# WHO SYBTEMS MAR/APR 1980 



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## 3-100 <br> MIBRD SY8TEMB

Volume 1 Number 2

Editorial Correspondence should be sent to: S-100 MICROSYSTEMS, BOX 1192, Mountainside, NJ 07092.

## STAFF

Sol Libes
publisher/editor
Russell Gorr
executive editor
Dennis Thovson
technical editor
Jacob Epstein
CP/M* editor
Jon Bondy
Pascal editor
Don Libes
assistant editor
Lennie Libes
Susan Libes
subscriptions/office manager

S-100 MICROSYSTEMS is seeking articles on S-100 software, hardware and applications. Program listings should be typed on white paper with a new ribbon. Articles should be typed 40 characters/inch at 10 pitch. Author's name, address and phone number should be included on first page of article and all pages should be numbered. Photos are desirable and should be black and white glossy.

Commercial advertising is welcomed. Write to S-100 MICROSYSTEMS, Box 1192, Mountainside, NJ 07092, or phone Sol Libes at 201-277-2063 after 4 PM EST.

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## The

 Editor's Page by Sol LibesThe following is reprinted from-the February 1980 issue of BYTE magazine and I feel requires no further comment here.

## Battle of the Buses

In the October 1979 "BYTE News," page 107, Sol Libes contends, in an item about the S-100 bus, that "those who wish to have a machine capable of getting the maximum benefits of microprocessors must go the S-100 route." While Mr Libes was comparing the S-100 bus to all-in-one systems, such as the TRS-80 and PET, his statement leaves out a number of computer systems with as much capability as S-100 systems, perhaps more in some cases. For example, the SwTPC S/09 and the Ohio Scientific Challenger III Series are two systems that come to mind. The former uses a 6809 processor with the SS-50 bus (see October BYTE, inside front cover), and the latter uses 6800,6502 , and $Z 80$ processors and apparently OSI's own bus (see back cover, same issue). Both of these systems have a 20 -bit address bus for large memories. SwTPC and several other companies make SS-50 bus systems using the 6800 . Other non-S-100 bus systems include the Heath H 8 and H11. Any of these systems, and probably others that I have left out, can be as good for serious personal computer users as any S-100 bus computer. The S-100 bus is not the only possible route.

Mr Libes also writes that "the S-100 bus is not processor dependent." This statement is debatable, in spite of the existence of S-100 boards for a number of microprocessors. Several signals on the S-100 bus are generated ONLY by the 8080. Any other processor must be "bent" into generating (or responding to) these 8080 -specific signals.

Personal computing could use a truly processor-independent bus. I feel that the S-100 bus will not be totally satisfactory in this role.

The mention of specific products in this letter does not necessarily constitute endorsement of these products. My point is simply that there are other buses besides the S-100, and that systems using these other buses can be just as capable as S-100 systems.

## Jim Howell

5472 Playa Del Rey
San Jose CA 95123

## Author Libes replies:

Thank you for your letter regarding my comments on S-100 systems in the October BYTE News column. Despite the views expressed in your letter, I still stand by my view that "those who wish to have a machine capable of getting maximum benefits of microprocessors must go the S-100 route." I agree with you that SS-50 and OSI Challenger III systems offer more power than integrated systems such as the TRS-80, Apple and PET. However, they still leave much to be desired compared to S-100. I will explain shortly.
Further, I also stand by my statement that "the S-100 bus is not processor dependent." The fact is that presently there are manufacturers selling six different 8-bit processor boards (8080, 8085, Z80, 6502, 6800 and 6809) and five different 16-bit processor boards (9900, LSI-11, 8086, Z8000 and Pascal Microengine) for S-100 systems. This means that eleven microprocessors have already been interfaced to the S-100. I do not know of any other system with this processor independence. Many of these microprocessors could not even be interfaced to buses such as the SS-50 or OSI without sacrificing performance.
When it comes to maximum power and flexibility the S-100 offers the following advantages over all other systems:

- More software available. There are several times more languages, operating systems and applications packages for S-100 systems than for any other system.
- There are currently close to two dozen different manufacturers of S-100 mainframes and about fifty manufacturers of over 400 S-100 plug-in boards. This is many times more than for any other system.
- There is greater computer power capability with S-100. What other system has direct addressing of up to 16 megabytes of memory ( 24 address lines) and 64 K input/output ports (16 address lines), up to eleven vectored interrupt lines, up to sixteen masters on the bus (with priority), up to twenty-three plug-in slots on the motherboard, up to 10 MHz clock on the bus, plug-in operator front panel, and more.
- The S-100 bus is now standardized by the Institute of Electrical and Electronic Engineers (IEEE) assuring conformance among manufacturers.

Regarding your reference to the H8 bus, note that Heath has discontinued production of this unit. Besides, it was dedicated exclusively to the 8080 and therefore was destined to an early death. The Heath H11 is essentially the same as and uses the same bus specifications as a Digital Equipment Corp LSI-11. Few other firms support the LSI-11 with products within the price range of the typical hobbyist. The hardware and software facilities, compared to the S-100, are limited and expensive.

Again, thank you for reading my column and I welcome any further comments you wish to make regarding my opinions.

## Sol Libes

# At Intersystems, "dump" is an Instruction. Not a way of life. (Or, when you're ready for IEEE S-100, will your computer be ready for you?) 



We're about to be gadflies again.
While everyone's been busy trying to convince you that large buses housed in strong metal boxes will guarantee versatility and ward off obsolescence, we've been busy with something better. Solving the real problem with the first line of computer products built from the ground up to conform to the new IEEE S-100 Bus Standard. Offering you extra versatility in 8 -bit applications today. And a full 16 bits tomorrow.

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ance, flexibility and economy they offer. Whether you're looking at a new mainframe, expanding your present one or upgrading your system with an eye to the future. (Series II boards are compatible with most existing S-100 systems and all IEEE S-100 Standard cards as other manufacturers get around to building them.)

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economical 2 -serial, 4-parallel Multiple I/O board. 8-bit A/D-D/A converter. Our Double-Density High-Speed Disk Controller. And what is undoubtedly the most flexible front panel in the business. Everything you need for a complete IEEE S-100 system. Available separately, or all together in our new DPS-1 Mainframe!

Whatever your needs, why dump your money into obsolete products labelled "IEEE timing compatible" or other words people use to make up for a lack of product. See the future now, at your Intersystems dealer or call/ write for our new catalog. We'll tell you all about Series II and the new IEEE S-100 Bus we helped pioneer. Because it doesn't make sense to buy yesterday's products when tomorrow's are already here.

## Tunteralysitumas"

Ithaca Intersystems Inc.,
1650 Hanshaw Road/P.O. Box 91, Ithaca, NY 14850
607-257-0190/TWX: 5102554346

# S-100 MICROSYSTEM NEWS by Sol Libes 

## UCSD REVOKES LICENSES TO DISTRIBUTE PASCAL

The university of California, San Diego (UCSD) has sent out a letter to all computer clubs that are licensed to distribute the UCSD Pascal package that their license is terminated effective April 2, 1980. These organizations paid $\$ 250$ for the license and some of these organizations are considering filing suit against UCSD for breach of contract.

The clubs have distributed the UCSD Pascal package to their members at charges that ranged from $\$ 5$ to $\$ 50$. A UCSD Pascal user must now pay $\$ 250$ to Softech, the new UCSD licensee, to obtain a copy of the package.

## UCSD PASCAL NEWSLETTER PUBLISHED

The first i $\frac{1}{\text { ssue of the UCSD Pascal }}$ Hobby Newsletter has appeared in print. It is 9 pages and is chock full of a lot of valuable informatin that users of UCSD Pascal will find invaluable. To get on the mailing list send $\$ 2$ to: Jim McCord, 330 Vereda, Leyenda, CA 93017. It's an absolute bargain.

INTERNATIONAL PASCAL USER'S GROUP
If you are into pascal you will also want to join the "International Pascal User's Group (PUG). You will get an occassional newsletter that is a couple of hundred pages that is a compendium of all the news thats fit to print about pascal from all over the world. It contains a lot of useful programs, too! To join, send $\$ 6$ to PUG c/o Dick Shaw, Digital Equipment Corp., 5775 Peachtree Dunwoody Road, Atlanta, Ga. 30342.

$$
\begin{aligned}
& 28000 \text { S-100 CPU BOARDS TO BE OUT SOON }
\end{aligned}
$$

(formerly Ithaca Audio) will soon
introduce a Z8000 CPU board for S-100
based systems. It will meet IEEE S-100
specs and contain an operating system in
ROM. II will also have a version of their
Pascal/Z software which can compile to
either 280 or 28000 code. Therefore any
present $Z 80$ Pascal/Z software can be
immediately recompiled to the new 28000 .
II also plans to introduce a 88000
assembler.
National Multiplex Corp., 78 Oliver

Ave, Edison, NJ 08817 has also disclosed that they plan to soon have available a Z8000 CPU card for $\mathrm{S}-100$ systems. No further details are available at this time.

Ithaca Intersystems has also disclosed that they have a prototype CPU card running using the new Motorola 68000, 16-bit, microprocessor IC. Since the 68000 is not yet in production there are no immediate production plans.

NEW $\frac{\text { S-100 }}{\text { Godbo }} \frac{8-B I T}{\text { El }} \frac{\text { CPU }}{\text { ecto }} \frac{\text { CARDS }}{\text { ics }} \frac{\text { TO }}{\text { Oak }} \frac{\text { BE }}{1 a n} \frac{\text { AVAILABLE }}{\text { Airport }}$
Godbout Electonics, Oakland Airport, CA, plans to produce an $S-100$ card with dual microprocessors. The card will contain both 8085 A and 8088 microprocessors. Both are Intel ICs. The 8085A is the new improved version of the 8080A and hence can execute all the present 8080 software. It will be clocked at 8 MHz . The 8088 is the new Intel 16 -bit micro with 8 -bit I/O. Hence, it executes Intel's 8086 code. Thus a user can run standard CP/M written in 8080 code and software (e.g. 8086 BASIC) written in 8086 code. The board will also contain a memory manager circuit to provide extended memory addressing.

Tarbell Electronics, Carson, CA, will soon introduce a Z 80 CPU card for $\mathrm{S}-100$ systems. Don Tarbell said that he feels that the Z80, with its huge software base, will continue to be the dominant microprocessor for the next few years. The software base for 16 -bit micros, he feels, will take at least a year or two to develop and hence he decided to introduce an 8-bit CPU card rather than a 16-bit CPU card.

## CP/M USER GROUP NEWS

After over $\frac{1}{\text { a }}$ year of no new disks, the CPM User Group will add eight new disks to the CPM User Group Software Library. The new disks were prepared by Bob Van Valzah of the Chicago Area Computer Hobbyist Exchange club (CACHE). This will bring the CP/M User Group library up to a total of 428 -inch floppy disks.

The CP/M User Group will distribute copies to clubs, as in the past, for copying. This software will be available for copying at the $C P / M$ User Group meeting
at the Trenton Computer Festival, Trenton, NJ, April 19-20 (donation of $\$ 1 / d i s k$ is asked for).

```
\(\frac{C P / M}{\text { The }} \frac{\text { SOFTWARE }}{\text { Small Systems Group, Box } 5429, ~}\) Santa Monica, CA 90405 has prepared a directory which lists "all" CP/M applications programs and classifies them by type with the name address of the vendor. To get a copy send \(\$ 2\) or include a large self-addressed stamped envelope with \$l.
```

XITAN/TDL NEWSLETTER PUBLISHED
All those TLD/Xitan system owners may no longer have TDL or Xitan for user support (what little there was of it) but they do have the XITAN NEWSLETTER published irregularly by: Dennis Thovson, 243 McMane Ave., Berkeley Heights NJ 07922. Send \$l for a sample copy.

MICROCOMPUTER SOFTWARE HONORED
For the first time a microcomputer software package has placed on the prestigous DATAMATION magazine "Honor Roll of Software Packages". Naturally it was CP/M, a product of Digital Research. Microsoft BASIC and UCSD PASCAL received honorable mention.

IMSAI LIVES $\frac{\text { ANEW }}{\text { One }}$ of the pioneers in the microcomputer field is alive and functioning as the the IMSAI COMPUTER Division of Fischer-Freitas Corp., San Leandro CA. FFC was the outfit handling IMSAI's production and warranty service at the time that IMSAI went into bankruptcy, last year. Actually, the new IC division is staffed with all former IMSAI employees from the top down. The company is housed in two buildings totalling almost $12,000 \mathrm{sq}$. ft. and has 12 employees.

Fischer-Freitas purchased almost 90\% of the finished IMSAI stock and the rights to use the IMSAI and IMDOS trademarks. The IC division is manufacturing and selling the full line of IMSAI Hardware products. They are presently selling copies of the IMDOS operating system which was purchased from IMSAI. However, they are still negotiating with Digital Research regarding updating and marketing of IMDOS, in the future. Todd Fischer, FFC director, said that he expects to have this resolved shortly. In the meantime, IC is supporting all present registered owners of the IMDOS software and is providing updates as needed.

Their address is: IMSAI COMPUTER Div, Fischer-Freitas Corp., 2175 Adams Ave, San Leandro CA.

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## LETTERS TO THE EDITOR

Sol:
First issue looks great! Thanks.
Ho pe

it goes well.<br>Jack Mathis<br>Union Carbide

Bound Brook NJ

## Dear Sol:

I received the first issue of MICROSYSTEMS and was quite pleased with it. The magazine is still quite rough around the edges but shows a great deal of promise.

There are quite a few topics I would like to see covered. One of these is graphics for the S100. I think this is its weakest point. Also, a comparison of the $\mathrm{S}-100$ with others, such as the Apple and TRS-80. Other topics might be a troubleshooting column, diagnostics, getting started with Pascal, software other than BASIC and many others.

Above all, keep it interesting.
Thanks.
Hector M. Smith
Fountain Valley CA

Dear Sirs:
Please find enclosed my check for $\$ 14.00$ for a two year subscription, as mentioned is the announcement in Infoworld.

I look forward with enthusiasm to the first issue, and hope you receive the widespread support a publication with your aids and audience deserves.

William T. Hole MD

To Sol Libes \& Russell Gorr:
Bless you!!
The S-100 hobbyist has been increasingly taking a back seat in the micro world. The stores and mags are catering to either the small business market or the Pet/TRS-80/Apple crowd. I can understand it, because that's where the big bucks are. But that doesn't mean that I like it. Your magazine is a welcome move back toward the group that started it all.

William C. Burns
Palo Alto CA

## Dear Sol:

Congratulations on the forthcoming S-100 MICROSYSTEMS. I am enclosing a check for a three year subscription. I'll await the first issue eagerly (but patiently).

I first heard that you were thinking about an S-100 publication from John Dilks around Christmas and it had been on my mind to write my best wishes. However, my clairvoyent card is not working so my printer never followed through. Then I read of S-100 MICROSYSTEMS in Computer World. You have your public relations act together -- an area that ON LINE could have paid more attention to. John suggested also that you might be looking for people and that $I$ might contact you in that regard. With studied sanity, I have avoided doing that -- I am currently enjoying retirement (to just a full time job). However, if $I$ can be of help in a pinch, let me know. When I come out of retirement, I may try my hand at writing. If so, you'll be hearing from me. If you're looking for software reviewers, keep me in mind.

Anyway, $I$ hope that $S-100$ MICROSYSTEMS turns out to be as successful as I think it will. My one suggestion is to put heavy emphasis on $C P / M$ as the 'software bus' that has made $\mathrm{S}-100$ systems so usable. But I'm sure that you are ahead of me on that thinking.

My very best wishes for success. And thanks for your help and support with ON-LINE.

Dave Beetle
Los Gatos CA
Editor: Dave is the former publisher of ON-LINE, which until recently sold to Computer Shopper, was the buy-line type publication for computer hobbyists. Dave did a great job with it.

Dear Mr. Libes:
As a long time reader I congratulate you on your new venture. With your extensive preparation and professional style I look forward to your great success.

I am writing to tell you of another CPU chip operating on the S-100 bus (the l2th by my count). The Signetics $8 \times 300$ is a very fast 8 -bit machine currently offered as a slave processor for BASIC (FORTRAN and PASCAL coming).

I'm not sure of the part number, but this board is also the first $\mathrm{S}-100$ slave processor that I know of. I heard a rumor about a slave device using Western Digital's micro but have nothing concrete.

The BASIC processor is called a DLX-10 and is marketed by Alaska Computer Systems, 12759 Poway Rd, Poway CA, 92064.

I offer this information for your

- CONTINUED ON PAGE 55 -



# NORTH STAR TOPICS <br> by <br> Randy Reitz <br> 26 Maple St. <br> Chatham Township, N.J. 07928 

The first of a regular column by and for users of North Star Disk Systems.

Welcome to what $I$ hope will be a regular feature in S-100 MICROSYSTEMS. With some support from North Star users I hope to bring you some interesting topics each month. Usually $I$ will draw from my experience with the North Star MDS I started using in July '78, but $I$ will welcome any contributions.

Some topics I can write about include Tiny-c, CP/M, several $N^{*}$ utility programs such as MARYELLN and finally $N *$ 's version of UCSD Pascal. I will write about these topics as long as $I$ have something to share. Of course I would like to hear about what is interesting to other $N$ * users so let me hear from you.

I will get started this month with my recent experience with $N * D O S$ release 5.1S. I picked up 5.1S (the "S" is for single density) in October 1979, but I didn't find the time to install it until recently. I also picked up a copy of the new SOFT-DOC System Software Manual that $N^{*}$ sells for $\$ 7.50$. So far I can say that the Manual is well worth it, but 5.lS is not much different from release 4. Of course if you are lucky enough to have a double density MDS, you will need release 5 to get on the air. Double density is no doubt the reason for release 5 .

So far, the major difference I have found in using 5.ls is the DOS is generally more "chatty" with you. Since the four DOS commands that require a buffer to operate were "kicked out" of release 5 (these commands are CF, CD, DT \& CO - which was "kicked out" in release 4), I suppose $\mathrm{N}^{*}$ used the extra space to put in more text to print out. For example, the hard disk error message is a bit more civilized (the type of error is displayed now) and when the system comes up a sign-on message is printed. This sign-on message caused me some trouble since $I$ use the "auto" command feature at bootstrap time
doubled.
As I was preparing this patch, I noticed that $\mathrm{N}^{*}$ fixed a problem with the directory list routine. The check for the number of lines to print before pausing for the RETURN key has been restrictd to the console (device 0) only. In release 4, the pause would occur no matter what output device was selected for the directory listing.

This patch also goes in the directory printing routine at the point where one directory entry has finished printing and the directory maintenance routine is about to be entered to check for more entries to print:

| ORG 2597H |  |
| :--- | :--- |
| CALLPRTBLK | ; blank after type |
| MOV A,M | ;get file type |
| INX H | ;start of usr area |
| CPI 1 | ;is file type ABS? |
| JNZ BLK8 | ;8 blks, CRLF? |
| DS 7 | ;skip 7 bytes |
| JMP CKCRLF | ;4 blks, CRLF? |
| RETURN DS 1 | ;is needed-return here |
|  |  |

The remainder of the code can go in any convenient spot such as DOS I/O area.

| *PATCH | FOR | DOUBLE | LUMN FILE DIR LISTING |
| :---: | :---: | :---: | :---: |
| BLK 4 | MVI | C, 4 | ;initialize counter |
|  | CALL | PRTBLK | ;output one blank |
|  | DCR | C |  |
|  | JNZ | BLK $4+2$ | ;output 4 blanks |
|  | RET |  |  |
| BLK8 | CALL | BLK 4 | ;output 8 blanks |
| CKCRLF | CALL | B LK4 | ;after each user area |
|  | LDA | DIRTG | ;check column toggle |
|  | XRI | 1 | ; to see if CRLF is |
|  | STA | DIRTG | ; needed |
|  | CZ | DOSCRLF | ; do CRLF |
|  | JMP | RETURN | ; and return from PATCH |

We have acquired the rights to all TDL software (\& hardware). TDL software has long had the reputation of being the best in the industry. Computer Design Labs will continue to maintain, evolve and add to this superior line of quality software.

- Carl Galletti and Roger Amidon. owners.

Software with Manual/Manual Alone

All of the software below is available on any of the following media for operation with a $\mathbf{Z 8 0} \mathbf{C P U}$ using the CP/W* or similar type disk operating system (such as our own TPM*).
for TRS-80* CP/M (Model I or II)
for 8" CP/M (soft sectored single density) for $51 / 4^{\prime \prime} \mathrm{CP} / \mathrm{M}$ (soft sectored single density) for $51 / 4$ " North Star CP/M (single density) for 51/4" North Star CP/M (double density)

## BASIC I

A powerful and fast Z80 Basic interpreter with EDIT, RENUMBER, TRACE, PRINT USING, assembly language subroutine CALL, LOADGO for "chaining", COPY to move text, EXCHANGE, KILL, LINE INPUT, error intercept, sequential file handling in both ASCII and binary formats, and much, much more. It runs in a little over 12 K. An excellent choice for games since the precision was limited to 7 digits in order to make it one of the fastest around. \$49.95/\$15.

## BASIC II

Basic I but with 12 digit precision to make its power available to the business world with only a slight sacrifice in speed. Still runs faster than most other Basics (even those with much less precision). \$99.95/\$15.

## BUSINESS BASIC

The most powerful Basic for business applications. It adds to Basic II with random or sequential disk files in either fixed or variable record lengths, simultaneous access to multiple disk flies, PRIVACY command to prohibit user access to source code, global editing, added math functions, and disk file maintenance capability without leaving Basic (list, rename, or delete). \$179.95/\$25.

## ZEDIT

A character oriented text editor with 26 commands and "macro" capability for stringing multiple commands together. Included are a complete array of character move, add, delete, and display function. \$49.95./\$15.

## ZTEL

Z80 Text Editing Language - Not just a text editor Actually a language which allows you to edit text and also write, save, and recall programs which manipulate text. Commands include conditional branching, subroutine calls, iteration, block move, expression evaluation, and much more. Contains 36 value registers and 10 text registers. Be creative! Manipulate text with commands you write using Ztel. \$79.95/\$25.

## TOP

A Z80 Text Output Processor which will do text formatting for manuals, documents, and other word processing jobs. Works with any text editor. Does justification, page numbering and headings, spacing, centering, and much more! \$79.95/\$25.

## MACRO I

A macro assembler which will generate relocateable or absolute code for the 8080 or Z 80 using standard Intel mnemonics plus TDL/Z80 extensions. Functions include 14 conditionals, 16 listing controls, 54 pseudoops, 11 arithmetic/logical operations, local and global symbols, chaining files, linking capability with optional linker, and recursive/reiterative macros. This assembler is so powerful you'll think it is doing all the work for you. It actually makes assembly language programming much less of an effort and more creative. $\$ 79.95 / \$ 20$.

## MACRO II

Expands upon Macro l's linking capability (which is useful but somewhat limited) thereby being able to take full advantage of the optional Linker. Also a time and date function has been added and the listing capability improved. \$99.95/\$25.

## LINKER

How many times have you written the same subroutine in each new program? Top notch professional programmers compile a library of these subroutines and use a Linker to tie them together at assembly time. Development time is thus drastically reduced and becomes comparable to writing in a high level language but with all the speed of assembly language. So, get the new CDL Linker and start writing programs in a fraction of the time it took before. Linker is compatible with Macro I \& II as well as TDL/Xitan assemblers version 2.0 or later. $\$ 79.95 / \$ 20$.

## DEBUG I

Many programmers give up on writing in assembly language even though they know their programs would be faster and more powerful. To them assembly language seems difficult to understand and follow, as well as being a nightmare to debug. Well, not with proper tools like Debug I. With Debug I you can easily follow the flow of any $Z 80$ or 8080 program. Trace the program one step at a time or 10 steps or whatever you like. At each step you will be able to see the instruction executed and what it did. If desired, modifications can then be made before continuing. It's all under your control. You can before continuing. It's all under your control. You can
even skip displaying a subroutine call and up to seven even skip displaying a subroutine call and up to seven
breakpoints can be set during execution. Use of Debug I can pay for itself many times over by saving you valuable debugging time. $\$ 79.95 / \$ 20$.

## DEBUG II

This is an expanded debugger which has all of the features of Debug I plus many more. You can "trap" (i.e. trace a program until a set of register, flag, and/or memory conditions occur). Also, instructions may be entered and executed immediately. This makes it easy to learn new instructions by examining registers/memory before and after. And a RADIX function allows changing between ASCII, binary, decimal, hex, octal, signed decimal, or split octal. All these features and more add up to give you a very powerful development tool. Both Debug I and II must run on a Z80 but will debug both Z80 and 8080 code. $\$ 99.95 / \$ 20$.

## ZAPPLE

A Z80 executive and debug monitor. Capable of search, ASCII put and display, read and write to I/O ports, hex math, breakpoint, execute, move, fill, display, read and write in Intel or binary format tape, and more! on disk $\$ 34.95 / \$ 15$.

## APPLE

8080 version of Zapple $\$ 34.95 /$ S 15

## NEW! TPM now available for TRS-80 Model

 II!
## TPM*

A NEW Z80 disk operation system! This is not CP/M*. It's better! You can still run any program which runs with $\mathrm{CP} / \mathrm{M}^{\star}$ but unlike $\mathrm{CP} / \mathrm{M}^{\star}$ this operating system was written specifically for the $Z 80^{*}$ and takes full advantage written specifically for the Z80* and takes full advantage
of its extra powerful instruction set. In other words its not warmed over 8080 code! Available for TRS-80* (Model I or II). Tarbell, Xitan DDDC, SD Sales "VERSAFLOPPY", North Star (SD\&DD), and Digital (Micro) Systems. \$79.95/\$25.

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DEALER INQUIRIES INVITED.
*CHECK IF A CRLF IS NEEDED AT END OF DIR RSTDIR POP $H$;get adr of dir toggle PUSH PSW isave flags XRA A ;prepare to zap toggle CMP M ; check first if zero MOV M,A ;see if another CRLF CNZ DOSCRLF ; is needed POP PSW ;restore flags LHLD ADRAVAL ; set HL->next disk adr RET

This patch makes three references to N*DOS routines:

| PRTBLK | EQU | $26 B 4 H$ | ;print one blank |
| :--- | :--- | :--- | :--- |
| DOSCRLF | EQU | 26 ClH | ;output CR and LF |
| ADRAVAL | EQU | 28A2H | ;next avail disk adr |

Also, the byte containing the number of console lines to display should be to load my spooling programs, a part of which is the console output driver. Needless to say that the console output drive had better be in memory to receive the $N$ *DOS greetings. This "feature" of release 5 caused much shifting of my I/O drivers.

My major disappointment with release 5 is that $\mathrm{N}^{*}$ did not take the opportunity to enhance the "LI" directory listing with double column output. The "PRESS RETURN TO CONTINUE" message for CRT consoles is nice, but $I$ think a double column directory listing is better. I have had a double column directory patch for some time so this "feature" is not new for N*DOS (for example, Dr. Dobb's Journal had a letter for such a patch in the October, '79 issue); but $I$ have never been completely satisfied with their performance. For example, the Dr. Dobb's patch will output spaces after each directory entry so that two entries will fit exactly on one line of your CRT. The CRT terminal is expected to provide an automatic $C R$ and LF. Hence there is a version for 64 and 80 column CRTs. The double column patch $I$ was using relied upon the fact that the directory maintenance routine stores the numbr of the directory entry being processed in the B-reg while the entry is being printed. My trick was to check for the B-reg to be odd before doing the CRLF. This worked except when the directory contained blank entries.

Here is the ultimate, perfect and completely fulfilling double column directory listing patch for version 5.1s (I will be happy to help figure out addresses for other versions). The only way to do this right is to allocate a flag to indicate which side of the CRT (or hard copy device) the directory entry is currently being displayed. The patch requires 43 bytes which can usually be found in the DOS I/O area. Here is the patch that goes at the end of the directory maintenance routine just before returning to DOS or the program calling

## entry DLOOK:

DIRTG

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- CONTINUED ON PAGE 55 -


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# LINEAR PROGRAMMING TECHNIQUES IN PASCAL <br> by <br> W. M. Yarnall <br> 19 Angus Lane <br> Warren, N.J. 07060 


#### Abstract

This is the first part of a two-part article on Linear Programming in PASCAL. Linear Programming provides a means for solving problems with numerous constraints that often make solutions non-obvious.


## INTRODUCTION

Linear Programming is defined as a technique for finding the optimum solution to a set of linear equations which are subject to constraints. It was first devised in 1947 under sponsorship of the U.S. Air Force and used to solve many problems in military planning.

Since that time, the main applications have been in the area of industrial planning.

Several algorithms have been developed for the solution of Linear Programming problems. Probably the most successful for general problems is the Revised Simplex method, first developed by George B. Dantzig (Stanford). The first successful solution of a Linear Programming problem on a high-speed electronic computer occured in January, 1962, on the National Bureau of Standards SEAC machine. Since that time, the simplex algorithm and its variations have been coded for most of the large general purpose computers in the US and England.

## APPLICATIONS

Linear programming problems are concerned with the efficient use or allocations of limited resources to meet desired objectives. There are four generally recognized classes of such problems:

1. The Product Mix Problems. A typical example is that of the manufacturer who must determine what combination of his available resource will enable him to manufacture his products in a way which not only satisfies his production schedule but also maximizes his profit.
2. The Transportation Problem.

This is the example of a situation
where material (of possibly different types) must be shipped from several sources to several destinations; the different routes from each of the sources to each of the destinations have differing costs. The problem is to satisfy the demands at each destination from the available stocks at a minimum cost for transportation.
3. The Diet Problem.

This problem is also called by Lowe the "activity analysis" problem. Here we are given the nutritional content of a variety of foods (and their costs). We must find the mix which will satisfy the minimum daily requirement for each of the nutrients at a minimum cost.
4. Gaming Problems.

This class of problems is one in which it is desired to maximize the payoff of a game (or contest) in which the variables are the strategies each player may use. Only the subclass of "Zero-Sum Two-Person" games have been solved so far.

Each of these problems have in common that:

* A multiplicity of choice exists in the solution,
* The choices are bounded by constraints, and
* Some objective must be optimized (maximized or minimized).

In this two-part article, an example of each of the classes will be taken up, and the steps in the formulation will be shown. The wide range of problems to which this technique can be applied is suggested by the recommended reading.

The Revised Simplex Algorithm has been coded for use on small computers, and is shown in Listing l. It has been implemented in UCSD Pascal, a language chosen for its wide acceptance and installation on small systems. It requires, as a minimum, 48 K of RAM and a single floppy drive. The author uses a dual single-density North Star drive system with the North Star implementation of the UCSD Pascal.

First, the program is keyed in, using the Pascal editor. Then the program is compiled and linked to the library functions needed.

Before the program can be run, a data file must be build; part 2 of this article will provide a data file editor program (also in UCSD Pascal), and several example problems).
$\frac{\text { An Example Problem }}{\text { In order to }}$
In order to solve a problem using Linear Programming techniques, it must first be stated in a "standard form". It usually takes some analysis to get a problem into this standard form (see the bibliography for suggested reading in this area). Our example will be presented in the standard form.

Problem:
Maximize

$$
\begin{equation*}
x_{1}+2 x_{2}+3 x_{3}-x_{4} \tag{1}
\end{equation*}
$$

subject to

$$
\left.\begin{array}{rl}
x_{1}+2 x_{2}+3 x_{3} & =15  \tag{2}\\
2 x_{1}+x_{2}+5 x_{3} & =20 \\
x_{1}+2 x_{2}+x_{3}+x_{4} & =10
\end{array}\right\}
$$

Here, equation (l) states an objective to be maximized (maybe a profit?), and equations (2) represent the constraints.

Since the program of Listing $l$ is coded to minimize the objective function, we convert equation (1) to:
minimize

$$
\begin{equation*}
-x_{1}-2 x_{2}-3 x_{3}+x_{4} \tag{la}
\end{equation*}
$$

subject to the constraints of equation (2).

This demonstrates that any objective function can be converted from a minimum to a maximum (and vice versa) by a simple operation. Our standard form is the minimizing function (la).

The representation of equations (2) used in the program is the matrix form. That is, the coefficients of the X's are kept in a two-dimensional array, called "ABAR" in the program. When this data is fed into the program (via reading a data file), the ABAR matrix will contain:

$$
A B A R=\left|\begin{array}{llll}
1.0 & 2.0 & 3.0 & 0.0 \\
2.0 & 1.0 & 5.0 & 0.0 \\
1.0 & 2.0 & 1.0 & 1.0
\end{array}\right|
$$

It can be seen that a column in this array corresponds to the coefficient of a variable (col $1=X_{1}$, etc), and the row corresponds to a constraint (there are 3).

The right-hand-side (RHS) of equation (2) is kept in an array, as are the coefficients of the objective vector.

After the data file is read in, the program echoes it out, and goes through several iterations, until it finds the required solution.

Listing 2 shows the data file, listed by the "LIST" function of the file editor. Each record (there are 22) is shown with a record sequence number. The file uses record variants; the first data is the "TAG". Tag 0 specifies

* an alphanumeric name for the problem (TEST),
* the number of rows (3) in the ABAR matrix, and
* the number of columns (4) (variables) in the problem.
This record must be the first in the file, since the use of the arrays is controlled by its data.


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Record 1 has a tag of l. It is optional and provides up to 64 characters to define the problem to which this data applies.

Records 2 through 20 can be in any order. It is best, however, to organize the data somewhat, so that it may be proof-read more easily. Here, records 2-4 have a tag of 2 , which defines the right-hand-side (RHS) data, along with a row name and row index. Records 4, 9, 13 and 17 define column names, column indices and the coefficients of the objective vector (array). They have a tag of 4 .

The rest of records 2-20 have a tag of 6, and define the ABAR coefficients, along with the row and column indices.

Records 18 and 19 are redundant, since the $A B A R$ matrix is filled with zeros before the data is read in; they are included here only for completeness.

Record 2l, with a tag of 99, is not necessary; it is, however, best to have it included to provide assurance that you have built and read the file correctly and completely.

Listing 3 show the results of a run with this data. The program first asks for the data file name; the initial data is listed. Note that row $\operatorname{ABAR}(M+1)$ is the objective function data.

The program then enters Phase 1, in which it finds a "basic feasible solution" which satisfies all the constraints. In this problem it goes through three iterations and lists the values of the X's at each stage. When a basic feasible solution is reached, Phase 2 is started. In this phase, the final optimization is done. The final results are printed:

$$
\begin{aligned}
& x_{1}=2.5 \\
& x_{2}=2.5 \\
& x_{3}=2.5
\end{aligned}
$$

The pointer to the variable is shown in the third column of the "LIST \& $X$ ARRAYS" printout; the last column shows the value of the variable. In this listing, the row labelled "M+1" shows, in the value columns, the negative of the cost; since we originally multiplied by -1, it also directly shows the profit. Note that none of $\mathrm{X}_{4}$ was used in finding the optimum. It was in the problem at the start of phase 2, but the "better way" was found.

## PROGRAM ERRORS

Several types of errors in data are tested for and reported (if found).

1. When the data file is read in, if the first record does not have a tag of 0 , the output "BAD FILE FORMAT" is output. The $r$ un is aborted.
2. In Phase 2, the output "ERROR IN ITERATION $N^{\prime \prime}$ indicates that a division by

0 was detected. The data has an error in it.
3. In Phase l, the output "NO FEASIBLE SOLUTION AFTER N ITERATIONS" indicates that the constraints, as supplied, are inconsistent.
4. In Phase 2, the output "UNBOUNDED SOLUTION" indicates that the problem as posed, has a minimum at Negative Infinity. This problem cannot be solved as posed.

There is a possibility that, during Phase 2, the program may "cycle" - go through iterations in which multiple solutions are found having the same minimum. This generally means that the problem is over-constrained or has redundant constraints.

## MORE TO COME

Part 2 of this article will provide an editor program to list and/or build or modify data files, and several problems as examples of the use of the Linear Programming Technique.

## Suggested Reading

1. GASS, S.I.: "Linear Programming Methods and Applications", McGraw-Hill, 1958.
2. SAATY, T.L.: "Mathematical Methods of Operations Research", McGraw-Hill, 1959.
3. HILLIER, F.S. \& LIEBERMAN, G.J.: "Operations Research", Holden-Day, Inc. 1974.

NOTICE: SOFTWARE ON DISK AVAILABLE
The linear programming software listed in this article is available on 5$1 / 4$ in. North Star disk for $\$ 20$. This covers the cost of disk, mailing and handing. Order it directly from Bill Yarnall - address at beginning of article.

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## LISTING I

The Linear Programming Algorithm. (Note: line numbers are not part of the program.)

```
ENTEF FILE NAME ---> LINEARF. TEXT
    FROGRAM LINEARF;
    (* MINIMIZE A COST FUNCTION SUEJECT TO CUNSTEAINTS.
            MRXIMIZE NEGHTIVE OF 'PRUFIT' FUNCTION. *)
C:ONST
    MAXROW = 32;
    MAXCOL = 64;
    TMPE
    ROW = ARRAY [ 1. . MFNROW] OF RERL;
    COL = FRRAY [1. . MFXCOL ] OF RERL;
    FREC = RECORD
            CASE TFG:INTEGER OF
                61: (NHME:STRING[6]; NUM1, NUM2:INTEGER);
                    1: (HEADER:STRING[64 ]);
                    2: <RNAME :STRING[G]; RINDEX:INTEGER; FHS:REAL);
                    4: \INAME:STRING[6]; CINOEX:INTEGER; OBJ:RERL);
                    6: (R,S:INTEGER; T:RERL);
                    99: ()
                END;
    VAR
        U : FRRHY [1. MAXROW 1. MRKROW] OF REFL;
```



```
        X, XIK : FRFRHY [ 1. MFNROW] UF REFL;
        LIST : HRRAY [1. MFKROW] OF INTEGER;
        FOUNNHME : HF'R'HY [1. . FHMF'OWN] OF STFING[E];
        COLNAME : ARR'GY [1. MEXCOL] OF STRING[E];
        FILEID : FILE OF FREC;
        FILNAM : STRING;
        F .: FREC;
        HEADING : STRING[ E4 ];
        HDRFLFG : EOOLEFN
        M,N,MP,M1 : INTEGER;
        FNHME : STRING[6];
        RESULT : INTEGER;
    FRUZEEDURE PRINTH;
    BEGIN
        WRITELN(" ');
        WRITELNG* FROG. NAME = ', FNAME);
        WRITELN(' NO. ROWSS = ',M:6);
        WRITELN('NO. COLS = ',N:6);
        WRITELN(' ')
            4E: END; (* FRINTH *)
```

47 :
48
49
FROCEDURE PRINTE;
EEGIN
WRITELN(' ');
WRITELN(' ERD FILE FURMAT');
URITELN(' ')
END; (* FRINTE *)
FRUCECUURE FRINTC(E : RUW; C : COL);

## VAR

I : INTEGER;
58 :
59
BEGIN
WRITELN(")
WRITELN( INITIAL DHTA');
WRITELN( ),
WRITELN(' OBJECTIVE VECTOR $>$;
WRITELN(" ";
FOR I:=1 TO N DO WRITELN(COLNAME[I]: $8, C[I]: 14: 8)$;
WRITELN(' ');
WRITELN(' RHS VECTOR');
WRITELN(');
FUR I:=1 TO M DO WRITELN‘ROWNAME[I]:8, B[I]:14:8);
WRITELN(' ')
END; (* FRINTC *)
PRUCEDURE PRINTD;
VAR
I, J : INTEGER:
$76:$

## 77: BEGIN

78: WRITELNK' ');
79: WRITELN(' AEAR GRRAY');
31: WRITELN(' ');
81: FOR J:=1 TO N DO
82: FGR I:=1 TOM DO
WRITELN(I:6, J:6, ROMNAME[I]:8, COLNAME[J]:8, ABFR[ I, J]:14:8);
WRITELN(');
WRITELN( $\operatorname{HEAR}(M+1)$, $\operatorname{HEAR}(M+2)$ ) ;
WRITELN(' ');
FOR I:=1 TO N DO
WRITELN(COLNRME[ I ]:8, FE:AR[M1, I ]:14:8, FGBR[MF, I ]:14:8);
WRITELN(")
END; (* FRINTD *)
FROCEDURE FRINTX;
VAR
94: I : INTEGER;
95: 5 : STRING[6];
96 :
97: EEGIN
98: WRITELN(' );
99: WRITELN(' LIST \& X ARRAYS');
106: WRITELN(');
101: FOR I:=1 TO MP DO

```
102: REGIN
103: S:='
104: IF (LIST[I]<=N) THEN S:=COLNFME[LIST[I]];
105: IF (I>M) THEN S:=ROWNNHE[I ];
106: WRITELN(I:8,S:8,LIST[I]:7,X[I]:18:8);
107: END;
108: WRITELN(' ')
109: END; (* FRINTX *)
110:
111: FFOCEDURE INITIAL;
112: VFR
113: I,J : INTEGER;
114: SUM : REAL;
115: FIRSTIN : EOOLEFN;
116: B : FRKA'Y [1. . MANEOW] UF REFL;
117: C : ARRFY [1. MFXCOL] OF REFL;
118:
119: BEGIN
120: FOR I:=1 TO MAXFOW DO
121: FOR J:=1 TO MHXCOL DO ABHR[I, J ]:=0. ©;
122: FIRSTIN:=FHLSE;
123: IF NOT EOF(FILEID) THEN F:=FILEIO*;
124: IF F.TAG = 0
125: THEN
126: BEGIN
127: FIRSTIN:=TRUE;
128: FNAME:=F. NAME;
129: M:=F. NUM1;
130: N:=F.NUM2;
131: }\quadMP:=M+2
132: M1:=M+1;
133: FRINTH
134: END
135: ELSE BEGIN PRINTB; RESULT:=2 END;
136: GET(FILEID);
137: WHILE (FIRSTIN` FND (NOT EOF{FILEID)` DO
138: BEGIN
139: F:=FILEID^;
140: WITH F [DO
141: CHSE THG OF
142: 1: BEGIN (*: HEADER *)
143: HEFDING:=HEADER;
144: HDRFLHG:=TRUE
145: END;
146: 2: EEGIN (* ROWNAMME & RHS *)
147: ROWNMME[RINDEX]:=RNFME;
148: B[RINDEX]:=FHS
149: END;
150: 4: EEGIN (* COLNFME & IEJ *)
151: COLNFME[CINDEX]:=CNNGME;
152: C[CINDEX]:=OBJ
153: END;
154: 6: FBAR[R, 5]:=T;
155: 99:;
156: END; <* CASE OF THG *)
```

```
157: GET(FILEID)
158: END; (* WHILE *)
159: IF FIRSTIN THEN
160: EEGIN
161: PRINTC(B,C);
162: FOR J:=1 TO N DO ABAR[M1,J]:=C[J];
16\Sigma: FUR I:=1 TO M DO
164: IF B[I]< 0. 0 THEN
165: BEGIN
166: 
167: FOR J:=1 TO N DO ABAR[I,J]:=-ABAR[I,J]
168: END;
169: FOR J:=1 TO N DO
170: BEGIN
171: SUM:=00, a;
172: FOR I:=1 TO M DO SUM:=SUM-HE:HE[I, J ];
17S: FEAR[MF, J]:=SUM
174: END;
175: E[M1]:=0.0.0;
176: SUM: =0. 0;
177: FOR I:=1 TO M DO SIMM:=SUM-E[I ];
178: E[MP]:=SUM;
179: FOR I:=1 TO MP DO
180: BEGIN
181: X[I]:=E[[I];
182: LIST[I]:=N+I;
183: FOR J:=1 TO MP DO U[I,J]:=0.0
184: END;
185: FOR I:=1 TO MP DO U[I,I]:=1. 胙
186: FRINTD;
187: ROWNRME[M1]:='M+1 ';
188: ROWNAME[MP ]:='M+2 ';
189: PRINTX
190: END; (* IF FIRSTIN *)
191: WRITELN(' ');
192: END; <* INITIRL *)
193:
194:
195:
196: PROCEDURE EXIT1(X:INTEGER);
197: EEGIN
198: FESULT:=1; (* NORMAL EXIT *)
199: WRITELN(' END OF FHASE 1 FUR ', FNFHE,' FFTER ', }x
200: ' ITERHTIONS');
201: FRINTX;
202: EXIT(FHHSE1)
2013: END; (* EXIT1 *)
204:
205: PROCEDURE EXIT2(X:INTEGER);
206: BEGIN
207: RESULT:=2; (* ERROR EXIT *)
208: WRITELNく' ERROR IN ITERHTION , }%\mathrm{ ;
269: FRINTX;
210: EXIT(FHASE1)
211: END; (* EXIT2 *)
```

212:
213: PROCEDURE EXIT3(X:INTEGER);
214: BEGIN
215: RESULT:=3; (* NO FEASIBLE SOLUTION *)
216: WRITELNく' NO FERSIBLE SOLUTION AFTER ', $x$, ITERATIONS');
217: PRINTX;
218: EXIT(PHASE1)
219: END; (* EXIT3 *)
$220:$
221: PROCEDURE PHRSE1;
222: CONST
223: TOL = 1. $10 \mathrm{E}-6$;
224:
225: VRR
226: ITER, I, J, L, KSFFE : INTEGER;
227: SUM, TEMP, THETA, 2 : REAL;
228: XL, XLK : REAL;
229: DEL, V, W : ARRA' [1. MAKROW] OF REAL;
230: TEST : BUOLEAN;
231:
232: EEGIN
233: WRITELN(' START PHASE 1');
234: WRITELN(' ');
235: ITER: $=0$;
236: WHILE TRUE DO (* LOOP HERE *)
237: BEGIN
238: $\operatorname{IF}(\mathrm{ABS}(\mathrm{X}[\mathrm{MP}]$ ) ( TOL) THEN EXIT1 (ITER);
239: IF (X[MP] > TOL) THEN EXITE(ITER);
249: ITER: $=$ ITER +1 ;
241: FOR J:=1 TO N DO
242: BEGIN
243: SUM: = ©. an;
244: FOR I:=1 TO MP DO
245: SUM:=SUM1+U[ MP, I ] 3 \&BERR[ I, J ];
246: DEL[J]:=SUM
247: END;
248: TEST:=TRUE;
249: FOR $\mathrm{J}:=1$ TO N DO IF(DEL[J] 0 . 0) THEN TEST:=FRLSE;
250: IF TEST THEN EXIT3(ITER);
251: TEMP:=1. 日QE +36 ;
252: KSAVE:=0;
253: FOR J:=1 TO N DO
254: IF (DEL[J]<TEMP) THEN
255: BEGIN
256: TEMF: =DEL[J];
257: KSAVE:=J
258: END;
259: FOR I:=1 TO MP DO
260: BEGIN
261: SUM: =0. 0 ;
262: FOR J:=1 TO MP DO SUM:=SUMT+U[I, J ]*ABAR[ J, KSAVE ];
263: XIK[I]:=SUM
264: END;
265: THETA: $=1$. 1 AE +36 ;
266: L: = O $^{\text {; }}$
267: FOR I:=1 TO M DO

```
    IF (XIK[ I ]>0. 0) THEN
    BEGIN
        Z:=X[ I ]/XIK[ I ];
        IF (Z=THETA) RND (LIST[I I>N)
        THEN L:=I
        ELSE
        IF(Z<THETA) THEN
        BEGIN
            THETA:=Z;
                L:=I
            END
            END;
        IF(L=0゙) THEN EXITご(ITER);
        LIST[L]:=KSAVE;
        FOR I:=1 TO MP DO
        BEGIN
            V[I]:=XIK[I ]/XIK[L];
            W[I]:=U[L,I]
        ENND;
        XL:=X[L];
        XLK:=XIK[L];
        FOR I:=1 TO MP DO
        EEGIN
        Z:=THETA;
        IF(LIST[I] < KSAVE) THEN Z:=X[I]-XL*V[I];
        X[I]:=Z;
        FOR J:=1 TO M DO
        BEGIN
            Z:=W[J]/XLK;
                IF(I<`L) THEN Z:=U[I, J]-W[J]*V[I ];
            U[I, J]:=2
        END
        END;
        WRITELN(' ITERATION ', ITER,' OF ', PNAME);
        PRINTX
        END (* WHILE *)
        END; (* PHASE1 *)
    PROCEDURE PHASE2; FORWFRD;
    PROCEDURE EXIT4(X:INTEGER);
        BEGIN
        RESULT:=1; (* NORMRL EXIT *)
        WRITELN(' END OF PHASE 2 FOR ', PNAME,' RFTER ', }X,'\mathrm{ ITERATIONS');
        FRINTX;
        EXIT(PHRSE2)
    END;
    PROCEDURE EXIT5(X:INTEGER);
    BEGIN
        RESULT:=2; <* UNBOLINDED SOLUTION *)
        WRITELN(' UNBOUNDED SOLUTION FOR ', PNAME);
        PRINTX;
        EXIT(PHRSE2)
        END;
```

323:

324 :
325 :
326: $\mathrm{TOL}=-1 . \mathrm{BE}-6$;
327
328:
329:
$330:$
331 :
352:
333:
334 :
$335:$
336: ITER: $=0$;
337: WRITELNく' STHRT PHRSE 2');
338: WRITELN(' $>$;
339: WHILE TRUE DO (* LOOP HERE *)
340: EEGIN
341: FOR J:=1 TO N DO
342: EEGIN
343: SUM: = $\overline{0}$, ;
344: FGR I:=1 TO MF DO SUM:=SUM+U[M1, I ]*FEAR[I, J];
345: DEL[J]:=5UM
346: END;
347: TEST:=TRUE;
348: FOR J:=1 TO N DO IF (DEL[J]<TOL) THEN TEST:=FRLSE;
349: IF TEST THEN EXIT4(ITER);
350: ITER: = ITER+1;
351: TEMF:=1. $\mathfrak{E E}+36$;
352: KSHVE:= $\overline{0}$;
こ5s: FGR J:=1 TO NDO
554: IF(DEL[J]<TEMP) THEN
355: EEGIN
356: TEMF:=0EL[J];
357: KSFVE:=J
558: END;
ड59: FOR I:=1 TO MP DO
S64. BEGIN
SE1: SUM: = © 0 ;
S62: FQR J:=1 TO MP DO SUM:=SUM+U[I, J ]*FEFRE[J, KSFVE ];
363: XIK[I]:=SUM
364: END;
365: TEST:=TRUE;
366: FOR I:=1 TO MP DO IF (XIK[I ] 0 0. 0) THEN TEST:=FFLSE;
367: IF TEST THEN EXIT5(ITER);
368: THETA:=1. $9 E+36$;
369: L: $=0$;
370: FOR I:=1 TO M DO
371: IF(XIK[I ]>0. 0) THEN
372: EEGIN
373: $\quad Z:=X[$ I $]$ XIK [ I ];
374: IF( $2<$ THETA) THEN
375: BEGIN
376: THETR:=Z;
377: L:=I
378: END
379: END;

```
380: LIST[L]:=KSRVE;
381: FOR I:=1 TO MP DO
382: BEGIN
383: V[I]:=XIK[I]/XIK[L];
384:W[ I]:=U[L, I]
385: END;
386: XL:=X[L];
387: XLK:=XIK[L];
388: FOR I:=1 TO MF DO
389: EEGIN
390: z:=THETA;
391: IF(LIST[I]<>KSAVE) THEN z:=X[I]-XL*V[I];
392: X[I]:=2;
393: FOR J:=1 TO M DO
394: BEGIN
395: Z:=W[J]/XLK;
396: IF (I<>L) THEN Z:=U[I, J]-W[.J ]*V[I ];
397: U[I, J]:=Z
398: END
399: END;
400: WRITELN(' ITERRTION ',ITER,' OF ',PNAME);
401: PRINTX
402: END (* WHILE *)
403: END; (* FHASE2 *)
404:
405: EEGIN (* MAIN *)
4616: WRITE(' ENTER DATG FILE NAME --->) ');
407: RERD (FILNAM);
408: RESET(FILEID,FILNAM);
469: WRITELN(' ');
410: INITIRL;
411: IF (RESULT <> 2) THEN FHASE1;
412: IF (RESULT = 1) THEN PHASE2;
413: IF HDRFLFG THEN WRITELN(' ',HEADING);
414: WRITELN(' ');
415: HRITELN(')
416: END.
LISTING 2
A sample data file named "LINTEST.DATA".
    EDIT: L(IST, BCUILD, M<UDIFY, QCUIT [1. a] L
    LIST WHAT FILE? LINTEST. DATA
    STARTING AT WHAT RECORD? 0
```

| 0 : | 0 | TEST |  |  | 3 |  | 4 |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1: | 1 | TEST |  | LINE | ERR |  | T. , 〈GASS | P. 95) |
| 2 : | 2 | ROW1 |  |  | 1 |  | 15. 0000 |  |
| 3 : | 2 | ROW2 |  |  | 2 |  | 20.0000 |  |
| 4 : | 2 | -ROW3 |  |  | 3 |  | 10.0000 |  |
| 5 : | 4 | COL1 |  |  | 1 |  | -1. 101010 |  |
| $6:$ | 6 | ROW | 1 | COL | 1 |  | 1. 101000 |  |
| 7 : | 6 | ROW | 2 | COL | 1 |  | 2. 10000 |  |
| 8 : | 6 | ROW | 3 | COL | 1 |  | 1. 80600 |  |
| $9:$ | 4 | COL2 |  |  | 2 |  | -2. 100100 |  |
| $10:$ | 6 | ROW | 1 | COL | 2 |  | 2. 010100 |  |
| 11: | 6 | ROW | 2 | COL | 2 |  | 1. 101008 |  |
| 12 : | 6 | ROW | 3 | COL | 2 |  | 2. 1101010 |  |
| 13: | 4 | COL3 |  |  | 3 |  | -3. 000100 |  |
| 14 : | 6 | ROW | 1 | COL | 3 |  | 3. 110100 |  |
| $15:$ | 6 | ROW | 2 | c CO | 3 |  | 5. 10160 |  |
| $16:$ | 6 | ROW | 3 | COL | 3 |  | 1. 100100 |  |
| 17 : | 4 | COL4 |  |  | 4 |  | 1. 101081 |  |
| $18:$ | 6 | ROW | 1 | COL | 4 |  | 6. 106006 |  |
| 19 : | 6 | ROW | 2 | col | 4 |  | -. 616100 |  |
| $20:$ | 6 | ROW | 3 | COL | 4 |  | 1. 16060 |  |
| $21:$ | 99 | LOGIIC | CAL | EOF |  |  |  |  |

EDIT: L(IST, BCUILD, MCODIFY, QCUIT [1. Q] Q

## THE MM-103 DATA MODEM AND COMMUNICATIONS ADAPTER <br> FCC APPROVED <br> Both the modem and telephone system interface are FCC approved, accomplishing all the required protective functions with a miniaturized, proprietary protective coupler. <br> WARRANTY <br> One year limited warranty. Ten-day unconditional return privilege. Minimal cost, 24-hour exchange policy for units not in warranty. <br> 

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MODEM: (703) 750-0930 (300 baud)


A sample run, using the data file
"LINTEST. DATA".

ENTER DATA FILE NAME ---> LINTEST. DATA

```
FROGG. NAME = TEST
NO. ROWS = 3
NO. COLS = 4
```

INITIAL DATA
OBJECTIVE VECTOR

| COL1 | -1. 00000 |
| :---: | :---: |
| COL2 | -2. 001010 |
| COL3 | -3. 100608 |
| COL4 | 1. 191080 |

RHS VECTOR

| ROW1 | 15. 0080 |
| :--- | :--- |
| RON2 | 20.0000 |
| ROW3 | 10.0000 |

GEAR ARRAY

| 1 | 1 | ROW1 | COL1 | 1. 96400 |
| :---: | :---: | :---: | :---: | :---: |
| 2 | 1 | ROW2 | COL1 | 2. 181400 |
| 3 | 1 | ROW3 | COL1 | 1. 00400 |
| 1 | 2 | ROW1 | COL2 | 2. 101810 |
| 2 | 2 | ROW2 | COL2 | 1. 00000 |
| 3 | 2 | ROW3 | COL2 | 2. 010060 |
| 1 | 3 | ROW1 | COL3 | 3. 00100 |
| 2 | 3 | ROW2 | COL3 | 5. 000100 |
| 3 | 3 | ROW3 | COL3 | 1. 10600 |
| 1 | 4 | ROW1 | COL4 | 0. 60040 |
| 2 | 4 | ROW2 | COL4 | 0. 00000 |
| 3 | 4 | ROW3 | COL4 | 1. 00000 |

HEAR $(M+1)$, $\operatorname{HEHR}(M+2)$

| COL1 | -1. 196180 | -4. 90000 |
| :---: | :---: | :---: |
| COL2 | -2. 10000 | -5. 1010100 |
| COL3 | -3. 00006 | -9. 010100 |
| COL4 | 1. 00968 | -1. 01010 |

LIST \& X ARRF'r'S
1 ROW1
2 ROW2
5 15. 0000
620.0000

| 3 | ROWS | 7 | 10.0000 |
| :--- | :--- | :--- | :---: |
| 4 | $1+1$ | 8 | 0.0101000 |
| 5 | $1+2$ | 9 | -45.0000 |

## STHRT FHASE 1

ITEFATIGN 1 OF TEST

LIST \& X RRFFHYS

| 1 | ROW1 | 5 | 3. 000090 |
| :---: | :---: | :---: | :---: |
| 2 | ROW2 | 3 | 4. 05065 |
| 3 | ROW3 | 7 | 6. 060006 |
| 4 | M+1 | 8 | 12. 00610 |
| 5 | $M+2$ | 9 | -9. 19514 |

ITEFATION 2 OF TEST
LIST \& X AREFTE

| 1 | ROW1 | 2 | 2.14286 |
| :--- | :--- | :--- | :---: |
| 2 | ROW2 | 3 | 3.57143 |
| 3 | ROW3 | 7 | 2.14286 |
| 4 | $M+1$ | 8 | 15.1061 |
| 5 | $M+2$ | 9 | -2.14266 |

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# AN INTRODUCTION TO CP/M* 

by
Jake Epstein
Innotronics Brooks Road
Lincoln, MA 01773

* Reg Tmk Digital Research

PART II<br>FILE STRUCTURE AND COMMAND SYNTAX

In my last article, I discussed the basic memory map of the $C P / M$ ver. 1.4 operating system plus the dialog that occurs when the system is initially "booted up". In this article I will be discussing the command syntax of $C P / M$, basic file structure, and FLOPPY DISKETTE mapping.

To begin, I will describe the layout of the FLOPPY DISKETTE so that terms and concepts that I use later will be clearer to those who are not yet familiar with this storage device. Also, for the sake of clarity and to prevent confusion, I will limit the discussion for now to 8 inch, single density diskettes. The diskette is a thin magnetic disk made up off material similar to that used in audio recording tape, and is housed in a square package that gives the disk both protection from dirt and support. When placed in a device known as a DRIVE, the disk rotates 360 times every minute and data is read from and written to the diskette via a read/record head that moves in and out depending on information supplied by host computer or device. What has made the floppy diskette so economically viable is that when actuated, the head comes in contact with the revolving magnetic material thus eliminating the mechanical difficulties associated with hard disks where heads have to be extremely close to the medium but can not touch due the the injurious effects of abrasion. This is not to say that ebrasion is not a problem for diskettes, for they can eventually wear out through normal use and dirt contamination.

Of special interest to $C P / M$ users is diskette layout. The 8 inch variety contains 77 TRACKS which are actually concentric circular areas on the disk. When disks are initially formated, these tracks are laid out, thus read/record head alignment will prove very important for accuracy of data transfer. When a
drive is commanded to read or write a certain track, the read/record head moves in or out (also known as SEEK) to find the specific track. In order to calibrate the drive, the head will perform a seek TRACK 0 upon system reset, and thus, all movement of the head can be monitored by software. The track at the outside edge of the diskette is track 0 whereas the innermost track is number 76. Each track is divided into 26 SECTORS thus the total number of sectors on the diskette is 77 * 26 or 2002 . In order to locate specific sectors, the first sector of each track, sector 01 , is indicated via a hardware indicator which senses a hole in the diskette, the INDEX. Other than TRACK 00 , and the INDEX, the type of diskette used by CP/M, SOFT-SECTORED FORMAT, has no other hardware indicators of sector and track location, and thus alignment of the read head coupled with data encoded on the disk during FORMATTING aid in sector location.

There is a type of diskette that has a series of 32 holes to indicate sector location. These HARD SECTORED FORMAT diskettes are incompatible with CP/M and should be avoided even though they can be made useable through formatting tricks. Each sector of a properly formated diskette contains both areas for data (as stored and retrieved by the user) and areas for identification and error checking. It is beyond the scope of this article to discuss formatting, but investigation of the references appended to this text will provide information on this subject. For now, all one needs to know is that each sector contains the track number, sector number, an area for the storage of 128 data characters (BYTES) and a CRC (CYCLIC REDUNDANCY CHECK) location to aid in detecting errors. Thus the available storage on each disk is:

77 TRKS * 26 SECTRS/TRK * 128 BYTES/SECTR
$=256256$ BYTES.

For the uninitiated, a byte is a standard data size of the computer industry which is 8 bits long and represents 256 different BINARY or base 2 numbers. These numbers as stored on diskette can either represent numeric data, special codes such as machine instructions, or alphanumeric characters in the form of the 7 bit ASCII code. Because diskettes are organized as discrete data structures, the 2002 sectors, special characters and/or identification headings are not needed within the data to aid in its exchange. Thus in contrast to sequential storage systems such as MAGTAPE or PAPER TAPE, the disk operating system can handle large streams of data without needing to check for beginnings and/or endings. The disadvantage of disk systems is that the minimum number of data bytes that can be read or written at any one time is dependent on sector size. In contrast, serial systems, such as MAGTAPE, can read/write one character at a time even though this is rarely done. A result of disk storage is that there will be times when storage space will be wasted, for the amount of data stored may not use up an entire sector. An advantage is that having discrete structure allows for random access to files and/or parts of files. In other words, to access a file on a tape, one has to read the entire tape to find the desired data or start of data whereas in disk systems, data can be seen as a series of physical locations. Finally, the term used for the data stored in a sector (128 bytes) is RECORD; thus a record is 128 bytes in CP/M. Also, the term FILE is used to describe a set of data. In other words, the binary data that comprises a program such as a text-editor would be stored on diskette as a file, or, the text to this article could be stored as a file.

At this point $I$ shall change the subject and discuss Command Syntax. First of all, when the entire operating system is in main memory, communication is usually accomplished via a CONSOLE device such as a CRT terminal or video-keyboard interface. Through some hardware and software manipulation, however, other devices such as modem boards for communication over telephone lines or card readers for BATCH style processing can be used. The subsystem of CP/M that directs and processes this dialog is the CCP (CONSOLE COMMAND PROCESSOR), and it does so by forming an interface between the user implemented INPUT-OUTPUT routines found in BIOS (BASIC I/O SYSTEM), and file handing routines found in BDOS, (BASIC

DISK OPERATING SYSTEM). In a future article, $I$ will discuss BIOS handler modification for different hadware configurations, and user access of BDOS routines for generation of hardware independent programs that can be run on any CP/M or equivalent system.

Conversation when $C P / M$ is in the COMMAND MODE (communicates via CCP with console device) occurs via uppercase alphanumeric characters and CONTROL functions. The lowercase alphabet is accepted, but it is converted to upper case before processing or storage, and thus user programs that leave lowercase characters in areas that are used by CCP, i.e. file names, can cause difficulties later. Also important to note is that all text is buffered until a carriage return function is received at which time a linefeed is sent to the console and then the command text is interpreted and executed. When a character is typed at the keyboard, it is read and then transmitted or ECHOED back to the screen or printer of the console. A BUFFER is an area of temporary storage that can be found any where information needs to be held, and the area of system memory in the CP/Mmemory map reserved for both console and disk file buffering is location 080 H to 0FFH. Thus a text string of 128 characters can be entered before the buffer overflows causing an error to occur. When this happens, the entire text as typed in will be automatically sent to the console followed by a "?", and when the CCP does not understand a command due to a mistake in syntax or other error, the same type of echo of text occurs.

The following list of special characters are reserved and are only used in certain situations:

```
< > . , : = ? [ ]
```

The functions implemented via the CCP are almost all file handing in nature, whereas other functions such as memory modify or single step program execution are provided for in utility programs that run under CP/M. Several of these are provided by DIGITAL RESEARCH on the original distribution diskette with the CP/M operating system, while others are available as seperate software packages. The standard CP/M utilities will be discussed in future articles.

As discussed above, since floppy diskettes are organized as a series of discrete records, then special characters such as SOH (START OF HEADING) or FS (FILE SEPARATOR) and/or identification headings that are required to monitor sequential or continuous data streams are not required. However, in certain file types, CP/M
utilizes procedures and conventions used in other systems. Below is a table of four type of files used by CP/M:

## B INARY

ASCII

RELOCATABLE
Used by certain assemblers or compilers. This is a special coding of machine programs that can be made to run at any memory location by using linking utility programs.

In the above file types, binary is the most compact and can include any kind of data. What is meant by memory image is that the file is a copy of a block of data as it appeared in computer memory thus giving the user the ability to replicate a memory state at a future time after the computer has been shut off or the memory changed. The reason why an EOF is used with ASCII files is that this gives a method for retrieving an exact copy of a stored file as to length, for without this feature, the file would have to include all of the data that was unused from the last sector as explained above. Hex format is a very versatile data type in that software generated checksums can be incorporated into the file giving a means for error checking and correction. This however, causes considerable software overhead and requires a great deal more storage space than other data types. Because Hex format can be confusing, below is an example of data as represented by different data types and hex format.

165 Decimal = 10100101 Binary = A5 Hex =

$$
\begin{aligned}
& 01000001 \text { Binary }=41 \mathrm{Hex}=\text { A ASCII } \\
& 00110101 \text { Binary }=35 \mathrm{Hex}=5 \text { ASCII }
\end{aligned}
$$

In the above the hexadecimal equivalent of the number, A5, is stored as ASCII codes so immediately it should be apparent that hex files will be at least twice as long as binary files. Checksums are generated in different ways, but usually all the data bytes in a BLOCK (a sub unit of a file usually associated with large storage devices like mag tape) are added together and then an adjusted number is stored in a non-data area of the file. When files are read and the stored checksum is different than the one generated, then an error has occured. Some software systems have the means to correct errors. At present, CP/M can only detect errors by using hardware generated CRC (CYCLIC REDUNDANCY CHECKS) which are generated in a similar manner to checksums, but error correction is not available. Thus, using file types that use checksums can prove useful for increasing reliability of disk storage. It should be mentioned that actually, CRC errors are detected by routines in BIOS, so that different disk controllers handle the errors differently.

CP/M uses a special file called the DIRECTORY to store pointers to file locations on the diskette. This directory files located in 16 sectors on track 02 . After a great deal of searching, I found that the following list of sectors contain the directory file on track 02:

```
sectors 1, 7, 13, 19, 25, 5, 11, 17, 23, 3
        9, 15, 21, 2, 8, 14 (decimal)
```

The reason that the sectors are not in order, i.e. $1,2,3$ etc., is that SKEWING is used to make reading and writing as efficient as possible. When two sectors are close together, hardware and software may not have time to identify and read/write the second one after doing so with the first before it slips by. In this situation, the system has to wait for one full revolution of the diskette for the second sector to come around again. The skewing for $C P / M 1.4$ is 6 but in other systems, it may be different causing incompatibility problems. In CP/M 2.0, this skewing can be modified by the user because this aspect of the system is handle in BIOS as opposed to version 1.4 where it is handled in BDOS. In a future article $I$ will discuss this and other enhancements found in version 2.0. A word of warning to CP/M 1.4 users. If you have a CP/M system for 5.25 in disks and wish to add 8 inch disks, you will have problems because of sector skewing and track size. The same is true for 8 inch users that want to add the smaller drives. CP/M 2.0 eliminates this problem.

The data structure that stores information about each file on the disk is
the FCB (FILE CONTROL BLOCK). Since I will be dedicating a lengthy discussion to the FCB in a future article on BDOS function calls, I will only describe a few concepts at this time. Each FCB has an area of 11 bytes in length that contains a PRIMARY and a SECONDARY name. The primary name can be any combination of up to 8 characters except for those that are reserved as mentioned above. Also, the primary name may be less than 8 characters, but when it is stored in the FCB, each empty position after the last character will be filled with ASCII "space" (20H). If a primary name is entered that is greater than 8 bytes, then it will be truncated upon storage. Names that are exactly the same as CCP function commands should not be used, for when files are accessed in the LOAD and EXECUTE function, the CCP will generate an error message ("?") because it will expect a command fucntion. Below is a list of possible primary names:

## FRTMARY NAME

FEFRESENTATON IN FCB

## TEXTEMTT

BASTC
FASCAL. 785
F- -80
1234
following names are not allowed.
PRIMARY NAME
REASON
LETTER?D
? IS RESERVED
JACOB E
SPACE IN NAME
REN IS A CPM FUNCTION

Secondary names are used to indicate certain types of files, and thus the CCP and/or utility programs can determine the data type of the file. For example, as I shall explain in next month's article, a file with the secondary name of COM can be LOADED into memory and then EXECUTED as a program. DIGITAL REASEARCH has reserved several secondary names for use in the operating system, but as software becomes more available and diverse, new reserved secondary names are added to the list. Of course, the user can use any secondary name that he/she desires, but if he/she uses a reserved name, then the file should fit the criteria of that file type. Below is a listing of the most used reserved secondary names:

NAME DATA TYPE

| COM | BINARY | PROGRAMS THAT CAN BE EXECUTED |
| :---: | :---: | :---: |
| HEX | HEX FORMAT | OBJECT OF ASSEMBLERS OR COMPILERS |
| ASM | ASCII | SOURCE CODE FOR ASSEMBLERS |
| MAC | ASCII | SOURCE CODE FOR MACRO-ASSEMBLERS |
| BAS | ASCII | SOURCE CODE FOR BASIC COMPILERS |
| FOR | ASCII | SOURCE CODE FOR FORTRAN COMPILERS |
| PRN | ASCII | PRINTOUTS OF TEXT FORMATORS, COMPILERS, ASSEMBLERS,ETC. |
| SUB | ASCII | SOURCE FOR SUBMIT UTILITY |
| \$\$\$ | - | TEMPORARY FILE MAY BE ANY FORMAT |
| LIB | ASCII | LIBRARY FILES |
| TEX | ASCII | ASCII FOR TEXT-FORMATORS |
| DOC | ASCII | MESSAGES OR DOCUMENTATION |
| MSG | ASCII | SAME AS DOC |
| TXT | ASCII | SAME AS DOC |
| REL | $\begin{aligned} & \text { RELOCAT- } \\ & \text { ABLE } \end{aligned}$ | OBJECT OF RELOCATING ASSEMBLERSSOURCE TO RELOCATING LINKERS |

There are several that $I$ left out, but in future articles $I$ will try to include as many as $I$ know of. In naming files, the primary and secondary names are written together but separated by a ".". This delimiter is not found in the FCB but represents the position between the eighth


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and ninth characters in the name block. Here are a few examples. Please remember that "-" represents space or ASCII 20H:

FTLEE NAME
FCB REFRESENTATION
BASTC.COM
FORTRANE. DOC
LETTER.I

The final aspect of CP/M file structure that $I$ wish to discuss this month is sector allocation and file size. Each file control block has space for 15 data bytes. This list is referred to as the DISK ALLOCATION MAP in the CP/M documentation, and it is used by BDOS to find the ordered list of sectors comprising the file pointed to by the FCB. After a great deal of analysis, I discovered that each position in this table represents a block of 8 sectors. Block numbers 00 and $0 l$ contain the directory as I listed above, but blocks 02 and up point to data files. BDOS also uses another byte in the $F C B$ that keeps track of the total number of records (sectors) in each file in the event that not every sector in a block is used. For example, if a file needs only 3 sectors, then the other 5 in the block pointed to by the disk allocation map are unused. This phenomemon is the same as that discussed above in reference to unused sector space. By comparing actual file size with total available space, CP/M has a means of managing disk space for small file lengths. Since there are $l 6$ positions in the disk allocation map, then the following figures can be calculated for the maximum storage capacity for one file pointed to by one FCB:

> 16 blks * 8 sectrs/blk * 128 bytes/sectr = 16384 bytes

When a FCB becomes full, the byte containing the number of records becomes equal to 80 H ( 128 decimal) and then BDOS creates (during write functions) or searches (during read functions) for a new FCB with the same file name as before. This new $F C B$ is called an extension and BDOS is able to create or read up to 15 extensions. Thus in CP/M ver. l.4, files can be created that are 16 * 16384 or 262 144 bytes in length. Of course as calculated above, this is impossible because there are only 256256 bytes of storage on a single density disk. The total amount of storage available for data is calculated below:

256256 bytes/disk

- 6656 bstes/tracks 0 and 1 (ssstem tracks)
-. . 1.24 bstes/directors rile


## 248576 bytes

Since each block in a disk access table points to 8 sectors, then this total length in bytes is 1024 (128 x 8). When files are shorter than 1024 bytes or the last block of a file is not full, then this unused space will be wasted. If
however, there were 64 files each a maximum of 16384 bytes in length as calculated above, then total storage would appear to be 64 * 16384 or 1048576 bytes. (The reason why I chose the number 64 is that the total length of a FCB as found in the directory is 32 bytes, thus each sector can contain $128 / 32$ or 4 FCBs. 16 sectors in the directory * 4 gives 64.) This is of course quite a bit more than the disk can store. Actually, the total storage space is determined by the fact that CP/M 1.4 supports 243 blocks, and since blocks 00 and 01 are used by the directory, the maximum storage on a disk when there is no unused space is 241 * 1024 or 246784 bytes. Finally, the CCP command STAT will give the unused space on a disk. This is given in 1000 byte increments thus a disk with no files in the directory will appear to have 241 K ( $K=1000$ ) instead of 246 K and the size of individual files will given to the nearest 1000 above the actual size. For example, STAT will give the size of a file of 25 6 bytes (2 sectors) as lK, (the full block).

In conclusion, I have included a great deal of information that may or may not prove useful at this time to all users of CP/M, but hopefully, it will help you to expand your knowledge of file structure and management. In next month's article, I will get into more practical matters. I hope to cover CCP command functions, list utilities and there basic functions, and space permitting, begin to discuss BIOS functions and I/O handler modifications.

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## LINEAR PROGRAMMING

## Continued from page 27

ITERATION 3 OF TEST
LIST \& $X$ RRRAYS

| 1 | ROW1 | 2 | 2.14286 |
| :--- | :--- | :--- | :---: |
| 2 | ROW2 | 3 | 3.57143 |
| 3 | ROW3 | 4 | 2.14286 |
| 4 | $M+1$ | 8 | 12. 8571 |
| 5 | $M+2$ | 9 | -0.001616188 |

END OF PHASE 1 FOR TEST AFTER 3 ITERATIONS
LIST \& X ARRA'Y'S

| 1 | ROW1 | 2 | 2. 14286 |
| :---: | :---: | :---: | :---: |
| 2 | ROW2 | 3 | 3. 57143 |
| 3 | ROWS | 4 | 2. 14286 |
| 4 | M+1 | 8 | 12. 8571 |
| 5 | $\mathrm{M}+2$ | 9 | -0.00060648 |

## START PHASE 2

ITERATION 1 OF TEST
LIST \& X ARFA'YS

| 1 | ROW1 | 2 | 2. 50060 |
| :---: | :---: | :---: | :---: |
| 2 | ROW2 | 3 | 2. 5614140 |
| 3 | ROW3 | 1 | 2. 501010 |
| 4 | M+1 | 8 | 15. 8000 |
| 5 | $\mathrm{M}+2$ | 9 | 0. 0.0661612 |

END OF PHASE 2 FÜR TEST AFTER 1 ITERATIONS

## LIST \& X RRR'H'T'S

| 1 | ROW1 | 2 | 2. 50000 |
| :---: | :---: | :---: | :---: |
| 2 | ROW2 | 3 | 2. 56101 |
| 3 | ROW3 | 1 | 2. 504010 |
| 4 | $\mathrm{M}+1$ | 8 | 15. 00610 |
| 5 | $\mathrm{M}+2$ | 9 | 9. 00806012 |

TEST OF LINEAR OPT., (GASS P. 95)

## WRITE FOR S-100 MICROSYSTEMS

We are looking for articles on serious microcomputer applications, reviews of software and hardware, and tutorials on S-100 based systems. Writing for S-100 MICROSYSTEMS is probably easier than you think. Remember, $\mathrm{S}-100$ MICROSYSTEMS is devoted to the serious microcomputer user. Since most of the other microcomputer magazines cater to beginners and game-oriented users, an S100 article will only appeal to a small fraction of the readers and may easily get lost between games, etc. junk. Therefore an article you write for SiOO-MICROSYSTEMS will get to the attention of the right people.

# ADDRESSING THE CURSOR 

## by <br> Larry Stein <br> Computer Mart of New Jersey 497 Lincoln Highway Iselin N.J. 08830

How to get more professional looking video I/O via software control in your BASIC programs. This is the first part of a two-part article.

After you have gotten over the hurdle of writing your first program, you should be able to understand that the difference between good programs and great programs is not what they do, but how easily they allow you to do it!!!

One of the things you can do to add quality to your program is to use the video screen to its fullest capabilities. This is possible if you own either a terminal with addressable cursor (e.g. Hazeltine l500, ADM3A, etc.) ,or a computer with a memorymapped video screen, such as a SOL,VDM, TRS-80, PET or APPLE.

If your computer falls into the above categories, making the video screen look very professional and easy to use is a simple matter. This is done with a technique called "direct cursor addressing".

To understand direct cursor addressing, first draw a picture of your video screen on graph paper, allowing as many boxes horizontally and vertically as your screen contains. For example, if you own a 24 by 80 terminal, make sure your graph paper screen has 80 boxes across and 24 boxes down. Number the upper-left hand box $1, l$ and the lowerright hand box 80,24. The upper-right hand box will therefore be 80,1 and the lower-left hand box will be 1,24. These numbers represent the column and row on your video screen. The 80 squares across the top are numbered l through 80 and the squares down the left are numbered 1 through 24. Therefore, any
box on the video screen can be identified by giving its coordinates: first the column (horizontal position) then the row (vertical position); i.e. when the cursor is at the lyth column across and the 9 th row down, the cursor is at $(15,9)$.

In the illustration at the top of the next page the screen columns are numbered from 1 to 80 and the video screen rows are numbered from 1 to 24 corresponding to an 80 by 24 terminal screen. Also, the following characters in the above layout are at the designated addresses.

|  | ADDRESS |  |
| :---: | ---: | :---: |
| CHARACTER | row | col umn |
|  |  | 18 |
| a | 3 | 3 |
| b | 6 | 72 |
| c | 18 | 45 |
| d | 9 | 45 |

Now we are ready to put the direct cursor addressing scheme to work in our program. I will use the BASIC language for all examples and the Hazeltine 1500, Lear Siegler $A D M-3 A$ and $S O L$ terminal computer for all the examples. If you have another terminal or computer, refer to the manual for the direct cursor addressing commands. Usually, in order to move the cursor around the screen,

11111111112222222222333333333344444444445555555555666666666677777777778 12345678901234567890123456789012345678901234567890123456789012345678901234567890

 3 4 XXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXX

 7XXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXX

 10 XxXxXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXX
 1 2 X XXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXX






 20 XXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXX 21 XXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXX
 23 24 XXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXX
you will supply some special characters and then the column and row of the location where you want the cursor to be placed.

The Hazeltine 1500 terminal requires the receipt of a lead-in character followed by an escape followed by the column and row. The Lear Siegler ADM-3A requires the receipt of a lead-in character, followed by an equal sign
followed by the column and row. The SOL terminal computer requires the receipt of a lead-in character followed by the character 2 followed by the column followed by another lead-in character followed by the character 1 and then the row.

The following table tells what decimal characters have to be sent to the 3 different terminals to place the cursor at the designated row and column.

| characters | 1st | 2nd | 3 rd |  | 4th | 5 th | 6 th |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Hazeltine 1500 | 126 | \| 17 | column | 1 | row | 1 |  |
| Lear Siegler ADM-3A | 27 | 161 | row | 1 | column | 1 |  |
| SOL terminal computer | 27 | 12 | row | 1 | 27 | 11 | column |

One additional factor in determining the row and column has to be considered. The terminals described here reference the screen relative to zero, which means that the upper left
hand position is 0,0 . The 1500 and the ADM-3A use a special formula to determine what character is output for the row and column to be addressed. Rather than go over the formula here, if
you own either of these terminals, refer to the manual, or look at the sample program there to see the technique used.

Now in order to use what we have learned about direct cursor adressing, let's look at a sample program written in Microsoft BASIC version 4.51. This program includes routines for the previously mentioned terminals and asks which one is being used. The program then sets a software switch to use the proper routine for that terminal.

In this program, the variable X is used to designate the column (Xcoordinate) and the variable $Y$ is used to designate the row (Y-coordinate) on the screen. If we want to position the cursor at screen location column 15 and row 7 we would set $X=15$ and $Y=7$ and then GOSUB to the cursor subroutine. The cursor subroutine will position the cursor at $(15,7)$, leave it there and return to the BASIC program where we will display something.

One special function is reserved for clearing the screen. If we set $X=0$ and $Y=0$ this will tell the cursor subroutine to clear the screen.

In the example program, line xxxx sets $X$ and $Y$ to 0 and then does a GOSUB to XXXX to clear the screen. The next statement in the program sets $Y$ to ll and $X$ to 14 and does a GOSUB to the cursor routine. On returning, the BASIC program prints the program title which will appear on the screen at column 14 and row 11.

Following the rest of the program, you can see that on all screens which require data input, the hoadings of the fields to be filled in are first printed on the screen and then the cursor is placed next to each field name for data input. This allows the user to see all of the fields required before entering any data. Seeing the questions before having to supply the answers is always of benefit!! If you have any doubts as to the value of this kind of screen for data input purposes, just type in the sample program and run it.

The non-computer data entry operator will find it invaluable to use this kind of screen and there is no limit to the aids your program can provide. One of these operator aids is showing the field delimiters for the data being entered on the screen as in statement XXXX. Another valuable aid is to re-display the data that was just typed in, edited into the correct format. Just re-displaying the data will help when the operator has deleted
any character in the field and BASIC has displayed the deleted characters surrounded by the familiar \ \. When you re-display the data, only the data will be printed on the screen as in statement $x x x x$.

I hope that this example and exercise will enable you to use your own imagination to make more use of the hardware you have at your fingertips. This is by no means a complete expose of all of the things that can be done with your screen on your microcomputer, but it will start you on your way away from the thinking that you have a teletype machine or other hardcopy device that can only scroll up the screen.

In the next issue of $\mathrm{S}-100 \mathrm{MICRO}-$ SYSTEMS I will discuss the following programs in detail.

If your computer has a video screen,

## USE IT!!!!!



$830 \mathrm{IF} \mathrm{I}=0$ Then 850
$840 \mathrm{Y}=\mathrm{N}: \mathrm{X}=11+\mathrm{M}$ : $\operatorname{GOSUB} 2160$
850 GOUB 2650

```
860 IF IN=13 THEN M=WD:GOTO 950 : REM CARRIAGE RETURN
890 IF IN=DT THEN 1990: REM DELETE CHARACTER
l
l
l
970 PRINT
980 NEXT N
990 PRINT % MD YOU WANT TO MAKE ANY CHANGES (Y/N) ";
*)
```



```
040 GOTO 630
S60 REM **
                                    ROUTINE to SAVE LabElS ON DISK
M0 REM **************************************************
lol
M30 PRINT "ENTER THE DRIVE ON WHICH LABEL IS TO BE STORED (A,B,C,D) n
*)
50 IF DS{"A."OR
80 PRINT
200 FILES FS
*)
*)
2300 FS=D$+2$+".LAB"
250 PRINT#1,WDS+","+LN$
1260 FOR N=1'TO LN 
1280 NEXT N
1290 CLOSE N
l300 PRINT
    320 FILES FS
    330 PRINT 
    350 REM *
    OREM *
    ROUTine to align label
    380 REM ********************************************************
#410 PRINT nDO YOU WRNT TO ALIGN THE LABELS (Y/N) ";
IF z$="N" OR z$="n" THEN 1570
*)
l
*)
500 PRINT "DO YOU NEED MORE ALIGNMENT (Y/N) ";
L/,
REM * RRM *
ROUTINE to Print Labels
l
1580 FOR N=1 TO LN
1590 LPRINT AS
1610 FOR N=1 TO SK
1630 NEXT N
```

| $\begin{aligned} & 1650 \\ & 1660 \end{aligned}$ | PRINT "DO YOU WANT TO PRINT MORE LABELS (Y/N) "; $z \$=1$ NPUT $\$(1):$ PRINT $2 \$$ |
| :---: | :---: |
| 1670 |  |
| 1680 | STOP |
| 1690 | GOтO 1680 |
| 1700 | PRINT "DO YOU WANT TO PRINT THE SAME LABEL (Y/N) "; |
| 1710 | z\$=INPUT\$ (1) : PRINT $2 \$$ |
| 1720 | IF $\mathrm{z} \$=$ "N" OR $\mathrm{z} \$=$ " n " THEN 350 |
| 1730 | gosub 2570 |
| 1740 | Gото 1000 |
| 1750 | REM ********** |
| 1760 | REM * |
| 1770 | REM * ROUTINE TO MOVE CURSOR RIGHT AND Left |
| 1780 | REM * |
| 1790 | REM ********************* |
| 1800 | I=1 |
| 1810 | IF IN=BS AND M $\left.\right\|^{1}$ THEN M=M-1: Сото 840 |
| 1820 | IF IN=FS AND MI WD THEN M=M+1:GOTO 840 |
| 1830 | $\mathrm{I}=0 \quad 1$ d |
| 1840 | GOTO 840 |
| 1850 | REM ******************** |
| 1860 | REM * |
| 1870 | REM * ROUTINE TO HANDLE CURSOR AT END OF FIELD |
| 1880 | REM * |
| 1890 | REM ******* |
| 1900 | gosub 2650 |
| 1910 | IF IN=BS OR IN=FS THEN I=1:GOTO 840 |
| 1920 | IF $\mathrm{TNS}=\mathrm{CHR}$ ( 13 ) THEN 950 |
| 1930 | Gото 1900 |
| 1940 | REM ****************** |
| 1950 | REM * |
| 1960 | REM * ROUTINE TO DELETE A CHARACTER |
| 1970 | REM * |
| 1980 | REM ************************************** |
| 1990 |  |
| 2000 |  |
| 2020 |  |
|  |  |
| 2040 | REM * ROUTINE TO INSERT A CHARACTER |
| 2050 | REM * |
| 2060 | REM ************************** |
| 2070 | Ll\$ $=$ MIID $(\mathrm{L}$, $\mathrm{M}, \mathrm{WD}-\mathrm{M}$ ) |
| 2080 2090 |  |
| 2100 | GOTO 840 (LS,M,WD-M+1) |
| 2110 | REM ****************************************************** |
| 2120 | REM * |
| 2130 | REM * UNIVERSAL CURSOR POSITIONING ROUTINE |
| 2140 | REM * |
| 2150 | REM ******************** |
| 2160 | ON AM GOTO 2220, 2360,2290 |
| 2170 |  |
| 2180 | REM * |
| 2190 | REM * CURSOR POSITIONING FOR ADM-3A terminal |
| 2200 | REM * |
| 2210 | REM ********************************** |
| 2220 | IF $\mathrm{Y}+\mathrm{X}=0$ THEN PRINT CBRS (26) ; : RETURN |
| 2230 | PRINT CHRS (27) + CHRS ( 61 ) + CHRS ( $31+\mathrm{Y}$ ) + CHRS ( 31+X) ; : RETURN |
| 2240 | REM ********************************************* |
| 2250 | REM * |
| 2260 | REM * CURSOR POSITIONING FOR THE SOL terminal computer |
| 2270 | REM |
| 2280 | REM ***************************************************** |
| 2290 | IF $\mathrm{Y}+\mathrm{X}=0$ THEN PRINT CHRS (11) ; :RETURN |
| 2300 | PRINT CHRS (27) +CHRS (2) +CHRS (Y-1) +CHRS (27) + CHRS (1) + CHRS ( $\mathrm{X}-1)$ |
| 2310 | REM **************************************************** |
| 2320 | REM |
| 2330 | REM * CURSOR POSITIONING FOR HAZELTINE 1500 terminal |
| 2340 | REM * |
| 2350 2360 | REM *\#*************************************** |
| 2360 2370 |  |
| 2370 2380 | $\mathrm{IF}_{\mathrm{y}=\mathrm{Y}+96} \mathrm{X} 32$ THEN $\mathrm{x}=\mathrm{X}+96$ |
| 2390 | $\mathrm{D}=\mathrm{D}+1: \mathrm{IF}$ D=5 THEN D=0:PRINT |
| 2400 | PRINT CHRS (126) + CHRS (17) + CHRS ( $\mathrm{X}-1)+$ CHRS (Y-1) ; : RETURN |
| 2410 | REM ************************* |
| 2420 | REM * |
| 2430 2440 | REM * ROUTINE TO DETERMINE LABEL SIze |
| 2450 | REM * |

```
4460 LINEINPUT "ENTER LABEL WIDTH (IN CHARACTERS) "; WD
480 LINEINPUT "ENTER NUMBER OF PRINT LINES PER LABEL "; LN
2490 LN=VAL(LN
510 DIM AS (LN)
\(5520 \mathrm{FOR} \mathrm{N}=1 \mathrm{TO}\) LN
\(530 \mathrm{AS}(\mathbb{N})=\mathrm{STR}\).
\(2530 \mathrm{AS}(\mathrm{N})=\)
2540
```



```
570 LINEINPUT "ENTER TOTAL NUMBER OF LABELS TO BE PRINTED \(n\); NB
\(2580 \mathrm{NB}=\mathrm{VAL}(\mathrm{NB}\) ) \()\)
2590 RETURN
630 REM * ROUTINE FOR DIRECT INPUT FROM TERMINAL
650 OUT 29,1
6660 WAIT \(29,1,0\)
670 WAIT \(29,1,0\)
IN =NP ( 28 )
680 INS \(=\) CHR (IN)
680 INS =CHR
```



```
REM * ROUTINE TO GENERATE PRODIGY DISKETTE LABELS * *
奖 \(=0: \mathrm{X}=0\) :GOSUB 2160
PRINT "ENTER DISKETTE NUMBER XXXX"; \(\operatorname{STRINGS}(4,8)\)
LINEINPUT DSS
```



```
2790 LINEINPUT UNS
2800 PRINT "ENTER DATE (MM/DD/YY) \(\mathrm{xx} / \mathrm{xx} / \mathrm{xx}\) ";STRINGS \((8,8)\);
2810
LINENPUT DTS
2820 PRINT "ENTER DEALER NAME \(\quad x x x x x x x x x x x x x x x x x x x x x x x x x x x x x x\) ";STRINGS(30,8);
840 IF LEN(DLS) 30 THEN PRINT "DEALER NAME TOO LONG ! ":GOTO 2820
2850 ERASE AS
```



```
2890 AS (3) \(=\) " " "
2900 AS (4) \(="\) DEALER: "
2910 AS (5) \(=\) " MASTER DISKETTE - RETURN IMMEDIATELY"
\(2930 \mathrm{AS}(7)="\) COPYRIGHT (C) 1979 PRODIGY SYSTEMS, INC.
2940 AS ( 8 ) =" ALL WORLDWIDE RIGHTS RESERVED"
950 LINEINPUT "ENTER TOTAL NUMBER OF LABELS TO BE PRINTED "; NB
\(2960 \mathrm{NB}=\mathrm{VAL}\) ( NB S \()\)
\(2980 \mathrm{WD}=\mathrm{VAL}\) (WDS)
\(3000 \mathrm{LN}=\mathrm{VAL}\)
\(3010 \mathrm{LNS})\)
\(3010 \mathrm{SN}=\mathrm{VAL}=1{ }^{11}(\mathrm{LNS})\)
020 SK=VAL (SKS)
030 GOTO 1390
3040 GTTOP
3040 STO
3050 REM *
ROUTINE FOR RETRIEVING LABELS FROM DISK
0.070
REM
```



```
3100 PRINT "ENTER DRIVE ON WHICH LABEL IS STORED (A, B, C,D) ";
2120 IF ZS["A" OR 2\$|"D" THEN 3100
\(130 \mathrm{FS}=2 \$+"\) : .LAB"
3140 PRINT
3150 FILES FS
3170 PRINT "ENTER A FILE NAME FROM THE ABOVE LIST"
180 LINEINPNTER "USE FILE NAME FROM THE ABOVE LIST"
\(3190 \mathrm{FS}=2 \mathrm{~F}+\mathrm{H}^{\prime \prime} \mathrm{LAB}\) "
3210 OPEN "III, 1,FS
3220 INPUT\#1,BS,CS
```



```
3240 ERASE AS
```


# THE S-100 BUSNEW VERSUS OLD <br> <br> by <br> <br> by <br> <br> Sol Libes 

 <br> <br> Sol Libes}

## A comparison between the old Altair S-100 bus standard and the new IEEE S-100 bus standard.

The IEEE S-100 bus standard is being voted on now and will no doubt be adopted by the time the next issue of |S-100 MICROSYSTEMS| appears in print. The complete proposed standard was reprinted in the previous issue of this magazine.

There are approximately 200,00 "old" S-100 systems presently in operation. Many owners of these systems will be upgrading these systems with new CPU and peripheral boards to keep in step with the changing technology. How compatible will these new "IEEE S-100 compatible" boards be with "old" $S-100$ mainframes?


#### Abstract

The IEEE $S-100$ standard has defined many bus lines which were not previously defined and redefined some pins. The roblem has been compounded by the fact that many $S-100$ board manufacturers took liberties with the loosely defined Altair bus and created some of their own pin functions. For example, pin 67 (now defined as PHANTHOM*) was used for as many as eight different functions by various $\mathrm{S}-100$ manufacturers.

First of all lets compare the difference between the features of the old $\mathrm{S}-100$ bus and the new IEEE $\mathrm{S}-100$ bus. They are shown in Table 1.


TABLE 1
S-100 FEATURES

| IEEE S-100 | ORIGINAL S-100 |
| :---: | :---: |
| 1. Designed to support more powerful 8-bit (z80) and 16-bit (e.g. Z8000) mpus | 1. Designed specifically for 8080 |
| 2. Operating speed up to 10 MHz | 2. Operating speed - 2 MHz |
| 3. 16 Megabyte direct memory addressing range | 3. 65 K direct memory addressing range |
| 4. Up to 65 K I/O ports | 4. Up to 256 I/0 ports |
| 5. Up to 16 Masters on bus | 5. Single Master operation |
| 6. 8 or 16 -bit data transfers | 6. Only 8-bit data transfers |
| 7. 10 vectored interrupts plus NMI | 7. 8 vectored interrupts |
| 8. Phantom |  |
| 9. Three more ground lines |  |
| 10. 3 undefined lines | 19 undefined lines |
| 11. 5 reserved for future use lines |  |

The Data Bus
The data bus has been changed so that it can support both 8-bit and 16-bit word transfers. This is shown in Table \#2. Thus the user can insert either 8 -bit or 16 -bit CPU cards into the bus. In fact, the user can operate with both 8-bit and 16-bit processors on the bus via the DMA protocol. For an explanation of the operation of the data bus refer to last months article on the IEEE S-100 standard.

The Address Bus
The address bus has been greatly expanded from 16 direct address lines for 55 K of memory and 256 ports for 8-bit processors to 24 direct address lines for up to 16 Megabytes of memory
and up to 65 K of $\mathrm{I} / 0$ ports for l6-bit processors. This is shown in Table \#3.


#### Abstract

The Status Bus A comparison of the old and new status bus is shown in Table \#4. The number of lines has been reduced from ll to 9 functions. Three status signals have been eliminated, namely STACK, RUN and SS (now used for ERROR*, RFU and NDEF, respectively). One new status signal has been added: sXTRQ*. Further, status signal labels now start with a lower-case "s" letter (except MWRT) while the old S-100 bus used and upper-case "S" lettr. Note also that an active low state is nowwindicated by an "*" compared to the over-bar in the old S-100. This was done because few printers can put an over-bar above a printed character.


|  | TAB | $\begin{array}{r} \# 2 \\ \text { BUS } \\ \hline \end{array}$ |
| :---: | :---: | :---: |
| IEEE S-100 |  | OLD S-100 |
|  | 16 lines | 16 lines |
| 8-bit Master | 16-bit master | 8-bit master only |
| 8-Data In | 16-Bidirectional | 8-Data In |
| 8-Data Out | \& | 8-Data Out |
|  | 8-Bidirectional |  |
|  | $8 \text {-Da } \frac{\text { or }}{\text { ta }} \text { In }$ |  |
|  | 8-Data In |  |

TABLE \#3

| ADDRESS BUS |  |
| :---: | :---: |
| IEEE S-100 | OLD S-100 |
| 16 or 24 lines | 16 lines |
| ( $\mathrm{A} 0-\mathrm{A} 23$ ) | (A0-A159) |
| 16 Megabytes memory | 65 K bytes memory |
| directly addressable | directly addressable |
| I/O addressing | I/O addressing |
| up to 65 K ports | up to 256 ports |
| (A0-A7 \& A8-A15) | (A0-A7 or A 8 -A15) |

TABLE \#4
STATUS BUS


Control Output Bus
The control output bus is shown in Table \#5. These signals determine timing and movement of data during a bus cycle. Two lines have been eliminated; PWAIT and PINTE (now defined as RFU). One new signal has been added; pSTVAL (previously this line was used for the Øl clock). A lower case "p" is now used instead of the upper case "P" to denote a control line.

Control Input Bus
The control input bus is compared in Table \#6. These signals synchronize slave to master operation.

Two new lines have been added; NMI* and SIXTN* previously these lines were not defined).

Vectored Interrupt Bus
\#7. This bus is compared in Table conjunction with the INT* signal and a vectored interrupt controller circuit. Two new signal lines have been added (using previously undefined lines).

TABLE \#5
CONTROL OUTPUT BUS


TABLE \#6
CONTROL INPUT BUS


TABLE \#7
VECTORED INTERRUPT BUS

| IEEE S-100 |  | 01d S-100 |
| :---: | :---: | :---: |
| 10 lines |  | 8 lines |
|  | (VI 0* | VIO |
| Inputs to an interrupt con- | VII* | $\overline{\text { VII }}$ |
| troller circuit which arbi- | VI 2* | VI2 |
| trates among 8 inputs and | VVI 3* | VI3 |
| asserts INT*. | VVI 4* | VI4 |
| VIO has highest priority | VI 5* | VI5 |
|  | VI 6* | VI6 |
|  | VI 7* | VIT |
| new lines |  |  |
| Asserted when error occurs | ERROR* |  |
| Asserted when impending power failure occurs | PWR FAIL* |  |

Utili.ty Bus
The utility lines, compared in Table \#9, have reduced in number from 34 to 22. This has occurred primarily through the use of many of the 19 previously "undefined" lines. The Øl clock signal was deleted since it was rarely used and has no meaning for most microprocessors which use only one clock signal. The PROT, UNPROTECT and PS functions have fallen into dis-use by present memory board manufacturers. The sense switch disable (SSW-DSB) has similary fallen into dis-use and has been replaced by GND and RFU lines.

On the other hand the ground lines have be increased from 2 to 5
lines to decrease the impedance of the ground circuit. Further, the location of the ground lines affords a small improvement in shielding between lines.

DMA Control Bus
This bus, compared in Table \#8, has been greatly expanded from 4 to 8 lines, to allow for multimaster operation. These lines (DMAD through DMA3) were previously undefined. The DMA control lines are used in conjunction with the HOLD* and pHLDA lines.

TABLE \#8
DMA CONTROL BUS

| IEEE $\mathrm{S}-100$ |  | Old S-100 |
| :---: | :---: | :---: |
| 8 signals lines |  | 4 signal lines |
| Address disable | ADSB* | ADD DSB |
| Data Out disable Disable PM | DODSB* | DO DSB |
| Status disable signal drivers | SDSB* | STA DSB |
| Control disable <br> new lines | CDDSB* | C/C DSB |
| arbitrate among up to 16 masters encoded priority requests are | $\left\{\begin{array}{l}\text { DMA 0* } \\ \text { DMA }\end{array}\right.$ |  |
| asserted \& lines contain priority no. of highest requestor | $\left\{\begin{array}{l}\text { DMA 2* } \\ \text { DMA 3* }\end{array}\right.$ |  |

TABLE \#9
UTILITY BUS

| IEEE S-100 |  | Old S-100 |
| :---: | :---: | :---: |
| 22 lines |  | 34 lines |
| Clock signal | CLOCK | CLOCK |
| Master clock signal | $\varnothing$ | $\emptyset 2$ |
| Resets all masters | RESET* | PRESET |
| Resets all slaves | SLAVE CLR* | EXT CLR |
| Power-on clear | POC* | POC |
| Overlays slaves | PHANTOM* |  |
| Not defined | NDEF (3 lines) | NDEF (19 lines) |
| Reserved for future use | RFU (5 lines) |  |
|  | +8V(2 lines) | +8V (2 lines) |
|  | +16V | +16V |
|  | -16V | $-16 \mathrm{~V}$ |
|  | GND (5 lines) | GND (2 lines) deleted |
|  |  | $\varnothing 1$ |
|  |  | PROT |
|  |  | UNPROT |
|  |  | PS - protect status |
|  |  | SSW-DSB |

## PRODUCT REVIEW

# THE CGS-808 INTELLIGENT COLOR GRAPHICS BOARD 

by<br>Jon Bondy<br>Box 148<br>Ardmore, Pa. 19003

Although most of the energy which I have applied to my computer system has been towards such mundane things as CRT's and printers, my main interests in having a home computer center around graphics and music. A few high density graphics boards have been available for the $S-100$ bus for some time (like the Matrox board), but they have been for the most part both inflexible and more expensive, so I devoted time to other projects and waited for the arrival of an inexpensive grahics board. In the fall of 1979, an ad for an Intelligent Color Graphics board appeared in BYTE from a firm of which $I$ had never heard, called Biotech Electronics. They offered a kit for $\$ 99$ which included the more expensive IC's along with a PC board and instructions. Somehow, this one seemed worth the risk, so I purchased it.

The board, called the CGS-808, uses the Motorola 6847 CRT controller chip to provide a wide range of low and high desnity color graphics, along with alphanumeric characters in a 16 by 32 character format. It employs an on-board 8085 processor to do the graphics bookkeeping, allowing the main processor in the system to attend to other matters. Up to two onboard 2708 (or 2716) EPROMs are used to store programs which make the CGS-808 an intelligent peripheral device. Unlike some other graphics boards, the video refresh memory resides entirely on the CGS808, making additional purchases of memory boards unnecessary, and leaving the address space of the main processor free for other uses. You can set up the 6847's modes, clear the screen, draw a dot, draw a line, and read or write the graphics refresh memory with simple commands from the central processor using their Firmware Pack I, and their other Firmware packages allow more complex graphics to be performed by the on-board 8085. The 8085 and the 6847 take turns accessing the
video refresh memory, so there is no "snow" on the graphics screen while it is in use.

My primary interest in graphics is to do high density line drawings, and the CGS-808 provides enough density to start to do some serious work in its highest density mode with 256 by 192 dots on the screen. At this density, you give up the range of colors which are offered with

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lower density modes (from 64 by 32 display elements on up), but I intended to use the board with a black and white monitor anyway, so this was not a great loss to me.

The bare board kit arrived quickly and consisted of a set of Hardware and Software manuals, the PC board, and four ICs (8085, 6847, 2708 and 1372). Biotech does not warrantee these ICs because they claim that they may be damaged by static by an inexperienced user. I have never had problems with static so far, and I doubt that this would be as large a problem as Biotech claims. In any event, most other manufactureres will replace ships which prove to be defective, whether they are MOS or not. The only defective chip which $I$ found in the entire process of building my kit was a 74 LS 74 which was purchased from another vendor. Total parts cost for the board (aside from the kit) was about $\$ 150$.

The board was as well made as I have seen in Hobby products, with IC locations and types silk-screened over a good solder mask. Complete discussions of board theory (both hardware and software) were contained in the manuals, along with step-by-step assembly instructions and sample programs. The discussions on board use and the 6847 were at times somewhat cryptic, but all of the information is there, and the sample programs were very useful. One surprise here was that their section on debugging the board was four pages long, since many vendors omit that section entirely. Complete schematics for the board were included, but there was no listing of the Firmware EPROM, so that when the board did not work immediately, it was a bit difficult to determine if it was a software or hardware problem. Biotech has informed me that a complete listing of the ROM will be made available at a nominal charge by the time this article is published.

Acquiring the hardware to complete the kit took a bit longer than had been anticipated, with some of the 74 LS 300 series ICs being the major problem. The Hardware manuals list the parts by schematic resistor, capacitor and IC number first, and then provides a summary for parts ordering, which was very useful. The RAM on the board consists of $2 l l 4 \mathrm{~s}$, so they were plentiful and reasonably priced.

Before I discuss my debugging problems, let me describe how the main processor communicates with the CGS-808. The CGS-808 reserves four input and four output ports for communication between the central processor and the on-board processor, and these ports can be placed anywhere in the range of 256 port addresses as long as they start on a port number which is a multiple of four. Of the four ports, only two are actually used by the board, with two unused. It turns out that it is possible to use one of
these unused ports for an interesting purpose, which I will get to later. Let's assume that we have set the ports up to start at port-0.

The two ports which the CGS-808 normally uses serve different purposes. Writing to port-0 serves to send both OP codes and parameters to the 8085. The OP codes are numeric values which indicate how the parameters which follow are to be processed, and they consist of 0 (clear screen), 1 (set mode), 2 (plot point), 3 (draw line), 4 (alphanumeric/semigrahic characters), 5 (read screen memory) and 6 (write screen memory). The parameters vary with each of the OP codes as appropriate; for instance, the draw line $O P$ code (3) requires starting $X$ and $Y$ coordinates, ending $X$ and $Y$ coordinates, and a line color.

When the OP code and its associated parameters have been sent to port 0 , writing any value at all to port-l starts the 8085 off executing the requrested operation. the data sent to this port is immaterial, and is ignored.

Reading port-0 provides you with some statusinformation such as whether an invalid OP code has been received, whether the board is so confused by the sequence of data that you should reset it (hardware reset), whether it thinks it is imputting parameters or not, and whether it is ready for a new command or is still executing the previous one. If you are executing a screen memory read, you read the screen data from this port also.

Reading port-l provides some more status, including whether the most recent OP code or parameter has been accepted by the 8085 yet or not, and whether there is output data in the data register (port-0) during a screen memory read operation.

When the board is first powered up, the 8085 is reset by on-board circuitry and it then sets the board's mode to high density graphics and clears the screen. My board did this when powered up, so I knew that the ROM, 8085, and video chips were functioning properly. Since the screen was an even shade of grey, it seemed likely that the refresh memory ws working properly also, so $I$ was encouraged. Unfortunately, I could not get the proper status bits to show up on port-0, so I suspected that the Biotech Rom was in error; perhaps if I had known the company better, I would not have wasted my time in this particular direction.

At times the video board would spontaneously give me a "light show" for minutes on end, while at other times it would do nothing at all no matter what I did with the ports. After setting up some software loops to write to ports 0 and 1 repeatedly, it became clear on the 'scope that not all of the port writes were setting through to the video board. I isolated it to a flip-flop which was not functioning properly, and the board worked
as soon as a new 7474 was installed.
I now came to another problem area; one which I had not expected. A diagonal line drawn across the screen was jagged and uneven. After some experimenting, it turned out that all of the odd pixels (picture elements, or dots) were less wide than the even pixels. Again, I first thought that it was a software error (in the line drawing algorithm) on Biotech's part. I tried to write a bit pattern into the refresh memory to see if the problem would also appear with this simple input. When it continued to produce pixels of uneven size, I assumed that it was a hardware problem. A call to Biotech revealed that sometimes the duty cycle of the color burst oscillator was not exactly 50\%, which caused some pixels to be clocked onto the screen more rapidly than others. The former pixels would then be less wide than the latter. Putting a lokohm resistor from pin 3 of the 1372 to ground fixed than problem.

During the debugging phases, the CGS808 had often thought that it needed a reset because the sequence of OP codes and parameters was not what it had expected (either because of a hardawre problem or my misunderstanding of the way the board worked). Unfortunately, I was debugging the board using my UCSD PASCAL Monitor program (which will appear in the next issue of this magazine), and pressing reset caused me to have to reboot. It seemed that if the software was aware that a board reset was required (see the status bits described for port-0 above), it should also be able to reset the board in software. What was required was a way to generate a low (0 volts) signal on the board on command from the central processor.

I looked for $I / O$ ports from which I could steal a signal, but they are all 'smart' ports which hold the data inside themselves until strobed by the 8085, and so could not be used. There were, however, strobe signals for each of the two ports which the CGS-808 had reserved but not used, and these were normally high but went low when the port was addressed. I connected the strobe for writing to port 2 (U48 pin 5) of the processor's reset circuitry, but it did not work. I then removed the 3.3 microfarad reset capacitor (C22) so that my short strobe would not have to fight the capacitor, and the trick worked. I now can reset the board by writing to port-2.

I did not want to lose the power-on reset and the ability to clear the board from the front panel reset switch, so I installed a diode between the $\mathrm{S}-100$ bus's reset line and the 8085's reset circuitry so that a board reset would not cause a system reset, but a system reset would cause a board reset. Addition of a small capacitor should also restore the power-on reset.

The board's speed is adequate but not high enough to allow something like real time 3-D rotations of complex objects. I asked Biotech if there was a way to increase the speed of the board and they said that if $I$ had purchased fast enough RAMs (about 300 nsec ) I could replace an RC network which drove the 8085's clock with a 6 MHz or 7 MHz crystal. This would increase the board's speed by about a factor of three. In addition, if $I$ was willing to have snow on the screen during processing, I could take pin 1 of $U 26$ (a flip-flop reset pin) and hold it always low, defeating the bus arbitration circuitry which prevents the 8085 and the 6847 from competing for the refresh memory bus. I have not done this yet, but they indicate that approximately a 40\% increase in speed could be expected from this modification. If you make this mod, be certain not to ground the other end of the run to U 26 pin-l, as it is an output pin on the 6847. Cut the lead before grounding U26 pin-l.

The 6847 has an annoying feature which is that it places a border around the graphics area of the screen. In high density mode, that border can only be white, making for some synch problems on my monitor and making it impossible to see lines drawn along the border in white. There can also be problems with monitor bloom

## THE S-100 MICROSYSTEMS SOFTWARE CATALOG

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# TARBELL DISK CONTROLLER MODS For Ithaca Audio Z80 CPU <br> by <br> George Lyons 280 Henderson St. Jersey City, N.J. 

To use the Tarbell floppy disk contoller in my system some modifications were required for it to work at all and others were desirable for better performance. The attached schematic shows the changes should anyone wish to use them. The diagram is purely schematic, so some circuit disconnections shown actually require more than one trace cut.

1. PWAIT compatibility--With the Ithaca Audio or similar Z80 cpus, PWAIT is not generated exactly like the 8080 and is mostly useless; it cannont be fed back to PRDY as on the Tarbell board, so U57 (diagram lower right) is disconnected and reused below.
2. Bootstrap PROM in Z80 systems--Z80 cpu cards which do not multiplex status on the data lines like the 8080 cannot use the standard Tarbell bootstrap PROM. One way it can be used is to put a new bus signal $\overline{S M E R}$ on an unused bus line "Xl" (top left). This signal is the standard line 47, SMER, inverted and not disabled by STATDSB. It allows enabling the PROM and disabling system memory only during memory read cycles, even though SMER itself becomes disabled. On the Ithaca Audio cpu $\overline{\text { SMER }}$ can be made by connecting ICl8 pin 14 to the bus with driver IC 39 pins 9-10, in which case it is conveniently disabled in the HOLD state even though insensitive to STATDSB, to prevent any unwanted reset of the boot flip-flop U34 from a HoLD. Status signal Ml at U43 pin 1 is replaced by reinverted SMER. PSYNC itself, instead of gated by 01 is used at U43 pin 1 for $Z 80$ cpus synthesizing PSYNC not precisely aligned with 01.

An alternative method eliminating the extra bus signal $\overline{\text { SMER }}$ is shown in fig. 2. A 7474 flip-flop is added at socket U58 and used to latch SMEMR on the rising edge of PSYNC, and to reset on the falling edge of PDBIN, "remembering" SMER while it is disabled by STATDSB. $\overline{\text { PHLD }}$ gates PSYNC at $U 28$ to clear U58 during DMA operations.
3. Separate Reset and Bootstrap
 bootstrap PROM when hitting RESET, Tarbell provides a dip swithch on the controller. To more conveniently do so, from a keyboard, the boot flip-flop U34 is rewired as shown with input BOOT placed on another unused bus line "X2" (high to engage the PROM). This also allows use of a ROM monitor with its own jump-on-reset circuit (such as on the Ithaca Audio cpu) by using X 2 to disable the ROM's jump circuit when high. A panel switch or keyboard key can then select which process is desired upon cpu reset. When a keyboard is interfaced through a MERLIN vide board inverting the active-low EDIT key is convenient.
4. Reset logic at $4 \mathrm{mhz}-$ The original Circuit pulses the Master Reset pin of the 1771 when starting the bootstrap. I find my 4 mhz cpu tends to execute the first bootstrap instruction, IN WAIT, before INTRQ has reset, causing Lost Data errors. This is partly due to MR being intiated on the trailing, rather than the leading edge of $\overline{\text { POC. Or gate }}$ U46 was therefore added to prevent Master Reset on booting and the reset circuit was simplified to use only POC instead of Preset and Extclr as well. U57 (upper left) provides a buffered POC to all functions. Master Reset could also be fixed with a one-shot but the spare one-shots are used below. Two strikes of the RESET key, one with and one without Master Reset are needed to boot in a 4 mhz system.
5. Merlin Video Compatibility-- This video bōard, no longer manufactured, has no on board display memory but uses DMA to refresh the screen--and must be deactivatted during disk transfers. To avoid doing so with software, the 1771 HLD signal is placed on a third spare bus line "X3" (lower right) and jumpered to the PRIORITY-IN pad on the Merlin. The Merlin Priority-In circuit may also be modified to disable immediately instead of waiting until completion of any DMA in progress.

6. Automatic Head-load Timing--The 1771 inserts a 10 ms deIa $\bar{y}$ in commands requesting the head to load so that for optimal efficiency the head load command bit must be omitted under software control when the head is already loaded from previous use. This loms delay is actually unnecessary because the head settling delay is really provided by the one-shot U4l and the HLT pin on the 1771, and is around 50 ms . Use of separate head-loading and non loading commands can be eliminated by activating the 1771 TEST pin; it eliminates the 10 ms delay so all commands can be of the head-loading type. However, TEST also speeds up head positioner step pulses to an unusable rate. This is cured by applying SEEKCOMPLETE to the TEST pin (lower left).
7. Dual Head Load--When switching back and forth bet ween two drives the Tarbell board sends the head load signal to only one drive at a time. By rewiring U6l and U63 as shown and replacing U63 by a NAND gate, both heads can remain down simultaneously. Neither head will load until it is selected, but once loaded all remain so until the 1771 HLD signal goes off. The spare one-shot U4l is added to the original HLT circuit to provide settling delay when engaging a second head. Note the reversal of the drive connections to pads E20, El9 resulting from the NAND gate replacement U63.
8. PerSci compatibility--PerSci drives with voice-coil head positioner can be operated in a "fast seek" mode which does not use the 1771 STEP signal but software generated high speed pulses
from pad El4 instead. The Tarbell board provides for permanently jumpering the controller in this fast-seek mode or in normal mode. Since the software for positioning the head is different in these two configurations modification was done to allow both types of software to use the board without changing jumpers. The PerSci drive accepts both types of STEP signal in any case. The seek mode is selected by bit Q4 of the function register U40, which also controls the source of the WATT state generated when reading the WAIT port (FC). In fast seek mode SEEK COMPLETE generates the WAIT state for detecting completion of a fast mult-track seek, while in slow-seek mode the 1771 signals completion of seeks with the usual INTRQ signal. The circuit shown lower right allows either pad El4 or 1771 STEP to drive the step line in nomral mode and disconnects the 1771 STEP signal in fast mode (necessary because a STEP pulse is generated in using the 1771 to set the DIRC signal level). Gates U28 and U29 shown are reused from other mods above where they were fed; alternatively spare AND gate at U36 and the added OR gate at U46 could be used.

Conclusion--Those are all the modifications. One more circuit might be added if the controller were entirely redesigned, but requires half a dozen gates. This would be to provide WAIT stat synchronization on the DATA port rather than on the WAIT port, permitting Z80 block I/O instructions to do disk transfers. It is messy because the 1771 is configured to use the DRQ and INTRQ signals as interrupts, not READY signals, and is not shown.

;THIS MODULE CONTAINS ALL THE INPUT/OUTPUT ROUTINES FOR THE CP/M SYSTEM, INCLUDING
THE DISK ROUTINES. THE DISK ROUTINES.
;THIS VERSION SUPPORTS THE FOUR DRIVE SYSTEM OF CP/M 1.4. IT CAN ALSO BE USED WITH ;THE ONLY CHANGE NEEDED IF YOU ARE USING 1.3 IS
;TO CHANGE THE BDOS EQUATE FROM "CBASE +3106 Hn ; TO CHANGE THE BDO
******** THIS BIOS REQUIRES 1 K MORE THAN THE 512 BYTES ; TO RUN THIS BIOS YOU WILL HAVE TO GENERATE A CPM THAT ; IS 1 KNLESS THAN THE AMOUNT OF MEMORY THAT YOU WANT IT
 ;TO MOVCPM OR RELOC ( 23 FOR 24K SYSTEM). USE THE IF YOU WANT VDM AS STD OUTPUT-"PTVDM" SHOULD BE TRUE. IF YOU WANT STANDARD OUTPUT ROUTINE-" PTVDM" SHOULD BE FALSE. >>>>>>>>>>>>>>>>>>>>>>>>>>>>>>>>>>>>>>>>>>>>>>>>>>>>>>>>>>>> "MYSYS" MUST BE FALSE <<<<<<<<<<<<<<<<<<<<<<<<<<<<<<<
$\qquad$

THE DISK DRIVER ROUTINES IN THIS BIOS ARE DEPENDENT ; ON EACH OTHER FOR CERTAIN INFORMATION AND TIMINGS. ONLY BE DONE AFTER CAREFUL STUDY

THIS VERSION CONTAINS SUPPORT FOR FAST READING AND
WRITING ON A NON-INTERLEAVED BASIS. THAT IS WREAD OR WRITE NON-INTERLEAVED BASIS. THAT IS, IT CAN ; THE DISK if the reads or writes are issued in ,
if you have only one disk change "onedsk" to true ; ERROR MESSAGES are printed in the following format ;xaAA twn snn $Z$

```
WHERE X = (R)READ, (W) WRITE
            WPY - NOT READY - BIT 7
            WPR - WRITE PROTECT - BIT 
            W, WRITE FAULT - BIT 5
            CRC - CRC ERROR - BIT 3
                DRQ - DATA RATA - DEST
            TNN = TSY - BUSY - BIT O
            SNN = SECTOR WHICH CONTANS EROROR (DECIMAL)
THE BIT POSITIONS ARE GIVEN SO THAT THE 1771 MANUAL
---WRITTEN BY M. D. NICHOLS--
```

The following program is an
improved BIOS (Basic Input/Output
System) for CP/M when used with a
Tarbell single density disk controller.
It offers several improvements over the
original BIOS provided with the Tarbell
Disk controller. For example, it will
support 4 disk drives, compared to 2
previously and provides faster disk I/O
operations. Further, it includes a五

100 Guy Street
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## SIOYDIN U!fIEW

$$
\begin{aligned}
& \text { VDM- } \\
& \text { source }
\end{aligned}
$$

compurir, inzinuas

momucom wos

| ;FREQUENTLY Changed parameters |  |  |  |
| :---: | :---: | :---: | :---: |
| True |  |  |  |
| TRUE | EQU | 0 FFFFF H | ; USED IN IF STATEMENTS |
| false | EQU | not true | ;USED IN If Statements |
| PTVDM | EQU | true | ;VDM FOR CONOT? |
| ONEDSK | EQU | false | ; USE If you only have 1 drive |
| TEST | EQU | true | ; TESTING AIDS |
| MSIZE | EQU | 47 | ;MEMORY SIZE |
| MODIF | EQU | 'G' | ; Modif level of driver |
| SUBmod | EQU | '4. | ;SUB MODIFICATION LEVEL |
| MYSYS | EQU | TRUE | ; SPECIAL STUFF FOR MY SYSTE |
| jbise | EOU | 0F8H | ; BASE FOR CONTROLLER PORTS |
| DCOM | EQu | DBASE | ;DISK COMMAND PORT |
| dStat | EQU | dBASE | ;DISK Status port |
| DTRK | EQU | DBASE+1 | ;DISK TRACK PORT |
| DSECT | EQU | DBASE+2 | ;DISK SECTOR PORT |
| dDATA | EQU | DBASE +3 | ;DISK DATA PORT |
| DCONT | EQU | DBASE+4 | ;DISK WAIT PORT |
|  | EQU | DBASE+4 | ;DISK CONTROL PORT |
|  | ORG | MSI2E*1024-512 | ;THIS SETS ADDR OF CBIOS |
| cbase | EQu | (MSI2E-15)*1024 | ; temporary calculation |
| CCP | EQU | CBASE +2900 H | ;START OF CP/M SYSTEM |
| bDOS | EQU | CBASE +3106 H | ; START OF bDOS FOR 1.4 |
| CPML | EQU | \$-CCP | ;LENGTH OF CP/M W/O CBIOS |
| NSECTS | EQU | CPML/128 | ; SECT CP/m excl. CBios |
| RTCNT | EQU | 10 | ;RETRY COUNT ON ERRORS |
| ; |  |  |  |
|  |  |  |  |
| ; | JMP | B00T | ; FRom cold start loader |
| Xwboot | JMP | WB 00 T | ; FROM WARM START |
| XCONST | JMP | CONST | ; CHECK CONSOLE STAT |
| XCONIN | JMP | CONIN | ; GET CONSOLE Character |
| XCONOT | JMP | CONOT | ; Write console character |
| XLIST | JMP | LIST | ; Write char on list dev |
| xpunc. | .TMP | PUNCH | ; PUNCH Char on paper tape |
| XRDR | JMP | READER | ; READ Char from paper tape |
| хноме | JMP | home | ; RESTORE HEAD TO TRK 0 |
| xDSKSEL | JMP | DSKSEL | ;SELECT DISK DRIVE |
| xsEttrk | JMP | SETTRK | ;SET TRACK NUMBER |
| xSETSEC | JMP | SETSEC | ; SET SECTOR NUMBER |
| XSETDMA | JMP | SETDMA | ; SET DMA ADDRESS |
| XREAD | JMP | Read | ; READ SElected sector |
| XWRITE | JMP | WRITE | ; Write selcected sector |
|  |  |  |  |
| ;GIVE SIGN ON MESSAGE |  |  |  |
| Bоot: |  |  |  |
|  | LXI | SP, 80H | ;SET STACK POinter |
|  | XRA |  | ; Clear a |
|  | STA | 4 | ;FORCE DISK A |
| ; |  |  |  |
|  | $\begin{aligned} & \text { IF } \\ & \text { out } \end{aligned}$ | MYSYS $0 \times 8 \mathrm{~S}$ OR PTVDM |  |
|  | MVT | 0 C 8 H OCOH | ;RESET VDM |
|  | call | XCONOT | ; ${ }^{\text {SEEND }}$ IT OUT |
|  | Endif |  |  |
| ; |  |  |  |
|  | LXI | ${ }^{\text {H, SMSG }}$ | ; PRINT OPENING MESSAGE |
|  | ${ }_{\text {JMP }}$ | ${ }_{\text {GMSG }}$ | ;NOW DR Housekeeping |

; Warm boot routine--activated when you control c
wBoot:

| LXI | SP, 80 ${ }^{\text {r }}$ | ;SET Stack pointer |
| :---: | :---: | :---: |
| mVI | c, 0 | ;USED BY DSKSEL AND SETTRK |
| CALL | xDSKSEL | ; SELECT DISK 0 |
| MVI | D, NSECTS | ; \# OF SECTORS IN D |
| MVI | B,2 | ; START WITH SECTOR |
| LXI | H,CCP | ; GET STARTING ADDR |
| Call | SETTRK | ; SElect track ( C ) |
| PUSH | B | ; SAVE BC |
| MOV | C, B | ; PUT SECTOR IN C |
| CALL | xsetsec | ;SET IT UP |
| mov | B, H | ;GET READY FOR SET DMA |
| mov | c, L | ; |


| call | xsEtdma | ; SET IT UP |
| :---: | :---: | :---: |
| POP | B | ; RESTORE BC |
| call | xread | ; READ Record |
| JNz | wbootx | ; ERROR ON READ |
| DCR | D | ; OF SECTORS TO READ - |
| JZ | GOCPM | ; DONE--GO TO CPM |
| INR | ${ }^{\text {B }}$ | ; POINT TO NEXT SECTOR |
| mov | A, B | ; MOVE IT TO A |
| CPI | 27 | ; END OF TRACK |
| JC | wboot 2 | ; NO---CONTINUE READING |
| INR | c | ;TRACK NOW EQUAL 1 |
| MOV | B, C | ; SET SECTOR BACK TO |
| JMP | WBOOT 2 | ; READ NEXT TRACK |
| LXI | H,BTMSG | ;GET ADDR OF ERROR MSG |
| CALL | PMSG | ;PRINT IT |
| CALL | XCONIN | ; WAIT FOR KB ENTRY |
| JMP | wBootl | ; DO WARM Boot again |

; $;$ this routine is the exit to CP/M System

GOCPM:

| MVI | A, OC3H | ; PUT JUMP TO WARM BOot |
| :---: | :---: | :---: |
| STA |  | ; AT ADDR 2 ERO |
| LXI | H, xwB00т | ;GET ADDR OF WARM Boot entry |
| SHLD | 1 | ;FINISH THE JUMP INSTUCTION |
| STA | 5 | ; PUT JUMP TO BDOS AT 5 |
| LXI | H,BDOS | ;GET ADDR OF bDOS ENTRY |
| SHLD |  | ;FINISH JUMP INSTRUCTION |
| STA | 38日 | ; SET UP INTERRUPT TRAP |
| LXI | H, TRAP |  |
| SHLD | 39 H | ; ARM IT |
| LXI | H,80H | ;SET default dma addr |
| SHLD | DMAADD | ; AND STORE IT |
| LDA | 4 | ; GET PREV DRIVE |
| mov | C, A | ; SHOULD BE IN C |
| ${ }^{\text {IF }}$ | TEST |  |
| XRA | A | ; Clear a |
| OuT | OFFH | ;SHOW IN THE FP LIGHTS |
| ENDIF |  |  |
| JMP | CCP | ;GO TO CP/M |

this routine is a null interrupt trap
;FOR NOW, JUST RE-ENABLE INTERRUPTS
trap:

| DB | 0 | ;LEAVE ROOM FOR JUMP OR CALL |
| :--- | :--- | :--- |
| DB | 0 | $;$ |
| DB | 0 | $;$ |
| EI |  | ;'GUR INTERRUPTS BACK ON |
| RET |  | BACK FROM WHENCE YOU CAME |

; Select disk given by reg c
'́SKSEL:
$\begin{array}{ll}\text { MOV } & \text { A, C } \\ \text { STA } \\ \text { RET }\end{array} \quad$ NXTDSK
; PUT NEW DISK IN A
;SAVE IT FOR REAL RTN
; SELECT DISK STORED in NXTDSK
SELDSK:

| PUSH | H |  |
| :---: | :---: | :---: |
| PUSH | D | ;SAVE REGS |
| PUSH | B |  |
| LDA | NXTDSK | ; GET NEW DISk \# |
| mov | C, A |  |
| ANI | H | ;LOOK AT 3 LSB'S |
| LxI | H, diskno | ; GET ADDR OF diskno |
| CMP | $\cdots$ | ; $\mathrm{NEW}=$ OLD |
| Jz | SElxit | ; IF SO,RETURN |
| IF | onedsk |  |
| PUSH | B | ;SAVE REGS |
| PUSH | A | ; SAVE REGS |
| LXI | H,MNTMSG | ; PRINT 'MOUNT' MSG |
| CALL | PMSG | ; TELL THEM |
| POP | A | ; Retrieve reg a |

SWGLSKSOYOIW OOT-S

| STA | DISKNO | ; Store it for later |
| :---: | :---: | :---: |
| ADI | 'A' | ; MAKE IT ASCII LEtTER |
| mov | C, A | ; PUT IN C FOR CONOT |
| call | xconot | ;PRINT IT |
| CALL | XCONIN | ; WAIt For go ahead |
| POP | B | ;RESTORE BC |
| XRA | A | ; Clear a |
| JMP | SELXIT | ; Exit Seldsk |
| endif |  |  |
| MOV | E, M | ; PUT OLD DISK \# IN DE |
| MVI | D,0 | ;Clear d for dad |
| LXI | H, TRKTB | ; GET ADDR of track table |
| DAD | D | ;ADD DISK \# TO ADDR |
| IN | DTRK | ;GET TRACK FROM OLD DRIVE |
| mov | M, A | ; STore it in track table |
| mov | E, C | ;GET NEW DRIvE \# |
| LXI | H, TRKTB | ; GET ADDR OF TRK table |
| DAD | D | ;GET HEAD LOC ON NEW DRIVE |
| MOV | A, M | ; And put in rec a |
| STA | HOLD | ;SAVE FOR SEEK Routine |
| out | DTRK | ;ADJUST 1771 |
| mov | A, C | ;GET DISK \# |
| STA | diskno | ; STORE IT FOR LATER USE |
| CMA |  | ; INVERT BIT FOR LATCH |
| ADD | A | ; PUT BITS 0-1 AT 4-5 |
| ADD | A | ; |
| ADD | A | ; |
| ADD | ${ }^{\text {a }}$ |  |
| ORI | 2 | ; MAKE LATCH COMMAND |
| OuT | DCONT | ; SET latch with code |
| IN | dStat | ; Unload head because 1771 does |
| ANI | 20H | ; NOT RECOGNIZE DRIVE SWITCH |
| JNz | Seldel |  |
| XRA | A | ;SAY EVERYTHING OK |
| POP | B | ;RESTORE REGS |
| POP | D | ; RESTORE REGS |
| POP | H | ;RESTORE REGS |
| RET |  | ; |

; SET the track given in reg c
SETTRK:
$\begin{array}{ll}\text { MOV } & \text { A, } \mathrm{C} \\ \text { STA } & \text { TRK } \\ \text { RET }\end{array}$
;TRK WAS IN REG C
; PUT IT WHERE IT CAN BE FOUND
;
;'SET DISK SECTOR NUMBER
SETSEC:
$\begin{array}{ll}\text { MOV } & \text { A, C } \\ \text { STA } & \text { SECT } \\ \text { RET }\end{array}$
;GET SECTOR NUMBER
; PUT AT SECTOR
; SET DISK DMA ADDRESS
SETDMA:

| mov | H,B | ; MOV bC To hl |
| :---: | :---: | :---: |
| MOV | L, C |  |
| SHLD | DMAADD | ;SAVE It |
| RET |  | ; |

; READ A SECTOR AT SECT, SEEK THE
;USE STARTING ADDRESS AT DMAADD
READ:

|  | Call | SEEK | ;MOVE HEAD TO TRACK |
| :---: | :---: | :---: | :---: |
|  | MVI | A, RTCNT | ; GET RETRY CNT |
| RRETRY: |  |  |  |
|  | STA | ERCNT | ; SAVE ERROR CNT |
|  | LDA | SECT | ; GEt sector number |
|  | LHLD | DMAADD | ;GET STARTING ADDRESS |
| READ 1: |  | DSECT | ; SET SECTOR IN 1771 |
|  | CALL | RDYCK | ; SEE IF DRIVE READY |
|  | DI |  | ;DO NOT ALLOW interrupts |
|  | IN | dstat | ;Get status |
|  | ANI | 20 H | ;Check if head loaded |


|  | MVI | A, 88H | ; SEtup read w/o head load |
| :---: | :---: | :---: | :---: |
|  | JNZ | READE | ; if loaded - then do it |
|  | ORI | 4 | ; ElSE Force head load |
| READE: | OUT | DCOM | ;SEND COMMAND TO 1771 |
| RLOOP: |  |  |  |
|  | call | fread | ; READ A SECtor |
| RDDONE: |  |  |  |
|  | EI |  | ;ALLOW INTERRUPTS AGAI |
|  | IN | DSTAT | ; READ DISK STATUS |
|  | ${ }^{\text {ANI }}$ | 9 DH | ; LOOR AT ERROR BITS |
|  | RZ |  | ; RETURN IF NON |
| CHECK: |  | ERRS |  |
| ; |  |  |  |
|  | IF | TEST |  |
|  | CMA |  | ; FP INVERTED LOGIC |
|  | OuT | OFFH | ; in the limelight |
|  | Endif |  | ; |
| ; | call | ERCHK | ;Check for seek error |
|  | LDA | ERCNT | ; GET ERROR CNT |
|  | DCR | A | ;DECREMENT COUNT |
|  | JN2 | RRETRY | ; TRY TO READ AGAIN |
|  | MVI | A, 'R' | ; SHOW READ ERROR |
|  | JMP | ERRMSG | ; tell someone |
| FREAD: |  |  |  |
|  | IN | DWAIT | ; WAIT FOR DRQ OR INTRQ |
|  | ORA | A | ; SET FLAGS |
|  | RP |  | ;DONE IF INTRQ |
|  | IN | DDATA | ;READ A BYTE FROM DISK |
|  | mov | M, A | ; PUT BYTE IN MEMORY |
|  | INX | H | ; INCR MEMORY POINTER |
|  | JMP | fread | ;KEEP READING |

; WRITE THE SECTOR AT SECT--LOAD hEAD FIRST ;USE STARTING ADDRESS AT DMAADD

|  | $\begin{aligned} & \text { CALL } \\ & \text { MVI } \end{aligned}$ | $\begin{aligned} & \text { SEEK } \\ & \mathrm{A}, \mathrm{RTCNT} \end{aligned}$ | ;GET ON RIGHT TRACK <br> ;GET RETRY COUNT |
| :---: | :---: | :---: | :---: |
| WRETRY: |  |  |  |
|  | STA | ERCNT | ; SAVE ERROR CNT |
|  | LDA | SECT | ; GET SECTOR \# |
|  | LHLD | DMAADD | ; GET Starting addr |
| WRITE1: | out | DSECT | ; SET SECT into 1771 |
|  | call | RDYCK | ; SEE IF DRIVE READY |
|  | DI |  | ; DO NOT ALLOW INTERRUPTS |
|  | IN | dStat | ;GET disk status |
|  | ANI | 20 H | ; Check for head loaded |
|  | MVI | A, OAB ${ }^{\text {H }}$ | ; SETUP WRITE W/O HEAD LO |
|  | JNZ | WRITEN | ; IF LOADED THEN DO IT |
|  | ORI | 4 | ;FORCE WRITE WITH HEAD L |
| WRITEN: | OUT | DCOM | ;issue command |
|  |  |  |  |
|  | CAll | FWRITE | ; WRite a sector |
| WDONE: | EI |  | ;ALLOW interrupts again |
|  | IN | DSTAT | ;READ DISK STATUS |
|  | ${ }_{\text {ANI }}$ | OFDH | ; MASK NON-ERROR BITS |
|  | R2 |  | ;RETURN IF NO ERRORS |
|  | STA | ERRS | ; SAve error flags |
|  | IF | TEST |  |
|  | CMA |  | ; INVERT THEM |
|  | $\begin{aligned} & \text { OUT } \\ & \text { ENDIF } \end{aligned}$ | OFFH | ; PUT THEM ON FP LEDS |
|  | call | ERCHK | ; CHK/CORRECT SEEK ERR |
|  | LDA | ERCNT | ; GET ERROR CNT |
|  | DCR | wr | ; DECREMENT COUNT |
|  | JNZ | WRETRY | ; TRY WRITE AGAIN |
|  | MVI | ${ }^{\text {A , 'W' }}$ | ; SHOW WRITE ERROR |
|  | JMP | ERRMSG | ;DO ERROR MESSAGES |
| WRITE: | IN | DWAIT | ; WAIT FOR READY |
|  | ORA | A | ; SET FLAGS |
|  | RP |  | ;GET OUT WHEN DONE |


; 'HECK to SEE IF ON CORRECT track-
; CHECK TO SEE IF ON CO
; CORRECT IF NECESSARY
' CHKSK:

| MVI | A, 0 C 4 H | ; SET UP READ addr |
| :---: | :---: | :---: |
| OUT | DCOM | ;COMMAND TO 1771 |
| DI |  | ;DO NOT ALLOW INTERRUPTS |
| IN | DWAIT | ; WAIT FOR 1ST DRQ (TRK) |
| IN | dDATA | ; READ THE TRACK ADDR |
| EI |  | ;ALlow interrupts again |
| PUSH | PSW | ; SAVE TRACK |
| CALL | SWAIT | ; WAIT FOR OPERATION TO FINISH |
| POP | PSW | ;GET TRACK BACK |
| STA | HOLD | ; USE IT TO SET UP SEEK RTN |
| OUT | DTRK | ;UPDATE TRACK REGISTER |

'track desired has already been stored by settrk
'

| SEEK: |  |  |  |
| :---: | :---: | :---: | :---: |
|  | ${ }_{\text {LDA }}^{\text {CALL }}$ | SELDSK <br> TRK | ;WILL DO NOTHING IF FROM CHKSK ;GET WHERE WE ARE GOING TO |
|  | PUSH | B | ; SAVE BC |
|  | mov | B, A | ;SAVE TRK IN b |
|  | LDA | HoLD | ;GET PREV TRACK |
|  | CMP | B | ; Compare to current |
|  | JZ | Forint | ;SAME SO FORCE InTERRUPT |
|  | mov | A, B | ; PUT CURRENT In A |
|  | STA | HoLD | ; SAVE FOR NEXT TIME |
|  | POP | B | ; REStore bC |
|  | CPI | 0 | ;SEE IF TRK $=0$ |
|  | JZ | homrtn | ; Do home routine instead |
|  | out | dDATA | ;GIVE DESIRED TRK TO 1771 |
|  | call | SWAIT | ; WAit till not busy |
|  | MVI | A, 1AH | ;SEEK-10MS- |
|  | call | SCMND | ;issue command |
|  | call | SLOOP | ; Give time for head to settle |
|  | RET |  | ; |
| S' ${ }_{\text {CMND }}$ |  |  |  |
|  | out | DCOM | ;issue command |
| SWAIT: |  |  |  |
|  | CALL | ${ }_{2 \text { IP }}^{2 \text { IP }}$ | ; WE NEED AT LEAST 12 US |
| SBUSY: |  |  |  |
|  | ${ }^{\text {IN }}$ | dStat | ; WAIT FOR NOT BUSY, IE INTRQ |
|  | ${ }^{\text {RRC }}$ |  |  |
| 2IP: | J | Sbus |  |
|  | RET |  | ; intre reset by reading status |
| ; Forint |  |  |  |
|  | MVI | $\mathrm{A}, \mathrm{OD} 0 \mathrm{H}$ | ;Clear any pending command |
|  | OUT | DCOM | ; AND FORCE TYPE 1 STATUS |
|  | POP | B | ; Clean up |
|  | RET |  | ;GO BACK |
| ŚLOOP: |  |  |  |
|  | PUSH | H | ; GIVE TIME FOR HEAD TO SEttle |
|  | LXI | H,9*256 | ; NEED AbOUT 10 MS |
| SLOOP1: | DCR |  |  |
|  | JN2 | SLOOP1 | ; |
|  | DCR | \% | ; |
|  | JN2 | SLOOP1 | ; |
|  | POP | H | ; |
|  | RET |  | ; |

orint :

| SEEK: |  |  |  |
| :---: | :---: | :---: | :---: |
|  | ${ }_{\text {LDA }}^{\text {CALL }}$ | SELDSK <br> TRK | ;WILL DO NOTHING IF FROM CHKSK ;GET WHERE WE ARE GOING TO |
|  | PUSH | B | ; SAVE BC |
|  | mov | B, A | ;SAVE TRK IN b |
|  | LDA | HoLD | ;GET PREV TRACK |
|  | CMP | B | ; Compare to current |
|  | JZ | Forint | ;SAME SO FORCE InTERRUPT |
|  | mov | A, B | ; PUT CURRENT In A |
|  | STA | HoLD | ; SAVE FOR NEXT TIME |
|  | POP | B | ; REStore bC |
|  | CPI | 0 | ;SEE IF TRK $=0$ |
|  | JZ | homrtn | ; Do home routine instead |
|  | out | dDATA | ;GIVE DESIRED TRK TO 1771 |
|  | call | SWAIT | ; WAit till not busy |
|  | MVI | A, 1AH | ;SEEK-10MS- |
|  | call | SCMND | ;issue command |
|  | call | SLOOP | ; Give time for head to settle |
|  | RET |  | ; |
| S' ${ }_{\text {CMND }}$ |  |  |  |
|  | out | DCOM | ;issue command |
| SWAIT: |  |  |  |
|  | CALL | ${ }_{2 \text { IP }}^{2 \text { IP }}$ | ; WE NEED AT LEAST 12 US |
| SBUSY: |  |  |  |
|  | ${ }^{\text {IN }}$ | dStat | ; WAIT FOR NOT BUSY, IE INTRQ |
|  | ${ }^{\text {RRC }}$ |  |  |
| 2IP: | J | Sbus |  |
|  | RET |  | ; intre reset by reading status |
| ; Forint |  |  |  |
|  | MVI | $\mathrm{A}, \mathrm{OD} 0 \mathrm{H}$ | ;Clear any pending command |
|  | OUT | DCOM | ; AND FORCE TYPE 1 STATUS |
|  | POP | B | ; Clean up |
|  | RET |  | ;GO BACK |
| ŚLOOP: |  |  |  |
|  | PUSH | H | ; GIVE TIME FOR HEAD TO SEttle |
|  | LXI | H,9*256 | ; NEED AbOUT 10 MS |
| SLOOP1: | DCR |  |  |
|  | JN2 | SLOOP1 | ; |
|  | DCR | \% | ; |
|  | JN2 | SLOOP1 | ; |
|  | POP | H | ; |
|  | RET |  | ; |

;
; READ OR WRITE ERROR DETECTEDANDLE NO REC FOUN
CONDITION ELSE NOPMAL RETRY LOOP
ERCHK:
MASK FOR NRF
; NOT NRF FAULT

OP1:

HOME ROUTINE-SET TRK TO ZERO AND SEEK
WILL PICK UP DURING READ OR WRITE
; 'оме:
$\begin{array}{ll}\text { XRA } & \text { A } \\ \text { STA } & \text { TRK } \\ \text { RET } & \end{array}$
; CLEAR A
; PUT Where it can be found
; this routine only called from seek
TO PERFORM ACTUAL HOME
omrtn

| MVI | A,ODOH | ;RESET ANY PENDING COMMAND |
| :--- | :--- | :--- |
| OUT | DCOM | ;ISSE COMMAD |
| CALL | SWAIT | ;WATT FOR NOT BUSY |
| MVI | A, OAH | ;10 MS SEER RATE-HOME |
| CALL | SCMND | ;ISSUE CEMMAND-WAIT FOR INTRQ |
| CALL | SLOOP | ;FOR HEAD SETTLING |
| RET |  | ;GO BACK |

; Check for drive ready-tf not tell operator ;AND WAIT FOR C
' RDYCK:

| IN | dStat | ;GEt disk status |
| :---: | :---: | :---: |
| ANI | 80 H | ; MASK For ready |
| R2 |  | ;ok we are ready |
| PUSH | H | ;SAVE ReGs |
| PUSH | B | ;SAVE ReGs |
| LxI | H, NRDYMS | ; POINT TO MESSAGE |
| Call | PMSG | ; PRINT IT OUT |
| LDA | dISKNO | ;GET CURRENT DRIVE |
| ADI | 'A' | ; Change to ascil |
| mov | C, A | ; SET UP TO PRINT |
| Call | XCONOT | ; PRINT IT |
| CALL | XCONIN | ;GET KEYBd Char |
| LXI | H, CRLF | ; SET UP CR AND LF |
| CALL | PMSG | ; PRINT IT |
| POP | B | ;RESTORE REGS |
| POP | н | ;RESTORE REGS |
| RET |  | ; |

;
; PRINT
PMSG:

; ERROR MESSAGE ROUTINES
ERRMSG:

| PUSH | H | ;SAVE HL |
| :---: | :---: | :---: |
| PUSH | B | ; SAve bC |
| MOV | B, A | ;SAVE TYPE indicator |
| Lxi | H,CRLF | ; DO CRLF |
| CALL | PMSG | ;PRINT IT |
| mov | C, B | ;GET TYPE INDIC |
| CALL | XCONOT | ; PRINT IT |
| LDA | ERRS | ;GET ERROR BYTE |
| LXI | H, TYPERR-3 | ; POINT TO ERROR table |
| INX | H | ; POint to next entry |
| INX | н | ; in error type table |
| INX | н |  |
| RLC |  | ;SHIFT bit into CARRY |
| JNC | LOCERR | ;NOT IN ERROR-KEEP LOOKING |
| MVI | B, 3 | ; SET B = 3 |
| mov | C, M | ; GET FIRST LETTER |
| CALL | XCONOT | ; PRINT IT |
| INX | н | ; POIN T TO NEXT LETtER |


| If | MYS YS |  |
| :---: | :---: | :---: |
| IN | 2 | ; REAd console stat |
| ani | 80H | ; LOok at bit 7 |
| JN2 | conin | ; KEEP WAIting |
| IN | 0 | ;GET DATA BYTE |
| ani | 7FH | ;TURN PARITY OfF |
| RET |  | ; |


|  | CPI | ODH | ; IS IT CR |
| :---: | :---: | :---: | :---: |
|  | $\mathrm{JZ}^{\text {I }}$ | CRTN | ; YES |
|  | CPI | 0AH | ; IS IT LF |
|  | J2 | CLF | ; YES |
|  | CPI | ОВ | ;is it home |
|  | J2 | ном | ; YES |
|  | CPI | OCH | ; IS IT FF |
|  | Jz | CLEAR | ; YES |
| CRTN: | JMP | EXIT | ; NOT WAnted-do nothing |
|  | mov | A, L | ;GET CURSOR LOC LSB |
|  | ANI | A, OCOH |  |
|  | mov | L, A | ; PUT BACK |
|  | JMP | EXIT | ;GET OUT |
| Delete: |  |  |  |
|  | MOV | A, L | ;GET CURSOR LOC LSb |
|  | ani | 63 | ; SEE If begin of line |
|  | $\mathrm{JZ}^{\text {J }}$ | EXIT | ; YES - DO NOTHING |
|  | DCX | H |  |
|  | MVI | M, ' ' | ;SET NEW LOC TO SPACE |
|  | JMP | EXIT | ; GET OUT |
| CLF: |  |  |  |
|  | PUSH PUSH | ${ }_{8}$ | ; SAVE Regs for scrol |
|  | LXI | B, 64 | ;put line length in bC |
|  | DAD | B | ; AdD to Current loc |
|  | mov | A, H | ;GET LOC MSB |
|  | CPI | vDBAS2+4 | ; is it past end of screen |
|  | POP | B | ; RESTORE BC |
|  | JNZ | XTRA | ; HL OK but gigo on stack |
|  | CALL | SCRL | ; SCROLL IT |
|  | POP | H | ; RESTORE HL |
|  | JMP | EXIT | ; GET OUT |
| XTRA : | XTHL |  | ; SWAP HL WITH tos |
|  | POP | H | ; CLEAN STACK AND CORRECT HL |
|  | JMP | EXIT | ;GET OUT |
| clear: |  |  |  |
|  | LXI | H, VDBASE | ; GEt Start of Screen in hl |
| CLEAR1: | MVI | A, VDBAS $2+4$ | ; PUT END OF SCREEN MSB IN A |
|  | MVI | M, ' | ; Put blank in mem |
|  | INX | H | ; POINT TO NEXT LOC |
|  | CMP | H | ; Check for end of screen |
|  | JNZ | Clear 1 | ; NO, STAY in loop |
| ном: |  |  |  |
|  | LXI | h, vDBASE | ; PUt Start of Screen in hl |
| EXIT: | JMP | EXIT | ; |
|  | mov | A, M | ;GET MEMORY BYTE |
|  | ORI | 80 H | ; TURN ON CURSOR |
|  | mov | M, A | ; PUT IT BACK |
|  | MOV | A, C | ; MAKE IT LOOK NORMAL |
|  | RET |  | ; GOODBYE |
|  |  |  |  |
| vDMP | DW | 0 | ; CURRENT CURSOR LOC |
|  | Endif |  | ; |
| ; |  |  |  |
| $\begin{aligned} & \text {;WRITE A } \\ & \text { LIST: } \end{aligned}$ | chara | ER ON Listi | vice |
|  |  |  |  |
|  | ${ }_{\text {IF }}^{\text {MVI }}$ | MYSYS |  |
|  | MVI | ${ }_{C}^{\text {A, OAH }}$ | ; CHECK FOR A CR <br> ; IF IT IS THEN DO NULL RTN |
|  | MVI | A, 25 | ; |
|  | Jz | ADDNL |  |
|  | MVI | A, OCH | ; CHECK FOR FF |
|  | CMP |  | ; DO FF RTN If Yes |
|  | MVI | A, 25 | ; |
|  | Jz | ADDNL | ; |
|  | mVI | A, ODH | ; |
|  | CMP | C | ; |
|  | MVI | A, 02 | ; |
|  | Jz | ADDNL | ; |
| LIST1: |  |  |  |
|  | IN | 2 | ; CHECK STATUS |
|  | ANI | 40H | ; MASK OFF TBE |


|  | J2 | LIST | ; WAIT FOR TBE |
| :---: | :---: | :---: | :---: |
|  | MOV | A, C | ;GET DATA ByTE |
|  | out | 03 | ; SEND IT OUT |
|  | RET |  | ; |
| ADDNL: |  |  |  |
|  | PUSH | B | ;SAVE BC |
|  | MOV | B,A | ; SAVE IT IN B |
| ADDNL 1: |  |  |  |
| DDDN 3 , | call | LIST1 | ; PRINT CR FIRST |
| ADDNL3: | : mVI | C,07FH | ;GET NULL Char |
|  | DCR | B | ;DECREMENT COUNTER |
|  | JNZ | ADDNL1 | ;IF <>0 THEN DO MORE NULLS |
| ADDNL 4: |  |  |  |
|  | ${ }_{\text {POP }}$ | , | ; RESTORE B |
|  | mov | A, C | ;RESTORE A |
|  | RET |  | ;RETURN FROM LIST |
|  | ENDIF |  | ; |
| ; |  |  |  |
| ; | IF | NOT MYSYS |  |
|  | CALL | CONOT | ; ROUTE TO CONSOLE FOR NOW |
|  | RET |  |  |
|  | Endif |  |  |
| ; NORMAL | LLY USED | TO PUNCH Pa | APE |
| ; CAN BE | E USED AS | Bit bucket | heck files |
|  |  |  |  |
|  |  |  |  |
|  | RET |  | ; |
| ; |  |  |  |
| ; NORMAL | lly used | to read pap |  |
| ; NOT IM | MPLEMENTED | but CPM | ES EOF |
| ; Reader: |  |  |  |
|  | MVI | A, 1AH | ; SET A=CTL-Z (EOF) |
|  | RET |  | ; |
|  |  |  |  |
| ; DATA A | areas for | CPM |  |
| ' ${ }^{\text {TRK }}$ |  |  |  |
|  | DB |  | ; TRK WANTED |
| TRKTB | ${ }^{\text {DB }}$ | 1,1,1,1 | ; TRACK TABLE |
| SECT | DB | 0 | ;SECTOR \# |
| DMAADD | DW | 0 | ; DMA ADDRESS |
| diskno | DB | 0 | ;SELECTED DISK |
| NXTDSK | DB | 0 |  |
| HoLD | DB | 0 | ;SAVE AREA FOR SEEK |
| ' ERCNT: | DB |  | ; ERROR COUNT |
| ERRS: | DB | 0 | ; ERROR HOLD AR |
| SKCNT: | ${ }_{\text {DB }}^{\text {DB }}$ | 0 | ; SEEK ERROR HOLD |
| ; |  |  |  |
| ; |  |  |  |

## CURSOR

Continued from Page 38
3250 DIM AS (LN)
3260 FOR $N=1$ TO
3260 FOR $\mathrm{N}=1$ TO LN
$3280 \mathrm{AS}(\mathrm{N})=\mathrm{DS}+\mathrm{STRINGS}(\mathrm{WD}-\mathrm{LEN}(\mathrm{D} \$), 32)$
3290 NEXT N
3290 NEXT N
3300 CLOSE
3310 GOSUB 2550
3320 GOTO 630
3320 GOTO 630
3330
3330 REM
3340 REM
3340 REM
3350 REM
3360 REM *
3370 REM *
3380 REM * IN. THIS LISTING, " 1 " MEANS "IS GREATER THAN
3380 REM *
3390 REM *
3400 REM ************
3410 PRINT "HI THERE"

## LETTERS

## Continued from Page 8

use and as a co-conspirator to foment S-100 computing until a better system evolves.

Again, congratulation and good luck.

Edward Lee
Riderwood Md

Dear Sol:
Hope you don't mind the "plug" in our newsletter. I was in Atlantic City last month and picked up one of your flyers at the Computer Store in Linwood. SOCCC is a small club (25-30 members) of mostly $S-100$ systems, so most of the members will probably subscribe. Next time you're in the southern California area we'd be interested in hearing from you - give me a call or drop a note.

I just got my first issue yesterday - read it from cover-to-cover, and in spite of the error (omission) on page $44 / 45$ (CBBS) I enjoyed the magazine very much. I like the format but most of all, the content - good, worthwhile and useful articles.

If there's anything $I$ (or the club) can do to help, don't hesitate to call or write.

Mel Hengen
South Orange County Computer $\mathrm{Cl} u \mathrm{~b}$
Fountain Valley, CA

To: Sol Libes
Your first issue of s-100 MICROSYSTEMS was very informative and interesting to read. I hope you will be able to keep the good work up. I enjoyed reading the IEEE $S-100$ information. However, I must confess that I will be studying it for some time to fully understand the bus. Even though I have a SYM-I (6502 based single board computer), I feel that sometime (a Year or two) in the future I would like to build an $S-100$ system. It would not have to be a 6502 based system (I have corresponded with you on that subject). To spend all that money for a mainframe, etc. I might very well be interested in a 16-bit machine. I am currently investigating the Motorola MC68000. Thus, I am very happy to see the IEEE S-100 enhancements to accommodate l6-bit machines. I am very interested in having a computer system that is both "powerful and has great flexibility". With those thoughts in mind, I am looking forward to more issues of $S-100$ MICROSYSTEMS and $I$ would be interested in articles that cover, or are about the
following:

* S-100 MAINFRAMES - good, bad, noise problems and their solutions, power supplies, etc.
* CPU card design principles to meet IEEE standards.
* Operating systems (Good start with the $C P / M$ article).
* IN GENERAL - a good mixture of hardware and software articles (as per Vol 1/ No 1).

George V. Wilder
Lisle, Ill

## CLUBS

## Continued from Page 12

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MAR/APRIL 1980

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## S-100 NEWS

## Continued from Page 7

AMRAD TO PUBLISH CBBS DIRECTORY
The Amateur Radio Research and Development Corporation (AMRAD) -- an amateur radio and computer society with headquarters in the Washington, DC area -is conducting a survey of computer message systems. AMRAD plans to publish a directory of Computerized Bulletin Board Systems, Apple Bulletin Board Systems, Forum-80's and similar systems in the near future. The directory is to be available to anyone at a nominal charge of \$1. AMRAD member and those who contribute first-hand information about existing message systems will receive the directory free of charge.

Individuals connected with existing message systems are asked to send their names and addresses to David W. Borden, Rt 2, Box 233B, Sterling, Virginia 22170, and request a copy of the AMRAD Computer Message System Questionnaire. This questionnaire contains all the elements of information needed for entries in the directory.

## IEEE DEVELOPING ASSEMBLY LANGUAGE

 STANDARD FOR MICROSThe Institute of Electrical and Electronic Engineers is developing a standard for assembly language on microprocessors (IEEE Task P694/D11). It is long overdue and will be of enormous value to all assembly language programmers who are struggling with different microprocessors. The group working on the standard has done some genuinely worthwhile things, such as showing that all the current major chips can be handled nicely by one standard. The problems at present are incredible. For example, on some chips mov A,B means move the contents of register $B$ to $A$, while on others it means just the opposite.

Right now, AMD is second sourcing the zilog z8000, and would you believe it, they are not using the zilog mnemonics! Hopefully, the new IEEE standard will cure problems such as this. In another example, Zilog did not use the Intel mnemonics for the Z 80 's instructions which were a subset of the 8080 .

The standard also covers Instruction Names, Address Modes, Operand Sequences, Expression Evaluation, Constants, Lables, Comments and Assembler Directives. The standard does not specify the syntax necessary to support macros or conditional assembly.

The IEEE Computer Society is to be congratulated for its activities in
developing computer standards. They are overcoming problems created by companies that all too often purposely create incompatibilities in order to protect their competitive position.

I predict that this standard will meet with the wide adoption that the other IEEE standards (IEEE-488 and IEEE-S-100 bus standards) are meeting with. Incidentally, you can obtain a copy of the Assembly Language Standard draft by sending a self-addressed $10 \times 13$ size envelope with 4 stamps on it to Dr. Robert G. Stewart, Chairman Computer Standards Committee, IEEE Computer Society, 1658 Belvoir Drive, Los Altos CA 94022.

Incidentally, the IEEE is also working on several other standards relevant to the microcomputer area. They are: Multibus (Task No. P696.2), Microbus (Task No. P696.3), Futurebus (Task No. P696.4), Floating Point (Task No. P754), High Level Languages (Task No. 755), Pascal (Task No. 770) and Relocatable Object Format (Task No. 695). I will try and report on the progress of these standards in a future S-100 MICROSYSTEMS column.

## SOL-LIKE COMPUTER TO BE PRODUCED

Lee Felsenstein, the wizard who created the Processor Tech SOL computer, the VDM, the $3 P+S$ and several other products is back with a new computer that is in effect a "SOL-II". (incidentally, before we go any further I must say that contrary to a statement made in a magazine recently, the $S O L$ was not named after me.)

The new computer, whose new name and manufacturer is still a secret will debut in April. It is being manufactured here in the U.S. by a Swedish concern. Physically it will look like a sol except that it will be housed in a vacuum formed plastic enclosure. It will operate with a TV monitor and have a built-in cassette, printer and spare parallel $I / O$ ports. The entire computer will be on a single board which includes the keyboard and will have as an option, a 4-slot $\mathrm{S}-100$ bus motherboard (IEEE compatible) for plug-ins such as disk controller and memory and $I / O$ expansion.

The computer will have Microsoft BASIC (V5.l) in ROM. They can be removed and the space converted to RAM by changing some jumpers. The computer will have a lot of powerful goodies such as: Real-time clock, interrupt driven keyboard, priority interrupt system, 16 K static RAM on the main board, use of cassette port for RS-232 modem I/O, options such as the ACE multitasking scheduler and best of all, the unit will sell for under $\$ 1 \mathrm{~K}$.

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Every issue of S-100 MICROSYSTEMS brings you the latest in the $\mathrm{S}-100$ world. Articles on applications, tutorials, software development, letters to the editor, newsletter columns, book reviews, new products, etc. Material to keep you on top of the ever changing microcomputer scene.
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[^1]

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## NEW PRODUCTS



## VIDEO DIGITIZER/MONITOR INTERFACE

TECMAR INC. has introduced a Real Time Video Digitizer and Monitor Interface (RT+MI) for $S-100$ systems, which digitizes video data from TV cameras and uses the digital data to reconstruct a picture on a TV moniotr. It digitizes the picture in $1 / 60$ second and deposits it into memory as a single operation using DMA. It displays pictures in 16 gray levels or black and white combination. It can simultaneously deposit and display allowing constant viewing of picture until desired image is seen on monitor, when it can be frozen on monitor. The last digital image displayed remains in memory and can be displayed at any time thereafter without erasing it from memory. The image can be processed or put on disk for later retrieval. Maximum resolution is $512 \times 240$ pixels. The complete RT+MI, includes Video $A / D$ (shown above), Video $D / A$ and DMA controller boards occupying $3 \mathrm{~S}-100$ slots and costing $\$ 850$. For information contact: TEC Mar Inc., 23414 Greenlawn, Cleveland Ohio 44122; tel: (216)382-7599.

## 64K RAM BOARD INTRODUCED

Measurement Systems \& Controls Inc., has released the DMB-6400, a 64 K Bank Switchable Dynamic Memory Board for S-100 systems. Output port addressing is used for bank selection of 4 totally independent 16 K banks of memory. Each bank can be turned ON or OFF at system
reset and PHANTOM can be used by any of the four banks. Four diagnostic LEDs indicate which banks of memory are on. The board will operate with all 8080, 8085 @ 3 MHz and most Z 80 A @ 4 MHz CPU boards. In addition, it will operate with the Marin Chip M9900 CPU. For further information contact: Measurement Systems \& Controls Inc., 867 North Main Street, Orange CA 92667 ; tel: (714)633-4460.


## NORTH STAR INTRODUCES HARD DISK

North Star Computers Inc. has announced a new Winchester-type l8mbyte enhancement for its Horizon computers and users of North Star floppy disk systems. Up to four hard disks and 2 mini-floppy disks may be accommodated on one system, providing up to 72 Mbytes of storage on the hard disks and over lmbyte on the floppy drives. Century Data Marksman hard disk drives are utilized with an average access time of 78 msec . Software supplied with the system includes a File Manager, a program for creating backup diskettes, a Command Processor (supporting all the North Star floppy disk DOS and Monitor commands, while adding others) and BASIC interpreter (modified to support hard disk files and run all previous North Star Basic programs with little or no change). For more information contact: North Star Computers, 1440 Fourth Street, Berkeley CA 94710; tel: (415) 527-6950.


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Our innovative Z-80A CPU board is truly the first of a new generation of $S$-100 bus equipment . . a a generation that's designed to accomodate multi-user setups and other high level industrial, scientific, and commercial applications.
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These advanced features give you the power you need for future expansion, as well as the system flexibility that comes from superior design . . . but perhaps best of all, the price is competitive with boards that do a whole lot less. The Z-80A CPU board is available for $\mathbf{\$ 2 2 5}$ unkit, $\mathbf{\$ 2 9 5}$ assembled, and $\$ 395$ for units qualified under our Certified System Component program.


Next month: The second new generation CPU board, featuring our 8085/8088 dual processing technique, as well as the Spectrum color graphics board. Please do not call for information on these products; sendan SASE and we'll send the information to you.

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| Dig Grp | $\$ 319$ | $\$ 379$ | n/a |
| S-100 | $\$ 449$ | $\$ 499$ | $\$ 599$ |
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| S-100 | $\$ 599$ | $\$ 689$ | $\$ 789$ |
| S-100 (2) | $\$ 649$ | $\$ 729$ | $\$ 849$ |
| H8 (3) | $\$ 649$ | $\$ 749$ | $\mathrm{n} / \mathrm{a}$ |
| Dig Grp | $\$ 599$ | $\$ 679$ | $\mathrm{n} / \mathrm{a}$ |
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