## PART 5

## DOS/BATCH DEVICE DRIVERS

## PART 5

## CHAPTER 1

## USING DEVICE DRIVERS OUTSIDE DOS/BATCH

Subroutines to handle I/O transfers between a PDP-Il and each of its peripheral devices are developed as required for use within the Disk Operating System DOS/BATCH. These subroutines are made available within an I/O Utilities Package for the benefit of PDP-ll users who have configurations unable to support DOS/BATCH or who wish to run programs outside DOS/BATCH control.

All the subroutines associated with one peripheral device form an entity known as a device driver. This part provides a general description of a driver and shows how it can be used in a stand-alone environment. The unique properties of each driver are discussed in separate documents, which are made available as part of the Device Driver Package. The I/O utilities package for any system is determined by the peripherals of that system. Thus, the full documentation for a particular package consists of this document and applicable supplements.

## PART 5

## CHAPTER 2

## DRIVER FORMAT

### 2.1 STRUCTURE

The basic principle of all drivers under the DOS/BATCH Monitor is that they must present a common interface to the routines using them in order to provide deviceindependent operation. The subroutines are structured to meet this end. Moreover, a driver can be loaded anywhere in memory under Monitor Control. Its code is always position-independent (PIC). ${ }^{1}$

A detailed description of a driver is found in Chapter 5-4. This section describes driver interfaces.

### 2.1.1 Driver Interface Table

The first section of each driver is a table which contains, in a standard format, information on the nature and capabilities of the device it represents and entry points to each of its subroutines. The calling program can use this table as required, regardless of the device being called. See Section 5-4.1 for a detailed description of the table.

### 2.1.2 Setup Routines

Each driver is expected to handle its device under the PDP-11 interrupt system. When called by a program, therefore, a driver subroutine merely initiates the required action by setting the device hardware registers appropriately. It returns to the calling program by a standard subroutine exit.

The main setup routines prepares for a data transfer to or from the device, using parameters supplied by the calling program. Normally, blocks of data are moved at each transfer. The driver returns control to the program only when the whole block has been transferred or when it is unable to continue because there is no more data available.

[^0]The driver can also contain subroutines by which the calling program can request (1) start-up or shut-down action, such as leader or trailer functions for a paper tape punch, or (2) some special function provided by the device hardware (or a software simulation of that for some similar device), e.g., rewind of a magnetic tape or DECtape.

### 2.1.3 Interrupt Servicing

The driver routine to service device interrupts is particularly dependent upon the device hardware provisions for controlling transfers. In general, the driver determines the cause of the interrupt and checks whether the last action was performed correctly or was prevented by some error condition. If more device action is needed to satisfy the program request, the driver again initiates that action and takes a normal interrupt exit. If the program request has been fully met, control is returned to the program at an address supplied at the time of the request.

### 2.1.4 Error Handling

Device errors can be handled in two ways. There are some errors for which recovery can be programmed; the driver, if appropriate, attempts this itself (as in the case of parity or timing failure on a bulk-storage device) or recalls the program with the error condition flagged (as at the end of a physical paper tape). Other errors normally require external action, perhaps by an operator. The driver calls a DOS/ BATCH error handler via an IOT call with supporting information on the processor stack.

### 2.2 INTERFACE TO THE DRIVER

### 2.2.1 Control Interface

The principal link between a calling program and any driver subroutine is the first word of the driver table (link word). In order to provide the control parameters for a device operation, the calling program prepares a list in a standardized form and places a pointer to the list in the link word. The called driver uses the pointer to access the parameters. The driver can place return status information (if needed) in the list area via the link word. The first word of the driver table can also act as a busy indicator; if it is $\varnothing$, the driver is not currently performing a task, but if it contains a listpointer, the driver can be assumed to be busy. Since most drivers support only one job at a time, the link word state is significant.

### 2.2.2 Interrupt Interface

Although the driver expects to use the interrupt system, it does not itself ensure that its interrupt vector in the memory area below $4 \varnothing \varnothing_{8}$ has been set up correctly; the Monitor takes care of this. However, the driver table contains the information required to initialize the appropriate vector.

## PART 5

## CHAPTER 3

## STAND-ALONE USE

Because each driver is designed for operation within the device-independent framework of the Monitor, it can be similarly used in other applications. Since the easiest way to use the driver is to assemble it with the program that requires it, this method will be described first. Other possible methods will be discussed later.

### 3.1 DRIVER ASSEMBLED WITH PROGRAM

### 3.1.1 Setting Interrupt Vector

As noted in Section 5-2.2.2, the calling program must initialize the device transfer vectọ within memory locations $\emptyset$ - 377 . The address of the driver's interrupt entry point can be identified on the source listing by the symbolic name which appears as the contents of the Driver Table Byte, DRIVER+5. The priority level at which the driver expects to process the interrupt is at byte DRIVER+6. For a program which can reference position-dependent code, the setup sequence,might be:

| MOV | \#DVRINT,VECTOR | ;SET INT. ADDRESS |
| :--- | :--- | :--- |
| MOVB | DRIVER+6,VECTOR +2 | ;SET PRIORITY |
| CLRB | VECTOR +3 | ;CLEAR UPPER STATUS BYTE |

where the Driver Table Byte (at DRIVER+5) shows the follwing instruction:
.BYTE DVRINT-DRIVER
If the program must be position-independent, it can take advantage of the fact that the Interrupt Entry address is stored as an offset from the start of the driver, as illustrated above. In this case, a sample sequence might be:

| MOV | PC, R1 | ;GET DRIVER START |
| :--- | :--- | :--- |
| ADD | \#DRIVER-.,R1 |  |
| MOV | \#VECTOR,R2 | ;...\& VECTOR ADDRESSED |
| CLR | @R2 | ;SET INT. ADDRESS |
| MOVB | $5(R 1), @ R 2$ |  |
| ADD | R1,(R2)+ |  |
| CLR | @R2 |  |
| MOVB | $6(R 1), @ R 2$ |  |

### 3.1.2 Parameter Table for Driver Call

For any call to the driver the program must provide a list of control arguments mentioned in Section 5-2.2.1. This list must adhere to the following format: ${ }^{1}$
[SPECIAL FUNCTION POINTER] ${ }^{2}$
[BLOCK NO.] ${ }^{3}$
STARTING MEMORY ADDRESS FOR TRANSFER
NO. OF WORDS to be transferred (2's complement) STATUS CONTROL showing in Bits:

```
\emptyset-2 Function (octally 2=WRITE, 4=READ) 4
8-1\varnothing Unit (if Device can consist of several units, e.g., DECtape)
11 Direction for DECtape travel ( }\varnothing=\mathrm{ Forward)
```

ADDRESS for RETURN ON COMPLETION [RESERVED FOR DRIVER USE] ${ }^{5}$

The list can be assembled in the required format since its content will not vary. The driver can return information in this area; this will not corrupt the program data.

On the other hand, most programs will probably use the same list area for several tasks or even for different drivers. In this case, the program must contain the necessary routine to set up the list for each task before making the driver call. The driver may refer to the list again when it is recalled by an interrupt or when returning information to the calling program. Therefore, the list must not be changed until any driver has completed a function requested; for concurrent operations, different list areas must be provided.

[^1]
### 3.1.3 Calling the Driver

To enable the driver to access the parameter list, the program must set the first word of the driver to an address six bytes less than that of the word containing the MEMORY START ADDRESS. It can then directly call the required driver subroutine by a normal JSR PC,xxxx call, where $x \times x x$ is the address of the driver subroutine.

As an example, the following position-independent code might appear in a program which wishes to read Blocks \#1øø-103 backward from DECtape unit 3 into a buffer starting at address BUFFER.

|  | MOV | PC, Rø | ; GET TABLE ADDRESS |
| :---: | :---: | :---: | :---: |
|  | ADD | \#TABLE+12-. , R $\varnothing$ |  |
|  | MOV | PC,@Rø | ; GET AND STORE... |
|  | ADD | \#RETURN-.,@Rø | ; ...RETURN ADDRESS |
|  | MOV | \#54ø4,-(Rø) | ; SET READ REV. UNIT 3 |
|  | MOV | \#-1ø24.,-(Rø) | ; 4 BLOCKS REQUIRED |
|  | MOV | PC, - (Rø) | ; GET AND STORE |
| - | ADD | \#BUFFER-., @Rø | ; . . BUFFER ADDRESS |
|  | MOV | \#1ø3,-(Rø) | ; START BLOCK |
|  | CMP | -(Rø), -(Rø) | ;SUBTRACT 4 FROM POINTER |
|  | MOV | $\mathrm{R} \varnothing, \mathrm{DT}$ | ; SET DRIVER LINK |
|  | JSR | PC, DT.TFR | ; GO TO TRANSFER ROUTINE |
| WAIT : | . |  | ; RETURNS HERE WHEN |
|  | $\cdot$ |  | ;...TRANSFER UNDER WAY ; RETURNS HERE WHEN |
|  | . |  | ; ...TRANSFER COMPLETE |
| TABLE: | . WORD $\varnothing$ |  | ;LIST AREA SET |
|  | -WORD $\varnothing$ |  | ;...by Above SEQUENCE |
|  | . WORD $\varnothing$ |  |  |
|  | .WORD $\varnothing$ |  |  |
|  | .WORD $\varnothing$ |  |  |
|  | .WORD $\varnothing$ |  |  |

### 3.1.4 User Registers

During its setup operations for the function requested, the driver assumes that Processor Registers $\emptyset-5$ are available for its use. If their contents are of value, the program must save them before the driver is called.

While servicing intermediate interrupts, the driver may need to save or restore its registers. It expects to have two subroutines available for the purpose (provided by the Monitor). It accesses them via addresses in memory locations $44{ }_{8}$ and $46_{8}$.

```
MOV @#44,-(SP) ;OR MOV @#46,-(SP)
JSR
R5,@(SP)+
```

The driver must also ensure that the start addresses are set into the correct locations ( $44_{8}$ and $46_{8}$ ).

At its final interrupt, the driver saves the contents of Registers $\varnothing$ - 5 before returning control to the calling program completion return.

### 3.1.5 Returns From Driver

As shown in the example in Section 5-3.1.3, the driver returns control to the calling program immediately after the JSR as soon as it has set the device in motion. The program can wait or carry out alternative operations until the driver signals completion by returning at the address specified (i.e., RETURN above). Prior to this, the program must not attempt to access the data being read in, nor refill a buffer being wxitten out.

The program routine beginning at address RETURN varies according to the device being used. In general, the driver has given control to the routine for one of two reasons; either the function has been satisfactorily performed, or it cannot be carried out due to some hardware failure with which the driver is unable to cope, though the program may be able to do so. In the latter case, the driver uses the STATUS word in the program list to show the cause:
$\begin{array}{ll}\text { Bit } 15=1 \quad \begin{array}{l}\text { indicates that a device or timing failure occurred } \\ \text { and the driver has not been able to overcome this, }\end{array} \\ & \text { perhaps after several attempts. }\end{array}$

Bit $14=1$ shows that the end of the available data has been reached.

The driver places in $R \varnothing$ the contents of its first word as a pointer to the parameter table (see Section 5-3.1.2).

Possibly, the driver has transferred only some of the data requested. In this case, it shows in the RESERVED word of the program list a negative count of the words not transferred in addition to setting Bit 14 of the STATUS word. As mentioned in the note in Section 5-3.1.2, this applies only to non-bulk storage devices. The drivers for DECtape or disks ${ }^{1}$ always endeavor to complete the full transfer, even beyond a parity failure, or they take more drastic action (see Section 5-3.1.6).

[^2]It is thus the responsibility of the program RETURN routine to check the information supplied by the driver in order to verify that the transfer was satisfactory and to handle the error situations appropriately.

In addition, the routine must contain a sequence to take care of the Processor Stack, Registers, etc. As noted earlier, the driver takes the completion return address after an interrupt and saves Registers $\varnothing-5$ on the stack above the Interrupt Return Address and Status. The program routine should, therefore, contain some sequence to restore the processor to its state prior to such interrupt, e.g., using the same Restore subroutine illustrated earlier:

| MOV | @\#46,-(SP) | ;CALL REGISTER RESTORE |
| :--- | :--- | :--- |
| JSR | R5,@(SP)+ |  |
| $\cdot$ |  |  |
| $\cdot$ |  |  |
| • |  |  |
| RTI |  |  |
|  |  |  |
|  |  |  |

All hardware errors other than those noted in the previous section cannot normally be overcome by the program or by the driver on its behalf. Some of these could be due to an operator fault, such as not turning on a paper tape reader or not setting the correct unit number on a DECtape transport. Once the operator has rectified the problem, the program could continue. Other errors, however, require hardware repair or even software repair, e.g., if the program asks for Block $2 \varnothing \varnothing \varnothing$ on a device having a maximum of $1 \varnothing \varnothing \varnothing$. In general, all these errors result in the driver placing identifying information on the processor stack and calling IOT to produce a trap through location $348^{\circ}$

Under DOS/BATCH, the Monitor provides a routine to print a teleprinter message when this occurs. In a stand-alone environment, the program using the driver must itself contain the routine to handle the trap (unless the user wishes to modify the driver error exits before assembly). The handler format depends upon the program. The following format takes advantage of the information supplied by the driver:

|  | (SP) : | Return Address |  |
| :---: | :---: | :---: | :---: |
| 2 | (SP) : | Return Status | Stored by IOT call |
| 4 | (SP) : | Error No. Code | Generally unique to driver |
| 5 | (SP) : | Error Type Code: | $1=$ Recoverable after Operator Action <br> 3 = No recovery |
| 6 | (SP) : | Additional <br> Information | Such as content of Driver, Control Register, Driver Identity, etc. |

As a rule, the driver expects a return following the IOT call in the case of recoverable errors but contains no provision for an IOT call following a return from irrecoverable errors.

### 3.1.7 General Comment

The source language of each driver has been written for use with DOS/BATCH and contains some code which is not accepted by the Paper Tape Software PAL-llR, in particular, .TITLE, .GLOBL, and Conditional Assembly directives. Such statements should be deleted before the source is used. Similarly, an entry in the driver table gives the device name as .RAD5ø 'DT' to obtain a specifically packed format used internally by DOS/BATCH. If the user wishes to keep the name, for instance, for identification purposes as discussed in Section 5-3.3, .RAD5 $\emptyset$ might easily be changed to .ASCII without detrimental effect, or it might be replaced with . WORD ø.

### 3.2 DRIVERS ASSEMBLED SEPARATELY

Rather than assemble the driver with every program requiring its availability, the user may wish to hold it in binary form and attach it to the program only when loaded. The only requirement is that the start address of the driver should be known or be determinable by the program.

The example in Section 5-3.1.2 showed that the Interrupt Servicing routine can be accessed through an offset stored in the Driver Table. The same technique can be used to call the setup routines, as these also have corresponding offsets in the Table, as follows:

| DRIVER +7 | Open $^{1}$ |
| ---: | :--- |
| +10 | Transfer $^{+11}$ |
| +12 | Close $^{1}$ |
|  | Special Functions ${ }^{1}$ |

The problem is the start address. There is the obvious solution of assembling the driver at a fixed location so that each program using it can immediately reference the location chosen. This ceases to be convenient when the program has to avoid the area occupied by the driver. A more general method is to relocate the driver as dictated by the program using it, thus taking advantage of the position-independent nature of the driver. The Absolute Loader, described in the Paper Tape Software Handbook (DEC-11-XPTSA-A-D), provides the capability to continue a load from the point at which it ended. Using this facility to enter the driver immediately following the program, the program might contain the following code to call the subroutine to perform the transfer illustrated in Section 5-3.1.3.

[^3]|  | MOV | PC; RI | ;GET DRIVER START ADDRESS |
| :---: | :---: | :---: | :---: |
|  | ADD | \#PRGEND-., Rl |  |
|  | MOV | PC, Rø | ; GET TABLE ADDRESS |
|  | ADD | \#TABLE+12-.,Rø | ; AND SET UP AS SHOWN |
|  | . |  | ; ...IN SECTION 5-3.1.3 |
|  | - |  |  |
|  | CMP | -(Rø), -(Rø) | ;FINAL POINTER ADJUSTMENT |
|  | MOV | R $\varnothing$, @R1 | ; STORE IN DRIVER LINK |
|  | CLR | -(SP) | ;GET BYTE SHOWING... |
|  | MOVB | $1 \varnothing$ (R1), @SP | ; ...tRANSFER OFFSET |
|  | ADD | (SP) +,R1 | ; COMPUTE ADDRESS |
|  | JSR | PC,@RI | ;GO TO DRIVER |
|  | - |  |  |
|  | - |  |  |
| PGREND: | - |  |  |
|  | .END |  |  |

This technique can be extended to cover situations in which several drivers are used by the same program, provided that it takes account of the size of each driver (known because of prior assembly) and that the drivers themselves are always loaded in the same order.

For example, to access the second driver, the above sequence would be modified to:

|  | MOV | PC, Rl | ; GET DRIVER 1 ADDRESS |
| :---: | :---: | :---: | :---: |
|  | ADD | \#PRGEND-.,R1 |  |
|  | ADD | \#DVR1S2,R1 | ;SET TO DRIVER 2 |
|  | - |  | * |
|  | - |  | , |
| DVR1SZ=n |  |  |  |
| PRGEND: |  |  |  |
|  | .END |  |  |

An alternative method may be to use the MACRO Assembler in association with the Linker program LINK, both of which are available through the DECUS Library. The start address of each driver is identified as a global. Any calling programs need merely include a corresponding .GLOBAL statement, e.g., .GLOBL DT.

### 3.3 DEVICE-INDEPENDENT USAGE

The drivers are assigned for use in a device-independent environment, i.e., one in which a calling program need not know in advance which driver has been associated with a table for a particular run. One application of this type might be to allow line printer output to be diverted to some other output medium because the line printer is not currently available. Another might be to provide a general program to analyze data samples although these on one occasion might come directly
from an Analog-to-Digital converter and on another be stored on a DECtape because the sampling rate was too high to allow immediate evaluation.

Programs of this type should be written to use all the facilities that any one device might offer, but not necessarily for each device. For instance, the program should ask for start-up procedures because it may sometime use a paper tape punch which provides them, even though it may normally use DECtape which does not. As noted in paragraph 5-2.2.1, the driver table contains an indication of its capabilities to handle this situation. The program can thus examine the appropriate item before calling the driver to perform some action. As an example, the code to request start-up procedures might be (assuming Rø already set to List Address) :

| MOV | \#DVRADD,RI | ; GET DRIVER ADDRESS |
| :---: | :---: | :---: |
| TSTB | 2 (RI) | ;BIT 7 SHOWS |
| BPL | NOOPEN | ; ...OPEN ROUTINE PRESENT |
| MOV | $\mathrm{R} \varnothing$,@R1 | ; STORE TABLE ADDRESS |
| CLRB | -(SP) | ; BUILD ADDRESS |
| MOVB | 7(RI) , @SP | ; ...OF THIS ROUTINE |
| ADD | (SP) +, RI |  |
| JSR | PC,@R1 | ; ...AND GO TO IT |
|  |  | ;FOLLOWED POSSIBLY BY |
|  |  | ; WAIT AND COMPLETION |
|  |  | ; PROCESSING |
|  |  | ;RETURN TO COMMON OPERATION |

NOOPEN:

Similarly, the indicators show whether the device is capable of performing input or output, or both; whether it can handle ASCII or binary data; whether it is a bulk storage device capable of supporting a directory structure or is a terminaltype device requiring special treatment. Other table entries show the device name as identification and the number of words the device might normally expect to transfer at a time (in l6-word units). All of the information can be readily examined by the calling program, thus enabling the use of a common call sequence for any I/O operation, as illustrated in the example on the following page.

|  | MOV | \#DVRADR,R5 | ; SET DRIVER START |
| :---: | :---: | :---: | :---: |
|  | JSR | R5,IOSUB | ; CALl SET UP SUB |
|  | BR | WAIT | ; SKIP table following on return |
|  | . WORD | $1 \varnothing$ | ;TRANSFER REQUIRED |
|  | .WORD | 183 | ; BLOCK NO. |
|  | -WORD | BUFFER | ; BUFFER ADDRESS |
|  | . WORD | -256. | ; WORD COUNT |
|  | - WORD | $4 \varnothing 4$ | ; READ FROM UNIT 1 |
|  | .WORD | RETURN | ; EXIT ON COMPLETION |
|  | . WORD | $\emptyset$ | ; RESERVED |
| WAIT: |  |  | ; CONTINUE HERE... |
|  | - |  |  |
|  | - | - |  |
| IOSUB: | MOV | @SP,Rø | ;PICK UP DRIVER ADDR |
|  | MOV | R5, R1 | ; SET UP POINTER TO LIST |
|  | TST | (RI) + | ; BUMP TO COLLECT CONTENT |
|  | . |  | ; ROUTINE CHECKS ON DEVICE |
|  | . | . | ; ...CAPABILITY USING Rl |
|  | - |  | ;...TO ACCESS LIST AND |
|  | - |  | ; . . R $\varnothing$ THE DRIVER TABLE |
|  | - |  | ; IF O.K... |
|  | MOV | @R1,R1 | ;GET ROUTINE OFFSET |
|  | ADD | R $\varnothing, \mathrm{RI}$ |  |
|  | CLR | -(SP) | ;USE IT TO BUILD |
|  | MOVB | @RI,@SP | ;...ENTRY POINT |
|  | ADD | $R \varnothing$,@SP |  |
|  | JSR | PC, @ (SP) + | ; CALL DRIVER |
|  | RTS | R5 | ;EXIT TO CALLER |

The calling program, or a subroutine of the type just illustrated, may take advantage of a feature mentioned earlier; the fact that when a driver is in use, its first word is non-zero. The driver itself does not clear this word except in special cases shown in the description for the driver concerned. If the program itself always ensures that the first word of the driver is set to zero between driver tasks, then this word forms a suitable driver-busy flag. Under DOS, the program parameter list is extended to allow additional words to provide linkage between lists as a queue in which the list indicated in the driver's first word is the first link.

The preceding paragraphs indicate possible ways of incorporating the available drivers into the type of environment for which they were designed. The user should carefully read the more detailed description of the driver structure in Chapter 5-4, and the individual driver specifications before determining the final form of his program.

A word of warning is appropriate here. Although most drivers set up an operation and then wait for an interrupt to produce a completion state, there are some cases in which the driver can finish its required task without an interrupt, e.g., "opening" a paper tape reader involves only a check on its status. Moreover, where "Special Functions" are concerned, the driver routine may determine from the code specified that the function is not applicable to its device, and therefore, have nothing to do. In such cases, the driver clears the intermediate return address from the processor stack and immediately takes the completion return. Special problems can arise, however, if the driver concerned is servicing several tasks, any of which can cause a queue for the driver's services under DOS/BATCH. To overcome these problems, the driver expects to be able to refer to flags outside the scope of the list so far described. This can mean that a program using such a driver may also need to extend the list range to cover such possibilities. Particular care should be exercised in such cases.

## PART 5

## CHAPTER 4

## I/O DRIVERS WITHIN THE DOS/BATCH OPERATING SYSTEM

The principal function of an $I / O$ driver is to satisfy a Monitor processing routine's requirement for the transfer of a block of data in a standard format to or from the device it services. This involves setting up the device hardware registers to cause the transfer and gaining control under the interrupt scheme of PDP-11, making allowance for peculiar device characters (e.g., conversion to or from ASCII if some special code is used).

The I/O driver must also include routines for handing device start-up or shutdown such as punching leader or trailer, and for making available to the user certain special features of the device, such as rewind of magtape.

### 4.1 DRIVER STRUCTURE

In order to provide a common interface to the Monitor, all drivers must begin with a table of identifying information as follows:

DVR:

| BUSY FLAG (initially $\varnothing$ ) <br> FACILITY INDICATOR (expanded below) |
| :--- |
| Offset to <br> Interrupt Routine* |
| Offset to <br> OPEN Routine* |
| Standard Buffer Size <br> in l6-word Units. |
| Offset to <br> CLOSE Routine* |
| SpacePriority for <br> Interrupt Service* |
| Offset to <br> Transfer Routine* |

The table should be extended as follows if the device is file-structured:


The driver routines that set up the transfer and control under the interrupt follow the table.

Bits in the Facility Indicator Word define the device for Monitor reference:

SPECIAL STRUCTURES
GENERAL STRUCTURE

*Multi-unit System
type devices (i.e., RK disk).

### 4.2 MONITOR CALLING

When a Monitor I/O processing routine needs to call the driver, it first sets up the parameters for the driver operation in relevant words of the appropriate DDB ${ }^{1}$, as illustrated in the following table.

[^4]DDB:

| SPECIAL FUNCTION CODE |
| :---: |
| DEVICE BLOCK NUMBER |
| MEMORY BLOCK ADDRESS |
| WORD COUNT (2'S COMPLEMENT) |
| TRANSFER FUNCTIONS (EXPanded below) |
| COMPLETION RETURN ADDRESS |
| (DRIVER WORD-COUNT RETURN) Set to Zero |

The relevant content of the Transfer Function word is as follows:


Provided that the Facility Indicator in the Driver Table described above shows that the driver is able to satisfy the request, according to the direction and mode and the service required, the Monitor routine places in Register 1 the relative By'te address of the entry in the Driver Table containing the offset to the routine to be used (e.g., for the Transfer routine, this would be $1 \varnothing$ ). The Monitor routine then calls the Driver Queue Manager, using a JSR PC, S.CDB instruction.

The Driver Queue Manager refers to the Busy Flag (Word $\varnothing$ of the driver table) to assure that the driver is free to accept the request. If the Busy Flag contains $\varnothing$, the Queue Manager inserts the address of the DDB from Register $\varnothing$ and jumps to the start of the routine in the driver using Register 1 content to evaluate the address required. If the driver is already occupied, the new request is placed in a queue linking the appropriate DDB's for datasets waiting for the driver's services. It is taken from the queue when the driver completes its current task. (This is done by a recall to the Queue Manager from the routine just serviced, using JSR PC,S.CDQ).

On entry to the Driver Routine, therefore, the address following the Monitor routine call remains as the "top" element of the processor stack. It can be used by the driver in order to make an immediate return to the Monitor (having initiated the function requested), using RTS PC. It should also be noted that the Monitor routine saves register contents if it needs them after the device action. The driver may thus freely use the registers for its own operations.

When the driver has completly satisfied the Monitor request, it should return control to the Monitor using the address set into the DDB. On such return, Register $\emptyset$ must be set to contain the address of the DDB just serviced and since the return will normally follow an interrupt, Registers $\varnothing-5$ at the interrupt must be stored on top of the stack.

### 4.3 DRIVER ROUTINES

### 4.3.1 TRANSFER

The sole purpose of the TRANSFER routine is to set the device in motion. The information needed to load the hardware registers is available in the DDB, whose address is contained in the first word of the driver. Conversion of the stored values is the function of the routine. It must also enable the interrupt; however, it need not set the interrupt vectors as these are preset by the Monitor when the driver is brought into core. After the TRANSFER'routine has activated the device, the routine returns to the calling processor by an RTS PC instruction.

### 4.3.2 Interrupt Servicing

The form of this routine depends upon the nature of the device. In most drivers it falls into two parts, one for handling the termination of a normal transfer and the other to deal with reported error conditions.

For devices which are word or byte-oriented, the routine must provide for individual word or byte transfers, with appropriate treatment of certain characters (e.g., TAB or Null) and for their conversion between ASCII or binary and any special device coding scheme, until either the word count in the DDB is satisfied or an error prevents this. On these devices, the most likely case for such error is the detection of the end of the physical medium; the treatment: for the error varies according to whether the device is providing input or accepting output. The calling program usually needs to take action in the former case and the driver should merely indicate the error by returning the unexpired portion of the word count in DDB Word 7 on exit to the Monitor. Output End of Data requires operator
action. To obtain this, the driver should call the Error Diagnostic Print routine within the Monitor by:

| MOV | DEVNAM, $-(S P)$ | ;SHOW DEVICE NAME |
| :--- | :--- | :--- |
| MOV | $\# 4 \varnothing 2,-(S P)$ | ;SHOW DEVICE NOT READY |
| IOT |  | ;CALL ERROR DIAGNOSTIC PRINT ROUTINE |

On the assumption that the operator will reset the device for further output and request continuation, the driver must follow the above sequence with a Branch or Jump to resume the transfer.

Normal transfer handling on blocked devices (or those like RFll Disk which are treated as such) is simpler since the hardware takes care of individual words or bytes and the interrupt only occurs on completion.

Errors that indicate definite hardware malfunctions must generate diagnostic messages to the Operator, The only recourse is to start the program over, after the malfunction has been corrected.

There are some errors which the driver can attempt to overcome by restarting the transfer. Device parity failure on input is a common example. If one or more retries are unsuccessful, the driver should normally allow programmed recovery and indicate the error by Bit 15 of DDB word 5. Nevertheless, because the program may try to process the data despite the error, the driver should attempt to transfer the whole block requested if this has not already been effected. The remaining forms of errors must be processed according to the type of recovery deemed desirable.

Whether the routine uses processor registers for its operation depends on considerations of the core space saved against the time taken to save the user's contents. However, on completion (or error return to the Monitor), the calling routine expects the top of the stack to contain the contents of Registers $\varnothing$-5 and Register $\varnothing$ to be set to the address of the DDB just serviced. The driver must, therefore, provide for this.

### 4.3.3 OPEN

This routine need be provided only for those devices that require some hardware initialization. It should not normally appear in drivers for devices used in a file oriented manner. The presence of the routine must be indicated by Bit 7 in the driver table Facility Indicator.

The OPEN routine may vary according to the transfer direction of the device. For output devices, the probable action required is the transmission of appropriate data, e.g., CR/LF at a keyboard terminal, form-feed at a printer, or null characters as punched leader code, and for this a return interrupt is expected. The OPEN routine should then be somewhat similar to the TRANSFER routine in that it sets the device going and makes an interim return via RTS PC, waiting until completion of the whole transmission before taking the final return address in the DDB.

An input OPEN may consist of just a check on the readiness of the device to provide data when requested. In this case, the desired function can be effected without any interrupt wait. The routine should, therefore, take the completion return immediately. Nevertheless, it must ensure that the saved PC value on top of the stack from the call to S.CDB is appropriately removed before exit. In the case of drivers which can service only one dataset at a time (i.e., Bit $\emptyset$ of their Facility Pattern word is set to $\varnothing$ ) and can never be queued, a TST (SP)+ instruction can effect this. However, a multiuser driver must allow for the possibility that it may be recalled to perform some new task waiting in a queue. This condition exists if the byte at DDB-3 is non-zero. In this case, the driver must simulate the interrupt expected by the completion process. This is accomplished by inserting a PS word on the stack above the return address supplied by the JSR of the Open request. A possible sequence for the interrupt simulation is illustrated below.

|  | MOV | DRIVER,Rø | ;PICK UP DDB ADDRESS |
| :---: | :---: | :---: | :---: |
|  | MOV | (SP) +, R5 | ;SAVE INTERIM RETURN |
|  | TSTB | -3 (Rø) | ; COME FROM QUEUE? |
|  | BEQ | EXIT |  |
|  | MOV | @\#177776, - (SP) | ; IF SO, STORE STATUS |
|  | MOV | R5,-(SP) | ; ...\& RETURN |
|  | SUB | \#14, SP | ; DUMMY SAVE REGS |
| EXIT: | JMP | @10(Rø) |  |

### 4.3.4 CLOSE

The CLOSE routine is like the OPEN routine, in that it should provide for the possibility of some form of hardware shut-down, such as the punching of trailer code and that it is not necessary for file-structured devices. Moreover, it is likely to be a requirement for output devices only. If it is provided, Driver Table Facility Indicator (Bit 6) must be set.

Again, the probable form is initialization of the hardware action required, with immediate return via RTS PC and eventual completion return via the DDB-stored address.

### 4.3.5 SPECIAL

This routine may be included if either the device itself contains the hardware to perform some special function or there is a need for software simulation of each hardware on other devices, e.g., tape rewind; it should not be provided otherwise. Its presence must be indicated by Bit 5 of the Facility Indicator.

The function itself is stoxed by the Monitor as a code in the DDB. When called, the driver routine must determine whether such function is appropriate in its case. If not, the completion return should be taken immediately with prior stack clearance, as discussed under OPEN. For a recognized function, the necessary routine must be provided. Its exit method depends upon the necessity for an interrupt wait.

### 4.4 DRIVERS FOR TERMINALS

The rate of input from terminal devices normally reflects the typing skill of the operator. For both input and output, the amount of data to be transferred on each occasion may be a varying length, i.e., a line rather than a block of standard size. Furthermore, echoing input may conflict with interrupting output. As a result, drivers for such devices demand special treatment.

Normal output operation, i.e., .WRITE by the program, is handled by the Monitor Processor. On recognizing that the device being used is a terminal, as shown by Bit 8 of the facility indicator, this routine always causes a driver transfer at the end of the user line, even though the internal buffer has not been filled. The driver, however, is given the whole of a standard buffer, padded as necessary with nulls. Provided the driver can ignore these, the effect is the suppression of trailing nulls. Input control remains the driver's responsibility since overcoming the rate problem requires circular buffering within the driver. This circular buffering feature allows the user type-ahead facilities. A subsequent input request may then be satisfied by data already in core. If the data is sufficient to fill the Monitor buffer, the driver awaits the next request before further transfer. If this is insufficient, the driver should operate as any other device and use subsequent interrupts to satisfy the Monitor's requests. Since the driver must stop any transfer at the end of a line in normal operation, in order to allow the Monitor to continue, the driver must simulate the filling of the buffer by null padding. If the user requests .TRAN's which are not line oriented, the buffer size varies from the standard and the driver assumes the program requires a complete buffer before return.

## PART 5

## CHAPTER 5

## SAMPLE LINE PRINTER DRIVER LISTING

The following is a sample listing of a DOS/BATCH Device Driver. The actual driver is the LPll Line Printer Driver (for device name LP:).






$$
5-26
$$

| 1 | 040322 | $\begin{aligned} & 120127 \\ & 000615 \end{aligned}$ | LP．143： | CMPB | R1，\＃15 | ； | CARRIAGETRETURN（15）？ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 2 | 0.06326 | 003016 |  | $86 T$ | LP．I14 | 1 | NO，ABOVE |
| 3 | 000330 | 001014 |  | BNE | LP．I15 | 1 | NO，BELOW |
| 4 | 000352 | $\begin{aligned} & 005767 \\ & 177464 \end{aligned}$ |  | TST | GVPRNT | 1 | PRINT THE CARRIAGE－RETURN ？ |
| 5 | 040336 | 001021 |  | BNE | LP． 116 | 1 | YES |
| 6 | 000340 | $\begin{aligned} & 046703 \\ & 177452 \end{aligned}$ |  | mov | LP．SLZ，R3 | 1 | R3－（ PRINTERIS WIDTH） |
| 7 | 000344 | 064403 |  | NEG | R3 | 1 |  |
| 8 |  |  |  | －IFOF | LSIIXSPREAD |  |  |
| 9 |  |  |  | TST | LP．FLG | 1 | ELONGATIUN ENABLED？ |
| 14 |  |  |  | BEQ | LP．IXX | $f$ | NO |
| 11 |  |  |  | ASR | R3 | ， | halve printeris wIdTh |
| 12 |  |  |  | MOV | R3．LP．FLG | 1 | REFINITIALIZE THE FLAG |
| 13 |  |  |  | －ENOC |  |  |  |
| 14 | 400346 |  | LP．IXX： |  |  |  |  |
| 15 | 10346 | 000732 |  | 日R | LP． 106 | 1 | SUPPRESS CARRIAGE－RETURN |
| 16 | 00390 |  | LP．114： | －IFDF | LSIIKSPREAO |  |  |
| 17 |  |  |  | TST | LP．FLG |  |  |
| 18 |  |  |  | BE゙Q | LP．IYY |  |  |
| 19 |  |  |  | CMPB | R1，${ }^{\text {R16 }}$ |  |  |
| 20 |  |  |  | BEQ | LP．IVA |  |  |
| 21 |  |  | LP．IYY： |  |  |  |  |
| 22 |  |  |  | －ENOC |  |  |  |
| 23 | 000350 | 120127 |  | CMPB | R1，\＃22 |  |  |
|  |  | 000022 |  |  |  |  |  |
| 24 | 406354 | 001616 |  | BNE | LP． 117 | 1 | NO |
| 25 | 00356 | 012701 |  | MOV | \＃SKIP2，R1 | 1 | SUBSTITUTE APPROPRIATE CHAR |
|  |  | 000012 |  |  |  |  |  |
| 26 | 00302 | $\begin{aligned} & 120127 \\ & 000012 \end{aligned}$ | LP．I15： | CMPG | R1，\＃12 | 1 | LINEFEED（12）？ |
| 27 | 00306 | 002411 |  | BLT | LP．I17 | ； | NO，BELOW |
| 28 | 00370 | 001404 |  | BEQ | LP．I16 | 1 | YE゙S |
| 29 | 06372 | 120127 |  | CMPB | R1， 113 | 1 | VERTICAL TAB（13）？ |
|  |  | 000013 |  |  |  |  |  |
| 30 | 00376 | 001717 |  | BEQ | LP．DNP | $\dagger$ | YES，IGNORE IT ！ |
| 31 | 06400 | 000400 |  | BR | LP．Id6 | 1 | NO，FURMFEED（14）ISOLATED |
| 32 | 04402 |  | LP．I16： |  |  |  |  |
| 33 | 300402 | $\begin{aligned} & 016703 \\ & 177410 \end{aligned}$ |  | MOV | LP，SIZ，R3 | ； | R3－－PRINTERIS WIDTH ） |
| 34 | 400406 | 005403 |  | NEG | R3 | ＇ |  |
| 35 |  |  | － | －IFOF | LSIL\＆SPREAD |  |  |
| 36 |  |  |  | TST | LP．FLG | 1 | ELONGATION ENABLED？ |
| 37 |  |  |  | BEQ | LP．IVA | 1 | NO，PRINT CHARACTER |
| 38 |  |  |  | ASR | 83 | 1 | HALVE PRINTERIS WIDTH |
| 39 |  |  |  | MOV | RS，LP．FLG | 1 | REOINITIALIZE THE FLAG |
| 40 |  |  |  | －ENDC |  |  |  |
| 41 | 00410 | 040674 |  | BR | LP．IVA | 1 | PRINT THE CHARACTER |


| 1 | 000412 | 012701 <br> 000040 | LP.117: | mov | W 40.12 |  | UNPRINTABLE, BLANK SUBSTITUTIO |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 2 | 000416 | 000667 |  | BR | LP.IV3 | 1 | PRINT A GLANK |
| 3 | 000420 | 120127 | LP.II8: | CMPB | R1, 1172 | 1 | LOWER CASE ALPHABET ? |
|  |  | 000172 |  |  |  |  |  |
| 456 | 006424 | 003003 |  | BGT | LP.I19 | 1 | EXCEEDS |
|  |  |  | 1 |  |  |  |  |
|  |  |  | 1 |  | LOWEK CASE TO | UPPE | H CASE CUNVERSION PERFORMED |
|  |  |  | 1 |  |  |  |  |
| 8 | 000426 | 042701 |  | BIC | \# $40, R 1$ | 1 | CONVERSION PERFORMEO |
|  |  | 000040 |  |  |  |  |  |
| 9 | 0100432 | 000661 |  | 8R | LP. 103 | $!$ | Print character RUBOUT (177)? |
| 10 | 00434 | 120127 | LP.I198 | CMPB | R1,*177 | 1 | RUEOUT (177) ? |
|  |  | 000177 |  |  |  |  |  |
| 11 | 00440 | 001676 |  | 8EQ | LP.DNP | 1 | YES, IGNQRED |
| 12 | 00442 | 126727 |  | CMPB | UPPCAS, 137 | 1 | UPPER CASE PERMITTED ? |
|  |  | 177352 |  |  |  |  |  |
|  |  | 000137 |  |  |  |  |  |
| 13 | 06450 | 101252 |  | BHI | LP.Ib3 | 1 | YES, PRINT CHARACTER |
| 14 | 04452 | 000757 |  | BR | LP.117 | 1 | UNPKINTABLE, BLANK SUBSTITUTIO |
| 15 |  |  |  |  |  |  |  |
| 16 | 00434 | 005303 | LP.120: | DEC | RS | 1 | BACKUP PRINT POSITION |
| 17 | 00450 | 005002 |  | DEC | $R 2$ | 1 | BACKUP BUFFER POSITION |
| 18 | 00400 | 004567 | LP.121: | JSR | RS,LF.SET | 1 | RESTORE TEMPORARIES |
|  |  | 000052 |  |  |  |  |  |
| 19 | 00464 | 052737 |  | BIS | H100,0\#LP.CSR | ; | ENABLE INTERRUPT |
|  |  | 000100 |  |  |  |  |  |
|  |  | 177514 |  |  |  |  |  |
| 20 | 00472 | 00ub02 |  | RTI |  | 1 | EXIT FROM INTERRUPT |
| 21 |  |  |  |  |  |  |  |
| 22 | 00474 | 005303 | LP.L22: | DEC | $R 3$ | \% | GACKUP PRINT POSITION |
| 23 | 00476 | 005302 |  | OEC | R2 | 1 | BACKUP BUFFER POSITION |
| 24 | 00500 | 016746 | LPAERK: | MOV | LP.NAM, - (SP) | 1 | DEVICE DRIVERIS MNEMONIC |
|  |  | 177316 |  |  |  |  |  |
| 25 | 00504 | 112746 |  | MOV |  | 1 | MeSSage code |
|  |  | 000402 |  |  |  |  |  |
| 26 | 00510 | 000004 |  | 107 |  |  |  |
| 27 | 00512 | ט00167 |  | JMP | LP.INT | 1 | TRY AgAIN |
|  |  | 177372 |  |  |  |  |  |



| ADO2 - | 000402 | 010000 | Blank | 000040 |
| :---: | :---: | :---: | :---: | :---: |
| BSLSH | 000134 | 000015 | DDBADR: | 000906 |
| DUBGLK* | 000004 | 000010 | DOBCRT: | 000014 |
| DOBDVA | 177776 | 000012 | DDBULA | 000002 |
| DOBUNT: | 000013 | 000004 | DITBMP | 000016 |
| DITBSY: | 000000 | 000002 | DITINT: | 000005 |
| DITMFD: | 000014 | 000012 | DIIOPN: | 000007 |
| DITPRI* | 000006 | 000011 | DITXFR: | 000010 |
| EMTINT: | 0000u6 | 000014 | EMTVAL | 104000 |
| EMTVEC: | 000030 | 000001 | FTOOS | 000001 |
| FTMUD | 000001 | 00001 | FTKPO3: | 000001 |
| F00] | 001401 | 001462 | F003 | 001403 |
| F005 | 001465 | 001407 | F011 | 001411 |
| Fu12 | 001412 | 001417 | F024 | 001424 |
| F042 | 011442 | 001450 | F052 | 001452 |
| KSBSIZ: | 060400 | 000012 | LP | OOUODOKG |
| LPTYP | 000000 | $000052 R$ | LP.BKS | 0000268 |
| LP.CLS | VUQU36R | 177514 | LP. DGR: | 177516 |
| LP.DNE | $000252 R$ | 000236 R | LP.DON | 000256 R |
| LP.ERR | $000500 R$ | 000110R | LP.IXX | 000346 R |
| LP.IV | $000124 R$ | OUOIO6R | LP.IOL | 000162 R |
| LP.IU2 | OOD 170 O | 0001768 | LP.IO4 | $000202 R$ |
| LP.IDS | 0002302 | $0 \cup 02 S 4 R$ | LP:I10 | $000274 R$ |
| LP.III | $000366 R$ | $000310 R$ | LP.I13 | $000322 R$ |
| LP.II4 | $000350 R$ | $000362 R$ | LP.I16 | $000402 R$ |
| LP.İ 7 | VOU412R | $000420 R$ | LP. 119 | 0004342 |
| LP.I20 | $000454 R$ | OUV460R | LP.I22 | 000474 R |
| LP.LIN | $000024 R$ | 000040 | LP.NAM | $000014 R$ |
| LP.OPN | U000S6R | 000536 R | LP.SIZ | $000016 R$ |
| LP.STS | $000516 R$ | OODOSOR | LP.TOF | 0000342 |
| LP,TRN | U00000R | 000200 | LP.TRT | 000236 R |
| LPil | 000061 | 001000 | OVL006: | 000002 |
| OVLO10: | U00006 | 000022 R | OV1061: | 000012 |
| OV2061: | 000012 | 000030 | PRI4 | 000200 |
| PRI7 | 000340 | 177776 | PSPRIO: | 177437 |
| RPGIT | 004000 | 000020 | RUBOUT: | 000177 |
| SKIP2 | 000012 | 000040 | STMASK= | 107070 |
| S.RSAVA | 000044 | 000011 | UPPCAS | 000020 R |
| V.CDB | 000050 | 000052 | $V . G T B=$ | 000054 |
| V.RLB | Qu0056 | 000046 | V.RSAV | 000044 |
| V.SVT $=$ | 000040 | 000042 | WIDTH | 000120 |
| XFTCOM: | buboue | 000000 | XFTMUOE | 000000 |
| XFTRPG | 000000 | 000000 |  |  |
| - ABS. | $0 \cup 0000$ |  |  |  |
|  | 0010572 |  |  |  |
| ERRORS DETECTEO: $v$ <br> FREE CURE: 15039. WOKUS |  |  |  |  |
|  |  |  |  |  |
| -LP:LPM/ | /CKFaSYif | FEATSW, | LPO[200,200] | I/LI:ME |




| rence table S-3 |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | 110 160 |  |  |  |  |  |  |
| R2 | 109 | 4-42* | 6=13 | 7-11 | 7* 150 | 7-16 | 100170 |
|  | 10. 230 | 11. 60 | 11-7 | 11. 80 | 11-13 | 11-170 |  |
| R3 | 1-10\# | 4-43* | 7-10 | 7-140 | 70270 | 8-23 | 9-60 |
|  | 9-70 | 9-330 | 9m 340 | 10. 100 | 10= $22^{\circ}$ | 110. 12 | 11. 180 |
| R4 | 1012\% | 4. 44 \# | 7-9 | 7.130 | 7-320 | 7-370 | 8-29 |
|  | 8-310. | 11-11 | 11-140 |  |  |  |  |
| R5 | 1012 | 4-454 | 8 m 60 | 8.80 | 10-180 | 11-190 |  |
| SKIP2 | 4-24* | 9. 25 |  |  |  |  |  |
| $\begin{aligned} & \text { SMBSLZ } \\ & \text { SP } \end{aligned}$ | 2-158* | 2-159 | 2-160 |  |  |  |  |
|  | 1-13\# | 4.46\% | 7-90 | 7-100 | 7-110 | 7-120 | 8-70 |
|  | 8-8 | 8.170 | 8-230 | 8. 290 | 8-300 | 8-31 | 10. 240 |
|  | 10. 250 | 110 4 | 11-50 | 11. 7 ? | $11=14$ | 11-150 | $11-16$ |
|  | 11-17 | 11918 |  |  |  |  |  |
| SPACES <br> SPREAD | 7-20 |  |  |  |  |  |  |
|  | $5=8$ | 5-11 | 5-47 | 6- 15 | 6"20 | 8-18 | 8* 24 |
|  | 908 | 9.16 | 9. 35 |  |  |  |  |
| STMASK | 1-93 ${ }^{\text {c }}$ |  |  |  |  |  |  |
| S.RSAV | 4-51\% | 8. 7 |  |  |  |  |  |
| TABCHUPPCAS | 1-122\# |  |  |  |  |  |  |
|  | 5-34 | 7.25 | 10-12 |  |  |  |  |
| V.CDB | 1-25 |  |  |  |  |  |  |
| V.CDU | 1-20* |  |  |  |  |  |  |
| V.GTB | 1\% 27 \# |  |  |  |  |  |  |
| V.RLE | 1-28\# | ,* |  |  |  |  |  |
| V.RRES | 1-24\# |  |  |  |  |  |  |
| V.RSAV | 1-23* |  | - |  |  |  |  |
| V.SVT | 1*21\% |  |  |  |  |  |  |
| V.XIT | 1-22* |  |  |  |  |  |  |
| WIOTH | 4-36 | 6-15 | 5-33 |  |  |  |  |
| XFTCOM | 3-7* |  |  |  |  | . |  |
| XFTOOS | 3-13* |  |  |  |  |  |  |
| XFTMUO | $3-9$ \# |  |  |  |  |  |  |
| XFTRPG | 3-11\% |  |  |  |  |  |  |
| SSPASS | 2-177\# |  |  |  |  |  |  |
|  | $6=7$ |  |  |  |  |  |  |




[^0]:    ${ }^{1}$ See Part 6 for information on PIC.

[^1]:    ${ }^{1}$ In some cases, it can be further extended as discussed in later sections.
    ${ }^{2}$ Required only if Driver is being called for Special Function; addresses a Special Function Block.
    ${ }^{3}$ Required only if the device is bulk storage (e.g., Disk or DECtape).
    ${ }^{4}$ Most devices transfer words regardless of their content, i.e., ASCII or Binary. Some devices (e.g., Card Reader) may be handled differently depending on the mode for these, Bit $\varnothing$ must also be set to indicate ASCII= $\varnothing$, Binary $=1$. In these cases, the driver always produces or accepts ASCII even though the device itself uses some other code.
    ${ }^{5}$ This word may be omitted if the device is bulk storage.

[^2]:    ${ }^{1}$ This includes RFll Disk; although this is basically word-oriented, it is assumed to be subdivided into 64 -word blocks.

[^3]:    ${ }^{1}$ If the routine is not provided, these are $\varnothing$.

[^4]:    ${ }^{1}$ Dataset Data Block - a 16-word table which provides the main source of communication between the Monitor drivers and a particular set of data being processed on behalf of a using program.

