being DBG\$GL_MOVED_DST_LIST_HEAD.

Define the fields of the Moved DST Record Entry. Also define the entry length and the declaration macro.

```

FIELD MDE$FIELDS =
  SET
  MDESL_FLINK        = [ 0, L_ ],    ! Forward link

```

```
MDESL_ORIGINAL_DSTPTR = [ 1, L_ ], ! The Original DST pointer
MDESL_REAL_DSTPTR     = [ 2, L_ ], ! The new, moved, real DST pointer
MDESL_REAL_LENGTH     = [ 3, L_ ], ! The moved DST's length
MDESL_NEXT_DSTPTR     = [ 4, L_ ], ! The next DST pointer
TES;
```

```
LITERAL MDESK_LENGTH = 5;           ! Length of an entry in longwords
```

```
MACRO MDE$RECORD = BLOCK[,LONG] FIELD(MDE$FIELDS) %;
```

```
! NOTE: MDESL_NEXT_DSTPTR is a simple address calculation to the next DST record
! with all the continuation records consumed. At the time the calculation is
! made it is not known if the record will be relocated. Thus any access to it
! must also check it for being relocated by calling DBG$RST_DST_PTR.
```

NUMBER SCANNER STATE TABLE

The Lexical Number Scanner in routine DBGSLEXICAL_SCANNER operates off a Finite-State Machine state table which defines the allowed formats of numeric constants in the current language. There is thus one Number State Table for each language supported by DEBUG (although some languages can share state tables). Each state in the state table is represented by a list of two or more State Table Transition Entries, each of which specifies a state transition to be taken if the current input character is of a class specified in the Transition Entry. The format of a Number Scanner State Table Entry is as follows:

```

3 3 2 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1
1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0

```

0	-----
	First State Table Transition Entry
1	-----
..	Additional State Table Transition Entries
..	-----
..	State Table Transition Entry for NUMST\$K_CLASS_OTHER

The individual State Table Transition Entry has the following format:

```

3 3 2 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1
1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0

```

0	-----
	NUMST\$W_NEXT_STATE : NUMST\$B_ACTION : \$B_CHAR_CLASS :

The whole Number Scanner State Table for a language simply consists of a sequence of concatenated State Table Entries, one entry per state in the Finite-State Machine. The whole State Table is thus a vector of State Table Transition Entries. Element zero of the state table is the beginning of the state table's Start State, and each transition's NUMST\$W_NEXT_STATE pointer is a state table index which points to the beginning of the next state's State Table Entry.

A pointer to a Number Scanner State Table is declared as follows:

```
STATE_TABLE: REF NUMST$TABLE
```

This declares STATE_TABLE to be a block vector which is referenced as .STATE_TABLE[.STATE_INDEX, field-name] where STATE_INDEX is the

vector index which points to a specific Number Scanner State Table Transition Entry.

Define the fields in the Number Scanner State Table Transition Entry.
Also define the declaration macro.

```
FIELD NUMST$FLD_DEF =
  SET
  NUMST$B_CHAR_CLASS = [ 0, B0_ ], Character class that activates
                                this state transition
  NUMST$B_ACTION      = [ 0, B1_ ], CASE index for semantic action to be
                                taken during state transition
  NUMST$W_NEXT_STATE = [ 0, W1_ ], Index of next state after transition
  TES;
```

```
MACRO
  NUMST$TABLE = BLOCKVECTOR[1, LONG] FIELD(NUMST$FLD_DEF) %;
```

Define the allowed values of the character class field NUMST\$B_CHAR_CLASS.
These are the different classes of characters which are significant when
scanning numeric constants.

```
LITERAL
  NUMST$K_MIN_CLASS      = 0,    --- Minimum character class code
  NUMST$K_CLASS_OTHER    = 0,    All other characters
  NUMST$K_CLASS_DIGIT     = 1,    The decimal digits, '0' - '9'
  NUMST$K_CLASS_HEXDIGIT = 2,    The hex digits 'A', 'C', and 'F'
  NUMST$K_CLASS_DOT       = 3,    The decimal point
  NUMST$K_CLASS_PLUS      = 4,    The plus sign '+' (for exponents)
  NUMST$K_CLASS_MINUS     = 5,    The minus sign '-' (for exponents)
  NUMST$K_CLASS_B         = 6,    The letter 'B' (for binary)
  NUMST$K_CLASS_D         = 7,    The letter 'D' (for exponents)
  NUMST$K_CLASS_E         = 8,    The letter 'E' (for exponents)
  NUMST$K_CLASS_Q         = 9,    The letter 'Q' (for exponents)
  NUMST$K_CLASS_X         = 10,   The letter 'X' (for hexadecimal)
  NUMST$K_CLASS_UNDERSCORE = 11,  ' ' is allowed in numbers in ADA
  NUMST$K_CLASS_POUND     = 12,  ' ' is allowed in numbers in ADA
  NUMST$K_CLASS_G         = 13,   The letter 'G' (for exponents)
  NUMST$K_MAX_CLASS      = 13,    --- Maximum character class code
```

Define the valid values for the NUMST\$B_ACTION field. These values are CASE
indices which select the semantic action routine to be executed for each
state transition.

```
LITERAL
  NUMST$K_MIN_ACTION      = 1,    --- Minimum action index
  NUMST$K_ACT_DO_NOthing  = 1,    No semantic action needed
  NUMST$K_ACT_GO_PAST_DIGIT = 2,   Go past digit in integer part
  NUMST$K_ACT_MARK_DEC_PT  = 3,    Mark that a decimal point was found
  NUMST$K_ACT_GO_PAST_FRAC = 4,    Go past digit in fraction part
  NUMST$K_ACT_MARK_E_EXP   = 5,    Mark that E-exponent field was found
  NUMST$K_ACT_MARK_D_EXP   = 6,    Mark that D-exponent field was found
```

NUMSTSK_ACT_MARK_Q_EXP	= 7,	Mark that Q-exponent field was found
NUMSTSK_ACT_BACKUP_PTRS	= 8,	Backup pointers and return token
NUMSTSK_ACT_GOT_NUMBER	= 9,	We got a number--return token for it
NUMSTSK_ACT_NOT_NUMBER	= 10,	Not a numeric constant
NUMSTSK_ACT_GIVE_ERROR	= 11,	We should never get here--give error
NUMSTSK_ACT_COB_CKNUM	= 12,	In COBOL this could be part of a name
NUMSTSK_ACT_COB_CKHEX	= 13,	In COBOL this could be part of a name if it is not in hex mode
NUMSTSK_ACT_GO_PAST_PACK	= 14,	Go past digit in Pack decimal
NUMSTSK_ACT_GO_PAST_PACK_FRAC=15,	!	Go past digit in Pack decimal
NUMSTSK_ACT_GOT_PACK_NUMBER=16,		We got a Packed decimal number-- return token for it
NUMSTSK_ACT_SAVE_BASE	= 17,	base#number in ADA
NUMSTSK_ACT_MARK_G_EXP	= 18,	Mark that G-exponent field was found
NUMSTSK_MAX_ACTION	= 18;	--- Maximum action index

Literals that are used in the new style operator evaluation. These values are stored in the ORTSW_ROUT field of the Operator Routine Table entry for a given operator operating on a given pair of types. The literal is used as a case index in the DBGSPERFORM_OPERATOR routine in order to actually do operations.

 Note - the order of the case indices in DBGPERMOP must correspond exactly to the ordering of the literals here. One must be very careful in changing this list.

This situation is of course a maintenance problem and we should eventually figure out a better way to get these literals shared.

In the meantime, the list below must be kept dense and must correspond exactly to the case labels in DBGPERMOP.

!***

LITERAL

ORTSK_MIN_ROUT = 1,

! Add.

ORTSK_ADD_B_B = 1,
 ORTSK_ADD_BU_BU = 2,
 ORTSK_ADD_W_W = 3,
 ORTSK_ADD_WU_WU = 4,
 ORTSK_ADD_L_L = 5,
 ORTSK_ADD_LU_LU = 6,
 ORTSK_ADD_F_F = 7,
 ORTSK_ADD_D_D = 8,
 ORTSK_ADD_G_G = 9,
 ORTSK_ADD_H_H = 10,
 ORTSK_ADD_FC_FC = 11,
 ORTSK_ADD_DC_DC = 12,
 ORTSK_ADD_GC_GC = 13,
 ORTSK_ADD_HC_HC = 14,

! Subtract.

ORTSK_SUB_B_B = 15,
 ORTSK_SUB_BU_BU = 16,
 ORTSK_SUB_W_W = 17,
 ORTSK_SUB_WU_WU = 18,
 ORTSK_SUB_L_L = 19,
 ORTSK_SUB_LU_LU = 20,
 ORTSK_SUB_F_F = 21,
 ORTSK_SUB_D_D = 22,
 ORTSK_SUB_G_G = 23,
 ORTSK_SUB_H_H = 24,
 ORTSK_SUB_FC_FC = 25,
 ORTSK_SUB_DC_DC = 26,
 ORTSK_SUB_GC_GC = 27,
 ORTSK_SUB_HC_HC = 28,

: Divide.

ORTSK_DIV_B_B	= 29.
ORTSK_DIV_BO_BU	= 30.
ORTSK_DIV_W_W	= 31.
ORTSK_DIV_WO_WU	= 32.
ORTSK_DIV_L_L	= 33.
ORTSK_DIV_LO_LU	= 34.
ORTSK_DIV_F_F	= 35.
ORTSK_DIV_D_D	= 36.
ORTSK_DIV_G_G	= 37.
ORTSK_DIV_H_H	= 38.
ORTSK_DIV_FC_FC	= 39.
ORTSK_DIV_DC_DC	= 40.
ORTSK_DIV_GC_GC	= 41.
ORTSK_DIV_HC_HC	= 42.

: Multiply

ORTSK_MUL_B_B	= 43.
ORTSK_MUL_BO_BU	= 44.
ORTSK_MUL_W_W	= 45.
ORTSK_MUL_WO_WU	= 46.
ORTSK_MUL_L_L	= 47.
ORTSK_MUL_LO_LU	= 48.
ORTSK_MUL_F_F	= 49.
ORTSK_MUL_D_D	= 50.
ORTSK_MUL_G_G	= 51.
ORTSK_MUL_H_H	= 52.
ORTSK_MUL_FC_FC	= 53.
ORTSK_MUL_DC_DC	= 54.
ORTSK_MUL_GC_GC	= 55.
ORTSK_MUL_HC_HC	= 56.

: Remainder.

ORTSK_MOD_L_L	= 57.
ORTSK_MOD_LO_LU	= 58.
ORTSK_REM_L_L	= 59.
ORTSK_REM_LO_LU	= 60.

: Shift operations.

ORTSK_SHIFT_LEFT_L_L	= 61.
ORTSK_SHIFT_RT_L_L	= 62.
ORTSK_SHIFT_RT_LO_LU	= 63.

: Exponentiation.

ORTSK_POWER_W_W	= 64.
ORTSK_POWER_L_L	= 65.
ORTSK_POWER_F_L	= 66.
ORTSK_POWER_D_L	= 67.
ORTSK_POWER_G_L	= 68.
ORTSK_POWER_H_L	= 69.
ORTSK_POWER_FC_L	= 70.

ORTSK_POWER_DC_L	= 71,
ORTSK_POWER_GC_L	= 72,
ORTSK_POWER_F_F	= 73,
ORTSK_POWER_D_F	= 74,
ORTSK_POWER_F_D	= 75,
ORTSK_POWER_D_D	= 76,
ORTSK_POWER_G_G	= 77,
ORTSK_POWER_H_H	= 78,
ORTSK_POWER_FC_FC	= 79,
ORTSK_POWER_DC_DC	= 80,
ORTSK_POWER_GC_GC	= 81,

! String operations.

ORTSK_CONCAT_T_T	= 82,
ORTSK_EQL_VT_VT	= 83,
ORTSK_EQL_T_T	= 84,
ORTSK_GEQ_VT_VT	= 85,
ORTSK_GEQ_T_T	= 86,
ORTSK_GTR_VT_VT	= 87,
ORTSK_GTR_T_T	= 88,
ORTSK_LEQ_VT_VT	= 89,
ORTSK_LEQ_T_T	= 90,
ORTSK_LSS_VT_VT	= 91,
ORTSK_LSS_T_T	= 92,
ORTSK_NEQ_VT_VT	= 93,
ORTSK_NEQ_T_T	= 94,

! Relational.

ORTSK_EQL_B_B	= 95,
ORTSK_EQL_W_W	= 96,
ORTSK_EQL_L_L	= 97,
ORTSK_EQL_F_F	= 98,
ORTSK_EQL_D_D	= 99,
ORTSK_EQL_G_G	= 100,
ORTSK_EQL_H_H	= 101,
ORTSK_EQL_FC_FC	= 102,
ORTSK_EQL_DC_DC	= 103,
ORTSK_EQL_GC_GC	= 104,
ORTSK_EQL_HC_HC	= 105,

ORTSK_NEQ_B_B	= 106,
ORTSK_NEQ_W_W	= 107,
ORTSK_NEQ_L_L	= 108,
ORTSK_NEQ_F_F	= 109,
ORTSK_NEQ_D_D	= 110,
ORTSK_NEQ_G_G	= 111,
ORTSK_NEQ_H_H	= 112,
ORTSK_NEQ_FC_FC	= 113,
ORTSK_NEQ_DC_DC	= 114,
ORTSK_NEQ_GC_GC	= 115,
ORTSK_NEQ_HC_HC	= 116,

ORTSK_GEQ_B_B	= 117,
ORTSK_GEQ_W_W	= 118,

```

ORTSK_GEQ_L_L      = 119,
ORTSK_GEQ_LO_LU    = 120,
ORTSK_GEQ_F_F      = 121,
ORTSK_GEQ_D_D      = 122,
ORTSK_GEQ_G_G      = 123,
ORTSK_GEQ_H_H      = 124,

```

```

ORTSK_GTR_B_B      = 125,
ORTSK_GTR_W_W      = 126,
ORTSK_GTR_L_L      = 127,
ORTSK_GTR_LO_LU    = 128,
ORTSK_GTR_F_F      = 129,
ORTSK_GTR_D_D      = 130,
ORTSK_GTR_G_G      = 131,
ORTSK_GTR_H_H      = 132,

```

```

ORTSK_LEQ_B_B      = 133,
ORTSK_LEQ_W_W      = 134,
ORTSK_LEQ_L_L      = 135,
ORTSK_LEQ_LO_LU    = 136,
ORTSK_LEQ_F_F      = 137,
ORTSK_LEQ_D_D      = 138,
ORTSK_LEQ_G_G      = 139,
ORTSK_LEQ_H_H      = 140,

```

```

ORTSK_LSS_B_B      = 141,
ORTSK_LSS_W_W      = 142,
ORTSK_LSS_L_L      = 143,
ORTSK_LSS_LO_LU    = 144,
ORTSK_LSS_F_F      = 145,
ORTSK_LSS_D_D      = 146,
ORTSK_LSS_G_G      = 147,
ORTSK_LSS_H_H      = 148,

```

```

! Bitwise logical operations.
!

```

```

ORTSK_BIT_AND_B_B  = 149,
ORTSK_BIT_AND_W_W  = 150,
ORTSK_BIT_AND_L_L  = 151,
ORTSK_BIT_EQV_B_B  = 152,
ORTSK_BIT_EQV_W_W  = 153,
ORTSK_BIT_EQV_L_L  = 154,
ORTSK_BIT_NOT_B    = 155,
ORTSK_BIT_NOT_W    = 156,
ORTSK_BIT_NOT_L    = 157,
ORTSK_BIT_OR_B_B   = 158,
ORTSK_BIT_OR_W_W   = 159,
ORTSK_BIT_OR_L_L   = 160,
ORTSK_BIT_XOR_B_B  = 161,
ORTSK_BIT_XOR_W_W  = 162,
ORTSK_BIT_XOR_L_L  = 163,

```

```

! Logical operations.
!

```

```

ORTSK_AND_B_B      = 164,
ORTSK_AND_W_W      = 165,

```

ORTSK_AND_L_L	= 166.
ORTSK_AND_D_D	= 167.
ORTSK_NOT_B	= 168.
ORTSK_NOT_W	= 169.
ORTSK_NOT_L	= 170.
ORTSK_NOT_D	= 171.
ORTSK_OR_B_B	= 172.
ORTSK_OR_W_W	= 173.
ORTSK_OR_L_L	= 174.
ORTSK_OR_D_D	= 175.
ORTSK_XOR_C_L	= 176.

! Unary plus and minus.

ORTSK_UNARY_MINUS_B	= 177.
ORTSK_UNARY_MINUS_W	= 178.
ORTSK_UNARY_MINUS_L	= 179.
ORTSK_UNARY_MINUS_LU	= 180.
ORTSK_UNARY_MINUS_F	= 181.
ORTSK_UNARY_MINUS_D	= 182.
ORTSK_UNARY_MINUS_G	= 183.
ORTSK_UNARY_MINUS_H	= 184.
ORTSK_UNARY_MINUS_FC	= 185.
ORTSK_UNARY_MINUS_DC	= 186.
ORTSK_UNARY_MINUS_GC	= 187.
ORTSK_UNARY_MINUS_HC	= 188.
ORTSK_UNARY_PLUS_B	= 189.
ORTSK_UNARY_PLUS_W	= 190.
ORTSK_UNARY_PLUS_L	= 191.
ORTSK_UNARY_PLUS_F	= 192.
ORTSK_UNARY_PLUS_D	= 193.
ORTSK_UNARY_PLUS_G	= 194.
ORTSK_UNARY_PLUS_H	= 195.
ORTSK_UNARY_PLUS_FC	= 196.
ORTSK_UNARY_PLUS_DC	= 197.
ORTSK_UNARY_PLUS_GC	= 198.
ORTSK_UNARY_PLUS_HC	= 199.

! Absolute Value.

ORTSK_ABS_B	= 200.
ORTSK_ABS_W	= 201.
ORTSK_ABS_L	= 202.
ORTSK_ABS_F	= 203.
ORTSK_ABS_D	= 204.
ORTSK_ABS_G	= 205.
ORTSK_ABS_H	= 206.

! Indirection.

ORTSK_INDIRECT_LU	= 207.
ORTSK_INDIRECT_TPTR	= 208.

! BLISS bit-select.

ORTSK_BITSELECT	= 209.
-----------------	--------

! Set operations.

ORTSK_DIFFERENCE_SET_SET = 210,
 ORTSK_EQL_SET_SET = 211,
 ORTSK_GEQ_SET_SET = 212,
 ORTSK_IN_SET_SET = 213,
 ORTSK_INTERSECT_SET_SET = 214,
 ORTSK_LEQ_SET_SET = 215,
 ORTSK_NEQ_SET_SET = 216,
 ORTSK_UNION_SET_SET = 217,

! C operations:

! ADDRESS OF (&), sizeof

! Also, addition or subtraction involving typed pointers
 ! is special-cased in C because there is scaling to be done.

ORTSK_ADDRESS_L = 218,
 ORTSK_SIZEOF_C = 219,
 ORTSK_ADD_TPTR_L = 220,
 ORTSK_SUB_TPTR_L = 221,
 ORTSK_SUB_TPTR_TPTR = 222,

! BASIC implication (A imp B)

ORTSK_BIT_IMP_B_B = 223,
 ORTSK_BIT_IMP_W_W = 224,
 ORTSK_BIT_IMP_L_L = 225,

! C pre- and post- increment, decrement (++X X++ --X X--)

ORTSK_PRE_INCR_L = 226,
 ORTSK_PRE_INCR_LU = 227,
 ORTSK_PRE_INCR_D = 228,
 ORTSK_PRE_INCR_TPTR = 229,
 ORTSK_POST_INCR_L = 230,
 ORTSK_POST_INCR_LU = 231,
 ORTSK_POST_INCR_D = 232,
 ORTSK_POST_INCR_TPTR = 233,
 ORTSK_PRE_DECR_C = 234,
 ORTSK_PRE_DECR_LU = 235,
 ORTSK_PRE_DECR_D = 236,
 ORTSK_PRE_DECR_TPTR = 237,
 ORTSK_POST_DECR_L = 238,
 ORTSK_POST_DECR_LU = 239,
 ORTSK_POST_DECR_D = 240,
 ORTSK_POST_DECR_TPTR = 241,

! Packed Decimal Operations.

ORTSK_ADD_P_P = 242,
 ORTSK_SUB_P_P = 243,
 ORTSK_MUL_P_P = 244,
 ORTSK_DIV_P_P = 245,
 ORTSK_UNARY_PLUS_P = 246,
 ORTSK_UNARY_MINUS_P = 247,

```
ORT$K_EQL_P_P      = 248,
ORT$K_NEQ_P_P      = 249,
ORT$K_GTR_P_P      = 250,
ORT$K_GEQ_P_P      = 251,
ORT$K_LSS_P_P      = 252,
ORT$K_LEQ_P_P      = 253,
```

```
! PL/I Bit-String Operations.
```

```
ORT$K_CONCAT_TF_TF = 254,
ORT$K_EQL_TF_TF    = 255,
ORT$K_NEQ_TF_TF    = 256,
ORT$K_GTR_TF_TF    = 257,
ORT$K_GEQ_TF_TF    = 258,
ORT$K_LSS_TF_TF    = 259,
ORT$K_LEQ_TF_TF    = 260,
ORT$K_BIT_NOT_TF   = 261,
ORT$K_BIT_AND_TF   = 262,
ORT$K_BIT_OR_TF    = 263,
```

```
ORT$K_EQL_RFA_RFA  = 264,
ORT$K_NEQ_RFA_RFA  = 265,
```

```
ORT$K_UNARY_PLUS_Q  = 266,
ORT$K_UNARY_MINUS_Q = 267,
ORT$K_UNARY_PLUS_O  = 268,
ORT$K_UNARY_MINUS_O = 269,
```

```
! Built-in Functions
```

```
ORT$K_SUCC_ENUM     = 270,
ORT$K_PRED_ENUM     = 271,
```

```
! Fixed binary operations.
```

```
ORT$K_UNARY_PLUS_FIXED = 272,
ORT$K_UNARY_MINUS_FIXED = 273,
ORT$K_ABS_FIXED        = 274,
ORT$K_ADD_FIXED_FIXED  = 275,
ORT$K_SUB_FIXED_FIXED  = 276,
ORT$K_MUL_FIXED_FIXED  = 277,
ORT$K_DIV_FIXED_FIXED  = 278,
ORT$K_EQL_FIXED_FIXED  = 279,
ORT$K_NEQ_FIXED_FIXED  = 280,
ORT$K_LSS_FIXED_FIXED  = 281,
ORT$K_GTR_FIXED_FIXED  = 282,
ORT$K_LEQ_FIXED_FIXED  = 283,
ORT$K_GEQ_FIXED_FIXED  = 284,
ORT$K_MAX_ROUT         = 284;
```

```
!*** ALSO CHANGE the corresponding
!*** definition in DBGPERMOP.MAR if
!*** you change this.
```

```
! Literals that are defined for type id check in DBG$PERFORM_TYPEID_CHECK.
```

```
! LITERAL
```



```
ORTSK_TYPEID_MIN_ROOT = 1,  
ORTSK_TYPEID_ENUM_ENUM = 1,  
ORTSK_TYPEID_SET_SET = 2,  
ORTSK_TYPEID_TPTR_TPTR = 3,  
ORTSK_TYPEID_SUBRNG_SUBRNG = 4,  
ORTSK_TYPEID_MAX_ROOT = 4;
```

!***-

TYPE GRAPH

The type conversions that are done for a given operator in a given language are specified by a set of directed acyclic graphs. For example, part of the graph that specifies the implicit type conversions to be done for FORTRAN addition might look like:

```

      L -> F -> D -> H
          \   \
          FC -> DC
  
```

Each graph is represented by listing its edges. Each edge is just a pair of Type IDs.

Below we define the data structure used to declare the edges of these graphs.

```

3 3 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1
1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0
  
```

```

+-----+
| HIGHER_TYPE | LOWER_TYPE |
+-----+
  
```

Define the fields in an entry in an edge of a Type Graph.

```
FIELD TYPE_GRAPH$FLD_DEF =
```

```

SET
TYPE_GRAPH$B_LOWER_TYPE = [0,B0_],
TYPE_GRAPH$B_HIGHER_TYPE = [0,B1_],
TYPE_GRAPH$W_BOTH_TYPES = [0,W0_]
TES;
  
```

Define the macro which is used to declare an entry in a Type Graph

```
MACRO
```

```
TYPE_GRAPH$ENTRY = BLOCK [1,WORD] FIELD (TYPE_GRAPH$FLD_DEF) %;
```


TYPE CONVERSION INFORMATION TABLE

The Type Conversion Information Table is a vector of longwords used by a language which gives pointers to tables which contain the information necessary to select the type conversions and then the routine invocations which does the type conversion.

3 3 2 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1
1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0

0	CVTINFO\$MPTBL
1	CVTINFO\$CVTBL
2	flags

A pointer to an Type Conversion Information Table is declared as follows:

CVTINFO_TABLE: REF CVTINFO\$TABLE

Define the fields of the Type Conversion Information Table entry. Also define the declaration macro.

```
FIELD CVTINFO$FLD_DEF =
SET
  CVTINFO$MPTBL    = [ 0, L_ ], | Address of Type Map Table
  CVTINFO$CVTBL    = [ 1, L_ ], | Address of Type Conversion Table
  CVTINFO$V_ROUND  = [ 2, V_(0) ], | Flag to indicate conversion result
                                     | is rounded
TES;
```

```
MACRO
  CVTINFO$TABLE = BLOCK[3, LONG] FIELD (CVTINFO$FLD_DEF) %;
```

TYPE CONVERT ENTRY

A Type Convert Entry contains a type conversion routine index which identifies the routine to be invoked, from data type, and to data type.

```

3 3 2 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1
1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0

```

```

+-----+
| TYPE_CVT$W_ROUT      | :  HIGHER_TYPE  :  LOWER_TYPE  |
+-----+

```

A pointer to an Type Map Entry is declared as follows:

```
CVTPTR : REF TYPE_CVT$ENTRY
```

Define the fields of the Type Map Table entry. Also define an entry declaration macro.

FIELD

```
TYPE_CVT$FLD_DEF =
```

```
SET
```

```

TYPE_CVT$B_LOWER_TYPE = [0, B0-],
TYPE_CVT$B_HIGHER_TYPE = [0, B1-],
TYPE_CVT$W_MAP_PAIR    = [0, W0-],
TYPE_CVT$W_ROUT        = [0, W1-]
TES;

```

MACRO

```
TYPE_CVT$ENTRY = BLOCK[1, LONG] FIELD (TYPE_CVT$FLD_DEF) %;
```

```

! Literals that are used in the new style operator evaluation.
! The literal is used as a case index in the DBG$LANGUAGE_TYPE_CONV
! routine in order to actually do the conversion.

```

```
!***
```

LITERAL

```

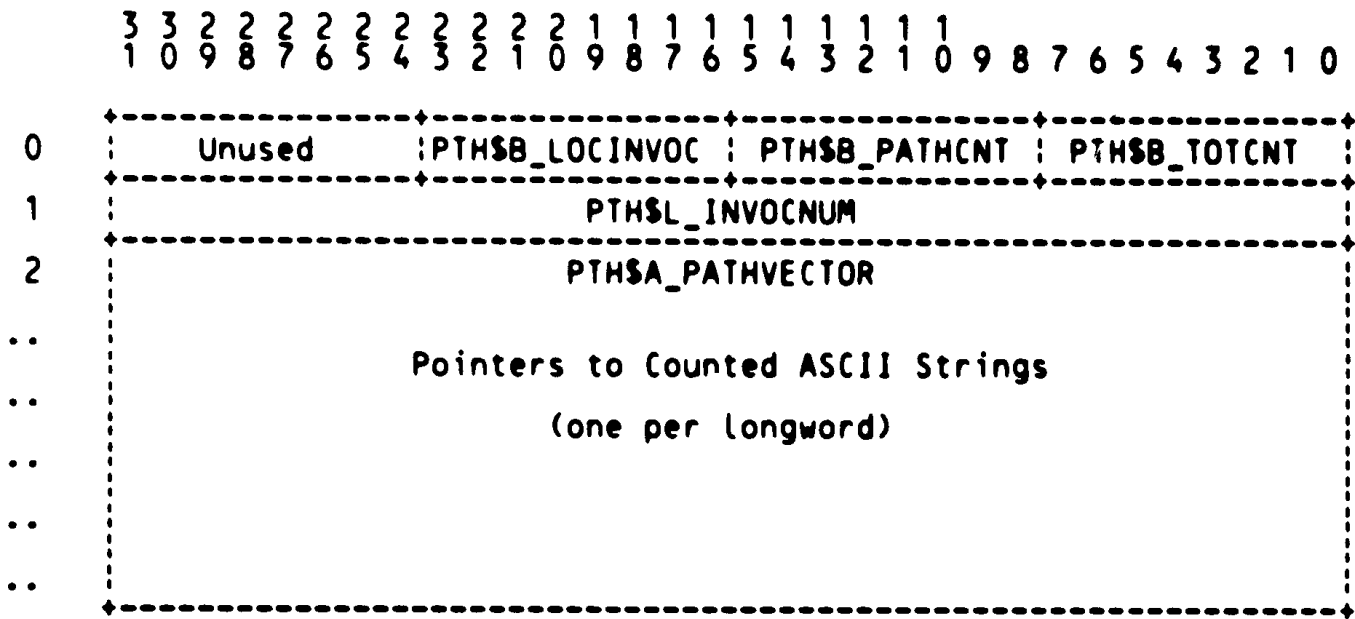
CVT$K_MIN_ROUT      = 1,
CVT$K_PLI_CVT       = 1,
CVT$K_COB_PICT      = 2,
CVT$K_MAX_ROUT      = 2;

```

```
!***
```

PATHNAME DESCRIPTORS

The pathname descriptor is used internally in the Debugger to represent a pathname, including data qualification. Thus 'MOD\ROUT\X' is a path-name in this sense, and so is 'MOD\ROUT\A.B.C'. The format of a path-name descriptor is shown here:



The Pathname Descriptor contains two counts fields which give the total number of names in the pathname (five for M\R\A.B.C) and the total number of names before the first dot (three for M\R\A.B.C). It also contains an invocation number (which is zero if none was explicitly specified) and an index into PTHSA_PATHVECTOR which indicates which pathname component was qualified by the invocation number (1 for A 2\X and 4 for M\R1\R2\R3 5\X). This index is zero if no invocation number was given. The rest of the descriptor consists of longwords containing pointers to the individual names in the pathname, each represented by a Counted ASCII String.

A pointer to a Pathname Descriptor is declared as follows:

```
PATHPTR: REF PTHSPATHNAME
```

Define the Pathname Descriptor fields.

```
FIELD PTHSFLD_DEF =
SET
PTHSB_TOTCNT = [ 0, 80_ ], ! Total number of names in pathname
```

```

PTHSB_PATHCNT    = [ 0, B1_ ],
PTHSB_LOCINVOC   = [ 0, B2_ ],
! ---
PTHSL_INVOCNUM   = [ 0, B3_ ],
PTHSA_PATHVECTOR= [ 1, L_ ],
PTHSA_PATHVECTOR= [ 2, A_ ]

TES;

```

including all data qualification
Number of names in pathname proper
before any data qualification
The location in the vector below of
the name with an invocation num-
ber (first name is 1) or zero
Unused
The invocation number or zero
The start of the pathname vector--has
one name pointer (pointing to a
Counted ASCII name) per longword.

```

! Define the maximum allowed size of a pathname, i.e. the maximum number of
! individual names, including data components. Also define the size of the
! fixed part of the descriptor and the declaration macro.

```

```

LITERAL

```

```

DBG$K_MAX_PATHNAME = 50,
DBG$K_PATHDESCSIZE = 2,
DBG$K_PATHNAME_SIZE = DBG$K_MAX_PATHNAME + DBG$K_PATHDESCSIZE;

```

Maximum size of a pathname
Size of fixed part of Pathname
Descriptor

```

MACRO

```

```

PTH$PATHNAME = BLOCK[DBG$K_PATHNAME_SIZE] FIELD(PTH$FLD_DEF) %;

```

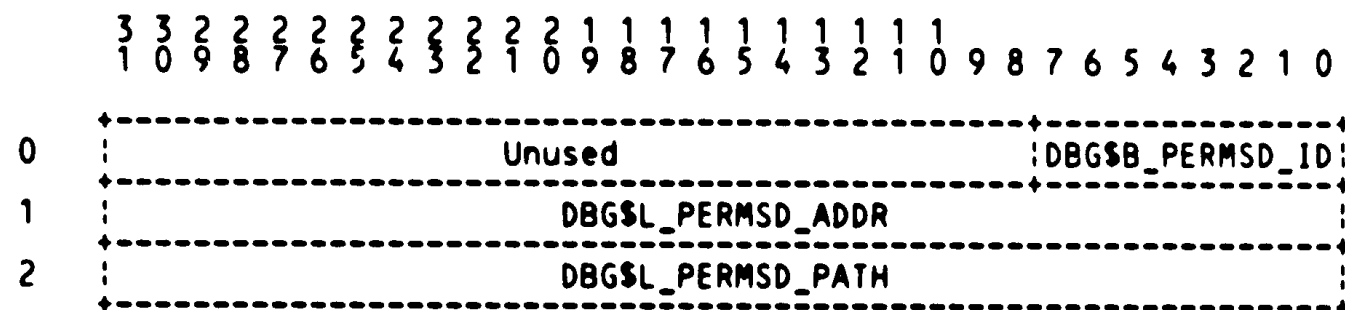
```

!***

```

P E R M A N E N T S Y M B O L D E S C R I P T O R S

The Permanent Symbol Descriptor is used to describe Debugger "permanent symbols" in Value Descriptors, Primary Descriptors, and address expressions. Permanent symbols name the VAX registers; "R0" is thus a permanent symbol. Permanent Symbol Descriptors are only used for languages COBOL and PL/I. A Permanent Symbol Descriptor has this format:



A pointer to a Permanent Symbol Descriptor is declared as follows:

```
PSDPTR: REF DBG$PERMSD
```

Define the Permanent Symbol Descriptor fields and declaration macro.

```
FIELD DBG$PERMSD_FIELDS =
```

```
SET
DBG$B_PERMSD_ID   = [ 0, B0_ ], : Symbol ID (register number)
DBG$L_PERMSD_ADDR = [ 1, L_ ],  : Address of the register contents
DBG$L_PERMSD_PATH = [ 2, L_ ],  : Address of symbol's pathname descr.
TES;
```

```
LITERAL
DBG$K_PERMSD_SIZE = 3;           : Size of Permanent Symbol Descriptor
                                : in longwords
```

```
MACRO
DBG$PERMSD = BLOCK[DBG$K_PERMSD_SIZE] FIELD(DBG$PERMSD_FIELDS) %;
```

These are the possible values of the DBG\$B_PERMSD_ID field.

```
LITERAL
DBG$K_R0      = 200,           : Register R0
DBG$K_R1      = 201,           : Register R1
DBG$K_R2      = 202,           : Register R2
DBG$K_R3      = 203,           : Register R3
DBG$K_R4      = 204,           : Register R4
DBG$K_R5      = 205,           : Register R5
```

DBG\$K_R6	= 206,	! Register R6
DBG\$K_R7	= 207,	! Register R7
DBG\$K_R8	= 208,	! Register R8
DBG\$K_R9	= 209,	! Register R9
DBG\$K_R10	= 210,	! Register R10
DBG\$K_R11	= 211,	! Register R11
DBG\$K_AP	= 212,	! Argument Pointer (AP)
DBG\$K_FP	= 213,	! Frame Pointer (FP)
DBG\$K_SP	= 214,	! Stack Pointer (SP)
DBG\$K_PC	= 215,	! Program Counter (PC)
DBG\$K_PSL	= 216;	! Processor Status Longword (PSL)

!***-


```
PRID$ENTRY = BLOCK[,LONG] FIELD(PRID$FLD_DEF) %;
```

```
! Define the valid values of the PRID$B_KIND field.
```

```
LITERAL PRID$K_CONSTANT = 1;          ' Predeclared ID constant
```


The possible values of the DBG\$B_DHEAD_LANG field are the language codes defined in the section entitled "Miscellaneous Literals" above. The possible value of the DBG\$B_DHEAD_TYPE field, which are defined in the same section, are the following:

DBG\$K_PRIMARY_DESC	! This is a Primary Descriptor
DBG\$K_VALUE_DESC	! This is a non-volatile Value Descriptor
DBG\$K_V_VALUE_DESC	! This is a volatile Value Descriptor

PRIMARY DESCRIPTOR DEFINITIONS

Primary Descriptors describe Primary Symbols parsed in DEBUG commands. A Primary Symbol may be a simple name, such as 'ABC', or a more complex name including pathname qualification, data qualification, subscripting, and dereferencing, such as 'MOD\ROUT\A.B[2,3]^C(4).D'.

The Primary Descriptor described here is the Primary Descriptor built by the language-independent parser in module DBGPARSER.

A Primary Descriptor consists of a Root Node followed by a doubly linked list of Sub-Nodes which describe the individual components of the symbol described by the Primary Descriptor. Each Sub-Node thus describes an instance of data qualification, subscripting, dereferencing, or the like. The last Sub-Node describes the final object named by the Primary Symbol.

Consider the Primary Symbol 'A\B\C.D[2]^E' for example. The Primary Descriptor Root Node would point to a Sub-Node for A\B\C which requires a Record Sub-Node. It points to another Sub-Node for component D of record A\B\C. This component is an array and requires an Array Sub-Node. The next node is the Sub-Node for the array element A\B\C.D[2]. This element is of a pointer type. The next Sub-Node is a node for the pointed-to object which happens to be a record and therefore calls for a Record Sub-Node. Finally, the last Sub-Node on the chain is for component E of that record. The type of Sub-Node required depends on the data type of E, but would be a Normal Sub-Node for all data types other than arrays and records.

! those routines which can accept either a Primary Descriptor or a Value
! Descriptor as input.

MACRO

DBG\$PRIMARY = BLOCK [,LONG]
FIELD(DBG\$DHDR_FIELDS, DBG\$PRIM_FIELDS, DBG\$VALUE_FIELDS) %;

! Define the size in longwords of the Primary Descriptor Root Node.

LITERAL

DBG\$K_PRIMARY_SIZE = 9; ! Size of Root Node in longwords

PRIMARY DESCRIPTOR NORMAL SUB-NODE

Primary Descriptor Sub-Nodes are attached to the Primary Descriptor Root Node via a doubly linked list. The first Sub-Node represents the first data item in the Primary Symbol (for example, A/B in 'A/B.C[[2],D') and subsequent nodes represent subsequent items (such as .C, .C[[2], and .D). All Primary Descriptor Sub-Nodes include a 'common core' of information, with other fields added for certain specific sub-nodes. The 'Normal Sub-Node' described here is used for all 'normal' symbols, meaning all symbols for which no specific Primary Descriptor Sub-Node has been defined below. This Sub-Node has only the 'common core' fields and no type-specific fields:

```

3 3 2 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1 1
1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0

```

0	----- DBG\$P_NODE_FLINK -----
1	----- DBG\$P_NODE_BLINK -----
2	----- Unused : \$B_PNODE_FLAGS : \$B_PNODE_FCODE : \$B_PNODE_KIND -----
3	----- DBG\$P_NODE_TYPEID -----
4	----- DBG\$P_NODE_SYMID -----
5	----- DBG\$P_NODE_RELOC -----

A pointer to a Primary Descriptor Sub-Node is declared as follows:

```
PNODEPTR: REF DBG$PRIM_NODE
```

The following are the meanings of the fields in the Normal Sub-Node.

DBG\$B_PNODE_KIND	-- May be any RST kind, although it usually is RST\$K_DATA or RST\$K_TYPCOMP.
DBG\$B_PNODE_FCODE	-- Can be any FCODE other than those for which separate Primary Descriptor Sub-Nodes are described below.
DBG\$P_NODE_TYPEID	-- The Type ID for the object's type.
DBG\$P_NODE_SYMID	-- The SYMID for the object. This field may be zero, as it is for an array element or a pointed-to element.
DBG\$P_NODE_RELOC	-- A 'relocation constant' for the address. This is normally zero, but can contain a user address (in the case of anonymous references) or a byte offset (for strings)

Define the size of the Primary Descriptor Normal Sub-Node.

```
LITERAL    DBGSK_PRIM_SIZE_NORMAL    = 6;      ! Size of Normal Sub-Node in longwords
```

```
! Define the fields of the Primary Descriptor Sub-Node.  Also define the
! declaration macro.
```

```

FIELD DBG$PNODE_FIELDS =
SET
DBG$PNODE_FLINK = [ 0, L_ ], : Link to following primary sub-node
DBG$PNODE_BLINK = [ 1, L_ ], : Link to preceding primary sub-node
DBG$PNODE_KIND = [ 2, B0_ ], : The RST "kind" of this component
DBG$PNODE_FCODE = [ 2, B1_ ], : FCODE for this component
DBG$PNODE_FLAGS = [ 2, B2_ ], : Flags byte
DBG$PNODE_EVAL = [ 2, V2_(0) ], : Evaluate subscripting, dot
                                : qualification, or dereferencing
                                : operation on this component
DBG$PNODE_PNARR_COLUMN = [ 2, V2_(1) ], : Set if array in 'column major' order
DBG$PNODE_PNARR_BITREF = [ 2, V2_(2) ], : Set if this is a BIT array
DBG$PNODE_PNARR_RANGE = [ 2, V2_(3) ], : Set if array has subscript ranges
DBG$PNODE_PNVAR_VALID = [ 2, V2_(4) ], : Set if tag has been validated
DBG$PNODE_IGNORE = [ 2, V2_(5) ], : Says to ignore SYMID in address
                                : computation.
DBG$PNODE_TYPEID = [ 3, L_ ], : TYPEID for this component
DBG$PNODE_SYMID = [ 4, L_ ], : Generic SYMID for this component
DBG$PNODE_RELOC = [ 5, L_ ], : Address relocation factor

DBG$PNODE_PNREC_INDEX = [ 6, W0_ ], : Record Component Index
DBG$PNODE_PNREC_NCOMPS = [ 6, W1_ ], : Number of components in record

DBG$PNODE_PNVAR_INDEX = [ 6, W0_ ], : Variant Component Index
DBG$PNODE_PNVAR_NCOMPS = [ 6, W1_ ], : Number of components in variant
DBG$PNODE_PNVAR_TAGID = [ 7, L_ ], : SYMID of tag variable (or 0)
DBG$PNODE_PNVAR_COMPLST = [ 8, L_ ], : Pointer to list of record components
                                : in this record variant
DBG$PNODE_PNVAR_DSTPTR = [ 9, L_ ], : Pointer to Variant Value DST Record
                                : for this record variant

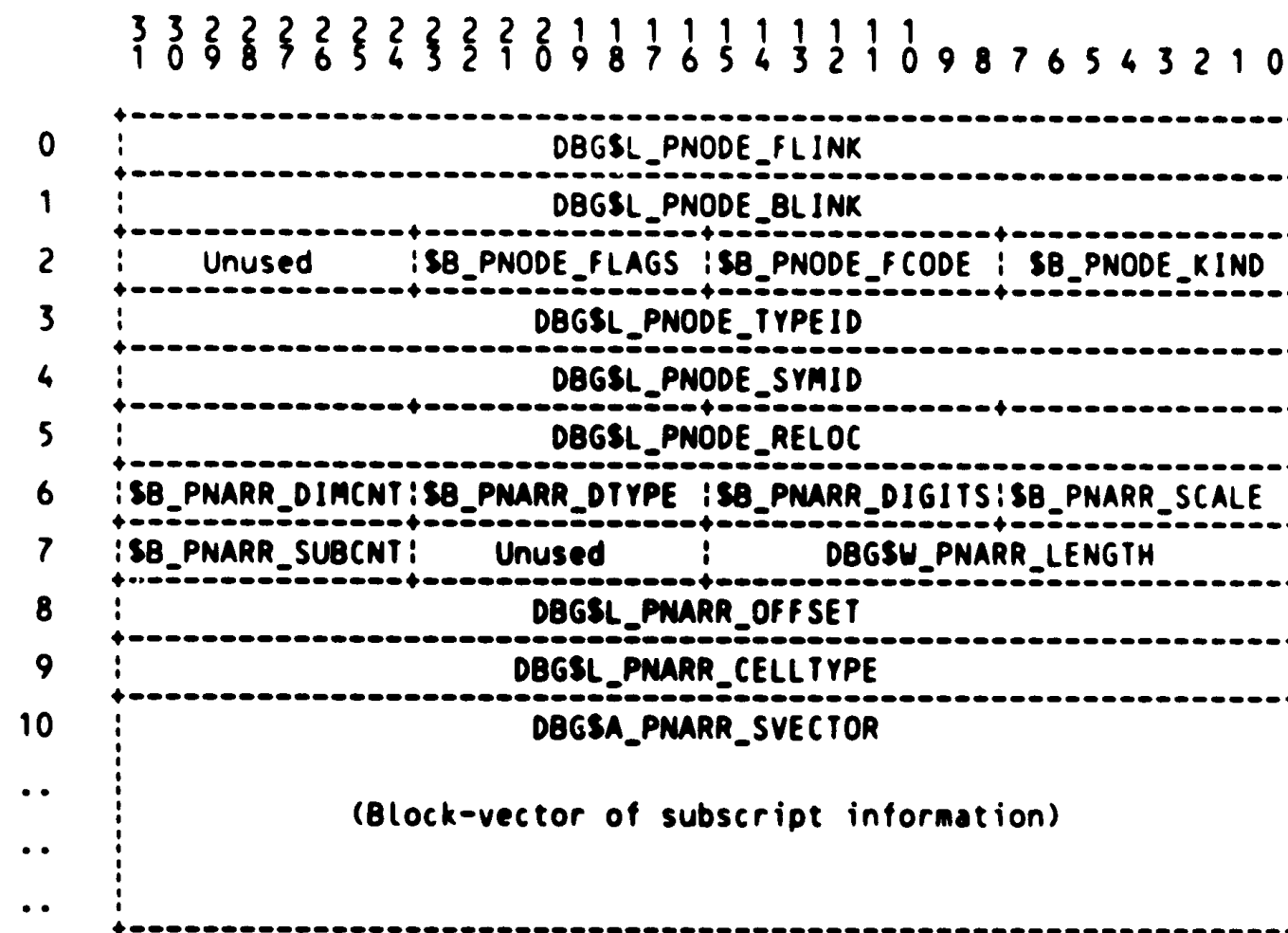
DBG$PNODE_PNARR_SCALE = [ 6, B0_ ], : Decimal scale factor for array element
DBG$PNODE_PNARR_DIGITS = [ 6, B1_ ], : Decimal digits for array element or 0
DBG$PNODE_PNARR_DTYPE = [ 6, B2_ ], : DTYPE of an array element
DBG$PNODE_PNARR_DIMCNT = [ 6, B3_ ], : Number of array dimensions
DBG$PNODE_PNARR_LENGTH = [ 7, W0_ ], : Length of each array element
DBG$PNODE_PNARR_SUBCNT = [ 7, B3_ ], : Number of actual subscripts supplied
DBG$PNODE_PNARR_OFFSET = [ 8, L_ ], : Offset to element [0,0,...,0] (bit or
                                : byte depending on BITREF flag)
DBG$PNODE_PNARR_CELLTYPE = [ 9, L_ ], : Type ID of array element data type
DBG$PNODE_PNARR_SVECTOR = [ 10, A_ ], : BLOCKVECTOR of subscripts
TES:

```

```
MACRO DBG$PRIM_NODE = BLOCK [ ,LONG] FIELD(DBG$PNODE_FIELDS)%;
```


PRIMARY DESCRIPTOR ARRAY SUB-NODE

This Primary Descriptor Sub-Node is used for array types (i.e., for Primary Symbol components whose FCODE is Array). The node contains all information needed later about the array type itself, including the dimension count, all subscript types and bounds, etc. This is the format:



The following are the meanings of the fields in the Array Type Sub-Node.

DBG\$B_PNODE_KIND	-- Must be RST\$K_DATA or RST\$K_TYPCOMP.
DBG\$B_PNODE_FCODE	-- Must be RST\$K_TYPE_ARRAY.
DBG\$B_PNODE_TYPEID	-- The Type ID for the array type.
DBG\$B_PNODE_SYMID	-- The SYMID for the array (or zero).
DBG\$B_PNARR_SCALE	-- The element scale factor or zero.
DBG\$B_PNARR_DIGITS	-- The element digit count or zero.
DBG\$B_PNARR_DTYPE	-- The element VAX standard type code (if applicable) or zero.

```

DBG$B_PNARR_DIMCNT -- The number of array dimensions.
DBG$B_PNARR_SUBCNT -- The number of actual subscripts supplied.
                      In some languages this may be less than
                      the number of array dimensions.
DBG$B_PNARR_LENGTH -- The length of each array element. This
                      length is in bytes except for DTYPEs
                      V (bits) and P (packed decimal).
DBG$L_PNARR_OFFSET -- The offset to element ARR[0,0,...,0]
                      from the start of the array.
DBG$L_PNARR_CELLTYPE -- The Type ID of the array element type.
DBG$A_PNARR_SVECTOR -- Start of the subscript block-vector.

```

The block-vector of subscript information consists of a vector whose elements are blocks of the following format. There is once such block for each dimension in the array.

```

3 3 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1 1
1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0

```

0	DBG\$L_PNSUB_SVALUE
1	DBG\$L_PNSUB_STRIDE
2	DBG\$L_PNSUB_LBOUND
3	DBG\$L_PNSUB_UBOUND
4	DBG\$L_PNSUB_TYPEID

Define the fields in the subscript information block-vector.

FIELD DBG\$PNSUB_FIELDS =

```

SET
DBG$L_PNSUB_SVALUE = [ 0, L_ ], | Subscript Value
DBG$L_PNSUB_STRIDE = [ 1, L_ ], | Subscript Stride (bit or byte
                                | depending on BITREF flag)
DBG$L_PNSUB_LBOUND = [ 2, L_ ], | Subscript Lower Bound
DBG$L_PNSUB_UBOUND = [ 3, L_ ], | Subscript Lower Bound
DBG$L_PNSUB_TYPEID = [ 4, L_ ], | Subscript type's Type ID. If the
                                | subscript type is longword
                                | integer, this field is zero.

```

TES;

MACRO

DBG\$PRIM_NODE_SUBS = BLOCKVECTOR[,5, LONG] FIELD(DBG\$PNSUB_FIELDS) %;

Define the size of the Primary Descriptor Array Sub-Node.

LITERAL

```
DBG$K_PRIM_SIZE_ARRAY = 10,  ! Size of fixed part in longwords
DBG$K_PRIM_SIZE_SUBS   = 5;  ! Size of variable part per subscript,
                           ! also in longwo ds
```

PRIMARY DESCRIPTOR RECORD SUB-NODE

This Primary Descriptor Sub-Node is used for record types (i.e., for Primary Symbol components whose FCODE is Record). The only extra field in this sub-node is the record (or "structure") component index which indicates which record component is selected from the defined components of this record type.

3 3 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1 1
1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0

0	DBG\$L_PNODE_FLINK												
1	DBG\$L_PNODE_BLINK												
2	Unused	: \$B_PNODE_FLAGS :				: \$B_PNODE_FCODE :				: \$B_PNODE_KIND :			
3	DBG\$L_PNODE_TYPEID												
4	DBG\$L_PNODE_SYMID												
5	DBG\$L_PNODE_RELOC												
6	DBG\$W_PNREC_NCOMPS					:	DBG\$W_PNREC_INDEX						

The following are the meanings of the fields in the Record Type Sub-Node.

DBG\$B_PNODE_KIND -- Must be RST\$K_DATA or RST\$K_TYPCOMP.
 DBG\$B_PNODE_FCODE -- Must be RST\$K_TYPE_RECORD.
 DBG\$L_PNODE_TYPEID -- The Type ID for the record type.
 DBG\$L_PNODE_SYMID -- The SYMID for the record (or zero).
 DBG\$W_PNREC_INDEX -- The index of the selected record component in the set of record components for this record type. This is meaningful only if the DBG\$V_PNODE_EVAL is set--otherwise this field is always initialized to 1.
 DBG\$W_PNREC_NCOMPS -- The number of components in this record

Define the size of the Primary Descriptor Record Sub-Node.

LITERAL
 DBG\$K_PRIM_SIZE_RECORD = 7; ! Size of Record Sub-Node in longwords

! Define the size of the Primary Descriptor Variant Sub-Node.

LITERAL
DBG\$K_PRIM_SIZE_VARIANT = 10; ! Size of Variant Sub-Node in longwords

PRIMARY PARSER STATE TABLE

The Primary Parser (DBG\$PRIMARY_PARSER) operates off a Finite-State Machine (FSM) state table which defines the allowed formats of Primary Symbols in the current language. A Primary Symbol consists of a symbol name which may include pathname qualification, data qualification, subscripting, dereferencing, and possibly other qualification depending on the language. In PASCAL, for example, M[R\A.B[2,3]^.[4].D is a valid Primary Symbol. The operators in a Primary Symbol may include '[', '^', '.', ']', and so on. Exactly which operators are allowed and in which order they are allowed is language-dependent, and this is what is specified by the Primary Parser State Table. Each state in the state table is represented by zero or more Primary Parser State Table Transition Entries, each of which specifies a transition to be taken if the current operator is the one specified in the transition entry. The format of a Primary Parser State Table Entry is as follows:

3 3 2 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1
1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0

0	First Primary Parser State Table Transition Entry
1	Additional State Table Transition Entries
..	
..	Longword with value zero (0)

The individual Primary Parser State Table Transition Entry has this format:

3 3 2 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1
1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0

0	PRIMARY\$W_NEXTSTATE	SB_ACTION	SB_OPCODE
---	----------------------	-----------	-----------

The whole Primary Parser State Table for a language simply consists of a sequence of concatenated State Table Entries, with one entry per state in the Finite-State Machine. The whole state table is thus a vector of transition entries and is indexed as such. Element zero of the state table is the beginning of the state table's Start State, and each transition's next state pointer (PRIMARY\$W_NEXT STATE) is a state table index which points to the beginning of the next state's state table entry.

A pointer to a Primary Parser State Table is declared as follows:

PRIMARY_TABLE: REF PRIMARY\$TABLE

This declares PRIMARY_TABLE to be a block vector which is referenced as .PRIMARY_TABLE[.STATE_INDEX, field-name] where STATE_INDEX is the vector index pointing to the current Primary Parser State Table Transition Entry.

Define the fields in the Primary parser State Table Transition Entry. Also define the declaration macro.

FIELD PRIMARY\$FLD_DEF =

```

SET
PRIMARY$B_OPCODE = [ 0, B0_ ],  | Operator code which activates this
                                | state transition
PRIMARY$B_ACTION = [ 0, B1_ ],  | CASE index for semantic action to be
                                | taken during state transition
PRIMARY$W_NEXTSTATE=[ 0 , W1_ ] | Index of next state after transition
TES;
```

MACRO

PRIMARY\$TABLE = BLOCKVECTOR[.1, LONG] FIELD(PRIAMRY\$FLD_DEF) %;

Define the allowed values of the PRIMARY\$B_ACTION field. These values are CASE indices which select the semantic action routine to be executed for each state transition.

LITERAL

```

PRIMARY$K_MIN_ACTION          = 1,  | --- Minimum action index
PRIMARY$K_ACT_START_GBL       = 1,  | Global scope backslash found
PRIMARY$K_ACT_GBL_TERM        = 2,  | Terminator after global symbol
PRIMARY$K_ACT_START_SLASH     = 3,  | Backslash after first operand
PRIMARY$K_ACT_START_DOT       = 4,  | Dot qualification after first operand
PRIMARY$K_ACT_START_DOT_PLI   = 5,  | " in PL/I
PRIMARY$K_ACT_START_DOT_COB   = 6,  | " in COBOL
PRIMARY$K_ACT_START_SUBSCR     = 7,  | Subscripting after first operand
PRIMARY$K_ACT_START_SUBSCR_BLI = 8,  | Bliss subscript
PRIMARY$K_ACT_START_SUBSCR_PLI = 9,  | PL/I subscript
PRIMARY$K_ACT_START_DEREF     = 10,  | PASCAL dereference ^
PRIMARY$K_ACT_START_DEREF_PLI = 11,  | PL/I "->" operator
PRIMARY$K_ACT_START_TERM      = 12,  | Terminator after first operand
PRIMARY$K_ACT_START_BIF_CALL  = 13,  | Built-in function call
PRIMARY$K_ACT_START_TICKR     = 48,  | Ada tick operator

PRIMARY$K_ACT_SLASH_SLASH     = 14,  | Backslash after previous backslash
PRIMARY$K_ACT_SLASH_INVOCNUM  = 15,  | Invocation Number in pathname
PRIMARY$K_ACT_SLASH_DOT       = 16,  | Dot qualification after backslash
PRIMARY$K_ACT_SLASH_DOT_PLI   = 17,  | " in PL/I
PRIMARY$K_ACT_SLASH_SUBSCR     = 18,  | Subscripting after backslash
PRIMARY$K_ACT_SLASH_SUBSCR_PLI = 19,  | " in PL/I
PRIMARY$K_ACT_SLASH_SUBSCR_BLI = 20,  | Bliss subscript after backslash
PRIMARY$K_ACT_SLASH_DEREF     = 21,  | PASCAL ^ after backslash
PRIMARY$K_ACT_SLASH_DEREF_PLI = 22,  | PL/I "->" operator after "/"
PRIMARY$K_ACT_SLASH_TERM      = 23,  | Terminator operator after backslash
```


PRIMARYSK_ACT_SLASH_TICK	= 49, !	Ada tick operator
PRIMARYSK_ACT_DOT_SLASH_COB	= 24, !	A of B\C in COBOL
PRIMARYSK_ACT_DOT_DOT	= 25, !	Dot qualification after previous dot
PRIMARYSK_ACT_DOT_DOT_PLI	= 26, !	" in PL/I
PRIMARYSK_ACT_DOT_DOT_COB	= 27, !	" in COBOL
PRIMARYSK_ACT_DOT_SUBSCR	= 28, !	Subscripting after dot qualification
PRIMARYSK_ACT_DOT_SUBSCR_PLI	= 29, !	" in PL/I
PRIMARYSK_ACT_DOT_SUBSCR_COB	= 30, !	" in COBOL
PRIMARYSK_ACT_DOT_DEREF	= 31, !	PASCAL ^ after dot
PRIMARYSK_ACT_DOT_DEREF_PLI	= 32, !	PL/I "->" operator after "."
PRIMARYSK_ACT_DOT_TERM	= 33, !	Terminator operator after dot qualif.
PRIMARYSK_ACT_DOT_TERM_PLI	= 34, !	" in PL/I
PRIMARYSK_ACT_DOT_TERM_COB	= 35, !	" in COBOL
PRIMARYSK_ACT_DOT_TICK	= 50, !	Ada tick operator
PRIMARYSK_ACT_SUBSCR_DOT	= 36, !	Dot qualification after subscripting
PRIMARYSK_ACT_SUBSCR_DOT_PLI	= 37, !	" in PL/I
PRIMARYSK_ACT_SUBSCR_SUBSCR	= 38, !	Subscripting after subscripting
PRIMARYSK_ACT_SUBSCR_SUBSCR_PLI	= 39, !	" in PL/I
PRIMARYSK_ACT_SUBSCR_DEREF	= 40, !	PASCAL ^ after []
PRIMARYSK_ACT_SUBSCR_DEREF_PLI	= 41, !	PLI "->" operator after subscript
PRIMARYSK_ACT_SUBSCR_TERM	= 42, !	Terminator operator after subscripting
PRIMARYSK_ACT_SUBSCR_TERM_PLI	= 43, !	" in PL/I
PRIMARYSK_ACT_DEREF_DOT	= 44, !	Dot after dereferencing
PRIMARYSK_ACT_DEREF_SUBSCR	= 45, !	Subscripting after dereferencing
PRIMARYSK_ACT_DEREF_DEREF	= 46, !	Dereferencing after dereferencing
PRIMARYSK_ACT_DEREF_TERM	= 47, !	Terminator after dereferencing
PRIMARYSK_MAX_ACTION	= 50; !	--- Maximum action index

P R I N T I N F O R M A T I O N T A B L E

The Print Information Table is a blockvector indexed by FCODE which gives pointers to tables and print routine indexes necessary to print primary symbols for a given language.

```

3 3 2 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1
1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0

```

0	----- PRTINFO\$ROUT_INDEX -----
1	----- PRTINFO\$CHARTBL -----

A pointer to an Print Information Table is declared as follows:

```
PRINT_INFO_TABLE: REF PRTINFO$TABLE
```

Define the fields of the Print Information Table entry. Also define the declaration macro.

```
FIELD PRTINFO$FLD_DEF =
```

```

SET
PRTINFO$ROUT_INDEX   = [ 0, L_ ],
PRTINFO$CHARTBL      = [ 1, L_ ],
TES;

```

```
MACRO
```

```

PRTINFO$TABLE = BLOCKVECTOR[RST$K_TYPE_MAXIMUM + 1, 2, LONG]
FIELD(PRTINFO$FLD_DEF) %;

```

```
! Define valid PRTINFO$ROUT_INDEX literals.
```

```
LITERAL
```

```

PRT$K_MIN_ROUT      = 0,
PRT$K_OTHER         = 0,
PRT$K_ARRAY         = 1,
PRT$K_POINTER       = 2,
PRT$K_RECORD        = 3,
PRT$K_RECORD_COB    = 4,
PRT$K_POINTER_C     = 5,
PRT$K_RECORD_C_PL1  = 6,
PRT$K_VARIANT       = 7,
PRT$K_MAX_ROUT      = 7;

```

```
! Define valid Print Character Table Indexes.
```

```
LITERAL
```

PRTSK_BEGIN_CHAR = 0;
PRTSK_SEPARATOR_CHAR = 1;
PRTSK_END_CHAR = 2;
PRTSK_RECPTR_CHAR = 3;
PRTSK_MAX_PRTCHAR = 4;

REGISTER DESCRIPTORS

A Register Descriptor is used to describe the absolute address of a register or of a byte address within a register. It thus contains a register number, a byte offset within that register (0, 1, 2, and 3 are the only possible values), and a scope number which indicates to which call frame in the call stack this register belongs. The descriptor also contains a "sentinel" value which is used for validity checking and for ensuring that a Register Descriptor always has a non-zero value.

Register Descriptors are built by routine DBG\$STA_ADDRESS_TO_REGDESCR and are later used for various purposes when registers are examined. They are also required by routine DBG\$STA_REGISTER_NAME which generates the register name for printing.

A Register Descriptor fits in a single longword. This is the format:

```

 3 3 2 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1
 1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0
+-----+-----+-----+-----+
|          DBG$W_REGD_SCOPENUM          |SB_REGD_REGNUM|_SENTINEL|OFF|
+-----+-----+-----+-----+

```

A Register Descriptor is declared as follows:

```
REGDESCR: DBG$REGDESCR
```

Define the fields of the Register Descriptor. Also define the declaration macro.

```
FIELD DBG$REGD_FLD_DEF =
```

```

SET
DBG$V_REGD_OFFSET   = [ 0, V_(0,2) ],    ! Byte offset from start of register
DBG$V_REGD_SENTINEL = [ 0, V_(2,6) ],    ! Sentinel value--must be %X'2D'
DBG$B_REGD_REGNUM   = [ 0, BT_ ],        ! The actual register number
DBG$W_REGD_SCOPENUM = [ 0, W1_ ],        ! The register's scope number
TES;

```

```
MACRO
```

```
DBG$REGDESCR = BLOCK[1, LONG] FIELD(DBG$REGD_FLD_DEF) %;
```

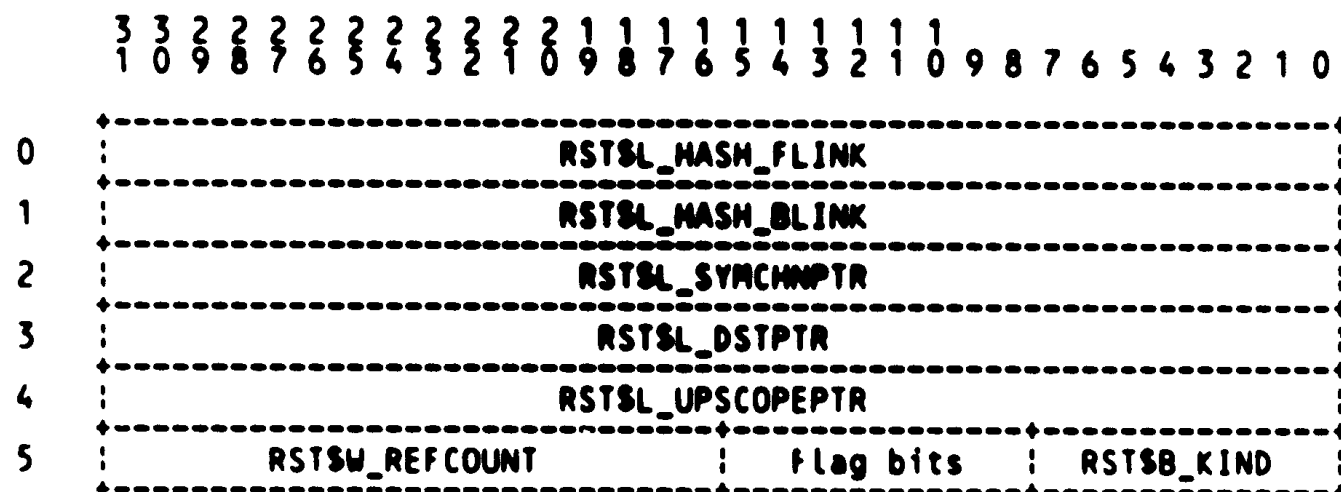
RUN - TIME SYMBOL TABLE DEFINITIONS

This section contains all definitions related to the structure and contents of the Run-Time Symbol Table (RST) through which the Debugger accesses the Debug Symbol Table (DST) in the user's executable image. The RST itself is built and accessed entirely through the Symbol-Table Access Routines in modules RSTCNTRL and RSTACCESS.

The RST is the structure through which the DST is accessed, one module at a time. This means that for each "active" module in the user's program, there is an RST data structure which describes that module and all lexical entities and data items in it. This structure, which is scattered through the Debugger's free memory, contains one entry for each lexical entity and data item in the module's DST. Each RST entry contains a pointer to the corresponding DST entry so all the DST entry information can be accessed rapidly. In addition, the RST itself is hashed by symbol name so that a symbol's RST and DST entries can be located rapidly given the symbol name.

RST ENTRY COMMON CORE

All entries in the Run-Time Symbol Table have the same information in their first parts. This RST entry "common core" is described here:



The hash chain forward and backward links link together all RST entries for symbols whose names have the same hash value. The hash chain is doubly linked so that RST entries can be removed from the hash chains when the whole RST for a module is removed to free up its space. Each hash chain has a permanent list head; RST entries in the RST Hash Table therefore never have zero hash links.

Some RST entries are not in the RST Hash Table, however. Such entries are "temporary" in the sense that they are created when needed and are thrown away when no longer referenced. Examples include many Data Type RST Entries and RST entries with invocation numbers. Such RST entries are put on the Temporary RST Entry List. This is a singly linked list which uses the RSTSL_HASH_FLINK field for the links; RSTSL_HASH_BLINK is left zero in this case.

All RST entries for a given module are linked together through the RSTSL_SYMNPTR pointer in a singly linked list. The list starts at the Module RST Entry and is terminated by a zero forward link. This list is traversed when the Module's whole RST is released to make room for another module.

The RSTSL_DSTPTR pointer gives the address of the symbol's DST entry. The RSTSL_UPSCOPEPTR pointer generally points to the RST entry of the symbol's "containing" entity. This usually means the containing lexical entity, but for data record components it means the containing record.

The RSTSB_KIND field identifies what kind of symbol and what kind of RST entry this is. The possible values are listed at the end of this section. The RSTSW_REFCOUNT field contains a reference count which speci-

ties how many references (pointers) to this RST entry have been passed to the rest of the Debugger. When this count is non-zero, this RST entry cannot be released to free storage. The reference count is thus used to prevent a "dangling pointer" problem.

A pointer to a Run-Time Symbol Table Entry is declared as follows:

RSTPTR: REF RST\$ENTRY

This declaration macro collects together all the individual FIELD sets declared below for the various kinds of RST entries.

MACRO

RST\$ENTRY = BLOCK[,LONG] FIELD(RST\$FLD_CORE,RST\$FLD_MOD,RST\$FLD_LEX,RST\$FLD_DATA) %;

! These are the RST entry common core field definitions.

FIELD RST\$FLD_CORE =

SET		
RST\$FL_HASH_FLINK	= [0, L_]	Symbol name hash chain forward link
RST\$FL_HASH_BLINK	= [1, L_]	Symbol name hash chain backward link
RST\$FL_SYMCNPTR	= [2, L_]	Pointer to the next RST entry belong- to the same module as this one
RST\$FL_DSTPTR	= [3, L_]	Pointer to this symbol's DST entry
RST\$FL_UPSCOPEPTR	= [4, L_]	Pointer to the RST entry one level up in scope from this one
RST\$B_KIND	= [5, B0_]	The kind of RST entry this is
RST\$V_GLOBAL	= [5, V1_(0)]	Flag set if this is a global symbol
RST\$V_NONZLENGTH	= [5, V1_(1)]	Flag set to TRUE if this symbol has a non-zero length
RST\$V_INVOCNUM	= [5, V1_(2)]	Flag set if symbol has invocation num.
RST\$V_SET_TYPEPTR	= [5, V1_(3)]	Flag set to indicate that the Data RST Entry RST\$FL_TYPEPTR field should be set during RST build.
RST\$V_MARKBIT	= [5, V1_(4)]	Mark bit available for temporary use (such as stopping recursion)
RST\$V_COBOLGBL	= [5, V1_(5)]	Flag bit set to indicate this symbol has the COBOL "global" attribute
RST\$V_REGISTER	= [5, V1_(6)]	Flag set if SYMID represents a register
RST\$W_REFCOUNT	= [5, W1_]	Reference count for storage management

TES;

!***

The following are the possible values for the RST\$B_KIND field. This field specifies what kind of RST entry this is; it thus indicates both what kind of symbol the RST entry represents and what the format of the RST entry is after the "common core" part.

LITERAL

RST\$K_INVALID	= 0,	Invalid code--cannot occur in RST en- try but can be returned by some Symbol Table Access routines
----------------	------	--

RST\$K_NOTUNIQUE = 9,	:	Symbol not unique--cannot occur in RST
	:	entry but can be returned by the
	:	DBG\$STA_GETSYMBOL routine.
RST\$K_MODULE = 1,	:	Module RST entry
RST\$K_ROUTINE = 2,	:	Routine RST entry
RST\$K_BLOCK = 3,	:	Block RST entry
RST\$K_ENTRY = 8,	:	Entry Point RST entry
RST\$K_LABEL = 4,	:	Label RST entry
RST\$K_LINE = 5,	:	Line number RST entry
RST\$K_DATA = 6,	:	Data item RST entry
RST\$K_TYPCOMP = 10,	:	Type component RST entry
RST\$K_TYPE = 7,	:	Data type RST entry
RST\$K_VARIANT = 11,	:	Record variant set RST entry
RST\$K_INVOCNUM = 12,	:	Invocation number RST Entry
RST\$K_OVERLOAD = 13,	:	Overloaded Symbol RST Entry
RST\$K_KIND_MINIMUM = RST\$K_INVALID,	:	Minimum possible kind value
RST\$K_KIND_MAXIMUM = RST\$K_OVERLOAD;	:	Maximum possible kind value

!***-

RST ENTRY FOR A MODULE

A module is represented by a Module RST Entry at all times, whether the rest of the module is in the RST or not. This Module Entry then points to a linked list of RST entries for the remaining symbols in the module via the RSTSL_SYMCHNPTR pointer; the end of this list is indicated by a zero pointer. The RSTSL_UPSCOPEPTR pointer is not used as an up-scope pointer in a Module Entry--there is nothing up-scope from a module. Instead RSTSL_NXTMODPTR, which occupies the same location in the RST entry as RSTSL_UPSCOPEPTR, is used to link all Module RST Entries in the whole program into a singly linked list. This list is scanned during RST initialization and when processing the SHOW MODULES command.

```

3 3 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1 1
1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0

```

0	RSTSL_HASH_FLINK		
1	RSTSL_HASH_BLINK		
2	RSTSL_SYMCHNPTR		
3	RSTSL_DSTPTR		
4	RSTSL_NXTMODPTR		
5	RSTSW_REFCOUNT	Flag bits	RST\$B_KIND
6	RSTSL_SAT_PTR		
7	RSTSL_PCTBL_BASE		
8	RSTSL_MODRSTSIZ		
9	RSTSL_MODSRCTBL		
10	Unused	RST\$B_LANGUAGE	Flag bits
11	RSTSL_MODPCTBL		
12	RSTSL_BASEVA		
13	Unused	SB_IMGFILCHAN	

The RSTSL_MODPCTBL field points to a table of pointers to the PC-Correlation Table DST records for this module (used for PC to line number correlation and vice versa). The first cell of this table (namely MOD_PC_TBL[0]) gives the number of PC-Correlation Table DST records in the module's DST, and the remaining cells (i.e., MOD_PC_TBL[1] through MOD_PC_TBL[MOD_PC_TBL[0]]) contain pointers to those DST records. Each

cell is one longword. If there are no PC-Correlation Table DST records for a given module, the RST\$\$_MODPCTBL field is zero.

PC-Correlation Table DST records need not have any special order or nesting with respect to the DST records for routines and lexical blocks and are therefore viewed as belonging to the module as a whole. The PC-Correlation Tables for a module are scanned by searching them in the order of their pointers in the RST\$\$_MODPCTBL table. The addresses given by the PC-Correlation Table DST records are relative to the value of the RST\$\$_PCTBL_BASE field, which contains the address of the lowest address routine in the module.

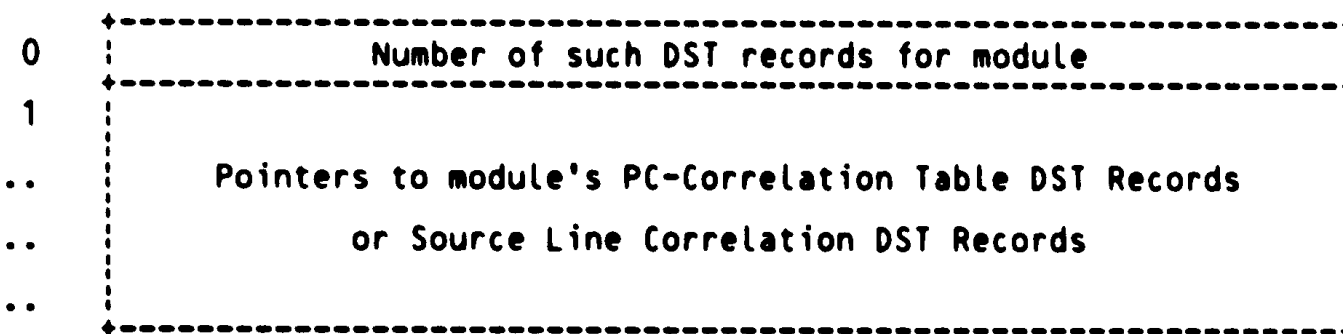
The RST\$\$_MODSRCTBL field points to a table of pointers to the Source Line Correlation DST Records for this module. The first cell of this table (i.e., MODSRCTBL[0]) contains the number of Source Line Correlation DST Records in this module's DST and the remaining cells (i.e., MODSRCTBL[1] through MODSRCTBL[MODSRCTBL[0]]) contain pointers to those DST records. Each cell in the table is a longword. If there are no such DST records at all, the RST\$\$_MODSRCTBL field is zero.

The memory blocks pointed to by RST\$\$_MODPCTBL and RST\$\$_MODSRCTBL both have the following format:

```

3 3 2 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1
1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0

```



When registers are examined using a numbered scope which does not correspond to any routine in the RST, a dummy Module RST Entry is created to represent that scope. In such an RST entry, the RST\$\$_MODNUMSCP bit is set to distinguish it from other Module RST Entries. In addition, the scope number (the number of call frames from the top of the call stack) of the desired scope is stored in the RST\$\$_MODSCPNUM field. This field overlays the RST\$\$_MODRSTSI7 field, which is not needed in this case. The name of such a "module" is just the scope number in decimal ASCII. These special Module RST Entries are built by the GET_REGISTER_SYMID routine and are recognized by the DBG\$\$_STA_SETCONTEXT routine, both in RSTACCESS.

This set of field declarations includes all fields outside the RST common core

! used in the Module RST Entry.

FIELD RST\$FLD_MOD =

SET		
RST\$L_NXTMODPTR = [4, L_],	Pointer to the next Module RST Entry	
RST\$L_SAT_PTR = [6, L_],	Pointer to this module's Static Address Table (SAT)	
RST\$L_PCTBL_BASE = [7, L_],	The base address for the offsets given by the PC Correlation Tables	
RST\$L_MODRSTSIZ = [8, L_],	The total size in bytes of the RST for this module--used only by the SHOW MODULE command.	
RST\$L_MODSCPNUM = [8, L_],	The scope number if the MODNUMSCP bit described below is set	
RST\$L_MODSRCTBL = [9, L_],	Pointer to table of pointers to the module's Source Line Correlation DST Records (or zero)	
RST\$V_MODSET = [10, V_(0)],	Flag indicating this module is SET	
RST\$V_MOD_IN_RST = [10, V_(1)],	Flag indicating module is in the RST	
RST\$V_ANONMOD = [10, V_(2)],	Flag indicating this is the "anonymous module" containing all unclaimed global symbols.	
RST\$V_MODNUMSCP = [10, V_(3)],	Flag indicating that this is a special RST entry for a numbered scope	
RST\$V_SHARE_IMAGE = [10, V_(4)],	Flag indicating that this is a shared image	
RST\$V_OLDPLIFLAG = [10, V_(5)],	Flag indicating this is a PL/I module compiled with an old compiler which generates incorrect VAX standard array descriptors	
RST\$B_LANGUAGE = [10, B1_],	The language of this module	
RST\$L_MODPCTBL = [11, L_],	Pointer to a table of pointers to the module's PC-Correlation Table DST records (or zero).	
RST\$L_BASEVA = [12, L_],	The base address of a shared image	
RST\$B_IMGFIChan = [13, B_]	The executable shared image file channel number	

TES;

! Define a symbolic name for the size of the Module RST Entry in longwords.

LITERAL

RST\$K_MODENTSIZ = 12;	Size of Module RST Entry
RST\$K_SHARED_MODENTSIZ = 14;	Size of Shared Image Module RST Entry


```
RST$L_BREAKADDR = [ 10, L_ ]    ! Routine breakpoint address (start
                                ! address, start address + 2, or
                                ! address from Prolog DST Record)
TES;
```

```
! Define symbolic names for the lengths of the lexical entity RST entries.
! All lengths are expressed in longwords.
```

```
LITERAL
```

```
RST$K_ROUTENTSIZ = 11,          ! Size of Routine RST Entry
RST$K_LEXENTSIZ   = 8,          ! Size of Lexical Block RST Entry
RST$K_EPTENTSIZ   = 11,         ! Size of Entry Point RST Entry
RST$K_LBLENTSIZ   = 7,          ! Size of Instruction Label RST Entry
RST$K_LINENTSIZ   = 8;          ! Size of Line Number RST Entry
```


The RST entry for an entry point (i.e., an alternate point through which a routine can be called) is shown here. The up-scope pointer points to the RST entry of the containing lexical entity. An entry point has no independent extent and therefore its RST entry has only a start address. The Entry Point RST Entry is always built when the whole module's RST is built.

3 3 2 2 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1 1
1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0

0	RSTSL_HASH_FLINK
1	RSTSL_HASH_BLINK
2	RSTSL_SYMCHNPTR
3	RSTSL_DSTPTR
4	RSTSL_UPSCOPEPTR
5	RSTSW_REFCOUNT Flag bits RST\$B_KIND
6	RSTSL_STARTADDR
7	unused
8	unused
9	unused
10	RSTSL_BREAKADDR

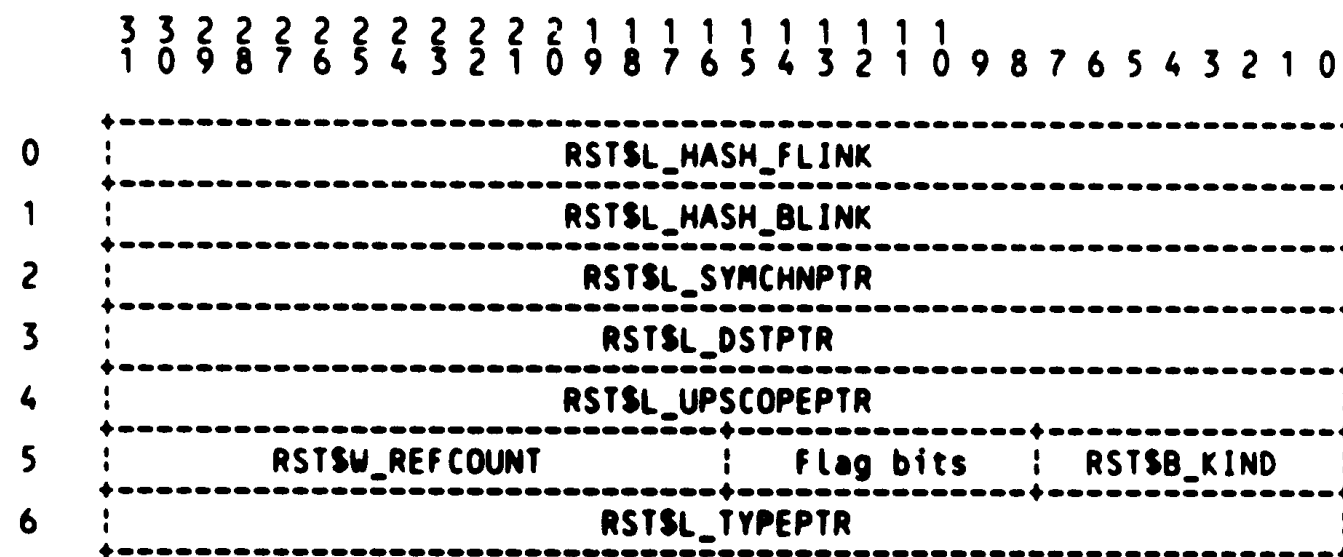
The RST entry for a label (which labels a point in the user's code) is shown here. The up-scope pointer points to the RST entry of the containing lexical entity. Since a label does not have extent, there is only a start address in the RST entry; no end address is given. The Label RST Entry is always built when the whole module's RST is built.

3 3 2 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1 1 1 0 9 8 7 6 5 4 3 2 1 0

0	RST\$L_HASH_FLINK		
1	RST\$L_HASH_BLINK		
2	RST\$L_SYMCHNPTR		
3	RST\$L_DSTPTR		
4	RST\$L_UPSCOPEPTR		
5	RST\$W_REFCOUNT	Flag bits	RST\$B_KIND
6	RST\$L_STARTADDR		

RST ENTRY FOR A DATA SYMBOL

The RST entry for a data symbol has the format shown here. The up-scope pointer points to the RST entry for the containing lexical entity (i.e., module, routine, or lexical block) or, if this is a data component, to the RST entry for the containing data object. The Data Symbol RST Entry also has a Type Pointer (RSTSL_TYPEPTR). For simple data types (those defined by one-byte type codes or standard descriptors) this pointer is zero, but for more complex data types (records, variants, and enumeration types) this pointer points to the RST entry for that data type. If the type pointer is zero, the type can be found directly in the data symbol's DST entry.



This set of field declarations includes the fields outside the RST common core used in the Data Symbol, Data Type Component, Data Type, Variant Set, and Invocation Number RST Entries. The individual illustrations show which RST entries use which fields.

FIELD RST\$FLD_DATA =

SET		
RSTSL_TYPEPTR	= [6, L_],	A pointer to the Data Type RST Entry which gives the data type of this data object. For simple types, this pointer is zero.
RSTSB_FCODE	= [6, B0_],	The "Format Code" of a Data Type RST Entry--indicates nature of type: array, atomic, record, etc.
RSTSW_TYPREFCNT	= [6, W1_],	The number of Data RST Entries that reference this Type RST Entry
RSTSL_TYPREFTBL	= [7, L_],	Pointer to Type Reference Table
RSTSL_BITSIZE	= [8, L_],	The length in bits of data items of this data type

```

RSTSL_DST_TYP_REC_PTR = [ 9, L_ ],
RSTSL_TYPCOMPNT= [ 10, L_ ],
RSTSA_TYPCOMPLST= [ 11, A_ ],
RSTSL_VARSETCNT = [ 2, L_ ],
RSTSL_VARTAGPTR = [ 4, L_ ],
RSTSA_VARSETTBL = [ 6, A_ ],
RSTSL_INVOCNUM = [ 6, L_ ]
TES;

```

A pointer to the DST record containing the embedded type-spec described by the DEBUG built DST record pointed to by RSTSL_DSTPTR.
 The number of record components if this is a data record type
 The start of a table of record component RST pointers
 The number of distinct variants in this record variant set
 A pointer to the RST entry of the tag variable for this variant set
 The start of a table of variant set members, giving the tag value and component list for each variant.
 The desired invocation number

! Declare symbolic names for the lengths of the various data, type, and invocation number RST entries. All lengths are in longwords.

LITERAL

```

RSTSK_DATENTSIZ = 7,
RSTSK_OLENTSIZ = 6,
RSTSK_TYSENTSIZ = 11,
RSTSK_VARENTSIZ = 6,
RSTSK_INVENTSIZ = 7;

```

! Size of the Data Symbol RST Entry and the Type Component RST Entry
 ! Size of the Overloaded Symbol RST Entry
 ! Size of the Data Type RST Entry
 ! Size of the Variant Set RST Entry
 ! Size of Invocation Number RST Entry

The RST entry for a data type component has the format shown here. A "component" in this sense is a record component, one of a set of variants, or one of the enumeration literals of an enumeration type. The interpretation of a data type component is thus dependent on the "type kind" of the containing data type.

Like the Data Item RST Entry, the this RST entry also has a type pointer which specifies the data type of the component. This is zero for simple types or a pointer to a Data Type RST Entry for a more complex data type (such as records).

3 3 2 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1 1 0 9 8 7 6 5 4 3 2 1 0

0	RSTSL_HASH_FLINK		
1	RSTSL_HASH_BLINK		
2	RSTSL_SYNCHNPTR		
3	RSTSL_DSTPTR		
4	RSTSL_UPSCOPEPTR		
5	RSTSW_REFCOUNT	Flag bits	RSTSB_KIND
6	RSTSL_TYPEPTR		

The RST entry for a data type has the format shown here. Not all data types have RST entries--those described by one-byte type codes or standard descriptors do not. Data Type RST Entries are used for all record, variant, and enumeration type data types, however. The up-scope pointer always points to the RST entry for the lexical entity (module, routine, or lexical block) in which the data type was declared. A data type need not have a name, however--it can be anonymous as it is in a PL/I or COBOL record declaration. Data Type RST Entries for records, variants, and enumeration types are built when the RST as a whole is built. All other Data Type RST Entries are built only when needed.

3 3 2 2 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1 1
1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0

0	RST\$L_HASH_FLINK		
1	RST\$L_HASH_BLINK		
2	RST\$L_SYMNPNTR		
3	RST\$L_DSTPTR		
4	RST\$L_UPSCOPEPTR		
5	RST\$W_REFCOUNT	Flag bits	RST\$B_KIND
6	RST\$W_TYPREFCNT	Unused	RST\$B_FCODE
7	RST\$L_TYPREFTBL		
8	RST\$L_BITSIZE		
9	RST\$L_DST_TYP_REC_PTR		
10	RST\$L_TYPCOMP CNT		
11	RST\$A_TYPCOMPLST		
..	Pointers to the Type Component RST Entries of		
..	the components of this record		
..	or to the Data Item RST Entries of the elements		
..	of this enumeration type		
..	(Not used for other kinds of data types)		
..			

!***

These are the possible values for the RST\$B_FCODE field.

LITERAL

RST\$K_TYPE_MINIMUM = 1,		--- Minimum possible FCODE value</td
RST\$K_TYPE_ARRAY	= 1,	Array type
RST\$K_TYPE_ATOMIC	= 2,	Atomic VAX standard type
RST\$K_TYPE_DESCR	= 3,	VAX standard descriptor type
RST\$K_TYPE_ENUM	= 4,	Enumeration type
RST\$K_TYPE_PICT	= 5,	Picture type (as in Cobol and PL/I)
RST\$K_TYPE_PTR	= 6,	Typed pointer type
RST\$K_TYPE_RECORD	= 7,	Record data type
RST\$K_TYPE_SET	= 8,	Set type
RST\$K_TYPE_SUBRNG	= 9,	Subrange data type
!	= 10,	Unused--available for future use
!	= 11,	Unused--available for future use
RST\$K_TYPE_COBHACK	= 12,	Cobol Hack data item
RST\$K_TYPE_BLIDATA	= 13,	Bliss data item
RST\$K_TYPE_BLIFLD	= 14,	Bliss field
RST\$K_TYPE_FILE	= 15,	File data type (as in Pascal)
RST\$K_TYPE_PTR	= 16,	Untyped pointer data type
RST\$K_TYPE_AREA	= 17,	Area type
RST\$K_TYPE_OFFSET	= 18,	Offset type
RST\$K_TYPE_VARIANT	= 19,	Variant Set (as in Pascal and ADA)
RST\$K_TYPE_RFA	= 20,	Record file address type
RST\$K_TYPE_SELF_REL_LAB	= 21,	Self relative label
RST\$K_TYPE_TASK	= 22,	Task type (as in ADA)
RST\$K_TYPE_MAXIMUM = 22;		--- Maximum possible FCODE value</td

!***

This is the format of a Variant Entry as pointed to by the pointers in the Variant-Set RST Entry. This entry gives the list of record components which comprise this particular record variant. It also gives a pointer to the Variant-Value DST record for this variant.

3 3 2 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1
1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0

0	RST\$L_VAR_DSTPTR
1	RST\$L_VAR_COMPCNT
2	RST\$A_VAR_COMPLST
..	List of RST pointers to the record
..	components in this variant
..	(One per longword)

Declarations for the fields in the Variant Entry. Also a declaration macro.

```
FIELD RST$FLD_VARENT =  
  SET  
  RST$L_VAR_DSTPTR = [ 0, L_ ], ! Pointer to Variant-Value DST record  
  RST$L_VAR_COMPCNT = [ 1, L_ ], ! Number of components in this variant  
  RST$A_VAR_COMPLST = [ 2, A_ ], ! Start of component RST pointer list  
  TES;  
  
MACRO  
  RST$VAR_ENTRY = BLOCK[,LONG] FIELD(RST$FLD_VARENT) %; ! Declaration macro
```


These definitions have something to do with variant record tag variables. Unfortunately, no comments were put in to say exactly what they are used for.

Define something or another.

FIELD RST\$TAG_BLOCK_FIELDS =

SET
RST\$L_TAG_NUMVALS = [0, L_], ! Number of values to follow (1 or 2)
RST\$L_TAG_LOWBOUND = [1, L_], ! Lower bound value
RST\$L_TAG_HIGHBOUND = [2, L_], ! Upper bound value
TES;

LITERAL

RST\$K_TAG_BLOCK_SIZE = 3; ! The size of the tag variable block

MACRO

RST\$TAG_LIST = BLOCKVECTOR[,RST\$K_TAG_BLOCK_SIZE]
FIELD(RST\$TAG_BLOCK_FIELDS) %;

!***

RST ENTRY FOR AN OVERLOADED SYMBOL

The RST entry for an Overloaded Symbol has the format shown below.
An Overloaded Symbol RST Entry is created for each Overloaded Symbol
DST record (DST\$K_OVERLOAD). The RST entry is hashed normally so that
DBG\$STA_GETSYMBOL can find it.

3 3 2 2 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1 1
1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0

0	-----+----- RST\$L_MASH_FLINK -----+-----																			
1	-----+----- RST\$L_MASH_BLINK -----+-----																			
2	-----+----- RST\$L_SYMCHNPTR -----+-----																			
3	-----+----- RST\$L_DSTPTR -----+-----																			
4	-----+----- RST\$L_UPSCOPEPTR -----+-----																			
5	-----+----- RST\$W_REFCOUNT -----+-----										: Flag bits : -----+-----					RST\$B_KIND -----+-----				

SCOPE LIST DEFINITIONS

The scope list maintains the list of scopes declared with the SET SCOPE command. It is a singly linked list pointed to by the global variable SCOPE\$LIST in module RSTCNTRL. Each scope entry on the list has this format:

```

3 3 2 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1
1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0

```

0	-----SCOPE\$L_FLINK-----
1	-----SCOPE\$L_STATE-----
2	-----SCOPE\$L_RSTPTR-----
3	-----SCOPE\$L_MODPTR-----

A pointer to a scope list entry is declared as follows:

```
SCOPEPTR: REF SCOPE$ENTRY
```

Declare the scope entry fields and the definition macro.

```

FIELD SCOPE$FLD_DEF =
SET
SCOPE$L_FLINK   = [ 0, L_ ],      ! Scope list forward link or zero
SCOPE$L_STATE  = [ 1, L_ ],      ! The scope "state" or kind
SCOPE$L_RSTPTR = [ 2, L_ ],      ! Pointer to scope's own RST entry
SCOPE$L_MODPTR = [ 3, L_ ],      ! Pointer to scope's Module RST Entry
TES;

```

```

LITERAL
SCOPE$K_ENTSIZE = 4;           ! Size of scope entry in longwords

```

```

MACRO
SCOPE$ENTRY = BLOCK[SCOPE$K_ENTSIZE] FIELD(SCOPE$FLD_DEF) %;

```

These are the possible values of the SCOPE\$L_STATE field. Each of these values indicates what kind of scope is to be searched to match a given pathname.

```

LITERAL
SCOPE$K_NORMAL   = 1,          ! Normal named scope
SCOPE$K_NUMBERED = 2,          ! Numbered scope (PC in CALL stack)
SCOPE$K_GLOBAL   = 3,          ! Global Symbol Table
SCOPE$K_SETMODS  = 4,          ! All SET modules

```

THE SCREEN DISPLAY ENTRY

The Screen Display Entry contains all information needed to represent and output a screen display on the terminal screen. All Display Entries are linked together by forward and backward links. Each Display Entry contains all information about where the corresponding display is placed on the screen, what attributes it has, and what the contents (the text) of the display is. The text is represented by pointers to the Screen Display Line Entries for the text lines that make up the display.

3 3 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1 1
1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0

0	-----	DBGSL_DISP_FLINK	-----
1	-----	DBGSL_DISP_BLINK	-----
2	-----	DBGSW_DISP_FLAGS	SB_DISP_REND : SB_DISP_KIND
3	-----	Unused	DBGSW_DISP_SCROLL
4	-----	DBGSW_DISP_RLEN	DBGSW_DISP_RBEG
5	-----	DBGSW_DISP_CLEN	DBGSW_DISP_CBEG
6	-----	DBGSW_DISP_DCOL	DBGSW_DISP_DROW
7	-----	DBGSW_DISP_LINECNT	DBGSW_DISP_MAX_LINECNT
8	-----	DBGSL_DISP_START_LINE_PTR	-----
9	-----	DBGSL_DISP_END_LINE_PTR	-----
10	-----	DBGSL_DISP_WINDOW_PTR	-----
11	-----	DBGSL_DISP_ERROR_PTR	-----
12	-----	DBGSL_DISP_OLDTXT_PTR	-----
13	-----	DBGSL_DISP_MODPTR	-----
14	-----	DBGSL_DISP_CENTER	-----
15	-----	DBGSL_DISP_MARKLINE	-----
16	-----	DBGSL_DISP_MINLINE	-----
17	-----	DBGSL_DISP_MAXLINE	-----
18	-----	DBGSL_DISP_CMDLIST	-----

19

DBG\$A_DISP_NAME

..

The display name in Counted ASCII

..

..

All Screen Display Entries are linked together on a doubly linked list for which the OWN block DBG\$SCR_DISPLAY_LIST is the list head. Also, the DBG\$L_DISP_START_LINE_PTR and DBG\$L_DISP_END_LINE_PTR fields of the Display Entry constitute the list head for the doubly linked list of Screen Display Line Entries associated with this display.

A pointer to a Display Entry is declared as follows:

DISP_PTR: REF DBG\$DISP_ENTRY

Define the fields of the Screen Display Entry.

FIELD DBG\$DISP_FLD =

SET

DBG\$L_DISP_FLINK	= [0, L_]	:	Screen display list forward link
DBG\$L_DISP_BLINK	= [1, L_]	:	Screen display list backward link
DBG\$B_DISP_KIND	= [2, B0_]	:	Screen display kind
DBG\$B_DISP_REND	= [2, B1_]	:	Screen display rendition bits
DBG\$V_DISP_REND_BLD	= [2, V1_ (0)]	:	Bolding rendition bit
DBG\$V_DISP_REND_RV	= [2, V1_ (1)]	:	Reverse-video rendition bit
DBG\$V_DISP_REND_BLK	= [2, V1_ (2)]	:	Blinking rendition bit
DBG\$V_DISP_REND_UND	= [2, V1_ (3)]	:	Underline rendition bit
DBG\$W_DISP_FLAGS	= [2, W1_]	:	Screen display flag bits
DBG\$V_DISP_REMOVE	= [2, V2_ (0)]	:	Remove display from pasteboard
DBG\$V_DISP_INVSCR	= [2, V2_ (1)]	:	Invalidate DISP_SCROLL field
DBG\$V_DISP_ATTOP	= [2, V2_ (2)]	:	Display at top of source file
DBG\$V_DISP_ATBOT	= [2, V2_ (3)]	:	Display at bottom of source file
DBG\$V_DISP_MARKFLG	= [2, V2_ (4)]	:	Mark changed lines in display
DBG\$V_DISP_NEWDISP	= [2, V2_ (5)]	:	New display--suppress marking
DBG\$W_DISP_SCROLL	= [3, SW0_]	:	Screen display scrolling count
DBG\$W_DISP_RBEG	= [3, W1_]	:	Unused (available for future use)
DBG\$W_DISP_RLEN	= [4, W0_]	:	Screen window beginning row location
DBG\$W_DISP_CBEG	= [4, W1_]	:	Screen window row length (height)
DBG\$W_DISP_CLEN	= [5, W0_]	:	Screen window beginning column location
DBG\$W_DISP_DROW	= [5, W1_]	:	Screen window column length (width)
DBG\$W_DISP_DCOL	= [6, W0_]	:	Screen window display row location
DBG\$W_DISP_MAX_LINECNT	= [6, W1_]	:	Screen window display column location
DBG\$W_DISP_LINECNT	= [7, W0_]	:	Screen display's maximum line count
DBG\$W_DISP_LINECNT	= [7, W1_]	:	(maximum lines saved in memory)
DBG\$L_DISP_START_LINE_PTR	= [8, L_]	:	Screen display's actual line count
		:	Pointer to first line of display text
		:	(points to display line entry)


```

DBGSL_DISP_END_LINE_PTR = [ 9, L_ ], : Pointer to last line of display text
                                : (points to display line entry)
DBGSL_DISP_WINDOW_PTR = [ 10, L_ ], : Pointer to first line in screen window
                                : (points to a Display Line Entry)
DBGSL_DISP_ERROR_PTR = [ 11, L_ ], : Pointer to list of error message
                                : Display Line Entries for display
DBGSL_DISP_OLDTXT_PTR = [ 12, L_ ], : Pointer to list of old Display Line
                                : Entries--used to mark changes
DBGSL_DISP_MODPTR = [ 13, L_ ], : Pointer to Module RST Entry--used for
                                : source line displays only
DBGSL_DISP_CENTER = [ 14, L_ ], : Central line number for source display
DBGSL_DISP_MARKLINE = [ 15, L_ ], : Line number to be marked with "-->" in
                                : source line display (or -1)
DBGSL_DISP_MINLINE = [ 16, L_ ], : Minimum source line number in module
DBGSL_DISP_MAXLINE = [ 17, L_ ], : Maximum source line number in module
DBGSL_DISP_CMDLIST = [ 18, L_ ], : Pointer to display's DEBUG command
                                : list entry (or zero)
DBGSA_DISP_NAME = [ 19, A_ ] : The display's name in Counted ASCII
                                : TES;

```

MACRO

```
DBG$DISP_ENTRY = BLOCK[,LONG] FIELD(DBG$DISP_FLD) %;
```

LITERAL

```
DBG$K_DISP_ENTSIZE = 19; : Size of the fixed part of the Screen
                        : Display Entry in longwords
```

```
! Define the literals used to indicate the display kind. The display kind
! determined how the contents of the display are generated.
```

LITERAL

```

DBG$K_DISP_MINKIND = 0, : ---Minimum kind value
DBG$K_DISP_NOKIND = 0, : No kind--used in calls to indicate no
                        : change to the display's kind
DBG$K_DISP_NORMAL = 1, : Normal display kind--contents deter-
                        : mined by direct user input
DBG$K_DISP_DO = 2, : Contents specified by DO command list
DBG$K_DISP_SOURCE = 3, : Source display with contents speci-
                        : fied by a source command list
DBG$K_DISP_REGISTER = 4, : Contents is machine register display
DBG$K_DISP_MAXKIND = 4; : ---Maximum kind value

```

```
! Define literals for the display rendition bits.
```

LITERAL

```

DBG$M_DISP_REND_BLD = 1, : Bolding rendition mask
DBG$M_DISP_REND_RV = 2, : Reverse-video rendition mask
DBG$M_DISP_REND_BLK = 4, : Blinking rendition mask
DBG$M_DISP_REND_UND = 8; : Underline rendition mask

```

```
! Define the literals used to indicate the scrolling direction when scrolling
! a screen display.
```

LITERAL

DBG\$K_SCROLL_UP	= 1,	! Scroll the display up
DBG\$K_SCROLL_DOWN	= 2,	! Scroll the display down
DBG\$K_SCROLL_LEFT	= 3,	! Scroll the display to the left
DBG\$K_SCROLL_RIGHT	= 4,	! Scroll the display to the right

THE SCREEN DISPLAY LINE ENTRY

The Screen Display Line Entry holds the text and other attributes of a line within a screen display. This entry contains forward and backward links which link this Line Entry to the Line Entries for the previous line and the next line of the same display. The Line Entry also contains the line's rendition attributes (such as reverse video, blinking, etc.) which are to be passed to the Screen Management Package. Finally it contains the actual text of the line, including the length of that text, and the length of the text buffer.

Screen Display Line Entries come in two variants, the Normal and the Source Display Line Entry. Source Display Line Entries are used for source displays as the name implies, and Normal Display Line Entries are used for all other kinds of displays. Normal Display Line Entries can optionally include rendition information on a per-character basis.

NORMAL DISPLAY LINE ENTRY

The Normal Display Line Entry is used for all screen displays except Source screen displays. This is the format of the Normal Screen Display Line Entry:

3 3 2 2 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1
1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0

```

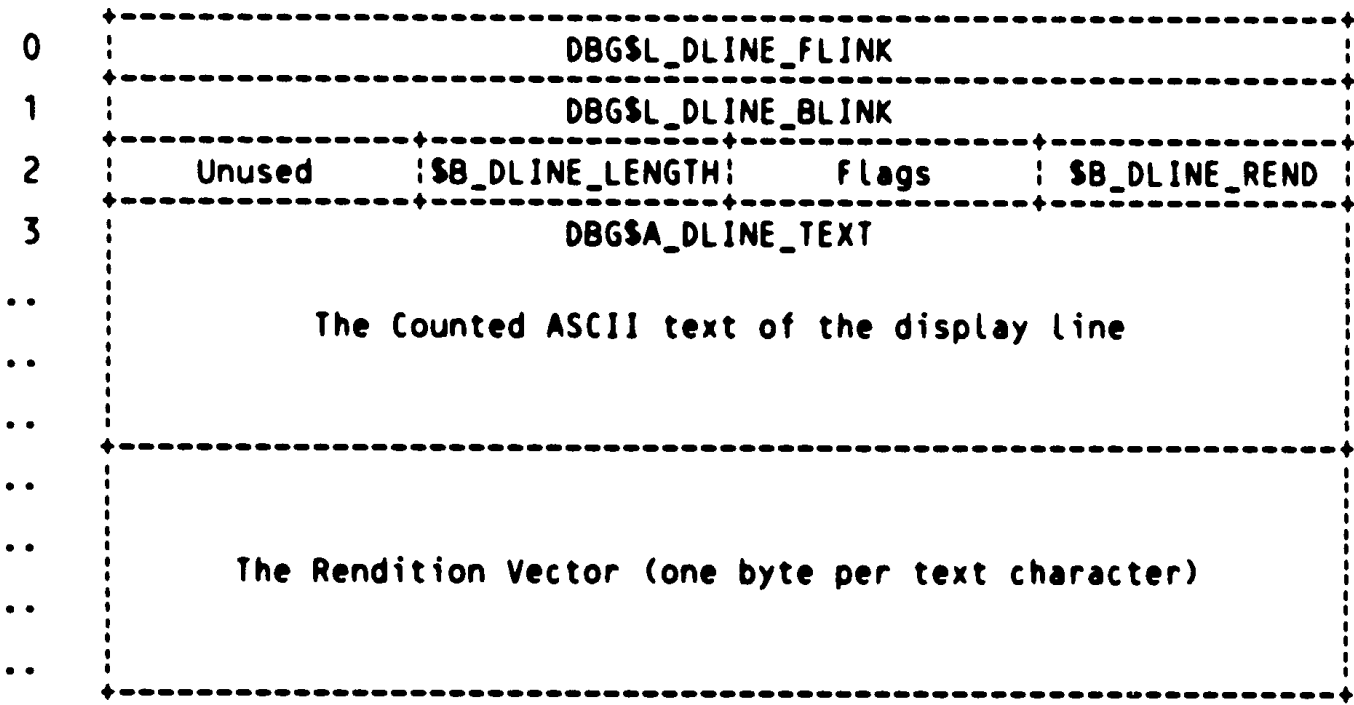
0  |-----+-----+-----+-----+
   |          DBG$DLINE_FLINK
1  |-----+-----+-----+-----+
   |          DBG$DLINE_BLINK
2  |-----+-----+-----+-----+
   |  Unused  |$B_DLINE_LENGTH|  Flags  |$B_DLINE_REND
3  |-----+-----+-----+-----+
   |          DBG$A_DLINE_TEXT
..
..
..
..

```

NORMAL DISPLAY LINE ENTRY WITH RENDITION VECTOR

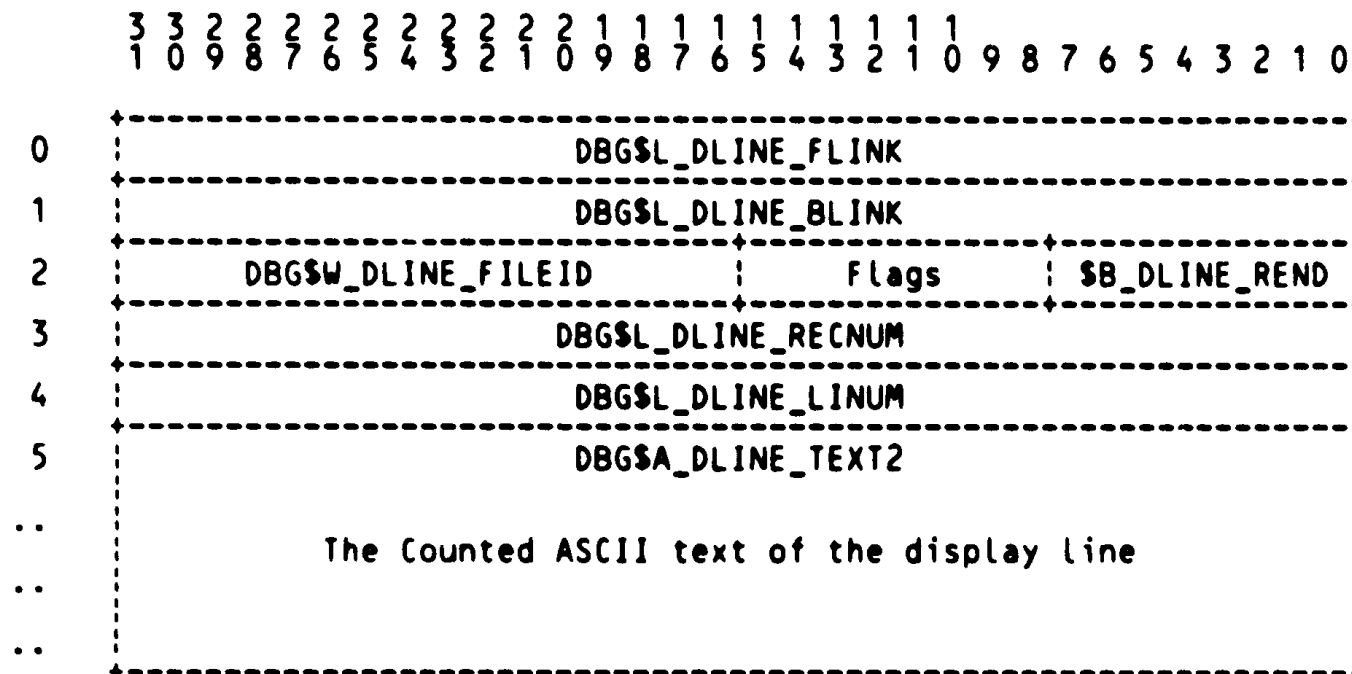
The Normal Display Line Entry may include a "rendition vector" if the screen rendition is not the same for the entire line. In this case, the DBG\$V_DLINE_RENDFLG is set in the Display Line Entry's flag field and a "rendition vector" is present at the end of the Display Line Entry, immediately following the Counted ASCII text of the display line. The rendition vector contains one byte of rendition codes for each character in the Counted ASCII text of the line. The rendition vector has the same length as the Counted ASCII text.

3 3 2 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1
1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0



SOURCE DISPLAY LINE ENTRY

The Source Display Line Entry is used for source displays only, and contains extra line number information which is not needed for other kinds of displays. The DBG\$V_DLINE_SOURCEFLG bit is always set in a Source Display Line Entry. This is the format of the Source Display Line Entry:



Each Display Line Entry is part of a doubly linked list of such entries associated with some display. The fields DBG\$L_DISP_START_LINE_PTR and DBG\$L_DISP_END_LINE_PTR in the display's Screen Display Entry constitute the list head for this list of Display Line Entries.

The Display Line Entries associated with a display are reused as new text is added to the display contents. For this reason, each entry contains not only the length of the current text line but the length of the entry's text buffer as well. The text buffer will be longer than the current text line if it contained a longer text line earlier.

A pointer to the Screen Display Line Entry is declared as follows:

```
DLIN_PTR: REF DBG$DLIN_ENTRY
```

Define the fields of the Screen Display Line Entry.

FIELD DBG\$DLINE_FLD =

```

SET
DBG$DLINE_FLINK = [ 0, L_ ], ! Forward link to next line entry
DBG$DLINE_BLINK = [ 1, L_ ], ! Backward link to previous line entry
DBG$DLINE_REND = [ 2, B0_ ], ! Rendition bits for this line
DBG$DLINE_REND_BLD = [ 2, V0_(0) ], ! Bolding rendition bit
DBG$DLINE_REND_RV = [ 2, V0_(1) ], ! Reverse-video rendition bit
DBG$DLINE_REND_BLK = [ 2, V0_(2) ], ! Blinking rendition bit
DBG$DLINE_REND_UND = [ 2, V0_(3) ], ! Underline rendition bit
! The flag byte
DBG$DLINE_RENDFLG = [ 2, V1_(0) ], ! Line includes rendition vector
DBG$DLINE_SOURCEFLG = [ 2, V1_(1) ], ! Set for Source Disp Line Entry
!
DBG$DLINE_LENGTH = [ 2, B2_ ], ! The length of the text buffer (bytes)
DBG$DLINE_TEXT = [ 3, A_ ], ! The line's text in Counted ASCII
!
DBG$DLINE_FILEID = [ 2, W1_ ], ! Source File ID of source line
DBG$DLINE_RECNUM = [ 3, L_ ], ! Source file record number of line
DBG$DLINE_LINUM = [ 4, L_ ], ! Line number of source line
DBG$DLINE_TEXT2 = [ 5, A_ ], ! Source line's text in Counted ASCII
TES;

```

MACRO

DBG\$DLINE_ENTRY = BLOCK[,LONG] FIELD(DBG\$DLINE_FLD) %;

LITERAL

```

DBG$K_DLINE_ENTSIZE = 3, ! Size of fixed part of Normal Display
                        ! Line Entry in longwords
DBG$K_DLINE_ENTSIZE2 = 5; ! Size of fixed part of Source Display
                        ! Line entry in longwords

```

THE SCREEN PASTEBOARD ENTRY

The screen pasteboard, on which all displays to be output on the terminal are pasted to determine what occludes what, is represented by a blockvector. The vector is indexed by line number and has one entry (block) for each line on the screen. The length of the vector is actually 21 (not 24) so that the last three lines on the screen are always reserved for user program output and DEBUG input prompting. For each line, the vector contains one Pasteboard Entry of the following format:

	3	3	2	2	2	2	2	2	2	2	2	2	1	1	1	1	1	1	1	1	1	1	0	9	8	7	6	5	4	3	2	1	0
0	DBG\$W_PASTE_SCROLL										SB_PASTE_REND					SB_PASTE_KIND																	
1	DBG\$L_PASTE_DISPID																																
2	DBG\$L_PASTE_DLINE																																

Define the fields of the individual Pasteboard Entry. Also declare the declaration macro and the size of the vector and of each block.

```
FIELD DBG$PASTE_FLD =
SET
DBG$B_PASTE_KIND = [ 0, B0_ ], ! The kind of line this entry holds
DBG$B_PASTE_REND = [ 0, B1_ ], ! The rendition bits for this line
DBG$W_PASTE_SCROLL = [ 0, SW_ ], ! The scrolling amount for this line
DBG$L_PASTE_DISPID = [ 1, L_ ], ! Pointer to Display Entry for line
DBG$L_PASTE_DLINE = [ 2, L_ ], ! Pointer to Display Line Entry
TES;
```

```
MACRO
DBG$PASTEBOARD =
BLOCKVECTOR(DBG$K_PASTE_SIZE, DBG$K_PASTE_ENTSIZE, LONG)
FIELD(DBG$PASTE_FLD) %;
```

```
LITERAL
DBG$K_PASTE_SIZE = 21, ! Number of lines in pasteboard
DBG$K_PASTE_ENTSIZE = 3; ! Size of one Pasteboard Entry in longwords
```

Define the allowed values of the DBG\$B_PASTE_KIND field.

```
LITERAL
DBG$K_PASTE_NULL = 1, ! Null line--no display covers it
DBG$K_PASTE_TEXT = 2, ! Text line from a display
DBG$K_PASTE_BLANK = 3, ! Blank line within a display
```

DBGSK_PASTE_LABEL = 4; ! Top border line with label
DBGSK_PASTE_BORDER = 5; ! Bottom border line on screen

! Screen Window Entry in longwords

S O U R C E D I R E C T O R Y S E A R C H L I S T

The Source Directory Search List is a list structure which keeps track of all source directory names specified by the user via SET SOURCE and SET SOURCE/MODULE=xxx commands. This structure is thus searched when a source file must be opened so that the file is opened in the proper directory.

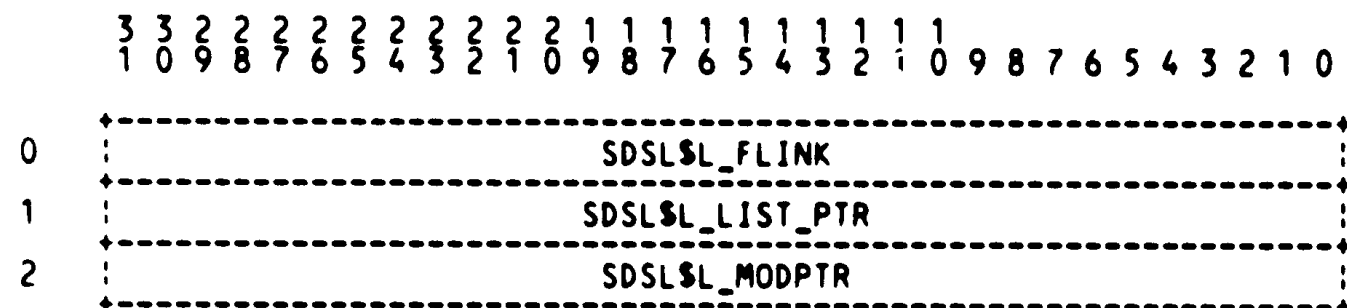
The structure of this list is as follows. Variable DBF\$SRC_DIR_LIST in module DBG\$SOURCE points to a singly linked list of Source Directory Search List Header Blocks. Each Header Block contains a Module RST Entry pointer (which may be zero) and points to a singly linked list of Source Directory Search List Entries. Each such entry contains a directory name as a Counted ASCII string.

The list is searched by first searching all the Header Blocks for the Module RST Entry pointer for the current module, i.e. the module from which source lines are to be displayed. If the desired module is not found, the Header Block with the zero Module RST Entry pointer is used instead. The found Header Block then points to the linked list of directory names to use when searching for the desired source file. If there is no Header Block with either the desired module pointer or the zero module pointer, no Source Directory Search List Entries are used--the file name from the Declare Source File DST command is used as is.

The command SET SOURCE/MODULE=modname dir1,...,dirN is represented by a Header Block with a pointer to the Module RST Entry for module 'modname' and a pointer to a linked list of Entries. There is one Entry for each of dir1, ..., dirN. The command SET SOURCE dir1,...,dirN is represented by a Header Block with a zero Module RST Entry pointer and of course a pointer to a linked list of Entries for dir1, ..., dirN.

SOURCE DIRECTORY SEARCH LIST HEADER BLOCK

The Source Directory Search List Header Block, as described on the previous page, has the following format:



A pointer to a Source Directory Search List Header Block is declared as follows:

```
SDSL_PTR: REF SDSL$HEADER
```

Declare the fields of the Source Directory Search List Header Block. Also declare the declaration macro.

```
FIELD SDSL$FLD_HEADER =
  SET
    SDSL$FLINK      = [ 0, L_ ],      ! Forward link to next Header Block
    SDSL$LIST_PTR   = [ 1, L_ ],      ! Pointer to list of SDSL Entries
    SDSL$MODPTR     = [ 2, L_ ],      ! Pointer to Module RST Entry of module
                                         to which search list applies; is
                                         zero if for all other modules
  TES;

LITERAL SDSL$K_HDR_SIZE = 3;          ! Size of SDSL Header Block in longwords

MACRO SDSL$HEADER = BLOCK[SDSL$K_HDR_SIZE] FIELD(SDSL$FLD_HEADER) %;
```

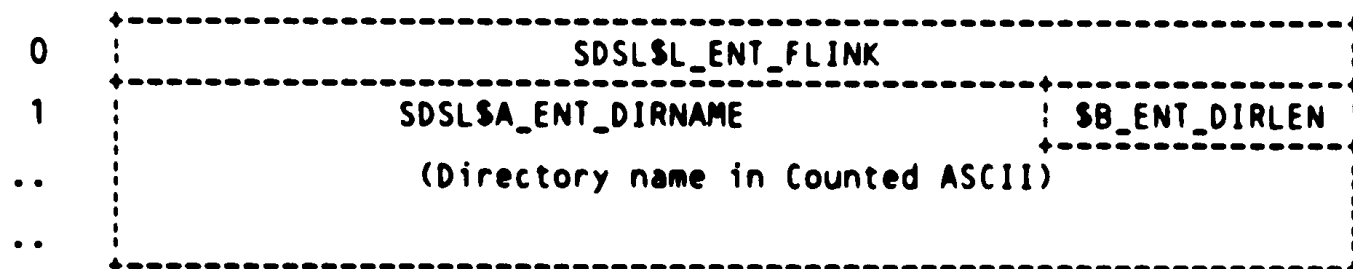
SOURCE DIRECTORY SEARCH LIST ENTRY

There is one Source Directory Search List Entry for each directory name specified on a SET SOURCE command. The "directory name" may in fact contain a full file name, including extension and version number, but in most cases it is used only to override the directory name as given in the Declare Source File DST command. This is the format of a Source Directory Search List Entry:

```

3 3 2 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1 1
1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0

```



A pointer to a Source Directory Search List Entry is declared like this:

```
SDSL_ENT_PTR: REF SDSL$ENTRY
```

Declare the fields of the Source Directory Search List Entry. Also declare the declaration macro.

```

FIELD SDSL$FLD_ENTRY =
SET
SDSL$ENT_FLINK = [ 0, L_ ], ! Forward link to next SDSL Entry
SDSL$B_ENT_DIRLEN = [ 1, B0_ ], ! Length in characters of directory name
SDSL$A_ENT_DIRNAME = [ 1, A1_ ], ! First character of directory name
TES;

```

```

LITERAL
SDSL$K_ENT_SIZE = 1;           ! Size of fixed portion of SDSL Entry
                                ! in longwords

```

```
MACRO SDSL$ENTRY = BLOCK[,LONG] FIELD(SDSL$FLD_ENTRY) %;
```

SOURCE FILE CONTROL BLOCK

This is the format of the Source File Control Block:

	3	3	2	2	2	2	2	2	2	2	2	1	1	1	1	1	1	1	1	1	0	9	8	7	6	5	4	3	2	1	0	
0	SFCBSL_FLINK																															
1	SFCBSL_BLINK																															
2	Unused (MBZ)																SFCBSB_KIND															
3	SFCBSW_RFACURLN										SFCBSW_RF SPACING																					
4	SFCBSL_RFATBLPTR																															
5	SFCBSL_DSTPTR																															
6	SFCBSL_FABPTR																															
7	SFCBSL_RABPTR																															
8	SFCBSL_NAMPTR																															
9	SFCBSL_XABDATPTR																															
10	SFCBSL_XABFHCPTR																															
11	SFCBSL_NAMBUFFER																															
12	SFCBSL_CURRECNUM																															
13	SFCBSL_CUR_RFA0																															
14	Unused										SFCBSW_CUR_RFA4																					
15	SFCBSL_LBRINDEX																															

The Source File Control Blocks (SFCBs) keep track of all currently open source files. This includes modules within source libraries as well as normal RMS files. Each such block contains all information needed to read source records from the corresponding file. These blocks form a circular doubly linked list, where variable DBG\$SRC_SFCB_PTR in module DBG\$SOURCE points to the first block on the list. The blocks are always maintained in order so that the Most Recently Used SFCB is first on the list and the Least Recently Used block is last. This ordering allows DEBUG to close the Least Recently Used source file when a file must be closed in order to open a new file.

DEBUG always keeps a fixed number of Source File Control Blocks on the list. This number is given by variable DBG\$SRC MAX FILES in module DBG\$SOURCE. It can be changed by the user with the SET MAX_SOURCE_FILES command. This number limits the number of source files which are open at the same time. Such a limit is required to prevent DEBUG from using up too many of the limited number of channels the user can have open at any one time. Thus, when all SFCBs are in use and a new source file must be opened, the Least Recently Used file is closed and its SFCB is reused for the new file.

A pointer to a Source File Control Block is declared as follows:

SFCB_PTR: REF SFCB\$BLOCK

Declare the fields of the Source File Control Block. Also declare the declaration macro.

```
FIELD SFCB$FLD_DEF =
SET
SFCB$FLINK      = [ 0, L_ ],      ! Forward link to next SFCB
SFCB$BLINK      = [ 1, L_ ],      ! Backward link to previous SFCB
SFCB$B_KIND     = [ 2, BO_ ],      ! The kind of source file this is
SFCB$RFA$PACING = [ 3, WO_ ],      ! The allocated length of the RFA table
SFCB$RFA$CURLN  = [ 3, WI_ ],      ! The used length of the RFA table
SFCB$RFA$BLPTR  = [ 4, L_ ],      ! Pointer to Record File Address (RFA)
                                     ! table for this source file
SFCB$DSTPTR     = [ 5, L_ ],      ! Pointer to DST Declare Source File
                                     ! command for this source file
SFCB$FABPTR     = [ 6, L_ ],      ! Pointer to the FAB for this file
SFCB$RABPTR     = [ 7, L_ ],      ! Pointer to the RAB for this file
SFCB$NAMPTR     = [ 8, L_ ],      ! Pointer to the NAM block for this file
SFCB$XABDATPTR  = [ 9, L_ ],      ! Pointer to the XAB Date and Time block
                                     ! for this file
SFCB$XABFHCPTR  = [ 10, L_ ],     ! Pointer to XAB File Header Character-
                                     ! istics block for this file
SFCB$NAMBUFFER  = [ 11, L_ ],     ! Pointer to buffer for NAM block file
                                     ! name strings for this file
SFCB$CURRECNUM  = [ 12, L_ ],     ! Record number at which the source file
                                     ! is currently positioned
SFCB$CUR_RFA0   = [ 13, L_ ],     ! RFA at which the source file is cur-
                                     ! rently positioned (first 4 bytes)
SFCB$CUR_RFA4   = [ 14, WO_ ],     ! RFA at which the source file is cur-
                                     ! rently positioned (last 2 bytes)
SFCB$LBRINDEX   = [ 15, L_ ]      ! Library file index used by librarian
TES;
```

```
LITERAL
SFCB$K_SIZE      = 16;           ! Size of SFCB in longwords
```

```
MACRO SFCB$BLOCK = BLOCK[SFCB$K_SIZE] FIELD(SFCB$FLD_DEF) %;
```

```
! Declare the possible values of the SFCB$B_KIND field.
```

LITERAL

SFCBSK_NOTUSED = 1,
SFCBSK_RMSFILE = 2,
SFCBSK_LBRFILE = 3,

SFCBSK_FILE_UNAVAIL = 4;

! This SFCB is not used (is available)
! This SFCB is used for an RMS file
! This SFCB is used for a module within
! a source library
! This SFCB is occupied by a file which
! is presently not available

SOURCE FILE ID TABLE

The Source File ID Table is used during the decoding of the Source Line Correlation DST Record to keep track of the File IDs used to reference the current module's source files. The Source File ID Table consists of a singly linked list of Source File ID Table Entries, where each entry gives a File ID, a pointer to the corresponding Declare Source File command, (This command gives all needed information about the identity and nature of the source file.), and the current source record number. This table is used locally in routine DBG\$SRC_TYPE_LNUM_SOURCE in module DBG\$SOURCE. This is the format of each Source File ID Table Entry:

```

3 3 2 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1
1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0

```

0	-----SFIT\$L_FLINK-----
1	-----SFIT\$L_FILE_ID-----
2	-----SFIT\$L_DSTPTR-----
3	-----SFIT\$L_CURRECNUM-----

A pointer to a Source File ID Table Entry is declared as follows:

```
SFIT_PTR: REF SFIT$ENTRY
```

Declare the fields of the Source File ID Table Entry. Also declare the declaration macro.

```

FIELD SFIT$FLD_DEF =
SET
SFIT$L_FLINK      = [ 0, L_ ], ! Forward link to next SFIT entry
SFIT$L_FILE_ID   = [ 1, L_ ], ! File ID value for this source file
SFIT$L_DSTPTR    = [ 2, L_ ], ! Pointer to DST Declare Source File
                                command for this source file
SFIT$L_CURRECNUM = [ 3, L_ ]   ! Current source record number for
                                this source file
TES;

```

```

LITERAL
SFIT$K_SIZE      = 4;          ! Size of SFIT entry in longwords

```

```
MACRO SFIT$ENTRY = BLOCK[SFIT$K_SIZE] FIELD(SFIT$FLD_DEF) %;
```

S O U R C E F I L E R E C O R D F I L E A D D R E S S T A B L E

This is the format of each Record File Address Table entry:

3 3 2 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1 1
1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0

0	RFATBL\$RFA0																			
1											RFATBL\$W_RFA4									

The Record File Address Table correlates record numbers in a source file with RFA's for that file. It is used during the TYPE command and other commands that cause source lines to be displayed, to help locate the desired record more efficiently.

There is one RFA table for each open file. The table is allocated during DBG\$SRC_INIT and a pointer to the table is placed in the Source File Control Block. The table consists of a fixed number of entries (RFATBL_SIZE_ENTRIES, defined in DBGSOURCE).

For every N records in the file, there is one RFA entry. That is, the first entry is the RFA for record 1, the next for record (1+N), the next for record (1+2N), and so on. The spacing N can be adjusted dynamically and is kept in the Source File Control Block.

A pointer to an RFA table is declared as follows:

```
RFATBL_PTR : REF RFATBL$BLOCKVECTOR(RFATBL_SIZE_ENTRIES)
```

Declare the fields of an RFA table entry. Also declare the declaration macro.

```
FIELD RFATBL$FLD_DEF =
```

```
SET
RFATBL$RFA0  = [ 0 , L_ ] , ! Block number part of RFA
RFATBL$W_RFA4 = [ 4 , W0_ ] , ! Byte offset part of RFA
TES;
```

```
LITERAL
```

```
RFATBL$K_SIZE = 6; ! Table entry size in bytes
```

```
MACRO
```

```
RFATBL$BLOCKVECTOR(N) =
BLOCKVECTOR [N, RFATBL$K_SIZE, BYTE] FIELD(RFATBL$FLD_DEF) %;
```


S T A T I C A D D R E S S T A B L E D E F I N I T I O N S

The Static Address Table (SAT) consists of a two-level structure. There is a Program Static Address Table which specifies what static addresses (code addresses and static data addresses) belong to what modules. This table consists of a linked list of SAT entries where each entry specifies an address range and a pointer to the corresponding Module RST Entry. The list is ordered by start address so that low start addresses come before higher start addresses. It should be noted that a given address may be covered by more than one SAT entry because different modules can have global PSECTs in common--Fortran COMMON blocks cause this situation to arise. Given an address, the Program Static Address Table can thus be searched to find the module that contains the address.

On a second level, there is a Module Static Address Table for each module that is SET. This SAT is a linked list of exactly the same structure as the Program SAT except that the RST pointer in the SAT entry points to the RST entry of the symbol, not the module, containing the address range. This list is also ordered by start address and again more than one symbol may contain the same address (due to Fortran EQUIVALENCEing for example). If two entries have the same start address, the entry with the larger end address is placed first in the Module SAT; this is crucial to the correct behavior of the DBG\$STA_SETCODE routine where routine addresses are compared to PC values from the VAX call stack. The Module RST Entry contains a pointer (RST\$SAT_PTR) which points to the first SAT entry on the module's SAT chain.

Note (R. Title)- for a future release of DEBUG, I intend to turn the SAT chains into binary trees, instead of linked lists. The reason for this is that for large programs, the SAT chains can grow quite long (e.g., 1000 entries), and it becomes quite expensive to search this list linearly.

The tree search algorithm will work something like this: If the desired address is greater than the address in the SAT\$END field of the current node, follow the "right son" pointer. If it is less than the address in the SAT\$START field of the current node, then follow the "left son" pointer. If it is in the range, then we have found the desired SAT entry and we are done.

During the transition to this new data structure, the same four fields are being kept in the SAT entry in the same place (so that the old code still works). The "right son" field is overlaid with the FLINK field, so that the tree search algorithm will work on the old linked lists (a linear list is just an extremely unbalanced tree, with no left branches). The "left son" field will occupy the fifth longword.

The individual Static Address Table entry has this format:

```

  3 3 2 2 2 2 2 2 2 2 2 1 1 1 1 1 1 1 1 1
  1 0 9 8 7 6 5 4 3 2 1 0 9 8 7 6 5 4 3 2 1 0

```

0	SAT\$L_FLINK (also SAT\$L_RIGHTSON)
1	SAT\$L_START
2	SAT\$L_END
3	SAT\$L_RSTPTR
4	SAT\$L_LEFTSON

A pointer to a Static Address Table entry is declared as follows:

SATPTR: REF SAT\$ENTRY

Field definitions for the SAT entry.

FIELD SAT\$FLD_DEF =

```

SET
SAT$L_FLINK      = [ 0, L_ ],      ! Forward link to next SAT entry on the
                                     chain--zero terminate the chain.
SAT$L_RIGHTSON   = [ 0, L_ ],      ! Link to right son in binary tree.
SAT$L_START      = [ 1, L_ ],      ! Start address of address range
SAT$L_END        = [ 2, L_ ],      ! End address of address range
SAT$L_RSTPTR     = [ 3, L_ ],      ! Pointer to RST entry of module or sym-
                                     bol containing this address range
SAT$L_LEFTSON    = [ 4, L_ ]      ! Link to left son in binary tree.
TES;
```

LITERAL

SAT\$K_ENTSIZE = 5; ! Size of one SAT entry in longwords

MACRO

SAT\$ENTRY = BLOCK[SAT\$K_ENTSIZE] FIELD(SAT\$FLD_DEF) %;

DBG\$W_VALUE_TOKENCODE=[4, W0],	:	Value of TOKEN\$W_CODE for constants
DBG\$W_VALUE_SIGN_CODE=[4, W1],	:	Sign code (Value of Unary -/+
	:	TOKEN\$W_CODE) for constants
DBG\$A_VALUE_VMSDESC = [5, A],	:	Address of VAX/VMS descriptor
DBG\$B_VALUE_CLASS = [5, B3],	:	DSC\$B_CLASS sub-field
DBG\$B_VALUE_DTYPE = [5, B2],	:	DSC\$B_DTYPE sub-field
DBG\$W_VALUE_LENGTH = [5, W0],	:	DSC\$W_LENGTH sub-field
DBG\$L_VALUE_POINTER = [6, L],	:	DSC\$A_POINTER sub-field
DBG\$B_VALUE_SCALE = [7, B0],	:	DSC\$B_SCALE sub-field
DBG\$B_VALUE_DIGITS = [7, B1],	:	DSC\$B_DIGITS sub-field
DBG\$V_VALUE_FL BINSIZE = [7, 19, 1, 0],	:	DSC\$V_FL BINSIZE sub-field
DBG\$L_VALUE_POS = [7, L],	:	DSC\$L_POS sub-field
DBG\$A_VALUE_ADDRESS = [8, A],	:	First byte of value
DBG\$L_VALUE_VALUE0 = [8, L],	:	First longword of value
DBG\$L_VALUE_VALUE1 = [9, L],	:	Second longword of value
TES;		

MACRO

```
DBG$VALDESC = BLOCK [,LONG] FIELD(DBG$DHDR_FIELDS, DBG$VALUE_FIELDS) %;
```

```
: Define the "base size" of a value descriptor in longwords. For the cases
: where the value is actually stored in the descriptor, the total size
: is the base size plus the number of longwords needed to hold the actual
: value.
```

LITERAL

```
DBG$K_VALDESC_BASE_SIZE = 8;
```

O L D D E B U G G E R D E F I N I T I O N S

These definitions are used by old Debugger modules which also need to use the new RST (and other) definitions.

LITERAL

```
! constants used to identify individual rst flags
rst$f_global = 1,
rst$f_nonzlength = 2,
rst$f_modset = 3;
```

MACRO

```
! Get the srm type address from a dst with address calculation commands
! The type is following the name string
! 7 - size of fixed portion of dst
! 1 - byte to hold count of name string
! dst[dstr_name] - contains count of characters in name string
!
! add this distance to base address to get address of byte
! containing the srm type
```

```
CAD_SRM_TYPEA(DST) =
( .DST + 7 + 1 + .DST[dst$b_name] ) %,
```

```
! Get the srm type from a dst with address calculation commands
```

```
CAD_SRM_TYPE(DST) =
( .(CAD_SRM_TYPEA(DST))<0,8,0> ) %,
```

```
! Get the address of the calculation commands in a dst which contains
! them. They start one byte past the srm type.
```

```
CAD_COMMANDS(DST) =
( CAD_SRM_TYPEA(DST) + 1 ) %;
```

MACRO

```
RST_UNITS( bytes ) =
( ((bytes) + %upval-1)/%upval )
%;
```

MACRO

```
! DEBUG tells the RST module about ASCII
! strings by passing a counted string pointer.
CS_POINTER = REF VECTOR[1,BYTE] %;
```

LITERAL

```
! We will never print "symbol+offset" when the
! upper bound for "symbol" is 0 and when
! the offset is greater than RST_MAX_OFFSET
```

RST\$K_MAX_OFFSET = %X'100';

!+
! Since scope definitions are recursive, we must
! stack ROUTINE BEGINS in the routine ADD_MODULE.
! It is no coincidence that this stack limit is the
! same as the limit on the length (in elements) of
! symbol pathnames.
!-

LITERAL
MAX_SCOPE_DEPTH = dbg\$K_max_pathname; ! Routines can be nested to a maximum depth.

FIELD

```
    VALU_FIELD_SET =  
    SET  
        VALU_NT_PTR    = [ 0,0,32,0 ],      ! Associated NT pointer.  
        VALU_VALUE     = [ 4,0,32,0 ]      ! The actual value.  
    TES;
```

```
!+  
! Declare an occurrence or REF to a VALUE_DESCRIPTOR  
! via the following macros.  
!-
```

LITERAL

```
    VALU_DESC_SIZE = 8;           ! Each one is 2 longwords long.
```

MACRO

```
    VALU_DESCRIPTOR = BLOCK[ VALU_DESC_SIZE, BYTE ] FIELD( VALU_FIELD_SET ) %;
```

!+
!+ Array Bounds Descriptor

An array bounds Descriptor is used to pass around all needed information about an array and its associated dimensions. Like VALU_DESCRIPTORs, they are simply 2-longword blocks, but this might change.

!-----longword-----!

address of array
length of array

Such Descriptors must be accessed via the following field names.

FIELD
SET ARRAY_BOUNDS_SET =
ARRAY_ADDRESS = [0,0,32,0], ! Beginning address of array.
ARRAY_LENGTH = [4,0,32,0] ! Size, in bytes, of array.
TES;

!+
!+ Declare an occurrence or REF to an array bounds
!+ descriptor via the following macros.

LITERAL ARRAY_BOUNDS_SIZE = 8; ! Each one is 2 longwords long.

MACRO ARRAY_BOUNDS_DESC = BLOCK[ARRAY_BOUNDS_SIZE, BYTE] FIELD(ARRAY_BOUNDS_SET) %;

LITERAL
SL_ACCE_INIT = 0. ! See above. "SL" --> SAT/LVT
SL_ACCE_RECS = 1.
SL_ACCE_SORT = 2.
SL_ACCE_FREE = 3;

↑
You declare an occurrence or REF of an SAT datum via
the macro, SAT_RECORD.
-

MACRO
SAT_RECORD = BLOCK[] FIELD(sat\$fld_def) %;

♦♦

Now we bring in the GST definitions directly from STARLET.REQ.

The format of one of these concatenated GSD records is a single leading byte containing the value 1, indicating that the record is indeed a GSD record.

Each entry in the record has a fixed number of overhead bytes followed by a symbol name that is a variable number of bytes. The entries we are interested in processing are the global symbol definitions and entry point symbol and mask definitions. The other defined type, PSECT definition, is noted only because it must be successfully passed over. The format of each of these types is illustrated below:

Global symbol definition:

0	! GSD type 1 !	
1	! data type !	ignored for now
2	! flag	bit 1 set means that this is a definition. ignore bit 0.
3	! bytes !	
4	! psect index !	ignored.
5	! value	4 bytes
9	! symbol name !	
		stock counted character string.

The entry point symbol and mask definition entry is identical to the global symbol definition illustrated above, with the addition of a two byte field for the procedure's register save mask. This two byte field is located after the symbol value field (which is an entry point address).

0	! GSD type 2 !	
1	! data type !	ignored for now
2	! flag	not relevant for entry point def.
3	! bytes !	
4	! psect index !	ignored
5	! value	4 bytes

9 10	register save mask	ignored, 2 bytes
11	symbol name	stock counted character string

The procedure definition with formal argument descriptions is identical to the "GSD type 2" definition above with the addition of fields that describe each formal argument. Following the symbol name there is one byte containing the minimum number of arguments allowed, and one byte containing the maximum number of arguments. These bytes are followed by a series of records of 2 - 257 bytes that describe each of the formal arguments (the number of these records is equal to the maximum number of arguments for the procedure.

0	GSD type 3	
1	data type	ignored for now
2 3	flag bytes	bit 1 set indicates that this is a definition. Bit 0 ignored.
4	psect index	ignored
5	value	4 bytes
9 10	register save mask	ignored, 2 bytes
11	symbol name	stock counted character string
	min # of args	1 byte
	max # of args	1 byte
	formal arg #1 description	
	:	
	formal arg #n description	

 format of each formal argument description.

0	! arg value ctl !	1 byte
1	! remaining byte!	1 byte (0 - 255)
	! detailed argument description !	

PSECT definition:

0	! GSD type 0 !	
1	! alignment !	
2	! flag	
3	! bytes !	
4	! allocation !	4 bytes
8	! symbol name !	stock counted character string.

The format of GST records is defined as a BLOCK
 (even though the records are variable sized)
 with the following fields:

--

FIELD

GST_FIELD_SET =

SET

GST_ENTRY_TYPE	= [0,0, 8,0],	! Type of GST record.
GST_IS_DEFN	= [2,1, 1,0],	! This flag implies whether or no the record is an entry definition.
GST_VALUE	= [5,0,32,0],	! The value of the global. (This won't work for PSECTs)
GST_P_NAME_CS	= [8,0, 8,0],	! Character count of psect name
! The following 2 pairs are mutually exclusive. The first one is for GST records of type GST_GLOBAL_DEFN, the second is for GST_ENTRY_DEFN or GST_PROC_DEFN.		
GST_G_NAME_CS	= [9,0, 8,0],	! The symbol name is a counted string.

```
GST_G_NAME_ADDR = [10,0, 8,0 ],
```

```
! A dotted reference to this field
! picks up the count, an undotted
! one addresses the counted string.
! The name string itself. An undotted
! reference is the address of the name,
! a dotted one is the 1st character.
```

```
GST_E_NAME_CS   = [11,0, 8,0 ],
```

```
GST_E_NAME_ADDR = [12,0, 8,0 ],
```

```
! The entry name is a counted string.
! A dotted reference to this field
! picks up the count, an undotted
! one addresses the counted string.
! The name string itself. An undotted
! reference is the address of the name,
! a dotted one is the 1st character.
! Maximum number of formal arguments.
! Remaining byte count of argument
! descriptor.
```

```
GST_P_MAX_ARG   = [ 1,0, 8,0 ],
GST_P_REM_CNT   = [ 1,0, 8,0 ]
```

```
TES;
```

```
+ You declare an occurrence or REF of a GST datum via:
```

```
LITERAL
```

```
GST_RECORD_SIZE = 43;    ! Each GST record is at most 43 bytes long.
```

```
MACRO
```

```
GST_RECORD = BLOCK[ GST_RECORD_SIZE, BYTE] FIELD( GST_FIELD_SET ) %;
```

```
LITERAL
```

```
GST_PSECT_OVERHEAD = 9,
GST_GLOBAL_OVERHEAD = 10,
GST_ENTRY_OVERHEAD = 12,
```

```
! Minimum number of bytes in psect def
! Minimum number of bytes in a global entry
! Minimum number of bytes in entry
! point symbol and mask definition
! Minimum size of formal argument descriptor
! Size of min. and max. overhead in GST
! Type of GST record
! Record Type is GST
```

```
GST_ARGDSC_OVERHEAD = 2,
GST_MINMAX_OVERHEAD = 2,
GST_RECORD_TYPE = 0,
GST_TYPE = 1;
```

```
+ The GST record types are defined as:
```

```
LITERAL
```

```
! GST types:
```

```
GST_LOWEST      = 1,
```

```
! We don't support global PSECTs now.
```

```
GST_PSECT_DEFN  = 0,
```

```
! P-SECT record.
```

```
GST_GLOBAL_DEFN = 1,
```

```
! A global symbol definition record.
```

```
GST_ENTRY_DEFN  = 2,
```

```
! An entry point definition.
```

```
GST_PROC_DEFN   = 3,
```

```
! A procedure with formal argument desc.
```

```
GST_HIGHEST     = 3;
```

```
! Highest one we support.
```

```

!+
BLISS uses 'non-standard' DST records to encode
most of its local symbol information. These records
are like most DST records except that the TYPE
information is variable-sized.
--

```

FIELD

```

BLZ_FIELD_SET =
SET
BLZ_SIZE      = [ 0,0, 8,0 ],      ! First byte is record size in bytes.
! The next byte contains DSC$K_DTYPE_2, or we
! wouldn't be applying this structure to a given
! DST record.

BLZ_TYP_SIZ   = [ 2,0, 8,0 ],      ! Type info takes up this
!                                     many bytes.
BLZ_TYPE      = [ 3,0, 8,0 ],      ! Which type of type Zero
!                                     this corresponds to.
BLZ_ACCESS    = [ 4,0, 8,0 ],      ! Access field.

! Sub fields of _ACCESS are offset from beginning
! of the BLZ record.
        BLZ_ACCES_TYPE = [ 4,0, 2,0 ],      ! Type of access,
        BLZ_ACCES_BASD = [ 4,2, 2,0 ],      ! based or not,
        BLZ_ACCES_BREG = [ 4,4, 4,0 ],      ! associated register.

BLZ_STRUCT    = [ 5,0,3,0 ],      ! Type of STRUCTURE reference.
BLZ_REF       = [ 5,7,1,0 ],      ! Indicates whether symbol has REF attribute

! **** The following only work when BLZ_TYP_SIZ is 3.

BLZ_VALUE     = [ 6,0,32,0 ],      ! DST VALUE field.
BLZ_NAME_CS   = [ 10,0, 8,0 ],     ! The symbol name is a counted string.
!                                     A dotted reference to this field
!                                     picks up the count, an undotted
!                                     one addresses the counted string.
BLZ_NAME_ADDR = [ 11,0, 8,0 ],     ! The name string itself. An undotted
!                                     reference is the address of the name,
!                                     a dotted one is the 1st character.

! The following fields are in the variable length part of the DST record.
! They should only be applied once the structure type has been determined
! or errors could result.

U_ALLOC_STRUC = [ 6, 0, 32, 0 ],    ! no of units alloc (except BLOCKVECTOR)
U_ALLOC_BVEC  = [ 10, 0, 32, 0 ],   ! no of units alloc for BLOCKVECTOR
SIGN_EXT_VEC  = [ 10, 7, 1, 0 ],    ! sign ext field for VECTOR
UNIT_SIZE_BLOCK = [ 10, 0, 8, 0 ],  ! unit size for BLOCK
UNIT_SIZE_BVEC  = [ 14, 0, 8, 0 ],  ! unit size for BLOCKVECTOR
UNIT_SIZE_VEC   = [ 10, 0, 4, 0 ],  ! unit size for VECTOR
TES;

```

```

!+

```


! You declare a REF to a BLZ_DST datum via:
!_

LITERAL
BLZ_REC_SIZ = 54; ! Each DST record is at most 54 bytes long.

MACRO
BLZ_RECORD = BLOCK[BLZ_REC_SIZ, BYTE] FIELD(BLZ_FIELD_SET) %;

!+
! The type zero sub types,
! as defined in CPO021.MEM,
! must be within the following
! range.
!_

LITERAL
! Type Zero Sub-Types:
BLZ_LOWEST = 1, ! Lowest variable type we support.
BLISS_Z_FORMAL = 1, ! Description of a ROUTINE formal.
BLISS_Z_SYMBOL = 2, ! A BLISS LOCAL symbol.
BLZ_HIGHEST = 2; ! Highest variable type we support.

*
Dummy descriptors

Whenever a type DSC\$K_DTYPE_CAD item is being accessed by descriptor we must build a dummy one in free storage and put the current address in it. The address of this dummy descriptor is then used by the rest of the debugger in referencing the item. To be able to free up all this space, the items are kept on a linked list. At the end of command processing this list is walked down and all the space is returned to the free storage manager.

Each entry on the linked list is a 3 longword vector

pointer to next item on the linked list

pointer to the dummy descriptor

size in bytes of the dummy descriptor

The global variable DBG\$GL_DLISHEAD if nonzero points to the first item on the list.

LITERAL

DLIS_ENTRY = 3 ;	! Size in longwords of entry on list
DLIS_LINK = 0 ;	! Pointer to next item on list
DLIS_POINTER = 1 ;	! Pointer to dummy descriptor
DLIS_DESCSIZE = 2 ;	! Size in bytes of descriptor

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! E N D O F D B G L I B . R E Q

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