

CCCCCCCCCCCC	DDDDDDDDDDDD	DDDDDDDDDDDD		
CCCCCCCCCCCC	DDDDDDDDDDDD	DDDDDDDDDDDD		
CCCCCCCCCCCC	DDDDDDDDDDDD	DDDDDDDDDDDD		
CCC	DDD	DDD	DDD	DDD
CCC	DDD	DDD	DDD	DDD
CCC	DDD	DDD	DDD	DDD
CCC	DDD	DDD	DDD	DDD
CCC	DDD	DDD	DDD	DDD
CCC	DDD	DDD	DDD	DDD
CCC	DDD	DDD	DDD	DDD
CCC	DDD	DDD	DDD	DDD
CCC	DDD	DDD	DDD	DDD
CCC	DDD	DDD	DDD	DDD
CCC	DDD	DDD	DDD	DDD
CCC	DDD	DDD	DDD	DDD
CCC	DDD	DDD	DDD	DDD
CCC	DDD	DDD	DDD	DDD
CCC	DDD	DDD	DDD	DDD
CCC	DDD	DDD	DDD	DDD
CCC	DDD	DDD	DDD	DDD
CCC	DDD	DDD	DDD	DDD
CCCCCCCCCCCC	DDDDDDDDDDDD	DDDDDDDDDDDD		
CCCCCCCCCCCC	DDDDDDDDDDDD	DDDDDDDDDDDD		
CCCCCCCCCCCC	DDDDDDDDDDDD	DDDDDDDDDDDD		

```

CCCCCCCC DDDDDDDD DDDDDDDD LL      IIIIII  BBBB8888
CCCCCCCC DDDDDDDD DDDDDDDD LL      IIIIII  BBBB8888
CC        DD        DD        LL      II       BB      BB
CC        DD        DD        LL      II       BB      BB
CC        DD        DD        LL      II       BB      BB
CC        DD        DD        LL      II       BB      BB
CC        DD        DD        LL      II       BBBB8888
CC        DD        DD        LL      II       BBBB8888
CC        DD        DD        LL      II       BB      BB
CC        DD        DD        LL      II       BB      BB
CC        DD        DD        LL      II       BB      BB
CC        DD        DD        LL      II       BB      BB
CCCCCCCC DDDDDDDD DDDDDDDD LLLLLLLLLL IIIIII  BBBB8888
CCCCCCCC DDDDDDDD DDDDDDDD LLLLLLLLLL IIIIII  BBBB8888

```

```

88888888 333333 222222
88888888 333333 222222
BB      BB 33      33 22      22
BB      BB 33      33 22      22
BB      BB 33      33 22      22
BB      BB 33      33 22      22
88888888 33      22
88888888 33      22
BB      BB 33      33 22      22
BB      BB 33      33 22      22
BB      BB 33      33 22      22
88888888 33      22
88888888 33      22
88888888 333333 2222222222
88888888 333333 2222222222

```

```

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*****

```

```

**
TITLE: CDDLIB                CDD Bliss Library
VERSION: 'V04-000'
FACILITY: Common Data Dictionary'
ABSTRACT:
    This module is the library file used for compiling all other
    modules in the CDD facility.
ENVIRONMENT:
AUTHOR: Jeff East, 22-Jan-80
MODIFIED BY:
    7-May-81 (JAE) Added $IO_SYNC macro.

```

```

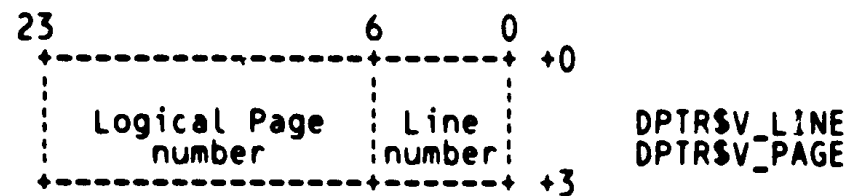
--
XTITLE 'CDD Bliss Library'

```

```
%SBTTL 'STRUCTURE DEFINITIONS'
```

```
STRUCTURE DEFINITIONS
```

```
On-disk Pointer Structure
```

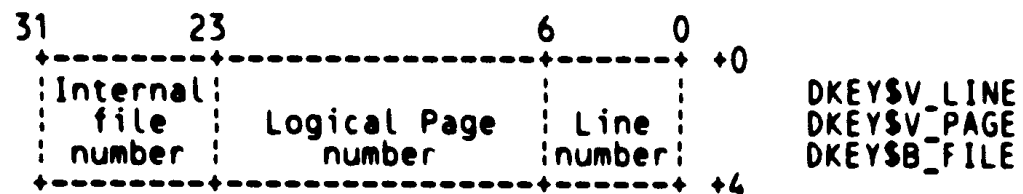


```
LITERAL
DPTRSS_BLOCK_LENGTH = 3;
```

```
MACRO
$DPTR = BLOCK[DPTRSS_BLOCK_LENGTH, BYTE] FIELD (DPTRSZ_FIELDS)
%;
```

```
FIELD DPTRSZ_FIELDS =
SET
DPTRSV_VALUE      = [0, 0, 24, 0],
DPTRSV_LINE       = [0, 0, 7, 0],
DPTRSV_PAGE       = [0, 7, 17, 0]
TES;
```

In-core Disk Pointer Structure



When a disk pointer is being passed around routines, it is passed as a Disk Key (DKEY), rather than just a disk pointer.

Disk Keys include all the information of a disk pointer, but also include a pointer to their file's FCB.

LITERAL

DKEY\$S_BLOCK_LENGTH = 4;

MACRO

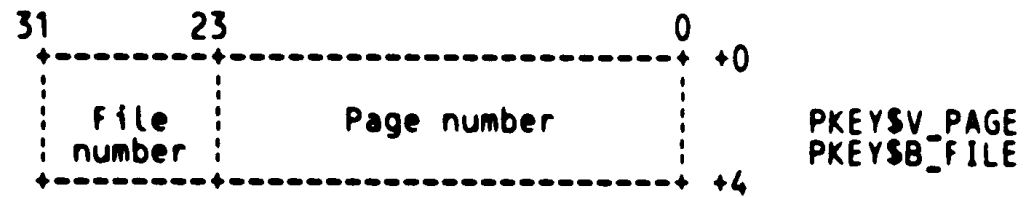
\$DKEY = BLOCK[DKEY\$S_BLOCK_LENGTH, BYTE] FIELD (DKEY\$Z_FIELDS)
%;

FIELD DKEY\$Z_FIELDS =

SET
 DKEY\$V_LINE = [0, 0, 7, 0],
 DKEY\$V_PAGE = [0, 7, 17, 0],
 DKEY\$B_FILE = [0, 24, 8, 0]

TES;

In-core Page Number Structure



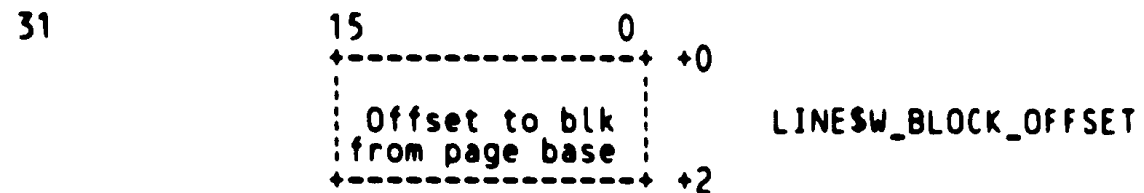
Dictionary page numbers that are passed around the hash table use the Page Key (PKEY) structure to provide both the page and file numbers.

```
LITERAL
PKEY$$_BLOCK_LENGTH = 4;
```

```
MACRO
$PKEY = BLOCK[PKEY$$_BLOCK_LENGTH, BYTE] FIELD (PKEY$Z_FIELDS)
%;
```

```
FIELD PKEY$Z_FIELDS =
SET
PKEY$V_PAGE           = [0, 0, 24, 0],
PKEY$B_FILE           = [0, 24, 8, 0]
TES;
```

Line Index Format



Each block on a dictionary page is pointed to by a line index on that page. Disk pointers (DPTR) between the various blocks actually point to the block's line index. The line index is then used to locate the physical dictionary block.

The PAGESINDEX macro allows the programmer to access a line index while it resides on a dictionary page. It also uses the LINESZ_FIELDS to access the individual fields within a line index.

LITERAL

```
LINESW_BLOCK_LENGTH = 2;
```

MACRO

```
$LINE = BLOCK[LINESW_BLOCK_LENGTH, BYTE] FIELD (LINESZ_FIELDS)
```

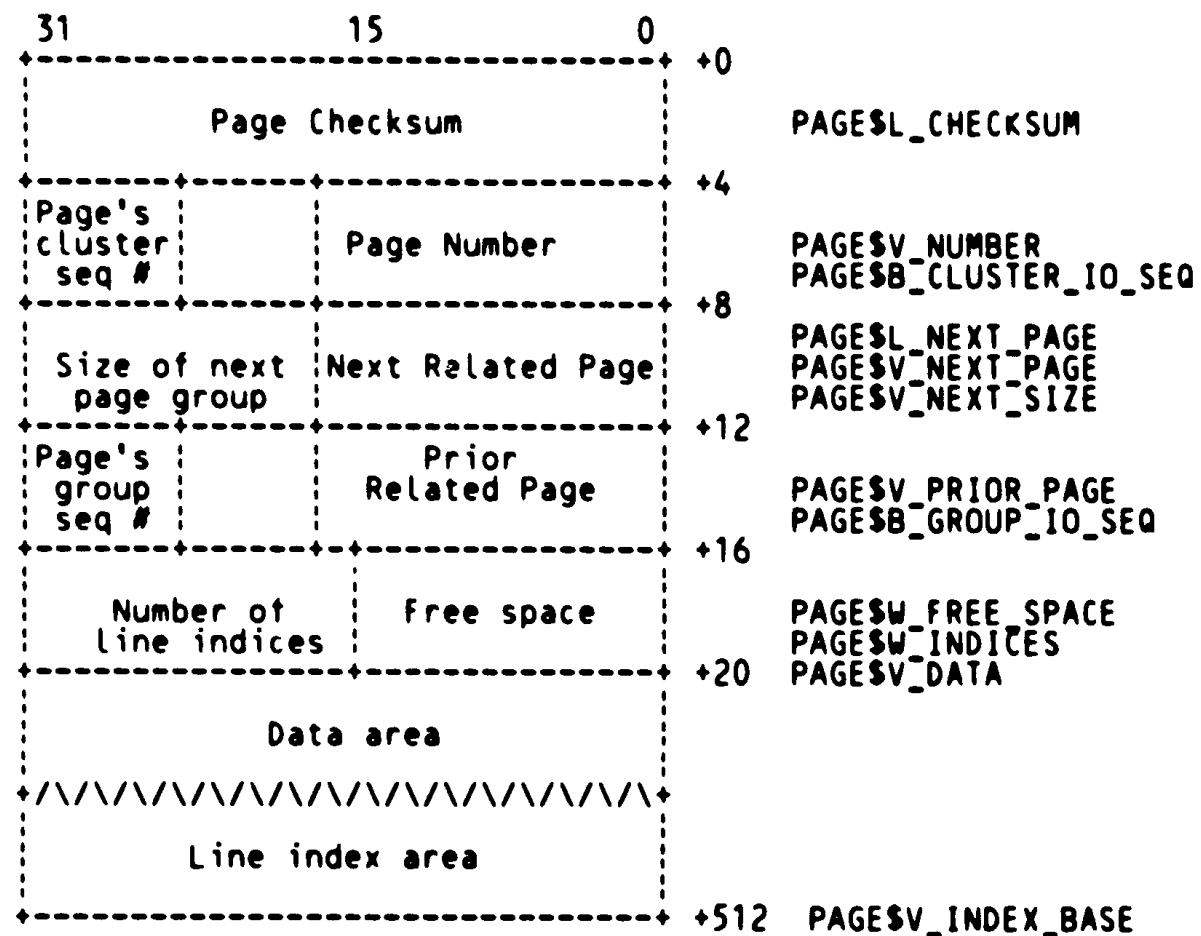
```
%;
```

FIELD LINESZ_FIELDS =

```
SET LINESW_BLOCK_OFFSET = [0, 0, 16, 0]
```

```
TES;
```

Page Format



Each page in the dictionary file starts with a page header. The next and prior related page pointers are used for linking page groups. The following types of page groups exist:

- 1) Lockable page group.
Each named entity (and history list) owns exactly one lockable page group. These groups contain all unnamed offspring of the name entry (or history list).

A page group may only be accessed through its portal page. A group's portal page is the page on which the named entity (or history list head) resides. If a group's portal page is locked, then no pages in the group may be accessed.

The PAGESINDEX structure is used to access a line index entry on a page.

A page whose checksum is zero is a locked page, and indicates that

the sub-tree below it is incomplete and in a transient state. Such pages may never be read.

Each page has two page sequence numbers. The cluster sequence number must be the same for all pages in the cluster. The group sequence number must be the same for all pages in an I/O group.

The cluster I/O sequence number is bumped whenever more than one group in the cluster is written. The group I/O sequence number is bumped whenever the group is written. These sequence numbers enable the detection of incomplete cluster unstage operations.

```

LITERAL
  PAGESS_BLOCK_LENGTH = 512;

MACRO
  $PAGE = BLOCK[PAGESS_BLOCK_LENGTH, BYTE] FIELD (PAGE$Z_FIELDS)
  %:

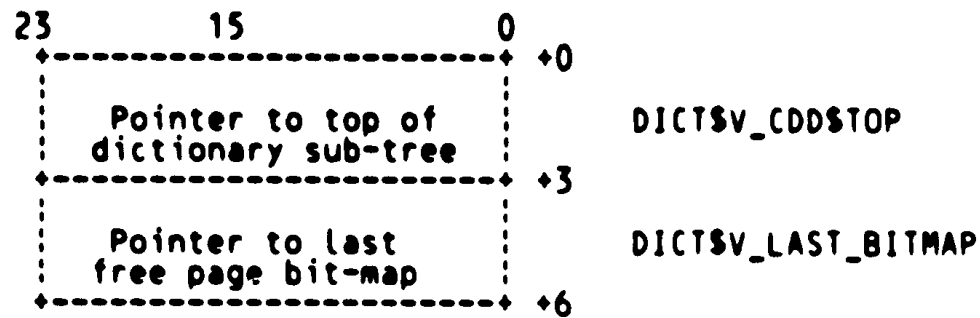
FIELD PAGE$Z_FIELDS =
  SET
    PAGE$L_CHECKSUM          = [0, 0, 32, 0],
    PAGE$V_NUMBER           = [4, 0, 17, 0],
    PAGE$B_CLUSTER_IO_SEQ   = [7, 0, 8, 0],
    PAGE$L_NEXT_PAGE        = [8, 0, 32, 0],
    PAGE$V_NEXT_PAGE        = [8, 0, 17, 0],
    PAGE$V_NEXT_SIZE        = [8, 17, 15, 0],
    PAGE$V_PRIOR_PAGE       = [12, 0, 17, 0],
    PAGE$B_GROUP_IO_SEQ     = [15, 0, 8, 0],
    PAGE$W_FREE_SPACE       = [16, 0, 16, 0],
    PAGE$W_INDICES          = [18, 0, 16, 0],
    PAGE$V_DATA              = [20, 0, 0, 0],
    PAGE$V_INDEX_BASE       = [512, 0, 0, 0]
  TES;

LITERAL
  PAGE$K_BASE = BLOCK[0, PAGE$V_DATA,
    PAGESS_BLOCK_LENGTH, BYTE],
  PAGE$K_FREE_SPACE = PAGESS_BLOCK_LENGTH - PAGE$K_BASE,
  PAGE$S_INDEX_BASE = BLOCK[0, PAGE$V_INDEX_BASE,
    PAGESS_BLOCK_LENGTH, BYTE];

STRUCTURE
  PAGE$INDEX[I, 0, P, S, E] =
    (PAGE$INDEX+PAGE$S_INDEX_BASE+0-((I)*LINESS_BLOCK_LENGTH)) <P,S,E>;

```

Dictionary Header Block



Each dictionary uses page 1 as its dictionary header page. The dictionary header block immediately follows the page header on page 1.

`DICTSV_CDDSTOP` points to the root of the tree/sub-tree contained in the dictionary file.

`DICTSV_LAST_BITMAP` points to the last free page bit-map entry.

LITERAL

`DICTSS_BLOCK_LENGTH = 6 + PAGESK_BASE;`

MACRO

`%DICT = BLOCK[DICTSS_BLOCK_LENGTH, BYTE] FIELD (PAGESZ_FIELDS, DICTSZ_FIELDS)`

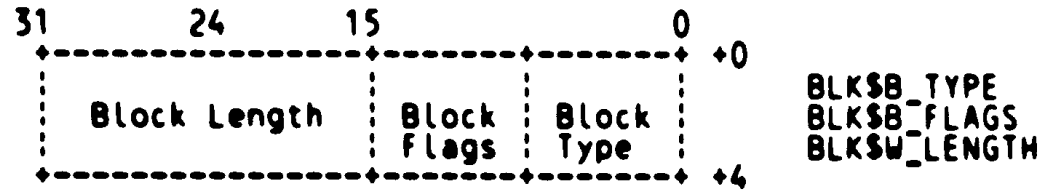
`%;`

FIELD `DICTSZ_FIELDS =`

`SET`
`DICTSV_CDDSTOP = [0+PAGESK_BASE, 0, 24, 0],`
`DICTSV_LAST_BITMAP = [3+PAGESK_BASE, 0, 24, 0]`

`TES;`

Universal Block Header



This block header is used in every non-fixed block on the dictionary pages. Fixed blocks (one which always occur in the same place whenever they exist) do not have this block header.

```
LITERAL
  BLK$S_BLOCK_LENGTH = 4;
```

```
MACRO
  SBLK = BLOCK[BLK$S_BLOCK_LENGTH, BYTE] FIELD (ACL$Z_FIELDS,
  ACLC$Z_FIELDS,
  ATT$Z_FIELDS,
  BLK$Z_FIELDS,
  ELST$Z_FIELDS,
  ENT$Z_FIELDS,
  FPCB$Z_FIELDS,
  LIST$Z_FIELDS,
  LST$Z_FIELDS,
  NAME$Z_FIELDS,
  NAT$Z_FIELDS,
  NNAM$Z_FIELDS,
  NODE$Z_FIELDS,
  SEGS$Z_FIELDS,
  SLST$Z_FIELDS,
  SSAS$Z_FIELDS,
  STR$Z_FIELDS,
  TEXT$Z_FIELDS)

  X;
```

```
FIELD BLK$Z_FIELDS =
  SET
    BLK$B_TYPE           = [0, 0, 8, 0],
    BLK$B_FLAGS         = [1, 0, 8, 0],
    BLK$W_LENGTH        = [2, 0, 16, 0]
  TES;
```

```
LITERAL
  BLK$K_FLAGS           = BLOCK[0, BLK$B_TYPE; , BYTE],
  BLK$K_BASE            = BLK$S_BLOCK_LENGTH;
```

!*

Block types

The ordering and clustering of these block types is significant.

All block types must be contiguous without any holes and bounded by the symbols BLKSK_TYPE_FIRST and BLKSK_TYPE_LAST.

All node and NAME block types must be contiguous and bounded by the symbols BLKSK_TYPE_NODE_FIRST and BLKSK_TYPE_NODE_LAST.

To add more node or NAME block types, insert them at the front of the table.

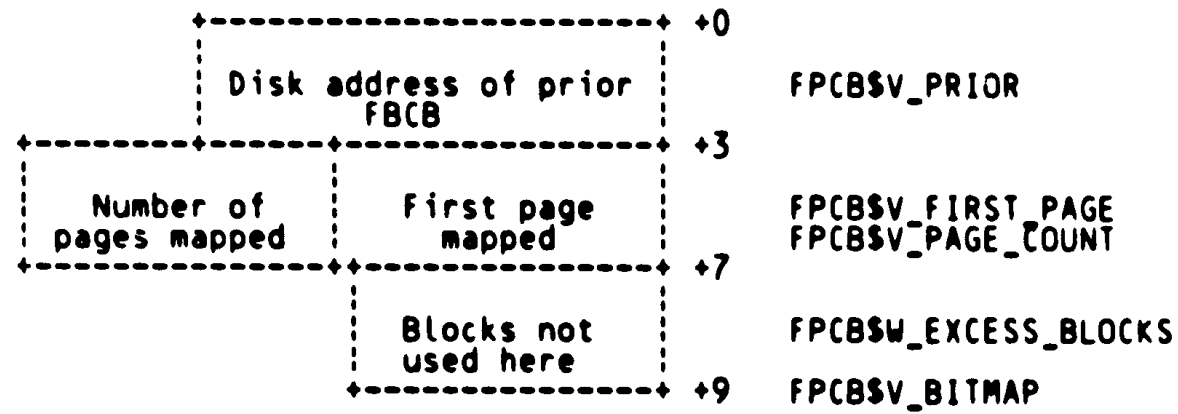
To add any other block type, append them to the end of the table.

CHANGING THE VALUES OF ANY OF THE ACTUAL BLOCK TYPES WILL INVALIDATE EXISTING DICTIONARIES.

LITERAL

BLKSK_TYPE_FIRST	= 101,	! Lowest block type
BLKSK_TYPE_NODE_FIRST	= 101,	! First node or NAM block
BLKSK_TYPE_DIR_NAM	= 101,	! Directory NAM block
BLKSK_TYPE_FIL_NAM	= 102,	! File NAM block
BLKSK_TYPE_TERM_NAM	= 103,	! Terminal NAM block
BLKSK_TYPE_DIR_NODE	= 104,	! Directory node block
BLKSK_TYPE_TERM_NODE	= 105,	! Terminal node block
BLKSK_TYPE_NODE_LAST	= 105,	! Last node or NAM block
BLKSK_TYPE_BITMAP	= 106,	! Free page bitmap
BLKSK_TYPE_ENTITY_ATT	= 107,	! Entity attribute block
BLKSK_TYPE_ENTITY_LIST	= 108,	! Entity list block
BLKSK_TYPE_ENTITY_LIST_ATT	= 109,	! Entity list attribute block
BLKSK_TYPE_NULL_ATT	= 110,	! Null attribute block
BLKSK_TYPE_NUM_ATT	= 111,	! Numeric attribute block
BLKSK_TYPE_SHORT_ATT	= 112,	! Short string attribute block
BLKSK_TYPE_STRING_ATT	= 113,	! String attribute block
BLKSK_TYPE_STRING_LIST	= 115,	! String list block
BLKSK_TYPE_STRING_LIST_ATT	= 116,	! String list attribute block
BLKSK_TYPE_STRING_SEG	= 117,	! String segment block
BLKSK_TYPE_TEXT	= 118,	! Text block
BLKSK_TYPE_ACL	= 119,	! Access control list entry
BLKSK_TYPE_ACLC	= 120,	! Access control list criterion
BLKSK_TYPE_LAST	= 120,	! Highest block type

Free Page Control Block (FPCB)
Free Page Bit Map (FPBM)



Free pages in the dictionary file are controlled by the Free Page Control Blocks (FPCB) and the Free Page Bit Map (FPBM). Each free page in the file has its corresponding bit set on.

The disk's cluster size may make it impossible to use all of the allocated blocks. If so, the un-used blocks are tallied in FPCBSW_EXCESS_BLOCKS, and will be used when the file is next extended.

The FPCBs are chained as the dictionary file grows in size.

LITERAL

FPCBS_BLOCK_LENGTH = BLK\$K_BASE + 9;

STRUCTURE

FPCB[0, P, S, E; BLOCKS] = [FPCBS_BLOCK_LENGTH+(BLOCKS+7)/8]
(FPCB+0)<P,S,E>;

FIELD FPCBSZ_FIELDS =

SET
 FPCBSV_PRIOR = [0+BLK\$K_BASE, 0, 24, 0],
 FPCBSV_FIRST_PAGE = [3+BLK\$K_BASE, 0, 17, 0],
 FPCBSV_PAGE_COUNT = [3+BLK\$K_BASE, 17, 15, 0],
 FPCBSW_EXCESS_BLOCKS = [7+BLK\$K_BASE, 0, 16, 0],
 FPCBSV_BITMAP = [9+BLK\$K_BASE, 0, 0, 0]

TES;

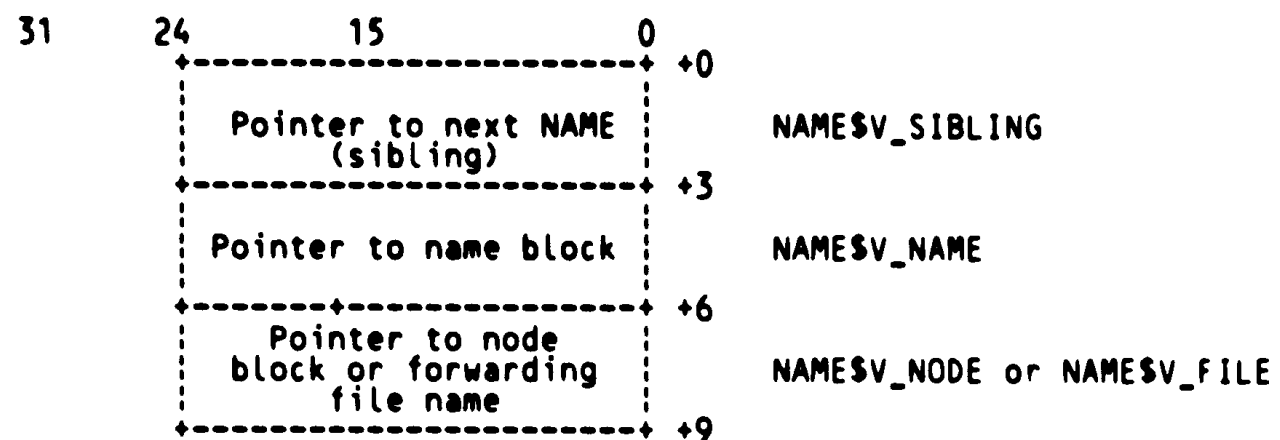
LITERAL

FPCBK_BASE = FPCBS_BLOCK_LENGTH,
 FPCBS_BITMAP = FPCB[0, FPCBSV_BITMAP];

STRUCTURE

FPBM[1] = (FPBM+FPCBS_BITMAP)<1-1, 1, 0>;

Name Block



There are three types of NAM Blocks, Directory, File, and Terminal.
 Directory NAM Block is made up of a Block Header + NAM Block.
 File NAM Block is made up of a Block Header + NAM Block.
 Terminal NAM Block is made up of a Block Header + NAM Block + Terminal NAM Block.

LITERAL

NAMES_BLOCK_LENGTH = 9 + BLK_BASE;

MACRO

\$NAME = BLOCK[NAMES_BLOCK_LENGTH, BYTE] FIELD (BLK_FIELDS, NAMES_FIELDS)

%;

FIELD NAMES_FIELDS =

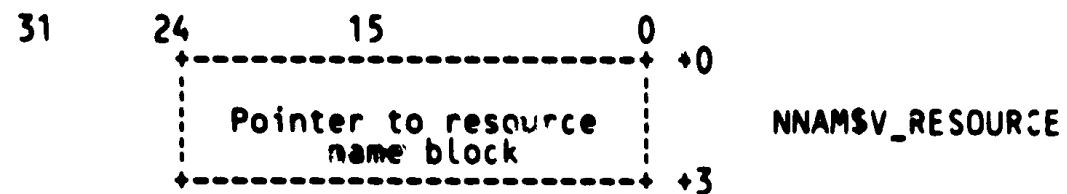
SET
 NAMESV_SIBLING = [0+BLK_BASE, 0, 24, 0],
 NAMESV_NAME = [3+BLK_BASE, 0, 24, 0],
 NAMESV_NODE = [6+BLK_BASE, 0, 24, 0],
 NAMESV_FILE = [6+BLK_BASE, 0, 24, 0]

TES.

LITERAL

NAMESK_BASE = NAMES_BLOCK_LENGTH;

Node Nam Block



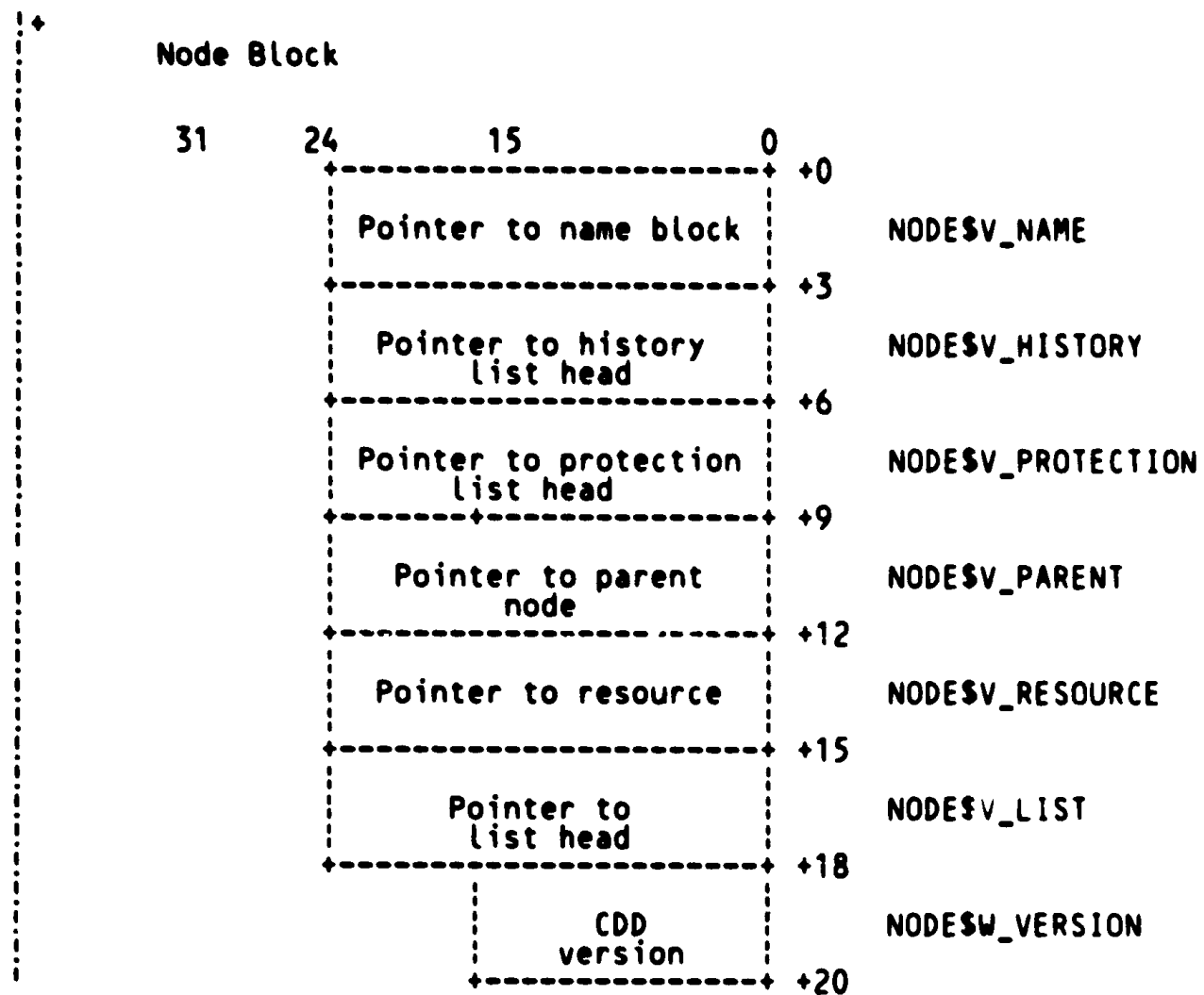
Node NAM Block is made up of a Block Header + NAM Block + Node NAM Block.

```
LITERAL
NNAMSS_BLOCK_LENGTH = 3 + NAME$K_BASE;
```

```
MACRO
$NNAM = BLOCK[NNAMSS_BLOCK_LENGTH, BYTE] FIELD (BLK$Z_FIELDS,
NAME$Z_FIELDS,
NNAM$Z_FIELDS)
%;
```

```
FIELD NNAM$Z_FIELDS =
SET
NNAMSV_RESOURCE = [0+NAME$K_BASE, 0, 24, 0]
TES;
```

```
LITERAL
NNAM$K_BASE = NNAMSS_BLOCK_LENGTH;
```



There are two types of Node Blocks Directory and Terminal.
 Directory Node Block is made up of a Block Header + Node Block
 Terminal Node Block is made up of a Block Header + Node Block

```
LITERAL
NODE$S_BLOCK_LENGTH = 20+ BLK$K_BASE;
```

```
MACRO
$NODE = BLOCK[NODE$S_BLOCK_LENGTH, BYTE] FIELD (BLK$Z_FIELDS,
NODE$Z_FIELDS)
X;
```

```
FIELD NODE$Z_FIELDS =
SET
NODE$V_ORDERED = [1, 0, 1, 0],
NODE$V_NAME = [0+BLK$K_BASE, 0, 24, 0],
NODE$V_HISTORY = [3+BLK$K_BASE, 0, 24, 0].
```



```
NODE$V_PROTECTION    = [6+BLK$K_BASE, 0, 24, 0],
NODE$V_PARENT        = [9+BLK$K_BASE, 0, 24, 0],
NODE$V_RESOURCE      = [12+BLK$R_BASE, 0, 24, 0],
NODE$V_LIST          = [15+BLK$K_BASE, 0, 24, 0],
NODE$W_VERSION       = [18+BLK$K_BASE, 0, 16, 0]
```

TES;

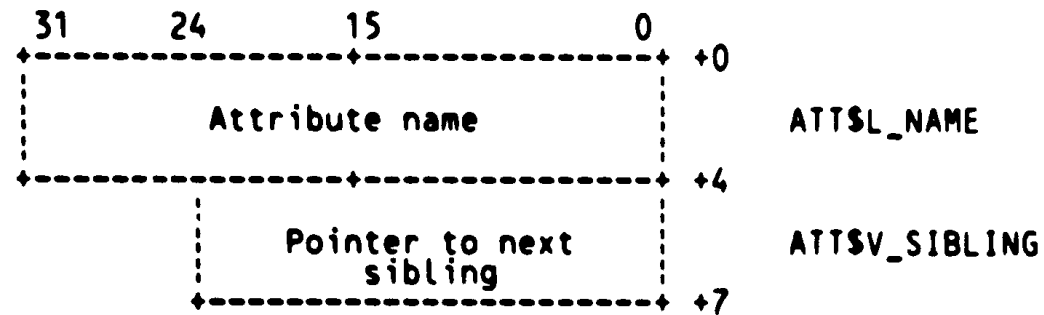
```
LITERAL
NODE$K_BASE          = NODE$S_BLOCK_LENGTH;
```

!+
!-
!-

The following flag occurs in the BLK\$B_FLAGS field of
a directory node.

```
LITERAL
NODE$M_ORDERED      = 1^1 - 1^0;    ! List elements are sorted by name
```

Common Attribute Block



Each attribute has the common attribute block immediately following the universal block header.

There are six types of Attribute Blocks Entity Attribute Block, List Attribute Block, Numeric Attribute Block, Null Attribute Block, Short String Attribute Block, and String Attribute Block.

Entity Attribute Block is made up of a Block Header + Attribute Block + Entity Attribute Block.

List Attribute Block is made up of a Block Header + Attribute Block + List Attribute Block.

Numeric Attribute Block is made up of a Block Header + Attribute Block + Numeric Attribute Block.

Null Attribute Block is made up of a Block Header + Attribute Block.

Short String Attribute Block is made up of a Block Header + Attribute Block + Short String Attribute.

String Attribute Block is made up of a Block Header + Attribute Block + String Attribute Block.

```
LITERAL
ATTSS_BLOCK_LENGTH = 7 + BLK$K_BASE;
```

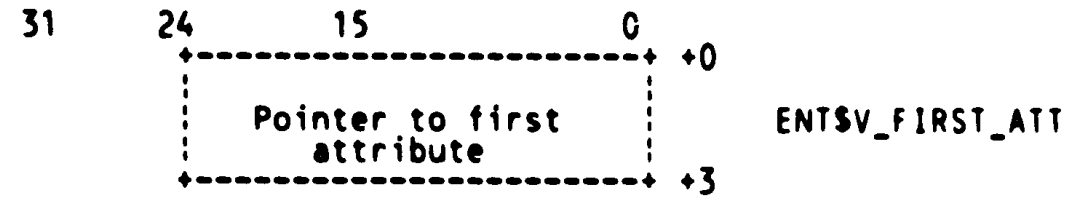
```
MACRO
$ATT = BLOCK[ATTSS_BLOCK_LENGTH, BYTE] FIELD (ATT$Z_FIELDS,
BLK$Z_FIELDS)
%;
```

```
FIELD ATT$Z_FIELDS =
SET
ATTSL_NAME           = [0+BLK$K_BASE, 0, 32, 0],
ATTSV_SIBLING        = [4+BLK$K_BASE, 0, 24, 0]
TES;
```

```
LITERAL
```

ATT\$K_BASE = ATT\$S_BLOCK_LENGTH;

Entity Attribute Block



Entity Attribute Block is made up of a Block Header + Attribute Block + Entity Attribute Block.

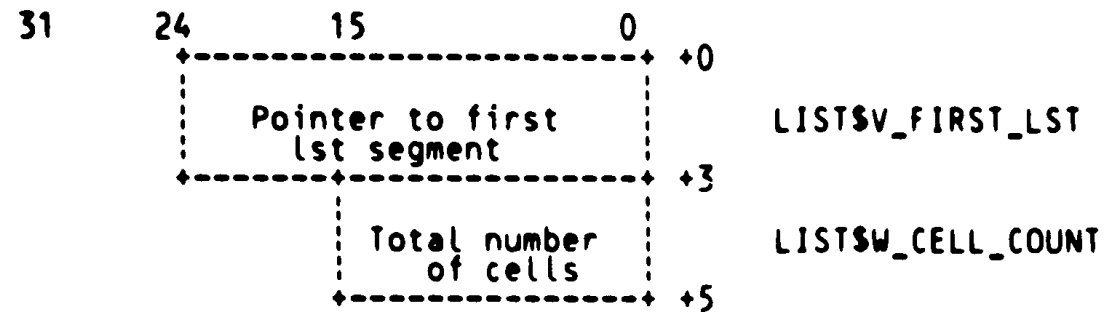
```
LITERAL
ENTSS_BLOCK_LENGTH = 3 + ATTSK_BASE;
```

```
MACRO
SENT = BLOCK[ENTSS_BLOCK_LENGTH,BYTE] FIELD (ATTSZ_FIELDS,
BLKSZ_FIELDS,
ENTSZ_FIELDS)
%;
```

```
FIELD ENTSZ_FIELDS =
SET
ENTSV_FIRST_ATT = [0+ATTSK_BASE, 0, 24, 0]
TES;
```

```
LITERAL
ENTSK_BASE = ENTSS_BLOCK_LENGTH;
```

List Attribute Block



List Attribute Block is made up of a Block Header + Attribute Block + List Attribute Block.

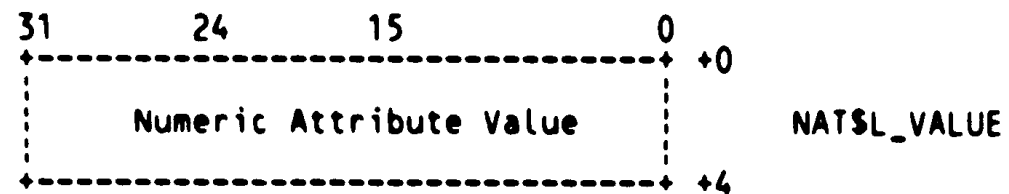
```
LITERAL
LIST$$BLOCK_LENGTH = 5 + ATT$K_BASE;
```

```
MACRO
$LIST = BLOCK[LIST$$BLOCK_LENGTH, BYTE] FIELD (ATT$Z_FIELDS,
                                                BLK$Z_FIELDS,
                                                LIST$Z_FIELDS)
%;
```

```
FIELD LIST$Z_FIELDS =
SET
LIST$V_FIRST_LST      = [0+ATT$K_BASE, 0, 24, 0],
LIST$W_CELL_COUNT    = [3+ATT$K_BASE, 0, 16, 0]
TES;
```

```
LITERAL
LIST$K_BASE          = LIST$$BLOCK_LENGTH;
```

Numeric Attribute Block



Numeric Attribute Block is made up of a Block Header + Attribute Block + Numeric Attribute Block.

Numeric attribute block contains the value of the numeric attribute.

```
LITERAL
NAT$$BLOCK_LENGTH = 4 + ATT$K_BASE;
```

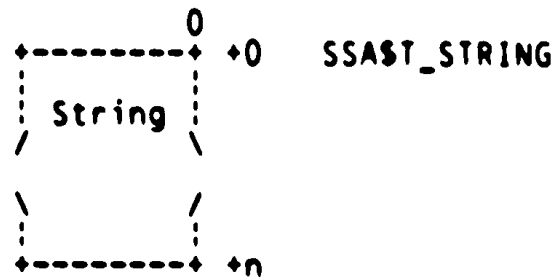
```
MACRO
$NAT = BLOCK[NAT$$BLOCK_LENGTH, BYTE] FIELD (ATT$Z_FIELDS,
                                             BLK$Z_FIELDS,
                                             NAT$Z_FIELDS)
%;
```

```
FIELD NAT$Z_FIELDS =
SET
NAT$L_VALUE          = [0+ATT$K_BASE, 0, 32, 0]
TES;
```

```
LITERAL
NAT$K_BASE          = NAT$$BLOCK_LENGTH;
```

Short String Attribute Block

31



Short String Attribute Block is made up of a Block Header + Attribute Block + Short String Attribute.

Strings whose total length is between 0 and 255 bytes are usually stored using the short string attribute block. If a string is too long to be stored in an SSA, then it is stored using a normal string attribute block (STR).

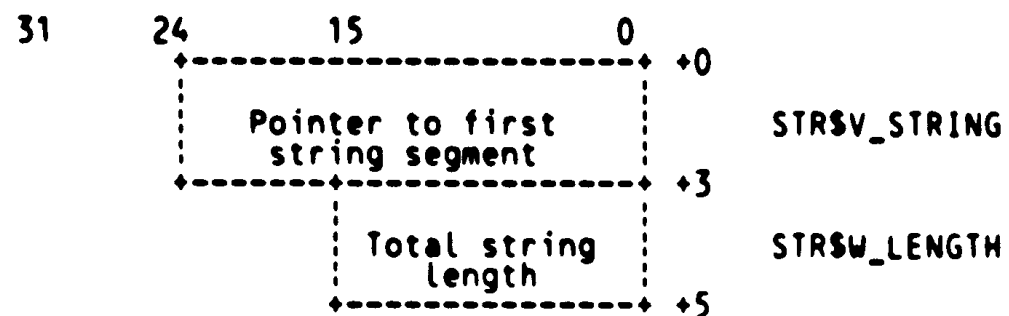
```
LITERAL
SSASS_BLOCK_LENGTH = 0 + ATT$K_BASE;
```

```
MACRO
$SSA = BLOCK[SSASS_BLOCK_LENGTH, BYTE] FIELD (ATT$Z_FIELDS,
                                               BLK$Z_FIELDS,
                                               SSASZ_FIELDS)
%;
```

```
FIELD SSASZ_FIELDS =
SET
SSAST_STRING          = [0+ATT$K_BASE, 0, 0, 0]
TES;
```

```
LITERAL
SSASK_BASE           = SSASS_BLOCK_LENGTH;
```

String Attribute Block



String Attribute Block is made up of a Block Header + Attribute Block + String Attribute Block.

The string is broken into one or more segments. The string segments are stored in SEG blocks.

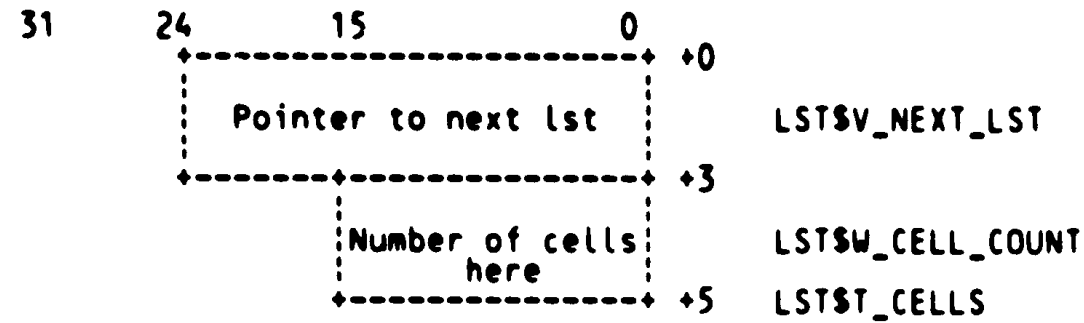
```
LITERAL
STRSS_BLOCK_LENGTH = 5 + ATT$K_BASE;
```

```
MACRO
$STR = BLOCK[STRSS_BLOCK_LENGTH, BYTE] FIELD (ATT$Z_FIELDS,
                                               BLK$Z_FIELDS,
                                               STR$Z_FIELDS)
%;
```

```
FIELD STR$Z_FIELDS =
SET
STRSV_STRING          = [0+ATT$K_BASE, 0, 24, 0]
STRSW_LENGTH          = [3+ATT$K_BASE, 0, 16, 0]
TES;
```

```
LITERAL
STR$K_BASE           = STRSS_BLOCK_LENGTH;
```


LST Segment Block



There are two types of LST Segment Blocks Entity List Block and String List Attribute Block.

Entity List Block is made up of a Block Header + LST Block + Entity List Block.

String List Attribute Block is made up of a Block Header + LST Block + String List Attribute Block.

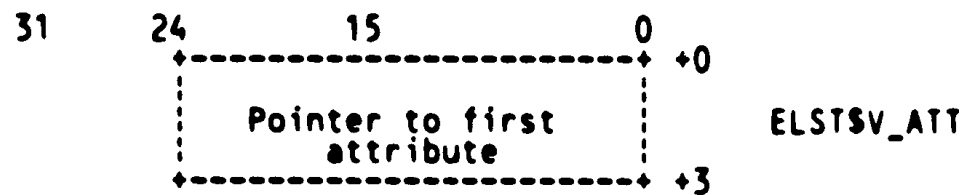
```
LITERAL
LSTSS_BLOCK_LENGTH = 5 + BLK$K_BASE;
```

```
MACRO
$LST = BLOCK[LSTSS_BLOCK_LENGTH, BYTE] FIELD (BLK$Z_FIELDS,
LST$Z_FIELDS)
%;
```

```
FIELD LST$Z_FIELDS =
SET
LSTSV_NEXT_LST      = [0 + BLK$K_BASE, 0, 24, 0],
LSTSW_CELL_COUNT   = [3 + BLK$K_BASE, 0, 16, 0],
LSTST_CELLS        = [5 + BLK$K_BASE, 0, 0, 0]
TES;
```

```
LITERAL
LST$K_BASE          = LSTSS_BLOCK_LENGTH;
```

Entity List Block



Entity List Block is made up of a Block Header + LST Block + Entity List Block.

```
LITERAL
ELSTSS_BLOCK_LENGTH = 3;
```

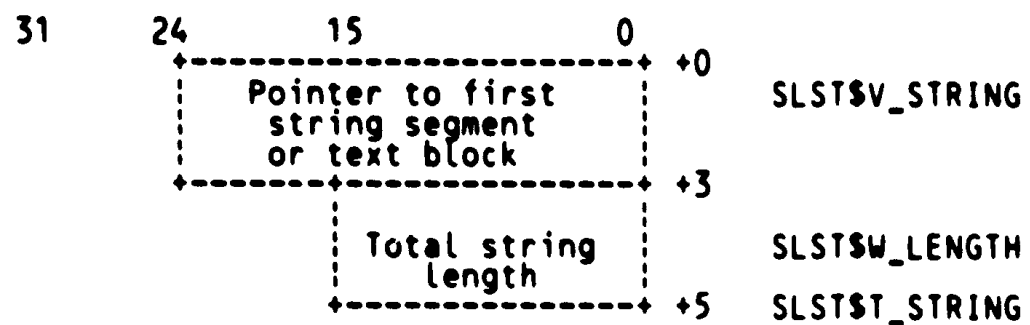
```
MACRO
SELST = BLOCK[ELSTSS_BLOCK_LENGTH, BYTE] FIELD (ELST$Z_FIELDS)
%;
```

```
FIELD ELST$Z_FIELDS =
SET
ELSTSV_ATT = [0, 0, 24, 0]
TES;
```

```
LITERAL
ELST$K_BASE = ELSTSS_BLOCK_LENGTH;
```

```
STRUCTURE
ELST$ELM[I] =
(ELST$ELM + LST$K_BASE + (I * ELSTSS_BLOCK_LENGTH));
```

String List Attribute Block



String List Attribute Block is made up of a Block Header + LST Block + String List Attribute Block.

```
LITERAL
SLST$$BLOCK_LENGTH = 5;
```

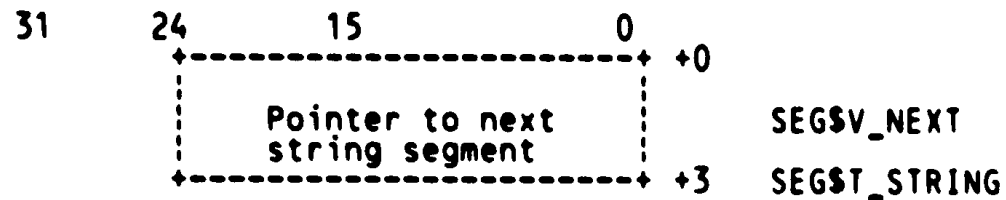
```
MACRO
SSLST = BLOCK[SLST$$BLOCK_LENGTH, BYTE] FIELD (SLST$Z_FIELDS)
%;
```

```
FIELD SLST$Z_FIELDS =
SET
    SLST$V_STRING      = [0, 0, 24, 0],
    SLST$W_LENGTH     = [3, 0, 16, 0],
    SLST$T_STRING     = [5, 0, 0, 0]
TES;
```

```
LITERAL
SLST$K_BASE          = SLST$$BLOCK_LENGTH;
```

```
STRUCTURE
SLST$ELM[I] =
(SLST$ELM + LST$K_BASE + (I * SLST$$BLOCK_LENGTH));
```

String Segment Block



String Segment Block is made up of a Block Header + String Segment Block.

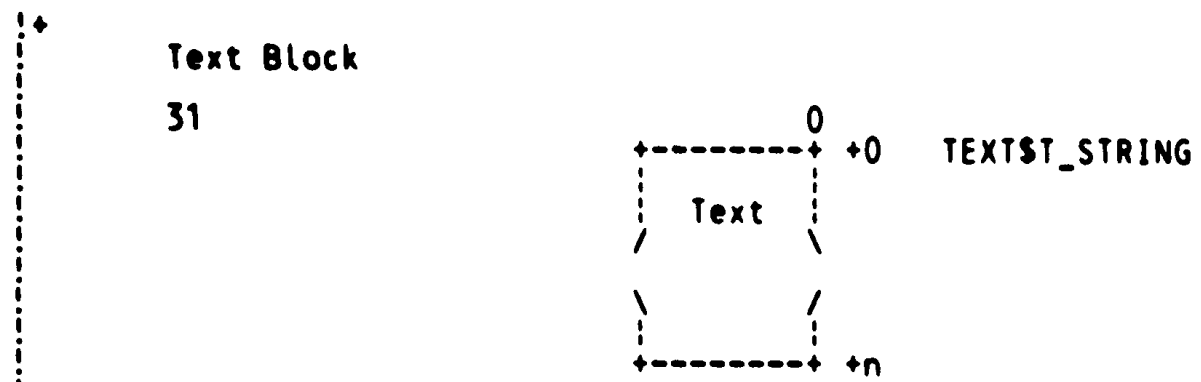
Each string attribute points to zero or more string segment blocks. Each block contains a portion of the whole string. The final string is found by catenating all the string segments together.

```
LITERAL
SEG$$_BLOCK_LENGTH = 3 + BLK$K_BASE;
```

```
MACRO
$SEG = BLOCK[SEG$$_BLOCK_LENGTH, BYTE] FIELD (BLK$Z_FIELDS,
                                                SEG$Z_FIELDS)
X;
```

```
FIELD SEG$Z_FIELDS =
SET
    SEG$V_NEXT           = [0+BLK$K_BASE, 0, 24, 0],
    SEG$T_STRING         = [3+BLK$K_BASE, 0, 0, 0]
TES;
```

```
LITERAL
SEG$K_BASE = SEG$$_BLOCK_LENGTH;
```



Text Block is made up of a Block Header + Text Block.

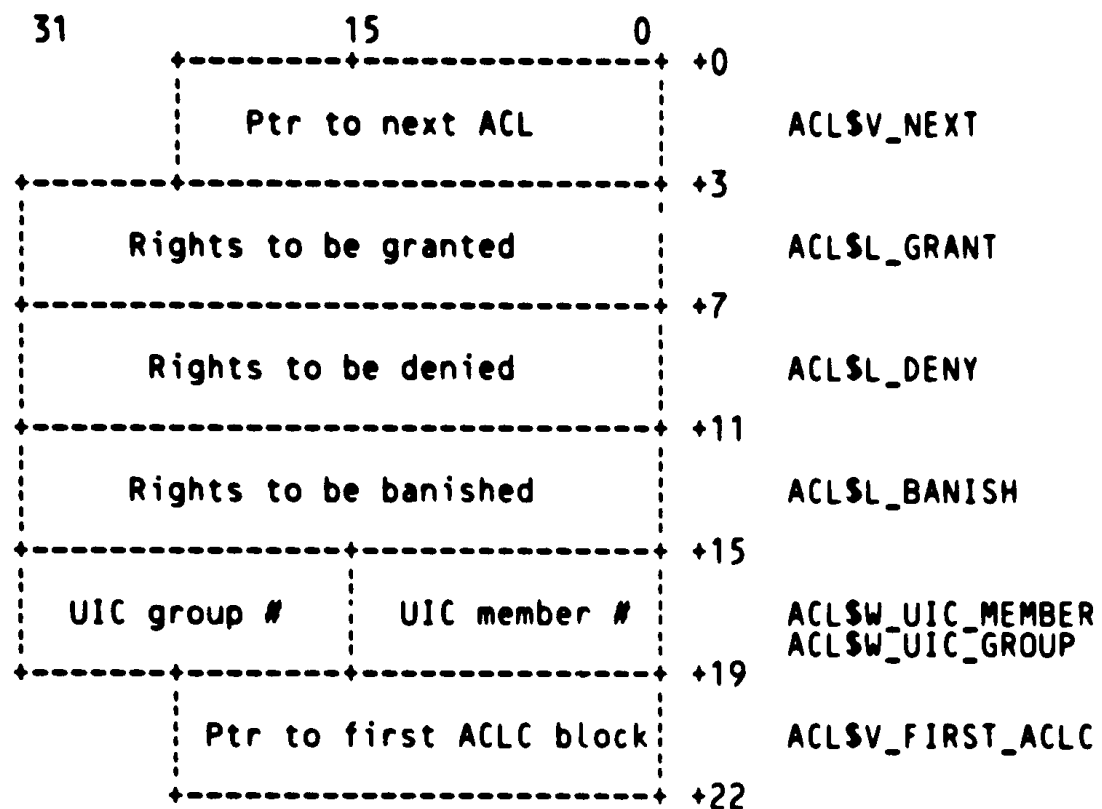
```
LITERAL
TEXT$$BLOCK_LENGTH = 0 + BLK$K_BASE;
```

```
MACRO
$TEXT = BLOCK[TEXT$$BLOCK_LENGTH, BYTE] FIELD (BLK$Z_FIELDS,
TEXT$Z_FIELDS)
%;
```

```
FIELD TEXT$Z_FIELDS =
SET
TEXT$T_STRING . = [0+BLK$K_BASE, 0, 0, 0]
TES;
```

```
LITERAL
TEXT$K_BASE = TEXT$$BLOCK_LENGTH;
```

Access Control List Entry (ACL)



Each node has an Access Control List made up of zero or more Access Control List Entries (ACL).

The ACL block is appended to a BLK to make an access control list entry.

```
LITERAL
ACLSS_BLOCK_LENGTH = 22+BLK$K_BASE;
```

```
MACRO
$ACL = BLOCK[ACLSS_BLOCK_LENGTH, BYTE] FIELD (BLK$Z_FIELDS,
ACL$Z_FIELDS)
%;
```

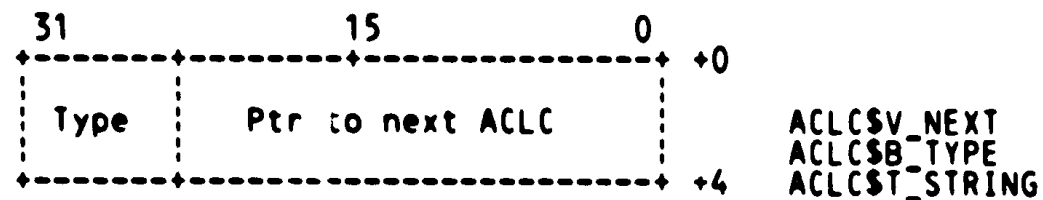
```
FIELD ACL$Z_FIELDS =
SET
ACL$V_NEXT = [0+BLK$K_BASE, 0, 24, 0],
ACL$SL_GRANT = [3+BLK$K_BASE, 0, 32, 0],
ACL$SL_DENY = [7+BLK$K_BASE, 0, 32, 0],
ACL$SL_BANISH = [11+BLK$K_BASE, 0, 32, 0],
ACL$SW_UIC_MEMBER = [15+BLK$K_BASE, 0, 16, 0],
ACL$SW_UIC_GROUP = [17+BLK$K_BASE, 0, 16, 0],
ACL$V_FIRST_ACLC = [19+BLK$K_BASE, 0, 24, 0]
```

TES;

LITERAL
ACLSK_BASE

= ACLSS_BLOCK_LENGTH;

Access Control List Criterion (ACLC)



Each ACL may have one or more Access Control List Criterion (ACLC) chained from it.

Each ACLC specifies one user identification criterion for the ACL entry. The user must match all the criteria for the ACL entry to apply to him.

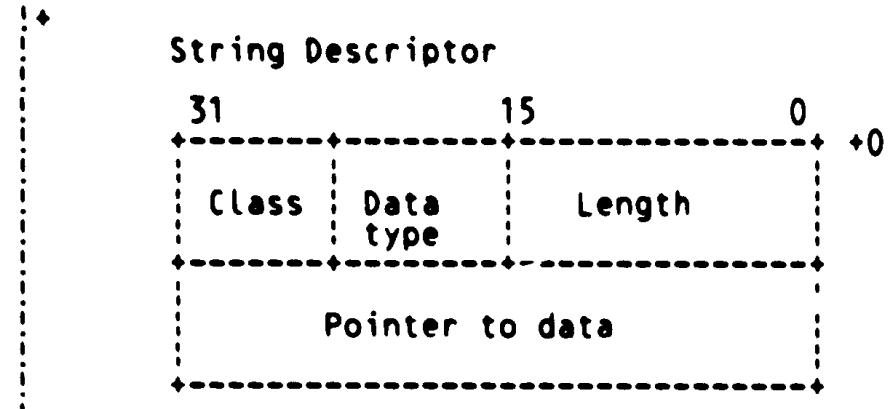
An ACLC is appended to a BLK to form the criterion block.

```
LITERAL
ACLC$S_BLOCK_LENGTH = 4+BLK$K_BASE;
```

```
MACRO
$ACLC = BLOCK[ACLC$S_BLOCK_LENGTH, BYTE] FIELD (BLK$Z_FIELDS,
ACLC$Z_FIELDS)
%;
```

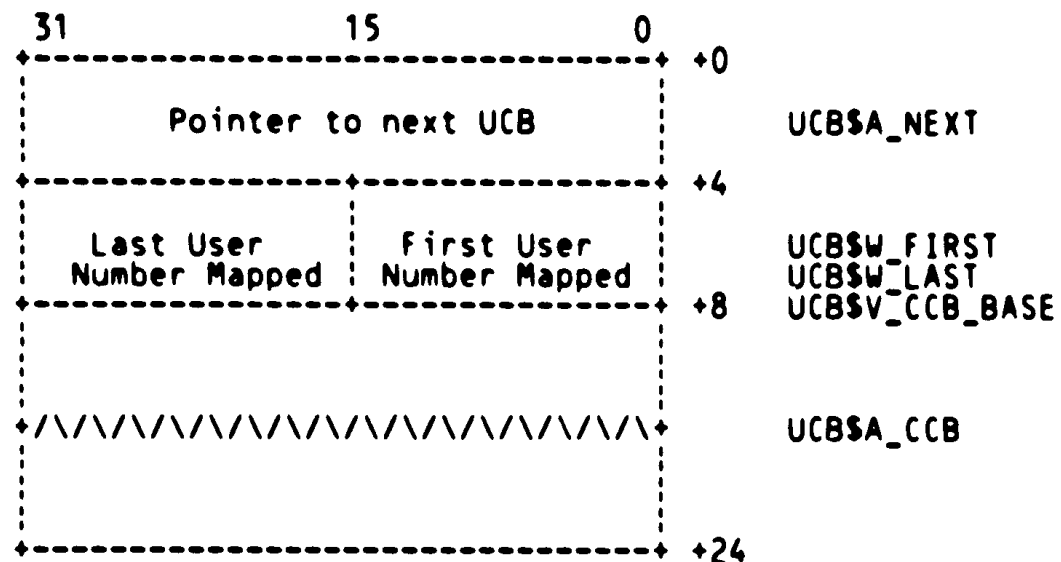
```
FIELD ACLC$Z_FIELDS =
SET
ACLC$V_NEXT = [0+BLK$K_BASE, 0, 24, 0],
ACLC$B_TYPE = [3+BLK$K_BASE, 0, 8, 0],
ACLC$T_STRING = [4+BLK$K_BASE, 0, 0, 0]
TES;
```

```
LITERAL
ACLC$K_BASE = ACLC$S_BLOCK_LENGTH;
```

```
MACRO  
  $DSC = BLOCK[8,BYTE]  
  %;
```

User Control Block (UCB)



A user's context pointer is an origin 1 index into the CCBs in the UCB list. Each UCB can point to UCB\$K_CCB number of CCBs. If the user passes a context number that doesn't map to an active CCB, we tell him it's an invalid context pointer.

Except in some extreme case, any image will probably never have more than 4 users active at any one time. But we can handle more!

```
LITERAL
UCB$K_CCB_BASE      = 8,
UCB$K_CCB           = 4,      ! Number of CCBs/UCB
UCB$$BLOCK_LENGTH  = UCB$K_CCB_BASE + UCB$K_CCB * 4;
```

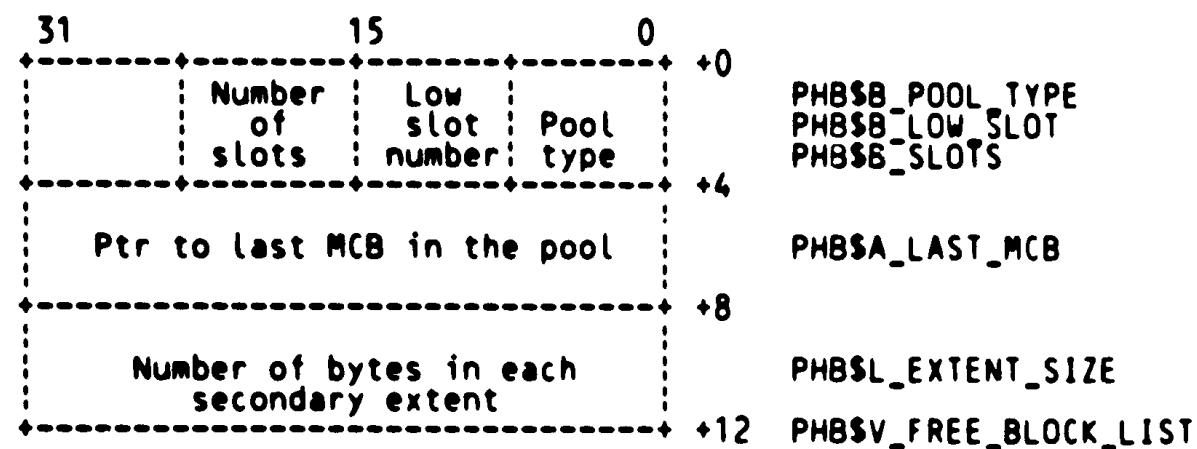
```
MACRO
%UCB = BLOCK[UCB$$BLOCK_LENGTH, BYTE] FIELD (UCB$Z_FIELDS)
%;
```

```
FIELD   UCB$Z_FIELDS =
SET
UCB$A_NEXT      = [0, 0, 32, 0],
UCB$W_FIRST     = [4, 0, 16, 0],
UCB$W_LAST     = [6, 0, 16, 0],
UCB$V_CCB      = [8, 0, 0, 0],
UCB$A_CCB      = [0, 0, 32, 0]
TES;
```

```
LITERAL
UCB$C_FIRST = BLOCK[0, UCB$W_FIRST; UCB$$BLOCK_LENGTH, BYTE];
```

```
STRUCTURE
UCB$CCB[I, O, P, S, E] = (UCB$CCB+UCB$K_CCB_BASE+
(I-(UCB$CCB+UCB$C_FIRST)<0,16,0>)*4+O)<P,S,E>;
```

Pool Header Block (PHB)



Pools are used to provide space for various in-core block types. Blocks that are related are usually allocated from the same pool. This is done because pools provide good locality of reference and this allocation scheme reduces page faults.

Each pool consists of one or more extents. Each extent consists of a Memory Control Block and the rest of the space allocated to the extent. Each pool has a Pool Header Block associated with it. This block identifies the MCB list, as well as contains the pool's free block list.

The free block list is a vector of linked lists. Each slot in the vector corresponds to a block type. When a block is freed, it is linked into its associated free list. When a block is requested, its free list is checked to see if it is non-empty. If so, then a block is allocated from it. Otherwise, a block is allocated from one of the pool's extents.

```

LITERAL
  PHB$S_BLOCK_LENGTH      = 12;

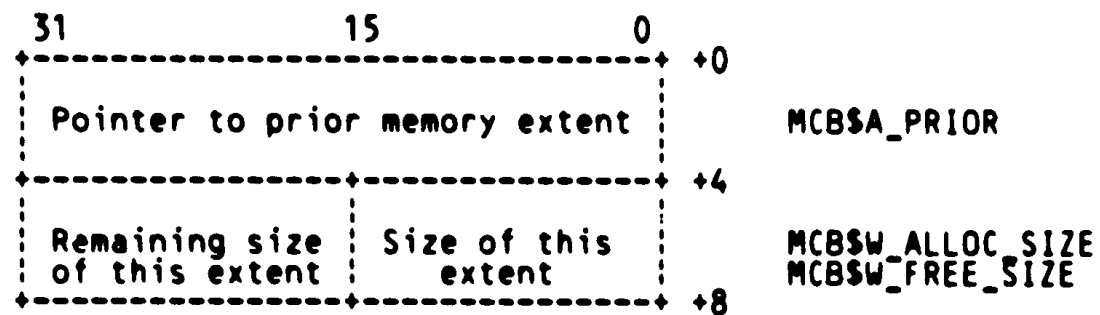
MACRO
  $PHB = BLOCK[PHB$S_BLOCK_LENGTH, BYTE] FIELD (PHB$Z_FIELDS)
  %:

FIELD  PHB$Z_FIELDS =
SET
  PHB$B_POOL_TYPE      = [0, 0, 8, 0],
  PHB$B_LOW_SLOT      = [1, 0, 8, 0],
  PHB$B_SLOTS         = [2, 0, 8, 0],
  PHB$A_LAST_MCB     = [4, 0, 32, 0],
  PHB$L_EXTENT_SIZE  = [8, 0, 32, 0],
  PHB$V_FREE_BLOCK_LIST = [12, 0, 0, 0]

TES;

```

Memory Control Block (MCB)



Each pool has one or more Memory Control Blocks associated with it.

This list serves two purposes:

- 1) It points to each memory extent we asked LIB\$GET_VM for.
- 2) It keeps track of memory that has been allocated for the pool, but never used.

LITERAL

MCBSS_BLOCK_LENGTH = 8;

MACRO

\$MCB = BLOCK[MCBSS_BLOCK_LENGTH, BYTE] FIELD (MCBSZ_FIELDS)

%;

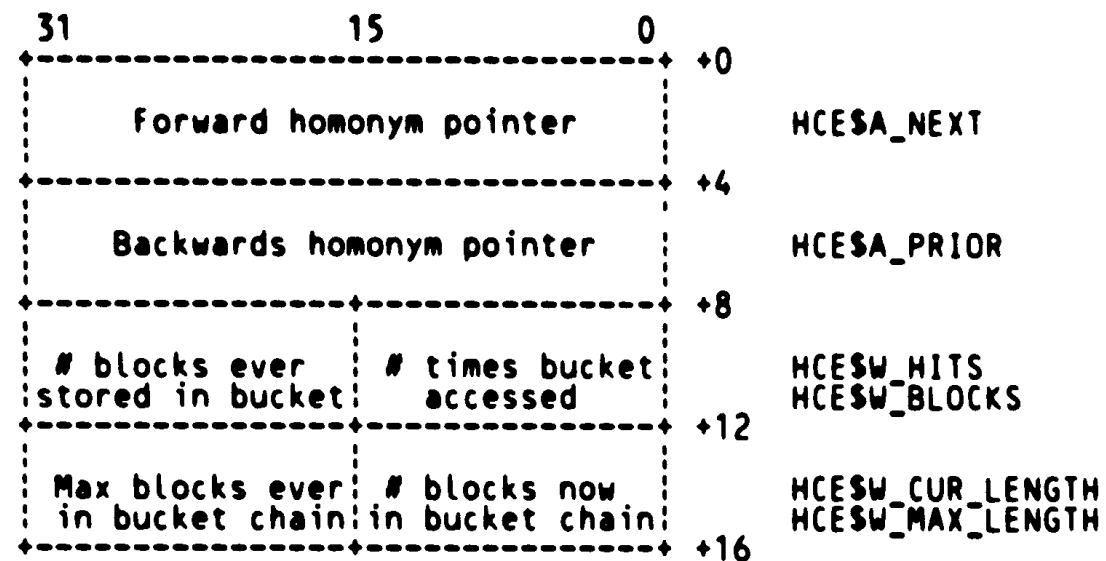
FIELD MCB\$Z_FIELDS =

SET

MCBSA_PRIOR = [0, 0, 32, 0],
 MCB\$W_ALLOC_SIZE = [4, 0, 16, 0],
 MCB\$W_FREE_SIZE = [6, 0, 16, 0]

TES;

Hash Control Entry (HCE)



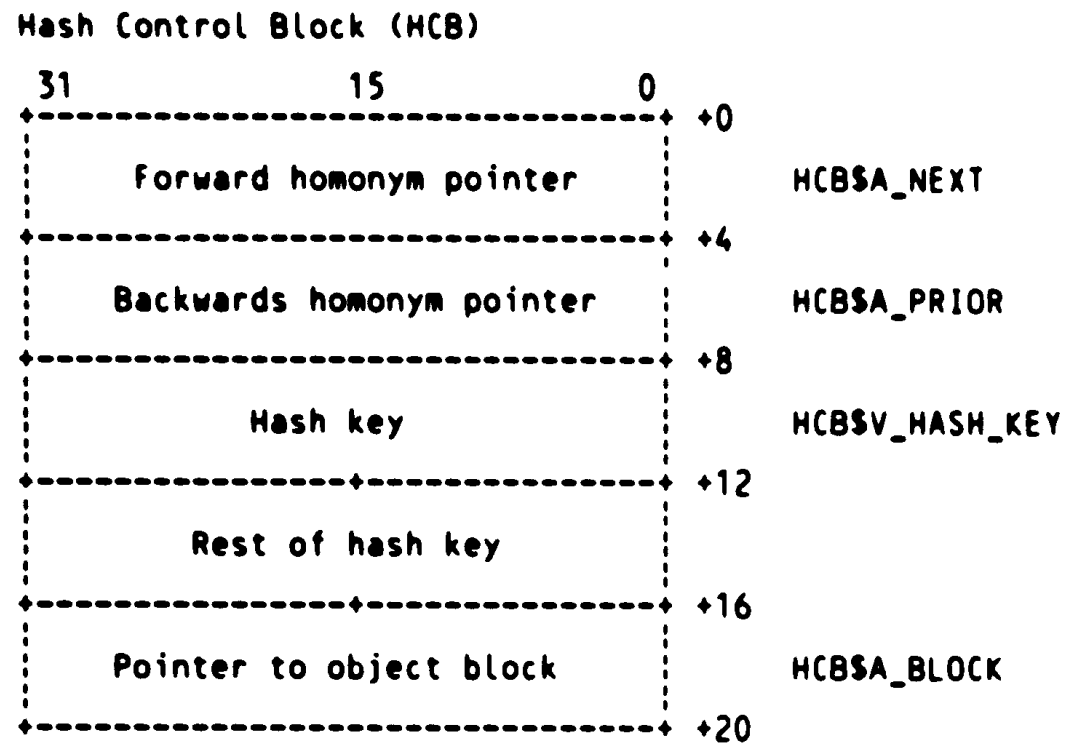
HCEs are found in each user's CCB. Each HCE forms one hash bucket. We keep some statistics on the buckets in an attempt to tune the hash function.

```
LITERAL
HCESS_BLOCK_LENGTH = 16;
```

```
MACRO
$HCE = BLOCK[HCESS_BLOCK_LENGTH, BYTE] FIELD (HCE$Z_FIELDS)
X;
```

```
FIELD HCE$Z_FIELDS =
SET
HCE$A_NEXT           = [0, 0, 32, 0],
HCE$A_PRIOR         = [4, 0, 32, 0],
HCE$W_HITS           = [8, 0, 16, 0],
HCE$W_BLOCKS         = [10, 0, 16, 0],
HCE$W_CUR_LENGTH     = [12, 0, 16, 0],
HCE$W_MAX_LENGTH     = [14, 0, 16, 0]
TES;
```

```
LITERAL
HCE$K_HOMONYM_LIST = BLOCK[0, HCE$A_NEXT; .BYTE];
```



LITERAL

```

HCBSK_PCB_NUMBER    = 1,      ! Hash type is page number
HCBSK_LCCC_CODE     = 2,      ! Hash type is location code
HCBSK_LCCB_ADDRESS  = 3,      ! Hash type is entity disk address
HCBSK_KEY_LENGTH    = 8,      ! Size of a hash key
HCBS$BLOCK_LENGTH  = 20;
    
```

MACRO

```

$HCB = BLOCK[HCBS$BLOCK_LENGTH, BYTE] FIELD (HCBSZ_FIELDS)
%:
    
```

FIELD HCBSZ_FIELDS =

```

SET
  HCBSA_NEXT      = [0, 0, 32, 0],
  HCBSA_PRIOR     = [4, 0, 32, 0],
  HCBSV_HASH_KEY  = [8, 0, 0, 0],
  HCBSA_BLOCK     = [16, 0, 32, 0]
    
```

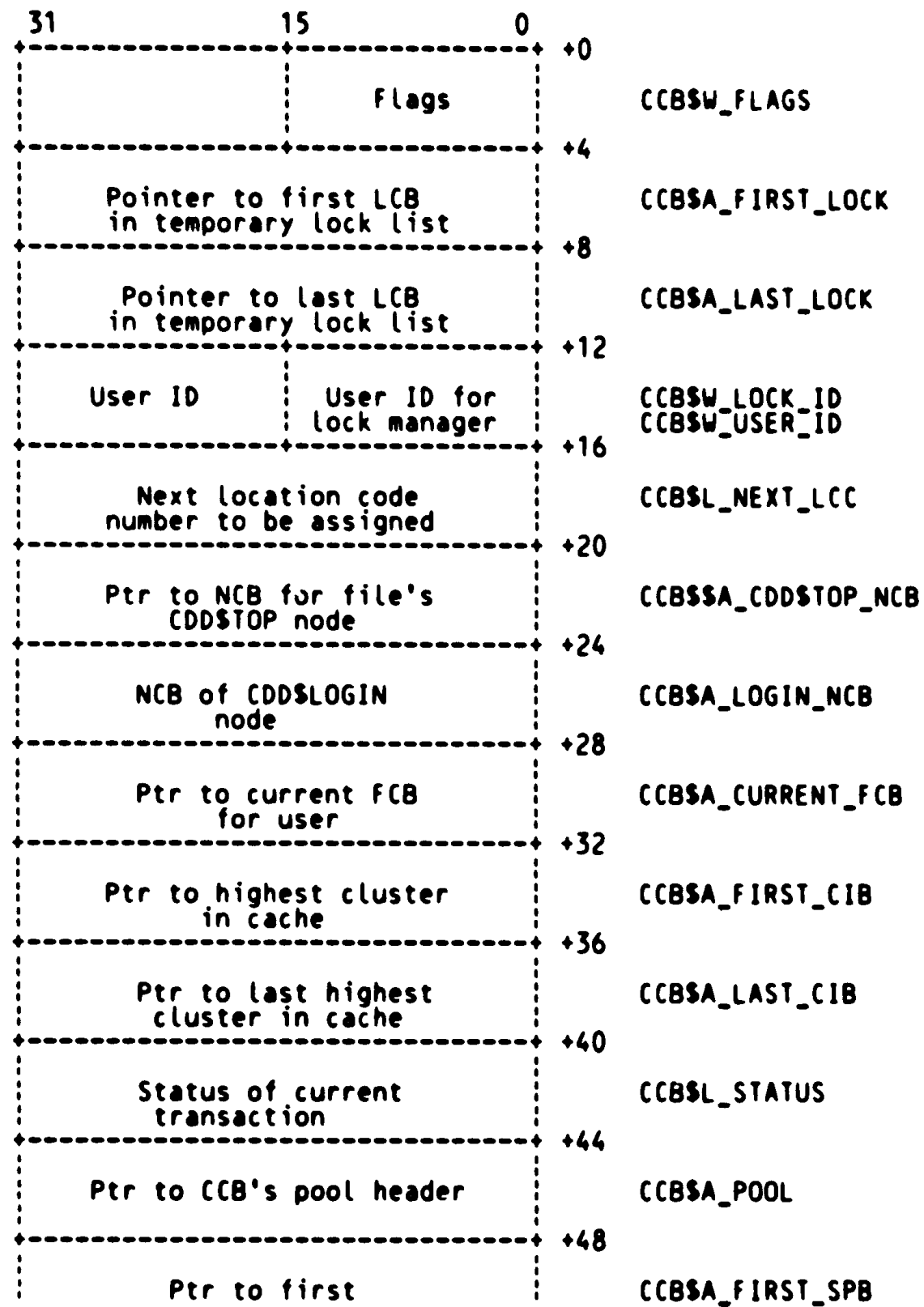
TES;

LITERAL

```

HCBSK_HOMONYM_LIST = BLOCK[0, HCBSA_NEXT; , BYTE];
    
```

Context Control Block (CCB)



CCBSW_FLAGS

CCBSA_FIRST_LOCK

CCBSA_LAST_LOCK

CCBSW_LOCK_ID
CCBSW_USER_ID

CCBSL_NEXT_LCC

CCBSA_CDD\$TOP_NCB

CCBSA_LOGIN_NCB

CCBSA_CURRENT_FCB

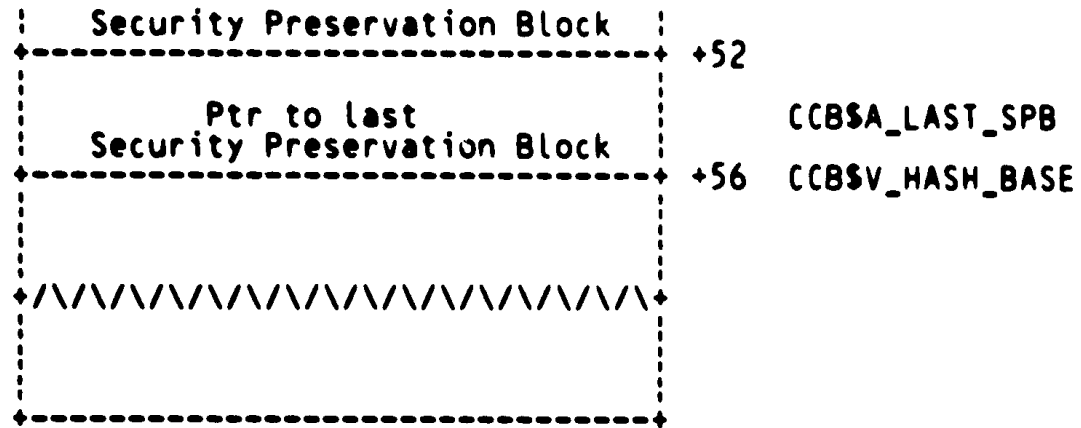
CCBSA_FIRST_CIB

CCBSA_LAST_CIB

CCBSL_STATUS

CCBSA_POOL

CCBSA_FIRST_SPB



The hash table consists of a number of hash entries.
See the HCE description for the format of these entries.

```

LITERAL
CCBSK_HASH_TABLE_SIZE = 151,
CCBS$BLOCK_LENGTH    = 56 + (CCBSK_HASH_TABLE_SIZE * HCESS_BLOCK_LENGTH);
    
```

```

MACRO
$CCB = BLOCK[CCBS$BLOCK_LENGTH, BYTE] FIELD (CCBSZ_FIELDS)
%:
    
```

```

FIELD CCBSZ_FIELDS =
SET
  CCBSW_FLAGS           = [0, 0, 16, 0],
  CCBSV_CORRUPT        = [0, 0, 1, 0],           ! Stream is corrupt
  CCBSA_FIRST_LOCK     = [4, 0, 32, 0],
  CCBSA_LAST_LOCK      = [8, 0, 32, 0],
  CCBSW_LOCK_ID        = [12, 0, 16, 0],
  CCBSW_USER_ID        = [14, 0, 16, 0],
  CCBSL_NEXT_LCC       = [16, 0, 32, 0],
  CCBSA_CDD$STOP_NCB   = [20, 0, 32, 0],
  CCBSA_LOGIN_NCB      = [24, 0, 32, 0],
  CCBSA_CURRENT_FCB    = [28, 0, 32, 0],
  CCBSA_FIRST_CIB      = [32, 0, 32, 0],
  CCBSA_LAST_CIB       = [36, 0, 32, 0],
  CCBSL_STATOS         = [40, 0, 32, 0],
  CCBSA_POOL           = [44, 0, 32, 0],
  CCBSA_FIRST_SPB      = [48, 0, 32, 0],
  CCBSA_LAST_SPB       = [52, 0, 32, 0],
  CCBSV_HASH_BASE      = [56, 0, 0, 0]
TES:
    
```

```

LITERAL
CCBSK_LOCK_LIST      = BLOCK[0, CCBSA_FIRST_LOCK; , BYTE],
CCBSK_CLUSTER_LIST  = BLOCK[0, CCBSA_FIRST_CIB; , BYTE],
CCBS$HASH_BASE       = BLOCK[0, CCBSV_HASH_BASE; , BYTE],
CCBSK_SPB_LIST       = BLOCK[0, CCBSA_FIRST_SPB; , BYTE];
    
```

```

STRUCTURE
CCB$HASH[1, 0, P, S, E] =
    
```

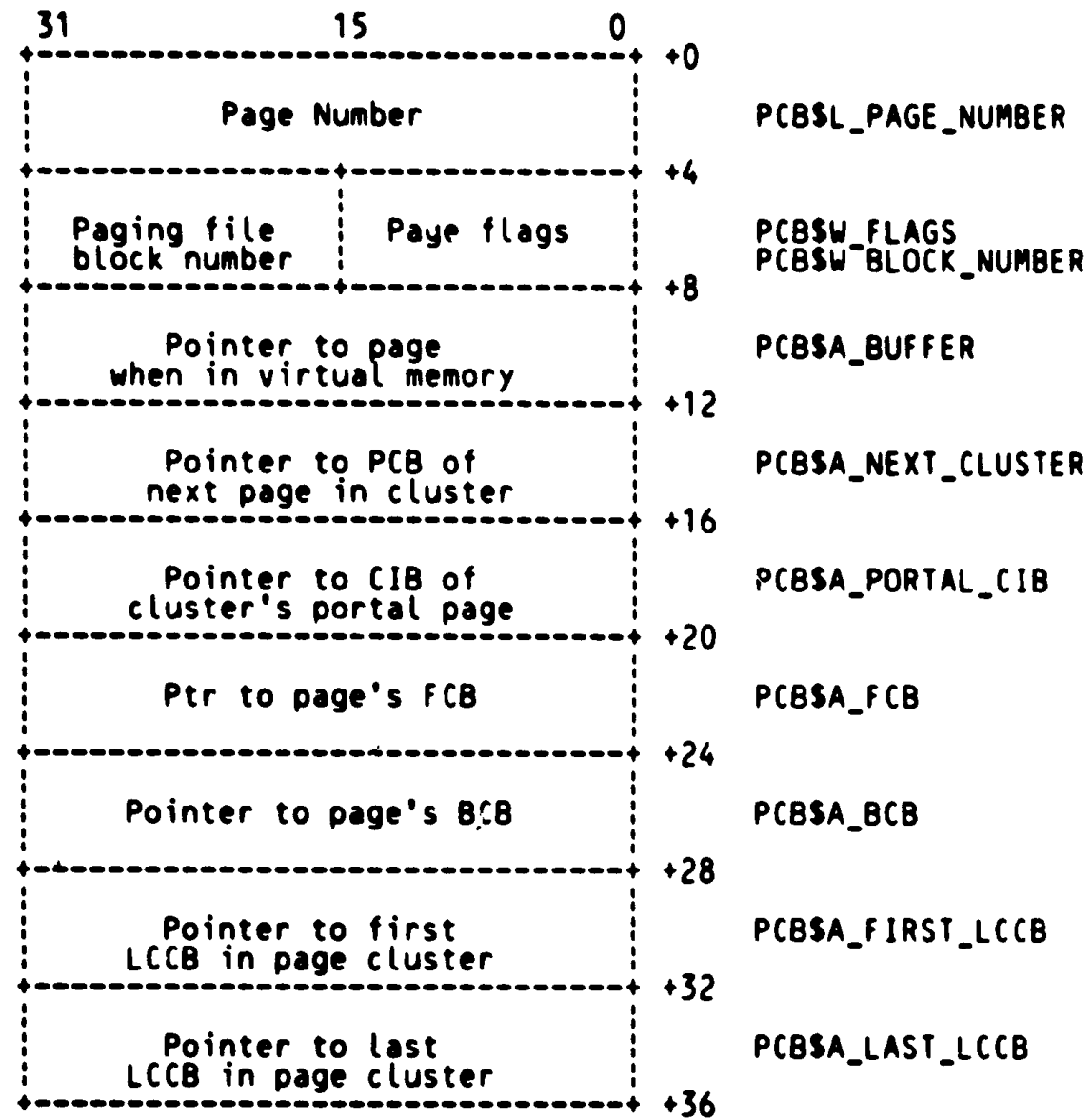

(CCB\$HASH+CCB\$\$_HASH_BASE+0+(I+ACCESS_BLOCK_LENGTH))<P,S,E>;

LITERAL

CCB\$M_CORRUPT = 1^1 - 1^0; ! Stream is corrupt

↑

Page Control Block (PCB)



A Page Control Block (PCB) exists for every page in the cache, and for pages that only have presence locks on them (thus they're not in the staging cache). Portal pages have a Cluster Information Block (CIB) appending to their PCB.

LITERAL
PCBS\$BLOCK_LENGTH = 36;

MACRO
\$PCB = BLOCK[PCBS\$BLOCK_LENGTH, BYTE] FIELD (PCBS\$Z_FIELDS)
X;

```

FIELD  PCBSZ_FIELDS =
SET
PCBSL_PAGE_NUMBER      = [0, 0, 32, 0],
PCBSW_FLAGS            = [4, 0, 16, 0],
PCBSV_MODIFIED         = [4, 0, 1, 0],      ! Must write back to dict.
PCBSV_NEW_PAGE         = [4, 3, 1, 0],      ! Page was in free chain
PCBSV_READ_ONLY        = [4, 4, 1, 0],      ! Page is read only
PCBSV_FREE_PAGE        = [4, 5, 1, 0],      ! Page is a free page
PCBSV_AVAILABLE       = [4, 6, 1, 0],      ! Page available in cache
PCBSV_PORTAL_PAGE      = [4, 7, 1, 0],      ! Page is portal page, CIB follows
PCBSW_BLOCK_NUMBER     = [6, 0, 16, 0],
PCBSA_BUFFER           = [8, 0, 32, 0],
PCBSA_NEXT_CLUSTER    = [12, 0, 32, 0],
PCBSA_PORTAL_CIB       = [16, 0, 32, 0],
PCBSA_FCB              = [20, 0, 32, 0],
PCBSA_BCB              = [24, 0, 32, 0],
PCBSA_FIRST_LCCB      = [28, 0, 32, 0],
PCBSA_LAST_LCCB       = [32, 0, 32, 0]
TES:

```

```

LITERAL
PCBSK_LCCB_LIST        = BLOCK[0, PCBSA_FIRST_LCCB; , BYTE];

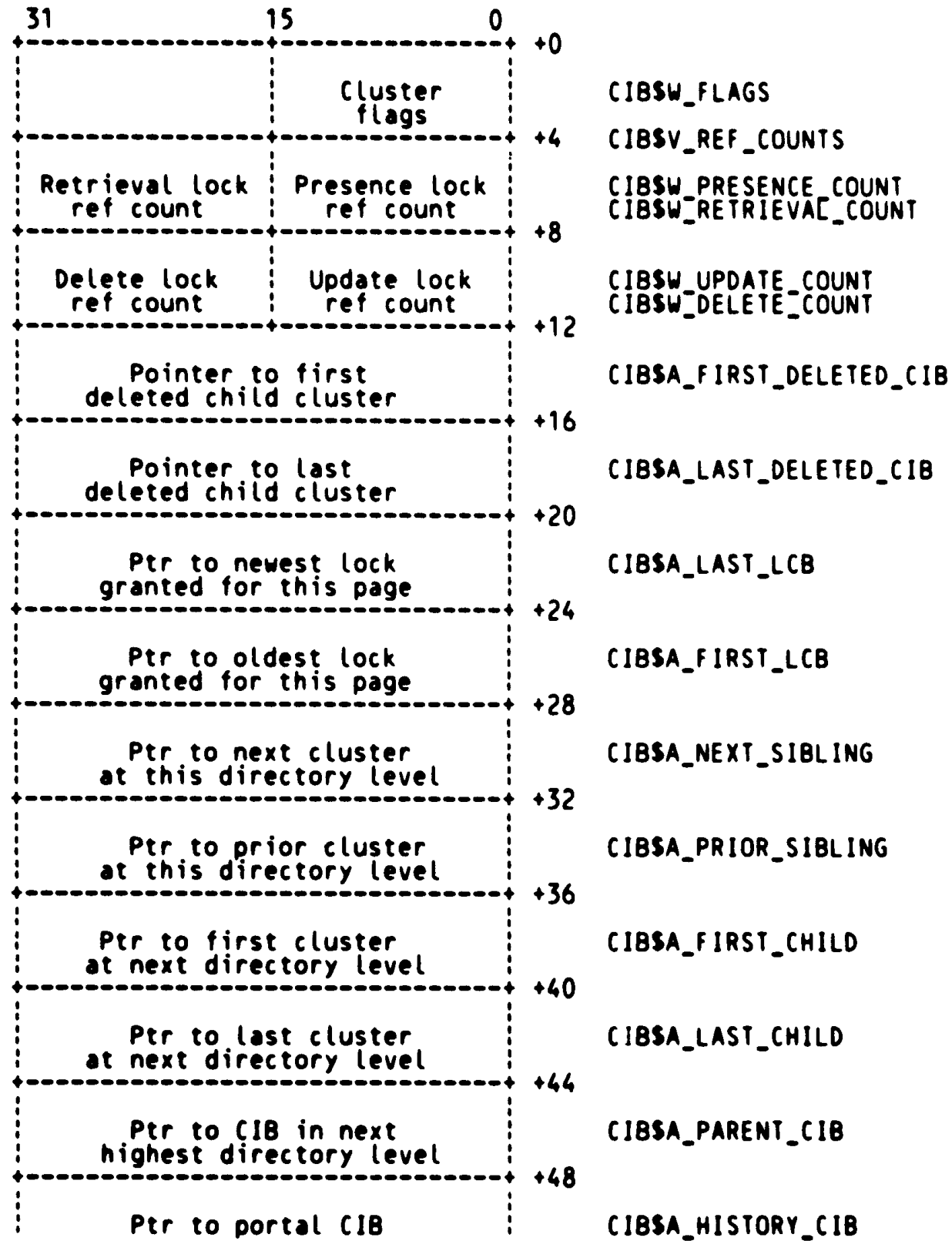
```

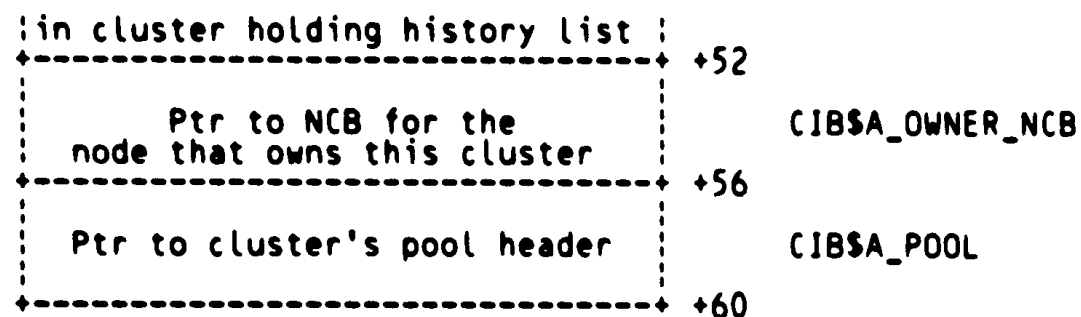
```

LITERAL
PCBSM_MODIFIED         = 1^1 - 1^0,      ! Must write back to dict.
PCBSM_NEW_PAGE         = 1^4 - 1^3,      ! Page was in free chain
PCBSM_READ_ONLY        = 1^5 - 1^4,      ! Page is read only
PCBSM_FREE_PAGE        = 1^6 - 1^5,      ! Page is a free page
PCBSM_AVAILABLE       = 1^7 - 1^6,      ! Page available in cache
PCBSM_PORTAL_PAGE      = 1^8 - 1^7;      ! Portal page, CIB follows

```

Cluster Information Block (CIB)





Each clusters' portal page has a CIB associated with it. This block is physically appended to the PCB of the cluster's portal page.

The CIB also points to the cluster's pool header.

The CIB\$V_REF_COUNTS lists must have each of the lock ref counts in the same order as the LOCK\$K_xxx lock request symbols.

LITERAL

```
CIB$$BLOCK_LENGTH = 60 + PCB$$BLOCK_LENGTH;
```

MACRO

```
%CIB = BLOCK[CIB$$BLOCK_LENGTH, BYTE] FIELD (CIB$Z_FIELDS, PCB$Z_FIELDS)
```

```
%;
```

```
FIELD SET CIB$Z_FIELDS =
```

```

CIB$W_FLAGS           = [0+PCB$$BLOCK_LENGTH, 0, 16, 0],
CIB$V_LOCKED_DELETE  = [0+PCB$$BLOCK_LENGTH, 0, 1, 0],
CIB$V_LOCKED_UPDATE  = [0+PCB$$BLOCK_LENGTH, 1, 1, 0],
CIB$V_LOCKED_RETRIEVAL = [0+PCB$$BLOCK_LENGTH, 2, 1, 0],
CIB$V_LOCKED_PRESENCE = [0+PCB$$BLOCK_LENGTH, 3, 1, 0],
CIB$V_LOCKED         = [0+PCB$$BLOCK_LENGTH, 0, 4, 0],
CIB$V_LOGIN          = [0+PCB$$BLOCK_LENGTH, 4, 1, 0],
CIB$V_COMPLETE       = [0+PCB$$BLOCK_LENGTH, 5, 1, 0],
CIB$V_NEW            = [0+PCB$$BLOCK_LENGTH, 6, 1, 0],
CIB$V_HISTORY        = [0+PCB$$BLOCK_LENGTH, 7, 1, 0],
CIB$V_REF_COUNTS     = [4+PCB$$BLOCK_LENGTH, 0, 0, 0],
CIB$W_PRESENCE_COUNT = [4+PCB$$BLOCK_LENGTH, 0, 16, 0],
CIB$W_RETRIEVAL_COUNT = [6+PCB$$BLOCK_LENGTH, 0, 16, 0],
CIB$W_UPDATE_COUNT   = [8+PCB$$BLOCK_LENGTH, 0, 16, 0],
CIB$W_DELETE_COUNT   = [10+PCB$$BLOCK_LENGTH, 0, 16, 0],
CIB$A_FIRST_DELETED_CIB = [12+PCB$$BLOCK_LENGTH, 0, 32, 0],
CIB$A_LAST_DELETED_CIB = [16+PCB$$BLOCK_LENGTH, 0, 32, 0],
CIB$A_LAST_LCB       = [20+PCB$$BLOCK_LENGTH, 0, 32, 0],
CIB$A_FIRST_LCB      = [24+PCB$$BLOCK_LENGTH, 0, 32, 0],
CIB$A_NEXT_SIBLING   = [28+PCB$$BLOCK_LENGTH, 0, 32, 0],
CIB$A_PRIOR_SIBLING  = [32+PCB$$BLOCK_LENGTH, 0, 32, 0],
CIB$A_FIRST_CHILD    = [36+PCB$$BLOCK_LENGTH, 0, 32, 0],
CIB$A_LAST_CHILD     = [40+PCB$$BLOCK_LENGTH, 0, 32, 0],
CIB$A_PARENT_CIB     = [44+PCB$$BLOCK_LENGTH, 0, 32, 0],

```

```

CIBSA_HISTORY_CIB      = [48+PCB$$_BLOCK_LENGTH, 0, 32, 0],
CIBSA_OWNER_NCB       = [52+PCB$$_BLOCK_LENGTH, 0, 32, 0],
CIBSA_POOL             = [56+PCB$$_BLOCK_LENGTH, 0, 32, 0]
TES;

```

```

LITERAL
CIBSK_LOCK_LIST       = BLOCK[0, CIBSA_LAST_LCB; , BYTE],
CIBSK_SIBLING_LIST   = BLOCK[0, CIBSA_NEXT_SIBLING; , BYTE],
CIBSK_CHILD_LIST     = BLOCK[0, CIBSA_FIRST_CHILD; , BYTE];

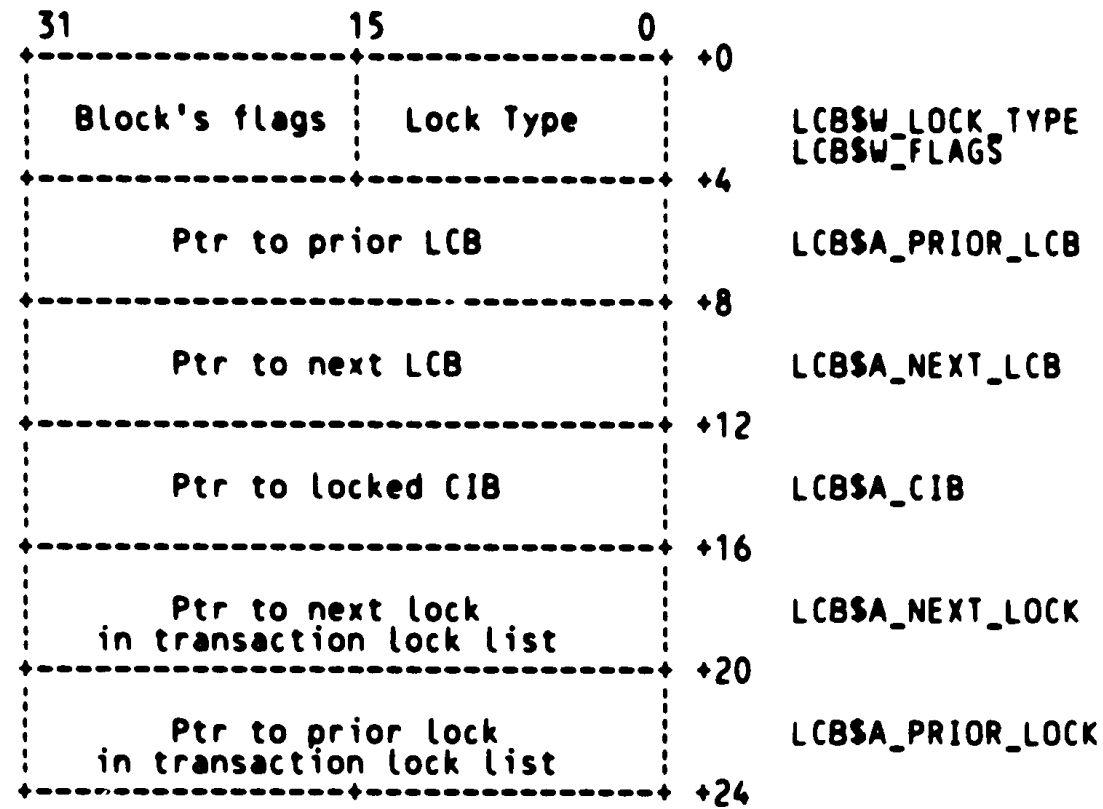
```

```

LITERAL
CIBSM_LOCKED_DELETE   = 1^1 - 1^0,
CIBSM_LOCKED_UPDATE   = 1^2 - 1^1,
CIBSM_LOCKED_RETRIEVAL = 1^3 - 1^2,
CIBSM_LOCKED_PRESENCE = 1^4 - 1^3,
CIBSM_LOCKED          = 1^4 - 1^0,
CIBSM_LOGIN           = 1^5 - 1^4,      ! Cluster is on login path
CIBSM_COMPLETE        = 1^6 - 1^5,      ! Cluster is complete in cache
CIBSM_NEW              = 1^7 - 1^6,      ! Cluster read in this transaction
CIBSM_HISTORY         = 1^8 - 1^7;      ! History list is in cluster

```

Lock Control Block (LCB)



LCBs are used to keep track of which locks exist on a cluster (CIB).

Each LCB is linked to its CIB, and to other LCBs for that cluster. When an LCB is granted, it is placed in the transaction's lock list until the transaction terminates.

LITERAL

LCBSS_BLOCK_LENGTH = 24;

MACRO

\$LCB = BLOCK[LCBSS_BLOCK_LENGTH, BYTE] FIELD (LCBSZ_FIELDS)

%;

FIELD LCBSZ_FIELDS =

SET

LCBSW_LOCK_TYPE	=	[0, 0, 16, 0],	
LCBSW_FLAGS	=	[2, 0, 16, 0],	
LCBSV_CURRENT	=	[2, 1, 1, 0],	! Allocated in current transaction
LCBSA_PRIOR_LCB	=	[4, 0, 32, 0],	
LCBSA_NEXT_LCB	=	[8, 0, 32, 0],	
LCBSA_CIB	=	[12, 0, 32, 0],	
LCBSA_NEXT_LOCK	=	[16, 0, 32, 0],	
LCBSA_PRIOR_LOCK	=	[20, 0, 32, 0],	

TES;

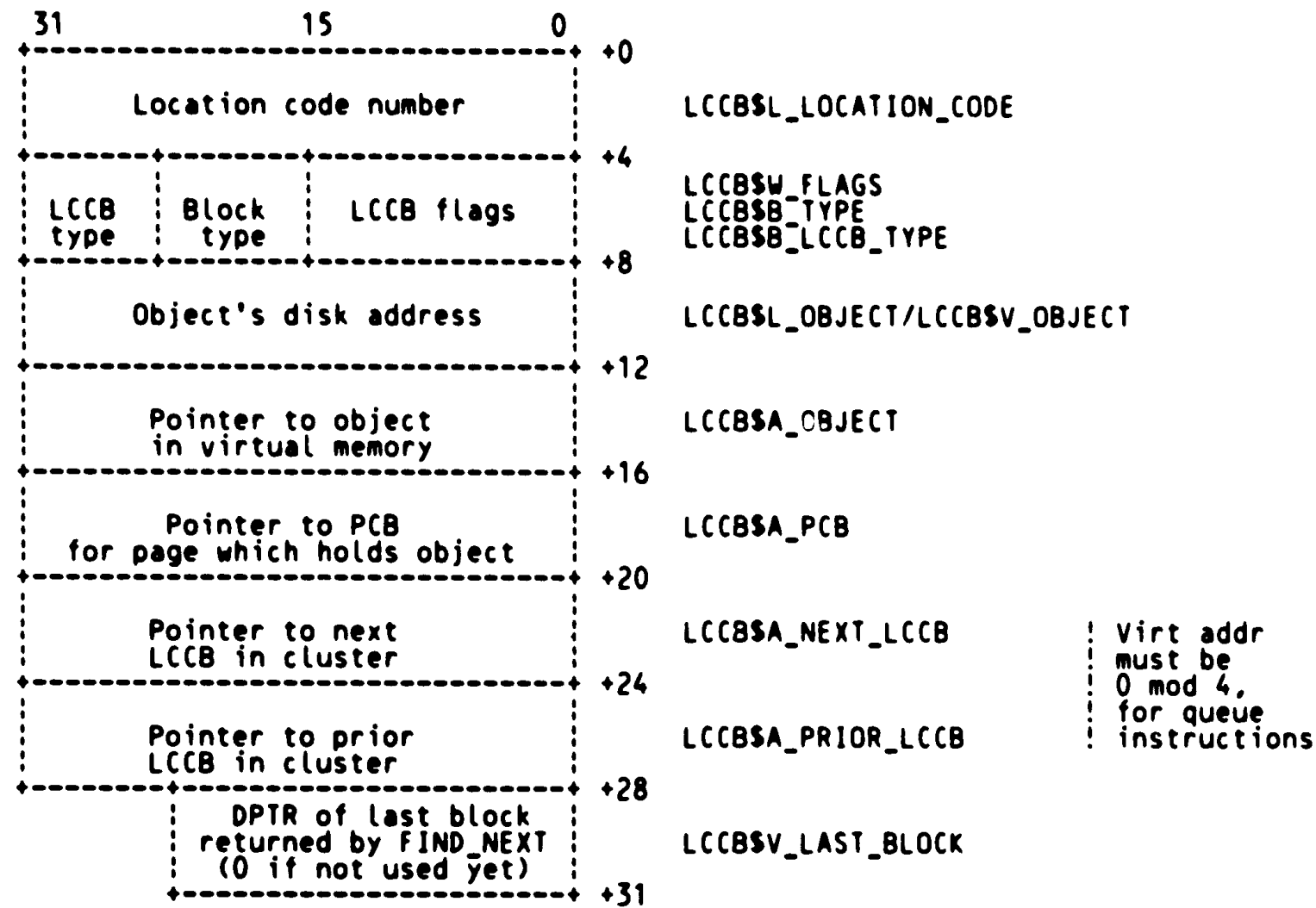
LITERAL
LCBSK_LOCK_LIST
LCBSK_TEMP_LOCK_LIST

= BLOCK[0, LCBSA_PRIOR_LCB; , BYTE],
= BLOCK[0, LCBSA_NEXT_LOCK; , BYTE];

LITERAL
LCBSM_CURRENT

= 1^2 - 1^1; ! Allocated in current transaction

Location Code Control Block (LCCB)



Location codes are used to provide a convenient means for the user to identify a particular object. Each location code is associated with an LCCB. The following objects may be assigned location code:

- 1) Entities
- 2) Lists
- 3) Nodes

LITERAL
LCCBS_BLOCK_LENGTH = 31;

MACRO

```
%LCCB = BLOCK[LCCBS$ _BLOCK_LENGTH, BYTE] FIELD (LCCBS$Z_FIELDS)
%;
```

```
FIELD LCCBS$Z_FIELDS =
SET
```

```
LCCBSL_LOCATION_CODE = [0, 0, 32, 0],
LCCBSW_FLAGS = [4, 0, 16, 0],
LCCBSV_AVAILABLE = [4, 0, 1, 0],
LCCBSV_MUST_SCAN = [4, 1, 1, 0],
LCCBSV_GHOST = [4, 2, 1, 0],
LCCBSV_DIRECTORY = [4, 3, 1, 0],
LCCBSV_TERMINAL = [4, 4, 1, 0],
LCCBSV_ENTITY_ATT = [4, 5, 1, 0],
LCCBSV_ENTITY_LIST_ATT = [4, 6, 1, 0],
LCCBSV_ENTITY_LIST = [4, 7, 1, 0],
LCCBSV_STRING_LIST_ATT = [4, 8, 1, 0],
LCCBSB_TYPE = [6, 0, 8, 0],
LCCBSB_LCCB_TYPE = [7, 0, 8, 0],
LCCBSL_OBJECT = [8, 0, 32, 0],
LCCBSV_OBJECT = [8, 0, 24, 0],
LCCBSA_OBJECT = [12, 0, 32, 0],
LCCBSA_PCB = [16, 0, 32, 0],
LCCBSA_NEXT_LCCB = [20, 0, 32, 0],
LCCBSA_PRIOR_LCCB = [24, 0, 32, 0],
LCCBSV_LAST_BLOCK = [28, 0, 24, 0]
```

```
TES;
```

```
LITERAL
```

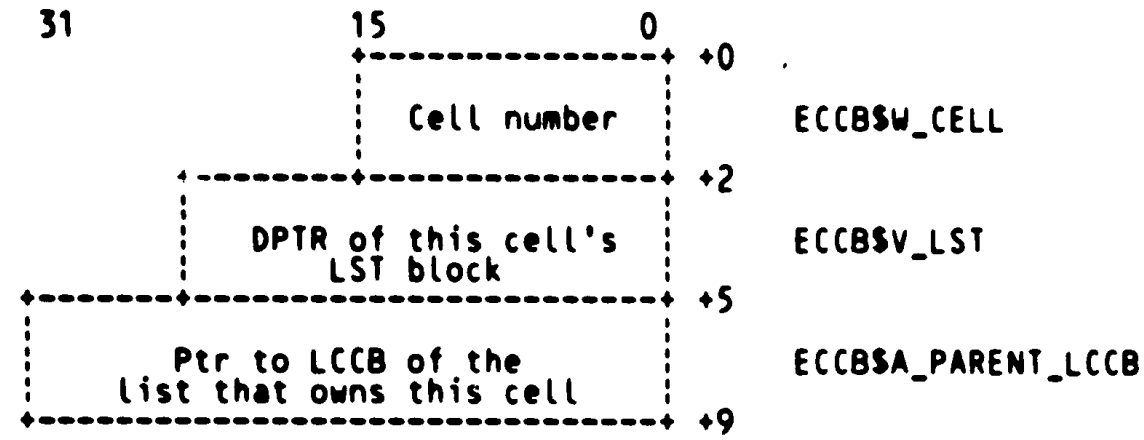
```
LCCBSK_LCCB_LIST = BLOCK[0, LCCBSA_NEXT_LCCB; , BYTE];
```

```
LITERAL
```

```
LCCBSM_AVAILABLE = 1^1 - 1^0, ! Object's virtual addr is known
LCCBSM_MUST_SCAN = 1^2 - 1^1, ! Protection tree must be scanned
LCCBSM_GHOST = 1^3 - 1^2, ! LCCB may not be fetched by LCC
LCCBSM_DIRECTORY = 1^4 - 1^3, ! Directory NCB
LCCBSM_TERMINAL = 1^5 - 1^4, ! Terminal NCB
LCCBSM_ENTITY_ATT = 1^6 - 1^5, ! Entity attribute LCCB
LCCBSM_ENTITY_LIST_ATT = 1^7 - 1^6, ! Entity list attribute LCCB
LCCBSM_ENTITY_LIST = 1^8 - 1^7, ! Entity list ECCB
LCCBSM_STRING_LIST_ATT = 1^9 - 1^8, ! String list LCCB
```

```
LCCBSK_LCCB_TYPE_FIRST = 1,
LCCBSK_NCB = 1, ! LCCB includes NCB
LCCBSK_ECCB = 2, ! LCCB includes ECCB
LCCBSK_LCCB = 3, ! LCCB
LCCBSK_LCCB_TYPE_LAST = 3;
```

Entity Cell Control Block (ECCB)



Each cell in an entity list may be assigned a location code.

This block is appended to the location code's LCCB to name the specific cell represented by the location code.

LITERAL

ECCBSS_BLOCK_LENGTH = LCCBSS_BLOCK_LENGTH + 9;

MACRO

SECCB = BLOCK[ECCBSS_BLOCK_LENGTH, BYTE] FIELD (ECCBSZ_FIELDS, LCCBSZ_FIELDS)

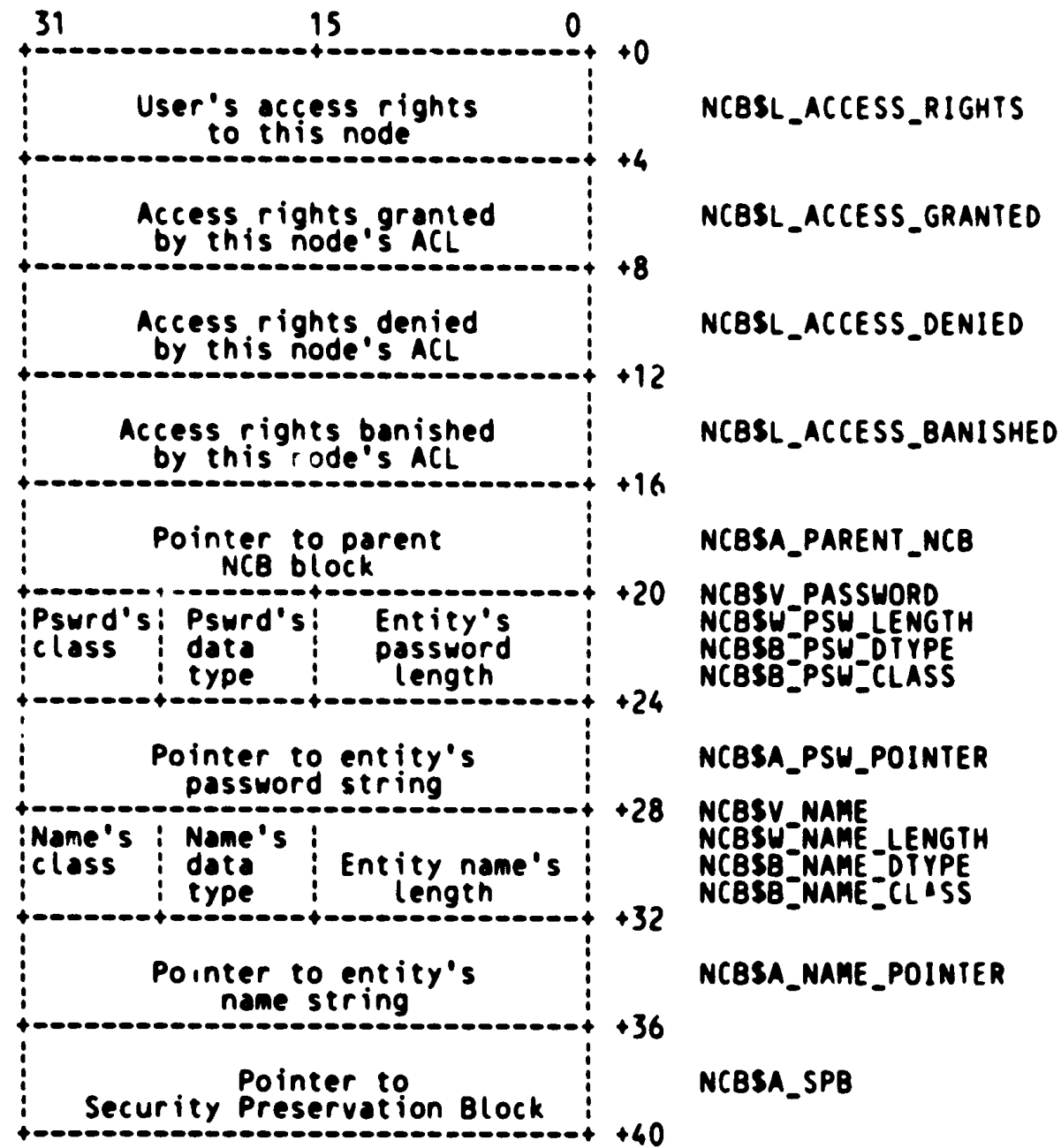
%;

FIELD ECCBSZ_FIELDS =

SET
 ECCBSW_CELL = [0+LCCBSS_BLOCK_LENGTH, 0, 16, 0],
 ECCBSV_LST = [2+LCCBSS_BLOCK_LENGTH, 0, 24, 0],
 ECCBSA_PARENT_LCCB = [5+LCCBSS_BLOCK_LENGTH, 0, 32, 0]

TES;

Node Control Block (NCB)



The LCCB for a node is followed by an NCB. This gives additional information needed for the node.

LITERAL

NCBS_BLOCK_LENGTH = 40+LCCBS_BLOCK_LENGTH;

MACRO

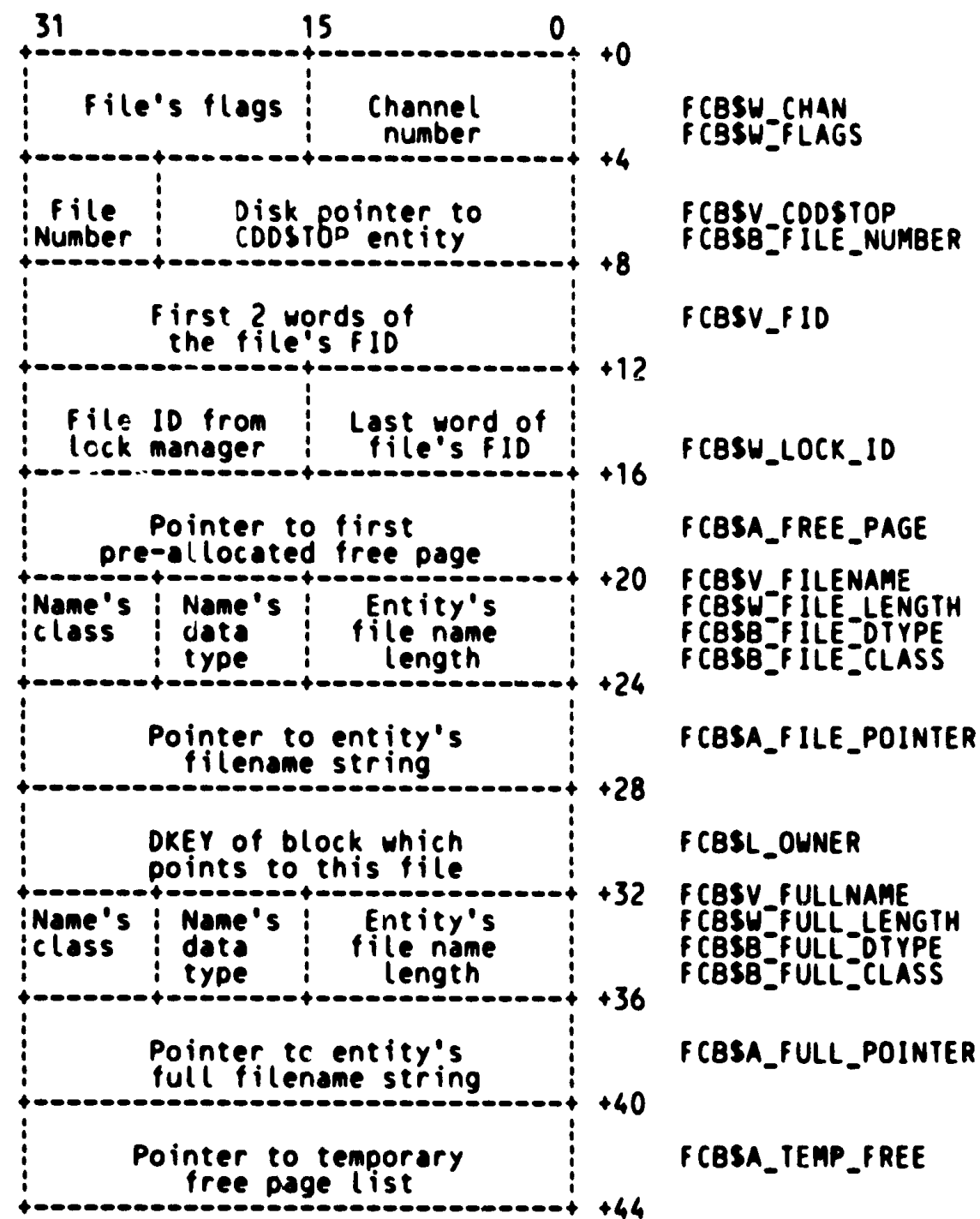
\$NCB = BLOCK[NCB\$\$_BLOCK_LENGTH, BYTE] FIELD (LCCB\$\$_FIELDS,
NCB\$\$_FIELDS)

%;

```
FIELD SET NCB$$_FIELDS =
NCB$$_ACCESS_RIGHTS = [0+LCCB$$_BLOCK_LENGTH, 0, 32, 0],
NCB$$_ACCESS_GRANTED = [4+LCCB$$_BLOCK_LENGTH, 0, 32, 0],
NCB$$_ACCESS_DENIED = [8+LCCB$$_BLOCK_LENGTH, 0, 32, 0],
NCB$$_ACCESS_BANISHED = [12+LCCB$$_BLOCK_LENGTH, 0, 32, 0],
NCB$$_PARENT_NCB = [16+LCCB$$_BLOCK_LENGTH, 0, 32, 0],
NCB$$_PASSWORD = [20+LCCB$$_BLOCK_LENGTH, 0, 0, 0],
NCB$$_PSW_LENGTH = [20+LCCB$$_BLOCK_LENGTH, 0, 16, 0],
NCB$$_PSW_DTYPE = [22+LCCB$$_BLOCK_LENGTH, 0, 8, 0],
NCB$$_PSW_CLASS = [23+LCCB$$_BLOCK_LENGTH, 0, 8, 0],
NCB$$_PSW_POINTER = [24+LCCB$$_BLOCK_LENGTH, 0, 32, 0],
NCB$$_NAME = [28+LCCB$$_BLOCK_LENGTH, 0, 0, 0],
NCB$$_NAME_LENGTH = [28+LCCB$$_BLOCK_LENGTH, 0, 16, 0],
NCB$$_NAME_DTYPE = [30+LCCB$$_BLOCK_LENGTH, 0, 8, 0],
NCB$$_NAME_CLASS = [31+LCCB$$_BLOCK_LENGTH, 0, 8, 0],
NCB$$_NAME_POINTER = [32+LCCB$$_BLOCK_LENGTH, 0, 32, 0],
NCB$$_SPB = [36+LCCB$$_BLOCK_LENGTH, 0, 32, 0]
```

TES;

File Control Block (FCB)



Each active (open) dictionary file has an FCB.

The FCBSL_OWNER field holds the DKEY of the attribute block (FIL) which pointed to the file. The primary dictionary file has key 0, while files that are not currently pointed to have a value of -1

! in the owner field.
!+

LITERAL

FCBSK_FILE_LIMIT = 255,
FCBSS_BLOCK_LENGTH = 44;

MACRO

\$FCB = BLOCK[FCBSS_BLOCK_LENGTH, BYTE] FIELD (FCBSZ_FIELDS)
X;

FIELD FCBSZ_FIELDS =
SET

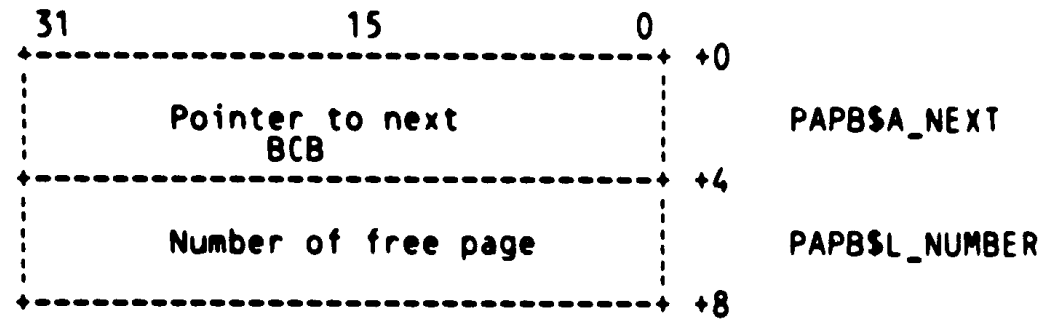
FCBSW_CHAN	=	[0, 0, 16, 0],
FCBSW_FLAGS	=	[2, 0, 16, 0],
FCBSV_READ_ONLY	=	[2, 0, 1, 0],
FCBSV_ROOT	=	[2, 1, 1, 0],
FCBSV_CDDSTOP	=	[4, 0, 24, 0],
FCBSB_FILE_NUMBER	=	[7, 0, 8, 0],
FCBSV_FID	=	[8, 0, 0, 0],
FCBSW_LOCK_ID	=	[14, 0, 16, 0],
FCBSA_FREE_PAGE	=	[16, 0, 32, 0],
FCBSV_FILENAME	=	[20, 0, 0, 0],
FCBSW_FILE_LENGTH	=	[20, 0, 16, 0],
FCBSB_FILE_DTYPE	=	[22, 0, 8, 0],
FCBSB_FILE_CLASS	=	[23, 0, 8, 0],
FCBSA_FILE_POINTER	=	[24, 0, 32, 0],
FCBSL_OWNER	=	[28, 0, 32, 0],
FCBSV_FULLNAME	=	[32, 0, 0, 0],
FCBSW_FULL_LENGTH	=	[32, 0, 16, 0],
FCBSB_FULL_DTYPE	=	[34, 0, 8, 0],
FCBSB_FULL_CLASS	=	[35, 0, 8, 0],
FCBSA_FULL_POINTER	=	[36, 0, 32, 0],
FCBSA_TEMP_FREE	=	[40, 0, 32, 0]

TES;

LITERAL

FCBSM_READ_ONLY = 1^1 - 1^0, ! File can only be opened for read
FCBSM_ROOT = 1^2 - 1^1; ! FCB is the root dictionary file

Pre-Allocated Page Block (PAPB)



Pre-allocated pages are represented by PAPBs linked onto the FCB in which the pages reside. Each PAPB has the page number of the free page it represents.

```

LITERAL
  PAPBS_BLOCK_LENGTH = 8;
  
```

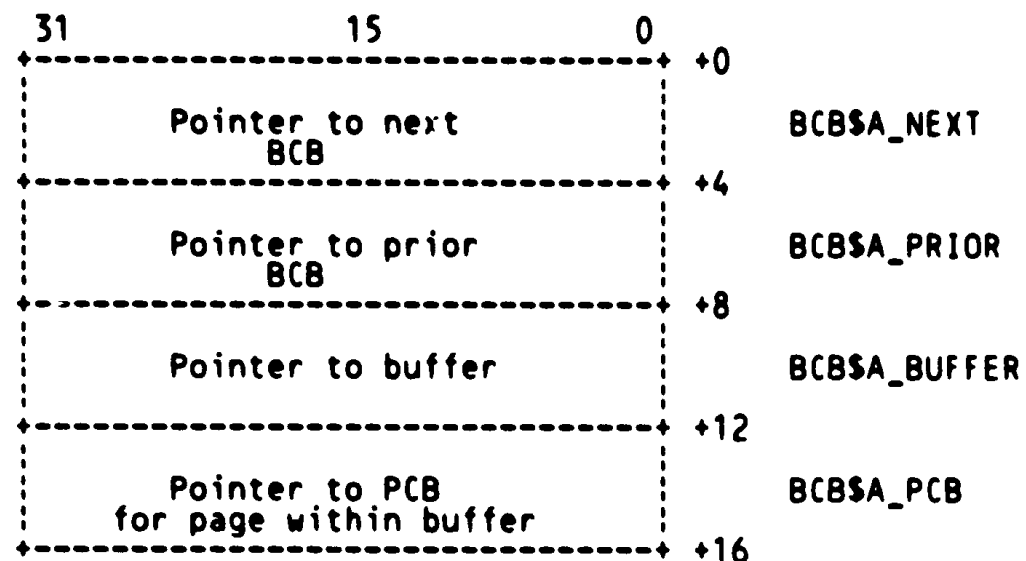
```

MACRO
  $PAPB = BLOCK [PAPBS_BLOCK_LENGTH, BYTE] FIELD (PAPBSZ_FIELDS)
  %;
  
```

```

FIELD  PAPBSZ_FIELDS =
  SET
    PAPBSA_NEXT          = [0, 0, 32, 0],
    PAPBSL_NUMBER       = [4, 0, 32, 0]
  TES;
  
```


Buffer Control Block (BCB)



The buffer control blocks (BCBs) are used to control the pages that are in memory.

The BCBs are linked into a queue. When a buffer is needed, the last buffer in the queue is assigned.

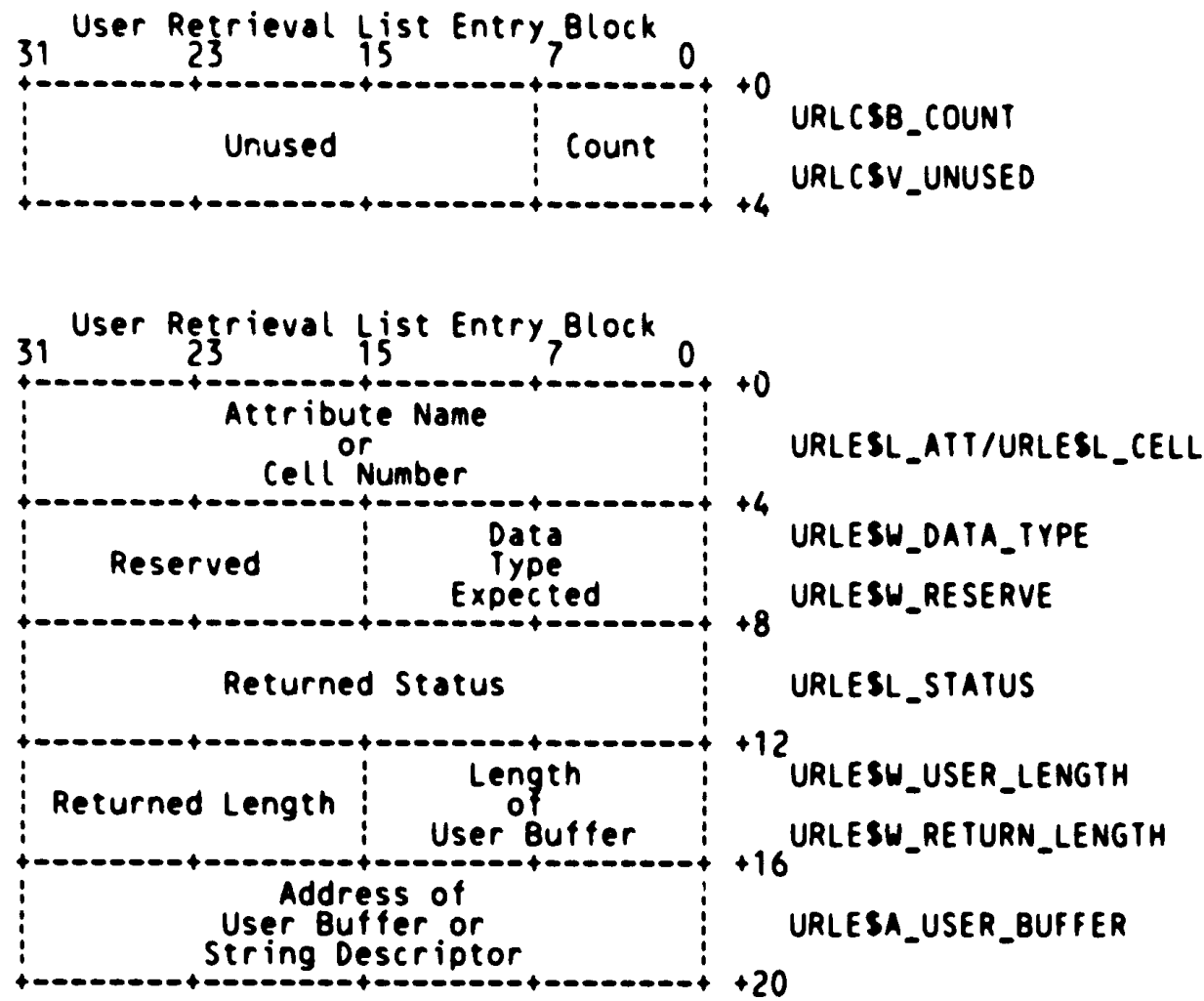
The associated PCB contains information about the page and the buffer. If a page is modified while in the buffer, the page is written back to the work file before the buffer is reused.

```
LITERAL
BCB$K_NUMBER      = 16,      ! Number of in-core buffers
BCB$S_BLOCK_LENGTH = 16;
```

```
MACRO
$BCB = BLOCK[BCB$S_BLOCK_LENGTH, BYTE] FIELD (BCB$Z_FIELDS)
%;
```

```
FIELD BCB$Z_FIELDS =
SET
BCB$A_NEXT      = [0, 0, 32, 0],
BCB$A_PRIOR    = [4, 0, 32, 0],
BCB$A_BUFFER    = [8, 0, 32, 0],
BCB$A_PCB      = [12, 0, 32, 0]
TES;
```

```
LITERAL
BCB$K_BUFFER_LIST = BLOCK[0, BCB$A_NEXT; , BYTE];
```



LITERAL
 URLC\$\$BLOCK_LENGTH = 4;
 URLE\$\$BLOCK_LENGTH = 20;

MACRO
 \$URLC = BLOCK[URLC\$\$BLOCK_LENGTH,BYTE] FIELD (URLC\$Z_FIELDS)
 %;

FIELD URLC\$Z_FIELDS =
 SET
 URLC\$B_COUNT = [0, 0, 8, 0],
 URLC\$V_UNUSED = [1, 0, 24, 0]
 TES;

MACRO
 \$URLE = BLOCK[URLE\$\$BLOCK_LENGTH,BYTE] FIELD (URLE\$Z_FIELDS)
 %;

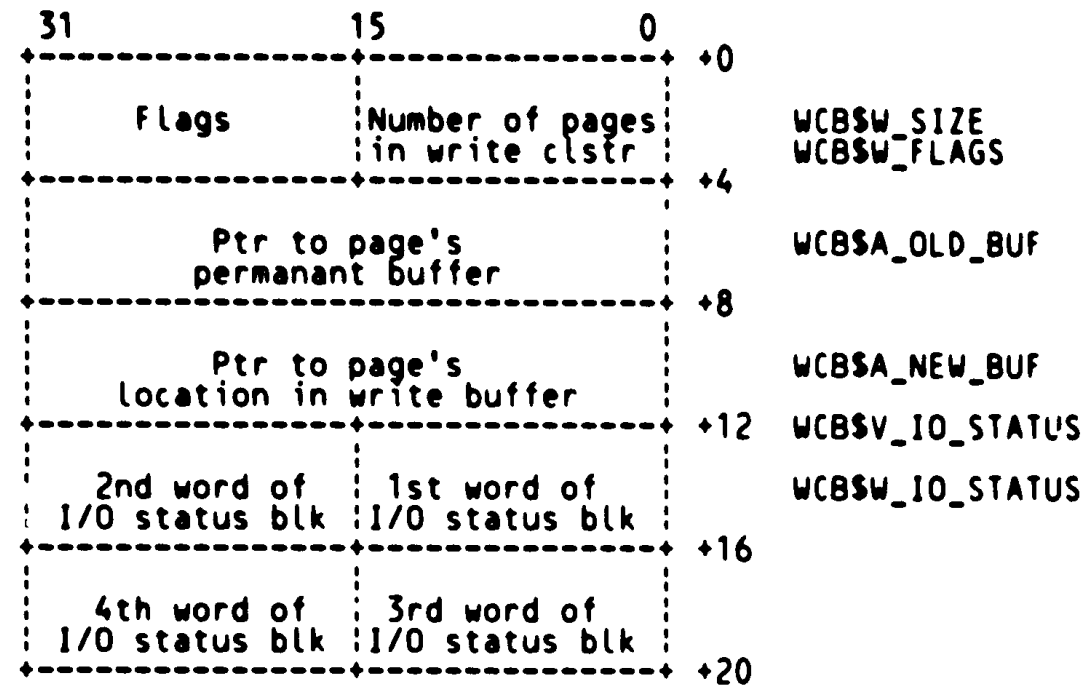
FIELD URLE\$Z_FIELDS =
 SET

```
URLESL_ATT           = [0, 0, 32, 0],
URLESL_CELL          = [0, 0, 32, 0],
URLE$W_DATA_TYPE     = [4, 0, 16, 0],
URLE$W_RESERVE       = [6, 0, 16, 0],
URLESL_STATUS        = [8, 0, 32, 0],
URLE$W_USER_LENGTH   = [12, 0, 16, 0],
URLE$W_RETURN_LENGTH = [14, 0, 16, 0],
URLE$A_USER_BUFFER   = [16, 0, 32, 0]
```

TES:

```
STRUCTURE
URL$BLOCK [I] = (URL$BLOCK + URLE$$_BLOCK_LENGTH +
(I * URLE$$_BLOCK_LENGTH));
```

Write Control Block (WCB)



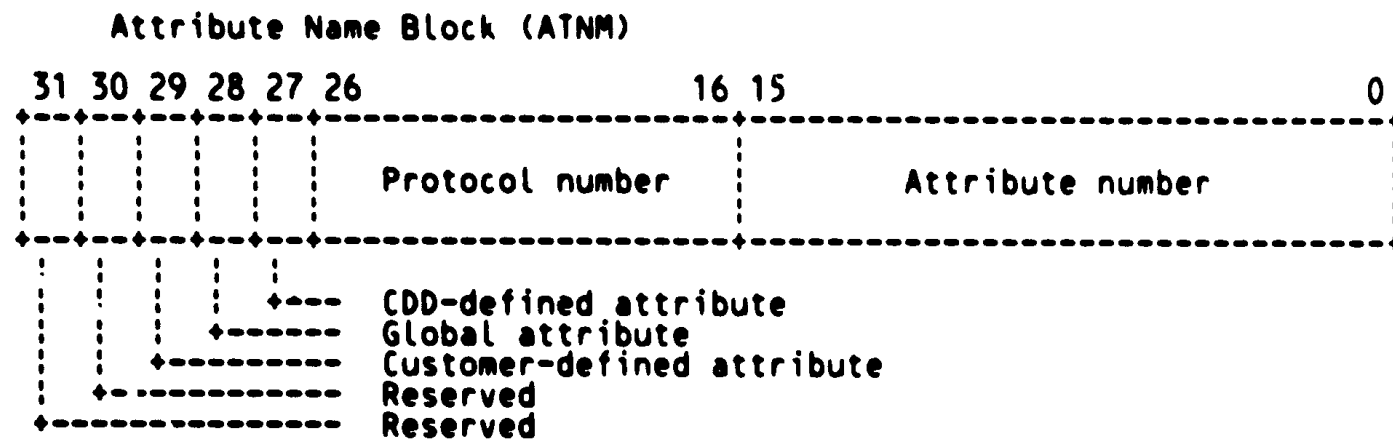
LITERAL
WCBS_BLOCK_LENGTH = 20;

MACRO
\$WCB = BLOCK[WCBS_BLOCK_LENGTH, BYTE] FIELD (WCBSZ_FIELDS)
%;

FIELD WCBSZ_FIELDS =
SET
 WCBSW_SIZE = [0, 0, 16, 0],
 WCBSW_FLAGS = [2, 0, 16, 0],
 WCBSV_WRITABLE = [2, 0, 1, 0], ! Write group to file
 WCBSV_GROUP = [2, 1, 1, 0], ! Group of pages
 WCBSA_OLD_BUF = [4, 0, 32, 0],
 WCBSA_NEW_BUF = [8, 0, 32, 0],
 WCBSV_IO_STATUS = [12, 0, 0, 0],
 WCBSW_IO_STATUS = [12, 0, 16, 0]

TES;

LITERAL
 WCBSM_WRITABLE = 1^1 - 1^0, ! Write group to file
 WCBSM_GROUP = 1^2 - 1^1; ! Group of pages



The attribute name block defines the break down of the attribute name. Bits 0 to 15 contains the number. Bits 16 to 26 contains the protocol. The remaining bits are flags define the type of attribute. Bit 27 on implies a system defined attribute. Bit 28 on implies a global defined attribute. Bit 29 on implies a customer defined attribute.

NOTE: If the high order word of the attribute name block is 0 the block is describing a cell. That is the protocol will equal 0 and the flag bits will equal 0.

LITERAL

```
ATNMSS_BLOCK_LENGTH = 4;
```

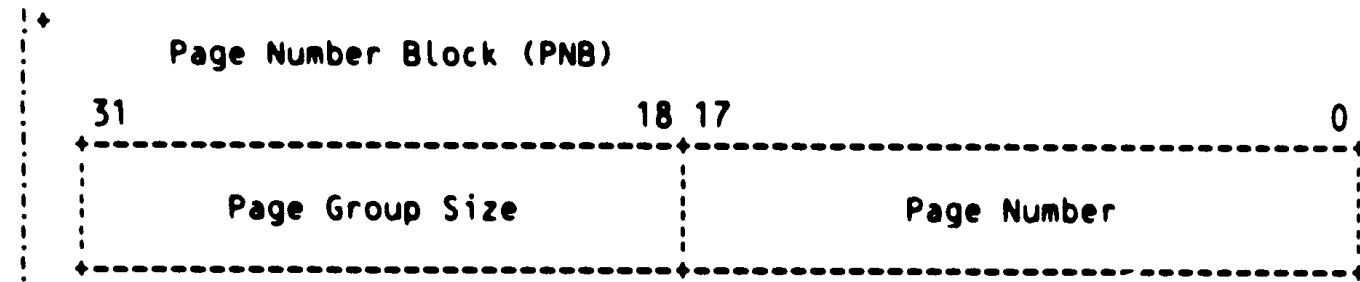
MACRO

```
$ATNM = BLOCK[ATNMSS_BLOCK_LENGTH, BYTE] FIELD (ATNMSZ_FIELDS)
%;
```

FIELD ATNMSZ_FIELDS =

```
SET
  ATNMSW_NUMBER = [0, 0, 16, 0],
  ATNMSW_HIGH_ORD = [0, 16, 16, 0],
  ATNMSV_PROTOCOL = [0, 16, 11, 0],
  ATNMSV_FLAGS = [0, 27, 5, 0],
  ATNMSV_SYSTEM = [0, 27, 1, 0],
  ATNMSV_GLOBAL = [0, 28, 1, 0],
  ATNMSV_CUSTOMER = [0, 29, 1, 0]
```

```
TES;
```



The page number block define the break down of the page number. Bits 0 to 17 is the page number. Bits 18 to 31 de ine the number of consecutive pages that can be read or written at one time.

LITERAL

PNBSS_BLOCK_LENGTH = 4;

MACRO

\$PNB = BLOCK[PNBSS_BLOCK_LENGTH, BYTE] FIELD (PNB\$Z_FIELDS)

%;

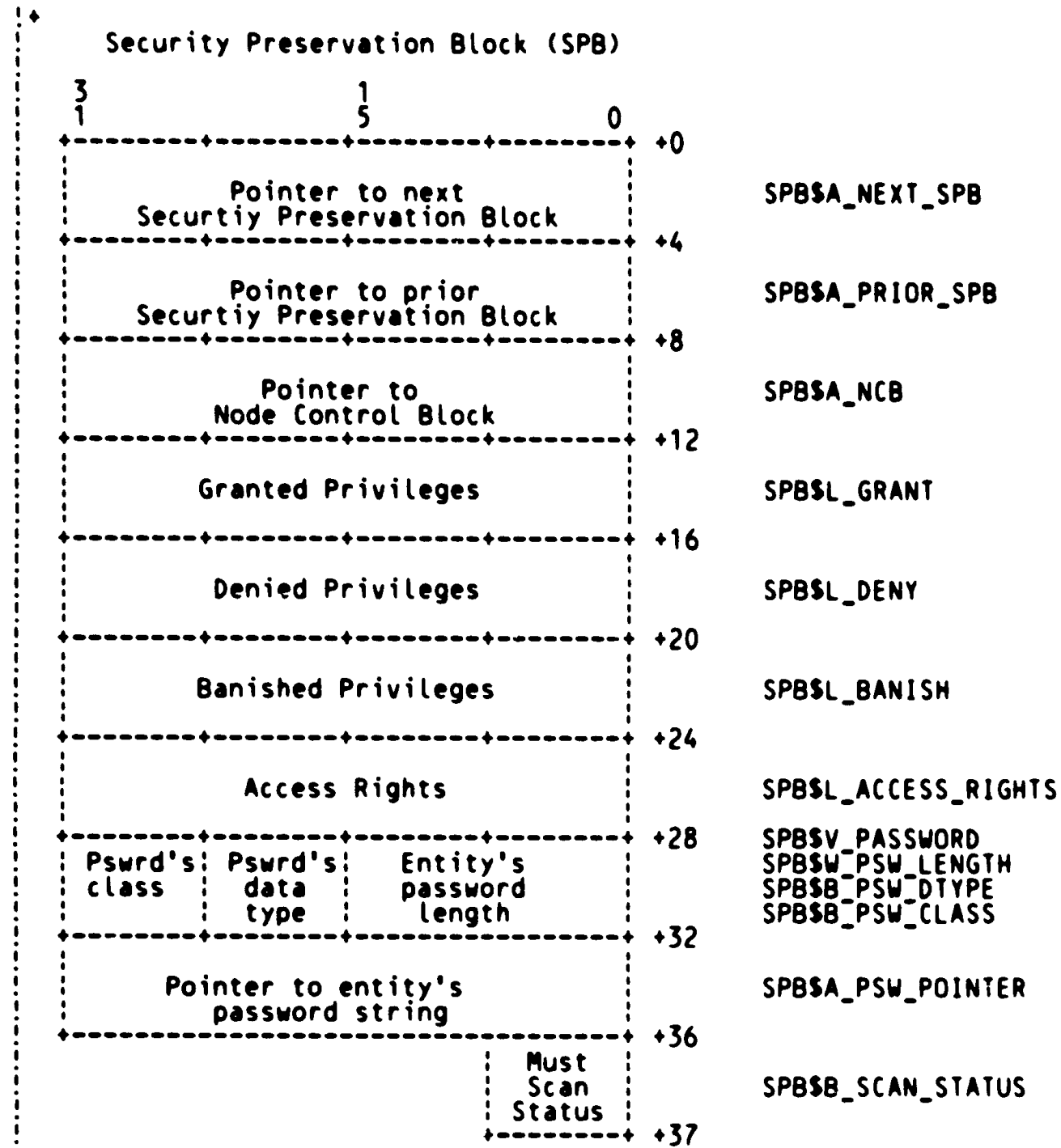
FIELD PNB\$Z_FIELDS =

SET

PNB\$V_PAGE_NUM = [0, 0, 17, 0],

PNB\$V_GROUP_SIZ = [0, 17, 15, 0]

TES;



LITERAL
 SPB\$S_BLOCK_LENGTH = 37;

MACRO

```
SSPB = BLOCK[SPBSS_BLOCK_LENGTH, BYTE] FIELD (SPB$Z_FIELDS)  
%;
```

```
FIELD SPB$Z_FIELDS =  
SET
```

```
SPB$A_NEXT_SPB      = [0, 0, 32, 0],  
SPB$A_PRIOR_SPB    = [4, 0, 32, 0],  
SPB$A_NCB           = [8, 0, 32, 0],  
SPB$L_GRANT         = [12, 0, 32, 0],  
SPB$L_DENY          = [16, 0, 32, 0],  
SPB$L_BANISH        = [20, 0, 32, 0],  
SPB$L_ACCESS_RIGHTS = [24, 0, 32, 0],  
SPB$V_PASSWORD      = [28, 0, 0, 0],  
SPB$W_PSW_LENGTH    = [28, 0, 16, 0],  
SPB$B_PSW_DTYPE     = [30, 0, 8, 0],  
SPB$B_PSW_CLASS     = [31, 0, 8, 0],  
SPB$A_PSW_POINTER   = [32, 0, 32, 0],  
SPB$B_SCAN_STATUS  = [36, 0, 8, 0]
```

```
TES;
```

```
LITERAL
```

```
SPB$K_QUE_HEADER    = BLOCK[0, SPB$A_NEXT_SPB; , BYTE];
```



```

%SBTTL      'System Literal Definitions'

```

```

SYSTEM LITERAL DEFINITIONS

```

```

These literals are only used internally.

```

```

CDD Implementation Version

```

```

LITERAL

```

```

CDD$K_FACILITY      = 43,
CDD$K_LOWEST_VERSION = 201, ! Lowest compatible version
CDD$K_VERSION       = 201; ! Present version

```

```

Boolean literals

```

```

LITERAL

```

```

TRUE  = 1,      ! Boolean TRUE value
FALSE = 0;      ! Boolean FALSE value

```

```

The following literals are used to validate routines' parameter
lists.

```

```

LITERAL

```

```

ARG$K_OPTIONAL      = 1,      ! Parameter is optional
ARG$K_REQUIRED      = 2,      ! Parameter is required
ARG$K_SYNC          = 3,      ! See CDD$SU_VALIDATE documentation
ARG$K_SYNCIF        = 4,      ! See CDD$SU_VALIDATE documentation
ARG$K_MARK           = 0,      ! Mark parameter as missing
ARG$K_DEFAULT       = 1,      ! Use default value
ARG$K_STRING        = 1,      ! Default value is a null string
ARG$K_LONG           = 2,      ! Default value is a longword
ARG$K_WORD          = 3,      ! Default value is a word value
ARG$K_REF            = 0,      ! Parameter passed by reference
ARG$K_VALUE         = 1;      ! Parameter passed by value

```

```

Access Lock Constants

```

```

NOTE:  The order of the LOCK$K_NORMAL lock constants MUST be
the same as the order of the CIB$V_REF_COUNTS lock
ref count fields in the CIB.

```

```

LITERAL
LOCK$K_NULL           = 0,           ! Not locked
LOCK$K_NORMAL         = 1,           ! Base of partially queued locks
  LOCK$K_PRESENCE     = 1,
  LOCK$K_RETRIEVAL    = 2,
  LOCK$K_UPDATE       = 3,
  LOCK$K_DELETE       = 4,
LOCK$K_NORMAL_TYPES   = 4,           ! Number of partially queued locks
!!!
!!! LOCK$K_QUEUED      = 5,           ! Base of fully queued locks
!!! LOCK$K_FULL_RET    = 5,
!!! LOCK$K_FULL_UPD    = 6,
!!! LOCK$K_QUEUED_TYPES = 2,           ! Number of fully queued locks

LOCK$K_XXX_GET        = 1^17-1^16,   ! Flag indicates must establish lock
LOCK$K_XXX_GET_IF     = 1^18-1^17,   ! Get lock if not already present
LOCK$K_OPTIONS        = 1^32-1^16;   ! Option bits for locks
    
```

```

!+
!-
Types of page purge requests
    
```

```

LITERAL
PURGE$K_ALL           = 1,           ! Complete purge
PURGE$K_PRESENCE     = 2,           ! Purge but keep portal page
PURGE$K_RETRIEVAL    = 3,           ! Checkpoint & keep retrieval locks
PURGE$K_UPDATE       = 4,           ! Checkpoint & keep all locks
PURGE$K_SUBTREE      = 100,         ! Purge whole subtree
PURGE$K_CLUSTER      = 101;        ! Purge this cluster only
    
```

```

!+
!-
Types of blocks that can be allocated in a pool.
    
```

```

!+
!-
Types of pools
    
```

```

LITERAL
MEM$K_LOWEST_POOL     = 1,           ! Lowest pool type
MEM$K_CCB_POOL        = 1,           ! Pool for CCB and HCBs
MEM$K_CIB_POOL        = 2,           ! Pool for cluster blocks
MEM$K_HIGHEST_POOL    = 2;          ! Highest pool type
    
```

```

!+
!-
CCB Pool block types
    
```

```

LITERAL
MEM$K_CCB_LOWEST      = 1,           ! Lowest block type in CCB pool
MEM$K_CCB_CCB         = 1,           ! CCB block
    
```

```

MEMSK_CCB_LOW_SLOT      = 2,      ! Lowest block type in free block list
MEMSK_CCB_HCB          = 2,      ! HCB block
MEMSK_CCB_HIGHEST     = 2;      ! Highest block type in CCB pool

```

```

!+
Cluster Pool block types
!-

```

```

LITERAL
MEMSK_CIB_LOWEST       = 1,      ! Lowest block type in cluster pool
MEMSK_CIB_CIB         = 1,      ! CIB block
MEMSK_CIB_LOW_SLOT    = 2,      ! Lowest block type in free block list
MEMSK_CIB_PCB         = 2,      ! PCB,
MEMSK_CIB_LCB        = 3,      ! LCB,
MEMSK_CIB_LCCB       = 4,      ! LCCB,
MEMSK_CIB_ECCB       = 5,      ! ECCB,
MEMSK_CIB_NCB        = 6,      ! NCB,
MEMSK_CIB_SPB        = 7,      ! SPB,
MEMSK_CIB_HIGHEST     = 7;      ! Highest block type in cluster pool

```

```

!+
These flags tell the deletion routine how it is to handle the
following cases:
!-

```

```

DELSK_FAST           says that pointers do not have to be cleaned up,
                    as the block they reside in is going to be deleted.

DELSK_PRESERVE       indicates that a directory node is merely to be
                    emptied, and that its cluster is not to be deleted.

DELSK_SUBDICTIONARY says that sub-files are to have their contents
                    deleted.

```

```

LITERAL
DELSK_FAST           = 1^1 - 1^0,
DELSK_PRESERVE       = 1^2 - 1^1,
DELSK_SUBDICTIONARY = 1^3 - 1^2;

```

```

!+
User Identification Criteria
!-

```

```

LITERAL
CDDSK_ACL_LOWEST     = 1,
CDDSK_ACL_PASSWORD   = 1,      ! PASSWORD
CDDSK_ACL_TERMINAL   = 2,      ! TERMINAL name or class
CDDSK_ACL_UIC        = 3,      ! UIC
CDDSK_ACL_USERNAME   = 4,      ! USERNAME
CDDSK_ACL_HIGHEST    = 4;

```

%SBTTL 'Security Masks'

SECURITY MASKS

CDD security bits

LITERAL

CDD\$K_PROT_C	= 1^1 - 1^0,	! CONTROL access
CDD\$K_PROT_D	= 1^2 - 1^1,	! LOCAL DELETE access
CDD\$K_PROT_G	= 1^3 - 1^2,	! GLOBAL DELETE access
CDD\$K_PROT_H	= 1^4 - 1^3,	! HISTORY list entry creation access
CDD\$K_PROT_P	= 1^5 - 1^4,	! PASS THRU access
CDD\$K_PROT_S	= 1^6 - 1^5,	! SEE (read) access
CDD\$K_PROT_U	= 1^7 - 1^6,	! UPDATE terminal node access
CDD\$K_PROT_X	= 1^8 - 1^7,	! EXTEND directory node access
CDD\$K_PROT_F	= 1^9 - 1^8,	! FORWARDing directory creation allowed

Macro-security values

CDD\$K_PROT_ANY	= 1^9 - 1^0,
CDD\$K_PROT_DELETE	= CDD\$K_PROT_D OR CDD\$K_PROT_G,
CDD\$K_PROT_EXTEND	= CDD\$K_PROT_F OR CDD\$K_PROT_X,
CDD\$K_PROT_UPDATE	= CDD\$K_PROT_C OR CDD\$K_PROT_D OR CDD\$K_PROT_G OR CDD\$K_PROT_H OR CDD\$K_PROT_U OR CDD\$K_PROT_X OR CDD\$K_PROT_F,

Other processor security bits

VAX-11 Datatrieve

CDD\$K_DTR_PROT_E	= 1^17 - 1^16,	! EXTEND file
CDD\$K_DTR_PROT_R	= 1^18 - 1^17,	! READ file
CDD\$K_DTR_PROT_M	= 1^19 - 1^18,	! MODIFY file
CDD\$K_DTR_PROT_W	= 1^20 - 1^19,	! WRITE file

%SBTTL 'User Literal Definitions'

USER LITERAL DEFINITIONS

These symbols are needed by users of the program interface.

System Defined Attribute Names

LITERAL

CDD\$K_SYSNAM_FLAGS = 1^28 OR 1^27 OR 0^16; ! Global/System-defined/Protocol=0

LITERAL

CDD\$K_FIRST_SYSNAM	= 1 OR CDD\$K_SYSNAM_FLAGS,	! Lowest system defined attribute name value
CDD\$K_FILE	= 1 OR CDD\$K_SYSNAM_FLAGS,	! Node's file name
CDD\$K_HISTORY	= 2 OR CDD\$K_SYSNAM_FLAGS,	! History list head
CDD\$K_NAME	= 3 OR CDD\$K_SYSNAM_FLAGS,	! Node's name
CDD\$K_PROTOCOL	= 5 OR CDD\$K_SYSNAM_FLAGS,	! Node's protocol name
CDD\$K_TYPE	= 6 OR CDD\$K_SYSNAM_FLAGS,	! Type of object pointed to by location code
CDD\$K_PATHNAME	= 7 OR CDD\$K_SYSNAM_FLAGS,	! Node's complete pathname
CDD\$K_SHORT_PATHNAME	= 8 OR CDD\$K_SYSNAM_FLAGS,	! Node's path to CDD\$DEFAULT directory
CDD\$K_ORDER	= 9 OR CDD\$K_SYSNAM_FLAGS,	! Directory's order
CDD\$K_LAST_SYSNAM	= 9 OR CDD\$K_SYSNAM_FLAGS;	! Highest system defined attribute name value

Attribute and Entity Types

LITERAL

CDD\$K_FIRST_TYPE	= 1.
CDD\$K_ENTITY	= 1.
CDD\$K_ENTITY_LIST	= 2.
CDD\$K_NULL	= 3.
CDD\$K_NUMERIC	= 4.
CDD\$K_STRING	= 5.
CDD\$K_STRING_LIST	= 6.
CDD\$K_DIRECTORY	= 7.
CDD\$K_TERMINAL	= 8.
CDD\$K_LAST_TYPE	= 8.

User's entity purge options

LITERAL

CDD\$K_ALL	= 1^1 - 1^0.
CDD\$K_ABORT	= 1^2 - 1^1.
CDD\$K_CHECKPOINT	= 1^3 - 1^2.

!+
:-
User's node creation options

LITERAL
CDD\$K_NOHISTORY = 1^1 - 1^0, ! Doesn't want history list cluster
CDD\$K_NOACL = 1^2 - 1^1, ! Don't create default ACL entry
CDD\$K_CREATE = 1^3 - 1^2, ! Create dictionary file if needed
CDD\$K_FIRST = 1^4 - 1^3, ! Insert as first node
CDD\$K_LAST = 1^5 - 1^4, ! Insert as last node

!+
:-
User's node deletion options

LITERAL
CDD\$K_CHECK = 1^1 - 1^0, ! Fail if directory has children
CDD\$K_SUBDICTIONARY = 1^2 - 1^1, ! Delete contents of subdictionaries

!+
:-
Values of the CDD\$K_ORDER attribute

LITERAL
CDD\$K_SORTED = 1, ! Directory is sorted
CDD\$K_NONSORTED = 2, ! Directory is not sorted

```
%SBTTL      'LINKAGE DEFINITIONS'
```

```
LINKAGE DEFINITIONS
```

```
CDDCALL
```

```
This linkage uses the CALLG/CALLS linkage convention, except  
that it allows for one global register to be used in  
parameter passing.
```

```
R11 -- used to pass the user's context pointer.
```

```
LINKAGE  
CDDCALL = CALL : GLOBAL (USER_CONTEXT = 11);
```

```
SYS_JSB
```

```
This linkage provides us with a general JSB routine linkage.
```

```
LINKAGE  
SYS_JSB = JSB;
```

```

ZSBTTL      'MACRO DEFINITIONS'

```

```

MACRO DEFINITIONS

```

```

$ACTIVE
$INACTIVE

```

These macros declare that we have started, and finished, respectively, a CDD transaction. They abort the transaction if another transaction is in progress.

```

MACRO
$ACTIVE =
  BEGIN
    EXTERNAL
      CDD$GB_INUSE:          BYTE;

    EXTERNAL LITERAL
      CDD$_NOTASTREE;

    BUILTIN
      TESTBITSS;

    IF TESTBITSS (CDD$GB_INUSE) THEN
      SIGNAL (CDD$_NOTASTREE);
  END
Z,

$INACTIVE =
  BEGIN
    EXTERNAL
      CDD$GB_INUSE:          BYTE;

    CDD$GB_INUSE = FALSE;
  END
Z,

```

```

$BITCLEAR

```

This macro checks to see if any bit in a mask is set in the target area. If not, it returns TRUE.

```

$BITCLEAR(target, mask) =
  (target AND mask) EQLU 0
Z,

```


♦
 SBITSET

This macro checks to see if any bit in a mask is set in the target area. If so, it returns TRUE.

```
SBITSET(target, mask) =
(target AND mask) NEQU 0
%,
```

♦
 \$DONE_TRANS

This macro is used to terminate a transaction.

Call:

```
$DONE_TRANS [(dsc1 [, dsc2] ...)]
```

Where:

dsci ::= the names of dynamic descriptors which are to have their strings returned to the string pool.

```
$DONE_TRANS (dsc1) =
BEGIN
EXTERNAL ROUTINE
CDD$$$SN_DONE_TRANS      : CDDCALL      NOVALUE;

CDD$$$SN_DONE_TRANS (dsc1
%IF NOT %NULL(%REMAINING) %THEN , %REMAINING %FI );
.USER_CONTEXT[CCBSL_STATUS]
END
%,
```

♦
 \$FIND_ENTITY

This macro returns the virtual address of the LCCB associated with a location code. It also checks to make certain that the cluster is locked as requested, and that the cluster's node allows the requested security access.

Call:

```
lccb-block.wa.v = $FIND_ENTITY (valid-arg.ra.v ,
( CHECK      { RETRIEVAL      READ
              { UPDATE        } ) , { MODIFY    } );
              DELETE          DELETE
```

ANY

Where:

valid-arg ::= the address of the calling routine's validated argument list.
The argument list must have the following format:

valid-arg[0] ::= address of longword holding context #
valid-arg[1] ::= address of descriptor holding path name, or zero.

The first set of keywords names the desired lock state of the entity's cluster.

The last set of keywords names the intended access to the cluster.

```

SFIND_ENTITY (valid_arg, locking, security) =
  BEGIN
    EXTERNAL ROUTINE
      CDDSSN_FIND_ENTITY      : CDDCALL;

    CDDSSN_FIND_ENTITY (..valid_arg[1],
      SFIND_XXX_LOCKING (%REMOVE(locking)),
      %NAME('CDDSK_PROT_', security))
  END
%,

SFIND_XXX_LOCKING (class, type) =
  %IF %IDENTICAL (class, %QUOTE LOCK) %THEN
    LOCKSK_XXX_GET OR
  %ELSE
    %IF %IDENTICAL (class, %QUOTE LOCKIF) %THEN
      LOCKSK_XXX_GET_IF OR
    %ELSE
      %IF NOT %IDENTICAL (class, %QUOTE CHECK) %THEN
        %ERROR ('Invalid locking keyword: ', class)
      %FI
    %FI
  %FI
  %NAME ('LOCKSK_', type)
%,

```

SFIND_NODE

This macro returns the virtual address of the NCB associated with a path name or location code. It also checks to make certain that the cluster is locked as requested, and that the target node allows the requested security access.

Call:

```
ncb-block.wa.v = $FIND NODE (valid-arg.ra.v ,
  LOCK      RETRIEVAL  READ
  ( { LOCKIF } { UPDATE  } ) , { MODIFY } );
  CHECK     DELETE     DELETE
  ANY
```

Where:

valid-arg ::= the address of the calling routine's validated argument list.
The argument list must have the following format:

```
valid-arg[0] ::= address of longword holding
               context #
valid-arg[1] ::= address of descriptor holding
               path name, or zero.
valid-arg[2] ::= address of longword holding
               location code, or zero.
```

The first set of keywords names the desired lock state of the node's cluster.

The last set of keywords names the intended access to the cluster.

```
$FIND NODE (valid_arg, locking, security) =
BEGIN
  EXTERNAL ROUTINE
  CDD$$SN_FIND_NODE      : CDDCALL;

  CDD$$SN_FIND NODE (valid_arg,
    $FIND_XXX_LOCKING (%REMOVE(locking)),
    %NAME?'CDD$K_PROT_', security))
END
%.
```

\$FIND_PARENT

This macro returns the virtual address of the NCB associated with a location code. It also checks to make certain that the cluster is locked as requested, and that the cluster's node allows the requested security access.

Call:

```
status.wlc.v = $FIND_PARENT (valid-arg.ra.v ,
  LOCK      RETRIEVAL  READ
  ( { LOCKIF } { UPDATE  } ) , { MODIFY } , ncb-block.wa.r,
  CHECK     DELETE     DELETE
  ANY
  name.wt.ds);
```

Where:

valid-arg ::= the address of the calling routine's validated argument list.

The argument list must have the following format:

```

valid-arg[0] ::= address of longword holding
                context #
valid-arg[1] ::= address of descriptor holding
                path name, or zero.
valid-arg[2] ::= address of longword holding
                location code, or zero.

```

The first set of keywords names the desired lock state of the target cluster.

The last set of keywords names the intended access to the cluster.

```

SFIND_PARENT (valid_arg, locking, security, ncb_block, name) =
BEGIN
  EXTERNAL ROUTINE
    CDDSSN_FIND_PARENT      : CDDCALL;

  CDDSSN_FIND_PARENT (valid_arg,
    SFIND_XXX_LOCKING (%REMOVE(locking)),
    %NAME('CDDSK_PROT_', security), ncb_block, name)
END
%,

```

\$INIT_DSC

This macro is used to initialize dynamic string descriptors.

Call:

```
$INIT_DSC (dsc1 [, dsc2] ...)
```

```

$INIT_DSC[DSC_NAM] =
BEGIN
  DSC_NAM[DSC$B_DTYPE] = DSC$K_DTYPE_T;
  DSC_NAM[DSC$B_CLASS] = DSC$K_CLASS_D;
  DSC_NAM[DSC$W_LENGTH] = 0;
  DSC_NAM[DSC$A_POINTER] = 0;
END
%,

```

\$IO_SYNC (ef, iosb)

This macro waits for the event flag (ef) to be set and for the I/O status parameter (iosb) to be filled in (non-zero).

ef must be the target event flag number.

iosb must be the address of the I/O status block. This must be defined as a VECTOR[,WORD] VOLATILE structure.

```

$IO_SYNC (ef, iosb) =
  DO
    BEGIN
      LOCAL
        STATUS:          LONG;

      STATUS = $WAITFR (efn = ef);
      IF NOT .STATUS THEN
        SIGNAL_STOP (.STATUS);
      STATUS = $CLREF (efn = ef);
      IF NOT .STATUS THEN
        SIGNAL_STOP (.STATUS);
    END WHILE .iosb[0] EQLU 0
  %

```

SMARK_PAGE (page-block)

This macro marks a page as modified. Such a page must be written back to the dictionary file when it is purged from the staging buffers.

```

SMARK_PAGE (page_block) =
  page_block[P[BSV_MODIFIED] = TRUE
  %

```

\$PARAMETERS (arg1 [, arg2] ...)

This macro is used to build the control vector for the parameter list validate routine (CDD\$\$U_VALIDATE).

There is one entry (arg1, arg2, etc) for every formal parameter in the routine. Each entry has the following format:

```

( REQUIRED [, [n] [, [REF : VALUE] :
  { SYNC : $SYNCIF }, n :
  OPTIONAL [, MARK : , DEFAULT [, STRING : , LONG : , WORD] ] )

```

```

$PARAMETERS[] =
  UPLIT (%LENGTH, $PARM_PROCESS (%REMOVE(%REMAINING)))

```

```

%,
SPARM_PROCESS[arg] =
  BYTE ($SPARM_DECIDE (%REMOVE(arg)))
%,
SPARM_DECIDE(status)[] =
  %IF %IDENTICAL (status, %QUOTE REQUIRED) %THEN
    ARG$K_REQUIRED, 0,
    %IF %NULL (%REMAINING) %THEN
      ARG$K_REF, 0
    %ELSE
      $SPARM_REQUIRED (%REMAINING)
    %FI
  %ELSE
    %IF %IDENTICAL (status, %QUOTE SYNC) %THEN
      ARG$K_SYNC, 0, 0, %REMAINING
    %ELSE
      %IF %IDENTICAL (status, %QUOTE SYNCIF) %THEN
        ARG$K_SYNCIF, 0, 0, %REMAINING
      %ELSE
        %IF %IDENTICAL (status, %QUOTE OPTIONAL) %THEN
          ARG$K_OPTIONAL, $SPARM_OPTIONAL (%REMAINING), 0
        %ELSE
          %ERROR ('Invalid parameter status: ', status)
          0, 0, 0, 0
        %FI
      %FI
    %FI
  %FI
%,
SPARM_REQUIRED(number, pass) =
  %IF %NULL (pass) %THEN
    0
  %ELSE
    %NAME ('ARG$K_', pass)
  %FI
  %IF %NULL (number) %THEN
    0
  %ELSE
    , number
  %FI
%,
SPARM_OPTIONAL(action)[] =
  %IF %IDENTICAL (action, %QUOTE MARK) %THEN
    ARG$K_MARK, 0
  %ELSE
    %IF %IDENTICAL (action, %QUOTE DEFAULT) %THEN
      ARG$K_DEFAULT, %NAME ('ARG$K_', %REMAINING)
    %ELSE
      %ERROR ('Invalid action keyword: ', action)
      0, 0
    %FI
  %FI

```

%,

\$PRESENT

This macro checks to see if a parameter is present (non-zero) in a list.

```
$PRESENT (name) =
.name NEQA 0
```

%,

\$RECOVERY (RESET ; ENABLE ; DISABLE)

This macro determines the ability of the exit handler to perform a purge of the cache if the task aborts.

\$RECOVERY (DISABLE)

declares that the internal data structures or external disk structure is in an indeterminant state and cannot be recovered by the exit handler.

\$RECOVERY (ENABLE)

declares that the data structure manipulation is completed.

\$RECOVERY (RESET)

specifies that the data structures are in a recoverable state.

Note that these can be nested. Recovery is only possible if the recovery counter is zero (reset).

\$RECOVERY (option) =

BEGIN

EXTERNAL

CDD\$GW_RECOVERY: WORD;

%IF %IDENTICAL (option, %QUOTE DISABLE) %THEN

CDD\$GW_RECOVERY = .CDD\$GW_RECOVERY + 1;

%ELSE

%IF %IDENTICAL (option, %QUOTE ENABLE) %THEN

CDD\$GW_RECOVERY = .CDD\$GW_RECOVERY - 1;

%ELSE

%IF %IDENTICAL (option, %QUOTE RESET) %THEN

CDD\$GW_RECOVERY = 0;

%ELSE

%ERROR ('Illegal recovery option: ', option)

%FI

%FI

```

%      END %FI
%,

```

```

$RELEASE_LOCK

```

```

This macro calls the CDD$$SN_RELEASE routine to release one
or more locks.

```

```

Call:

```

```

    $RELEASE_LOCK ( { RETRIEVAL
                     UPDATE } , ncb-block1 ...);
                     DELETE

```

```

$RELEASE_LOCK (locking)[] =

```

```

BEGIN

```

```

    EXTERNAL ROUTINE
    CDD$$SN_RELEASE          : CDDCALL          NOVALUE;

```

```

    CDD$$SN_RELEASE (%NAME ('LOCK$K_', locking), %REMAINING)

```

```

END

```

```

%,

```

```

$SIGNAL_SEVERE (error)

```

```

This routine signals a severe error.

```

```

$SIGNAL_SEVERE (ERROR) =

```

```

    SIGNAL (ERROR OR ST$$K_SEVERE)

```

```

%,

```

```

$STATIC_DSC

```

```

This macro is used to initialize static string descriptors.

```

```

Call:

```

```

    $STATIC_DSC (dsc1 [, dsc2] ...)

```

```

Where:

```

```

    dsci := name ! ( name , source )

```

```

Source is the name of another descriptor. The named descriptor
is initialized to point to the same string as source.

```



```

%,
$STRING_PTR_SETUP[PAIR] =
  $STRING_PTR_INIT (%REMOVE(PAIR))
%,
$STRING_PTR_INIT (STR_NAME, STR_VAL) =
  STR_NAME[DSC$A_POINTER] = UPLIT BYTE(%REMOVE(STR_VAL));
%,

```

↑

STEXTC

This macro is used to define counted strings. The first byte of such strings is a count of the number of characters in the string.

Each string is defined to be a VECTOR[,BYTE] structure, with the 0 element being the character count, and the actual string starting at STRING[1].

STEXTC ((name1, 'str1') [, (name2, 'str2')] ...);

↓

```

$STEXTC[PAIR] =
  $STEXTC_STR(%REMOVE(PAIR))
%,
$STEXTC_STR(NAME, TSTR) =
  BEGIN
    NAME = UPLIT BYTE (%CHARCOUNT(%REMOVE(TSTR)), %REMOVE(TSTR)) :
    VECTOR[%CHARCOUNT(%REMOVE(TSTR))+1, BYTE];
%,

```

↑

\$VALIDATE (cntl, [arg-list])

This macro generates a call to the general transaction setup routine.

cntl -- the address of the control vector for CDD\$\$U_VALIDATE.

arg-list-- is optional. If present, it is the address of the vector which is to receive the verified argument list from CDD\$\$U_VALIDATE.

↓

```

$VALIDATE(cntl, arg_list) =
  BEGIN
    EXTERNAL ROUTINE
      CDD$$SN_START_TRANS      : CDDCALL NOVALUE;

```

BUILTIN
AP;

```
%IF %NULL (arg_list) %THEN  
  CDDSSN_START_TRANS (.AP, cntl)  
%ELSE  
  CDDSSN_START_TRANS (.AP, cntl, arg_list)  
%FI
```

%;
END

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VAX/VMS V4.0

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The image displays a grid of 100 computer terminal screens, arranged in 10 rows and 10 columns. Each screen shows a different output from a VAX/VMS system, including:

- System status and configuration screens.
- Code listings and program outputs.
- Graphical displays, such as bar charts and histograms.
- Specific application screens with titles like:
 - COO
 - COOSHR MAP
 - W.C. Johnson
 - COOEXC2 LIS
 - COOLIB LIS
 - COOEXC LIS
 - WRITEBOOT LIS
 - COOLIB B32