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FOR USE WITH C D C® C Y B E R 2 0 0 OPERATING SYSTEM VERSION 1

# REFERENCE MANUAL



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# **LIST OF EFFECTIVE PAGES**

New features, as well as changes, deletions, and additions to information in this manual are indicated by bars in the margins or by a dot near the page number if the entire page is affected. A bar by the page number indicates pagination rather than content has changed.





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nower case<br>indicate<br>strings.

indicate actual keywords.

are hexadecimal numbers.

UPPERCASE

lowercase

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 $\label{eq:2.1} \mathcal{L}_{\mathcal{A}}(\mathcal{A}) = \mathcal{L}_{\mathcal{A}}(\mathcal{A}) = \mathcal{L}_{\mathcal{A}}(\mathcal{A})$ 

The ELSE statement can be used with a block IF statement to provide an alternate path of execution for a block IF statement. An ELSE statement can have a statement label, but the label cannot be referenced in any other statement.

#### **ELSE IF**

The ELSE IF statement has the following form:

ELSE IF (expr) THEN

expr Any logical expression.

The ELSE IF statement can be used with a block IF statement to provide an alternate path of execution for a block IF statement or another ELSE IF statement, and to perform a conditional test. An ELSE IF statement can have a statement label, but the label cannot be referenced in any other statement. The effect of execution of an ELSE IF statement is the same as for the block IF statement.

#### **END IF**

The END IF statement has the following form:

END IF

The END IF statement terminates a block IF structure. Each block IF statement must have a corresponding END IF statement.

#### **BLOCK IF STRUCTURES**

Block IF structures provide . for alternate execution of blocks of statements. A block IF structure begins with a block IF statement and ends with an END IF statement; it can contain an ELSE statement or one or more ELSE IF statements. Each IF, ELSE, or ELSE IF statement can be followed by a block of executable statements called an If-block, else-block, and elseif-block respectively.

elseif-block. Control can<br>else-block, or elseif-block. An if-block, else-block, and elseif-block can contain any number of executable statements or no statements. Control can transfer out of an if-block, else-block, or. else if-block. Control cannot transfer into an if-block,

A simple block IF structure is shown in figure 5-1. If the expression in the block IF statement is true, execution continues with the first statement in the if-block. If the expression is false, control transfers to the statement following the END IF statement.

IF (expr). THEN

if-block

END IF

Figure 5-1. Simple Block IF Structure

A block IF structure that contains an ELSE statement is shown in figure 5-2. If the expression in the block IF statement is true, execution continues with the first executable statement in the if-block. If a statement in the A block IF structure that contains an ELSE statement is shown in figure 5-2. If the expression in the block IF statement is true, execution continues with the first executable statement in the if-block. If a statement in t if-block does not transfer control elsewhere, control transfers to the statement following the END IF statement after execution of the if-block.



Figure 5-2. Block IF Structure With ELSE Statement

If the expression in the block IF statement is false, control transfers to the first statement in the else-block. If a statement in the else-block does not transfer control elsewhere, control transfers to the statement following the END IF statement after execution of the else-block.

A block IF statement can have no more than one associated ELSE statement.

A block IF structure that contains ELSE IF statements is shown in figure 5-3. If the expression in the block IF statement is true, execution continues with the first executable statement in the if-block. If a statement in the if~block does not transfer control elsewhere, control transfers to the statement following the END IF statement after execution of the if-block.



#### Figure 5-3. Block IF Structure With ELSE IF Statements

If the expression in the block IF statement is false, control transfers to the first ELSE IF statement that is associated with the block IF statement. The expression in this ELSE IF statement is evaluated. If the expression is true, execution continues with the first executable statement in elseif-block-1. If a statement in elseif-block-1 does not transfer control elsewhere, control transfers to the statement following the END IF statement after execution of elseif-block-1.

If the expression in the first ELSE IF statement is false, control transfers to the second ELSE IF statement that is associated with the block IF statement. The expression in the second ELSE IF statement is evaluated in the same manner as in the first ELSE IF statement. Any number of ELSE IF statements can appear in a block IF structure.

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An ELSE statement can also appear in this structure;<br>however, it must follow the last ELSE IF statement. The else-block associated with the ELSE statement is executed if all of the logical expressions in the block IF statement and ELSE IF statements are false.

### **NESTING BLOCK IF STRUCTURES**

A nested block IF structure is a block IF structure that appears in an if-block, else-block, or elseif-block of<br>another block IF structure. A nested block IF structure must appear entirely within an if-block, else-block, or elseif-block. Control can transfer from an if-block, else-block, or elseif-block of a nested block IF structure to the if-block, else-block, or elself-block of the outer block<br>IF structure in which the nested block IF structure appears. Control cannot transfer from an if-block, else-block, or elseif-block of an outer block IF structure to an if-block, else-block, or elseif-block of a nested block IF structure, however. Nested block IF structures are shown in figure 5-4.



Figure 5-4. Nested Block IF Structure

A block IF structure can appear within the range of a DO loop, but the entire block IF structure must appear in the DO loop range. An END IF statement cannot be the terminal statement of a DO loop. A DO loop can appear in an if-block, else-block, or elseif-block but the entire range of the DO loop must appear in the if-block, else-block, or elseif-block.

## **DO STATEMENT**

Execution of a group of statements can be repeated a specified number of times through use of the DO statement. The range of a DO statement is the set of executable statements beginning with the first executable<br>statement following the DO and ending with the terminal statement associated with the DO. A DO statement along with its range is referred to as a DO loop.

#### **DEFINING A DO LOOP**

The DO statement has the following form:

DO  $n i = m_1, m_2, m_3$ 

- $\mathbf n$ The label of the terminal statement.
- i The control variable, a simple integer variable.
- The initial value parameter of i, an integer  $m_1$ constant or a simple integer variable with a value greater than zero.
- The terminal value parameter of i, an integer  $m<sub>2</sub>$ constant or a simple integer variable with a value greater than zero.
- Optional. **The** incrementation value  $m<sub>3</sub>$ parameter for i. an integer constant or a simple integer variable with a value greater than zero. Default value is 1.

The terminal statement of a DO loop can be any assignment statement and almost any input or output statement. However, any flow control statement other than a CONTINUE is either highly restricted or must not appear as the terminal statement of a DO. The terminal statement must not be any of the following:

- A RETURN, STOP, or PAUSE statement
- A GO TO statement of any form
- A block IF, ELSE, ELSE IF, or END IF statement ×
- A block WHERE or OTHERWISE statement
- A special call statement
- A DO statement
- A READ statement containing an ERR or END branch
- A CALL statement that passes a return label
- An arithmetic IF statement
- A logical IF statement containing any of these restricted forms

The terminal statement must physically follow and be in the same program unit as the DO statement that refers to it.

Example:

 $\bullet$ 

DO 10 I=1,11,3  $IF(ALIST(I)-ALIST(I+1))15.10.10$ 15 ITEMP=ALIST(I) 10  $ALIST(I)=ALIST(I+1)$ 300 WRITE(6,200) ALIST

The statements following DO up to and including<br>statement 10 are executed four times. The DO loop is executed with I equal to 1, 4, 7, 10. Statement  $300$  is then executed.

A DO loop can be initially entered only through the DO<br>statement. That is, the group of statements in figure 5-5 are incorrect. The GO TO statement in figure 5-5 transfers control into the range of the DO before the DO statement has been executed.

> GO TO 100 DO 100 1=1,50 100  $A(I)=I$

Figure 5-5. Incorrect: Entering Range of DO Before DO Execution

Execution of a DO statement causes the following sequence of operations:

- i is assigned the value of  $m_1$ .  $1.$
- The range of the DO statement is executed.

- 3. i is incremented by the value of  $m_3$ .
- 4. i is compared with m<sub>2</sub>. If the value of i is less than or equal to the value of m<sub>2</sub>, the sequence of operations starting at step 2 is repeated. If the value of i is greater than the value of  $m_2$  the DO is said to have been satisfied, the control variable becomes undefined (has an unpredictable value), and control passes to the statement following the statement labeled n. If  $m_1$  is greater than  $m_2$ , the range of the DO is still executed once.

A transfer out of the range of a DO loop is allowed at any time. When such a transfer occurs, the control variable remains defined at its most recent value in the loop. If control eventually is returned to the same range without entering at the DO statement, the statements executed while control is out of the range are said to define the extended range of the DO. The extended range of a DO must not contain a DO that has its own extended range.

The control variable, initial parameter, terminal parameter, and incrementation parameter of a DO must not be redefined during the execution of the range of that DO. However, the group of statements in figure 5-6 are correct. If ever an element of the array RA is zero or negative, it is set to 1 and the DO statement is reentered, which reinitializes the control variable I.

> $K=0$ GO TO 300  $200$  RA(I)=1. 300 DO 100 1=1,50 K=K+1 IF (RA(l).LE.O.) GO TO 200 100 RA(l)=K

Figure 5-6. DO Control Variable Reinitialization

#### **NESTING DO LOOPS**

When a DO loop contains another DO statement, the grouping is called a DO nest. DO loops can be nested to any number of levels. The range of a DO statement can include other DO statements only if the range of each inner DO is entirely within the range of the containing DO statement. When DO loops are nested, each must have a different control variable.

The terminal statement of an inner DO loop must be either the same statement as the terminal statement of the containing DO loop or must occur before it. If more than one DO loop has the same terminal statement, a branch to that statement can be made only from within the range or extended range of the innermost DO. Figure 5-7 gives an example of an incorrect transfer into the range of an inner DO. Since statement 500 ih figure 5-7 is the terminal statement for more than one DO loop, if the first element of any row in array A is less than or equal to zero, the consequent branch to the CONTINUE statement will be an entrance into the range of the inner DO.

If the nested loops in figure 5-7 did not share a terminal statement or if the outer loop did not reference the terminal statement, the loops would be correctly nested.

DO 500 1=1.5 IF (A(l,1).LE.O.) GO TO 500 DO 500 K=1,10 A(l,K)=SQRT(A(l ,K)) 500 CONTINUE

#### Figure 5-7. Example of Incorrect Sharing of Terminal Statement

The range of a DO loop can contain a block IF structure, but the entire block IF structure must appear in the DO loop range. An END IF statement cannot be the terminal statement of a DO loop. A DO loop can appear in an if-block, else-block, or elseif-block, but the entire range of the DO loop must appear in the if-block, else-block, or else if-block.

The range of a DO loop can contain a block WHERE structure, but the entire block WHERE structure must appear in the DO loop range. An END WHERE statement can be the terminal statement of a DO loop.

## **CONTINUE STATEMENT**

The CONTINUE statement has the following form:

**CONTINUE** 

The CONTINUE statement performs no operation. It is an executable statement that can be placed anywhere in a program without interrupting the flow of control. The CONTINUE statement is generally used to carry a CONTINUE statement is generally used to carry a statement label. For example, it can provide DO loop termination when a GO TO or IF would otherwise be the last statement of the range of the DO.

## **PAUSE STATEMENT**

The PAUSE statement has the following form:

PAUSE n

n Optional. A string of one to five decimal digits,  $or a$  character constant.

If a string is given, it is displayed in the job dayfile or at the terminal. The string is also placed in the output file<br>for the job. Program execution then continues with the for the job. Program execution then continues with the rest executable statement . following the PAUSE statement. If no string is given, instead of n being displayed and output, the string PAUSE is displayed and output before program execution continues.

## **STOP STATEMENT**

The STOP statement has the following form:

STOP n

n Optional. A string of one to five decimal digits, or a character constant.

Upon execution of the STOP statement, program execution<br>unconditionally terminates and control is returned to the operating system. If a string is given, it is displayed in the job dayfile or at the terminal. The string is also placed in the output file for the job. If no string is given, instead of n being displayed and output, the string STOP is displayed and output.

## **RETURN STATEMENT**

Subroutine and function subprograms contain one or more RETURN statements that when executed cause immediate return of control to the referencing program unit. The RETURN statement must not appear in a main program.

#### Form:

RETURN $\vert n \vert$ 



In a function subprogram, execution of a RETURN causes the function value to be returned to the referencing program unit and to be substituted for the most recently executed function reference in that program unit. Evaluation of the expression that contained the function reference continues. The integer n must not appear after a RETURN statement in a function subprogram.

In a subroutine subprogram, when n is not given, execution of a RETURN returns control to the first executable statement following the CALL statement last executed in the calling program unit. When n is given, control returns instead to a statement indicated in the argument list of the CALL statement. The statement label to which control. returns is given by the actual argument corresponding to the nth asterisk dummy argument in the SUBROUTINE or<br>ENTRY statement of the called subroutine. If there are<br>fewer than n such statement label arguments or if n 0,<br>the return is as if n had not been specified (that is, con returns to the first executable statement following the appropriate CALL statement).

## **CALL STATEMENT**

The CALL statement is used to transfer control to a subroutine subprogram, System Input/Output (SIO) module, System Request Language (SRL) module, assembly language subroutine, or any other external subroutine. The execution of a CALL statement is not complete until the subroutine designated in the statement completes execution and returns control to the calling program unit.

Form:

CALL s 
$$
(a_1, a_2, \ldots, a_n)
$$

s The symbolic name of a subroutine, or an entry point name in a subroutine.

Optional. An actual argument which can be an expression, vector, descriptor, array, external procedure name, or the label of an<br>executable statement in the same program<br>unit (the label is prefixed by an ampersand). When the argument list is omitted, the parentheses and commas must also be omitted. n must equal the number of dummy arguments in the SUBROUTINE or ENTRY statement for s.

Execution of the CALL statement transfers control to entry point name s. See the heading Passing Arguments Between Subprograms in section 7 for a further description of actual arguments in CALL statements.

Control normally returns to the first executable statement following the CALL statement. However, control can be made to return to some other statement in the program unit by appropriate selection of the CALL statement's actual arguments. If the dummy argument list in the called subroutine contains at least n asterisks, and if the called subroutine contains a RETURN n statement, then upon execution of the RETURN n statement, control returns to the statement having the nth statement label in the CALL statement actual argument list.

For example, the program in figure 5-8 uses both the RETURN n and the RETURN statement formats. If the data read with the READ statement in the subroutine is less than 1.0 or greater than 10.0, control transfers back to th message is printed out and the program terminates. On the other hand, if the data is within the appropriate range, the.<br>subroutine ' continues' executing until the RETURN statement is reached, at which time control transfers back to the main program statement that immediately follows the call to the subprogram.



## **VECTOR PROGRAMMING 1 1**

Detailed in this section are the ways that a user can introduce machine vector instructions into the object code for a FORTRAN program. Any of the forms described here can be used in the same program with the previously described FORTRAN features.

## **AUTOMATIC VECTORIZATION**

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Automatic vectorization is a process by which the FORTRAN compiler translates an iterative, sequential procedure into parallel procedures. The aim of the process is to utilize the capabilities of the CYBER 200 hardware to produce optimal object code, without requiring alteration of FORTRAN programs that do not use the extensions of CYBER 200 FORTRAN, and without necessitating that a problem be reconceptualized in terms of parallel processes. Automatic vectorization of a FORTRAN program is selected by including the V compile option in the FORTRAN system control statement that requests compilation of the program.

Under the V option, CYBER 200 vector instructions are generated for DO loops that have certain characteristics. The object code generated for a loop that is accepted by the vectorizer consists of vector instructions rather than scalar instructions. If a loop is rejected by the vectorizer, the compiler attempts to transform the loop into a call to one of the supplied STACKLIB routines.

Automatic vectorization can be used with any FORTRAN program, including FORTRAN programs that do not use any of the extensions of CYBER 200 FORTRAN. However, because of the restrictiveness of the conditions for vectorization, summarized in table 11-1, it might not be possible for the vectorizer alone to achieve the degree of vectorization desired. As an alternative, the programmer can elect to use other methods, in conjunction with the V compile option or not, to specify vector operations explicitly.

#### **GENERAL CHARACTERISTICS OF VECTORIZABLE DO LOOPS**

A simple vectorizable DO loop is shown  $at(3)$  in figure 11-1. The range of a vectorizable loop can contain assignment statements, CONTINUE statements, and DO statements. An input/output statement or IF statement, for example, is not acceptable in a loop that is to be vectorized.

The initial, terminal, and incrementation parameters of the DO statement of a vectorizable loop must have certain characteristics. The incrementation parameter, if present, must be 1; an incrementation value of 2, for example, causes the loop not to be vectorized. Secondly, FORTRAN allows the parameters to be constants or variables; however, a variable initial, terminal, or incrementation parameter does prohibit the vectorization of any containing DO loop. For instance, the vectorizable loop defined at  $(3)$  has a variable terminal parameter. Loop  $(2)$  contains loop  $(3)$  and, consequently, cannot be vectorized. Thirdly, the iterative count of a loop or entire



Figure 11-1. Form of Vectorizable DO Loops

nest of loops must be less than or equal to 2<sup>16</sup>–1 (that is, 65535). By this criterion, loops *(T*) and *(*6) in part B of figure 11–1 can be vectorized, depending on the range of the innermost loop; but loop  $(5)$  cannot be vectorized (because  $30 * 200 * 11 = 66000$ ).

When the initial or terminal parameter of a loop is a variable, the dimensions of the loop-dependent array references within the loop are used to determine the largest possible iterative count through which the loop can pass, and this count is used to decide if the loop can be .<br>vectorized.

The U compile option can be selected for unsafe vectorization. When U is selected, the compiler vectorizes loops that contain dummy arrays, even if the terminal value of the loop is variable. The optimization is considered unsafe because the presence of a variable dimension might cause the iterative loop count to exceed 65535.

The U compile option also enables vectorization of loops that contain an equivalenced data element on the left side of an assignment statement.

If a loop cannot be vectorized (loop  $(2)$  in figure 11-1, for instance), then a loop containing the nonvectorizable loop cannot be vectorized either. By this criterion, loop  $(1)$  is nonvectorizable.

### TABLE 11-1. CRITERIA FOR VECTORIZABLE LOOPS



### **ASSIGNMENT STATEMENTS IN VECTORIZABLE DO LOOPS**

Operators in assignment statements in a vectorizable loop can be any of the arithmetic or logical operators. The use of relational operations within a loop causes the loop not to be vectorized.

The type of an operand occurring in the range of a vectorizable loop can be integer, real, or logical. A vectorizable loop containing a logical assignment statement is shown in figure 11-2.

LOGICAL A, C, R DIMENSION A(50000), C(50000), R(49999) DO 999 X=2,50000  $R(X-1) = (A(X-1)$ . AND.  $A(X)$ ) .OR. (C(X-1) .AND. C(X)) 999 CONTINUE

Figure 11-2. Vectorizable Loop #1

For complex vector arithmetic expressions, the following restrictions apply:

- Operands can be integer, real, or complex. Double-precision operands are not allowed.
- Exponentiation is not allowed; the operators in a complex vector expression can be only  $+, -, *,$  and  $/$ .

For double-precision vector arithmetic expressions, the following restrictions apply:

- The expression must consist of either a double-precision vector or a reference to a FORTRAN-supplied double-precision vector function. No operators are allowed.
- The expression can appear only in a vector arithmetic assignment statement of type double-precision.

Given the declarations:

DESCRIPTOR Dl, SCRP, RZLT DIMENSION SCRP $(3,3)$ , VR $(100)$ , R $(100)$ DATA Dl/VR(1;50)/, SCRP(3,l)/VR(l;100)/

the following are examples of vector arithmetic expressions:

**e** VR(l;lOO)

Current values of the 100 consecutive elements in the array VR.

**e** Dl

Current values of the first 50 consecutive elements in the array VR.

 $D1 + N$ 

A vector formed by adding the value of the scalar N to each element of VR(1;50).

 $-(D1 + N)/2.**M$ 

A vector formed by adding N to Dl, negating, then performing a divide by 2\*\*M on each element of  $V_{R(1;50)}$ .

 $SCRP(3,1)$ 

Current values of the 100 consecutive elements in the array VR.

•  $VEXP(R(10; 52); RZLT)$ 

Vector function reference.

#### **VECTOR RELATIONAL EXPRESSIONS**

A vector relational expression consists of a relational operator flanked by two expressions. The relational opera tors are:



The periods are part of the operators and must appear.

A vector relational expression has one of the following forms:



vaej

sae A scalar arithmetic expression of type real or integer, but not of type complex or double-precision.

op Cone of the relational operators.

A vector arithmetic expression of type real or integer, but not of type complex or double-precision.

A vector relational expression, which al ways contains one or more vector data elements, evaluates to a bit vector of truth values represented by bits 0 and 1. (In contrast, evaluation of a scalar relational expression results in a. single logical value.)

When both operands for a relational operation are vectors, the operation compares successive elements of one vector operand with corresponding elements of the other vector operand. If the specified relation holds between the pair of elements, the operation sets (assigns 1 to) the operation sets (assigns 1 to) corresponding bit in the result bit vector. If the relation does not hold, the operation clears (assigns 0 to) the corresponding bit in the result bit vector. When one operand is a vector and the other a scalar, the scalar is compared with each element of the vector during evaluation of the expression.

Given the declaration;

DESCRIPTOR Dl

the following are examples of vector relational expressions:

**e** X+Y/3.\*Z .LT~ VR(l;lOO)

Bit vector having a length of 100 bits, where the i<sup>th</sup> bit is 1 if the i<sup>th</sup> element of the real vector  $VR(1;100)$  is greater than or equal to the value of the scalar arithmetic expression X+Y/3.\*Z, and is O otherwise.

**D1.NE. D1\*2** 

Bit vector having the length of the vector that D1 points to, where the i<sup>th</sup> bit is 1 if an element of the vector is nonzero, and is 0 otherwise.

**e** VR(1;89) ~EQ. VR(2;89)

Bit vector having a length of 89 bits, where the i<sup>th</sup> bit is 1 if the  $i$ <sup>th</sup> element of the real vector  $VR(1;89)$  equals the  $i$ <sup>th</sup> element of the real vector  $VR(2,89)$ , and is 0 otherwise.

 $R(10;20)$ .GE. 0.34

Bit vector having a length of 20 bits, where the i<sup>th</sup> bit is 1 if the i<sup>th</sup> element of the real vector R(10;20) is greater than or equal to  $0.34$ , and is  $0$  otherwise.

#### $VR(1;25)*3.$ .EQ.  $(VC(1;25)-ST+REM)$

Bit vector having a length of 25 bits, where the j<sup>th</sup><br>bit is 1 if the j<sup>th</sup> element of the real vector<br> $VR(1;25)$ , multiplied by 3, is equal to the j<sup>th</sup><br>element of the real vector  $VC(1;25)$  plus the value of REM-ST, and is 0 otherwise.

### **BIT EXPRESSIONS**

A bit expression can be a vector relational expression, bit vector, bit descriptor, bit vector function reference, bit descriptor array element, bit constant, bit variable, bit array element, or a bit expression enclosed in parentheses. If B and C are vector bit expressions, then B followed by a<br>logical operator followed by C is also a vector bit expression.

The operators used in vector bit expressions are the logical operators interpreted so that truth is the bit value 1 and falsity is the bit value 0. The mathematical definitions of the logical operators are given in section 3, the precedences of the operators are the same as for logical operators in logical expressions.

A vector relational expression evaluates to an entire bit vector of logical results. Logical operations on bit vectors are performed on either corresponding elements of two vector operands or a scalar operand paired with successive elements of a vector operand.

Given the declarations:

BIT VC(100), BV, BV2, BVA, B1, B2 DESCRIPTOR BV, BV2, BVA

the following are examples of vector bit expressions:

 $\overline{R}V$ 

Value of the bit vector that BV points to.

BV.OR.BV2

Bit vector in which the  $i<sup>th</sup>$  element is 1 unless the  $i<sup>th</sup>$  elements of the vector values of BV and BV2 are 0, in which case the  $i<sup>th</sup>$  element is 0.

 $((BV.OR. BY2).OR. BYA)$ 

Bit vector in which the  $i<sup>th</sup>$  element is 1 unless the i<sup>th</sup> elements of the values of (BV .OR. BV2) and BVA are 0, in which case the i<sup>th</sup> element is 0.

(VC(1,25).GE. 5.431).AND. (VC(1,25).LT. 5.931)

Bit vector in which the i<sup>th</sup> element is 1 if the i<sup>th</sup> elements of the values of the two vector relational expressions are 1, and is 0 otherwise.

B "101"

Bit constant.

 $B1$ 

Bit variable.

.NOT. (B1.AND.B2)

Bit expression.

## **EXECUTABLE STATEMENTS**

CYBER 200 FORTRAN vector programming statements include two assignment statements, one ASSIGN<br>statement, and one specification statement. The types of vector expressions appearing in the assignment statements described here are shown in table  $11-2$ .





<sup>T</sup>A vector arithmetic expression can be a primary in a vector relational expression which in turn can be a primary in a vector bit expression.

#### **DESCRIPTOR ASSIGN STATEMENT**

At execution time, the descriptor ASSIGN statement can be used to associate a vector with a descriptor. The descriptor ASSIGN statement has two forms.

First form:

ASSIGN p,q

A descriptor, or a descriptor array element. Þ

A vector, descriptor, or descriptor array element.  $\mathbf{Q}$ 

The type of p must agree with the type of q. The descriptors and vectors can be of type integer, real, bit, or complex: double-precision vectors cannot be assigned to a descriptor.

Execution of this form of the descriptor ASSIGN statement causes  $p$  to point to the vector  $q$  (if  $q$  is a vector) or to the vector that q points to (if q is a descriptor or descriptor<br>array element). After the ASSIGN has been executed, the operand p points to a vector.

The use of the first form of the descriptor ASSIGN statement is illustrated in part A of figure 11-8. In the figure, a bit array is declared and given initial values. The rows of the bit array are then indicated to be vectors in the first descriptor ASSIGN statement within the DO.<br>After the descriptor array element Y(I) has been given a<br>bit vector value, the second ASSIGN statement assigns the same bit vector value to the descriptor array element Y(16-I). The result of completing the DO loop is the configuration shown in part B of figure 11-8.

Memory for a vector need not be allocated at compile time as happens when, for example, a vector name is used in the first form of the descriptor ASSIGN statement. The following descriptor ASSIGN statement form can be used to allocate dynamic space for a vector at execution time, and to cause a descriptor or descriptor array element to point to it.



I

After execution of this second form of the descriptor ASSIGN statement, p points to a vector consisting of dynamic space having a length of e words if p is real or integer, e bits if p is bit, or 2e words if e is complex.<br>Before the value of the dynamic space pointer is incremented, the current value of the dynamic space pointer is assigned as the base address of the vector pointed to by p.

Dynamic space is managed as a stack. After an assignment to dynamic space has been made with a descriptor ASSIGN statement, the dynamic space pointer (the location of the top of the stack) is incremented. The next space available is made to begin at the first doubleword boundary. Subsequently, the execution of a FREE statement or a RETURN statement releases all dynamic space allocated by descriptor ASSIGN statements in the program unit.

#### **FREE STATEMENT**

Execution of the FREE statement (or completion of program unit execution) reverses the effect of a descriptor ASSIGN statement in which a reference .DYN. to the dynamic space pointer appears. The FREE statement resets the dynamic space pointer to the value it had before execution of the first descriptor ASSIGN statement in the program unit. All space assigned through the use of descriptor ASSIGN statements is released; if more than one such assignment was made, all are reversed.

Form:

FREE

## **VECTOR ARITHMETIC .ASSIGNMENT STATEMENT**

A vector arithmetic assignment statement has the following form:

- v=e
- v A vector ·of type integer, real, complex, or double-precision; or a descriptor or descriptor array element of type integer, real, or complex.
- e A vector arithmetic expression, or a scalar arithmetic expression.

The value of  $e$  is assigned to  $v$  . When  $v$  is double-precision, e can only· be a double-precision vector or a reference to a predefined CYBER 200 FORTRAN double-precision vector function (listed in appendix E).

If  $e$  evaluates to a scalar, that scalar is stored into every element of v; but if e evaluates to a vector, the first element of  $e$  is stored into the first element of  $v$ , the second element of e is stored into the second element of v, and so on. If the type of v differs from that of e, type conversion takes place, during assignment, to the type of *v* • Type conversion rules are given in table 11-3.

Examples, given the following declarations:

DOUBLE PRECISION RES(30), DPN, EXN(30)<br>DESCRIPTOR D1, DEX DATA DEX, D1/RES(1;20), EXN(1;15)/

**<sup>e</sup>**RES(1;30) = DSQRT(DPN)

Replace the value of the  $i<sup>th</sup>$  element of RES(1;30) with the value returned by the predefined function reference DSQRT(DPN). The sequential equivalent is:

 $DO 3I = 1,30$  $3$  RES(I) = DSQRT (DPN)

 $RES(1;30) = EXN(1;30)$ 

Replace the value of the ith element of RES(1;30) with the ith element of EXN(1;30). The sequential equivalent is:

DO  $3I = 1,30$  $3$  RES(I) = EXN(I)



#### $DEX = -(D1+N)/2$ .\*\*M

Cause DEX to point to the vector value of the vector arithmetic expression  $-(D1+N)/2.**M$ . The sequential equivalent is:

 $DO 3 I = 1,15$  $EXN(I) = -(EXN(I)+N)/2.**M$  $\boldsymbol{\mathcal{X}}$ ASSIGN DEX, EXN(1;15)

 $D1 = -30.0$ 

Replace the value of the ith element of the vector<br>pointed to by D1 with -30.0. The sequential equivalent is:

DO 3 I =  $1,15$ <br>EXN(I) = -30.0

 $\bf{3}$ 

ASSIGN D1, EXN(1;15)

#### **BIT ASSIGNMENT STATEMENT**

A bit assignment statement has the following form:

 $v = e$ 

- A bit vector, bit descriptor, bit descriptor array  $\mathbf{v}$ element, bit variable, or bit array element.
- A bit expression.  $\ddot{\textbf{e}}$

Execution of the statement causes the complete evaluation of e and the storing of the result into  $v$ . If  $v$  is longer than the result, the remaining rightmost bits of  $v$  are padded with zeros.

Examples, given the following declarations:

**BIT BV(60), BD DESCRIPTOR BD** 

 $BV(1;60) = BD$ 

Replace the value of the ith element of  $BV(1;60)$ with the value of the ith element of the vector pointed to by BD.

## $BD = BD$ .OR.  $BV(1;60)$

Cause BD to point to the vector value of the vector bit expression BD .OR. BV(1;60).

 $BV = B'1'$ 

Set each element of the array BV to 1. The sequential equivalent is:

 $DO 4I = 1,60$  $\overline{a}$  $BV(I) = B'1'$ 

 $BV = B'100'$ 

Set each element of the array to BV 1. (Note that the bit constant B '100' is truncated to B '1' since each element of the bit array BV is assigned separately.) The sequential equivalent is:

$$
DO 5 I = 1,60
$$
  
5 B V(I) = B'100'

 $BV(1:60) = B'1'$ 

Set the bit vector to 1 one and 59 zeros.

 $BV(1;60) = B'101'$ 

Set the bit vector BV to 101 followed by 57 zeros.

#### **WHERE STATEMENT**

The WHERE statement has the following form:

WHERE  $(e)$  st

A vector bit expression. ė

A vector assignment statement. st

All vector operands in the vector bit expression and in the vector expression that appears in the vector assignment statement must have the same length. All vectors and vector expressions that appear in the vector assignment statement must be of type integer or real. The vector expression that appears in the vector assignment statement

П

can contain only addition, subtraction, multiplication, and division operations, and references to the vector functions VFLOAT, VIFIX, VINT, VAINT, VSQRT, VABS, and VIABS.

When the WHERE statement is executed, the vector bit expression is evaluated. The evaluation produces a control vector. A control vector is a bit vector that controls the storing of values into a vector. This control vector is used for the vector assignment statement that appears in the WHERE statement.

The vector expression that appears in the vector assignment statement is evaluated. Each value of the result vector is assigned to the corresponding vector element on the left side of the vector assignment statement only if the corresponding element in the control vector contains a 1 bit.

A value is not assigned to the corresponding vector element on the left side of the vector assignment statement if the corresponding element in the control vector contains. a O bit. If a value is not assigned to a vector element, data flag branches are disabled for the operations that compute that value.

: Given the declarations:

REAL A(5), B(5), C(5) BIT CA(5), CB(5) DATA A/3.0, 9.0, 12.0, 2.0, 16.0/ DATA B /6.0, 6.0, 10.0; 5.0, 4.0/ DATA C/9.0, 3.0, 0.0,7.0, 7.0/ DATA CA /B'11101'/, CB /B'11000'/

the following are examples of the WHERE statement:

WHERE  $(A(1;5),LT.B(1;5)) C(1;5)=B(1;5)-A(1;5)$ 

Causes the value the value 3.0 to be assigned to the first and fourth elements of vector  $C(1;5)$ . All other elements of vector C(l;5) are unchanged.

**WHERE (CA(1;5).AND.CB(1;5)) C(1;5)=A(1;5)+B(1;5)** 

Causes the values  $9.0$  and  $15.0$  to be assigned to the  $first$  and second elements of vector  $C(1;5)$ respectively. All other elements of vector C(1;5) are unchanged.

**e** WHERE (C(l;5).NE.O.O) A(1;5)=B(l;5)/C(l;5)

Causes the values .66667, 2.0, .71429, and .57143 to be assigned to the first, second, fourth, and fifth elements of vector. A(1;5) respectively. The division of the third element of vector B(1;5) by the third element of vector C(1;5) is a division by zero; however, because the third element of the control vector contains a O bit, no data flag branch occurs.

#### **BLOCK WHERE STATEMENT**

The block WHERE statement has the following form:

WHERE(e)

e A vector bit expression.

When the block WHERE statement is executed, the vector bit expression is evaluated. The evaluation produces a control vector. A control vector is a bit vector that controls the storing of values into a vector. This control vector is used for all of the vector assignment statements

that appear between the block WHERE statement and the next END WHERE statement.

#### **OTHERWISE STATEMENT**

The OTHERWISE statement has the following form:

**OTHERWISE** 

The OTHERWISE statement can be used with a block WHERE statement to reverse the effect of the control vector established in the block WHERE statement. Reversing the effect of the control vector causes a value to be assigned to a vector element only if the corresponding element in the control vector contains a 0 bit, rather than a 1 bit. An OTHERWISE statement affects all vector assignment statements that appear between the OTHERWISE statement and the next END WHERE statement.

#### **END WHERE STATEMENT**

The END WHERE statement has the following form:

END WHERE

The END WHERE statement terminates a block WHERE structure. Each block WHERE statement must have one corresponding.END WHERE statement.

### **BLOCK WHERE STRUCTURES**

Block WHERE structures provide for the execution of any number of vector assignment statements using a single control vector. A control vector is a bit vector that controls the storing of values into a vector.

A block WHERE structure begins with a block WHERE contain one OTHERWISE statement. The block WHERE statement can be followed by a block of vector assignment statements called a where-block. An OTHERWISE statement can be followed by a block of vector assignment statements called an otherwise-block.

A where-block or an otherwise-block can contain any no statements. No other types of statements can appear in a where-block or otherwise-block. All vector operands that appear in a where-block or otherwise-block must have the same length as the control vector that is established in the block> WHERE statement. All vectors and vector expressions that appear in a where-block or otherwise-block must be of type integer or real. All vector<br>expressions that appear in a where-block or expressions that appear in a where-block or otherwise-block can contain only addition, subtraction, multiplication, and division operations, and references to the vector functions VFLOAT, VIFIX, VINT, VAINT, VSQRT, VABS, and VIABS.

Control must not transfer into a where-block or otherwise-block.

A simple block WHERE structure is shown in figure 11-9. When the block WHERE statement is executed, the vector bit expression in the block WHERE statement is evaluated. The evaluation produces a control vector. A control vector is a bit vector that controls the storing of values into a vector. This control vector is used for all of the vector assignment statements that appear between the block WHERE statement and the next END WHERE statement.



Figure 11-9. Simple Block WHERE Structure

The vector assignment statements in the where-block are then executed in order as follows:

- The vector expression that appears in a vector A. assignment statement in the where-block is evaluated.
- Each value of the result vector is assigned to the corresponding vector element on the left side of the  $2.$ vector assignment statement only if the corresponding element in the control vector contains a 1 bit. A<br>value is not assigned to the corresponding vector element on the left side of the vector assignment statement if the corresponding element in the control vector contains a 0 bit. If a value is not assigned to a<br>vector element, data flag branches are disabled for the operations that compute that value.
- $3<sup>2</sup>$ Steps 1 and 2 are repeated for each vector assignment statement in the where-block.

A block WHERE structure that contains an OTHERWISE statement is shown in figure 11-10. When the block WHERE statement is executed, the vector bit expression in the block WHERE statement is evaluated. The evaluation<br>produces a control vector. A control vector is a bit vector produces a control that controls the storing of values into a vector. This control vector is used for all of the vector assignment statements that appear between the block WHERE statement and the next END WHERE statement.

**WHERE** (a) where-block **OTHERWISE** otherwise-block **END WHERE** Figure 11-10. Block WHERE Structure With **OTHERWISE Statement** 

The vector assignment statements in the where-block are then executed in order as follows:

The vector expression that appears in a vector 1. assignment statement in the where-block is evaluated.

 $2.$ Each value of the result vector is assigned to the corresponding vector element on the left side of the vector assignment statement only if the corresponding element in the control vector contains a 1 bit. A value is not assigned to the corresponding vector element on the left side of the vector assignment statement if the corresponding element in the control<br>vector contains a 0 bit. If a value is not assigned to a vector element, data flag branches are disabled for the operations that compute that value.

Steps 1 and 2 are repeated for each vector assignment statement in the where-block.

The vector assignment statements in the otherwise-block are then executed in order as follows:

- The vector expression that appears in a vector  $\mathbf{1}$ assignment statement in the otherwise-block is evaluated.
	- Each value of the result vector is assigned to the corresponding vector element on the left side of the vector assignment statement only if the corresponding element in the control vector contains a  $0$  bit.  $\tilde{A}$ value is not assigned to the corresponding vector element on the left side of the vector assignment<br>statement if the corresponding element in the control vector contains a 1 bit. If a value is not assigned to a vector element, data flag branches are disabled for the operations that compute that value.
	- Steps 1 and 2 are repeated for each vector assignment statement in the otherwise-block.

Given the declarations:

 $3.$ 

 $2.$ 

 $3.$ 

REAL A(5), B(5), C(5), D(5), E(5) BIT CA(5), CB(5) DATA A/3.0, 1.0, 6.0, 12.0, 12.0/ DATA B/4.0, 1.0, 8.0, 12.0, 16.0/ DATA C/5\*0.0/, D/5\*0.0/, E/5\*0.0/ DATA CA/B'10101'/, CB/B'11111'/

the following are examples of block WHERE structures:

WHERE  $(A(1;5), \text{NE.B}(1;5))$ E(1;5)=A(1;5)\*A(1;5)+B(1;5)\*B(1;5)<br>C(1;5)=VSQRT(E(1;5);C(1;5))  $D(1;5)=A(1;5)*B(1;5)/2.0$ **END WHERE** 

The statements in the where-block cause the values 25.0, 100.0, and 400.0 to be assigned to the first, third, and fifth elements of vector E(1;5) respectively. They also cause the values 5.0, 10.0, and 20.0 to be assigned to the first, third, and fifth elements of vector  $C(1;5)$ respectively, and cause the values 6.0, 24.0, and 96.0 to be assigned to the first, third, and fifth elements of vector D(1;5) respectively. All other elements of vectors  $C(1;5)$  and  $D(1;5)$  are unchanged.

WHERE (CA(1;5).AND.CB(1;5))  $E(1;5)=A(1;5)*A(1;5)+B(1;5)*B(1;5)$  $C(1;5) = VSQRT(E(1;5); C(1;5))$  $D(1;5)=A(1;5)*B(1;5)/2.0$ **END WHERE** 

Has the same effect as the previous example.

WHERE  $(A(1;5), \text{NE.B}(1;5))$ E(1;5)=A(1;5)\*A(1;5)+B(1;5)\*B(1;5)<br>C(1;5)=VSQRT(E(1;5);C(1;5))  $D(1;5)=A(1;5)*B(1;5)/2.0$ **OTHERWISE**  $E(1;5)=16.0$ <br>C(1;5)=16.0  $D(1;5)=16.0$ **END WHERE** 

The statements in the where-block cause the values 25.0, 100.0, and 400.0 to be assigned to the first, third, and fifth elements of vector  $E(1,5)$  respectively. They also cause the values 5.0, 10.0, and 20.0 to be assigned to the first, third, and fifth elements of vector C(l;5) respectively, and cause the values 6.0, 24.0, and 96.0 to be assigned to the first, third, and fifth elements of vector  $D(1;5)$  respectively.

The statements in the otherwise-block cause the value 16.0 to be assigned to the second and fourth elements of vectors  $C(1;5)$  and  $D(1;5)$ .

### **NESTING BLOCK WHERE STRUCTURES**

A block WHERE structure can appear in an if-block, else-block, or elseif-block of a block IF structure, but the entire block WHERE structure must appear in the if-block, else-block, or elseif-block.

A block WHERE structure can appear in the range of a DO loop, but the entire block WHERE structure must appear in the range of the DO loop. An. END WHERE statement can be the terminal statement of a DO loop.

## **DECLARATIONS**

Vector programming adds one specification statement to the list of nonexecutable statements that can appear at the beginning of a CYBER 200 FORTRAN program unit.

#### **DESCRIPTOR STATEMENT**

The DESCRIPTOR statement has the following form:

DESCRIPTOR  $v_1, v_2, \ldots, v_n$ 

A variable, array declarator, or array name,  $v_i$ of type real, integer, bit, or complex.

All variables in the DESCRIPTOR statement list are declared to be descriptors, and any array or array declarator list element specifies a descriptor array. For example, the statement pair:

DESCRIPTOR A,B,C(3,4) REAL A,B(6,2),C

specifies A to be a real descriptor, and B and C to be. descriptor arrays having 12 type real descriptor array elements each.

The type of  $v_i$  must be established with an explicit type declaration statement, or by the first-letter rule. Although vectors can be double-precision, a descriptor cannot be double-precision.

## **INITIALIZING DESCRIPTORS AND VECTORS**

The nonexecutable DATA statement, described in section 6, can be used to place initial values in vectors and descriptors before the program begins executing. Double-precision vectors cannot be initialized in a DATA statement, although the double-precision array or individual array elements can be so initialized.

As described in section  $6$ , a data initialization statement consists of pairs of lists; a list of variables is paired with a list of constants used as the initial values for the variables. Besides scalar list elements, the list of variables can include vectors, descriptors, descriptor arrays, and descriptor array elements.

For vectors, a vector name in the variable list must contain only integer constant subscript expressions and vector length specification. The number of constant list elements corresponding to the name must be equal to the length of the vector. For example, if a vector name in the variable list is  $A(1,10)$ , then  $10$  consecutive constant list elements must correspond to the vector name. (This is similar to the way that arrays can be initialized in a DATA statement.)

For descriptors and descriptor array elements, a descriptor in the variable list must correspond only to a vector, which must contain only integer constant subscript expressions and vector length specification.

The repeat count specification in a DATA statement (section 6) can be used to specify the repeated use of a vector for initialization of more than one descriptor or descriptor array element. The data. types of corresponding variable list and constant list items must, in the above cases, be the same.

Examples of initializing vector descriptors are given in section 16.

## **'VECTOR FUNCTION SUBPROGRAMS**

Vector function subprograms are defined in almost the same way that other function subprograms are defined. The differences lie in the argument list form, the number of data types available for vector function results, and the fact that the function name must appear in a<br>DESCRIPTOR statement in the function subprogram.

Form:

 $\mathbf{t}$ 

 $\ddot{\mathbf{f}}$ 

t FUNCTION  $f(a_1, a_2, \ldots, a_n; *)$ 

Optional declaration of the type of f. When present, t can be INTEGER, REAL, BIT, or COMPLEX but cannot be DOUBLE PRECISION.

- The function's symbolic name.
- $a_i$ Dummy argument. The possible dummy arguments here are the same as for scalar functions, n must be greater than or equal to 1.

The function name f must appear in a DESCRIPTOR statement within the function subprogram. If t is not specified, f can appear in a type statement or be typed implicitly. The semicolon in the dummy argument list is required to separate the input list from the dummy output argument, which is represented by the asterisk.

Refer to Function Subprograms in section 7 for a more<br>detailed. discussion. of function. names .and function. subprogram program units.

## **REFERENCING VECTOR FUNCTIONS**

A vector function is referenced when the name of the function, followed by an actual argument list enclosed in<br>parentheses, appears in an arithmetic expression in an arithmetic assignment statement. The actual arguments that can correspond to a dummy argument are shown in section 7.

The actual argument list in the function reference is divided into two parts by a semicolon. Input arguments precede the semicolon and are separated by commas; they can be scalar expressions, vectors, descriptors, or descriptor array elements. A single output argument of the function follows the semicolon and can be a vector, descriptor, or descriptor array element. The output argument must be the same data type as the function. CYBER 200 FORTRAN permits double-precision output arguments to be used in references only to some predefined<br>CYBER 200 FORTRAN vector functions (listed in appendix E).

## **SECONDARY ENTRY POINTS**

Vector function subprograms, like scalar function subprograms, can have multiple entry points defined for<br>them. The ENTRY statement (described in section 7) specifies that the first executable statement following the ENTRY statement is a secondary entry point. More than one entry point can be declared in a subprogram; also, a scalar or vector function subprogram can have both scalar and vector secondary entry points.

Form:

ENTRY  $e(a_1, a_2, \ldots, a_n; *)$ 

The symbolic name of the entry point.  $\bullet$ 

Dummy argument. The possible dummy a: arguments here are the same as for scalar ENTRY statements, n must be greater than or equal to 1.

Like the function name, the entry point name must appear in a DESCRIPTOR statement within the scalar or vector function subprogram. Again, the semicolon separates the<br>dummy input argument list from the dummy output<br>argument which is represented by the asterisk.

The statements made in section 7 with respect to referencing secondary entry point names apply to the referencing of the entry point names defined in a vector function subprogram.

 $\overline{\phantom{a}}$ 

 $\overline{\phantom{a}}$ 

#### TABLE 14-1. FORTRAN-SUPPLIED FUNCTIONS (Contd)



<sup>11</sup> Provides the same effect as the implied conversion in assignment statements.

tttEach control vector c is type bit, and each index vector i is type integer.

## **FUNCTION DESCRIPTIONS**

The following descriptions are listed in strict alphabetical order. However, since a naming convention uses the letter V as a prefix to scalar function names to produce the corresponding vector function names, all functions with vector results can be found under  $V$  and  $Q8V$ . If a vector function input argument can be a vector, it is implied that it can also be a descriptor or descriptor array element. Also, except for some of the vector functions, none of the functions alters the values of its arguments. The mathematical values of some of the mathematical functions can be indefinite.

A generic function generates a real or integer result, depending on the mode of the argument. For instance, the Q8SSUM function is a generic function:.

A typeless function generates a result that is not converted for use as an argument or for assignment. For instance, the Q8SINSB function is a typeless function.

#### **ABS( a)**

For a real number x, ABS(x) computes the absolute value 1x1.

See ASIN for a description of the ACOS function.

### **AIMAG(a)**

This returns the imaginary part of a complex number as a real number; if  $x+iy$  is the complex number, AIMAG returns y.

### **AINT(a)**

For a real number x,  $\text{AINT}(x)$  computes  $[x]$ , where  $[A]$  is the sign of A times the largest integer less than or equal to 1A1. AINT returns a real result even though its value is always integral.

#### **ALOG(a)**

This computes the natural logarithm of a real number greater than zero. The result is a real number accurate to approximately 45 bits.

For a given real number x, ALOG(x) is calculated as follows.

For x outside the range:

 $((2)^{1/2}/2 \leq x < (2)^{1/2})$ 

let:

 $x = 2^n * w$ 

where:

1/2Sw<l

and n is an integer that satisfies the equation.

Also, let:

 $t = (w - (2)1/2/2)/(w + (2)1/2/2)$ 

Then:

 $\log_e(x) = (n - 1/2) * \log_e(2) + \log_e((1 + t)/(1 - t))$ 

For x in the range:

 $((2)^{1/2}/2 \leq x < (2)^{1/2})$ 

let:

 $t = (x - 1) / (x + 1)$ 

Then:

 $log_e(x) = log_e((1 + t) / (1 - t))$ 

In either case:

$$
\log_e((1+t)/(1-t)) = 2t \sum_{n=0}^{6} c_n t 2n
$$

where:

 $c_0 = 1.0000000000000000172016224 * 10^0$ 

 $c_1 = 3.3333333333327618176885283 * 10^{-1}$ 

 $c_2$  = 2.000000003098077890899307  $*$  10-1

 $c_3 = 1.428570799460827347261398 * 10^{-1}$ 

 $c_4$  = 1.111171831154342806719000 \* 10<sup>-1</sup>

 $c_5 = 9.060935658179353717214254 * 10^{-2}$ 

 $c_6 = 8.419186575863053137534817 * 10^{-2}$ 

If a zero or negative argument is received, a data flag branch occurs inside the routine.

#### **AlOGlO(a)**

This computes the logarithm of a real number. The result is a real number that is accurate to approximately 45 bits.

For a given real number x greater than zero:

 $log_{10}(x) = log_{10}(e) * log_{0}(x)$ 

where the natural logarithm is computed as-described for the function ALOG.

If a zero or negative argument is received, a data fiag branch occurs inside the routine.

## **AMAX0(a<sub>1</sub>,a<sub>2</sub>, ...)**

This searches a list of integer numbers for the list element having the maximum value. The integer found is returned as a real number.

### **AMAX1(a<sub>1</sub>,a<sub>2</sub>, ...)**

This searches a list of real numbers for the list element having the maximum value and returns that value.

#### $AMINO(a_1, a_2, \ldots)$

This searches a list of integer numbers for the list element having the minimum value. The integer found is returned as a real number.

#### **AMIN1**  $(a_1, a_2, \ldots)$

This searches a list of real numbers for the list element having the minimum value and returns the number when found.

### AMOD $(\alpha_1, \alpha_2, \ldots)$

This computes one real number modulo a second real number and produces a real result. AMOD(x,y) is defined as  $x-[x/y]^* y$ , where [A] is the sign of A times the largest integer less than or equal to IAI.

The system control statements accompanying a CYBER 200 FORTRAN program must include a call to the FORTRAN compiler. The parameters for this call optionally declare files for input and output, and optionally include instructions to the compiler to. (for example) output storage maps. Additional control statements are required to load and to execute the compiled program, and can be used to change at run time the file declarations made in a PROGRAM statement.

## **FORTRAN STATEMENT**

The FORTRAN system control statement is used to execute the CYBER 200 FORTRAN compiler. In the statement parameter descriptions that follow, underlining indicates the minimum number of characters that can be used in specifying the parameter.

Forms:

FORTRAN.

#### $\tt FORTRAN (NPUT=f_{1,BINARY}=f_{2/l2}$ , LIST=f3/l3/d3,0PTIONS=olist)

 $INPUT=f<sub>1</sub>$ 

Optional;  $f_1$  is the name of the file containing the FORTRAN source program to be compiled. When the parameter is omitted, the default file name INPUT is used.

 $BINARY=f_2/l_1$ 

Optional;  $f_2$  is the name of the file that is to receive the compiler-generated object modules. 12 is a specification of the length of  $f_2$ , and can be either an integer constant or a hexadecimal number prefixed with a #. 12 can be omitted along with the slash. When the entire parameter is omitted, the default file name BINARY is used. When  $1_2$  or the entire parameter is omitted, the default file length of 16 small pages is used.

 $LIST = f_3/l_3/d_3$ 

Optional;  $f_3$  is the name of the file that is to receive the<br>compiler-generated listings compiler-generated and program output. 13 is a specification of the length of  ${\rm f}_3$ . Like  ${\rm l}_2$ ,  ${\rm l}_3$  can be either an integer constant or a hexadecimal number prefixed with  $a \#$ .  $d_3$  is the routing disposition of f3 and must be PR (the line printer) or can be omitted (in which case no

OPTIONS=olist

routing is performed). 13 and d3 can occur in either order. When 13 is omitted, the default file size of 336 small pages is used. When the entire parameter is omitted, the default is OUTPUT.

Optional; olist is some logical combination of the compile option letters ABCEIKLMOR SUVYZ12, with the restriction that Y must not occur with any other option except L.<br>Default olist is B. When Default olist is  $B$ . O=olist is omitted, or when B is included in olist, the object file for the program is built. The object file is not built when the O=olist parameter without the B option appears in the parameter list for the<br>FORTRAN system control FORTRAN statement.

Alternative delimiters for the parameter list are a comma or blank instead of the left parenthesis, and a period instead of the right parenthesis. When communicating interactively with the system, the user can replace a period with a carriage return.

The FORTRAN system control statement parameters must be separated by commas or blanks. Partial parameter lists are acceptable, with default values used for the omitted parameters. The first form of the FORTRAN statement selects all defaults for the parameters. The  $I=$ ,  $B=$ , and  $L=$ parameters can be interchanged without consequence; the O= parameter must occur last.

The object and output files (specified by the  $B=$  and  $L=$ parameters of the FORTRAN system control statement) do not have to exist when the control statement is executed.<br>If the file does not exist, it is automatically created on a unit assigned by the operating system and with the length specified in the control statement. If the file does exist and has write access, it is automatically destroyed and recreated on the same unit with the length specified in the control statement. If the file does exist but does not have write access, a request is made to interactive users for permission to destroy the file. If permission is granted, the procedure followed is the same as for files that exist with write access. If permission is not granted, or if the user is in batch mode, the job is aborted.

When a compile option letter appears in the O=olist parameter, certain actions are performed during parameter, certain actions are performed during compilation that would not be performed otherwise. The L option is an exception in that the listing of the source program is inhibited rather than initiated by its appearance in olist.

## **A - ASSEMBLY LISTING**

An assembly listing of the object code can be placed in the output file by selecting the A option.

#### **B - BUILD OBJECT FILE**

An object file is required for the loading and execution of the FORTRAN program. A request that the file be built is made by selecting the B option.

#### **C - CROSS-REFERENCE LISTING**

All mentions in the source program to labels and symbolic names are listed in tabular form in the output file by selecting the C option.

### **E - EXTENDED BASIC BLOCK OPTIMIZATION**

The E<sub>.</sub>option selects optimization of extended basic blocks. This optimization involves compile-time computable result propagation, redundant code elimination, and instruction scheduling. The E option is included in the 0 option. The E option effectively selects options P, R, and I.

### I - **INSTRUCTION SCHEDULING**

The I option selects optimization of object instructions according to the results of a critical path analysis. The I option is included in the 0 and E options.

#### **K - 64-BIT COMPARE**

This option enables fullword (64-bit} integer compares for .EQ. and .NE. operators in logical IF statements. Otherwise, 48-bit compares are performed for the .EQ. and .NE. operations (integers are 48 bits).

#### **L - SOURCE LISTING SUPPRESSION**

The first part of the output file for a CYBER 200 FORTRAN program is normally the source program listing. This can be omitted from the file by selecting the L option.

### **M - MAP OF REGISTER FILE AND STORAGE ASSIGNMENTS**

A listing in the output file of all variables, constants, externals, arrays, and descriptors, along with a map of the contents of the register file, is produced when the M option is selected.

### **0 - OPTIMIZATION**

The 0 option selects all available optimization of scalar object code. More efficient object code is produced at the expense of increased compilation time. The 0 option effectively selects options Z, E, R, I, and P.

#### **P - PROPAGATION**

The P option selects compile-time-computable result propagation.

### **R - REDUNDANT CODE ELIMINATION**

The R option selects elimination of redundant code. The R option is included in the 0 and E options.

#### **S - SUPPRESS DEBUG SYMBOL TABLE CREATION**

The effect of this option is to suppress generation in the binary output of a debug symbol table for each program unit. The symbol table makes it possible for the<br>system-provided debugging utility DEBUG to recognize names in the FORTRAN program and for a FORTRAN run-time routine to identify the source line in a user routine at which a run-time error occurred. The user must not select this option if DEBUG is going to have to interpret variables, names, and symbolic addresses; if only absolute addresses will be used in commands to DEBUG, the S option can be selected.

#### **U - UNSAFE VECTORIZATION**

The U option enables unsafe vectorization of certain DO loops. If the terminal value of a DO loop is variable and the loop contains any references to dummy arrays, the compiler cannot determine the number of iterations of the loop. Vectorization of such loops is considered unsafe because the loop count might exceed 65535, which is the maximum length of a vector. If a DO loop contains an assignment statement that has an equivalenced data element on the left side, the loop can be vectorized only if the U compile option is selected.

## **V - VECTORIZATION AND AUTOMATIC RECOGNITION OF STACKLIB LOOPS**

Vectorization of certain CYBER 200 FORTRAN language constructs and automatic recognition and conversion of certain DO loops into calls to a stacklib routine are requested with the V compile option. The language constructs that fall under these categories are described in section 11.

#### **Y - SYNTAX CHECK**

A partial compilation can be performed to check the syntax of a FORTRAN program and any resulting diagnostics can be produced by selecting the Y compile option. The Y option can appear alone or with the L or S options (such as LY or SY); all other option combinations using Y are invalid compile option lists and produce an error accompanied by a dayfile message.

messages: compiler failure messages, compilation error messages, run-time error messages, and vectorizer messages. This appendix describes the four groups of diagnostic messages.

## **COMPILER FAILURE AND COMPILATION ERRORS**

Compiler failure messages are messages generated because of compiler failure. Compilation error messages are messages generated because of errors in the program. The serioumess of the error is indicated by the error type.

### **COMPILER FAILURE**

Error messages produced when the compiler fails are listed in table B-1. The compiler failure error type is:

A (abort) Compilation was terminated because of compiler failure. The return code is 8  $(RC=8)$ 

#### **COMPILATION ERRORS**

Error messages produced when the compiler detects errors in the source program are listed in table B-2. Some of the error numbers have no messages currently assigned to them. These error numbers are reserved for future use by CDC. Compilation error types are:

- W (warning) The statement in error was compiled. Compilation continued, but part of the statement might not have been processed. The return code is 4 (RC=4).
- F (fatal) The statement in error was not compiled. Object code generation is inhibited. The return code is 8 (RC=8).



#### TABLE B-1. COMPILER FAILURE MESSAGES







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## TABLE B-2. COMPILATION ERROR MESSAGES (Contd)

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#### **RETURN CODES**

The user has control over the execution of a batch job in that the user can determine whether to initiate error exit processing or to allow batch job processing to continue. The TV control statement allows a termination value to be entered with the program to be executed. The termination value is used to determine when error exit processing is to occur. All return codes having a value less than or equal to the termination value are ignored and job processing continues. All return codes having a value greater than the termination value cause error processing specified by the EXIT control statement to take place.

For example, a termination value of 8 would allow all warning and fatal errors to be ignored, andcause error exit processing to occur for abort errors. A termination value of 0 would trap all errors, including warning codes. The termination value control statement is discussed in the Operating System reference manual.

## **RUN-TIME ERRORS**

Error messages listed. in table B-3 are produced when error conditions are detected during the execution of a previously compiled program. Some of the error numbers have no messages currently assigned to them. currently assigned to

These error numbers are reserved for future use by CDC. The system error processor (SEP) can be called upon to change the attributes of certain run-time errors. Run-time error types are:

W (waming) Nonfatal error. A warning is issued and execution continues. The return code is 4 (RC=4).

F (fatal) Execution is terminated abnormally when this error condition<br>exists. The return code is 8 The return code is 8  $(RC=8)$ .

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Ĥ

C (catastrophic) Condition is nonalterable by SEP and not subject to user control,

other than replacement of the standard message. The return code is  $8$  (RC=8).

All errors having a warning classification can be made fatal. Those errors which are designated as fatal can be altered to warning level. Catastrophic errors cannot be altered to fatal or warning level; however, the standard message can be replaced.

Error messages for mathematical routines have the CYBER 200 FORTRAN library function name appended to the message. In like manner, input/output error messages have the file name appended to the message.

The form of a run-time error message is:

ERROR xxx IN subr AT LINE nn

## TABLE B-3. RUN-TIME ERRORS

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## TABLE B-3. RUN-TIME ERRORS (Contd)

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## TABLE B-3. RUN-TIME ERRORS (Contd)



## **GLOSSARY** c

Terms used in the main text of this manual are described in this section. The definitions give the general meanings of the terms. Precise definitions are given in the main text. Also, most general terms regarding computers and terms defined in the American National Standards documents regarding the FORTRAN language have been excluded.

Array -<br>An ordered set of variables identified by a single<br>symbolic name. Referencing a single element of an array requires the array name plus a subscript that specifies the element's position in the array.

Array Declarator -<br>Specifies the dimensions of an array. It consists of an array name followed by a parenthesized list of integer constants or simple integer variables that specify the largest value of each dimension.

ASCII Data - Characters, each of which has a standard internal representation. One byte (8 bits) is required for each character.

ASCII File - A type of file that can be manipulated with formatted READ statements, formatted WRITE statements, PRINT statements, and PUNCH statements.

#### Binary File -

A type of file that can be manipulated by unformatted input/output routines.

Bit Data -<br>A binary value represented in a FORTRAN program as<br>a binary number in the format B'bb ••• b' where each b is a 0 or a 1. Each 0 or 1 becomes a 0 bit or a 1 bit in the internal representation for the binary value.

Buffer Input/Output -<br>Input and output statements that cause data to be transferred between binary files and a buffer area in main memory.

#### Character Data -

An ASCII value represented in a FORTRAN program by a character string in the format 'cc... $c^{\dagger}$  where each c is in ASCII. Each character becomes a byte of ASCII data in the internal representation for the ASCII value.

Colon Notation - The notation used to express implied DO subscript expressions in a subarray. The colons separate the initial, terminal, and incrementation values for the implied DO.

Columnwise -<br>The ordering of the elements in an array declared in a DIMENSION, COMMON, or explicit type statement (the other ordering is rowwise). The succession of subscripts corresponding to the elements of a columnwise array is with the value of the leftmost subscript expression varying the fastest.

#### Compile Time -

The period of time during which the FORTRAN compiler is reading with the user's program and producing the relocatable module for the program. Compilation is initiated by the FORTRAN system control statement.

#### Conformability -

Determines whether two subarrays can occur in the same expression. Two subarrays are conformable if they contain the same number of implied DO subscripts and if corresponding implied DO subscript expressions are identical.

#### Control Vector -

A bit vector that controls the storing of values into a vector. The control vector elements are set to a configuration of Os and ls. Control vectors are used in WHERE statements, block WHERE structures, and some FORTRAN-supplied functions.

#### Controllee File -

A file that consists of object code generated by the loader. The loader builds a controllee file from relocatable object code produced by a compiler, plus relocatable object code of any externally-defined routines.

Data Element -<br>A constant, variable, array, or array element.

Data Flag Branch Manager (DFBM) -<br>A FORTRAN run-time and CYBER 200 library routine<br>that processes data flag branches when they occur in an executing program. A data flag branch is a hardware function of the CYBER 200 computers.

Data Flag Branch (DFB) Register - Part of the data flag branch hardware. It is a 64-bit register located in the CYBER 200 central processor.

Declaration - A specification statement that declares attributes of variables, arrays, or function names.

Defining -<br>Process whereby a variable or array element acquires a predictable or meaningful value. Definition can take place through data initialization, parameter place through data initialization, parameter association, DO statement execution, input statement execution, or assignment statement execution. Defining contrasts with naming and referencing.

#### Descriptor -

A pointer to a vector. In several FORTRAN forms, the descriptor can be used instead of the vector.

Dominance - A conventional data type hierarchy determining the data type of the result of expression evaluation. Dominated operands are converted during evaluation to the dominant type. The type complex dominates all other types, with types double-precision, real, and integer following in order of decreasing dominance.

Drop File - A file that is created and maintained for each executing program. Contains any modified pages of the program file, any free space attached, and any read-only data space defined to have temporary write access.

Dynamic Space - Virtual memory space available for allocation and deallocation at execution time. In particular, space for vectors can be assigned in the dynamic space area by using the descriptor ASSIGN statement.

Explicit Typing - Specification of the data type of a variable or array by means of one of the explicit type statements (the INTEGER, REAL, COMPLEX, DOUBLE PRECISION, BIT, CHARACTER, and LOGICAL statements). Explicit typing overrides any implicit typing.

External Function - A function that is defined outside of the program unit that references it. A reference to an external function generates code in the user's object program that causes control to transfer to the external function during program execution. External functions contrast with in-line functions.

File -

A collection of information that can be defined by output statements, or referenced by input statements. Depending on the type of output used to create it, a file can be either implicit or explicit.

First-Letter Rule - Default type association for data names according to the first letter of the name. Type assignment made is type integer to any name beginning with the letter I, J, K, L, M, or N, and type real to all others. The IMPLICIT statement is used to alter these defaults.

Floating-Point - Refers to the internal representation for real, double-precision, and complex data.

Generic Function -<br>A function whose result mode depends on the mode of the argument.

Hexadecimal Data - A value represented in a FORTRAN program as a hexadecimal number in the format  $X<sup>th</sup>$   $\ldots$  h' where each h is a hexadecimal digit (one of the digits 0 through 9 or one of the letters A through F). Each digit becomes the 4-bit binary equivalent in the internal representation for the value.

Implicit Typing - Specification of the data type of a variable or array by means of the first-letter rule for data names.

Index Vector - An integer vector whose elements are indexes into another vector. An index is an ordinal number indicating element position in a vector. Some of the FORTRAN-supplied functions use index vectors.

In-Line Function -

A type of predefined function. Referencing an in-line function causes the function's object code to be inserted directly into the relocatable object code of the user's program during compilation. In-line functions contrast with external functions.

Input  $-$ 

The name of the file read with FORTRAN READ statements that do not specify a unit number. To be used, INPUT must be declared in the PROGRAM statement or in the execution line.

Large Page -<br>A block of 65536 words in memory starting on a large page boundary. A loader call parameter can be used to tell the operating system that the specified modules are to be placed within a large page loaded on a large page boundary.

Loader -<br>A utility that links relocatable object modules, together with modules from user libraries or the system library as needed to satisfy external references. It then converts external references and relocatable addresses into the virtual address constants. Thus, relocatable modules are transformed into a virtual code controllee file with the (default) name of GO.

Logical Unit Number -<br>Integer between 1 and 99 associated with a file by<br>means of the PROGRAM statement declarations or<br>execution line declarations, and used to refer to the file when performing FORTRAN input/output.

Loop-Dependent -<br>Describes a variable whose value changes as the value of the control variable of a DO loop passes through the range specified in the DO statement. A specified in the  $D\overline{O}$ loop-dependent variable is defined within the range of the loop, while a loop-independent variable is defined (or could be defined with the same effect) outside the range of the loop.

Loop-Independent -<br>Describes a variable whose value remains constant within the range of a DO loop.

Naming -<br>Identifying data (or a procedure) without necessarily implying that its current value is to be made available (or, for procedures, that the procedure actions are to be made available) during the execution of the statement in which it is identified. Naming contrasts with referencing and defining.

Object Module -<br>The relocatable representation of a program unit created by compilation of the program unit.· Consists of object code.

Output - The name of the file to which all run-time error messages and records output with PRINT statements are written. WRITE statements can also be used to write on OUTPUT if OUTPUT is given a logical unit number in the PROGRAM statement.

Precedence - A conventional arithmetic, relational, and logical operator hierarchy determining the order in which<br>operations are performed during expression operations are performed during expression evaluation. Operator precedence in FORTRAN corresponds to the mathematical notion of the precedence of mathematical operations.

Predefined Function - FORTRAN-supplied code that performs common manipulations. Predefined functions can be in-line functions, external functions, or both in-line and external functions.

This appendix contains a list of the functions that are<br>available for reference for any CYBER 200 FORTRAN<br>program, as discussed in section 14. For each function,

table E-1 indicates what type of code (in-line, external, or both) is generated during compilation as a result of referencing the function.







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TABLE E-1. SUPPLIED FUNCTIONS (Contd)

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# **CYBER 200 FORTRAN STATEMENT SUMMARY F**

This appendix contains a summary of the statement forms described in the main text. Given are the entities that compose each statement; refer to the main text for the detailed specifications for these entities. Abbreviations used in this appendix are the following:



Brackets around an item indicates that the item is optional.

## **ASSIGNMENT STATEMENTS**



**And International** 

# **FLOW CONTROL STATEMENTS**



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