


## COMPUTER SYSTEMS META/MASTER REFERENCE MANUAL



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Address comments concerning this manual to:

Control Data Corporation Software Documentation 4201 North Lexington Avenue St. Paul, Minnesota 55112
or use Comment Sheet in the back of this manual.

## PREFACE

This manual is directed at programmers using the $3300 / 3500$ Meta-Assembler. It discusses the principles, features, methods, rules, and techniques of producing a META language program.

The reader is assumed to be familiar with the CONTROL DATA ${ }^{\circledR} 3300$ Computer System or the CONTROL DATA ${ }^{\circledR} 3500$ Computer System. In addition, familiarity with the $3300 / 3500$ MASTER Multiprogramming Executive Operating System and the $3300 / 3500$ COMPASS Assembly Language is helpful.

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The $3300 / 3500 \mathrm{META} / \mathrm{MASTER}$, a Meta Assembler executing on a CONTROL DATA ${ }^{\circledR} 3300$ Computer System or CONTROL DATA ${ }^{(®)} 3500$ Computer System, provides a versatile, extensive, and self-extending language for directing the generation of object code.
1.1 FEATURES

Using Meta-Assembler (META), the programmer can choose a relocatable binary output format acceptable for loading and executing under the 3300/ 3500 MASTER Multiprogramming Executive Operating System, or define output as a byte stream not restricted to a 24 -bit object word. MetaAssembler is an ideal language in which to code compilers and assemblers, or to produce code for an alternate computer system. The object computer, real or simulated, may have a word size up to 48 bits.

Source statements to META include directives that control the assembler in much the same way machine language instructions control the computer, procedure definitions and references, and function definitions and references.

The Meta-Assembler language allows simple, brief notation, nested functions and procedures, and complex expressions involving sets.

Procedures and functions provide extensive parameterization of source statements. For example, META includes standard procedures for the $3300 / 3500$ machine language instructions and for generating equivalent code. While these mnemonics resemble the $3300 / 3500$ COMPASS repertoire, differences in syntax and in notation used for operand fields and modifiers cause incompatibilities between the two languages. In addition, META does not recognize COMPASS macros, most pseudo instructions, or numeric operation codes. For example, the representation of an octal number in the 3300 COMPASS language is a string of octal digits followed by the letter B. The representation of an octal mumber in the META language is the letter $O$ followed by a string of octal digits enclosed in apostrophes.

A complete list of $3300 / 3500$ mnemonic instructions is given in appendix B .
A group of Meta-Assembler directives makes it possible for the programmer to assign his program and data to as many as 15 relocatable control sections, as well as one absolute control section. The assembler main-
tains a location counter for each section so that data locations within a control section are relative to the beginning of that section. The programmer can increment these counters by words or bytes. (He can also: define the size of words and bytes.)

## 1.2 <br> HARDWARE CONFIGURATION

The requirements for executing Meta-Assembler on the 3300/3500 are the minimum requirements for executing MASTER.

MASTER minimum core memory (about 16 K should be available for META)
One CONTROL DATA ${ }^{\circledR} 3304$ or 3504 Processor
One CONTROL DATA ${ }^{\circledR} 3311$ Multiprogramming Module ( 3300 only)
One CONTROL DATA ${ }^{\circledR} 405$ Card Reader and buffered controller
One CONTROL DATA ${ }^{\circledR} 501,505$, or 3254 Line Printer and buffered controller
One CONTROL DATA ${ }^{\circledR} 415$ Card Punch and buffered controller
Two CONTROL DATA ${ }^{\circledR} 3306$ or 3307 (3507) Communication (Data) Channels
2.5 million words ( 10 million characters) of mass storage. $90 \mathrm{~K}-100 \mathrm{~K}$ words or about 9 scratch segments of mass storage should be available to META for its temporary files.

## 2.1

CHARACTER SET

Programs written for Meta-Assembler may use alphabetic characters A-Z, numeric characters $0-9$, blank spaces, and the special characters listed below.

| $+$ | plus | ' | apostrophe |
| :---: | :---: | :---: | :---: |
| - | minus | $\leq$ | less than or equal |
| * | multiply | $\geq$ | greater than or equal |
| / | divide | [ | left bracket |
| $=$ | equal | ] | right bracket |
| $<$ | less than | $\dagger$ | decimal exponent |
| > | greater than | $\downarrow$ | binary exponent |
| - | period | 7 | NOT |
| , | comma | ; | semicolon |
| $($ | left parenthesis | $\rightarrow$ | right arrow |
| ) | right parenthesis | = | identity |
| \% | percent | : | colon |
| \$ | dollar ${ }^{\text {c }}$ | V | OR |
|  |  | $\wedge$ | AND |

The relationship of these characters to printer graphic characters, internal octal codes, and card codes is shown in appendix A. Characters that have special significance as operators are given in Table 2-1.

## 2.2 <br> STATEMENT FORMAT

> A Meta-Assembler statement consists of a label field, a command field, an operand field, and a comments field. Each field is terminated by two or more consecutive blanks.

## Format:

Leabel command operand comments ...

Statements can begin in character position 1 and continue through character position 71. A semicolon in character position 72 indicates card continuation. Any information beyond character position 72 is not interpreted by Meta-Assembler but does appear on the assembly listing. Thus, columns 72-80 can be used for sequencing.

Within a field, a single blank can separate elementary items, operators, and delimiters. A blank is optional for separating a symbolic operator, such as ${ }^{* *}$, from its operands, but is required for separating a mnemonic operator (AND) from its operands.
2.2.1

LABEL FIELD

The label field begins in character position 1 or 2, and is terminated by two consecutive blanks. If character positions 1 and 2 are blank, the statement has no label.

A label field may contain a symbol, set element reference, or SYM attribute function reference (section 6.5). A set element reference is legal only with an RDEF directive (section 4.3.2).

The definition of a symbol in a label field depends on the content of the command field. Throughout this manual, unless stated otherwise, a label field symbol is optional and, if present, is the value of the control counter prior to processing the command field. This value is either a word address or a byte address, depending on the command (section 3.1).

### 2.2.2

COMMAND FIELD

The command field of a statement begins with the first nonblank character following the label field and is terminated by two consecutive blanks. If the label field is blank, the command may start in character position 3.

The following are legal command field entries.

## An assembler directive

A mnemonic machine instruction code followed by its modifiers which are separated from the instruction code by a comma

The name of a previously defined FORM
The name of a previously defined procedure which may be followed by a set separated from the procedure name by a comma

A SYM attribute function reference

### 2.2.3

OPERAND FIELD

The command field entry of a Meta-Assembler statement determines if an entry is required in the operand field. If present, the operand field begins with the first nonblank character following the command field. The overand field contains either an expression or a set that supplies information for the command field. For sets, see section 2.7.

Two consecutive blanks terminate the operand field.
2.2.4

COMMENTS FIELD

### 2.2.5

STATEMENT
CONTINUATION
Normally, character position 71 terminates a source statement that has not yet terminated. However, a line of code that cannot be contained in the first 71 character positions can be continued to the next line by placing a semicolon in character position 72 and continuing the field at character position 2 of the next line; character position 1 is ignored.

Any character other than a semicolon in character position 72 is ignored.

### 2.2.6

EXAMPLES
The following line contains all four fields.


The following line has a blank label field and does not contain comments.

command operand
The following line is continued.


The following line contains a command and a comment.
END *C

The following line is a comment line.


The following line is not continued; character position 71 terminates the operand field.


## 2.3 <br> ELEMENTARY <br> ITEMS

The basic representation of data for META are elementary items. An elementary item is self-defining|and its meaning is immediately obvious; no additional information is needed for its interpretation. Meta-Assembler recognizes the following as elementary items.

Delimiters
Decimal integers
Binary coded decimal (BCD) integers
Octal integers

Floating-point real numbers
BCD character strings, left or right adjusted
ASCII character strings
Arithmetic and logical operators (symbolic and mnemonic)
2.3.1

DELIMITERS
2.3.2

DECIMAL INTEGER

META recognizes the following characters as delimiters.

| Comma | Delimits subfields of a source statement field, <br> elements of a set or subset, and subscripts of a set <br> element reference. |
| :--- | :--- |
| Parentheses $\quad$Enclose and delimit function arguments and nested <br> expressions. |  |
| Brackets | Enclose and delimit nested subsets and set element <br> references. |
| Blank $\quad$Separates elementary items for visual clarity or de- <br> limits them when required. |  |
| Apostrophes $\quad$Terminate fields of a source statement. |  |
| Enclose and delimit character and numeric strings. <br> Within a character string, only the single apostrophe <br> is a delimiter; any other delimiter is accepted as a <br> valid character in the string. A pair of apostrophes <br> signifies a valid BCD or ASCII apostrophe. |  |

A decimal integer is a string of numeric characters from the character set 0-9. Meta-Assembler converts a decimal integer to its binary equivalent. If the resulting binary number exceeds 48 bits, META truncates it with the loss of the most significant bits and sets an error flag. META also sets an error flag and truncates the resulting binary number if it exceeds a specified field size during data generation.

Examples:
2.3.3

BCD INTEGER

To specify a BCD integer, write the letter D followed by a string of one to eight numeric characters from the character set 0-9, enclosed in apostrophes.

Examples:

D'078'
D'123'
-D'123'

A BCD integer is not converted to its binary equivalent, but is represented as a string of 6 -bit $B C D$ characters (appendix A). If the number of $B C D$ characters is greater than eight, truncation causes loss of the most significant characters and META sets an error flag. During data generation, if the field into which the integer is to be placed is too small, META truncates the most significant characters and sets an error flag (section 4.5).

If during data generation, the field size is greater than required, the BCD integer is right adjusted with leading zeros.

Expressions containing BCD integers are evaluated using 6-bit BCD arithmetic. The sign of a $B C D$ integer is placed in the left bit of the rightmost digit of the number.

Examples:

> D'123'
-D'123'

| 00 | 01 | 02 | 03 |
| :--- | :--- | :--- | :--- |


| 00 | 01 | 02 | 43 |
| :--- | :--- | :--- | :--- |

An octal integer is noted with the letter $O$ followed by a string of numeric characters from the character set 0-7 enclosed in apostrophes.

Examples:
O'77'
O'123'

Meta-Assembler converts an octal integer to its binary equivalent. If the resulting binary number exceeds 48 bits, truncation causes loss of the most significant bits and META sets an error flag. If during data generation, the field into which the integer is to be placed is too small, META truncates the most significant digits and sets an error flag (section 4.5).

### 2.3.5 <br> REAL NUMBER

If the number of characters in a $B C D$ string exceeds eight, truncation causes loss of the leftmost characters and META sets an error flag. During data generation, if the field into which the character string is to be placed is too small, META truncates the leftmost characters and sets an error flag.

If the field size is greater than that required to hold a left-adjusted string, the string used in data generation is left adjusted with trailing blanks. If the field size is not a multiple of 6 bits, the extraneous bits are on the left and are 0. The remainder of the field is used for characters and is blank filled.

Example:
$L^{\prime} A B^{\prime}$ is stored in 21-bit field as


A right-adjusted string used in data generation is right adjusted with leading zeros if the field size is greater than that required to hold the string.

An ASCII character string is written as the letter A followed by an apostrophe, a string of one to six ASCII characters (appendix A), and an apostrophe. Each ASCII character occupies eight bits.

Because an apostrophe is a delimiter, an apostrophe as a character in the string is represented as two consecutive apostrophes.

Example:

$$
\begin{array}{ll}
A^{\prime} A B^{\prime} ' C D ' & \text { A string of five characters } A B^{\prime} C D \\
A^{\prime \prime \prime} A B C ' \prime & \text { A string of five characters 'ABC' }
\end{array}
$$

If the number of characters exceeds six, truncation causes loss of the leftmost characters and META sets an error flag. During data generation, if the field into which the character string is to be placed is too small, META truncates the leftmost characters and sets an error flag.

An ASCII string used in data generation is right adjusted with leading zeros if the field size is greater than that required to hold the string.
2.3.8

OPERATORS

The following table summarizes legal operators and their hierarchies and meanings.

Table 2-1. Legal Operators

| Operator | Alternate Mnemonic | Meaning | Hierarchy |
| :---: | :---: | :---: | :---: |
|  |  | Unary plus Unary minus | 1 |
| $1$ | $\begin{gathered} \text { DS } \\ \text { BS } \end{gathered}$ | Decimal scaling Binary scaling | 2 |
|  |  | Arithmetic product Arithmetic quotient | 3 |
|  |  | Arithmetic addition Arithmetic subtraction | 4 |
|  | LT EQ <br> NE <br> GT <br> LE <br> GE | Less than (compare) <br> Equal (compare) <br> Not equal (compare) <br> Greater than (compare) <br> Less than or equal (compare) <br> Greater than or equal (compare) | 5 |
| ** | AND | Logical product (AND) | 6 |
| $++$ | $\begin{aligned} & \mathrm{XOR} \\ & \mathrm{OR} \end{aligned}$ | Logical difference (exclusive OR) Logical addition (inclusive OR) | 7 |
| = |  | Unary equal; 1-word literal <br> Unary double equal; 2-word literal | 8 |

2.4 SYMBOLS

## 2.5

LOCATION COUNTERS

A symbol is an alphabetic character from the set A-Z followed by 0-11 alphabetic or numeric characters from the sets A-Z, 0-9.

Examples:

Legal Symbols
P
R3
PROGRAM

## Illegal Symbols

5A
ST\$RT
ABC-1

A unique location counter is associated with each of the 16 control sections available under Meta-Assembler. META interprets a reference to a control section name as a reference to the current value of the location counter (a word address) within that control section.

In addition, META interprets the character $\$$ as the value of the current location counter, a word address, prior to processing the line containing $\$$. Location counters are discussed in detail in section 4.4.

## 2.6

EXPRESSIONS
A combination of one or more elementary items, symbols, set element references, or function references makes up an expression. The programmer can form subexpressions by using parentheses in the normal role of arithmetic grouping. Thus an expression may contain subexpressions which in turn are made up of operators and other subexpressions or elementary items.

## Examples:

$$
\begin{aligned}
& \$ \\
& A+2 \\
& (A+2) * B
\end{aligned}
$$

Rules for evaluating expressions are:
Expressions are evaluated left to right with lower numbered hierarchies evaluated first.

Parenthetical subexpressions are expanded from the inside and are performed first.

Operators of equal hierarchy are evaluated left to right.
If a mnemonic operator is used in lieu of a special symbol (e.g., DS instead of 1 ), it must be preceded and followed by a single blank.

The value of a compare operation is 1 if the expression is true, 0 if it is false.

For the < or LT and the > or GT operators, 0 is greater than -0. For the $\leq$ or LE, the $7=$ or $N E$, and the $=$ or EQ operators, 0 is equal to -0 .

In expressions used for data generation, META performs the arithmetic operation and places the vulue in the specified field. If the resultant value exceeds the specified field size, META truncates the most significant bits and flags the error.

Examples:

Expression
$A+B>C$
$\mathrm{A}+\mathrm{B} * \mathrm{C}$
$(\mathrm{A}+\mathrm{B}) * \mathrm{C}$
$(\mathrm{A}<\mathrm{B})++(\mathrm{C}>\mathrm{D})$

## Evaluation

Add A to B ; compare the result to C .
Multiply B by C; add A to the product.
Add A to B; multiply the sum by $C$.
Compare $A$ to $B$; compare $C$ to $D$; perform a logical OR on the two subexpressions. If either or both inequalities are true, the value is 1 ; if both are false, the value is 0 .

If an expression contains relocatable symbolic addresses, its value must be relative to a single location counter, or not related to a location counter and thus nonrelocatable.

## Examples:

In the following examples, $\mathrm{P}_{\mathrm{i}}, \mathrm{D}_{\mathrm{i}}$, and $\mathrm{C}_{\mathrm{i}}$ refer to relocatable addresses in the program, data, and common areas.

The following are relocatable addresses.
D +1

Subtracting one relocatable address from another in the same program control section produces an absolute nonrelocatable result.

$$
\begin{aligned}
& P_{1}-P_{2} \\
& -C_{1}+C_{2} \\
& D_{1}-P_{1}+P_{2}-D_{2}+C_{1}-C_{2}
\end{aligned}
$$

The result of an expression cannot be the sum of two or more relocatable addresses in the same or different control sections. The following are illegal.

$$
\begin{array}{ll}
P_{1}+\left(P_{2}+5\right) & \text { Relocated twice relative to } P \\
P+D & \text { Relocated to both } P \text { and } D \\
-P_{1}-P_{2}+P_{3}-P_{4} & \text { Relocated twice relative to } P
\end{array}
$$

Single relocation or an absolute value can legally result from a complex expression.

$$
\begin{aligned}
& P_{1}-P_{2}+P_{3} \\
& -P_{1}+P_{2}-P_{3} \\
& P+D_{1}-\left(D_{2}+2\right)-C_{1}+\left(C_{2}-6\right)
\end{aligned}
$$

$$
\mathrm{P}_{1}+\left(\mathrm{P}_{2}+5\right)+\mathrm{D}_{1}-\left(\mathrm{D}_{2}+2\right)-\mathrm{C}_{1}+\left(\mathrm{C}_{2}+6\right) \quad \text { Result }+9 \text { is not relocatable }
$$

### 2.6.1

ATTRIBUTES

### 2.6.2

MODES OF
EXPRESSIONS

Single positive relocation
Single negative relocation
Result P-8 is single positive relocation
relative to any control section. relative to any control section.

A mode associated with each elementary item defines how META is to interpret the data when it performs an arithmetic operation on the item. MetaAssembler recognizes 11 modes accessible through the mode attribute function (section 6.2).

Expressions are evaluated using either integer, real, or binary-codeddecimal arithmetic. META permits mixed-mode arithmetic on real and integer values, converting the integer to a real value and performing the operation in floating-point arithmetic. The mode of the result is real. With any combination other than real and integer, if all elements of the expression are not of the same arithmetic type, META flags an error and sets the value of the expression to 0.

In arithmetic and relational expressions, META treats character strings and addresses that are not external as integers.

META performs logical operations on a bit-by-bit basis without regard to mode. The result of a logical or compare operation is in integer mode.

The following table shows legal combinations of operators and operands. For,,$+- *$, and /, interchanging the first two columns does not affect the result. The mode of the second value must not be external.

Table 2-2. Combinations of Operators

| Operator | Mode 1st Value | Mode 2nd Value | Mode of Result |
| :---: | :---: | :---: | :---: |
| 1 | Integer <br> Real <br> Decimal | Integer <br> Integer <br> Integer | Integer <br> Real <br> Decimal |
| $\downarrow$ | Integer Real | Integer <br> Integer | Integer <br> Real |
| +, - | Integer <br> Integer <br> Integer <br> Real <br> Decimal <br> Word Addr <br> Word Addr <br> Byte Addr <br> Ext Wrd Addr <br> Ext Byte Addr | Integer <br> Real <br> Word Addr <br> Real <br> Decimal <br> Word Addr <br> Byte Addr <br> Byte Addr <br> Integer <br> Integer | Integer <br> Real <br> Word Addr <br> Real <br> Decimal <br> Word Addr <br> Byte Addr <br> Byte Addr <br> Ext Wrd Addr $\dagger$ <br> Ext Byte Addr $\dagger$ |
| *, / | Integer <br> Integer <br> Real <br> Decimal | Integer <br> Real <br> Real <br> Decimal | Integer <br> Real <br> Real <br> Decimal |
| **, ++, -- | Any | Any | Integer |
| $>,=, 7 i=,<, \leq, \geq$ | Mode 1, 3, 5, 7, $9,11 \dagger \dagger$ <br> Real <br> Decimal | ```Modc 1, 3,5,7, 9,11\dagger\dagger Real Decimal``` | Integer <br> Integer <br> Integer |

$\dagger$ External word addresses and external byte addressed cannot be interchanged.
$\dagger \dagger$ Section 6.2

Scale factors, both decimal and binary, must be integer.

Examples:

| $1.5 * 3$ | Legal; value is 4.5 real. |
| :--- | :--- |
| $D^{\prime} 15^{\prime}+D^{\prime} 17^{\prime}$ | Legal; both items are decimal integers. |
| $D^{\prime} 15^{\prime}+1.5$ | Illegal; conflicting modes. |
| $1.5 \downarrow 2.5$ | Illegal; scaling factor is not an integer. |

2.6 .3

A literal is an expression beginning with an equal or a double equal sign depending on whether the value is to occupy one or two words.

Examples:

$$
\begin{array}{ll}
=O^{\prime} 77700077 \\
=\mathrm{A}+\mathrm{B}-\$ & ==1.2 \\
=1 & ==A^{\prime} \mathrm{ABCDEF}
\end{array}
$$

META places the value of the expression in a literal table. If the value exceeds the specified number of object computer words, META truncates it and flags the error. If the object computer word size is greater than 24 bits, use of a 2 -word literal causes truncation because the maximum precision allowed is 48 bits. By using one or more LIT directives (section 4.4.5), the programmer can designate which control sections are to contain literal tables. If the program contains no LIT directive, the literal table is appended to the program section. The address of a literal is the address of the literal table entry relative to the beginning of the control section. Literals with identical expression values are entered into a single literal table only once.

An attempt to place a literal in a numbered common area is flagged as an error; numbered common cannot be preset.

A set is one or more set elements separated by commas. A set element is an expression, a set name, or a subset. A subset is a set enclosed in brackets.

The NSET directive (section 4.3.3) assigns a set name to a set. Set names can also be assigned through the PROC and FUNC directives (sections 5.1.1 and 5.1.2).

Examples:


A is a set of two elements.
$B$ is a set of two elements. The first element is an expression; the second is a set name.
k. NSET, $[1,2],[3,[3,2,33],$,
[6],$A \ldots$
C is a set of three alements. The first is a subset which is a set of two elements. The second element is a subset which is a set of two elements, the first of which is itself a subset. The third alemont is a set name.

In the preceding example, the first element of set $C$ could have been written as A.

To refer to a set element, write the name of the set followed by a left bracket and one or more expressions separated by commas and a rightbracket. The values of expressions represent the ordinal location of the set element referenced. From left to right, they represent the level of the elemont in a set containing subsets. To refer to an entire set, write the name of the set.

If the reference is to a nonexistent element, META uses zero.

Example:

The symbol A is defined as the set $5, \mathrm{C},[9,[3,4]]$. The set has three elements. The third element [9, [3, 4] ] contains two elements, the second of which also contains two elements [3,4].

| Reference | Element | Value |
| :---: | :---: | :---: |
| A | All | ¢́, C, [9, [3, 4]] |
| A[1] | First element of A | 5 |
| A[2] | Second element of A | C |
| A[3] | Third element of $A$ | 9, [3,4] |
| $\mathrm{A}[3,1]$ | First element of subset of third element of $A$ | 9 |
| A $[3,2]$ | Second element of subset of third element of $A$ | 3,4 |
| $\mathrm{A}[3,2,1]$ | First element of subset of second element of subset of third element of $A$ | 3 |
| A $[15,33]$ | Nonexistent element | 0 |

In the preceding example, if $C$ is a set name for a set consisting of the list elements 7, 8, 6, elements of $C$ could be referred to as follows:

| Reference | Element | Value |
| :--- | :--- | :---: |
| $A[2,1]$ or $C[1]$ | First element of $C$ | 7 |
| $A[2,2]$ or $C[2]$ | Second element of $C$ | 8 |
| $A[2,3]$ or $C[3]$ | Third element of $C$ | 6 |

The Meta-Assembler maintains information about a set and its elements together with the symbol defining the set. The programmer can access this information for use by the assembler through attribute function references. For example, the NUM attribute function (section 6.3) supplies the number of elements in the set.

Meta-Assembler provides location control by making available one absolute and up to 15 relocatable control sections, each with an associated location counter. The counters can be incremented in word or byte increments.

## 3.1 RELOCATABLE ADDRESSES

A relocatable address is either a word address (mode 9) or a byte address (mode 11). Mode is specified implicitly by the directive. Word-oriented directives cause definition of relocatable word addresses. Byte-oriented directives cause definition of relocatable byte addresses.

A label field symbol is a word address for the following directives.

| RES | TEXT |
| :--- | :--- |
| GEN | TEXTA |
| GEND |  |

Also, literals and control section names are word addresses. A reference to a control section name returns, as a word address, the current value of the location counter in use prior to processing the line. Use of the $\$$ returns the word value of the current location counter prior to processing of the line.

For the following directives, a label field symbol is a byte address.

| NOLIST | LIT | DETAIL | ENTRY |
| :--- | :--- | :--- | :--- |
| LIST | RESB | SECP | EXT |
| SPACING | GENB | SECD | GOTO |
| EJECT | TITLE | SECA | ENDS |
| ORG | BRIEF | TEXTC | TREF |
|  |  |  | LIBS |

A label field symbol on a FORM reference line or a procedure reference line is a byte address. This means that a mnemonic instruction (which is actually a procedure reference) does not cause the counter to be rounded to the nearest word address.

A word-oriented directive that follows a byte-oriented address causes the control counter for the section to be rounded up to the nearest word address. A byte-oriented directive always uses the next available byte.

Use of the $\$$ returns the word value of the current location counter prior to processing the line.

Examples:


RES in the first example causes the control counter to be rounded up to the next word boundary prior to definition of the symbol B1. The control counter is not rounded up in the second example.


RES and RESB are discussed in sections 4.4.6 and 4.4.7.

Location counters are designated $0-15$, corresponding to the 16 control sections a programmer can define using SECA, SECD, and SECP' directives (chapter 4).

Location counter 0 is reserved for the absolute control section (defined by SECA).

Location counter 1 is reserved for the first program control directive. If the program has no SECP directive defining a program control section name, location counter 1 is still used for the program.

Location counters 2 through 15 are used for either program control sections or data control sections. As META encounters each SECP or SECD directive, it assigns the next available location counter.

Directives control the operation of Meta-Assembler much the same as machine mnemonic codes direct the computer. The programmer can use directives to:

- Control the content and format of the Meta-Assembler listing.
- Define word and byte size when the object computer is not a 3300 or 3500.
- Define a symbol and assign it a value or set of values.
- Assign up to 15 relocatable and one absolute location counters for address assignment.
- Generate code to be loaded and executed on the object computer.
- Specify field sizes for the object code.
- Specify that certain symbols are entry points to separately assembled subprograms, or that symbols used within the current subprogram are external to it.
- Repeat or skip source statements conditionally.
- Terminate assembly of a subprogram or group of subprograms.
- Define a procedure and assign it one or more names for subsequent reference.
- Define a function and assign it one or more names.


## 4.1

LISTING CONTROL

Through listing control directives, the programmer suppresses portions of the output listing, selects spacing, places a title at the top of any page of the listing, and requests the level of detail he wants to appear in the listing. For all listing control directives, a label is optional; if present, it has the current location counter value.
4.1.1

NOLIST
NOLIST suppresses generation of the output listing until the assembler en- • counters a LIST directive. The NOLIST line is suppressed from the listing. Format:

Lobere NoLIs.T comments. -

LIST causes resumption of the normal assembler listing following a NOLIST directive. LIST appears on the output listing.

Format:

## Rabel LIST, comments..

4.1 .3

SPACING
SPACING allows the programmer to select single, double, or triple spacing in the output listing.

Format:
leabel SPACIMG exp comments..
$\exp \quad$ Expression evaluated as 1, 2, or 3 corresponding to single, double or triple spacing, respectively. Otherwise, directive is ignored.

The specified spacing remains in effect until another SPACING directive appears. If no SPACING directives appear in a program, the listing is single spaced.

EJECT terminates the current page of the output listing and causes listing to resume at the top of the following page. EJECT is printed as the first line of the next page.

Format:
Rabel EJECT comments. .-

If EJECT is already the first line of a page, it is printed but has no other effect.
4.1.5

TITLE

TITLE causes the current page to be ejected and the TITLE directive line itself to be printed on the first line of the new page. Until another TITLE directive is processed, all succeeding pages begin with this title.

Format:
label TITLE 'Gharacterstringícomments con
character string $\quad 1-56$ characters that appear as title at top of each page of output listing (section 8.1)
4.1.6

BRIEF
4.1.7 DETAIL

BRIEF causes listing of source lines and lines of code generated by data generating directives only. BRIEF remains in effect until a DETAIL direclive occurs. The default mode of listing is BRIEF.

Format:
Kabeil BRIEE comments -

DETAIL causes listing of all lines of code other than library procedure definations in subsequent LIBS directives (section 5.1.6) and causes listing of procedure expansions. DETAIL remains in effect until a BRIEF directive is processed. A NOLIST directive takes precedence over a DETAIL direclive.

Format:
Rabeili, DETAIL comment is.

## 4.2

OBJECT MACHINE DEFINITION (UNIT)

## 4.3 <br> SYMBOL AND SET DEFINITION

The Meta-Assembler running on a Control Data 3300 or 3500 Computer Sysfem to assemble programs for other computers must have certain information about the object computer to generate the proper binary information. The UNIT directive defines the byte size and word size of the object computer. Word size of the object computer must not be less than 8 bits nor greater than 48 bits.

Format:

$$
\text { Label UNII } I_{1} \text { exploit, exp } p_{2} \text { comments. }
$$

| label | Optional |
| :--- | :--- |
| $\exp _{1}$ | Evaluatable nonrelocatable expression defining the byte <br> size of the object computer in bits. During assembly, the <br> location counter is incremented by 1 for each $\exp _{1}$ bits. |
| $\exp _{2}$ | Evaluatable nonrelocatable expression specifying the nom- <br> der of bytes per word. |

In the absence of a UNIT directive, META uses the host computer unit size of 6 bits per byte and 4 bytes per word. Binary output is in the form acceptable to the $3300 / 3500$ MASTER relocatable loader.

UNIT, if used, must precede all lines of code other than listing control directives and comment cards. Use of UNIT causes binary output to be in the alternate form (appendix C).
*

A symbol that appears in the label field of an EQU or RDEF directive has a defined value. Whenever the symbol is used in an expression, this defined value rather than the address of the symbol is used in evaluating the expression.

A symbol that appears in the label field of an NSET directive or the label field, command field (as a modifier), or the operand field of a PROC or FUNC directive (chapter 5), becomes the set name for a list of set elements. Whenever the subscripted set name is used in an expression, the value of the set element is used in evaluating the expression.

EQU assigns the value and attributes of the operand field expression to the label field symbol.

Format:
Label Fou exp commentis.-1

The label field must contain a symbol. A symbol defined by EQU cannot be redefined later in the program.

Example:
HB. ERS. 7.1
4.3.2 REF

RDEF assigns the value and attributes of the operand field expression to the symbol or set element named in the label field.

Format:

> Label RDEF exp comments.

The label field must contain a symbol or a set element reference. The value and attributes assigned to this symbol or set element remain in effect until an RDEF with an identical label field symbol or set element is processed or until an RPT (section 4.7.1) with an identical label field symbol is processed. If the operand field is blank, the symbol or set element has a value of 0 .

Example:


A has value of 15
$B$ has value of current location counter
$C$ has value $A+3$, or 18
Illegal; A is doubly defined
Illegal; B is doubly defined
Legal; C changed from 18 to 20
D has value 0
Illegal; D may not be redefined
Define set E
Redefine element two of set E

NSET assigns the label field symbol as the set name of the operand field set. The label field must contain a symbol which is the name by which the set or set elements can be referenced. If the operand field is blank, the set consists of one element which has a value of 0 .

Format:

## Label NSET set comments.

Example:


A is a set of three elements.
$B$ is a set of two elements, the second of which is a set of two edemints.

C is a set of two elements, the second of which is a set of three elements.
$D$ is a set of one zero element.
A[2] is redefined to be 9 in the last line of the examples; thus the final set A is defined as though the following had been written.

$$
A, N S E T \quad 3,9, \$
$$

An entire set can be redefined through use of NSET.
Example:


A forward reference is a reference to a symbol or set element before it is defined. The Meta-Assembler processes forward references in two passes. On the first pass, a reference to a symbol before it is defined is not given a value; a reference to the symbol after it is defined is given the most recently assigned value. On the second pass, the forward reference is given the most recent value assigned.

An expression cannot contain a forward reference if:

1. The value affects location counting.
2. The undefined symbol is defined subsequently by an EQU directive that contains a second forward reference.
3. The undefined symbol or set element is not defined subsequently.
4. The expression is not evaluatable.

A forward reference to a symbol or set element redefined subsequently by RDEF or NSET directives that contain forward references yields the final value assigned to the symbol or set element.

Examples:

Legal use of forward reference


First Pass Second Pass
A undefined
$A=2.5$
$B=2.5$
$B=2.5$
$A A$ undefined $\quad A A=7$
$X X$ undefined $\quad X X=2$
BB undefined $\quad \mathrm{BB}=2$
$Y Y=2 \quad \bullet Y Y=2$
$X X=7 \quad X X=7$
$\mathrm{AA}[2]$ undefined $\mathrm{AA}[2]=9$
$\mathrm{B} B=8 \quad \mathrm{BB}=8$
$\mathrm{AA}[1]=1 ; \mathrm{AA}[2] \quad \mathrm{AA}[1]=1 ; \mathrm{AA}[2]$
$=8 ; \mathrm{AA}[3]=7 \quad=8 ; \mid \mathrm{AA}[3]=7$
$\mathrm{AA}[2]=9$

Replacing the last line of the previous example with the following would achieve the same result.
|AA $\ldots \ldots$

Illegal Forward References:
 $1, R E S_{1} A A, \ldots, 1$,


GEN directive contains forward reference to AA which is defined by an EQU containing a forward reference.

Value affecting location counting must be defined on first pass. RES affects location counting.

Value affecting location counting must be defined on first pass.

RES affects location counting.

Value affecting location counting must be defined on first pass.

RES affects location counting.
4.4

LOCATION CONTROL

Meta-Assembler provides one absolute and 15 relocatable control sections, each of which has an associated location counter. Any program can use one or more control sections.

Meta-Assembler directives described in this section assign names to control sections and address values to location counters.
4.4.1 SEC

The first SECP directive defines a program control section.

Format:

> label SEC, P. sym comments . -

$$
\begin{aligned}
& \text { label } \begin{array}{l}
\text { Optional; if present, label has the value of the location } \\
\text { counter after the SECP directive is processed. } \\
\text { sym } \\
\text { 1-8 character name of program control section (subpro- } \\
\text { gram name). A reference to sym later in the program } \\
\text { returns the current value of the associated location } \\
\text { counter. }
\end{array} . \quad \begin{array}{l}
\text {. }
\end{array} \text {. } \quad \text {. }
\end{aligned}
$$

After the first SECP naming a specific sym, successive SECP directives using this sym indicate that the code that follows is an extension of the presviously declared program control section. A programmer coding for MASTER may use only one program control section; any additional SECP directive naming a new sym is flagged with an informative $\mathbf{D}$ error.
4.4.2

SECD

The first SECD in a program defines a blank common, numbered common, or labeled common control section.

Format:

- Label SECR sym, exp commentisi i
label Optional; if present, label has the value of the location counter after the SECD directive is processed.
sym Optional; name of control section defined or referenced. zero or Control section defined or referenced blank

1-8-character symbol

1-4 decimal digits is zero or blank common.

Control section is labeled common block.

Control section is numbered common block.

For the $3300 / 3500$ relocatable output, if sym is blank or 0 , the block name is 1 MWUN for chapter 1 and 2 MWM for chapter 2. For other than the $3300 / 3500$, the block name depends on the item type (appendix C).

If sym is a symbol, a reference to the symbol later in the program returns the value of the associated location counter.
$\exp \quad$ Optional; if execution is under $3300 / 3500 \mathrm{MASTER}$, exp is an evaluatable expression with value 1 or 2 designating the chapter to which the section is assigned. If exp is absent, chapter one is assigned.

Each new sym on an SECD directive causes creation of a new control section starting at relative address 0 . If a sym appears on a subsequent SECD directive, exp is ignored and code following the subsequent $S E C D$ directive down to the next location control directive is an extension of the previously declared control section.
4.4.3

SECA
A program can have an absolute control section declared by a SECA directive.

Format:
Rabel SECA sym comments.

| label | Optional; if present, label has the value of the location counter after the SECA is processed. |
| :---: | :---: |
| sym | 1-8 character name of the absolute control section. A subsequent reference to sym returns the current value of the absolute location counter. |

Any SECA directive after the first one in a program indicates that the code following it is an extension of the originally defined absolute control section. If SECA is preceded by an ORG directive setting the absolute location counter, the code following the SECA extends the absolute control section.

SECA cannot be used when coding for MASTER.

ORG sets the specified control counter to a specified address.
Format:
Label aRG exp comments.-
label Optional; if present, the label has the value of the location counter after the ORG directive is processed.
$\exp \quad$ Evaluatable expression. The expression indicates the controd counter to be selected and the address to which it is to be set. Lines of code following ORG are placed in the contron section indicated.

Examples:
*


Defines program control section ALPHA.

Specifies labeled common block of COMM in chapter one.

Specifies numbered common block. Chapter one is implied.

Specifies resumption of program control section. (Here, ORG has the same effect as SECP ALPHA).

Specifies resumption of labeled common block COMM.

Location within COMM.
Selects absolute location counter and sets its value to 50 .

Specifies blank common. For the 3300/ 3500 the block name is $1 \wedge \wedge \wedge \wedge \wedge \wedge \wedge$; otherwise, the block name is 00000000 .

Selects the location counter for COMM and sets the location counter value to the value of $D$.

Specifies resumption of program controll section ALPHA.
4.4 .5

Format:
Label i LIT sym icommenitisi_1

| label | Optional; if present, label has the value of the location <br> counter. |
| :--- | :--- |
| sym | Name of a previously defined control section. . |

META places literals (section 2.6.3) in the control section specified by a LIT directive, regardless of which control section contains the reference, until it encounters another LIT directive designating a different control section for literals. In any given literal table, only one entry is made for identical literals. However, a literal table can have entries that duplicate entries in other literal tables. A literal results in the generation of object code.

In the absence of a LIT directive, literals are appended to the first program control section.
4.4.6

RES
RES adds the value of the expression in the operand field to the current loca- tion counter value as a word increment to reserve storage.

Format:
label, RES exp comments,
exp Evaluatable nonrelocatable expression (must not contain a forward symbolic reference or reference to an externally defined symbol).

Examples:


Increment location counter by two words. Increment location counter by ten words.
Decrement location counter by five words.
4.4.7 RES

RESB adds the value of the expression in the operand field to the current value of the location counter as a byte increment to reserve storage.

Format:

## label RESB exp comments -

$\exp \quad$ Evaluatable nonrelocatable expression by which to increment the counter.

Examples:


## 4.5 <br> DATA <br> GENERATION

Data generating directives define data formats and generate words or bytes of information to be loaded into the computer at execution time.
4.5.1

GEN
GEN places the values of expressions in the operand field set in successive words, one word for each expression.

Format:
Nobel GEN set comment is.
set Set of expressions to be generated. A set of sets is not permitted.

Examples:
For the following examples, the object computer word size is 24 bits.
A. GEN $5, C^{\prime} A B C D^{\prime}$...

Generate two words, the first containing 5, the second contraining the internal BCD representation of $A B C D$.


Results in the same values as the above.





Illegal; values must be single precision.

Legal; reference to external symbol.
4.5.2

GEND

GEND generates the values of expressions in the operand field set, two object computer words per expression. Maximum precision for a value is 48 bits. If the object computer word size exceeds 24 bits, META truncates the value to 48 bits and flags the error.

Format:
label GEND set comments i-
label Optional
set Set of expressions to be generated. A set of sets is not permitted.

Example:
In the following example, the object computer word size is 24 bits.
A. GENA $2.44,25, C^{\prime} A B C^{\prime}$..

The code generates six words. The first two words contain the floatingpoint representation of 2.4 . The next two contain the binary integer representation of 25 . The last two words contain the internal BCD representation of ABC right-justified with leading zeros.

GENB evaluates the values of expressions in the operand field set and places the values in successive bytes. If the value of an expression exceeds the byte size specified in the UNIT directive, META truncates the value to the byte size and flags the error.

## Format:

## Llobeil GENB set commentsi i

label Optional
set Set of expressions to be placed in successive bytes. A set of sets is not permitted.

## Example:

For the following example, the object computer byte size is 6 bits.
TOE GENB $5,9,63,14,-2 \ldots$
The above code generates five 6-bit bytes. The last byte contains the one's complement of -2 truncated to 6 bits (111101).

FORM defines a data format by specifying field sizes, left to right, in one or more object computer bytes.

## Format:

> Kabiel FARM set comments -

| label | Required; label is the name referring to FORM |
| :--- | :--- |
| set | A set of expressions, each of which defines a field size |
|  | in bits. A set of sets is not allowed. |

## Examples:

For the following examples, the object computer byte size is 6 bits and the object computer word size is 24 bits.

|  |  |
| :---: | :---: |
| NOPD $\ldots$, FORM, 24, | One field, four bytes |
| WORD , FORM, 4,8, | One field, eight bytes |
| CHARS, FSRM, $6, ~ 6 a, b, 6,1$ | Four fields, four bytes |
| $A D R \ldots F Q R M, Z_{y} / \mathbf{Z} \ldots \ldots$ | Two fields, four bytes |
| INST 1 FQRM, $I_{1 \ldots \ldots \ldots}$ | Four fields, four bytes |

To refer to a format defined by a FORM directive, place the label of the FORM directive line in the command field of a line. Supply a set of expressions, corresponding to the fields, in the operand field of the referencing line. A form reference generates code starting with the next available byte.

A label on the line referring to a FORM directive has the value of the location counter prior to processing the line. If a value exceeds the specified field size or if the field size exceeds 48 bits, high-order bits are truncated ,and an error flag is generated. For a negative value, the one's complement of the absolute value is used unless the value is in BCD decimal mode. For a BCD decimal value, the sign is inserted in the leftmost bit of the least significant character position of the field.

If the field contains a 6-bit character type value and the field size is not a multiple of 6 -bits, the characters are placed in the rightmost bits of the field with the leftmost extraneous bits zero.

References to FORM directives can be circular.
Examples:


Generates a single word with value $\$+3$ right justified in the 24-bit field.


Generates the 48-bit floating-point value of 1.59.


Generates one word containing the following octal value: $||21|| 77||17|| 12$


Generates a word with zero in the leftmost 7 bits and the byte value of the location counter in the rightmost 17 bits.


Generates a word with value 12 right adjusted in the leftmost 10 bits, zeros in the next 3 bits, and the current word address plus 2 in the rightmost 15 bits.

The following example illustrates circularity of forms.


Generates 2 bytes
Generates 2 bytes identical to last two

Generates 4 bytes filling last with zeros
Generates 5 bytes

The 4 bytes generated by G are:
The 5 bytes generated by G are:

| 04 10 20 40 <br> 06 14 30 61 |
| :--- |

In the following example, BCD characters XY are to be stored in a 19-bit field.


BCD characters $X, Y$, and blank are placed in the rightmost 18 bits of the field. The leftmost bit is 0 .

TEXT generates an integral number of object computer words containing the specified BCD character string.

Format:

$$
\text { Rabel TEX } T_{1} \text { 'striing'_ comments }
$$

The last word is padded with blanks as needed. If the object computer word size is not a multiple of 6 bits, as many characters as fit are placed in each word, right adjusted with upper bits zero.
4.5.6

TEXTC

TEXTC is identical to TEXT except that the BCD character string generated is placed in consecutive words without padding the last word.

Format:

$$
\text { Rabe } l_{1} \text {, TEXTTC 'string'_ comments. }
$$

TEXTC generates code starting with the next available byte.

## 4.6 <br> PROGRAM LINKING

TEXTA generates 8-bit ASCII characters in the same way TEXT generates BCD characters. Padding of the last word, if needed, is with the internal representation of ASCII blanks.

Format:

The directives ENTRY and EXT do not define symbols, but either classify symbols defined within the subprogram as being known outside the subprogram, or classify symbols referenced in a subprogram as being defined outside of the subprogram.
4.6.1

ENTRY
4.6.2 EXT

## 4.7

REPEAT
AND SKIP

The ENTRY directive specifies which symbols defined may be referenced by subprograms compiled or assembled independently. That is, ENTRY directives list entry points to the current subprogram.

Format:

sym $_{i} \quad$ Entry point symbols, 1-8 BCD characters

The EXT directive lists symbols which are defined as entry points in ingependently compiled or assembled subprograms, but for which references appear in the subprogram being assembled.

Format:

$\operatorname{sym}_{\mathrm{i}} \quad$ External symbols, $1-8 \mathrm{BCD}$ characters

Source statements can be processed repeatedly or skipped conditionally through use of the RPT and GOTO directives.
4.7.1

RPT
RPT specifies processing a portion of code a given number of times.

Format:
Rebel, RPT exp, linid comments..

| label | Optional; if present, the original value is 0 . The <br> value of the label is tested and incremented by 1 |
| :--- | :--- |
| prior to each processing of the lines of code, to a |  |
| final value that is the value of exp. |  |$\quad$| Absolute evaluatable nonrelocatable expression (con- |
| :--- |
| tains no forward or external references) indicating |
| the number of times the following lines are to be |
| processed. If exp is less than or equal to 0 , the |
| following lines are not processed and the RPT acts |
| as a skip. |

RPTs may be nested to a level of at least six and possibly more depending on available table space. Space not required for processing functions and procedures could be used for additional levels of RPTs. Processing of repeated statements is from innermost to outermost. Every inner RPT range must lie totally within the range of the next outer RPT.

The programmer can redefine the RPT label within the repeated statements to terminate a repetition prematurely.

Examples:
The following sequence generates a 10 -word table of even numbers, $0-18$. Because linid is absent, only one line is processed.


Generates one word for each value $0,2,4,6,8, \ldots, 18$

The following example illustrates two levels of repeats; the nested repeats produce 10 words.


| $\mathbf{Q}=1, \mathbf{R}=1$ | 2 | $\mathbf{Q}=3, \mathbf{R}=2$ | 5 |
| :--- | :--- | :--- | :--- |
| $\mathbf{Q}=1, \mathbf{R}=2$ | 3 | $\mathbf{Q}=4, \mathbf{R}=1$ | 5 |
| $\mathbf{Q}=2, \mathbf{R}=1$ | 3 | $\mathbf{Q}=4, \mathbf{R}=2$ | 6 |
| $\mathbf{Q}=2, \mathbf{R}=2$ | 4 | $\mathbf{Q}=5, \mathbf{R}=1$ | 6 |
| $\mathbf{Q}=3, \mathbf{R}=1$ | 4 | $\mathbf{Q}=5, \mathbf{R}=2$ | 7 |

In the following example, lines 5-8 are processed three times.


In this example, the elements of set $S$ are initially zero. On the first processing of lines $5-8, C$ is 1 , and $S[1]$ is redefined as $A+C-1$, or 4 . On the second RPT directive, the test $S[1]=B$ is not true ( 0 ); the GEN line is skipped. When lines $5-8$ are repeated, $C$ is 2 and $S[2]$ is redefined as 5 . The test $\mathrm{S}[2]=\mathrm{B}$ is true (1) so the GEN line is processed; it generates one word with a value of 20 . On the final iteration, C is $3, \mathrm{~S}[3]$ is redefined as 6 , and the test $\mathrm{S}[3]=\mathrm{B}$ is not true ( 0 ); the GEN statement is skipped. Without the use of repeats, this example would be:


GOTO specifies a conditional skip.
Format:
Rabel, GOTA exp, linid $d_{1}, \ldots$, linid ${ }_{n}$ comments
$\begin{array}{ll}\exp & \text { Evaluatable nonrelocatable expression } \\ \text { linid }_{i} & \text { Line identifiers defined as labels on lines following } \\ & \text { GO'TO. }\end{array}$ GOTO.

Expression exp is evaluated and used as an index to the list of line identifiers. The line containing the label identified by the indexed line identifier is the next line assembled. For example, if exp has value 2 , the second line identifier is the label of the next line to be assembled. If $0 \geq \exp >n$, where $n$ is the number of line identifiers, assembly continues with the next line.

Example:
For the following lines of code, since $(B-A) * B=2$, the next line assembled after GOTO is the line identified by the second line identifier, the line labeled BILL. Lines between GOTO and the line labeled BILL are skipped.

4.7.3

LNID
LNID inserts a dummy label for line identification purposes. The label has no value and is not entered in the Meta-Assembler symbol table. As long as no ambiguity exists, the same label may appear on more than one LNID line, or on any non-LNID line, or on both LNID and non-LNID lines.

## Format:

Rabel, LNID comments --

There is no operand; comments can be entered immediately after the command without the use of an asterisk.

LNID is particularly useful for defining the range of an RPT, since the use of normal labels may sometimes result in duplicate symbol definitions.
4.7.4

RPT AND GOTO PROCESSING

When META encounters an RPT directive, it compresses lines of code within the RPT range by removing comments and redundant blanks, and stores the lines in an internal table of definitions.

In the process of saving the lines of code within the RPT range, the assembler examines the command field of each line to ensure that the RPT range does not include an END or FINIS directive. The assembler also recognizes procedure and function definitions (chapter 5) which are within the range of an RPT.

When a procedure or function definition appears within an RPT range, label field symbols within the procedure or function definition are lucal to the procedure or function definition and are not considered in determining the RPT range.

Example:


Not end of RPT range

End of RPT range

A GOTO directive may appear within the range of an RPT. The object of the GOTO may be either within or outside the range of the RPT. If the object of a GOTO is outside the range of an RPT, the RPT is terminated.

Within a procedure or function definition, the object of a GOTO or an RPT must be within the procedure or function definition, and must be át the same level as the GOTO or RPT directive line. (Level of definition is discussed in section 5.4.)

Examples:

$B$ is within the procedure definition and is at the same level as the RPT directive line.

| KL. PRDC |  |
| :---: | :---: |
| $\begin{aligned} & C_{L} \quad P R \phi C_{1} \\ & A, \quad \text { NAME } \end{aligned}$ |  |
|  |  |
|  |  |
| н-1 |  |
| ENDS. |  |
|  |  |
| $3 \ldots \ldots$ LSIO $\ldots \ldots \ldots \ldots$ |  |
|  | 12 |

Illegal; $B$ is not in the procedure definition. If the procedure is referenced, the GOTO is terminated on encountering ENDS.
4.8

ASSEMBLY
TERMINATION

The directives END and FINIS specify the end of a subprogram and of a set of subprograms, respectively.
4.8.1

END

END terminates a subprogram. The symbol in the operand field is optional but, if present, must be a symbol of eight characters or fewer declared as an entry point in some subprogram. The symbol specifies the symbolic location at which execution is to begin.

Format:

## flabel END symbol comments.

4.8.2

FINIS
FINIS causes termination of assembly. Normally, an assembly is a set of subprograms, each of which ends with an END directive. The FINIS directive should follow the END directive for the final subprogram.

Format:

$$
\text { Rabe }_{1} 1_{1} \text { FINIS } S_{1} \text {, comments }
$$

Procedure and function definitions are bodies of code resembling subroutines but processed during assembly rather than object-time execution. They provide programmers with a means of conditionally generating sequences of code. A procedure reference consists of the appearance of the procedure name in the command field of a statement; the referenced procedure generates object code each time it is referenced according to parameters supplied with the reference. A function reference consists of the function name and its argument appearing in a statement; the function generates a value or set of values dependent on the argument.

A procedure or function definition begins with a PROC or FUNC directive, respectively, and terminates with an ENDS directive. The definition must precede a reference to it.

A functinn or procedure definition can wholly contain other definitions and references to yet other definitions. Such definitions are nested. Each nested definition is considered one 'evel higher than the definition that contains it. Nesting can occur to a level of 14. Levels of nesting are discussed more fully in section 5.4 .

Examples of nesting:


If the procedure being defined contains a forward reference to a locally defined symbol, proper data generation cannot result in a single pass. An optional parameter on the PROC directive indicates a two-pass procedure to permit local forward references. The Meta-Assemblert then makes a preliminary symbol defining pass through the procedure similar to the first assembly pass of a program.

## 5.1 DIRECTIVES

5.1.1 PROC

META provides directives specifically related to use of procedures and functions.

A PROC directive declares the beginning of a procedure definition.
Format:
Rabiel PROC, set namey set nome nexp comments

| label | Optional; if present, label becomes the name of sets <br> given on NAME lines in the procedure. <br> setname $1_{1}$ <br> Optional; set name that identifies the set in the <br> command field of the procedure reference. This <br> setname is in the command field and is separated <br> from PROC by a comma. |
| :--- | :--- |
| setname 2 | Optional; set name that identifies the set in the <br> operand field of the procedure reference. |
| exp $\quad$Optional; if value of expression is nonzero, procedure <br> requires two passes. Note: This option requires <br> core for expression building and causes a reduction in <br> assembly speed. It should not be used unless the <br> procedure contains a forward reference. |  |

When defining a two-pass procedure, the user should take care to prevent the inadvertent doubling of expression values. For the following lines of code, after a reference to procedure TWO, A has value 1 because it was initialized to zero each pass; B has value 2 because it was not initialized and was incremented once each procedure pass.

Example:


Operand field set has name $P$ TWO is entry name to procedure

TWO is reference to procedure
5.1. 2

FUNC
Format:

## Raber FUNC setname comments..

label | Optional; if present, label becomes the name of the |
| :--- |
| sets appearing on NAME lines in the function when the |
| function is referenced. |

setname $\quad$| Setname becomes the name of a set of parameters |
| :--- |
| passed to the function. |

A function should not include directives that generate code or affect counters.

Example:

| \|-1. FlUNC_ | Begin function FX |
| :---: | :---: |
| EX_工 NAME_, |  |
| $A, \ldots N S E T+H_{y}, 5,6$ |  |
|  |  |
| $\ldots \ldots$ | End function FX |
| + |  |
| , |  |
|  | Set B has two elements, 4 and 5 |

5.1.3 NAME

FUNC declares the beginning of a function definition. sets appearing on NAME lines in the function when the function is referenced. passed to the function.

NAME directives define entry names by which a function or procedure can be referenced. They must be between the PROC or FUNC directive and its associated ENDS directive. The label field symbol of the NAME directive is used as the command field of the statement referencing the function or procedure. Any number of NAME directives can appear within a definition.

Format:
Kabeil NAME set comments . -

| label | Required symbol; an entry name to the procedure or <br> function. |
| :--- | :--- |
| set | Optional set of expressions or sets that are to be <br> associated with this NAME. The name associated with <br> this set is in the label field of the PROC or FUNC <br> directive preceding this NAME. If the PROC or FUNC <br> label field contains a set name and the operand field <br> of the NAME directive is blank, the set consists of one <br> element having a value of 0. |

Example:


The procedure can contain references to a set named E. When the procedure is referred to by name ENTER1, elements 12 and I are assigned set name $E$ as if the following line had been written:
E. NSET I $2, I, I$

If, instead, the procedure is referred to by name ENTER2, elements 13 and $J$ are assigned set name $E$ as if the following line had been written:
E.NSEILI $13, J_{1}$

ENDS terminates a procedure or function definition.
Format:

## Kabel ENDS exp comments ...

When ENDS terminates a procedure definition, META expects no operand field entry. However, an asterisk must precede comments.

When ENDS terminates a function definition, exp is either an expression that defines the function value, a set name for a set of values, or set elements enclosed by brackets. A function reference that returns a set or a set name may be used instead of a subset. That is, to return a set, exp must be one of the forms:
(set)
setname
func (set)
Examples:


Begin outer procedure.

Begin inner procedure.

End inner procedure.

End outer procedure. Begin first function.

End first function. The value of the function is the sum of the first two values of the calling set. Begin second function.

End second function. The function returns a set of values rather than a single value.

The TREF directive terminates processing of a reference to a procedure or function definition before the ENDS directive.

Format:
label TREF exp comments -

For a function reference, control returns to the statement containing the reference and passes to it the value or set defined by the expression in the operand field of the TREF. Exp is either an expression that defines the function value, a set name for a set of values, or set elements enclosed by brackets. A function reference that returns a set or a set name may be used instead of a subset. That is, to return a set, exp can be one of the forms:
(set)
setname
fund (set)
Example:


A reference to IDENT terminates at the statement before the LDA NAME directive. References to the procedure by names LDA or STA terminate at the ENDS directive.

The LIBS directive enables the user to retireve procedure definitions from a file. It must not appear within a procedure or function definition.

Format:

label
dsi
sym $_{i}$

Optional symbol.
Data set identifier of the file containing the procedure definitions. This file, if it is not the system library file, must have been allocated and opened through use of MASTER control cards before META executes (3300/3500 MASTER Reference Manual Pub. No. 60213600). Procedures are searched for by name; they can be in any order on the file. If no dsi is given, META.uses MM.

Label field symbol of each NAME directive line for every outer procedure to be retrieved.

Function definitions can be obtained from a file through nesting of definitions and through externalization (section 5.4).

Procedures are stored on the system library by GLIB, the MASTER library generation program, and can be placed on some other file through use of XFER, the MASTER transfer routine (MASTER Reference Manual). They cannot be on an auxiliary library.

Examples:
The following procedure definition appears in-a procedure library on file DSI.


Procedure P1 is obtained by LIBS as follows:
LLIBS L'DSII', Plum

After P1 has been obtained, function names F1 and F2 are defined by writing P1 as a command field entry.


A procedure with names $A$ and $B$ is on the system library, *LIB.


By using the following LIBS directive, both A and B are defined and may be referenced. The user needs to specify only the first procedure name to obtain the entire definition.


If a user has no use for the $A$ entry name, he can save core during assembly by obtaining only the portion of the definition following the $B$ entry name.


## 5.2 DEFINITION PROCESSING

When META encounters a procedure or function definition, it compresses the lines of code representing the procedure by removing comments and redundant blanks, and stores the lines in core.

Meta-Assembler removes the NAME lines of outer level procedures and functions and inserts the labels of these lines into the symbol table. These labels are procedure or function entry names, and contain the location of the definition and the values of any sets associated with the NAMEs.

Entry names of inner definitions are not processed. Meta-Assembler stores these in the procedure definitions area as part of the lines of code comprising the definition. When procedure or function definitions are nested, entry points to the inner definitions are not known until the outer procedure is referenced. META does not save outer level PROC and FUNC lines, but instead creates a PROC or FUNC symbol table entry for each such line.

When an outer procedure or function is referenced, META processes only PROC, FUNC, NAME, and ENDS lines of the next level of procedures or functions. Unless the inner procedure name is externalized (section 5.4) subsequent reference to an inner procedure may occur only within the next outer procedure.

Each procedure and function definition may contain several NAME directive lines. The position of a NAME directive determines the first line of code to be processed when the procedure is referenced.

Example:


If the procedure is called by name $X$, the first line of code processed is:


If the procedure is called by name $Y$, the first line processed is:


The position of NAME directive lines within a procedure affects LIBS directive processing. If the following line is written the entire procedure is retrieved from the file.


If LIBS is written as below, the only line preceding the NAME line with label $Y$ retrieved is the PROC directive line.


To cefer to a proceduy, write the laber of anyan ME directivemnetn the dcfinition as a command. The label ficld can be blank or can contain a symbol that is assigned the value of the current location counter. To $\mathrm{su}_{\mathrm{i}}$. ly parameters to the procedure, place a set in the operan! field of the procedure call line, append a set to the procedure name in the command field, or do both. Within the procedure definition, the sets are referred to as if they were defined by NSET directives. If set names are provided in the command and operand fields of the PROC directive or the operand field of the FUNC directive and the corresponding field of the procedure or function reference is blank, the set used consists of one zero element.

Example:


When the procedure is referred to by name ENTER, elements A, B, \$, $[\mathrm{C}-3,5]$ are associated with name JOE as if the following line had been written.
$\mid T A E$, NSEI $A, B_{1} \$,[C-3,5] \ldots$

JOE [1] refers to A, JOE [2] refers to B, JOE [3] refers to the value of $\$$ at the time the reference occurs, and JOE [4] consists of a subset of two elements, C-3,5.

Set $X, Y$ has set name $M$ and is referred to as if the following line had been written.


Thus

$$
\begin{aligned}
& \operatorname{JOE}[3]=\$ \\
& \operatorname{JOE}[4,1]=\mathrm{C}-3 \\
& \mathrm{M}[2]=\mathrm{Y}
\end{aligned}
$$

The label appearing on the PROC directive line assigns a name to the set in the operand field of the NAME line. In the preceding example, $E$ is the set 12, I.

To refer to a function, write the label of a NAME directive appearing in the function definition followed by an argument enclosed in parentheses. A function reference must include the parentheses.

Example:


In the above reference, FU [1] is 15 and FU [2] is 4; A has value $(15+4-1) / 4$. If the reference had been|CQUOT(), the set FU would have been a single element set with value of zero and would have been illegal because FU [2] is a divisor with value 0 .

Parameters are referenced within a function in exactly the same way as they are referenced within a procedure.

A reference to a function that returns a value may appear as an operand in an expression. Reference to a function that returns a set may appear anywhere a set name may appear.

When a procedure is referred to, META forms as many as three sets in the symbol table. The set in the operand field of the procedure reference line, the set appearing in the command field of the procedure reference line, and the NAME directive set associated with the procedure reference.

The set in the operand field of the procedure reference line is evaluated and entered in the assembler symbol table. Its set name is the symbol that appeared in the first operand subfield of the PROC directive line for the procedure. The level of definition of the set is one greater than the level in effect for the procedure reference line.

A set appearing in the command field of the procedure reference line is processed in the same manner as the operand field set of the procedure reference line. The name of this set is the entry in the second subfield of the command field of the PROC directive.

The NAME directive set associated with the procedure reference is treated differently. At the time of procedure reference, the elements of the NAME directive set are already in the assembler symbol table but have no set name. META forms the NAME directive set in the assembler symbol table by copying the elements of the NAME directive set from one point in the symbol table to another and by assigning them the set name (the symbol from the label field of the PROC directive line). The level of definition is the same as for the other two sets previously described.

When META encounters the ENDS line for the procedure, it removes local symbols and sets from the symbol table. Externalized symbols are saved.

Meta-Assembler processes lines of code between a NAME line and a TREF or an ENDS line as if these lines appeared on the source input file. The lines are read from core storage rather than from the source input file. When there is nesting of definition, it also reads PROC, NAME, and ENDS lines from core storage. Again, processing is similar to that for lines on the source input file. Information is extracted from the first encountered PROC line and all associated NAME lines; other lines are skipped until a corresponding ENDS line. Had these lines been on the source input file, the assembler would have saved them. However, since the lines are already in core, it is unnecessary to save them again.

META allows nesting of function and procedure references as well as nesting of definitions. A definition can contain a reference to another procedure and, within that procedure, there can be a reference to still another procedure. Nesting of references, as with nesting of definitions, can continue to 14 levels.

Meta-Assembler recognizes 16 levels of symbol definition. Symbols defined at a given level are always available at the given level and all higher (inner) levels, but cannot be referred to at lower (outer) levels.

Symbols external to the program (i.e., those appearing as operands in an EXT directive) are defined at level 0. Symbols defined in the program but outside of procedures or functions are at level one. Symbols defined within procedures or functions are at level two or higher, the level being raised by one for each nesting of the reference.

Except for labels of NAME directives, which are available to the next outer level, labels within a procedure or function definition are local to the procedure or function; they are not available to outer procedures or to the program.

To make a label defined within a procedure or function available outside that procedure or function, the programmer can append one or more dollar signs to the symbol. Each dollar sign lowers the definition of the symbol one level to a minimum level of 1.

## Examples:



Define A one level lower. Define $B$ two levels lower. Define P[1] one level lower (see section 6.5 for ${ }^{\circ} \mathrm{SYM}$ ).

Thus, by lowering the procedure level of a symbol definition, the definition is available at a lower level outside the procedure or function.

Example:
In the following example, procedure $C$ is defined at level 2 when referenced by the main program (second line from bottom). Its entry name ( $C$ ) is known at levels 1 and 2. Within $C$, a call to procedure A defines A one level higher (level 3) causing its entry name (A) to be known at levels 2 and 3. Label E is local to procedure A. Label B is known at levels 1, 2, and 3. Label D is known at levels 2 and 3. Consequently, labels D and E are not known when they are referenced at level 1 by the GEN directive following the reference to procedure $C$. If the reference were to A instead of C, A would be defined at level 2 making labels $B$ and D available to the GEN directive.


A known at levels 3 and 2
E local to level 3
B known down to level 1
D known at levels 3 and 2
B,D, and E all known
$C$ known at levels 2 and 1

Level 2; E not known

Level 1; C known
Level 1; D and E not known

In the process of assembling source programs, Meta-Assembler constructs tables of information about elements of the source program. Attribute functions provide the user with information about expressions and sets.

The implicit attribute of a symbol or a set element is its value. Within: Meta-Assembler, the value attribute of a symbol is synonymous with the symbol; no further notation is needed to obtain that information.

Example:
Let A and B be defined as follows:


Within META, $A * B$ and $3 * 4$ are identical expressions.
Attribute functions are used to obtain information about attributes other than value. As with a symbolic reference, an attribute function reference results in a value. To refer to an attribute function, write the attribute name followed by an expression or set enclosed in parentheses. An attribute function reference can be an operand in an expression.

## 6.1 RELOCATION <br> (REL)

The relocation attribute function, REL, returns value zero if the expression within the parentheses is not a value or is an absolute value. If the expression is relocatable relative to a control section origin, REL returns the internal location counter designation (1-15) of the control section containing the expression.

Example:
$\bullet$


If A is in control section using location counter 1, go to $\mathbf{C}$.

Assume the program contains only one program control section and that B is an expression in that section.
$\operatorname{REL}(B)=1 \quad$ The first program control section is always assigned location counter 1.
$\operatorname{REL}(15)=0 \quad$ The argument is absolute.

The mode attribute function, MDE, returns the mode of the argument.

| Mode | Type of Expression |
| :--- | :--- |
| 0 | Not a value; for example, a set or function name |
| 1 | Integer (decimal or octal) value |
| 2 | Real- or floating-point value |
| 3 | BCD character string, right adjusted |
| 4 | BCD decimal integer |
| 5 | BCD character string, left adjusted |
| 7 | ASCII character string |
| 9 | Relocatable word address (includes literals, control |
| 10 | section names, and specia' character \$) |
| 11 | External word address |
| 12 | Relocatable byte address |
| 12 | External byte address |

Examples: Let $\mathrm{A}, \mathrm{B}$, and C be defined as follows.


$$
\begin{aligned}
& \mathrm{D}[1]=\operatorname{MDE}(\mathrm{A})=2 \\
& \mathrm{D}[2]=\operatorname{MDE}(\mathrm{B})=1 \\
& \mathrm{D}[3]=\operatorname{MDE}(\mathrm{C})=4
\end{aligned}
$$

NUM returns the number of elements in a set. If the symbolic item is not a set, NUM returns value 0 .

Examples: Let A and B be defined as shown.
$A, N S E I_{1}, 4,5,5,1,10, \ldots \ldots$
$B_{1}, E Q U, 13, \ldots, \ldots, \ldots$

$$
\begin{array}{ll}
\operatorname{NUM}(A)=3 & \text { Set A has three elements. } \\
\operatorname{NUM}(\mathrm{A}[1])=0 & \mathrm{~A}[1] \text { is a value, not a set. } \\
\operatorname{NUM}(\mathrm{A}[3])=2 & \mathrm{~A}[3] \text { is a set of two elements. } \\
\operatorname{NUM}(\mathrm{A}[3,1])=0 & \mathrm{~A}[3,1] \text { is value } 7, \text { not a set. } \\
\operatorname{NUM}(\mathrm{B})=0 & \text { B is not a set. } \\
\text { NET } &
\end{array}
$$

$$
\operatorname{NUM}(\mathrm{C})=1 \quad \text { Set } \mathrm{C} \text { has one element (zero) }
$$

The following example tests for number of elements in a set and tests elements of a set for subsets.

6.4 SIZE OF DATA

SZE returns either the number of object machine bytes needed to contain the value of an expression or the number of characters, depending on the mode of the expression. If the item is not a value, SZE returns value zero. SZE considers an address to be a one-word value.

SZE returns values depending on mode.

| Mode | Size |
| :---: | :--- |
| 0 | Zero |
| 1 or 2 | Number of bytes |

3, 4, 5, or 7 Number of characters

9, 10, 11
or 12 One word expressed as a byte count

## Examples:

Let $A, B, C, D$, and E bedefined as shown for an object computer word size of 24 bits and byte size of 6 bits.


| SZE $(A)=3$ | Three characters |
| :--- | :--- |
| SZE $(B)=8$ | Two words or eight bytes |
| SZE $(C)=4$ | One word or four bytes |
| SZE $(D)=4$ | One word or four bytes |
| SZE $(E)=0$ | E is a set, not a value |

SYM causes the value of the argument expression to be treated as a symbol. A SYM attribute function reference can appear in the label, command, or operand field. By using SYM, the programmer creates a symbol which is the value of the argument expression. The assembler represents the symbol as either 24 or 48 bits.

One use of the SYM attribute function is to refer to a symbol that is otherwise illegal. SYM can also be used for symbol concatenation.

Another use is to move a symbolic parameter into any field of a procedure or function. In this way, symbols supplied as parameters oan be defined within a procedure or function.

Examples:


The following example illustrates symbol concatenation. It generates symbol XY by scaling parameters X and Y into appropriate bit positions to form the value of the argument expression.


A reference to the symbol Q8Q. XYZ is ordinarily illegal because of the decimal point. It can, however, be referred to through use of the SYM attribute.


A reference to a procedure can be a SYM -defined name


The above code causes the procedure JFR to be interpreted as if it had been written:


After the inner reference to procedure $A B$, the EQU line becomes:


## 6.6 <br> WORD ADDRESS (WRD)

If the mode of the argument expression is either 9 (word) or 11 (byte), WRD returns the value of the argument as a word address. If the mode of the argument expression is 12 (external byte address), WRD changes the mode to 10 (external word address). If the mode of the argument expression is not $9,10,11$, or 12 , WRD returns the argument expression unchanged. If the argument expression is a byte address that does not correspond to a word address, truncation occurs.

Examples:

|Computer word 6 bits per byte, four bytes per word. A has mode 11, value 0.
AA has mode 9 , value 0 .
$B$ has mode 11, value 4.
BB has mode 9 , value 1.
C has mode 9 , value 2.
CC has mode 9 , value 2 .
D has mode 1 , value 10.
DD has mode 1, value 10. E has mode 11, value 12. $F$ has mode 11, value 13. |FF has mode 9, value 3 truncated. $G$ has mode 10 , value 0 .
GG has mode 12 , value 0 . GGG has mode 10 , value 0 .

## 6.7 <br> BYTE ADDRESS <br> (BYT)

If the mode of the argument expression is 9 or 11 (word or byte), BYT returns the value of the argument expression as a byte address. If the mode of the argument is 10 (external word address), BYT changes the mode to 12 (external byte address). If the mode is not $9,10,11$ or 12, BYT returns the argument expression unchanged.

Examples:


Computer word 6 bits per byte, four bytes per word. A has mode 9 , value 0 . AA has mode 11, value 0 . $B$ has mode 9, value 1. $B B$ has mode 11, value 4.
$C$ has mode 11, value 8. CC has mode 11, value 8. D has mode 1, value 10. DD has mode 1 , value 10. Ethas mode 10, value $0_{a}$ EE has mode 12, value 0 .

META can be called either by a MASTER task name control card or by a task already in execution.

When called by control card, META is loaded and placed in multiprogrammed execution as soon as its class, core, and file requirements can be met. When called by a CALL macro, a copy of META is loaded, if the job making the call does not already have a copy of the task. If it has a copy, the call is queued; that is, the caller must wait for the existing copy. Since META reinitializes itself, a job may make multiple calls to the MetaAssembler. Parameters ordinarily specified on a META control card (including parentheses) are passed as secondary parameters of a CALL macro. For use of CALL macro, see MASTER Reference Manual.

When the object deck is to be executed, it must be called by a task name control card or another task. The job monitor then calls the loader which loads relocatable binary information, links independently assembled subprograms, and loads and links library routines referenced by the loaded program. The program then executes multiprogrammed with all other active tasks.

## 7.1 <br> CONTROL CARDS

Assembly of META source programs under MASTER and execution of $3300 / 3500$ binary object decks require MASTER control cards identifiable by a $\$$ in column 1 (except for the end-of-file card). The name of the control card followed by any necessary parameters begins in column 2. The name and parameters must be contained on an 80 -column card.

MASTER control cards optionally accompanied by source and data decks are read serially from the input card reader. Cards required for META jobs are described in sections 7.1.1 through 7.1.5.

## 7.1:1

\$JOB
A JOB card must appear in a job deck either as the first card or, if a DIRECT card is used, as the second card.

```
$JOB, c,i,t\ell,\ell,p
```

BCD account number; required
$1 \quad \mathrm{BCD}$ job identifier; required
tl Time limit in minutes; optional
$\ell \quad$ Printer line limit (1-99999); optional
p
Punched card limit (0-99999); optional

Example:
$\$$ JOB, 639, DJ, 15, 150, 100, COMMENTS
7.1.2
\$SCHED

A SCHED card, immediately follows the JOB card in the job deck and provides the system with core and scratch mass storage requirements.
\$SCHED, CORE=qp, SCR=seg,.

Oth $r$ SCHED card parameters, not normally required by the META assembler, are described in the MASTER Reference Manual.

CORE $=$ qp $\quad$ Estimate of maximum amount of core, in quarter pages, required for assembly or execution, whichever has the higher core requirement. The estimate for the META assembler is a minimum of 32 quarter pages. Add four quarter pages if MASTER mnemonic instruction set is required and allow for any other procedures or functions.

If the loader determines that the estimate is below that required by the job, the job is terminated with a message on the OUT file.

When the CORE field is omitted, qp is set by an installation parameter.

$\mathrm{SCR}=\mathrm{seg}$

Number of segments of mass storage scratch area required by the job. The segment size is determined when the operating system is installed.

If the length of a segment is 10,000 words, the file for executable output (usually LGO) requires roughly one segment for each 400 source statements. Normally, LGO needs only one segment.
META uses at least one and sometimes three system scratch files in addition to files indicated on the META card. All are in standard MASTER blocked format with a block size of 1280 characters. META always |uses a file with the dsi INT for source card images of the subprogram being assembled. The SCR field must schedule sufficient segments for this file to contain the largest subprogram or a set of subprograms to be assembled.

If the X or F option is requested, META uses a scratch file having the dsi BIN. Normally, one segment is sufficient; the file contains most of the binary output for one subprogram.
If a cross reference table is requested, META writes reference information on a scratch file with the dsi INTP. Normally, one segment is sufficient for INTP.

If the sum of the mass storage requirements indicated by the JOB card line and punch limits and the SCR and ABORT requests exceeds the storage reserved for these files, the job is not initiated.
When the SCR field is omitted, seg is set to an installation parameter.
7.1.3
\$META

The MASTER task name control card that causes META to be called, loaded, and executed (multiprogrammed) has the following format.

The optional parameters, $p_{i}$, are separated by commas and may appear in any order within the parentheses. Parameters have the format:

$$
\begin{aligned}
& \text { assembly option }=\mathbf{d s i} \\
& \text { or }
\end{aligned}
$$

assembly option
The assembly options are character strings, beginning with $\mathrm{I}, \mathrm{L}, \mathrm{X}, \mathrm{F}, \mathrm{P}$, or $R$. The dsi's are MASTER data set identifiers of 1-4 alphanumeric characters; 0000 may not be used for a dsi.

The options, and the data set identifier assigned for each if none is given on the META card, are listed below:

| Option | Significance | dsi |
| :---: | :--- | :--- |
| I | Source input | INP |
| L | Listable output | OUT |
| X | Load-and-go output | LGO |
| F | Load-and-go output with <br> forced execution | LGO |
| P | Punchable output | PUN |
| R | Cross reference table <br> (selectable only in con- <br> junction with $L$ ) | Same dsi as for L |

The $X$ and $F$ options are mutually exclusive. If the $X$ option is used and assembly errors occur, META issues a SUPPRESS request (MASTER Reference Manual) so that the object program is not executed. Under the $X$ option, assembly errors do not prevent generation of the executable output, just its loading and execution in the same job. The $F$ option causes execution of the $3300 / 3500$ object program despite assembly errors.

The Meta-Assembler source deck can be on the standard input card reader (INP) or a file, such as a magnetic tape file, specified by the programmer. If it is on the card reader, the MASTER input preprocessor transfers the deck from the card reader to the INP file. The programmer has the option of bypassing this transfer by placing a DIRECT card in front of his deck.

MASTER either accumulates Meta-Assembler printer output on the mass storage standard output file (OUT) for automatic post-job processing, prints output directly during job execution, or places the output on some other file specified by the user and for which printing is not automatic.

Similarly, MASTER either accumulates Meta-Assembler binary output on a punch file (PUN) for automatic post-job punching, punches output directly during job execution, or places the output on some other file specified for the user and for which punching is not automatic. •

For all output options, META assigns a system scratch file if the user does not specify either a standard file (OUT, PUN, or LGO) or a permanent file. All scratch files are automatically released at job end. The SCR parameter on the SCHED card must allow for all scratch files.

Use of permanent files is described in the MASTER Reference Manual.

Example:
\$META(LIST, XCUTE, PUNCH)

META is loaded from MASTER library file *LIB. Source statements are read by META from the INP file. Statements and assembly listings are written on the job OUT file and automatically printed. The punchable output is written on the job PUN file and automatically punched. Executable output is written on the LGO file.
\$META(IN=SRCE, LIST=OUT , FORSX=GOGO)

META is loaded from MASTER library file *LIB. It reads source statements from file SRCE. Printer output goes to the OUT file and is automatically printed. The job does not have any punch output. Executable output goes to user file GOGO. Because of the F option, the program on GOGO can be loaded and executed despite errors. occurring during assembly.

A task name control card directs MASTER to call and load the object-time - program from the specified file and to begin execution of the task.

If the object-time program is to be executed following assembly, a task name card of the following form must follow the source deck (if it is on the standard input file) or the META card (if the source deck is elsewhere).
\$name, dsi
name
1-4 alphanumeric characters; name is required.
dsi dsi of an opened file from which the named task is to be loaded. When the dsi is zero or the field is omitted, MASTER looks for the task on the system library. Normally, dsi is LGO.

For execution of a previously assembled program, the task name card for the object deck immediately follows the SCHED card. The object deck follows the task name card or is on the named file.
7.1.5

END-OF-FILE
A job is terminated with an end-of-file card characterized by 7,8 punches in columns one and two. Columns 3-80 may contain comments.

88 END OF File

! !







 55 5








## 7.2

SAMPLE DECKS

The following sample deck structures illustrate the use of MASTER control cards in job decks.

Assemble, list, and execute


Assemble and list


This job does not|include execution of an object deck because the source program on file SRCE contains a UNIT directive describing a computer system other than the 3300 or 3500 . Output is to permanent file BEN. In this example, SRCE and BEN are on magnetic tape. For use of 607 parameter on SCHED card and for use of *DEF cards, refer to the MASTER Reference Manual.

Execute only


This example illustrates execution of a $3300 / 3500$ deck assembled previously by META.

List only


This job assembles the source deck but produces only a listing as output.

## 8.1 LIST FORMAT

When the L option is selected on the META control card, META generates list output. Each page of list output is in the following format:

title
date
source statement number
error code
relocation section
word address .
byte position
operand relocation
Object computer word
source statement

Characters supplied by TITLE directive.
Date of computer run.
Position of source statement in the source deck (00000-99999).
Code if source statement is erroneous (section 8.2).
Control section (00-15) containing object computer word.
Address of object computer word.
On byte-oriented source lines, position of byte in word from left to right. $00-\mathrm{n}$, respectively, where $n$ is the number of bytes per word.

Control section (00-15) containing operand; $X$ indicates operand is external symbol.
Object computer word generated by META (3-16 octal digits).

1-80 characters of source input line, including sequence number if provided.

## Example:

| METAMMSTER 000001 00002 | VER 1.0 <br> 010000000000 010000000000 |  |  | $\begin{aligned} & \text { LIBS } \\ & \text { IDENT } \end{aligned}$ | $\begin{aligned} & \text { LFMIBA, IDENT } \\ & \text { ARTAF } \end{aligned}$ | 09/06/68 | PAGE | 1 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| meta/master | VER 1.0 | FuMct | TEST |  |  | 09/06/68 | PAGE | 2 |
| 00003 | O1 0000000000 |  |  | TITLE | \#FWNCTIOM DIR |  |  |  |
| 00004 | O1 00000000000 |  |  | Emtry | BEEIM |  |  |  |
| 00005 | 010000000000 |  |  | EXT |  |  |  |  |
| ${ }^{000006}$ |  |  | ${ }_{51}$ | FUNC | \$2 |  |  |  |
| 00007 00008 |  |  | fukl | MNAF | 1,4 |  |  |  |
| 00008 00009 |  |  | funt | WNVE |  |  |  |  |
| 00010 |  |  |  | ENOS | si[2] + S2[2] | S2[3]), 5 |  |  |
| 00011 | 010000000000 | 0101000000 | begin | WP |  |  |  |  |
| 00012 | 01 00000000100 |  |  | EnI | 0.1 |  |  |  |
| 00013 | O1 01000000000100 | 14100000 |  | EMI | Funl $(1,2,3)$ |  |  |  |
|  | 010000000200 | 14000014 |  |  |  |  |  |  |
| 00014 | 0100000003300 |  |  | EM | $\pi$ |  |  |  |
| 00015 | 010000000400 | 1. |  | ISE | 24.1 |  |  |  |
| 00016 | O1 0000000000500 | 01100030 |  | EM | Tif |  |  |  |
|  | O1 00000000500 | 0114600042 |  |  |  |  |  |  |
| 00017 | 010000000600 | 0100700024 |  | RTJ | resut |  |  |  |
| 00018 | 01 0000000700 |  |  | EnI | 0.2 |  |  |  |
| 00019 | O100000000700 | 14200000 |  | EMI | Funk $(3,1,5)$ |  |  |  |
|  | 010000001000 |  |  |  |  |  |  |  |
| 00020 | 0100000011.00 | 01 14600050 |  | ENA | 12 |  |  |  |
| 00021 | 010000001200 |  |  | ISE | 48,2 |  |  |  |
|  | 010000001200 | 04200060 |  | E4 | T2F |  |  |  |
| 00022 | 010000001300 | 0114600056 |  | ena | T2F |  |  |  |
| 00023 | 010000001400 |  |  | RTJ | RESUL |  |  |  |
| 00024 | 010000001400 | 010070002 | SEIt | MSET |  |  |  |  |
| 00025 | 010000001500 |  |  | ENI | 0,3 |  |  |  |
| E\#* |  | 14300000 | ENOS | $\mathrm{si}[2]+$ |  |  |  |  |
| 00026 | 010000001600 <br> 010000001600 | 14000000 |  |  | Fun3(SEN) |  |  |  |
| 00027 | 010000001700 | 1 14600064 |  | ENA | T3 |  |  |  |
| 00028 | -1100000020 00 | 0114600064 |  | ISE | 30,3 |  |  |  |
| 00029 | 010000002000 | 04300036 |  | EMA | T3F |  |  |  |
| 00029 | 010000002100 | 01 14600072 |  |  |  |  |  |  |
| 00030 | 01.00000022 00 |  |  | RTJ | RESUT |  |  |  |
| 00031 | 010000002200 0100000023 00 | 01 00700024 |  | WP,I | BEGIM |  |  |  |
|  | 010000002300 | 0101400000 |  |  |  |  |  |  |
| ${ }_{0}^{00033}$ | 010000002400 |  | resut | UPP |  |  |  |  |
| 00034 | 010000002400 | 0101000024 |  |  |  |  |  |  |

8.2 ERROR CODES

Meta-Assembler flags each detected error with a single-character error code and 3 asterisks on the line of the source statement in error.

| Code | Meaning |
| :---: | :---: |
| C*** | Common error. An attempt was made to assemble information into numbered common. |
| $\mathrm{D}^{* * *}$ | Double definition. 1) A symbol has two values at the same level, or 2) A subprogram that does not contain a UNIT directive contains more than one SECP directive. |
| E*** | Expression error. The expression is syntactically correct, but an error, such as an illegal combination of modes, exists. |
| F*** | Forward reference error. A forward reference appeared in an expression which must be evaluatable. |
| I*** | Illegal instruction. The command field contains a symbol that is neither a directive nor the name of a procedure or FORM. The command field contains a misplaced directive. |
| N*** | Nesting error. More than 14 procedure levels or six RPT nests were encountered, or an RPT, procedure, or function is improperly nested. |
| $\mathrm{R}^{* * *}$ | Relocation error. The relocation associated with an expression is neither absolute, nor singularly positive, nor singularly negative, nor an external plus or minus a constant. |
| S*** | Syntax error. The syntax is unrecognizable or illegal. For example, a symbol has more than 12 characters. |
| T*** | Truncation error caused by 1) A value larger than the receiving field can accept. Note: No error is flagged when all the truncated bits are the same as the most significant bit (sign) of the value placed in the field. 2) A word-oriented statement following a byte-oriented statement. 3) Mixing of wordoriented and byte-oriented operations. |
| U*** | Undefined symbol. An operand contains a reference to a symbol that is neither defined in the program nor declared as external. |

## 8.3

SUPPLEMENTARY INFORMATION

Following the source program listing, META provides supplementary information as a standard part of the Meta-Assembler output listing. The supplementary information is identified as follows:

| Message | Meaning |
| :--- | :--- |
| LITERALS | $\begin{array}{l}\text { Identifies the list of literals. The loca- } \\ \text { tion and control section designator (0-15) } \\ \text { are given for each literal. }\end{array}$ |
|  | $\begin{array}{l}\text { Begins new page. Identifies list of con- } \\ \text { trol section names, octal length of section }\end{array}$ |
| in words, and location counter designator |  |$\}$

## 8.4

 CROSS REFERENCE TABLE
## 8.5 <br> MESSAGES <br> <br> ON OUT

 <br> <br> ON OUT}After detecting an error, META writes one of the following messages on the OUT file for the job.

Message
**META request ERROR code DSI dsi LINE line

Cause

Input/output error occurred. If other than read error (PICK reject code 04000000 or 050 xxxxx ), run is abnormally terminated. Message appears as voluntary abort code on accounting information as well as in listing.
request Blocker/deblocker or system OCARE request name
code Reject code for request: (Q) for blocker/deblocker. (A) for system OCARE
dsi Data set identifier for request
line - Number of META source input line
**META BAD LIBRARY
**META FINIS GENERATED
**META ILLEGAL \$META CARD
**META \$SCHED MORE CORE

The overlays of META are not in task directory. Library generation is incorrect. The run is abnormally terminated and message also appears as voluntary abort code.

FINIS directive generated because of end-of-file condition encountered on source input file. Execution continues.
\$META card contains illegal parameter such as illegal option or data set identifier. The run is abnormally terminated and message also appears as voluntary abort code.

Request for additional core rejected. The run is abnormally terminated and message also appears as voluntary abort code. Resubmit job with more core specified on \$SCHED card.

Examples:
**META SEXPAND ERROR 30000000 DSI INT LINE 10422
**META PICK ERROR 05000000 DSI INP LINE 00012

## APPENDIX SECTION

## 0

## 0

$\qquad$


| Type of Character | 501 <br> Printer <br> Graphic |  | Internal Code Octal | Card <br> Code |
| :---: | :---: | :---: | :---: | :---: |
| Blank | blank |  | 60 | space |
|  | $1+$ | plus | 20 | 12 |
|  |  | minus | 40 | 11 |
|  | * | times | 54 | 11,4, 8 |
|  |  | divide | 61 | 0,1 |
|  | $=$ | equals | 13 | 3,8 |
|  | $<$ | less than | 32 | 12,0 |
|  | > | greater than | 57 | 11, 7, 8 |
|  |  | period | 33 | 12, 3, 8 |
|  | - | comma | 73 | 0,3,8 |
|  | ( | left parenthesis | 74 | 0,4,8 |
|  | ) | right parenthesis | 34 | 12,4,8 |
|  | \% | percent | 16 | 6,8 |
| Special | $\$$$\neq$$\leq$ | dollar | 53 | 11, 3, 8 |
|  |  | not equal (apostrophe on keypunch) | 14 | 4,8 |
|  |  | less or equal | 15 | 5,8 |
|  | $\geq$ | greater or equal | 35 | 12,5,8 |
|  | [ | left bracket | 17 | 7,8 |
|  | 1 | right bracket | 72 | 0,8,2 |
|  | 1 | decimal exponent | 55 | 11,5,8 |
|  | $\downarrow$ | binary exponent | 56 | 11,6,8 |
|  | 7 | NOT | 36 | 12,6,8 |
|  | ; | semicolon | 37. | 12, 7, 8 |
|  | $\overrightarrow{\mathrm{I}}$ | right arrow | 75 | 0,5,8 |
|  |  | identity | 76 | 0,6,8 |
|  | : | colon | 12 | 2,8 |
|  | $V$$\wedge$ | OR | 52 | 11, 0 |
|  |  | AND | 77 | 0,7,8 |

TABLE A-1. BCD/ASCII Conversion Table

| 6-bit <br> BCD Code | $\begin{gathered} \text { 8-bit } \\ \text { ASCII } \end{gathered}$ <br> Character | Binary Status of ASCII Character (bit positions) |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | 7* | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| 00 | 0 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | 0 |
| 01 | 1 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | 1 |
| 02 | 2 | 0 | 0 | 1 | 1 | 0 | 0 | 1 | 0 |
| 03 | 3 | 0 | 0 | 1 | 1 | 0 | 0 | 1 | 1 |
| 04 | 4 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 0 |
| 05 | 5 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 1 |
| 06 | 6 | 0 | 0 | 1 | 1 | 0 | 1 | 1 | 0 |
| 07 | 7 | 0 | 0 | 1 | 1 | 0 | 1 | 1 | 1 |
| 10 | 8 | 0 | 0 | 1 | 1 | 1 | 0 | 0 | 0 |
| 11 | 9 | 0 | 0 | 1 | 1 | 1 | 0 | 0 | 1 |
| 12 | : | 0 | 0 | 1 | 1 | 1 | 0 | 1 | 0 |
| 13 | $=$ | 0 | 0 | 1 | 1 | 1 | 1 | 0 | 1 |
| 14 | ' | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 1 |
| 15 | \& | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 0 |
| 16 | \% | 0 | 0 | 1 | 0 | 0 | 1 | 0 | 1 |
| 17 | [ | 0 | 1 | 0 | 1 | 1 | 0 | 1 | 1 |
| 20 | + | 0 | 0 | 1 | 0 | 1 | 0 | 1 | 1 |
| 21 | A | 0 | 1 | 0 | 0 | 0 | 0 | 0 | 1 |
| 22 | B | 0 | 1 | 0 | 0 | 0 | 0 | 1 | 0 |
| 23 | C | 0 | 1 | 0 | 0 | 0 | 0 | 1 | 1 |
| 24 | D | 0 | 1 | 0 | 0 | 0 | 1 | 0 | 0 |
| 25 | E | 0 | 1 | 0 | 0 | 0 | 1 | 0 | 1 |
| 26 | F | 0 | 1 | 0 | 0 | 0 | 1 | 1 | 0 |
| 27 | G | 0 | 1 | 0 | 0 | 0 | 1 | 1 | 1 |
| 30 | H | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 0 |
| 31 | 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 1 |
| 32 | $<$ | 0 | 0 | 1 | 1 | 1 | 1 | 0 | 0 |
| 33 | - | 0 | 0 | 1 | 0 | 1 | 1 | 1 | 0 |
| 34 | ) | 0 | 0 | 1 | 0 | 1 | 0 | 0 | 1 |
| 35 | $\wedge$ | 0 | 1 | 0 | 1 | 1 | 1 | 1 | 0 |
| 36 | " | 0 | 0 | 1 | 0 | 0 | 0 | 1 | 0 |
| 37 | ; | 0 | 0 | 1 | 1 | 1 | 0 | 1 | 1 |

*ASCII bit 7 is unassigned and 0 for all codes.

TABLE A-1. BCD/ASCII Conversion Table

| 6-bit <br> BCD Code | 8-bitASCIICharacter | Binary Status of ASCII Character (bit positions) |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | 7* | 6 | 5 | 4 | 3 | 2 | 1 | 0 |
| 40 | - | 0 | 0 | 1 | 0 | 1 | 1 | 0 | 1 |
| 41 | J | 0 | 1 | 0 | 0 | 1 | 0 | 1 | 0 |
| 42 | K | 0 | 1 | 0 | 0 | 1 | 0 | 1 | 1 |
| 43 | L | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 0 |
| 44 | M | 0 | 1 | 0 | 0 | 1 | 1 | 0 | 1 |
| 45 | N | 0 | 1 | 0 | 0 | 1 | 1 | 1 | 0 |
| 46 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | 1 | 1 |
| 47 | P | 0 | 1 | 0 | 1 | 0 | 0 | 0 | 0 |
| 50 | Q | 0 | 1 | 0 | 1 | 0 | 0 | 0 | 1 |
| 51 | R | 0 | 1 | 0 | 1 | 0 | 0 | 1 | 0 |
| 52 | ! | 0 | 0 | 1 | 0 | 0 | 0 | 0 | 1 |
| 53 | \$ | 0 | 0 | 1 | 0 | 0 | 1 | 0 | 0 |
| 54 | * | 0 | 0 | 1 | 0 | 1 | 0 | 1 | 0 |
| 55 | \# | 0 | 0 | 1 | 0 | 0 | 0 | 1 | 1 |
| 56 | 1 | 0 | 1 | 0 | 0 | 0 | 0 | 0 | 0 |
| 57 | > | 0 | 0 | 1 | 1 | 1 | 1 | 1 | 0 |
| 60 | Blank | 0 | 0 | 1 | 0 | 0 | 0 | 0 | 0 |
| 61 | / | 0 | 0 | 1 | 0 | 1 | 1 | 1 | 1 |
| 62 | S | 0 | 1 | 0 | 1 | 0 | 0 | 1 | 1 |
| 63 | T | 0 | 1 | 0 | 1 | 0 | 1 | 0 | 0 |
| 64 | U | 0 | 1 | 0 | 1 | 0 | 1 | 0 | 1 |
| 65 | V | 0 | 1 | 0 | 1 | 0 | 1 | 1 | 0 |
| 66 | W | 0 | 1 | 0 | 1 | 0 | 1 | 1 | 1 |
| 67 | X | 0 | 1 | 0 | 1 | 1 | 0 | 0 | 0 |
| 70 | Y | 0 | 1 | 0 | 1 | 1 | 0 | 0 | 1 |
| 71 | Z | 0 | 1 | 0 | 1 | 1 | 0 | 1 | 0 |
| 72 | ] | 0 | 1 | 0 | 1 | 1 | 1 | 0 | 1 |
| 73 | , Comma | 0 | 0 | 1 | 0 | 1 | 1 | 0 | 0 |
| 74 | 1 | 0 | 0 | 1 | 0 | 1 | 0 | 0 | 0 |
| 75 | $\sim$ | 0 | 1 | 0 | 1 | 1 | 1 | 0 | 0 |
| 76 | - | 0 | 1 | 0 | 1 | 1 | 1 | 1 | 1 |
| 77 | $?$ | 0 | 0 | 1 | 1 | 1 | 1 | 1 | 1 |

*ASCII bit 7 is unassigned and 0 for all codes.

A 3300/3500 META mnemonic instruction is a procedure reference in which the label field optionally contains a symbolic address, the command field contains a mnemonic instruction and modifiers, and the operand field contains operands that depend on the mnemonic.

META assembles $3300 / 3500$ mnemonic instructions through the use of three standard sets of procedures on the system library. The sets are identified through their primary entry names as IDENT, MONITOR, and BDP.

## IDENT

IDENT includes procedures for the $3300 / 3500$ mnemonic instructions executable in program state, for the HLT instruction, and for the following 3300,'3000 COMPASS/MASTER pseudo instructions.

```
IDENT
BSS
BSS, C
DEC
DECD
```

9

Capabilities paralleling those provided by the following pseudo instructions are available through Meta-Assembler directives (Chapter 4).

| END | NOLIST |
| :--- | :--- |
| FINIS | LIST |
| ENTRY | EJECT |
| EXT | TITLE |
| EQU |  |

Of these, TITLE is the only directive that does not correspond to the COMPASS psuedo instruction.
META does not recognize the following 3300/3500 COMPASS/MASTER pseudo instructions.

| REM | IFZ | IFF | BCD, C |
| :--- | :--- | :--- | :--- |
| COMMON | PRG | IFN | ENDM |
| OCT | ORGR | DATA | LIBM |
| VFD | BCD | EQU,C | SPACE |
| IFT | MACRO |  |  |

## MONITOR

MONITOR includes procedures for assembling 3300/3500 mnemonic instructions executable in the monitor state only.

BDP
BDP provides for assembly of $3300 / 3500$ BDP instructions only.

## USE OF LIBS

Before they can be used, the $3300 / 3500$ mnemonic instruction procedures must be obtained from the library through use of the LIBS directive.

Examples:


Program state instructions
All but monitor state instructions
All but BDP instructions

## MASTER PROGRAM TASK

A META-Assembler program to be executed as a task under the MASTER multiprogramming operating system must include a copy of the user interrupt control routine (UIC) that provides the task with an entrance and an exit. Each subprogram must begin with a LIBS directive.

When loading and execution of the assembled output is called for by the task name card (section 7.1.4), the call connects with the UIC routine which contains a return jump to the task primary entry point. The return address is inserted into the operand field for the UJP as a normal function of a return jump execution. To obtain a copy of UIC, the program must declare UIC as an external symbol.

## Example:



Call for library procedures. First subprogram named JOE.

START is the task primary entry point.

Begin second subprogram named XY.

## PROCEDURE SETS

Three tables present brief descriptions of procedure references and resultant object code assembled by the DENT, MONITOR, and BDP procedure sets. For a complete description of the actual machine instructions, refer to the 3300 or 3500 Computer System Reference Manual.

Because the $3300 / 3500$ instructions are assembled through procedures, operation code modifiers must be defined as symbols having values. A reference to each of the sets IDENT, MONITOR, and BDP, causes the symbols for operation code modifiers to be defined. No other definition can be given these symbols. Thus, a group of words is reserved for each set of procedures.

The following list of terms defines modifiers, operands, registers, and nonstandard symbols that appear in the tables.

In some instructions, the execution address $m$ or $r$, or the shift count $k$ may be modified by adding to them the contents of an index register, $B^{b}$. The 2 -bit designator $b$ specifies which of the three index registers is to be used. Symbols representing the respective modified quantities are $M, R$, and $K$.

Term

A
b

B
.

n

D

MONITOR operation modifier: Conversion (alter the characters transmitted). Other: 24-bit A register or word count control for INAC, and INAW.
The $b$ subfield designates an index register. The $b$ subfield may be represented by a digit; a symbol; or an expression with a nonrelocatable value of 12 , or 3 .
MONITOR operation|modifier: Backward read or write. Other: Index register defined by $\mathrm{B}^{\mathrm{b}}$.

Examples:


Delimiting character is K .
V has mode 3.
Delimiting character is V .

48-bit E register.
IDENT and BDP operation modifier: Indicates equal.
IDENT operation modifier: Indicates greater than or equal.
MONITOR operation modifier: Indicates half assembly or disassembly.
$B D P$ operation modifier: Indicates $(B C R)=01_{2}$ jump condition.

IDENT operation modifier: Indicates indirect addressing.
Increment or decrement. The i address subfield may contain a symbol, constant, or expression which results in a nonrelocatable value from 0 to 7 .
MONITOR operation modifier: Indicates interrupt on completion.
Shift count
Field length of block. $0-177_{8}$. The $\ell$ address subfield may be a symbol or an expression which results in a nonrelocatable value from 1 to $177_{8}$.
BDP operation modifier: Indicates $(B C R)=10_{2}$ jump condition.
BDP operation modifier: Indicates left-to-right scan.
IDENT operation modifier: Indicates less than.
Number of characters in field R.
Number of characters in field $S$.
15 -bit word address, first operand, or jump address. The m address subfield may contain a symbol, $\$$, a constant, an expression, or a literal.
Actual operand or jump address as modified; $M=m+\left(B^{b}\right)$.
MONITOR operation modifier: Indicates no assembly or disassembly.
Same as $m$, second operand address.
IDENT and BDP operation modifier: Indicates not equal.
15 (or 17)-bit P register.
24-bit Q register.
17-bit character address. The $r$ address subfield may contain a symbol, literal, constant, external symbol, expression, or $\$$.
Actual character address as modified; $R=r+\left(B^{b}\right)$.
BDP operation modifier: Indicates right-to-left scan.
Abbreviation for read next instruction at. For example, RNI P+1 means read the next instruction at the current location plus 1 of the $P$ register.
Same as $r$, second operand address.
IDENT operation modifier: Sign extension if S present; no sign extension if $S$ omitted.
Other: Same as $R$, second operand address; $S=s+\left(B^{b}\right)$.
Scan character
6-bit address in register file. The $v$ address subfield may contain a symbol, constant, or expression which results in a nonrelocatable value 0 to $63_{10}$.

Term
w Page index file address.
$\mathbf{x}$
Connect code or interrupt mask. The $x$ address subfield may contain a symbol, constant, or expression that results in a nonrelocatable value $0 \leq x \leq 2^{12}-1$.
y
15 -bit operand. The y address subfield may contain a symbol, * or **, constant, an expression, or a literal.

ZRO . BDP operation modifier: Indicates (BCR)=0 jump condition.
()

Operation analysis symbol indicating the contents of. For example, (A) means the contents of the A register.

Operation analysis symbol indicating replace. For example, ( $M$ )-(A) means replace the contents of the A register with the contents of the $M$ operand field.

Procedures for COMPASS pseudo instructions precede the tables.
IDENT procedures are grouped according to instruction types as:

Transfers
Arithmetic operatiens
Character operations
Decisions
Jumps, pauses, and stops
Interrupt operations
No-operation instruction
Shift instructions
Logical instructions
MONITOR procedures are grouped according to instruction types as:
Transfers
Decisions
Jumps, pauses, and stops
Input/output operations
Interrupt operations
BDP procedures are not divided into subgroups.

IDENT sym The IDENT procedure names a subprogram and provides control information for META. The operand field contains a 1-8 character symbol naming the subprogram. The procedure contains a SECP directive that places the name on the IDC card of the relocatable object subprogram deck. The label field is defined as the value of the location counter.

The subprogram name is not an entry point name and cannot be referred to within the source subprogram. Each subprogram must have a SECP directive or IDENT instruction preceding all but the LIBS, UNIT, or list control directives.

Lines of code following IDENT are assembled, using the location control counter, until the next SECP, SECA, SECD, or ORG directive.

BSS reserves and labels a block of words in any area. The label field is blank or contains a symbol defined as the 15 -bit relocatable word address of the first word in the block.

The operand field specifies the number of words to be reserved. It must contain a constant, a symbol, or an address expression that results in a nonrelocatable value.

## Example:

ABLE BSS. $12 \ldots$


A double asterisk is illegal in the operand field. A symbol in the operand field must be defined in the label field of a preceding instruction.

A negative operand field such as -O '2' is interpreted as $\mathrm{O}^{\prime} 77777775^{\prime}$. META reserves $77777775_{8}$ words.

If the operand field is in error or is zero, no storage is reserved but the label field symbol is defined. If the operand field is zero, and a byte-oriented instruction immediately precedes the BSS, the next instruction that uses space begins with a new word.

BSS, C reserves and labels a block of bytes. The label field is blank or contains a symbol defined as a 17-bit relocatable address of the first byte (BCD character position) in the block to be reserved. The operand field specifies the number of bytes reserved. It must contain a constant, a symbol, or an address expression that results in a nonrelocatable value.

A negative operand field such as $-\mathrm{O}^{\prime} \mathbf{2 '}^{\prime}$ is interpreted as $\mathrm{O}^{\prime} 77777775^{\prime}$. META reserves $77777775_{8}$ bytes.

A zero operand does not reserve space but the label field symbol is defined.

Example:



DEC $d_{1}, d_{2}, \ldots, d_{n}$

DECD $d_{1}, d_{2}, \ldots, d_{n}$

DEC generates one computer word for each decimal value in the operand field. The label field is blank or contains a symbol defined as a 15-bit relocatable address of the first word generated. The operand field contains values, symbols, or expressions that result in decimal values.

Example:

$$
\text { DEC }-38,4.29,18 \ldots
$$

Generates three words.
DECD generates two computer words in 48-bit internal floating-point format for each real (floating-point) value in the operand field. The label field is blank or contains a symbol defined as the 15 -bit relocatable address of the first word generated. The operand field contains values, symbols, or expressions that result in real or floating-point values.

TABLE B-1. DENT PROCEDURE REFERENCES


TABLE B-1. IDENT PROCEDURE REFERENCES

| T y p e | Command Field | Operand Field | Operation | Object Code |
| :---: | :---: | :---: | :---: | :---: |
|  | LCAQ, 1 | m, b | Complement of $(\mathrm{M}) \rightarrow \mathrm{A}$; complement of $(M+1) \rightarrow Q$ | $\|$23 17   14 <br> 26 A b  m |
|  | LDA, I | m, b | $(\mathrm{M}) \rightarrow \mathbf{A}$ |  |
|  | LDAQ, I | m, b | $(\mathrm{M}) \rightarrow \mathrm{A}, \quad(\mathrm{M}+1) \rightarrow \mathrm{Q}$ |  |
|  | LDI, I | m, b | $\left(\mathrm{M}_{14-00}\right) \rightarrow \mathrm{B}^{\mathrm{b}}$ |  |
|  | LDQ, I | m, b | $(\mathrm{M}) \rightarrow \mathrm{Q}$ |  |
|  | QEL |  | $(Q) \rightarrow E_{23-00}$ |  |
|  | RIS |  | Relocate to instruction state |  |
|  | ROS |  | Relocate to operand state |  |
|  | STA, I | m, b | $(\mathrm{A}) \rightarrow \mathrm{M}$ | $\left.\begin{array}{\|llll}23 & 17 & 14\end{array}\right] 00$ |
|  | STAQ, I | m, b | $(\mathrm{A}) \rightarrow \mathrm{M},(\mathrm{Q}) \rightarrow \mathrm{M}+1$ |  |
|  | STI, I | m, b | $\left(B^{b}\right) \rightarrow M_{14-00}$ | 23 17 14 00 |
|  |  |  |  | 47 ab m |

TABLE B-1. IDENT PROCEDURE REFERENCES


TABLE B-1. IDENT PROCEDURE REFERENCES

| T $\mathbf{y}$ p e | Command Field | Operand Field | Operation | Object Code |
| :---: | :---: | :---: | :---: | :---: |
|  | ADA, I | m, b | $(\mathrm{A})+(\mathrm{M}) \rightarrow \mathrm{A}$ | $\|$23 17 14  <br> 30 $a b$  00 |
|  | ADAQ, I | m, b | $(A, Q)+(M, M+1) \rightarrow A, Q$ |  |
|  | AIA | b | $(A)+\left(B^{b}\right) \rightarrow A$, sign of $\left(B^{b}\right)$ is extended prior to addition |  |
|  | AQA |  | $(\mathrm{A})+(\mathrm{Q}) \rightarrow \mathrm{A}$ |  |
|  | DVA, I | m, b | $(\mathrm{A}, \mathrm{Q}) /(\mathrm{M}) \rightarrow \mathrm{A}$, remainder $\rightarrow \mathrm{Q}$ | $\|$23 17 14  <br> 51 a/bil   |
|  | DVAQ, I | m, b | $(\mathrm{A}, \mathrm{Q}, \mathrm{E}) /(\mathrm{M}, \mathrm{M}+1) \rightarrow \mathrm{A}, \mathrm{Q},$ <br> remainder with sigaic extended $\rightarrow E$ | $\left.\right\|_{$23 17 14  <br> 57 $a$ $b$ $}$ |
|  | FAD, I | m, b | Floating-point addition of $(M, M+1)$ to $(A, Q) \rightarrow A, Q$ |  |
|  | FDV, 1 | m, b | Floating-point division of (A, Q) by ( $M, M+1$ ) $\rightarrow \mathrm{A}, \mathrm{Q}$; remainder with $\operatorname{sign}$ extended $\rightarrow(E)$ |  |
|  | FMU, I | m, b | Floating-point multiplication of $(A, Q)$ and $(M, M+1) \rightarrow A, Q$ |  |
|  | FSB, I | m, b | Floating-point subtraction of $(M, M+1)$ from ( $A, Q$ ) $\rightarrow A, Q$ |  |

TABLE B-1. IDENT PROCEDURE REFERENCES


TABLE B-1. IDENT PROCEDURE REFERENCES


TABLE B-1. IDENT PROCEDURE REFERENCES


TABLE B-1. DDENT PROCEDURE REFERENCES


TABLE B-1. IDENT PROCEDURE REFERENCES


TABLE B-1. IDENT PROCEDURE REFERENCES

| T <br> y <br> p <br> e | Command Field | Operand Field | Operation | Object Code |
| :---: | :---: | :---: | :---: | :---: |
|  | NOP |  | No operation (assembled NOP), RNI P+1 |  |
|  | SCAQ <br> SHA | $k, b$ $k, b$ | Shift ( $A, Q$ ) left end around until upper 2 bits of $A$ are unequal; residue $K=k$ - shift count; if $b=1,2$, or $3, K \rightarrow B^{b}$; if $b=0$, $K$ is discarded <br> Shift (A); shift count $K=k+\left(B^{b}\right)$ (signs of $k$ and $B^{b}$ extended); if bit 23 of $K=1$, shift right; complement of lower 6 bits equals shift magnitude; if bit 23 of $K=0$, shift left; lower 6 bits equal shift magnitude; leit shitts end around; right shifts end off | 23 17   <br> 13 14 00  |
|  | SHAQ | k, b | Shift (A, Q) as or register; shift count $K=k+\left(B^{b}\right)$ (signs of $k$ and $\mathrm{B}^{\mathrm{b}}$ extended); if bit 23 of $\mathrm{K}=1$, shift right and complement of lower 6 bits equals shift magnitude; if bit 23 of $\mathrm{K}=0$, shift left and lower 6 bits equal shift magnitude; left shifts end around; right shifts end off |  |
|  | SHQ | k, b | Shift (Q); shift count $K=k+\left(B^{b}\right)$ (signs of $k$ and $B^{b}$ extended); if bit 23 of $K=1$, shift right, complement of lower 6 bits equals shift magnitude; if bit 23 of $K=0$, shift left, lower 6 bits equal shift magnitude; left shifts end around; right shifts end off |  |

TABLE B-1. IDENT PROCEDURE REFERENCES


TABLE B-1. LDENT PROCEDURE REFERENCES


TABLE B-2. MONITOR PROCEDURE REFERENCES


TABLE B-2. MONITOR PROCEDURE REFERENCES


TABLE B-2. MONITOR PROCEDURE REFERENCES


Pause Sensing Mask

| Mask Bits | Mask Codes | Condition | Notes |
| :---: | :---: | :---: | :--- |
| 00 | 0001 | I/O channel 0 busy | Channel read or write operation in |
| 01 | 0002 | 1 | progress, the External MC logic <br> 02 |
| 03 | 0004 | 2 | within the channel is set, or a Reply |
| 04 | 0010 | 3 | or Reject from a previous operation |
| 05 | 0020 | 4 | is still present at the channel |
| 06 | 0040 | 6 |  |
| 07 | 0100 | 7 |  |
| 08 | 0200 |  |  |
| 09 | 0400 | Typewriter busy | Typewriter I/O in progress |
| 10 | 1000 | Typewriter NOT finish | Finish logic not set |
| 11 | 2000 | Typewriter NOT repeat | Repeat logic not set |
|  | 4000 | Search/Move control | Search or Move operation in |
|  |  | busy | progress |

TABLE B-2. MONITOR PROCEDURE REFERENCES


INTERRUPT MASK REGISTER BIT ASSIGNMENTS

| Mask Bit <br> Positions | Mask Codes (x) | Interrupt Conditions Represented |
| :---: | :---: | :---: |
| 00 | 0001 | I/O Channel 0 (includes interrupts generated within |
| 01 | 0002 | 1 the channel and external equipment |
| 02 | 0004 | 2 interrupts) |
| 03 | 0010 | 3 |
| 04 | 0020 | 4 |
| 05 | 0040 | 5 |
| 06 | 0100 | 6 |
| 07 | 0200 | 7 |
| 08 | 0400 | Real-time clock |
| 09 | 1000 | Exponent overflow/underflow \& BCD faults |
| 10 | 2000 | Arithmetic overflow \& divide faults |
| 11 | 4000 | Search/Move completion |

TABLE B-2. MONITOR PROCEDURE REFERENCES


TABLE B-2. MONITOR PROCEDURE REFERENCES


TABLE B-2. MONITOR PROCEDURE REFERENCES

| $\mathbf{T}$ <br> $\mathbf{y}$ <br> $\mathbf{p}$ <br> e | Command Field | Operand Field | Operation | Object Code |
| :---: | :---: | :---: | :---: | :---: |
|  | OTAC, INT <br> OTAW,INT <br> OUTC, INT, B, $\mathrm{H}, \mathrm{A}$ <br> OUTW, INT, B, N, A <br> SEL | ch ch $\mathrm{ch}, \mathrm{r}, \mathrm{s}$ ch, m, n $\mathrm{x}, \mathrm{ch}$ | Character from ( $\mathrm{A}_{05-00}$ ) is sent to peripheral device, <br> (A) retained <br> Transfers ( $\mathrm{A}_{11-00}$ ) or $\mathrm{A}_{23-00}$, depending on type of I/O channel, to a peripheral device <br> Storage words assembled into 6 - or 12 -bit characters and sent to a peripheral device <br> Transfer 12- or 24 -bit words from storage to a peripheral device <br> If channel ch is busy, read reject instruction from $P+1$; if not busy, send a 12 -bit function code on channel ch with a function enable, RNI P+2 |  |
|  | CILO | cm | Lockout external interrupt on masked channels, cm, until channel is not busy |  |

TABLE B-2. MONITOR PROCEDURE REFERENCES


Internal Status Sensing Mask

| Masked Bit Positions | Mask Codes (x) | Interrupt Conditions Represented |
| :---: | :---: | :---: |
| 00 | 0001 | Parity error on channel ch |
| 01 | 0002 | Channel ch busy reading |
| 02 | 0004 | Channel ch busy writing |
| 03 | 0010 | External reject active on channel ch |
| 04 | 0020 | No-response reject active on channel ch |
| 05 | 0040 | $\dagger$ Illegal write |
| 06 | 0100 | Channel ch preset by CON or SEL, but no reading or writing in progress |
| 07 | 0200 | Internal I/O channel interrupt on channel ch upon: <br> 1) completion of read or write operation, or <br> 2) end-of-record |
| 08 | 0400 | $\dagger$ Exponent overflow/underflow fault (floating-point) |
| 09 | 1000 | $\dagger$ Arithmetic overflow fault (adder) |
| 10 | 2000 | $\dagger$ Divide fault |
| 11 | 4000 | $\dagger$ BCD fault |

$\dagger$ Peripheral Equipment Reference Manual, Pub. No. 60108800

TABLE B-2. MONITOR PROCEDURE REFERENCES

$\dagger$ Internal faults are cleared when sensed.

TABLE B-3. BDP PROCEDURE REFERENCES


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TABLE B-3. BDP PROCEDURE REFERENCES


TABLE B-3. BDP PROCEDURE REFERENCES


TABLE B-4. OCTAL CODE INDEX TO MNEMONICS

| Octal Code | Mnemonic | Octal Code | Mnemonic | Octal Code | Mnemonic |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 00.0 | HLT | 05.5 | QSG, S | 15.5 | INQ, S |
| 00.1-6 | SJ1-SJ6 | 05.6 | ASG | 15.6 | INA |
| 00.7 | RTJ | 05.7 | QSG | 15.7 | INQ |
| 01 | UJP, I | 06.0-7 | MEQ | 16.0 | No-op |
| 02.0 | No-op | 07.0-7 | MTH | 16.1-3 | XOI |
| 02.1-3 | IJI | 10.0 | SSH | 16.4 | XOA, S |
| 02.4 | No-op | 10.1-3 | ISI | 16.5 | XOQ, S |
| 02.5-7 | IJD | 10.4-7 | ISD | 16.6 | XOA |
| 03.0 | AZJ, EQ | 11.0 | ECHA | 16.7 | XOQ |
| 03.1 | AZJ, NE | 11.4 | ECHA, S | 17.0 | No-op |
| 03.2 | AZJ, GE | 12.0-3 | SHA | 17.1-3 | ANI |
| 03.3 | AZJ, LT | 12.4-7 | SHQ | 17.4 | ANA, S |
| 03.4 | AQJ, EQ | 13.0-3 | SHAQ | 17.5 | ANQ, S |
| 03.5 | AQJ, NE | 13.4-7 | SCAQ | 17.6 | ANA |
| 03.6 | AQJ, GE | 14.0 | NOP | 17.7 | ANQ |
| 03.7 | AQJ, LT | 14.1-3 | ENI | 20 | LDA, I |
| 04.0-3 | ISE | 14.4 | ENA, S | 21 | LDQ, I |
| 04.4 | ASE, S | 14.5 | ENQ, S | 23 | LOCH |
| 04.5 | QSE, S | 14.6 | ENA | 24 | LCA, I |
| 04.6 | ASE | 14.7 | ENQ | 25 | LDAQ, I |
| 04.7 | QSE | 15.0 | No-op | 26 | LCAQ, I |
| 05.0-3 | ISG | 15.1-3 | INI | 27 | LDL, I |
| 05.4 | ASG | 15.4 | INA, S | 30 | ADA, I |

TABLE B-4. OCTAL CODE INDEX TO MNEMONICS

| Octal Code | Mnemonic | Octal Code | Mnemonic | Octal Code | Menmonic |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 31 | SBA, I | 53. $(0+\mathrm{b}) 4$ | AIA | 64.4-7 | $\left\lvert\, \begin{aligned} & \text { MVE, dc } \\ & \text { MVZS, dc } \end{aligned}\right.$ |
| 32 | ADAQ, I | 53.41 | TQM |  |  |
|  |  |  |  | 65.0-3 | SCAN, LR, EQ |
| 33 | SBAQ, I | 53.42 | TAM |  | SCAN, LR, NE |
|  |  |  |  |  | $\begin{aligned} & \text { SCAN, RL, EQ } \\ & \text { SCAN. RL, EQ } \end{aligned}$ |
| 34 | RAD, I | 53. $(4+\mathrm{b}) 0$ | TAI |  | SCAN, RL, EQ |
| 35 | SSA, I | 53. (4+b)3 | TIM | 65.4-7 | SCAN, RL, EQ, dc SCAN, LR, NE, dc |
| 36 | SCA, I | 53. $(4+\mathrm{b}) 4$ | IAI |  | $\begin{array}{l\|l} \mathrm{SCAN}, & \mathrm{RL}, \mathrm{EQ}, \mathrm{dc} \\ \mathrm{SCAN} & \mathrm{RL}, \mathrm{NE}, \mathrm{dc} \end{array}$ |
| 37 | LPA, I | 54 | LDI, I |  |  |
|  |  |  |  | 66.0-3 | ATD |
| 40 | STA, I | 55.0 | RIS |  | CVBD |
|  |  |  |  |  | CVDB |
| 41 | STQ, I | 55.1 | ELQ |  | DTA |
|  |  |  |  |  | PAK |
| 42 | SACH | 55.2 | EUA |  | UPAK |
| 43 | SQCH | 55.3 | EAQ | 66.4-7 | $\int_{\mathrm{DTA}, \mathrm{dc}}^{\mathrm{ATD}}, \mathrm{dc}$ |
| 44 | SWA I | 55.4 | ROS |  |  |
|  |  |  |  | 67.0-3 | ADM |
| 45 | STAQ, I | 55.5 | Q |  | CMP |
|  | SCHA, I | 55.6 |  |  | SBM |
| 46 |  |  | AEU |  | TST |
|  |  |  |  |  | ZADM |
| 47 | STI, I | 55.7 | AQE |  |  |
|  |  |  |  | 67.4-7 | CMP, de |
| 50 | MUA, I | 56 | MUAQ, I |  | TSTN |
| 51 | DVA, I | 57 | DVAQ, I | 70.0 | JMP, HI |
| 52 | CPR, I | 60 | FAD, I | 70.1 | JMP, Z RO |
| 53.01 | TMQ | 61 | FSB, I | 70.2 | JMP, LOW |
| 53.02 | TMA | 62 | FMU, I | 70.6 | LBR |
| 53.04 | AQA | 63 | FDV, I | 70.7 | SBR |
| 53. (0+b) 0 | TIA | 64.0-3 | $\begin{aligned} & \text { MVBF } \\ & \text { MVE } \\ & \text { MVZ F } \\ & \text { MVZS } \end{aligned}$ | 71 | SRCE, INT SRCN, INT |
|  |  |  |  |  |  |
| 53. (0+b) 3 | TMI |  |  |  |  |

TABLE B-4. OCTAL CODE INDEX TO MNEMONICS

| Octal Code | Mnemonic | Octal Code | Mnemonic |
| :---: | :---: | :---: | :---: |
| 72 | MOVE, INT | 77.61 | PRP <br> TMAV |
| 73 | $\begin{aligned} & \text { INPC, INT, B, H, A } \\ & \text { INAC, INT } \end{aligned}$ | 77.62 | SBJP |
| 74 | OUTC, INT, B, H, A INAW, INT | 77.624 | SDL |
|  | INPW, INT, B, N, A | 77.63 | $\begin{aligned} & \text { CRA } \\ & \text { SRA } \end{aligned}$ |
| 75 | $\begin{aligned} & \text { OUTC, INT, B, H, A } \\ & \text { OTAC, INT } \end{aligned}$ | 77.634 | ACR |
|  |  |  | RCR |
| 76 | OUTW, INT, B, N, A OTAW, INT | 77.64 | APF |
| 77.0 | CON | 77.65 | PFA |
| 77.1 | SEL | 77.66 | AOS |
| 77.2 | $\begin{aligned} & \text { EXS } \\ & \text { COPY } \end{aligned}$ | 77.664 | AIS |
|  |  | 77.67 | OSA |
| 77.3 | INS <br> CINS |  |  |
|  |  | 77.674 | ISA |
| 77.4 | INTS | 77.70 | SLS |
| 77.50 | INCL | 77.71 | SFPF |
| 77.51 | CILO CLCA IOCL | 77.72 | SBCD |
|  |  |  |  |
|  |  | 77.73 | DINT |
| 77.52 | SSIM | 77.74 | EINT |
| 77.53 | SCIM | 77.75 | CTI |
| 77.54 | ACI | 77.76 | CTO |
| 77.55 | CLA | 77.77 | UCS |
| 77.56 | JAA |  |  |
| 77.57 | LAPR |  |  |
| 77.60 | PAUS |  |  |

```
0
```



0

When the META source deck contains a UNIT directive, the object computer is not the 3300 or 3500 , and binary output (if requested) is in an alternate form. Information is written as binary card images, that is, in 40 -word logical records in standard MASTER blocked format (MASTER Reference Manual).

Each 40 -word logical record consists of a set of 1606 -bit bytes. Binary output is in the form of a byte stream. The first four bytes of each logical record are:

| Byte | Value |
| :---: | :--- |
| 1 | Unused; 0 |
| 2 | $05_{8}$ |
| 3 | Unused; 0 |
| 4 | Unused; 0 |

The byte stream consists of multibyte items. The first byte of an item is its item type, indicating the class of information. The number and contents of the bytes in the item vary according to item type.

Type

## Information



Type

4

6

7

8

Item type
Contents of a word ( n bytes)
Item type; item contains relocation information associated with preceding type 4 item
(Leftmost bit position of field in ward ( 7 bits)
Field size (7 bits)
Positive or negative relocation (1 bit)
0 Positive
1 Negative
Word or Byte relocation (1 bit)
0 Word
1 Byte
Unused; 0 (2 bits)
Relocation counter
Item type; item contains external reference information associated with preceding type 4 item
(Bit position of field in ward (7 bits)
Field size (7 bits)
Positive or negative relocation (1 bit)
0 Positive
1 Negative
Word or byte relocation (1 bit)
0 Word
1 Byte
External symbol table ordinal (14 bits)
Item type
Entry point symbol
Entry point byte address
Relocation counter
Item type
Transfer symbol
Item type; end of stream on a logical record
Item type; end of stream

The number of bytes in a type 4 item is a function of the object computer word size. A value is right justified in the number of bytes required. For example, if the object computer word size is 19 bits, $n$ equals 4.

All symbols are left justified and blank filled in eight bytes. The collection of type 2 items forms the external symbol table. Type 7 items refer to this table.

For type 6 and 7 items, bit positions are numbered from right to left in ascending order, beginning with zero. Thus, for a word address reference on the 3300 , the following is true.

Leftmost bit position of field in word
Field size 15

When word size is 12 , the leftmost bit position of a 13 -bit field is 0 .
Example:
The following program results in the binary (byte) stream shown.
Program:


Binary stream of 6-bit bytes:

Card 1

Card 2


The first four bytes cause rows 7 and 9 to be punched in column 1 of a binary card; column 2 is blank. Successive bytes consist of items and their associated information. A space is indicated by $\wedge$.

For a subprogram, all external symbol items (item type 2) form a table of external symbols that immediately follows the table of control section name items.

Normally, a load address item (item type 3) immediately follows the last external symbol item. A load address item appears in the stream as necessary and always precedes the first contents-of-word item (item type 4). If a load address is more than one greater than the address associated with the previous contents-of-word item, META generates a load address item.

Example:


The binary output stream for the above is as follows.

Card 1


For a subprogram, all control section name items (item type 1) form a table of control section names. This table is first in the binary output stream. The entries in the table are in order according to their associated location counters. The first entry is for counter 0 or 1 depending on whether or not the program uses 0 , the absolute location counter.

Example:
First entry in the control section name table is for location counter 0.


Location counter 2
Location counter 0
Location counter 1
Location counter 3
Location counter 4

Binary stream for above program:

Card 21

| 00 | 05 | 00 | 00 |  |  |  |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 01 | A | B | $\wedge$ | $\wedge$ | $\wedge$ | $\wedge$ | $\wedge$ | $\wedge$ | 00 | 00 | 00 | 00 | 10 |
| 01 | X | Y | $\wedge$ | $\wedge$ | $\wedge$ | $\wedge$ | $\wedge$ | $\wedge$ | 00 | 00 | 00 | 00 | 11 |
| 01 | J | 0 | E | $\wedge$ | $\wedge$ | $\wedge$ | $\wedge$ | $\wedge$ | 00 | 00 | 00 | 00 | 12 |
| 01 | K | A | $\wedge$ | $\wedge$ | $\wedge$ | $\wedge$ | $\wedge$ | $\wedge$ | 00 | - | 00 | 00 | 11 |
| 01 | K | B | $\wedge$ | $\wedge$ | $\wedge$ | $\wedge$ | $\wedge$ | $\wedge$ | 00 | 00 | 00 | 00 | 12 |
| 11 | X | Y | Z | $\wedge$ | $\wedge$ | $\wedge$ | $\wedge$ | $\wedge$ |  |  |  |  |  |
| 00 |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 00 | 05 | 00 | 00 |  |  |  |  |  |  |  |  |  |  |
| 77 |  |  |  |  |  |  |  |  |  |  |  |  |  |

It is possible to change control sections at points other than at word boundaries and resume the control sections.

Example:


Causes item 4 to be 5 bytes
Causes item type 1
Causes item types 3 and 4

Causes item type 1
Causes item types 3 and 4
Returns to address A (counter 1)
Causes item types 3 and 4
Causes item type 9 (118)

The binary output stream for the above is as follows.


For the previous example, note that in the control section named JOB, two contents-of-word items (item type 4) are generated for the contents-of-word location zero.

GENB 1

| 04 | 01 | 00 | 00 | 00 |
| :--- | :--- | :--- | :--- | :--- |

GENB 3

| 04 | 00 | 03 | 00 | 00 |
| :--- | :--- | :--- | :--- | :--- |

Contents-of-word

| 01 | 03 | 00 | 00 |
| :--- | :--- | :--- | :--- |

A field with a size greater than that of the object computer word may contain a relocatable value or an external symbol plus or minus a constant. More than one contents-of-word item (type 4) result, but they are not consecutive. A relocatable reference item (type 6) or an external reference item (type 7) immediately follows the first contents-of-word item (type 4). The condition can be detected when the leftmost bit in the word and the field size indicate a position beyond the preceding computer word.

Example:


Causes item 4 to be 3 bytes
Causes item 1

48-bit field (four 12-bit words)
Causes items 3, 4, 6, 4, 4, 4

Causes items 4, 6, 4
Causes items 4, 6, 4
Causes item 9


The $3300 / 3500$ MASTER relocatable loader accepts relocatable binary object decks produced by the Meta-Assembler when there is no UNIT directive. During assembly, the X or F option on the META control card uses card images of the relocatable deck to be written on the LGO file (or some other file optionally specified). The P option on the META card causes the binary deck to be punched. A binary deck is comprised of the following types of cards.

Subprogram identification card (IDC)
Block common table cards (BCT)
Subprogram entry point cards (EPT)
Relocatable information cards (RIF)
External name and linkage cards (XNL)
Transfer cards (TRA)
End loading card (ELD)
These cards are described in the MASTER Reference Manual. Information on the cards is related to directives as shown in table D-1.

Table D-1 Loader Cards

| Card | W | Source of Information |
| :---: | :---: | :---: |
| IDC | ${ }^{41} 8$ | Name taken from SECP directive; length of subprogram calculated by META. |
| BCT | 478 | Names of labeled and numbered common blocks taken from SECD directives. |
| EPT | 428 | Entry points taken from ENTRY directives. |
| RIF | $1-368$ | Relocatable information generated by mnemonic instructions, GEN, GEND, GENB, LIT, TEXT, TEXTC, TEXTA. RES or RESB causes start of new RIF card. Relocation factor set for character addressing if symbol generated is defined in bytes. Increment/ decrement count and base depend on relocation counter used by Meta-Assembler. |
| XNL | 438 | External symbols taken from EXT directives. |
| TRA | 448 | Transfer point symbol taken from END directive. |
| ELD | 778 | Card generated upon encountering FINIS. |

## -

## GLOSSARY OF TERMS

## Absolute program

A program that must be loaded into specific core storage locations.

## Assemble

To prepare an object language program for the 3300/3500 Computer System or for some other computer system from a symbolic source language program.

ASCII

American Standard Code for Information Interchange
Attribute
A characteristic of a symbol (value), such as its size in words or bytes and its mode of representation (decimal, octal, character, etc.).

Byte
A subdivision of a word as defined by a UNIT directive, if the source program contains one; otherwise, a byte is 6 bits.

Byte stream
Output from the Meta-Assembler when the source program contains a UNIT directive. Each 40 -word record ( 1606 -bit bytes) consists of 11 types of multibyte items.

Command
The field in the source statement that specifies the operation to be performed by the Meta-Assembler.

## Control Section

The portion of object code generated under a single location counter.

## Definition

1. A group of source statements comprising a procedure or function. 2. The association of a symbol with a value and its other attributes so that use of the symbol causes its value or the address of its value to be used.

## Delimiter

Character or characters that limit a string of characters and therefore cannot be a member of the string.

## Directive

A source statement that instructs a Meta-Assembler.

## Elementary item

A self defining component of an expression.

## Entry point

A label of a source statement at which execution or processing can begin.

## Expression

A valid series of values, symbols, and functions that may be connected by mnemonic or symbolic operators as required to cause a desired computation.

## External symbol

A label defined in a subprogram other than the subprogram or at a level othe: than a level currently being assembled and used as an operand in the program or at the level being assembled.

## Forward reference

A label that is referenced in the operand field and has not been defined previously.

## Function

A series of source statements that, when referenced, provides a single value or a set to be used in the source statement containing the reference.

## Label

1. A string of alphanumeric characters used to identify or describe an item or placed at any location for informational and instructional purposes. 2. To assign a symbol as a means of identifying a source statement or a location in an object deck.

## Literal

An item of data having a constant value.

## Location counter

A counter for the 16 control sections controlled by the assembler.

## Meta-Assembler

An assembler that transcends the capabilities of a conventional assembler by allowing extensive programmer control of the assembly process.

Mnemonic instruction
Use of symbolic notation in place of actual machine code. A mnemonic instruction must be translated to actual operation codes by META procedure references.

## Operand

A piece of data upon which an operation is performed; the contents of the operand field of a source statement.

## Operator

The symbol or mnemonic that tells what to do with two operands, e.g., * is the operator for multiplication of the two operands as in $A * B$.

## Procedure

A subset of source statements meeting a specific purpose that can be repeatedly referenced to cause parameterized code generation.

## Processing

The interpretation by the Meta-Assembler of a source statement or group of source statements.

Real number
A value written with a decimal point, using decimal digits. The sign is a unary operator. An integer exponent preceded by $E$ may follow the real number.

Symbolic referencing
The assembler allows mnemonic symbols to be used in place of instruction codes, modifiers, addresses, formats, procedures, and functions. The assembler interprets the symbol and determines where to find specific information.

A collection of elements that bear a relationship to one another and have a common name. An element may be a set; i.e., a subset of a set. A reference to an element consists of the set name followed by one or more integers enclosed in brackets indicating the location of the element.

## Source program

A program written in META language that must be translated into machine language before it can be executed.

## Statement

An instruction interpreted by an assembler.

## Subprogram

A part of a program that can be assembled independently.

## Subscript

One or more integers enclosed by brackets used to specify a particular element in a set.
Unary operator
An operator such as the sign of a value ( + or - ) that operates on one operand only rather than causing an addition or a subtraction.

## Word

A group of bytes as defined by the UNIT directive if the source program contains one; otherwise, 24 bits, the standard $3300 / 3500$ word size.

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## COMMENT AND EVALUATION SHEET

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